Practical Binary Analysis #11

Seminar @ Gondow Lab.

What you will learn

- 1. Internals of libdft
 - a. data structure of libdft
 - b. how libdft works
- 2. How to use libdft to build DTA-tools
 - a. a tool that prevents remoto control-hijacking attacks
 - b. a tool that automatically detects information leaks

- 1. Internals of libdft
- 2. Using DTA to Detect Remote Control-Hijacking
- 3. Circumeventing DTA with implicit Flows
- 4. A DTA-Based Data Exfiltration Detector

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What is libdft

- Open source binary-level taint tracking library
- Byte-granularity taint-tracking system built on Intel Pin
- Supports only 32-bit x86
 - o althogh you can use it on a 64-bit platform
- Relies on legacy versions of Pin
- Supports only for "Regular" x86 instructions,
 - not for extended instruction sets like MMX or SSE.



- taint: the effect of something bad or unpleasant.(OALD)
- libdft is based on Pin (between 2.11 and 2.13)
- 64-bit version of libdft: https://github.com/AngoraFuzzer/libdft64

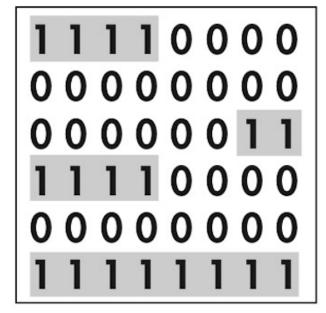
- 1. Internals of libdft
 - Shadow Memory
 - how to store taint info
 - Virtual CPU
 - how to propagete taint info
 - The libdft API and I/O interface
 - how to instrument
 - Taint Policy
- 2. Using DTA to Detect Remote Control-Hijacking
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Internals of libdft - Shadow Memory

Bitmap (1 color)

- Supports only 1 taint color.
- Slightly faster and use less memory than STAB.

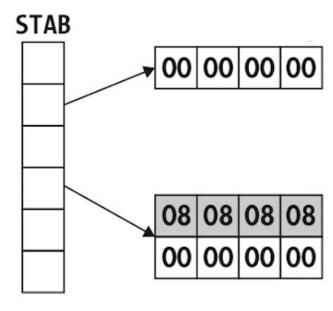
Bitmap (1 color)

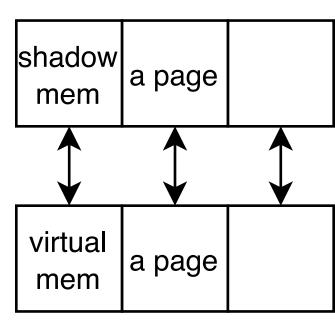


Internals of libdft - Shadow Memory

STAB: Segment Translation Table (8 color)

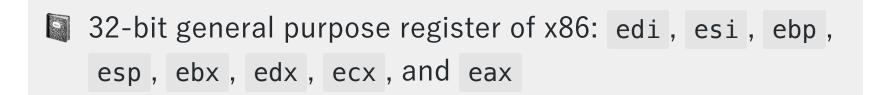
- Supports 8 taint color (because we use 8 bits to contain color).
- Contains one entry for every memory page.
 - and is allocatd in page-sized chunks
- Input and Output of STAB
 - o input: some upper bits of virtual memory address
 - 16 bit if each page size is 2^{16} B(= 64KB)
 - output: some upper bits of shadow memory address
 - 16 bit if each page size is 2^{16} B(= 64KB)
- Lower bits can be used to access each shadow memory.
- Shadow memory pages is adjacent if the corresponding virutual memory is adjacent.





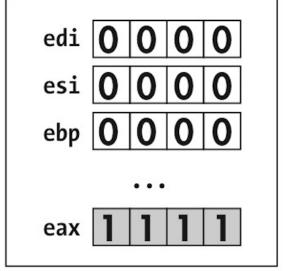
Internals of libdft - Virtual CPU

- Virtual CPU keeps track of the taint status of CPU register.
 - This is stored in memory as a special data structure.
- Virtual CPU is a kind of shadow memory.
 - for each of 32-bit general-purpose CPU registers.

