

A tutorial on high availability clusters

Fundamentals



**Public businesses
Unacceptable outage
Uninterrupted services**



Fundamentals



Negative impact in companies' profit and business
Unacceptable in E-commerce and E-business companies
Colateral damages in global companies and economies

Business Continuity and Disaster Recovery

Fundamentals



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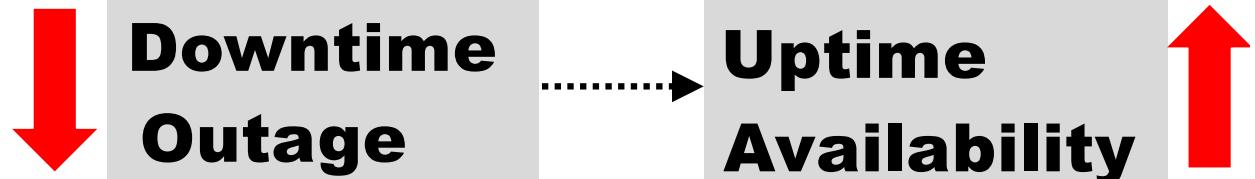
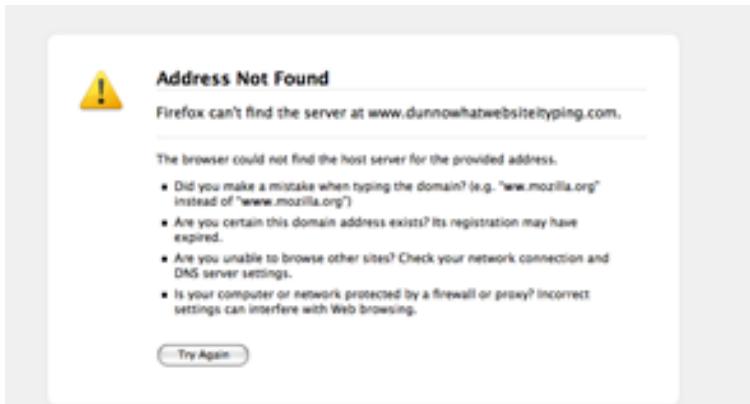
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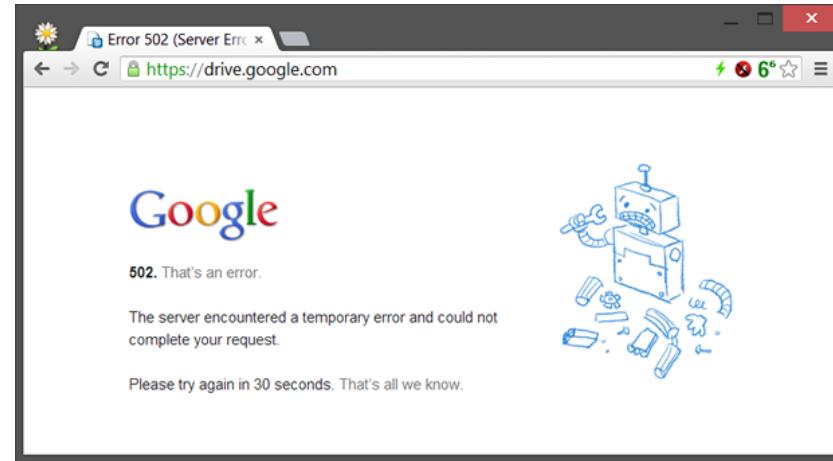
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Fundamentals



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http://www.google.com/appsstatus

Fundamentals – a definition for availability

- Measure the percentage of time (hour, day, month) in which the service is available for the end-user.
- Should be globally analysed and not by the components separately.
- Should be measured in a user perspective.
- Its value has positive (or negative) impact in the business.

Fundamentals – a definition for availability



Fundamentals – a definition for availability

$$D = \frac{\text{Uptime}}{\text{Operation}} \quad \text{with } D \in [0, 1]$$

Uptime = duration of uptime period (h, min, week, day, month,...)

Operation = duration of total period of operation (h, min, week, ...)

Example: Availability in a week (168h) for a uptime of 165h.

$$D = \frac{165}{168} = 0.9821 = 98.21\%$$

Fundamentals – a definition for availability

	Per Hour	Per Day	Per Week	Per Year
99.999%	.0006	.01	.10	5
99.98%	.012	.29	2	105
99.95%	.03	.72	5	263
99.90%	.06	1.44	10	526
99.70%	.18	4.32	30	1577

Values for *downtime* (minutes)

- 24x60x365 operation:**
- 525600 minutes per year
- 24x7x60 operation:**
- 10080 minutes per week

An availability of **99.999%** means 5 minutes of *downtime* in a year!

At least three decimal digits should be used ("The myth of nines")

A common metric widely used to assess infrastructure effectiveness

Fundamentals – mean availability

- **Measures for availability (mean time)**

Mean Time Between Failures (MTBF) = $\frac{\text{Operation}}{\text{Total failures}}$
Mean duration of component (hardware) downtime.

Mean Time to Repair (MTTR) = $\frac{\sum \text{Repair}}{\text{Total failures}}$
Mean time to repair a component.

When applied to
network services

Mean Time Between Service Outage (MTBSO)
Mean Time To Service repair (MTTSR)

Fundamentals – mean availability

$$\text{Mean availability} = \frac{\text{MTBF}}{\text{MTBF} + \text{MTTR}}$$

MTBF=1000h



MTTR=1h

$$D_{\text{mean}} = \frac{1000}{1001} = 99,90\%$$

$$\text{Mean availability} = \frac{\text{MTBSO}}{\text{MTBSO} + \text{MTTSR}}$$

MTBSO=1000 h



MTTSR=3h

$$D_{\text{mean}} = \frac{1000}{1003} = 99,70\%$$

$$\text{Annualized Failure Rate (AFR)} = \frac{365 \times 24}{\text{MTBF}} = \frac{8760}{\text{MTBF}}$$

Example: Standard S.M.A.R.T for disks

S.M.A.R.T. = Self-Monitoring, Analysis and Reporting Tech.

- Available for HDD and SSD devices
- It detects and reports several disks' fiability indicators
- Goal: to antecipate hardware failures

In Linux:

```
smartctl -i /dev/sda
```

```
smartctl -a /dev/sda
```

Fundamentals – common downtime causes



Planned

Proactive maintenance

Hardware replacement

Hardware upgrade

Software upgrade

Users are usually warned.



Unplanned

Reactive response

Hardware, Natural, Power and network failures

Web and DB server down

DNS server down

Vulnerabilities exploits

No time to warn users.
Backup strategies should start

Fundamentals – downtime costs

Considering:

- 24x7 operation
- $D_{mean} = 98.1\%$
- Cost of downtime per hour = 10,000€



Then:

- Downtime = 3.192 hours per week
- Total cost of downtime $\approx \underline{31,920\text{€}}$ per week

If $D_{mean} = 99.9\%$ then downtime ≈ 10 minutes (0.166h) = **1660€**

Gain = 31.920€ - 1660€ = **30.260€ per week!**

Fundamentals – downtime costs



Direct

Business activities

Productivity

Intranet and internal processes



Indirect

Clients' dissatisfaction

Stock options

Negative publicity

Legal processes

Company reliability

External reputation

Worker's performance impact



Fundamentals – downtime costs

- How does it cost to reduce downtime from 3.2h to 10 minutes?
- What if we reduce from 3.2h to 1 h?
- The benefits worth the investment?

What is the ROI on applying the measures?

- Reduce the risk (**R**) before (**b**) and after (**a**) the investment
- Assess potential gains (**G**) after the investment (**I**)

$$G = R_b - R_a$$

$$ROI = \frac{G - I}{I} , \text{ being } I \text{ the cost of preventive measures}$$

Fundamentals – calculating the risk

- Probability (P) of an event
- Duration (D) of the downtime event
- Impact (I), means the percentage of affected users

The effect of an event (E_x) is measured by its value before (b) and after (a) of applying the preventive measures.

$$E_{ax} = P_{ax} \times D_{ax} \times I_{ax} \quad \text{Downtime cost}$$

$$E_{dx} = L_{dx} \times D_{dx} \times I_{dx} \quad \text{Implementation cost}$$

$$R_{\text{before}} = C_D \times (E_{b1} + E_{b2} + \dots + E_{bx})$$

$$R_{\text{after}} = I + (C_D \times (E_{a1} + E_{a2} + \dots + E_{ax}))$$

Fundamentals – calculating the risk

$$G = R_a - R_d$$

$$R_a > R_d$$

$$G = C_D \times (E_{a1} + E_{a2} + \dots + E_{ax}) - I + (C_D \times (E_{d1} + E_{d2} + \dots + E_{dx}))$$

$$G > 0$$

$$ROI = \frac{G - I}{I} \%$$

(annual basis or per cycle)

Examples:

Marcus E, Stern H., “*Blueprints for high availability*”; 2003; Wiley; ISBN: 0471430269;
pp. 42-46

Fundamentals – calculating the risk

Calculating the effect before applying the measures:

Table 3.2 Effects of Outages before Clustering Software Is Installed

OUTAGE TYPE	BEFORE DURATION (D)	BEFORE LIKELIHOOD (L)	BEFORE IMPACT (I)	BEFORE EFFECT (D × L × I)
Crash and reboot	60 minutes*	10 [†]	100%	$E_{B1} = 600$ (during the day)
Crash and reboot (off-hours)	120 minutes [‡]	10 [†]	75% [§]	$E_{B2} = 900$
Scheduled reboot	30 minutes	60	50% [§]	$E_{B3} = 900$
Motherboard or other major hardware failure	24 hours (1,440 minutes)*	2	100%	$E_{B4} = 2,880$
Network card failure	4 hours (240 minutes)**	2	100%	$E_{B5} = 480$
Application failure	60 minutes	20 ^{††}	100%	$E_{B6} = 1,200$
Scheduled maintenance	4 hours (240 minutes) [#]	20	50%	$E_{B7} = 2,400$
Failover testing	0 ^{§§}	0	0	$E_{B8} = 0$

Total effect of outages:

9,360 minutes (99.644 percent availability over 5 years)

Marcus E. Stern H., “Blueprints for high availability”; 2003; Wiley; ISBN: 0471430269;

Fundamentals – calculating the risk

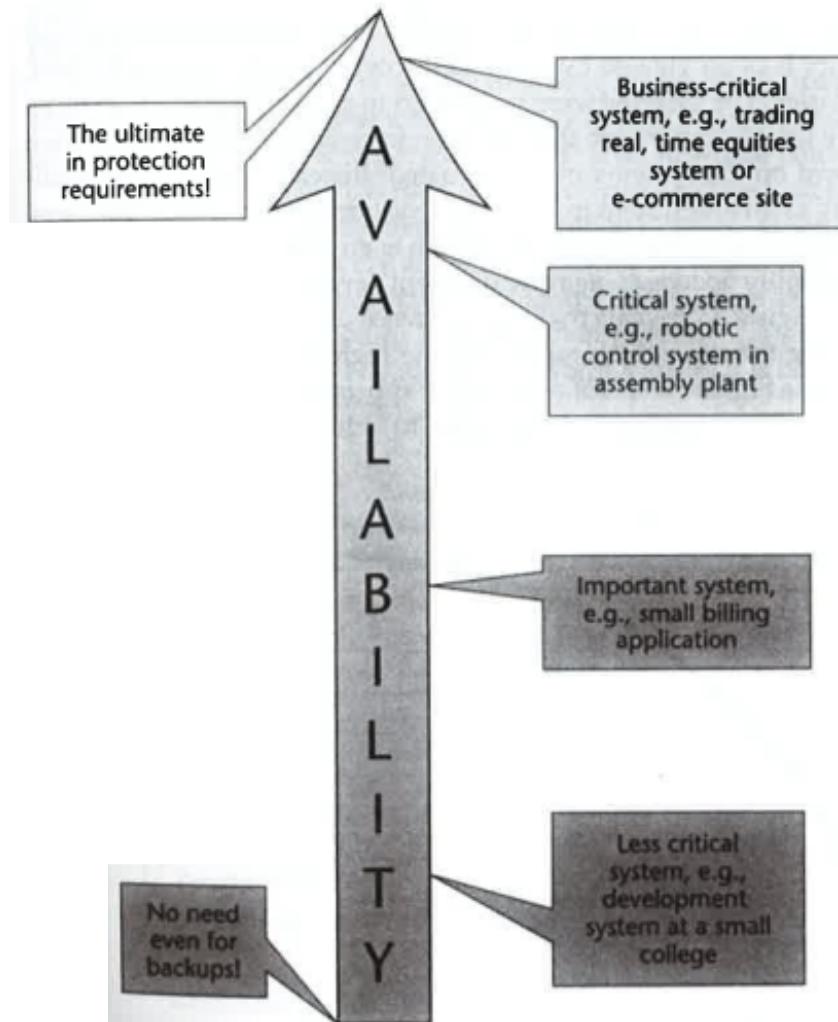
Calculating the effect after applying the measures:

Table 3.3 Effects of Outages after Clustering Software Is Installed

OUTAGE TYPE	AFTER DURATION (D)	AFTER LIKELIHOOD (L)	AFTER IMPACT (I)	AFTER EFFECT (D × L × I)
Crash and reboot	5 minutes [†]	10	100%	$E_{A1} = 50$ (during the day)
Crash and reboot (off-hours)	5 minutes [†]	10	75%	$E_{A2} = 37.5$
Scheduled reboot	5 minutes	60	50%	$E_{A3} = 150$
Motherboard or other major hardware failure	5 minutes	2	100%	$E_{A4} = 10$
Network card failure	2 minutes, then 5 minutes [‡]	2	100% then 50%	$E_{A5} = 9$
Application failure	3 minutes [§]	20	100%	$E_{A6} = 60$
Scheduled maintenance	5 minutes	20	50%	$E_{A7} = 50$
Failover testing	5 minutes	20	50%	$E_{A8} = 50$
Total effect of outages:				416.5 minutes (99.984 percent availability over 5 years)

Marcus E, Stern H., "Blueprints for high availability"; 2003; Wiley; ISBN: 0471430269;

Fundamentals – calculating the risk



Hospitals
Aviation systems
Other highly critical systems

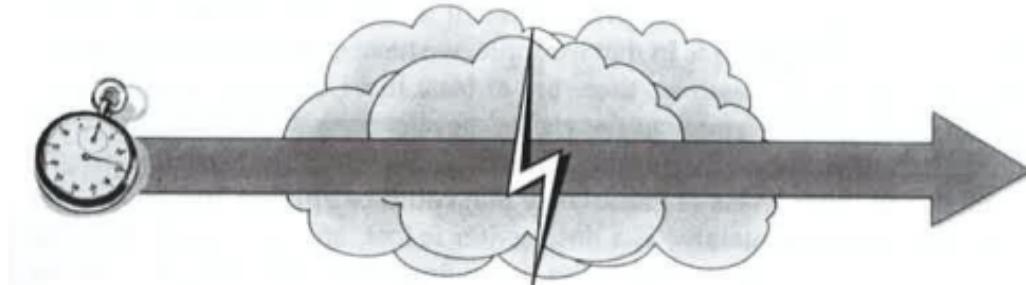
Web-based businesses
Industrial-based businesses

Universities and schools
Other less critical businesses

Small businesses
Inactive computers

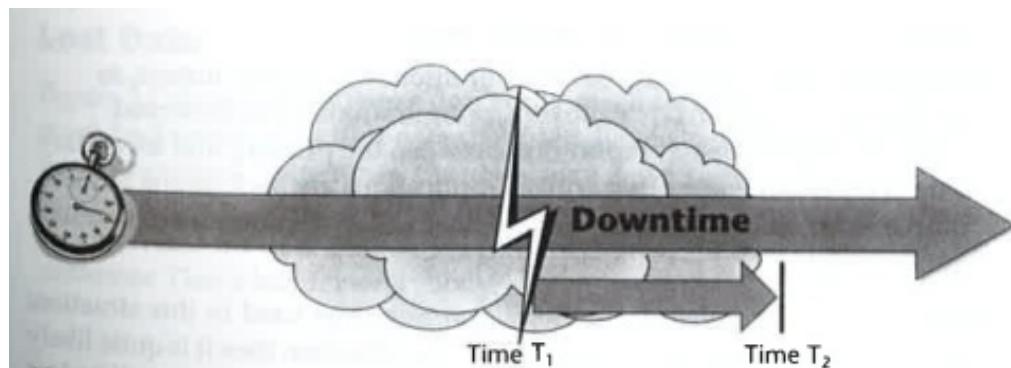
Marcus E. Stern H., "Blueprints for high availability"; 2003; Wiley; ISBN: 0471430269;

Outage lifecycle



Begin of an outage

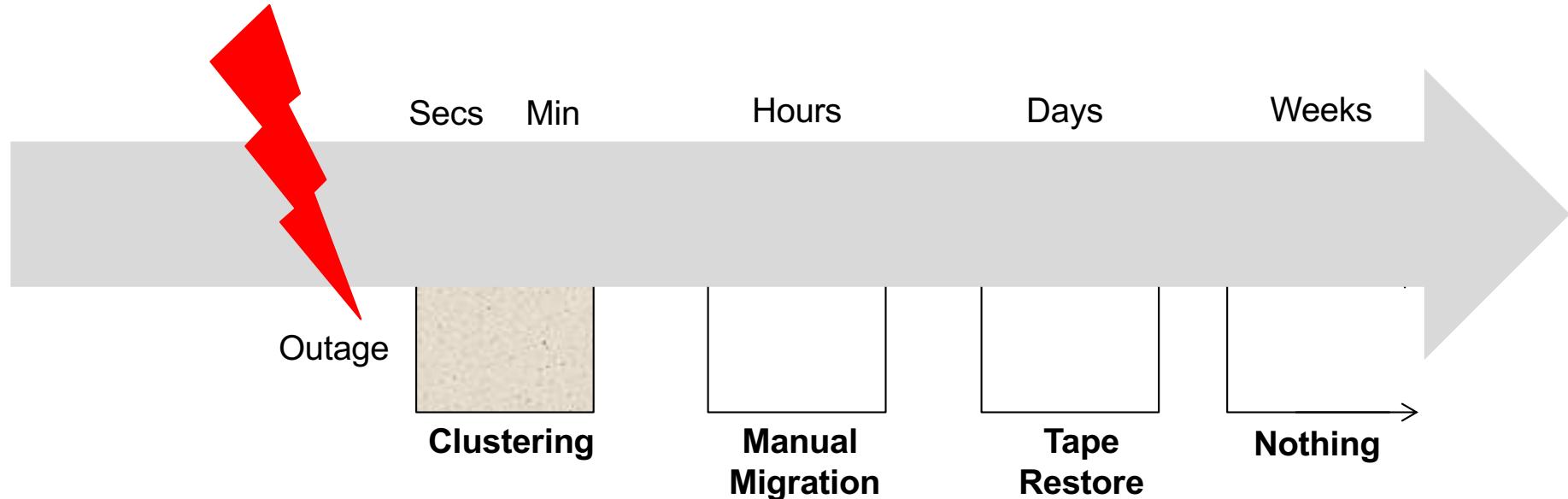
Marcus E, Stern H., "Blueprints for high availability"; 2003; Wiley; ISBN: 0471430269;



