

## 1 Syntax of lambda calculus

Assume a countable set of variable names, denoted by  $a, b, c, \dots, x, y, z, a_0, a_1, \dots$

**Definition.** A  $\lambda$ -term is defined by the following context-free grammar:

$$\begin{aligned} \langle term \rangle &:= \langle name \rangle \\ &| (\lambda \langle name \rangle . \langle term \rangle) \\ &| (\langle term \rangle \langle term \rangle) \end{aligned}$$

**Conventions.**

1. **Function application** is left-associative, so  $((A_1 A_2) A_3) \dots A_k$  can be abbreviated as  $A_1 A_2 A_3 \dots A_k$
2. Nested **abstractions**  $(\lambda x_1. (\lambda x_2. (\dots \lambda x_k. A) \dots))$  can be abbreviated as  $\lambda x_1 x_2 \dots x_k. A$

**Example.**

$\lambda xy. FAB$  means  $((\lambda x. (\lambda y. F)) A) B$

## 2 Free variables

1.  $\langle name \rangle$  is free in  $\langle name \rangle$
2.  $\langle name \rangle$  is free in  $\lambda \langle name' \rangle . \langle term \rangle$  if  $\langle name \rangle \neq \langle name' \rangle$  and  $\langle name \rangle$  is free in  $\langle term \rangle$
3.  $\langle name \rangle$  is free in  $\langle term' \rangle \langle term'' \rangle$  if  $\langle name \rangle$  is free in  $\langle term' \rangle$  or  $\langle name \rangle$  is free in  $\langle term'' \rangle$

## 3 Bound variables