Heap store Objects: Where? Objects are more permanent "objects"
method called includes arrays
smethod returns
Sobject lives on "objects") .. can't store objects in frames The memory where objects are stored is called the heap.

Example

public class Date private into private Strand private into private int	ing month;	instance var (non-str - stored in bject) of	atic freds	s) ate
Each instance	re of Da	te laid or	it as -	
	•		> a Striv	na
		(		
object header	Lay	month	year	

Memory contains: chain of linked frames is called a - frames - seen already Das objects (including arrays) stacks method area method bytecode constant pools class variables (static fields)

Stack
One record (frame)
per welhood call
Parameters, local variables
operand stacks

One record per object (instance) Instance variables (non-static fields)

Method area

One record per class

Constant data: constant pool, method bytecode

Class variables (static fields)

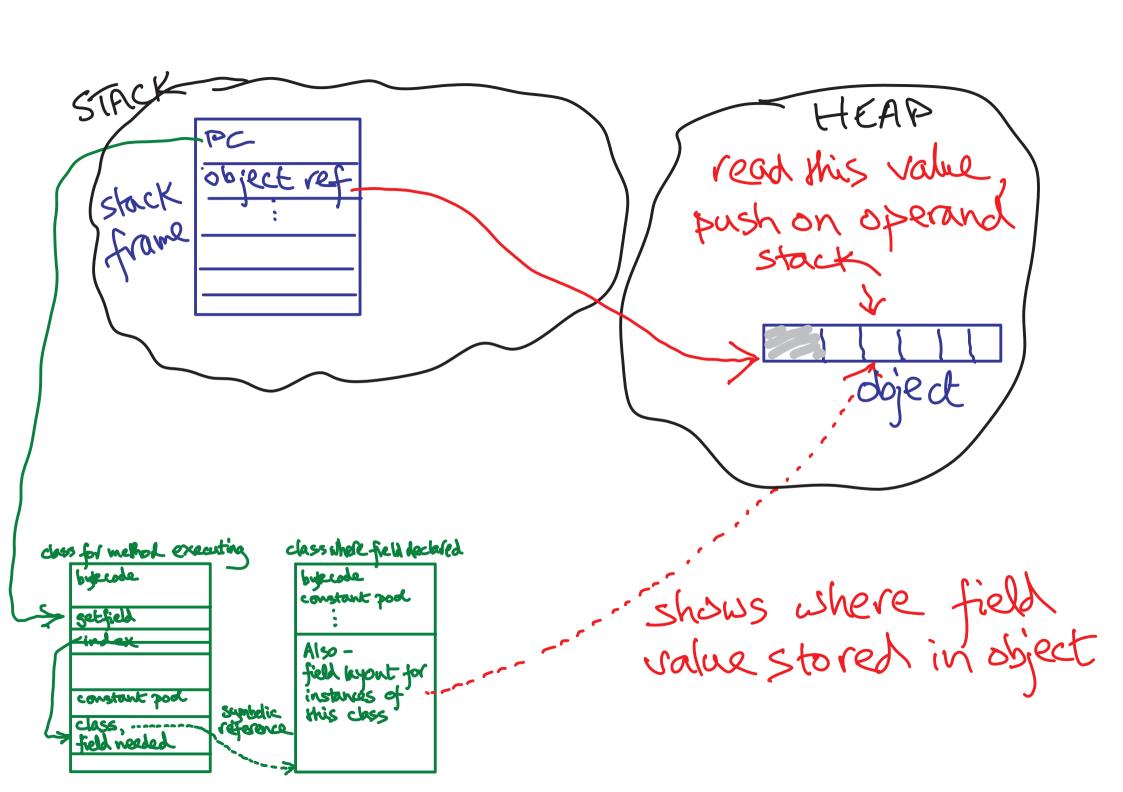
Govill see these later in shi workshop

Exact layout depends An object Number of references to this object « Explained later count class address of values of instance variables (fields) method area for class of shis object object header includes these

