Sistemas Operativos
fork, exec & exit

Patrício Domingues
ESTG/IPLeiria

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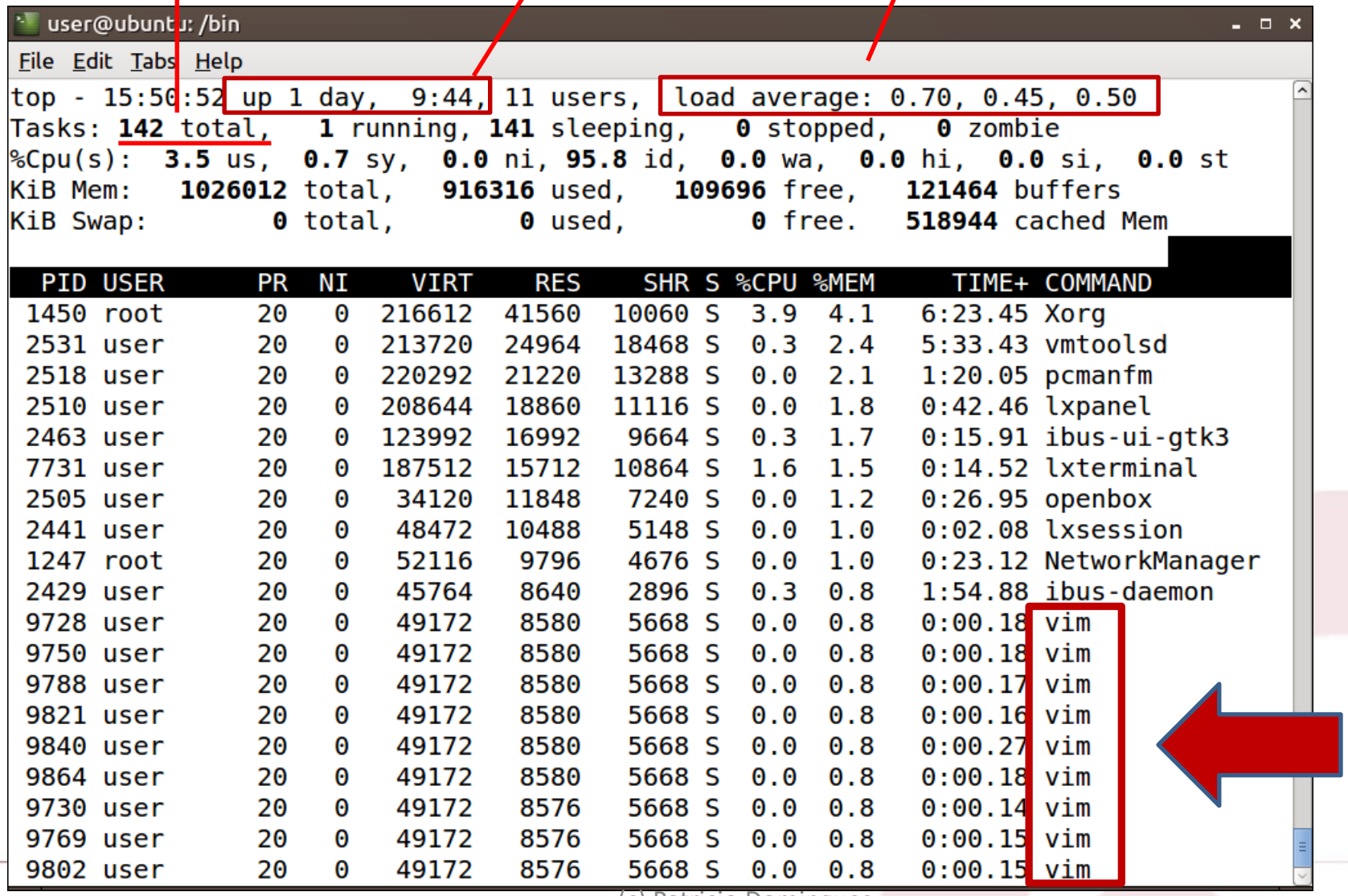
"top" / linux

