# teambrbr002 UFMG

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## Índice

### 1 Grafos

#### 1.1 AGM Directionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
//
// O(m log(n))
3c9 struct node {
        pair < ll, int > val;
        ll lazy;
4e4
b19
        node *1, *r;
f93
        node() {}
c53
        node(pair<int, int> v) : val(v), lazy(0), l(NULL), r(NULL) {}
        void prop() {
a9c
768
             val.first += lazy;
b87
            if (1) 1->lazy += lazy;
d3b
            if (r) r \rightarrow lazy += lazy;
c60
            lazv = 0:
05b
296 }:
de5 void merge(node*& a, node* b) {
        if (!a) swap(a, b);
c11
802
        if (!b) return;
626
        a->prop(), b->prop();
d04
        if (a->val > b->val) swap(a, b);
```

```
4b0
        merge(rand()%2 ? a->1 : a->r, b);
b82 }
d01 pair <11, int > pop(node *& R) {
        R->prop();
22 e
        auto ret = R->val;
        node* tmp = R;
af0
3f3
        merge(R->1, R->r);
6 c 9
        R = R - > 1;
3 e 4
        if (R) R->lazy -= ret.first;
7 c 3
        delete tmp;
edf
        return ret;
c4e }
6f6 void apaga(node * R) { if (R) apaga(R->1), apaga(R->r), delete R; }
f13 ll dmst(int n, int r, vector<pair<int, int>, int>, int>>& ar) {
94 e
        vector<int> p(n); iota(p.begin(), p.end(), 0);
a23
        function < int(int) > find = [&](int k) { return
    p[k] == k?k:p[k] = find(p[k]); };
2d7
        vector < node *> h(n);
56f
        for (auto e : ar) merge(h[e.first.second], new node({e.second,
    e.first.first}));
fd1
        vector < int > pai(n, -1), path(n);
        pai[r] = r;
66 e
        11 \text{ ans} = 0;
04b
603
        for (int i = 0: i < n: i++) { // vai conectando todo mundo
2a3
             int u = i, at = 0;
             while (pai[u] == -1) {
cae
                 if (!h[u]) { // nao tem
daa
947
                     for (auto i : h) apaga(i);
77 c
                     return LINF:
dd1
167
                 path[at++] = u, pai[u] = i;
55 e
                 auto [mi, v] = pop(h[u]);
```

```
64c
                ans += mi;
                if (pai[u = find(v)] == i) { // ciclo
5 e 2
                     while (find(v = path[--at]) != u)
86f
621
                         merge(h[u], h[v]), h[v] = NULL, p[find(v)] = u;
                     pai[u] = -1;
57a
                }
0 d8
            }
ce8
5 df
947
        for (auto i : h) apaga(i);
ba7
        return ans;
e02 }
```

### 1.2 Articulation Points

```
// Computa os pontos de articulação (vertices criticos) de um grafo
// art[i] armazena o numero de novas componentes criadas ao deletar
// se art[i] >= 1. entao vertice i eh ponto de articulação
// O(n+m)
1a8 int n;
789 vector < vector < int >> g;
4ce stack < int > s:
b66 vector < int > id, art;
3e1 int dfs_art(int i, int& t, int p = -1) {
cf0
        int lo = id[i] = t++;
18e
        s.push(i);
        for (int j : g[i]) if (j != p) {
9 a 3
            if (id[j] == -1) {
206
                int val = dfs_art(j, t, i);
                lo = min(lo, val);
0c3
                if (val >= id[i]) {
588
66a
                    art[i]++;
                    while (s.top() != j) s.pop();
bd9
2eb
                    s.pop();
1 f 3
                // if (val > id[i]) aresta i-j eh ponte
238
328
            else lo = min(lo, id[j]);
762
3bd
        if (p == -1 and art[i]) art[i]--;
```

```
253     return lo;
8e1 }

d79 void compute_art_points() {
597     id = vector < int > (n, -1);
a62     art = vector < int > (n, 0);
6bb     int t = 0;
d41     for (int i = 0; i < n; i++) if (id[i] == -1)
625          dfs_art(i, t, -1);
379 }</pre>
```

### 1.3 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
14e int n, m;
248 int d[MAX];
e93 vector<pair<int, int>> ar; // vetor de arestas
9e2 vector < int > w;
                                 // peso das arestas
6be bool bellman ford(int a) {
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
8a8
        d[a] = 0:
4e3
        for (int i = 0; i \le n; i++)
891
            for (int j = 0; j < m; j++) {
6 e 4
                if (d[ar[i].second] > d[ar[i].first] + w[i]) {
705
                    if (i == n) return 1;
e 93
                     d[ar[j].second] = d[ar[j].first] + w[j];
84b
                }
            }
a 82
bb3
        return 0;
6eb }
```

### 1.4 Block-Cut Tree

```
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloracao da arvore eh tal que uma cor sao
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
//
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
//
// Para todo i < blocks.size()</pre>
// blocks[i] eh uma componente 2-vertice-conexa maximal
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
// pos[i] responde a qual vertice da arvore vertice i pertence
// Arvore tem no maximo 2n vertices
//
// O(n+m)
d10 struct block_cut_tree {
        vector < vector < int >> g, blocks, tree;
d8e
43b
        vector < vector < pair < int , int >>> edgblocks;
4ce
        stack<int> s:
6 c 0
        stack<pair<int, int>> s2;
2bb
        vector < int > id, art, pos;
763
        block_cut_tree(vector < vector < int >> g_) : g(g_) {
af 1
            int n = g.size();
37a
            id.resize(n, -1), art.resize(n), pos.resize(n);
6f2
            build();
        }
6bd
        int dfs(int i, int& t, int p = -1) {
df6
cf0
            int lo = id[i] = t++;
18e
            s.push(i);
827
            if (p != -1) s2.emplace(i, p);
53f
            for (int j : g[i]) if (j != p and id[j] != -1)
   s2.emplace(i, j);
            for (int j : g[i]) if (j != p) {
cac
                if (id[j] == -1) {
9a3
121
                    int val = dfs(j, t, i);
                    lo = min(lo, val);
0 c3
```

```
588
                    if (val >= id[i]) {
                         art[i]++:
66a
483
                         blocks.emplace_back(1, i);
                         while (blocks.back().back() != j)
110
138
                             blocks.back().push_back(s.top()), s.pop();
128
                         edgblocks.emplace_back(1, s2.top()), s2.pop();
47 e
                         while (edgblocks.back().back() != pair(j, i))
bсе
                             edgblocks.back().push_back(s2.top()),
   s2.pop();
870
                    // if (val > id[i]) aresta i-i eh ponte
85 с
328
                else lo = min(lo, id[j]);
            }
344
3bd
            if (p == -1 and art[i]) art[i]--;
253
            return lo:
726
        }
        void build() {
0a8
6bb
            int t = 0:
            for (int i = 0; i < g.size(); i++) if (id[i] == -1) dfs(i,
abf
   t, -1);
56 c
            tree.resize(blocks.size());
            for (int i = 0; i < g.size(); i++) if (art[i])</pre>
f7d
                pos[i] = tree.size(), tree.emplace_back();
965
973
            for (int i = 0; i < blocks.size(); i++) for (int j :</pre>
   blocks[i]) {
403
                if (!art[j]) pos[j] = i;
                else tree[i].push_back(pos[j]),
101
   tree[pos[j]].push_back(i);
3df
        }
c03
056 };
1.5 Blossom
// Matching maximo em grafo geral
//
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
```

```
042 vector < int > g[MAX];
128 int match[MAX]; // match[i] = com quem i esta matchzado ou -1
1f1 int n, pai[MAX], base[MAX], vis[MAX];
26a queue < int > q;
107 void contract(int u, int v, bool first = 1) {
165
        static vector < bool > bloss;
fbe
        static int 1:
418
        if (first) {
            bloss = vector <bool > (n. 0);
a 47
042
            vector < bool > teve(n, 0);
ddf
            int k = u: l = v:
31e
            while (1) {
297
                teve[k = base[k]] = 1;
                if (match[k] == -1) break;
116
dfa
                k = pai[match[k]];
68b
            while (!teve[l = base[1]]) l = pai[match[1]];
d31
5 d 6
2 e 9
        while (base[u] != 1) {
            bloss[base[u]] = bloss[base[match[u]]] = 1;
e29
8fa
            pai[u] = v:
            v = match[u]:
0b0
a51
            u = pai[match[u]];
58e
71c
        if (!first) return:
        contract(v, u, 0):
95e
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
6ee
594
            base[i] = 1;
ca7
            if (!vis[i]) q.push(i);
            vis[i] = 1:
29a
857
        }
e35 }
f10 int getpath(int s) {
88f
        for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1, vis[i] =
   0;
ded
        vis[s] = 1; q = queue < int > (); q.push(s);
402
        while (q.size()) {
be 1
            int u = q.front(); q.pop();
bdc
            for (int i : g[u]) {
7a2
                if (base[i] == base[u] or match[u] == i) continue;
e35
                if (i == s or (match[i] != -1 and pai[match[i]] != -1))
4 f 2
                     contract(u. i):
e2e
                else if (pai[i] == -1) {
545
                    pai[i] = u;
                    if (match[i] == -1) return i;
f6a
```

```
818
                     i = match[i]:
29 d
                     vis[i] = 1; q.push(i);
90 e
                }
            }
0b5
634
daa
        return -1;
a16 }
83f int blossom() {
        int ans = 0:
315
        memset(match, -1, sizeof(match));
2 e 3
        for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
f76
            for (int i : g[i]) if (match[i] == -1) {
1bc
                 match[i] = i:
f1d
                 match[j] = i;
0df
                 ans++;
c2b
                 break;
723
da8
        for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
7 e 3
            int j = getpath(i);
5f2
            if (j == -1) continue;
0df
            ans++;
3a0
            while (j != -1) {
ef0
                 int p = pai[j], pp = match[p];
348
                 match[p] = j;
fe9
                 match[j] = p;
55 d
                 j = pp;
797
            }
f70
        }
ba7
        return ans;
fcd }
```

#### 1.6 Centro de arvore

```
d47
            if (d[v] > df) f = v, df = d[v];
            for (int u : g[v]) if (u != par[v])
e68
1 a 5
                d[u] = d[v] + 1, par[u] = v, dfs(u);
90d
        };
        f = df = par[0] = -1, d[0] = 0;
1 b 0
41e
        dfs(0):
        int root = f;
c2d
0f6
        f = df = par[root] = -1, d[root] = 0;
        dfs(root):
14e
761
        vector < int > c:
87e
        while (f != -1) {
            if (d[f] == df/2 \text{ or } d[f] == (df+1)/2) \text{ c.push_back}(f);
999
19c
            f = par[f];
3bf
00f
        return {df, c};
9c7 }
1.7 Centroid
// Computa os 2 centroids da arvore
// O(n)
97a int n, subsize[MAX];
042 vector < int > g[MAX];
98f void dfs(int k, int p=-1) {
bd2
        subsize[k] = 1:
6 e 5
       for (int i : g[k]) if (i != p) {
801
            dfs(i, k):
            subsize[k] += subsize[i];
2 e 3
1 b 2
        }
5a5 }
2e8 int centroid(int k, int p=-1, int size=-1) {
        if (size == -1) size = subsize[k]:
e73
        for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
8df
bab
            return centroid(i, k, size);
839
        return k:
```

b6a }

051

f20 pair < int, int > centroids (int k=0) {

dfs(k);

```
909
         int i = centroid(k), i2 = i;
        for (int j : g[i]) if (2*subsize[j] == subsize[k]) i2 = j;
8dd
0cb
         return {i, i2};
cf4 }
1.8 Centroid decomposition
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
// O(n log(n))
042 vector < int > g[MAX];
ba8 int sz[MAX], rem[MAX];
747 void dfs(vector<int>& path, int i, int l=-1, int d=0) {
         path.push_back(d);
75 f
         for (int j : g[i]) if (j != 1 and !rem[j]) dfs(path, j, i,
    d+1):
3e9 }
071 int dfs_sz(int i, int l=-1) {
        sz[i] = 1:
е5с
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
 i):
191
         return sz[i]:
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
735
             return centroid(j, i, size);
d9a
        return i;
96e }
d79 ll decomp(int i, int k) {
        int c = centroid(i, i, dfs_sz(i));
106
a67
        rem[c] = 1;
        // gasta O(n) aqui - dfs sem ir pros caras removidos
        11 \text{ ans} = 0;
04b
        vector < int > cnt(sz[i]):
020
878
         cnt[0] = 1:
0a8
        for (int j : g[c]) if (!rem[j]) {
5 b 4
             vector < int > path;
baf
             dfs(path, j);
1 a 1
             for (int d : path) if (0 \le k-d-1 \text{ and } k-d-1 \le sz[i])
```

### 1.9 Centroid Tree

```
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
// O(n log(n)) de tempo e memoria
845 vector < int > g[MAX], dist[MAX];
c1e int sz[MAX], rem[MAX], p[MAX];
071 int dfs_sz(int i, int l=-1) {
        sz[i] = 1:
02 c
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
е5с
   i):
191
        return sz[i]:
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
735
            return centroid(j, i, size);
d9a
        return i;
96e }
324 void dfs_dist(int i, int 1, int d=0) {
541
        dist[i].push_back(d);
5 a 1
        for (int j : g[i]) if (j != l and !rem[j])
82a
            dfs_dist(j, i, d+1);
645 }
27e void decomp(int i, int l = -1) {
        int c = centroid(i, i, dfs_sz(i));
1 b 9
        rem[c] = 1, p[c] = 1;
534
        dfs dist(c, c);
a2a
       for (int j : g[c]) if (!rem[j]) decomp(j, c);
ebd }
```

```
76c void build(int n) {
235         for (int i = 0; i < n; i++) rem[i] = 0, dist[i].clear();
867         decomp(0);
96b         for (int i = 0; i < n; i++) reverse(dist[i].begin(),
         dist[i].end());
a78 }</pre>
```

### 1.10 Dijkstra

```
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF.
// entao x nao alcanca i
//
// O(m log(n))
eff ll d[MAX]:
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
1a8 int n:
abc void diikstra(int v) {
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
a7f
        d[v] = 0;
88 c
        priority_queue < pair < ll , int >> pq;
b32
        pq.emplace(0, v);
265
        while (pg.size()) {
            auto [ndist, u] = pq.top(); pq.pop();
a 25
            if (-ndist > d[u]) continue;
953
cda
            for (auto [idx, w] : g[u]) if (d[idx] > d[u] + w) {
331
                d[idx] = d[u] + w:
a 84
                pq.emplace(-d[idx], idx);
c56
еБс
        }
fec }
```

### 1.11 Dinitz

```
// O(min(m * max_flow, n^2 m))
// Grafo com capacidades 1: O(min(m sqrt(m), m * n^(2/3)))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
```

```
472 struct dinitz {
        const bool scaling = false; // com scaling -> 0(nm
61f
   log(MAXCAP)),
206
        int lim:
                                      // com constante alta
670
        struct edge {
358
            int to, cap, rev, flow;
7f9
            bool res;
d36
            edge(int to_, int cap_, int rev_, bool res_)
                 : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
a 94
f70
        };
002
        vector < vector < edge >> g;
216
        vector < int > lev, beg;
a71
        11 F:
190
        dinitz(int n) : g(n), F(0) {}
        void add(int a, int b, int c) {
087
            g[a].emplace_back(b, c, g[b].size(), false);
bae
            g[b].emplace_back(a, 0, g[a].size()-1, true);
4 c 6
5 c 2
123
        bool bfs(int s, int t) {
            lev = vector < int > (g.size(), -1); lev[s] = 0;
90f
64c
            beg = vector < int > (g.size(), 0);
8 b 2
            queue < int > q; q.push(s);
402
            while (q.size()) {
be 1
                 int u = q.front(); q.pop();
bd9
                for (auto& i : g[u]) {
                     if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
dbc
b4f
                     if (scaling and i.cap - i.flow < lim) continue;</pre>
                     lev[i.to] = lev[u] + 1;
185
8ca
                     q.push(i.to);
                }
f97
e87
0 de
            return lev[t] != -1;
742
dfb
        int dfs(int v, int s, int f = INF) {
50b
            if (!f or v == s) return f;
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
88f
027
                 auto& e = g[v][i];
                if (lev[e.to] != lev[v] + 1) continue;
206
                 int foi = dfs(e.to, s, min(f, e.cap - e.flow));
ee0
                if (!foi) continue:
749
                e.flow += foi, g[e.to][e.rev].flow -= foi;
3 c 5
                return foi;
45с
            }
618
bb3
            return 0:
```

```
4b1
ff6
        ll max_flow(int s, int t) {
a86
            for (lim = scaling ? (1<<30) : 1; lim; lim /= 2)</pre>
9d1
                 while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
4ff
            return F:
8b9
        }
86f }:
// Recupera as arestas do corte s-t
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
f07
        g.max_flow(s, t);
68 c
        vector<pair<int, int>> cut;
1b0
        vector < int > vis(g.g.size(), 0), st = {s};
321
        vis[s] = 1:
3c6
        while (st.size()) {
            int u = st.back(); st.pop_back();
b17
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)</pre>
322
c17
                vis[e.to] = 1, st.push_back(e.to);
d14
        for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])</pre>
481
            if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i,
942
   e.to):
d1b
        return cut:
1e8 }
```

### 1.12 Dominator Tree

```
// Codigo do Kawakami. Se vira pra usar ai
//
// build - O(m log(n))
// dominates - O(1)
1a8 int n;
bbf namespace d_tree {
042
        vector < int > g[MAX];
        // The dominator tree
b39
        vector < int > tree[MAX];
        int dfs_l[MAX], dfs_r[MAX];
5af
        // Auxiliary data
        vector < int > rg[MAX], bucket[MAX];
a2e
3ef
        int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
44b
        int ancestor[MAX], label[MAX];
563
        vector < int > preorder;
```

```
void dfs(int v) {
76a
6a1
            static int t = 0;
            pre[v] = ++t;
db6
767
            sdom[v] = label[v] = v:
            preorder.push_back(v);
a3d
d08
            for (int nxt: g[v]) {
                 if (sdom[nxt] == -1) {
56c
                     prv[nxt] = v;
eed
900
                     dfs(nxt):
f48
2 b 5
                 rg[nxt].push_back(v);
5 a 1
            }
d6a
62e
        int eval(int v) {
c93
            if (ancestor[v] == -1) return v:
a75
            if (ancestor[ancestor[v]] == -1) return label[v];
f33
            int u = eval(ancestor[v]);
b49
            if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
            ancestor[v] = ancestor[u]:
66e
            return label[v]:
c24
0b9
        void dfs2(int v) {
4b2
6a1
            static int t = 0;
330
            dfs l[v] = t++:
5 e 0
            for (int nxt: tree[v]) dfs2(nxt);
            dfs_r[v] = t++:
8 e 2
cfa
с2 с
        void build(int s) {
603
            for (int i = 0; i < n; i++) {</pre>
                 sdom[i] = pre[i] = ancestor[i] = -1;
e6f
2 e 1
                 rg[i].clear();
                 tree[i].clear();
50a
                 bucket[i].clear();
666
3ba
772
            preorder.clear();
с6с
            dfs(s):
12b
            if (preorder.size() == 1) return;
3 c 7
            for (int i = int(preorder.size()) - 1; i >= 1; i--) {
6 c 6
                 int w = preorder[i];
                for (int v: rg[w]) {
a52
                     int u = eval(v);
5 c 1
                     if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] = sdom[u];</pre>
a17
018
680
                 bucket[sdom[w]].push_back(w);
                 ancestor[w] = prv[w];
ea7
                for (int v: bucket[prv[w]]) {
b99
```

```
5 c 1
                    int u = eval(v):
977
                    idom[v] = (u == v) ? sdom[v] : u;
aff
2cc
                bucket[prv[w]].clear();
0a3
            for (int i = 1; i < preorder.size(); i++) {</pre>
d0c
6.6
                int w = preorder[i]:
14b
                if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
32 f
                tree[idom[w]].push_back(w);
c58
8ac
            idom[s] = sdom[s] = -1;
1b6
            dfs2(s):
d09
       }
        // Whether every path from s to v passes through u
        bool dominates(int u, int v) {
490
c75
            if (pre[v] == -1) return 1; // vacuously true
            return dfs l[u] <= dfs l[v] && dfs r[v] <= dfs r[u]:
2ea
332
ce9 }:
```

### 1.13 Euler Path / Euler Cycle

```
// Para declarar: 'euler < true > E(n); ' se guiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice, id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get_path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
//
// O(n+m)
63f template <bool directed=false > struct euler {
1a8
        int n:
4c0
        vector < vector < pair < int , int >>> g;
d63
        vector < int > used;
30f
        euler(int n_) : n(n_), g(n) {}
        void add(int a, int b) {
50f
```

```
4 \, cd
            int at = used.size();
c51
            used.push_back(0);
74e
            g[a].emplace_back(b, at);
            if (!directed) g[b].emplace_back(a, at);
fab
411
d41 #warning chamar para o src certo!
        pair < bool, vector < pair < int, int >>> get_path (int src) {
eed
            if (!used.size()) return {true, {}};
baf
b25
            vector < int > beg(n, 0);
            for (int& i : used) i = 0;
4 e.c.
            // {{vertice, anterior}, label}
            vector<pair<int, int>, int>> ret, st = {{src, -1},
   -1}};
3 c 6
            while (st.size()) {
8ff
                int at = st.back().first.first;
                int& it = beg[at];
002
                while (it < g[at].size() and used[g[at][it].second])</pre>
8a1
   it++:
8 e 4
                if (it == g[at].size()) {
                     if (ret.size() and ret.back().first.second != at)
9 d d
b82
                         return {false, {}};
                     ret.push_back(st.back()), st.pop_back();
420
2c0
                } else {
                     st.push_back({{g[at][it].first, at},
daa
   g[at][it].second});
                     used[g[at][it].second] = 1;
eb8
396
                }
b3a
            if (ret.size() != used.size()+1) return {false, {}};
a19
f77
            vector<pair<int, int>> ans;
            for (auto i : ret) ans.emplace_back(i.first.first,
fdf
   i.second);
459
            reverse(ans.begin(), ans.end());
997
            return {true, ans}:
844
9b6
        pair < bool, vector < pair < int, int >>> get_cycle() {
            if (!used.size()) return {true, {}};
baf
ad1
            int src = 0;
34b
            while (!g[src].size()) src++;
687
            auto ans = get_path(src);
            if (!ans.first or ans.second[0].first !=
   ans.second.back().first)
                return {false, {}};
b82
            ans.second[0].second = ans.second.back().second;
350
8 b 8
            ans.second.pop_back();
ba7
            return ans:
48f
        }
```

```
711 };
```

### 1.14 Floyd-Warshall

```
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
//
// O(n^3)
1a8 int n;
ae5 int d[MAX][MAX];
73c bool floyd_warshall() {
e 22
        for (int k = 0; k < n; k++)
        for (int i = 0; i < n; i++)</pre>
830
f90
        for (int j = 0; j < n; j++)
0ab
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
830
        for (int i = 0; i < n; i++)</pre>
             if (d[i][i] < 0) return 1;</pre>
753
bb3
        return 0;
192 }
```

### 1.15 Functional Graph

```
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (O sse ta no ciclo)
// pos[i] fala a posicao de i no array que eh a concat. dos ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir de i
// Se quiser outra operacao, da pra alterar facil o codigo
// Codigo um pouco louco, tenho que admitir
//
// build - O(n)
// f_k - O(log(min(n, k)))
// path - O(log(min(n, k)))
```

```
6ef namespace func_graph {
1 a 8
        int n;
        int f[MAX], vis[MAX], d[MAX];
ce2
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
ebd
        int sz[MAX], comp;
6a9
        vector < vector < int >> ciclo;
405
        11 val[MAX], jmp[MAX], seg[2*MAX];
97c
        11 op(11 a, 11 b) { return a+b; }; // mudar a operacao aqui
        void dfs(int i, int t = 2) {
27b
9 c 9
            vis[i] = t;
f09
            if (vis[f[i]] >= 2) \{ // comeca ciclo - f[i] eh o rep.
e0a
                d[i] = 0, rt[i] = comp;
                sz[comp] = t - vis[f[i]] + 1;
74c
                p[i] = pp[i] = i, jmp[i] = val[i];
97b
                ciclo.emplace_back();
15c
bfb
                ciclo.back().push_back(i);
a22
            } else {
                if (!vis[f[i]]) dfs(f[i], t+1);
c16
                rt[i] = rt[f[i]];
8 c 0
                if (sz[comp]+1) { // to no ciclo
195
d0f
                    d[i] = 0:
                    p[i] = pp[i] = i, jmp[i] = val[i];
97b
bfb
                    ciclo.back().push_back(i);
c20
                } else { // nao to no ciclo
00d
                    d[i] = d[f[i]]+1, p[i] = f[i]:
                    pp[i] = 2*d[pp[f[i]]] == d[pp[pp[f[i]]]]+d[f[i]] ?
   pp[pp[f[i]]] : f[i];
                    jmp[i] = pp[i] == f[i] ? val[i] : op(val[i],
114
   op(jmp[f[i]], jmp[pp[f[i]]]));
               }
db8
003
            if (f[ciclo[rt[i]][0]] == i) comp++; // fim do ciclo
e4a
29a
            vis[i] = 1:
0ba
        void build(vector<int> f_, vector<int> val_ = {}) {
1da
bcb
            n = f_size(), comp = 0;
527
            if (!val_.size()) val_ = f_;
830
            for (int i = 0; i < n; i++)</pre>
998
                f[i] = f_[i], val[i] = val_[i], vis[i] = 0, sz[i] = -1;
            ciclo.clear();
e74
            for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
158
6bb
            int t = 0:
            for (auto& c : ciclo) {
daa
                reverse(c.begin(), c.end());
336
                for (int j : c) {
ea5
```

```
85 b
                     pos[i] = t;
948
                    seg[n+t] = val[j];
c82
                    t++;
25е
                }
cbc
dc1
            for (int i = n-1; i; i--) seg[i] = op(seg[2*i],
   seg[2*i+1]);
90b
       }
        int f k(int i. ll k) {
283
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
ded
                if (big <= k) k -= big, i = pp[i];</pre>
584
                else k--, i = p[i];
            }
09с
77 e
            if (!k) return i;
a 19
            return ciclo[rt[i]][(pos[i] - pos[ciclo[rt[i]][0]] + k) %
   sz[rt[i]]];
       }
f34
        ll path(int i, ll k) {
047
3cf
            auto query = [&](int 1, int r) {
3 e 4
                11 q = 0:
47a
                for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2) {
27е
                    if (1\%2 == 1) q = op(q, seg[1]);
1f2
                    if (r\%2 == 0) q = op(q, seg[r]);
598
                }
bef
                return q;
6e1
            }:
b73
            11 ret = 0;
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
327
                if (big \leq k) k -= big, ret = op(ret, imp[i]), i =
   pp[i];
f9e
                else k--, ret = op(ret, val[i]), i = p[i];
7 e 3
еЗс
            if (!k) return ret;
a9e
            int first = pos[ciclo[rt[i]][0]], last =
   pos[ciclo[rt[i]].back()];
            // k/sz[rt[i]] voltas completas
            if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] * query(first,
430
   last));
            k %= sz[rt[i]];
9af
еЗс
            if (!k) return ret;
            int l = pos[i], r = first + (pos[i] - first + k - 1) %
8ea
   sz[rt[i]];
```

```
982
            if (1 <= r) return op(ret, query(1, r));</pre>
            return op(ret, op(query(1, last), query(first, r)));
687
380
51f }
1.16 HLD - aresta
// SegTree de soma
// query / update de soma das arestas
// Complexidades:
// build - 0(n)
// query_path - O(log^2 (n))
// update_path - 0(log^2(n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
c0d
        vector<pair<int, int> > g[MAX];
e65
        int pos[MAX], sz[MAX];
        int sobe[MAX], pai[MAX];
7c0
096
        int h[MAX], v[MAX], t;
        void build hld(int k, int p = -1, int f = 1) {
Осе
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
180
            for (auto& i : g[k]) if (i.first != p) {
418
                auto [u. w] = i:
dd2
a76
                sobe[u] = w; pai[u] = k;
0 c 1
                h[u] = (i == g[k][0] ? h[k] : u);
da7
                build_hld(u, k, f); sz[k] += sz[u];
865
                if (sz[u] > sz[g[k][0].first] or g[k][0].first == p)
9 a 3
                    swap(i, g[k][0]);
            }
804
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
4 dd
1f8
        void build(int root = 0) {
            t = 0;
a34
295
            build hld(root):
            seg::build(t, v);
c83
ea2
3fc
        ll query_path(int a, int b) {
2d5
            if (a == b) return 0;
aa1
            if (pos[a] < pos[b]) swap(a, b);</pre>
```

```
if (h[a] == h[b]) return seg::query(pos[b]+1, pos[a]);
29b
            return seg::query(pos[h[a]], pos[a]) +
fca
   query_path(pai[h[a]], b);
87 f
        void update_path(int a, int b, int x) {
920
d54
            if (a == b) return;
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            if (h[a] == h[b]) return (void) seg::update(pos[b]+1,
   pos[a], x);
701
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
   b, x);
dbf
d0a
        11 query_subtree(int a) {
            if (sz[a] == 1) return 0;
b9f
2f6
            return seg::query(pos[a]+1, pos[a]+sz[a]-1);
77 f
        void update_subtree(int a, int x) {
acc
            if (sz[a] == 1) return;
a5a
9 cd
            seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
a 46
7be
        int lca(int a, int b) {
aa1
            if (pos[a] < pos[b]) swap(a, b);
ca5
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
219
599 }
1.17 HLD - vertice
// SegTree de soma
// query / update de soma dos vertices
//
// Complexidades:
// build - O(n)
// query_path - 0(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
042
        vector < int > g[MAX];
e 65
        int pos[MAX], sz[MAX];
```

bd4

int peso[MAX], pai[MAX];

```
096
        int h[MAX], v[MAX], t;
        void build_hld(int k, int p = -1, int f = 1) {
Осе
            v[pos[k] = t++] = peso[k]; sz[k] = 1;
b18
b94
            for (auto& i : g[k]) if (i != p) {
                pai[i] = k:
78d
                h[i] = (i == g[k][0] ? h[k] : i);
26e
                build_hld(i, k, f); sz[k] += sz[i];
193
                if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
   g[k][0]);
d94
667
            if (p*f == -1) build hld(h[k] = k, -1, t = 0):
a83
1f8
        void build(int root = 0) {
            t = 0;
a34
295
            build_hld(root);
            seg::build(t, v);
c83
ea2
        11 query_path(int a, int b) {
3fc
            if (pos[a] < pos[b]) swap(a, b);</pre>
            if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
4bf
            return seg::query(pos[h[a]], pos[a]) +
fca
   query_path(pai[h[a]], b);
c17
920
        void update_path(int a, int b, int x) {
aa1
            if (pos[a] < pos[b]) swap(a, b);</pre>
198
            if (h[a] == h[b]) return (void) seg::update(pos[b], pos[a],
   x):
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
701
   b, x);
421
        11 query_subtree(int a) {
d0a
b3e
            return seg::query(pos[a], pos[a]+sz[a]-1);
        }
ba2
        void update_subtree(int a, int x) {
acc
a22
            seg::update(pos[a], pos[a]+sz[a]-1, x);
480
        }
        int lca(int a, int b) {
7be
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
ca5
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
219
        }
de3 }
```

### 1.18 HLD sem Update

```
// querv de min do caminho
//
// Complexidades:
// build - 0(n)
// query_path - O(log(n))
826 namespace hld {
c0d
        vector<pair<int, int> > g[MAX];
e 65
        int pos[MAX], sz[MAX];
7 c 0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
ea2
        int men[MAX], seg[2*MAX];
        void build_hld(int k, int p = -1, int f = 1) {
Осе
180
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
418
            for (auto& i : g[k]) if (i.first != p) {
                 sobe[i.first] = i.second; pai[i.first] = k;
1f5
                h[i.first] = (i == g[k][0] ? h[k] : i.first);
6fa
                 men[i.first] = (i == g[k][0] ? min(men[k], i.second) :
87b
    i.second):
4b2
                build_hld(i.first, k, f); sz[k] += sz[i.first];
bc3
                if (sz[i.first] > sz[g[k][0].first] or g[k][0].first
    == p)
9a3
                     swap(i, g[k][0]);
ea4
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
667
8ec
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
3ae
            for (int i = 0; i < t; i++) seg[i+t] = v[i];
            for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
8db
    seg[2*i+1]);
ea5
f04
        int query_path(int a, int b) {
            if (a == b) return INF;
490
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
98f
            if (h[a] != h[b]) return min(men[a], query_path(pai[h[a]],
    b));
46b
            int ans = INF, x = pos[b]+1+t, y = pos[a]+t;
646
            for (; x \le y; ++x/=2, --y/=2) ans = min({ans, seg[x],
    seg[y]});
ba7
            return ans;
```

```
3a9 }
ee6 };
```

### 1.19 Hopcroft Karp

```
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// O(|E| * sqrt(|V|)) com constante baixa
// Para grafos esparsos gerados aleatoriamente, roda em O(|E| *
   log(|V|))
// com alta probabilidade
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
958 struct hopcroft_karp {
14e
        int n, m;
        vector < vector < int >> g;
789
        vector < int > dist, nxt, ma, mb;
5 e a
605
        hopcroft_karp(int n_, int m_) : n(n_), m(m_), g(n),
ac5
            dist(n), nxt(n), ma(n, -1), mb(m, -1) {}
ba6
        void add(int a, int b) { g[a].push_back(b); }
        bool dfs(int i) {
caf
            for (int &id = nxt[i]; id < g[i].size(); id++) {</pre>
32b
d9b
                int j = g[i][id];
dd2
                if (mb[j] == -1 or (dist[mb[j]] == dist[i] + 1 and
   dfs(mb[i]))) {
bfe
                     ma[i] = j, mb[j] = i;
8a6
                     return true;
                }
c96
cf0
d1f
            return false:
0 de
838
        bool bfs() {
            for (int i = 0; i < n; i++) dist[i] = n;</pre>
85d
26a
            queue < int > q;
ad2
            for (int i = 0; i < n; i++) if (ma[i] == -1) {</pre>
d6b
                 dist[i] = 0;
3f2
                q.push(i);
```

```
030
43 f
            bool rep = 0;
402
            while (q.size()) {
379
                int i = q.front(); q.pop();
48e
                for (int j : g[i]) {
096
                     if (mb[j] == -1) rep = 1;
395
                     else if (dist[mb[j]] > dist[i] + 1) {
998
                         dist[mb[j]] = dist[i] + 1;
a21
                         q.push(mb[j]);
                    }
040
36 e
                }
fc5
            }
d14
            return rep;
ad7
bf7
        int matching() {
7 c 9
            int ret = 0;
5a8
            for (auto& i : g) shuffle(i.begin(), i.end(), rng);
6d4
            while (bfs()) {
c79
                for (int i = 0; i < n; i++) nxt[i] = 0;
                for (int i = 0; i < n; i++)
830
475
                     if (ma[i] == -1 and dfs(i)) ret++;
939
edf
            return ret:
b77
        }
cd2 }:
```

#### 1.20 Isomorfismo de arvores

```
// thash() retorna o hash da arvore (usando centroids como vertices
    especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
//
// O(|V|.log(|V|))
91f map < vector < int >, int > mphash;
df6 struct tree {
1a8
         int n;
789
         vector < vector < int >> g;
         vector < int > sz, cs;
1 b 5
         tree(int n_{-}): n(n_{-}), g(n_{-}), sz(n_{-}) {}
76b
         void dfs_centroid(int v, int p) {
588
             sz[v] = 1;
fa7
             bool cent = true;
```

```
18e
            for (int u : g[v]) if (u != p) {
365
                dfs_centroid(u, v), sz[v] += sz[u];
e90
                if(sz[u] > n/2) cent = false;
            }
есе
            if (cent and n - sz[v] <= n/2) cs.push_back(v);</pre>
1f6
368
784
        int fhash(int v, int p) {
544
            vector<int> h:
332
            for (int u : g[v]) if (u != p) h.push_back(fhash(u, v));
1 c.9
            sort(h.begin(), h.end());
3ac
            if (!mphash.count(h)) mphash[h] = mphash.size();
bbc
            return mphash[h];
748
        }
38f
        11 thash() {
23a
            cs.clear();
3 a 5
            dfs_centroid(0, -1);
16d
            if (cs.size() == 1) return fhash(cs[0], -1);
772
            11 h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
fae
            return (min(h1, h2) << 30) + max(h1, h2);
138
        }
4dd };
// Versao mais rapida com hash, ideal para hash de floresta.
// subtree_hash(v, p) retorna o hash da subarvore enraizada em v com
// tree_hash() retorna o hash da arvore.
// forest hash() retorna o hash da floresta.
// use o vetor forb[] para marcar vertices que nao podem ser visitados.
//
// O(|V|.log(|V|))
c8a mt19937
   rng(chrono::steady_clock::now().time_since_epoch().count());
426 int uniform(11 1, 11 r) {
969
        uniform_int_distribution < 11 > uid(1, r);
f54
        return uid(rng);
5fc }
3e4 const int MOD = 1e9 + 7;
58c const int H = 13;
db1 const int P = uniform(1, MOD-1);
325 const int P2 = uniform(1, MOD-1);
df6 struct tree {
2 e 2
        int fn:
```

```
vector < vector < int >> g;
789
347
        vector < int > sz, cs;
268
        vector < bool > forb;
f73
        tree(int n_): fn(n_{-}), g(n_{-}), sz(n_{-}), forb(n_{-}) {}
762
        void dfs_size(int v, int p) {
588
            sz[v] = 1:
d8a
            for (int u : g[v]) if (u != p and !forb[u]) {
db5
                 dfs_size(u, v), sz[v] += sz[u];
896
            }
156
        }
301
        void dfs_centroid(int v, int p, int n) {
fa7
            bool cent = true:
            for (int u : g[v]) if (u != p and !forb[u]) {
d8a
e1f
                dfs_centroid(u, v, n);
e 90
                if(sz[u] > n/2) cent = false;
235
1f6
            if (cent and n - sz[v] \le n/2) cs.push_back(v);
188
1fc
        int subtree_hash(int v, int p) {
3a7
            int h = H:
d8a
            for (int u : g[v]) if (u != p and !forb[u]) {
d36
                h = 11(h) * (P + subtree_hash(u, v)) % MOD;
cf2
81 c
            return h;
d83
126
        int tree hash(int v=0) {
23a
            cs.clear():
575
            dfs_size(v, -1);
8a5
            dfs_centroid(v, -1, sz[v]);
d8d
            if (cs.size() == 1) return subtree_hash(cs[0], -1);
098
            assert (cs.size() == 2);
403
            int h1 = subtree hash(cs[0], cs[1]):
            int h2 = subtree_hash(cs[1], cs[0]);
c49
1ae
            return 11(P + h1) * (P + h2) % MOD;
ad7
        }
c50
        int forest_hash() {
5d7
            fill(sz.begin(), sz.end(), 0);
eb4
            int hash = 1;
778
            for (int v = 0; v < fn; v++) if (!sz[v] and !forb[v]) {</pre>
                hash = hash * 11(P2 + tree_hash(v)) % MOD;
f 21
cc2
34 e
            return hash:
2ce
        }
2fe }:
```

### 1.21 Johnson

```
// funciona igual ao Floyd-Warshall
// encontra o menor caminho entre todo
// par de vertices e retorna 1 sse tem
// ciclo negativo no grafo
// O(nm log(m))
1b8 vector<pair<int, ll>> g[MAX]; // {vizinho, peso}
1a5 ll d[MAX][MAX];
e33 bool johnson(int n) {
        vector <11> h(n, 0);
61e
        for (int i = 0; i <= n; i++)</pre>
4 e 3
84d
            for (int v = 0; v < n; v++)
                for (auto [u, w] : g[v]) if (h[u] > h[v] + w) {
aa2
705
                    if (i == n) return 1;
e47
                    h[u] = h[v] + w;
c36
                }
603
        for (int i = 0; i < n; i++) {
            for (int j = 0; j < n; j++) d[i][j] = LINF;</pre>
bb0
682
            d[i][i] = 0:
88c
            priority_queue < pair < ll, int >> pq;
99c
            pq.emplace(0, i);
            while (pq.size()) {
265
d82
                auto [ndist, v] = pq.top(); pq.pop();
                if (-ndist > d[i][v]) continue;
a3f
209
                for (auto [u, w] : g[v]) {
5 a 5
                    w += h[v] - h[u]:
                    if (d[i][u] > d[i][v] + w) {
f18
db9
                         d[i][u] = d[i][v] + w;
558
                         pq.emplace(-d[i][u], u);
ef1
                    }
                }
f7f
a90
f90
            for (int j = 0; j < n; j++)
                d[i][j] += h[j] - h[i];
fa1
8f9
        }
bb3
        return 0;
d3d }
```

### 1.22 Kosaraju

```
// 0(n + m)
1a8 int n;
042 vector < int > g[MAX];
58d vector<int> gi[MAX]; // grafo invertido
c5a int vis[MAX];
ee6 stack<int> S:
a52 int comp[MAX]; // componente conexo de cada vertice
1ca void dfs(int k) {
59a
        vis[k] = 1;
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
            if (!vis[g[k][i]]) dfs(g[k][i]);
8d5
58f
        S.push(k);
89c }
436 void scc(int k, int c) {
59a
        vis[k] = 1;
52 c
        comp[k] = c;
ff0
        for (int i = 0; i < (int) gi[k].size(); i++)</pre>
            if (!vis[gi[k][i]]) scc(gi[k][i], c);
bf6
088 }
db8 void kosaraju() {
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
991
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);
991
        for (int i = 0; i < n; i++) vis[i] = 0;
d32
        while (S.size()) {
70b
            int u = S.top();
7de
            S.pop();
f43
            if (!vis[u]) scc(u, u);
207
        }
e21 }
1.23 Kruskal
// Gera e retorna uma AGM e seu custo total a partir do vetor de
   arestas (edg)
// do grafo
//
```

 $// O(m \log(m) + m a(m))$ 

```
// 864875
1b9 vector < tuple < int, int >> edg; // {peso,[x,y]}
// DSU em O(a(n))
4a6 void dsu_build();
d78 int find(int a);
369 void unite(int a, int b);
c67 pair<ll, vector<tuple<int, int, int>>> kruskal(int n) {
8 d 2
        dsu_build(n);
        sort(edg.begin(), edg.end());
e31
854
        11 cost = 0;
        vector < tuple < int , int , int >> mst;
979
        for (auto [w,x,y] : edg) if (find(x) != find(y)) {
fea
9de
            mst.emplace_back(w, x, y);
45f
            cost += w:
05a
            unite(x,y);
ca2
5 df
        return {cost, mst};
b6a }
1.24 Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinitz
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
6c6 struct kuhn {
```

int n. m:

vector < vector < int >> g;

vector < int > vis, ma, mb;

kuhn(int n\_, int m\_) : n(n\_), m(m\_), g(n),

789

d3f

40e

```
8af
             vis(n+m), ma(n, -1), mb(m, -1) {}
 ba6
         void add(int a, int b) { g[a].push_back(b); }
         bool dfs(int i) {
 caf
 29a
             vis[i] = 1;
29b
             for (int j : g[i]) if (!vis[n+j]) {
 8 c 9
                 vis[n+j] = 1;
 2cf
                 if (mb[j] == -1 or dfs(mb[j])) {
                      ma[i] = j, mb[j] = i;
 bfe
 8a6
                      return true;
b17
                 }
82a
             }
 d1f
             return false;
 4ef
         }
bf7
         int matching() {
 1ae
             int ret = 0, aum = 1;
 5 a 8
             for (auto& i : g) shuffle(i.begin(), i.end(), rng);
 392
             while (aum) {
                 for (int j = 0; j < m; j++) vis [n+j] = 0;
 618
 c.5d
                 aum = 0:
 830
                 for (int i = 0; i < n; i++)
 0.1 f
                      if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
 0.85
             }
 edf
             return ret;
 2ee
b0d }:
 ebf pair<vector<int>, vector<int>> recover(kuhn& K) {
 e80
         K.matching();
 50 c
         int n = K.n, m = K.m;
 9d0
         for (int i = 0; i < n+m; i++) K.vis[i] = 0;</pre>
         for (int i = 0; i < n; i++) if (K.ma[i] == -1) K.dfs(i);</pre>
 bde
 8ad
         vector < int > ca. cb:
 576
         for (int i = 0; i < n; i++) if (!K.vis[i]) ca.push_back(i);</pre>
 f 24
         for (int i = 0; i < m; i++) if (K.vis[n+i]) cb.push_back(i);
         return {ca, cb};
 aad
 55f }
1.25 LCA com binary lifting
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
//
// Complexidades:
```

```
// build - O(n log(n))
// lca - O(log(n))
677 vector < vector < int > > g(MAX);
41c int n. p:
e75 int pai[MAX2][MAX];
999 int in [MAX], out [MAX];
1ca void dfs(int k) {
fdf
        in[k] = p++:
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
9b7
            if (in[g[k][i]] == -1) {
ba6
                pai[0][g[k][i]] = k;
c38
                dfs(g[k][i]);
            }
e2d
26f
        out[k] = p++;
691 }
c11 void build(int raiz) {
        for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
        p = 0, memset(in, -1, sizeof in);
c63
       dfs(raiz):
ecb
        // pd dos pais
511
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < n; i++)
d38
            pai[k][i] = pai[k - 1][pai[k - 1][i]];
530 }
00f bool anc(int a, int b) { // se a eh ancestral de b
bfe
        return in[a] <= in[b] and out[a] >= out[b];
2 d 6 }
7be int lca(int a. int b) {
       if (anc(a, b)) return a:
86d
       if (anc(b, a)) return b;
e 52
        // sobe a
f70
        for (int k = MAX2 - 1; k >= 0; k--)
acf
            if (!anc(pai[k][a], b)) a = pai[k][a];
847
        return pai[0][a];
5 c 4 }
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
```

```
// build - 0(n)
// kth, lca, dist - O(log(n))
9c6 int d[MAX], p[MAX], pp[MAX];
d40 void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
e9d void add_leaf(int i, int u) {
        p[i] = u. d[i] = d[u]+1:
b15
        pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
33f }
 c37 int kth(int i, int k) {
4 e 3
         int dd = max(0, d[i]-k);
935
         while (d[i] > dd) i = d[pp[i]] >= dd ? pp[i] : p[i];
d9a
         return i;
f3c }
7be int lca(int a, int b) {
        if (d[a] < d[b]) swap(a, b);</pre>
6cd
         while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] : p[a];
984
        while (a != b) {
932
             if (pp[a] != pp[b]) a = pp[a], b = pp[b];
e7c
             else a = p[a], b = p[b];
4ea
3f5
         return a;
21d }
4fe int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
042 vector < int > g[MAX];
3ab void build(int i, int pai=-1) {
         if (pai == -1) set_root(i);
5cf
15 f
        for (int j : g[i]) if (j != pai) {
d31
             add_leaf(j, i);
b21
             build(j, i);
43b
        }
74a }
1.26 LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
```

```
// anc(a, b) responde se 'a' eh ancestral de 'b'
// Complexidades:
// build - O(n)
// lca - O(log(n))
// anc - 0(1)
042 vector < int > g[MAX];
713 int pos[MAX], h[MAX], sz[MAX];
ff1 int pai[MAX], t;
8bf void build(int k, int p = -1, int f = 1) {
        pos[k] = t++; sz[k] = 1;
        for (int& i : g[k]) if (i != p) {
e26
78d
            pai[i] = k;
26e
            h[i] = (i == g[k][0] ? h[k] : i);
cb8
            build(i, k, f); sz[k] += sz[i];
            if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k][0]);
cd1
        }
917
3da
        if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
1b9 }
7be int lca(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
ca5
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
219 }
00f bool anc(int a, int b) {
db5
        return pos[a] \le pos[b] and pos[b] \le pos[a] + sz[a] - 1;
272 }
1.27 LCA com RMQ
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b. lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
//
// Complexidades:
// build - O(n)
// lca - 0(1)
// dist - 0(1)
```

1a5 template < typename T > struct rmq {

int n; static const int b = 30;

vector <T> v;

517

fcc

```
70e
        vector < int > mask, t;
        int op(int x, int y) { return v[x] < v[y] ? x : y; }</pre>
18 e
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
6ad
        rma() {}
43 c
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) {
2e5
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
a 61
                at = (at << 1) & ((1 << b) -1);
76a
                 while (at and op(i, i-msb(at&-at)) == i) at ^= at&-at;
53 c
243
            for (int i = 0; i < n/b; i++) t[i] =</pre>
   b*i+b-1-msb(mask[b*i+b-1]):
39 d
            for (int i = 1; (1<<i) <= n/b; i++) for (int i = 0;
   i+(1<< j) <= n/b; i++)
ba5
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
   t[n/b*(j-1)+i+(1<<(j-1))]);
2d3
        int small(int r. int sz = b) { return
c92
   r-msb(mask[r]&((1<<sz)-1));}
        T query(int 1, int r) {
b7a
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
27b
7bf
            int ans = op(small(1+b-1), small(r));
            int x = 1/b+1, y = r/b-1;
e80
e 25
            if (x <= y) {
                int j = msb(y-x+1);
a4e
002
                 ans = op(ans, op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
4b6
ba7
            return ans:
        }
6bf
021 };
065 namespace lca {
042
        vector < int > g[MAX];
8ec
        int v[2*MAX], pos[MAX], dep[2*MAX];
8bd
        int t;
2de
        rmq<int> RMQ;
4cf
        void dfs(int i, int d = 0, int p = -1) {
c97
            v[t] = i, pos[i] = t, dep[t++] = d;
cac
            for (int j : g[i]) if (j != p) {
8ec
                dfs(j, d+1, i);
                 v[t] = i, dep[t++] = d;
cf2
843
d6a
789
        void build(int n, int root) {
a 34
            t = 0:
            dfs(root);
14e
```

```
3f4
             RMQ = rmq < int > (vector < int > (dep, dep + 2 * n - 1));
657
        }
7be
        int lca(int a, int b) {
             a = pos[a], b = pos[b];
ab7
9 c 0
             return v[RMQ.query(min(a, b), max(a, b))];
5 db
b5d
        int dist(int a, int b) {
670
             return dep[pos[a]] + dep[pos[b]] - 2*dep[pos[lca(a, b)]];
5 b 7
645 }
```