### For local variables: use iload, fload etc.) ne variables using bytecode

Accessing instance variables using bytecode

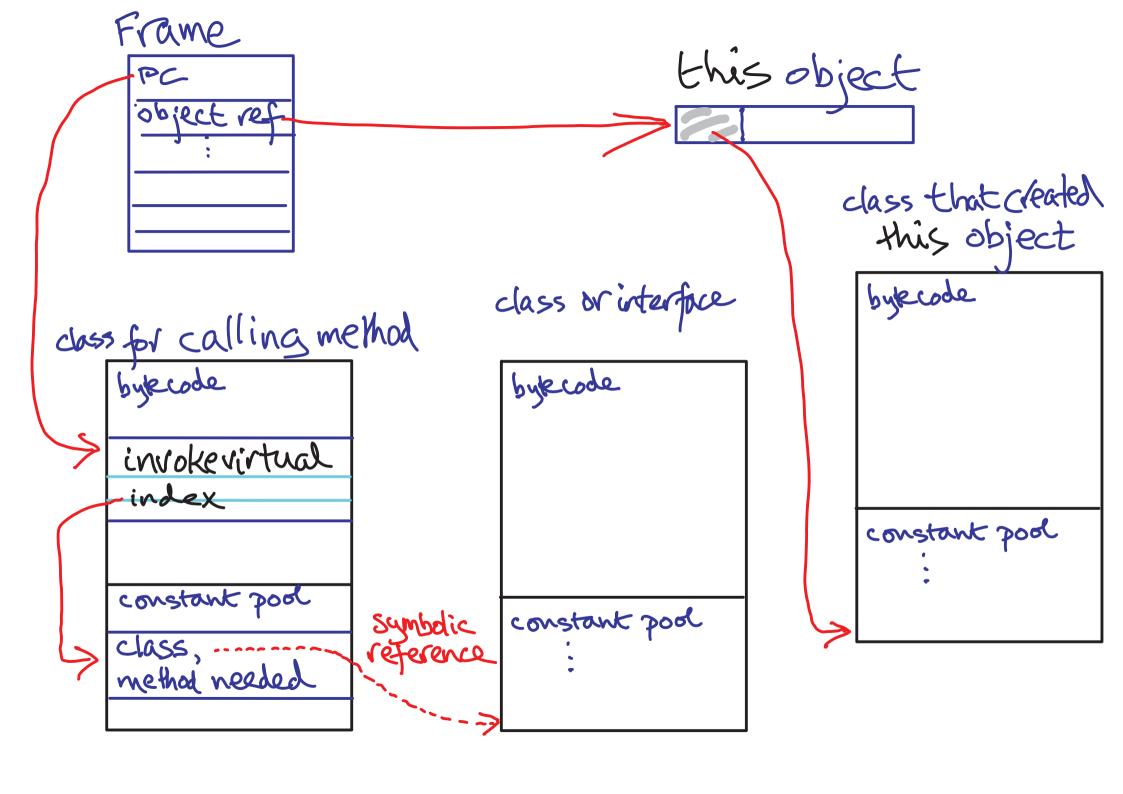
getfield 2-byte index ---, object ref => ---, value of variable pushes variable value onto operand stack index is index into constant pool - object reference is address of object - constant pool entry is symbolic reference - names of instance variable + class declaring it