142 processes

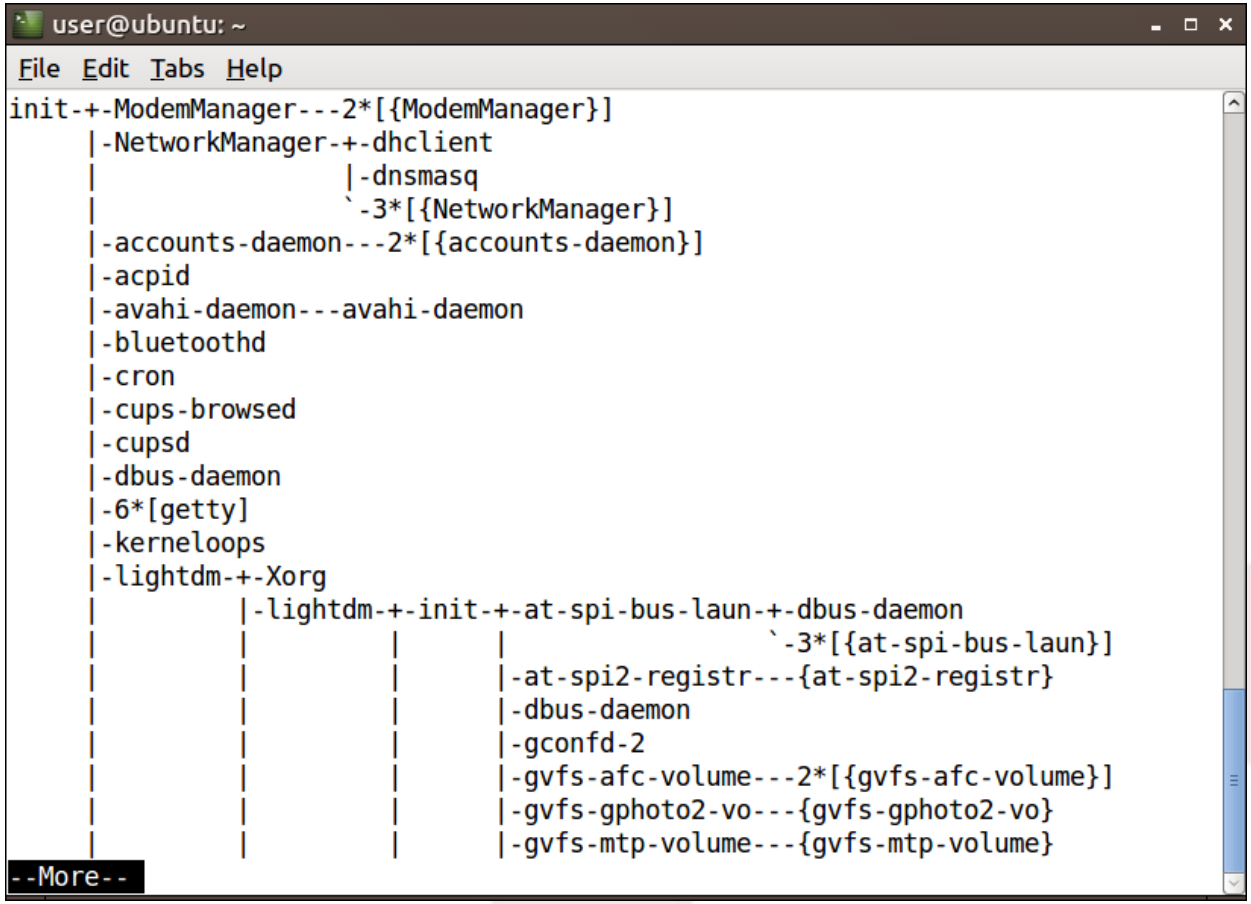
uptime

Average load of the system



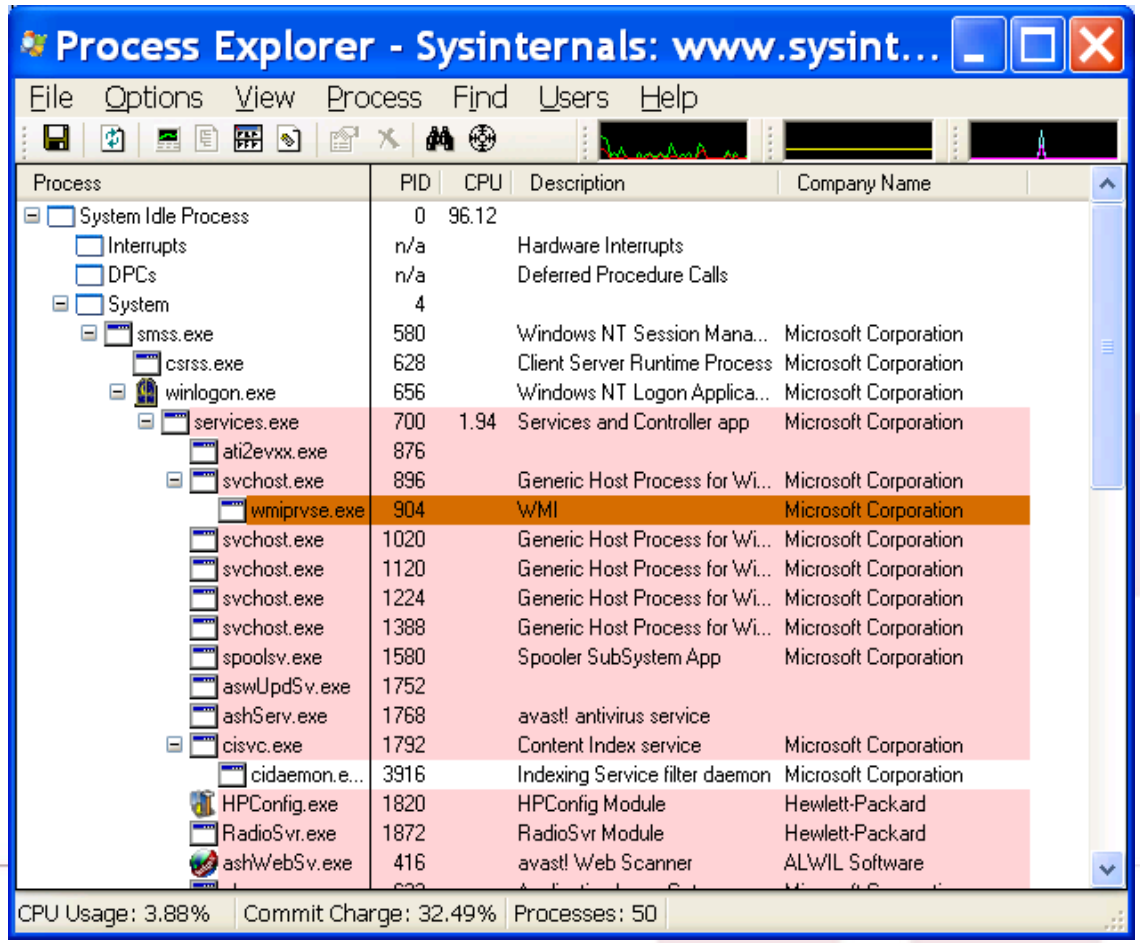
Processes (1)

- ✓ Processes in Unix
  - All are descendant from the *init* (PID=1) process
  - A process is created…by another process
    - Parent / son relationship
    - Tree of processes



