



```
#include <stdlib.h>
#include <string.h>
#include <ctype.h>

#define MAXPAROLA 30
#define MAXRIGA 80

int main(int argc, char *argv[])
{
    int freq[MAXPAROLA]; /* vettore di contatori
delle frequenze delle lunghezze delle parole */
    char riga[MAXRIGA];
    int i, inizio, lunghezza;
    FILE *f;

    for(i=0; i<MAXPAROLA; i++)
        freq[i]=0;

    if(argc != 2)
    {
        fprintf(stderr, "ERRORE: serve un parametro con il nome del file\n");
        exit(1);
    }
    f = fopen(argv[1], "r");
    if(f==NULL)
    {
        fprintf(stderr, "ERRORE: impossibile aprire il file %s\n", argv[1]);
        exit(1);
    }

    while( fgets( riga, MAXRIGA, f ) != NULL )
```

Recursion

Theory Aspects of Recursion

Stefano Quer

Dipartimento di Automatica e Informatica
Politecnico di Torino

The stack

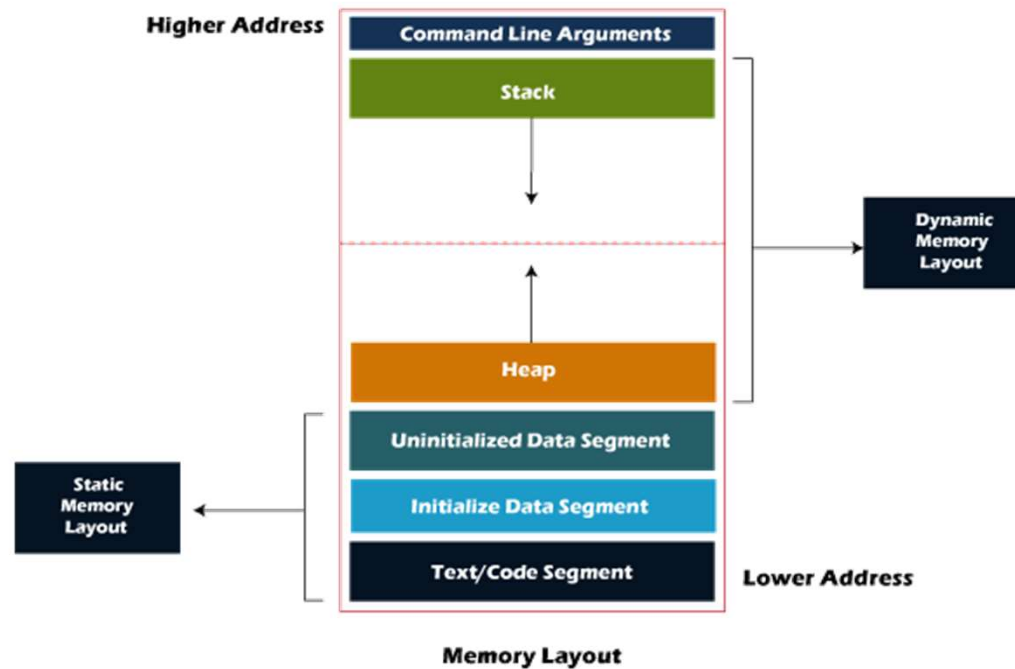
- ❖ In computer science a stack is an Abstract Data Type (ADT) that serves as a collection of elements

The **stack** was introduced in the unit on lists

- ❖ A stack supports the following operations
 - Push: Inserts objects on top
 - Pop: Extracts objects from the top
 - As both insertions and extractions are performed at the same end of the ADT, the stack follows a LIFO (Last-In First-Out) strategy

The stack

- ❖ A programmer can implement its own stacks
- ❖ The operating system (or any application) can use its own stack as well



The stack

- ❖ For a C compiler, the stack is the data structure containing at least
 - Formal parameters
 - Local variables
 - The return address when the function execution is over
 - The pointer to the function's code

