

Algorithms Review





Introduction

- **Insertion Sort Analysis**
 - Proving correctness using loop-invariants
 - RAM model of computation
 - We discussed how we are normally only interested in growth of running time:
 - » **Best-case linear in $O(n)$, worst-case quadratic in $O(n^2)$**
- **Divide and Conquer**
 - General Method
 - Merge Sort
- **Analysis of divide-and-conquer algorithms**
 - Recurrence Equation
 - Merge Sort Analysis
 - » **Proof by Picture of Recursion Tree, Telescoping, Induction**
 - » **$\Theta(n \lg n)$ for all cases**



Asymptotic Growth

- **Asymptotic Growth**
 - **O-notation**
 - **Ω -notation**
 - **Θ -notation**
- **Recurrences**
 - **Substitution method**
 - **Recursion-tree method**
 - **Master method**



Sorting

- **Heap Sort**
 - **Heaps**
 - **MAX-HEAPIFY**
 - **BUILD-MAX-HEAP**
 - **HEAPSORT**
- **Priority Queues**
 - **MAXIMUM(S)**
 - **EXTRACT-MAX(S)**
 - **INCREASE-KEY(S,x,k)**
 - **INSERT(S, x)**
 - **Applications**



Sorting

- **QuickSort**
 - Divide and Conquer
 - PARTITION
 - Analysis of QuickSort
 - » Worst case $\Theta(n^2)$
 - » Expected running time: $\Theta(n \lg n)$
 - Randomized quicksort
- **The best worst-case running time for comparison sorting**
 - Decision-tree Model
 - **Theorem.** Any comparison sorting algorithm requires $\Omega(n \lg n)$ comparisons in the worst case



Sorting in linear time

- **Counting sort**

- How if there are *17* elements not greater than *x* in A?
 - » Put the *last one* in position *17*, the *penultimate one* in position *16*,...
- If $k = O(n)$, then counting sort takes $\Theta(n)$ time.

- **Radix sort**

- Sort on **least-significant digit** first with auxiliary **stable** sort.
- $\Theta(d T(n))$, if we use **counting sort** and d is constant, $\Theta(n)$

- **Bucket sort**

- Divide the interval $[0,1)$ into n equal-sized subintervals, or bucket, and then distribute the n input number into the buckets.
- $\Theta(n)$ under uniform distribution



Medians and Order Statistics

- **Order Statistics**

- Expected linear time selection

- » **Main idea: PARTITION**

- » **Worst-case $\Theta(n^2)$**

- Worst-case linear time selection

- » **Generate a good pivot recursively.**



Other D&C

- **Square Matrix Multiplication**
 - Simple Divide and Conquer algorithm
 - » $\Theta(n^3)$
 - Strassen's algorithm
 - » 8 recursive multiplications of $n/2 \times n/2$ matrices to only 7.
 - » $\Theta(n^{\lg 7})$
- **The maximum-subarray problem**
 - Stock Investing Problem
 - Divide and Conquer algorithm
 - » 3 sub-problems, one of which is not recursive
 - » $\Theta(n \lg n)$



Dynamic Programming

- **Assembly Lines**

$$f_1[j] = \begin{cases} e_1 + a_{1,1} & \text{if } j = 1 \\ \min(f_1[j-1] + a_{1,j}, f_2[j-1] + t_{2,j-1} + a_{1,j}) & \text{if } j \geq 2 \end{cases}$$
$$f_2[j] = \begin{cases} e_2 + a_{2,1} & \text{if } j = 1 \\ \min(f_2[j-1] + a_{2,j}, f_1[j-1] + t_{1,j-1} + a_{2,j}) & \text{if } j \geq 2 \end{cases}$$

- **Matrix-chain multiplication**

$$m[i, j] = \begin{cases} 0 & \text{if } i = j \\ \min_{i \leq k < j} \{ m[i, k] + m[k + 1, j] + p_{i-1}p_kp_j \} & \text{if } i < j \end{cases}$$

- **Elements of DP Algorithms**

- Optimal Substructure
- Overlapping Subproblem
- Reconstructing an optimal solution



Dynamic Programming

- Memoization
- Longest Common Subsequence

$$c[i, j] = \begin{cases} c[i-1, j-1] + 1 & \text{if } i, j > 0 \text{ and } x_i = y_j \\ \max(c[i, j-1], c[i-1, j]) & \text{otherwise} \end{cases}$$

- Max Sum
 - $b[j] = \max(b[j-1] + a[j], a[j])$, $1 \leq j \leq n$.



