- 1. HEREDITARY G-COMPACTNESS, HEREDITARY BOREL CARDINALITY
- **Definition 1.1.** We say that a theory T is G-compact if  $\equiv_{\mathbf{L}}$  is type-definable, or equivalently, if the Galois group  $\operatorname{Gal}(T)$  is a Hausdorff group.
- **Definition 1.2.** We say that a structure is *G-compact* if its theory is *G-compact*.
- **Definition 1.3.** If M and  $N = (S_i, R_j, f_k, c_l)_{i,j,k,l}$  (where  $S_i$  are sorts,  $R_j$  are predicates,  $f_k$  are functions and  $c_l$  are constants) are structures (possibly in different languages), then we say that M interprets N if we can in  $M^{eq}$  definable sets  $S'_i$ , definable relations  $R'_j$ , definable functions  $f'_k$  and definable points  $c'_l$ , such that  $(S'_i, R'_i, f'_k, c'_l)_{i,j,k,l} \cong N$ .
- **Definition 1.4.** A theory T is said to be *hereditarily G-compact* if for every model  $M \models T$ , and every structure N interpreted by M (with parameters), N is G-compact.
- Remark 1.5. It is enough to consider any single  $|T|^+$ -saturated  $M \models T$  (the only way the choice of the model matters is in the realised types of the parameters used in the interpretation). Furthermore, one can check that T is hereditarily G-compact if and only if its restriction to every countable sublanguage is hereditarily G-compact. (?)
  - Fact 1.6. Every simple theory is G-compact, and every theory interpreted by a simple theory is simple. Therefore, every simple theory is hereditarily G-compact.
- **Example 1.7.** Let T be any non-G-compact theory. Then  $T^{Sk}$ , the Skolemization of T, is G-compact but not hereditarily G-compact (because it interprets T).
  - **Example 1.8.** Any o-minimal expansion of a group (with at least two definable points) is G-compact (because it has definable Skolem functions). On the other hand, we will see below that no theory which interprets dense linear orderings is hereditarily G-compact (this includes many o-minimal theories, such as all expansions of the real field).
  - Question 1.9. Is simplicity equivalent to hereditary G-compactness?
  - Question 1.10. Is NSOP equivalent to hereditary G-compactness?
- **Definition 1.11.** Given two posets  $P = (P, <_P), (Q, <_Q)$ , the linear sum  $P \oplus Q$  is defined as  $(P \sqcup Q, <)$  where a < b if:
  - $a \in P$  and  $b \in Q$ , or
  - $a, b \in P$  and  $a <_P b$ ,
  - or  $a, b \in Q$  and  $a <_Q b$ .

Informally,  $P \oplus Q$  is the disjoint union of P and Q with Q put after P.

35 Remark 1.12. The linear sum  $\oplus$  is clearly associative.

**Theorem 1.13.** Suppose  $(P, \leq)$  is a partially ordered set such that the three natural embeddings of P into  $P \oplus P \oplus P$  are elementary. Then P is not hereditarily G-compact (even without parameters, and without imaginary sorts if one has sorts with arbitrarily many definable constants).

Proof. Let  $c_n^i$ ,  $n \in \mathbb{N}^+$ ,  $i = 0, \ldots, n-1$  be any (possibly imaginary) pairwise distinct, definable constants. Let  $P_n = \bigcup_{i=0}^{n-1} P \times \{c_n^i\}$ . Furthermore, define a cyclic order  $S_n$  on  $P_n$  by saying that  $S_n((p_1, c_n^i), (p_2, c_n^j), (p_3, c_n^k))$  whenever S(i, j, k) (where S is the standard cyclic ordering on  $\{1, \ldots, n\}$ ) or i = j and  $p_1 < p_2$  or j = k and  $p_2 < p_3$  or k = i and  $p_3 < p_1$ . Finally, put  $R_n(p, a_n^i) := (p, a_n^{i+1})$  (with addition modulo n in the upper index).