#### 1.28 Line Tree

```
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - O(log(n))
1a8 int n;
3ae namespace linetree {
        int id[MAX], seg[2*MAX], pos[MAX];
f37
43f
        vector < int > v [MAX]. val[MAX];
        vector < pair < int , pair < int , int > > ar;
430
dc6
        void add(int a, int b, int p) { ar.push_back({p, {a, b}}); }
0a8
        void build() {
b09
            sort(ar.rbegin(), ar.rend());
            for (int i = 0; i < n; i++) id[i] = i, v[i] = {i},
    val[i].clear();
8bb
            for (auto i : ar) {
                int a = id[i.second.first], b = id[i.second.second];
c91
                 if (a == b) continue;
f6f
c58
                if (v[a].size() < v[b].size()) swap(a, b);</pre>
                 for (auto j : v[b]) id[j] = a, v[a].push_back(j);
fb8
                 val[a].push_back(i.first);
482
                for (auto j : val[b]) val[a].push_back(j);
78b
e39
                 v[b].clear(), val[b].clear();
            }
012
8 e 8
            vector<int> vv;
            for (int i = 0; i < n; i++) for (int j = 0; j < n
    v[i].size(); j++) {
e52
                pos[v[i][j]] = vv.size();
```

```
941
                if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
1cb
                else vv.push_back(0);
475
bb4
            for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];
            for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
69 e
   seg[2*i+1]);
9fe
       }
        int query(int a, int b) {
4ea
596
            if (id[a] != id[b]) return 0; // nao estao conectados
ab7
            a = pos[a], b = pos[b];
d11
            if (a > b) swap(a, b);
199
            b--:
38 a
            int ans = INF;
513
            for (a += n, b += n; a <= b; ++a/=2, --b/=2) ans =
   min({ans, seg[a], seg[b]});
ba7
            return ans:
952
        }
00f };
```

#### 1.29 Link-cut Tree

```
// Link-cut tree padrao
// Todas as operacoes sao O(\log(n)) amortizado
1ef namespace lct {
        struct node {
3c9
19f
            int p, ch[2];
            node() \{ p = ch[0] = ch[1] = -1; \}
062
f43
        };
5f3
        node t[MAX];
971
        bool is root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
657
    t[t[x].p].ch[1] != x);
cf1
        }
        void rotate(int x) {
ed6
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
8fa
            t[x].p = pp, t[p].p = x;
49b
        }
07 с
        void splay(int x) {
```

```
18c
            while (!is_root(x)) {
                                                                             4e4
                                                                                         ll lazy;
497
                 int p = t[x].p, pp = t[p].p;
                                                                             f93
                                                                                         node() {}
0 c 5
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
                                                                             7a8
                                                                             546
   == x) ? x : p):
                rotate(x):
                                                                             b07
64f
            }
                                                                             53b
d8d
4fa
        }
                                                                             6e0
                                                                                     }:
f16
        int access(int v) {
0eb
            int last = -1;
                                                                             c53
            for (int w = v; w+1; last = w, splay(v), w = t[v].p)
                                                                             99 e
01a
024
                 splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                                                                             e4d
                                                                                     int sz;
3 d3
            return last:
0 a 4
        }
                                                                             95 a
                                                                                     void prop(int x) {
                                                                                         if (t[x].lazy) {
e89
        int find_root(int v) {
                                                                             dc1
5 e 3
            access(v):
                                                                             25е
3 de
            while (t[v].ch[0]+1) v = t[v].ch[0];
                                                                             2ab
f05
            return splay(v), v;
                                                                             edc
                                                                             942
ee7
142
        void link(int v, int w) { // v deve ser raiz
                                                                             1ba
5 e 3
            access(v):
                                                                                         if (t[x].rev) {
                                                                             aa2
                                                                             f 95
10d
            t[v].p = w;
c56
                                                                             379
        void cut(int v) { // remove aresta de v pro pai
4 e 6
                                                                             c3d
5 e 3
            access(v):
                                                                             50e
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
                                                                             230
5f5
                                                                             f9d
        int lca(int v, int w) {
bbb
                                                                             564
                                                                                     void update(int x) {
948
            return access(v). access(w):
                                                                             1a3
b6d
                                                                             8ca
e4e }
                                                                             621
                                                                             c.4f
                                                                             269
                                                                             400
                                                                                         }
1.30 Link-cut Tree - aresta
                                                                                     }
                                                                             28b
                                                                             971
                                                                                     bool is_root(int x) {
// Valores nas arestas
                                                                             657
// rootify(v) torna v a raiz de sua arvore
                                                                                t[t[x].p].ch[1] != x);
// query(v, w) retorna a soma do caminho v--w
                                                                             cf1
// update(v, w, x) soma x nas arestas do caminho v--w
                                                                             ed6
                                                                                     void rotate(int x) {
//
                                                                             497
// Todas as operacoes sao O(log(n)) amortizado
                                                                             fc4
                                                                             251
1ef namespace lct {
                                                                             461
        struct node {
3 c.9
                                                                             a 76
19f
            int p, ch[2];
                                                                             8fa
810
            ll val, sub;
```

aa6

04a

bool rev;

int sz, ar;

```
node(int v, int ar_) :
            p(-1), val(v), sub(v), rev(0), sz(ar_{-}), ar(ar_{-}), lazy(0) {
                ch[0] = ch[1] = -1:
        node t[2*MAX]; // MAXN + MAXQ
        map<pair<int, int>, int> aresta;
                if (t[x].ar) t[x].val += t[x].lazy;
                t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
                swap(t[x].ch[0], t[x].ch[1]);
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
            t[x].lazy = 0, t[x].rev = 0;
            t[x].sz = t[x].ar, t[x].sub = t[x].val;
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
                prop(t[x].ch[i]);
                t[x].sz += t[t[x].ch[i]].sz;
                t[x].sub += t[t[x].ch[i]].sub;
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
            int p = t[x].p, pp = t[p].p;
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
            update(p), update(x);
444
f31
       }
```

```
238
        int splay(int x) {
            while (!is root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
                rotate(x);
64f
72c
            return prop(x), x;
aab
08f
        }
f16
        int access(int v) {
0eb
            int last = -1:
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
294
9f1
        void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar);
  }
        int find root(int v) {
e89
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
16a
82f
        bool conn(int v. int w) {
            access(v), access(w):
2 cf
b9b
            return v == w ? true : t[v].p != -1;
ec0
277
        void rootify(int v) {
            access(v):
5 e 3
a02
            t[v].rev ^= 1;
a05
        11 auerv(int v. int w) {
971
            rootify(w), access(v);
b54
249
            return t[v].sub;
652
        }
3fa
        void update(int v, int w, int x) {
b54
            rootify(w), access(v);
12c
            t[v].lazv += x;
74f
204
        void link_(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
523
        void link(int v, int w, int x) { // v--w com peso x
6b8
379
            int id = MAX + sz++;
```

```
110
            aresta[make_pair(v, w)] = id;
            make_tree(id, x, 1);
a88
c88
            link_(v, id), link_(id, w);
58 c
        void cut (int v. int w) {
e 63
            rootify(w), access(v);
b54
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
7cd
        }
031
        void cut(int v, int w) {
            int id = aresta[make_pair(v, w)];
b0f
a4a
            cut_(v, id), cut_(id, w);
840
        }
bbb
        int lca(int v. int w) {
            access(v):
5 e 3
a8b
            return access(w);
524
9ce }
```

#### 1.31 Link-cut Tree - vertice

```
// Valores nos vertices
// make_tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v. w. x) soma x nos vertices do caminho v--w
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
3 c 9
        struct node {
            int p, ch[2];
19 f
810
            ll val, sub;
aa6
            bool rev;
            int sz;
e4d
4 e 4
            ll lazv:
f 93
            node() {}
aa0
            node(int v) : p(-1), val(v), sub(v), rev(0), sz(1),
   lazv(0) {
b07
                ch[0] = ch[1] = -1;
с4е
2b7
        };
5 f 3
        node t[MAX]:
95 a
        void prop(int x) {
```

```
dc1
            if (t[x].lazy) {
9 f 7
                t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].sz;
edc
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
e26
            }
            if (t[x].rev) {
aa2
                swap(t[x].ch[0], t[x].ch[1]);
f95
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
            }
50e
230
            t[x].lazy = 0, t[x].rev = 0;
c62
564
        void update(int x) {
            t[x].sz = 1. t[x].sub = t[x].val:
ec2
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
                t[x].sub += t[t[x].ch[i]].sub;
269
            }
400
        }
da7
971
        bool is root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
657
   t[t[x].p].ch[1] != x);
       }
cf1
ed6
        void rotate(int x) {
497
            int p = t[x].p, pp = t[p].p;
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
fc4
251
            bool d = t[p].ch[0] == x;
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
        }
f31
238
        int splav(int x) {
            while (!is_root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
0 c 5
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
64f
                rotate(x):
72c
            return prop(x), x;
aab
       }
08f
f16
        int access(int v) {
            int last = -1:
0eb
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
```

```
t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last:
294
        }
f 17
        void make_tree(int v, int w) { t[v] = node(w); }
        int find_root(int v) {
e89
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
        }
16a
f94
        bool connected(int v, int w) {
2cf
            access(v), access(w);
b9b
            return v == w ? true : t[v].p != -1:
ec6
277
        void rootify(int v) {
5 e 3
            access(v);
a02
            t[v].rev ^= 1;
        }
a 05
971
        11 query(int v, int w) {
            rootify(w), access(v);
b54
249
            return t[v].sub;
652
3fa
        void update(int v, int w, int x) {
b54
            rootify(w), access(v);
12 c
            t[v].lazv += x:
74 f
142
        void link(int v, int w) {
821
            rootifv(w):
389
            t[w].p = v;
8a8
        void cut(int v. int w) {
0.31
b54
            rootify(w), access(v);
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
264
d9a
        int lca(int v, int w) {
bbb
5e3
            access(v):
a8b
            return access(w);
524
        }
f9f }
```

### 1.32 Max flow com lower bound

```
// add(a, b, l, r):
// adiciona aresta de a pra b, onde precisa passar f de fluxo, l <= f
     <= r
// add(a, b, c):</pre>
```

```
// adiciona aresta de a pra b com capacidade c
// Mesma complexidade do Dinitz
cd5 struct 1b max flow : dinitz {
5 се
        vector < int > d;
d8 c
        lb max flow(int n): dinitz(n + 2), d(n, 0) {}
        void add(int a, int b, int l, int r) {
b12
c97
            d[a] -= 1;
            d[b] += 1:
f1b
4 c 0
            dinitz::add(a, b, r - 1);
ed4
087
        void add(int a, int b, int c) {
0f3
            dinitz::add(a, b, c);
039
7 a 1
        bool has_circulation() {
50c
            int n = d.size();
854
            11 cost = 0;
603
            for (int i = 0: i < n: i++) {
                if (d[i] > 0) {
c69
                    cost += d[i]:
f56
57a
                    dinitz::add(n, i, d[i]);
c72
                } else if (d[i] < 0) {</pre>
b76
                    dinitz::add(i, n+1, -d[i]);
b42
                }
            }
676
            return (dinitz::max_flow(n, n+1) == cost);
067
110
        bool has flow(int src, int snk) {
7 bd
387
            dinitz::add(snk, src, INF);
e40
            return has_circulation();
cc1
        11 max_flow(int src, int snk) {
4eb
ee8
            if (!has_flow(src, snk)) return -1;
            dinitz::F = 0;
4ad
fe5
            return dinitz::max_flow(src, snk);
619
e8f };
1.33 MinCostMaxFlow
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
```

```
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
// 0(nm + f * m log n)
123 template < typename T> struct mcmf {
        struct edge {
670
b75
            int to, rev, flow, cap; // para, id da reversa, fluxo,
    capacidade
7f9
            bool res; // se eh reversa
635
            T cost: // custo da unidade de fluxo
892
            edge(): to(0), rev(0), flow(0), cap(0), cost(0),
    res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T cost_, bool
    res )
                : to(to_), rev(rev_), flow(flow_), cap(cap_),
f8d
    res(res), cost(cost) {}
723
        };
002
        vector < vector < edge >> g;
168
        vector<int> par_idx, par;
f1e
        T inf:
a03
        vector <T> dist;
        mcmf(int n) : g(n), par_idx(n), par(n),
    inf(numeric_limits <T>::max()/3) {}
91 c
        void add(int u, int v, int w, T cost) { // de u pra v com cap
    w e custo cost
            edge a = edge(v, g[v].size(), 0, w, cost, false);
2fc
234
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
b24
            g[u].push_back(a);
c12
            g[v].push_back(b);
        }
0ed
        vector <T> spfa(int s) { // nao precisa se nao tiver custo
8bc
   negativo
871
            deque < int > q;
3d1
            vector < bool > is_inside(g.size(), 0);
            dist = vector<T>(g.size(), inf);
577
a 93
            dist[s] = 0:
a30
            q.push_back(s);
            is inside[s] = true:
e cb
```

```
14d
             while (!q.empty()) {
                 int v = q.front();
b1e
                 q.pop_front();
ced
                 is_inside[v] = false;
48d
                 for (int i = 0; i < g[v].size(); i++) {</pre>
76e
9d4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
943
                          dist[to] = dist[v] + cost;
ed6
                          if (is_inside[to]) continue;
                          if (!q.empty() and dist[to] > dist[q.front()])
   q.push_back(to);
b33
                          else q.push_front(to);
                          is_inside[to] = true;
b52
2 d 1
                     }
8 cd
                 }
f2c
8 d 7
             return dist;
96c
2 a 2
        bool dijkstra(int s, int t, vector < T > & pot) {
             priority_queue <pair <T, int>, vector <pair <T, int>>,
489
   greater<>> q;
577
             dist = vector <T>(g.size(), inf);
             dist[s] = 0:
a 93
115
             g.emplace(0, s);
             while (q.size()) {
402
91b
                 auto [d, v] = q.top();
833
                 q.pop();
68b
                 if (dist[v] < d) continue;</pre>
                 for (int i = 0; i < g[v].size(); i++) {</pre>
76e
9 d 4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
                     cost += pot[v] - pot[to];
e8c
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
                          dist[to] = dist[v] + cost;
943
441
                          q.emplace(dist[to], to);
88b
                          par_idx[to] = i, par[to] = v;
873
                     }
de3
                 }
9d4
             return dist[t] < inf:</pre>
1 d 4
c68
3 d 2
        pair < int , T > min_cost_flow(int s, int t, int flow = INF) {
3dd
             vector<T> pot(g.size(), 0);
             pot = spfa(s); // mudar algoritmo de caminho minimo aqui
9 e 4
```

```
d22
            int f = 0;
ce8
            T ret = 0:
4a0
             while (f < flow and dijkstra(s, t, pot)) {</pre>
                 for (int i = 0; i < g.size(); i++)</pre>
bda
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
d2a
71b
                 int mn flow = flow - f. u = t;
                 while (u != s){
045
                     mn_flow = min(mn_flow,
90f
07 d
                         g[par[u]][par_idx[u]].cap -
    g[par[u]][par_idx[u]].flow);
3d1
                     u = par[u]:
935
1f2
                 ret += pot[t] * mn_flow;
476
                 u = t;
                 while (u != s) {
045
e 09
                     g[par[u]][par_idx[u]].flow += mn_flow;
498
                     g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
                     u = par[u];
3d1
                 }
bcc
04 d
                 f += mn_flow;
36 d
            }
15b
             return make_pair(f, ret);
        }
ссЗ
        // Opcional: retorna as arestas originais por onde passa flow
182
        vector < pair < int , int >> recover() {
             vector<pair<int,int>> used;
24a
2a4
            for (int i = 0; i < g.size(); i++) for (edge e : g[i])
                 if (e.flow == e.cap && !e.res) used.push_back({i,
587
   e.to}):
f6b
            return used:
390
697 };
1.34 Prufer code
// Traduz de lista de arestas para prufer code
// e vice-versa
```

```
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de 0 a n-1
// Todo array com n-2 posicoes e valores de
```

```
// O a n-1 sao prufer codes validos
// O(n)
47d vector < int > to_prufer (vector < pair < int , int >> tree) {
        int n = tree.size()+1;
2 cf
        vector < int > d(n, 0):
4aa
        vector < vector < int >> g(n);
f87
        for (auto [a, b] : tree) d[a]++, d[b]++,
f60
             g[a].push_back(b), g[b].push_back(a);
c5a
        vector < int > pai(n, -1);
260
        queue < int > q; q.push(n-1);
402
        while (q.size()) {
             int u = q.front(); q.pop();
be 1
            for (int v : g[u]) if (v != pai[u])
34c
9 c 9
                 pai[v] = u, q.push(v);
70d
        }
399
        int idx, x:
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
4b8
        vector < int > ret;
b28
        for (int i = 0; i < n-2; i++) {</pre>
d4b
            int y = pai[x];
            ret.push_back(y);
e81
666
            if (--d[v] == 1 \text{ and } v < idx) x = v;
             else idx = x = find(d.begin()+idx+1, d.end(), 1) -
   d.begin();
5f9
edf
        return ret:
d3b }
4d8 vector <pair <int, int>> from_prufer (vector <int> p) {
455
        int n = p.size()+2;
126
        vector < int > d(n, 1);
        for (int i : p) d[i]++:
650
85b
        p.push_back(n-1);
399
        int idx, x;
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
1 df
        vector<pair<int, int>> ret;
b06
        for (int y : p) {
dab
            ret.push_back({x, y});
666
            if (--d[y] == 1 \text{ and } y < idx) x = y;
            else idx = x = find(d.begin()+idx+1, d.end(), 1) -
367
   d.begin();
c3b
edf
        return ret;
765 }
```

#### 1.35 Rooted Euler Tour Tree

```
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
// Chamar ETT E(n. w), passando n = numero de vertices
// e w = vector com os valores de cada vertice
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// get_root(v) retorna a raiz de v
// condense(v) remove v e faz seus filhos apontarem para pai(v)
// erase(v) remove v e corta suas arestas
// - Se usar condense ou erase, nao pode usar get_root!
// vtx_query e vtx_update operam sobre valor de um vertice
// subtree_query e subtree_update operam sobre sub-arvore
// is_ancestor(v, u) retorna se v eh ancestral de u
//
// Tudo O(log(n)) com alta probabilidade
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
0f9 template < typename T> struct rooted_ett {
        const static T ZERO = T();
83 c
3c9
        struct node {
            node *1, *r, *p;
ed1
            int pr, sz;
fa4
875
            T val, sub, lazy;
53e
            int id; // id, ou -1 se nao eh first
5ef
            int qt_f; // numero de 'first' na subarvore
3f6
            node(T v, int id_) : 1(NULL), r(NULL), p(NULL), pr(rng()),
                sz(1), val(v), sub(v), lazy(), id(id_), qt_f(id >= 0)
3ee
  {}
а9с
            void prop() {
9fc
                if (lazy != ZERO) {
1 e 2
                    if (id >= 0) val += lazy;
2cf
                    sub += lazy * qt_f;
b87
                    if (1) 1->lazy += lazy;
d3b
                    if (r) r->lazy += lazy;
f74
                }
520
                lazv = ZER0:
4e7
01e
            void update() {
c21
                sz = 1, sub = val, qt_f = id >= 0;
171
                if (1) 1->prop(), sz += 1->sz, sub += 1->sub, qt_f +=
```

```
1 -> qt_f;
                                                                                27b
                 if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub, qt_f +=
                                                                                d79
117
   r->qt_f;
            }
93e
                                                                                94 f
4f5
        }:
                                                                                d69
73c
        int size(node* x) { return x ? x->sz : 0; }
                                                                                8ea
bcf
        void join(node* 1, node* r, node*& i) {
                                                                                1bc
986
             if (!1 or !r) return void(i = 1 ? 1 : r);
                                                                                86 e
161
             1->prop(), r->prop();
                                                                                d72
ff5
             if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
                                                                                df9
982
             else join(1, r->1, r->1), r->1->p = i = r;
                                                                                cf1
bda
             i->update():
                                                                                f1a
84d
        }
                                                                                b2d
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
a20
                                                                                681
26a
             if (!i) return void(r = 1 = NULL);
                                                                                97 c
c89
             i->prop();
                                                                                ace
             if (kev + size(i->1) < v) {</pre>
d9e
                                                                                a28
                 split(i\rightarrow r, i\rightarrow r, r, v, kev+size(i\rightarrow l)+1), l = i;
                                                                                474
448
                 if (r) r -> p = NULL;
a21
                                                                                b38
6e8
                 if (i->r) i->r->p = i;
                                                                                ead
396
             } else {
                                                                                1f7
                 split(i->1, 1, i->1, v, key), r = i;
                                                                                275
98d
5 a 3
                 if (1) 1->p = NULL;
899
                 if (i->1) i->1->p = i;
                                                                                7af
18b
                                                                                c9f
bda
             i->update();
                                                                                b1f
134
3fa
        pair < node *, int > get_idx(node * i) { // {root, idx}}
                                                                                61 c
6cf
             int ret = size(i->1);
65d
             for (; i->p; i = i->p) if (i != i->p->1)
                                                                                1aa
fc8
                 ret += size(i -> p -> 1) + 1;
                                                                                031
             i->prop();
c89
                                                                                f07
e51
             return {i. ret}:
                                                                                c59
        }
fa4
                                                                                e 39
                                                                                4e6
                                                                                d05
4fb
        vector < node *> first, last;
                                                                                d06
96c
        rooted_ett(int n, vector<T> w) : first(n), last(n) {
                                                                                fcc
aab
             node * tmp;
                                                                                20 c
             for (int i = 0; i < n; i++) {
                                                                                f6e
603
761
                 first[i] = new node(w[i], i);
                                                                                43 c
45f
                 join(first[i], last[i] = new node(ZERO, -1), tmp);
                                                                                00b
            }
d89
b6c
562
        rooted_ett(const rooted_ett& t) { throw logic_error("Nao
                                                                                1ab
   copiar a ETT!"); }
                                                                                b73
```

```
\simrooted_ett() {
        for (int i = 0; i < first.size(); i++) delete first[i],</pre>
delete last[i];
   }
    tuple < node *, int, int > get_range(int i) {
        auto [root, l_idx] = get_idx(first[i]);
        return {root, l_idx, get_idx(last[i]).second};
    tuple < node *, node *, node *> split3(node * root, int 1, int r) {
        node *L, *M, *R;
        split(root, M, R, r+1);
        split(M. L. M. 1);
        return {L, M, R};
    node* extract_range(node*& root, int 1, int r) {
        auto [L, M, R] = split3(root, 1, r);
        join(L, R, root);
        return M;
    }
    void join3(node* L, node* M, node* R) {
        join(L, M, M), join(M, R, R);
    }
    void link(int v, int u) { // parent(v) <-- u</pre>
        auto [root_v, 1, r] = get_range(v);
        root_v = extract_range(root_v, 1, r);
        if (u == -1) return;
        auto [root, idx] = get_idx(first[u]);
        node *L, *R;
        split(root, L, R, idx+1);
        ioin3(L, root v, R);
    void cut(int v) { // corta de v pro pai
        link(v, -1);
    int get_root(int v) {
        auto [root, idx] = get_idx(first[v]);
        while (root->1) root = root->1;
        return root -> id;
    }
    // se usar condense ou erase, nao pode usar get_root!
    void condense(int v) { // remove v, pai dos filhos vira pai(v)
        auto [root, 1, r] = get_range(v);
```

```
a56
            auto E2 = extract_range(root, r, r);
            auto E1 = extract_range(root, 1, 1);
00d
468
            join(E1, E2, E2);
eb6
        }
930
        void erase(int v) { // remove v e arestas adjacentes
124
            cut(v):
f75
            condense(v):
а9с
        }
        T vtx_query(int v) {
c74
20c
            auto [root, idx] = get_idx(first[v]);
5a1
            auto [L, M, R] = split3(root, idx, idx);
d43
            T ans = M->sub:
20c
            join3(L, M, R);
ba7
            return ans;
010
8e6
        void vtx_update(int v, T val) {
20c
            auto [root, idx] = get_idx(first[v]);
5 a 1
            auto [L, M, R] = split3(root, idx, idx);
409
            M \rightarrow lazy += val;
6aa
            M ->prop();
20c
            join3(L, M, R);
385
c99
        T subtree_query(int v) {
b73
            auto [root, 1, r] = get_range(v);
ace
            auto [L, M, R] = split3(root, 1, r);
d43
            T ans = M->sub;
20c
            join3(L, M, R);
ba7
            return ans;
40f
        void subtree_update(int v, T val) {
80d
b73
            auto [root, 1, r] = get_range(v);
            auto [L, M, R] = split3(root, 1, r);
ace
409
            M \rightarrow lazv += val:
            join3(L, M, R);
20c
afb
        bool is_ancestor(int v, int u) { // se v eh ancestral de u
a8e
497
            auto [root_v, l_v, r_v] = get_range(v);
3 c 7
            auto [root_u, l_u, r_u] = get_range(u);
            return root v == root u and l v <= l u and r u <= r v:
b9c
786
7b9 };
```

### 1.36 Sack (DSU em arvores)

```
// Responde queries de todas as sub-arvores
// offline
//
// O(n log(n))
6bf int sz[MAX], cor[MAX], cnt[MAX];
042 vector < int > g[MAX];
6df void build(int k, int d=0) {
        sz[k] = 1:
e8f
01a
        for (auto& i : g[k]) {
30f
            build(i, d+1); sz[k] += sz[i];
925
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
011
        }
189 }
74f void compute(int k, int x, bool dont=1) {
        cnt[cor[k]] += x;
de9
828
        for (int i = dont; i < g[k].size(); i++)</pre>
            compute(g[k][i], x, 0);
b5c
896 }
dc4 void solve(int k, bool keep=0) {
32a
        for (int i = int(g[k].size())-1; i >= 0; i--)
b4c
            solve(g[k][i], !i);
4a0
        compute(k, 1);
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
830
        if (!keep) compute(k, -1, 0);
8bc }
1.37 Stable Marriage
// Emparelha todos os elementos de A com elementos de B
// de forma que nao exista um par x \in A, y \in B
// e x nao pareado com y tal que x prefira parear com y
// e y prefira parear com x.
// a[i] contem os elementos de B ordenados por preferencia de i
// b[j] contem os elementos de A ordenados por preferencia de j
// | A | <= | B |
//
// Retorna um vetor v de tamanho |A| onde v[i] guarda o match de i.
```

```
// O(|A| * |B|)
                                                                              48e
380 vector < int > stable_marriage (vector < vector < int >> &a,
                                                                              740
   vector < vector < int >> &b) {
                                                                              994
652
        int n = a.size(), m = b.size();
                                                                              d64
        assert(a[0].size() == m and b[0].size() == n and n <= m);
83e
017
        vector < int > match(m, -1), it(n, 0);
        vector inv_b(m, vector<int>(n));
e6f
                                                                              3de
a34
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
                                                                              3 c 3
9f2
                                                                              9 c 5
             inv_b[i][b[i][j]] = j;
                                                                              2ef
26a
        queue < int > q;
                                                                              266
5af
        for (int i = 0; i < n; i++) q.push(i);</pre>
        while (q.size()) {
402
                                                                              253
379
            int i = q.front(); q.pop();
                                                                              38a }
            int j = a[i][it[i]];
4 b 8
             if (match[j] == -1) match[j] = i;
57c
                                                                              6bb
02d
             else if (inv_b[j][i] < inv_b[j][match[j]]) {</pre>
                                                                              991
5 d 1
                 q.emplace(match[j]);
                 it[match[j]]++;
e7d
                                                                              3be
f1d
                 match[i] = i;
                                                                              eal }
            } else q.emplace(i), it[i]++;
bc4
258
825
        vector < int > ret(n);
d72
        for (int i = 0; i < m; i++) if (match[i] != -1) ret[match[i]]</pre>
   = i:
edf
        return ret;
Off }
                                                                             //
1.38 Tarjan para SCC
// O(n + m)
                                                                              //
042 vector < int > g[MAX];
4ce stack < int > s;
a42 int vis[MAX], comp[MAX];
3fd int id[MAX];
// se quiser comprimir ciclo ou achar ponte em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro pai da DFS tree
f32 int dfs(int i, int& t) {
cf0
        int lo = id[i] = t++;
18e
        s.push(i);
```

0 c 2

vis[i] = 2;

```
for (int j : g[i]) {
            if (!vis[j]) lo = min(lo, dfs(j, t));
             else if (vis[j] == 2) lo = min(lo, id[j]);
        }
        // aresta de i pro pai eh uma ponte (no caso nao direcionado)
        if (lo == id[i]) while (1) {
            int u = s.top(); s.pop();
            vis[u] = 1, comp[u] = i;
            if (u == i) break;
        }
         return lo;
f93 void tarjan(int n) {
         int t = 0:
         for (int i = 0; i < n; i++) vis[i] = 0;
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);</pre>
1.39 Top Tree
// lct<T> t(n. w) armazena uma floresta enraizada dinamica com 'n'
// vertices indexados de O a n-1 com peso w[i] do tipo T.
// Varias operacoes podem ser realizadas.
// Complexidades:
// Construcao - O(n) espaco e tempo
// Todas as outras operacoes - O(log(n)) amortizado, \Theta(n) pior
// O(n) memoria em qualquer momento
// 1. Consultas estruturais
// - find_root(v): raiz do vertice v (seguindo arestas pai)
// - lca(v, u): ancestral comum mais baixo entre v e u, ou -1 se nao
    evistir
// - connected(v, u): se v e u estao no mesmo componente conexo
// - parent(v, k): k-esimo pai de v (seguindo k arestas), ou -1 se nao
// - depth(v): numero de arestas entre v e sua raiz
// - dist(v, u): distancia entre v e u na arvore nao direcionada
// - subtree_size(v): numero de vertices na sub-arvore de v
// - adjacent(v, u): se existe uma aresta {v, u}
```

```
// 2. Atualizacoes estruturais
// - evert(v): rotaciona para que v seja a raiz de seu componente
// - link(v, u): adiciona aresta v -> u (pai de v eh u)
   . Pre-condicao: nao deve criar ciclo!
// . Se v nao eh raiz, eh feito evert(v)
// - cut(v): remove aresta de v para seu pai
// - cut(v, u): remove aresta {v, u}
   . Pre-condicao: aresta {v, u} deve existir!
// 3. Operacoes em vertices
// - t[v]: retorna w[v]
// . Pode ser usado para atualizar: t[v]++
// Para as seguintes operacoes que agregam w:
// Defina uma tag lazy associativa e comutativa para combinar
   atualizacoes em 'struct lazy'
// - Defina o que eh a tag lazy (int)
// - Defina o elemento neutro (0)
// - Defina como combinar tags lazy em 'operator+=' (+=)
// - Defina como aplicar a tag lazy em um valor em 'apply_lazy(val,
   1z) ' (+=)
//
// Defina um agregado associativo em 'struct data':
// - Defina o que serah mantido (soma/minimo/tamanho/...)
// . Manter tamanho eh obrigatorio!
// - Defina o elemento neutro (0/INF/...)
// - Defina como agregar em 'operator+=' (+=/min=/...)
// - Defina como aplicar tag lazy no agregado em 'apply_lazy(lz)'
// - Se a operacao nao for comutativa:
// . Defina como reverter um intervalo em reverse()
// 4. Operacoes em caminhos
// - path_query(v): retorna o agregado de w de v ate sua raiz
// . Nessa ordem, para operacoes nao comutativas!
// - path_query(v, u): retorna o agregado de w de v ateh u (nao
   direcionado)
// . Nessa ordem, para operacoes nao comutativas!
// - path_update(v, u, val): aplica lazy(val) em w de v ate u (nao
   direcionado)
// 5. Operacoes em sub-arvores
// - subtree_query(v, exclude_v): retorna o agregado de w na
   sub-arvore de v
// . Se exclude_v, exclui v do agregado
// . Ordem nao eh definida para operacoes nao comutativas!
```

```
// - subtree_update(v, val, exclude_v): aplica lazy(val) em w na
   sub-arvore de v
// . Se exclude_v, exclui v da atualizacao
// Para a proxima operacao que combina agregados de filhos:
// Defina uma combinação associativa e comutativa de filhos em
   'data child'
// - Defina o elemento neutro
// - Defina o valor de um filho, dado seu agregado
   (tamanho/tamanho^2/soma/...)
// - Defina como combinar valores de filhos (+=/min=/...)
// 6. Combinação de dados dos filhos
// - combine_children(v): retorna a combinacao dos filhos de v
   . Se voce definir data_child() = 1 e combinar com +=,
       returna numero de filhos
b5d template < typename T> struct lct {
07 с
        struct lazy {
f 48
            T val:
d26
            lazy(T val_ = 0) : val(val_) {}
            lazy& operator +=(const lazy& lz) {
be2
17 c
                val += lz.val;
357
                return *this:
fed
            }
2 e 2
        }:
64b
        void apply_lazy(T& val, const lazy& lz) {
17 c
            val += lz.val;
073
        }:
994
        struct data {
026
            T sum:
e4d
            int sz:
62 f
            data() : sum(0), sz(0) {}
ce5
            data(T val) : sum(val), sz(1) {}
382
            data& operator +=(const data& d) {
4b6
                sum += d.sum;
73 c
                sz += d.sz:
357
                return *this;
cbb
66a
            friend data operator +(data a, const data& b) { return a
   += b; }
c50
            void apply_lazy(const lazy& lz) {
```

```
977
                if (sz == 0) return:
2a5
                sum += lz.val * sz;
b76
            }
            void reverse() {}
ae9
5 a 9
        }:
97 c
        struct data_child {
0f9
            int cnt:
afd
            data_child() : cnt(0) {}
            data_child(data sub) { // 'sub' representa dado da
9 d 4
   sub-arvore
88d
                cnt = 1: // para numero de filhos
7ad
b76
            data_child& operator +=(const data_child& d) {
eb8
                cnt += d.cnt;
357
                return *this;
64a
            friend data_child operator +(data_child a, const
   data_child& b) { return a += b; }
52c
       }:
3 c 9
        struct node {
f48
            T val;
2fc
            int p, ch[4];
7 e 4
            data path, sub, all;
d9d
            data child data ch:
43e
            lazy lz_path, lz_sub;
533
            bool rev, fake;
373
            node(): p(-1), path(), sub(), all(), data_ch(),
   lz_path(), lz_sub(), rev(0), fake(1) {
dc5
                memset(ch, -1, sizeof ch);
            }
cf9
a94
            node(T val ) : node() {
                val = val :
d9b
                path = all = data(val);
446
                data_ch = data_child(all);
e 12
е7с
                fake = false;
168
            }
44d
        };
bbb
        vector < node > t;
08f
        vector < int > fakes;
ce0
        int free fake:
        lct(int n, vector < T > w = {}) : t(2*n), fakes(n), free_fake(n) {}
61e
9a9
            if (!w.size()) w = vector < T > (n, T()):
```

```
461
            for (int i = 0; i < n; i++) t[i] = node(w[i]);
2b9
            iota(fakes.begin(), fakes.end(), n);
1 b 0
        }
        void prop rev(int v) {
ec8
            if (v == -1) return;
501
953
            t[v].path.reverse();
b5a
            swap(t[v].ch[0], t[v].ch[1]);
            t[v].rev ^= 1;
a 02
a20
        }
f9b
        void prop_path(int v, const lazy& lz) {
53a
            if (v == -1 or t[v].fake) return;
906
            apply_lazy(t[v].val, lz);
6ff
            t[v].path.apply_lazy(lz);
1e8
            t[v].lz_path += lz;
361
            t[v].all = t[v].path + t[v].sub;
            t[v].data_ch = data_child(t[v].all);
4df
3df
        }
        void prop_sub(int v, const lazy& lz, bool virt) {
cbb
            if (v == -1) return;
501
db2
            t[v].sub.apply_lazy(lz);
8a0
            t[v].lz sub += lz:
91 b
            if (!t[v].fake and virt) prop_path(v, lz);
            else t[v].all = t[v].path + t[v].sub;
46e
58d
098
        void prop(int v) { // lazy nao inclui o vertice
501
            if (v == -1) return;
2d0
            if (t[v].rev) {
                for (int i = 0; i < 2; i++) prop_rev(t[v].ch[i]);</pre>
8 c 3
1e3
                t[v].rev = false:
ae5
dd1
            if (t[v].lz_path.val != lazy().val) {
492
                for (int i = 0; i < 2; i++) prop_path(t[v].ch[i],
   t[v].1z_path);
806
                t[v].lz_path = lazy();
048
b96
            if (t[v].lz_sub.val != lazy().val) {
                for (int i = 0; i < 4; i++) prop_sub(t[v].ch[i],</pre>
6a5
   t[v].lz_sub, i > 1);
65 e
                t[v].lz_sub = lazy();
aa2
            }
        }
db4
437
        node get_ch(int v, int c) {
```

```
17f
            return t[v].ch[c] == -1 ? node() : t[t[v].ch[c]];
                                                                            93 d
                                                                            e 07
                                                                                    void add_virt_ch(int v, int u) {
ee8
31f
        void update(int v) {
                                                                           61 c
                                                                                        if (u == -1) return;
                                                                                        for (int i = 2; i < 4; i++) if (t[v].ch[i] == -1) {</pre>
            if (!t[v].fake) t[v].path = get_ch(v, 1).path +
                                                                           b51
   data(t[v].val) + get_ch(v, 0).path;
                                                                            dab
                                                                                            attach(v. i. u):
            t[v].sub = get_ch(v, 0).sub + get_ch(v, 1).sub + get_ch(v,
                                                                            505
                                                                                            return:
                                                                                        }
   2).all + get ch(v, 3).all:
                                                                            019
            t[v].all = t[v].path + t[v].sub;
                                                                            680
                                                                                        int w = fakes[--free_fake];
361
                                                                            782
                                                                                        attach(w, 2, t[v].ch[2]);
f06
            if (t[v].fake) t[v].data_ch = get_ch(v, 2).data_ch +
                                                                                        attach(w, 3, u), attach(v, 2, w);
                                                                            6ef
   get_ch(v, 3).data_ch;
                                                                            999
                                                                                   }
bc2
            else t[v].data ch = data child(t[v].all);
                                                                            e 26
                                                                                    void prop_path(int v) {
e3b
                                                                            9b1
                                                                                        if (t[v].fake) prop_path(t[v].p);
                                                                            5 e 3
                                                                                        prop(v);
        int dir(int v) {
                                                                                   }
8b9
                                                                            a98
            if (t[v].p == -1) return -1;
4cb
                                                                            4f3
                                                                                    void rem_virt_ch(int v) {
c04
            for (int i = 0; i < 4; i++) if (t[t[v].p].ch[i] == v)
                                                                            672
                                                                                        int p = t[v].p;
                                                                            038
                                                                                        prop_path(p);
   return i:
            return -1;
                                                                            656
                                                                                        if (t[p].fake) {
daa
66b
                                                                            676
                                                                                            int pp = t[p].p;
                                                                                            attach(pp, dir(p), t[p].ch[dir(v)^1]);
679
        bool is root(int v) {
                                                                           f 2.3
252
            if (t[v].p == -1) return true;
                                                                            7 e 2
                                                                                            if (t[pp].fake) splay(pp);
            return t[v].fake ? !t[t[v].p].fake : (dir(v)&2);
                                                                           ffe
                                                                                            fakes[free_fake++] = p;
caf
287
        }
                                                                           df2
                                                                                        } else attach(p, dir(v), -1);
a5e
        void rotate(int v) {
                                                                            07f
                                                                                        t[v].p = -1;
078
            int p = t[v].p, pp = t[p].p, dir_v = dir(v), dir_p =
                                                                            a88
                                                                            7f6
   dir(p);
                                                                                    int real_par(int v) {
                                                                            31a
                                                                                        v = t[v].p:
                                                                                        if (!t[v].fake) return v;
8b6
            attach(p, dir_v, t[v].ch[dir_v^1]);
                                                                            ae6
a67
            attach(v, dir_v^1, p);
                                                                           bf5
                                                                                        splay(v);
6 c 7
            if (dir_p+1) attach(pp, dir_p, v);
                                                                            fe7
                                                                                        return t[v].p;
72e
            else t[v].p = pp;
                                                                           973
8 c 5
        }
                                                                           f 16
                                                                                    int access(int v) { // retorna ultimo vertice acessado
                                                                            35 b
ec1
        void splay(int v) {
                                                                                        update(v):
5 e 3
            prop(v);
                                                                           bf5
                                                                                        splay(v);
            while (!is_root(v)) {
                                                                                        add_virt_ch(v, t[v].ch[1]);
7 d3
                                                                            c4b
931
                int p = t[v].p, pp = t[p].p;
                                                                           b49
                                                                                        attach(v, 1, -1);
bf6
                if (pp+1) prop(pp);
                                                                                        int w = v;
                                                                            aee
a51
                prop(p), prop(v);
                                                                            7 c 5
                                                                                        while (t[v].p+1) {
c31
                if (!is_root(p)) rotate(dir(v) == dir(p) ? p : v);
                                                                            425
                                                                                            w = real_par(v);
e61
                rotate(v);
                                                                            50a
                                                                                            splay(w);
5 d2
            }
bed
                                                                            323
                                                                                            rem_virt_ch(v);
29f
        void attach(int v, int c, int u) {
                                                                           547
                                                                                            add_virt_ch(w, t[w].ch[1]);
908
            t[v].ch[c] = u;
                                                                           b4e
                                                                                            attach(w, 1, v);
c99
            if (u+1) t[u].p = v;
35b
            update(v):
                                                                           bf5
                                                                                            splay(v);
```

```
da0
df 1
            return w:
e96
        }
        // consultas estruturais
        int find root(int v) {
e89
5 e 3
            access(v):
9f0
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
f05
            return splay(v), v;
a26
c5d
        int lca(int v, int u) {
b77
            int ret = (access(v), access(u));
a2a
            if (v != u \text{ and } t[v]. p == -1) \text{ return } -1;
edf
            return ret:
872
05f
        bool connected (int v. int u) {
d16
            return lca(v, u) != -1;
1 a 1
ba2
        int parent(int v, int k = 1) {
5 e 3
            access(v):
            if (k >= t[v].path.sz) return -1;
217
667
            while (true) {
                 int sz_r = t[v].ch[1] == -1 ? 0 :
c0f
   t[t[v].ch[1]].path.sz;
                if (k < sz_r) v = t[v].ch[1], prop(v);</pre>
c4a
                 else {
4 e 6
733
                     k -= sz r:
                     if (k == 0) return splay(v), v;
3fd
01b
                     v = t[v].ch[0], prop(v);
9bc
                    k --;
acf
                }
077
            }
67d
        int depth(int v) {
baa
5 e 3
            access(v):
192
            return t[v].path.sz - 1;
ec2
419
        int dist(int v. int u) {
de8
            return depth(v) + depth(u) - 2*depth(lca(v, u));
1f5
363
        int subtree_size(int v) {
5 e 3
            access(v):
            return get_ch(v, 2).all.sz + get_ch(v, 3).all.sz + 1;
f37
eba
417
        bool adjacent(int v, int u) {
            return connected(v, u) and dist(v, u) == 1;
54a
```

```
eb0
        }
        // atualizacoes estruturais
        void evert(int v) { // rotaciona para que v seja a raiz
adb
5 e 3
            access(v):
            prop_rev(v);
533
        }
bd9
        void link(int v, int u) { // adiciona aresta v -> u (pai de v
7af
   eh u) - nao pode criar ciclo!
5 c 2
            evert(v). access(u):
7b7
            add_virt_ch(u, v);
d4a
4e6
        void cut(int v) { // corta aresta de v para o pai
5 e 3
            access(v):
            if (t[v].ch[0]+1) t[t[v].ch[0]].p = -1;
b1d
5ff
            t[v].ch[0] = -1;
35 b
            update(v);
8ad
        void cut(int v, int u) { // aresta {v, u} deve existir!
af7
            cut(depth(v) > depth(u) ? v : u);
8 c 6
        }
416
        // operacao de vertice
5f1
        T& operator[](int v) { // w[v] (pode ser alterado)
5 e 3
            access(v):
337
            return t[v].val:
492
        }
        // operacoes de caminho - vertices devem estar no mesmo
            componente!
7 f 4
        data path_query(int v) { // de v para a raiz
5 e 3
            access(v);
            return t[v].path:
1bb
2 c 5
        data path_query(int v, int u) { // de v para u
76 d
            int rt = find_root(v);
db2
708
            evert(u), access(v);
b4e
            auto ret = t[v].path;
e70
            evert(rt):
edf
            return ret;
267
        void path_update(int v, int u, T val) {
ee6
db2
            int rt = find root(v):
708
            evert(u), access(v);
715
            prop_path(v, lazy(val));
e70
            evert(rt):
cfe
        }
```

```
// operacoes de sub-arvore
        data subtree_query(int v, bool exclude_v = false) {
acc
            access(v):
5 e 3
            data ret = get_ch(v, 2).all + get_ch(v, 3).all;
ffe
            if (!exclude_v) ret += data(t[v].val);
827
edf
            return ret:
696
        }
56f
        void subtree_update(int v, T val, bool exclude_v = false) {
5 e 3
            access(v):
207
            prop_sub(t[v].ch[2], lazy(val), true);
bb9
            prop_sub(t[v].ch[3], lazy(val), true);
bf9
            if (!exclude_v) apply_lazy(t[v].val, lazy(val));
7b9
       }
        // combinação de dados das sub-arvores dos filhos
26e
        data_child combine_children(int v) {
5 e 3
            access(v):
            return get_ch(v, 2).data_ch + get_ch(v, 3).data_ch;
db4
303
       }
a22 };
```

### 1.40 Topological Sort

```
// Retorna uma ordenacaoo topologica de g
// Se g nao for DAG retorna um vetor vazio
// O(n + m)
042 vector < int > g[MAX];
b6a vector<int> topo_sort(int n) {
46e
        vector < int > ret(n, -1), vis(n, 0);
        int pos = n-1, dag = 1;
f51
        function < void(int) > dfs = [&](int v) {
36d
            vis[v] = 1;
cca
440
            for (auto u : g[v]) {
                if (vis[u] == 1) dag = 0;
152
532
                else if (!vis[u]) dfs(u);
e37
            ret[pos--] = v, vis[v] = 2;
d44
        };
57e
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);
```

