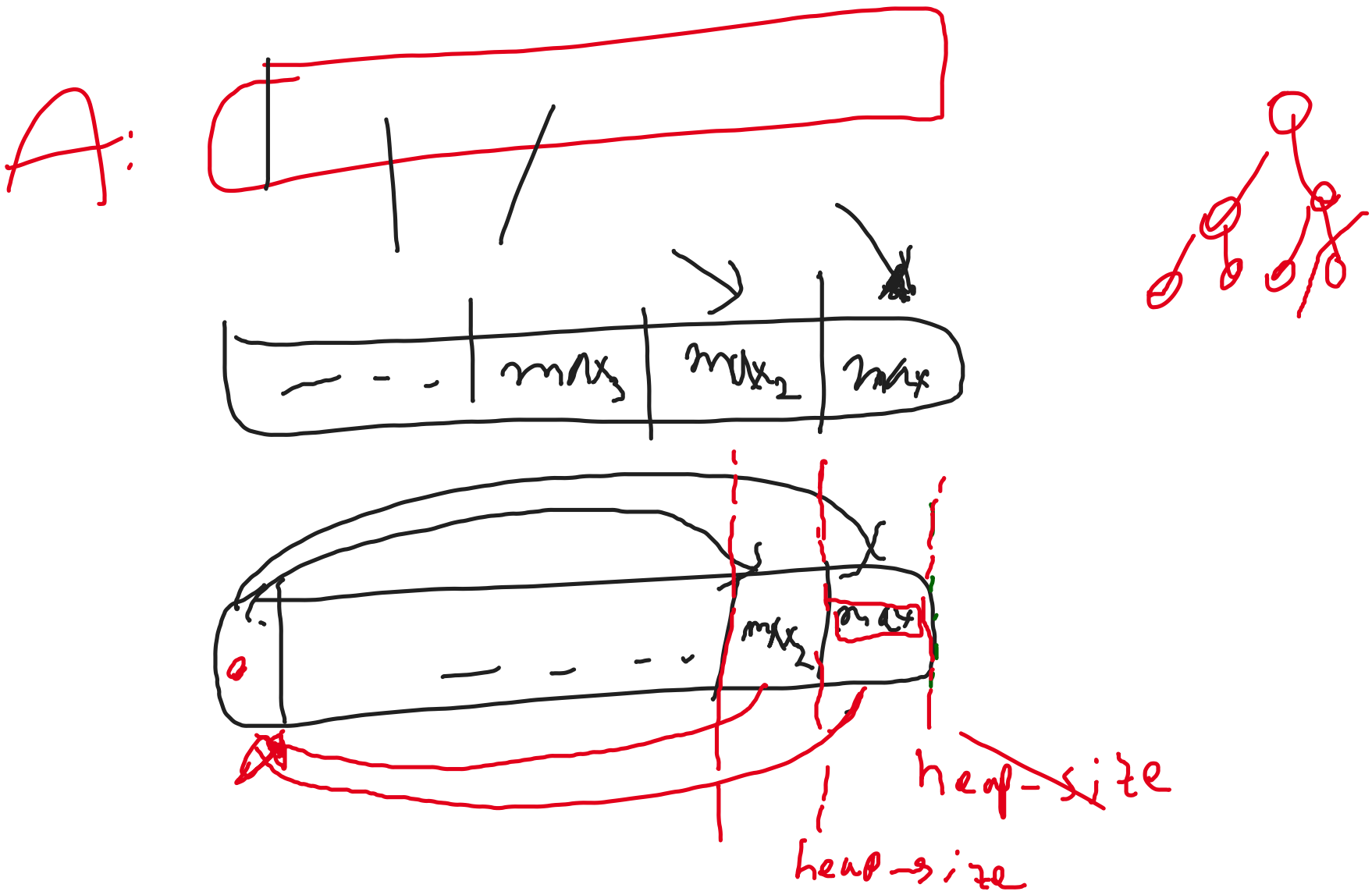
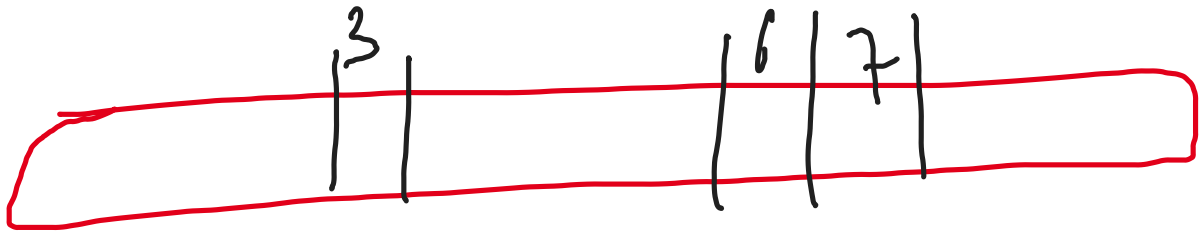
