

# Introduction to Computer Science

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# Recap

- Last time:
  - Representing numbers in the computer
    - Whole numbers in binary and hexadecimal notation
    - Positive real numbers in fixed-point binary
- This time:
  - Negative numbers
  - Arithmetic
  - Floating point

# Arithmetic

- We have a representation, need manipulations
- Again, useful to understand decimal arithmetic:

$$\begin{array}{r} 127 \\ + 84 \\ \hline \text{(Carry 1 1 )} \\ 211 \end{array}$$

- When a value is too large ( $>9$ ) for a column, part of it is carried to the next column
- Same in binary, but carry at  $>1$

# Binary Addition

$$\begin{array}{r} 1\ 0\ 0\ 1 \\ +\ 0\ 0\ 1\ 1 \\ \hline \text{(Carry } 1\ 1\ ) \\ 1\ 1\ 0\ 0 \end{array}$$

- Try:
  - $1011 + 0010$
  - $1101 + 0100$

# Overflow

- Numbers represented by fixed number of bits
- Large numbers “overflow” the maximum capacity of the number of allocated bits – computer will usually raise a hardware “error”
- Images typically use 8-bits to store whole number values representing image brightness
- 16, 32, 64 bit basic representations are now in common use (list in handout)

# Subtraction

- Back to the decimals:

$$\begin{array}{r} 211 \\ - 84 \\ \hline \end{array}$$

(Borrow 1 1 )

$$\begin{array}{r} 127 \end{array}$$

- Binary

$$\begin{array}{r} 1100 \\ - 1001 \\ \hline \end{array}$$

(Borrow 1 )

$$\begin{array}{r} 0011 \end{array}$$

# Negative numbers

- We need to be able to represent negative numbers purely in binary
- In decimal representation, we have ten symbols plus the “minus” sign
- In binary, only two symbols
- So how do we denote the “sign” of a number?

# Try something obvious

- Add a “sign” bit – assume on the left
- Then  $+5 = \mathbf{0101}$  and  $-5 = \mathbf{1101}$
- We know that  $+5 + -5 = 0$

$$\begin{array}{r} 0101 \\ + 1101 \\ \hline \text{(Carry } 1 \quad 1 \text{ )} \\ 10010 \end{array}$$

- We need to be a bit smarter about this...



# An observation about overflow

- Let's investigate overflow a bit more closely

$$\begin{array}{r} 0\ 1\ 1\ 1 \\ +\ 1\ 1\ 1\ 1 \\ \hline \text{(Carry } 1\ 1\ 1\ ) \\ (1)\ 0\ 1\ 1\ 0 \end{array}$$

- Interesting...try something else

$$\begin{array}{r} 1\ 0\ 1\ 0 \\ +\ 1\ 1\ 1\ 1 \\ \hline \text{(Carry } 1\ 1\ ) \\ (1)\ 1\ 0\ 0\ 1 \end{array}$$

# An observation about overflow

- And again

$$\begin{array}{r} 0\ 1\ 1\ 1 \\ +\ 1\ 1\ 1\ 0 \\ \hline \text{(Carry } 1\ 1\ 1\text{ )} \\ (1)\ 0\ 1\ 0\ 1 \end{array}$$

- And one more

$$\begin{array}{r} 1\ 0\ 1\ 0 \\ +\ 1\ 1\ 0\ 1 \\ \hline \text{(Carry } 1\ 1\text{ )} \\ (1)\ 0\ 1\ 1\ 1 \end{array}$$

# Two's Complement

- Addition of a “large” number plus overflow looks like subtraction

$$+111\dots111 \rightarrow -1$$

$$+111\dots110 \rightarrow -2$$

$$+111\dots101 \rightarrow -3$$

- This leads us to the **two's complement (2C)** representation

# Two's Complement Arithmetic

- In 2C the convention is:
  - All positive numbers start with 0
  - All negative numbers start with 1
  - Negation is achieved by:
    - Flipping all the bits
    - Adding 1 to the least significant (right-most)
- 011010 (positive) → 100110 (negative)

$$\begin{array}{r} 011010 \\ + 100110 \\ \hline \text{(Carry } 11111 \text{ )} \\ (1)000000 \end{array}$$

# Have a go

- 2C is used to implement subtraction
  - Easier to calculate 2C and add than to subtract
- $0001 - 0110$ ?
- $1110 - 1001$ ?

# Range of 2C

- Given N bits, what is the largest value of a 2C number?
  - Recap:  $111\dots111 = 2^N - 1$
  - Then  $011\dots111 = 2^{N-1} - 1$
- What is the smallest (most negative) value?
  - $100\dots000 = (-)100\dots000 = -2^{N-1}$
- Example: 8 bits
  - Maximum:  $2^7 - 1 = 127$
  - Minimum:  $-2^7 = -128$

# Fixed Point Arithmetic

- Everything is the same as for whole numbers
- Example:  $01001.010 - 00010.100$
- Take 2C and add:

$$\begin{array}{r} 01001010 \\ + 11101100 \\ \hline \text{(Carry } 1 \quad 1 \quad ) \\ (1)00110110 \end{array}$$

- $00110.101 - 10110.010?$

# Floating Point

- Fixed point trades accuracy for range
- A possible alternative:

$$V = M \times 2^E$$

- Why? Powers of 2 are easy to compute in binary!
- M is the *mantissa*: 2C fixed point, one integer bit
- E is the *exponent*: 2C integer

$$0101\ 0110 = 0.101 \times 2^{0110} = 0.625 \times 64 = 04$$

- 1101 1110 = ?



# Summary

- Binary arithmetic is like decimal arithmetic
  - But we are not practised so find it hard
- Negative numbers are tricky things
  - But we can use a few of our own tricks – 2C.
- Floating point is an alternative, but is very unnatural for us

Next time, we will begin to study how computer are organised, and how they execute programs