Put  $\mathcal{P} = (P_n, S_n, R_n)_{n \in \mathbb{N}^+}$ . Clearly P interprets  $\mathcal{P}$  (without parameters). We will show that  $\mathcal{P}$  is not G-compact.

Denote by  $\mathcal{P}^*$  a model constructed analogously to  $\mathcal{P}$ , only with P replaced by its monster model extension  $P^*$ . It is easy to see that then  $\mathcal{P} \leq \mathcal{P}^*$  (because they are interpreted in the same way from P and  $P^*$ , respectively).

Choose arbitrary  $p_0 \in P$ . We will show that for n > 1,

$$n/2 - 1 < d_L((p_0, c_n^0), (p_0, c_n^{\lfloor n/2 \rfloor}) < \infty.$$

This will complete the proof, as the diameter of  $[(p_0, c_n^0)_{n \in \mathbb{N}^+}]_{\equiv \mathbb{L}}$  will be unbounded. For the first inequality, note that if  $\bar{p}_1, \bar{p}_2 \in P_n^*$  have the same type over some  $\mathcal{M} \leq \mathcal{P}^*$ , we must have  $S_n(\bar{p}_1, \bar{p}_2, R_n(\bar{p}_1))$  or  $S_n(\bar{p}_1, \bar{p}_2, R_n^{-1}(\bar{p}_1))$  (because there are some points in  $\mathcal{M}$  between  $\bar{p}_1$  and each of  $R_n(\bar{p}_1)$  and  $R_n^{-1}(\bar{p}_2)$ ). In particular, the second coordinate can change by at most one step between  $\bar{p}_1$  and  $\bar{p}_2$ . Therefore, since  $c_n^0$  and  $c_n^{\lfloor n/2 \rfloor}$  differ by  $\lfloor n/2 \rfloor$ , we obtain the first inequality.

since  $c_n^0$  and  $c_n^{\lfloor n/2 \rfloor}$  differ by  $\lfloor n/2 \rfloor$ , we obtain the first inequality. For the second, more substantial step, by the assumption and saturation of  $P^*$ , we can find some  $P^-, P^+ \leq P^*$  such that  $P^- < P < P^+$  and  $P^- \cup P \cup P^+ \leq P^*$ , and such that there exists an  $f: P^- \cup P \cup P^+$  such that restrictions to  $P^-, P$  and  $P^+$  are isomorphisms with  $P, P^+$  and  $P^-$ , respectively.

Now, define partial functions  $\sigma_-, \sigma, \sigma_+ \colon \mathcal{P}^* \to \mathcal{P}^*$  by the formulas:

$$\sigma_{-}((p, c_n^i)) := \begin{cases} (f(p), c_n^i) & \text{if } p \in P^- \\ (p, c_n^i) & \text{if } p \in P^+ \end{cases}$$

$$\sigma((p, c_n^i)) := \begin{cases} (f(p), c_n^i) & \text{if } p \in P \\ (p, c_n^i) & \text{if } p \in P^- \end{cases}$$

$$\sigma_{+}((p, c_n^i)) := \begin{cases} (f(p), c_n^{i+1}) & \text{if } p \in P^+ \\ (p, c_n^i) & \text{if } p \in P \end{cases}$$

It is easy to check that of these is elementary partial maps in  $\mathcal{P}^*$  (because their domains and ranges are models of  $\text{Th}(\mathcal{P})$ , and they are clearly isomorphisms between those), so they extend to automorphisms of  $\mathcal{P}^*$ . Moreover, it is easy to see

that the extensions are all Lascar strong automorphisms, and  $\sigma_- \circ \sigma_+ \circ \sigma((p, c_n^i)) = (p, c_n^{i+1}).$ 

**Proposition 1.14.** If P is a model complete poset such that for any  $Q \equiv P$  we have  $P \oplus Q \equiv P$ , then P satisfies the hypothesis of Theorem 1.13.

