



C & x86 Review



CS5435 - Spring 2024
Marina Bohuk



The C Language

- Barebones language → Very fast, versatile
- Does not come with any “extras” like Java and Python
 - No classes or objects
- Does not do boundary or index checks
- Does not abstract anything away

Pointers

- Store a memory address of an object
 - `char *c;`
 - `int *x;`
 - `Object *o;`

variable	address	memory
	1000	5
x	1016	1000

Arrays

- Arrays are pointers!
 - `int intArray[3] = {2, 4, 6};`
- The length allocates space, but there is no index checking
- Why do the addresses jump by 4?
- Strings are just arrays

variable	address	memory
intArray	???	1000
	...	
&intArray[0]	1000	2
&intArray[1]	1004	4
&intArray[2]	1008	6
	1012	???

Example program

```
#include <stdio.h>
#include <string.h>
```

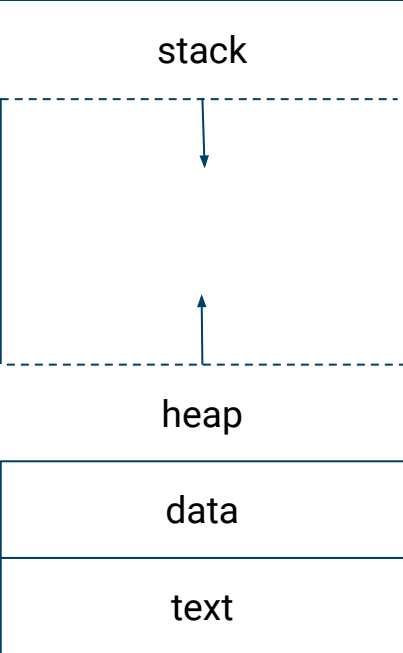
```
void greeting( char* temp1 )
{
    char name[400];
    memset(name, 0, 400);
    strcpy(name, temp1);
    printf( "Hi %s \n", name );
}
```

```
int main(int argc, char* argv[] )
{
    greeting( argv[1] );
    printf( "Bye %s\n", argv[1] );
}
```

Compiling

```
gcc -ggdb -mpreferred-stack-boundary=2 -zexecstack -fno-stack-protector -no-pie  
-fno-pie -m32 meet.c -o meet
```

- ggdb: debugger symbols
- mpreferrred-stack-boundary=2: align the stack in 4 byte chunks
- zexecstack: use an executable stack
- fno-stack-protector: disable stack protections
- no-pie/fno-pie: no ASLR
- m32: use 32-bit x86, not 64



Memory allocation

Static

- Amount of memory needed already known at compile time
 - `int array[10];`
- This memory is allocated on the **stack**.

Dynamic

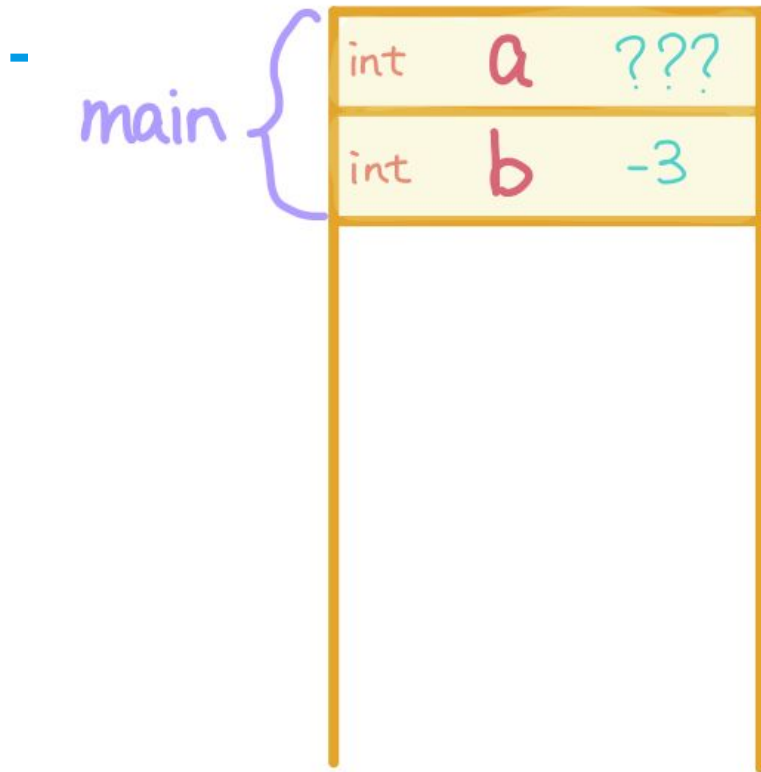
- Amount of memory needed not known
 - `int *array = (int*)malloc(n * sizeof(int));`
- This memory is allocated on the **heap**.