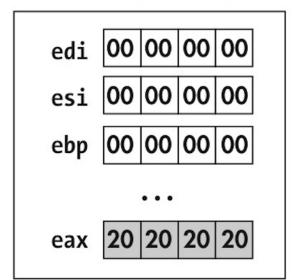


Virtual CPU

VCPU (1 color)



VCPU (8 colors)



Internals of libdft - libdft API and I/O interface

- libdft provides a taint tracking API.
- Two import tools for building DTA tools is those that
 - manipulate shadow memory (Tagmap API)
 - add callbacks and instrument code



tagmap: shadow memory

Internals of libdft - libdft API and I/O interface

- libdft provides a taint tracking API.
- Two import tools for building DTA tools is those that
 - o manipulate tagmap (Tagmap API) ←
 - add callbacks and instrument code

Tagmap API

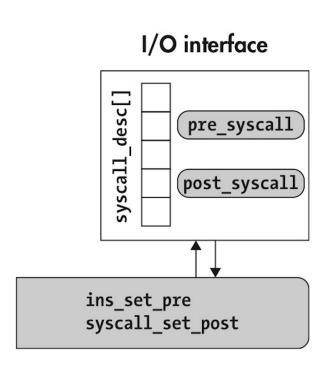
- tagmap_setb() : sets status of tagmap in byte granularity.
 - tagmap_setn(): taints arbitary number of bytes.
- tagmap_getb(): gets status of tagmap in byte granularity.
 - tagmap_getn() : checks arbitary number of bytes.

libdft API and I/O interface

- libdft provides a taint tracking API.
- Two import tools for building DTA tools is those that
 - manipulate tagmap (Tagmap API)
 - o add callbacks and instrument code ←

API for adding callbacks and instrumentation code

- syscall_set_pre(): register callbacks for syscall events.
- syscall_set_post() : same as above
- syscall_desc[] : store syscall pre- and post-handlers.
 - use *syscall number* to index this array.
- ins_set_pre : register callbacks for instructions.
- ins_set_post : same as above



Taint Policy

- libdft taint policy defines the following classes of instructions.
- Each of these classess define how to propagete and merge taint.
- 1. ALU: arithmetic and logic instruction such as add, sbb, and, xor, div, and imul
- 2. XFER: instructions that copy a value such as mov. Simply copies the taint info.
- 3. CLR: (=clear), instructions that reset taint info of output opperands
- 4. SPECIAL: instructions that require special rules
- 5. FPU, MMX, SSE: instructions libdft doesn't supports. Doesn't propagete taint info, causing undertainting.

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- 2. Using DTA to Detect Remote Control-Hijacking
 - Examples of exec family
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Remote Control-Hijacking Attack

- Goal is to detect attacks where data received from network is used to control the argument of execve call (-> arbitary code execution).
- taint source: the network receive functions, recv and recvfrom.
- tiant sink : execve

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Examples of exec family

	引数の渡し方(I or v)	環境変数(e)	パス検索(p)
exec <u>l</u>	list	引き継ぐ	しない
exec <u>v</u>	vector	引き継ぐ	しない
exec <u>le</u>	list	渡す	しない
exec <u>ve</u>	vector	渡す	しない
exec <u>lp</u>	list	引き継ぐ	する

```
#include <unistd.h>
int main(void){
  char *argv[] = {"ls", "-l", NULL};
  char *envp[] = {NULL};
  execve("/bin/ls", argv, envp);
}
```

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Header files

• all libdft tools are just Pin tools ilnked with the libdft library.

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Functions

- syscall_desc: Index this array with syscall number of syscall you're installing
 such as __NR_socketcall or __NR_execve.
- Details of these functions will be explained later.

