

# Exact model reduction for Quantum Systems

- An algebraic approach -

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DI PADOVA

What are the **minimal resources** needed to reproduce a target quantum process?

**Resources?** Mathematical or physical degrees of freedom (“memory”).

# Why is it interesting?

- Model reduction for **quantum dynamics**;
- Efficient quantum simulation (quantum systems are notoriously difficult to simulate);
- Efficient implementations of **controllers, error suppression schemes and quantum filters** ;
- Easier models to study;
- Proving optimality of quantum algorithms;
- Probing “quantumness” of processes;
- Efficient generation of quantum trajectories (Montecarlo methods).

# The model (discrete-time quantum walk)

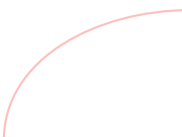
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$$\mathcal{B}(\mathcal{H}) = \mathbb{C}^{n \times n}$$

$\rho$  are density operators:

$$\rho \in \mathbb{C}^{n \times n}, \rho = \rho^\dagger \geq 0, \text{tr}[\rho] = 1$$


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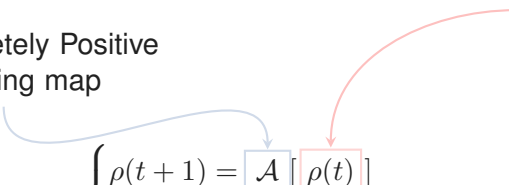
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A blue arrow points from the text " $\mathcal{A}$  is a Completely Positive Trace Preserving map" to the box around  $\mathcal{A}$  in the equation. A red arrow points from the text " $\rho$  are density operators:" to the box around  $\rho(t)$  in the equation.

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$\mathcal{S}$  is a set of initial conditions

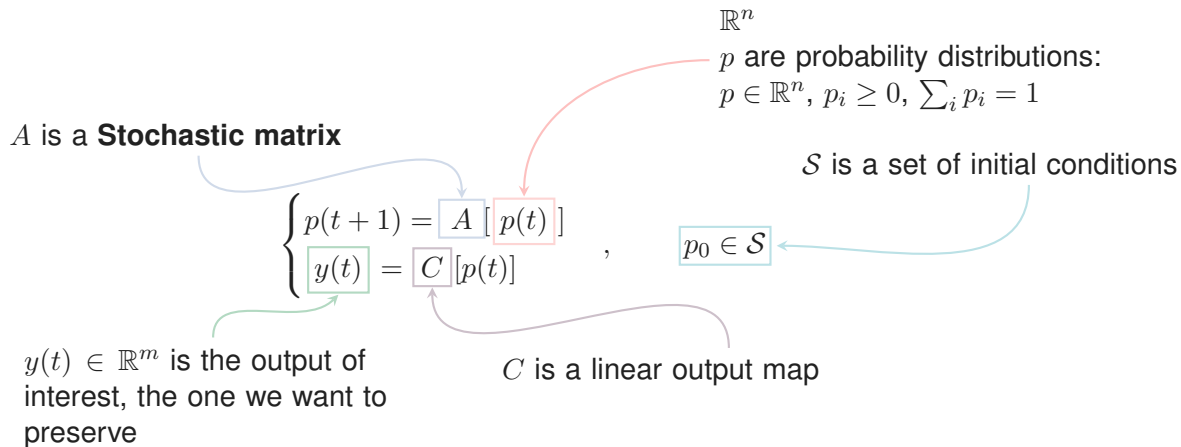
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# The classical counterpart (Markov Chains)



# The problem: quantum model reduction

Given a QS  $(\mathcal{A}, \mathcal{S}, \mathcal{C})$  find **another QS**  $(\check{\mathcal{A}}, \check{\mathcal{S}}, \check{\mathcal{C}})$  and a linear map  $\Phi : \mathbb{C}^{n \times n} \rightarrow \mathbb{C}^{r \times r}$

such that for all  $t \geq 0$  and  $\rho_0 \in \mathcal{S}$ ,  $\check{\rho}_0 = \Phi[\rho_0]$

- **exact model reduction**

$$\mathcal{C}[\mathcal{A}^t[\rho_0]] = \check{\mathcal{C}}[\check{\mathcal{A}}^t[\check{\rho}_0]];$$

- approximate model reduction (future work)

$$\mathcal{C}[\mathcal{A}^t[\rho_0]] \approx \check{\mathcal{C}}[\check{\mathcal{A}}^t[\check{\rho}_0]].$$

# First step. Assume:

$$\ker[\mathcal{C}] \neq \{0\},$$

$$\rho_0 \in \mathcal{S} = \mathfrak{D}(\mathcal{H}).$$

We are only interested in reproducing trajectories of certain values of interest starting from any initial condition.