Stack

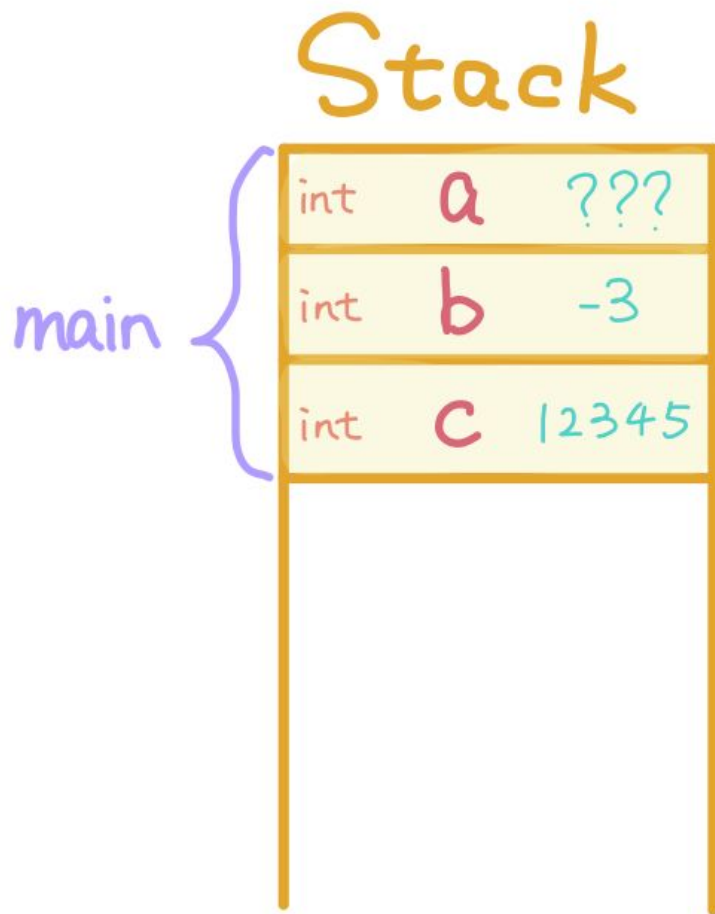


```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```

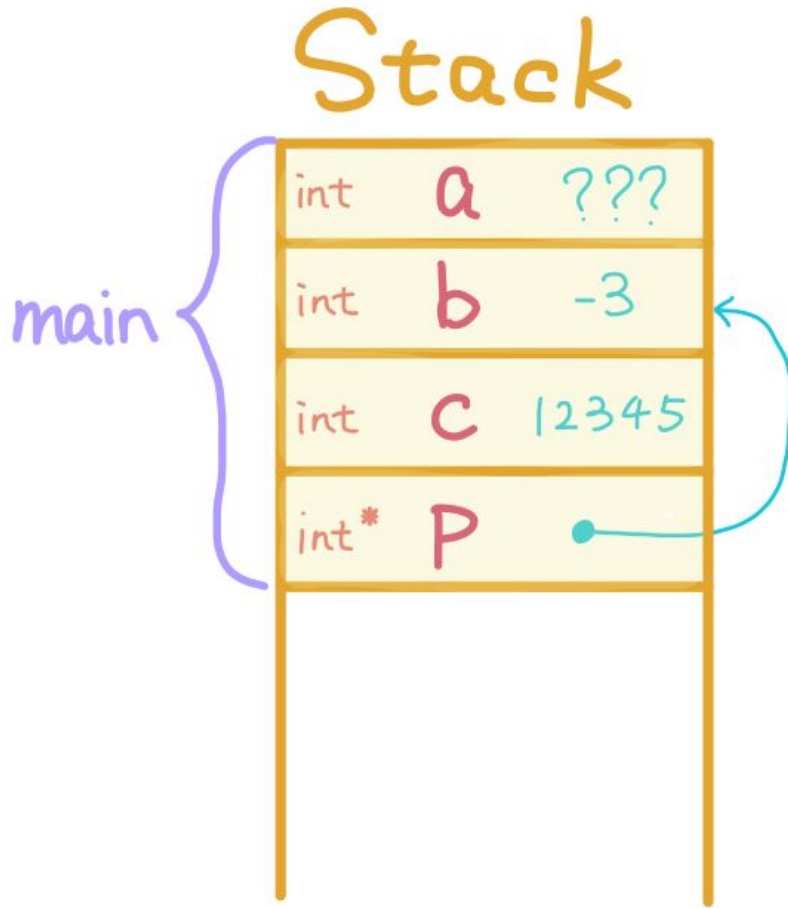
Stack



