

MEET TREES STUFF

1. SOME OBSERVATIONS ABOUT TREES

$\mathcal{O} = \{x_0, x_1, \dots\}$ is an orbit of a finite partial automorphism p of a \wedge -tree (so $x_j = p^j(x_0)$).

Remark 1.1. If for some i, j we have $x_i \geq x_j$, then whenever (for some $k \in \mathbf{Z}$) x_{i+k}, x_{j+k} exist, then $x_{i+k} \geq x_{j+k}$.

In particular, if x_{2j-i} exists, then $x_j \geq x_{2j-i}$, so $x_i \geq x_{2j-i}$, and likewise, if $x_{j+k(j-i)}$ exists, then $x_i \geq x_{j+k(j-i)}$. \diamond

Proposition 1.2. *Suppose (in a \wedge -tree) $x_1 \leq x_2$ and $y \not\geq x_1$. Then $x_1 \wedge y = x_2 \wedge y$.*

Proof. Trivially, $x_1 \wedge y \leq x_2 \wedge y$.

On the other hand, $x_2 \wedge y < x_2$, so $x_2 \wedge y \not\geq x_1$. Furthermore, since $y \geq x_2 \wedge y$ and $y \not\geq x_1$, we have $x_2 \wedge y \not\geq x_1$. Thus $x_2 \wedge y \leq x_1$, whence $x_2 \wedge y \leq x_1 \wedge y$. \square

The following proposition describes the “spiral” behaviour of orbits.

Proposition 1.3. *Suppose $n > 0$ is minimal such that x_0 is comparable to x_n . Then:*

- (1) x_i is comparable to x_j if and only if $i \equiv j \pmod{n}$,
- (2) if $i_1 \equiv i_2 \not\equiv j \pmod{n}$, then $x_{i_1} \wedge x_j = x_{i_2} \wedge x_j$.

Proof. (1): \Leftarrow is an easy consequence of Remark 1.1. For \Rightarrow , suppose this is not true. Then we can choose $i \not\equiv j \pmod{n}$ such that x_i is comparable to x_j . Choose i, j such that $|i - j|$ is minimal.

Note that we may assume without loss of generality that $i < j$, and then furthermore, that $i = 0$. Thus j is the smallest number such that $n \nmid j$ and x_j is comparable to x_0 . Necessarily (by choice of n), $j > n$.

Notice also that x_n and x_j cannot be comparable: otherwise, x_0 and x_{j-n} are comparable, which contradicts minimality of j .

Now, we have one of the three:

- $x_n \geq x_0$, in which case $x_{n+j} \geq x_n, x_j$,
- $x_n < x_0$ and $x_0 \leq x_j$, or
- $x_n, x_j < x_0$.

In all three cases, x_n and x_j are comparable, a contradiction.

(2): Since x_0 and x_n are comparable, it follows that x_{i_1} and x_{i_2} are also comparable. Suppose without loss of generality that $x_{i_1} \geq x_{i_2}$.

Since of course $x_{i_1} \wedge x_j \leq x_{i_1}$, we have that x_{i_2} and $x_{i_1} \wedge x_j$ are comparable.

If $x_{i_2} \geq x_{i_1} \wedge x_j$, then $x_{i_2} \wedge x_j \geq x_{i_1} \wedge x_j$, and since $x_{i_1} \geq x_{i_2}$, we also have $x_{i_1} \wedge x_j \geq x_{i_2} \wedge x_j$, so we are done.

Otherwise, $x_{i_2} < x_{i_2} \wedge x_j$, so in particular, $x_{i_2} < x_j$, which contradicts (1), so we are done. \square

Proposition 1.4. *Suppose that neither $x_0 \wedge x_1 < x_1 \wedge x_2$ nor $x_0 \wedge x_1 > x_1 \wedge x_2$.*

Then $x_0 \wedge x_1$ is a fixed point, and whenever x_{i+1} exists, $x_0 \wedge x_1 = x_i \wedge x_{i+1}$. Moreover, $x_0 \wedge x_1 = \bigwedge \mathcal{O}$, i.e. it is the smallest element of the semilattice generated by \mathcal{O} .