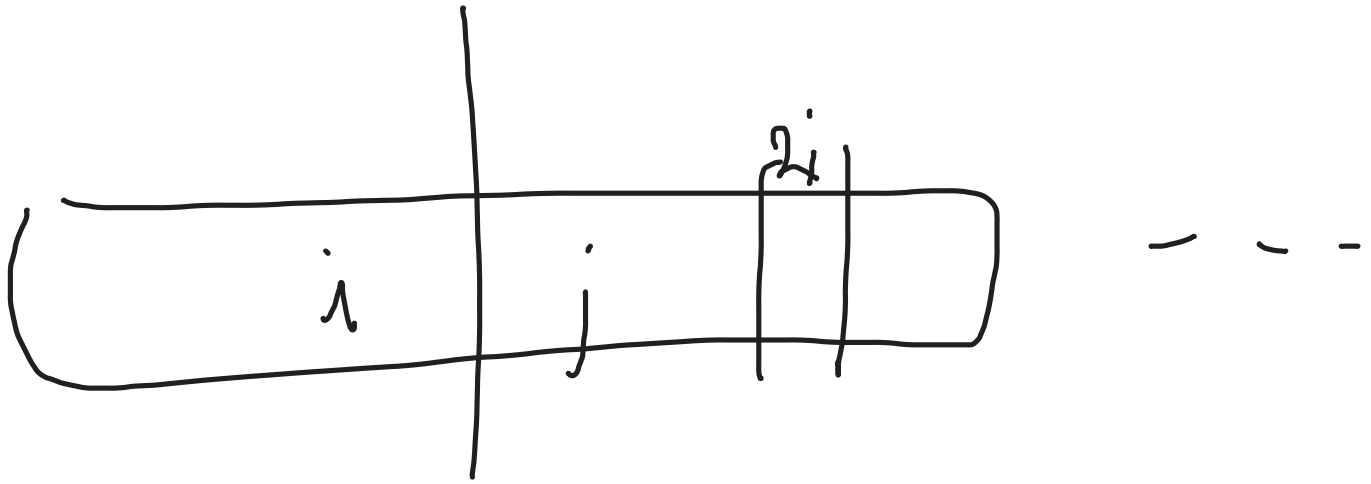

