Middleware and Web Services Lecture 6: High Availability and Performance

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Good Performance

- What influences good performance?
 - Number of users and concurrent connections
 - Number of messages and messages' sizes
 - Number of services
 - Infrastructure capacity, availability, configuration, ...
- How can we achieve good performance?
 - Infrastructure
 - → Scalability, failover, cluster architectures
 - Performance tuning
 - → Application Server, JVM memory, OS-level tuning, Work managers configuration
 - Service configuration
 - → Parallel processing, process optimization

Overview

- Infrastructure
 - Load Balancers
 - Cluster Architecture
- Performance Tuning

Definitions

Scalability

- server scalability
 - \rightarrow ability of a system to scale when input load changes
 - → users should not feel a difference when more users access the same application at the same time
 - → horizontal scaling
 - → adding new instances of applications/servers
 - → vertical scaling
 - → adding new resources (CPU, memory) to a server instance
- network traffic
 - → bandwidth capacity influences performance too
 - → service should limit the network traffic through caching

Availability

- probability that a service is operational at a particular time
 - \rightarrow e.g., 99.9987% availability downtime ~44 seconds/year

Definitions (Cont.)

High Availability

- When a server instance fails, operation of the application can continue
- Failures should affect application availability and performance as little as possible

• Application Failover

- When an application component performing a job becomes unavailable, a copy of the failed object finishes the job.
- Issues
 - \rightarrow A copy of the failed object must be available
 - \rightarrow A location and operational status of available objects must be available
 - \rightarrow A processing state must be replicated

Load Balancing

- Distribution of incoming requests across server instances

Performance Metrics

- Latency
 - A client-side metric

- CPU intensive service or a bad configuration of a service
 - → consider asynchronous processing when CPU intensive
- Writing to a data store

Performance Metrics

- Queries/Requests per Second (QPS)
 - A server-side metric

- Caching may improve performance
 - → even if data changes often, with high QPS caching improves a lot

Infrastructure Example – Weblogic

Best Configuration Practices

Domain configuration

- A server is an admin server or a managed server
- Each server is running on a separated JVM
- A physical machine may run one or more servers
- There should be at least two managed servers running on one machine
 - → This provides a better performance (as opposed to one server running on one machine)
- A domain can have clustered or unclustered servers

• Load balancers (LB)

- Load Balancers are not part of the domain
 - → They are external to Weblogic server
 - → There is usually one HW LB and several SW LBs
 - \rightarrow Software LB
 - → Realized by the Web Tier (Apache HTTP server)
 - → Redirects requests too all managed servers in a domain (across multiple machines)

Overview

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Load Balancing

- Distributes a load to multiple app/object instances
 - App instances run on different machines
 - Load sharing: equal or with preferences
 - Health checks
- Types
 - DNS-based load balancer
 - → DNS Round Robin
 - NAT-based load balancer (Layer-4)
 - Reverse-proxy load balancer (Layer-7)
 - \rightarrow application layer
 - → Sticky sessions
 - → JSession, JSession-aware load balancer
 - Client-side load balancer
 - \rightarrow LB run by a client
 - → a client uses a replica-aware stub of the object from the server

DNS-based Load Balancer

DNS Round Robin

- A DNS record has multiple assigned IP addresses
- DNS system delivers different IP addresses from the list
- Example DNS A Record: company.com A 147.32.100.71 147.32.100.72 147.32.100.73
- Advantages
 - Very simple, easy to implement
- Disadvantages
 - IP address in cache, could take hours to re-assign
 - No information about servers' loads and health

Reverse Proxy Load Balancer

HTTP Sticky Sessions Example

- How to identify a server that hosts the session state
 - Passive cookie persistence LB uses a cookie from the server
 - Active cookie persistence LB adds its own cookie

Types of Load Balancers

- Software
 - Apache mod_proxy_balancer
 - → HTTP Session persistence sticky sessions
 - WebLogic proxy plug-in

/soa-infra is a first part of an URL path that rules in this Location will be applied (this is a standard Apache configuration mechanism) czfmwapp{N} is a hostname that corresponds to a virtual IP to which the managed server JVM processes is bounded (using the tcp port 8001).