5 *Proof.* By applying assumption for Q = P, we have  $P \equiv P \oplus P$ , and by applying it again for  $Q = P \oplus P$ , we get  $P \equiv P \oplus P \oplus P$ . The result follows immediately by model completeness.

**Proposition 1.15.** ( $\mathbf{Q}$ , <) and ( $\mathbf{Z}$ , <) satisfy the assumptions of Theorem 1.13, so they are not hereditarily G-compact (without parameters).

Consequently, every infinite linear ordering without endpoints which is discrete or dense is not hereditarily G-compact (without parameters).

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*Proof.* In both cases, the conclusion follows from a straightforward application of the Ehrenfeucht-Fraïssé games. For example, to check that the middle copy of  $\mathbf{Z}$  is elementarily embedded in  $\mathbf{Z} \oplus \mathbf{Z} \oplus \mathbf{Z} = \mathbf{Z}^{\oplus 3}$ , choose some an arbitrary finite sequence  $k_1, \ldots, k_n \in \mathbf{Z}$ , and enumerate  $\mathbf{Z}^{\oplus 3}$  as  $\mathbf{Z} \times \{1, 2, 3\}$ . We need to show that  $(\mathbf{Z}, <, k_1, \ldots, k_n) \equiv (\mathbf{Z}^{\oplus 3}, (k_1, 2), \ldots, (k_n, 2))$ . For an Ehrenfeucht-Fraïssé game of length N,

**Proposition 1.16.** If (L, <) is an  $\aleph_1$ -saturated infinite linear order, then there is an infinite definable set D such that  $(D, < \upharpoonright_D)$  is dense or discrete.

20 Proof. Notice that since (L, <) is  $\aleph_1$ -saturated, it follows immediately that given a definable  $D \subseteq L$ ,  $(D, < \upharpoonright_D)$  is also  $\aleph_1$ -saturated.

If L defines an infinite discrete linear order, we are done. So suppose it does not. We aim to show that there is a definable dense linear order.

Given a linear order (K, <), denote by S(K) the set of all immediate successors in K, and denote by P(K) the set of all immediate predecessors. It is clear that S(K) and P(K) are both definable in (K, <).

Claim 1: Let (K, <) be an arbitrary linear order. If I is a convex component of P(K) in K (i.e. a maximal subset of P(K) which is convex in K), then either I is finite or I contains arbitrarily long finite intervals.

Proof. If I is finite, we are done. Suppose, then, that I is infinite.

For brevity, given  $a \in S(K)$ , write a + (-1) or a - 1 for its predecessor, and given  $a \in P(K)$ , write a + 1 for its successor in K. In the same way, let us write a + k for all  $k \in \mathbf{Z}$  for the k-th successor (when it exists).

Take any  $a \in I$ , and consider the set  $S_a$  of a + k for all  $k \in \mathbf{Z}$  (for which this makes sense).

Note that  $S_a$  is convex and all elements of  $S_a$ , except the last one (if it exists) are contained in I.

Thus, if  $S_a$  is infinite, we are done (because we can find arbitrarily long finite intervals in  $I \cap S_a$ ). So suppose  $S_a$  is finite. Then it has a smallest element  $a_-$ 

and a largest element  $a^+$ . Note that this implies that  $a^+ \notin P(K)$ , so  $I \setminus S_a = I \setminus [a_-, a_+] = I \cap (-\infty, a_-)$ . Thus, because I is infinite, there is some  $b \in I$  such that  $b < a_-$ . But then for each  $n \in \mathbb{N}_{>0}$ , if b + n exists, then  $b + n \in (b, a_-) \subseteq I$  (otherwise, for some n we would have  $b + n = a_-$ , which is impossible, because  $a_-$  is not a successor). This implies that b + n + 1 exists. Thus by induction, for all  $n \in \mathbb{N}$ , b + n exists, and (b, b + n + 1) is an interval with n elements.  $\square$ (claim)

Now, if some convex component of P(L) contains arbitrarily long finite intervals, then by compactness, there is an infinite discrete interval, which contradicts the assumption from the second paragraph.