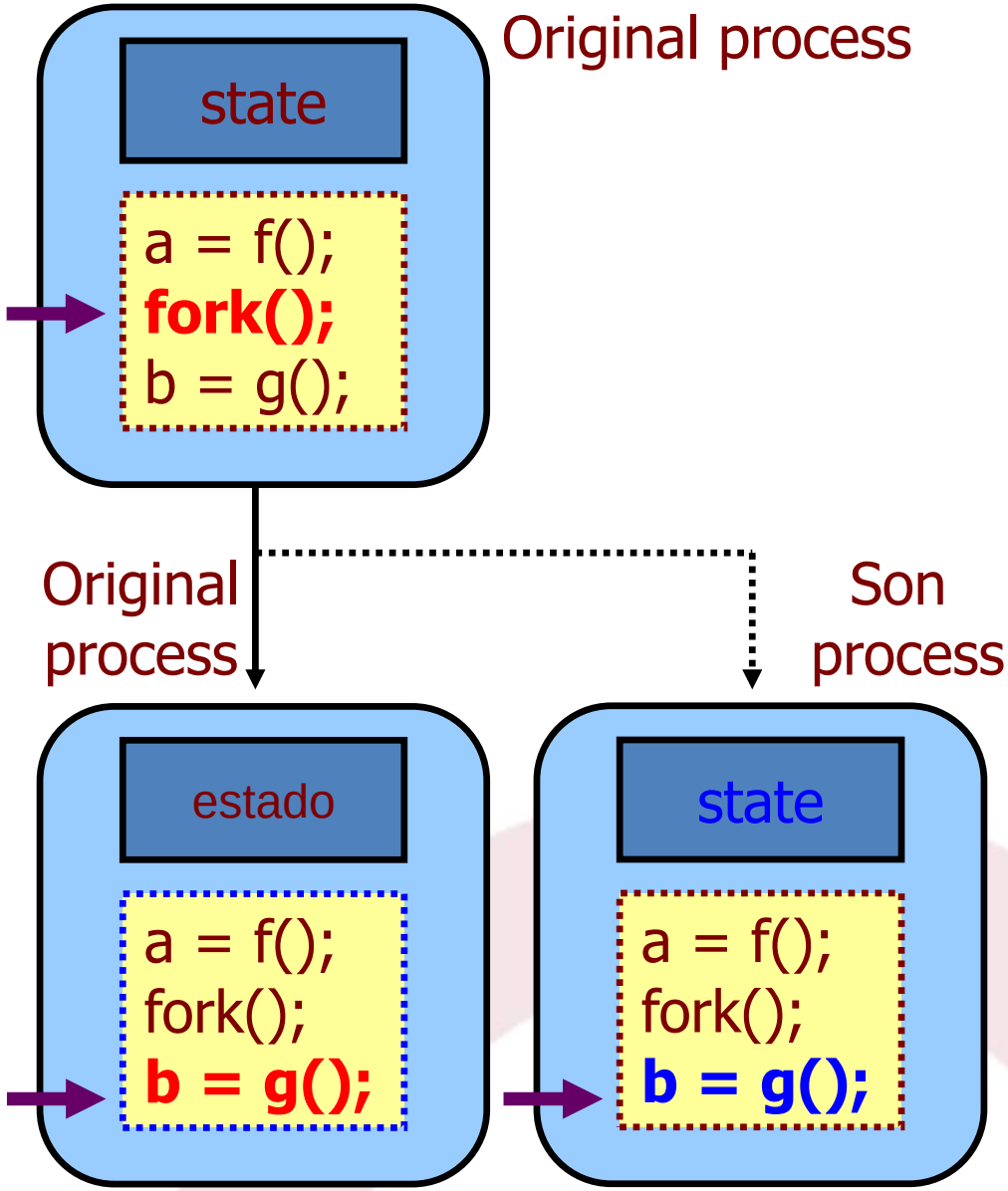
Processes (2)

- ✓There is also a process hierarchy in Microsoft Windows
  - "process explorer" (sysinternals.com)



Process model in UNIX

- `fork()`
  - system call to create a process
  - It is called by the parent process
- The son process inherits all characteristics of the parent process
  - Variables, program counter, open files, allocated memory, etc.
  - The son is a snapshot of the parent
- After "fork", each process executes separately
  - The change of a variable in one process does **not** reflect on the other one



`fork` system call

- ✓ `pid_t fork(void);`
- ✓ The fork system call returns an integer:
  - `0` to the newly created son process
  - `> 0` to the calling parent process
    - The return value corresponds to the PID of the newly created process
- ✓ It can also returns `-1` if an error has ocurred

Example – `fork` system call

```c
#include <stdio.h>
#include <unistd.h>
#include <sys/wait.h>
#include <sys/types.h>

int main()
{
  pid_t id;

  id = fork(); /* returns 0: son process; > 0 to the parent */
  if (id == 0)
    {  /* code only executed by the son process */
      printf("[%d] I'm the son!\n", getpid());
      printf("[%d] My parent is: %d\n", getpid(), getppid());
    }
  else if (id > 0 )
    {/* code only executed by the parent process */
      printf("[%d] I'm the father!\n", getpid());
      wait(NULL);
    }
  return 0;
}
```

```
user@ubuntu: ~
File  Edit  Tabs  Help
user@ubuntu:~$ ./fork.exe
[10567] I'm the father!
[10568] I'm the son!
[10567] My parent is: 10567
```

Example – `fork` in a for loop

- How many processes are created by the following code?

```
#include <…>
int main(void){
    int i;
    for(i=0;i<3;i++){
        if( fork() == 0 ){
            printf("PID=%u\n", getpid());
            fflush(stdout);
        }
    }
    return 0;
}
```

Only newly created processes print their PID
Answer: 7
2^n -1, with n=3

Process ID (PID) – #1

- ✓ Process ID (PID)
  - Integer identifier of a process
- ✓ The Linux kernel allocates process IDs to processes in a strictly linear fashion.
  - If pid 37 is the highest number currently allocated, pid 38 will be allocated next, even if the process last
- ✓ For compatibility with old UNIX, the max value for PID is 32768 (16-bit signed int)
- ✓ This value can be changed
  - `/proc/sys/kernel/pid_max`