#### 1.41 Unrooted Euler Tour Tree

```
// Mantem uma floresta dinamicamente
// e permite queries/updates em componentes
//
// Chamar ETT E(n, w), passando n = numero de vertices
// e w = vector com os valores de cada vertice
//
// link(v. u) cria uma aresta entre v e u
// cut(v, u) corta aresta entre v e u
// get_root(v) retorna a raiz de v
// branc_query(v, u) e branch_update(v, u) operam sobre o componente
// conexo de v se arestas {v, u} fosse removida do grafo
// comp_query(v) e comp_update(v, val) operam sobre o compoente conexo
// point_update(v, val) atualiza o valor do vertice v
// point_update(v, val, u) atualiza o valor da aresta {v, u}
//
// Tudo O(log(n)) com alta probabilidade
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
9c9 template < typename T> struct unrooted_ett {
        const static T ZERO = T();
83 c
3 c 9
        struct node {
ed1
            node *1, *r, *p;
1 c 1
            bool is_vtx;
eba
            int pr, sz, vtx_cnt;
875
            T val, sub, lazy;
            node(T v, bool is_vtx_ = false) : 1(NULL), r(NULL),
b5f
    p(NULL), is_vtx(is_vtx_),
                pr(rng()), sz(1), vtx_cnt(is_vtx), val(v), sub(v),
4ef
   lazy(ZERO) {}
a9c
            void prop() {
9fc
                if (lazy != ZERO) {
094
                    if (is_vtx) val += lazy;
489
                    sub += lazy * vtx_cnt;
b87
                    if (1) 1->lazy += lazy;
d3b
                    if (r) r->lazy += lazy;
f2b
                }
520
                lazy = ZER0;
```

```
d34
01e
             void update() {
1eb
                 sz = 1, vtx_cnt = is_vtx, sub = val;
                 for (auto& i : {1, r}) if (i) {
ac2
c89
                     i->prop();
                      sz += i->sz, vtx_cnt += i->vtx_cnt;
fe4
ea7
                      sub += i-> sub:
                 }
ac7
3 e 7
eaf
        };
73c
        int size(node* x) { return x ? x->sz : 0; }
bcf
        void join(node* 1. node* r. node*& i) {
986
             if (!1 or !r) return void(i = 1 ? 1 : r);
161
            1->prop(), r->prop();
ff5
             if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
982
             else join(1, r - > 1, r - > 1), r - > 1 - > p = i = r;
bda
             i->update();
84d
a20
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
             if (!i) return void(r = 1 = NULL);
26a
c89
             i->prop():
             if (key + size(i->1) < v) {
d9e
448
                 split(i\rightarrow r, i\rightarrow r, r, v, kev+size(i\rightarrow l)+1), l = i;
a 2 1
                 if (r) r -> p = NULL;
6e8
                 if (i->r) i->r->p = i;
396
             } else {
98d
                 split(i\rightarrow 1, 1, i\rightarrow 1, v, key), r = i;
5 a 3
                 if (1) 1->p = NULL;
899
                 if (i->1) i->1->p = i;
18b
bda
             i->update();
134
3fa
        pair < node *, int > get_idx(node * i) { // {root, idx}}
             int ret = size(i->1);
6cf
             for (; i->p; i = i->p) if (i != i->p->1)
65 d
fc8
                 ret += size(i->p->1) + 1;
c89
             i->prop();
e51
             return {i, ret};
fa4
        }
        // as treaps sao disjuntas!
        map < pair < int , int > , node *> mp;
Зсс
        unrooted_ett(int n, const vector < T > & w) { // w = valor de cada
72a
   vertice
98f
             for (int i = 0; i < n; i++) mp[pair(i, i)] = new</pre>
```

```
node(w[i], true);
045
        unrooted_ett(const unrooted_ett& t) { throw logic_error("Nao
8d2
   copiar a ETT!"); }
057
        ~unrooted_ett() { for (auto [__, ptr] : mp) delete ptr; }
57a
        void link(int v. int u. T val = ZERO) {
            auto [root_v, idx_v] = get_idx(mp[pair(v, v)]);
0e6
a60
            node *L_v, *R_v;
f8c
            split(root_v, L_v, R_v, idx_v);
f7a
            join(R_v, L_v, root_v);
a2b
            auto [root, idx] = get_idx(mp[pair(u, u)]);
031
            node *L. *R:
680
            split(root, L, R, idx);
1 d 6
            join(L, mp[pair(u, v)] = new node(val), L); // w[(u, v)]
7ce
            join(L, root_v, L);
37b
            join(L, mp[pair(v, u)] = new node(ZERO), L); // w[(v, u)]
0 c 7
            join(L, R, R);
        }
1b4
        node* cut(int v, int u) { // retorna a raiz de v
85 с
4 c 2
            auto [root, 1] = get_idx(mp[pair(u, v)]);
4a4
            int r = get_idx(mp[pair(v, u)]).second;
541
            bool rev = false;
19 f
            if (1 > r) swap(1, r), rev = true;
bae
            node *L, *E1, *M, *E2, *R;
сбе
            split(root, E2, R, r+1);
a8c
            split(E2, M, E2, r);
b7a
            split(M, E1, M, 1+1);
            split(E1, L, E1, 1);
1aa
739
            delete E1:
            delete E2:
ef6
            mp.erase(pair(v, u));
ac8
717
            mp.erase(pair(u, v));
0 c 7
            join(L, R, R);
b31
            return rev ? R : M;
        }
ca8
        // opera sobre component(v) em F \ {v, u}
13f
        T branch_query(int v, int u) {
4da
            auto root = cut(v, u);
e 38
            T ret = root->sub;
32 e
            link(v. u):
```

```
edf
            return ret;
0 e 3
2bf
        void branch_update(int v, int u, T val) {
4da
            auto root = cut(v, u);
fc6
            root -> lazy += val;
32e
            link(v, u);
827
       }
        // opera sobre component(v)
        T comp_query(int v) {
b3d
da0
            auto root = get_idx(mp[pair(v, v)]).first;
6f0
            return root -> sub;
4fd
623
        void comp_update(int v, T val) {
da0
            auto root = get_idx(mp[pair(v, v)]).first;
fc6
            root -> lazy += val;
0 e 4
       }
        // w[v] += val, ou w[(v, u)] += val
        void point_update(int v, T val, int u = -1) {
e21
            if (u == -1) u = v;
e4a
            auto [root, idx] = get_idx(mp[pair(v, u)]);
ca9
df9
            node *L, *M, *R;
d26
            split(root, M, R, idx+1);
70b
            split(M, L, M, idx);
409
            M->lazy += val;
675
            join(L, M, M);
087
            join(M, R, R);
3f0
       }
9b3 };
1.42 Vertex cover
// Encontra o tamanho do vertex cover minimo
// Da pra alterar facil pra achar os vertices
// Parece rodar com < 2 s pra N = 90
// 0(n * 1.38^n)
76a namespace cover {
```

5 a 4

042

823

const int MAX = 96;

vector < int > g[MAX];

bitset < MAX > bs [MAX];

```
1a8
        int n;
        void add(int i, int j) {
697
bd0
             if (i == j) return;
78 c
             n = \max(\{n, i+1, j+1\});
200
             bs[i][j] = bs[j][i] = 1;
        }
203
6c0
        int rec(bitset < MAX > m) {
             int ans = 0:
1 a 4
25 b
             for (int x = 0; x < n; x++) if (m[x]) {
002
                 bitset < MAX > comp;
4bf
                 function < void(int) > dfs = [&](int i) {
b96
                     comp[i] = 1, m[i] = 0;
0 c 3
                     for (int j : g[i]) if (m[j]) dfs(j);
815
                 };
963
                 dfs(x);
d34
                 int ma, deg = -1, cyc = 1;
                 for (int i = 0; i < n; i++) if (comp[i]) {</pre>
417
                     int d = (bs[i]&comp).count();
d0b
18a
                     if (d <= 1) cvc = 0;
                     if (d > deg) deg = d, ma = i;
c1f
d8e
                 }
269
                 if (deg <= 2) { // caminho ou ciclo</pre>
340
                     ans += (comp.count() + cyc) / 2;
5 e 2
                     continue;
702
3f9
                 comp[ma] = 0;
                 // ou ta no cover, ou nao ta no cover
1dd
                 ans += \min(1 + rec(comp), deg + rec(comp & \sim bs[ma]));
6e6
            }
ba7
             return ans:
2ec
f5c
        int solve() {
3 c 5
             bitset < MAX > m;
603
             for (int i = 0; i < n; i++) {</pre>
939
                 m[i] = 1:
f90
                 for (int j = 0; j < n; j++)
741
                     if (bs[i][j]) g[i].push_back(j);
13e
4f9
             return rec(m);
708
        }
9c5 }
```

#### 1.43 Virtual Tree

```
// Comprime uma arvore dado um conjunto S de vertices, de forma que
// o conjunto de vertices da arvore comprimida contenha S e seja
// minimal e fechado sobre a operacao de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o vizinho
// Retorna a raiz da virtual tree
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
// O(k log(k))
b36 vector <pair <int, int>> virt[MAX];
d41 #warning lembrar de buildar o LCA antes
c14 int build virt(vector<int> v) {
        auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
   lca::pos[j]; };
074
        sort(v.begin(), v.end(), cmp);
        for (int i = v.size()-1; i; i--) v.push_back(lca::lca(v[i],
   v[i-1]));
074
        sort(v.begin(), v.end(), cmp);
d76
        v.erase(unique(v.begin(), v.end()), v.end());
37 c
        for (int i = 0; i < v.size(); i++) virt[v[i]].clear();</pre>
        for (int i = 1; i < v.size(); i++) virt[lca::lca(v[i-1],</pre>
197
   v[i])].clear();
        for (int i = 1; i < v.size(); i++) {</pre>
ad7
            int parent = lca::lca(v[i-1], v[i]);
51b
290
            int d = lca::dist(parent, v[i]);
d41 #warning soh to colocando aresta descendo
4d0
            virt[parent].emplace_back(v[i], d);
fe5
832
        return v[0];
142 }
```

- 2 math
- 3 geometry
- 4 Matematica

#### 4.1 2-SAT

```
// solve() retorna um par, o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
// O(|V|+|E|) = O(\#variaveis + \#restricoes)
138 struct sat {
e6c
         int n, tot;
789
         vector < vector < int >> g;
0ca
         vector<int> vis, comp, id, ans;
4ce
         stack < int > s:
141
         sat() {}
         sat(int n_{-}) : n(n_{-}), tot(n), g(2*n) {}
172
f32
         int dfs(int i, int& t) {
cf0
             int lo = id[i] = t++:
efc
             s.push(i), vis[i] = 2;
             for (int j : g[i]) {
48e
740
                 if (!vis[j]) lo = min(lo, dfs(j, t));
994
                 else if (vis[j] == 2) lo = min(lo, id[j]);
d64
3de
             if (lo == id[i]) while (1) {
3 c 3
                 int u = s.top(); s.pop();
9 c 5
                 vis[u] = 1, comp[u] = i;
                 if ((u>1) < n \text{ and } ans[u>1] == -1) ans[u>1] = \sim u & 1;
91 d
2ef
                 if (u == i) break;
            }
60 d
253
             return lo;
        }
dec
74a
         void add_impl(int x, int y) { // x -> y = !x ou y
26a
             x = x >= 0 ? 2*x : -2*x-1;
2b8
             v = v >= 0 ? 2*v : -2*v-1;
             g[x].push_back(y);
a1e
1 e 2
             g[y^1].push_back(x^1);
ef0
        }
```

```
e85
         void add_cl(int x, int y) { // x ou y
0 b 5
             add_impl(\simx, y);
254
487
         void add_xor(int x, int y) { // x xor y
0b7
             add_cl(x, y), add_cl(\simx, \simy);
9a1
978
         void add_eq(int x, int y) { // x = y
c86
             add_xor(\simx, y);
b91
         void add_true(int x) { // x = T
b10
18b
             add_impl(\sim x, x);
9 e 2
d14
         void at most one(vector<int> v) { // no max um verdadeiro
54d
             g.resize(2*(tot+v.size()));
             for (int i = 0; i < v.size(); i++) {</pre>
f14
8 c 9
                 add_impl(tot+i, \simv[i]);
                 if (i) {
a8f
                      add_impl(tot+i, tot+i-1);
b6a
                      add_impl(v[i], tot+i-1);
3 d3
                 }
0f7
084
258
             tot += v.size();
        }
b00
        pair < bool, vector < int >> solve() {
a8e
27b
             ans = vector < int > (n, -1);
6bb
             int t = 0:
             vis = comp = id = vector \langle int \rangle (2*tot, 0);
0 de
             for (int i = 0; i < 2*tot; i++) if (!vis[i]) dfs(i, t);</pre>
53c
f88
             for (int i = 0; i < tot; i++)</pre>
                 if (comp[2*i] == comp[2*i+1]) return {false, {}};
4 c.9
997
             return {true, ans};
7b3
        }
ef6 }:
4.2 Avaliação de Interpolação
// Dado 'n' pontos (i, v[i]), i \in [0, n),
```

```
// Dado 'n' pontos (i, y[i]), i \in [0, n),
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x'
// Tudo modular, precisa do mint
//
// O(n)

ee8 mint evaluate_interpolation(int x, vector<mint> y) {
80e    int n = y.size();
```

```
vector \langle mint \rangle sulf (n+1, 1), fat (n, 1), if at (n);
184
        for (int i = n-1; i >= 0; i--) sulf[i] = sulf[i+1] * (x - i);
6fa
29h
        for (int i = 1; i < n; i++) fat[i] = fat[i-1] * i;</pre>
        ifat[n-1] = 1/fat[n-1]:
0da
3db
        for (int i = n-2; i >= 0; i--) ifat[i] = ifat[i+1] * (i + 1);
ca1
        mint pref = 1, ans = 0;
5ea
        for (int i = 0; i < n; pref *= (x - i++)) {
            mint num = pref * sulf[i+1];
42 f
b4e
            mint den = ifat[i] * ifat[n-1 - i];
0bd
            if ((n-1 - i)\%2) den *= -1;
03f
            ans += y[i] * num * den;
        }
ce6
ba7
        return ans;
4fe }
4.3 Berlekamp-Massey
// guess_kth(s, k) chuta o k-esimo (0-based) termo
// de uma recorrencia linear que gera s
// Para uma rec. lin. de ordem x, se passar 2x termos
// vai gerar a certa
// Usar aritmetica modular
// Pro fast_evaluate, precisa de ntt e divmod (powerSeries.cpp)
//
// Complexidades: (n = |s|)
// evaluate: 0(n^2 log k)
// fast_evaluate: O(n log n log k)
// berlekampMassey: 0(n^2 + 0(evaluate))
b7c template < typename T> T evaluate (vector < T> c, vector < T> s, 11 k) {
ff2
        int n = c.size();
9ee
        assert(c.size() <= s.size());
```

auto mul =  $\lceil \& \rceil \pmod{\text{const vector}} \& a$ , const vector< T > & b) {

for (int i = 0; i < a.size(); i++) for (int j = 0; j <

for (int i = ret.size()-1; i >= n; i--) for (int j = n-1;

vector <T> ret(a.size() + b.size() - 1);

ret[i+j] += a[i] \* b[j];

ret[i-j-1] += ret[i] \* c[j];

d09

564

d75

cff

112

b.size(); j++)

i >= 0; i--)

```
16d
            ret.resize(min<int>(ret.size(), n));
edf
            return ret:
3 b 9
        };
        vector < T > a = n == 1 ? vector < T > ({c[0]}) : vector < T > ({0, 1}),
1a6
   x = \{1\}:
95f
        while (k) {
            if (k\&1) x = mul(x, a);
7 f 1
b28
            a = mul(a, a), k >>= 1;
8ea
dd6
        x.resize(n);
ce8
        T ret = 0:
        for (int i = 0; i < n; i++) ret += x[i] * s[i];</pre>
e72
edf
        return ret;
7e2 }
83e mint fast_evaluate(poly c, poly s, ll k) {
        if (k < s.size()) return s[k];</pre>
fb0
        int n = c.size();
ff2
        assert(c.size() <= s.size());</pre>
9ee
79f
        auto f = poly(n + 1, 1);
bcc
        for (int i = 0; i < n; i++) f[i] = -c[n-i-1];
7 ce
        poly a = n == 1 ? poly(\{c[0]\}) : poly(\{0, 1\}), x = \{1\};
        while (k) {
95f
3df
            if (k&1) x = divmod(convolution(x, a), f).second:
            a = divmod(convolution(a, a), f).second, k >>= 1;
cac
        }
11b
eb1
        mint ret = 0;
        for (int i = 0; i < n; i++) ret += x[i] * s[i];
e72
edf
        return ret:
95d }
192 template < typename T > vector < T > berlekamp_massey(vector < T > s) {
ce8
        int n = s.size(), l = 0, m = 1;
222
        vector < T > b(n). c(n):
46e
        T \ 1d = b[0] = c[0] = 1;
        for (int i = 0; i < n; i++, m++) {
620
793
            T d = s[i];
            for (int j = 1; j <= 1; j++) d += c[j] * s[i-j];</pre>
ab6
            if (d == 0) continue;
5f0
8b4
            vector < T > temp = c;
            T coef = d / ld:
369
ba6
            for (int j = m; j < n; j++) c[j] -= coef * b[j-m];
```

```
88f
            if (2 * 1 \le i) 1 = i + 1 - 1, b = temp, 1d = d, m = 0;
76 a
90 с
        c.resize(1 + 1);
844
        c.erase(c.begin());
        for (T\& x : c) x = -x:
0dc
807
        return c:
4d9 }
2cf template < typename T> T guess_kth(const vector < T>& s, ll k) {
        auto c = berlekamp_massey(s);
96a
        return evaluate(c, s, k);
697 }
4.4 Binomial Distribution
```

```
// binom(n, k, p) retorna a probabilidade de k sucessos
// numa binomial(n, p)
361 double logfact[MAX];
9e4 void calc() {
7a0
        logfact[0] = 0;
152
        for (int i = 1; i < MAX; i++) logfact[i] = logfact[i-1] +</pre>
   log(i);
67a }
94c double binom(int n, int k, double p) {
        return exp(logfact[n] - logfact[k] - logfact[n-k] + k * log(p)
   + (n-k) * log(1 - p));
587 }
```

# 4.5 Convolucao de GCD / LCM

```
// O(n log(n))
// multiple_transform(a)[i] = \sum_d a[d * i]
bbe template < typename T> void multiple_transform(vector < T>& v, bool
   inv = false) {
64a
        vector < int > I(v.size()-1);
        iota(I.begin(), I.end(), 1);
847
674
        if (inv) reverse(I.begin(), I.end());
dad
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
a8a
            v[i] += (inv ? -1 : 1) * v[i*i];
338 }
```

```
// \gcd_{convolution(a, b)[k]} = \sum_{gcd(i, j)} = k} a_i * b_j
fe2 template < typename T > vector < T > gcd_convolution(vector < T > a,
   vector <T> b) {
        multiple_transform(a), multiple_transform(b);
bdf
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
dea
        multiple_transform(a, true);
3f5
        return a;
984 }
// divisor_transform(a)[i] = \sum_{d|i} a[i/d]
be7 template < typename T > void divisor_transform(vector < T > & v, bool inv
   = false) {
64a
        vector < int > I(v.size()-1):
847
        iota(I.begin(), I.end(), 1);
        if (!inv) reverse(I.begin(), I.end());
5ea
dad
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
            v[i*j] += (inv ? -1 : 1) * v[i];
14f
aa7 }
// lcm_convolution(a, b)[k] = \sum_{lcm(i, j)} = k} a_i * b_j
b1b template < typename T > vector < T > lcm_convolution (vector < T > a,
   vector <T> b) {
        divisor_transform(a), divisor_transform(b);
3af
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
d8f
        divisor_transform(a, true);
3f5
        return a:
f5a }
    Coprime Basis
// Dado um conjunto de elementos A constroi uma base B
// de fatores coprimos tal que todo elemento A[i]
```

```
// pode ser fatorado como A[i] = \prod B[j]^p_ij
//
// Sendo n o numero de inserts, a complexidade esperada fica
// O(n*(n*loglog(MAX) + log(MAX)^2))
//
// No pior caso, podemos trocar n*loglog(MAX) por
// 8n, se MAX <= 1e6
// 10n, se MAX <= 1e9
// 16n. se MAX <= 1e18
// 26n, se MAX <= 1e36
ebc template <typename T> struct coprime_basis {
a00
        vector <T> basis:
```

```
60 e
        coprime_basis() {}
        coprime_basis(vector<T> v) { for (T i : v) insert(i); }
055
        void insert(T z) {
845
             int n = basis.size();
сЗс
efe
             basis.push_back(z);
43 с
             for (int i = n; i < basis.size(); i++) {</pre>
21 c
                 for (int j = (i != n) ? i+1 : 0; j < basis.size();</pre>
   j++) {
4се
                     if (i == j) continue;
                     T &x = basis[i]:
024
                     if (x == 1) {
c91
fac
                          i = INF;
5 e 2
                          continue;
                     }
6e0
544
                     T &v = basis[i];
                     T g = gcd(x, y);
3c9
                     if (g == 1) continue;
e 10
                     y /= g, x /= g;
15 b
8 c 6
                     basis.push_back(g);
069
                 }
422
             }
             basis.erase(remove(basis.begin(), basis.end(), 1),
fe8
    basis.end());
1a5
        }
4ba
        vector<int> factor(T x) {
21 d
             vector < int > fat(basis.size()):
6fd
             for (int i = 0; i < basis.size(); i++) {</pre>
25 c
                 while (x % basis[i] == 0) x /= basis[i], fat[i]++;
8de
            }
6a7
             return fat:
b5d
        }
671 };
4.7 Crivo de Eratosthenes
// "O" crivo
```

```
// Encontra maior divisor primo
// Um numero eh primo sse divi[x] == x
// fact fatora um numero <= lim
// A fatoração sai ordenada
//
// \text{ crivo } - O(n \log(\log(n)))
```

```
// fact - O(log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
        for (int i = 2; i <= lim; i++) if (divi[i] == 1)
d46
018
            for (int j = i; j <= lim; j += i) divi[j] = i;
349 }
d41 #warning A funcao fact ira adicionar o 1 no vetor se voce tentar
   fatorar especificamente o numero 1
470 void fact(vector<int>& v, int n) {
        if (n != divi[n]) fact(v, n/divi[n]);
        v.push_back(divi[n]);
ab4
1db }
// Crivo linear
// Mesma coisa que o de cima, mas tambem
// calcula a lista de primos
//
// O(n)
f12 int divi[MAX]:
fd3 vector < int > primes;
fb9 void crivo(int lim) {
d5a
        divi[1] = 1;
f70
        for (int i = 2; i <= lim; i++) {</pre>
3eb
            if (divi[i] == 0) divi[i] = i, primes.push_back(i);
3ba
            for (int j : primes) {
522
                if (j > divi[i] or i*j > lim) break;
00b
                divi[i*j] = j;
            }
491
85a
        }
519 }
// Crivo de divisores
// Encontra numero de divisores
// ou soma dos divisores
//
// O(n log(n))
f12 int divi[MAX];
```

```
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
f53
        for (int i = 2: i <= lim: i++)
424
594
            for (int j = i; j <= lim; j += i) {</pre>
                // para numero de divisores
                 divi[j]++;
9e0
                // para soma dos divisores
278
                divi[j] += i;
c58
            }
fc1 }
// Crivo de totiente
// Encontra o valor da função
// totiente de Euler
//
// O(n log(log(n)))
5f4 int tot[MAX];
fb9 void crivo(int lim) {
a27
        for (int i = 1; i <= lim; i++) {</pre>
bc9
            tot[i] += i:
feb
            for (int j = 2*i; j <= lim; j += i)</pre>
837
                 tot[i] -= tot[i];
678
        }
212 }
// Crivo de função de mobius
//
// O(n log(log(n)))
4e1 char meb[MAX];
fb9 void crivo(int lim) {
649
        for (int i = 2; i <= lim; i++) meb[i] = 2;
        meb[1] = 1:
ace
842
        for (int i = 2; i <= lim; i++) if (meb[i] == 2)
8b8
            for (int j = i; j <= lim; j += i) if (meb[j]) {</pre>
686
                if (meb[j] == 2) meb[j] = 1;
                 meb[j] *= j/i\%i ? -1 : 0;
ae1
            }
97 f
9bc }
// Crivo linear de funcao multiplicativa
```

```
// Computa f(i) para todo 1 <= i <= n, sendo f
// uma funcao multiplicativa (se gcd(a,b) = 1,
// entao f(a*b) = f(a)*f(b)
// f_prime tem que computar f de um primo, e
// add_prime tem que computar f(p^{(k+1)}) dado f(p^k) e p
// Se quiser computar f(p^k) dado p e k, usar os comentarios
//
// O(n)
fd3 vector < int > primes;
623 int f[MAX], pot[MAX];
//int expo[MAX];
5c4 void sieve(int lim) {
        // Funcoes para soma dos divisores:
fc9
        auto f_prime = [](int p) { return p+1; };
31c
        auto add_prime = [](int fpak, int p) { return fpak*p+1; };
        //auto f_pak = [](int p, int k) {};
02d
        f[1] = 1;
f70
        for (int i = 2; i <= lim; i++) {</pre>
e6b
            if (!pot[i]) {
e74
                primes.push_back(i);
f05
                f[i] = f_prime(i), pot[i] = i;
                //\exp [i] = 1;
            }
b71
3 b 9
            for (int p : primes) {
                if (i*p > lim) break;
b9f
569
                if (i\%p == 0) {
                    f[i*p] = f[i / pot[i]] * add_prime(f[pot[i]], p);
b97
                    // se for descomentar, tirar a linha de cima tambem
                    //f[i*p] = f[i / pot[i]] * f_pak(p, expo[i]+1);
                    //\exp(i*p) = \exp(i)+1;
                    pot[i*p] = pot[i] * p;
51f
c2b
                    break;
643
                } else {
9ef
                    f[i*p] = f[i] * f[p];
638
                    pot[i*p] = p;
                    //\exp [i*p] = 1;
                }
6f7
f31
            }
1 bb
        }
350 }
```

#### 4.8 Deteccao de ciclo - Tortoise and Hare

```
// Linear no tanto que tem que andar pra ciclar,
// 0(1) de memoria
// Retorna um par com o tanto que tem que andar
// do fO ate o inicio do ciclo e o tam do ciclo
58d pair<11, 11> find_cycle() {
        11 \text{ tort} = f(f0):
273
b2b
        ll hare = f(f(f0));
b1b
        11 t = 0:
683
        while (tort != hare) {
b4d
            tort = f(tort):
4b2
            hare = f(f(hare));
c82
            t++;
        }
93 d
0e8
        11 \text{ st} = 0:
        tort = f0;
909
683
        while (tort != hare) {
            tort = f(tort);
b4d
1a2
            hare = f(hare);
397
            st++;
c91
        }
73 d
        ll len = 1;
        hare = f(tort);
3cd
683
        while (tort != hare) {
1a2
            hare = f(hare);
040
            len++;
f1a
ebd
        return {st, len};
899 }
```

#### 4.9 Division Trick

#### 4.10 Equação Diofantina Linear

```
// Encontra o numero de solucoes de a*x + b*y = c,
// em que x \in [lx, rx] e y \in [ly, ry]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
// O(log(min(a, b)))
c5e template < typename T > tuple < ll, T, T > ext_gcd(ll a, ll b) {
        if (!a) return {b, 0, 1};
3bd
c4b
        auto [g, x, y] = ext_gcd < T > (b%a, a);
c59
        return \{g, y - b/a*x, x\};
8a8 }
// numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
14c template < typename T = 11> // usar __int128 se for ate 1e18
2a4 ll diophantine(ll a, ll b, ll c, ll lx, ll rx, ll ly, ll ry) {
        if (1x > rx \text{ or } 1y > ry) return 0;
        if (a == 0 and b == 0) return c ? 0 : (rx-lx+1)*(ry-ly+1);
a98
        auto [g, x, y] = ext_gcd < T > (abs(a), abs(b));
8се
9 c 3
        if (c % g != 0) return 0;
        if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <= ry);
249
4се
        if (b == 0) return (ry-ly+1)*(lx <= c/a and c/a <= rx);
fb1
        x *= a/abs(a) * c/g, y *= b/abs(b) * c/g, a /= g, b /= g;
        auto shift = [\&](T qt) \{ x += qt*b, y -= qt*a; \};
b20
        auto test = [&](T& k, ll mi, ll ma, ll coef, int t) {
efa
866
            shift((mi - k)*t / coef);
            if (k < mi) shift(coef > 0 ? t : -t);
79d
74d
            if (k > ma) return pair < T, T > (rx+2, rx+1);
41f
            T x1 = x;
633
            shift((ma - k)*t / coef);
            if (k > ma) shift (coef > 0 ? -t : t);
c5b
4a9
            return pair <T, T>(x1, x);
8 e 1
        };
639
        auto [11, r1] = test(x, lx, rx, b, 1);
38e
        auto [12, r2] = test(v, lv, rv, a, -1);
c43
        if (12 > r2) swap(12, r2);
50a
        T l = max(11, 12), r = min(r1, r2);
339
        if (1 > r) return 0:
        ll k = (r-1) / abs(b) + 1;
42f
839
        return k; // solucoes: x = 1 + [0, k)*|b|
98e }
```

#### 4.11 Euclides estendido

```
// Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
// Assume a, b >= 0
//
// O(log(min(a, b)))

2be tuple<11, l1, l1> ext_gcd(l1 a, l1 b) {
3bd     if (!a) return {b, 0, 1};
550     auto [g, x, y] = ext_gcd(b%a, a);
c59     return {g, y - b/a*x, x};
354 }
```

# 4.12 Exponenciacao rapida

```
// (x^y mod m) em O(log(y))
03c ll pow(ll x, ll y, ll m) \{ // \text{ iterativo} \}
c85
        ll ret = 1:
1b8
        while (y) {
895
            if (y & 1) ret = (ret * x) % m;
23b
            v >>= 1:
             x = (x * x) % m;
сс5
        }
020
edf
        return ret;
12b }
03c ll pow(ll x, ll y, ll m) { // recursivo
        if (!y) return 1;
426
        11 ans = pow(x*x\%m, y/2, m);
88 d
        return v%2 ? x*ans%m : ans;
7d4 }
```

# 4.13 Fast Walsh Hadamard Transform

```
// FWHT<'!'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
//
// O(n log(n))

382 template<char op, class T> vector<T> FWHT(vector<T> f, bool inv = false) {
```

```
b75
        int n = f.size();
d78
        for (int k = 0; (n-1) >> k; k++) for (int i = 0; i < n; i++) if
   (i>>k&1) {
29e
            int j = i^{(1 << k)};
            if (op == '\^') f[i] += f[i], f[i] = f[i] - 2*f[i];
627
            if (op == ', ') f[i] += (inv ? -1 : 1) * f[j];
a38
93c
            if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
1 bb
        }
578
        if (op == ', ', and inv) for (auto& i : f) i /= n;
abe
        return f;
50e }
// Generalizacao de FWHT de Xor
//
// Convolucao de soma mod B, usar tamanho potencia de B!!
// Precisa definir o tipo T e a raiz primitiva g
// satisfazendo g^b == g
// Se possivel, hardcodar a multiplicacao de matriz
// feita em cada iteracao faz ficar bem mais rapido
//
// O(n b log_b(n))
4fc template < class T>
811 vector <T > FWHT (vector <T > f, int b, T g, bool inv = false) {
        int n = f.size();
929
        vector <T> w(b):
a7a
        w[1] = g;
        for (int i = 2; i < b; i++) w[i] = w[i - 1] * g;
aab
        w[0] = w[b - 1] * g;
dec
d59
        if (inv) reverse(w.begin() + 1, w.end());
c31
        for (int pot = 1; pot < n; pot *= b) {</pre>
339
            for (int i = 0; i < n; i++) if (!(i / pot % b)) {
4 e 1
                vector <T> res(b);
                for (int j = 0; j < b; j++) {
сОс
                     for (int k = 0; k < b; k++)
a41
a32
                         res[j] = res[j] + w[j * k % b] * f[i + k *
   pot];
9 c 8
                    if (inv) res[j] = res[j] / b;
dc3
e0e
                for (int j = 0; j < b; j++) f[i + j * pot] = res[j];
7f8
            }
        }
25e
```

```
abe
        return f;
902 }
// Exemplos da FWHT Generalizada:
//
// mod 7, resposta mod 998244353:
// T = mint, g = 14553391
//
// mod 3, resposta cabe em um long long:
// T = array<11, 2>, g = {0, 1};
//
// using T = array<11, 2>;
    // T operator +(const T& a, const T& b) {
    // return T{a[0] + b[0], a[1] + b[1]};
// }
    // T operator *(const T& a, const T& b) {
    // return T{a[0] * b[0] - a[1] * b[1],
                 a[0] * b[1] + a[1] * b[0] - a[1] * b[1];
    //
// };
   // T operator /(const T& a, const int& b) {
    // return T{a[0] / b, a[1] / b};
// }
4.14 FFT
// Chamar convolution com vector < complex < double >> para FFT
// Precisa do mint para NTT
//
// O(n log(n))
// Para FFT
488 void get_roots(bool f, int n, vector<complex<double>>& roots) {
f 26
        const static double PI = acosl(-1);
71 a
        for (int i = 0; i < n/2; i++) {
            double alpha = i*((2*PI)/n);
b1e
1a1
            if (f) alpha = -alpha;
069
            roots[i] = {cos(alpha), sin(alpha)};
804
       }
de5 }
// Para NTT
9f7 template < int p>
97b void get_roots(bool f, int n, vector<mod_int<p>>& roots) {
1 e 6
        mod_int  r;
de9
        int ord;
```

57a

if (p == 998244353) {

```
9b6
            r = 102292;
81b
            ord = (1 << 23):
121
        } else if (p == 754974721) {
            r = 739831874;
43a
f0a
            ord = (1 << 24):
d48
        } else if (p == 167772161) {
a2a
            r = 243;
033
            ord = (1 << 25);
5 a 4
        } else assert(false);
547
        if (f) r = r^(p - 1 - ord/n);
ee2
        else r = r^(ord/n):
be4
        roots[0] = 1:
        for (int i = 1; i < n/2; i++) roots[i] = roots[i-1]*r;
078
63f }
8a2 template < typename T > void fft(vector < T > & a, bool f, int N,
   vector<int>& rev) {
        for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
bc7
   a[rev[i]]):
        int 1, r, m;
12b
        vector < T> roots(N);
cb4
        for (int n = 2; n <= N; n *= 2) {</pre>
192
0f4
            get_roots(f, n, roots);
5dc
            for (int pos = 0; pos \langle N; pos += n \rangle {
432
                1 = pos + 0, r = pos + n/2, m = 0;
                while (m < n/2) {
a88
297
                     auto t = roots[m] * a[r];
254
                     a[r] = a[1] - t;
b8f
                    a[1] = a[1] + t;
2c9
                    1++, r++, m++;
d89
                }
1fd
            }
185
        }
        if (f) {
235
            auto invN = T(1) / T(N);
1 c 5
557
            for (int i = 0; i < N; i++) a[i] = a[i] * invN;
256
        }
1b1 }
bf5 template < typename T > vector < T > convolution (vector < T > & a,
   vector < T > & b) {
        vector <T> l(a.begin(), a.end()), r(b.begin(), b.end());
87a
        int N = 1.size()+r.size()-1;
e0a
f03
       int n = 1, log_n = 0;
        while (n \le N) n *= 2, log_n++;
0 a 4
808
        vector < int > rev(n);
```

```
603
        for (int i = 0; i < n; i++) {</pre>
434
             rev[i] = 0;
f 44
             for (int j = 0; j < log_n; j++) if (i >> j & 1)
                 rev[i] = 1 << (log_n-1-i);
4ff
256
143
         assert(N <= n);
fa4
        l.resize(n):
7 e 4
        r.resize(n);
56e
        fft(1, false, n, rev);
        fft(r, false, n, rev);
fcf
917
        for (int i = 0; i < n; i++) l[i] *= r[i];
88b
        fft(l, true, n, rev);
5e1
        l.resize(N):
        return 1;
792
bd6 }
// NTT
6c8 template < int p, typename T > vector < mod_int < p >> ntt(vector < T > & a,
    vector < T > \& b) {
         vector < mod int < p >> A(a.begin(), a.end()), B(b.begin(),
d52
    b.end());
d29
         return convolution(A, B);
3bf }
// Convolucao de inteiro
// Precisa do CRT
// Tabela de valores:
// [0,1]
             - <int, 1>
// [-1e5, 1e5] - <11, 2>
// [-1e9, 1e9] - <__int128, 3>
b3c template < typename T, int mods >
eec vector<T> int convolution(vector<int>& a. vector<int>& b) {
         static const int M1 = 998244353, M2 = 754974721, M3 =
    167772161;
bf5
         auto c1 = ntt < M1 > (a, b);
221
         auto c2 = (mods >= 2 ? ntt < M2 > (a, b) : vector < mod_int < M2 >> ());
f9b
         auto c3 = (mods >= 3 ? ntt < M3 > (a, b) : vector < mod_int < M3 >> ());
         vector <T> ans;
2da
5 c 5
        for (int i = 0; i < c1.size(); i++) {</pre>
             crt < T > at (c1[i].v, M1);
c09
316
             if (mods >= 2) at = at * crt<T>(c2[i].v, M2);
987
             if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
b2b
             ans.push back(at.a);
```

```
26d
            if (at.a > at.m/2) ans.back() -= at.m;
                                                                             ec9
                                                                                     }
b9f
ba7
                                                                                      for (int i = 0; i < m; i++) if (where[i] == -1)</pre>
        return ans;
                                                                             12 e
                                                                             018
5e8 }
                                                                                          return pair(INF, ans);
                                                                             280
                                                                                      return pair(1, ans);
                                                                             292 }
4.15 Gauss
                                                                             4.16 Gauss - Z2
// Resolve sistema linear
// Retornar um par com o numero de solucoes
// e alguma solucao, caso exista
                                                                             // D eh dimensao do espaco vetorial
                                                                             // add(v) - adiciona o vetor v na base (retorna se ele jah pertencia
// O(n^2 * m)
                                                                                 ao span da base)
                                                                             // coord(v) - retorna as coordenadas (c) de v na base atual (basis^T.c
67a template < typename T>
                                                                             // recover(v) - retorna as coordenadas de v nos vetores na ordem em
728 pair <int, vector <T>> gauss (vector <vector <T>> a, vector <T> b) {
        const double eps = 1e-6;
                                                                                 que foram inseridos
6ca
        int n = a.size(), m = a[0].size();
                                                                             // coord(v).first e recover(v).first - se v pertence ao span
f92
2f0
        for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre>
                                                                             //
                                                                             // Complexidade:
3cb
        vector < int > where (m. -1);
                                                                             // add, coord, recover: O(D^2 / 64)
        for (int col = 0, row = 0; col < m and row < n; col ++) {
237
f05
            int sel = row;
                                                                             cd4 template<int D> struct gauss_z2 {
b95
            for (int i=row: i<n: ++i)</pre>
                                                                             3 c 1
                                                                                      bitset <D > basis[D], keep[D];
e55
                 if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
                                                                             b16
                                                                                      int rk, in;
            if (abs(a[sel][col]) < eps) continue;</pre>
                                                                                      vector < int > id:
2c4
                                                                             482
1ae
            for (int i = col: i <= m: i++)
                 swap(a[sel][i], a[row][i]);
dd2
                                                                             cf2
                                                                                      gauss_z2 (): rk(0), in(-1), id(D, -1) {};
2 c 3
            where [col] = row;
                                                                             04 e
                                                                                      bool add(bitset<D> v) {
0 c 0
            for (int i = 0; i < n; i++) if (i != row) {</pre>
                                                                             42 c
                                                                                          in++;
96c
                T c = a[i][col] / a[row][col]:
                                                                             fb0
                                                                                          bitset <D> k:
                                                                                          for (int i = D - 1; i >= 0; i--) if (v[i]) {
d5c
                for (int j = col; j <= m; j++)</pre>
                                                                             659
c8f
                     a[i][j] -= a[row][j] * c;
                                                                             189
                                                                                              if (basis[i][i]) v ^= basis[i], k ^= keep[i];
            }
490
                                                                             4 e 6
                                                                                              else {
b70
                                                                                                  k[i] = true, id[i] = in, keep[i] = k;
            row++;
                                                                             ea6
        }
                                                                                                  basis[i] = v, rk++;
3 d 8
                                                                             6се
                                                                             8a6
                                                                                                  return true;
b1d
        vector < T > ans(m, 0);
                                                                             b34
                                                                                              }
        for (int i = 0; i < m; i++) if (where[i] != -1)</pre>
                                                                             09 с
e1a
            ans[i] = a[where[i]][m] / a[where[i]][i];
                                                                             d1f
12a
                                                                                          return false;
603
        for (int i = 0; i < n; i++) {
                                                                             58b
            T sum = 0:
501
                                                                             0f6
                                                                                      pair < bool, bitset < D >> coord(bitset < D > v) {
a75
            for (int j = 0; j < m; j++)
                                                                             944
                                                                                          bitset <D> c;
5 a 9
                 sum += ans[j] * a[i][j];
                                                                             659
                                                                                          for (int i = D - 1; i >= 0; i--) if (v[i]) {
b1f
            if (abs(sum - a[i][m]) > eps)
                                                                             a39
                                                                                              if (basis[i][i]) v ^= basis[i], c[i] = true;
6cd
                 return pair(0, vector <T>());
                                                                             8af
                                                                                              else return {false, bitset<D>()};
```

```
a08
5 db
             return {true, c};
a08
330
        pair < bool , vector < int >> recover (bitset < D > v) {
             auto [span, bc] = coord(v);
22e
             if (not span) return {false, {}};
af8
f79
            bitset < D > aux:
            for (int i = D - 1; i >= 0; i--) if (bc[i]) aux ^= keep[i];
5 a 0
ea9
            vector < int > oc;
            for (int i = D - 1; i >= 0; i--) if (aux[i])
   oc.push_back(id[i]);
             return {true. oc}:
001
b75
688 };
```

## 4.17 Integracao Numerica

```
// Metodo de Simpson 3/8
// Integra f no intervalo [a, b], erro cresce proporcional a (b - a)^5

676 const int N = 3*100; // multiplo de 3
287 ld integrate(ld a, ld b, function<ld(ld)> f) {
    ld s = 0, h = (b - a)/N;
    for (int i = 1; i < N; i++) s += f(a + i*h)*(i%3 ? 3 : 2);
    da     return (f(a) + s + f(b))*3*h/8;
    c7e }</pre>
```

#### 4.18 Inverso Modular

```
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^(b-2)

f0a ll inv(ll a, ll b) {
    ae1      return a > 1 ? b - inv(b%a, a)*b/a : 1;
    cf9 }

// computa o inverso modular de 1..MAX-1 modulo um primo
    a88 ll inv[MAX]:
    Of2 inv[1] = 1;
    Ofa for (int i = 2; i < MAX; i++) inv[i] = MOD - MOD/i*inv[MOD%i]%MOD;</pre>
```

#### 4.19 Karatsuba

```
// Os pragmas podem ajudar
// Para n \sim 2e5, roda em < 1 s
//
// O(n^1.58)
//#pragma GCC optimize("Ofast")
//#pragma GCC target ("avx,avx2")
77a template < typename T > void kar(T* a, T* b, int n, T* r, T* tmp) {
d4c
        if (n <= 64) {
510
             for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)
212
                 r[i+j] += a[i] * b[j];
505
             return;
        }
bb8
194
        int mid = n/2;
2d7
        T * atmp = tmp, *btmp = tmp+mid, *E = tmp+n;
        memset(E, 0, sizeof(E[0])*n);
4 f 1
c 65
        for (int i = 0; i < mid; i++) {</pre>
             atmp[i] = a[i] + a[i+mid];
c.72
4b9
             btmp[i] = b[i] + b[i+mid];
a3f
38a
        kar(atmp, btmp, mid, E, tmp+2*n);
b1e
        kar(a, b, mid, r, tmp+2*n);
229
        kar(a+mid, b+mid, mid, r+n, tmp+2*n);
c65
        for (int i = 0; i < mid; i++) {</pre>
735
            T \text{ temp} = r[i+mid];
de7
            r[i+mid] += E[i] - r[i] - r[i+2*mid];
f1e
            r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
f72
28f }
e38 template<typename T> vector<T> karatsuba(vector<T> a, vector<T> b)
   {
ba3
        int n = max(a.size(), b.size());
        while (n&(n-1)) n++;
a84
ca9
        a.resize(n), b.resize(n);
        vector \langle T \rangle ret (2*n), tmp (4*n);
ae0
644
        kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
edf
        return ret:
f87 }
```

## 4.20 Logaritmo Discreto

// Resolve logaritmo discreto com o algoritmo baby step giant step

```
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem. retorna -1
//
// O(sqrt(m) * log(sqrt(m))
d41
da8 int dlog(int b, int a, int m) {
        if (a == 0) return b ? -1 : 1; // caso nao definido
d41
a6e
        a \%= m, b \%= m;
        int k = 1. shift = 0;
31e
        while (1) {
6e3
           int g = gcd(a, m);
d47
           if (g == 1) break:
d41
9bc
           if (b == k) return shift;
642
           if (b % g) return -1;
c36
            b /= g, m /= g, shift++;
            k = (11) k * a / g % m;
9ab
515
d41
af7
        int sq = sqrt(m)+1, giant = 1;
975
        for (int i = 0; i < sq; i++) giant = (11) giant * a % m;</pre>
d41
0 b 5
        vector < pair < int , int >> baby;
33f
        for (int i = 0, cur = b; i <= sq; i++) {
496
            babv.emplace back(cur, i);
            cur = (11) cur * a % m;
16c
622
eb4
        sort(baby.begin(), baby.end());
d41
        for (int j = 1, cur = k; j <= sq; j++) {
9 c 9
ace
            cur = (11) cur * giant % m;
            auto it = lower_bound(baby.begin(), baby.end(), pair(cur,
   INF));
            if (it != baby.begin() and (--it)->first == cur)
d26
                return sq * j - it->second + shift;
ac3
        }
b9d
d41
daa
        return -1:
739 }
4.21 Miller-Rabin
// Testa se n eh primo. n \le 3 * 10^18
```

```
// O(log(n)), considerando multiplicacao
```

```
// e exponenciacao constantes
d8b ll mul(ll a, ll b, ll m) {
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
        return ret < 0 ? ret+m : ret;</pre>
074
2f3 }
03c ll pow(ll x, ll y, ll m) {
13a
        if (!v) return 1;
dbc
        ll ans = pow(mul(x, x, m), y/2, m);
7fa
        return v%2 ? mul(x, ans, m) : ans;
539 }
1a2 bool prime(ll n) {
       if (n < 2) return 0;
237
       if (n <= 3) return 1;
9de
     if (n % 2 == 0) return 0;
       ll r = __builtin_ctzll(n - 1), d = n >> r;
       // com esses primos, o teste funciona garantido para n <= 2^64
        // funciona para n <= 3*10^24 com os primos ate 41
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
dd1
   1795265022}) {
            ll x = pow(a, d, n);
da0
709
            if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
4a2
            for (int j = 0; j < r - 1; j++) {
10f
                x = mul(x, x, n):
                if (x == n - 1) break:
df0
1ff
            if (x != n - 1) return 0:
e1b
5b0
        }
6a5
        return 1:
9a1 }
```

# Multipoint Evaluation And Interpolation

```
// Evaluation:
// Avalia o polinomio f(x) nos pontos p[0], p[1], ..., p[n-1]
// Interpolation:
// Retorna o polinomio f(x) de grau n que
// satisfaz f(x) = y pra o conjunto de pontos x, y
// Precisa do ntt e
// - do divmod pro evaluate
```

```
// - da derivada pro interpolate
// O divmod e a derivada estao no arquivo powerSeries.cpp
// O(n log^2(n))
55c namespace multipoint {
        vector < poly > tree;
204
415
        void build(vector<mint>& p) {
2fa
            int n = p.size();
ab2
            tree.resize(2*n);
f43
            for (int i = 0; i < n; i++) tree [n + i] = \{-p[i], 1\};
917
            for (int i = n - 1; i > 0; i - -)
                 tree[i] = convolution(tree[2*i], tree[2*i + 1]);
e 64
1f2
d21
        vector < mint > evaluate(poly& f, vector < mint > & p) {
734
            build(p);
2fa
            int n = p.size();
            vector<poly> ans(2 * n);
305
            ans [1] = divmod(f, tree[1]).second:
66a
            for (int i = 2; i < 2 * n; i++)
4cb
                 ans[i] = divmod(ans[i/2], tree[i]).second;
67d
512
            vector<mint> results(n);
8f6
            for (int i = 0; i < n; i++) results[i] = ans[n + i][0];</pre>
238
            return results:
084
        poly prod(vector<mint>& p, int 1, int r) {
c25
8f7
            if (l == r) return {-p[l], 1};
            int m = (1 + r) / 2;
ee4
6df
            return convolution(prod(p, 1, m), prod(p, m + 1, r));
b0b
        poly interpolate(vector<mint>& x, vector<mint>& y) {
ec4
            int n = x.size();
34a
12e
            polv p = D(prod(x, 0, n - 1)):
            auto d = evaluate(p, x);
beb
305
            vector<poly> ans(2 * n);
            for (int i = 0; i < n; i++) ans [n + i] = {y[i] / d[i]};
5d0
0f2
            for (int i = n - 1; i > 0; i--) {
                 poly p1 = convolution(tree[2*i], ans[2*i + 1]);
ccd
5 c 7
                poly p2 = convolution(tree[2*i + 1], ans[2*i]);
6ea
                 ans[i] = p1:
                for (int j = 0; j < p1.size(); j++) ans[i][j] += p2[j];</pre>
9 d 8
5 df
887
            return ans[1]:
1 a 1
        }
44e }
```

#### 4.23 NTT

```
// Precisa do mint (primitivas de aritmetica modular)
//
// O(n log (n))
4e9 const int MOD = 998244353;
Of4 typedef mod_int<MOD> mint;
c4b void ntt(vector < mint > & a, bool rev) {
        int n = a.size(); auto b = a;
6f1
479
        assert(!(n&(n-1)));
        mint g = 1;
513
459
        while ((g^{(MOD / 2)}) == 1) g += 1;
        if (rev) g = 1 / g;
574
        for (int step = n / 2; step; step /= 2) {
e55
             mint w = g^{(MOD)} / (n / step), wn = 1;
306
41 e
             for (int i = 0; i < n/2; i += step) {
                 for (int j = 0; j < step; <math>j++) {
c29
673
                     auto u = a[2 * i + j], v = wn * a[2 * i + j +
    step];
                     b[i+j] = u + v; b[i + n/2 + j] = u - v;
464
09e
c39
                 wn = wn * w;
ade
            }
257
             swap(a. b):
c50
        if (rev) {
1bb
b18
             auto n1 = mint(1) / n;
b28
             for (auto& x : a) x *= n1;
        }
eaa
574 }
7c4 vector<mint> convolution(const vector<mint>& a, const
    vector<mint>& b) {
03a
        vector<mint> l(a.begin(), a.end()), r(b.begin(), b.end());
d71
        int N = 1.size()+r.size()-1, n = 1;
        while (n \le N) n *= 2;
73b
fa4
        l.resize(n);
7 e 4
        r.resize(n);
156
        ntt(1, false);
557
        ntt(r, false);
917
        for (int i = 0; i < n; i++) l[i] *= r[i];
de9
        ntt(1. true):
5 e 1
        l.resize(N);
792
        return 1:
```