[T₁; T₂] - Downtime

Marcus E, Stern H., "Blueprints for high availability"; 2003; Wiley; ISBN: 0471430269;

Outage lifecycle – outage mitigation

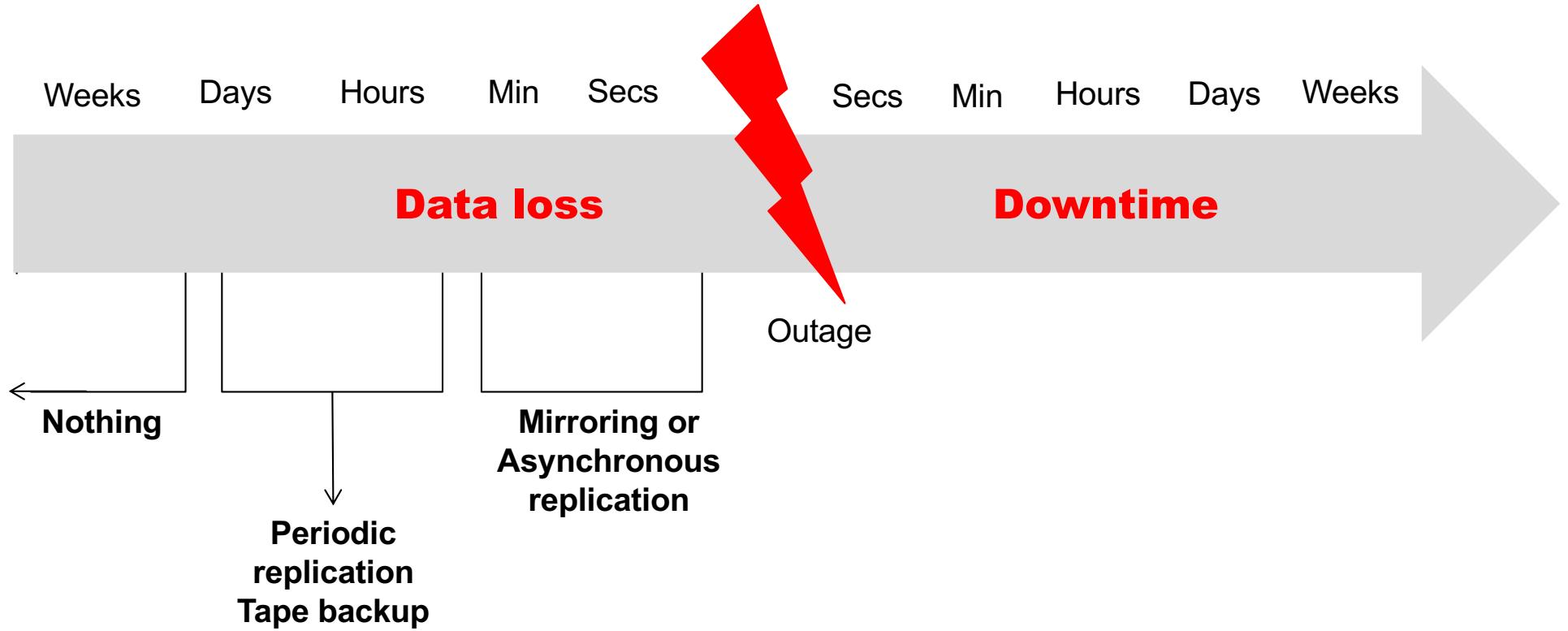


Time is an important variable

Downtime will happen in the presence of an outage

Goal: to mitigate outage in a minimum period of time

Outage lifecycle – data loss mitigation

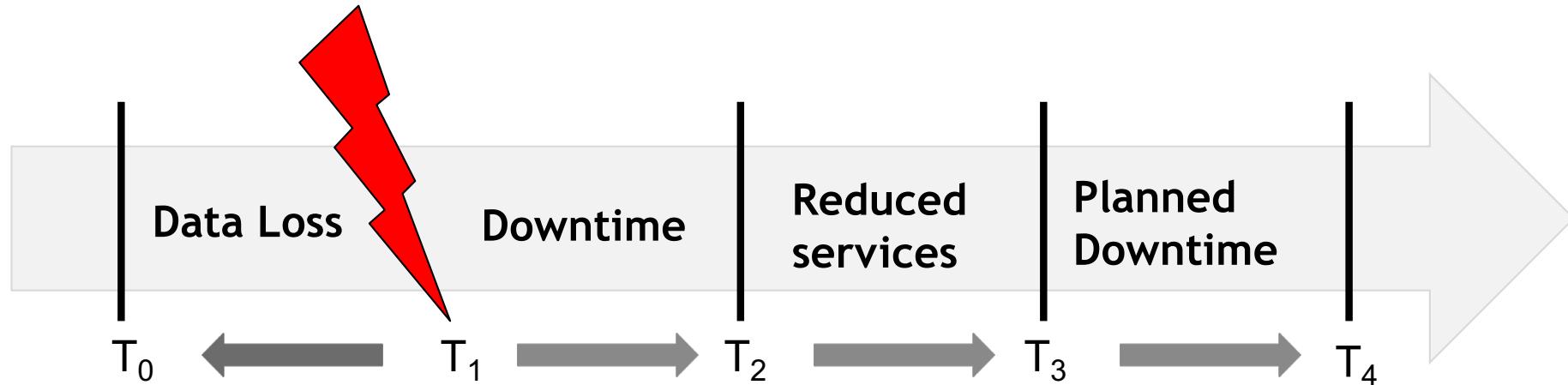


Time is again an important variable

Outage will cause data loss

Goal: to minimize data loss and to assure system consistency

Outage lifecycle



HA implementation is a business decision.

How much does it cost the infrastructure downtime?

How much does it cost to implement HA procedures to mitigate downtime?

Impact on the organization

- To define clear and tangible goals
- To define a physical environment
- To automate processes
- To define environment for testing and QA
- To create a stock of physical components
- To adjusts contracts for critical components
- To schedule processes
- To plan catastrophic scenarios
- To train
- To document everything!

Impact on the organization - vulnerabilities

VULNERABILITY	LIKELIHOOD (1-3)	IMPACT (1-3)	LEVEL OF CONCERN (LIKELIHOOD × IMPACT)	COMMENTS
Failed disk	3	1	3	Critical systems already have mirrored disks.
Blown CPU	2	2	6	Systems are not clustered; a failed CPU will result in major downtime for one server.
Database corruption	2	3	6	Corruption in a critical database could shut down the web site for hours.
Unreadable backups	1	2	2	Only an issue if we lose data from our disks. Then it could be quite serious.
Network component failure	2	2	4	We have a very complex network, with many single points of failure.
Data center fire	1	3	3	Could cause the loss of our entire computing infrastructure.
Extended power outage	2	2	4	Won't damage data, but could keep us down for a long time.
Flooding	2	3	6	This area has a history of flooding. With the data center in a low floor, results could be catastrophic.
Chemical spill	1	3	3	An interstate highway goes within 200 yards of the front door. Always a small risk.
Tornado	1	2	2	Would have to hit the building to be a major problem.
Earthquake	2	3	6	Can, of course, be a major disaster.
Bioterrorism	1	3	3	Unlikely, but if it happened, could be serious.

Marcus E, Stern H., "Blueprints for high availability"; 2003; Wiley; ISBN: 0471430269;

HA design – 20 goals

1. Don't be cheap (€ £ \$)
2. Assume nothing from constructors, trades, etc...
3. Remove single points of failure (SPOF)
4. Enforce security (physical, logical)
5. Redefine number of servers
6. Performance measurements
7. Enforce changes to the configuration

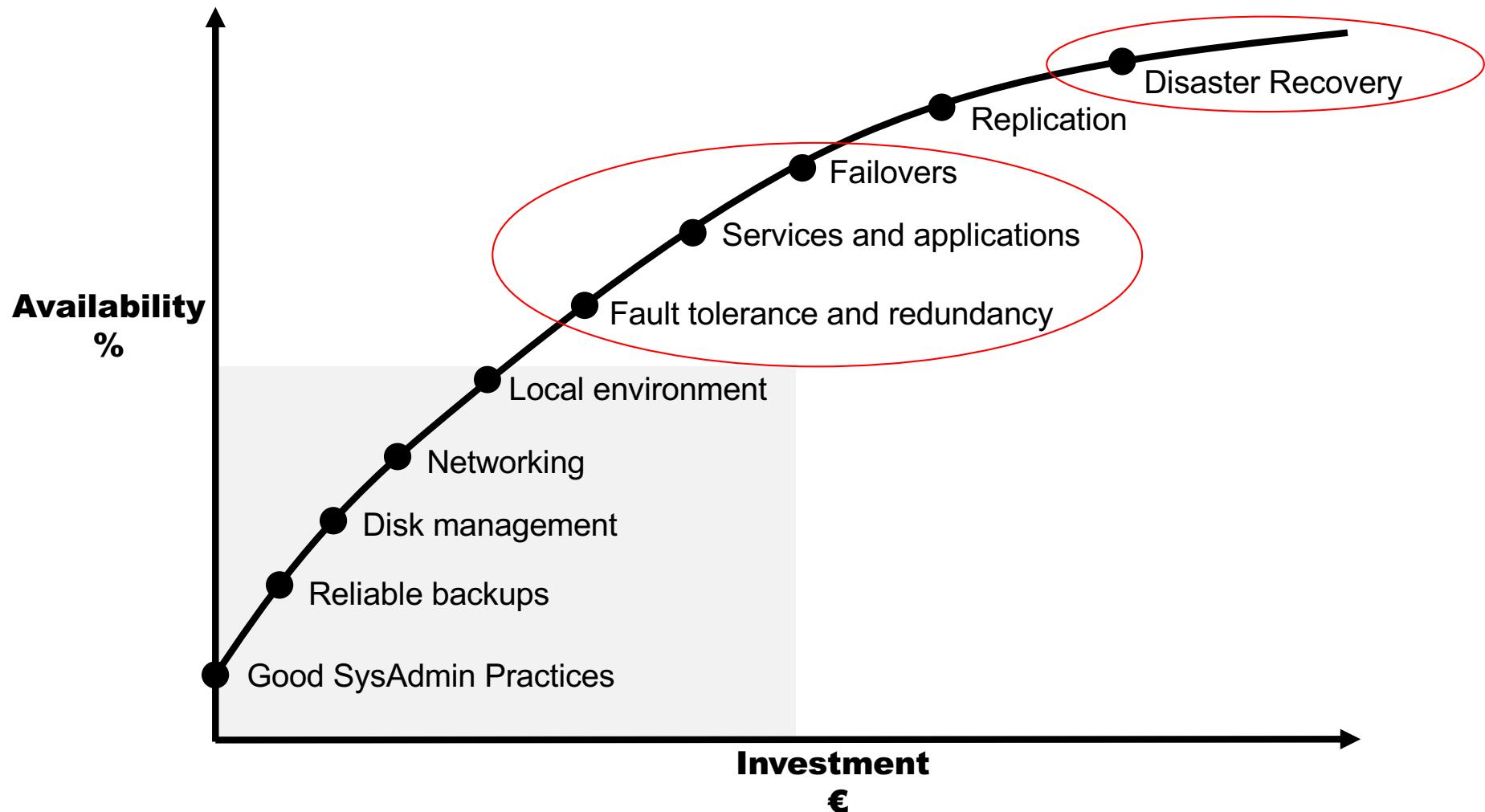
HA design – 20 goals

8. Document everything
9. Define and apply SLA
10. Proactive planning
11. Test everything
12. Separate environments (production, test, QA, dev,...)
13. Learn from history
14. Design for growth

HA design – 20 goals

15. Choose mature software with support
16. Choose mature and reliable hardware
17. Reuse configurations
18. Resources (papers, BCPs, bibliography, ...)
19. One problem, one solution!
20. “Keep it simple!”

Availability index

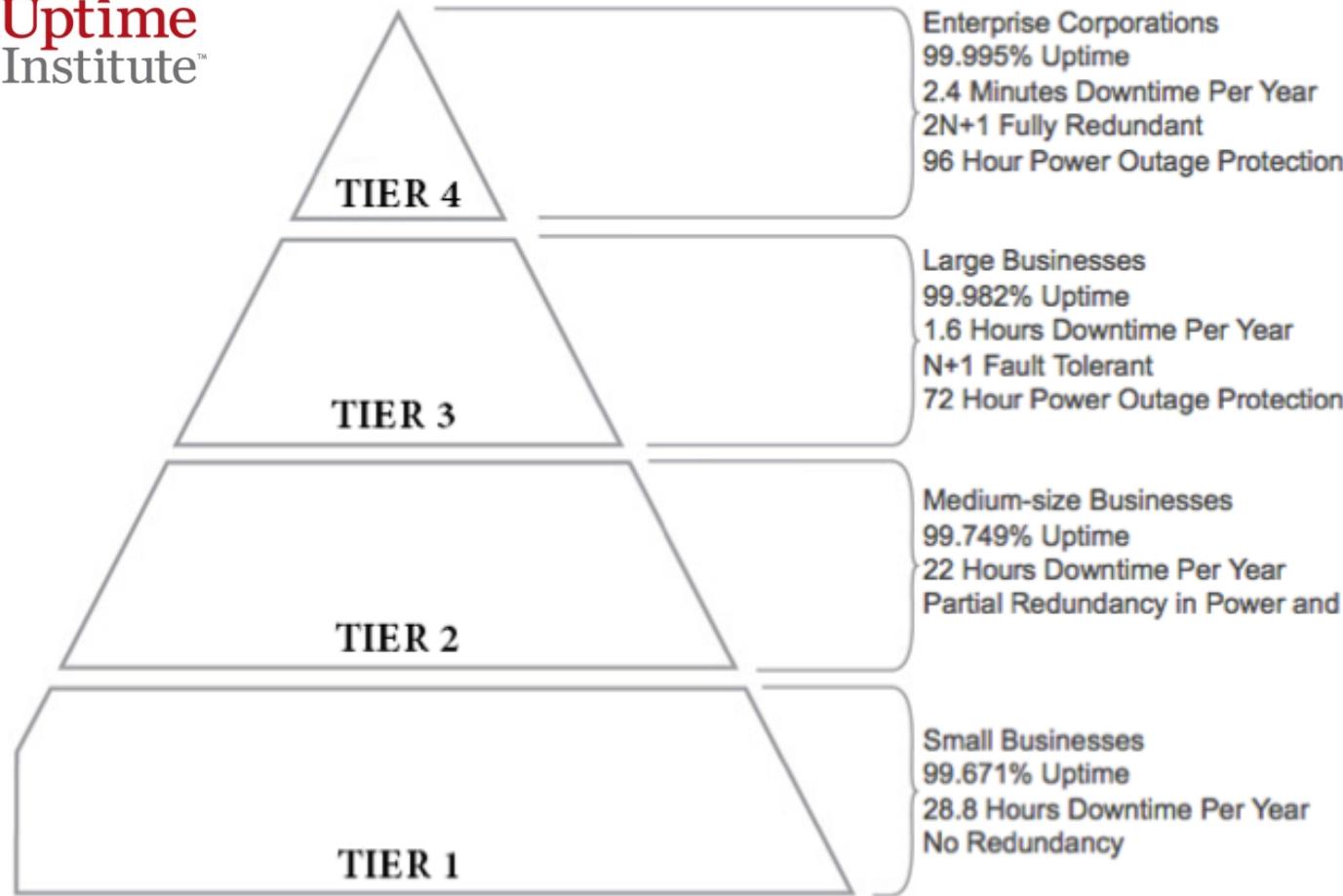


Adapted by Marcus E, Stern H., "Blueprints for high availability"; 2003; Wiley; ISBN: 0471430269;

Availability index

Availability measure is used to evaluate datacenters

Uptime
Institute™



Source: www.datacenters.com

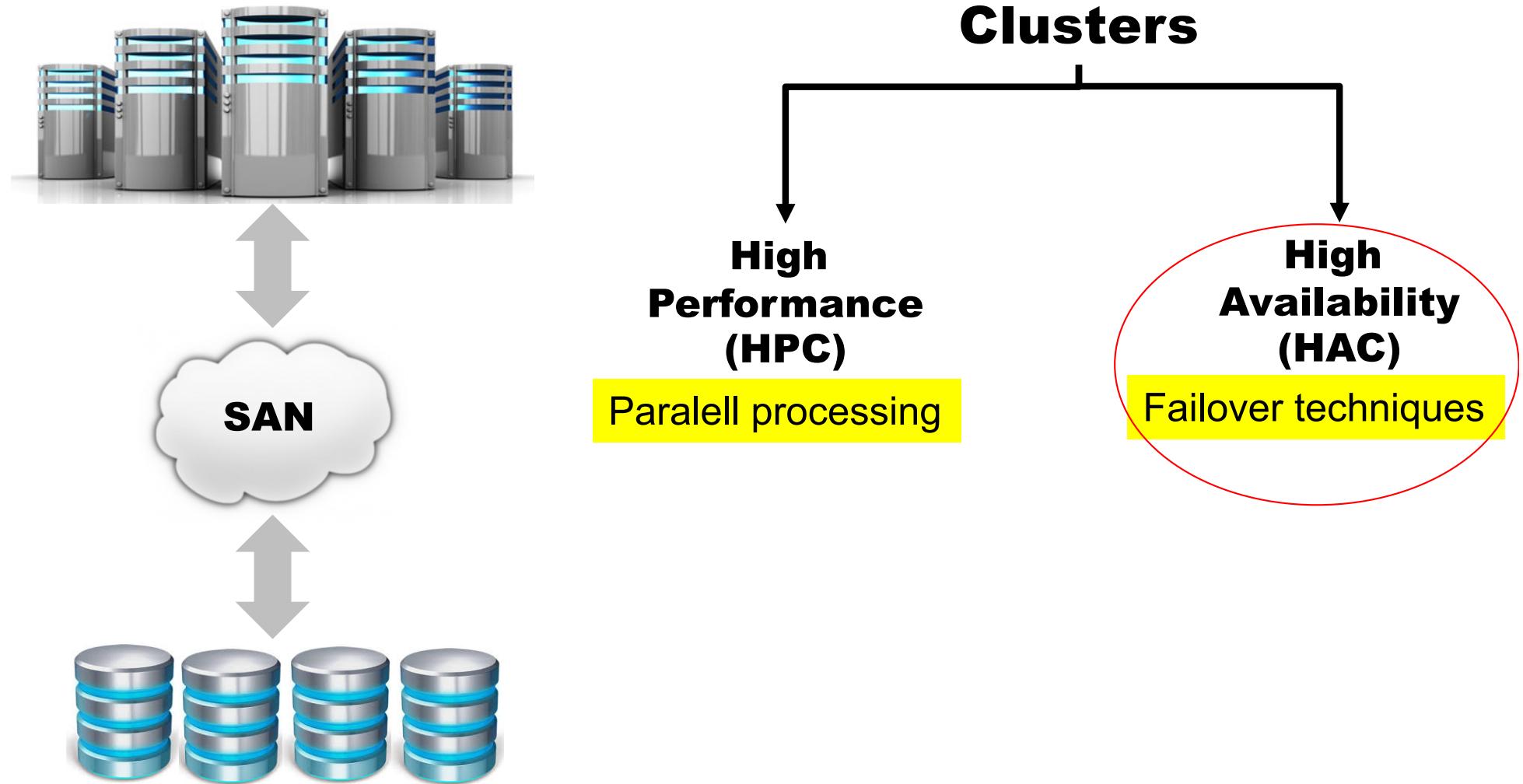
Fault tolerance
and continuous
availability