Process ID (PID) - #2

- ✓ Within a C program, the PID of the calling process is returned with `getpid()`
  - `pid_t getpid(void);`
- ✓ The PID of the parent process is available through `getppid()`
  - `pid_t getppid(void);`

```
printf ("My pid=%jd\n", (intmax_t) getpid ());
printf ("Parent's pid=%jd\n", (intmax_t) getppid());
```

`printf` format specifiers (2)

- ✓Source: http://www.pixelbeat.org/programming/gcc/int_types/

```
uint32_t uint32=0xffffFFFF;
uintmax_t uintmax=UINTMAX_MAX;
off_t offset=TYPE_MAX(off_t); /* Depends on _FILE_OFFSET_BITS */
time_t time=TYPE_MAX(time_t); /* May be float! */
size_t size=TYPE_MAX(size_t); /* Depends on int size */

printf("native int bits %20zu %16x\n"
       "native long bits%20zu %16lx\n"
       "uint32_t max    %20"PRIu32" %16"PRIx32"\n"
       "uintmax_t max   %20ju %16jx\n" /* try PRIuMAX if %ju unsupported */
       "off_t max       %20jd %16jx\n" /* try PRIdMAX if %jd unsupported */
       "time_t max      %20jd %16jx\n"
       "size_t max      %20zu %16zx\n",
       sizeof(int)*CHAR_BIT, UINT_MAX,
       sizeof(long)*CHAR_BIT, ULONG_MAX,
       uint32, uint32,
       uintmax, uintmax,
       (intmax_t)offset, (intmax_t)offset,
       (intmax_t)time, (intmax_t)time,
       size, size);
```

Running applications

- ✓ But…
  - If all new processes execute the code of their parents, how can new applications be run?
    - The fork system call creates a clone of the parent process
- ✓ How do we run an application?
  - vim, ps, ls, find, firefox,...
- ✓ Answer
  - The “exec” family of system calls
    - These syscalls replace the image of the calling process

The *exec* family of system calls

- ✓ "exec" system calls

```
int execl(const char *path, const char *arg, ...);
int execlp(const char *file, const char *arg, ...);
int execle(const char *path, const char *arg, ..., char *const envp[]);
int execv(const char *path, char *const argv[]);
int execvp(const char *file, char *const argv[]);
```

- ✓ Usage of an exec system call is
  - exec...(application_to_be_run)
- ✓ "exec"
  - Replaces the image of the current process by another one from a given executable
    - Functions with "p" are "PATH"-aware
    - Functions with "v" get their parameters from a vector of strings
    - Functions with "l" get their parameters from a list, where itens are separated by "," and the list ends with NULL
- ✓ Example: **execl("/bin/ps", "ps", "aux",NULL);**

Example - running “ls” (1)

- ✓ Running “ls -a” resorting to “execlp”
- ✓ Question
  - ✓ Why the “This cannot happen!” message?