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main function

```
int main(int argc, char **argv) {
  PIN_InitSymbols();
  if (unlikely(PIN_Init(argc, argv))) {
    return 1;
  if (unlikely(libdft_init() != 0)) {
    libdft_die();
    return 1;
  syscall_set_post(&syscall_desc[__NR_socketcall], post_socketcall_hook);
  syscall_set_pre(&syscall_desc[__NR_execve], pre_execve_hook);
  PIN_StartProgram();
  return 0;
```

main function

```
int main(int argc, char **argv) {
  PIN_InitSymbols(); // in case symbols are available
  if (unlikely(PIN_Init(argc, argv))) {
    return 1;
  if (unlikely(libdft_init() != 0)) { // init data structures such as tagmap
    libdft_die(); // deallocat any resources libdft may have allocated
    return 1;
```

- unlikely(hoge) tells compiler that hoge is unlikely to be ture (= likely to be false).
 - Better branch prediction and less cycles.

main function

```
// socketcall events as the taint sources
syscall_set_post(&syscall_desc[__NR_socketcall], post_socketcall_hook);
// taint sink
syscall_set_pre(&syscall_desc[__NR_execve], pre_execve_hook);

PIN_StartProgram(); // never returns

return 0; // never reached
} // end of main
```

- socketcall events include recv and recvfrom events
 - On some architectures—for example, x86-64 and ARM—there is no socketcall() system call; instead socket(2), accept(2), bind(2), and so on really are implemented as separate system calls. (from man socketcall)

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Details of func - alert

- This alert function
 - i. prints an alert message with the details about the tainted address.
 - ii. exit from the application

Details of func - check_string_taint

```
void check_string_taint(const char *str, const char *source) {
                                         // to store "color"
 uint8 t tag;
 uintptr_t end = (uintptr_t)str + strlen(str); // end of string
 for (uintptr_t addr = start; addr <= end; addr++) {</pre>
   tag = tagmap_getb(addr);
                                        // get the "color" of addr
   if (tag != 0)
     alert(addr, source, tag);
                            // alert if tag is tainted
 fprintf(stderr, "OK\n");
} // end of check_string_taint
```

Details of func - post_socketcall_hook

```
static void post_socketcall_hook(syscall_ctx_t *ctx) {
  int fd;
 void *buf;
  size t len;
  int call = (int)ctx->arg[SYSCALL_ARG0];
  unsigned long *args = (unsigned long *)ctx->arg[SYSCALL_ARG1];
  switch (call) {
  case SYS_RECV:
  case SYS_RECVFROM:
   // . . .
    // omitted
    // ...
    break;
  default:
    break;
```

29/82

Details of func - post_socketcall_hook

- ctx : contains
 - i. the arguments that were passed to the syscall
 - ii. the return value of syscall.

Details of func - post_socketcall_hook

```
switch (call) {
  case SYS_RECV:
  case SYS_RECVFROM:
    // ...
    // omitted
    // ...
    break;
  default:
    break;
} // end of post_socketcall_hook
```

- Ignores any cases other than SYS_RECV or SYS_RECVFROM.
 - Thus, catching all the socketcall does work.

Details of func - post_socketcall_hook (omitted part)

```
... // start of omitted part
    if (unlikely(ctx->ret <= 0)) {</pre>
      return;
    fd = (int)args[0];
    buf = (void *)args[1];
    len = (size_t)ctx->ret;
    for (size_t i = 0; i < len; i++) {
      if (isprint(((char *)buf)[i]))
        fprintf(stderr, "%c", ((char *)buf)[i]);
      else
        fprintf(stderr, "\\x%02x", ((char *)buf)[i]);
    fprintf(stderr, "\n");
    tagmap_setn((uintptr_t)buf, len, 0x01);
... // end of omitted part
```

Details of func - post_socketcall_hook (omitted part)

Details of func - post_socketcall_hook (omitted part)

```
for (size_t i = 0; i < len; i++) { // print each char
    if (isprint(((char *)buf)[i]))
        fprintf(stderr, "%c", ((char *)buf)[i]);
    else
        fprintf(stderr, "\\x%02x", ((char *)buf)[i]);
}
fprintf(stderr, "\n");
tagmap_setn((uintptr_t)buf, len, 0x01);
... // end of omited part</pre>
```

