# CS241 Coursework Report – 2008321

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# 1 – Implementation

#### 1.1 - Design

The application's runtime is split into three main phases: setup, packet processing and cleanup. In the setup phase, data structures (a queue for packets, and a binary tree for unique Ipv4 addresses) are initialised and allocated onto the heap, and the pool of worker threads is created. The threads are initialised to wait idly on a condition variable signal from the main thread. During packet processing, the main thread continually executes  $pcap\_loop()$  [1]. Packets are intercepted at the network interface, each triggering the callback function dispatch(). In a thread-safe manner, this adds a packet to the queue and signals the condition variable. A freshly awoken worker thread then dequeues and processes the packet. The final phase of runtime serves to deallocate all memory allocated to the heap, and to join all worker threads, collating the statistics from each to produce a final report to the user.

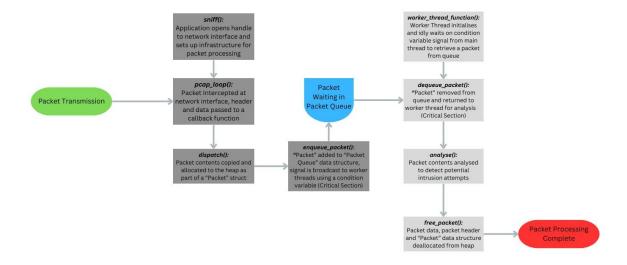


Figure 1: Packet Processing Flow

# 1.2 - Choice of Multithreading Model

Although a One Thread per X model is easier to implement, necessitating far less synchronisation, it presents two key issues. Firstly, if processing a large number of packets, the computational cost brought about by the continuous creation and destruction of threads grows to a problematic size, slowing application performance significantly. The other issue is the wasted CPU utilisation caused by constant context switches as the CPU must schedule time slices for a vast number of threads. The Threadpool Model avoids both these issues. It is far more scalable, enabling steady performance during the high-traffic bursts typical of internet traffic. The drawback, however, is the added implementation complexity due to the greater need for synchronisation and need for a queue data structure to manage packets. The flowchart below illustrates the workflow of a worker thread.

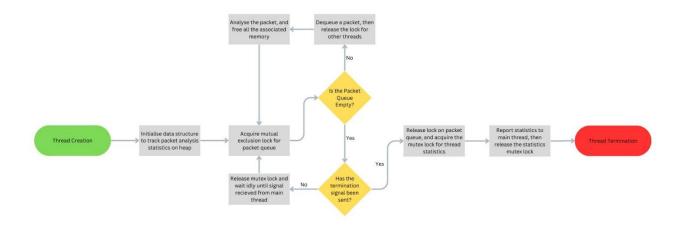


Figure 2: Worker Thread Execution Flow

## 1.3 – Synchronisation

A major design philosophy of the application was to minimise the need for synchronisation, as threads waiting on each other are likely to slow performance [2]. However, synchronisation primitives were implemented in four key places. Firstly, both a mutual exclusion lock and a condition variable are used for the packet queue. The mutex ensures packets are processed exactly once and are not lost while being added, avoiding potential memory leaks. The condition variable saves CPU resources by preventing busy waiting when the packet queue is empty. The second lock is used to synchronise printing to the terminal.

Another lock was used to synchronise search and insertion operations on the IP tree. A tree per thread design was considered, however, this would have necessitated the merging of many binary search trees into one while ensuring nodes were unique, a process which would have been both difficult to implement and computationally costly. Furthermore, each individual insertion operation has only logarithmic time complexity in average case, and analysis for most packets will not enter this critical section. For these reasons, a global tree was used. The final lock was used to safely report statistics to the main thread, preventing race conditions and consequent errors in the final report.

Deadlock is prevented by design, and therefore thread progress is ensured. The execution flow of the program ensures each thread is only waiting on, or allocated, at most one mutually excludable resource at any given time, preventing a circular wait condition and therefore deadlock from ever

occurring. More formally, in terms of a resource allocation graph, the implementation ensures that all vertices denoting processes are either isolated or pendant (have degree 0 or 1 respectively), and therefore cannot be part of a cycle, in which all vertices must have degree at least two. It follows that no cycles are present in the graph. Bounded waiting is achieved by internal fairness mechanisms in the Linux implementation of the *pthreads* library [3], which attempt to grant a mutex lock to the process which has been waiting the longest.

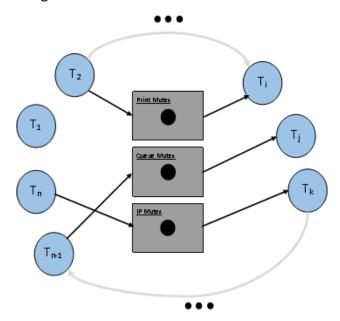


Figure 3: Resource allocation graph, with an arbitrary enumeration of vertices, representative of a typical state during packet processing

### 1.4 - Detecting Malicious Packets

The analysis of packets is done by using the layers of the OSI model successively to examine packet headers in detail. Firstly, the link layer ethernet header is inspected to determine the protocol for the network layer above. In the case of an ARP packet, the operation type is checked to spot and count ARP Replies. Alternatively, if the network layer protocol is IP and the transport layer protocol is TCP, the control flags in the TCP header are inspected to detect SYN packets. If one is found, a tree search/insert method is used to track unique IP addresses SYN packets have been sent from. After this, we check the destination port of the packet, to spot HTTP transmissions bound for port 80, and check the http request for attempts to access blacklisted URLs.