A:
heap-size
heap-size

A: 

$A[i] \geq A[2i]$

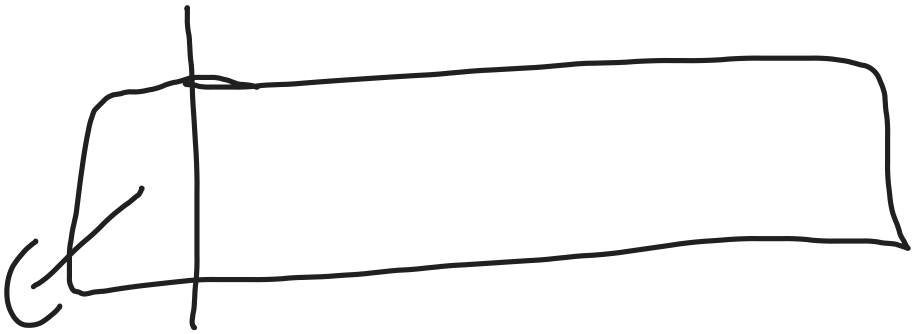
$A[i] \geq A[2i+1]$

$parent(7) = \left\lfloor \frac{7}{2} \right\rfloor = 3$



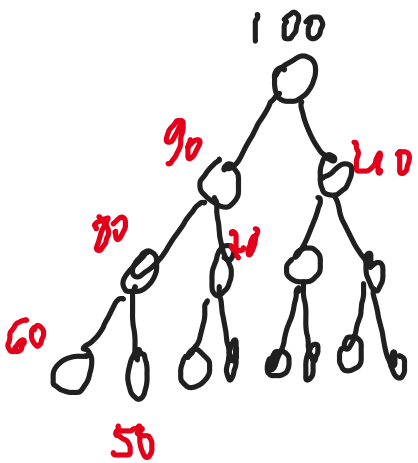
$i < \frac{n}{2} \Rightarrow 2i < n$

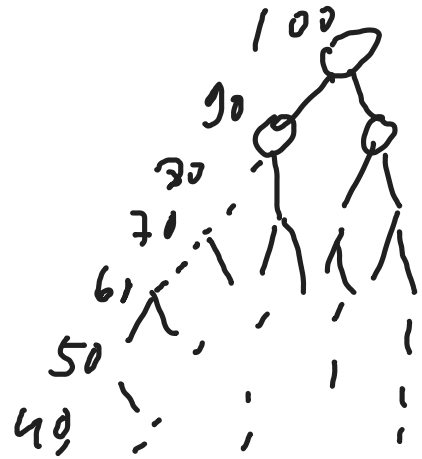
$j > \frac{n}{2} \Rightarrow 2j > n$

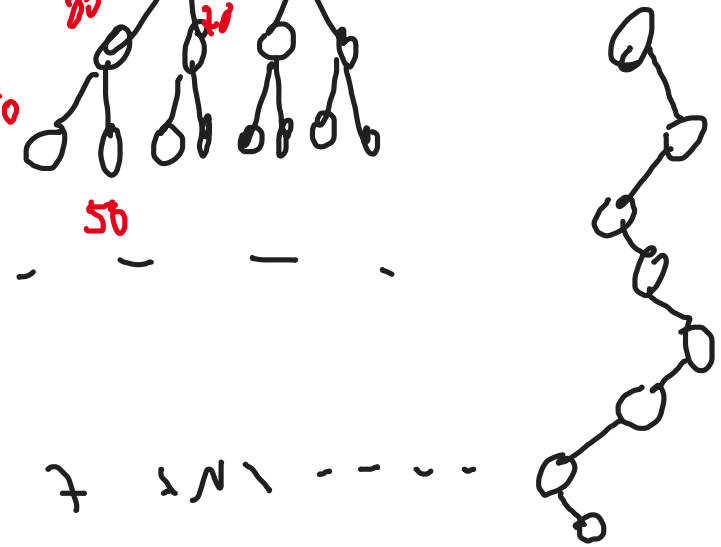


מה הגובה המינימלי של עץ בינארי עם 100 צמתים?