- isprint(hoge) returns
 - none-0 if hoge can be printed
 - 0 if hoge can't be printed

- buf: first address that will be tainted
- len: # of bytes to taint
- 0x01 : taint color

Details of func - post_socketcall_hook (repeated)

```
static void post_socketcall_hook(syscall_ctx_t *ctx) {
  int fd;
 void *buf;
  size t len;
  int call = (int)ctx->arg[SYSCALL_ARG0];
  unsigned long *args = (unsigned long *)ctx->arg[SYSCALL_ARG1];
  switch (call) {
  case SYS_RECV:
  case SYS_RECVFROM:
   // ...
    // omitted, just explained
    // . . .
    break;
  default:
    break;
```

Details of func - pre_execve_hook

```
static void pre_execve_hook(syscall_ctx_t *ctx) {
  const char *filename = (const char *)ctx->arg[SYSCALL_ARG0];
  char *const *args = (char *const *)ctx->arg[SYSCALL_ARG1];
  char *const *envp = (char *const *)ctx->arg[SYSCALL_ARG2];
  check_string_taint(filename, "execve command");
 while (args && *args) {
    fprintf(stderr, "(dta-execve) arg: %s (@%p)\n", *args, *args);
    check_string_taint(*args, "execve argument");
    args++;
 while (envp && *envp) {
    fprintf(stderr, "(dta-execve) env: %s (@%p)\n", *envp, *envp);
    check_string_taint(*envp, "execve environment parameter");
   envp++;
```

Details of func - pre_execve_hook

```
static void pre_execve_hook(syscall_ctx_t *ctx) {
  const char *filename = (const char *)ctx->arg[SYSCALL_ARG0];
  char *const *args = (char *const *)ctx->arg[SYSCALL_ARG1];
  char *const *envp = (char *const *)ctx->arg[SYSCALL_ARG2];
  check_string_taint(filename, "execve command");
...
```

- ctx : contains
 - i. the arguments that were passed to the syscall
 - ii. the return value of syscall.
- Check whether filename is tainted or not.

Details of func - pre_execve_hook

```
while (args && *args) {
  fprintf(stderr, "(dta-execve) arg: %s (@%p)\n", *args, *args);
  check string taint(*args, "execve argument");
  args++;
while (envp && *envp) {
  fprintf(stderr, "(dta-execve) env: %s (@%p)\n", *envp, *envp);
  check_string_taint(*envp, "execve environment parameter");
  envp++;
```

Check whether args and envp are taited or not.

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- 2. Using DTA to Detect Remote Control-Hijacking
 - Examples of exec family
 - Header files
 - Functions
 - main function
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Test of Control-Flow Hijacking - main

```
int main(int argc, char *argv[]) {
  char buf [4096];
  struct sockaddr_storage addr;
  int sockfd = open_socket("localhost", "9999");
  socklen t addrlen = sizeof(addr);
  recvfrom(sockfd, buf, sizeof(buf), 0, (struct sockaddr *)&addr, &addrlen);
  int child_fd = exec_cmd(buf);
  FILE *fp = fdopen(child_fd, "r");
 while (fgets(buf, sizeof(buf), fp)) {
    sendto(sockfd, buf, strlen(buf) + 1, 0, (struct sockaddr *)&addr, addrlen);
  return 0;
```

• Some error-handling codes are omitted.

Test of Control-Flow Hijacking - main

```
int main(int argc, char *argv[]) {
  char buf[4096];
  struct sockaddr_storage addr;

int sockfd = open_socket("localhost", "9999"); // 1

  socklen_t addrlen = sizeof(addr);
  recvfrom(sockfd, buf, sizeof(buf), 0, (struct sockaddr *)&addr, &addrlen); // 2
...
```

- 1. Opens a socket.
- 2. Receives a meassage from the socket.

Test of Control-Flow Hijacking - main

```
int child_fd = exec_cmd(buf); // 3
FILE *fp = fdopen(child_fd, "r");
while (fgets(buf, sizeof(buf), fp)) {
    sendto(sockfd, buf, strlen(buf) + 1, 0, (struct sockaddr *)&addrlen); // 4
}
return 0;
}
```

- 3. Executes a command.
- 4. Writes the output of the command output to network socket.
- exec_cmd is vulnerable function.
 - The args of execve can be influenced by an attacker.