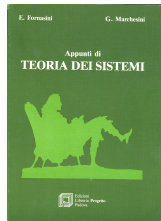
# Linear model reduction

$$\begin{cases} \rho(t+1) = \mathcal{A}[\rho(t)] \\ Y(t) = \mathcal{C}[\rho(t)] \end{cases}$$

$$\mathcal{N} = \{X \in \mathbb{C}^{n \times n} \mid \mathcal{C}[\mathcal{A}^t[X]] = 0, \forall t \geq 0\}$$

$\mathcal{N}$  is the largest  $\mathcal{A}$ -invariant operators subspace contained in  $\ker[\mathcal{C}]$ .

$$\begin{cases} \begin{bmatrix} \rho_{\mathcal{N}} \\ \rho_{\mathcal{N}^\perp} \end{bmatrix} (t+1) = \begin{bmatrix} \mathcal{A}_{\mathcal{N},\mathcal{N}} & \mathcal{A}_{\mathcal{N},\mathcal{N}^\perp} \\ 0 & \mathcal{A}_{\mathcal{N}^\perp,\mathcal{N}^\perp} \end{bmatrix} \begin{bmatrix} \rho_{\mathcal{N}} \\ \rho_{\mathcal{N}^\perp} \end{bmatrix} (t) \\ Y(t) = \begin{bmatrix} 0 & C_{\mathcal{N}^\perp} \end{bmatrix} \begin{bmatrix} \rho_{\mathcal{N}} \\ \rho_{\mathcal{N}^\perp} \end{bmatrix} (t) \end{cases}$$



# Linear model reduction pt. 2

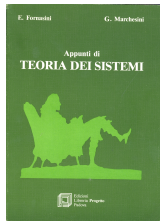
This is equivalent to finding factors  $\mathcal{R}$  and  $\mathcal{J}$  such that  $\mathcal{J}\mathcal{R} = \Pi_{\mathcal{N}^\perp}$  and  $\mathcal{R}\mathcal{J} = \mathcal{I}$  so that

$$\check{\mathcal{A}} = \mathcal{R}\mathcal{A}\mathcal{J}, \quad \check{\mathcal{C}} = \mathcal{C}\mathcal{J}, \quad \Phi = \mathcal{R}.$$

**Note for later:**

Restricting onto  $\mathcal{N}^\perp$  provides the minimal linear equivalent model.

Nonetheless, we can include some variables from  $\mathcal{N}$  and the reduced model is still equivalent.



**Problem:**

$\check{A}$  is not necessarily CPTP

How do we ensure that physical (probability) constraints are satisfied?

**The protagonist of this work:**  
**\*-algebras**



# Definition

We define a  $*$ -algebra  $\mathcal{A}$  as an operator space closed under matrix multiplication and adjoint action.

$$X, Y \in \mathcal{A} \quad \Rightarrow \quad X + Y \in \mathcal{A} \quad X^\dagger, Y^\dagger \in \mathcal{A} \quad \text{and} \quad XY \in \mathcal{A}$$

It is the fundamental mathematical structure that supports a **quantum probability space**.

Classically you can see an algebra as the set of random variables that are measurable w.r.t. a given  $\sigma$ -algebra.

# Conditional expectations & state extensions

A *conditional expectation*  $\mathbb{E}_{\mathcal{A},\rho}[\cdot]$  is a **CP unital proj. onto a \*-algebra  $\mathcal{A}$**  (and such that  $\text{tr}[\rho \mathbb{E}_{\mathcal{A},\rho}[X]] = \text{tr}[\rho X]$  for all  $X \in \mathcal{B}(\mathcal{H})$ ).