More generally, if $x_0 \wedge x_k$ is not strictly greater nor strictly smaller than $x_k \wedge x_{2k}$, then, then $x_k \wedge x_{2k}$ is a fixed point of p^k and it is minimal in the \wedge -lattice generated by $\{x_{nk} \mid n \in \mathbf{N}\}$.

Proof. Note that $x_0 \wedge x_1, x_1 \wedge x_2 \leq x_1$, so they are comparable. By the assumption, $x_0 \wedge x_1 = x_1 \wedge x_2$. It follows immediately that $x_0 \wedge x_1$ is a fixed point, so for every i , $x_i \wedge x_{i+1} = x_0 \wedge x_1$ (as long as x_{i+1} exists).

But then for $i > 1$ we have $x_i \wedge x_0 \wedge x_1 = x_i \wedge x_i \wedge x_{i-1} = x_i \wedge x_{i-1} = x_0 \wedge x_1$. Since $\bigwedge \mathcal{O} = \bigwedge_{i=0}^m x_i$, the conclusion follows.

The “more generally” part is a special case with the partial automorphism p^k in place of p . \square

For each $j, k \in \mathbf{N}$, put $y_j^k := x_j \wedge x_{j+k}$. Notice that $p^l(y_j^k) = y_{j+l}^k$.

Proposition 1.5. *Suppose $n > 0$ is minimal such that y_0^k is comparable to y_n^k (in particular, suppose that such n exists). Then n divides k .*

Proof. We have the following claim.

Claim. y_0^k and y_k^k are comparable.

Using Claim, by Proposition 1.3(1) applied to y_i^k with $i = 0$ and $j = k$, we conclude that $0 \equiv k \pmod{n}$, i.e. n divides k . Thus, we only need to prove the claim.

Proof of claim. Suppose that $y_0^k \geq y_n^k$ (the other case is analogous). Then we have:

$$y_0^k \geq y_n^k \geq y_{2n}^k \geq \dots \geq y_{kn}^k. \quad (1)$$

Note that immediately by definition of y_j^k , $x_k \geq y_0^k, y_k^k$. Thus y_0^k and y_k^k are comparable. As in the preceding paragraph, we conclude that if $y_0^k < y_k^k$, then $y_0^k < y_{nk}^k = y_{kn}^k$. This contradicts (1), so we must have $y_0^k \geq y_k^k$. \square (claim)

\square

Proposition 1.6. *Suppose $k \in \mathbf{N}$ is minimal such that for some $n > 0$, y_0^k is comparable but not equal to y_n^k . Suppose in addition that $k > 0$.*

Then k divides every such n , and the minimal n is equal to k .

Proof. We have the following claim.

Claim. If $n = 1$, then $k = 1$.

Proof. The proof is by contraposition. Suppose that $k > 1$. Then for every n , $y_0^1 = x_0 \wedge x_1$ is either equal or incomparable to $y_n^1 = x_n \wedge x_{n+1}$.

In particular, y_0^1 is neither greater nor smaller than y_1^1 . By Proposition 1.4, it follows that $y_0^1 \wedge y_1^1 = y_1^1 \wedge y_2^1$. In particular, $n > 1$. \square

Suppose for now that n is minimal. Then by Proposition 1.5, n divides k . Put $k' := k/n$, consider the partial automorphism p^n , and put $x'_j = x_{nj}$, and for $y'_j = x_{nj} \wedge x_{nj+k} = x'_j \wedge x'_{j+k'}$.

Then we have that $y'_0 = x'_0 \wedge x'_{k'} = x_0 \wedge x_k$ is comparable but not equal to $y'_1 = x'_1 \wedge x'_{1+k'} = x_n \wedge x_{n+k}$. Moreover, by minimality of k , it follows that k' is minimal such that for some n' the meets $x'_0 \wedge x'_{k'}$ and $x'_{n'} \wedge x'_{n'+k'}$ are comparable but not equal. By Claim it follows that $k' = 1$, i.e. the minimal n is k .