Continuous
maintenance

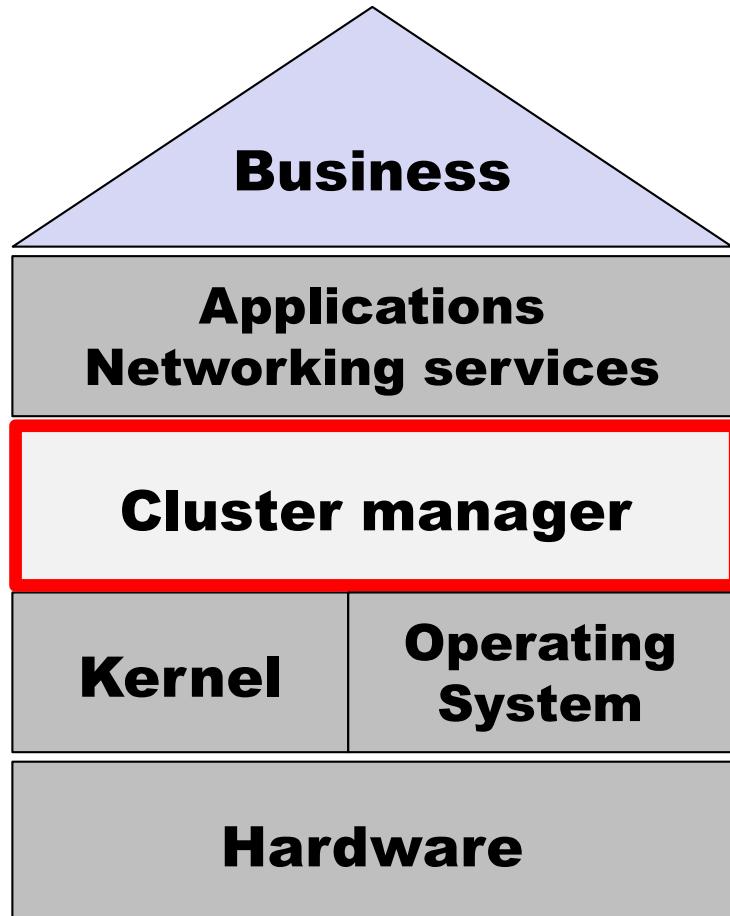
Hardware
redundancy

Basic
facilities

Clustering technologies



Clustering technologies - HA



Cluster managers

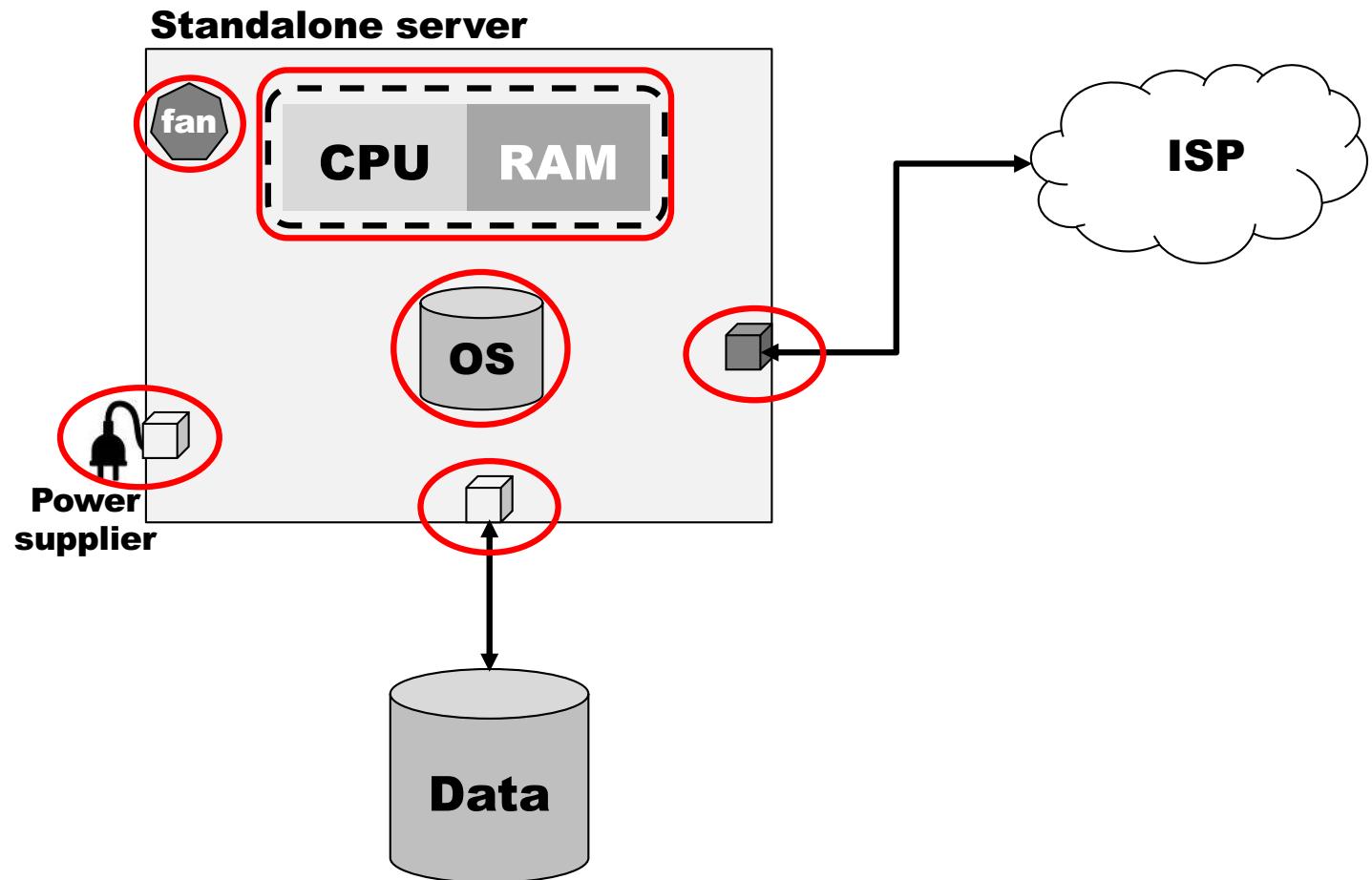
- OS dependent
- Mandatory to be in the market!
- Manufacturers awareness

Opensource cluster managers

- Mainly for Linux OS
- Linux Virtual Server (LVS)
- Heartbeat: www.linux-ha.org

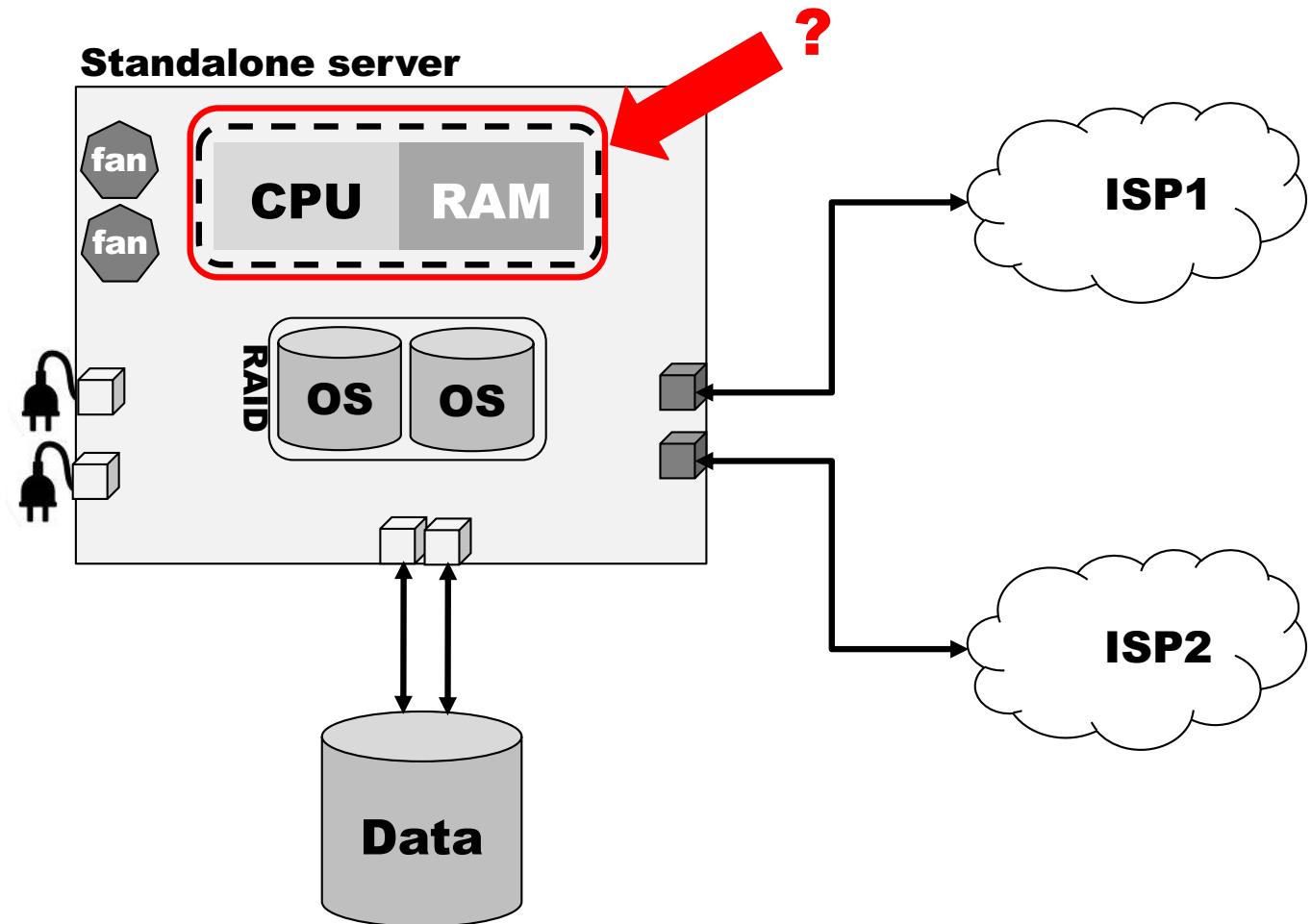
Clustering technologies – HA

Identify Single Points of Failure (SPoF)



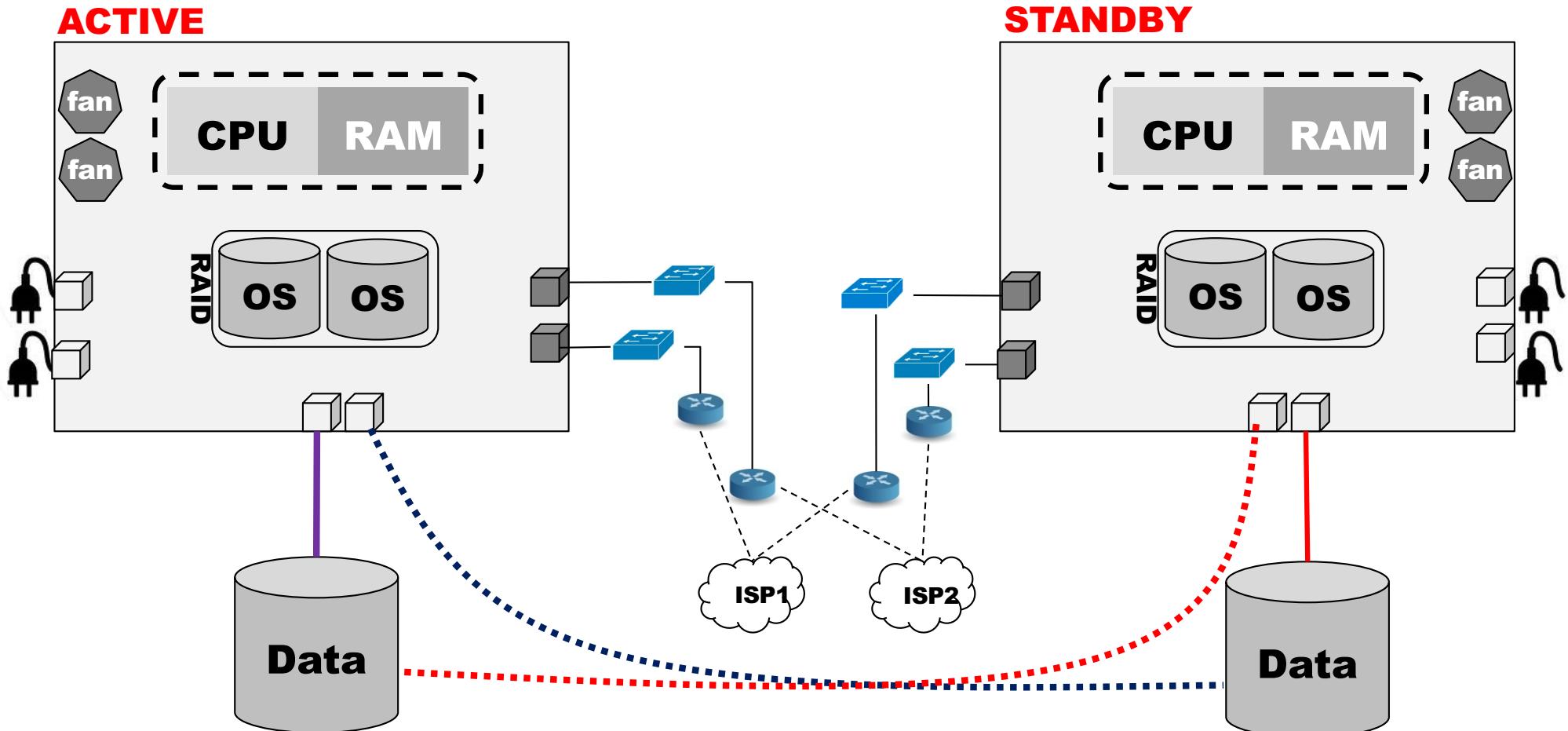
Clustering technologies - HA

Eliminate (mitigate) SPoF (1)



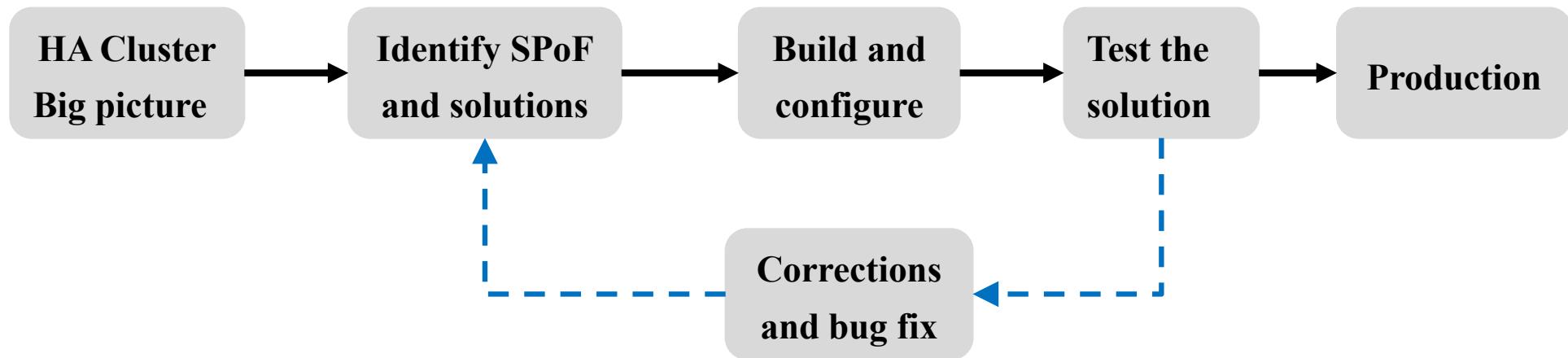
Clustering technologies - HA

Eliminate (mitigate) SPoF (2)



Clustering technologies – HA - planning

General methodology to setup a HA cluster



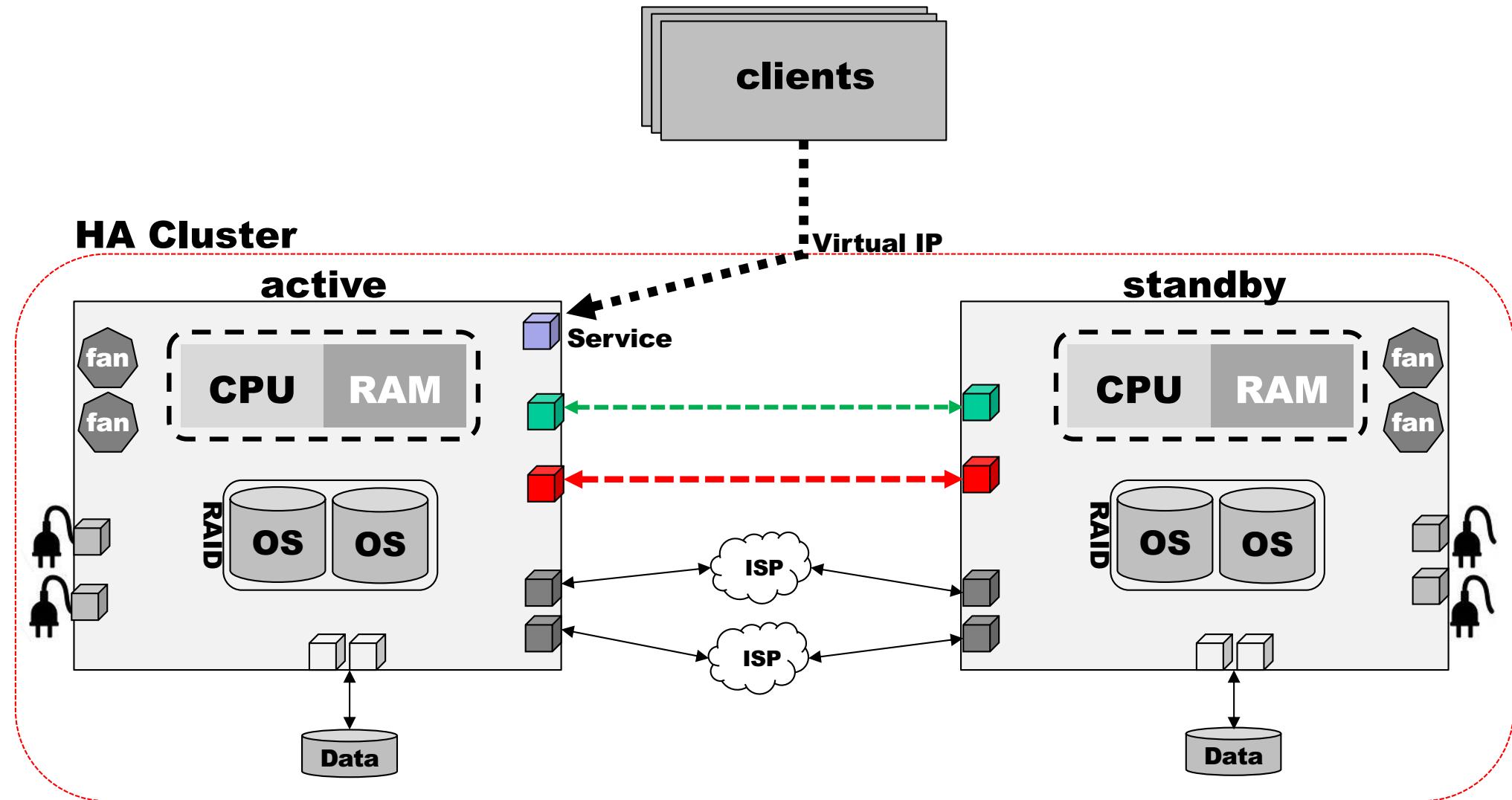
- Planning is crucial
- Level of availability may required mixed solutions
- There is not one solution that fit all problems
- Practice ... practice and more practice!