Test of Control-Flow Hijacking - cmd

- cmd contains:
 - prefix for the command output
 - datefmt for the output of date command
 - o cmd, date itself

Test of Control-Flow Hijacking - exec_cmd

```
int exec_cmd(char *buf) {
  int pid; int p[2]; char *argv[3];
  for (i = 0; i < strlen(buf); i++) { // [1]</pre>
    if (buf[i] == '\n') {
      cmd.prefix[i] = '\0';
      break;
    cmd.prefix[i] = buf[i]; // ** Buffer overflow **
  argv[0] = cmd.cmd;
  argv[1] = cmd.datefmt;
  argv[2] = NULL;
  pipe(p) // omitted error-handling
```

• [1] lacks propper bound check, allowing attackers to overwrite the cmd field.

Test of Control-Flow Hijacking - exec_cmd

```
switch (pid = fork()) {
\cdot \cdot \cdot \cdot //  case -1:
case 0: /* Child */
  printf("(execve-test/child) execv: %s %s\n", argv[0], argv[1]);
  fflush(stdout);
  close(1);
  dup(p[1]);
  close(p[0]);
  printf("%s", cmd.prefix);
  fflush(stdout);
  execv(argv[0], argv);
  ... // error-handling
default: /* Parent */
  close(p[1]);
  return p[0];
return -1;
```

45/82

Test of Control-Flow Hijacking - main (repeated)

```
int main(int argc, char *argv[]) {
  char buf [4096];
  struct sockaddr_storage addr;
  int sockfd = open_socket("localhost", "9999");
  socklen t addrlen = sizeof(addr);
  recvfrom(sockfd, buf, sizeof(buf), 0, (struct sockaddr *)&addr, &addrlen);
  int child_fd = exec_cmd(buf);
  FILE *fp = fdopen(child_fd, "r");
 while (fgets(buf, sizeof(buf), fp)) {
    sendto(sockfd, buf, strlen(buf) + 1, 0, (struct sockaddr *)&addr, addrlen);
  return 0;
```

Test of Control-Flow Hijacking - test of overflow

No Buffer-Overflow

- 1. Start server. (localhost(=127.0.0.1), port 9999)
- 2. Send prefix to localhost:9999 using nc.



Test of Control-Flow Hijacking - test of overflow

Buffer-Overflow

```
$ ./execve-test-overflow &
[1] 2061
$ nc -u 127.0.0.1 9999
(execve-test/child) execv: /home/binary/code/chapter11/echo bb...bb/home/binary/.../echo
aa...aabb...bb/home/binary/code/chapter11/echo bb...bb/home/binary/code/chapter11/echo
                        ./execve-test-overflow
^C[1]+ Done
```

```
static struct __attribute__((packed)) {
} cmd;
```



-u option: Use UDP instead of TCP.

Test of Control-Flow Hijacking - Detect Hijacking

Using DTA to Detect Hijacking Attempt

- 1. Check whether argments of execve are taitned.
- 2. dta-execve notices that the command is tainted with 0x01.
- 3. Raises an alert and then stops the child process.

Overview

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Circumeventing DTA with Implicit Flows

```
int exec_cmd(char *buf) {
  for (i = 0; i < strlen(buf); i++) {</pre>
    if (buf[i] == '\n') {
      cmd.prefix[i] = '\0';
      break;
    c = 0;
    while (c < buf[i]) c++; // increment c until c == b[i]</pre>
    cmd.prefix[i] = c; // c == b[i], but c is not tainted
   // Set up argv and continue with execv
```

- Without explicitly copying b[i] to c, c has the same value as b[i].
- We call this *implicit flow*.
 - libdft cannot track this kind of flow, causing undertainting.
- Malware may contain implicit flow to confuse taint analysis.