4bb }

# 4.24 Operacoes em Polinomios e Series de Potencias

```
// Precisa do NTT
// O exp nao foi bem testado
//
// Fonte:
   github.com/celiopassos/competitive-programming/blob/master/algorithms
// D, I: O(n)
// inv, divmod, log e exp: O(n log(n))
0d8 using poly = vector<mint>;
c8d const int MAGIC = 512;
e47 poly D(poly p) {
cd1
       if (p.empty()) return p;
73c
        for (int i = 0; i + 1 < p.size(); i++)
            p[i] = (i + 1) * p[i + 1];
9 c 7
087
        p.pop_back();
74e
        return p;
b66 }
62b polv I(polv p) {
       int n = p.size();
2fa
ef6
        p.push_back(0);
        for (int i = n - 1; i \ge 0; i - -)
056
5fe
            p[i + 1] = p[i] / (i + 1);
481
        p[0] = 0;
74e
        return p;
809 }
3ef poly inv(poly p) {
640
        assert(!p.empty() && p[0] == 1);
253
        poly q = {mint(1) / p[0]};
        int n = p.size(), k = 1;
ee3
d20
        while (k < n) {
            k *= 2:
539
d93
            q.resize(2 * k):
d01
            ntt(q, false);
0ac
            poly p0(2 * k);
f39
            copy_n(p.begin(), min(k, n), p0.begin());
697
            ntt(p0, false);
818
            for (int i = 0; i < 2 * k; i++)
```

```
eef
                 q[i] *= 2 - p0[i] * q[i];
ff8
             ntt(q, true);
             q.resize(k);
afe
        }
cd0
        q.resize(n);
ba3
bef
         return q;
60b }
18f pair <poly, poly > divslow(const poly& a, const poly& b) {
bea
         poly q, r = a;
/meathematwings/eformaslizeowerses or isoiszeh(p)p) {
d01
             q.push_back(r.back() / b.back());
e4d
             if (q.back() != 0)
c06
                 for (int i = 0; i < b.size(); i++)</pre>
7d9
                     r.end()[-i-1] -= q.back() * b.end()[-i-1];
515
            r.pop_back();
733
539
        reverse(q.begin(), q.end());
442
         return {q, r};
9b9 }
// retorna (q, r) : a(x) = b(x) * q(x) + r(x)
06b pair <poly, poly > divmod(const poly& a, const poly& b) {
884
         if (a.size() < b.size()) return {{}, a};</pre>
         if (max(b.size(), a.size() - b.size()) < MAGIC) return</pre>
    divslow(a. b):
c83
         poly ra = poly(a.rbegin(), a.rend());
b35
        poly rb = poly(b.rbegin(), b.rend());
        int k = a.size() - b.size() + 1;
b8b
b8b
        rb.resize(k);
        poly irb = inv(move(rb)), q = convolution(ra, irb);
864
57b
        q = poly(q.rend() - k, q.rend());
         poly r = convolution(move(q), b);
fe8
f63
        for (int i = 0; i < r.size(); i++) r[i] = a[i] - r[i];
         while (r.size() > 1 && r.back() == 0) r.pop_back();
46f
442
         return {q, r};
8af }
53b poly log(poly p) {
640
         assert(!p.empty() && p[0] == 1);
2fa
         int n = p.size();
        auto d = D(p), i = inv(p);
983
25 f
        auto r = convolution(d, i);
7 c.9
        r.resize(n - 1):
c7b
         return I(move(r));
35 c }
```

```
84d poly exp(poly p) {
        assert(p.empty() || p[0] == 0);
380
5 a 8
        polv q = \{1\};
ee3
        int n = p.size(), k = 1;
d20
        while (k < n) {
539
            k *= 2;
afe
            q.resize(k);
0a7
            poly b = log(q);
2 d3
            for (int i = 0; i < k; i++) b[i] *= -1;
f89
            b[0] += 1:
45 d
            for (int i = 0; i < min(n, k); i++) b[i] += p[i];</pre>
b0d
            q = convolution(q, b);
afe
            a.resize(k);
a3f
ba3
        q.resize(n);
bef
        return q;
f78 }
```

# 4.25 Pollard's Rho Alg

```
// Usa o algoritmo de deteccao de ciclo de Floyd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
d8b 11 mul(11 a, 11 b, 11 m) {
e7a
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
       if (! y) return 1;
13a
dbc
        11 ans = pow(mul(x, x, m), y/2, m);
        return y%2 ? mul(x, ans, m) : ans;
7fa
539 }
1a2 bool prime(ll n) {
       if (n < 2) return 0;
1aa
237
        if (n <= 3) return 1:
9de
       if (n % 2 == 0) return 0;
```

```
f6a
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
dd1
         for (int a: {2, 325, 9375, 28178, 450775, 9780504,
    1795265022}) {
da0
            ll x = pow(a, d, n);
             if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue:
709
4a2
             for (int j = 0; j < r - 1; j++) {
                 x = mul(x, x, n):
10f
df0
                 if (x == n - 1) break;
1ff
e1b
             if (x != n - 1) return 0;
5b0
        }
6a5
         return 1:
9a1 }
9cf ll rho(ll n) {
0f9
         if (n == 1 or prime(n)) return n;
f7c
         auto f = [n](11 x) \{ return mul(x, x, n) + 1; \};
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
8a5
         while (t % 40 != 0 or gcd(prd, n) == 1) {
533
8a0
             if (x==y) x = ++x0, y = f(x);
e13
             q = mul(prd, abs(x-y), n);
21f
            if (q != 0) prd = q;
450
            x = f(x), y = f(f(y)), t++;
379
002
         return gcd(prd, n);
523 }
5b7 vector<ll> fact(ll n) {
        if (n == 1) return {};
1 b 9
0ec
        if (prime(n)) return {n};
        11 d = rho(n):
0ed
1de
         vector < 11 > 1 = fact(d), r = fact(n / d):
3af
        1.insert(1.end(), r.begin(), r.end());
792
         return 1;
902 }
4.26 Produto de dois long long mod m
// 0(1)
d8b 11 mul(11 a, 11 b, 11 m) { // a*b % m
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
```

#### 4.27 Simplex

```
// Maximiza c^T x s.t. Ax <= b. x >= 0
// O(2^n), porem executa em O(n^3) no caso medio
395 const double eps = 1e-7;
493 namespace Simplex {
69c
        vector < vector < double >> T;
14e
        int n. m:
43e
        vector < int > X, Y;
c51
        void pivot(int x, int y) {
8e6
            swap(X[v], Y[x-1]);
d03
            for (int i = 0; i <= m; i++) if (i != y) T[x][i] /=
   T[x][y];
33c
            T[x][y] = 1/T[x][y];
38b
            for (int i = 0; i <= n; i++) if (i != x and abs(T[i][v]) >
   eps) {
774
                for (int j = 0; j <= m; j++) if (j != y) T[i][j] -=
   T[i][y] * T[x][i];
                T[i][y] = -T[i][y] * T[x][y];
3 d8
            }
a7d
e05
        }
        // Retorna o par (valor maximo, vetor solucao)
6f8
        pair < double , vector < double >> simplex(
e9d
                 vector < vector < double >> A, vector < double >> b,
   vector <double > c) {
            n = b.size(), m = c.size();
5 bb
002
            T = vector(n + 1, vector < double > (m + 1));
2d9
            X = vector < int > (m);
0 c 2
            Y = vector < int > (n):
115
            for (int i = 0; i < m; i++) X[i] = i;
            for (int i = 0; i < n; i++) Y[i] = i+m;</pre>
51f
            for (int i = 0; i < m; i++) T[0][i] = -c[i];
5 b 5
603
            for (int i = 0; i < n; i++) {
                 for (int j = 0; j < m; j++) T[i+1][j] = A[i][j];
ba6
                T[i+1][m] = b[i];
eca
            }
07с
667
            while (true) {
714
                int x = -1, y = -1;
2 db
                double mn = -eps;
                for (int i = 1; i <= n; i++) if (T[i][m] < mn) mn =
   T[i][m], x = i;
af2
                if (x < 0) break;
```

```
882
                for (int i = 0; i < m; i++) if (T[x][i] < -eps) { y = }
   i: break: }
4a6
                if (y < 0) return \{-1e18, \{\}\}; // sem solucao para Ax
   \leq b
7fb
                pivot(x, y);
            }
472
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
562
                for (int i = 0; i < m; i++) if (T[0][i] < mn) mn =
   T[0][i], y = i;
9Ъ0
                if (v < 0) break:
034
                mn = 1e200:
5af
                for (int i = 1; i \le n; i++) if (T[i][y] > eps and
   T[i][m] / T[i][y] < mn
48 f
                     mn = T[i][m] / T[i][y], x = i;
53b
                if (x < 0) return {1e18, {}}; // c^T x eh ilimitado
7fb
                pivot(x, y);
81 e
            }
290
            vector < double > r(m);
            for (int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]] =
32 f
   T[i+1][m];
            return {T[0][m], r};
e 5.9
7 a 4
a64 }
      Teorema Chines do Resto
// Combina equacoes modulares lineares: x = a (mod m)
```

```
// O m final eh o lcm dos m's, e a resposta eh unica mod o lcm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
3bd
        if (!a) return {b, 0, 1};
550
        auto [g, x, y] = ext_gcd(b\%a, a);
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11> struct crt {
        Ta, m;
        crt() : a(0), m(1) {}
5f3
        crt(T a_, T m_) : a(a_), m(m_) {}
7eb
```

```
911
        crt operator * (crt C) {
             auto [g, x, y] = ext_gcd(m, C.m);
238
             if ((a - C.a) % g) a = -1;
dc0
4f9
             if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);
             T lcm = m/g*C.m;
d09
             T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
             return crt((ans % lcm + lcm) % lcm, lcm);
1 f 2
        }
0d9 };
```

#### 4.29 Totiente

```
// O(sqrt(n))
a7e int tot(int n) {
0f6
        int ret = n;
505
        for (int i = 2; i*i <= n; i++) if (n % i == 0) {
b0c
            while (n \% i == 0) n /= i;
125
            ret -= ret / i:
34a
af4
        if (n > 1) ret -= ret / n;
edf
        return ret;
fae }
```

# 5 Strings

## 5.1 Aho-corasick

```
// query retorna o somatorio do numero de matches de
// todas as stringuinhas na stringona
//
// insert - O(|s| log(SIGMA))
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - O(|s|)

eal namespace aho {
807          map<char, int> to[MAX];
c87          int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
bfc         void insert(string& s) {
```

```
05е
             int at = 0;
 b4f
             for (char c : s) {
b68
                 auto it = to[at].find(c);
                 if (it == to[at].end()) at = to[at][c] = ++idx;
 1 c 9
 361
                 else at = it->second:
 ff4
142
             term[at]++. sobe[at]++:
 6eb
 d41 #warning nao esquece de chamar build() depois de inserir
 0a8
         void build() {
26a
             queue < int > q;
 537
             q.push(0);
 dff
             link[0] = exit[0] = -1:
 402
             while (q.size()) {
 379
                 int i = q.front(); q.pop();
 3 c 4
                 for (auto [c, j] : to[i]) {
 5da
                     int 1 = link[i];
                     while (l != -1 and !to[l].count(c)) l = link[l];
 102
 7a5
                     link[j] = 1 == -1 ? 0 : to[1][c];
                     exit[j] = term[link[j]] ? link[j] : exit[link[j]];
 3ab
                     if (exit[j]+1) sobe[j] += sobe[exit[j]];
6f2
 113
                     q.push(j);
f1d
                }
 367
             }
 768
bc0
         int query(string& s) {
             int at = 0, ans = 0;
 86 d
b4f
             for (char c : s){
                 while (at != -1 and !to[at].count(c)) at = link[at];
1ca
 5b9
                 at = at == -1 ? 0 : to[at][c];
 2 b 1
                 ans += sobe[at]:
b85
ba7
             return ans:
 038
        }
 a30 }
5.2 eertree
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
//
// O(n) amortizado, considerando alfabeto O(1)
```

```
8eb struct eertree {
7сс
        vector < vector < int >> t;
42e
       int n, last, sz;
        vector < int > s, len, link, qt;
745
d36
        eertree(int N) {
            t = vector(N+2, vector(26, int()));
ec8
            s = len = link = qt = vector < int > (N+2);
cee
            s[0] = -1:
cd1
288
            link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
688
            sz = 2, last = 0, n = 1;
8dc
        void add(char c) {
244
692
            s[n++] = c -= 'a':
34f
            while (s[n-len[last]-2] != c) last = link[last];
289
            if (!t[last][c]) {
                int prev = link[last];
dab
553
                while (s[n-len[prev]-2] != c) prev = link[prev];
                link[sz] = t[prev][c];
fb2
3f5
                len[sz] = len[last]+2;
                t[last][c] = sz++;
1 f 8
f8b
344
            qt[last = t[last][c]]++;
b1d
f17
        int size() { return sz-2; }
2af
        11 propagate() {
b73
            11 \text{ ret} = 0:
            for (int i = n; i > 1; i--) {
ebb
                qt[link[i]] += qt[i];
fd3
db5
                ret += qt[i];
074
            }
edf
            return ret:
ef6
a2e };
5.3 KMP
// matching(s, t) retorna os indices das ocorrencias
// de s em t
// autKMP constroi o automato do KMP
// Complexidades:
// pi - O(n)
```

// match - 0(n + m)

```
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
ea8 template < typename T > vector < int > pi(T s) {
         vector < int > p(s.size());
019
725
         for (int i = 1, j = 0; i < s.size(); i++) {</pre>
a51
             while (j \text{ and } s[j] != s[i]) j = p[j-1];
             if (s[j] == s[i]) j++;
973
f8c
             p[i] = i;
e0a
74 e
         return p;
f50 }
c10 template < typename T> vector <int> matching(T& s, T& t) {
658
         vector<int> p = pi(s), match;
a1b
         for (int i = 0, j = 0; i < t.size(); i++) {</pre>
6be
             while (j \text{ and } s[j] != t[i]) j = p[j-1];
             if (s[j] == t[i]) j++;
c.4d
310
             if (j == s.size()) match.push_back(i-j+1), j = p[j-1];
028
ed8
         return match:
c82 }
a2d struct KMPaut : vector < vector < int >> {
47 c
         KMPaut(){}
6 c 7
         KMPaut (string& s) : vector < vector < int >> (26,
    vector < int > (s.size()+1)) {
503
             vector < int > p = pi(s);
04b
             auto& aut = *this:
4fa
             aut[s[0]-'a'][0] = 1;
             for (char c = 0; c < 26; c++)
19a
5 d 3
                 for (int i = 1; i <= s.size(); i++)</pre>
42 b
                      aut[c][i] = s[i] - a' == c ? i+1 : aut[c][p[i-1]]:
4bb
        }
79b };
5.4 Manacher
// manacher recebe um vetor de T e retorna o vetor com tamanho dos
   palindromos
```

```
// manacher recebe um vetor de T e retorna o vetor com tamanho dos
   palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// ret[2*i+1] = tamanho maior palindromo centrado em i e i+1
//
// Complexidades:
// manacher - O(n)
// palindrome - <O(n), O(1)>
```

```
// pal_end - 0(n)
28a template < typename T> vector < int> manacher (const T& s) {
        int l = 0, r = -1, n = s.size();
18f
fc9
        vector < int > d1(n), d2(n):
        for (int i = 0; i < n; i++) {</pre>
603
            int k = i > r ? 1 : min(d1[l+r-i], r-i);
821
            while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++;
61a
61e
            d1[i] = k--;
9f6
            if (i+k > r) l = i-k, r = i+k;
950
        }
e03
       1 = 0. r = -1:
603
       for (int i = 0: i < n: i++) {
            int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
a 64
            while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k]) k++;
2 c 6
            d2[i] = --k:
eaa
26d
            if (i+k-1 > r) l = i-k, r = i+k-1;
4fe
        vector < int > ret(2*n-1);
c41
e6b
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
e1d
edf
        return ret:
ebb }
// verifica se a string s[i..j] eh palindromo
cac template < typename T > struct palindrome {
        vector < int > man;
f97
        palindrome(const T& s) : man(manacher(s)) {}
b2d
9 d 7
        bool query(int i, int j) {
            return man[i+j] >= j-i+1;
bad
1 e 7
60c }:
// tamanho do maior palindromo que termina em cada posicao
7cb template < typename T > vector < int > pal_end(const T& s) {
        vector < int > ret(s.size());
e57
fde
        palindrome <T> p(s);
d51
        ret[0] = 1:
        for (int i = 1; i < s.size(); i++) {</pre>
88e
a32
            ret[i] = min(ret[i-1]+2, i+1);
            while (!p.query(i-ret[i]+1, i)) ret[i]--;
6ea
78e
edf
        return ret:
8bd }
```

# 5.5 Min/max suffix/cyclic shift

```
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)
016 template < typename T > int max_suffix(T s, bool mi = false) {
        s.push_back(*min_element(s.begin(), s.end())-1);
476
1a4
        int ans = 0:
88 e
        for (int i = 1; i < s.size(); i++) {</pre>
eec
            int j = 0;
            while (ans+j < i and s[i+j] == s[ans+j]) j++;
708
7a2
            if (s[i+j] > s[ans+j]) {
b52
                if (!mi or i != s.size()-2) ans = i;
e 51
            } else if (j) i += j-1;
69 c
ba7
        return ans;
f2a }
a1a template < typename T> int min_suffix(T s) {
76b
        for (auto& i : s) i *= -1;
09 d
        s.push_back(*max_element(s.begin(), s.end())+1);
        return max_suffix(s, true);
925
ec0 }
97c template < typename T > int max_cyclic_shift(T s) {
        int n = s.size();
163
1ad
        for (int i = 0; i < n; i++) s.push_back(s[i]);</pre>
        return max_suffix(s);
20a
d34 }
08a template < typename T> int min_cyclic_shift(T s) {
        for (auto& i : s) i *= -1;
76b
7be
        return max_cyclic_shift(s);
c7a }
5.6 String Hashing
// Complexidades:
// construtor - \Omega(|s|)
// operator() - 0(1)
```

```
// Complexidades:
// construtor - O(|s|)
// operator() - O(1)

878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
```

```
463 int uniform(int 1, int r) {
a7f
        uniform_int_distribution < int > uid(1, r);
f54
        return uid(rng);
d9e }
9e0 template <int MOD> struct str_hash { // 116fcb
        static int P;
dcf
        vector<11> h, p;
        str_hash(string s) : h(s.size()), p(s.size()) {
ea8
7a2
            p[0] = 1, h[0] = s[0];
ad7
            for (int i = 1; i < s.size(); i++)</pre>
84c
                p[i] = p[i - 1]*P%MOD, h[i] = (h[i - 1]*P + s[i])%MOD;
1ef
af7
        11 operator()(int 1, int r) { // retorna hash s[1...r]
749
            11 hash = h[r] - (1 ? h[1 - 1]*p[r - 1 + 1]%MOD : 0);
dfd
            return hash < 0 ? hash + MOD : hash;</pre>
        }
3ba
116 };
217 template < int MOD > int str_hash < MOD > :: P = uniform (256, MOD - 1); //
   1 > | sigma|
```

# 5.7 String Hashing - modulo 2<sup>61</sup> - 1

```
// Quase duas vezes mais lento
// Complexidades:
// build - 0(|s|)
// operator() - 0(1)
9d0 const ll MOD = (111<<61) - 1;
e38 ll mulmod(ll a, ll b) {
ff3
        const static ll LOWER = (111 << 30) - 1, GET31 = (111 << 31) - 1;
410
        11 \ 11 = a\&LOWER, h1 = a>>30, 12 = b\&LOWER, h2 = b>>30;
        11 m = 11*h2 + 12*h1, h = h1*h2;
d54
        11 \text{ ans} = 11*12 + (h>>1) + ((h&1)<<60) + (m>>31) +
   ((m\&GET31) << 30) + 1;
1 dd
        ans = (ans\&MOD) + (ans >> 61), ans = (ans\&MOD) + (ans >> 61);
c0f
        return ans - 1:
f98 }
798 mt19937 64
   rng(chrono::steady_clock::now().time_since_epoch().count());
f89 ll uniform(ll l, ll r) {
        uniform_int_distribution<11> uid(1, r);
969
```

```
f54
        return uid(rng);
cac }
d7d struct str_hash {
        static 11 P:
c20
dcf
        vector <11> h, p;
ea8
        str_hash(string s) : h(s.size()), p(s.size()) {
7a2
            p[0] = 1, h[0] = s[0];
ad7
            for (int i = 1; i < s.size(); i++)</pre>
                p[i] = mulmod(p[i-1], P), h[i] = (mulmod(h[i-1], P), h[i])
632
   P) + s[i])%MOD;
507
af7
        11 operator()(int 1. int r) { // retorna hash s[1...r]
538
            11 hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
dfd
            return hash < 0 ? hash + MOD : hash;</pre>
544
        }
148 };
6c5 ll str_hash::P = uniform(256, MOD - 1); // l > |sigma|
5.8 Suffix Array - O(n log n)
// kasai recebe o suffix array e calcula lcp[i],
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],...,n-1]
//
// Complexidades:
// suffix_array - O(n log(n))
// kasai - O(n)
733 vector<int> suffix_array(string s) {
b38
        s += "$";
043
        int n = s.size(), N = max(n, 260);
2f3
        vector < int > sa(n), ra(n);
29b
        for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];
0a2
        for(int k = 0; k < n; k ? k *= 2 : k++) {
5 ce
            vector < int > nsa(sa), nra(n), cnt(N);
            for (int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
fae
   cnt[ra[i]]++:
            for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
4 c 4
368
            for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]] = nsa[i];
            for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=
28f
   ra[sa[i]] !=
f86
                ra[sa[i-1]] or ra[(sa[i]+k)\%n] != ra[(sa[i-1]+k)\%n];
```

26b

ra = nra;

```
d5e
             if (ra[sa[n-1]] == n-1) break;
                                                                              39 d
11e
                                                                                  i+(1<<j) <= n/b; i++)
057
        return vector<int>(sa.begin()+1, sa.end());
                                                                              ba5
ff3 }
                                                                              41a
481 vector < int > kasai (string s, vector < int > sa) {
                                                                              e 34
232
        int n = s.size(), k = 0;
                                                                              27b
        vector < int > ra(n), lcp(n);
                                                                              e80
408
676
        for (int i = 0; i < n; i++) ra[sa[i]] = i;
                                                                              fd3
                                                                              a4e
740
        for (int i = 0; i < n; i++, k -= !!k) {
                                                                              ea3
199
            if (ra[i] == n-1) { k = 0; continue; }
                                                                                  t[n/b*j+y-(1<<j)+1]));
1 de
             int i = sa[ra[i]+1]:
                                                                              be6
             while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
891
                                                                              62 a
d98
            lcp[ra[i]] = k;
                                                                              093
a07
                                                                              bab };
5ed
        return lcp;
fbe }
                                                                              9d7 struct suffix_array {
                                                                              ac0
                                                                                       string s;
                                                                                       int n:
                                                                              1a8
                                                                              5b4
    Suffix Array - O(n)
                                                                              2de
                                                                                      rmq<int> RMQ;
// Rapidao
                                                                              d6e
// Computa o suffix array em 'sa', o rank em 'rnk'
                                                                              91 d
// e o lcp em 'lcp'
                                                                              82 d
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
                                                                                  int k) {
// Complexidades
                                                                              c17
// O(n) para construir
                                                                              bac
// query - 0(1)
                                                                              703
                                                                              000
1a5 template < typename T > struct rmq {
                                                                                      }
                                                                              6f3
517
        vector <T> v;
                                                                              d66
fcc
        int n; static const int b = 30;
                                                                              a 76
70e
        vector < int > mask, t;
                                                                              3a9
```

int op(int x, int y) { return  $v[x] \leftarrow v[y]$  ? x : y; }

int small(int r, int sz = b) { return

at = (at <<1) & ((1 << b) -1);

r-msb(mask[r]&((1<<sz)-1));}

rmq() {}

}

int msb(int x) { return \_\_builtin\_clz(1)-\_\_builtin\_clz(x); }

 $rmq(const\ vector < T > \&\ v_) : v(v_), n(v.size()), mask(n), t(n)$ 

while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;

for (int i = 0, at = 0; i < n;  $\max[i++]$  = at |= 1) {

for (int i = 0; i < n/b; i++) t[i] = small(b\*i+b-1);

183

ee1

6ad

43c

2 e 5

a61

c00

c2f

ea4

```
for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
   t[n/b*(j-1)+i+(1<<(j-1))]);
        int index_query(int 1, int r) {
            if (r-l+1 \le b) return small(r, r-l+1);
            int x = 1/b+1, y = r/b-1;
            if (x > y) return op(small(1+b-1), small(r));
            int j = msb(y-x+1);
            int ans = op(small(1+b-1), op(t[n/b*j+x],
            return op(ans. small(r)):
        T query(int 1, int r) { return v[index_query(1, r)]; }
        vector < int > sa, cnt, rnk, lcp;
        bool cmp(int a1, int b1, int a2, int b2, int a3=0, int b3=0) {
            return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3);
        template < typename T> void radix (int* fr, int* to, T* r, int N,
            cnt = vector < int > (k+1, 0);
            for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;</pre>
            for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
            for (int i = N-1; i+1; i--) to[--cnt[r[fr[i]]]] = fr[i];
        void rec(vector<int>& v. int k) {
            auto &tmp = rnk, &m0 = lcp;
            int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
7f8
            vector < int > R(sz2+3);
74 f
            for (int i = 1, j = 0; j < sz2; i += i/(3) R[j++] = i;
b30
            radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
207
            radix(&tmp[0], &R[0], &v[0]+1, sz2, k);
5f1
            radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
af5
            int dif = 0:
            int 10 = -1, 11 = -1, 12 = -1;
ed9
d81
            for (int i = 0; i < sz2; i++) {
                if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
8de
```

```
v[tmp[i]+2] != 12)
                                                                              5a4
                                                                                                   k = 0:
                     10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
                                                                              5 e 2
b43
                                                                                                   continue;
                                                                              9df
   v[tmp[i]+2], dif++;
                 if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
                                                                              39a
                                                                                               int j = sa[rnk[i]+1];
                 else R[tmp[i]/3+sz] = dif;
                                                                              891
                                                                                               while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
1f5
            }
                                                                              825
                                                                                              lcp[rnk[i]] = k;
d18
                                                                              a3e
                                                                                          }
47f
            if (dif < sz2) {
                                                                                          RMQ = rmq<int>(lcp);
                                                                              9ff
146
                 rec(R, dif);
                                                                              9a8
                                                                                      }
                 for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;
746
105
            } else for (int i = 0; i < sz2; i++) sa[R[i]-1] = i;</pre>
                                                                              588
                                                                                      int query(int i, int j) {
                                                                              d97
                                                                                          if (i == j) return n-i;
6f4
            for (int i = 0, i = 0; i < sz2; i++) if (sa[i] < sz)
                                                                              223
                                                                                          i = rnk[i], i = rnk[i];
   tmp[j++] = 3*sa[i];
                                                                              c3a
                                                                                          return RMQ.query(min(i, j), max(i, j)-1);
7 ce
            radix(&tmp[0], &m0[0], &v[0], sz, k);
                                                                              940
74d
            for (int i = 0; i < sz2; i++)</pre>
                                                                              71 c
                                                                                      pair<int, int> next(int L, int R, int i, char c) {
                 sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2;
                                                                              024
                                                                                          int 1 = L, r = R+1;
с9е
                                                                                          while (1 < r) {
                                                                              40 c
            int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
                                                                                              int m = (1+r)/2;
332
                                                                              ee4
            while (p >= 0 \text{ and } p2 >= 0) {
                                                                                               if (i+sa[m] >= n \text{ or } s[i+sa[m]] < c) l = m+1:
1 c 9
                                                                              e7e
                if ((sa[p2]%3==1 and cmp(v[m0[p]], v[sa[p2]],
                                                                              ef3
                                                                                               else r = m:
   R[m0[p]/3],
                                                                              ebe
                     R[sa[p2]/3+sz])) or (sa[p2]%3==2 and cmp(v[m0[p]],
                                                                              575
                                                                                          if (1 == R+1 \text{ or } s[i+sa[1]] > c) return \{-1, -1\};
0 ce
                                                                              eb7
                                                                                          L = 1;
   v[sa[p2]],
                     v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
   R[sa[p2]/3+1]))
                                                                              9e2
                                                                                          1 = L \cdot r = R+1;
                                                                                          while (1 < r) {
300
                     sa[at--] = sa[p2--];
                                                                              40 c
cb0
                 else sa[at--] = m0[p--]:
                                                                              ee4
                                                                                              int m = (1+r)/2;
                                                                                              if (i+sa[m] >= n \text{ or } s[i+sa[m]] <= c) l = m+1;
214
                                                                              1a1
f2b
            while (p >= 0) sa[at--] = m0[p--];
                                                                              ef3
                                                                                              else r = m;
            if (N\%3==1) for (int i = 0; i < N; i++) sa[i] = sa[i+1];
                                                                                          }
eb6
                                                                              b5b
        }
                                                                              56a
                                                                                          R = 1-1;
ee6
                                                                                          return {L. R}:
                                                                              e13
938
        suffix_array(const string \& s_) : s(s_), n(s.size()), sa(n+3),
                                                                              71b
                                                                                      // quantas vezes 't' ocorre em 's' - 0(|t| log n)
                 cnt(n+1), rnk(n), lcp(n-1) {
e62
9fe
            vector < int > v(n+3);
                                                                              66 d
                                                                                      int count_substr(string& t) {
f9b
            for (int i = 0; i < n; i++) v[i] = i;
                                                                              b2b
                                                                                          int L = 0, R = n-1;
eba
            radix(&v[0], &rnk[0], &s[0], n, 256);
                                                                              c9d
                                                                                          for (int i = 0; i < t.size(); i++) {</pre>
e6d
            int dif = 1:
                                                                              de0
                                                                                               tie(L, R) = next(L, R, i, t[i]);
830
            for (int i = 0; i < n; i++)
                                                                              4fc
                                                                                               if (L == -1) return 0;
                 v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i-1]]);
419
                                                                              cff
            if (n >= 2) rec(v, dif);
7cf
                                                                              fbf
                                                                                          return R-L+1;
            sa.resize(n):
                                                                                      }
fb9
                                                                              aaa
            for (int i = 0; i < n; i++) rnk[sa[i]] = i;</pre>
76f
                                                                                      // exemplo de f que resolve o problema
            for (int i = 0, k = 0; i < n; i++, k -= !!k) {
892
668
                 if (rnk[i] == n-1) {
                                                                                          https://codeforces.com/edu/course/2/lesson/2/5/practice/contes
```

```
57e
        ll f(ll k) \{ return k*(k+1)/2; \}
        11 dfs(int L, int R, int p) { // dfs na suffix tree chamado em
e68
   pre ordem
            int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
c54
            // Tem 'ext - p' substrings diferentes que ocorrem 'R-L+1'
                vezes
            // O LCP de todas elas eh 'ext'
            11 \text{ ans} = (ext-p)*f(R-L+1);
f80
            // L eh terminal, e folha sse L == R
63c
            if (sa[L]+ext == n) L++:
            // se for um SA de varias strings separadas como s#t$u&,
                usar no lugar do if de cima
            // (separadores < 'a', diferentes e inclusive no final)
            // while (L \leq R && (sa[L]+ext == n || s[sa[L]+ext] \leq
               'a')) {
            // L++;
            // }
            while (L <= R) {
add
5 a 8
                int idx = L != R ? RMQ.index_query(L, R-1) : -1;
5ef
                if (idx == -1 or lcp[idx] != ext) idx = R;
478
                ans += dfs(L, idx, ext);
28d
                L = idx+1:
590
ba7
            return ans;
        }
e 2.1
        // sum over substrings: computa, para toda substring t
           distinta de s.
        // \sum f(# ocorrencias de t em s) - O (n)
        ll sos() { return dfs(0, n-1, 0); }
ca8
6fa };
5.10 Suffix Array Dinamico
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
//
// Complexidades:
```

```
// Construir sobre uma string de tamanho n: O(n log n)
// push_front e pop_front: O(log n) amortizado
2fe struct dyn_sa {
3 c 9
        struct node {
             int sa, lcp;
1d4
ed1
             node *1, *r, *p;
f0d
             int sz, mi;
17b
             node(int sa_, int lcp_, node* p_) : sa(sa_), lcp(lcp_),
543
                 1(NULL), r(NULL), p(p_), sz(1), mi(lcp) {}
01e
            void update() {
58f
                 sz = 1, mi = lcp;
bd7
                 if (1) sz += 1->sz, mi = min(mi, 1->mi);
                 if (r) sz += r \rightarrow sz, mi = min(mi, r \rightarrow mi);
a 54
            }
27 c
574
        };
bb7
        node* root:
295
        vector<ll> tag; // tag of a suffix (reversed id)
        string s; // reversed
ac0
        dyn_sa() : root(NULL) {}
cf4
        dyn_sa(string s_) : dyn_sa() {
e 45
            reverse(s_.begin(), s_.end());
ae4
519
            for (char c : s_) push_front(c);
2a7
a 86
        \sim dvn_sa() {
609
            vector < node *> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
                 if (!x) continue:
ee9
1 c 7
                 q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x:
653
            }
        }
8 c 1
        int size(node* x) { return x ? x->sz : 0; }
73 c
08e
        int mirror(int i) { return s.size()-1 - i; }
580
        bool cmp(int i, int j) {
a29
            if (s[i] != s[j]) return s[i] < s[j];</pre>
5 b 4
             if (i == 0 or j == 0) return i < j;</pre>
988
            return tag[i-1] < tag[j-1];</pre>
9fd
919
        void fix_path(node* x) { while (x) x->update(), x = x->p; }
245
        void flatten(vector<node*>& v, node* x) {
8 c 8
             if (!x) return:
e 96
            flatten(v. x->1);
```

```
2 a 2
            v.push_back(x);
            flatten(v, x->r);
42d
01f
        void build(vector<node*>& v, node*& x, node* p, int L, int R,
   11 1, 11 r) {
            if (L > R) return void(x = NULL);
04c
331
            int M = (L+R)/2;
3 e 3
            11 m = (1+r)/2;
7 e 5
            x = v[M];
63e
            x - p = p;
bb3
            tag[x->sa] = m;
            build(v, x->1, x, L, M-1, 1, m-1), build(v, x->r, x, M+1,
ae0
   R. m+1. r):
ca8
            x->update();
a3a
        void fix(node*& x, node* p, ll l, ll r) {
82f
7f0
            if (3*max(size(x->1), size(x->r)) \le 2*size(x)) return
   x->update();
            vector<node*> v;
3 d 1
            flatten(v, x);
0 c c
            build(v, x, p, 0, v.size()-1, 1, r);
ea9
b86
        node* next(node* x) {
b19
728
            if(x->r)
a 9 1
                x = x -> r:
347
                 while (x->1) x = x->1:
ea5
                 return x;
e7d
402
            while (x->p \text{ and } x->p->r == x) x = x->p;
137
            return x->p;
48b
b68
        node* prev(node* x) {
            if (x->1) {
e41
a26
                 x = x -> 1:
93c
                 while (x->r) x = x->r;
ea5
                 return x;
9be
            }
6a1
            while (x->p \text{ and } x->p->1 == x) x = x->p;
137
            return x->p;
73e
        }
        int get_lcp(node* x, node* y) {
4f7
75a
            if (!x or !y) return 0; // change defaut value here
            if (s[x->sa] != s[y->sa]) return 0;
e51
843
            if (x->sa == 0 \text{ or } y->sa == 0) return 1;
            return 1 + query(mirror(x->sa-1), mirror(y->sa-1));
4 d 0
8 d 6
        }
```

```
ad6
        void add_suf(node*& x, node* p, int id, ll l, ll r) {
91 e
             if (!x) {
8e3
                 x = new node(id, 0, p);
8e2
                 node *prv = prev(x), *nxt = next(x);
65 d
                 int lcp_cur = get_lcp(prv, x), lcp_nxt = get_lcp(x,
   nxt);
ca3
                 if (nxt) nxt->lcp = lcp_nxt, fix_path(nxt);
71 f
                 x \rightarrow lcp = lcp_cur;
7b4
                 tag[id] = (1+r)/2;
                 x->update();
ca8
505
                 return;
d0e
             }
4a3
             if (cmp(id, x->sa)) add_suf(x->l, x, id, l, tag[x->sa]-1);
сЗа
             else add_suf(x->r, x, id, tag[x->sa]+1, r);
3db
             fix(x, p, 1, r);
c98
        }
ec2
        void push_front(char c) {
cc7
             s += c:
493
             tag.push_back(-1);
             add_suf(root, NULL, s.size() - 1, 0, 1e18);
05 e
        }
1f2
7f3
        void rem suf(node*& x. int id) {
6cf
             if (x->sa != id) {
864
                 if (tag[id] < tag[x->sa]) return rem_suf(x->1, id);
e6f
                 return rem suf(x->r, id);
             }
2ae
2 cf
             node * nxt = next(x);
             if (nxt) nxt->lcp = min(nxt->lcp, x->lcp), fix_path(nxt);
09b
b20
             node *p = x->p, *tmp = x;
f3f
             if (!x->1 \text{ or } !x->r) {
2fd
                 x = x - > 1 ? x - > 1 : x - > r;
753
                 if (x) x - p = p:
696
            } else {
7f7
                 for (tmp = x->1, p = x; tmp->r; tmp = tmp->r) p = tmp;
f2a
                 x->sa = tmp->sa, x->lcp = tmp->lcp;
482
                 if (tmp -> 1) tmp -> 1 -> p = p;
14 c
                 if (p->1 == tmp) p->1 = tmp->1;
a 94
                 else p \rightarrow r = tmp \rightarrow 1;
97 c
            }
b5e
             fix_path(p);
7 c 3
             delete tmp;
510
        }
15 b
        void pop_front() {
             if (!s.size()) return;
abe
342
             s.pop_back();
```

```
436
             rem_suf(root, s.size());
сбе
             tag.pop_back();
987
        int guerv(node* x, 11 1, 11 r, 11 a, 11 b) {
530
             if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1) \text{ return}
e51
   s.size():
            if (a <= l and r <= b) return x->mi;
ef5
8eb
             int ans = s.size();
             if (a \le tag[x->sa] \text{ and } tag[x->sa] \le b) ans = min(ans,
   x \rightarrow lcp);
d99
             ans = min(ans, query(x->1, 1, tag[x->sa]-1, a, b));
261
             ans = min(ans, querv(x->r, tag[x->sa]+1, r, a, b)):
ba7
             return ans:
4 c8
588
        int query(int i, int j) { // lcp(s[i..], s[j..])
209
             if (i == j) return s.size() - i;
29e
             11 a = tag[mirror(i)], b = tag[mirror(j)];
710
             int ret = query(root, 0, 1e18, min(a, b)+1, max(a, b));
edf
             return ret:
84e
        // optional: get rank[i], sa[i] and lcp[i]
        int rank(int i) {
044
396
            i = mirror(i):
52f
             node * x = root;
7 c.9
             int ret = 0:
f4c
             while (x) {
33e
                 if (tag[x->sa] < tag[i]) {</pre>
f9d
                     ret += size(x->1)+1;
a91
                     x = x -> r;
                 } else x = x - > 1:
6dc
a19
edf
             return ret:
153
        pair < int , int > operator[](int i) {
649
52f
             node* x = root;
31e
             while (1) {
d4d
                 if (i < size(x->1)) x = x->1;
4 e 6
                 else {
85f
                     i = size(x->1);
                     if (!i) return {mirror(x->sa), x->lcp};
e03
040
                     i--, x = x->r;
b9b
7 a 2
             }
90c
        }
4c2 };
```

#### 5.11 Suffix Automaton

```
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
16e namespace sam {
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
c1a
0b8
        int nxt[2*MAX][26];
e6a
        void add(int c) {
            int at = cur:
17a
9a6
            len[sz] = len[cur]+1, cur = sz++;
500
            while (at != -1 and !nxt[at][c]) nxt[at][c] = cur, at =
   link[at]:
7ea
            if (at == -1) { link[cur] = 0; return; }
654
            int q = nxt[at][c];
fd9
            if (len[q] == len[at]+1) { link[cur] = q; return; }
31 f
            int qq = sz++;
2c3
            len[qq] = len[at]+1, link[qq] = link[q];
            for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];</pre>
9a9
            while (at != -1 and nxt[at][c] == q) nxt[at][c] = qq, at =
e76
   link[at]:
8d8
            link[cur] = link[q] = qq;
61a
94 e
        void build(string& s) {
889
            cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
            for (auto i : s) add(i-'a');
9fe
            int at = cur:
17a
            while (at) acc[at] = 1, at = link[at];
121
0e7
        }
        // coisas que da pra fazer:
28 c
        11 distinct_substrings() {
04b
            11 \text{ ans} = 0:
a1e
            for (int i = 1; i < sz; i++) ans += len[i] - len[link[i]];
ba7
            return ans;
0d7
        string longest_common_substring(string& S, string& T) {
a6c
419
            build(S):
111
            int at = 0, 1 = 0, ans = 0, pos = -1;
            for (int i = 0; i < T.size(); i++) {</pre>
d59
f2c
                while (at and !nxt[at][T[i]-'a']) at = link[at], l =
   len[at]:
                if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'], 1++;
efa
749
                else at = 0.1 = 0:
a1a
                if (1 > ans) ans = 1, pos = i;
2b3
            }
```

```
20f
            return T.substr(pos-ans+1, ans);
        }
930
46e
        11 dp [2 * MAX];
        11 paths(int i) {
455
2 a 8
            auto& x = dp[i];
            if (x) return x;
dee
            x = 1;
483
            for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
   paths(nxt[i][i]);
            return x;
ea5
d88
        void kth_substring(int k, int at=0) { // k=1 : menor substring
105
9d2
            for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {</pre>
d58
                if (paths(nxt[at][i]) >= k) {
                     cout << char('a'+i);</pre>
d02
c43
                     kth_substring(k-1, nxt[at][i]);
505
                     return;
69a
5 f 4
                k -= paths(nxt[at][i]);
ef6
            }
a13
c37 };
5.12 Trie
// trie T() constroi uma trie para o alfabeto das letras minusculas
// trie T(tamanho do alfabeto, menor caracter) também pode ser usado
// T.insert(s) - 0(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, -1 se nao achar - O(|s|)
// T.count_pref(s) numero de strings que possuem s como prefixo -
   0(|s|)
ab5 struct trie {
        vector < vector < int >> to;
e1a
450
        vector < int > end, pref;
af0
        int sigma; char norm;
        trie(int sigma_=26, char norm_='a') : sigma(sigma_),
   norm(norm) {
58a
            to = {vector < int > (sigma)};
86e
            end = \{0\}, pref = \{0\};
fe1
64e
        void insert(string s) {
c67
            int x = 0;
```

```
7e7
            for (auto c : s) {
                 int &nxt = to[x][c-norm];
800
dd7
                 if (!nxt) {
0aa
                     nxt = to.size();
526
                     to.push_back(vector<int>(sigma));
770
                     end.push_back(0), pref.push_back(0);
933
                 }
827
                 x = nxt, pref[x]++;
34 c
            end[x]++, pref[0]++;
421
e66
        }
6b2
        void erase(string s) {
c67
            int x = 0:
            for (char c : s) {
b4f
800
                 int &nxt = to[x][c-norm];
                 x = nxt, pref[x]--;
10 c
d8e
                 if (!pref[x]) nxt = 0;
885
104
             end[x]--, pref[0]--;
8bf
        }
        int find(string s) {
aee
c 67
            int x = 0:
7 e 7
            for (auto c : s) {
2ec
                 x = to[x][c-norm];
59b
                 if (!x) return -1;
42 d
ea5
            return x;
есс
        }
839
        int count_pref(string s) {
b25
             int id = find(s);
fc1
             return id >= 0 ? pref[id] : 0;
560
        }
8ca };
5.13 Z
// z[i] = lcp(s, s[i..n))
//
// Complexidades:
// z - O(|s|)
// \text{ match - } O(|s| + |p|)
a19 vector<int> get_z(string s) {
163
        int n = s.size();
2b1
        vector < int > z(n, 0);
```

# 6 graphs

### 7 Problemas

# 7.1 Angle Range Intersection

```
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
//
// Tudo 0(1)
32a struct angle_range {
        static constexpr ld ALL = 1e9, NIL = -1e9;
75e
395
        ld 1, r;
        angle_range() : 1(ALL), r(ALL) {}
c77
894
        angle_range(ld l_, ld r_) : l(l_), r(r_) { fix(l), fix(r); }
4ee
        void fix(ld& theta) {
da7
            if (theta == ALL or theta == NIL) return;
323
            if (theta > 2*pi) theta -= 2*pi;
868
            if (theta < 0) theta += 2*pi;</pre>
625
2ee
        bool empty() { return l == NIL; }
931
        bool contains(ld q) {
40f
            fix(a);
4d7
            if (1 == ALL) return true;
fec
            if (1 == NIL) return false;
            if (1 < r) return 1 < q and q < r;
6a6
075
            return q > 1 or q < r;
800
9 c 7
        friend angle_range operator &(angle_range p, angle_range q) {
743
            if (p.l == ALL or q.l == NIL) return q;
```

```
20f
             if (q.1 == ALL or p.1 == NIL) return p;
7 d 5
             if (p.l > p.r \text{ and } q.l > q.r) \text{ return } \{\max(p.l, q.l),
   min(p.r, q.r)};
             if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
aa6
             if (p.1 > p.r) {
8b8
249
                 if (q.r > p.l) return {max(q.l, p.l) , q.r};
6f7
                 else if (q.1 < p.r) return \{q.1, \min(q.r, p.r)\};
270
                 return {NIL, NIL};
337
             if (max(p.1, q.1) > min(p.r, q.r)) return {NIL, NIL};
5 a 8
bcb
             return {max(p.1, q.1), min(p.r, q.r)};
142
        }
5e1 }:
```

### 7.2 Area da Uniao de Retangulos

```
// O(n log(n))
// 5d8d2f
aa4 namespace seg {
        pair < int , 11 > seg[4*MAX];
b1b
        ll lazv[4*MAX], *v;
1a8
        int n:
        pair<int, 1l> merge(pair<int, 1l> 1, pair<int, 1l> r){
e 01
719
             if (1.second == r.second) return {1.first+r.first.
   1.second:
53b
             else if (l.second < r.second) return l;</pre>
             else return r;
aa0
d82
        }
6fc
        pair < int, ll > build(int p=1, int l=0, int r=n-1) {
3 c 7
             lazv[p] = 0;
             if (1 == r) return seg[p] = {1, v[1]};
bf8
             int m = (1+r)/2:
ee4
             return seg[p] = merge(build(2*p, 1, m), build(2*p+1, m+1,
432
   r));
f 94
        void build(int n2. 11* v2) {
d9e
680
             n = n2, v = v2;
6f2
             build():
f8a
        }
        void prop(int p, int l, int r) {
ceb
208
             seg[p].second += lazy[p];
2c9
             if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3 c 7
             lazv[p] = 0;
```

```
bf2
        pair < int, 11 > query (int a, int b, int p=1, int 1=0, int r=n-1)
693
   {
6 b 9
            prop(p, 1, r);
527
            if (a <= l and r <= b) return seg[p];</pre>
            if (b < 1 or r < a) return {0, LINF};</pre>
9b7
ee4
            int m = (1+r)/2;
            return merge (query (a, b, 2*p, 1, m), query (a, b, 2*p+1,
   m+1, r));
        }
786
        pair < int, 11 > update (int a, int b, int x, int p=1, int 1=0,
07с
   int r=n-1) {
6b9
            prop(p, 1, r);
9 a 3
            if (a \le 1 \text{ and } r \le b)
b94
                lazy[p] += x;
6b9
                prop(p, 1, r);
534
                return seg[p];
821
            if (b < l or r < a) return seg[p];
e9f
            int m = (1+r)/2;
ee4
086
            return seg[p] = merge(update(a, b, x, 2*p, 1, m),
579
                     update(a, b, x, 2*p+1, m+1, r));
        }
c65
043 };
eb5 ll seg_vec[MAX];
8be 11 area_sq(vector<pair<int, int>, pair<int, int>>> &sq){
        vector<pair<int, int>, pair<int, int>>> up;
28c
60a
        for (auto it : sq){
619
            int x1, y1, x2, y2;
ae0
            tie(x1, y1) = it.first;
68e
            tie(x2, y2) = it.second;
80f
            up.push_back(\{\{x1+1, 1\}, \{y1, y2\}\}\});
            up.push_back({{x2+1, -1}, {y1, y2}});
aee
6 c 3
        sort(up.begin(), up.end());
092
049
        memset(seg_vec, 0, sizeof seg_vec);
6fe
        11 H_MAX = MAX;
156
        seg::build(H_MAX-1, seg_vec);
7ba
        auto it = up.begin();
04b
        11 \text{ ans} = 0;
f 1 4
        while (it != up.end()){
07f
            11 L = (*it).first.first;
718
            while (it != up.end() && (*it).first.first == L){
127
                 int x, inc, y1, y2;
d35
                 tie(x. inc) = it->first:
```

```
d3d
                 tie(v1, v2) = it -> second;
5 d 1
                 seg::update(y1+1, y2, inc);
40 d
                 it++;
9b1
            }
852
            if (it == up.end()) break;
            11 R = (*it).first.first;
d8a
f59
            11 W = R - L;
efd
            auto jt = seg::querv(0, H_MAX-1);
91 a
            11 H = H MAX - 1;
e8a
            if (jt.second == 0) H -= jt.first;
8df
            ans += W*H;
5 c 8
ba7
        return ans;
385 }
```

# 7.3 Area Maxima de Histograma

```
// Assume que todas as barras tem largura 1.
// e altura dada no vetor v
//
// O(n)
15e ll area(vector<int> v) {
b73
        ll ret = 0;
4се
        stack < int > s;
        // valores iniciais pra dar tudo certo
        v.insert(v.begin(), -1);
447
        v.insert(v.end(), -1);
d56
1f8
        s.push(0);
0be
        for(int i = 0; i < (int) v.size(); i++) {</pre>
78 e
            while (v[s.top()] > v[i]) {
265
                11 h = v[s.top()]; s.pop();
de1
                ret = max(ret, h * (i - s.top() - 1));
40a
18 e
            s.push(i);
        }
020
edf
        return ret:
e43 }
```

#### 7.4 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
97c ll divi[MAX];
398 ll expo(ll a, ll b, ll m) {
1 c 1
       if (!b) return 1:
        11 ans = expo(a*a\%m, b/2, m);
399
751
        if (b\%2) ans *= a;
2 e 9
        return ans%m:
754 }
f0a ll inv(ll a. ll b){
bca
        return 1 < a ? b - inv(b%a,a)*b/a : 1;
041 }
153 template < typename T> tuple < T, T, T> ext_gcd(T a, T b) {
        if (!a) return {b, 0, 1};
3bd
        auto [g, x, y] = ext_gcd(b%a, a);
550
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11> struct crt {
627
        Ta, m;
5 f 3
        crt(): a(0), m(1) {}
7eb
        crt(T a_, T m_) : a(a_), m(m_) {}
911
        crt operator * (crt C) {
238
            auto [g, x, y] = ext_gcd(m, C.m);
dc0
            if ((a - C.a) \% g) a = -1;
            if (a == -1 or C.a == -1) return crt(-1, 0);
4 f 9
d09
           T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
            return crt((ans % lcm + lcm) % lcm, lcm);
d8d
        }
1f2
0d9 };
6f2 pair <11, 11> divide_show(11 n, int p, int k, int pak) {
4f7
        if (n == 0) return {0, 1};
d02
        11 blocos = n/pak, falta = n%pak;
        ll periodo = divi[pak], resto = divi[falta];
2 ce
        11 r = expo(periodo, blocos, pak)*resto%pak;
616
445
        auto rec = divide_show(n/p, p, k, pak);
a51
        ll y = n/p + rec.first;
        r = r*rec.second % pak;
bb9
```