Clustering technologies - HA

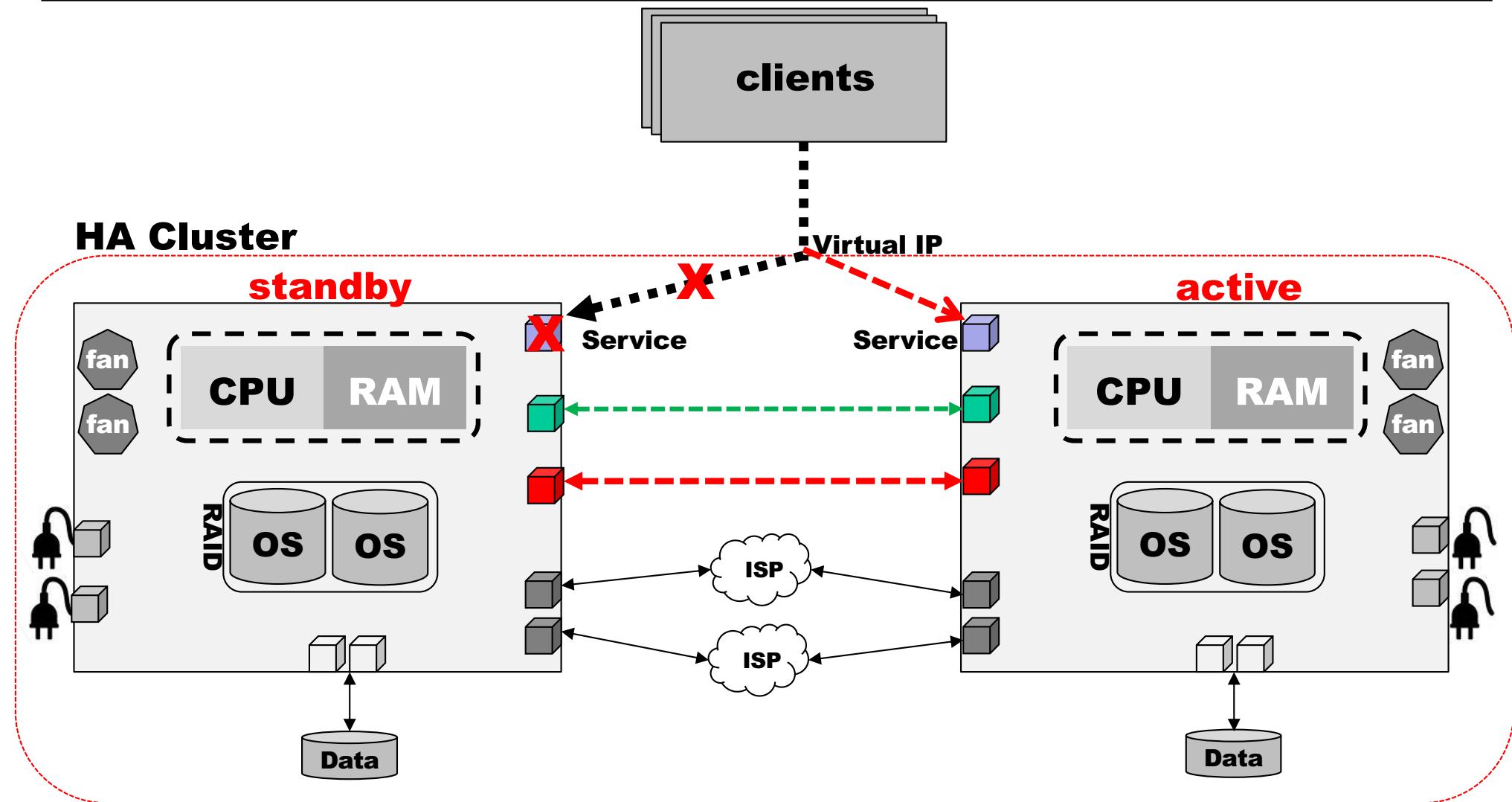
To have in mind:

- Define attainable goals
- Define and test different types of scenarios
- Different strategies: “*Active-Active*” or “*Active-Standby*”
- Network paths should also be replicated: racks, switches, routers, ISP...
- Applications should be “*HA ready!*”
- Organizations should be “*HA aware*”

Cluster Resource Manager

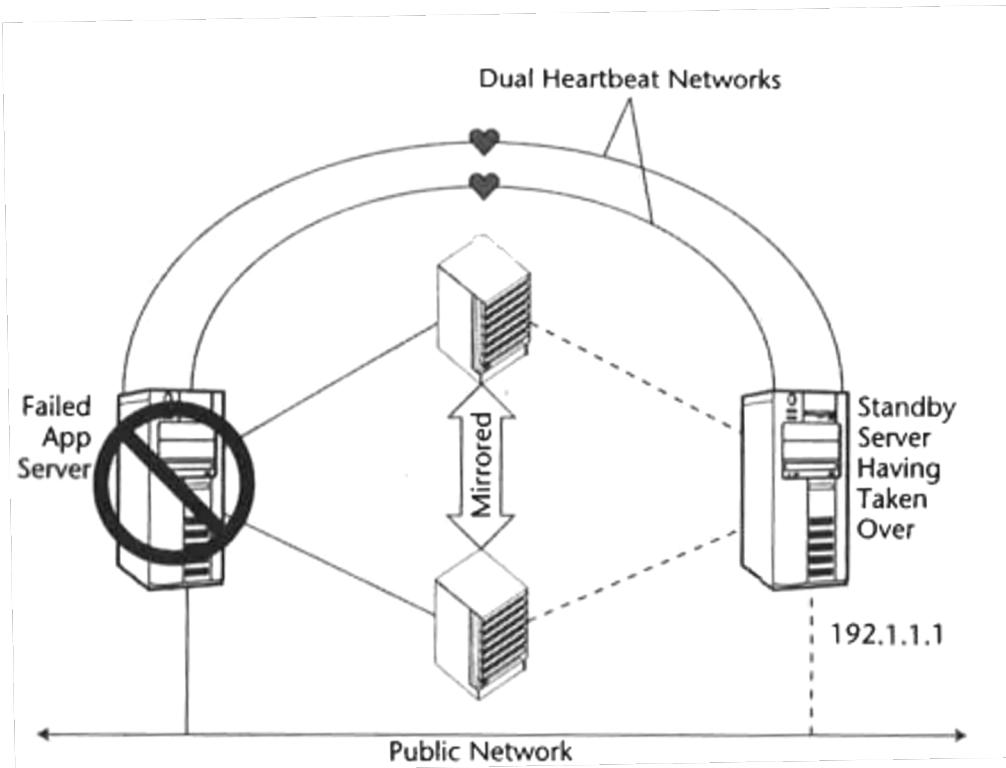


Cluster Resource Manager



Cluster Resource Manager

“Active-Passive” configuration



- A standby node
- Automatic recover after failure
- heartbeat between two nodes

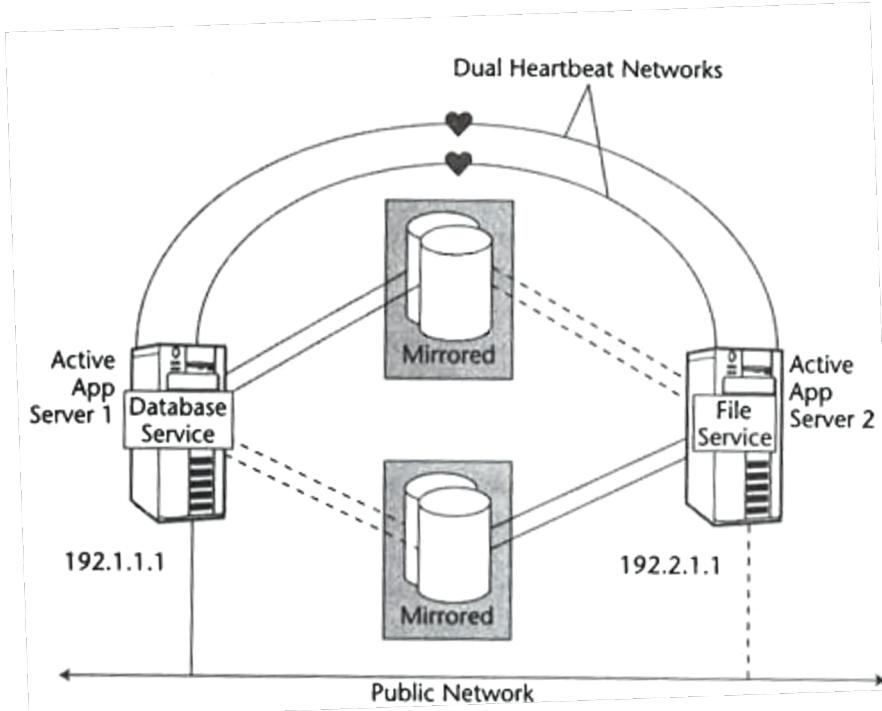
How to use standby nodes:

- ✓ software development
- ✓ Other applications
- ✓ QA and/or test system
- ✓ Solely standby node

Marcus E, Stern H., "Blueprints for high availability"; 2003; Wiley;
ISBN: 0471430269;

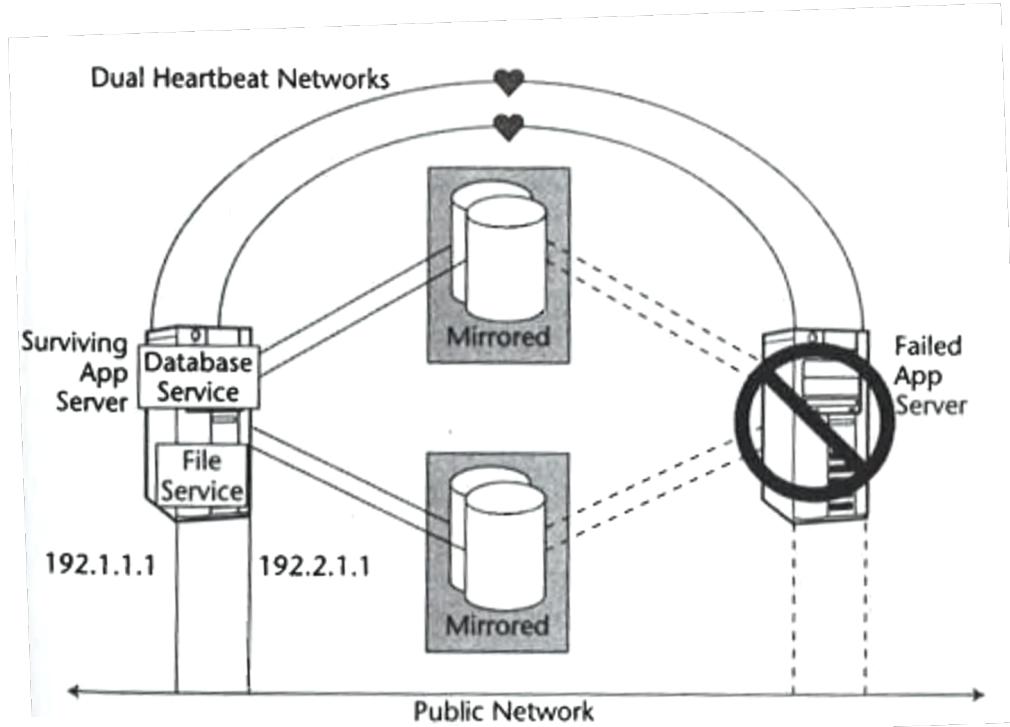
Cluster Resource Manager

“Active-Active” configuration



Better use of hardware

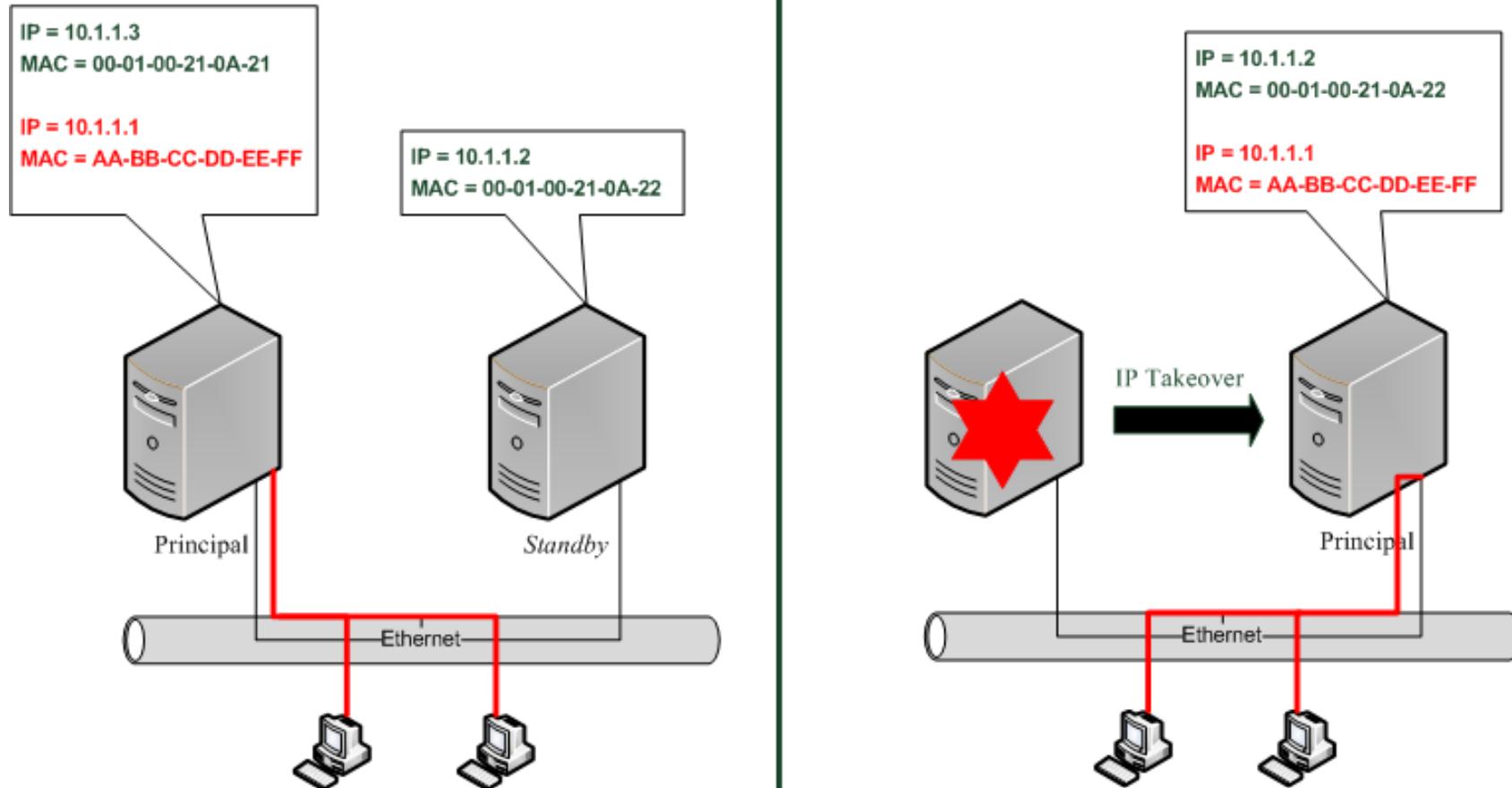
Marcus E, Stern H., “Blueprints for high availability”; 2003; Wiley;
ISBN: 0471430269;



Well tested and independent applications
Low performance after takeover
Management is more complex

Cluster Resource Manager

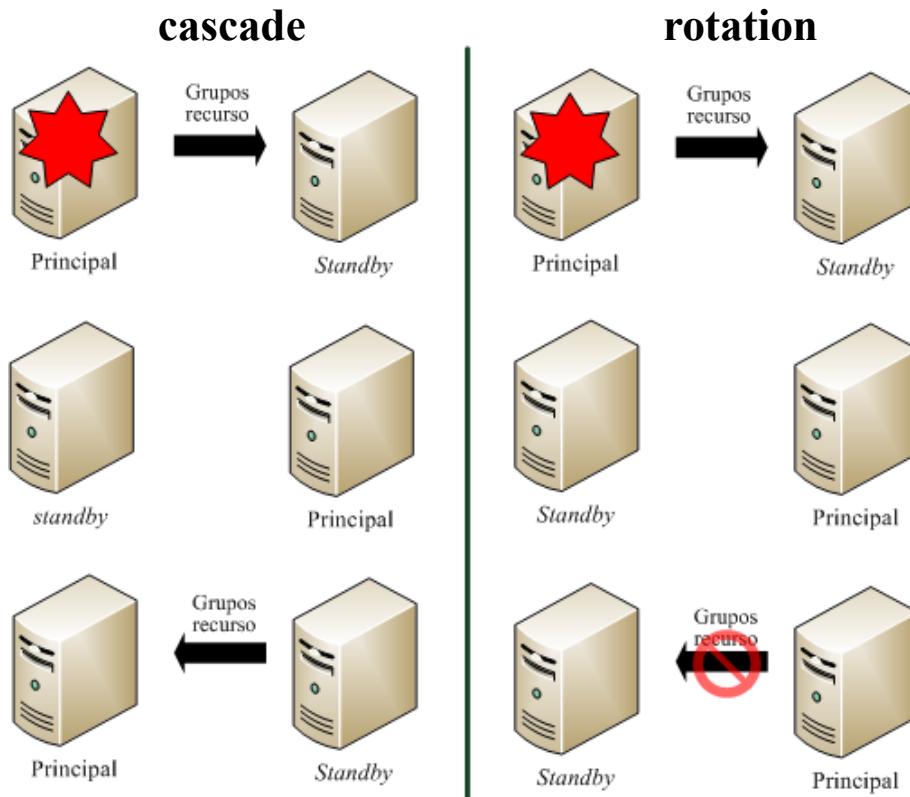
Address takeover



In "Gestão de Projetos TI e administração centralizada de sistemas e redes – cenários práticos em contexto empresarial";
Mário Antunes; IPLeiria; 2010