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 - Header files
 - Data structure
 - Functions
 - main function
 - Details of func
 - Test of Data Exfiltration

A DTA-Based Data Exfiltration Detector

- We use multiple taint color so that we can tell which file is leaking.
 - In the previous example, single taint color is enough to detect bytes are attackercontrolled or not.
- Taint settings
 - Taint source : open and read
 - Taint sink: socketcall, such as send, sendto

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Header files

```
#include "pin.H"

#include "branch_pred.h"
#include "libdft_api.h"
#include "syscall_desc.h"
#include "tagmap.h"
```

- all libdft tools are just Pin tools ilnked with the libdft library.
- This is same as the previous example.

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Data Structure

- fd2color: maps file discriptors to colors
- color2fname: maps taint colors to filenames

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Functions

```
void alert(uintptr_t addr, uint8_t tag);
static void post_open_hook(syscall_ctx_t *ctx);
static void post_read_hook(syscall_ctx_t *ctx);
static void pre_socketcall_hook(syscall_ctx_t *ctx);
```

- post_open_hook / post_read_hook runs after open / read syscall respectively.
- pre_socketcall_hook runs before the socketcall syscall such as recv or recvfrom.

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main function

```
int main(int argc, char **argv) {
  PIN_InitSymbols();
  if (unlikely(PIN_Init(argc, argv))) {
    return 1:
  if (unlikely(libdft_init() != 0)) {
    libdft_die();
    return 1;
  syscall_set_post(&syscall_desc[__NR_open], post_open_hook);
  syscall_set_post(&syscall_desc[__NR_read], post_read_hook);
  syscall_set_pre(&syscall_desc[__NR_socketcall], pre_socketcall_hook);
  PIN_StartProgram();
  return 0;
```

main func is almost identical to that of the previous example.

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Details of func alert

```
void alert(uintptr_t addr, uint8_t tag) {
  fprintf(stderr,
          "\n(dta-dataleak) !!!!!!! ADDRESS 0x%x IS TAINTED (tag=0x%02x), "
          "ABORTING !!!!!!\n",
          addr, tag);
  for (unsigned c = 0 \times 01; c \le MAX_COLOR; c \le 1) {
    if (tag & c) {
      fprintf(stderr, " tainted by color = 0x\%02x (%s)\n", c,
              color2fname[c].c_str());
  exit(1);
```

- Alert which address is tainted and with which color.
- It's possible that the data is tainted with multiple color (= multiple files).

```
static void post_open_hook(syscall_ctx_t *ctx) {
   static uint8_t next_color = 0x01;
   uint8_t color;
   int fd = (int)ctx->ret; // return value of syscall open
   const char *fname = (const char *)ctx->arg[SYSCALL_ARG0]; // filename to open
}
```

- ctx->ret : contains return value of syscall.
 - In this case, return value is the file discriptor that was opened.
- fname: filename that was opened.

```
static void post_open_hook(syscall_ctx_t *ctx) {
    if (unlikely((int)ctx->ret < 0)) {
        return;
    }
    if (strstr(fname, ".so") || strstr(fname, ".so.")) {
        return;
    }
    ...
}</pre>
```

- 1. Checks whether the return value is not smaller than 0.
 - You don't need to taint if open is failed.
- 2. Filters out uninteresting files such as shared libraries.
 - Shared libraries don't have any secret informations.
 - In a real-world DTA tool, you should filter out some more files.

```
static void post_open_hook(syscall_ctx_t *ctx) {
    if (!fd2color[fd]) {
        color = next_color;
        fd2color[fd] = color;
        if (next_color < MAX_COLOR) next_color <<= 1; // static variable
    } else {
        color = fd2color[fd]; // reuse color of a file with the same fd
    }
}</pre>
```

- "color" can be reused
 - if a file discriptor is closed and then the same file discriptor is reused.
- MAX_COLOR can be assigned to many fd
 - if we run out of "color".
 - We only supports 8 colors (because color is 8-bit wise).

```
static void post_open_hook(syscall_ctx_t *ctx) {
   if (color2fname[color].empty())
      color2fname[color] = std::string(fname);
   else
      color2fname[color] += " | " + std::string(fname);
}
```