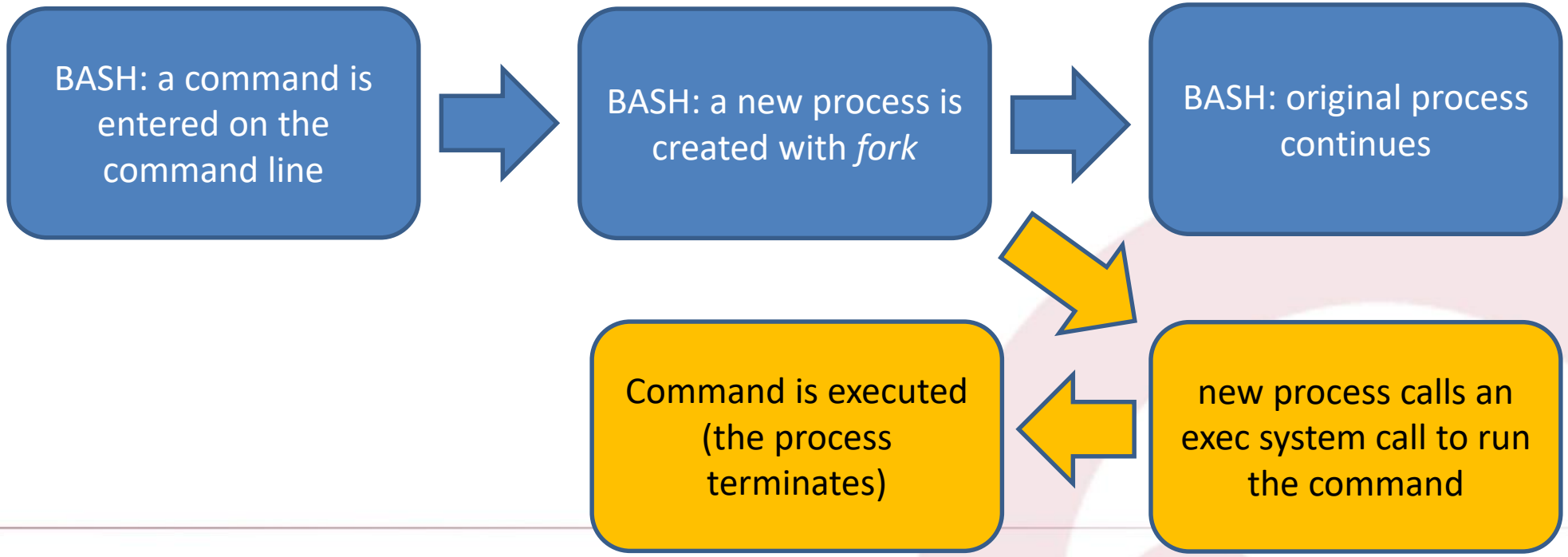
```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(){
  if (execlp("ls", "ls", "-a", NULL) == -1)
    perror("Error executing ls: ");
  else
    printf("This cannot happen!\n");
  return 0;
}
```

Example – running "ls" (2)

- ✓ Process launches "ls" via execlp
  - The exec system calls **NEVER** return when the execution is successful
  - The calling process image is replaced by the image of the executable called via "exec"
    - The calling process runs the executable
      - "ls" in our example
      - Therefore, *printf("This cannot happen!")* is removed from memory (as well as all the code of the calling process)
      - The code is replaced by the code of "ls -a"

Executing a command

- ✓ Executing a command
  - Example with bash
    - Applies to other shells (sh, zsh, etc.)
    - 1st – fork
    - 2nd – system call from the exec family



wait system call

- ✓ wait and waitpid

```
pid_t wait(int *status);
pid_t waitpid(pid_t pid, int *status, int options);
```

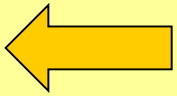
  - system calls used to synchronize a parent process with its children
  - Wait for state changes on their children processes
    - Child is stopped (SIGSTOP) or terminates
    - Child is resumed by a signal (SIGCONT)
  - Example
    - `wait(&status);` -- waits until a children process terminates
    - `waitpid(-1, &status, 0);` -- same as above

Zombie process

- ✓ A child that terminates, but has not been waited for becomes a "zombie“
- ✓ The kernel maintains a minimal set of information about the zombie process
  - PID, termination status, resource usage information
- ✓ As long as a zombie is not removed from the system via wait, it will consume a slot in the kernel process table
- ✓ If a parent process terminates, then its "zombie" children (if any) are adopted by *init*, which automatically performs a wait to remove the zombies

Creating zombies…

```c
#include <sys/types.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/wait.h>
void worker() {
  printf("[%d] Hi, I'm a worker process! Going to die...\n",
         getpid());
}
int main()
{ int i;
  for (i=0; i<5; i++) {
    if (fork() == 0) {
      worker();
      exit(0);
    }
  }
  system("ps aux | grep -i defunct");
  printf("[%d] Big father is sleeping!\n", getpid());
  sleep(10);
  return 0;
}
```



Question: how many processes are created?

Creating zombies…

✓ Results

```
user@ubuntu: ~/SO
File  Edit  Tabs  Help
user@ubuntu:~/SO$ ./zombie.exe
[17512] Hi, I'm a worker process! Going to terminate...
[17513] Hi, I'm a worker process! Going to terminate...
[17514] Hi, I'm a worker process! Going to terminate...
[17511] Hi, I'm a worker process! Going to terminate...
[17510] Hi, I'm a worker process! Going to terminate...
user     17467  0.2  0.4  12348  4616 pts/16   S+   12:59   0:00 vim zombie.c
user     17509  0.0  0.0   2024   276 pts/5    S+   13:02   0:00 ./zombie.exe
user     17510  0.0  0.0      0     0 pts/5    Z+   13:02   0:00 [zombie.exe] <defunct>
user     17511  0.0  0.0      0     0 pts/5    Z+   13:02   0:00 [zombie.exe] <defunct>
user     17512  0.0  0.0      0     0 pts/5    Z+   13:02   0:00 [zombie.exe] <defunct>
user     17513  0.0  0.0      0     0 pts/5    Z+   13:02   0:00 [zombie.exe] <defunct>
user     17514  0.0  0.0      0     0 pts/5    Z+   13:02   0:00 [zombie.exe] <defunct>
user     17515  0.0  0.0   2268   552 pts/5    S+   13:02   0:00 sh -c ps aux | grep -i zombie
user     17517  0.0  0.0   4680   832 pts/5    S+   13:02   0:00 grep -i zombie
[17509] Big father is sleeping!
user@ubuntu:~/SO$
```