Otherwise, all convex components of P(L) are finite, so  $L' = L \setminus P(L)$  is an infinite discrete linear order: indeed, if there are finitely many convex components, then P(L) is finite, in which case L' is clearly infinite (because L is infinite). Otherwise, if  $L \setminus P(L)$  was finite, then  $P(L) = L \setminus (L \setminus P(L))$  would have only finitely many convex components, all of them finite. But then L itself would be finite, a contradiction.

Similarly, we show that all convex components of S(L') are finite, so  $L_1 = L' \setminus S(L')$  is an infinite linear order.

In the same way, we define  $L_n$  for each  $n \in \mathbb{N}_{>0}$  recursively as  $L_{n+1} := L_n \setminus P(L_n) \setminus S(L_n \setminus P(L_n))$ , and they are all infinite linear orders. There are two cases: either  $L_n$  stabilises (i.e. eventually  $S(L_n) = P(L_n) = \emptyset$ ), or it does not. In the former case,  $L_n$  is clearly dense.

Corollary 1.17. Every o-minimal expansion of a group is G-compact but not hereditarily G-compact.

**Proposition 1.18.** By an analgous construction, the theory of (infinite) discrete linear orders is not hereditarily G-compact.

**Proposition 1.19.** If  $(L, \leq)$  is a totally ordered set, L interprets a dense linear order or L interprets a discrete linear order.

**Proposition 1.20** (?). If  $P = (P, \leq)$  is a poset such that  $P \equiv P \sqcup P$  ordered in such a way that the first copy goes before the second copy, and P is model complete, then an analogous construction works (probably).

Remark 1.21. For the "circles" construction to work for a general poset P, we certainly need P to be upwards and downwards directed, without least or largest element.

Remark 1.22. If P is a poset, and b is an element such that b is not the supremum of some two elements, then  $(-\infty, b)$  is upwards-directed.

**Question 1.23.** Suppose P has the property that every element is the supremum of some two elements. Does this imply IP (or some other general property)? (Note: atomless Boolean algebras have this property, as do binary ordered trees (growing down). They both have IP and SOP.)

Question 1.24. Are atomless Boolean algebras hereditarily G-compact?

Remark 1.25. In atomless Boolean algebras, we can always eliminate the parameters (up to a definable bijection). From now on, we will assume that there are no parameters.

**Lemma 1.26.** Let T be the theory of atomless boolean algebras. Let  $\varphi(x_1, \ldots, x_n)$  be any formula such that  $\varphi(x^1, \ldots, x^n) \vdash \bigwedge_i p(x^i)$  where p is a complete  $\emptyset$ -type. Then we can assume without loss of generality (up to a  $\emptyset$ -definable bijection) that p is the type of a partition and  $\varphi(x^1, \ldots, x^n)$  is equivalent to a formula of the form  $\bigvee \bigwedge x_i^i \cap x_{i'}^{i'} [=/\neq] \emptyset$  (i.e. a set of permissible intersection tables).

*Proof.* The first part is obvious. The second part is just calculation (and q.e.).  $\Box$ 

Corollary 1.27. Atomless boolean algebras don't interpret infinite linear orders.

*Proof.* Suppose a formula  $\varphi(x,y)$  defines a linear preorder with infinitely many equivalence classes. By  $\omega$ -categoricity, one of the |x|-types intersects infinitely many of these classes, so we can assume without loss of generality that  $\varphi$  implies this type, and so, by the preceding lemma, that x is a partition.

Now, let x and y be arbitrary partitions of the appropriate size. Let z be a partition of the same size, independent of x and y. Then by linearity, either  $\models \varphi(x,z)$  or  $\models \varphi(z,x)$ . But the type of x over z is the same as the type of z over x, so in fact  $\models \varphi(x,z) \land \varphi(z,x)$ . For the same reason,  $\models \varphi(z,y) \land \varphi(y,z)$ , and hence by transitivity  $\models \varphi(x,y) \land \varphi(y,x)$ , so x and y are equivalent and, as they were arbitrary,  $\varphi$  defines a total relation on its domain, a contradiction.  $\square$ 

Remark 1.28. Suppose atomless boolean algebras interpret a nontrivially directed poset. Let  $\varphi(x,y)$  be the formula defining it. Then we can restrict  $\varphi$  to a single complete type in such a way that it still defines a nontrivially directed poset.