```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```

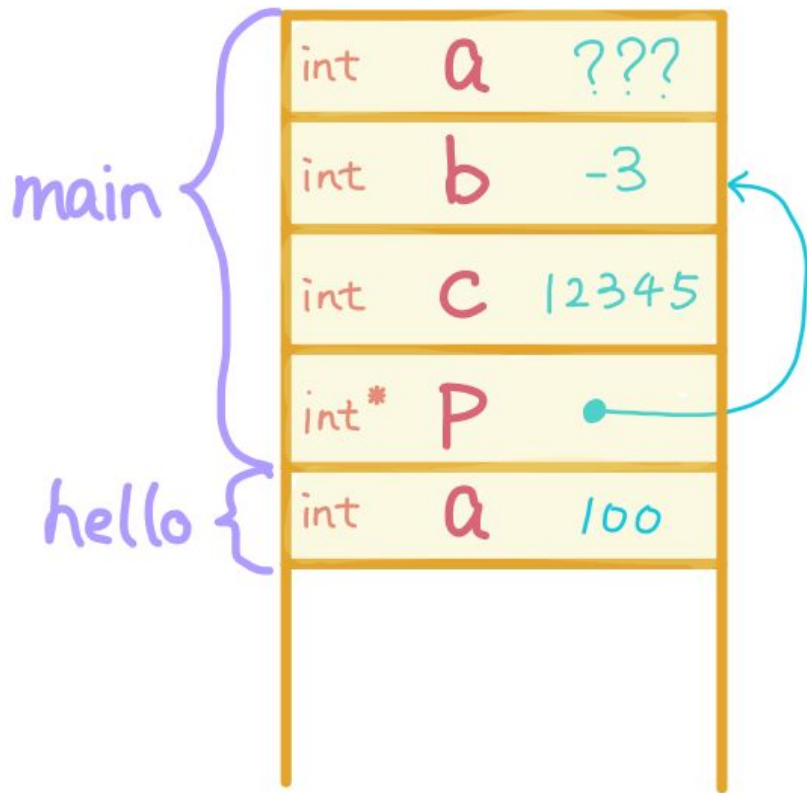


```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```

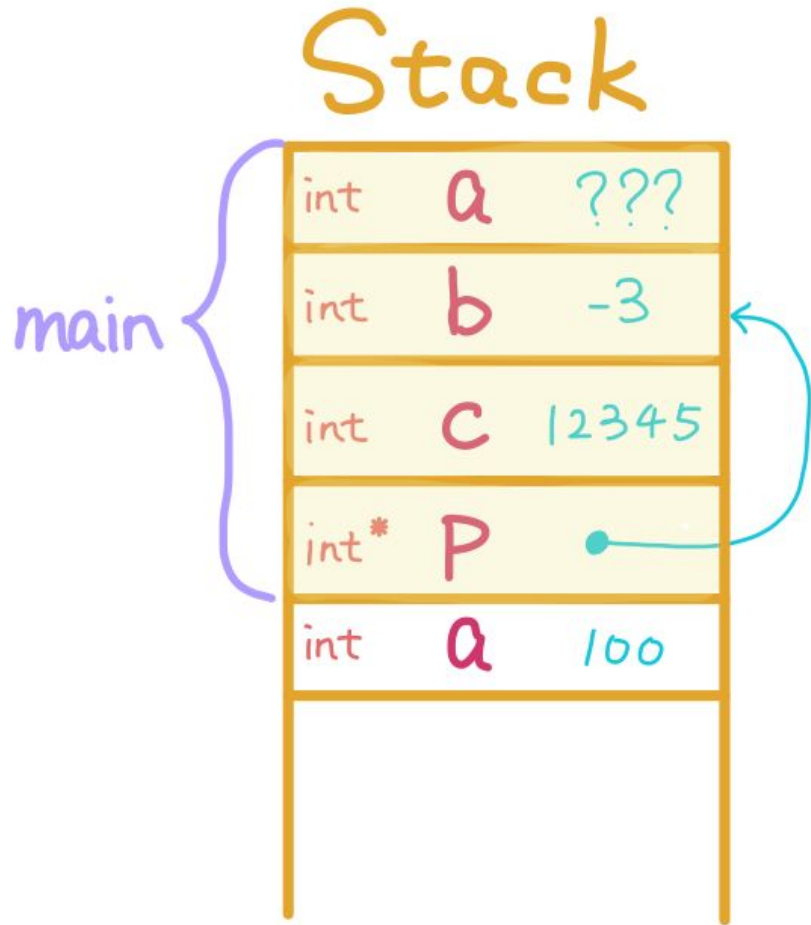