WebLogicCluster specifies the list of servers for load balancing

- Hardware
 - Cisco, Avaya, Barracude

Round-Robin Algorithm

Uses

```
request — client request with or without a cookie information

server_list — a list of servers that can process the request

rbinx — round robin index

sticky_sessions — associative array of pairs <session_id, server>

unhealthy_treshhold — a number of negative consecutive health checks before moving the server to the "unhealthy" state.
```

Round Robin Algorithm

- $-\mathit{if}$ session_id exist in the request and in sticky_sessions
 - → send the request to the server sticky_sessions[session_id]
- otherwise
 - → send the request to the rbinx server in the server_list
 - → extract session_id from the response from the server
 - \rightarrow if the session_id exist, add a pair <session_id;server_list[rbinx]> to sticky_sessions
 - → increase rbinx by one or reset it to 0 if it exceeds the length of server_list

Health Check

- Health Check
 - For each server in the server_list
 - → call the server's heatlhcheck endpoint
 - \rightarrow if a number of failed health checks for the server exceeds the unhealthy_threshold
 - → remove the server from the server_list
 - \rightarrow if the server was unhealthy and a there was a successful healthcheck
 - → add the server back to the server_list

Session State Persistence and Replication

- Session persistence
 - Session information is maintained in the database
 - Does not require sticky sessions
 - Implements HttpSession interface that writes data to the DB
- In-memory replication
 - A primary server holds a session state, the secondary server holds its replica.
 - Information about primary and secondary servers are part of JSession

In-Memory Replication

Session format

- It's a cookie
- JSESSIONID=SESSION_ID!PRIM_SERVER_ID!SEC_SERVER_ID!CREATION_TIME

 SESSION_ID session id, generated by the server to identify memory

 associated with the session on the server

 PRIM_SERVER_ID ID of the managed server holding the session data

 SEC_SERVER_ID ID of the managed server holding the session replica

 CREATION_TIME time the session data was created/updated

How LB uses this information

- LB has information whether the server is running or not (via healthchecks)
- if the primary server is running, it redirects the request there
- if the primary server is not running, it redirects the request to the secondary server directly
- if primary and secondary servers are not running, it redirect the request to any other server it has in the list this may cause side effects!

In-Memory Replication Scenarios

Overview

- Infrastructure
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Overview

• Cluster capabilities

- A group of servers (aka cluster members) act together to serve clients' requests
- Cluster is transparent to clients
- Servers can run on the same machines or on different machines
- Cluster's capacity can be increased by adding servers to the cluster
- Servers in a cluster may have the same copy of objects and they are aware of each other objects
 - → objects: applications, JMS destinations, RMI objects
 - → See Cluster-wide JNDI tree in Lecture 4

• Cluster Messaging Protocols

- When servers need to send messages to other members of the cluster
- *Unicast* one-to-one communication using TCP/IP sockets
- Multicast one-to-many communication
- Cluster services that rely on the cluster messaging protocol
 - → Cluster Membership
 - \rightarrow JNDI Replication
- There are services using persistent RMI connections between cluster members
 - → such as during in-memory replication of HTTP session information

Unicast

Overview

- Uses standard TCP/IP sockets to send messages across cluster members
- Uses a **group leader strategy** to limit the number of sockets required to send messages within the cluster

• Groups in the cluster

- Cluster is split into M groups with N servers, each group has a leader
- Servers communicate with the group leader, group leaders communicate with servers in the group and other group leaders
- When a group leader dies, the group elects another group leader
- There is up to MxN network messages for every cluster message

Cluster Services

Cluster Membership

- Cluster members maintain their view on what servers are in the cluster
- Each server sends heartbeat messages to the cluster
 - \rightarrow others know it is alive
- Each server receives heartbeat messages
 - → When a server misses a number of messages from a cluster, it removes the server from the list until it receives the next heartbeat message.