*Proof.* Note that a (countable) poset is nontrivially directed if and only if it has a cofinal chain. By pigeonhole principle, infinitely many members of the chain lie in a single type, and so we can apply the lemma.  $\Box$ 

Corollary 1.29. Suppose  $\varphi(x,y)$  defines a nontrivially directed poset on partitions of a given size. Let x be any partition, and let  $\bar{x}$  be the upper bound of all permutations of x. Then all row permutations of the intersection table of x and  $\bar{x}$  are permitted by  $\varphi$ .

Remark 1.30. Given a poset P, for some interval in P to be nontrivially directed, it is necessary that for all  $p_0, p_1 \in P$  there are some p, p' such that  $p_0 < p, p' < p_1$  and  $p_1$  is the supremum of p and p'.

**Example 1.31.** Define a poset P as follows. Let  $P_0 = \{p_{\min}, p_{\max}\}$ , where  $p_{\min} < p_{\max}$ . We construct  $P_{n+1}$  from  $P_n$  by adding, for each pair  $p_0, p_1 \in P_n$  of distinct points such that  $[p_0, p_1] = \{p_0, p_1\}$  in  $P_n$ , a pair of points p, p' such that  $p_0 < p, p' < p_1$  and p and p' are incomparable. Then put  $P = \lim_n P_n$  with inclusions.

Remark 1.32. The poset P in the preceding example has the property that no interval is nontrivially directed. In fact, all (open) intervals are isomorphic.

Remark 1.33. P is a non-distributive lattice.

5 Remark 1.34. (?) The poset P has q.e. in the language of bounded lattices (i.e.  $\cup, \cap, \leq, 0, 1$ ).

*Proof.* Let  $\varphi(x, \bar{a})$  be quantifier-free formula with parameters  $\bar{a}$ , where |x| = 1. We can assume without loss of generality that  $\bar{a}$  enumerates a sublattice of P.

Remark 1.35. Intervals in P have NIP.

10 Conjecture 1.36 (Well-known?). Every NIP unstable theory interprets an infinite linear order. (Note: every unstable weakly VC-minimal theory does.)

Question 1.37. Does an IP+SOP theory interpret an infinite linear order?

**Question 1.38.** For a poset  $(P, \leq)$ , how do the following properties relate:

- (1) P is directed (both ways?).
- (2) "circles with P" "work".
- (3) P has NIP.
- (4) P is nice with respect to VC-codimensions.

What if we only want them up to interpretation?

Question 1.39. Suppose P and Q are posets such that their induced circular orders are elementarily equivalent. Does the "circles" construction yield elementarily equivalent results?

**Example 1.40.**  $\omega \sqcup \omega^*$  yields the same thing as **Z**.

## 2. Adding parameters an G-compactness

**Example 2.1.** Consider an  $F = (F, \cdot, +, c)_{c \in \mathbf{Q}^{\text{alg}}}$  algebraically closed field of characteristic 0 with parameters for all algebraic numbers. Then the Shelah strong types and types over the empty set coincide, but for any transcendental t, the types over  $\{t\}$  do not coincide with strong types (e.g.  $\sqrt{t}$  and  $-\sqrt{t}$  have different strong types but the same type over t).

Question 2.2. Suppose T is G-compact (KP and L strong types coincide). If we add to the language some parameters for elements of a model of T, is the resulting theory still G-compact (i.e. do the KP and L strong types with parameters coincide)?