```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```

Stack

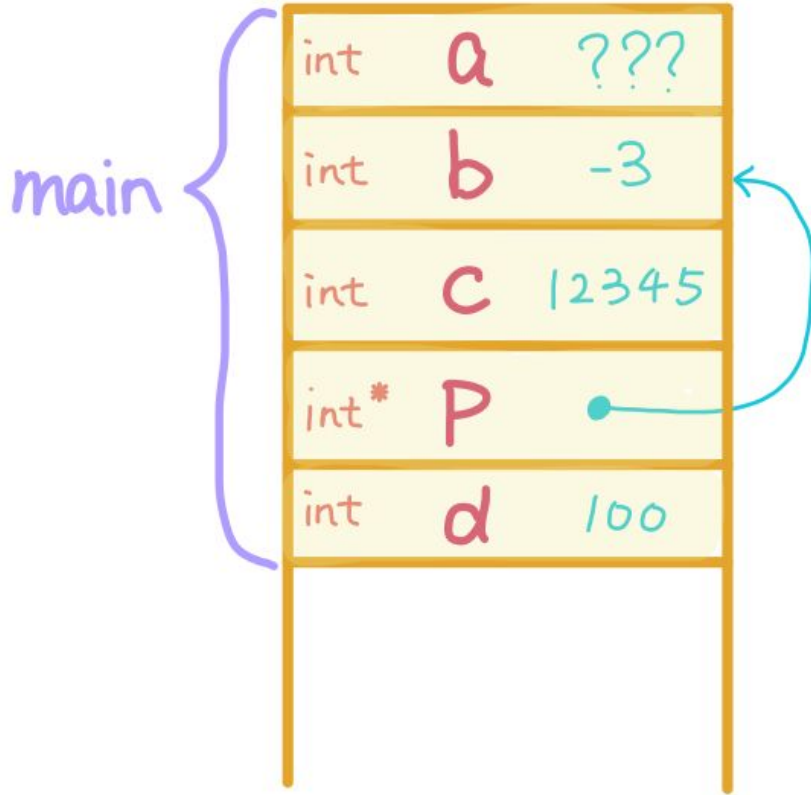


```
int hello() {  
    int a = 100;  
    return a;  
}  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```



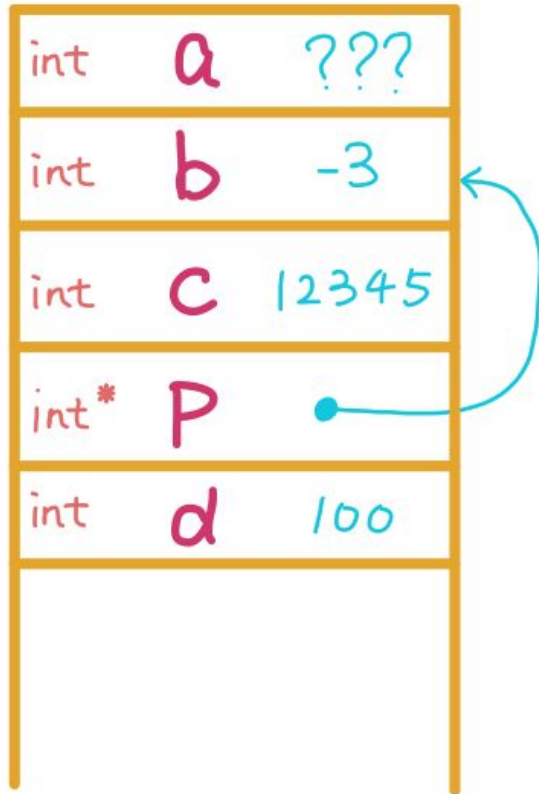
```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```