Greedy

- **Activity-selection problem**
 - **Optimal Substructure**
 - **Greedy algorithm**
- **Elements of the greedy strategy**
 - **Greedy-choice property and Optimal substructure**
 - **Fractional Knapsack Problem**
 - **0-1 Knapsack Problem (dynamic programming)**
- **Huffman codes**



Single-source shortest paths

- **Single-source shortest paths**
 - **Nonnegative edge weights**
 - » **Dijkstra's algorithm: $O(E + V \lg V)$**
 - **General**
 - » **Bellman-Ford: $O(VE)$**



All-pairs shortest paths

- All-pairs shortest paths (similar to matrix multiply)
 - $d^{(m)}_{ij} = \min_k \{d^{(m-1)}_{ik} + a_{kj}\}$
 - $(\min, +)$ multiplication
 - Improved matrix multiplication algorithm.
- Floy-Warshall algorithm
 - $c_{ij}^{(k)} = \min_k \{c_{ij}^{(k-1)}, c_{ik}^{(k-1)} + c_{kj}^{(k-1)}\}$
- Johnson's algorithm
 - Graph reweighting
 - » $h: V \rightarrow \mathbb{R}$, reweight $(u, v) \in E$ by $w_h(u, v) = w(u, v) + h(u) - h(v)$.
 - Algorithm:
 - » Find a function $h: V \rightarrow \mathbb{R}$, such that $w_h(u, v) \geq 0$ for all $(u, v) \in E$ by using Bellman-Ford
 - » using w_h from Run Dijkstra's algorithms each vertex $u \in V$
 - » For each $(u, v) \in V * V$, compute $\delta(u, v) = \delta_h(u, v) - h(u) + h(v)$



Back Tracking

- **Back Tracking Paradigm**

- A design technique, like divide-and-conquer.
- Useful for optimization problems and finding feasible solutions.
- Constraints
 - » **Explicit Constraints**
 - » **Implicit Constraints**
- General Method
 - » **Identify space state tree → Generate problem state → Is solution state? → Is answer state?**
 - » **Searching in DFS way**

- **N Queen problem**

- Algorithm
- Bound function



Branch and Bound Algorithms

- General Method
- Least Cost Search
 - $\hat{c}(X): \hat{c}(X) = f(h(X)) + \hat{g}(X)$



Computability

- **Computability**
 - **Undecidable Problem**
 - » **Hilbert's 10th Problem.**
 - » **Post's Correspondence Problem.**
 - » **Halting Problem.**
 - **NP-Complete**
 - » **P, EXP, NP, NP-Complete**
 - » **SAT is the first NP-Complete problem**
 - » **Other NP-Complete problems: 3-color, TSP, ...**



Sample Problem

- **Answer T/F for the following:**
 - Because inserting a key into a binary heap requires ($\lg n$) time in the worst case, building a heap of size n from scratch requires ($n \lg n$) time in the worst case.
- **Single Choice**
 - The worst-case running time of Insertion Sort is ()
 - A. $\Theta(n^2)$
 - B. $\Theta(n \lg n)$
 - C. $\Theta(n)$
 - D. $\Theta(n^3)$
- **Evaluate the following recursions**
 - $T(n) = T(7n/8) + n$



Sample Problem

- Find a maximum-length common subsequence of X and Y with sequence $X = \langle A, B, C, B, D, A, B \rangle$ and $Y = \langle B, D, C, A, B, A \rangle$.
- **Comprehensive**
 - Strategy?
 - Sorting?
- **Design?**