The stack

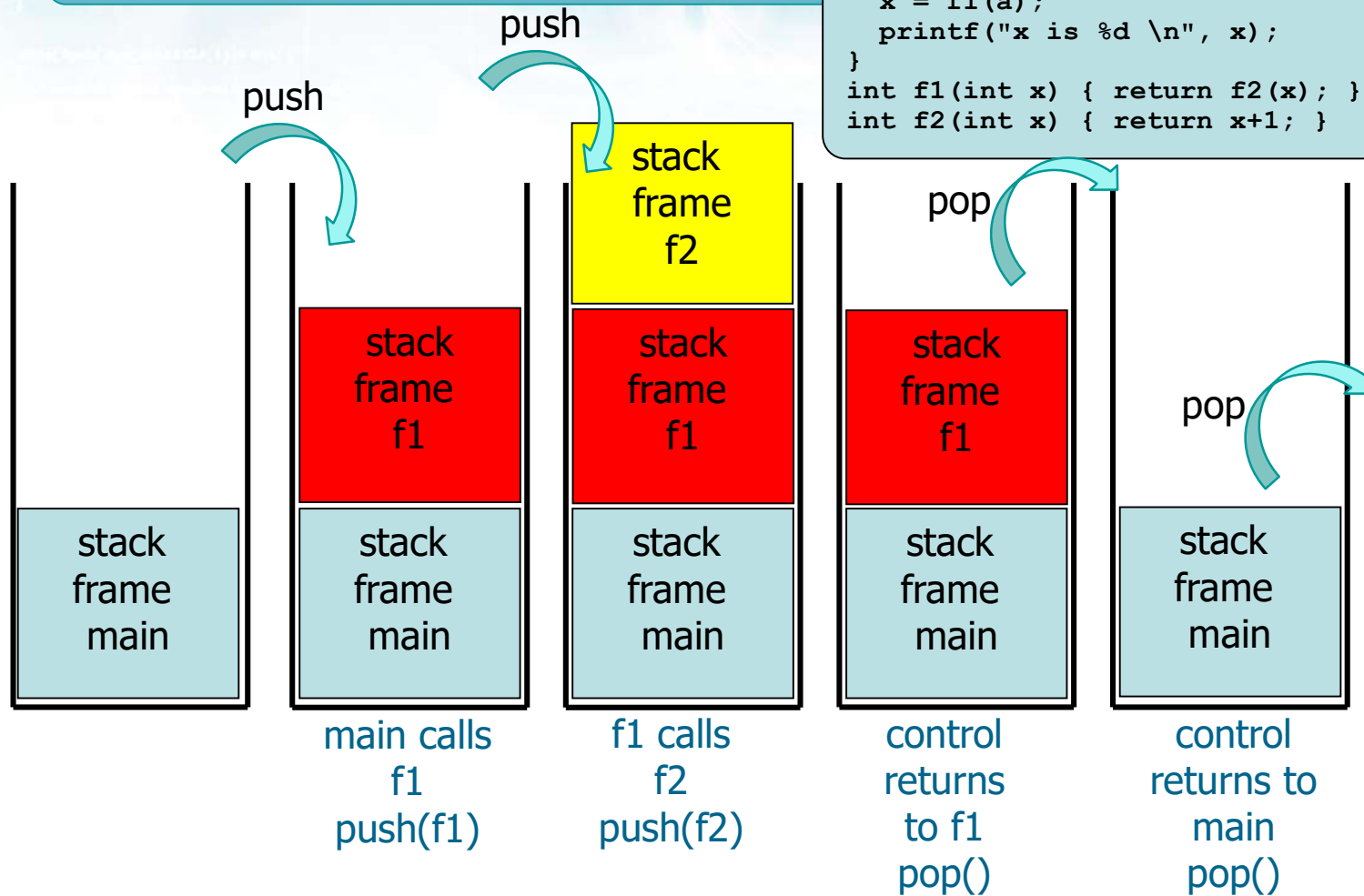
- ❖ All these pieces of data form a **stack frame**
 - A new stack frame is created when the function is called and the same stack frame is destroyed when the function is over
- ❖ Stack frames are stored in the system stack
 - The system stack has a predefined amount of memory available
 - When it goes beyond the space allocated to it, a **stack overflow** occurs
 - The stack grows from larger to smaller addresses (thus upwards)
 - The **stack pointer SP** is a register containing the address of the first available stack frame

Example

- ❖ Let us analyze the stack structure during the execution of the following program

```
int f1(int x);  
int f2(int x);  
  
main() {  
    int x, a = 10;  
    x = f1(a);  
    printf("x is %d \n", x);  
}  
int f1(int x) {  
    return f2(x);  
}  
int f2(int x) {  
    return x+1;  
}
```

```
int f1(int x);  
int f2(int x);  
main() {  
    int x, a = 10;  
    x = f1(a);  
    printf("x is %d \n", x);  
}  
int f1(int x) { return f2(x); }  
int f2(int x) { return x+1; }
```