Its dual is called a *state extension*  $\mathbb{J}_{\mathcal{A},\rho}[\cdot] = \mathbb{E}_{\mathcal{A},\rho}^\dagger[\cdot]$  is a **CPTP projection onto an operator subspace  $\mathcal{X}$**  (and is such that  $\mathbb{J}_{\mathcal{A},\rho}[\rho] = \rho$ ).

$\mathcal{X}$  is an operator subspace closed under adjoint action and a modified matrix product  $X \cdot_\rho Y = X\rho^{-1}Y$  and can be computed as  $\mathcal{X} = \rho^{1/2}\mathcal{A}\rho^{1/2}$ . We call this a **distorted algebra**.

# Properties of state extensions (new!)

$\mathbb{J}_{\mathcal{A},\rho}[\cdot]$  can be factorized in two non-square CPTP isometries  $\mathbb{J}_{\mathcal{A},\rho}[\cdot] = \mathcal{J}\mathcal{R}$ .

Given a CPTP map  $\mathcal{A}$ , its reduction onto the distorted algebra  $\mathcal{X}$ ,  $\mathcal{A}|_{\mathcal{X}} = \mathcal{R}\mathcal{A}\mathcal{J}$  is CPTP.

# Take home ideas

1. We need algebras to define a quantum probability space;
2. Projections onto distorted algebras provide CPTP reduction.

# Observable CPTP reduction

We leverage the knowledge of the quantities of interest to reduce the model.

1. Find the **non-observable subspace**

$$\mathcal{N} = \{X \in \mathbb{C}^{n \times n} \mid \mathcal{C}[\mathcal{A}^t[X]] = 0, \forall t \geq 0\}$$

2. Close its orthogonal complement to a **\*-algebra**  $\mathcal{A} = \text{alg}(\mathcal{N}^\perp)$ ;
3. Compute the **state extension**  $\mathbb{J}_{\mathcal{A}, I/n}[\cdot]$  and its CPTP factors  $\mathcal{R}$  and  $\mathcal{J}$ ;
4. Use the two factors to **reduce the dynamics**:

$$\check{\mathcal{A}} = \mathcal{R}\mathcal{A}\mathcal{J}, \quad \check{\mathcal{C}} = \mathcal{C}\mathcal{J}, \quad \Phi = \mathcal{R}.$$

**Main result:** The closure of  $\mathcal{N}^\perp$  (w.r.t.  $\langle \cdot, \cdot \rangle_{HS}$ ) to an algebra is “optimal”.

## Second step. Assume:

$$\mathcal{C} = \mathcal{I},$$

$$\rho_0 \in \mathcal{S} \subsetneq \mathfrak{D}(\mathcal{H}).$$

We are interested in reproducing the evolution of the entire state starting only from certain initial conditions.

# Reachable reduction: A naive approach

We leverage the knowledge of  $\rho_0 \in \mathcal{S}$  to reduce the model.

1. Find the **minimal subspace that contains the dynamics**

$$\mathcal{R} = \text{span}\{\mathcal{A}^t(\rho_0), t \geq 0\};$$

2. Close it to a **\*-algebra**  $\mathcal{A} = \text{alg}(\mathcal{R})$ ;
3. Compute the **state extension**  $\mathbb{J}_{\mathcal{A}, \rho}[\cdot]$  and its CPTP factors  $\mathcal{R}$  and  $\mathcal{J}$ ;
4. Use the two factors to **reduce the dynamics**:

$$\check{\mathcal{A}} = \mathcal{R}\mathcal{A}\mathcal{J}, \quad \check{\mathcal{C}} = \mathcal{C}\mathcal{J}, \quad \Phi = \mathcal{R}.$$

This approach works. **Is it the best we can do?**

# A simple example: Invariant State

Consider CPTP dynamics  $\mathcal{A}$  and  $\mathcal{S} = \{\rho_0 = \sum_j p_j \Pi_j\}$  such that

$$\rho_0 = \mathcal{A}[\rho_0].$$

Then

$$\text{alg}\{\mathcal{A}^t[\rho_0], t = 0, \dots, n^2 - 1\} = \text{alg}(\rho_0) = \text{span}\{\Pi_j, \forall j\}.$$