**Proposition 2.3.** Suppose G is a group definable in  $\mathfrak{C}$ . Let I be a large set (independent of  $\mathfrak{C}$ ). Consider a structure  $\mathfrak{C}' = (\mathfrak{C}, I \times G', I, \cdot, \pi)$ , where G' is a copy of G, the structure on  $\mathfrak{C}$  is standard, there is no internal structure on  $I \times G'$  or I,  $\cdot : G \times (I \times G') \to (I \times G')$  is the left action  $g \cdot (i, g') = (i, gg')$ , and  $\pi : I \times G' \to I$  is the standard projection.

Then  $\operatorname{Aut}(\mathfrak{C}') = G^I \rtimes (S_I \times \operatorname{Aut}(\mathfrak{C}))$ , where the action of  $S_I$  is obvious (and fixes G' coordinates and  $\mathfrak{C}$ ),  $\operatorname{Aut}(\mathfrak{C})$  acts on  $\mathfrak{C}$  in the usual way and permutes the G' coordinates in the standard manner (fixing I), while  $G^I$  permutes the G' coordinates by acting on the right coordinatewise (and fixes  $I, \mathfrak{C}$ ). Moreover,  $\operatorname{Autf}_L(\mathfrak{C}') = G^I \rtimes (S_I \times \operatorname{Autf}_L(\mathfrak{C}))$ , and similarly for Shelah and Kim-Pillay strong automorphism groups.

Meanwhile, for any  $i_0 \in I$  and  $g' \in G$  we have that the orbits of  $(i_0, g')$  via the action of automorphisms, strong Shelah, Kim-Pillay and Lascar automorphisms respectively are the same as their G,  $G^{00}$ ,  $G^{000}$  orbits, respectively.

*Proof.* The "meanwhile" part is standard, as once we add  $i_0$  as a parameter,  $\{i_0\} \times G'$  is just a principal homogenous space for G, which was analysed in detail in [?].

For the Aut( $\mathfrak{C}'$ ), first note that Aut( $\mathfrak{C}$ ) and  $S_I$  act naturally on  $\mathfrak{C}'$  (by acting on  $\mathfrak{C}$  part and permuting the G' coordinates in  $I \times G'$ , and by permuting the I coordinates in  $I \times G'$ , respectively), and the two actions commute. Thus it is easy to see that it is enough to show that Aut( $\mathfrak{C}'/I\mathfrak{C}$ ) =  $G^I$ . But from the same arguments as in the preceding paragraph, the automorphisms of  $\{i_0\} \times G'$  fixing  $\mathfrak{C}$  are the same as G acting on the right. Clearly, the actions on  $\{i\} \times G'$  and  $\{j\} \times G'$  are independent, and each of them must be preserved setwise by Aut( $\mathfrak{C}'/I\mathfrak{C}$ ), so we are done.

For the  $\operatorname{Autf}_L(\mathfrak{C}')$ , notice that clearly  $S_I, G^I \leq \operatorname{Autf}_L(\mathfrak{C}')$  (because given a small model  $M \leq \mathfrak{C}'$  such that  $M \cap I = I_0$  we have  $S_{I \setminus I_0}, G^{I \setminus I_0} \leq \operatorname{Aut}(\mathfrak{C}'/M)$ ), and clearly  $\operatorname{Autf}_L(\mathfrak{C}) \leq \operatorname{Autf}_L(\mathfrak{C}')$  (as we have no additional structure on  $\mathfrak{C}$ ), so we have  $G^I \rtimes (S_I \times \operatorname{Autf}_L(\mathfrak{C})) \leq \operatorname{Autf}_L(\mathfrak{C}')$ . The converse is trivial, because  $\operatorname{Autf}_L(\mathfrak{C}')$  acts on  $\mathfrak{C}$  by Lascar strong automorphisms.

