

Theory Assignment no. 3. CS 22203

Ex. 1. Refer to the recursive procedure that follows the Dynamic Programming (DP) backward approach discussed in class for solving the 0/1 Knapsack problem. Draw the recursion tree depicting the implementation of the procedure for the following instance. Give the solution along with the total profit.

$n=3$, $(p_1, p_2, p_3)=(10, 20, 15)$, $(w_1, w_2, w_3)=(3, 4, 7)$, $M=7$.

Ex. 2 Let $G=(V, E)$ be a directed graph with nodes v_1, v_2, \dots, v_n . We say that G is an ordered graph if it has the following properties:

Each edge goes from a node with a lower index to a node with a higher index. That is, every directed edge has the form (v_i, v_j) with $i < j$.

Each node except v_n has at least one edge leaving it. That is, for every node v_i , $i=1, 2, \dots, n-1$, there is at least one edge of the form (v_i, v_j) . The length of a path is the number of edges in it. The goal in this question is to solve the following problem:

Given an ordered graph G , find the length of the longest path that begins at v_1 and ends at v_n .

a) Show that the following algorithm does not correctly solve this problem by giving an example of an ordered graph on which it does not return the correct answer.

Set $w = v_1$

Set $L = 0$

While there is an edge out of the node w

 Choose the edge (w, v_j) for which j is as small as possible

 Set $w = v_j$

$L = L+1$

end While

Return L as the length of the longest path.

In your example, say what the correct answer is and also what the algorithm above finds.

b) Give an efficient Dynamic Programming algorithm that takes an ordered graph G and returns the length of the longest path that begins at v_1 and ends at v_n . [Hint: First, derive a solution (i.e., a recursive function) that exploits the principle of optimality (as was done for I/O knapsack in class) and build the algorithm from it].