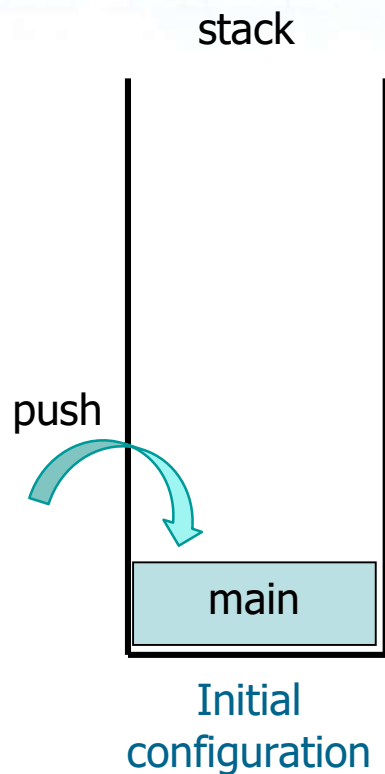


Recursive functions

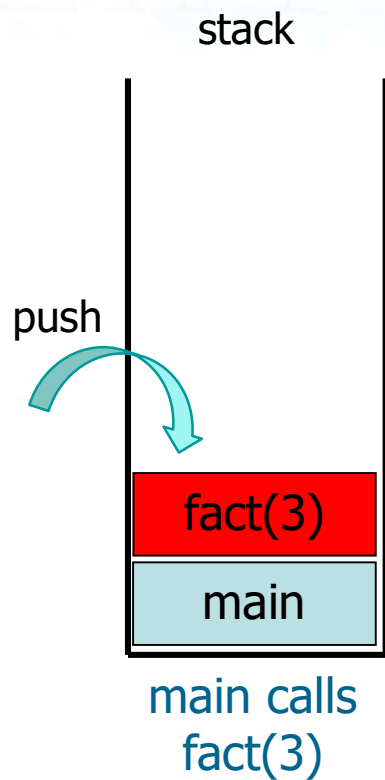
- ❖ With recursive functions
 - Calling and called functions coincide, but operate on different data
 - The system stack is used as in any other function call
- ❖ Too many recursive calls may result in a **stack overflow**

Example 1

```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```



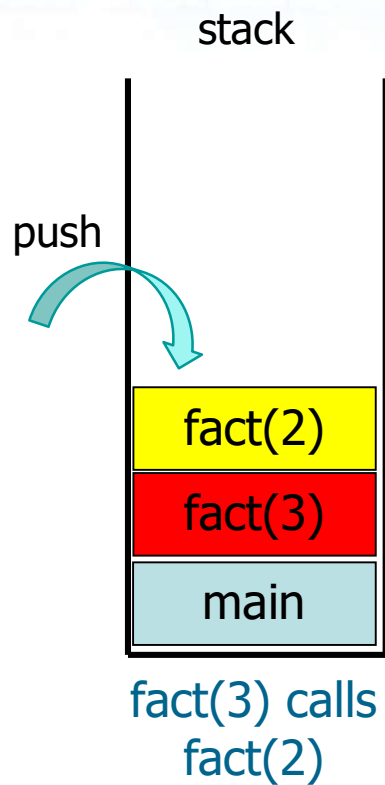
Example 1



```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

$3! = 3 * 2!$

Example 1



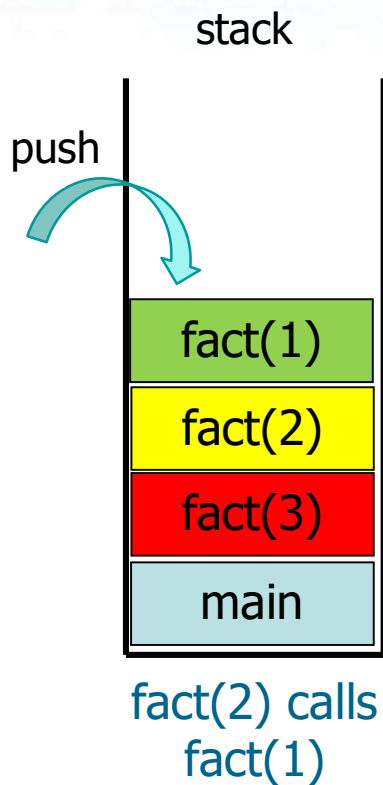
```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

$$3! = 3 * 2!$$



$$2! = 2 * 1!$$

Example 1



