

1. Explain what you think the worst-case, Big-O complexity and the best-case, Big-O complexity of merge sort is. Why do you think that?

The worst-case time complexity of merge sort is $O(n \log n)$, and the best-case time complexity is also $O(n \log n)$. This is because merge sort always divides the input into halves and then recursively sorts them. The merging step takes $O(n)$ time for each level of recursion, and there are $\log n$ levels of recursion. Therefore, the total time complexity is $O(n \log n)$ in both the worst and best cases.

2. Merge sort, as we have implemented it, has a recursive algorithm embedded in the code as this is the easiest way to think about it. It is also possible to implement merge sort iteratively:

```
// Iteratively sort subarray `A[low...high]` using a temporary array
void mergesort(int A[], int temp[], int low, int high)
{
    // divide the array into blocks of size `m`
    // m = [1, 2, 4, 8, 16...]

    for (int m = 1; m <= high - low; m = 2 * m)
    {
        // for m = 1, i = 0, 2, 4, 6, 8...
        // for m = 2, i = 0, 4, 8...
        // for m = 4, i = 0, 8...
        // ...

        for (int i = low; i < high; i += 2*m)
        {
            int from = i;
            int mid = i + m - 1;
            int to = min(i + 2*m - 1, high);
            merge(A, temp, from, mid, to);
        }
    }
}
```

Explain what you think the worst-case, Big-O complexity and the best-case, Big-O complexity is for this iterative merge sort. Why do you think that?

The worst-case time complexity of the iterative merge sort is $O(n \log n)$, which is the same as the recursive version. The best-case time complexity is also $O(n \log n)$, which occurs when the array is already sorted. This is because the algorithm always divides the array into halves, and then merges them in linear time. The outer loop iterates $\log n$ times, and the inner loop iterates n times, resulting in a total time complexity of $O(n \log n)$.