$$2^7 - 1 = 127$$







γ
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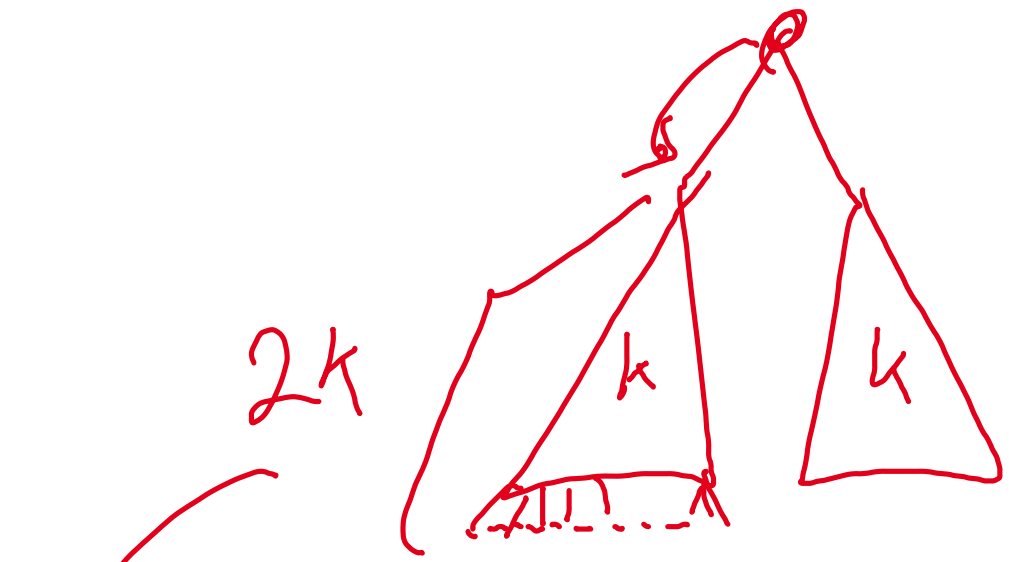
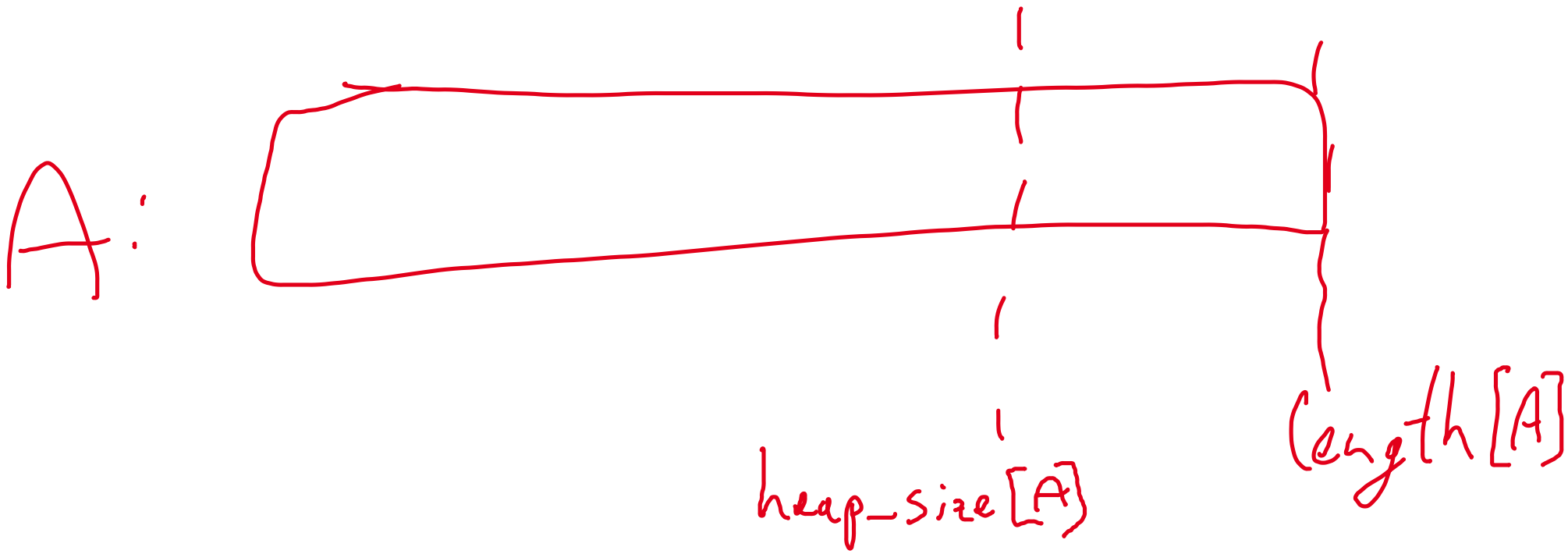
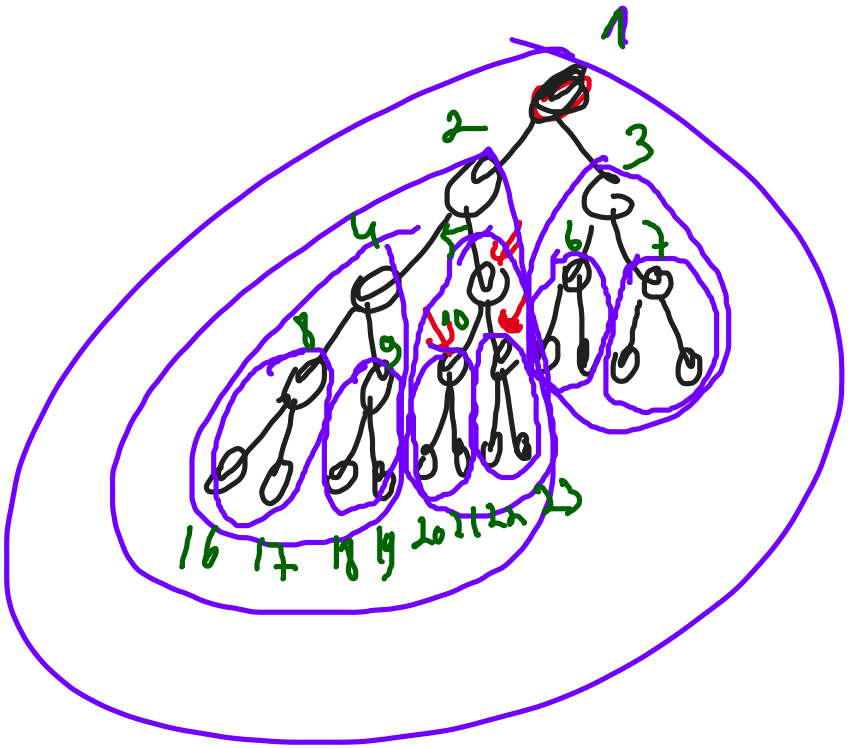