```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

$$3! = 3 * 2!$$



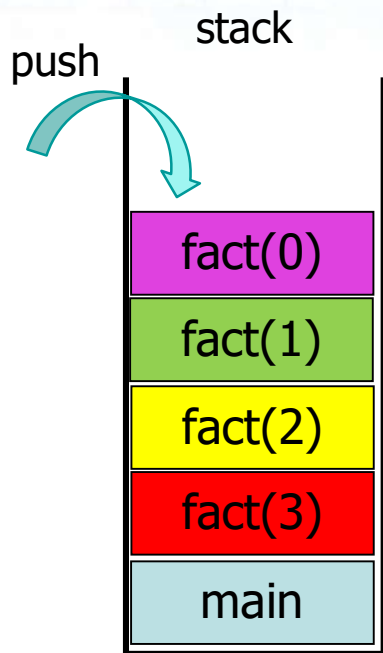
$$2! = 2 * 1!$$



$$1! = 1 * 0!$$

Example 1

Hit the termination condition



fact(1) calls
fact(0)

```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

$$3! = 3 * 2!$$



$$2! = 2 * 1!$$

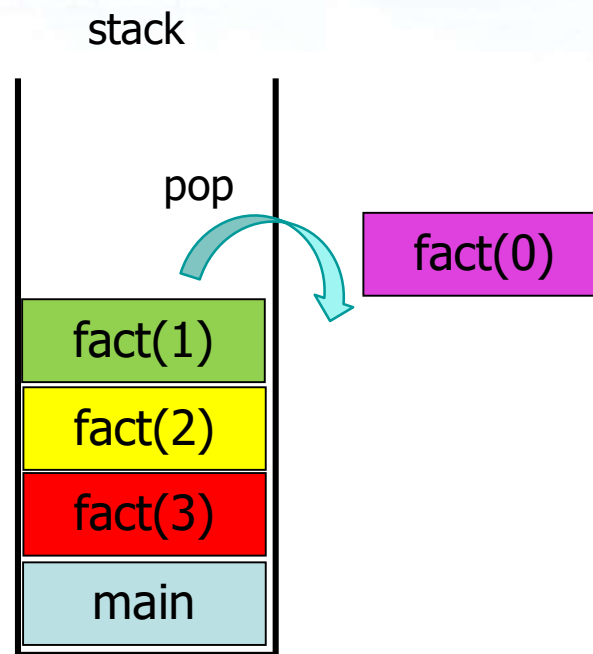


$$1! = 1 * 0!$$



$$0! = 1$$

Example 1



fact(0) terminates,
returns value 1 and
returns control
to fact(1)