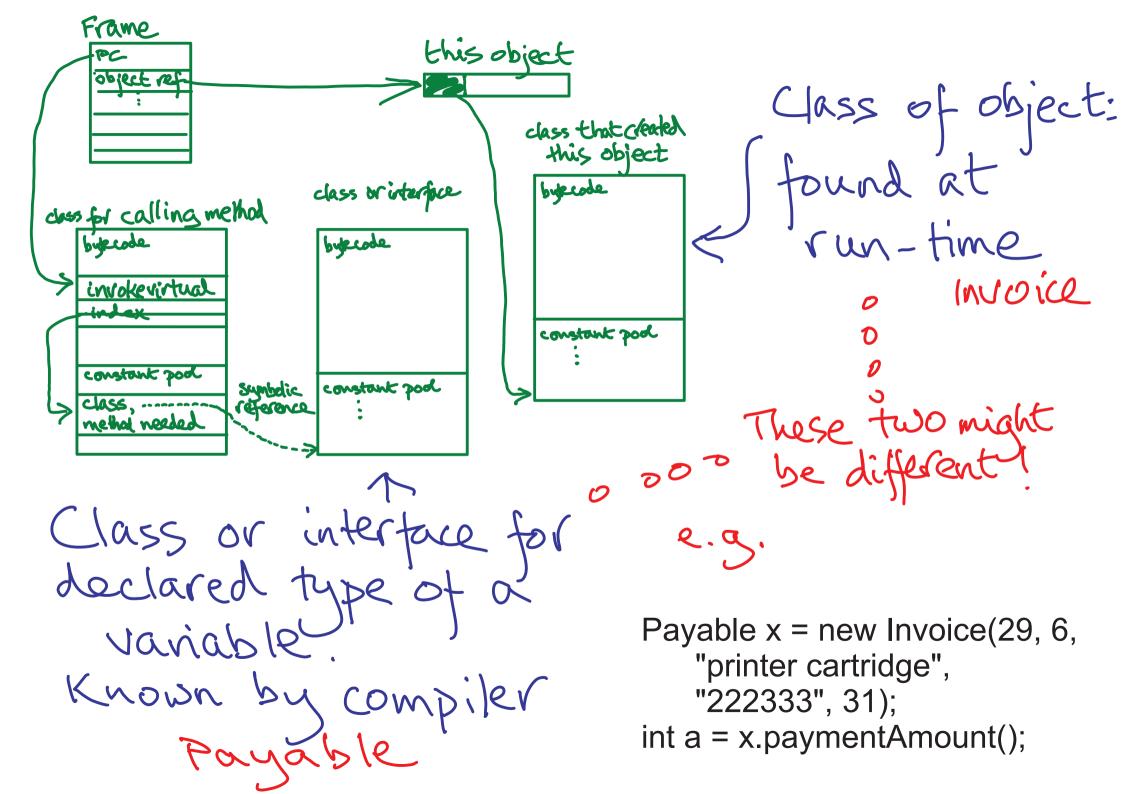


Accessing fields using bytecode putfield 2-byte index is similar, to pop value from operand stock: ...., object ref, new field value => Also getstatic, putstatico (understand these for static fields (class variables)

- but no object ref used for these

#### Calling an instance method 2-byte index invokevirtual Similar to invokestatic invokestatic 2-byte index · index is used as index into constant for static methods From it can be found but operand stack has - address of byterode for method ref to "this" object shat will execute method copied to variable #0 in frame





## Rule - of "method dispatch" It is always the class of the object that governs which method definition is chosen.

- . This makes a difference if the type of the variable is itself a class.
- . It means the method must be chosen at run-time the compiler can't work it out
- · C++ is différent! But she java rule is secure.

Which method definition? Every class extends

of you don't define to String:

Object

(in Object class)

Your class inherits the default of your definition overrides the default

Chains of overriding - anise from class extension (subclasses) invoke virtual must find correct definition (mot overrides all the others) Method must be Look at class of object defined somewhere otherwise class is abstract, so con't have instances. >Use it (compiler Superclass enforces this.)

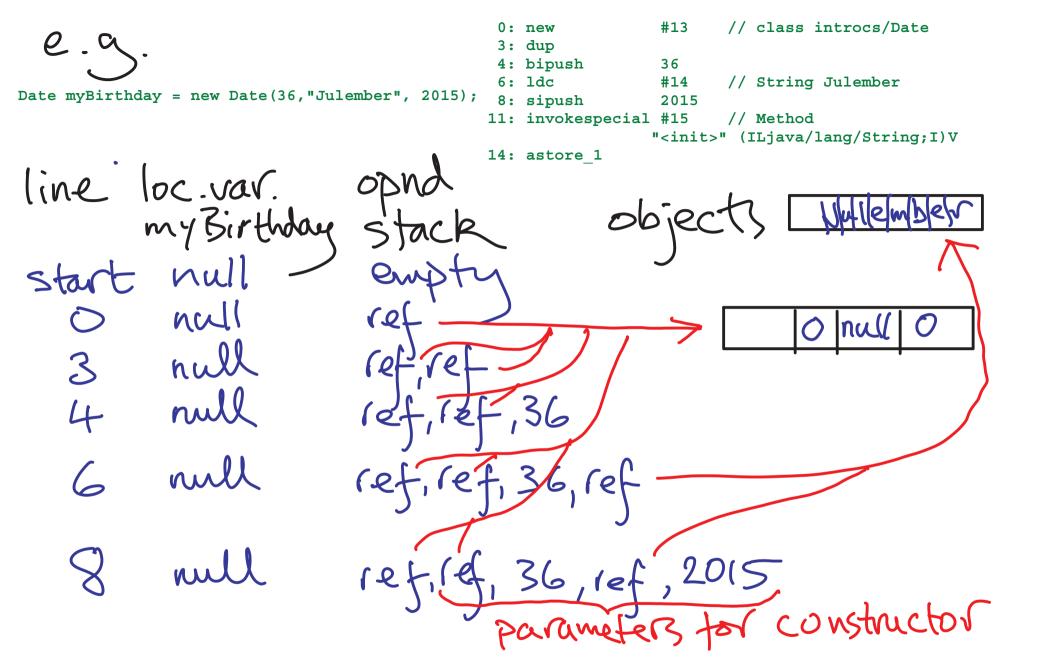
# Creating an object in bytecode Two steps D Allocate space, with default values (2) Execute constructor to do proper initialization

