Advanced Basics of *Mathematica* Grammar **2011** (easy/medium)

Replacements

Very important equations arising in the study on quantum integrable models are the so called Y-system equations:

$$Y_j(u+1) Y_j(u-1) = (1+Y_{j-1}(u))(1+Y_{j+1}(u))$$
 for $j=1, ..., n$ and $Y_j(u)=0$ for $j \le 0$ or $j \ge n+1$.

By considering a few different n's convince yourself that they imply that

$$Y_i(u + n + 3) = Y_{n-j+1}(u)$$

Hint: Understand

$$f[5] /. f[a_] \Rightarrow f[a-1] + a /; a > 0$$

$$f[5] //. f[a_] \Rightarrow f[a-1] + a$$

$$f[5] //. f[a_] \Rightarrow f[a-1] + a /; a > 0$$

Solution

Conjugate

Consider some generic complex function such as

```
s = (Exp[Ikx] + Exp[2Iky])
```

Sometimes we might want *Mathematica* to conjugate such expressions assuming all variables to be real unless told otherwise. That is we would like a function **conj** such that

```
conj[s] s // FullSimplify
```

should yield

$$2 (1 + Cos[kx - 2ky])$$

Create such function. Hint: Run

```
FullForm /@ \{ Exp[Ikx], A+BI, -2I \}
```

Solution

■ Tensor Products

How to proceed to get the following nice implementation of the tensor product:

```
 \begin{pmatrix} \mathbf{a} & \mathbf{b} \\ \mathbf{c} & \mathbf{d} \end{pmatrix} \otimes \begin{pmatrix} \alpha & \beta & \gamma \\ \mathbf{e} & \phi & \rho \end{pmatrix} // \mathbf{MatrixForm} 
\begin{pmatrix} \mathbf{a} & \alpha & \mathbf{a} & \beta & \alpha & \beta & \beta & b & \gamma \\ \mathbf{a} & \alpha & \alpha & \beta & \alpha & \beta & b & \beta & b & \gamma \\ \mathbf{a} & \alpha & \alpha & \beta & \beta & b & \beta & b & \beta & b & \beta \\ \mathbf{c} & \alpha & \alpha & \beta & \alpha & \beta & b & \beta & b & \beta & b & \beta \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \beta & \alpha & \beta & \alpha & \beta & \alpha \\ \mathbf{c} & \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \beta & \alpha & \beta & \alpha \\ \mathbf{c} & \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \beta & \alpha & \beta & \alpha \\ \mathbf{c} & \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha & \beta & \alpha & \beta & \alpha \\ \mathbf{c} & \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha & \alpha & \beta & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha \\ \mathbf{c} & \alpha & \alpha & \alpha & \alpha
```

Solution

Simplify

FullSimplify measures the complexity of the expression, which it tries to reduce then, in some very special way. It uses the function **LeafCount**. In some cases this leads to a really strange behavior:

Run

```
FullSimplify[Sin[3 (u + \pi)] + Cos[5 (u + \pi)]]
```

and

```
FullSimplify[Sin[3 (u + \pi)] - Cos[5 (u + \pi)]]
```

Apply **TreeForm** to both functions (with and without FullSimplify). Do it also for what you would like the output of the first line to be. Same for **LeafCount**. *Mathematica* wants to minimize **LeafCount**. Why doesn't *Mathematica* simplify the first line *as we would?*

Try to find some other examples where FullSimplify "fails".

The LeadCount function can be replaced by a user defined function. You can invent your own measure of complexity which you like. Hive several examples of your own **ComplexityCount** function and use it to simplify the expression you found.

Create your own function **Simpy** which simplifies expressions with 3π etc inside trigonometric functions in a more human way.

Hint: Functions which you might find useful are: TrigToExp, ExpandAll, ...

Solution

Sqrt killer

Define a function which will bring the expression with a single square roots in denominator to the canonical form, that is

$$\frac{\cdots + \sqrt{\cdots}}{(\cdots + \sqrt{\cdots}) \cdots (\cdots + \sqrt{\cdots})} \rightarrow \cdots + \cdots \sqrt{\cdots}$$

Solution

Memorizer

Define a function which will solve

$$F_n = F_{n-1} + 1 / F_{n-2}$$
 with $F_1 = F_2 = 1.2$

Plot the sequence $log(F_n)$ with n = 1, ..., 1000 (should take less than one second)

Solution

■ Puzzles*

*(from http://richardwiseman.wordpress.com/, solutions proposed here taken from the comment box and were given by Simon)

Puzzle 1: How can you place the arithmetical signs '+' and '-' between the consecutive numbers 123456789 so that the end result is 100?

Try to come up with a code to solve this problem. Also, decode the sollowing solution*

```
Do[If[ToExpression[str = StringJoin[Riffle[Characters["123456789"], i]]] == 100,
   Print["100=", str]], {i, Tuples[{"+", "-", ""}, 8]}]
```

Puzzle 2: Yesterday I went shopping and picked up four items. When I got to the till, the cashier added the price of the four items together and the bill was £7.11. I then noticed that I would get exactly the same total if I were to multiply the four prices. How much did each of the items cost?

Note: What the puzzle should say is: "Find integers a,b,c,d such that a+b+c+d=711 and a b c d = 711 000 000."

Again, try to come up with a code to solve this problem and/or decode the solution*

```
Select[IntegerPartitions[711, {4}], Times@@# == 711000000 &]
```

Solution given in the text