$$T(n) = T(n/2) + 1$$

$$T(n) = T(2n/3) + 1 \Rightarrow T(n) = \Theta(\log n)$$

$b = 1.5$


A:
heap_size[A]
length[A]



4 - 7

8 - 15

16 - 23

$$\left\lfloor \frac{23}{2} \right\rfloor = 11$$

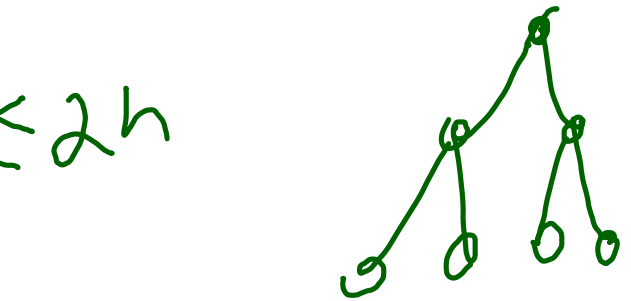
for $\frac{n}{2}$ downto 1

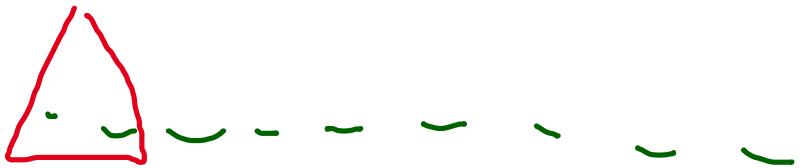
heapify(..)

$\frac{n}{2}$ קריאות לפחות היא קלט (גרוע)