Stack



```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```

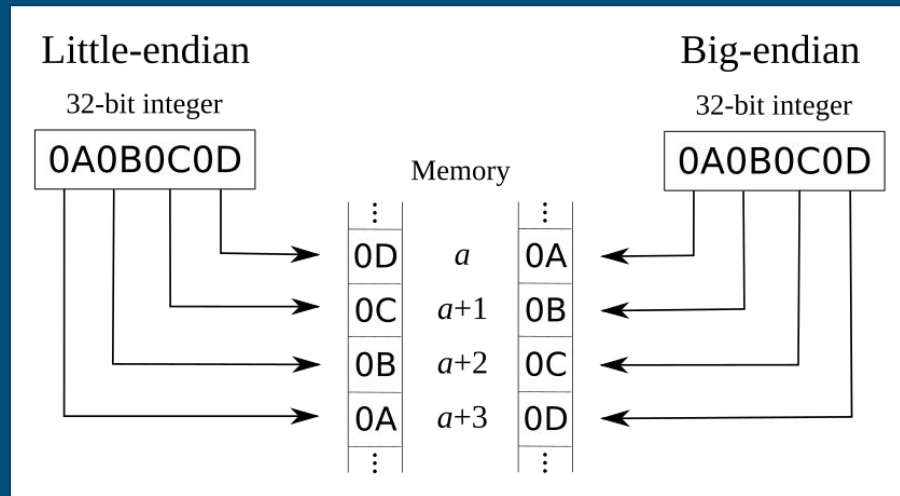
Stack



```
int hello() {  
    int a = 100;  
    return a;  
}  
  
int main() {  
    int a;  
    int b = -3;  
    int c = 12345;  
    int *p = &b;  
    int d = hello();  
    return 0;  
}
```


x86 assembly

- x86 first came around in the 70s and is still used without much change today
- C is compiled into x86 assembly
- Little endian (important when writing exploits!)
 - Endianness: how to order bytes in a unit
- Uses registers as a quick way to temporarily store information
 - Variables built into the processor
 - eax, ebx, ecx..



Example x86 instructions

`mov ecx, eax` // move value in eax into ecx

`mov eax, [ebx]` // move 4 bytes at the memory address pointed to by ebx into eax

`add eax, 0x4` // also sub, inc/dec

`push eax` // push the value in eax onto the stack

`pop eax` // pop the value on top of the stack into eax

`call function`

`jmp 0xffffffffb3` // jump to an address

x86 Calling Convention

A set of rules between functions:

- How are parameters passed? → in registers
- Where are registers preserved? → on the stack
- Where should local vars be stored? → on the stack
- How are values returned? → in register eax

Caller summary

- Prologue: Tasks to take care of BEFORE calling the subroutine
 - Save the caller-save registers
 - Place parameters on the stack (or in registers for 64-bit)
- Call the subroutine
 - The `call` instruction places the return address on the stack
- Epilogue: Tasks AFTER calling the subroutine
 - Remove params from the stack
 - Access return value in `eax`
 - Restore saved registers

Callee summary

- Prologue: Tasks to perform BEFORE executing the function body
 - Allocate space for local variables (on the stack): `sub 8, %esp`
 - Save callee-save registers
- Function body
- Epilogue: Tasks to perform AFTER executing the function body but BEFORE leaving the function
 - Save return value to `eax`
 - Restore callee-save registers (pop them from the stack)
 - Deallocate local variables
 - Return

GDB

- A debugger that allows you to see what's going on inside your C program
- PEDA: <https://github.com/longld/peda>
- Breakpoints:
 - You can set a breakpoint to tell GDB to stop when it reaches a certain location
 - `b hello` // function name
 - `b 123` // line number
- Disassemble:
 - Show the assembly code for a given function
 - `disas main` // function name
- Print out the current stack frame
 - `info frame`
- Print out memory
 - `x/64x 0xbffff5cc` // print out 64 bytes beginning at address 0xbffff5cc
- Step through code one line at a time
 - `step`

VM for HW 4

ssh your-net-id@35.226.118.201

password: your net id (change it as soon as you login)

Simple C program

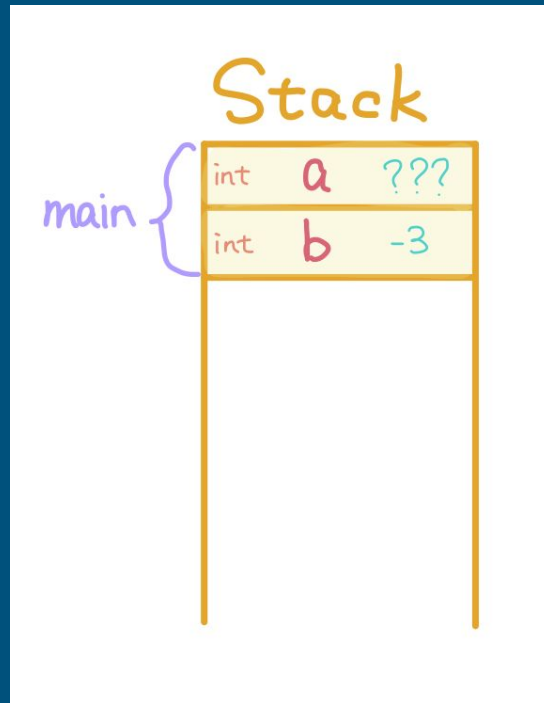
```
#include <stdio.h>
#include <string.h>
```