```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

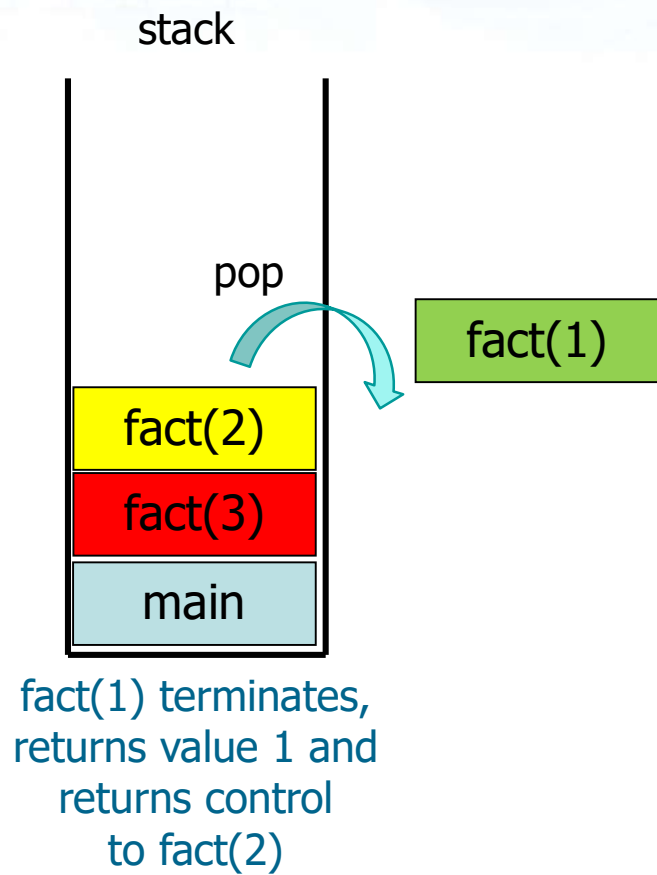
$$3! = 3 * 2!$$

$$2! = 2 * 1!$$

$$1! = 1 * 0!$$

$$0! = 1$$

Example 1



```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

$$3! = 3 * 2!$$

$$2! = 2 * 1!$$

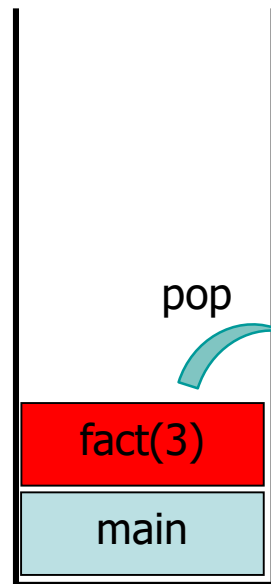
$$1! = 1 * 0! = 1$$

$$0! = 1$$

Example 1

```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

stack



fact(2)

fact(2) terminates,
returns value 2 and
returns control
to fact(3)

$$3! = 3 * 2!$$

$$2! = 2 * 1! = 2$$

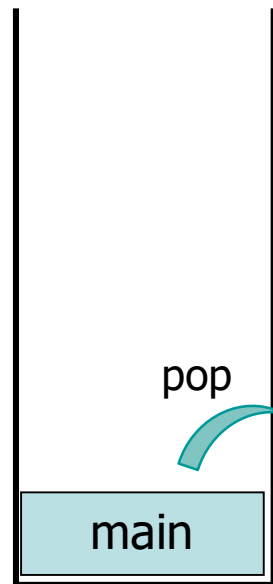
$$1! = 1 * 0! = 1$$

$$0! = 1$$

Example 1

```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

stack



pop



fact(3)

fact(3) terminates,
returns value 6 and
returns control
to main

$$3! = 3 * 2! = 6$$



$$2! = 2 * 1! = 2$$



$$1! = 1 * 0! = 1$$

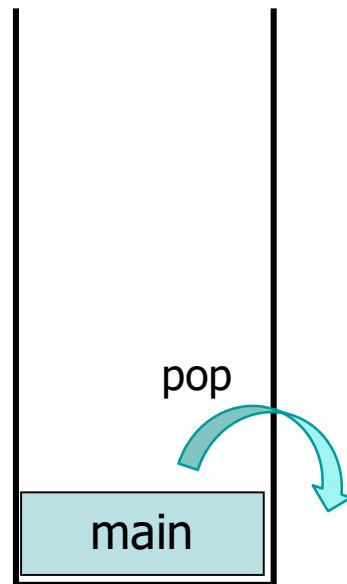


$$0! = 1$$

Example 1

```
main() {  
    long n;  
    printf("Input n: ");  
    scanf("%d", &n);  
    printf("%d %d \n", n, fact(n));  
}  
long fact(long n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

stack



fact(3) terminates,
returns value 6 and
returns control
to main

$$\begin{array}{c} 3! = 3 * 2! = 6 \\ \downarrow \uparrow \\ 2! = 2 * 1! = 2 \\ \downarrow \uparrow \\ 1! = 1 * 0! = 1 \\ \downarrow \uparrow \\ 0! = 1 \end{array}$$

Recursion versus iteration

- ❖ Recursion
 - May be memory-consuming
 - Is somehow equivalent to looping
- ❖ All recursive programs may be implemented in iterative form as well
 - There is a duality between recursion and iteration
- ❖ The best solution (efficiency and clarity of code) depends on the problem
- ❖ Try to remain at the highest possible abstraction level

Duality recursion - iteration

❖ Factorial iterative computation

- $5! = 1*2*3*4*5 = 120$
- The implementation may be iterative and recursive as well
- There is no need to use a stack

```
long fact(long n) {  
    long tot = 1;  
    int i;  
  
    for (i=2; i<=n; i++)  
        tot = tot * i;  
  
    return(tot);  
}
```

Duality recursion - iteration

❖ Fibonacci iterative computation

➤ 0 1 1 2 3 5 8 13 21 ...