# 2 - Testing

#### 2.1 - Correctness

The first round of testing had the principal aim of ensuring application correctness. Firstly, The VSCode debug feature and *gdb* were used for debugging. Following this, *Valgrind* was used to check for memory leaks, with the addition of *Helgrind* [4] for multithreading errors.

Figure 4: Valgrind report with helgrind enabled, after running packet processing for a short time on eth0

Terminal commands such as *hping* and *wget* were utilised to send packets over the *eth0* interface to initially test the software's capability to detect malicious packets. Another utilised tool was the packet analysis software *tcpdump*, which was used to confirm the absence of background SYN packets and validate the results of our own software.[5].

```
root@cs241:-/cs241/skeleton/src# tcpdump -w -/output.pcap -1 eth0
tcpdump: listening on eth0, link-type ENIOMB (Ethernet), snapshot length 262144 bytes

'C125 packets captured
126 packets received by filter
0 packets dropped by kernel
root@cs241:-/cs241/skeleton/src# tcpdump -r -/output.pcap 'tcp[13] = 2' | wc -l
reading from file /root/output.pcap, link-type ENIOMB (Ethernet), snapshot length 262144
100
root@cs241:-/cs241/skeleton/src# |

loo packets transmitted, 0 packets received, 100% packet loss
round-trip min/avg/max = 0.0/0.0/0.0 ms
root@cs241:-/cs241/skeleton/src# |

source NK: 32:35-00:00:02:02

Destination NAC: 32:35-00:12:34:56
Type: 0

source NK: 32:35-00:00:12:34:56
Type: 0

source NK: 32:35-00:00:12:
```

Figure 5: Terminal snippets using tcpdump to confirm correct SYN packet detection on eth0

For precise testing with no background packets, the loopback interface was utilised. To ensure all three attack types could be detected here, the etc/hosts file on the virtual machine was modified to redirect requests to the public IPs of the two blacklisted websites to 127.0.0.1 (localhost) [6]. While the application was being run concurrently, a custom python script was used to send a random barrage of intrusion attempts.

```
| [U200832]@viglab-28 cs241]$ ssh -p 3071 root@localhost |
| Linux cs241 5.10.0-8-and064 #1 SWP Debian 5.10.46-4 (2021-08-03) x86_64 |
| Linux cs241 5.10.0-8-and064 #1 SWP Debian 5.10.46-4 (2021-08-03) x86_64 |
| Linux cs241 5.10.0-8-and064 #1 SWP Debian 5.10.46-4 (2021-08-03) x86_64 |
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| Linux cs241 5.10.0-8-and064 #1 SWP Debian 5.10.46-4 (2021-08-03) x86_64 |
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| Linux cs241 5.10.0-8-and064 #1 SWP Debian 5.10.46-4 (2021-08-03) x86_64 |
| Linux cs241 5.10.0-8-and064 #1 SWP Debian 5.10.46-4 (2021-08-03) x86_64 |
| Linux cs2
```

Figure 6: Terminal output from testing with custom python script on loopback interface

# 2.2 - Optimisation

Secondly, care was taken to optimise the program with respect to runtime. An earlier design used many more mutual exclusion locks to update counts every time a packet was analysed, but it was discovered runtime could be saved by each thread having their own set of counts, which could then be collated after packet processing had concluded, where speed was less critical. Another question in this area related to the optimum size of the thread-pool [7]. By modifying one of the arguments of pcap\_loop(), and using the time.h library, we were able to analyse the runtime of processing 100 packets while varying thread-pool size in increments of five. The results suggested an optimum threadpool size of 10.

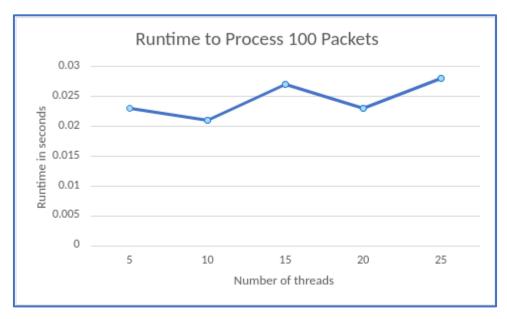


Figure 7: Experimental bivariate analysis on the effect of thread-pool size on runtime

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   <a href="https://documentation.agilepoint.com/8010/admin/sizingThreadPoolSizingGuidelines.html">https://documentation.agilepoint.com/8010/admin/sizingThreadPoolSizingGuidelines.html</a>

   (Accessed: 4 December 2023).

# Appendix: Python Script for testing on loopback interface