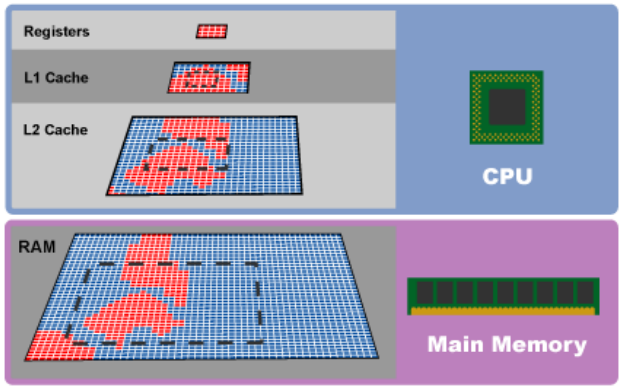
The `system` function

- ✓ In the zombie code, we have the following line of code

```
system("ps aux | grep –i zombie");
```

- ✓ system launches a shell that executes the command given as string
  - It executes **/bin/sh -c command_line** and waits for its termination
  - The Shell process is the one that actually executes the command line
  - It is costly, since it has to i) `fork` a process and then ii) `exec` its image to execute the Shell

CPU affinity (#1)



- ✓ On a multicore system, the OS scheduler needs to decide which processes runs on each CPU
- ✓ Once a process is running on a CPU, the OS scheduler tries to keep it there
  - Avoid the "cold cache" effect of moving a process to another CPU/core
    - When a process moves to another CPU/core, the cache(s) of the CPU/core do not have content of the process
      - The caches are *cold*

http://bit.ly/256cAlr

CPU affinity (#2)

- ✓ CPU affinity of a process can be controlled programatically
  - Hard affinity
- ✓ `int` **`sched_setaffinity`**`(pid_t pid, size_t setsize,const cpu_set_t *set);`
- ✓ `int` **`sched_getaffinity`**`(pid_t pid, size_t setsize,cpu_set_t *set);`
- ✓ void **CPU_SET** (unsigned long cpu, cpu_set_t *set);
- ✓ void **CPU_CLR** (unsigned long cpu, cpu_set_t *set);
- ✓ int **CPU_ISSET** (unsigned long cpu, cpu_set_t *set);
- ✓ void **CPU_ZERO** (cpu_set_t *set);

CPU affinity (#3)