Cluster Resource Manager

Resource groups – takeover strategies

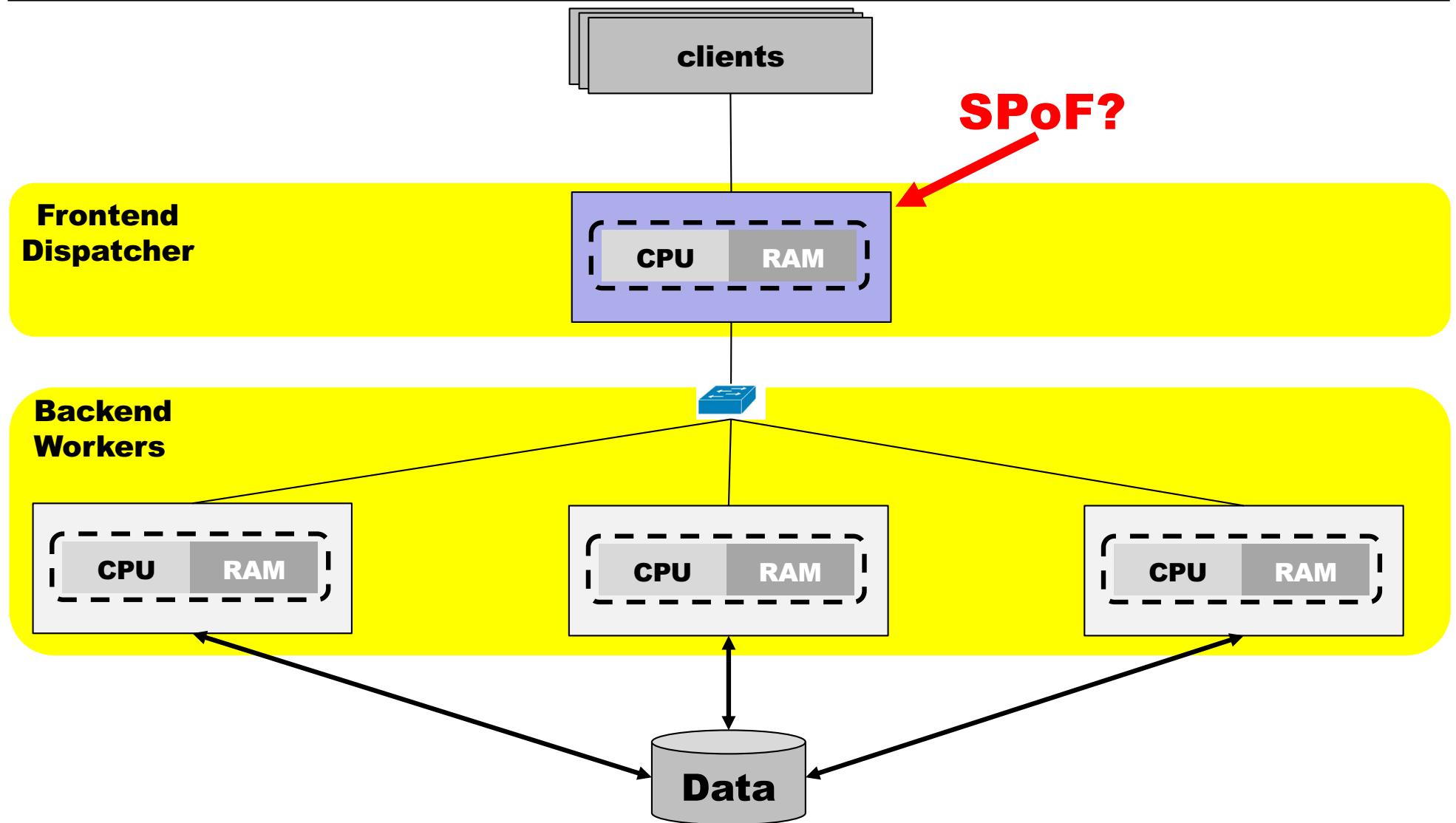


Resource group
critical applications
interfaces

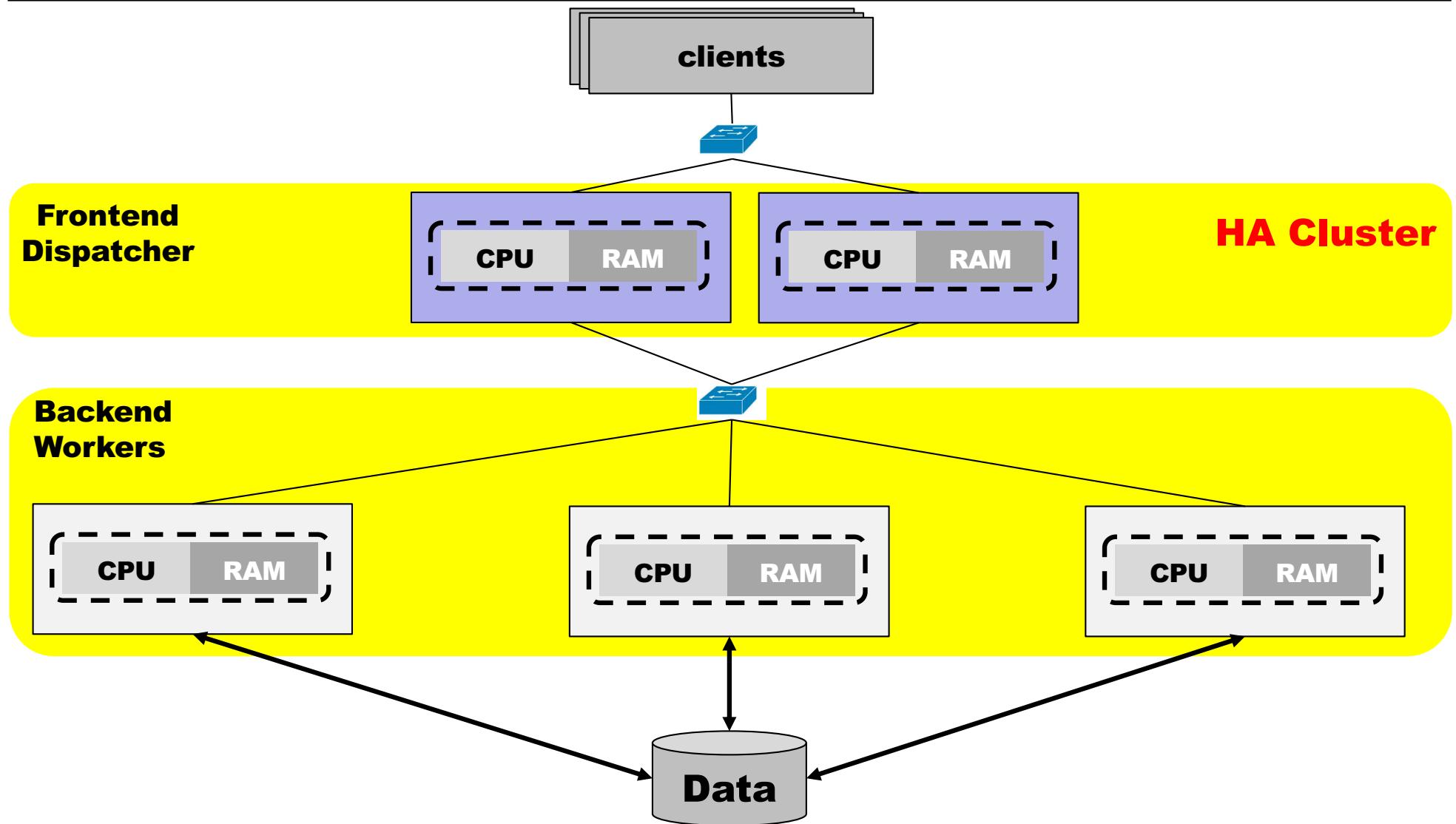
**Which method should
we choose?**

In "Gestão de Projetos TI e administração centralizada de sistemas e redes – cenários práticos em contexto empresarial";
Mário Antunes; IPLeiria; 2010

Cluster Resource Manager - Load balancing



Cluster Resource Manager - Load balancing



Cluster Resource Manager - Challenges

- Geographical distribution of the cluster nodes
- Infrastructure monitoring effectiveness
 - Cluster manager tools
 - Applications monitoring tools
 - Customized monitoring applications
- “*split-brain*” syndrome and how to mitigate it
- How to use backup node during idle period?

Final remarks

- Business are highly IT dependent
- Availability is a *must* and a *continuous challenge*
- Companies and businesses are too aware to HA thinking
- Basic concepts to cover core business and IT demands
- Mixed topologies to cover HA and LB features
- Keyword: **investment** (\$£€).

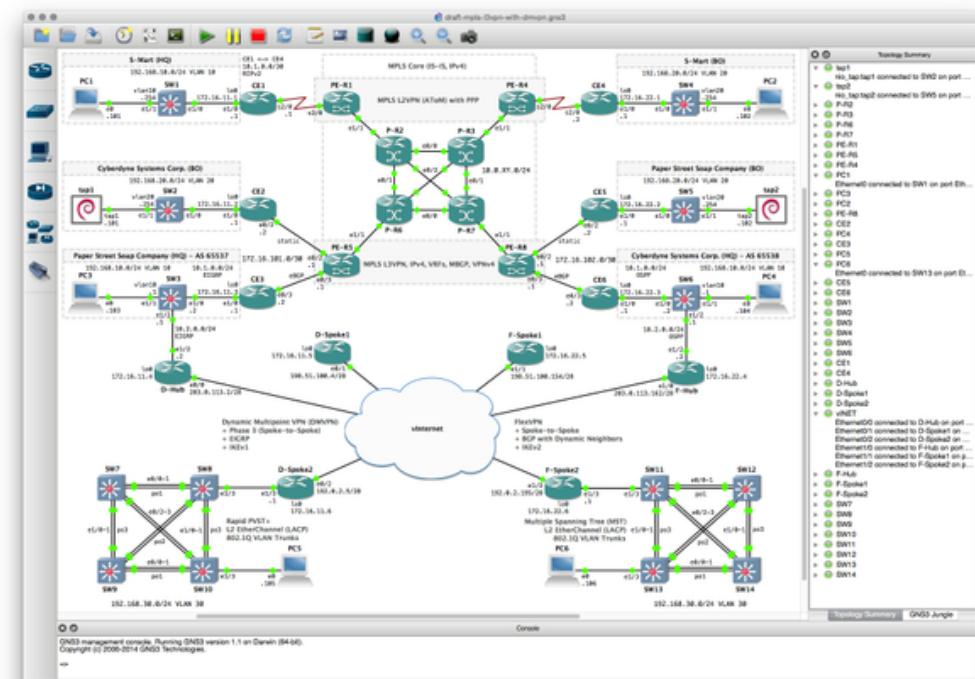
IT professionals should be *HA aware!*

High availability clusters

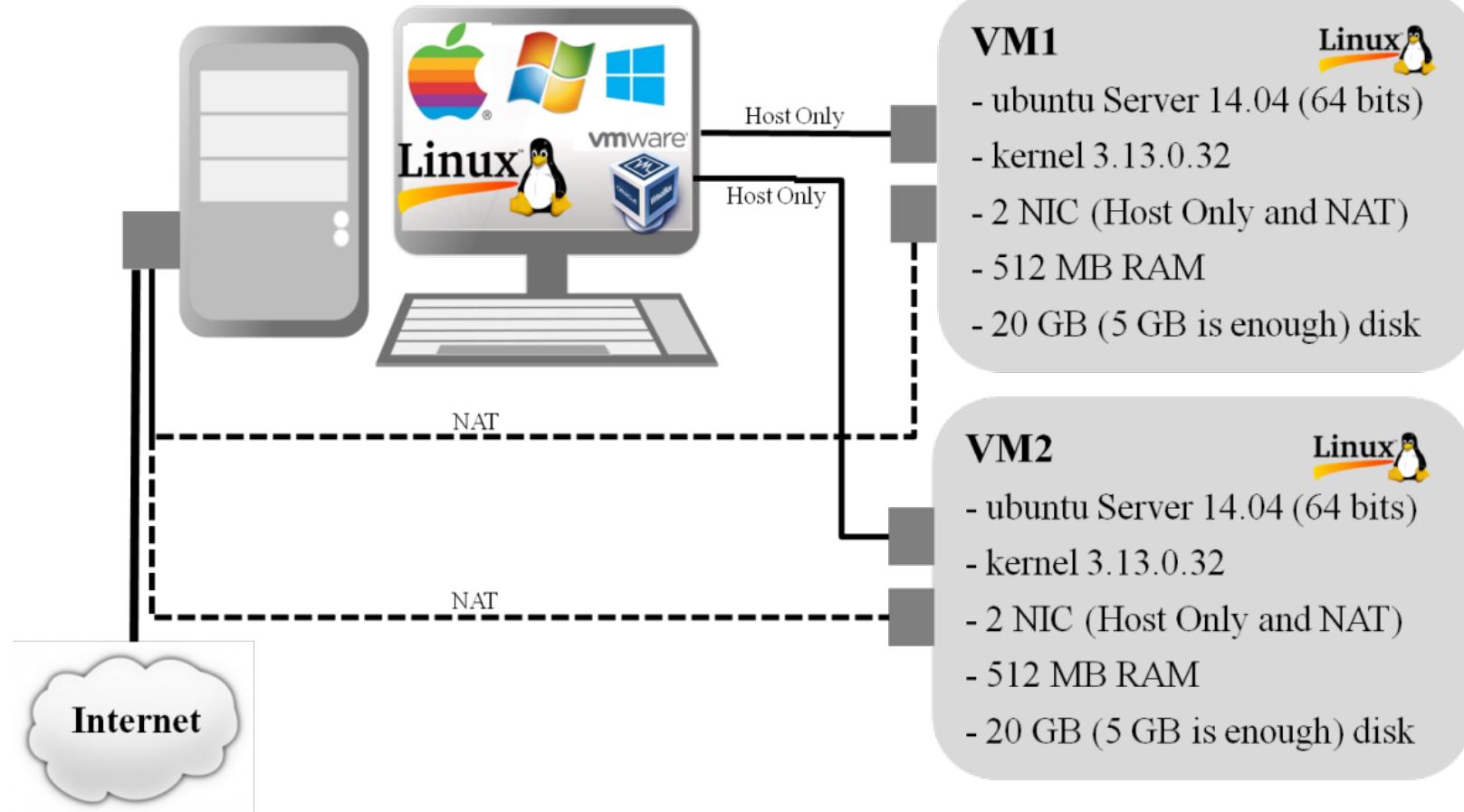
- case study with Heartbeat -

Lab setup – Visualization tool

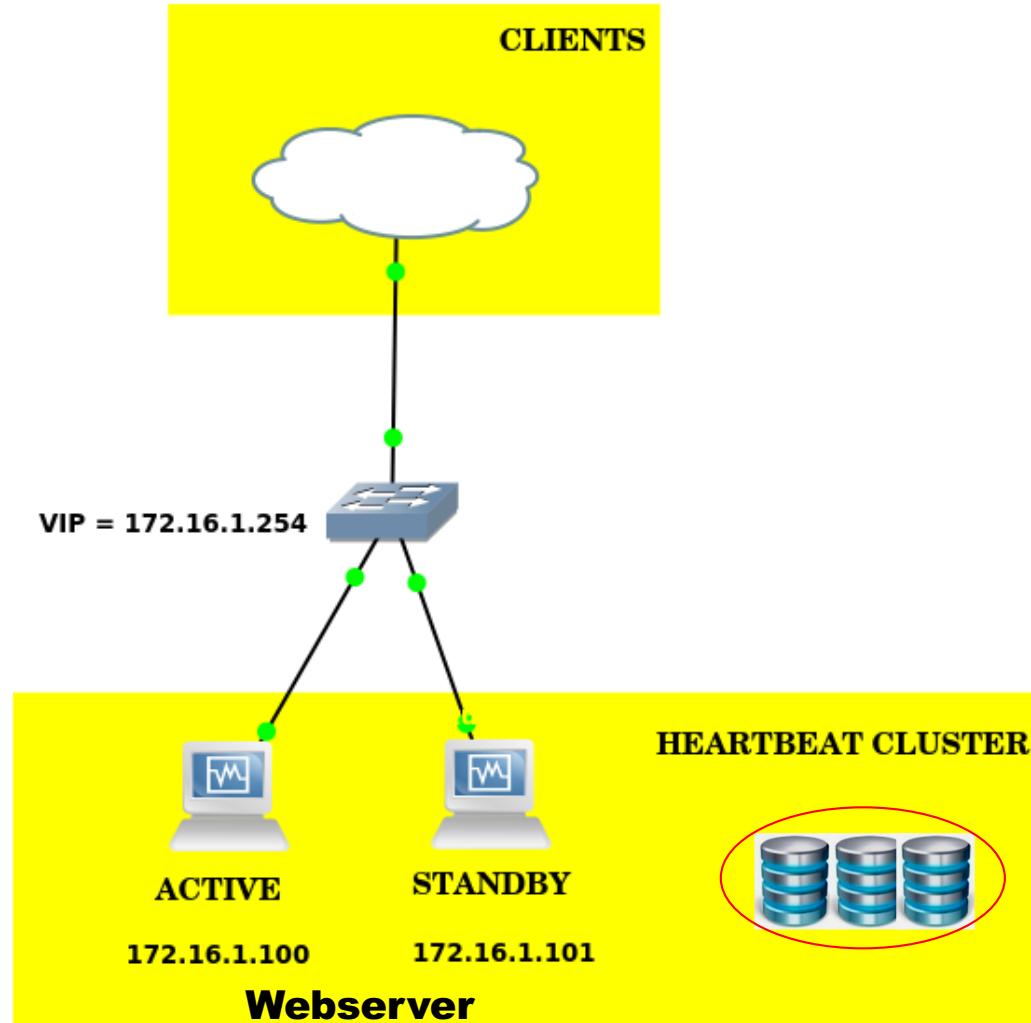
- Virtualization and emulation of network equipments and servers
- Integrates easily with Virtualbox
- Native hypervisor – run dynamips
- Graphical and visualization features
- Opensource and free!