The case of arbitrary n follows from Proposition 1.3(1) (applied to y_i^k). \square

Definition 1.7. Let $\mathcal{O} = \{x_0, \dots, x_m\}$ be the orbit of a finite partial automorphism p (such that $p(x_i) = x_{i+1}$).

Then:

- if for some positive integer $n \leq m$ we have that $x_0 = x_n$, and n is minimal, then we say that \mathcal{O} is a *loop of length n* ;
- if for some positive integer n we have that x_0 is comparable (but not equal) to x_n , and n is minimal, then we say that \mathcal{O} is a *spiral of length n* ;
- if for some positive integer $k \leq m/2$, we have that x_0 and x_k are incomparable, but $x_0 \wedge x_k$ is not equal to $x_k \wedge x_{2k}$, and k is minimal, then we say that \mathcal{O} is a *comb spiral of length k* ;
- otherwise (if \mathcal{O} is not a loop, a spiral, nor a comb loop), we say that \mathcal{O} is a *pseudoloop*.

Remark 1.8. Note that if x_0, x_1, \dots, x_m is of comb type k , then $y_0^k, y_1^k, \dots, y_{m-k}^k$ is of spiral type k .

Conversely, it is not hard to see that if y_0, \dots is an orbit of p of spiral type k , then for some $\bar{p} \supseteq p$, there is a \bar{p} -orbit x_0, \dots of comb type k , such that $y_j = x_j \wedge x_{j+k}$. \diamond

Proposition 1.9. If \mathcal{O} is of finite or spiral type, then for every k , either $x_0 \wedge x_k = x_k \wedge x_{2k}$ or x_0 and x_k are comparable.

Proof. Suppose that \mathcal{O} is of finite or spiral type n . If n divides k , then trivially x_0 and x_{nk} are comparable. Otherwise, by Proposition 1.3, we have that $x_0 \wedge x_k = x_{kn} \wedge x_{kn+k}$.

But if $x_0 \wedge x_k < x_k \wedge x_{2k}$, then $x_0 \wedge x_k < x_{kn} \wedge x_{kn+k}$, and likewise if $x_0 \wedge x_k > x_k \wedge x_{2k}$, then $x_0 \wedge x_k > x_{kn} \wedge x_{kn+k}$. In both cases, we have a contradiction. But since $x_0 \wedge x_k, x_k \wedge x_{2k} \leq x_k$, they are compatible, so they must be equal. \square

Note that Proposition 1.9 implies that orbits of finite or spiral type are not of comb type.

Remark 1.10. Proposition 1.6 implies that \mathcal{O} is of comb type k if and only if k is minimal such that for some n , the meets $x_0 \wedge x_k$ and $x_n \wedge x_{n+k}$ are comparable. \diamond

Conjecture 1.11. *If \mathcal{O} is of open type, then it can be extended to an orbit of any other type.*

If \mathcal{O} is not of open type, then it is determined.

2. PSUEDOLOOPS ANALYSIS

In this section, we have a blanket assumption that \mathcal{O} is a loop or a pseudoloop, i.e. \mathcal{O} is an antichain and for every k , we have $x_0 \wedge x_k = x_k \wedge x_{2k}$.

Remark 2.1. If $x_0 \wedge x_k = x_k \wedge x_{2k}$, then for every integer i we have $x_0 \wedge x_k = x_{ik} \wedge x_{(i+1)k}$. \diamond

Proposition 2.2. *Suppose \mathcal{O} is a pseudoloop and k, n are positive integers.*

Then $x_0 \wedge x_k \leq x_0 \wedge x_{kn}$.