$$T(n) = O(n \lg n)$$

$$\sum_i \frac{n}{2^i}(i-1) \le 2n$$





1

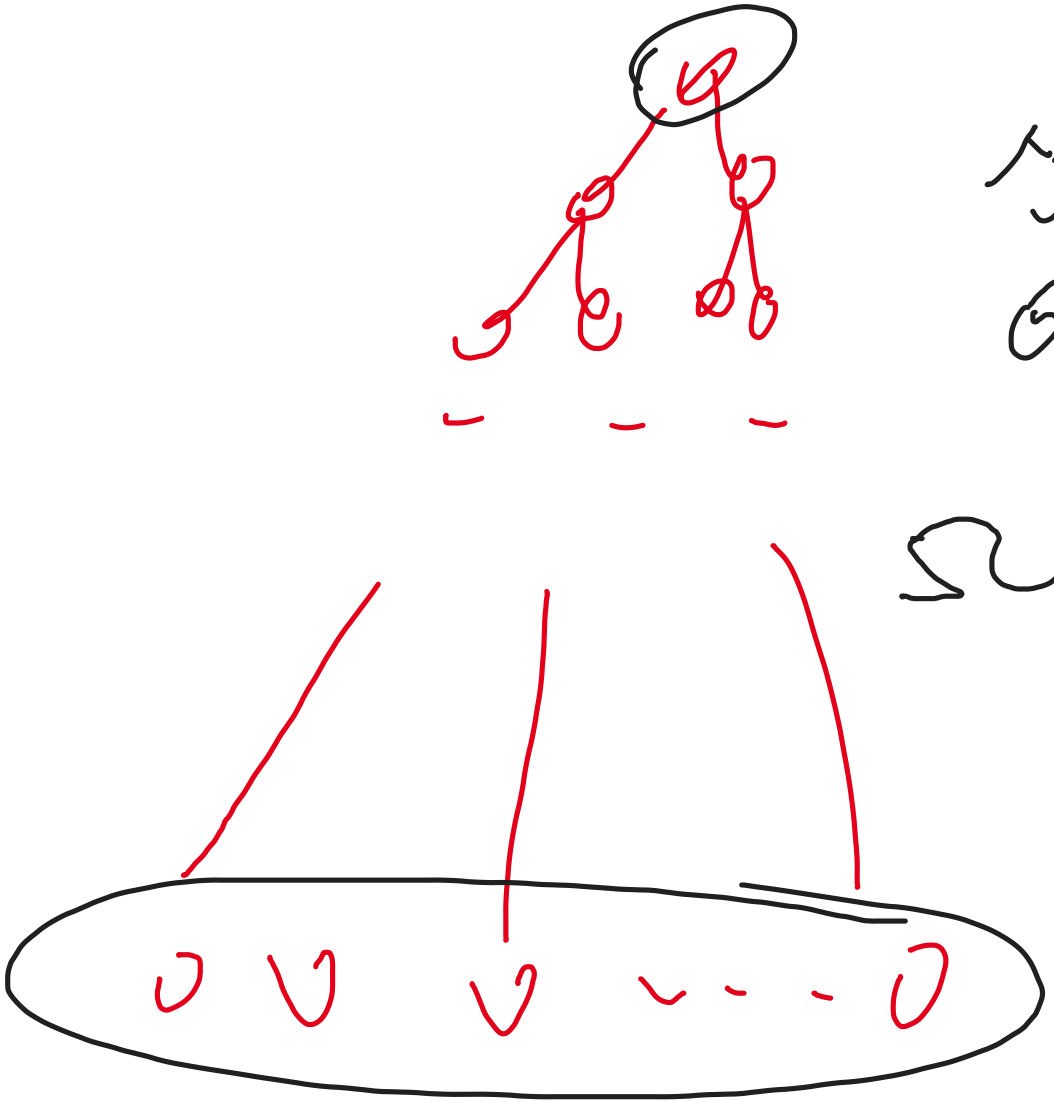
1

$\frac{n}{16} \cdot 3$

$\frac{n}{8} \cdot 2$

$\frac{n}{4} \cdot 1$

$\frac{n}{2} \cdot 0$

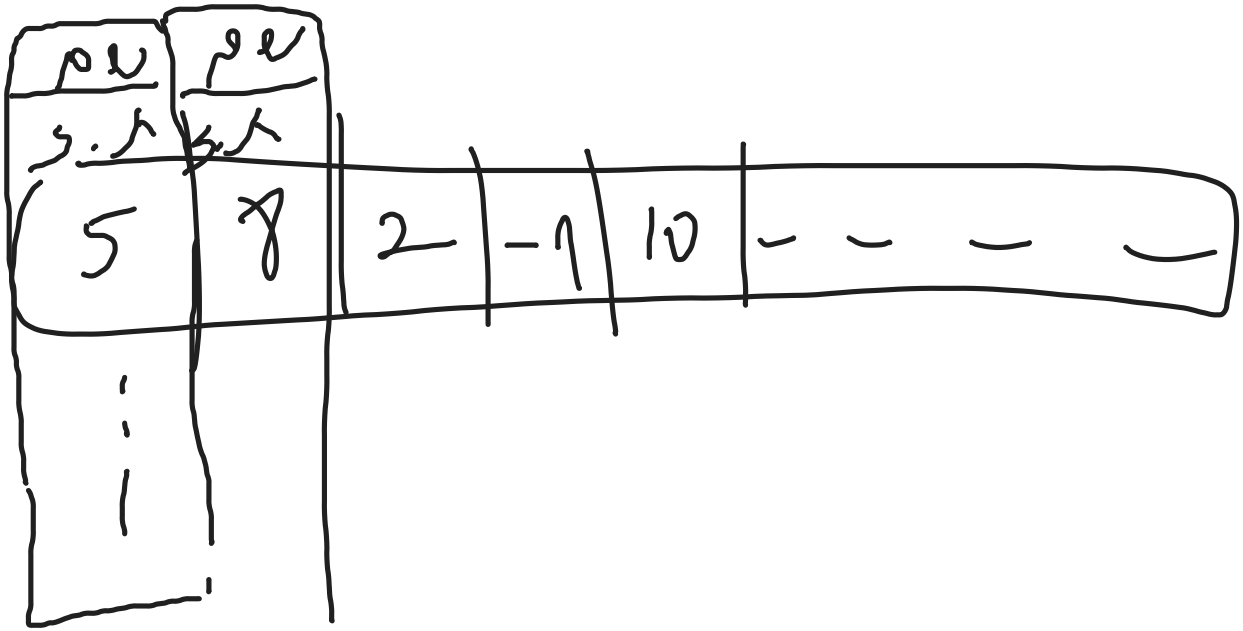


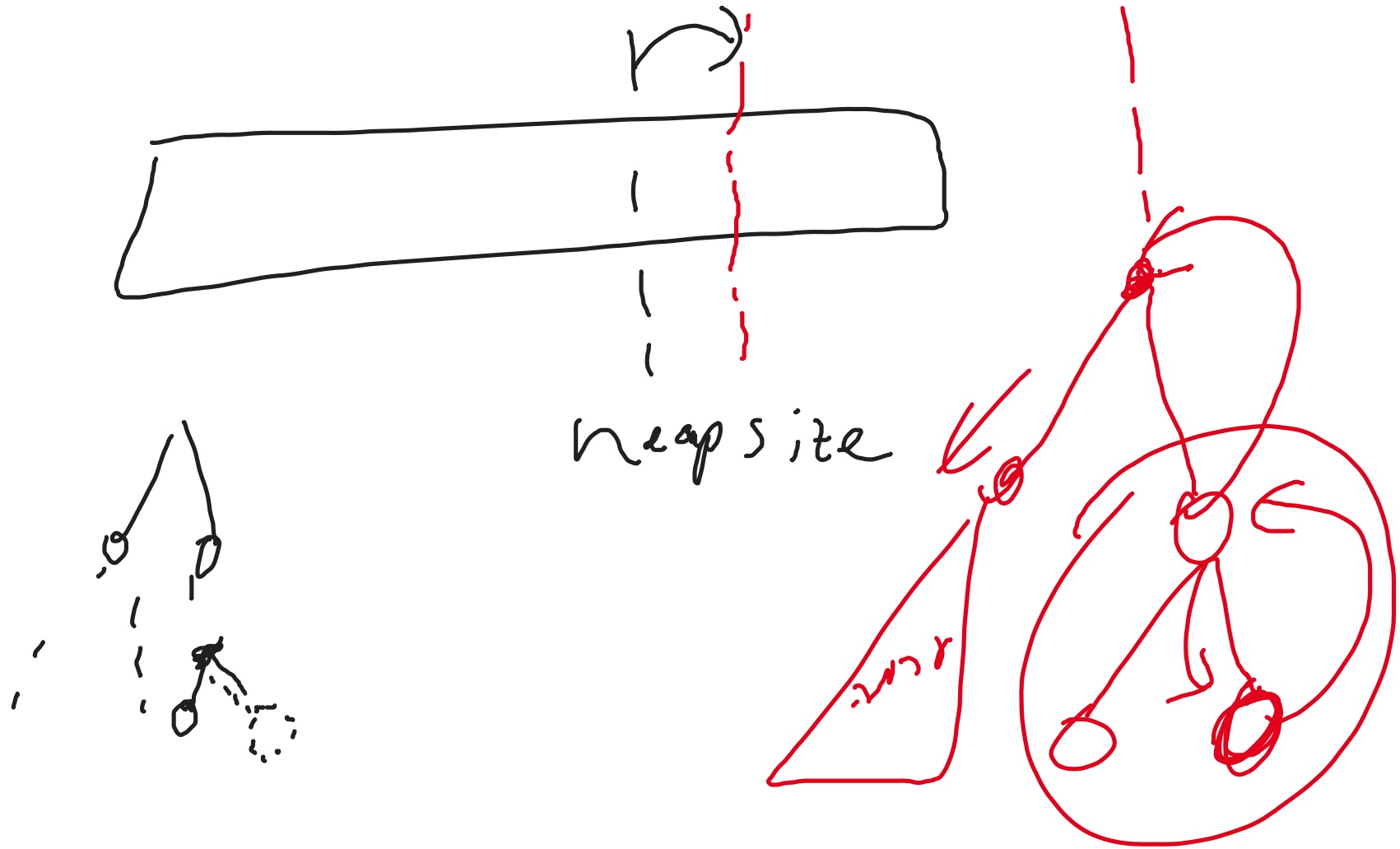
$\frac{h}{2}$ קירוב השורש

$\Theta(\log n)$ זקלו