## **Commutative (classical) algebra!**

However, a trivial (1-dim) dynamics  $\check{\rho}(t) = 1$  is sufficient to reproduce the states at each  $t$ : choose the CPTP reduction  $\check{\rho} = \mathcal{R}(\rho) = \text{Tr}(\rho)$ , and injection  $\mathcal{J}(\check{\rho}) = \rho_0 \check{\rho}$ , so that  $\rho_0 = \mathcal{J}(1)$  and

$$\check{\rho}(t+1) = \check{\rho}(t).$$



# Improving on algebras

Notice that  $\mathcal{R} = \text{span}\{\rho_0\}$  is a distorted algebra, closed under the product  $X \cdot_{\rho_0} Y = X\rho_0^{-1}Y$ .

**We need to close  $\mathcal{R}$  to a distorted algebra.**

In general, if  $\mathcal{R}$  is the reachable space we need to find a state  $\sigma$  such that  $\text{alg}_{\sigma}(\mathcal{R})$ :

- is the smallest distorted algebra that contains  $\mathcal{R}$ ;
- there exists a state extension  $\mathbb{J}$  that projects onto  $\text{alg}_{\sigma}(\mathcal{R})$  (Takesaki's theorem).

# How to “properly” close $\mathcal{R}$ to a distorted algebra

Given an operator space  $\mathcal{R}$ :

1. Compute  $\mathcal{A} = \text{alg}(\mathcal{R})$ ;
2. Compute its center  $\mathcal{Z} = \{X \in \mathcal{A} \mid [X, A] = 0, \quad \forall A \in \mathcal{A}\}$ ;
3. Pick a full-rank element  $\sigma \in \mathcal{Z}$ ;

**Main result:**  $\text{alg}_\sigma(\mathcal{R})$  is the **smallest distorted algebra** that allows for a **state extension** and contains  $\mathcal{R}$ .

# Reachable reduction

1. Find the **reachable subspace**

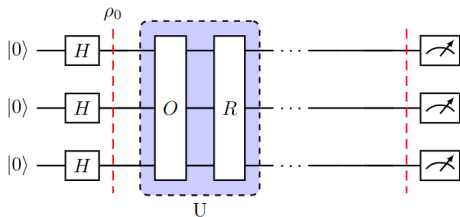
$$\mathcal{R} = \text{span}\{\mathcal{A}^t(\rho_0), t \geq 0\};$$

2. Close it to a **distorted algebra**  $\mathcal{A}_\sigma = \text{alg}_\sigma(\mathcal{R})$ ;
3. Compute the **state extension**  $\mathbb{J}_{\mathcal{A},\sigma}$  and its factors  $\mathcal{R}$  and  $\mathcal{J}$ ;
4. Use the two factors to **reduce the dynamics**:

$$\check{\mathcal{A}} = \mathcal{R}\mathcal{A}\mathcal{J}, \quad \check{\mathcal{C}} = \mathcal{C}\mathcal{J}, \quad \Phi = \mathcal{R}.$$

**Applications of interest:**

# Grover's algorithm



$$O|i\rangle = (-1)^{f(i)}|i\rangle$$

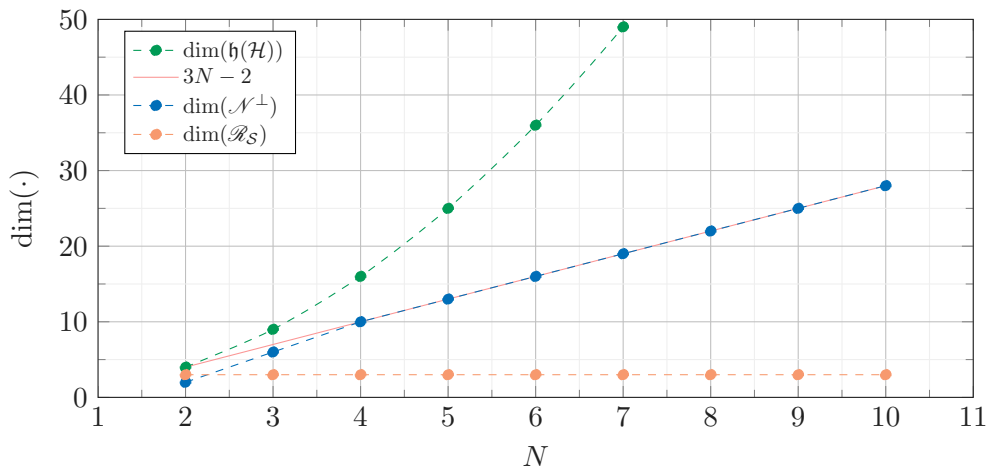
$$R = 2|+\rangle\langle+|^{\otimes N} - I_N$$