```
90f
         return {y, r};
533 }
6e6 ll solve_pak(ll n, ll x, int p, int k, int pak) {
         divi[0] = 1:
d34
f2b
         for (int i = 1; i <= pak; i++) {</pre>
901
             divi[i] = divi[i-1];
840
             if (i%p) divi[i] = divi[i] * i % pak;
51a
        }
         auto dn = divide_show(n, p, k, pak), dx = divide_show(x, p, k,
4ac
   pak),
162
              dnx = divide show(n-x, p, k, pak):
768
        11 y = dn.first-dx.first-dnx.first, r =
b64
             (dn.second*inv(dx.second, pak)%pak)*inv(dnx.second,
    pak)%pak;
035
         return expo(p, y, pak) * r % pak;
d78 }
9dd ll solve(ll n, ll x, int mod) {
         vector < pair < int , int >> f;
490
c3b
         int mod2 = mod;
        for (int i = 2; i*i <= mod2; i++) if (mod2%i==0) {</pre>
7b4
aff
             int c = 0;
75 b
             while (mod2\%i==0) mod2 /= i, c++;
2a1
             f.push_back({i, c});
fe7
Off
        if (mod2 > 1) f.push_back({mod2, 1});
e 96
         crt ans(0, 1);
a13
         for (int i = 0; i < f.size(); i++) {</pre>
702
             int pak = 1;
7 e 4
             for (int j = 0; j < f[i].second; j++) pak *= f[i].first;</pre>
             ans = ans * crt(solve_pak(n, x, f[i].first, f[i].second,
304
    pak), pak);
7fd
5fb
         return ans.a;
689 }
7.5 Closest pair of points
// O(nlogn)
915 pair <pt, pt > closest_pair_of_points(vector <pt > v) {
3d2
         int n = v.size();
fca
         sort(v.begin(), v.end());
```

```
31 c
        for (int i = 1; i < n; i++) if (v[i] == v[i-1]) return
   {v[i-1], v[i]};
        auto cmp_y = [&](const pt &1, const pt &r) {
c20
            if (1.v != r.v) return 1.v < r.v;</pre>
b53
920
             return 1.x < r.x;</pre>
55a
        };
62e
        set < pt, decltype(cmp_y) > s(cmp_y);
3 d 9
        int 1 = 0, r = -1;
6a2
        11 d2_min = numeric_limits<11>::max();
4d5
        pt pl, pr;
bd1
        const int magic = 5;
a55
        while (r+1 < n) {
7 f 1
             auto it = s.insert(v[++r]).first;
c92
            int cnt = magic/2;
773
             while (cnt-- and it != s.begin()) it--;
a01
d68
             while (cnt++ < magic and it != s.end()) {</pre>
f19
                 if (!((*it) == v[r])) {
                     11 d2 = dist2(*it, v[r]);
67e
74e
                     if (d2_min > d2) {
229
                         d2_min = d2;
                         pl = *it;
841
4f2
                         pr = v[r];
7 d 9
                     }
10a
                 }
40d
                 it++;
801
eb0
             while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min)
   s.erase(v[1++]);
de6
        return {pl, pr};
c74
f90 }
```

#### 7.6 Coloração de Grafo de Intervalo

```
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
//
// O(n log(n))

615 vector<int> coloring(vector<pair<int, int>>& v) {
3d2    int n = v.size();
c08    vector<pair<int, pair<int, int>>> ev;
603    for (int i = 0; i < n; i++) {</pre>
```

```
150
            ev.push_back({v[i].first, {1, i}});
cda
            ev.push_back({v[i].second, {0, i}});
6a4
49e
        sort(ev.begin(), ev.end());
        vector < int > ans(n), avl(n);
360
265
        for (int i = 0; i < n; i++) avl.push_back(n-i);</pre>
4bf
        for (auto i : ev) {
            if (i.second.first == 1) {
cbe
021
                 ans[i.second.second] = avl.back();
                 avl.pop_back();
a 0.0
e 98
            } else avl.push_back(ans[i.second.second]);
3a6
        }
ba7
        return ans:
83a }
     Conectividade Dinamica DC
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
8f2 typedef pair <int, int> T;
1cd namespace data {
553
        int n, ans;
573
        int p[MAX], sz[MAX];
        stack<int> S;
ee6
e5b
        void build(int n2) {
1 e 3
            n = n2;
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;</pre>
0b2
            ans = n;
cba
        }
1b1
        int find(int k) {
006
            while (p[k] != k) k = p[k];
839
            return k:
        }
c1e
072
        void add(T x) {
700
            int a = x.first. b = x.second:
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
            ans - - :
3c6
            if (sz[a] > sz[b]) swap(a, b);
4 c 2
            S.push(a):
582
            sz[b] += sz[a];
```

84b

p[a] = b;

```
e1a
5eb
        int query() {
ba7
            return ans;
35с
        }
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
61c
            if (u == -1) return:
270
            sz[p[u]] -= sz[u];
546
            p[u] = u;
Odf
            ans++:
456
        }
568 }:
357 int ponta[MAX]; // outra ponta do intervalo ou -1 se for query
4f0 int ans[MAX], n, q;
487 T qu[MAX];
47b void solve(int l = 0, int r = q-1) {
        if (1 >= r) {
0b1
            ans[1] = data::query(); // agora a estrutura ta certa
8 c 0
505
            return;
f77
        int m = (1+r)/2, qnt = 1;
962
fc7
        for (int i = m+1; i <= r; i++) if (ponta[i]+1 and ponta[i] < 1)</pre>
37d
            data::add(qu[i]), qnt++;
221
        solve(1. m):
593
        while (--qnt) data::rollback();
a2c
        for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i] > r)
37d
            data::add(qu[i]), qnt++;
37b
        solve (m+1, r);
        while (qnt --) data::rollback();
281
0d4 }
    Conectividade Dinamica LCT
```

```
cac
        };
        node t[2*MAX]; // MAXN + MAXQ
c53
99e
        map < pair < int , int > , int > aresta;
e4d
        int sz:
95 a
        void prop(int x) {
            if (t[x].rev) {
aa2
f 95
                 swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
            }
693
            t[x].rev = 0:
750
564
        void update(int x) {
            t[x].sub = t[x].val;
e8d
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
8ca
621
                prop(t[x].ch[i]);
78 d
                t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
3 e 4
        }
9bf
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
ed6
        void rotate(int x) {
            int p = t[x].p, pp = t[p].p;
497
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a 76
8fa
            t[x].p = pp, t[p].p = x;
            update(p), update(x);
444
f31
        }
238
        int splay(int x) {
18 c
            while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
0 c 5
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
                rotate(x);
64 f
72 c
aab
            return prop(x), x;
08f
f 16
        int access(int v) {
0eb
            int last = -1:
```

```
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last;
294
        void make_tree(int v, int w=INF) { t[v] = node(w); }
952
82f
        bool conn(int v. int w) {
2 cf
            access(v), access(w);
b9b
            return v == w ? true : t[v].p != -1;
ec0
277
        void rootify(int v) {
5 e 3
            access(v):
a02
            t[v].rev ^= 1:
a05
        int query(int v, int w) {
a1d
b54
            rootify(w), access(v);
249
            return t[v].sub;
c28
204
        void link_(int v, int w) {
821
            rootifv(w):
389
            t[w].p = v;
523
        void link(int v, int w, int x) { // v--w com peso x
6b8
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
ab6
            make tree(id, x):
c88
            link_(v, id), link_(id, w);
984
e63
        void cut_(int v, int w) {
b54
            rootify(w), access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
7cd
        void cut(int v. int w) {
031
            int id = aresta[make pair(v, w)];
b0f
            cut_(v, id), cut_(id, w);
a4a
        }
840
0d7 }
893 void dyn_conn() {
c5f
        int n, q; cin >> n >> q;
        vector <int> p(2*q, -1); // outra ponta do intervalo
d6e
b4f
        for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
fbf
        vector < pair < int , int >> qu(q);
        map < pair < int , int > , int > m;
139
abf
        for (int i = 0; i < q; i++) {
3 c 2
            char c: cin >> c:
ef6
            if (c == '?') continue:
```

```
602
            int a, b; cin >> a >> b; a--, b--;
d11
            if (a > b) swap(a, b);
8a1
            qu[i] = \{a, b\};
8d7
            if (c == '+') {
94b
                p[i] = i+q, p[i+q] = i;
906
                m[make_pair(a, b)] = i;
8a0
            } else {
412
                int j = m[make_pair(a, b)];
ac2
                p[i] = j, p[j] = i;
            }
0da
9e5
        }
447
        int ans = n;
abf
        for (int i = 0; i < a; i++) {
87 d
            if (p[i] == -1) {
                cout << ans << endl; // numero de comp conexos</pre>
886
5 e 2
                 continue;
b35
            }
            int a = qu[i].first, b = qu[i].second;
69 d
c4d
            if (p[i] > i) { // +
                if (lct::conn(a, b)) {
ac5
                     int mi = lct::query(a, b);
18f
993
                     if (p[i] < mi) {</pre>
dd3
                         p[p[i]] = p[i];
5 e 2
                         continue;
474
6f7
                     lct::cut(qu[p[mi]].first, qu[p[mi]].second), ans++;
                     p[mi] = mi;
6ea
9a9
d1d
                lct::link(a, b, p[i]), ans--;
9d0
            } else if (p[i] != i) lct::cut(a, b), ans++; // -
c.03
        }
56a }
     Conj. Indep. Maximo com Peso em Grafo de Intervalo
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento total
//
// O(n log(n))
31e vector<int> ind_set(vector<tuple<int, int, int>>& v) {
        vector < tuple < int , int , int >> w;
b27
f 14
        for (int i = 0; i < v.size(); i++) {</pre>
```

w.push\_back(tuple(get<0>(v[i]), 0, i));

w.push\_back(tuple(get<1>(v[i]), 1, i));

e 85

6f0

17 f

}

```
d1d
        sort(w.begin(), w.end());
844
        vector < int > nxt(v.size());
        vector < pair < ll, int >> dp(v.size());
c22
0eb
        int last = -1:
        for (auto [fim, t, i] : w) {
723
25a
            if (t == 0) {
4 ca
                nxt[i] = last;
5 e 2
                 continue;
5 f d
78b
            dp[i] = \{0, 0\};
cb8
            if (last != -1) dp[i] = max(dp[i], dp[last]);
            pair<11, int> pega = {get<2>(v[i]), -(get<1>(v[i]) -
911
   get<0>(v[i]) + 1)};
            if (nxt[i] != -1) pega.first += dp[nxt[i]].first,
5d3
   pega.second += dp[nxt[i]].second;
            if (pega > dp[i]) dp[i] = pega;
b08
            else nxt[i] = last;
7 cb
            last = i;
381
b7c
        }
977
        pair<11, int > ans = \{0, 0\};
919
        int idx = -1:
        for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans =
   dp[i], idx = i;
        vector < int > ret;
4 b 8
        while (idx != -1) {
fdd
d69
            if (get < 2 > (v[idx]) > 0 and
                (nxt[idx] == -1 or get<1>(v[nxt[idx]]) <</pre>
   get<0>(v[idx]))) ret.push_back(idx);
            idx = nxt[idx];
e4f
042
0ea
        sort(ret.begin(), ret.end());
        return ret;
edf
c4d }
```

#### 7.10 Convex Hull Dinamico

```
// insert - O(log n) amortizado
// is_inside - O(log n)

Ob9 struct upper {
    af8     set < pt > se;
    80b     set < pt >:: iterator it;

25c     int is_under(pt p) { // 1 -> inside ; 2 -> border
    fe0         it = se.lower_bound(p);
```

```
633
            if (it == se.end()) return 0:
            if (it == se.begin()) return p == *it ? 2 : 0;
a94
            if (ccw(p, *it, *prev(it))) return 1;
ca0
            return ccw(p, *prev(it), *it) ? 0 : 2;
402
dba
        void insert(pt p) {
eaa
712
            if (is_under(p)) return;
            if (it != se.end()) while (next(it) != se.end() and
a86
   !ccw(*next(it), *it, p))
316
                it = se.erase(it);
be3
            if (it != se.begin()) while (--it != se.begin() and
   !ccw(p. *it. *prev(it)))
                it = se.erase(it):
316
            se.insert(p);
0c8
5da
        }
750 }:
06f struct dyn_hull {
d93
        upper U, L;
333
        int is_inside(pt p) {
632
            int u = U.is_under(p), l = L.is_under({-p.x, -p.y});
4 c.c
            if (!u or !1) return 0:
fc0
            return max(u, 1);
478
eaa
        void insert(pt p) {
86 c
            U.insert(p);
925
            L.insert({-p.x, -p.y});
64 b
        }
285
        int size() {
7 c.2
            int ans = U.se.size() + L.se.size();
            return ans <= 2 ? ans/2 : ans-2;
1 c.9
ad5
65e };
```

# 7.11 Distancia maxima entre dois pontos

```
// max_dist2(v) - O(n log(n))
// max_dist_manhattan - O(n)

// Quadrado da Distancia Euclidiana (precisa copiar convex_hull, ccw e pt)
859 ll max_dist2(vector<pt> v) {
221     v = convex_hull(v);
```

```
a 14
        if (v.size() <= 2) return dist2(v[0], v[1%v.size()]);</pre>
04b
        11 \text{ ans} = 0:
323
        int n = v.size(), j = 0;
        for (int i = 0; i < n; i++) {
603
            while (! ccw(v[(i+1)\%n] - v[i], pt(0, 0), v[(j+1)\%n] - v[j])) j
057
   = (j+1) \%n;
            ans = \max(\{ans, dist2(v[i], v[j]), dist2(v[(i+1)%n],
e7a
   v[j])});
1f6
ba7
        return ans:
bda }
// Distancia de Manhattan
c51 template < typename T> T max_dist_manhattan(vector < pair < T, T>> v) {
8eb
        T min_sum, max_sum, min_dif, max_dif;
4 f 5
        min_sum = max_sum = v[0].first + v[0].second;
271
        min_dif = max_dif = v[0].first - v[0].second;
        for (auto [x, y] : v) {
c25
1cb
            min_sum = min(min_sum, x+y);
683
            max sum = max(max sum, x+v):
782
            min_dif = min(min_dif, x-y);
af7
            max_dif = max(max_dif, x-y);
e3a
9f0
        return max(max_sum - min_sum, max_dif - min_dif);
4e9 }
7.12 Distinct Range Query
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
789 namespace perseg { };
53d int qt[MAX];
edc void build(vector<int>& v) {
3 d2
        int n = v.size();
        perseg::build(n);
16b
        map < int , int > last;
663
        int at = 0;
05е
```

for (int i = 0; i < n; i++) {

perseg::update(i, 1);

at++;

if (last.count(v[i])) {

perseg::update(last[v[i]], -1);

603

817

a58

69a

d1f

4 f 2

# 7.13 Distinct Range Query com Update

```
// build - O(n log(n))
// query - 0(log^2(n))
// update - 0(log^2(n))
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
def
         using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update >;
3 a 1
042 int v[MAX], n, nxt[MAX], prv[MAX];
f60 map<int, set<int> > ocor;
e04 namespace bit {
         ord_set <pair <int, int>> bit[MAX];
0a8
         void build() {
3e1
             for (int i = 1; i <= n; i++) bit[i].insert({nxt[i-1],
   i-1});
             for (int i = 1; i <= n; i++) {</pre>
78a
edf
                 int j = i + (i\&-i);
d03
                 if (j <= n) for (auto x : bit[i]) bit[j].insert(x);</pre>
5cb
        }
af6
d3f
        int pref(int p, int x) {
7 c 9
             int ret = 0;
bbf
             for (; p; p -= p&-p) ret += bit[p].order_of_key({x, -INF});
edf
             return ret;
        }
0e1
d50
         int query(int 1, int r, int x) {
e 55
             return pref(r+1, x) - pref(1, x);
9b4
        }
ff3
         void update(int p, int x) {
f 17
             int p2 = p;
```

```
5 e d
            for (p++; p \le n; p += p\&-p) {
                 bit[p].erase({nxt[p2], p2});
ca8
f6b
                 bit[p].insert({x, p2});
            }
3df
151
        }
c63 }
0a8 void build() {
383
        for (int i = 0; i < n; i++) nxt[i] = INF;</pre>
7 b3
        for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
d07
        vector < pair < int , int >> t;
348
        for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
3fd
        sort(t.begin(), t.end());
        for (int i = 0; i < n; i++) {</pre>
603
b40
            if (i and t[i].first == t[i-1].first)
565
                 prv[t[i].second] = t[i-1].second:
            if (i+1 < n and t[i].first == t[i+1].first)</pre>
a8b
                 nxt[t[i].second] = t[i+1].second;
12f
        }
48d
        for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
a23
1 d7
        bit::build();
d44 }
aae void muda(int p, int x) {
        bit::update(p, x);
f92
c3d
        nxt[p] = x:
97c }
4ea int query(int a, int b) {
a0a
        return b-a+1 - bit::query(a, b, b+1);
511 }
ff3 void update(int p, int x) { // mudar valor na pos. p para x
        if (prv[p] > -INF) muda(prv[p], nxt[p]);
c0b
4ae
        if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
5bf
        ocor[v[p]].erase(p);
        if (!ocor[x].size()) {
4 b 4
19d
            muda(p, INF);
8 d 4
            prv[p] = -INF;
f6c
        } else if (*ocor[x].rbegin() < p) {</pre>
5 b 5
            int i = *ocor[x].rbegin();
f64
            prv[p] = i;
            muda(p, INF);
19d
5f2
            muda(i, p);
```

```
f36
        } else {
d46
            int i = *ocor[x].lower_bound(p);
33f
            if (prv[i] > -INF) {
f 17
                muda(prv[i], p);
8f9
                prv[p] = prv[i];
            } else prv[p] = -INF;
bc4
523
            prv[i] = p;
597
            muda(p, i);
2ac
c.96
        v[p] = x; ocor[x].insert(p);
38e }
```

#### 7.14 Dominator Points

```
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
// o multiset 'quina' guarda informação sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto nao dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
//
// Complexidades:
// is_dominated - O(log(n))
// insert - O(log(n)) amortizado
// query - 0(1)
e2a struct dominator_points {
baf
        set < pair < int , int >> se;
        multiset < int > quina;
4dd
        bool is_dominated(pair < int, int > p) {
a85
80f
             auto it = se.lower_bound(p);
633
             if (it == se.end()) return 0;
ab4
             return it->second >= p.second;
28f
        void mid(pair<int, int> a, pair<int, int> b, bool rem) {
99b
            pair < int , int > m = {a.first+1, b.second+1};
29a
b19
             int val = m.first + m.second;
638
             if (!rem) quina.insert(val);
731
             else quina.erase(quina.find(val));
241
        }
7 c 4
        bool insert(pair<int, int> p) {
```

```
fb4
            if (is_dominated(p)) return 0;
            auto it = se.lower_bound(p);
80f
ca9
            if (it != se.begin() and it != se.end())
                mid(*prev(it), *it, 1);
d4a
            while (it != se.begin()) {
1fa
049
                it--:
23c
                if (it->second > p.second) break;
                if (it != se.begin()) mid(*prev(it), *it, 1);
b86
316
                it = se.erase(it);
            }
acd
433
            it = se.insert(p).first;
69e
            if (it != se.begin()) mid(*prev(it), *it, 0);
96d
            if (next(it) != se.end()) mid(*it, *next(it), 0);
6a5
            return 1:
       }
688
5eb
        int query() {
956
            if (!quina.size()) return INF;
            return *quina.begin();
add
b8b
       }
09f };
```

#### 7.15 DP de Dominação 3D

```
// Computa para todo ponto i,
// dp[i] = 1 + max_{j dominado por i} dp[j]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
//
// O(n log^2 n), O(n) de memoria
c53 void lis2d(vector<vector<tuple<int, int, int>>>& v, vector<int>&
   dp, int 1, int r) {
        if (1 == r) {
893
            for (int i = 0; i < v[1].size(); i++) {</pre>
56f
8 b 5
                int ii = get<2>(v[1][i]);
                dp[ii] = max(dp[ii], 1);
1 ce
4b0
            }
505
            return;
3 e 4
        int m = (1+r)/2;
ee4
62b
        lis2d(v, dp, 1, m);
325
        vector < tuple < int , int , int >> vv[2];
d44
        vector < int > Z:
871
        for (int i = 1; i <= r; i++) for (auto it : v[i]) {
            vv[i > m].push_back(it);
2ef
```

```
042
             Z.push_back(get<1>(it));
0d1
e9f
        sort(vv[0].begin(), vv[0].end());
        sort(vv[1].begin(), vv[1].end());
9b5
0d1
        sort(Z.begin(), Z.end());
        auto get_z = [&](int z) { return lower_bound(Z.begin(),
573
   Z.end(), z) - Z.begin(); };
        vector < int > bit(Z.size());
        int i = 0:
181
e9a
        for (auto [v, z, id] : vv[1]) {
6bd
             while (i < vv[0].size() and get<0>(vv[0][i]) < v) {</pre>
397
                 auto \lceil v2, z2, id2 \rceil = vv \lceil 0 \rceil \lceil i++ \rceil;
                 for (int p = get_z(z2)+1; p <= Z.size(); p += p&-p)</pre>
ea0
                     bit [p-1] = max(bit [p-1], dp[id2]);
300
82 c
            }
d3b
             int q = 0;
             for (int p = get_z(z); p; p -= p\&-p) q = max(q, bit[p-1]);
fd9
614
             dp[id] = max(dp[id], q + 1);
acc
        }
        lis2d(v, dp, m+1, r);
c25
4d6 }
4de vector<int> solve(vector<tuple<int, int, int>> v) {
3d2
        int n = v.size();
        vector<tuple<int, int, int, int>> vv;
cd4
        for (int i = 0; i < n; i++) {</pre>
603
             auto [x, y, z] = v[i];
9be
             vv.emplace_back(x, y, z, i);
5bb
64 c
        sort(vv.begin(), vv.end());
bd3
e 11
        vector < vector < tuple < int , int , int >>> V;
        for (int i = 0: i < n: i++) {
603
             int j = i;
a5b
808
             V.emplace_back();
c 0 1
             while (j < n and get <0>(vv[j]) == get <0>(vv[i])) {
ba6
                 auto [x, y, z, id] = vv[j++];
cbb
                 V.back().emplace_back(y, z, id);
8bd
            i = j-1;
452
ac4
388
        vector < int > dp(n);
        lis2d(V, dp, 0, V.size()-1);
839
898
        return dp;
b0a }
```

## 7.16 Gray Code

```
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// 0(2^n)

df6 vector<int> gray_code(int n) {
    vector<int> ret(1<<n);
    for (int i = 0; i < (1<<n); i++) ret[i] = i^(i>>1);
edf return ret;
840 }
```

#### 7.17 Half-plane intersection

```
// Cada half-plane eh identificado por uma reta e a regiao ccw a ela
//
// O(n log n)
f4f vector <pt> hp_intersection(vector <line> &v) {
        deque<pt> dq = {{INF, INF}, {-INF, INF}, {-INF, -INF}, {INF,
   -INF}};
d41 #warning considerar trocar por compare_angle
        sort(v.begin(), v.end(), [&](line r, line s) { return
   angle(r.q-r.p) < angle(s.q-s.p); });
        for(int i = 0; i < v.size() and dq.size() > 1; i++) {
5 e 9
c69
            pt p1 = dq.front(), p2 = dq.back();
6 c 6
            while (dq.size() and !ccw(v[i].p, v[i].q, dq.back()))
                p1 = dq.back(), dq.pop_back();
47b
0a2
            while (dq.size() and !ccw(v[i].p, v[i].q, dq.front()))
7cf
                p2 = dq.front(), dq.pop_front();
4 d 9
            if (!dq.size()) break;
            if (p1 == dq.front() and p2 == dq.back()) continue;
606
c9b
            dq.push_back(inter(v[i], line(dq.back(), p1)));
            dq.push_front(inter(v[i], line(dq.front(), p2)));
65с
            if (dq.size() > 1 and dq.back() == dq.front())
fdd
   dq.pop_back();
4 d8
b2b
        return vector<pt>(dq.begin(), dq.end());
f56 }
```

# 7.18 Heap Sort

```
// O(n log n)
f18 void down(vector<int>& v, int n, int i) {
        while ((i = 2*i+1) < n) {
583
            if (i+1 < n and v[i] < v[i+1]) i++;</pre>
b27
            if (v[i] < v[(i-1)/2]) break;
322
            swap(v[i], v[(i-1)/2]);
170
        }
724 }
eb6 void heap_sort(vector<int>& v) {
3d2
        int n = v.size();
        for (int i = n/2-1; i \ge 0; i--) down(v, n, i);
61 d
917
        for (int i = n-1; i > 0; i--)
37 f
            swap(v[0], v[i]), down(v, i, 0);
b33 }
```

### 7.19 Hungaro

```
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
//
// O(n^3)
a6a template < typename T> struct hungarian {
1a8
a08
         vector < vector < T >> a;
f36
         vector <T> u, v;
5ff
         vector < int > p, way;
f1e
         T inf;
c3f
         hungarian(int n_): n(n_{-}), u(n+1), v(n+1), p(n+1), way(n+1) {
b2f
             a = vector < vector < T >> (n, vector < T > (n));
1f3
             inf = numeric_limits <T>::max();
78 f
d67
         pair < T, vector < int >> assignment() {
78a
             for (int i = 1; i <= n; i++) {
8 c 9
                 p[0] = i;
625
                 int j0 = 0;
                 vector <T> minv(n+1, inf);
ce7
241
                 vector < int > used(n+1, 0);
016
                 do {
```

```
472
                     used[j0] = true;
d24
                     int i0 = p[j0], j1 = -1;
7 e 5
                     T delta = inf;
9ac
                     for (int j = 1; j <= n; j++) if (!used[j]) {
                          T \text{ cur } = a[i0-1][j-1] - u[i0] - v[j];
7bf
                          if (cur < minv[j]) minv[j] = cur, way[j] = j0;</pre>
9f2
821
                          if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
4d1
f63
                     for (int j = 0; j <= n; j++)
                          if (used[j]) u[p[j]] += delta, v[j] -= delta;
2 c5
6ec
                          else minv[j] -= delta;
6 d 4
                     j0 = j1;
f4f
                 } while (p[j0] != 0);
016
                 do {
4 c 5
                     int j1 = way[j0];
0d7
                     p[j0] = p[j1];
6 d 4
                     j0 = j1;
886
                 } while (j0);
38d
306
             vector < int > ans(n);
             for (int j = 1; j <= n; j++) ans[p[j]-1] = j-1;
6db
             return make_pair(-v[0], ans);
da3
979
        }
64c };
```

## 7.20 Inversion Count

```
// Computa o numero de inversoes para transformar
// l em r (se nao tem como, retorna -1)
//
// O(n log(n))
37b template < typename T > 1l inv_count(vector < T > 1, vector < T > r = {}) {
bb6
        if (!r.size()) {
796
             r = 1:
             sort(r.begin(), r.end());
1bc
        }
dfb
874
        int n = 1.size();
        vector < int > v(n), bit(n);
8 c 0
4 e 9
        vector < pair < T, int >> w;
61c
        for (int i = 0; i < n; i++) w.push_back({r[i], i+1});</pre>
d1d
        sort(w.begin(), w.end());
603
        for (int i = 0; i < n; i++) {</pre>
bf3
             auto it = lower_bound(w.begin(), w.end(), make_pair(l[i],
    0));
```

```
1bf
             if (it == w.end() or it->first != l[i]) return -1; // nao
    da
             v[i] = it->second;
962
6c0
             it -> second = -1;
        }
964
04b
        11 \text{ ans} = 0:
        for (int i = n-1; i >= 0; i--) {
45 b
2d9
             for (int j = v[i]-1; j; j -= j k-j) ans += bit[j];
3 a 1
             for (int j = v[i]; j < n; j += j&-j) bit[j]++;
ebe
        }
ba7
        return ans;
eef }
7.21 LIS - recupera
// Calcula e retorna uma LIS
//
// O(n.log(n))
121 template < typename T > vector < T > lis(vector < T > & v) {
1fa
        int n = v.size(), m = -1;
f0c
        vector < T > d(n+1, INF);
aec
        vector < int > l(n);
007
        d[0] = -INF;
603
        for (int i = 0; i < n; i++) {
             // Para non-decreasing use upper_bound()
             int t = lower_bound(d.begin(), d.end(), v[i]) - d.begin();
4fd
             d[t] = v[i], l[i] = t, m = max(m, t);
3ad
        }
89 c
4ff
        int p = n;
        vector <T> ret;
5a9
        while (p--) if (l[p] == m) {
cdf
883
             ret.push_back(v[p]);
76b
             m - - ;
f83
969
        reverse (ret.begin(), ret.end());
edf
        return ret;
474 }
```

#### 7.22 LIS - tamanho

```
// Calcula o tamanho da LIS
// O(n log(n))
84b template < typename T > int lis(vector <T > &v){
        vector <T> ans:
5 e 0
        for (T t : v) {
            // Para non-decreasing use upper_bound()
            auto it = lower_bound(ans.begin(), ans.end(), t);
fe6
            if (it == ans.end()) ans.push_back(t);
d7f
b94
            else *it = t;
1f5
1eb
        return ans.size();
402 }
```

## 7.23 Minimum Enclosing Circle

```
// O(n) com alta probabilidade
22c const double EPS = 1e-12:
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
b2a struct pt {
662
        double x, v;
be7
        pt(double x_{=} = 0, double y_{=} = 0) : x(x_{=}), y(y_{=}) \{ \}
        pt operator + (const pt& p) const { return pt(x+p.x, y+p.y); }
7af
b23
        pt operator - (const pt& p) const { return pt(x-p.x, y-p.y); }
254
        pt operator * (double c) const { return pt(x*c, y*c); }
701
        pt operator / (double c) const { return pt(x/c, y/c); }
54d }:
2f9 double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
dd5 double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
e7c double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
3f4 pt center(pt p, pt q, pt r) {
5 d 9
        pt a = p-r, b = q-r;
        pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
e84
        return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x), c)) /
   cross(a. b):
fc8 }
aa8 struct circle {
f41
        pt cen;
c12
        double r;
```

```
898
        circle(pt cen_, double r_) : cen(cen_), r(r_) {}
83 c
        circle(pt a, pt b, pt c) {
13 d
            cen = center(a, b, c);
1f1
            r = dist(cen, a);
bc1
cd5
        bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
2a6 }:
806 circle minCirc(vector<pt> v) {
        shuffle(v.begin(), v.end(), rng);
ae0
        circle ret = circle(pt(0, 0), 0);
618
        for (int i = 0; i < v.size(); i++) if (!ret.inside(v[i])) {</pre>
16a
            ret = circle(v[i], 0):
f 1 1
            for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
                ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
881
b8c
                for (int k = 0; k < j; k++) if (!ret.inside(v[k]))
43 f
                     ret = circle(v[i], v[j], v[k]);
5f8
            }
6a1
        }
edf
        return ret:
eba }
     Minkowski Sum
```

```
// Computa A+B = \{a+b : a \setminus in A, b \setminus in B\}, em que
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
//
// O(|A|+|B|)
539 vector<pt> minkowski(vector<pt> p, vector<pt> q) {
051
        auto fix = [](vector<pt>& P) {
515
             rotate(P.begin(), min_element(P.begin(), P.end()),
   P.end());
018
            P.push_back(P[0]), P.push_back(P[1]);
f 24
        }:
889
        fix(p), fix(q);
        vector < pt > ret;
8af
692
        int i = 0, i = 0;
        while (i < p.size()-2 or j < q.size()-2) {
2ee
898
             ret.push_back(p[i] + q[j]);
732
             auto c = ((p[i+1] - p[i]) ^ (q[j+1] - q[j]));
ebc
             if (c >= 0) i = min<int>(i+1, p.size()-2);
81 e
             if (c <= 0) j = min<int>(j+1, q.size()-2);
9ff
        }
edf
        return ret;
```

```
d7c }
c3e ld dist_convex(vector<pt> p, vector<pt> q) {
        for (pt \& i : p) i = i * -1;
dc2
        auto s = minkowski(p, q);
44c
        if (inpol(s, pt(0, 0))) return 0;
95d
921
       ld ans = DINF:
073
        for (int i = 0; i < s.size(); i++) ans = min(ans,
f04
                disttoseg(pt(0, 0), line(s[(i+1)%s.size()], s[i])));
ba7
        return ans:
beb }
7.25 MO
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))
0d2 const int MAX = 1e5+10:
6ff const int SQ = sqrt(MAX);
b69 int v[MAX];
b65 int ans, freq[MAX];
9da inline void insert(int p) {
        int o = v[p];
ae0
591
       freq[o]++;
        ans += (freq[o] == 1);
992
21d }
a25 inline void erase(int p) {
ae0
       int o = v[p];
        ans -= (freq[o] == 1);
7ee
       freq[o]--;
ba2
dc7 }
e51 inline ll hilbert(int x, int y) {
        static int N = 1 << ( builtin clz(0) - builtin clz(MAX));
71e
100
        int rx, ry, s;
b72
       11 d = 0:
        for (s = N/2; s > 0; s /= 2) {
c95
           rx = (x \& s) > 0, ry = (y \& s) > 0;
e3e
            d += s * 11(s) * ((3 * rx) ^ ry);
d2e
           if (ry == 0) {
5aa
                if (rx == 1) x = N-1 - x, y = N-1 - y;
```

```
9dd
                 swap(x, y);
 e2d
 888
         }
 be2
         return d;
 7fa }
 bac #define HILBERT true
 617 vector<int> MO(vector<pair<int, int>> &q) {
 c3b
         ans = 0;
         int m = q.size();
 c23
 3f8
         vector < int > ord(m);
         iota(ord.begin(), ord.end(), 0);
 6a6 #if HILBERT
 8 c 4
         vector < 11 > h(m):
         for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first,
         sort(ord.begin(), ord.end(), [&](int 1, int r) { return h[1] <
    h[r]: }):
8c1 #else
         sort(ord.begin(), ord.end(), [&](int 1, int r) {
             if (q[1].first / SQ != q[r].first / SQ) return q[1].first
    < q[r].first;
 Odb
             if ((q[1].first / SQ) % 2) return q[1].second >
    q[r].second;
             return q[1].second < q[r].second;</pre>
 a 66
bec
f2e #endif
 435
         vector < int > ret(m);
         int 1 = 0, r = -1;
 3d9
         for (int i : ord) {
 8b0
 6 c 6
             int ql, qr;
 4f5
             tie(ql, qr) = q[i];
026
             while (r < qr) insert(++r);</pre>
 232
             while (l > ql) insert(--l);
75 e
             while (1 < q1) erase(1++);
             while (r > qr) erase(r--);
fe8
 381
             ret[i] = ans;
116
 edf
         return ret;
fb7 }
7.26 MO - DSU
// Dado uma lista de arestas de um grafo, responde
// para cada query(1, r), quantos componentes conexos
```

```
// o grafo tem se soh considerar as arestas l, l+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura rollbackavel
// O(m sqrt(q) log(n))
8d3 struct dsu {
553
        int n. ans:
2 e 3
        vector < int > p, sz;
        stack < int > S;
ee6
4b8
        dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) {
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;</pre>
aae
1 b 1
        int find(int k) {
            while (p[k] != k) k = p[k];
006
839
            return k;
c1e
        }
553
        void add(pair<int, int> x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
            ans - -:
            if (sz[a] > sz[b]) swap(a, b);
3 c 6
4 c2
            S.push(a);
582
            sz[b] += sz[a]:
84b
            p[a] = b;
720
        int query() { return ans; }
35 c
        void rollback() {
5cf
465
            int u = S.top(); S.pop();
            if (u == -1) return;
61 c
270
            sz[p[u]] -= sz[u];
            p[u] = u;
546
0 df
            ans++:
456
        }
9c1 };
1a8 int n;
e93 vector<pair<int, int>> ar;
// 9d242b
617 vector < int > MO (vector < pair < int , int >> &q) {
547
        int SQ = sqrt(q.size()) + 1;
       int m = q.size();
c23
3f8
       vector < int > ord(m);
be8
       iota(ord.begin(), ord.end(), 0);
d01
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
```

```
if (q[1].first / SQ != q[r].first / SQ) return
9 c 9
   q[l].first < q[r].first;
                return q[1].second < q[r].second;</pre>
a66
b90
                }):
435
        vector < int > ret(m);
3bd
        dsu small(n):
dd5
        for (int i = 0; i < m; i++) {</pre>
5ec
            auto [1, r] = q[ord[i]];
            if (1 / SQ == r / SQ) {
асс
00с
                for (int k = 1; k <= r; k++) small.add(ar[k]);</pre>
b99
                ret[ord[i]] = small.query();
64 e
                for (int k = 1; k \le r; k++) small.rollback();
259
        }
6b0
dd5
        for (int i = 0; i < m; i++) {
176
            dsu D(n):
ae9
            int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
e 25
            int last_r = fim;
ebc
            int j = i-1;
00с
            while (j+1 < m and q[ord[j+1]].first / SQ ==</pre>
   q[ord[i]].first / SQ) {
a0e
                auto [1, r] = q[ord[++j]];
f58
                if (1 / SQ == r / SQ) continue;
59b
                while (last r < r) D.add(ar[++last r]):
                for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
2cf
9b2
                ret[ord[j]] = D.query();
572
                for (int k = 1; k \le fim; k++) D.rollback();
9c8
            }
bdf
            i = j;
e99
edf
        return ret;
9d2 }
7.27 MO em Arvores
// Problema que resolve: https://www.spoj.com/problems/COT2/
// Complexidade sendo c = O(update) e SQ = sqrt(n):
// O((n + q) * sqrt(n) * c)
```

```
1bc const int MAX = 40010, SQ = 400;
042 vector < int > g[MAX];
c54 namespace LCA { ... }
249 int in [MAX], out [MAX], vtx [2 * MAX];
81b bool on [MAX];
4c3 int dif, freq[MAX];
9e2 vector < int > w;
d9a void dfs(int v. int p. int &t) {
        vtx[t] = v, in[v] = t++;
659
        for (int u : g[v]) if (u != p) {
18e
            dfs(u, v, t);
c53
e0f
        }
217
        vtx[t] = v, out[v] = t++;
42b }
e5f void update(int p) { // faca alteracoes aqui
        int v = vtx[p];
bbc
        if (not on[v]) { // insere vtx v
0ec
31c
            dif += (freq[w[v]] == 0);
b20
            freq[w[v]]++;
cf7
        else { // retira o vertice v
4 e 6
0a9
            dif -= (freq[w[v]] == 1);
            freq[w[v]]--;
fd3
2 c8
        on[v] = not on[v]:
73e
ea9 }
a3a vector < tuple < int, int, int >> build_queries (const vector < pair < int,
   int>>& q) {
       LCA::build(0);
ea6
        vector<tuple<int, int, int>> ret;
f77
aa9
        for (auto [1, r] : q){
d24
            if (in[r] < in[1]) swap(1, r);</pre>
6f9
            int p = LCA::lca(l, r);
            int init = (p == 1) ? in[1] : out[1];
826
07a
            ret.emplace_back(init, in[r], in[p]);
b0e
edf
        return ret:
8e6 }
f31 vector<int> mo_tree(const vector<pair<int, int>>& vq){
```

```
6bb
         int t = 0;
         dfs(0, -1, t);
dab
af1
         auto g = build_queries(vg);
f 48
         vector < int > ord(q.size());
be8
         iota(ord.begin(), ord.end(), 0);
         sort(ord.begin(), ord.end(), [&] (int 1, int r) {
d01
d8d
             int bl = get<0>(q[1]) / SQ, br = <math>get<0>(q[r]) / SQ;
596
             if (bl != br) return bl < br;</pre>
158
             else if (b1 % 2 == 1) return get<1>(q[1]) < get<1>(q[r]);
f1d
             else return get<1>(q[1]) > get<1>(q[r]);
0a8
        }):
80 e
         memset(freq, 0, sizeof freq);
bf6
         dif = 0:
ff2
         vector < int > ret(q.size());
         int 1 = 0, r = -1;
3d9
        for (int i : ord) {
8b0
             auto [ql, qr, qp] = q[i];
3 c 7
af7
             while (r < qr) update(++r);</pre>
             while (1 > q1) update(--1);
d6b
951
             while (1 < q1) update(1++);
             while (r > qr) update(r--);
6a1
3d8
             if (qp < 1 \text{ or } qp > r) \{ // \text{ se LCA estah entre as pontas} \}
74b
                 update(qp);
2 e 1
                 ret[i] = dif;
74b
                 update(qp);
e 83
0fe
             else ret[i] = dif;
0fd
        }
edf
         return ret:
48d }
7.28 Palindromic Factorization
// Precisa da eertree
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
// O(n log n), considerando alfabeto O(1)
```

070 struct eertree { ... };

```
0e7 ll factorization(string s) {
        int n = s.size(), sz = 2;
b19
580
        eertree PT(n);
147
        vector < int > diff(n+2), slink(n+2), sans(n+2), dp(n+1);
0ec
        dp[0] = 1:
        for (int i = 1; i <= n; i++) {
78a
c58
            PT.add(s[i-1]);
а7с
            if (PT.size()+2 > sz) {
6 c 4
                diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
                if (diff[sz] == diff[PT.link[sz]])
241
d6f
                    slink[sz] = slink[PT.link[sz]];
f53
                else slink[sz] = PT.link[sz];
eb9
                sz++:
f6a
            }
911
            for (int v = PT.last; PT.len[v] > 0; v = slink[v]) {
                sans[v] = dp[i - (PT.len[slink[v]] + diff[v])];
297
85d
                if (diff[v] == diff[PT.link[v]])
f20
                    sans[v] = (sans[v] + sans[PT.link[v]]) % MOD;
071
                dp[i] = (dp[i] + sans[v]) % MOD;
            }
е5е
fc0
5f0
        return dp[n];
3a7 }
```

# 7.29 Parsing de Expressao

```
// Operacoes associativas a esquerda por default
// Para mudar isso, colocar em r assoc
// Operacoes com maior prioridade sao feitas primeiro
cc1 bool blank(char c) {
f34
        return c == ' ';
ec3 }
8e4 bool is_unary(char c) {
        return c == '+' or c == '-';
f9c
b6b }
76d bool is_op(char c) {
010
        if (is_unary(c)) return true;
        return c == '*' or c == '/' or c == '+' or c == '-';
31c
4e4 }
fa3 bool r_assoc(char op) {
        // operator unario - deve ser assoc. a direita
cf0
        return op < 0;</pre>
```

```
c5c }
79d int priority(char op) {
        // operator unario - deve ter precedencia maior
        if (op < 0) return INF;</pre>
103
727
        if (op == '*' or op == '/') return 2:
        if (op == '+' or op == '-') return 1;
439
daa
        return -1;
966 }
c15 void process_op(stack<int>& st, stack<int>& op) {
        char o = op.top(); op.pop();
91 c
        if (o < 0) {
4e6
            o *= -1:
            int 1 = st.top(); st.pop();
1 e 2
Off
            if (o == '+') st.push(1);
7e9
            if (o == '-') st.push(-1);
320
        } else {
14 c
            int r = st.top(); st.pop();
            int 1 = st.top(): st.pop();
1 e 2
1 e 4
            if (o == '*') st.push(l * r);
f55
            if (o == ',') st.push(1 / r);
605
            if (o == '+') st.push(l + r);
c.40
            if (o == '-') st.push(1 - r);
0aa
        }
2b2 }
439 int eval(string& s) {
212
        stack < int > st, op;
d0c
        bool un = true:
1cf
        for (int i = 0; i < s.size(); i++) {</pre>
            if (blank(s[i])) continue:
68 d
139
            if (s[i] == '(') {
367
                op.push('(');
99 d
                un = true;
b88
            } else if (s[i] == ')') {
709
                while (op.top() != '(') process_op(st, op);
75 e
                op.pop();
ce2
                un = false:
003
            } else if (is_op(s[i])) {
                char o = s[i];
4d0
37 с
                if (un and is_unary(o)) o *= -1;
ae3
                while (op.size() and (
cd6
                             (!r_assoc(o) and priority(op.top()) >=
   priority(o)) or
```

```
c41
                             (r_assoc(o) and priority(op.top()) >
   priority(o)))
                    process_op(st, op);
c47
                op.push(o);
99d
                un = true;
196
            } else {
da8
                int val = 0:
                while (i < s.size() and isalnum(s[i]))
c2b
8a3
                    val = val * 10 + s[i++] - '0';
169
                i--:
25 d
                st.push(val);
                un = false:
442
            }
b19
        }
        while (op.size()) process_op(st, op);
7 f 6
123
        return st.top();
05c }
```

## 7.30 RMQ com Divide and Conquer

```
// Responde todas as queries em
// O(n log(n))
f74 typedef pair < pair < int , int > , int > iii;
7c6 #define f first
Oab #define s second
87d int n, q, v[MAX];
e3f iii qu[MAX];
aeb int ans[MAX], pref[MAX], sulf[MAX];
0e3 void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
8a3
        if (1 > r or q1 > qr) return;
        int m = (1+r)/2;
ee4
        int qL = partition(qu+ql, qu+qr+1, [=](iii x){return x.f.s <</pre>
1 b 1
   m; }) - qu;
        int qR = partition(qu+qL, qu+qr+1, [=](iii x){return x.f.f
   <=m;}) - qu;
3cd
        pref[m] = sulf[m] = v[m];
        for (int i = m-1; i >= 1; i--) pref[i] = min(v[i], pref[i+1]);
9f9
        for (int i = m+1; i <= r; i++) sulf[i] = min(v[i], sulf[i-1]);</pre>
ea8
b2a
        for (int i = qL; i < qR; i++)</pre>
f3a
            ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
```

```
364
        solve(l, m-1, ql, qL-1), solve(m+1, r, qR, qr);
13e }
7.31 Segment Intersection
// Verifica, dado n segmentos, se existe algum par de segmentos
// que se intersecta
//
// O(n log n)
6e0 bool operator < (const line& a, const line& b) { // comparador pro
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
    b.p.x))
780
            return ccw(a.p, a.q, b.p);
dc0
        return ccw(a.p, b.q, b.p);
e36 }
8e2 bool has_intersection(vector<line> v) {
576
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
            return interseg(a.first, b.first);
a 08
3e6
        };
        vector<pair<pt, pair<int, int>>> w;
e1b
f 14
        for (int i = 0; i < v.size(); i++) {
876
            if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
            w.push_back({v[i].p, {0, i}});
e 1 d
            w.push_back({v[i].q, {1, i}});
034
220
d1d
        sort(w.begin(), w.end());
7f2
        set < pair < line , int >> se;
        for (auto i : w) {
e 58
            line at = v[i.second.second];
bfd
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
d1e
                if (nxt != se.end() and intersects(*nxt, {at,
    i.second.second})) return 1;
257
                if (nxt != se.begin() and intersects(*(--nxt), {at,
    i.second.second})) return 1;
78 f
                se.insert({at. i.second.second}):
08b
                auto nxt = se.upper_bound({at, i.second.second}), cur
884
    = nxt, prev = --cur;
b64
                if (nxt != se.end() and prev != se.begin()
4fb
                     and intersects (*nxt, *(--prev))) return 1;
```

## 7.32 Sequencia de de Brujin

```
// Se passar sem o terceiro parametro, gera um vetor com valores
// em [0, k) de tamanho k^n de forma que todos os subarrays ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)</pre>
//
// Linear no tamanho da resposta
860 vector < int > de_brujin(int n, int k, int lim = INF) {
b55
        if (k == 1) return vector<int>(lim == INF ? 1 : n, 0);
5f6
        vector < int > 1 = \{0\}, ret; // l eh lyndon word
667
        while (true) {
c86
            if (1.size() == 0) {
1 b 9
                if (lim == INF) break;
daf
                1.push_back(0);
            }
bae
            if (n % 1.size() == 0) for (int i : 1) {
686
728
                ret.push_back(i);
                if (ret.size() == n+lim-1) return ret;
c99
            }
56e
630
            int p = 1.size();
            while (1.size() < n) 1.push_back(1[1.size()%p]);</pre>
905
            while (l.size() and l.back() == k-1) l.pop_back();
e7f
88a
            if (1.size()) 1.back()++;
2ef
edf
        return ret;
197 }
```

### 7.33 Shortest Addition Chain

```
// Computa o menor numero de adicoes para construir
// cada valor, comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore eh tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
```

```
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
3de pair < vector < int > , vector < int >> addition_chain() {
16f
        int MAX = 301:
875
        vector < int > dp(MAX), p(MAX);
        for (int n = 2; n < MAX; n++) {
1ab
7 c 0
             pair < int , int > val = {INF , -1};
212
             for (int i = 1; i < n; i++) for (int j = i; j; j = p[j])
94a
                 if (j == n-i) val = min(val, pair(dp[i]+1, i));
eb3
             tie(dp[n], p[n]) = val:
efe
             if (n == 9) p[n] = 8;
             if (n == 149 \text{ or } n == 233) \text{ dp}[n] --;
ba1
bcd
717
        return {dp, p};
84f }
7.34 Simple Polygon
// Verifica se um poligono com n pontos eh simples
// O(n log n)
6e0 bool operator < (const line& a, const line& b) { // comparador pro
    sweepline
191
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
2.31
    b.p.x))
780
             return ccw(a.p, a.q, b.p);
dc0
        return ccw(a.p, b.q, b.p);
e36 }
6f3 bool simple(vector<pt> v) {
576
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
e72
             if ((a.second+1)%v.size() == b.second or
80e
                 (b.second+1)%v.size() == a.second) return false;
a08
             return interseg(a.first, b.first);
1 c 5
        };
41a
        vector<line> seg;
e1b
        vector<pair<pt, pair<int, int>>> w;
f 14
        for (int i = 0; i < v.size(); i++) {</pre>
0a8
             pt at = v[i], nxt = v[(i+1)\%v.size()];
828
             if (nxt < at) swap(at, nxt);</pre>
937
             seg.push_back(line(at, nxt));
```