.כ"ג

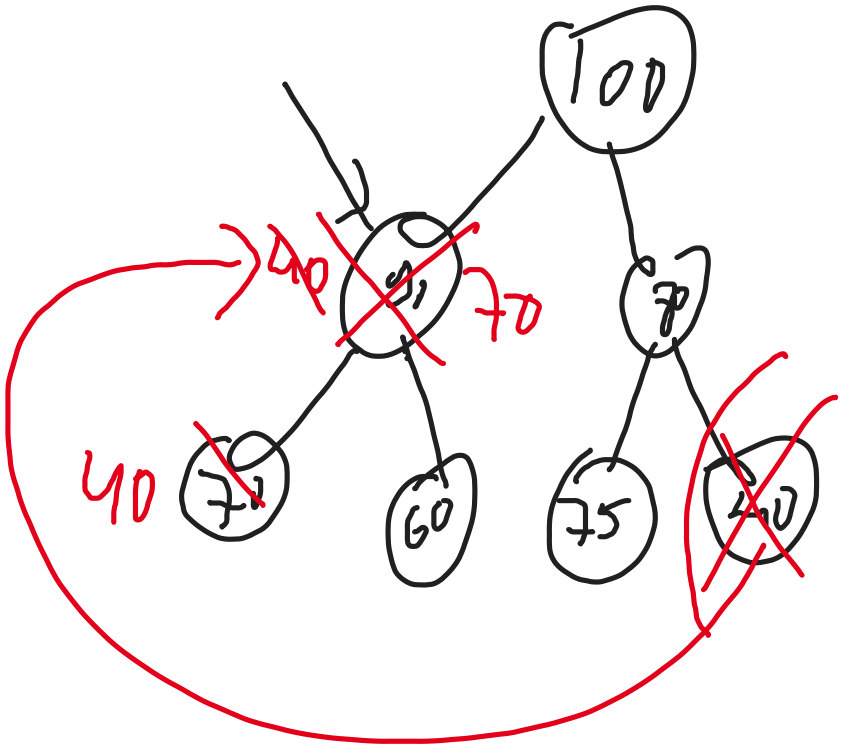
$\Omega(n \log n)$ סה"כ


5
8
2
-1
10

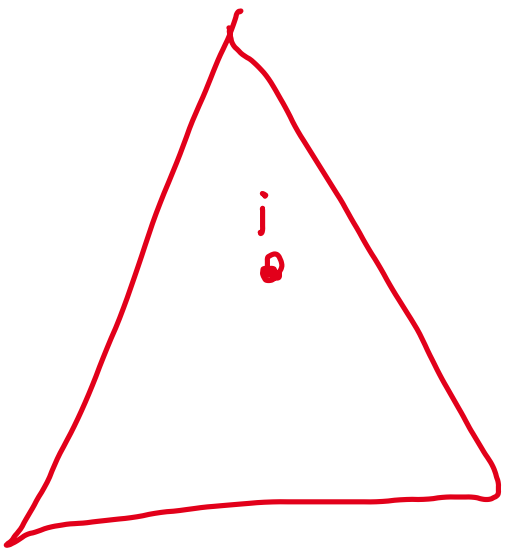

heap site

וכולה למעלה בפרה
טוב כלפי מעלה
טוב כלפי מעלה

כלפי מעלה - נתקן אם עיוקשת
כלפי מטה - דקדוקי להוסיף טוב, ולא רמתלול לכם ביותר
ורד למשוש ורעבות מיקלת לשונם.


100
40
70
40
60
75

Heap-Increase-key (A, i, key)