REPEAT  $U \approx \frac{\pi}{4}\sqrt{N}$  TIMES

$$\begin{cases} \rho(t+1) = U\rho(t)U^\dagger \\ \mathbf{p}(t) = \text{diag}[\rho(t)] \end{cases}$$

$$\rho_0 = |+\rangle\langle+|^{\otimes N}$$

# Numerical Analysis of Grover's algorithm



# Reachable reduction applied to Grover

Fix  $S = \{j | f(x) = 1\}$  the set of solutions to the search problem and  $M = |S|$  the number of solutions.

$$|\alpha\rangle := (N - M)^{-1/2} \sum_{j \notin S} |j\rangle \quad |\beta\rangle := (M)^{-1/2} \sum_{j \in S} |j\rangle$$

$$\mathcal{R} := \text{span}\{\rho(t), t \geq 0\} = \text{span}\{|\alpha\rangle\langle\alpha|, |\beta\rangle\langle\beta|, |\alpha\rangle\langle\beta| + |\beta\rangle\langle\alpha|\}$$

$$\mathcal{A} = \text{alg}(\mathcal{R}) = \text{span}\{|\alpha\rangle\langle\alpha|, |\beta\rangle\langle\beta|, |\alpha\rangle\langle\beta|, |\beta\rangle\langle\alpha|\} \simeq \mathbb{C}^{2 \times 2}$$

It's a qubit!

Regardless of the system's size  $N$



# Grover's reduced quantum model

$$\begin{cases} \check{\rho}(t+1) = \check{U} \check{\rho}(t) \check{U}^\dagger \\ \mathbf{p}(t) = \sum_{i \in S} |i\rangle\langle 1| \check{\rho}(t) |1\rangle\langle i| + \sum_{i \notin S} |i\rangle\langle 0| \check{\rho}(t) |0\rangle\langle i| \end{cases}$$

with

$$\check{\rho}_0 = \frac{N-M}{N} |0\rangle\langle 0| + \frac{M}{N} |1\rangle\langle 1| + \frac{\sqrt{(N-M)M}}{N} (|0\rangle\langle 1| + |1\rangle\langle 0|)$$

$$\check{U} = \frac{N-2M}{N} I_2 - i \frac{\sqrt{(N-M)M}}{N} \sigma_y$$

**NOTE:** Knowing  $\mathcal{R}$  “is equivalent” to finding the search problem solution.



# Continuous-time

Consider

$$\dot{\rho}(t) = \mathcal{L}[\rho(t)].$$

For  $e^{\mathcal{L}t}[\cdot]$  to be CPTP for all  $t \geq 0$ ,  $\mathcal{L}[\cdot]$  needs to be a GKSL generator:

$$\mathcal{L}(\rho) = -i[H, \rho] + \sum_k L_k \rho L_k^\dagger - \frac{1}{2} \{L_k^\dagger L_k, \rho\}.$$

## Main result:

Given a GKSL generator  $\mathcal{L}$ , its action onto the distorted algebra  $\mathcal{X}$ ,  $\mathcal{L}|_{\mathcal{X}} = \mathcal{R}\mathcal{L}\mathcal{J}$  is still a GKSL generator.