```
f7e
            w.push_back({at, {0, i}});
69c
            w.push_back({nxt, {1, i}});
            // casos degenerados estranhos
            if (isinseg(v[(i+2)%v.size()], line(at, nxt))) return 0;
ae8
            if (isinseg(v[(i+v.size()-1)%v.size()], line(at, nxt)))
b88
   return 0;
cba
        }
        sort(w.begin(), w.end());
d1d
7f2
        set <pair <line, int >> se;
        for (auto i : w) {
e58
ff8
            line at = seg[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
                if (nxt != se.end() and intersects(*nxt, {at,
7 c 4
   i.second.second})) return 0;
                if (nxt != se.begin() and intersects(*(--nxt), {at,
b34
   i.second.second})) return 0;
                se.insert({at, i.second.second});
78f
            } else {
537
                auto nxt = se.upper_bound({at, i.second.second}), cur
884
   = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                    and intersects(*nxt, *(--prev))) return 0;
403
                se.erase(cur);
cca
            }
7 be
d17
        }
6 a 5
        return 1;
af3 }
7.35 Steiner Tree
// steiner: retorna o peso da menor arvore que cobre os vertices S
// get_steiner: retorna o valor minimo e as arestas de uma solucao
// se nao tiver solucao retorna LINF
// grafo nao pode ter pesos negativos
// se so tiver peso nas arestas/vertices pode deletar os vw/w no codigo
//
//k = |S|
// 0(3^k * n + 2^k * m \log m)
// otimizacao: joga um vertice x do S fora e pegue a resposta em
```

dp[...][x] e reconstrua a arvore a partir dele

// ta comentado no codigo as mudancas necessarias

1a8 int n; // numero de vertices

```
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
920 11 d[1 << K][MAX]; // dp[mask][v] = arvore minima com o
   subconjunto mask de S e o vertice v
ObO 11 vw[MAX]; // peso do vertice
c8f ll steiner(const vector<int> &S) {
934
        int k = S.size(); // k--;
        for (int mask = 0; mask < (1 << k); mask++) for (int v = 0; v <
   n; v++) d[mask][v] = LINF;
        for (int v = 0; v < n; v++) d[0][v] = vw[v];
6b8
254
        for (int i = 0; i < k; ++i) d[1 << i][S[i]] = vw[S[i]];
042
        for (int mask = 1; mask < (1 << k); mask++) {</pre>
b5b
            for (int a = (mask - 1) & mask: a: a = (a - 1) & mask) {
638
                int b = mask ^ a:
6bf
                if (b > a) break;
84 d
                for (int v = 0; v < n; v++)
2 e 6
                     d[mask][v] = min(d[mask][v], d[a][v] + d[b][v] -
   vw[v]);
2ab
88 c
            priority_queue < pair < ll, int >> pq;
            for (int v = 0; v < n; v++) {
778
6ad
                if (d[mask][v] == LINF) continue;
5 ca
                pq.emplace(-d[mask][v], v);
f2e
            }
265
            while (pq.size()) {
a 25
                auto [ndist, u] = pq.top(); pq.pop();
                if (-ndist > d[mask][u]) continue;
dad
c38
                for (auto [idx, w] : g[u]) if (d[mask][idx] >
   d[mask][u] + w + vw[idx]) {
679
                    d[mask][idx] = d[mask][u] + w + vw[idx];
                    pq.emplace(-d[mask][idx], idx);
a2e
07е
            }
de5
a65
478
        return d[(1 << k) - 1][S[0]]; // S[k]
704 }
d41 #warning se k=1 a solucao eh a folha isolada e a funcao retorna
   edg = \{\}
d41 #warning se k=0 crasha
4d1 pair < ll, vector < pair < int, int >>> get_steiner (const vector < int > &S) {
        int k = S.size(); // k--;
        11 ans = steiner(S):
8ec
c8d
        vector<pair<int,int>> edg;
57f
        stack<pair<int,int>> stk;
09b
        stk.emplace((1 << k) - 1, S[0]); // S[k]
07 d
        while (!stk.empty()) {
```

```
c37
            bool cont = 0;
            auto [mask,u] = stk.top();stk.pop();
9c6
            if ((__builtin_popcount(mask) == 1 and u ==
de2
   S[__bit_width(mask) - 1])) continue;
            for (auto [idx, w] : g[u]){
851
                if (d[mask][u] == d[mask][idx] + w + vw[u]) {
bb4
fc7
                    edg.emplace_back(u, idx);
                    stk.emplace(mask, idx);
8ab
a04
                    cont = true;
c2b
                    break:
342
                }
ed9
            }
3 b 5
            if (cont) continue:
            for (int a = (mask - 1) & mask; a; a = (a - 1) & mask) {
b5b
638
                int b = mask ^ a;
                if (d[mask][u] == d[a][u] + d[b][u] - vw[u]) {
1 e 8
be8
                    stk.emplace(a, u);
c0a
                    stk.emplace(b, u);
                    cont = true;
a04
c2b
                    break:
                }
f52
d29
с5 с
            assert(!mask || cont);
2 b 1
be8
        return {ans, edg};
cf6 }
```

## 7.36 Sweep Direction

```
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 log n)
4b8 void sweep_direction(vector<pt> v) {
        int n = v.size();
3 d 2
163
        sort(v.begin(), v.end(), [](pt a, pt b) {
3 a 5
            if (a.x != b.x) return a.x < b.x;
572
            return a.y > b.y;
79a
b89
        vector < int > at(n);
516
        iota(at.begin(), at.end(), 0);
b79
        vector < pair < int , int >> swapp;
25e
        for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
95f
            swapp.push_back({i, j}), swapp.push_back({j, i});
```

```
269
        sort(swapp.begin(), swapp.end(), [&](auto a, auto b) {
            pt A = rotate90(v[a.first] - v[a.second]);
134
247
            pt B = rotate90(v[b.first] - v[b.second]);
            if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B)) return
615
   a < b:
224
            return compare_angle(A, B);
5e7
        }):
4 e 6
        for (auto par : swapp) {
e 24
            assert(abs(at[par.first] - at[par.second]) == 1);
            int l = min(at[par.first], at[par.second]),
a 96
0d3
                r = n-1 - max(at[par.first], at[par.second]);
            // l e r sao quantos caras tem de cada lado do par de
            // (cada par eh visitado duas vezes)
            swap(v[at[par.first]], v[at[par.second]]);
9cf
1 c 0
            swap(at[par.first], at[par.second]);
241
       }
6bb }
```

## 7.37 Triangulação de Delaunay

```
// Computa a triangulação de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as arestas
// sao as arestas da triangulação
// As arestas partindo de um vertice ja vem ordenadas por angulo.
// ou seja, se o vertice v nao esta no convex hull, (v, v_i, v_{i+1})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo vizinho
// Usa o alg d&c, precisa representar MAX_COOR^4, por isso __int128
// pra aguentar valores ateh 1e9
//
// Propriedades:
// 1 - 0 grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3 pontos
      nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
//
// O(n log n)
2ad typedef struct QuadEdge* Q;
ba5 struct QuadEdge {
53е
        int id;
114
        pt o:
41 e
        Q rot, nxt;
3 e 5
        bool used;
```

```
QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)):
3fc
4ba
            id(id_), o(o_), rot(nullptr), nxt(nullptr), used(false) {}
00f
        Q rev() const { return rot->rot; }
        Q next() const { return nxt; }
сЗс
        Q prev() const { return rot->next()->rot; }
188
        pt dest() const { return rev()->o; }
0d4
828 };
91b Q edge(pt from, pt to, int id_from, int id_to) {
        Q e1 = new QuadEdge(id_from, from);
        Q e2 = new QuadEdge(id_to, to);
61b
        Q e3 = new QuadEdge;
8f6
        Q e4 = new QuadEdge;
5 ca
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4, e2, e1\};
e69
f22
        tie (e1->nxt, e2->nxt, e3->nxt, e4->nxt) = \{e1, e2, e4, e3\};
1 a d
        return e1:
c70 }
d8d void splice(Q a, Q b) {
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
a6f
        swap(a->nxt, b->nxt);
da4
a58 }
167 void del_edge(Q& e, Q ne) { // delete e and assign e <- ne
        splice(e, e->prev());
cc0
        splice(e->rev(), e->rev()->prev());
eec
        delete e->rev()->rot, delete e->rev();
7ea
524
        delete e->rot; delete e;
6 b 2
        e = ne:
18b }
d08 Q conn(Q a. Q b) {
        Q = edge(a->dest(), b->o, a->rev()->id, b->id);
сс5
f2b
        splice(e, a->rev()->prev());
d37
        splice(e->rev(), b);
6bf
        return e;
f78 }
d64 bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf. (a, b, c) ?
        _{-}int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C = c*c - p2;
        return sarea2(p, a, b) * C + sarea2(p, b, c) * A + sarea2(p,
   c, a) * B > 0;
b54 }
540 pair < Q, Q > build_tr(vector < pt > & p, int 1, int r) {
```

```
09 d
        if (r-1+1 \le 3) {
2eb
             Q = edge(p[1], p[1+1], 1, 1+1), b = edge(p[1+1], p[r],
l+1, r);
912
             if (r-1+1 == 2) return \{a, a->rev()\};
             splice(a->rev(), b):
0ec
            11 ar = sarea2(p[1], p[1+1], p[r]);
сЗс
1af
             Q c = ar ? conn(b, a) : 0:
             if (ar >= 0) return {a, b->rev()};
021
9db
             return {c->rev(), c};
        }
bce
ee4
        int m = (1+r)/2;
328
        auto [la, ra] = build_tr(p, l, m);
b93
        auto [lb, rb] = build tr(p, m+1, r):
        while (true) {
667
b99
             if (ccw(lb->o, ra->o, ra->dest())) ra = ra->rev()->prev();
458
             else if (ccw(lb->o, ra->o, lb->dest())) lb =
   lb->rev()->next();
f 97
             else break:
24a
        Q b = conn(lb -> rev(), ra);
ca5
        auto valid = [&](Q e) { return ccw(e->dest(), b->dest(),
713
   b -> o): }:
        if (ra -> o == la -> o) la = b -> rev();
ee1
63 f
        if (1b->o == rb->o) rb = b;
667
        while (true) {
71 e
            0 L = b \rightarrow rev() \rightarrow next():
d11
             if (valid(L)) while (in_c(b->dest(), b->o, L->dest(),
   L->next()->dest()))
                 del_edge(L, L->next());
1 c 0
c76
            Q R = b -> prev();
             if (valid(R)) while (in_c(b->dest(), b->o, R->dest(),
2b0
   R->prev()->dest()))
541
                 del_edge(R, R->prev());
             if (!valid(L) and !valid(R)) break;
a3a
             if (!valid(L) or (valid(R) and in_c(L->dest(), L->o, R->o,
ccd
   R->dest())))
                 b = conn(R, b\rightarrow rev());
36 c
666
             else b = conn(b \rightarrow rev(), L \rightarrow rev());
94 d
a2b
        return {la, rb};
689 }
b58 vector < vector < int >> delaunay (vector < pt > v) {
3d2
        int n = v.size();
397
        auto tmp = v;
        vector<int> idx(n);
135
295
        iota(idx.begin(), idx.end(), 0);
```

```
fe9
        sort(idx.begin(), idx.end(), [&](int 1, int r) { return v[1] <</pre>
   v[r]; });
        for (int i = 0; i < n; i++) v[i] = tmp[idx[i]];</pre>
5 d8
        assert(unique(v.begin(), v.end()) == v.end());
        vector < vector < int >> g(n);
4aa
        bool col = true;
4ec
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1], v[i-2]))
   col = false;
        if (col) {
bf5
            for (int i = 1; i < n; i++)
aa4
                 g[idx[i-1]].push_back(idx[i]),
839
   g[idx[i]].push_back(idx[i-1]);
96b
            return g;
0ae
d36
        Q e = build_tr(v, 0, n-1).first;
        vector < Q > edg = {e};
113
5 d 1
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
            for (Q at = e; !at->used; at = at->next()) {
3ed
                 at->used = true;
60d
                 g[idx[at->id]].push_back(idx[at->rev()->id]);
cf8
15d
                 edg.push_back(at->rev());
9f2
            }
        }
d19
96b
        return g;
b43 }
```

## 7.38 Triangulos em Grafos

```
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
// O(m sqrt(m) log(n)) se chamar para todos os vertices
c0d vector<pair<int, int>> g[MAX]; // {para, peso}
d41 #warning o 'g' deve estar ordenado
9a5 vector < pair < int , pair < int , int >>> get_triangles (int i) {
771
        vector < pair < int , pair < int , int >>> tri;
        for (pair<int, int> j : g[i]) {
b23
2 b 3
            int a = i, b = j.first;
            if (g[a].size() > g[b].size()) swap(a, b);
6dd
            for (pair < int, int > c : g[a]) if (c.first != b and c.first
eb0
   > j.first) {
525
                auto it = lower_bound(g[b].begin(), g[b].end(),
   make_pair(c.first, -INF));
```

## 8 Primitivas

### 8.1 Aritmetica Modular

```
// O mod tem q ser primo
429 template <int p> struct mod_int {
c68
        ll expo(ll b, ll e) {
c85
             ll ret = 1;
c87
             while (e) {
cad
                 if (e % 2) ret = ret * b % p;
9d2
                 e /= 2, b = b * b % p;
c42
edf
             return ret;
734
        }
1f6
        11 inv(11 b) { return expo(b, p-2); }
4d7
        using m = mod_int;
d93
        int v:
fe0
        mod_int() : v(0) {}
e 12
        mod int(ll v ) {
019
             if (v_ >= p or v_ <= -p) v_ %= p;
bc6
             if (v_{-} < 0) v_{-} += p;
2e7
             v = v:
7f3
74 d
        m& operator +=(const m& a) {
2fd
             v += a.v;
             if (v >= p) v -= p;
ba5
357
             return *this;
c8b
eff
        m& operator -=(const m& a) {
8b4
             v -= a.v:
сс8
             if (v < 0) v += p;
357
             return *this;
f8d
        }
4c4
        m& operator *=(const m& a) {
```

```
8a5
            v = v * ll(a.v) % p;
357
            return *this;
d4c
        m& operator /=(const m& a) {
3f9
5 d 6
            v = v * inv(a.v) % p;
357
            return *this;
62d
d65
        m operator -(){ return m(-v); }
b3e
        m& operator ^=(11 e) {
06d
            if (e < 0) {
6 e 2
                v = inv(v);
00с
                e = -e:
275
284
            v = expo(v, e);
            // possivel otimizacao:
            // cuidado com 0^0
            // v = \exp(v, e''_{(p-1)});
357
            return *this;
6ed
        bool operator ==(const m& a) { return v == a.v; }
423
        bool operator !=(const m& a) { return v != a.v; }
69f
1 c 6
        friend istream& operator >>(istream& in, m& a) {
d1c
            11 val; in >> val;
d48
            a = m(val):
091
            return in:
870
44f
        friend ostream& operator <<(ostream& out, m a) {</pre>
5 a 0
            return out << a.v;</pre>
214
399
        friend m operator +(m a, m b) { return a += b; }
f9e
        friend m operator -(m a, m b) { return a -= b; }
9 c 1
        friend m operator *(m a, m b) { return a *= b; }
        friend m operator /(m a. m b) { return a /= b; }
51b
        friend m operator ^(m a, ll e) { return a ^= e; }
08f
b20 };
055 typedef mod_int<(int)1e9+7> mint;
8.2 Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
// Divisao, resto - O(n^2)
```

```
864 struct bint {
669
        static const int BASE = 1e9;
990
        vector < int > v;
3bd
        bool neg;
        bint() : neg(0) {}
609
d53
        bint(int val) : bint() { *this = val; }
        bint(long long val) : bint() { *this = val; }
e8f
        void trim() {
a0f
f 42
            while (v.size() and v.back() == 0) v.pop_back();
df8
            if (!v.size()) neg = 0;
8e3
        // converter de/para string | cin/cout
        bint(const char* s) : bint() { from_string(string(s)); }
294
548
        bint(const string& s) : bint() { from_string(s); }
        void from_string(const string& s) {
4ab
            v.clear(), neg = 0;
0a6
d72
            int ini = 0:
            while (ini < s.size() and (s[ini] == '-' or s[ini] == '+'
8e2
   or s[ini] == '0'))
71 d
                if (s[ini++] == '-') neg = 1;
883
            for (int i = s.size()-1; i >= ini; i -= 9) {
05е
                int at = 0:
5b1
                for (int j = max(ini, i - 8); j <= i; j++) at = 10*at
   + (s[i]-'0');
1fd
                v.push_back(at);
a5a
df8
            if (!v.size()) neg = 0;
e9a
2ff
        string to_string() const {
8be
            if (!v.size()) return "0";
793
            string ret;
            if (neg) ret += '-';
73 e
            for (int i = v.size()-1; i >= 0; i--) {
3e9
582
                string at = ::to_string(v[i]);
ced
                int add = 9 - at.size();
75 e
                if (i+1 < v.size()) for (int j = 0; j < add; j++) ret
   += '0':
f9f
                ret += at:
f 64
edf
            return ret;
770
d2f
        friend istream& operator>>(istream& in, bint& val) {
eb6
            string s; in >> s;
966
            val = s:
```

```
091
            return in;
328
99d
        friend ostream& operator << (ostream& out, const bint& val) {
8b9
            string s = val.to_string();
396
            out << s:
fe8
            return out;
ce1
        // operators
        friend bint abs(bint val) {
60a
c5f
            val.neg = 0;
d94
            return val:
44b
bee
        friend bint operator-(bint val) {
            if (val != 0) val.neg ^= 1;
815
d94
            return val;
326
        bint& operator=(const bint& val) { v = val.v, neg = val.neg;
41f
   return *this; }
        bint& operator=(long long val) {
249
            v.clear(), neg = 0;
0a6
            if (val < 0) neg = 1, val *= -1;</pre>
3a6
            for (; val; val /= BASE) v.push_back(val % BASE);
fdc
357
            return *this;
220
3bd
        int cmp(const bint& r) const { // menor: -1 | igual: 0 |
   maior: 1
b14
            if (neg != r.neg) return neg ? -1 : 1;
            if (v.size() != r.v.size()) {
0 bb
ff7
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
                return neg ? -ret : ret;
91b
1f6
            for (int i = int(v.size())-1; i >= 0; i--) {
478
                if (v[i] != r.v[i]) {
405
                    int ret = v[i] < r.v[i] ? -1 : 1;</pre>
2 e 5
91b
                    return neg ? -ret : ret;
9a9
                }
c32
bb3
            return 0;
07d
        friend bool operator < (const bint& l, const bint& r) { return
   1.cmp(r) == -1;}
        friend bool operator > (const bint& 1, const bint& r) { return
   1.cmp(r) == 1; }
        friend bool operator <= (const bint& 1, const bint& r) { return</pre>
   1.cmp(r) <= 0;}
        friend bool operator >= (const bint& 1, const bint& r) { return
954
```

```
1.cmp(r) >= 0; }
        friend bool operator == (const bint& 1, const bint& r) { return
   1.cmp(r) == 0: 
        friend bool operator!=(const bint& 1, const bint& r) { return
   1.cmp(r) != 0; }
38 e
        bint& operator +=(const bint& r) {
            if (!r.v.size()) return *this;
6bf
a 93
            if (neg != r.neg) return *this -= -r;
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
256
e 28
                if (i == v.size()) v.push_back(0);
08f
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
baa
                if ((c = v[i] >= BASE)) v[i] -= BASE;
8bb
357
            return *this;
ab1
54 c
        friend bint operator+(bint a, const bint& b) { return a += b; }
        bint& operator -=(const bint& r) {
9 c 8
6bf
            if (!r.v.size()) return *this;
524
            if (neg != r.neg) return *this += -r;
            if ((!neg and *this < r) or (neg and r < *this)) {
358
b10
                *this = r - *this;
a10
                neg ^= 1;
357
                return *this;
807
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
                v[i] -= c + (i < r.v.size() ? r.v[i] : 0);
9ef
                if ((c = v[i] < 0)) v[i] += BASE;
с8с
687
0eb
            trim();
357
            return *this;
f72
        friend bint operator-(bint a, const bint& b) { return a -= b; }
f44
        // operators de * / %
        bint& operator *=(int val) {
6b0
            if (val < 0) val *= -1, neg ^= 1;</pre>
bca
566
            for (int i = 0, c = 0; i < v.size() or c; i++) {
e 28
                if (i == v.size()) v.push_back(0);
352
                long long at = (long long) v[i] * val + c;
                v[i] = at % BASE;
6a3
b3d
                c = at / BASE;
cb1
0eb
            trim();
357
            return *this;
a 57
480
        friend bint operator *(bint a, int b) { return a *= b; }
```

```
d5c
        friend bint operator *(int a, bint b) { return b *= a; }
        using cplx = complex < double >;
13b
        void fft(vector<cplx>& a, bool f, int N, vector<int>& rev)
bfb
   const {
            for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
bc7
   a[rev[i]]);
bad
            vector < cplx > roots(N);
            for (int n = 2; n \le N; n *= 2) {
192
4 e 9
                 const static double PI = acos(-1);
                for (int i = 0; i < n/2; i++) {
71a
40d
                     double alpha = (2*PI*i)/n;
1a1
                     if (f) alpha = -alpha;
3f6
                     roots[i] = cplx(cos(alpha), sin(alpha));
f 16
                for (int pos = 0; pos \langle N; pos += n \rangle
3 e 9
                     for (int 1 = pos, r = pos+n/2, m = 0; m < n/2;
   1++, r++, m++) {
                         auto t = roots[m]*a[r];
297
                         a[r] = a[1] - t;
254
                         a[1] = a[1] + t:
b8f
                     }
b0d
e07
            if (!f) return:
3f1
08Ъ
            auto invN = cplx(1)/cplx(N);
873
            for (int i = 0; i < N; i++) a[i] *= invN;</pre>
c75
0 e 0
        vector < long long > convolution (const vector < int > & a, const
   vector < int > & b) const {
            vector < cplx > l(a.begin(), a.end()), r(b.begin(), b.end());
ff9
996
            int ln = l.size(), rn = r.size(), N = ln+rn+1, n = 1,
   log_n = 0;
            while (n <= N) n <<= 1, log_n++;
821
            vector<int> rev(n);
808
            for (int i = 0; i < n; i++) {
603
                rev[i] = 0:
434
f44
                for (int j = 0; j < log_n; j++) if (i >> j & 1)
4ff
                     rev[i] = 1 << (log_n-1-j);
256
230
            l.resize(n), r.resize(n);
a89
            fft(l, false, n, rev), fft(r, false, n, rev);
            for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
917
88b
            fft(1, true, n, rev);
            vector<long long> ret;
7ae
            for (auto& i : 1) ret.push_back(round(i.real()));
c14
edf
            return ret;
917
633
        vector<int> convert_base(const vector<int>& a, int from, int
```

```
to) const {
498
            static vector<long long> pot(10, 1);
671
            if (pot[1] == 1) for (int i = 1; i < 10; i++) pot[i] =
   10*pot[i-1];
4b8
            vector < int > ret;
            long long at = 0;
156
608
            int digits = 0;
941
            for (int i : a) {
412
                at += i * pot[digits];
                digits += from;
0.35
684
                while (digits >= to) {
0 c 8
                    ret.push_back(at % pot[to]);
cf9
                    at /= pot[to]:
fd4
                    digits -= to;
                }
122
87b
            }
944
            ret.push_back(at);
            while (ret.size() and ret.back() == 0) ret.pop_back();
384
edf
            return ret;
0.90
        }
edb
        bint operator*(const bint& r) const { // O(n log(n))
2af
            bint ret:
968
            ret.neg = neg ^ r.neg;
d5d
            auto conv = convolution(convert_base(v, 9, 4),
   convert_base(r.v, 9, 4));
a0e
            long long c = 0;
            for (auto i : conv) {
a74
f6d
                long long at = i+c;
4cb
                ret.v.push_back(at % 10000);
a25
                c = at / 10000;
773
            }
3cb
            for (; c; c /= 10000) ret.v.push_back(c%10000);
            ret.v = convert_base(ret.v, 4, 9);
0 e 2
            if (!ret.v.size()) ret.neg = 0;
25 c
            return ret:
edf
        }
c6b
359
        bint& operator*=(const bint& r) { return *this = *this * r; };
9a3
        bint& operator/=(int val) {
d9a
            if (val < 0) neg ^= 1, val *= -1;</pre>
            for (int i = int(v.size())-1, c = 0; i >= 0; i--) {
f 18
                long long at = v[i] + c * (long long) BASE;
2a7
e 02
                v[i] = at / val;
fb1
                c = at % val;
fdb
            }
0eb
            trim();
357
            return *this:
db6
        }
```

```
e74
        friend bint operator/(bint a, int b) { return a /= b; }
        int operator %=(int val) {
4a9
23b
            if (val < 0) val *= -1;
156
            long long at = 0;
            for (int i = int(v.size())-1; i >= 0; i--)
f31
                at = (BASE * at + v[i]) % val;
1 b 3
d22
            if (neg) at *=-1:
ce6
            return at;
4b4
2fb
        friend int operator % (bint a, int b) { return a % = b; }
        friend pair < bint, bint > divmod(const bint& a_, const bint& b_)
13b
   611
            if (a == 0) return {0, 0};
            int norm = BASE / (b_.v.back() + 1);
d8a
b4e
            bint a = abs(a<sub>_</sub>) * norm;
            bint b = abs(b_) * norm;
027
            bint q, r;
14d
            for (int i = a.v.size() - 1; i >= 0; i--) {
c.91
                r *= BASE, r += a.v[i];
b71
                long long upper = b.v.size() < r.v.size() ?</pre>
   r.v[b.v.size()] : 0:
                int lower = b.v.size() - 1 < r.v.size() ?</pre>
86d
   r.v[b.v.size() - 1] : 0;
                int d = (upper * BASE + lower) / b.v.back();
431
5d4
30f
                while (r < 0) r += b, d--; // roda 0(1) vezes
738
                q.v.push_back(d);
c6a
a48
            reverse(q.v.begin(), q.v.end());
            q.neg = a_.neg ^ b_.neg;
ae2
88b
            r.neg = a_.neg;
            q.trim(), r.trim();
8 e 5
            return {q, r / norm};
0ef
4fd
        bint operator/(const bint& val) { return divmod(*this,
   val).first: }
        bint& operator/=(const bint& val) { return *this = *this /
7 f 9
   val; }
        bint operator % (const bint & val) { return divmod(*this,
1f9
   val).second; }
        bint& operator%=(const bint& val) { return *this = *this %
   val; }
6c3 };
```

#### 8.3 Calendario

```
// Congruencia de Zeller
//
// Os dias da semana correspondem aos restos % 7
// Segunda=0, Terca=1, ..., Domingo=6
74e int get_id(int d, int m, int y) {
c5d
        if (m < 3) v--. m += 12;
        return 365 * y + y / 4 - y / 100 + y / 400 + (153 * (m - 3) +
   2) / 5 + d - 307;
ff5 }
ade tuple < int, int, int > date(int id) {
        int x = id + 1789995, n = 4 * x / 146097, i, j, d, m, y;
33 e
        x = (146097 * n + 3) / 4;
6fc
        i = (4000 * (x + 1)) / 1461001;
b1d
        x = 1461 * i / 4 - 31;
        j = 80 * x / 2447, d = x - 2447 * j / 80;
dd0
179
        x = i / 11:
e 85
        m = j + 2 - 12 * x, y = 100 * (n - 49) + i + x;
b86
        return {d, m, y};
0d9 }
8.4 Fracao
// Funciona com o Big Int
a4e template < typename T = int > struct frac {
a40
        T num, den;
e3f
        template < class U, class V>
61 d
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
bad
            assert(den != 0);
583
            if (den < 0) num *= -1, den *= -1;
a51
            T g = gcd(abs(num), den);
572
            num /= g, den /= g;
        }
fbf
        friend bool operator < (const frac& 1, const frac& r) {</pre>
51f
fa0
            return 1.num * r.den < r.num * 1.den;</pre>
a4e
4b5
        friend frac operator+(const frac& 1, const frac& r) {
b61
            return {1.num*r.den + 1.den*r.num, 1.den*r.den};
25 f
74 d
        friend frac operator-(const frac& 1, const frac& r) {
2 cd
            return {1.num*r.den - 1.den*r.num, 1.den*r.den};
8a7
        }
```

friend frac operator\*(const frac& 1, const frac& r) {

c80

```
510
            return {1.num*r.num, 1.den*r.den};
14b
a1b
        friend frac operator/(const frac& 1, const frac& r) {
            return {1.num*r.den, 1.den*r.num};
8f3
b2c
                                                                           8d7
        friend ostream& operator << (ostream& out, frac f) {</pre>
012
                                                                           4cb
            out << f.num << ',' << f.den:
37a
                                                                           858
fe8
            return out:
b49
       }
cdb }:
8.5 Geometria
c83 typedef double ld;
e3b const ld DINF = 1e18;
43a const ld pi = acos(-1.0);
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
        return abs(a - b) <= eps;</pre>
ba0
bfc }
                                                                           587
                                                                           6f8
b2a struct pt { // ponto
       ld x, y;
        pt(1d x_{-} = 0, 1d y_{-} = 0) : x(x_{-}), y(y_{-}) {}
3dd
        bool operator < (const pt p) const {</pre>
5bc
059
            if (!eq(x, p.x)) return x < p.x;
f98
            if (!eq(v, p.v)) return v < p.v;
bb3
            return 0;
f61
        bool operator == (const pt p) const {
a83
            return eq(x, p.x) and eq(y, p.y);
ed0
589
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
cb9
a24
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
        pt operator * (const ld c) const { return pt(x*c , y*c ); }
4 a 8
        pt operator / (const ld c) const { return pt(x/c , y/c ); }
a60
        ld operator * (const pt p) const { return x*p.x + y*p.y; }
3 b 6
```

ld operator ^ (const pt p) const { return x\*p.y - y\*p.x; }

friend istream& operator >> (istream& in, pt& p) {

return in >> p.x >> p.y;

6df

5ed

e37

e 45

a8b };

}

```
b3a struct line { // reta
730
        pt p, q;
0d6
        line() {}
        line(pt p_, pt q_) : p(p_), q(q_) {}
4b8
        friend istream& operator >> (istream& in, line& r) {
            return in >> r.p >> r.q;
        }
7ab }:
// PONTO & VETOR
364 ld dist(pt p, pt q) { // distancia
        return hypot(p.v - q.v. p.x - q.x);
c68 }
9d7 ld dist2(pt p, pt q) { // quadrado da distancia
        return sq(p.x - q.x) + sq(p.y - q.y);
80f }
483 ld norm(pt v) { // norma do vetor
490
        return dist(pt(0, 0), v);
cf7 }
589 ld angle(pt v) { // angulo do vetor com o eixo x
        1d \text{ ang} = atan2(v.v, v.x);
        if (ang < 0) ang += 2*pi;</pre>
19 c
        return ang;
404 }
298 ld sarea(pt p, pt q, pt r) { // area com sinal
606
        return ((q-p)^(r-q))/2;
1b1 }
e32 bool col(pt p. pt q. pt r) { // se p. q e r sao colin.
        return eq(sarea(p, q, r), 0);
e7d
98c }
Ocd bool ccw(pt p. pt q, pt r) \{ // \text{ se p, q, r sao ccw} \}
        return sarea(p, q, r) > eps;
85d }
1ef pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
е5с
        return pt(p.x * cos(th) - p.y * sin(th),
ff1
                p.x * sin(th) + p.y * cos(th));
41a }
ab1 pt rotate90(pt p) { // rotaciona 90 graus
```

```
a0d
        return pt(-p.y, p.x);
e4a }
// RETA
edc bool isvert(line r) { // se r eh vertical
87d
        return eq(r.p.x, r.q.x);
0fb }
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
        pt a = r.p - p, b = r.q - p;
b04
        return eq((a \hat{b}), 0) and (a * b) < eps;
726 }
98d ld get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
        return (r.p^r.q) / ((r.p-r.q)^v);
a0a }
256 pt proj(pt p, line r) { // projecao do ponto p na reta r
        if (r.p == r.q) return r.p;
bea
        r.q = r.q - r.p; p = p - r.p;
97a
9f8
        pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
        return proj + r.p;
232 }
d5c pt inter(line r, line s) { // r inter s
        if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
146
        r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
205
        return r.q * get_t(r.q, s) + r.p;
543
111 }
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
   de s
19b
        if (isinseg(r.p, s) or isinseg(r.q, s)
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
fcb ld disttoline(pt p, line r) { // distancia do ponto a reta
        return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
89a
1b7 }
bcc ld disttoseg(pt p, line r) { // distancia do ponto ao seg
73d
        if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
951
        if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
```

```
a19
        return disttoline(p, r);
367 }
11d ld distseg(line a, line b) { // distancia entre seg
        if (interseg(a, b)) return 0;
349
        ld ret = DINF;
        ret = min(ret, disttoseg(a.p, b));
341
        ret = min(ret, disttoseg(a.q, b));
ceb
        ret = min(ret, disttoseg(b.p. a));
093
448
        ret = min(ret, disttoseg(b.q, a));
edf
        return ret:
222 }
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
// ccw(r.p, r.q, p)
1a9 vector<pt> cut_polygon(vector<pt> v, line r) { // 0(n)
        vector<pt> ret;
8af
8a4
        for (int j = 0; j < v.size(); j++) {</pre>
            if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
dac
dce
            if (v.size() == 1) continue;
030
            line s(v[j], v[(j+1)%v.size()]);
ae3
            pt p = inter(r, s);
            if (isinseg(p, s)) ret.push_back(p);
a3d
d44
        }
8a1
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
        if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
24 d
edf
        return ret:
253 }
// distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
5f5 ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
080
        ld hor = 0, vert = 0;
34b
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
4 f d
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
80a
        return dist(pt(0, 0), pt(hor, vert));
96f
630 }
13d ld polarea(vector<pt> v) { // area do poligono
9 c 5
        ld ret = 0:
        for (int i = 0; i < v.size(); i++)</pre>
сбе
```

```
80f
            ret += sarea(pt(0, 0), v[i], v[(i + 1) % v.size()]);
d03
        return abs(ret);
5df }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector<pt>& v, pt p) { // O(n)
        int qt = 0;
8de
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            if (p == v[i]) return 2;
bda
6af
            int j = (i+1)%v.size();
e38
            if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
97f
                if ((v[i]-p)*(v[j]-p) < eps) return 2;</pre>
5 e 2
                continue:
            }
48b
388
            bool baixo = v[i].y+eps < p.y;</pre>
464
            if (baixo == (v[j].y+eps < p.y)) continue;</pre>
366
            auto t = (p-v[i])^{(v[j]-v[i])};
            if (eq(t, 0)) return 2;
1 b 4
            if (baixo == (t > eps)) qt += baixo ? 1 : -1;
839
d13
        }
b84
        return qt != 0;
a64 }
6ff bool interpol(vector<pt> v1, vector<pt> v2) { // se dois poligonos
   se intersectam - O(n*m)
        int n = v1.size(), m = v2.size();
7 d 1
        for (int i = 0; i < n; i++) if (inpol(v2, v1[i])) return 1;
c36
        for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return 1;</pre>
ab8
523
        for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
            if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[i],
   v2[(j+1)%m]))) return 1;
        return 0:
bb3
c58 }
494 ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
   poligonos
        if (interpol(v1, v2)) return 0;
f6b
349
        ld ret = DINF;
        for (int i = 0; i < v1.size(); i++) for (int j = 0; j <
   v2.size(); j++)
            ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
6 c 2
   v1.size()]),
                         line(v2[j], v2[(j + 1) % v2.size()])));
9 d 9
edf
        return ret:
```

```
125 }
138 vector<pt> convex_hull(vector<pt> v) { // convex hull - 0(n log(n))
        sort(v.begin(), v.end());
fca
d76
        v.erase(unique(v.begin(), v.end()), v.end());
        if (v.size() <= 1) return v;</pre>
52 d
526
        vector <pt> 1. u:
        for (int i = 0; i < v.size(); i++) {</pre>
f 14
fb2
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
   v[i]))
364
                l.pop_back();
c35
            l.push_back(v[i]);
58e
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
            while (u.size() > 1 \text{ and } ! ccw(u.end()[-2], u.end()[-1],
f 19
   v[i]))
7a8
                u.pop_back();
            u.push_back(v[i]);
a95
0b8
cfc
        1.pop_back(); u.pop_back();
        for (pt i : u) l.push_back(i);
82b
792
        return 1:
10d }
483 struct convex_pol {
f50
        vector < pt > pol;
        // nao pode ter ponto colinear no convex hull
        convex_pol() {}
d98
a04
        convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(log(n))
        bool is_inside(pt p) {
8af
            if (pol.size() == 0) return false;
b6e
            if (pol.size() == 1) return p == pol[0];
eae
            int 1 = 1, r = pol.size();
67 f
40 c
            while (1 < r) {
                int m = (1+r)/2;
ee4
48f
                if (ccw(p, pol[0], pol[m])) 1 = m+1;
ef3
                else r = m;
            }
91 c
00a
            if (l == 1) return isinseg(p, line(pol[0], pol[1]));
9e7
            if (1 == pol.size()) return false;
            return !ccw(p, pol[1], pol[1-1]);
1c0
6b0
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
```

```
719
        int extreme(const function < bool(pt, pt) > & cmp) {
             int n = pol.size();
b1c
4a2
             auto extr = [&](int i, bool& cur_dir) {
                 \operatorname{cur}_{\operatorname{dir}} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
22a
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
61a
364
            };
63d
             bool last dir. cur dir:
            if (extr(0, last_dir)) return 0;
a0d
993
            int 1 = 0, r = n;
            while (1+1 < r) {
ead
ee4
                 int m = (1+r)/2;
f29
                 if (extr(m, cur_dir)) return m;
44a
                 bool rel dir = cmp(pol[m], pol[l]);
                 if ((!last_dir and cur_dir) or
b18
                         (last_dir == cur_dir and rel_dir == cur_dir)) {
261
8a6
1 f 1
                     last_dir = cur_dir;
                 elser=m:
94a
606
792
            return 1:
56c
        }
316
        int max_dot(pt v) {
             return extreme([&](pt p, pt q) { return p*v > q*v; });
ec1
3b7
a54
        pair < int , int > tangents(pt p) {
ffb
             auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
             auto R = [\&](pt q, pt r) \{ return ccw(p, q, r); \};
8fd
            return {extreme(L), extreme(R)};
fa8
736
        }
3ec };
// CIRCUNFERENCIA
911 pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3 pontos
        b = (a + b) / 2;
174
        c = (a + c) / 2:
2ae
        return inter(line(b, b + rotate90(a - b)),
3f8
                line(c, c + rotate90(a - c)));
a12 }
4b3 vector<pt> circ_line_inter(pt a, pt b, pt c, ld r) { // intersecao
   da circunf (c, r) e reta ab
        vector <pt> ret;
8af
        b = b-a, a = a-c;
f2b
4 b 1
       1d A = b*b;
20a
        1d B = a*b:
2 e 9
        1d C = a*a - r*r;
```

```
1fa
        1d D = B*B - A*C:
818
        if (D < -eps) return ret;</pre>
dc5
        ret.push_back(c+a+b*(-B+sqrt(D+eps))/A);
        if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
20e
edf
        return ret:
cd8 }
ad2 vector<pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao da
    circunf (a, r) e (b, R)
        vector <pt> ret;
8af
b7e
        ld d = dist(a, b);
5 се
        if (d > r+R \text{ or } d+min(r, R) < max(r, R)) return ret;
398
        1d x = (d*d-R*R+r*r)/(2*d):
        1d y = sqrt(r*r-x*x);
183
        pt v = (b-a)/d;
325
        ret.push_back(a+v*x + rotate90(v)*y);
76 e
        if (y > 0) ret.push_back(a+v*x - rotate90(y)*y);
2cb
        return ret:
edf
fb1 }
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
f82
        if (!eq(angle(v1), angle(v2))) return angle(v1) < angle(v2);</pre>
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
27e }
b14 bool operator ==(const line& a. const line& b) {
        return !(a < b) and !(b < a);</pre>
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
        bool operator () (const line& a. const line& b) const {
            // assume que os segmentos tem p < q
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
            if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
231
   b.p.x))
780
                return ccw(a.p, a.q, b.p);
dc0
            return ccw(a.p, b.q, b.p);
243
367 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir;
5b0 struct cmp_sweepangle {
d80
        bool operator () (const line& a, const line& b) const {
```

```
522         return get_t(dir, a) + eps < get_t(dir, b);
653     }
97f };

8.6    Geometria - inteiro
2de #define sq(x) ((x)*(11)(x))
b2a struct pt { // ponto</pre>
```

```
int x, y;
e91
        pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
df1
5bc
        bool operator < (const pt p) const {</pre>
            if (x != p.x) return x < p.x;
95a
89c
            return y < p.y;</pre>
dcd
a83
        bool operator == (const pt p) const {
d74
            return x == p.x and y == p.y;
7 b 4
cb9
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
a24
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
        pt operator * (const int c) const { return pt(x*c, y*c); }
0 e f
        11 operator * (const pt p) const { return x*(11)p.x +
   y*(11)p.y; }
       11 operator ^ (const pt p) const { return x*(11)p.y -
   v*(11)p.x; }
        friend istream& operator >> (istream& in, pt& p) {
5 e d
            return in >> p.x >> p.y;
e37
        }
e45
840 }:
b3a struct line { // reta
730
       pt p, q;
0 d 6
       line() {}
4b8
     line(pt p_, pt q_) : p(p_), q(q_) {}
8 d 7
       friend istream& operator >> (istream& in, line& r) {
            return in >> r.p >> r.q;
4cb
858
       }
7ab }:
// PONTO & VETOR
ea8 11 dist2(pt p, pt q) { // quadrado da distancia
f24
        return sq(p.x - q.x) + sq(p.y - q.y);
515 }
5a2 11 sarea2(pt p, pt q, pt r) { // 2 * area com sinal}
```

```
586
        return (q-p)^(r-q);
bf4 }
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return sarea2(p, q, r) == 0;
034
a08 }
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
276
        return sarea2(p, q, r) > 0;
42b }
c31 int quad(pt p) { // quadrante de um ponto
        return (p.x<0)^3*(p.y<0):
fcf }
2df bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)
9fc
        if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
        return ccw(q, pt(0, 0), p);
ea1
771 }
ab1 pt rotate90(pt p) { // rotaciona 90 graus
a0d
        return pt(-p.v, p.x);
e4a }
// R.F.T.A
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
        pt a = r.p - p, b = r.q - p;
f65
        return (a ^ b) == 0 and (a * b) <= 0;
2ac
c9f }
676 \text{ bool} interseg(line r, line s) { // se o seg de r intersecta o seg
19b
        if (isinseg(r.p, s) or isinseg(r.q, s)
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
9e0 int segpoints(line r) { // numero de pontos inteiros no segmento
        return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
dd8 }
88a double get_t(pt v, line r) { // retorna t tal que t*v pertence a
        return (r.p^r.q) / (double) ((r.p-r.q)^v);
1ad
```

```
d27 }
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados paralelos
// assume que ta representado (inferior esquerdo, superior direito)
485 11 dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
c59
        int hor = 0, vert = 0;
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
34b
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
4fd
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
        return sq(hor) + sq(vert);
869
e13 }
9c3 ll polarea2(vector<pt> v) { // 2 * area do poligono
b73
        ll ret = 0:
        for (int i = 0; i < v.size(); i++)</pre>
с6е
            ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
532
        return abs(ret);
d03
d5f }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector<pt>& v, pt p) { // 0(n)
        int qt = 0;
8de
f14
        for (int i = 0; i < v.size(); i++) {
bda
            if (p == v[i]) return 2;
6af
            int j = (i+1)%v.size();
            if (p.y == v[i].y and p.y == v[j].y) {
cc6
547
                if ((v[i]-p)*(v[j]-p) <= 0) return 2;
5 e 2
                continue:
            }
b47
            bool baixo = v[i].y < p.y;</pre>
78c
            if (baixo == (v[j].y < p.y)) continue;</pre>
057
366
            auto t = (p-v[i])^(v[j]-v[i]);
2ad
            if (!t) return 2;
0bb
            if (baixo == (t > 0)) qt += baixo ? 1 : -1;
9cf
b84
        return qt != 0;
afd }
138 vector <pt > convex_hull(vector <pt > v) { // convex hull - 0(n log(n))
        sort(v.begin(), v.end());
fca
d76
        v.erase(unique(v.begin(), v.end()), v.end());
52d
        if (v.size() <= 1) return v;</pre>
```

```
526
        vector<pt> 1, u;
f 14
        for (int i = 0; i < v.size(); i++) {</pre>
fb2
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
   v[i]))
364
                l.pop_back();
            l.push_back(v[i]);
c35
        }
58e
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
f 19
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
   v[i]))
7a8
                 u.pop_back();
a95
            u.push_back(v[i]);
0b8
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
10d }
786 ll interior_points(vector<pt> v) { // pontos inteiros dentro de um
   poligono simples
        11 b = 0;
с4е
        for (int i = 0; i < v.size(); i++)</pre>
сбе
            b += segpoints(line(v[i], v[(i+1)\%v.size()])) - 1;
Осе
a1c
        return (polarea2(v) - b) / 2 + 1;
af2 }
483 struct convex_pol {
f50
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
        convex_pol() {}
d98
a04
        convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(log(n))
        bool is_inside(pt p) {
8af
            if (pol.size() == 0) return false;
b6e
            if (pol.size() == 1) return p == pol[0];
eae
67 f
            int l = 1, r = pol.size();
40 c
            while (1 < r) {
ee4
                 int m = (1+r)/2;
                 if (ccw(p, pol[0], pol[m])) l = m+1;
48 f
                 else r = m;
ef3
91 c
            if (l == 1) return isinseg(p, line(pol[0], pol[1]));
00a
9e7
            if (l == pol.size()) return false;
            return !ccw(p, pol[1], pol[1-1]);
1 c 0
6b0
        }
```

```
// ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
719
        int extreme(const function < bool(pt, pt) > & cmp) {
b1c
             int n = pol.size();
4 a 2
             auto extr = [&](int i, bool& cur dir) {
                 \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
22a
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
61a
364
            };
63d
             bool last_dir, cur_dir;
             if (extr(0, last dir)) return 0:
a0d
993
            int 1 = 0, r = n;
            while (1+1 < r) {
ead
ee4
                 int m = (1+r)/2;
                 if (extr(m, cur_dir)) return m;
f29
44a
                 bool rel_dir = cmp(pol[m], pol[l]);
b18
                 if ((!last_dir and cur_dir) or
261
                         (last_dir == cur_dir and rel_dir == cur_dir)) {
8a6
                     1 = m:
                     last_dir = cur_dir;
1 f 1
                elser=m:
94a
            }
606
792
            return 1;
56c
316
        int max_dot(pt v) {
ec1
             return extreme([&](pt p, pt q) { return p*v > q*v; });
3b7
a 54
        pair < int , int > tangents(pt p) {
ffb
             auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
             auto R = [\&](pt q, pt r) \{ return ccw(p, q, r); \};
8fd
fa8
            return {extreme(L), extreme(R)};
736
        }
3ec };
6e0 bool operator <(const line& a, const line& b) { // comparador pra
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
036
        bool b1 = compare_angle(v1, v2), b2 = compare_angle(v2, v1);
73c
        if (b1 or b2) return b1:
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
b61 }
b14 bool operator == (const line& a, const line& b) {
76c
        return !(a < b) and !(b < a);
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
```

```
08b
        bool operator () (const line& a, const line& b) const {
             // assume que os segmentos tem p < q
191
             if (a.p == b.p) return ccw(a.p, a.q, b.q);
             if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x < b.p.x))
614
780
                 return ccw(a.p, a.q, b.p);
             return ccw(a.p, b.q, b.p);
dc0
baf
        }
677 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir;
5b0 struct cmp_sweepangle {
        bool operator () (const line& a, const line& b) const {
261
             return get_t(dir, a) < get_t(dir, b);</pre>
dc5
f6d };
8.7 Geometria 3D
c83 typedef double ld;
e3b const ld DINF = 1e18:
107 const ld eps = 1e-9;
b32 #define sq(x)((x)*(x))
d97 \text{ bool} \text{ eq}(1d \text{ a. } 1d \text{ b})
             return abs(a - b) <= eps;</pre>
ba0
bfc }
```

 $pt(1d x_{-} = 0, 1d y_{-} = 0, 1d z_{-} = 0) : x(x_{-}), y(y_{-}), z(z_{-})$ 

return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);

pt operator + (const pt p) const { return pt(x+p.x, y+p.y,

pt operator - (const pt p) const { return pt(x-p.x, y-p.y,

if (!eq(x, p.x)) return x < p.x;

if (!eq(y, p.y)) return y < p.y;</pre>

if (!eq(z, p.z)) return z < p.z;

bool operator < (const pt p) const {</pre>

bool operator == (const pt p) const {

return 0;

2eb

a50

059

f98

44 c

bb3

6cd

a83

41 c

fb5

44b

392

z+p.z): }

z-p.z); }

{} 5bc

b2a struct pt { // ponto

ld x, y, z;