Lab setup – Visualization tool



Lab setup - heartbeat



Network setup

Webserver setup

Heartbeat setup

Logging

Config files (`/etc/ha.d/`)

authkeys

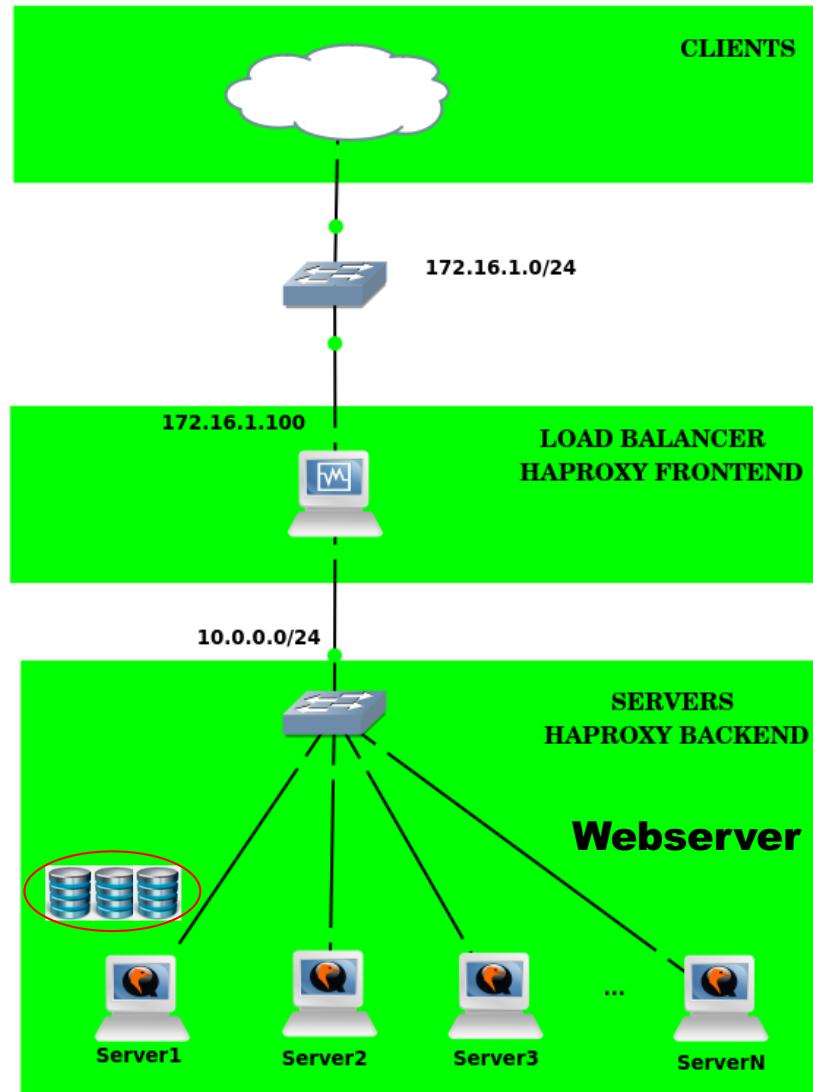
haresources

ha.cf

High availability clusters

- case study with HAProxy -

Lab setup - HAProxy



Web clients

- Browser launches on “host” system

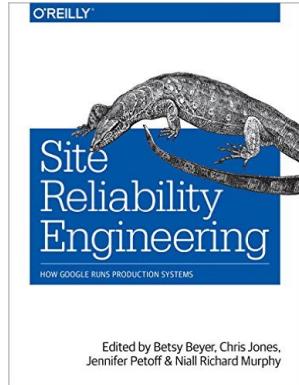
HAProxy Server

- Forward Web requests/replies
- Monitoring and configuration rules
- `/etc/haproxy/haproxy.conf`

Web servers

- Web server configuration rules
- Centralized storage access

Final remarks - bibliography



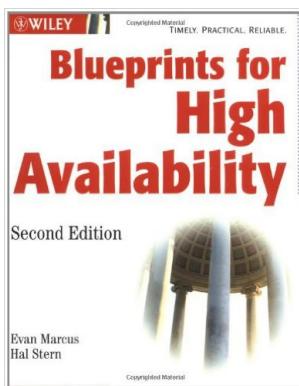
Site Reliability Engineering

Jennifer Petoff, Niall Richard Murphy, Betsy Beyer, Chris Jones
How Google Runs Production Systems
O'Reilly Media; 1 edition (April 16, 2016)
ISBN-13: 978-1491929124

<http://www.wired.com/2016/04/google-ensures-services-almost-never-go/>



<http://www.uptimeinstitute.com>
Reports and research publications



Blueprints for high availability

Marcus E. Stern H.
Wiley; 2003
ISBN: 0471430269

Redundancy

LB and HA in routing protocols

Dedicated HA and LB protocols

Mário Antunes
mario.antunes@ipleiria.pt

Redundancy

Hardware level



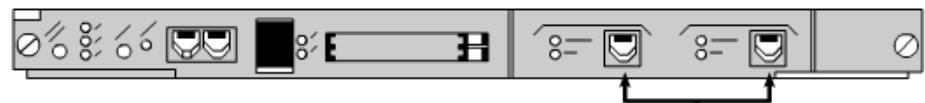
Catalyst 5000,5500, 6000, 6500



Uplinks redundantes

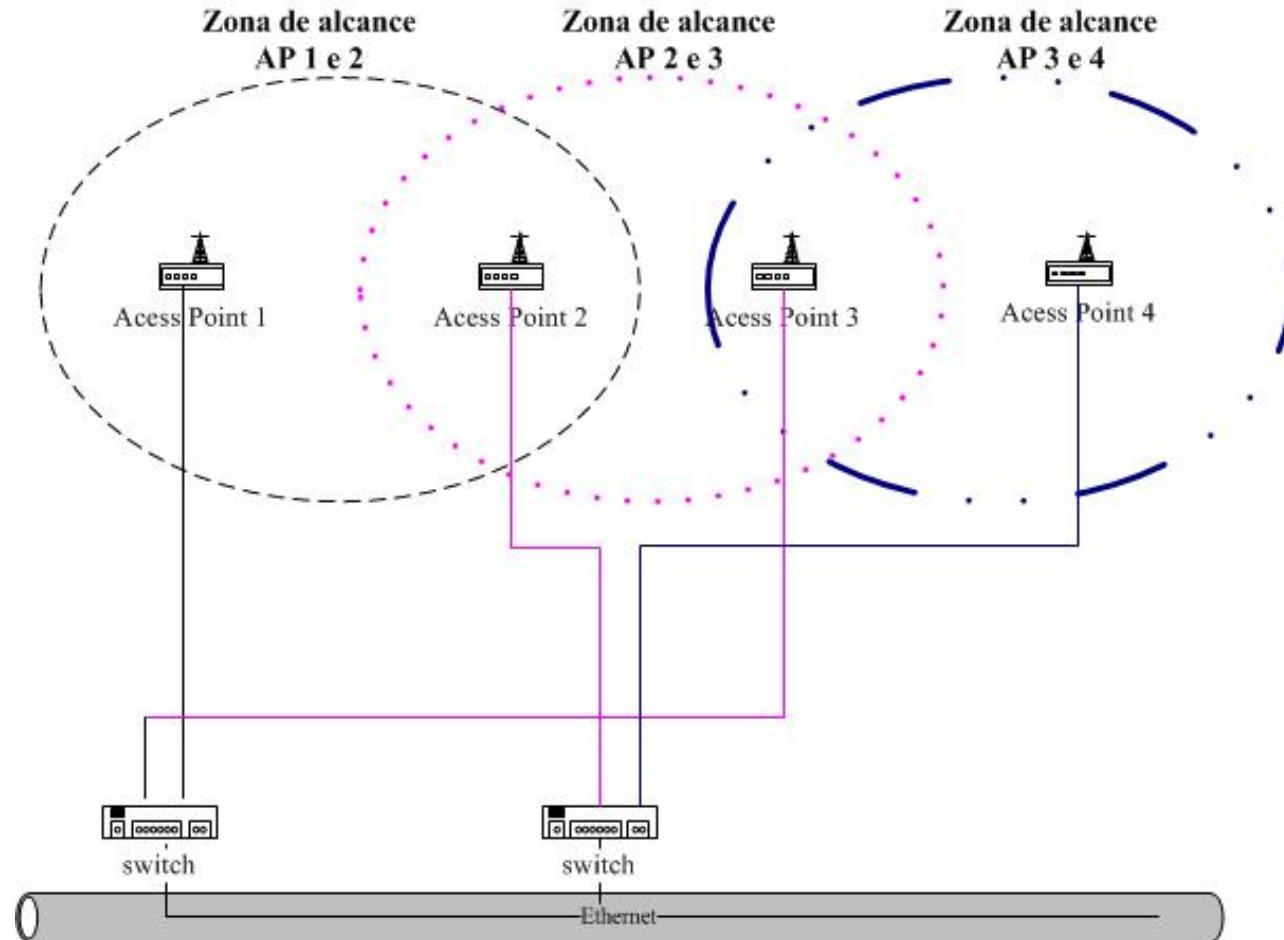


Power Supply



10/100BaseTX Fast EtherChannel
MDIX RJ-45 connections

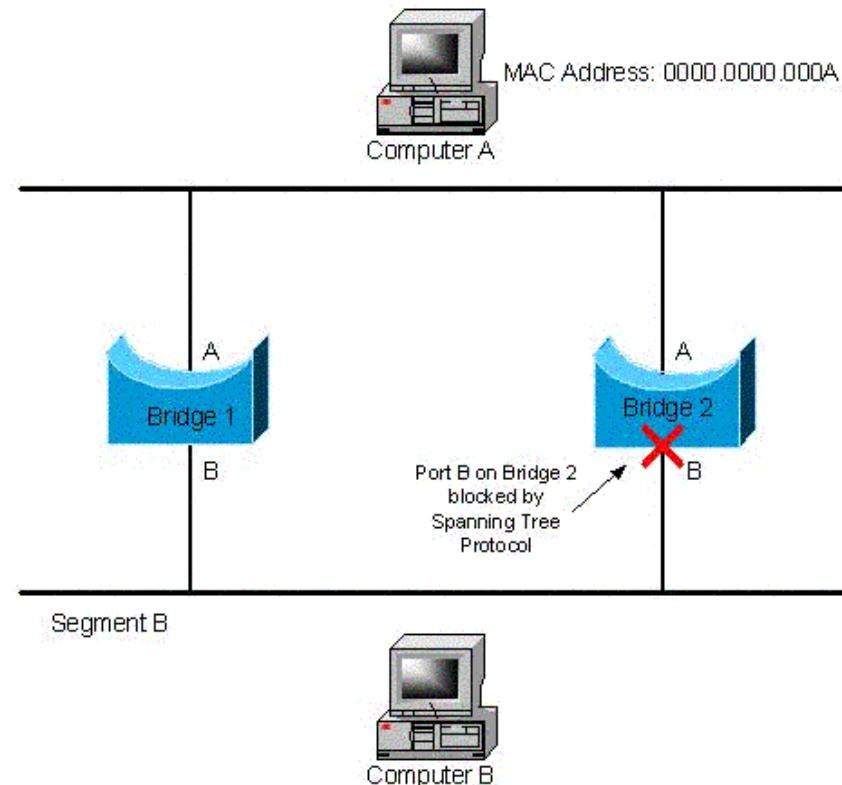
Redundancy - examples



In "Gestão de Projetos TI e administração centralizada de sistemas e redes – cenários práticos em contexto empresarial";
Mário Antunes; IPLEiria; 2010

Redundancy - examples

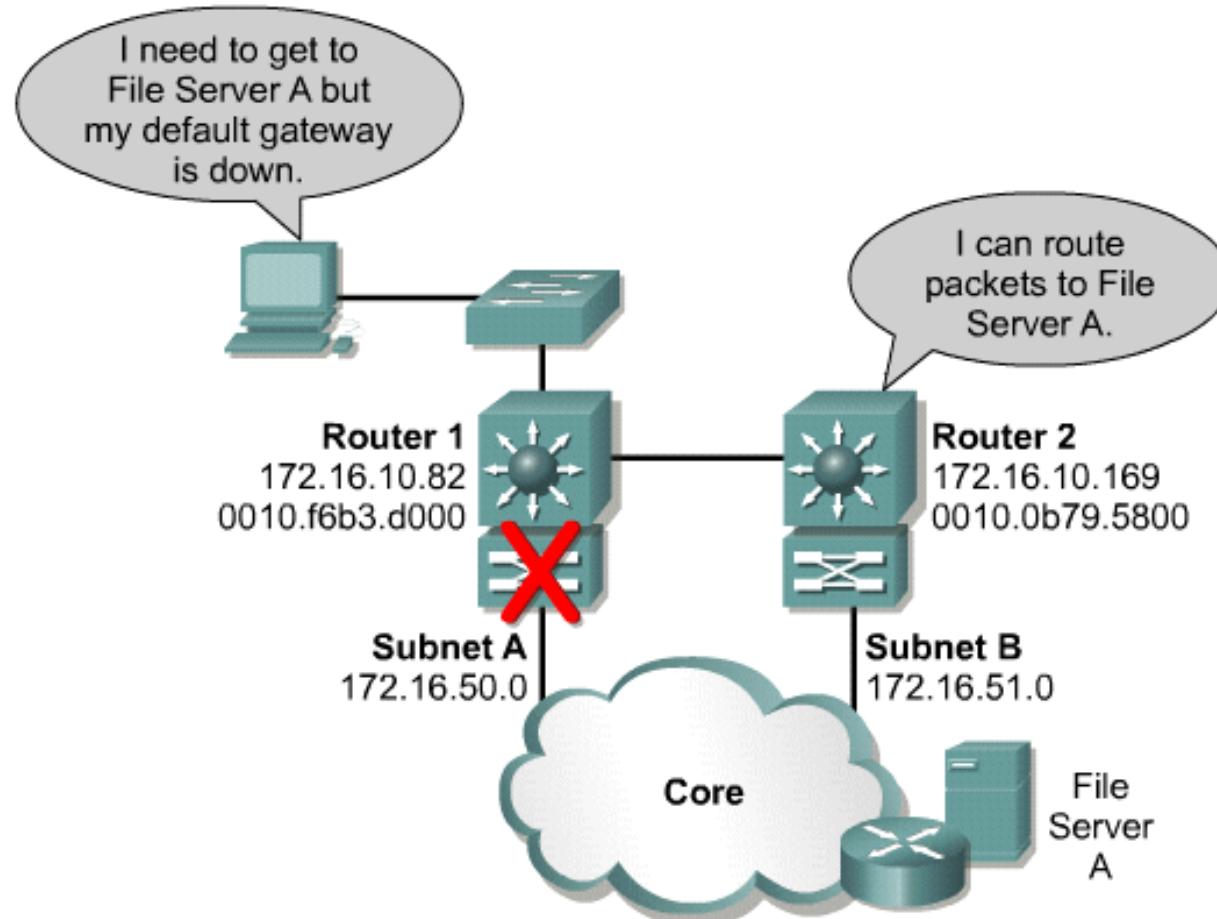
Redundancy L2 – Spanning Tree Protocol (STP)



<http://ipsit.bu.edu>

Case study – L3

The problem:



Case study – L3

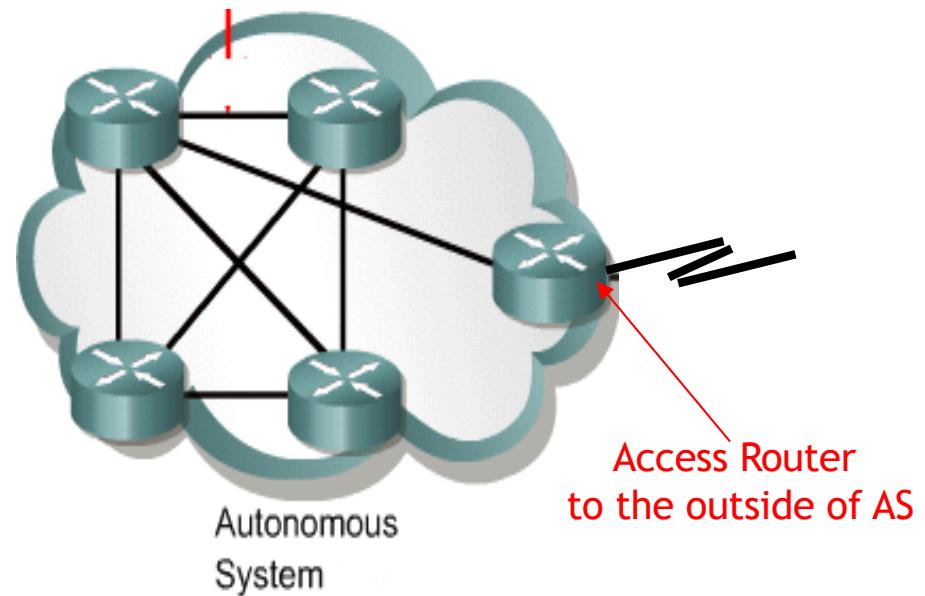
Main topics:

- Dynamic routing
- ICMP Redirect (IR)
- NAT with TCP Load Distribution
- Hot Standby Router Protocol (HSRP)
- Virtual Router Redundancy Protocol (VRRP)
- Gateway Load Balancing Protocol (GLBP)
- Single Router Mode (SRM)
- Server Load Balancing (SLB)

Routing fundamentals

Autonomous System (AS)

- Networks that share the same routing politics
- Usually under the same administrative control
- Identified by a unique identifier (AS Id):
 - 32 bits
 - Assigned by RIR