The argument for other two strong automorphism groups is analogous.  $\Box$ 

**Example 2.4.** Suppose  $\mathfrak{C} = (\mathbf{R}^*, \mathbf{Z}^*)$  (a large real closed field and a monster extension of the additive group of integers, with a constant for  $1 \in \mathbf{Z}$ ), and let G be the universal cover of  $SL_2(\mathbf{R})$ . Then the theory of  $\mathfrak{C}$  is G-compact (and, in fact,  $Autf_L(\mathfrak{C}) = Aut(\mathfrak{C})$ , because it has a pointwise definable model) and  $G^{00} \neq G^{000}$ . Consider  $\mathfrak{C}'$  as in the last proposition.

By the proposition  $\operatorname{Autf}_L(\mathfrak{C}') = \operatorname{Autf}_{KP}(\mathfrak{C}')$ , but given any  $i_0 \in I$  we have that  $\operatorname{Autf}_L(\mathfrak{C}'/i_0) \neq \operatorname{Autf}_{KP}(\mathfrak{C}'/i_0)$ .

**Example 2.5.** If we take for  $\mathfrak{C}$  a real closed field and  $G = (S^1)^*$ , we will have, as before,  $\operatorname{Autf}_L(\mathfrak{C}) = \operatorname{Aut}(\mathfrak{C})$ , whereas  $G^0 \neq G^{00}$  (because  $G/G^0$  is trivial and  $G/G^{00} = G(\mathbf{R}) = S^1$ ). Thus by the proposition we obtain a theory which has  $\equiv_{KP} = \equiv_{Sh}$ , but not so after adding a single parameter.

**Example 2.6.** We can combine the two previous example by taking  $\mathfrak{C}$  as in the first example, and  $G = \widetilde{\operatorname{SL}_2(\mathbf{R})}^* \times (S^1)^*$ , which yields an example with  $\operatorname{Autf}_L(\mathfrak{C}') = \operatorname{Aut}(\mathfrak{C}')$ , such that the automorphism group and all strong automorphism groups are pairwise distinct over any  $i_0 \in I$ .

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**Proposition 2.7.** The theory DLO of dense linear orderings (without endpoints) interprets without parameters a (structure analogous to) "circles" structure given by Casanovas et al, only without the "even" circles. Therefore, it is not hereditarily G-compact.

*Proof.* Let us denote by N the entire model of DLO. Let  $(c_{n,m})_{n,m}$  be dummy parameters (which can easily be defined in  $N^{eq}$ ), where n natural numbers, and  $1 \le m \le n$ .

Then we interpret the structure as follows:

- $M_n$  is the disjoint union  $\bigcup_{m \le n} \{c_{n,m}\} \times N$ .
- $g_n$  is defined by  $g_n \cdot (c_{n,m}, x) = (c_{n,m+1}, x)$ , where  $c_{n,n+1} = c_{n,1}$ .
- For each n, we have a natural circular order  $S'_n$  on  $\{c_{n,1}, \ldots, c_{n,n}\}$ . This allows us to define the (strict) circular ordering  $S_n$  on  $M_n$ .
  - If two elements  $a_1$ ,  $a_3$  have the same first coordinate, then an element  $a_2$  is between them if and only if it is also has the same first coordinate, and is between according to the ordering on the second coordinate.
  - If  $a_1, a_2, a_3$  have pairwise distinct first coordinates, then  $a_2$  is between them if and only if the first coordinate is between the first coordinates of  $a_1$  and  $a_3$  (in the sense of  $S'_n$ ).
  - If  $a_1, a_2$  have the same first coordinate, distinct from that of  $a_3$ , then  $a_2$  is between  $a_1$  and  $a_3$  if and only if the second coordinate of  $a_2$  is larger or smaller than the second coordinate of  $a_1$ , depending in the obvious way on the relationship between the first coordinates of  $a_1, a_3$ .
  - The interpretation in case when  $a_2$ ,  $a_3$  share the first coordinate, distinct from that of  $a_1$ , is analogous.

This  $(M_n, g_n, S_n)_{n \in 2\mathbb{N}+1}$  is clearly a model of the theory described in Casanovas et al, restricted to odd circles. Therefore, it is not G-compact, as was proven there, and DLO is not hereditarily G-compact.

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