```
pt operator * (const ld c) const { return pt(x*c , y*c ,
fb7
   z*c ); }
            pt operator / (const 1d c) const { return pt(x/c , y/c ,
7a1
   z/c ); }
            1d operator * (const pt p) const { return x*p.x + y*p.y +
   z*p.z; }
7 f 6
            pt operator ^ (const pt p) const { return pt(y*p.z -
   z*p.v, z*p.x - x*p.z, x*p.v - v*p.x); }
            friend istream& operator >> (istream& in, pt& p) {
5ed
                   return in >> p.x >> p.y >> p.z;
9bf
5 e 8
            }
3ee }:
b3a struct line { // reta
730
            pt p, q;
0 d 6
            line() {}
            line(pt p_, pt q_) : p(p_), q(q_) {}
4 b 8
            friend istream& operator >> (istream& in, line& r) {
8 d 7
                    return in >> r.p >> r.q;
4cb
858
            }
7ab }:
79b struct plane { // plano
7 e 1
            array<pt, 3> p; // pontos que definem o plano
29b
            array<ld, 4> eq; // equacao do plano
bb7
            plane() {}
fb0
            plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
            friend istream& operator >> (istream& in, plane& P) {
ca9
2ab
                    return in >> P.p[0] >> P.p[1] >> P.p[2];
70e
                    P.build():
544
            void build() {
0 a 8
da2
                    pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
                    eq = {dir.x, dir.y, dir.z, dir*p[0]*(-1)};
7 d 5
            }
41a
d5d }:
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de rotacao ao
   redor de z
2fb pt convert(ld rho, ld th, ld phi) {
           return pt(sin(phi) * cos(th), sin(phi) * sin(th),
   cos(phi)) * rho;
a4f }
```

```
// projecao do ponto p na reta r
256 pt proj(pt p, line r) {
            if (r.p == r.q) return r.p;
bea
            r.q = r.q - r.p; p = p - r.p;
97 a
9f8
            pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
            return proj + r.p;
232 }
// projecao do ponto p no plano P
bla pt proj(pt p, plane P) {
7b6
            p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2]
b69
            pt norm = P.p[1] ^ P.p[2];
6ab
            pt proj = p - (norm * (norm * p) / (norm*norm));
467
            return proj + P.p[0];
4a0 }
// distancia
a45 ld dist(pt a, pt b) {
            return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
2d0 }
// distancia ponto reta
137 ld distline(pt p, line r) {
ce1
            return dist(p, proj(p, r));
3c4 }
// distancia de ponto para segmento
d43 ld distseg(pt p, line r) {
73 d
            if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
951
            if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
200
            return distline(p, r);
42c }
// distancia de ponto a plano com sinal
7cc ld sdist(pt p, plane P) {
            return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z + P.eq[3];
150
d49 }
// distancia de ponto a plano
768 ld distplane(pt p, plane P) {
            return abs(sdist(p, P));
33d }
// se ponto pertence a reta
099 bool isinseg(pt p, line r) {
a32
            return eq(distseg(p, r), 0);
```

```
31a }
// se ponto pertence ao triangulo definido por P.p
cd2 bool isinpol(pt p, vector<pt> v) {
fad
            assert(v.size() >= 3);
            pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
bf4
8a4
            bool inside = true;
            int sign = -1;
cec
f14
            for (int i = 0; i < v.size(); i++) {</pre>
                    line r(v[(i+1)%3], v[i]);
834
2a9
                    if (isinseg(p, r)) return true;
4ef
                    pt ar = v[(i+1)\%3] - v[i];
                    if (sign == -1) sign = ((ar^(p-v[i]))*norm > 0);
320
                     else if (((ar^(p-v[i]))*norm > 0) != sign) inside
82b
   = false:
15e
            return inside:
aca
c81 }
// distancia de ponto ate poligono
361 ld distpol(pt p, vector <pt> v) {
            pt p2 = proj(p, plane(v[0], v[1], v[2]));
3 e 7
61a
            if (isinpol(p2, v)) return dist(p, p2);
349
            ld ret = DINF:
f14
            for (int i = 0; i < v.size(); i++) {</pre>
                    int j = (i+1)%v.size();
6af
5ee
                    ret = min(ret, distseg(p, line(v[i], v[j])));
7b2
            }
edf
            return ret;
a8d }
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
e51 enum RETCODE {BOTH, ONE, PARAL, CONCOR};
26b pair < RETCODE, pt > intersect(plane P, line r) {
fac
        1d d1 = sdist(r.p, P);
        1d d2 = sdist(r.q, P);
        if (eq(d1, 0) \text{ and } eq(d2, 0))
53a
504
                    return pair(BOTH, r.p);
72c
        if (eq(d1, 0))
847
                    return pair(ONE, r.p);
485
        if (eq(d2, 0))
168
                     return pair(ONE, r.q);
```

```
3fb
         if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0)) {}
             if (eq(d1-d2, 0)) return pair(PARAL, pt());
463
406
             return pair(CONCOR, pt());
91 c
        }
        1d frac = d1 / (d1 - d2);
c84
         pt res = r.p + ((r.q - r.p) * frac);
3ff
         return pair(ONE, res);
394
b92 }
// rotaciona p ao redor do eixo u por um angulo a
787 pt rotate(pt p, pt u, ld a) {
773
             u = u / dist(u, pt());
             return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) *
e6f
    sin(a);
7f0 }
8.8 Matriz
945 #define MODULAR false
5ed template<typename T> struct matrix : vector<vector<T>> {
         int n. m:
30 f
         void print() {
603
             for (int i = 0; i < n; i++) {</pre>
                 for (int j = 0; j < m; j++) cout << (*this)[i][j] << "</pre>
70 f
1fb
                 cout << endl;</pre>
            }
d98
        }
101
         matrix(int n_, int m_, bool ident = false) :
aa3
                 vector < vector < T >> (n_, vector < T > (m_, 0)), n(n_), m(m_)  {
b14
94 e
             if (ident) {
                 assert(n == m);
df7
                 for (int i = 0; i < n; i++) (*this)[i][i] = 1;
a89
359
             }
527
        }
         matrix(const vector<vector<T>>& c) : vector<vector<T>>(c),
b83
             n(c.size()), m(c[0].size()) {}
a3d
         matrix(const initializer_list<initializer_list<T>>& c) {
efc
f7e
             vector < vector < T >> val:
212
             for (auto& i : c) val.push_back(i);
303
             *this = matrix(val);
c50
        }
388
         matrix <T> operator*(matrix <T>& r) {
```

```
1 e 2
            assert(m == r.n):
82 c
            matrix < T > M(n. r.m):
d69
            for (int i = 0; i < n; i++) for (int k = 0; k < m; k++)
                for (int j = 0; j < r.m; j++) {
df4
                    T \text{ add} = (*this)[i][k] * r[k][j];
e34
f98 #if MODULAR
d41 #warning Usar matrix<ll> e soh colocar valores em [0, MOD) na
   matriz!
8b6
                    M[i][i] += add%MOD;
                     if (M[i][j] >= MOD) M[i][j] -= MOD;
983
8c1 #else
                    M[i][i] += add:
f2e #endif
620
                }
474
            return M;
394
528
        matrix <T> operator^(ll e){
f10
            matrix<T> M(n, n, true), at = *this;
c87
            while (e) {
2 e 2
                if (e\&1) M = M*at;
cc2
                e >>= 1:
c80
                at = at*at;
            }
eb6
474
            return M;
ca3
582
        void apply_transform(matrix M, ll e){
1 c3
            auto& v = *this;
c87
            while (e) {
                if (e & 1) v = M * v;
9ba
                e >>= 1;
cc2
419
                M = M * M;
d86
            }
        }
4 e 5
70d };
    Matroid
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
//
```

// Intersecao sem peso O(r^2 n)

```
// em que n eh o tamanho do conjunto e r eh o tamanho da resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma floresta
// build() : O(n)
// oracle() : 0(1)
fda struct graphic_matroid {
5da
        int n, m, t;
32 c
        vector<array<int, 2>> edges;
789
        vector < vector < int >> g;
62 e
        vector<int> comp, in, out;
        graphic_matroid(int n_, vector<array<int, 2>> edges_)
513
            : n(n_), m(edges_.size()), edges(edges_), g(n), comp(n),
a1f
   in(n), out(n) {}
        void dfs(int u) {
315
            in[u] = t++;
ab8
            for (auto v : g[u]) if (in[v] == -1)
17 d
                comp[v] = comp[u], dfs(v);
863
677
            out[u] = t;
        }
d83
945
        void build(vector<int> I) {
a 34
741
            for (int u = 0; u < n; u++) g[u].clear(), in[u] = -1;
667
            for (int e : I) {
0.05
                auto [u, v] = edges[e];
125
                g[u].push_back(v), g[v].push_back(u);
a8a
            for (int u = 0; u < n; u++) if (in[u] == -1)
809
a7d
                comp[u] = u, dfs(u);
207
f31
        bool is ancestor(int u. int v) {
            return in[u] <= in[v] and in[v] < out[u];</pre>
a 68
0 c 2
        bool oracle(int e) {
e6b
453
            return comp[edges[e][0]] != comp[edges[e][1]];
687
f75
        bool oracle(int e, int f) {
574
            if (oracle(f)) return true;
622
            int u = edges[e][in[edges[e][0]] < in[edges[e][1]]];</pre>
ff2
            return is_ancestor(u, edges[f][0]) != is_ancestor(u,
    edges[f][1]);
8a9
       }
691 };
```

```
// Matroid de particao ou cores
// Um conjunto eh independente se a quantidade de elementos
// de cada cor nao excede a capacidade da cor
// Quando todas as capacidades sao 1, um conjunto eh independente
// se todas as suas cores sao distintas
// build() : O(n)
// oracle() : 0(1)
994 struct partition_matroid {
501
        vector < int > cap, color, d;
608
        partition_matroid(vector<int> cap_, vector<int> color_)
04d
            : cap(cap_), color(color_), d(cap.size()) {}
        void build(vector<int> I) {
945
def
            fill(d.begin(), d.end(), 0);
e9d
            for (int u : I) d[color[u]]++;
c58
        }
514
        bool oracle(int u) {
0 a 1
            return d[color[u]] < cap[color[u]];</pre>
703
f7f
        bool oracle(int u. int v) {
            return color[u] == color[v] or oracle(v);
2 f 7
4 b 4
caa };
// Intersecao de matroid sem pesos
// Dadas duas matroids M1 e M2 definidas sobre o mesmo
// conjunto I, retorna o maior subconjunto de I
// que eh independente tanto para M1 quanto para M2
//
// \Omega(r^2*n)
// Matroid "pesada" deve ser a M2
132 template < typename Matroid1, typename Matroid2>
801 vector < int > matroid_intersection(int n, Matroid1 M1, Matroid2 M2) {
f5b
        vector < bool > b(n);
        vector < int > I[2];
a 64
a8b
        bool converged = false;
0 c 1
        while (!converged) {
742
            I[0].clear(), I[1].clear();
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
99d
09d
            M1.build(I[1]), M2.build(I[1]);
289
            vector < bool > target(n), pushed(n);
26a
            queue < int > q;
            for (int u : I[0]) {
5 c 5
                target[u] = M2.oracle(u);
2 b 2
```

```
c1b
                if (M1.oracle(u)) pushed[u] = true, q.push(u);
            }
0e6
3fe
            vector < int > p(n, -1);
07a
            converged = true;
402
            while (q.size()) {
                int u = q.front(); q.pop();
be1
5 c 6
                if (target[u]) {
101
                     converged = false;
c32
                     for (int v = u; v != -1; v = p[v]) b[v] = !b[v];
c2b
a80
                }
e78
                for (int v : I[!b[u]]) if (!pushed[v]) {
34 d
                    if ((b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v, u)))
                        p[v] = u, pushed[v] = true, q.push(v);
bae
533
                }
1d9
            }
5 e 7
b68
        return I[1];
381 }
// Intersecao de matroid com pesos
// Dadas duas matroids M1 e M2 e uma funcao de pesos w, todas
   definidas sobre
// um conjunto I retorna o maior subconjunto de I (desempatado pelo
// que eh independente tanto para M1 quanto para M2
// A resposta eh construida incrementando o tamanho conjunto I de 1 em
// Se nao tiver custo negativo, nao precisa de SPFA
//
// O(r^3*n) com SPFA
// O(r^2*n*log(n)) com Dijkstra e potencial
42a template < typename T, typename Matroid1, typename Matroid2>
2b5 vector < int > weighted_matroid_intersection (int n, vector < T > w,
   Matroid1 M1, Matroid2 M2) {
6c9
        vector < bool > b(n), target(n), is_inside(n);
563
        vector < int > I[2], from(n);
e 35
        vector<pair<T, int>> d(n);
        auto check_edge = [&](int u, int v) {
169
            return (b[u] and M1.oracle(u, v)) or (b[v] and
249
   M2.oracle(v, u));
253
        }:
667
        while (true) {
742
            I[0].clear(), I[1].clear();
99d
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
```

```
// I[1] contem o conjunto de tamanho I[1].size() de menor
                peso
            M1.build(I[1]), M2.build(I[1]);
09d
            for (int u = 0; u < n; u++) {
687
                 target[u] = false, is_inside[u] = false, from[u] = -1;
ea5
                 d[u] = {numeric_limits <T>::max(), INF};
961
            }
392
            deque < T > q;
8 d 3
476
            sort(I[0].begin(), I[0].end(), [&](int i, int j){ return
   w[i] < w[j]; \});
            for (int u : I[0]) {
5 c 5
2b2
                 target[u] = M2.oracle(u);
5 a 7
                 if (M1.oracle(u)) {
4ef
                     if (is_inside[u]) continue;
7cc
                     d[u] = \{w[u], 0\};
                     if (!q.empty() and d[u] > d[q.front()])
427
   q.push_back(u);
655
                     else q.push_front(u);
                     is_inside[u] = true;
4ae
                }
764
add
            }
            while (q.size()) {
402
97a
                 int u = q.front(); q.pop_front();
6f3
                is_inside[u] = false;
57a
                for (int v : I[!b[u]]) if (check_edge(u, v)) {
9de
                     pair <T, int > nd(d[u].first + w[v], d[u].second +
   1);
61b
                     if (nd < d[v]) {</pre>
                         from[v] = u, d[v] = nd;
6ac
bd7
                         if (is_inside[v]) continue;
                         if (q.size() and d[v] > d[q.front()])
   q.push_back(v);
                         else q.push_front(v);
275
587
                         is inside[v] = true:
b3f
                    }
                }
a3b
563
            pair < T, int > mini = pair (numeric_limits < T >:: max(), INF);
cc8
489
            int targ = -1;
259
            for (int u : I[0]) if (target[u] and d[u] < mini)</pre>
2b9
                 mini = d[u], targ = u;
            if (targ != -1) for (int u = targ; u != -1; u = from[u])
e14
                b[u] = !b[u], w[u] *= -1;
d89
f97
            else break:
c7d
b68
        return I[1];
8e7 }
```

## 9 DP

## 9.1 Divide and Conquer DP

```
// Particiona o array em k subarrays
// minimizando o somatorio das queries
//
// O(k n log n), assumindo quer query(l, r) eh O(1)
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
de6
        if (1 > r) return:
109
        int m = (1+r)/2, p = -1;
        auto& ans = dp[m][k&1] = LINF;
d2b
6e2
        for (int i = max(m, lk); i <= rk; i++) {</pre>
            ll at = dp[i+1][\sim k\&1] + query(m, i);
07b
57 d
            if (at < ans) ans = at, p = i;
        }
8f5
1ee
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
d3e }
cf1 ll DC(int n, int k) {
321
        dp[n][0] = dp[n][1] = 0;
f 27
        for (int i = 0; i < n; i++) dp[i][0] = LINF;</pre>
b76
        for (int i = 1; i <= k; i++) solve(i, 0, n-i, 0, n-i);
8e7
        return dp[0][k&1];
5e9 }
```

## 9.2 Longest Common Subsequence

```
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
//
// O(n*m), O(n+m) de memoria
eaf int lcs_s[MAX], lcs_t[MAX];
a6d int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
d12 void dp_top(int li, int ri, int lj, int rj) {
d13
        memset (dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
753
        for (int i = li; i <= ri; i++) {</pre>
9aa
            for (int j = rj; j >= lj; j--)
83b
                dp[0][j-1j] = max(dp[0][j-1j],
```

```
741
                (lcs_s[i] == lcs_t[j]) + (j > 1j ? dp[0][j-1 - 1j] :
   0));
            for (int j = lj+1; j <= rj; j++)
04с
                dp[0][i - 1i] = max(dp[0][i - 1i], dp[0][i-1 -1i]);
939
09f
       }
58f }
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
caO void dp_bottom(int li, int ri, int lj, int rj) {
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
3a2
        for (int i = ri; i >= li; i--) {
49с
            for (int j = lj; j <= rj; j++)</pre>
dbb
                dp[1][i - li] = max(dp[1][i - li].
                (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 - lj] :
4da
   0));
            for (int j = rj-1; j >= 1j; j--)
6ca
769
                dp[1][i - li] = max(dp[1][i - li], dp[1][i+1 - li]);
19b
       }
e8a }
93c void solve(vector<int>& ans, int li, int ri, int lj, int rj) {
        if (li == ri){
2ad
            for (int j = lj; j <= rj; j++)
49c
f5b
                if (lcs_s[li] == lcs_t[j]){
a66
                    ans.push_back(lcs_t[j]);
c2b
                    break:
                }
840
505
            return:
126
        }
534
        if (li == ri){
753
            for (int i = li: i <= ri: i++){
88f
                if (lcs_s[i] == lcs_t[li]){
531
                    ans.push back(lcs s[i]):
c2b
                    break:
                }
68a
            }
a03
505
            return;
76d
a57
        int mi = (li+ri)/2;
ade
        dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
d7a
        int i_{-} = 0, mx = -1;
        for (int j = lj-1; j <= rj; j++) {
aee
da8
            int val = 0;
2 bb
            if (j >= lj) val += dp[0][j - lj];
b9e
            if (j < rj) val += dp[1][j+1 - lj];</pre>
```

```
ba8
            if (val >= mx) mx = val, j_ = j;
14e
6f1
        if (mx == -1) return;
c2a
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri, j_+1, rj);
dd5 }
058 vector<int> lcs(const vector<int>& s, const vector<int>& t) {
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];</pre>
953
577
        for (int i = 0; i < t.size(); i++) lcs t[i] = t[i];
dab
        vector < int > ans;
599
        solve(ans, 0, s.size()-1, 0, t.size()-1);
ba7
        return ans:
17c }
9.3 Mochila
// Resolve mochila, recuperando a resposta
// O(n * cap). O(n + cap) de memoria
add int v[MAX], w[MAX]; // valor e peso
582 int dp[2][MAX_CAP];
// DP usando os itens [1, r], com capacidade = cap
Od6 void get dp(int x, int 1, int r, int cap) {
        memset (dp[x], 0, (cap+1)*sizeof(dp[x][0]));
        for (int i = 1; i <= r; i++) for (int j = cap; j >= 0; j--)
574
            if (j - w[i] >= 0) dp[x][j] = max(dp[x][j], v[i] + dp[x][j]
3a9
   - w[i]]);
b95 }
5ab void solve(vector<int>& ans, int 1, int r, int cap) {
893
        if (1 == r) {
9ff
            if (w[1] <= cap) ans.push_back(1);</pre>
5.05
            return:
13a
        }
        int m = (1+r)/2;
ee4
283
        get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
056
        int left_cap = -1, opt = -INF;
c94
        for (int j = 0; j \le cap; j++)
2f2
            if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
91 d
                opt = at, left_cap = j;
da3
        solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap - left_cap);
d75 }
```

```
0d7 vector<int> knapsack(int n, int cap) {
dab
        vector < int > ans;
1 e 0
        solve (ans, 0, n-1, cap);
ba7
        return ans;
e4d }
9.4 SOS DP
// O(n 2^n)
// soma de sub-conjunto
e03 vector<11> sos_dp(vector<11> f) {
       int N = __builtin_ctz(f.size());
6 c 0
        assert((1 << \mathbb{N}) == f.size());
e59
        for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
5 a 5
   mask++)
796
            if (mask>>i&1) f[mask] += f[mask^(1<<i)];</pre>
        return f;
abe
bec }
// soma de super-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
        int N = __builtin_ctz(f.size());
e59
        assert((1 << N) == f.size());
       for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
5 a 5
   mask++)
a3c
            if (\sim \max > i \& 1) f [\max k] += f [\max k^{(1 < i)}]:
abe
        return f;
dbd }
9.5 Subset sum
// Retorna max(x <= t tal que existe subset de w que soma x)
//
// O(n * max(w))
// O(max(w)) de memoria
efd int subset_sum(vector<int> w, int t) {
bb5
       int pref = 0, k = 0;
417
        while (k < w.size()) and pref + w[k] <= t) pref += w[k++];
1 e 7
       if (k == w.size()) return pref;
```

int W = \*max\_element(w.begin(), w.end());

444

```
vector < int > last, dp(2*W, -1);
44 d
d7b
        dp[W - (t-pref)] = k;
54 d
        for (int i = k; i < w.size(); i++) {</pre>
288
            last = dp;
            for (int x = 0; x < W; x++) dp[x+w[i]] = max(dp[x+w[i]],
15 f
   last[x]);
17b
            for (int x = 2*W - 1; x > W; x--)
303
                 for (int j = max(0, last[x]); j < dp[x]; j++)
595
                     dp[x-w[j]] = max(dp[x-w[j]], j);
        }
867
2fb
        int ans = t;
        while (dp[W - (t-ans)] < 0) ans --;
1 c 1
ba7
        return ans:
d88 }
    strings
10
```

### 11 Estruturas

### 11.1 BIT

```
// BIT de soma O-based
// upper_bound(x) retorna o menor p tal que pref(p) > x
//
// Complexidades:
// build - O(n)
// update - 0(log(n))
// query - 0(log(n))
// upper_bound - O(log(n))
8eb struct Bit {
1a8
         int n:
         vector<1l> bit;
406
e 86
         Bit(int _{n=0}) : n(_{n}), bit(n + 1) {}
70 f
         Bit(vector \langle int \rangle \& v) : n(v.size()), bit(n + 1) {
78a
             for (int i = 1; i <= n; i++) {</pre>
671
                 bit[i] += v[i - 1]:
edf
                 int j = i + (i \& -i);
b8a
                 if (j <= n) bit[j] += bit[i];</pre>
806
             }
e89
        }
625
         void update(int i, ll x) { // soma x na posicao i
```

```
b64
            for (i++; i <= n; i += i & -i) bit[i] += x;
d67
462
        11 pref(int i) { // soma [0, i]
b73
            11 \text{ ret} = 0:
4d3
            for (i++; i; i -= i & -i) ret += bit[i];
edf
            return ret;
0ef
        }
02a
        11 query(int 1, int r) { // soma [1, r]
89b
            return pref(r) - pref(l - 1);
ca8
014
        int upper_bound(ll x) {
1ba
            int p = 0;
0af
            for (int i = _-lg(n); i+1; i--)
                 if (p + (1 << i) <= n \text{ and } bit[p + (1 << i)] <= x)
6f5
                     x -= bit[p += (1 << i)];
68e
74e
            return p;
fdd
        }
502 };
```

#### 11.2 BIT 2D

```
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar update)
// Complexidades:
// construir - O(n log(n))
// update e query - O(\log^2(n))
a6b template < class T = int > struct bit2d {
acf
        vector < T > X;
a84
        vector < vector < T >> Y, t;
709
        int ub(vector<T>& v, T x) {
            return upper_bound(v.begin(), v.end(), x) - v.begin();
dde
9сс
5 cb
        bit2d(vector<pair<T, T>> v) {
            for (auto [x, y] : v) X.push_back(x);
2 e 1
            sort(X.begin(), X.end());
fd4
1ee
            X.erase(unique(X.begin(), X.end()), X.end());
d56
            t.resize(X.size() + 1);
d12
            Y.resize(t.size()):
3 d 0
            sort(v.begin(), v.end(), [](auto a, auto b) {
e8f
                return a.second < b.second; });</pre>
```

```
961
            for (auto [x, y] : v) for (int i = ub(X, x); i < t.size();</pre>
   i += i\&-i
                 if (!Y[i].size() or Y[i].back() != y)
b75
   Y[i].push_back(y);
7 c 7
            for (int i = 0; i < t.size(); i++) t[i].resize(Y[i].size()</pre>
   + 1);
8cc
        void update(T x, T y, T v) {
e 78
2a9
            for (int i = ub(X, x); i < t.size(); i += i\&-i)
                 for (int j = ub(Y[i], y); j < t[i].size(); j += j&-j)</pre>
   t[i][j] += v;
533
        }
        T query(T x, T y) {
5 d 2
966
            T ans = 0;
            for (int i = ub(X, x); i; i -= i&-i)
c54
                 for (int j = ub(Y[i], y); j; j -= j&-j) ans += t[i][j];
4fb
ba7
            return ans:
        }
62 d
        T query(T x1, T y1, T x2, T y2) {
46 d
fcf
            return query (x2, y2) - query (x2, y1-1) - query (x1-1, y2)
   v2) + querv(x1-1, v1-1);
232
6a7 };
```

## 11.3 BIT com update em range

```
// Operacoes O-based
// query(1, r) retorna a soma de v[1..r]
// update(1, r, x) soma x em v[1..r]
//
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
e04 namespace bit {
        11 bit[2][MAX+2];
1a8
        int n:
61 c
        void build(int n2, int* v) {
1e3
            n = n2:
535
            for (int i = 1; i <= n; i++)
edd
                bit [1] [\min(n+1, i+(i\&-i))] += bit [1][i] += v[i-1];
```

```
db0
        11 get(int x, int i) {
637
b73
           11 \text{ ret = 0};
360
            for (; i; i -= i&-i) ret += bit[x][i];
edf
            return ret:
99c
20c
        void add(int x. int i. ll val) {
            for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
503
bf6
        11 get2(int p) {
162
с7с
            return get(0, p) * p + get(1, p);
153
02a
        11 querv(int 1, int r) {
ff5
            return get2(r+1) - get2(1);
633
089
        void update(int 1, int r, ll x) {
e5f
            add(0, 1+1, x), add(0, r+2, -x);
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
e5f
f91 };
```

#### 11.4 BIT-Sort Tree

```
// Tipo uma MergeSort Tree usando Bit
// Apesar da complexidade ser pior, fica melhor na pratica.
// query(1, r, k) retorna o numero de elementos menores que k
// no intervalo [1, r]
// Usa O(n log(n)) de memoria
//
// Complexidades:
// construir - O(n log^2(n))
// query - O(log^2(n))
6fa template < typename T> struct ms_bit {
1 a 8
        int n;
b2f
        vector < vector < T >> bit;
        ms_bit(vector < T > \& v) : n(v.size()), bit(n+1) {
899
830
            for (int i = 0; i < n; i++)
                for (int j = i+1; j \le n; j += j \& -j)
d51
dad
                     bit[j].push_back(v[i]);
535
            for (int i = 1; i <= n; i++)
                 sort(bit[i].begin(), bit[i].end());
eec
b4d
        }
```

```
int p_query(int i, T k) {
257
7c9
            int ret = 0;
            for (i++; i; i -= i&-i)
be8
                ret += lower_bound(bit[i].begin(), bit[i].end(), k) -
1bd
   bit[i].begin();
edf
            return ret:
6f9
690
        int query(int 1, int r, T k) {
            return p_query(r, k) - p_query(l-1, k);
83 d
bcc
       }
8d0 };
```

### 11.5 Convex Hull Trick Dinamico

```
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
// O(log(n)) amortizado por insercao
// O(log(n)) por query
72c struct Line {
073
         mutable ll a. b. p:
         bool operator<(const Line& o) const { return a < o.a; }</pre>
8e3
         bool operator<(ll x) const { return p < x; }</pre>
abf
469 }:
326 struct dynamic_hull : multiset <Line, less <>> {
33a
        11 div(ll a, ll b) {
             return a / b - ((a ^ b) < 0 and a % b);
a20
        }
a8a
         void update(iterator x) {
bbb
b2a
             if (next(x) == end()) x -> p = LINF;
             else if (x->a == next(x)->a) x->p = x->b >= next(x)->b?
772
   LINF : -LINF;
424
             else x \rightarrow p = div(next(x) \rightarrow b - x \rightarrow b, x \rightarrow a - next(x) \rightarrow a);
0 c 4
        }
71 c
         bool overlap(iterator x) {
f18
             update(x);
cfa
             if (next(x) == end()) return 0;
             if (x->a == next(x)->a) return x->b>= next(x)->b;
a4a
```

```
d40
            return x -> p >= next(x) -> p;
901
        void add(ll a, ll b) {
176
1 c7
            auto x = insert({a, b, 0});
4ab
            while (overlap(x)) erase(next(x)), update(x);
dbc
            if (x != begin() and !overlap(prev(x))) x = prev(x),
   update(x);
            while (x != begin() and overlap(prev(x)))
0fc
                x = prev(x), erase(next(x)), update(x);
4 d 2
48f
        }
4ad
        ll querv(ll x) {
229
            assert(!empty());
7 d 1
            auto l = *lower_bound(x);
d41 #warning cuidado com overflow!
            return 1.a * x + 1.b;
aba
        }
3f5
8f2 };
```

### 11.6 Convex Hull Trick Estatico

```
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
//
// add O(1) amortizado, get O(1) amortizado
4b5 struct CHT {
942
        int it:
        vector<ll> a, b;
ac1
45e
        CHT(): it(0){}
0bb
        ll eval(int i, ll x){
93d
            return a[i]*x + b[i];
        }
b2a
63a
        bool useless(){
a20
            int sz = a.size();
35f
            int r = sz-1, m = sz-2, l = sz-3;
d41 #warning cuidado com overflow!
d71
            return (b[1] - b[r])*(a[m] - a[1]) <
                (b[1] - b[m])*(a[r] - a[1]);
413
a0c
        }
bf4
        void add(ll A, ll B){
7 f 5
            a.push_back(A); b.push_back(B);
565
            while (!a.empty()){
                if ((a.size() < 3) || !useless()) break;</pre>
233
ecb
                a.erase(a.end() - 2);
```

```
568
                b.erase(b.end() - 2);
b21
            }
d27
            it = min(it, int(a.size()) - 1);
6df
        }
        ll get(ll x){
81b
46a
            while (it+1 < a.size()){
3 c 4
                if (eval(it+1, x) > eval(it, x)) it++;
f 97
                else break;
fe9
420
            return eval(it, x);
b44
        }
450 }:
11.7 DSU
// Une dois conjuntos e acha a qual conjunto um elemento pertence por
//
// find e unite: O(a(n)) \sim = O(1) amortizado
8d3 struct dsu {
        vector < int > id, sz;
b33
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
        int find(int a) { return a == id[a] ? a : id[a] = find(id[a]);
0cf
   }
440
        void unite(int a, int b) {
605
            a = find(a), b = find(b);
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);
6d0
            sz[a] += sz[b], id[b] = a;
ea7
        }
8e1 };
// DSU de bipartido
//
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
//
// find e unite: O(log(n))
8d3 struct dsu {
6f7
        vector<int> id, sz, bip, c;
```

```
5 b 4
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
            iota(id.begin(), id.end(), 0);
db8
f25
        int find(int a) { return a == id[a] ? a : find(id[a]); }
ef0
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
f30
   color(id[a]); }
440
        void unite(int a, int b) {
            bool change = color(a) == color(b);
263
605
            a = find(a), b = find(b);
a89
            if (a == b) {
4 e d
                if (change) bip[a] = 0;
505
                return;
           }
32d
956
            if (sz[a] < sz[b]) swap(a, b);
            if (change) c[b] = 1;
efe
2cd
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
22b
       }
118 };
// DSU Persistente
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
// find e unite: O(log(n))
8d3 struct dsu {
33c
        vector < int > id, sz, ti;
733
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) 
db8
            iota(id.begin(), id.end(), 0);
       }
aad
5e6
        int find(int a, int t) {
6ba
            if (id[a] == a or ti[a] > t) return a;
ea5
            return find(id[a], t);
       }
6cb
        void unite(int a, int b, int t) {
fa0
            a = find(a, t), b = find(b, t);
84f
            if (a == b) return;
d54
            if (sz[a] < sz[b]) swap(a, b);
956
35 d
            sz[a] += sz[b], id[b] = a, ti[b] = t;
```

```
513 }
6c6 };
// DSU com rollback
//
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
// Sempre que uma variavel muda de valor, adiciona na stack
//
// find e unite: O(log(n))
// checkpoint: O(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
8d3 struct dsu {
        vector < int > id, sz;
825
        stack<stack<pair<int&, int>>> st;
27 с
98 d
        dsu(int n) : id(n), sz(n, 1) 
            iota(id.begin(), id.end(), 0), st.emplace();
1cc
8 cd
bdf
        void save(int &x) { st.top().emplace(x, x); }
30 d
        void checkpoint() { st.emplace(); }
5cf
        void rollback() {
ba9
            while(st.top().size()) {
                auto [end, val] = st.top().top(); st.top().pop();
6bf
149
                end = val;
f9a
            }
25 a
            st.pop();
3c6
        }
        int find(int a) { return a == id[a] ? a : find(id[a]); }
ef0
440
        void unite(int a, int b) {
605
            a = find(a), b = find(b);
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);
803
            save(sz[a]), save(id[b]);
            sz[a] += sz[b], id[b] = a;
0.66
1b9
       }
c6e }:
```

### 11.8 Li-Chao Tree

```
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
// O(log(MA-MI)), O(n) de memoria
5b0 template < 11 MI = 11(-1e9), 11 MA = 11(1e9) > struct lichao {
b3a
        struct line {
12d
            ll a, b;
cef
            array < int, 2 > ch;
            line(ll a_{-} = 0, ll b_{-} = LINF):
fdf
                a(a_{-}), b(b_{-}), ch(\{-1, -1\}) \{\}
423
888
            11 operator ()(11 x) { return a*x + b; }
d1d
        }:
17b
        vector < line > ln;
df8
        int ch(int p, int d) {
e85
            if (ln[p].ch[d] == -1) {
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
ef2
            return ln[p].ch[d];
86a
021
        lichao() { ln.emplace_back(); }
c33
        void add(line s, ll l=MI, ll r=MA, int p=0) {
            11 m = (1+r)/2:
3 e 3
911
            bool L = s(1) < ln[p](1);
d37
            bool M = s(m) < ln[p](m);
03ъ
            bool R = s(r) < ln[p](r);
            if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
825
            if (s.b == LINF) return;
cac
f6d
            if (L != M) add(s, l, m-1, ch(p, 0));
898
            else if (R != M) add(s, m+1, r, ch(p, 1));
76e
        }
092
        11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
            11 m = (1+r)/2, ret = ln[p](x);
11b
9db
            if (ret == LINF) return ret;
529
            if (x < m) return min(ret, query(x, 1, m-1, ch(p, 0)));
81a
            return min(ret, query(x, m+1, r, ch(p, 1)));
fba
        }
59b };
```

## 11.9 Li-Chao Tree - Lazy

```
// Sendo N = MA-MT:
// insert(\{a, b\}) minimiza tudo com ax+b - O(log N)
// insert(\{a, b\}, 1, r) minimiza com ax+b no range [1, r] - 0(\log^2 N)
// shift({a, b}) soma ax+b em tudo - 0(1)
// shift({a, b}, 1, r) soma ax+b no range [1, r] - O(log^2 N)
// query(x) retorna o valor da posicao x - O(\log N)
//
// No inicio eh tudo LINF, se inserir {0, 0} fica tudo 0
// O(n log N) de memoria ; O(n) de memoria se nao usar as operacoes de
   range
41c template < int MI = int(-1e9), int MA = int(1e9) > struct lichao {
b3a
        struct line {
12 d
            ll a. b:
            11 la, lb; // lazy
158
cef
             array < int, 2 > ch;
fdf
             line(ll a_{-} = 0, ll b_{-} = LINF):
b09
                 a(a_{-}), b(b_{-}), la(0), lb(0), ch(\{-1, -1\})  {}
888
             11 operator ()(11 x) { return a*x + b; }
92 e
        };
17b
        vector<line> ln:
df8
        int ch(int p, int d) {
e 85
             if (\ln \lceil p \rceil, \cosh \lceil d \rceil == -1) {
9af
                 ln[p].ch[d] = ln.size();
cdc
                 ln.emplace_back();
bc1
ef2
             return ln[p].ch[d];
86a
        }
021
        lichao() { ln.emplace_back(); }
ceb
        void prop(int p, int l, int r) {
ff8
             if (\ln[p]. \ln == 0 \text{ and } \ln[p]. \ln == 0) return;
1d3
             ln[p].a += ln[p].la, ln[p].b += ln[p].lb;
579
             if (1 != r) {
b9e
                 int pl = ch(p, 0), pr = ch(p, 1);
0d7
                 ln[pl].la += ln[p].la, ln[pl].lb += ln[p].lb;
fa8
                 ln[pr].la += ln[p].la, ln[pr].lb += ln[p].lb;
77 f
01e
             ln[p].la = ln[p].lb = 0;
89b
        }
c06
        11 query(int x, int p=0, int l=MI, int r=MA) {
6b9
             prop(p, 1, r);
```

```
6f3
            ll ret = ln[p](x);
            if (ln[p].ch[0] == -1 and ln[p].ch[1] == -1) return ret;
33b
            int m = 1 + (r-1)/2;
90d
da9
            if (x \le m) return min(ret, query(x, ch(p, 0), 1, m));
c55
            return min(ret, query(x, ch(p, 1), m+1, r));
953
        void push(line s, int p, int l, int r) {
5 df
6b9
            prop(p, 1, r);
90d
            int m = 1 + (r-1)/2:
911
            bool L = s(1) < ln[p](1);
d37
            bool M = s(m) < ln[p](m);
03b
            bool R = s(r) < ln[p](r):
            if (M) swap(ln[p].a, s.a), swap(ln[p].b, s.b);
c3f
cac
            if (s.b == LINF) return;
            if (L != M) push(s, ch(p, 0), 1, m);
c49
29e
            else if (R != M) push(s, ch(p, 1), m+1, r);
ceb
        void insert(line s, int a=MI, int b=MA, int p=0, int l=MI, int
a8e
   r = MA) {
6b9
            prop(p, l, r);
2d3
            if (a \le 1 \text{ and } r \le b) \text{ return push}(s, p, l, r);
            if (b < l or r < a) return;
1 d d
90d
            int m = 1 + (r-1)/2;
f1e
            insert(s, a, b, ch(p, 0), 1, m);
952
            insert(s, a, b, ch(p, 1), m+1, r);
        }
375
        void shift(line s, int a=MI, int b=MA, int p=0, int l=MI, int
97a
   r = MA) {
6 b 9
            prop(p, l, r);
90d
            int m = 1 + (r-1)/2;
            if (a <= 1 and r <= b) {</pre>
9a3
                ln[p].la += s.a. ln[p].lb += s.b;
ada
                return;
505
            }
570
1 dd
            if (b < l or r < a) return;
fdd
            if (ln[p].b != LINF) {
751
                push(ln[p], ch(p, 0), 1, m);
ade
                push(ln[p], ch(p, 1), m+1, r);
                ln[p].a = 0, ln[p].b = LINF;
c2f
199
a04
            shift(s, a, b, ch(p, 0), 1, m);
            shift(s, a, b, ch(p, 1), m+1, r);
e7d
d43
        }
285 };
```

## 11.10 MergeSort Tree

```
// Se for construida sobre um arrav:
//
         count(i, j, a, b) retorna quantos
//
         elementos de v[i..j] pertencem a [a, b]
//
         report(i, i, a, b) retorna os indices dos
//
         elementos de v[i..j] que pertencem a [a, b]
//
         retorna o vetor ordenado
// Se for construida sobre pontos (x, v):
//
         count(x1, x2, y1, y2) retorna quantos pontos
//
         pertencem ao retangulo (x1, y1), (x2, y2)
//
        report(x1, x2, y1, y2) retorna os indices dos pontos que
//
         pertencem ao retangulo (x1, y1), (x2, y2)
//
        retorna os pontos ordenados lexicograficamente
//
        (assume x1 \le x2, y1 \le y2)
//
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
// x dentre os pontos que possuem y em [v1, v2] (0 based)
// Se quiser usar para achar k-esimo valor em range, construir
// com ms_tree t(v, true), e chamar kth(l, r, k)
//
// Usa O(n log(n)) de memoria
//
// Complexidades:
// construir - O(n log(n))
// count - O(log(n))
// report - O(log(n) + k) para k indices retornados
// kth - O(log(n))
c6c template <typename T = int> struct ms_tree {
6f7
         vector < tuple < T, T, int >> v;
1a8
         int n:
         vector < vector < tuple < T, T, int >>> t; // {v, idx, left}
5ee
6ae
         vector <T> vv;
78 c
         ms_tree(vector < pair < T, T >> \& vv) : n(vv.size()), t(4*n), vv(n) {
e80
             for (int i = 0; i < n; i++) v.push_back({vv[i].first,</pre>
    vv[i].second, i});
fca
             sort(v.begin(), v.end());
224
             build(1, 0, n-1);
01a
             for (int i = 0; i < n; i++) vy[i] = get<0>(t[1][i+1]);
45 e
         ms_tree(vector<T>& vv, bool inv = false) { // inv: inverte
    indice e valor
8e8
             vector < pair < T, T >> v2;
e1e
             for (int i = 0; i < vv.size(); i++)</pre>
196
                 inv ? v2.push_back({vv[i], i}) : v2.push_back({i,
```

```
vv[i]});
cca
             *this = ms_tree(v2);
f23
2 c 6
        void build(int p, int 1, int r) {
             t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0}); //
1 d 2
   {min_x, max_x, 0}
5 c8
             if (1 == r) return t[p].push_back({get<1>(v[1]),
   get <2 > (v[1]), 0});
             int m = (1+r)/2;
ee4
             build(2*p, 1, m), build(2*p+1, m+1, r);
bd9
32d
             int L = 0, R = 0;
a03
             while (t[p].size() \le r-l+1) {
                 int left = get <2>(t[p].back());
68e
                 if (L > m-1 \text{ or } (R+m+1 \le r \text{ and } t[2*p+1][1+R] \le r
   t[2*p][1+L])) {
                      t[p].push_back(t[2*p+1][1 + R++]);
8cf
                      get < 2 > (t[p].back()) = left;
da0
5 e 2
                      continue;
ce0
249
                 t[p].push_back(t[2*p][1 + L++]);
339
                 get < 2 > (t[p].back()) = left + 1;
             }
208
        }
2eb
        int get_l(T y) { return lower_bound(vy.begin(), vy.end(), y) -
   vv.begin(); }
ebb
        int get_r(T y) { return upper_bound(vy.begin(), vy.end(), y) -
   vy.begin(); }
f62
        int count(T x1, T x2, T y1, T y2) {
902
             function < int (int, int, int) > dfs = [&] (int p, int 1, int
   r) {
7 c 6
                 if (1 = r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or } get < 1 > (t[p][0])
   < x1) return 0:
2bb
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) <= x2)
   return r-1;
                 int nl = get < 2 > (t[p][1]), nr = get < 2 > (t[p][r]);
784
eb6
                 return dfs(2*p, nl, nr) + dfs(2*p+1, l-nl, r-nr);
122
7cb
             return dfs(1, get_l(y1), get_r(y2));
f65
002
        vector<int> report(T x1, T x2, T y1, T y2) {
4 b 8
             vector<int> ret;
             function < void (int, int, int) > dfs = [&] (int p, int 1, int
85e
   r) {
                 if (1 == r \text{ or } x2 < get<0>(t[p][0]) \text{ or } get<1>(t[p][0])
882
```

```
< x1) return;
                if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) <= x2) {
8da
                    for (int i = 1; i < r; i++)
e 00
   ret.push_back(get<1>(t[p][i+1]));
505
                    return:
067
784
                int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
                dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
194
12b
8ad
            dfs(1, get_l(v1), get_r(v2));
edf
            return ret;
668
        }
985
        int kth(T y1, T y2, int k) {
902
            function < int (int, int, int) > dfs = [&](int p, int 1, int
   r) {
150
                if (k >= r-1) {
941
                    k = r-1;
                     return -1;
daa
b8d
                if (r-l == 1) return get <1>(t[p][1+1]);
8da
                int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
784
072
                int left = dfs(2*p, nl, nr);
3b6
                if (left != -1) return left;
04 d
                return dfs(2*p+1, 1-n1, r-nr);
a1b
7cb
            return dfs(1, get_l(y1), get_r(y2));
635
        }
1ce }:
11.11 Min queue - deque
// Tudo O(1) amortizado
1dc template < class T> struct minqueue {
        deque<pair<T, int>> q;
2d8
        void push(T x) {
3fc
```

```
56e
             int ct = 1;
             while (q.size() and x < q.front().first)</pre>
953
75 f
                  ct += q.front().second, q.pop_front();
987
             q.emplace_front(x, ct);
        }
e 8 d
42 d
         void pop() {
aa2
             if (q.back().second > 1) q.back().second--;
c51
             else q.pop_back();
5 f d
        }
```

### 11.12 Min queue - stack

```
// Tudo O(1) amortizado
557 template < class T> struct minstack {
        stack <pair <T, T>> s;
3fc
        void push(T x) {
12b
            if (!s.size()) s.push({x, x});
9d9
            else s.emplace(x, std::min(s.top().second, x));
f8d
4f0
       T top() { return s.top().first; }
94a
       T pop() {
1f2
            T ans = s.top().first;
2eb
            s.pop();
ba7
            return ans:
013
614
       int size() { return s.size(); }
        T min() { return s.top().second; }
13b
4c0 };
1dc template < class T> struct minqueue {
        minstack <T> s1, s2;
        void push(T x) { s1.push(x); }
7cd
c96
        void move() {
d4d
            if (s2.size()) return;
d92
            while (s1.size()) {
7ae
                T x = s1.pop();
489
                s2.push(x);
            }
656
ef1
787
        T front() { return move(), s2.top(); }
        T pop() { return move(), s2.pop(); }
23a
7 f 3
        int size() { return s1.size()+s2.size(); }
19c
        T min() {
cd6
            if (!s1.size()) return s2.min();
            else if (!s2.size()) return s1.min();
58e
31d
            return std::min(s1.min(), s2.min());
9c7
       }
6d3 };
```

#### 11.13 Order Statistic Set

```
// Funciona do C++11 pra cima
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
def
3a1
        tree order statistics node update >:
// para declarar:
// ord_set<int> s;
// coisas do set normal funcionam:
// for (auto i : s) cout << i << endl:
// cout << s.size() << endl;
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
// cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
// cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int>> com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k
```

# 11.14 Priority Queue DS

```
// Mantem updates aplicados em uma estrutura de dados
// que permita rollback e nao seja amortizada.
// Cada update possui uma prioridade,
// sendo possivel remover o update com maior prioridade.
// Os updates devem ser comutativos, ou seja, o estado
// da estrutura deve ser o mesmo independente da ordem
// que eles sejam aplicados.
//
// Complexidades:
// update - O(log(n) + T(n))
// query - T(n)
// pop - O(log(n) * T(n)) amortizado
//
// onde T(n) eh a complexidade do update
```

```
// assumes all priorities are distinct
945 template < typename DS, typename UPD > struct priority_queue_ds {
df4
        vector<tuple<UPD, int, int>> upd; // {u, p, idx_in_pos}
a7e
866
        set < pair < int , int >> st;
927
        vector < int > pos;
        priority_queue_ds(int n) : D(n) {}
cf0
6af
        void update(UPD u, int p) {
9ab
            D.update(u);
d07
            st.emplace(p. pos.size());
6ca
            upd.emplace_back(u, p, pos.size());
e3d
            pos.push_back(upd.size() - 1);
6af
        int query(int a) {
427
            return D.find(a);
aa3
        }
2 d3
42d
        void pop() {
25f
            int k = 1, min_p; // k = number of pops we will do
43e
            vector<tuple<UPD, int, int>> small, big;
639
            auto it = st.end();
231
            for (int qt = 0; qt++ < (k+1)/2;) {
049
                it--:
3ab
                min_p = it->first;
                int i = pos[it->second];
80f
e82
                if (qt > 1) big.push_back(upd[i]);
                k = \max < int > (k, upd.size() - i);
84b
            }
b9a
b3d
            for (int i = 0; i < k; i++) {
                D.rollback():
a62
6 d 8
                auto [u, p, idx] = upd.rbegin()[i];
                if (p < min_p) small.emplace_back(u, p, idx);</pre>
86d
588
            }
            st.erase(prev(st.end()));
23e
            upd.erase(upd.end() - k, upd.end());
623
            small.insert(small.end(), big.rbegin(), big.rend());
a 25
            for (auto [u, p, idx] : small) {
06f
9ab
                D.update(u);
                upd.emplace_back(u, p, idx);
c8e
                pos[idx] = upd.size() - 1;
a7d
```