✓Example

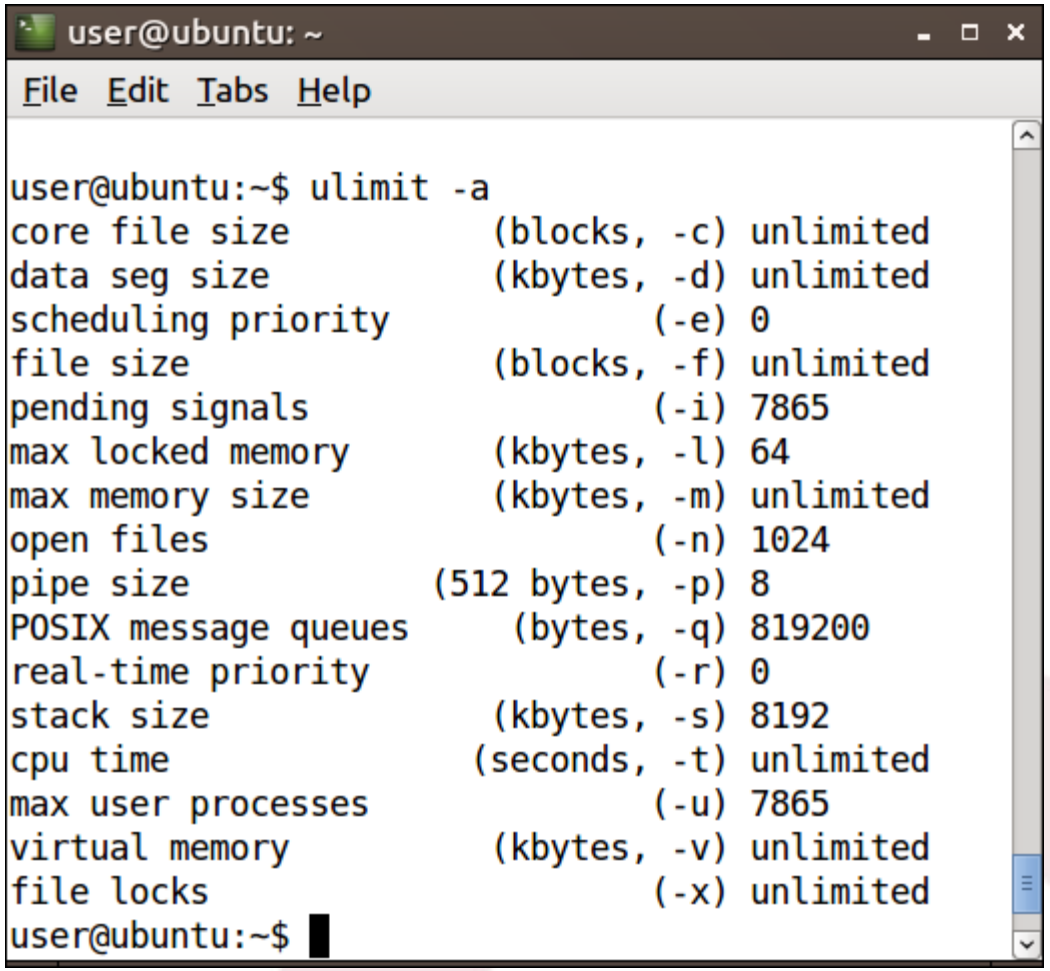
```
#define _GNU_SOURCE
#include <sched.h>
#include <stdio.h>cpu_set_t set;
int ret, i;
CPU_ZERO (&set);
ret = sched_getaffinity(0, sizeof (cpu_set_t), &set);
if (ret == -1){
        perror ("sched_getaffinity");
}
for (i=0; i < CPU_SETSIZE; i++) {
        int cpu;
        cpu = CPU_ISSET(i,&set);
        printf ("cpu=%i is %s\n", i, cpu?"set":"unset");
}
```

Termination of a process

- ✓ Reason for a process to terminate
  - Regular termination
    - exit, return of main function, etc.
  - Process has exceeded maximum CPU time (e.g., “ulimit” from bash)
  - Not enough memory
  - I/O failure
  - Invalid instruction (e.g., “divide by zero”)
  - OS action
    - Deadlock or OOM (Out of Memory Killer)
  - User action
    - Kill -9 PID or killall -9 process_name
  - …

ulimits (bash)

- ✓ The bash shell has an internal set of limits
  - `ulimit`
    - internal command
      - Processed by bash, there is no ulimit executable
    - There is no man for `ulimit`
      - help ulimit
  - `ulimit -a`
    - List all limits for the current session



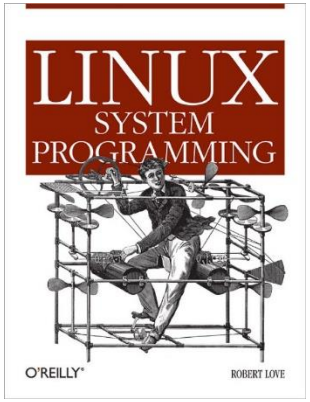
ulimits and core files

- ✓ ulimits has a parameter that controls the size of core file
  - `ulimit -c 99999999` ➔ unlimited
- ✓ With unlimited, the crash of an executable generates a core dump, that is, a file named "`core`"
- ✓ This core can be examined with a debugger, namely gdb
  - `gdb a.exe core`
  - `The program needs to be compiled with the right options`
    - `-g`
    - `-ggdb`

Bibliography

- Man pages
  - man 2 fork
  - man 2 exec
  - man system
  - man bash
  - help ulimit

- *Chapter 5 – Process management*, “Linux System Programming”, Robert Love, 2013



- printf format in C99 and C11 http://bit.ly/1OvGzGl