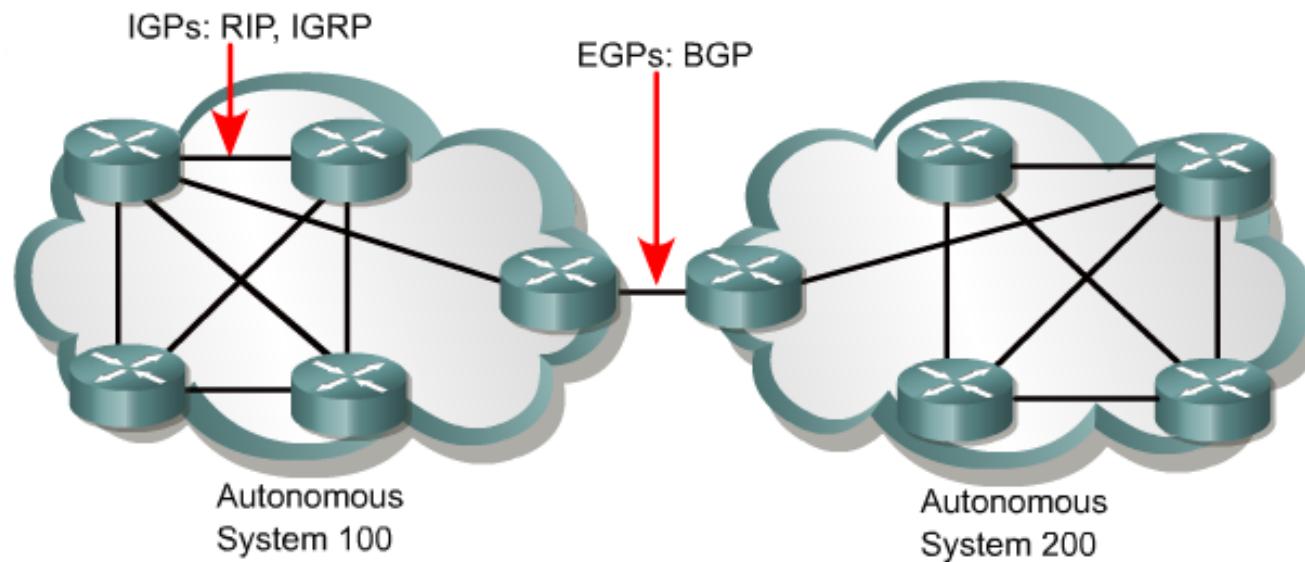
Routing fundamentals

IGP

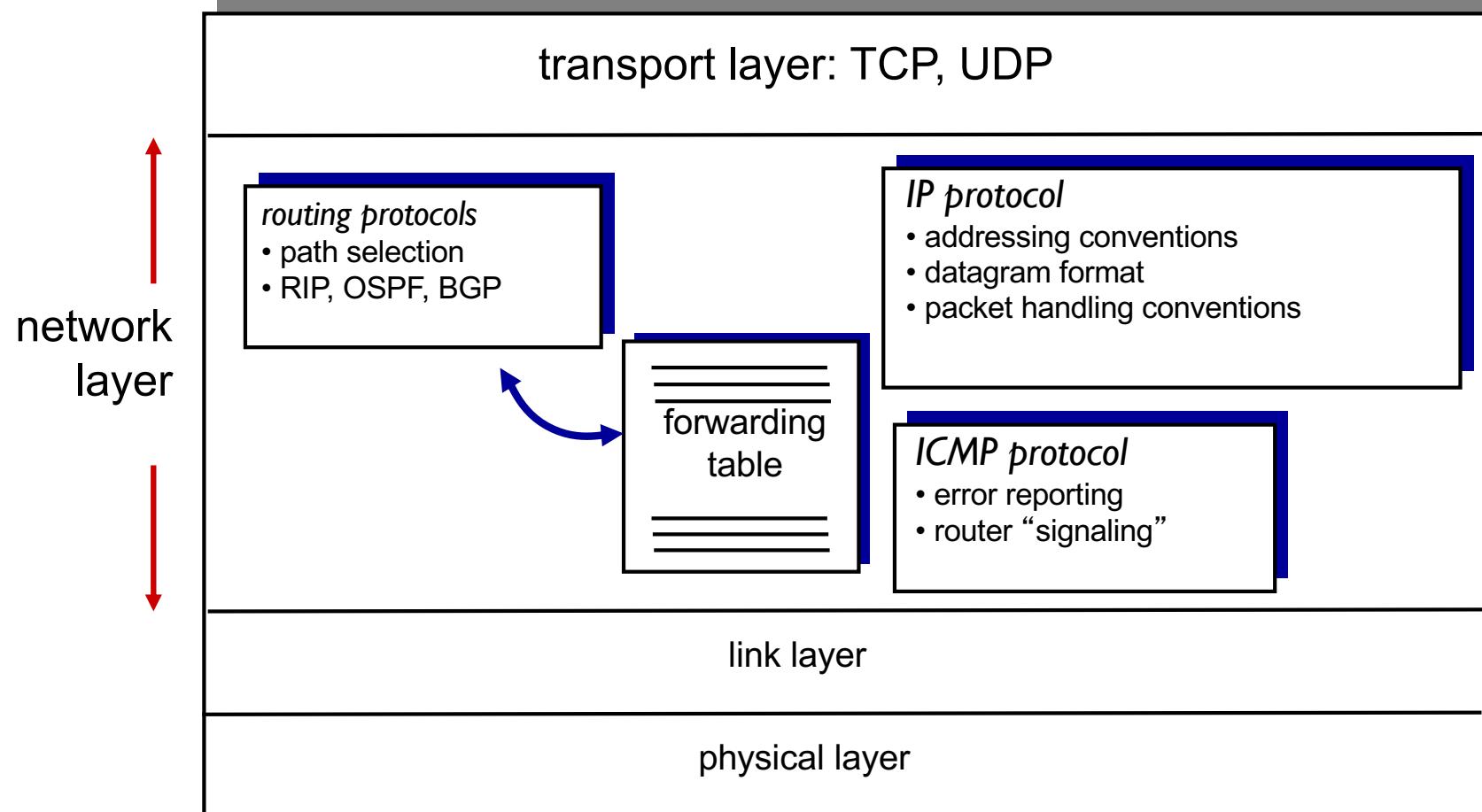
- *Routing inside AS*
- Scalability in the AS
- RIP, IGRP, EIGRP, OSPF

EGP

- *Routing between two AS*
- Hierarchy in large networks
- Management policies
- BGP



Routing fundamentals

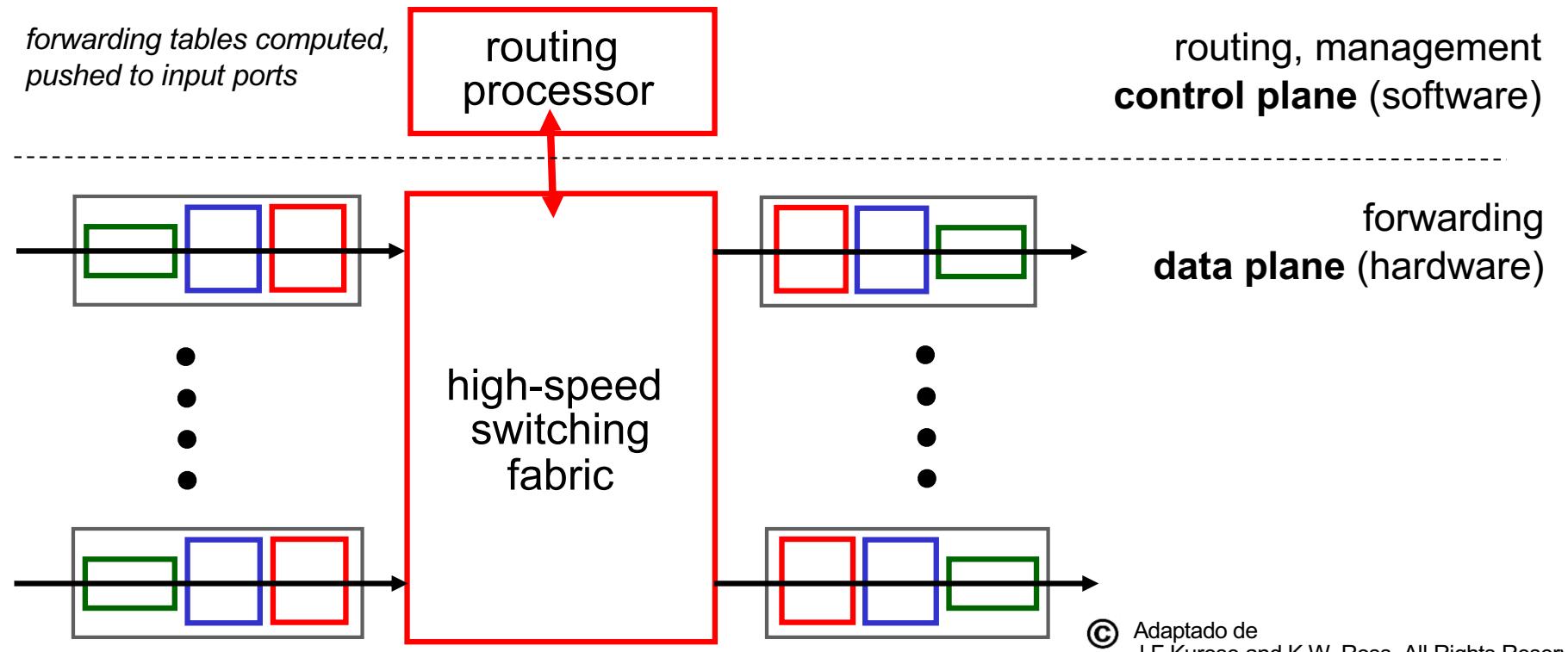


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Routing fundamentals

Main functions:

- Execute routing algorithms (RIP, OSPF, BGP)
- Route (*forwarding*) packets to the destination IP by an interface



Routing fundamentals

- Non-adaptive (static)
- Adaptive (dynamic)
 - Link State
 - Distance Vector

Routing – RIP protocol

RIP

- Through routes with the **same cost** or alternative routes learned by RIP
- Selective adjust of existing timers: *update, invalid, holddown, flush*
- Load balance by routes with the same cost
- Maximum number of connections:

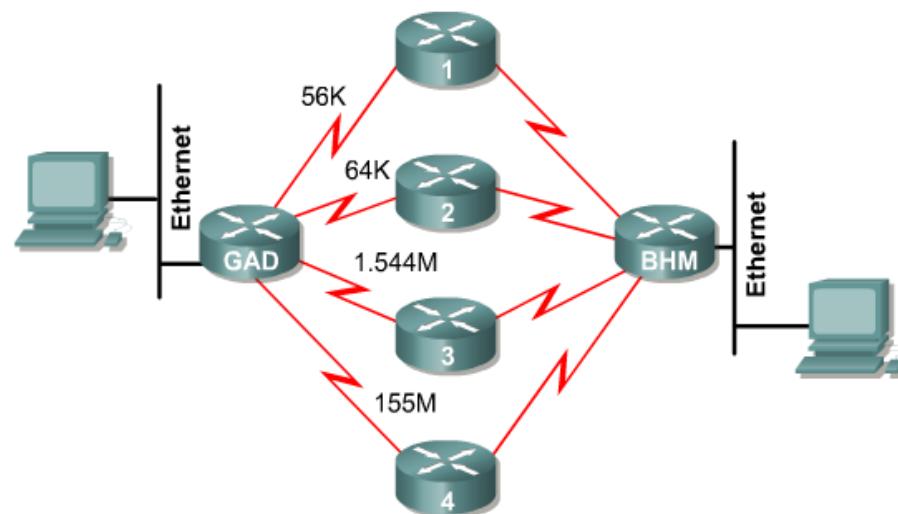
```
R1 (config)#router rip
```

```
R1 (config-router)#maximum-path 2
```

Routing – RIP protocol

Load balancing

- Load balancing between connections with the better equal cost
- RIP uses up to 32 connections (default=4)
(router(config-router) #**maximum-paths N**)
- Method used is “*round-robin*”



```
RouterC#show ip route 192.168.2.0
Routing entry for 192.168.2.0/24
Known via "rip", distance 120, metric 1
Redistributing via rip
Last update from 192.168.4.2 on FastEthernet0/0,
00:00:18 ago
Routing Descriptor Blocks:
192.168.4.1, from 192.168.4.1, 00:02:45 ago, via
FastEthernet0/0
Route metric is 1, traffic share count is 1
* 192.168.4.2, from 192.168.4.2, 00:00:18 ago,
via FastEthernet0/0
Route metric is 1, traffic share count is 1
```

Routing – EIGRP protocol

- Enhanced Interior Gateway Routing Protocol
- Distance vector
- Cisco proprietary
- Periodic updates: 90 seconds
- Main features:
 - Operates well in complex networks
 - Flexibility in networks with links that have distinct characteristics
 - Scalability in large scale networks

Routing – EIGRP protocol

Metric

- Default: BW e DLY
- Composed metric:
 - BW
 - DLY
 - LOAD
 - Fiability
 - MTU

$$\text{Metric} = [K1 * \text{BW} + (K2 * \text{BW}) / (256 - \text{LOAD}) + K3 * \text{DLY}]$$

Se $K5 \neq 0 \rightarrow \text{Metric}^* = [K5 / (\text{RELIABILITY} + K4)]$

Default = **BW + DLY**

```
Router>show ip protocols
Routing Protocol is igrp 300
  Sending updates every 90 seconds, next due in 55
  seconds
  Invalid after 270 seconds, hold down 280, flushed
  after 360
  Outgoing update filter list for all interfaces is
  not set
  Incoming update filter list for all interfaces is
  not set
  Default networks flagged in outgoing updates
  Default networks accepted from incoming updates
  IGRP metric weight K1=1, K2=0, K3=1, K4=0, K5=0
  IGRP maximum hopcount 100
  IGRP maximum metric variance 1
  Redistributing igrp 300
```

“K” parameters (default)

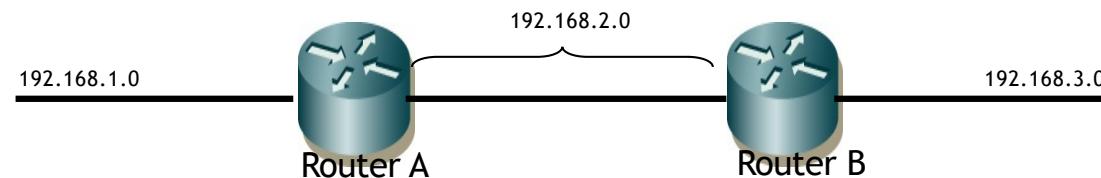
- **K1 = BW = 1**
- **K2 = fiability = 0**
- **K3 = DLY = 1**
- **K4 = LOAD = 0**
- **K5 = MTU = 0**

Routing – EIGRP protocol

Configuration:

EIGRP on AS 101
RouterA(config)#**router igrp 101**
Interfaces RouterA(config-router)#**network 192.168.1.0**
RouterA(config-router)#**network 192.168.2.0**

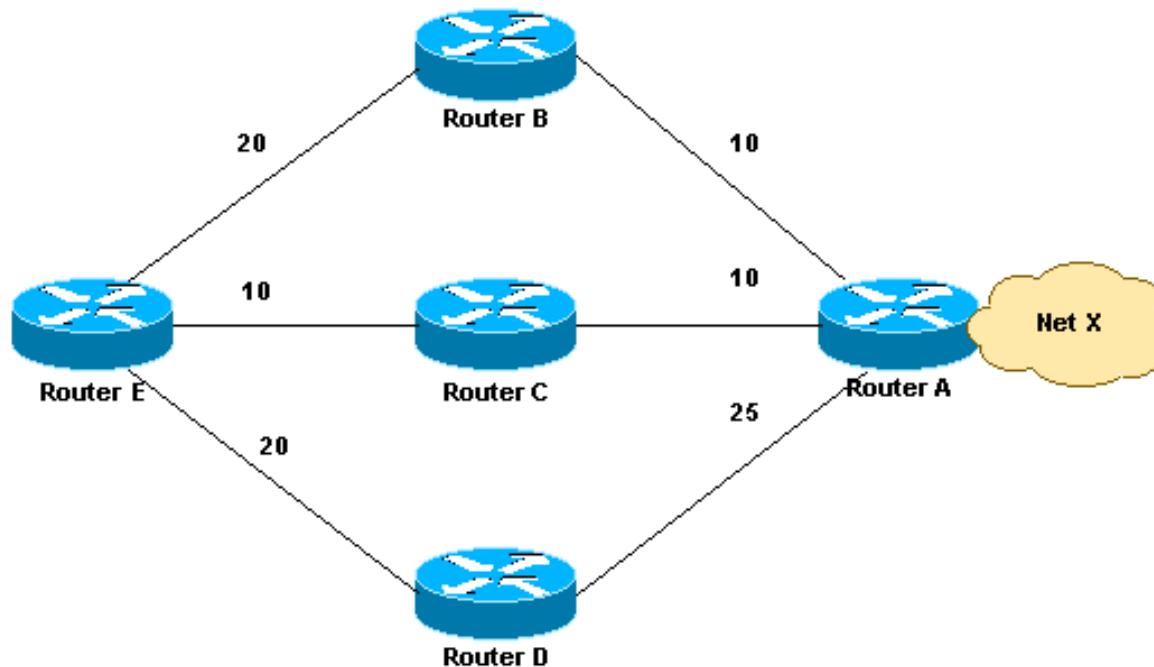
RouterB(config)#**router igrp 101**
RouterB(config-router)#**network 192.168.2.0**
RouterB(config-router)#**network 192.168.3.0**



Routing – EIGRP protocol

Load balance:

- Routes with equal costs (maximum-paths)
- Routes with different costs (variance)



E-B-A → 30
E-C-A → 20
E-D-A → 45

router eigrp 1
network *x.x.x.x*
variance 2

E-B-A → 30
E-C-A → 20

Routing – other protocols

OSPF

- Hierarchy with several areas
- Designated router (DR) and Backup Designated Router (BDR)
- Faster convergence after failure. Routers learn the whole network topology.

BGP

- Hierarchical organization
- Use “route reflectors” and confederations.

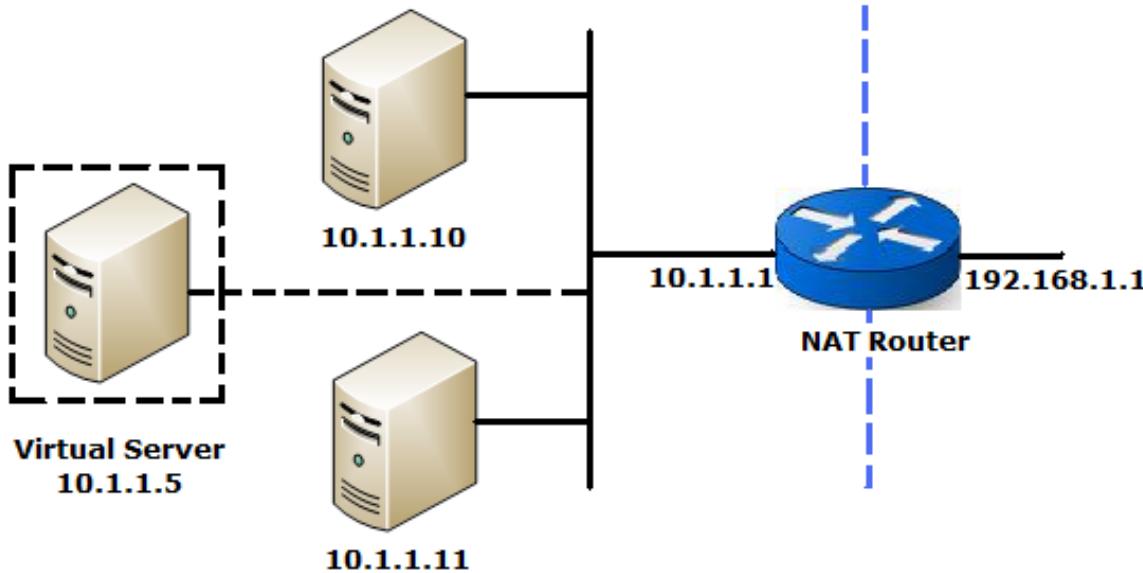
Routing – other protocols

- By default, router caches the routes previously used
- Load balancing → deactivate cache!

```
[no] ip route-cache [same-interface | flow | distributed | cef | policy]  
[no] ip cef
```

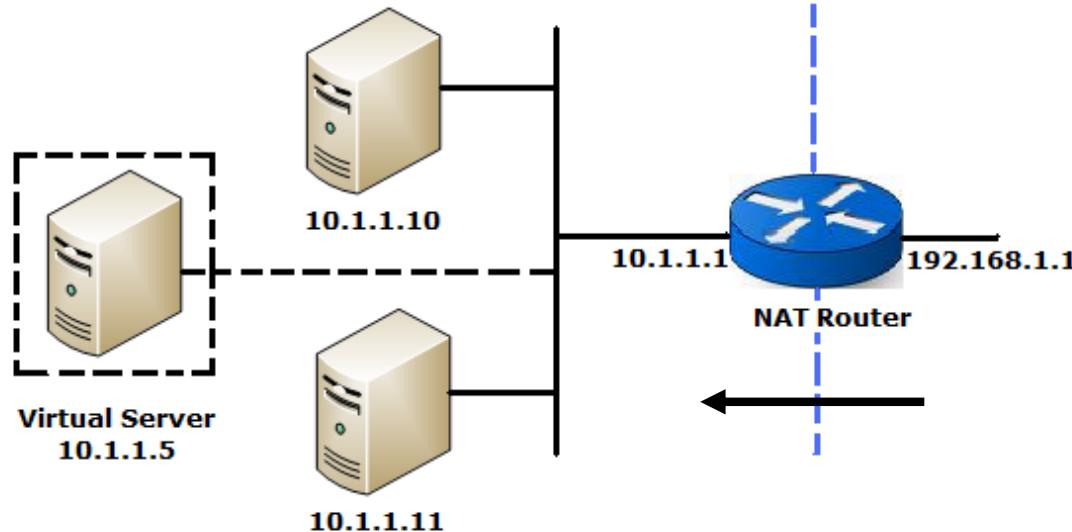
- fast switching
- Load balancing by destination

NAT with TCP Load Balancing



1. External requests are routed alternately between servers (p.e. web, telnet)
2. Assure balanced and rotative balancing between servers.