- Update the color2fname map with just opened filename.
- Filename is concatinated with " | "
 - if taint color is reused for multiple files.

```
static void post_read_hook(syscall_ctx_t *ctx) {
  int fd = (int)ctx->arg[SYSCALL_ARG0];
  void *buf = (void*)ctx->arg[SYSCALL_ARG1];
  size_t len = (size_t)ctx->ret;
  uint8_t color;
}
```

- fd: file discriptor that's being read
- buf: buffer into which bytes are read
- len: length of buffer that was read.

```
static void post_read_hook(syscall_ctx_t *ctx) {
    if(unlikely(len <= 0)) {
        return;
    }
    fprintf(stderr, "(dta-dataleak) read: %zu bytes from fd %u\n", len, fd);
}</pre>
```

- You don't need to taint if nothing was read, that is when...
 - i. 0 byte was read(, when return value is 0).
 - ii. read failed (, when return value is negative).

```
static void post_read_hook(syscall_ctx_t *ctx) {
    color = fd2color[fd];
    if(color) {
       tagmap_setn((uintptr_t)buf, len, color);
    } else {
       tagmap_clrn((uintptr_t)buf, len);
    }
}
```

- Taint bytes using tagmap_setn()
 - o if the fd is colored.
- Clear taint on bytes using tagmap_clrn()
 - o if the fd is not colored.

Details of func - pre_socketcall_hook

```
static void pre_socketcall_hook(syscall_ctx_t *ctx) {
 int fd:
 void *buf;
 size_t i, len;
 uint8_t tag;
 uintptr_t start, end, addr;
          = (int)ctx->arg[SYSCALL_ARG0];
 int call
 unsigned long *args = (unsigned long*)ctx->arg[SYSCALL_ARG1];
 switch(call) {
 case SYS SEND:
 case SYS_SENDTO:
   break;
 default:
   break;
```

 $\frac{71/82}{}$

Details of func - pre_socketcall_hook

- call : type(number) of socketcallsuch as recv, recvfrom, send, sendto.
- args: arguments that was passed to socketcall

Details of func - pre_socketcall_hook

```
static void pre_socketcall_hook(syscall_ctx_t *ctx) {
  switch(call) {
  case SYS_SEND:
  case SYS_SENDTO:
    ... // omitted
    break;
  default:
    break;
```

• Check the tagmap, if the socketcall is send or sendto.

Details of func - pre_socketcall_hook (omitted part)

```
static void pre_socketcall_hook(syscall_ctx_t *ctx) {
    fd = (int)args[0];
    buf = (void*)args[1];
    len = (size_t)args[2];
    start = (uintptr_t)buf;
    end = (uintptr_t)buf+len;
    for(addr = start; addr <= end; addr++) {</pre>
      tag = tagmap_getb(addr);
      if(tag != 0) alert(addr, tag);
```

- Loops over all of bytes in the send buffer and check whether they are tainted of not.
- If tainted, alert and exit the application.

Overview

- 1. about libdft
- 2. Using DTA to Detect Remote Control-Hijacking
- 3. Circemeventing DTA with implicit Flows
- 4. A DTA-Based Data Exfiltration Detector
 - Header files
 - Data structure
 - Functions
 - main function
 - Details of func
 - Test of Data Exfiltration

Test of Data Exfiltration - dataleak-test-xor - main

```
int main(int argc, char *argv[]) {
  size_t i, j, k;
  FILE *fp[10];
  char buf1[4096], buf2[4096], *filenames[10];
  struct sockaddr_storage addr;
  srand(time(NULL)); // set seed
  int sockfd = open socket("localhost", "9999"); // 1
  socklen t addrlen = sizeof(addr);
  recvfrom(sockfd, buf1, sizeof(buf1), 0, (struct sockaddr*)&addr, &addrlen); // 2
```

- 1. Open a socket.
- 2. Read filenames from the socket.