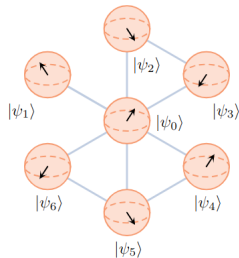
# Dephasing central spin model coupled to a Markovian environment

$$H = \sum_{k=0}^{N-1} J_{k,0} \sigma_z^{(k)} + \sum_{l>k} J_{k,l} \sigma_z^{(k)} \sigma_z^{(l)}$$

$$\{L_k\} = \{\sigma_x^{(j)}, \sigma_y^{(j)}, \sigma_z^{(j)}, \sigma_+^{(j)}, \sigma_-^{(j)}\}_{j=0,\dots,N-1}$$

$$\begin{cases} \dot{\rho} = -i[H, \rho] + \sum_k L_k \rho L_k^\dagger - \frac{1}{2} \{L_k^\dagger L_k, \rho\} \\ \rho_S(t) = \text{tr}_E[\rho(t)] \end{cases}$$

$$\rho_0 \in \mathcal{D}(\mathcal{H})$$



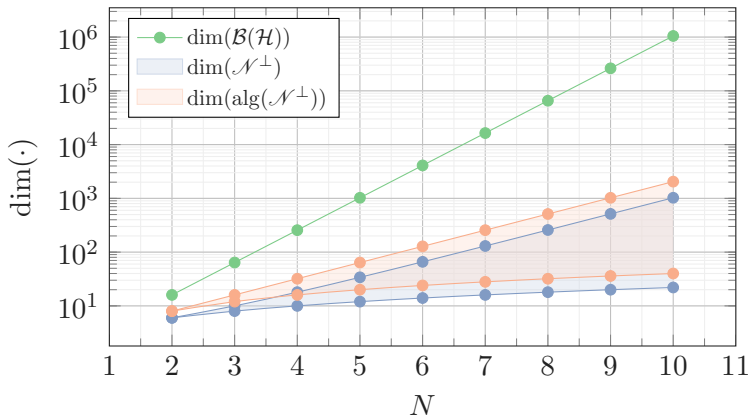
# Observable space and algebra

$$\mathcal{N}^\perp \subseteq \text{span}\{\sigma_0 \otimes I_{2^{N-1}}, \sigma_z \otimes I_{2^{N-1}}, \sigma_x \otimes |j\rangle\langle j|, \sigma_y \otimes |j\rangle\langle j|, \quad j = 0, \dots, 2^{N-1}\}$$

$$\text{alg}(\mathcal{N})^\perp \subseteq \text{span}\{\sigma_k \otimes |j\rangle\langle j|, \quad j = 0, \dots, 2^{N-1}, \quad k = 0, x, y, z\} \simeq \bigoplus_{j=0}^{2^{N-1}-1} \mathbb{C}^{2 \times 2}$$

The bath can be modeled classically!

# How much are we reducing?



$4N \leq \dim(\mathcal{A}) \leq 2^{N+1}$  (depending on the parameters) instead of  $2^{2N}$ .

# Conclusion

- **Framework for model reduction of statistical dynamics (CPTP)**
- Motivated by quantum walks/HMM;
- The Framework is very general and has been extended to:
  - Discrete-time case;
  - Continuous-time case;
  - Open systems, homogeneous;
  - Quantum trajectories.
- Solves an open problem in HMM (since '92);
- Explores quantumness of processes;
- Explores optimality of Grover's algorithm and other quantum walks;

# Outlook

- Approximate method;
- Asymptotic model reduction (ongoing w/ Viola);
- Physically relevant examples (ongoing w/ Viola);
- Positive linear systems (ongoing w/ Cortese and Ferrante);
- Infinite-dimensional case (e.g. bosons);
- Locality constraints, networks (ongoing w/ Peruzzo);
- Continuous-time with measurement, SME (ongoing w/ Pellegrini);
- Adiabatic elimination techniques (ongoing w/ Sarlette);

## Related publications

- TG and FT, "Minimal resources for exact simulation of quantum walks," IEEE CDC 2022;
- TG and FT, "Algebraic Reduction of Hidden Markov Models," IEEE TAC 2023;
- TG and FT, "Model Reduction for Quantum Systems: Discrete-time Quantum Walks and Open Markov Dynamics," (submitted to) IEEE TIT 2023;
- TG and FT, "Exact model reduction for conditional quantum dynamics," (submitted to) IEEE CDC and L-CSS 2024;
- TG, Yukuan Tao, FT, and Lorenza Viola, "Exact Model Reduction for Continuous-Time Open Quantum Dynamics," (to be submitted to) PRXQ 2024;

# Thanks for your attention!



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