NAT with TCP Load Balancing



```
ip nat pool teste 10.1.1.10 10.1.1.11 prefix-length 24 type  
    rotary  
ip nat inside destination list ALL_TCP pool teste  
ip alias 10.1.1.5 23  
ip alias 10.1.1.5 80  
!
```

NAT with TCP Load Balancing

```
ip access-list extended ALL_TCP  
permit tcp any host 10.1.1.5 eq telnet  
permit tcp any host 10.1.1.5 eq www
```

```
ip nat inside  
ip nat outside
```

Alternative in Cisco IOS:

To explore Server Load Balancing (SLB) features.

First Hop Redundancy Protocols

- Hot Standby Redundancy Protocol (HSRP)
- Virtual Router Redundancy Protocol (VRRP)
- Global Load Balancing Protocol (GLBP)

Hot Standby Router Protocol (HSRP)



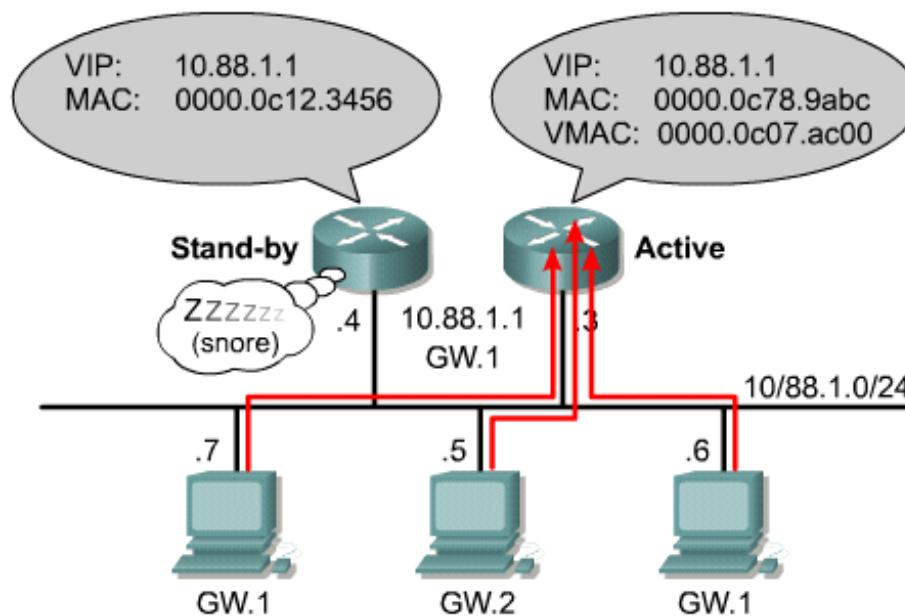
- Packets are redirected automatically to a *standby* router
- Transparently to the end-user
- Functioning:
 - A routers set share a MAC and IP addresses (Virtual Router)
 - One *router* is elected as active
 - Routers exchange control messages of HSRP state
 - ARP assignes MAC to the MAC Virtual
 - If active *router* fails, standby router functions as active router.

Hot Standby Router Protocol (HSRP)

HSRP Operation

FIGURE

1
2
3
4



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Hot Standby Router Protocol (HSRP)

Designating an Active Router

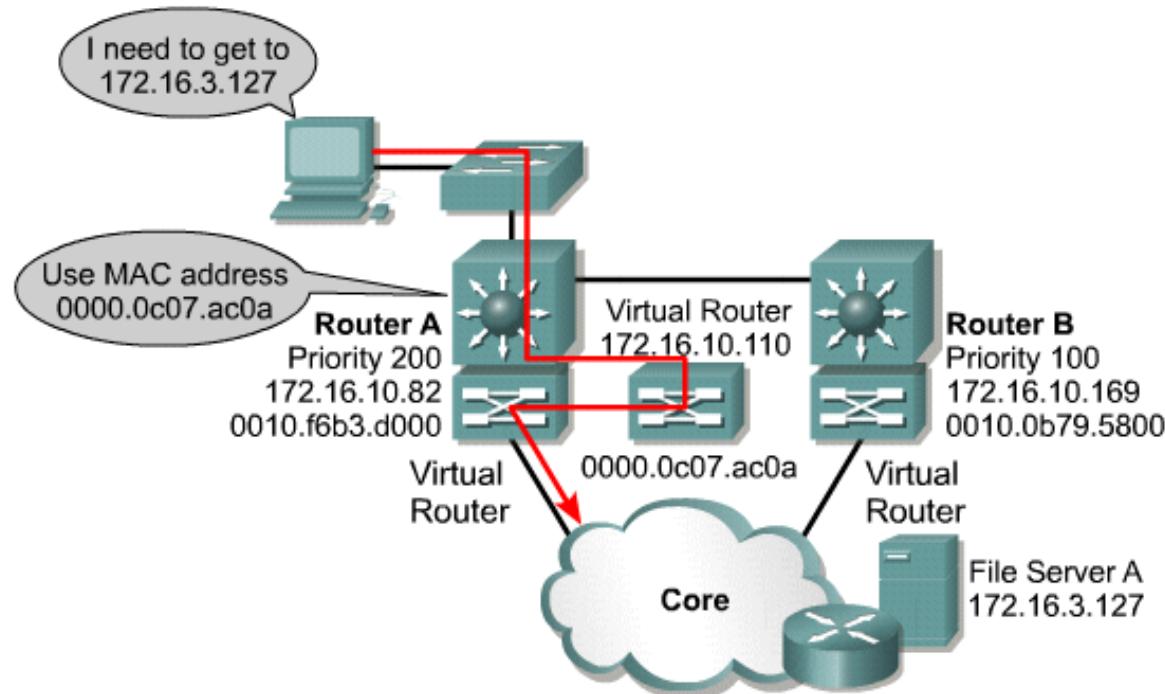
FIGURES

1

2

3

4



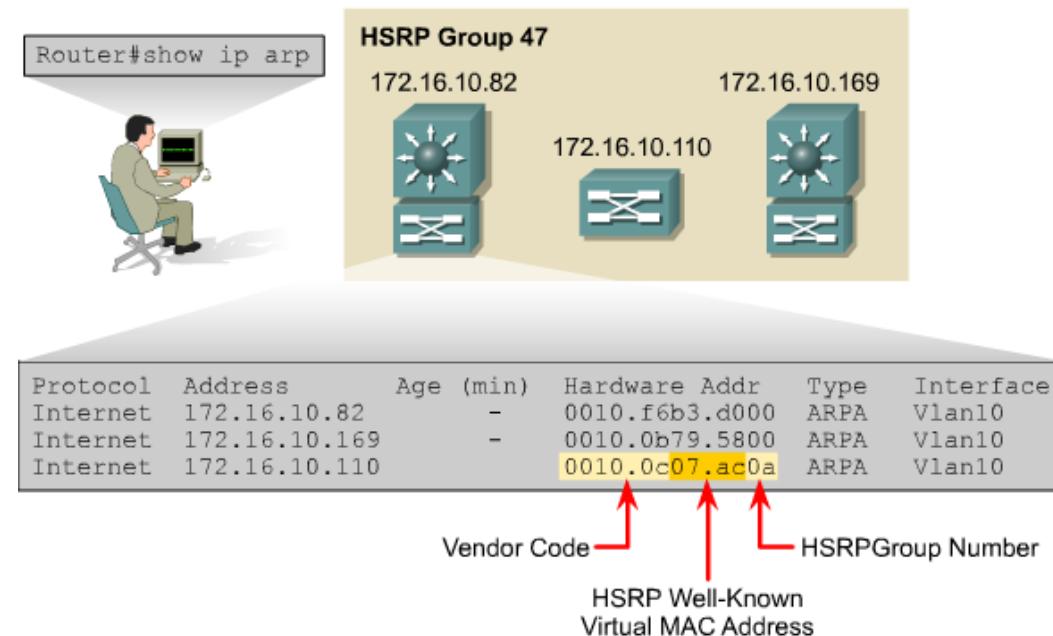
The router with the highest HSRP priority becomes the active router. The active router responds to ARP requests with the MAC address of the virtual router.

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Hot Standby Router Protocol (HSRP)

Virtual MAC address used at HSRP:

- Vendor ID – MAC address (24 bits)
- HSRP Code – 16 bits (“07.ac”)
- Group ID – last 8 bits of MAC endereço



Hot Standby Router Protocol (HSRP)

HSRP messages

1 Octet	1 Octet	1 Octet	1 Octet
Version	Op Code	State	HelloTime
Holdtime	Priority	Group	Reserved
Authentication Data			
Authentication Data			
Virtual IP Address			

Eleição de routers activo e standby

Initial
Stanby
Active

0 – Hello
1 – Coup
2 – Resign

[0 ... 255]

Intervalo de Hello

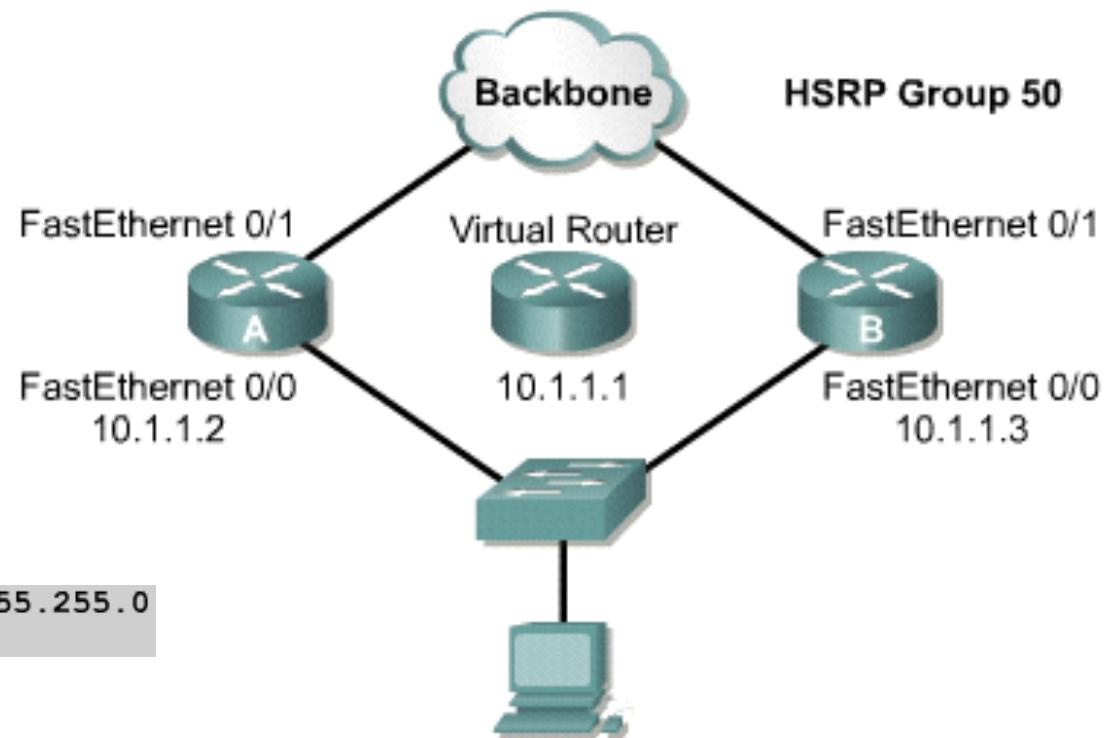
Transport protocol: UDP
Port: 1985
Destination IP : endereço 224.0.0.2
TTL = 1

Possible states:

- Initial, Learn, Listen, Speak, Stanby, Active

Hot Standby Router Protocol (HSRP)

Configuration



Router A

```
A(config-if)#ip address 10.1.1.2 255.255.255.0
A(config-if)#standby 50 ip 10.1.1.1

A(config-if)#standby 50 priority 150

A(config-if)#standby 50 preempt

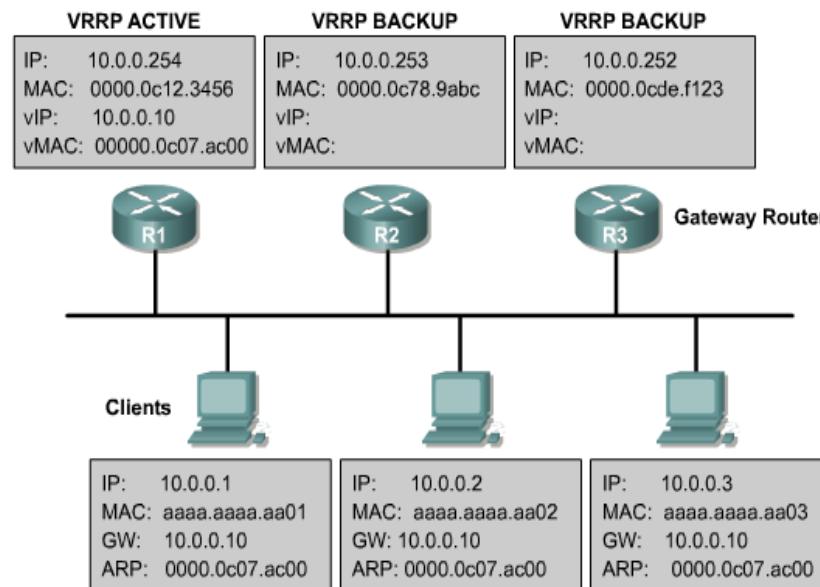
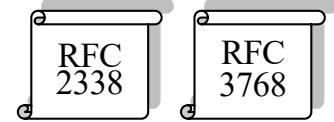
A(config-if)#standby 50 timers 5 15

A(config-if)#standby 50 track fastethernet 0/1 55
```

IP Address: 10.1.1.50
Default Gateway 10.1.1.1

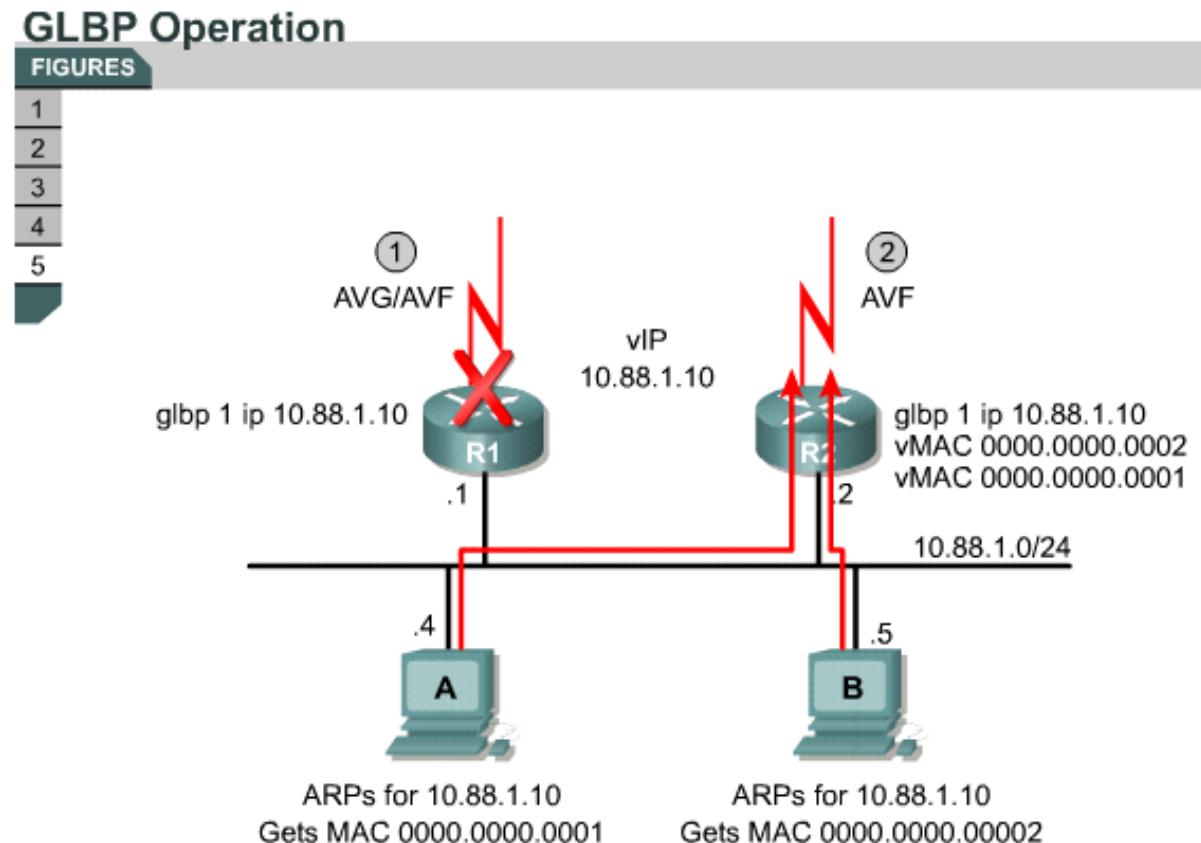
Virtual Router Redundancy Protocol (VRRP)

- Similar to HSRP
- Interoperability between equipments of distinct vendors
- A “Master” *router* and several “Backup” *routers*
- Periodic updates sent by Master *router*
- VRRP should be used when ot all routers are Cisco.



Gateway Load Balance Protocol (GLBP)

- Functions: Redundancy + Load Balance



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Bibliografia

- RFCs 2281, 2378 e 3768
- Mário Antunes; “*Gestão de Projetos TI e administração centralizada de sistemas e redes – cenários práticos em contexto empresarial*”; IPLeiria; 2010