Proof. This is immediate from Proposition 1.4. \square

Proposition 2.3. *Let $d = (k_1, k_2)$ be the greatest common divisor of k_1 and k_2 . Then either $x_0 \wedge x_{k_1} = x_0 \wedge x_d$ or $x_0 \wedge x_{k_2} = x_0 \wedge x_d$.*

Proof. Let a, b be positive integers such that $ak_1 = bk_2 + d$. Then (by Remark 2.1)

$$(x_0 \wedge x_{ak_1}) \wedge (x_0 \wedge x_{bk_2}) \leq x_{ak_1} \wedge x_{bk_2} = x_0 \wedge x_d.$$

On the other hand, by Proposition 1.4, $x_0 \wedge x_d \leq x_0 \wedge x_{ak_1}$ and $x_0 \wedge x_d \leq x_0 \wedge x_{bk_2}$, thus

$$(x_0 \wedge x_{ak_1}) \wedge (x_0 \wedge x_{bk_2}) = x_0 \wedge x_d$$

Since $x_0 \geq x_0 \wedge x_{ak_1}, x_0 \wedge x_{bk_2}$, the latter two are comparable, so

$$x_0 \wedge x_d = (x_0 \wedge x_{ak_1}) \wedge (x_0 \wedge x_{bk_2}) = \min(x_0 \wedge x_{ak_1}, x_0 \wedge x_{bk_2}),$$

which completes the proof. \square

Define a sequence $(n_i)_i$ in the following way: $n_0 = 1$, while n_{i+1} is the smallest positive integer n such that $x_0 \wedge x_n > x_0 \wedge x_{n_i}$ (if it exists).

Remark 2.4. Immediately by the definition, for every k , there is some i such that $x_0 \wedge x_k = x_0 \wedge x_{n_i}$. \diamond

Proposition 2.5. *For each i , $x_0 \wedge x_k \geq x_0 \wedge x_{n_i}$ if and only if n_i divides k .*

It follows that $(n_i)_i$ is strictly \downarrow -increasing and (by Remark 2.4) that for $n(k)$ being the largest n_i such that n_i divides k , we have $x_0 \wedge x_k = x_0 \wedge x_{n(k)}$.

Proof. \Leftarrow is immediate by Proposition 1.4.

For \Rightarrow , note that if $d = (k, n_i)$, then by Proposition 2.3, we have $x_0 \wedge x_d = x_0 \wedge x_{n_i}$. By minimality of n_i , it follows that $d = n_i$, so n_i divides k . \square

Proposition 2.6. *For each i and positive $m < n_i$, we have $x_0 \wedge x_{n_i} \neq x_m \wedge x_{m+n_i}$.*

Proof. Suppose $x_0 \wedge x_{n_i} = x_m \wedge x_{m+n_i}$. Then $x_0 \wedge x_{n_i} \leq x_m$, so $x_0 \wedge x_{n_i} \leq x_0 \wedge x_m$. But then by Proposition 2.5, n_i divides m , so $m \geq n_i$. \square

Corollary 2.7. *If y is in the \wedge -lattice generated by \mathcal{O} , $k \leq 2n$ and $y < x_0$, then $y \geq x_0 \wedge x_k$ if and only if $p^k(y) = y$.*

Proposition 2.8. *Consider a partial orbit x_0, \dots, x_{n+1} .*

Let i_0 be such that n_{i_0} is maximal not greater than n .

If n_{i_0} divides $n + 1$, then there are two ways to amalgamate x_0, \dots, x_n and x_1, \dots, x_{n+1} : either $x_0 \wedge x_{n+1} > x_0 \wedge x_{n_{i_0}}$ (so $n_{i_0+1} = n + 1$), or $x_0 \wedge x_{n+1} = x_0 \wedge x_{n_{i_0}}$.

Given any $k < n$, we have $x_{n+1} \wedge x_{n+1-k} = x_{n+1} \wedge x_{n+1-n(k)}$. If $n(k) < \sqrt{n}$, then $x_{n+1} \wedge x_{n+1-n(k)}$

Otherwise, if n_i does not divide $n + 1$, then