- $F(0) = 0$
- $F(1) = 1$
- $F(2) = F(0) + F(1) = 1$
- $F(3) = F(1) + F(2) = 2$
- $F(4) = F(2) + F(3) = 3$
- $F(5) = F(3) + F(4) = 5$

➤ The implementation may be iterative and recursive as well

➤ There is no need to use a stack

```
long fib(long int n) {  
    long int f1p=1, f2p=0, f;  
    int i;  
    if(n == 0 || n == 1)  
        return(n);  
    f = f1p + f2p;  
    for(i=3; i<= n; i++) {  
        f2p = f1p;  
        f1p = f;  
        f = f1p+f2p;  
    }  
    return(f);  
}
```

Duality recursion - iteration

❖ Binary search

- The implementation may be iterative and recursive as well
- There is no need to use a stack

```
int BinarySearch (  
    int v[], int l, int r, int k) {  
    int c;  
  
    while (l<=r){  
        c = (int) ((l+r) / 2);  
        if (k == v[c]) {  
            return(c);  
        }  
        if (k < v[c]) {  
            r = c-1;  
        } else {  
            l = c+1;  
        }  
    }  
    return(-1);  
}
```

Emulating recursion

- ❖ Recursion may be emulated explicitly dealing with a stack
 - Recursion is realized using the system stack to store and restore the local status
 - It is always possible to emulate recursion through iterations using a user-defined stack
 - The user stack mimics the system stack
 - It is manipulated by the programmer to store and restore information (function stack frames) as the system does into the system stack

Emulating recursion

```
long fact(long int n) {  
    if(n == 0)  
        return(1);  
    return(n * fact(n-1));  
}
```

Original recursive
function

Non recursive
(iterative) function
emulating recursion
using a user stack

```
long fact(long int n) {  
    long fact = 1; int status;  
    list_t *top;  
    top = NULL;  
    while (n>0) {  
        top = push (top, n);  
        n--;  
    }  
    do {  
        top = pop (top, &n, &status);  
        if (status!=FAILURE)  
            fact = n * fact;  
    } while (status!=FAILURE);  
    return fact;  
}
```

The ADT stack is
a user-defined
list_t object

Tail-recursive functions

- ❖ In the traditional recursive functions
 - Recursive calls are performed first
 - Then, the return value is used to compute the result
 - The final result is obtained after all calls have terminated, i.e., the program has returned from every recursive call
- ❖ Tail-recursion (or tail-end recursion) is a particular case of recursion

Tail-recursive functions

- ❖ In tail recursive function, the recursive call is the last operation to be executed, except for return

```
long fact(long int n) {  
    if (n == 0)  
        return(f);  
    return fact(n*fact(n-1));  
}
```

This function is not tail-recursive because the product can be executed only after returning from the recursive call

The system stack is required

```
fact(3)  
3 * fact(2)  
3 * (2 * fact(1))  
3 * (2 * (1 * fact(0)))  
3 * (2 * (1 * 1))
```

Tail-recursive functions

❖ Tail-recursive version of the factorial function

```
long fact(long int n) {  
    if (n == 0)  
        return(f);  
    return fact(n*fact(n-1));  
}
```

This function is tail-recursive because the product is executed before the recursive call

The system stack is **not** required

```
fact_tr(3,1)  
fact_tr(2,3)  
fact_tr(1,6)  
fact_tr(0,6)
```

The caller sets $n = 1$

```
long fact_tr(long n, long f){  
    if (n==0)  
        return f;  
    return fact_tr(n-1,n*f);}
```

Tail-recursive functions

❖ In tail recursive functions

- Calculations are performed first
- Recursive calls are done after
- Current results are passed to future calls
 - The return value of any given recursive step is the same as the return value of the next recursive call
 - The consequence of this is that once you are ready to perform your next recursive step, you do not need the current stack frame any more

Tail-recursive functions

- Current stack frame is not needed anymore
 - Recursion can be substituted by a simple jump (**tail call elimination**)
 - A proper compiler or language (Prolog, Lisp, etc.) may recognize tail recursive functions and it may optimize their code
 - Stack overflows does not happen anymore and there is no limit to the number of recursive calls that can be made
- Tail recursion is essentially equivalent to looping
- Tail recursion only applies if there are no instructions that follow the recursive call

Solution

```
void print (char *s) {  
    if (*s == '\0') {  
        return;  
    }  
    printf ("%c", *s);  
    print (s+1);  
    return;  
}
```

Printing a string:
There are no instructions that follow the recursive call. The compiler may understand this and it may avoid the stack. This function is tail recursive.

```
void reverse_print (char *s) {  
    if (*s == '\0') {  
        return;  
    }  
    reverse_print(s+1);  
    printf ("%c", *s);  
    return;  
}
```

Reverse printing a string:
There are instructions that follow the recursive call. The stack cannot be avoided. This function is not tail recursive.

Limits of the recursion

❖ Disadvantages

- The number of recursions is limited by the stack size
 - The stack consume memory
- Sub-problems may not be independent, and recomputations may occur leading to inefficiency