• JNDI Replication

- Provides each server with a cluster-wide view of the JNDI tree.
- Servers send JNDI update messages to the cluster when an object is bound or removed to their local tree.
- When a server leaves the cluster
 - \rightarrow other members remove its JNDI bindings from their three.
- When a server joins the cluster
 - \rightarrow The server asks other server for a JNDI state dump.

Overview

- Infrastructure
- Performance Tuning

Performance Limiting Factors

Monitoring

- Important to understand performance
 - DevOps monitoring trends
- What you need
 - Collect \rightarrow Filter \rightarrow Store \rightarrow View \rightarrow Tune
 - Metrics, dashboards, alerting, log management, reporting, tracing capabilities
 - It is necessary to organize metrics well in order to understand what is going on
 - Start from a high-level process, detail to technical components
- Source
 - Application server
 - → usually management beans with JMX interfaces
 - \rightarrow log files (access logs, server logs, etc.)
 - OS
 - \rightarrow many utilities available out of the box
 - → open sockets, memory, context switches, I/O performance, CPU usage
 - Database
 - \rightarrow applications may write metrics to the DB
 - \rightarrow SQL scripts to collect metrics

Monitoring Tools

- Commercial Monitoring Solutions
 - Application server vendor usually offers a monitoring solution
 - AppDynamics, Oracle Enterprise Manager, Splunk
 - Google stackdriver, Amazon AWS CloudWatch
- Open source examples
 - Elasticsearch + LogStash + Kibana
 - -InfluxDB + Telegraph + DataGraph

Tuning – A Layered Approach

- Application server can be tuned at multiple layers
 - Service configuration optimization
 - Transport-level tuning
 - Application Server Tuning
 - JVM Memory Tuning
 - OS Tuning
- Lower levels are cheaper to tune

Memory Allocations

Drawing not available

Generations

- Young objects get allocated in this space initially
- − Old − objects get promoted to old from young
- Perm space for permanent allocations, e.g. objects describing classes and methods

Garbage Collection

- Steps to move objects around
 - 1. Objects are created in young
 - 2. When young is full, the live objects are copied to old, dead are discarded
 - lightweight GC
 - 3. When young is full and no space in old \rightarrow the full GC frees the old space
 - Full GC nothing is running in JVM, the application stops
 - Too frequent full GC has an impact on performance
- A memory leak or inadequate heap allocation
 - Old is out of space \rightarrow full GC will run often (or continously)
 - High CPU utilization, The server will not be able to process/respond to requests

Heap Size and GC Runs

- Heap Size and GC runs
 - Wrong heap size allocation too small or memory leaks
 - GC full runs too often, this has a negative impact on performance

JVM Memory Tuning

- JVM Memory Parameters
 - -Xms initial java heap size
 - -Xmx maximum java heap size
 - -XX:NewSize the initial size of the heap for young generation
 - -XX: MaxNewSize the maximum size of the heap for young generation
- General recommendations
 - -Xms and -Xmx should be set to the same value (do not allow the heap to grow \rightarrow limit the overhead)
 - -XX:NewSize and -XX:MaxNewSize should be set to the one half of maximum heap
 - Example, 1GB heap size
 - -Xms1024m -Xmx1024m -XX:NewSize=500m -XX:MaxNewSize=500m

Asynchronous I/O: Recall

- Connections maintained by the OS, not the Web app
 - The Web app registers events, OS triggers events when occur

- Characteristics
 - Event examples: new connection, read, write, closed
 - The app may create working threads, but controls the number!
 - → much less number of working threads as opposed to blocking I/O

Work Manager Configuration

Work Manager

- Controls the number of thread allocated to processing of requests
- In WLS is called a dispatch policy
 - → Can be assigned to OSB proxy services
- Parameters
 - → *maximum threads* (max) *maximum number of working threads*
 - \rightarrow *capacity* (cap) *maximum number of connections*
- maximum connections waiting to be processed: cap max
- refused connections: when number of connections is > cap

Inbound throtling

- A dispatch policy applied to a single proxy service
- Rejected connections will not be processed