

Laboratory practice No. 5

Divide to conquer and Dynamic programming.

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3) Practice for final project defense presentation

3.1 The data structure that the program use is a matrix of costs and another matrix of visited vertices and the algorithm travel through the matrix, and the horizontal index of the vertex will represent the power set of the vertices not including the origin vertex, which is the final vertex too. While the algorithm creates previous values of the costs of going from one vertex to another going through other. So, a lot of cost depends on the previous values, so it can return and take them without doing the process again.

3.2 The program is going to have to do n times the power set of a set made up of $n-1$ elements, so it is going to be $n \cdot (2^{n-1})$

3.3 The algorithm works searching the shortest path that the delivery man has to cross for going by every vertex in the graph but obviously in the reality the delivery man took shortcuts, so the algorithm must have as vertices the points of delivery and the distance between two points of delivery is calculated directly if there isn't another vertex between them, or searching which path is the shortest going by another point.

3.4 The data structure used is a matrix of Strings, that will be fill by "K" that is the position of Karolina, "R" is the position of the radioactive waste, and numbers that are the distance from Karolina to that exactly point. So, the algorithm completes that number considering the position of Karolina while she gets every waste and accumulate the distance that Karolina cross.

3.5 The complexity of the algorithm is $O((x \cdot y) \cdot R)$

3.6 Being x and y the dimensions of the board and R the number of radioactive.

4) Practice for midterms

4.1.1

			C	A	L	L	E
		0	1	2	3	4	5
C	1	0	1	2	3	4	
A	2	1	0	1	2	3	
S	3	2	1	1	2	3	
A	4	3	2	2	2	3	

4.1.2

			M	A	D	R	E
		0	1	2	3	4	5
M	1	0	1	2	3	4	
A	2	1	0	1	2	3	
M	3	2	1	1	2	3	
A	4	3	2	2	2	3	

4.2.1 $O(X*Y)$, siendo X y Y la longitud de cada cadena

4.2.2 return table[lenx][leny]

4.3.1 a)

4.3.2 a)

4.4 c)

4.5.1 c)

4.5.2 a[mitad]

4.5.3 a, mitad+1, der, z

4.6.1 scm[i] = arr[i]

4.6.2 max++

4.6.3 max=scm[i]

4.6.4 $O(n^2)$

4.7.1 d[i][j]

4.7.2 d[l, k-1]

4.7.3 d[l-1, j]

4.7.4 $O(n^2)$

5) Recommended reading (optional)

6) Team work and gradual progress (optional)

Im alone ahah :(

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