i — מיקום הצומת שבו

key — הערך שבו

Heap Delete(A, i)

Increase(i, ∞)

Extract Max(..)

כנ"ל עבור $\Theta(m)$ –

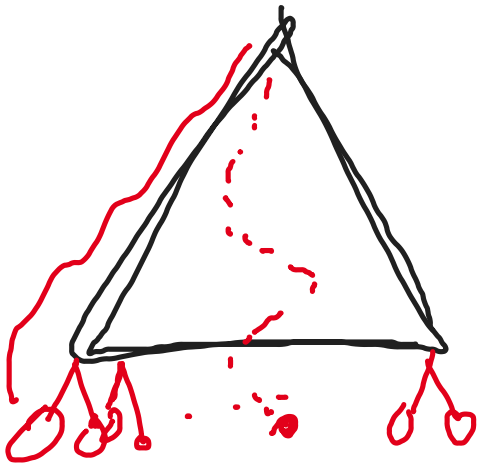
מדובר בקלט בגודל $n-m$:

$$\Theta\left[(n-m)+m\right]$$

הערה 1.9

טענה (משפט קנטור) לכל קבוצה A

ב. ניתן שפשוף הנקרא $\frac{2}{3}$ מהקודם:



בהדגשה A בגודל $\mathcal{P}(A)$ גדולה מ־A