Test of Data Exfiltration - dataleak-test-xor - main

```
int main(int argc, char *argv[]) {
  size_t fcount = split_filenames(buf1, filenames, 10); // 1
  for(i = 0; i < fcount; i++) {</pre>
    fp[i] = fopen(filenames[i], "r"); // 2
  i = rand() % fcount; // 3
  do { j = rand() % fcount; } while(j == i);
  memset(buf1, '\0', sizeof(buf1)); // initialize buffer
  memset(buf2, '\0', sizeof(buf2));
```

- 1. Gets each filenames.
- 2. Opens all the requested files.
- 3. Choses two of the opened files at random.

Test of Data Exfiltration - dataleak-test-xor - main

```
int main(int argc, char *argv[]) {
  while(fgets(buf1, sizeof(buf1), fp[i]) && fgets(buf2, sizeof(buf2), fp[j])) {
    for(k = 0; k < sizeof(buf1)-1 && k < <math>sizeof(buf2)-1; k++) {
      buf1[k] ^= buf2[k];
    sendto(sockfd, buf1, strlen(buf1)+1, 0, (struct sockaddr*)&addr, addrlen);
  return 0;
```

- Reads each files line by line, concatinating each pair of lines by operating XOR and sending to the socket.
- buf[sizeof(buf) 1] would be a NULL character, so you should loop over from 0
 to sizeof(buf) 2.

```
$ ./pin.sh -follow-execv -t ~/code/chapter11/dta-dataleak.so -- ~/code/chapter11/dataleak-test-xor &
(dta-dataleak) read: 512 bytes from fd 4
(dta-dataleak) clearing taint on bytes 0xffb4aa80 -- 0xffb4ac80

$ nc -u 127.0.0.1 9999
/home/binary/code/chapter11/dta-execve.cpp .../dta-dataleak.cpp .../date.c .../echo.c
```

- 1. Runs dataleak-test-xor server with dta-dataleak as Pin tool,o immeediately loading dataleak-test-xor itself.
- 2. Starts netcat session to connect to the server.
- 3. Sends a list of filenames to open.

```
$ nc -u 127.0.0.1 9999
/home/binary/code/chapter11/dta-execve.cpp .../dta-dataleak.cpp .../date.c .../echo.c
(dta-dataleak) opening /home/binary/code/chapter11/dta-execve.cpp at fd 5 with color 0x01
(dta-dataleak) opening /home/binary/code/chapter11/dta-dataleak.cpp at fd 6 with color 0x02
(dta-dataleak) opening /home/binary/code/chapter11/date.c at fd 7 with color 0x04
(dta-dataleak) opening /home/binary/code/chapter11/echo.c at fd 8 with color 0x08
```

1. Assigns each of files with a taint color.

```
$ nc -u 127.0.0.1 9999
/home/binary/code/chapter11/dta-execve.cpp .../dta-dataleak.cpp .../date.c .../echo.c
...
(dta-dataleak) read: 4096 bytes from fd 6
(dta-dataleak) tainting bytes 0x9b775c0 -- 0x9b785c0 with color 0x2
(dta-dataleak) read: 155 bytes from fd 8
(dta-dataleak) tainting bytes 0x9b785c8 -- 0x9b67663 with color 0x8
(dta-dataleak) send: 20 bytes from fd 4
...
```

- 1. Randomly chooses two files to leak,
 - that is, 6 and 8.
- 2. Intercepts the server's attempt to send the contents of files.

```
$ nc -u 127.0.0.1 9999
/home/binary/code/chapter11/dta-execve.cpp .../dta-dataleak.cpp .../date.c .../echo.c
...
(dta-dataleak) checking taint on bytes 0xffb48f7c -- 0xffb48f90...
(dta-dataleak) !!!!!! ADDRESS 0xffb48f7c IS TANTED (tag=0x0a), ABORTING !!!!!!!
  tainted by color = 0x02 (/home/binary/code/chapter11/dta-dataleak.cpp)
  tainted by color = 0x08 (/home/binary/code/chapter11/echo.c)
^C
```

- 1. Checks the taint color of the contents,
 - detecting that they're tainted.