```
void greeting( char* temp1 )
{
    char name[400];
    memset(name, 0, 400);
    strcpy(name, temp1);
    printf( "Hi %s \n", name );
}
```

```
int main(int argc, char* argv[] )
{
    greeting( argv[1] );
    printf( "Bye %s\n", argv[1] );
}
```


Simple Stack Smashing Example

- Preview of stack smashing before Thursday
- Capture the Flag (CTF) Challenge
- Goal: get the flag! (a_flag_looks_like_this)
- Files:
 - Source: <https://tinyurl.com/stacksmash5435>
 - Binary: <https://tinyurl.com/stacksmash5435bin>
- Connect to the challenge:
 - `nc kubenode.mctf.io 30005`



Simple Stack Smashing Example

How Cool Are You? (solved by 49 teams)

300

Are you cool enough to get in? Connect via `nc kubenode.mctf.io 30005` and find out.

You can grab the [binary here](#), and the [source here](#). Good luck.

Submit!

```

#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <time.h>

int main() {
    // Declare variables
    int your_estimated_coolness;
    char your_name[64];

    // Disable buffering for stdin and stdout
    setbuf(stdin, NULL);
    setbuf(stdout, NULL);

    // Seed the random number generator with the current time
    srand(time(NULL));

    // Talk to user
    printf("Whoa there, my dude. Only the truly cool are allowed in here.\n");

    // Use the standard coolness formula to determine how cool the user is
    your_estimated_coolness = 500 + 10 * (rand() % 1000);

    // Read in the user's name
    printf("What's your name? ");
    gets(your_name);

    // Evaluate them
    printf("%s, it looks like your coolness value is %d.\n", your_name, your_estimated_coolness)
    if (your_estimated_coolness < 1500000001) {
        printf("I'm sorry, you're just not cool enough. Get lost!\n");
    } else if (your_estimated_coolness > 1500000001) {
        printf("Whoa, you're... you're actually a bit too cool for us. Sorry...\n");
    } else if (your_estimated_coolness == 1500000001) {
        printf("Wow! Radical... You're in. Have fun.\n");
        printf("Oh, right, here's the flag: %s\n", getenv("FLAG"));
    }

    return 0;
}

```

Source:

<https://tinyurl.com/stacksmash5435>

Binary:

<https://tinyurl.com/stacksmash5435bin>

What does this program do? Where is the vulnerability?

1500000001 == 0x59682F01

More security practice

- MetaCTF:
 - Next monthly Flash CTF: <https://mctf.io/apr2024>
 - For an int overflow challenge, check out February's Flash CTF
- PicoCTF