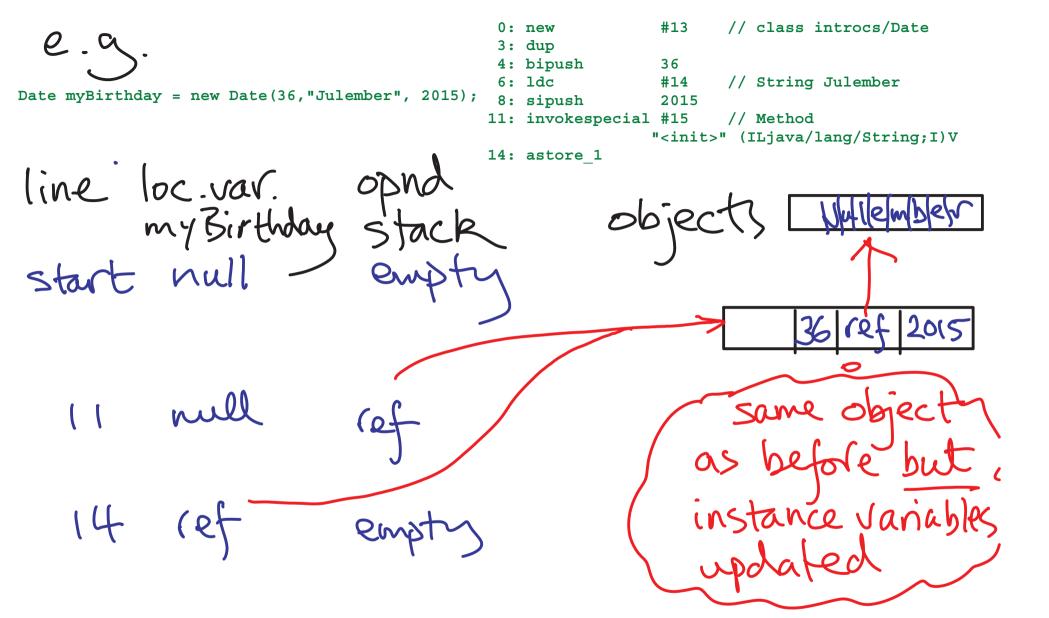
Creating an object in bytecode 1) Allocate space new 2-byte index · index is to constant pool entre with name of class for new object · finds class de finition · allocates heap space for object with correct fields, initialized to default of rets pushes reference to new object etc. .... => ...., object ref

Creating an object in bytecode
2) Execute constructor
invokespecial 2-byte index
· index into constant pool gives entry
· expects ref to new object on stack
D2 usually done as cf. invokevirtual
new duplicates top of stack dup duplicates top of stack invokespecial so object ref still invokespecial on stack at end

e.s.

```
Date myBirthday = new Date(36, "Julember", 2015);
                  #13 // class introcs/Date
 0: new
 3: dup
                  36
 4: bipush
 6: ldc
                  #14 // String Julember
 8: sipush
                 2015
 11: invokespecial #15
                         // Method
                  "<init>" (ILjava/lang/String; I) V
 14: astore 1
```





### Verifier checks on object creation

- new always followed by appropriate invokespecial
- 2) extra checks associated with class extension

> object always created consistent with what its designers intend

object	free object	objects, some parts free
free	) [ obj	ect
When new ob in free an Whom object	eas. is destroyed:	its space made

Garbage collection In some languages (e.g. C++) you destroy objects explicitly when you have finished with shem. Then their storage area can be recycled In Java you don't. The garbage collector is part of the JVM. It works out when objects are finished with a shen recycles storage area — marks it "free"

When is an object finished with? Each object is like Isla de Muerter in Pirates of the Caribbean: "It's an island that cannot be found, except by those who already know where it is." To know where an object is: have a reference to it. When there are no references to an object: it cannot be used any more. .. garbage collector can safely delete it.