```
ec7
bd1
54a };
11.15 Range color
// update(l, r, c) colore o range [l, r] com a cor c,
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
//
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - O(log(q))
df6 template<typename T> struct color {
        set < tuple < int , int , T >> se;
071
        vector<tuple<int, int, T>> update(int 1, int r, T val) {
9 c.4
            auto it = se.upper_bound({r, INF, val});
753
            if (it != se.begin() and get<1>(*prev(it)) > r) {
e 91
                auto [L, R, V] = *--it;
3f0
                se.erase(it):
bfd
                 se.emplace(L, r, V), se.emplace(r+1, R, V);
683
d9e
            it = se.lower bound({1, -INF, val});
            if (it != se.begin() and get<1>(*prev(it)) >= 1) {
516
e91
                auto [L, R, V] = *--it;
                se.erase(it):
3f0
75 a
                 se.emplace(L, l-1, V), it = se.emplace(1, R, V).first;
b65
d7b
            vector<tuple<int, int, T>> ret;
            for (; it != se.end() and get<0>(*it) <= r; it =</pre>
7 a 1
   se.erase(it))
8 c 0
                ret.push_back(*it);
b4a
            se.emplace(1, r, val);
            return ret;
edf
        }
b6c
        T query(int i) {
ff9
            auto it = se.upper_bound({i, INF, T()});
c31
8e7
            if (it == se.begin() or get<1>(*--it) < i) return -1; //</pre>
   nao tem
53d
            return get <2 > (*it);
daf
9e9 };
```

#### 11.16 RMQ < O(n), O(1) > - min queue// O(n) pra buildar, query O(1) // Se tiver varios minimos, retorna // o de menor indice 1a5 template < typename T > struct rmq { 517 vector < T> v; int n; static const int b = 30; fcc 70e vector < int > mask, t: 183 int op(int x, int y) { return $v[x] \le v[y]$ ? x : y; } ee1 int msb(int x) { return \_\_builtin\_clz(1)-\_\_builtin\_clz(x); } c92 int small(int r, int sz = b) { return r-msb(mask[r]&((1<<sz)-1));} rmq() {} 6ad $rmq(const\ vector < T > \&\ v_) : v(v_), n(v.size()), mask(n), t(n)$ 43c for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) { 2 e 5 at = (at << 1) & ((1 << b) -1);a61 c00while (at and op(i-msb(at&-at), i) == i) at ^= at&-at; c2f for (int i = 0; i < n/b; i++) t[i] = small(b\*i+b-1); ea4 for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0; 39d i+(1<< j) <= n/b; i++)t[n/b\*j+i] = op(t[n/b\*(j-1)+i],ba5 t[n/b\*(j-1)+i+(1<<(j-1))]);41a e34 int index\_query(int 1, int r) { 27b if (r-l+1 <= b) return small(r, r-l+1);</pre> e80 int x = 1/b+1, y = r/b-1; fd3 if (x > y) return op(small(1+b-1), small(r)); int j = msb(y-x+1); a4e int ans = op(small(l+b-1), op(t[n/b\*j+x], t[n/b\*j+y-(1<<j)+1]));be6 return op(ans, small(r)); 62a T query(int 1, int r) { return v[index\_query(1, r)]; } 093 bab }: 11.17 SegTreap // Muda uma posicao do plano, e faz query de operacao // associativa e comutativa em retangulo // Mudar ZERO e op

// Esparso nas duas coordenadas, inicialmente eh tudo ZERO

```
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
//
// Valores no X tem que ser de O ateh NX
// Para q operacoes, usa O(q log(NX)) de memoria, e as
// operacoes custa O(log(q) log(NX))
55b const int ZERO = INF:
560 const int op(int 1, int r) { return min(1, r); }
878 mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T> struct treap {
3c9
        struct node {
b19
            node *1. *r:
ee1
             int p;
850
            pair <11, 11> idx; // {y, x}
36 d
            T val, mi;
bc2
             node(ll x, ll y, T val_) : l(NULL), r(NULL), p(rng()),
1b5
                 idx(pair(y, x)), val(val_), mi(val) {}
01e
            void update() {
d6e
                 mi = val:
182
                if (1) mi = op(mi, 1->mi);
b68
                 if (r) mi = op(mi, r \rightarrow mi);
282
            }
        };
6e1
bb7
        node* root;
84b
        treap() { root = NULL; }
        \simtreap() {
cec
            vector < node *> q = {root};
609
402
             while (q.size()) {
                 node* x = q.back(); q.pop_back();
e5d
ee9
                 if (!x) continue;
1c7
                 q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x;
            }
653
        }
50e
        treap(treap&& t) : treap() { swap(root, t.root); }
225
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
             if (!l or !r) return void(i = 1 ? 1 : r);
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80 e
```

```
fa0
             else join(1, r->1, r->1), i = r;
bda
             i->update();
671
        void split(node* i, node*& l, node*& r, pair<ll, 11> idx) {
c82
26a
             if (!i) return void(r = 1 = NULL);
13c
             if (i->idx < idx) split(i->r, i->r, r, idx), l = i;
d26
             else split(i\rightarrow 1, l, i\rightarrow 1, idx), r = i;
             i->update();
bda
143
d3b
        void update(ll x, ll y, T v) {
df9
             node *L, *M, *R;
8b2
             split(root, M, R, pair(y, x+1)), split(M, L, M, pair(y,
   x));
1 e 4
             if (M) M \rightarrow val = M \rightarrow mi = v;
9e5
             else M = new node(x, y, v);
             join(L, M, M), join(M, R, root);
69d
58e
        T query(ll ly, ll ry) {
91b
df9
             node *L, *M, *R;
             split(root, M, R, pair(ry, LINF)), split(M, L, M, pair(ly,
1 c 0
   0));
0f7
             T \text{ ret} = M ? M \rightarrow mi : ZERO;
             join(L, M, M), join(M, R, root);
69d
edf
             return ret;
1ae
        }
bdf }:
46a template < typename T > struct segtreap {
        vector < treap < T >> seg;
6e7
        vector < int > ch[2];
        11 NX:
e4e
        segtreap(11 NX_{-}) : seg(1), NX(NX_{-}) { ch[0].push_back(-1),}
   ch[1]. push back (-1): }
a71
        int get_ch(int i, int d){
e51
             if (ch[d][i] == -1) {
2d6
                 ch[d][i] = seg.size();
23e
                 seg.emplace_back();
842
                 ch[0].push_back(-1), ch[1].push_back(-1);
3e1
968
             return ch[d][i];
bb6
        }
        T query(11 lx, 11 rx, 11 ly, 11 ry, int p, 11 l, 11 r) {
10c
             if (rx < 1 or r < 1x) return ZERO;</pre>
003
fOf
             if (lx <= l and r <= rx) return seg[p].query(ly, ry);</pre>
```

```
11 m = 1 + (r-1)/2;
e6a
354
            return op(query(lx, rx, ly, ry, get_ch(p, 0), l, m),
060
                    query(lx, rx, ly, ry, get_ch(p, 1), m+1, r));
a5e
f 48
        T query(11 lx, 11 rx, 11 ly, 11 ry) { return query(1x, rx, ly,
   ry, 0, 0, NX); }
249
        void update(ll x, ll y, T val, int p, ll l, ll r) {
            if (1 == r) return seg[p].update(x, y, val);
73 c
e6a
            11 m = 1 + (r-1)/2;
            if (x <= m) update(x, y, val, get_ch(p, 0), 1, m);</pre>
сс5
5a2
            else update(x, y, val, get_ch(p, 1), m+1, r);
980
            seg[p].update(x, y, val);
cc2
517
        void update(11 x, 11 y, T val) { update(x, y, val, 0, 0, NX); }
40a };
11.18 SegTree
```

```
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Pode usar a seguinte funcao para indexar os nohs:
// f(1, r) = (1+r) | (1!=r), usando 2N de memoria
// Complexidades:
// build - 0(n)
// query - 0(log(n))
// update - 0(log(n))
aa4 namespace seg {
005
        11 seg[4*MAX], lazy[4*MAX];
052
        int n, *v;
d22
        11 build(int p=1, int l=0, int r=n-1) {
3c7
            lazv[p] = 0;
            if (1 == r) return seg[p] = v[1];
6cd
            int m = (1+r)/2;
ee4
            return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1, r);
193
c71
        }
850
        void build(int n2. int* v2) {
680
            n = n2, v = v2;
6f2
            build():
acb
        }
ceb
        void prop(int p, int l, int r) {
```

```
cdf
            seg[p] += lazv[p]*(r-l+1);
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
2 c 9
3 c 7
            lazv[p] = 0;
c10
        }
2 c 3
        11 query(int a, int b, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
            if (a <= 1 and r <= b) return seg[p]:
527
            if (b < 1 or r < a) return 0;</pre>
786
ee4
            int m = (1+r)/2;
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
b1f
4 c 5
        11 update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
cfb
6 b 9
            prop(p, 1, r):
            if (a \le 1 \text{ and } r \le b)
9a3
b94
                lazy[p] += x;
6b9
                prop(p, 1, r);
534
                return seg[p];
821
e9f
            if (b < l or r < a) return seg[p];
            int m = (1+r)/2:
ee4
            return seg[p] = update(a, b, x, 2*p, 1, m) +
fdb
7fd
                update(a, b, x, 2*p+1, m+1, r);
75 c
       }
0af };
// Se tiver uma seg de max, da pra descobrir em O(log(n))
// o primeiro e ultimo elemento >= val numa range:
// primeira posicao >= val em [a, b] (ou -1 se nao tem)
119 int get_left(int a, int b, int val, int p=1, int l=0, int r=n-1) {
6b9
        prop(p, 1, r);
        if (b < l or r < a or seg[p] < val) return -1;</pre>
f38
       if (r == 1) return 1:
205
       int m = (1+r)/2;
ee4
753
       int x = get_left(a, b, val, 2*p, 1, m);
50e
       if (x != -1) return x;
сЗс
        return get_left(a, b, val, 2*p+1, m+1, r);
68c }
// ultima posicao >= val em [a, b] (ou -1 se nao tem)
992 int get_right(int a, int b, int val, int p=1, int l=0, int r=n-1) {
6b9
        prop(p, 1, r);
       if (b < l or r < a or seg[p] < val) return -1;</pre>
f38
205
       if (r == 1) return 1;
       int m = (1+r)/2;
ee4
       int x = get_right(a, b, val, 2*p+1, m+1, r);
1 b 1
```

```
50e
        if (x != -1) return x;
6a7
        return get_right(a, b, val, 2*p, 1, m);
1b7 }
// Se tiver uma seg de soma sobre um arrav nao negativo v. da pra
// descobrir em O(\log(n)) o maior j tal que v[i]+v[i+1]+...+v[j-1] <
   val
6a9 int lower_bound(int i, 11 % val, int p, int l, int r) {
6b9
        prop(p, 1, r);
        if (r < i) return n;
6e8
b5d
        if (i \le l \text{ and } seg[p] \le val) {
bff
            val -= seg[p];
041
            return n:
634
       if (1 == r) return 1;
Зсе
        int m = (1+r)/2;
ee4
514
        int x = lower_bound(i, val, 2*p, 1, m);
        if (x != n) return x:
ee0
8b9
        return lower_bound(i, val, 2*p+1, m+1, r);
2b8 }
```

### 11.19 SegTree 2D Iterativa

```
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em seg[x+n][y+n]
// Query: soma do retangulo ((x1, v1), (x2, v2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
// Para query com distancia de manhattan <= d, faca
// nx = x+v, nv = x-v
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
//
// Complexidades:
// build - 0(n^2)
// query - 0(log^2(n))
// update - 0(log^2(n))
731 int seg[2*MAX][2*MAX], n;
0a8 void build() {
919
       for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
            if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
c81
```

```
fe9
             if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
                                                                                 526
                                                                                              int tam;
d51
        }
                                                                                 125
                                                                                              ll sum, lazy; // lazy pra soma
499 }
                                                                                 4f3
                                                                                              ll mi1, mi2, mi; // mi = #mi1
                                                                                              ll ma1, ma2, ma; // ma = #ma1
                                                                                 c61
251 int querv(int x1, int v1, int x2, int v2) {
         int ret = 0, y3 = y1 + n, y4 = y2 + n;
                                                                                 426
                                                                                              node(11 x = 0) {
827
83e
        for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
                                                                                 ba6
                                                                                                  sum = mi1 = ma1 = x:
0f2
             for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
                                                                                 b29
                                                                                                  mi2 = LINF, ma2 = -LINF;
554
                 if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
                                                                                 62 c
                                                                                                  mi = ma = tam = 1;
6 b 0
                 if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
                                                                                 c.60
                                                                                                  lazv = 0:
c01
                 if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
                                                                                 b00
5 d4
                 if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
                                                                                 770
                                                                                              node(const node& 1. const node& r) {
2 d 0
             }
                                                                                 a95
                                                                                                  sum = 1.sum + r.sum, tam = 1.tam + r.tam;
                                                                                 c60
                                                                                                  lazv = 0;
edf
         return ret;
                                                                                 797
                                                                                                  if (l.mi1 > r.mi1) {
ff1 }
                                                                                                       mi1 = r.mi1, mi = r.mi;
                                                                                 230
                                                                                                       mi2 = min(1.mi1, r.mi2);
                                                                                 ea2
767 void update(int x, int y, int val) {
                                                                                 f1e
                                                                                                  } else if (l.mi1 < r.mi1) {</pre>
                                                                                                       mi1 = 1.mi1, mi = 1.mi;
66a
        int v2 = v += n;
                                                                                 e 34
        for (x += n; x; x /= 2, y = y2) {
192
                                                                                 4b3
                                                                                                       mi2 = min(r.mi1, l.mi2);
970
             if (x >= n) seg[x][y] = val;
                                                                                 ef2
                                                                                                  } else {
             else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
                                                                                 a39
                                                                                                       mi1 = 1.mi1, mi = 1.mi+r.mi;
ba9
                                                                                 83 d
                                                                                                       mi2 = min(1.mi2, r.mi2);
3b1
             while (y /= 2) \text{ seg}[x][y] = \text{seg}[x][2*y] + \text{seg}[x][2*y+1];
                                                                                 a92
                                                                                                  }
68b
        }
                                                                                 cd0
                                                                                                  if (1.ma1 < r.ma1) {
62e }
                                                                                 6a0
                                                                                                       ma1 = r.ma1, ma = r.ma;
                                                                                 96 d
                                                                                                       ma2 = max(1.ma1, r.ma2);
                                                                                 3c0
                                                                                                  } else if (1.ma1 > r.ma1) {
                                                                                                       ma1 = 1.ma1, ma = 1.ma;
                                                                                 ae0
11.20 SegTree Beats
                                                                                 2ca
                                                                                                       ma2 = max(r.ma1, 1.ma2);
                                                                                 da8
                                                                                                  } else {
// \text{ query}(a, b) - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
                                                                                 db2
                                                                                                       ma1 = 1.ma1, ma = 1.ma+r.ma;
// updatemin(a, b, x) faz com que v[i] \leftarrow min(v[i], x),
                                                                                 c 05
                                                                                                       ma2 = max(1.ma2. r.ma2);
// para i em [a, b]
                                                                                 11 c
                                                                                                  }
// updatemax faz o mesmo com max, e updatesum soma x
                                                                                 1ba
// em todo mundo do intervalo [a, b]
                                                                                 4b4
                                                                                              void setmin(ll x) {
                                                                                 55e
                                                                                                  if (x \ge ma1) return:
// Complexidades:
                                                                                 463
                                                                                                  sum += (x - ma1)*ma;
// build - O(n)
                                                                                 be5
                                                                                                  if (mi1 == ma1) mi1 = x;
// query - 0(log(n))
                                                                                 0a0
                                                                                                  if (mi2 == ma1) mi2 = x;
// update - O(log^2 (n)) amortizado
                                                                                 b81
                                                                                                  ma1 = x:
// (se nao usar updatesum, fica log(n) amortizado)
                                                                                 0 c 3
                                                                                 6cb
                                                                                              void setmax(ll x) {
7c6 #define f first
                                                                                 e 25
                                                                                                  if (x <= mi1) return;</pre>
Oab #define s second
                                                                                 7 e 8
                                                                                                  sum += (x - mi1)*mi;
                                                                                 Obb
                                                                                                  if (ma1 == mi1) ma1 = x;
f39 namespace beats {
                                                                                 c32
                                                                                                  if (ma2 == mi1) ma2 = x;
```

3 c 9

struct node {

```
1 f f
                 mi1 = x;
            }
a86
4cf
            void setsum(ll x) {
fe8
                 mi1 += x, mi2 += x, ma1 += x, ma2 += x;
620
                 sum += x*tam:
c46
                 lazv += x;
            }
b53
47f
        };
62b
        node seg[4*MAX];
052
        int n, *v;
93b
        node build(int p=1, int l=0, int r=n-1) {
d84
            if (1 == r) return seg[p] = {v[1]};
ee4
            int m = (1+r)/2;
3 d 6
            return seg[p] = {build(2*p, 1, m), build(2*p+1, m+1, r)};
444
0 d8
        void build(int n2. int* v2) {
            n = n2, v = v2;
680
6f2
            build():
acb
        }
        void prop(int p, int 1, int r) {
ceb
            if (1 == r) return;
8 се
            for (int k = 0; k < 2; k++) {
abd
d07
                 if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
843
                 seg[2*p+k].setmin(seg[p].ma1);
                 seg[2*p+k].setmax(seg[p].mi1);
f79
585
431
             seg[p].lazy = 0;
7ee
        pair <pair <11, 11>, 11> query (int a, int b, int p=1, int 1=0,
   int r=n-1) {
            if (b < 1 or r < a) return {{LINF, -LINF}, 0};</pre>
e07
            if (a <= 1 and r <= b) return {{seg[p].mi1, seg[p].ma1},
9be
   seg[p].sum};
6b9
            prop(p, 1, r);
            int m = (1+r)/2;
ee4
            auto L = query(a, b, 2*p, 1, m), R = query(a, b, 2*p+1,
   m+1. r):
96d
            return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)}, L.s+R.s};
e9d
2 c8
        node updatemin(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
            if (b < l or r < a or seg[p].ma1 <= x) return seg[p];</pre>
744
309
            if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
                 seg[p].setmin(x);
ccd
                 return seg[p];
534
```

```
bbf
6b9
             prop(p, 1, r);
             int m = (1+r)/2;
ee4
96a
             return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
                              updatemin(a, b, x, 2*p+1, m+1, r)};
4db
aad
044
        node updatemax(int a. int b. ll x. int p=1. int l=0. int
   r=n-1) {
b59
             if (b < 1 or r < a or seg[p].mi1 >= x) return seg[p];
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
a9e
e8a
                 seg[p].setmax(x);
534
                 return seg[p];
e9b
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2;
ee3
             return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
98b
                              updatemax(a, b, x, 2*p+1, m+1, r)};
323
aee
        node updatesum(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
e9f
             if (b < l or r < a) return seg[p];</pre>
9a3
             if (a <= 1 and r <= b) {</pre>
8f4
                 seg[p].setsum(x);
534
                 return seg[p];
596
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2:
7b6
             return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
                              updatesum(a, b, x, 2*p+1, m+1, r)};
483
111
        }
0d2 }:
```

## 11.21 SegTree Colorida

```
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
//
// Complexidades:
// construir - O(n log(n)) espaco e tempo
```

```
// query - 0(log(n))
// update - 0(log(n))
// paint - O(log(n)) amortizado
04f struct seg_color {
        struct node {
3 c 9
b19
            node *1. *r:
0f9
            int cnt;
9ca
            ll val, lazy;
            node(): 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
277
01e
            void update() {
d0a
                 cnt = 0, val = 0;
bc4
                for (auto i : {1, r}) if (i) {
c89
                     i->prop();
281
                     cnt += i->cnt, val += i->val;
                }
68d
554
            }
a9c
            void prop() {
2 dd
                if (!lazy) return;
3 f 7
                 val += lazy*(ll)cnt;
b64
                for (auto i : {1, r}) if (i) i->lazy += lazy;
                lazv = 0;
c60
            }
e 24
514
        };
1a8
        int n:
9b0
        vector < node *> seg;
6e0
        seg_color(vector<pair<int, int>>& v, int c) : n(v.size()),
   seg(c, NULL) {
            for (int i = 0; i < n; i++)</pre>
830
9b7
                 seg[v[i].second] = insert(seg[v[i].second], i,
   v[i].first, 0, n-1);
94a
        }
3 c 7
        \simseg_color() {
dde
            queue < node *> q;
3a6
            for (auto i : seg) q.push(i);
402
            while (q.size()) {
20b
                 auto i = q.front(); q.pop();
dab
                if (!i) continue;
7 c.7
                q.push(i->1), q.push(i->r);
5 се
                 delete i;
            }
c60
        }
139
40b
        node* insert(node* at, int idx, int val, int l, int r) {
            if (!at) at = new node();
1 a 4
```

```
232
            if (1 == r) return at->cnt = 1, at->val = val, at;
ee4
            int m = (1+r)/2;
137
            if (idx \le m) at->1 = insert(at->1, idx, val, 1, m);
            else at->r = insert(at->r, idx, val, m+1, r);
3e6
cff
            return at -> update(), at;
d6e
        11 query(node* at, int a, int b, int l, int r) {
870
61b
            if (!at or b < l or r < a) return 0;
d9f
            at->prop();
            if (a <= l and r <= b) return at->val;
cb2
ee4
            int m = (1+r)/2;
4c4
            return query (at \rightarrow 1, a, b, 1, m) + query (at \rightarrow r, a, b, m+1,
   r);
8 c 3
        11 query(int c, int a, int b) { return query(seg[c], a, b, 0,
e54
   n-1): }
91 c
        void update(node* at, int a, int b, int x, int l, int r) {
fba
            if (!at or b < l or r < a) return;
d9f
            at->prop();
            if (a <= 1 and r <= b) {</pre>
9a3
e9a
                 at -> lazy += x;
cb2
                 return void(at->prop());
051
            }
ee4
            int m = (1+r)/2;
0b0
            update(at->1, a, b, x, 1, m), update(at->r, a, b, x, m+1,
   r);
7b4
            at->update();
9fd
a 40
        void update(int c, int a, int b, int x) { update(seg[c], a, b,
   x, 0, n-1); }
        void paint(node*& from, node*& to, int a, int b, int l, int r)
70 c
   {
            if (to == from or !from or b < 1 or r < a) return;
10f
e 85
            from -> prop():
889
            if (to) to->prop();
9a3
            if (a <= 1 and r <= b) {
24 d
                if (!to) {
38f
                     to = from;
140
                     from = NULL:
505
                     return;
                }
e5f
ee4
                 int m = (1+r)/2;
1cb
                 paint(from->1, to->1, a, b, 1, m), paint(from->r,
   to->r, a, b, m+1, r);
72 d
                 to ->update();
270
                 delete from:
140
                 from = NULL:
```

```
505
                return:
            }
a0e
019
            if (!to) to = new node();
            int m = (1+r)/2;
ee4
            paint(from->1, to->1, a, b, 1, m), paint(from->r, to->r,
1cb
   a, b, m+1, r);
45a
            from->update(), to->update();
4aa
471
        void paint(int c1, int c2, int a, int b) { paint(seg[c1],
   seg[c2], a, b, 0, n-1); }
293 };
      SegTree Esparsa - Lazy
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
//
// Complexidades:
// build - 0(1)
// query - 0(log(n))
// update - 0(log(n))
aa4 namespace seg {
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
6de
e9a
        int get l(int i){
            if (L[i] == 0) L[i] = ptr++;
3db
a96
            return L[i];
b6e
        }
943
        int get_r(int i){
71b
            if (R[i] == 0) R[i] = ptr++;
283
            return R[i];
43a
        }
```

void build() { ptr = 2; }

lazy[get\_r(p)]^=lazy[p];

lazv[p] = 0:

prop(p, 1, r);

void prop(int p, int 1, int r) {

if (l != r) lazy[get\_l(p)]^=lazy[p],

int query(int a, int b, int p=1, int l=0, int r= $\mathbb{N}$ -1) {

if (!lazy[p]) return;
seg[p] = r-l+1 - seg[p];

e71

ceb

b77

76c

213

3c7

20b

158

6 b 9

```
786
             if (b < 1 \text{ or } r < a) \text{ return } 0;
527
             if (a <= l and r <= b) return seg[p];</pre>
ee4
             int m = (1+r)/2;
             return query(a, b, get_l(p), l, m)+query(a, b, get_r(p),
818
    m+1, r);
0d9
        }
51f
         int update(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
             prop(p, 1, r);
e9f
             if (b < 1 or r < a) return seg[p];</pre>
9a3
             if (a <= 1 and r <= b) {</pre>
ab6
                 lazv[p] ^= 1:
6b9
                 prop(p, 1, r);
534
                 return seg[p];
8 e 4
ee4
             int m = (1+r)/2;
             return seg[p] = update(a, b, get_l(p), l, m)+update(a, b,
43a
    get_r(p), m+1, r);
        }
1dc
dc3 };
        SegTree Esparsa - O(q) memoria
11.23
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
//
// Complexidades:
// query - 0(log(n))
// update - 0(log(n))
13d template < typename T > struct seg {
         struct node {
3 c 9
             node* ch[2]:
d53
970
             char d:
ca0
             T v;
с4е
             T mi:
d4e
             node(int d_, T v_, T val) : d(d_), v(v_) {
                 ch \lceil 0 \rceil = ch \lceil 1 \rceil = NULL:
e 71
d6e
                 mi = val;
065
b32
             node (node * x) : d(x->d), v(x->v), mi(x->mi) {
```

ch[0] = x -> ch[0], ch[1] = x -> ch[1];

c99

```
cb7
01e
            void update() {
909
                 mi = numeric_limits < T > :: max();
                 for (int i = 0; i < 2; i++) if (ch[i])
151
b5a
                     mi = min(mi, ch[i]->mi);
            }
fe3
530
        };
bb7
        node * root;
9c5
        char n;
ba7
        seg() : root(NULL), n(0) {}
512
        \simseg() {
4 c 0
             std::vector<node*> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
73f
                 delete x;
bf0
d3e
            }
        }
d8c
        char msb(T v, char l, char r) { // msb in range (l, r]
1a6
8 e 4
             for (char i = r; i > l; i--) if (v>>i&1) return i;
daa
             return -1:
688
430
        void cut(node* at, T v, char i) {
677
             char d = msb(v \cdot at -> v \cdot at -> d \cdot i);
23b
             if (d == -1) return; // no need to split
ebf
             node* nxt = new node(at);
d43
            at -> ch \lceil v >> d&1 \rceil = NULL:
34f
            at -> ch[!(v>>d&1)] = nxt;
            at -> d = d:
150
0b3
        }
6e5
        node* update(node* at, T idx, T val, char i) {
с8 с
             if (!at) return new node(-1, idx, val);
d67
             cut(at, idx, i);
1 a 2
             if (at -> d == -1) { // leaf}
792
                 at->mi = val;
ce6
                 return at;
a6f
b29
            bool dir = idx>>at->d&1;
             at->ch[dir] = update(at->ch[dir], idx, val, at->d-1);
c8f
7b4
            at->update();
ce6
            return at:
76d
        }
```

```
85 с
        void update(T idx, T val) {
8f4
            while (idx >> n) n++;
61 e
            root = update(root, idx, val, n-1);
79 d
        }
        T query(node* at, T a, T b, T l, T r, char i) {
9d8
df0
            if (!at or b < l or r < a) return numeric limits <T>::max():
            if (a <= l and r <= b) return at->mi;
fd3
841
            T m = 1 + (r-1)/2;
            if (at->d < i) {</pre>
c.85
c59
                if ((at->v>>i\&1) == 0) return query(at, a, b, 1, m,
   i-1);
                else return query(at, a, b, m+1, r, i-1);
ca4
934
373
            return min(query(at->ch[0], a, b, 1, m, i-1),
   query(at->ch[1], a, b, m+1, r, i-1));
2db
        T query(T 1, T r) { return query(root, 1, r, 0, (T(1) << n)-1,
034
   n-1); }
d7f }:
11.24 SegTree Iterativa
// Consultas O-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a. b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
6a4 int seg[2 * MAX];
1a8 int n;
0a8 void build() {
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] + seg[2*i+1];
9a8 }
4ea int query(int a, int b) {
        int ret = 0:
7 c.9
```

for(a += n, b += n; a <= b; ++a /= 2, --b /= 2) {

if (a % 2 == 1) ret += seg[a]:

if (b % 2 == 0) ret += seg[b];

728

4ea

244

ac0

}

```
edf
        return ret;
24a }
ff3 void update(int p, int x) {
37d
        seg[p += n] = x;
с8 с
        while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
02d }
11.25 SegTree Iterativa com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
// Complexidades:
// build - 0(n)
// query - 0(log(n))
// update - 0(log(n))
aa4 namespace seg {
6db
        11 seg[2*MAX], lazy[2*MAX];
        int n:
1 a 8
        ll junta(ll a, ll b) {
9b3
534
            return a+b:
e26
        // soma x na posicao p de tamanho tam
        void poe(int p, ll x, int tam, bool prop=1) {
1b4
517
            seg[p] += x*tam;
6ae
            if (prop and p < n) lazy[p] += x;</pre>
8bc
        }
        // atualiza todos os pais da folha p
        void sobe(int p) {
b1e
d5a
            for (int tam = 2; p /= 2; tam *= 2) {
                seg[p] = junta(seg[2*p], seg[2*p+1]);
4 ca
                poe(p, lazy[p], tam, 0);
388
            }
acd
        }
b76
        // propaga o caminho da raiz ate a folha p
        void prop(int p) {
a0a
076
            int tam = 1 << (LOG-1);</pre>
            for (int s = LOG; s; s--, tam /= 2) {
```

0 a 8

```
4 b 1
                 int i = p \gg s;
27 c
                 if (lazy[i]) {
860
                     poe(2*i, lazy[i], tam);
                     poe(2*i+1, lazy[i], tam);
e 38
b97
                     lazy[i] = 0;
de8
            }
3ed
         }
e 29
         void build(int n2, int* v) {
61 c
1e3
             n = n2;
95 f
             for (int i = 0; i < n; i++) seg[n+i] = v[i];
c41
             for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
    seg[2*i+1]);
             for (int i = 0; i < 2*n; i++) lazy[i] = 0;</pre>
f4c
8bb
4f3
         11 query(int a, int b) {
b73
             ll ret = 0;
             for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2) {
b48
                 if (a%2 == 1) ret = junta(ret, seg[a]);
a8e
 c58
                 if (b%2 == 0) ret = junta(ret, seg[b]);
510
             }
edf
             return ret;
38b
        }
a28
         void update(int a, int b, int x) {
c2d
             int a2 = a += n, b2 = b += n, tam = 1;
             for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
Off
32 a
                 if (a\%2 == 1) poe(a, x, tam);
                 if (b\%2 == 0) poe(b, x, tam);
9da
9bc
0 f 7
             sobe(a2), sobe(b2);
adc
        }
6dc };
11.26 SegTree PA
// Segtree de PA
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
// update_add soma PA(A, R) em [1, r]
// query(1, r) retorna a soma de [1, r]
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
//
// Complexidades:
```

```
// construir - O(n)
// update_set, update_add, query - O(log(n))
                                                                            06с
                                                                                                 seg[2*p+1].add_a += add_a + add_r * (m-l+1);
                                                                                                 seg[2*p+1].add_r += add_r;
                                                                            a6d
                                                                            8af
dc7 struct seg_pa {
350
        struct Data {
                                                                            953
                                                                                             add a = add r = 0:
8f5
            11 sum:
                                                                            ab7
                                                                                    }
662
            ll set_a, set_r, add_a, add_r;
                                                                            07f
            Data() : sum(0), set_a(LINF), set_r(0), add_a(0), add_r(0)
9b7
   {}
                                                                            0b7
                                                                                    int inter(pair<int, int> a, pair<int, int> b) {
                                                                            98 c
eb6
        }:
                                                                                         if (a.first > b.first) swap(a, b);
        vector < Data > seg;
                                                                            eef
                                                                                         return max(0, min(a.second, b.second) - b.first + 1);
16a
1 a 8
                                                                            628
        int n:
                                                                            be1
                                                                                    11 set(int a, int b, ll aa, ll rr, int p, int l, int r) {
d45
        seg_pa(int n_) {
                                                                            6b9
                                                                                         prop(p, 1, r);
                                                                                        if (b < l or r < a) return seg[p].sum;</pre>
e95
            n = n_{-};
                                                                            457
fc3
                                                                            9a3
                                                                                        if (a <= l and r <= b) {</pre>
            seg = vector < Data > (4*n);
ce0
        }
                                                                            91 c
                                                                                             seg[p].set_a = aa;
                                                                            774
                                                                                             seg[p].set_r = rr;
        void prop(int p, int 1, int r) {
                                                                            6b9
                                                                                             prop(p, 1, r);
ceb
            int tam = r-l+1:
                                                                                             return seg[p].sum;
d5a
                                                                            254
                                                                                        }
c3f
            11 &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r =
                                                                            8ee
                                                                                        int m = (1+r)/2;
   seg[p].set_r,
                                                                            ee4
                                                                                         int tam_l = inter({1, m}, {a, b});
                &add_a = seg[p].add_a, &add_r = seg[p].add_r;
                                                                            963
a1b
                                                                            c34
                                                                                         return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m) +
            if (set a != LINF) {
                                                                            365
                                                                                             set(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
c02
660
                set_a += add_a, set_r += add_r;
                                                                            8e2
                                                                                    }
                sum = set_a*tam + set_r*tam*(tam+1)/2;
06e
                                                                            f55
                                                                                    void update_set(int 1, int r, ll aa, ll rr) {
579
                if (1 != r) {
                                                                            6f7
                                                                                         set(1, r, aa, rr, 1, 0, n-1);
                    int m = (1+r)/2;
                                                                            913
ee4
                                                                            5f6
                                                                                    11 add(int a, int b, ll aa, ll rr, int p, int l, int r) {
886
                    seg[2*p].set_a = set_a;
                                                                            6b9
                                                                                        prop(p, 1, r);
358
                    seg[2*p].set_r = set_r;
                                                                            457
                                                                                        if (b < l or r < a) return seg[p].sum;</pre>
                    seg[2*p].add_a = seg[2*p].add_r = 0;
                                                                                         if (a <= 1 and r <= b) {</pre>
ed6
                                                                            9a3
                                                                            359
                                                                                             seg[p].add a += aa:
                     seg[2*p+1].set_a = set_a + set_r * (m-l+1);
                                                                                             seg[p].add_r += rr;
f0c
                                                                            1ee
471
                     seg[2*p+1].set_r = set_r;
                                                                            6b9
                                                                                             prop(p, 1, r);
d48
                    seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
                                                                            254
                                                                                             return seg[p].sum;
a97
                                                                            d19
                                                                                        }
823
                set_a = LINF, set_r = 0;
                                                                                        int m = (1+r)/2;
                                                                            ee4
953
                add_a = add_r = 0;
                                                                            963
                                                                                         int tam_l = inter({1, m}, {a, b});
            } else if (add a or add r) {
                                                                                        return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m) +
da7
                                                                            586
                sum += add_a*tam + add_r*tam*(tam+1)/2;
                                                                            695
                                                                                             add(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
18b
579
                if (1 != r) {
                                                                            904
ee4
                    int m = (1+r)/2;
                                                                            848
                                                                                    void update_add(int 1, int r, 11 aa, 11 rr) {
                                                                                         add(1, r, aa, rr, 1, 0, n-1);
                                                                            afa
                                                                            81 e
ff0
                    seg[2*p].add_a += add_a;
                    seg[2*p].add_r += add_r;
ec0
                                                                            f 45
                                                                                    ll query(int a, int b, int p, int l, int r) {
```

```
6b9
            prop(p, 1, r);
            if (b < 1 or r < a) return 0;
786
e9a
            if (a <= 1 and r <= b) return seg[p].sum;</pre>
            int m = (1+r)/2;
ee4
b1f
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
f6e
bfc
        11 query(int 1, int r) { return query(1, r, 1, 0, n-1); }
bc4 }:
11.27 SegTree Persistente
// SegTree de soma, update de somar numa posicao
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t. criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - O(log(n))
// update - 0(log(n))
```

```
6de const int MAXS = 2*MAX+UPD*LOG;

f6e namespace perseg {
bd6    ll seg[MAXS];
```

54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;

rt[0] = cnt++;

build(0, 0, n-1);

int m = (1+r)/2;

856

c50 a2e

f45 786

527

ee4

1 e d

}

```
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
       int n. *v:
3 c 4
        11 build(int p, int l, int r) {
6cd
            if (1 == r) return seg[p] = v[1];
855
           L[p] = cnt++, R[p] = cnt++;
ee4
            int m = (1+r)/2;
275
            return seg[p] = build(L[p], 1, m) + build(R[p], m+1, r);
39d
0 d8
        void build(int n2, int* v2) {
680
            n = n2, v = v2;
```

11 query(int a, int b, int p, int l, int r) {

if (a <= l and r <= b) return seg[p];</pre>

return query(a, b, L[p], 1, m) + query(a, b, R[p], m+1, r);

if (b < 1 or r < a) return 0;

```
if (a \le m)
ab8
b48
                return seg[p] = update(a, x, L[lp], L[p]=cnt++, l, m)
   + seg[R[p]=R[lp]];
            return seg[p] = seg[L[p]=L[lp]] + update(a, x, R[lp],
8a9
   R[p] = cnt ++, m+1, r);
788
6f6
        int update(int a, int x, int tt=t) {
            update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
ab3
e0d
            return t:
d63
        }
26f }:
       SegTree Persistente com Lazy
11.28
// Nao propaga, meio estranho de mexer, mas da
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, b, x, t) faz um update v[a..b]+=x a partir da
// versao de t. criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - \Omega(n)
// query - O(log(n))
// update - 0(log(n))
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
ab3 const int MAXS = 2*MAX + 4*UPD*LOG;
f6e namespace perseg {
9eb
        int seg[MAXS];
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
adf
        int build(int p, int l, int r) {
6cd
            if (1 == r) return seg[p] = v[1];
855
            L[p] = cnt++, R[p] = cnt++;
ee4
            int m = (1+r)/2;
            return seg[p] = max(build(L[p], 1, m), build(R[p], m+1,
01 d
```

11 query(int a, int b, int tt) {

int m = (1+r)/2;

return query(a, b, rt[tt], 0, n-1);

if (1 == r) return seg[p] = seg[lp]+x;

ll update(int a, int x, int lp, int p, int l, int r) {

4 d 2

182

c13

726

bb3

747

ee4

```
ffd
                                                                            6ef
                                                                                     priority_queue < T, vector < T>, greater <>> R;
        void build(int n2. int *v2) {
0 d8
                                                                            a20
                                                                                    T add_l, add_r;
680
            n = n2, v = v2;
856
            rt[0] = cnt++;
                                                                            055
                                                                                    T top_R() {
c50
            build(0, 0, n-1);
                                                                                         if (R.empty()) return inf;
                                                                            a34
a2e
                                                                            ffe
                                                                                         else return R.top() + add_r;
976
        int querv(int a, int b, int p, int l, int r) {
                                                                            074
                                                                                    }
27b
            if (b < l or r < a) return -INF;
793
            if (a <= l and r <= b) return lazy[p] + seg[p];</pre>
                                                                            70 с
                                                                                    T pop_R() {
            int m = (1+r)/2:
ee4
                                                                            66 f
                                                                                         T \text{ val} = top_R();
7a2
            int ret = lazy[p] + max(query(a, b, L[p], l, m), query(a,
                                                                            8e0
                                                                                         if (R.size()) R.pop();
   b, R[p], m+1, r));
                                                                            d94
                                                                                         return val:
edf
            return ret:
                                                                            21 d
                                                                                    }
9a7
442
        int query(int a, int b, int tt) {
                                                                            d9d
                                                                                    T top_L() {
            return query(a, b, rt[tt], 0, n-1);
                                                                                         if (L.empty()) return -inf;
c13
                                                                            b7b
a05
        }
                                                                            470
                                                                                         else return L.top() + add_1;
        int update(int a, int b, int x, int lp, int p, int l, int r) {
                                                                                    }
bc1
                                                                            31 d
            tie(seg[p], lazy[p], L[p], R[p]) = {seg[lp], lazy[lp],
3f6
   L[lp], R[lp]};
                                                                            821
                                                                                    T pop_L() {
847
            if (b < l or r < a) return seg[p] + lazy[p];</pre>
                                                                            66a
                                                                                         T val = top_L();
            if (a <= 1 and r <= b) return seg[p] + (lazy[p] += x);</pre>
                                                                            1e0
                                                                                         if (L.size()) L.pop();
32a
                                                                            d94
                                                                                         return val:
            int m = (1+r)/2;
                                                                            dfd
                                                                                    }
ee4
24a
            seg[p] = max(update(a, b, x, L[lp], L[p] = cnt++, l, m),
bdb
                          update(a, b, x, R[lp], R[p] = cnt++, m+1, r));
                                                                            86 d
                                                                                     size t size() {
1 e d
            lazy[p] = lazy[lp];
                                                                            7ff
                                                                                         return L.size() + R.size();
1 b 7
            return seg[p] + lazy[p];
                                                                            c4b
                                                                                    }
877
cbf
        int update(int a, int b, int x, int tt=t) {
                                                                                    SlopeTrick() : min_f(0), add_1(0), add_r(0) {};
                                                                            0e8
            assert(tt <= t);
aa8
661
            update(a, b, x, rt[tt], rt[++t]=cnt++, 0, n-1);
                                                                                    // return {min f(x), lx, rx}
                                                                                    // Em que [lx, rx] eh o intervalo que atinge o minimo
e0d
            return t:
aad
        }
                                                                            5ee
                                                                                     arrav<T. 3> querv() {
f27 };
                                                                                         return {min_f, top_L(), top_R()};
                                                                            e8a
                                                                            14f
                                                                                    }
                                                                                    // f(x) += a
11.29 SlopeTrick
                                                                            ad4
                                                                                     void add all(T a) {
                                                                            f8c
                                                                                         min_f += a;
// Armazena uma estrutura convexa piecewise linear
                                                                            78a
                                                                                    }
// Permite adicionar slopes sem peso e realizar query de minimo
// Comentarios acima das funcoes para explicar o que cada uma faz
                                                                                    // add \_
                                                                                    // f(x) += max(a - x, 0)
2f8 template < typename T> struct SlopeTrick {
                                                                                     void add_a_minus_x(T a) {
                                                                            60 a
64e
        T inf = numeric_limits < T > :: max() / 3;
                                                                            8 c 6
                                                                                         min_f += max(T(0), a - top_R());
acc
        T min_f;
                                                                                         R.push(a - add_r);
                                                                            cdb
```

f32

priority\_queue <T, vector <T>, less <>> L;

```
416
           L.push(pop_R() - add_l);
44c
        // add /
       // f(x) += max(x - a, 0)
        void add_x_minus_a(T a) {
7a9
            min_f += max(T(0), top_L() - a);
b36
988
           L.push(a - add_l);
           R.push(pop_L() - add_r);
e5a
f3a
       // add \/
       // f(x) += abs(x - a)
        void add_abs(T a) {
825
9сс
            add_a_minus_x(a);
e55
            add_x_minus_a(a);
639
       }
       // \/ -> \_
       // f_{new} (x) = min f(y) (y <= x)
        void clear_right() {
73b
b8e
            while (R.size()) R.pop();
2 b 3
       }
       // \/ -> /
       // f_{new} (x) = min f(y) (y >= x)
        void clear_left() {
fd5
e21
            while (L.size()) L.pop();
bc4
       }
       // \/ -> \ /
       // f_{new} (x) = min f(y) (x-b <= y <= x-a)
        void shift(T a, T b) {
564
25b
           assert(a <= b);
b95
           add_1 += a;
165
            add_r += b;
29a
       }
       // \/. -> .\/
       // f_{new} (x) = f(x - a)
        void shift(T a) {
5 d 6
a77
            shift(a, a);
       }
af 1
        // Retorna f(x)
        // O(size)
        T get(T x) {
c2a
```

```
7ce
             auto L2 = L;
202
             auto R2 = R:
bf4
            T ret = min_f;
6a9
             while (L.size()) {
efd
                 ret += max(T(0), pop_L() - x);
e50
886
             while (R.size()) {
97b
                 ret += max(T(0), x - pop_R());
8ef
            L = L2, R = R2;
98 a
edf
             return ret;
093
        }
        // O(min(size, st.size))
        void merge(SlopeTrick &st) {
9e9
             if (st.size() > size()) {
f 68
079
                 swap(*this, st);
788
1a3
             while (st.R.size()) {
85 b
                 add_x_minus_a(st.pop_R());
2 c 5
8c6
             while (st.L.size()) {
897
                 add_a_minus_x(st.pop_L());
b31
            }
eaf
             min_f += st.min_f;
3df
495 };
11.30 Sparse Table
// Resolve RMO
// MAX2 = log(MAX)
// Complexidades:
// build - O(n log(n))
// query - 0(1)
cca namespace sparse {
710
        int m[MAX2][MAX], n;
        void build(int n2, int* v) {
61 c
1 e 3
            n = n2:
78 e
            for (int i = 0; i < n; i++) m[0][i] = v[i];
            for (int j = 1; (1<<j) <= n; j++) for (int i = 0; i+(1<<<math>j)
a1c
\langle = n : i++ \rangle
5d5
                 m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
cae
        }
```

### 11.31 Sparse Table Disjunta

```
// Resolve qualquer operacao associativa
// MAX2 = log(MAX)
// Complexidades:
// build - 0(n log(n))
// query - 0(1)
cca namespace sparse {
        int m[MAX2][2*MAX], n, v[2*MAX];
9bf
5f7
        int op(int a, int b) { return min(a, b); }
0 d 8
        void build(int n2, int* v2) {
1 e 3
            n = n2;
df4
            for (int i = 0; i < n; i++) v[i] = v2[i];
a84
            while (n&(n-1)) n++;
3d2
            for (int j = 0; (1<<j) < n; j++) {
                int len = 1<<i:
1 c.0
d9b
                for (int c = len; c < n; c += 2*len) {</pre>
                    m[j][c] = v[c], m[j][c-1] = v[c-1];
332
668
                    for (int i = c+1; i < c+len; i++) m[j][i] =
   op(m[j][i-1], v[i]);
                    for (int i = c-2; i >= c-len; i--) m[i][i] =
432
   op(v[i], m[j][i+1]);
eda
f4d
ce3
9 e 3
        int query(int 1, int r) {
f13
            if (1 == r) return v[1];
            int j = __builtin_clz(1) - __builtin_clz(1^r);
e6d
d67
            return op(m[j][1], m[j][r]);
a7b
        }
fd8 }
```

## 11.32 Splay Tree

```
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO // NODE MAIS PROFUNDO VISITADO
```

```
// Todas as operacoes sao O(log(n)) amortizado
// Se quiser colocar mais informacao no node,
// mudar em 'update'
538 template < typename T > struct splaytree {
3 c 9
         struct node {
183
             node *ch[2]. *p:
e4d
             int sz:
f 48
             T val;
da0
             node(T v) {
696
                 ch[0] = ch[1] = p = NULL;
a26
                 sz = 1:
250
                 val = v:
2d0
01e
             void update() {
a 26
                 sz = 1:
с7с
                 for (int i = 0; i < 2; i++) if (ch[i]) {
d5f
                      sz += ch[i]->sz:
486
                 }
             }
f 45
aa3
        };
bb7
         node* root;
fbc
         splaytree() { root = NULL; }
214
         splaytree(const splaytree& t) {
cbf
             throw logic_error("Nao copiar a splaytree!");
1f1
891
        \simsplaytree() {
609
             vector < node *> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
73 f
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                 delete x;
d3e
            }
837
        }
94f
         void rotate(node* x) { // x vai ficar em cima
d9b
             node *p = x->p, *pp = p->p;
             if (pp) pp -> ch[pp -> ch[1] == p] = x;
ecf
286
             bool d = p \rightarrow ch[0] == x;
d63
             p -> ch[!d] = x -> ch[d], x -> ch[d] = p;
bad
             if (p->ch[!d]) p->ch[!d]->p = p;
fc2
             x \rightarrow p = pp, p \rightarrow p = x;
             p->update(), x->update();
1ea
007
        }
```

```
3fa
        node* splay(node* x) {
a39
            if (!x) return x;
4ea
            root = x;
             while (x->p) {
3cf
d9b
                 node *p = x->p, *pp = p->p;
359
                 if (!pp) return rotate(x), x; // zig
еЗс
                 if ((pp -> ch[0] == p)^(p -> ch[0] == x))
                     rotate(x), rotate(x); // zigzag
a2b
4b2
                 else rotate(p), rotate(x); // zigzig
            }
028
ea5
             return x;
379
313
        node* insert(T v. bool lb=0) {
b64
             if (!root) return lb ? NULL : root = new node(v);
002
             node *x = root, *last = NULL;;
31e
             while (1) {
                 bool d = x -> val < v;
5 d7
                 if (!d) last = x:
0fd
                 if (x->val == v) break;
с2е
                 if (x->ch[d]) x = x->ch[d];
c16
4 e 6
                 else {
                     if (lb) break;
dea
                     x -> ch [d] = new node(v);
055
99c
                     x \rightarrow ch[d] \rightarrow p = x;
30e
                     x = x -> ch[d]:
c2b
                     break:
                 }
68a
            }
1ab
0b6
             splay(x);
61c
             return lb ? splay(last) : x;
622
сОс
        int size() { return root ? root->sz : 0; }
2ca
        int count(T v) { return insert(v, 1) and root->val == v; }
        node* lower bound(T v) { return insert(v, 1); }
111
        void erase(T v) {
26b
446
            if (!count(v)) return;
             node *x = root, *1 = x -> ch[0];
bсе
268
            if (!1) {
8b1
                 root = x -> ch[1];
                 if (root) root->p = NULL;
32e
8f3
                 return delete x;
a86
5 e 7
             root = 1, 1 \rightarrow p = NULL;
             while (1-> ch[1]) 1 = 1-> ch[1];
902
bab
             splay(1);
            1 -> ch[1] = x -> ch[1];
f0e
            if (1->ch[1]) 1->ch[1]->p = 1;
7 d 9
```

```
bf0
             delete x;
62 a
             1->update();
007
24a
         int order_of_kev(T v) {
62b
              if (!lower_bound(v)) return root ? root->sz : 0;
1cc
              return root -> ch[0] ? root -> ch[0] -> sz : 0;
b00
         }
db6
         node* find_by_order(int k) {
084
             if (k >= size()) return NULL;
52 f
             node * x = root:
31 e
             while (1) {
20f
                  if (x \rightarrow ch[0] \text{ and } x \rightarrow ch[0] \rightarrow sz >= k+1) x = x \rightarrow ch[0];
4e6
a1c
                       if (x->ch[0]) k -= x->ch[0]->sz;
1dc
                       if (!k) return splay(x);
                       k--, x = x-> ch[1];
eb8
                  }
aca
             }
e01
         }
0de
         T min() {
19 c
52 f
             node * x = root;
6f6
              while (x -> ch[0]) x = x-> ch[0]; // max -> ch[1]
3e9
             return splay(x)->val;
70e
         }
4ff }:
```

### 11.33 Splay Tree Implicita

```
// vector da NASA
// Um pouco mais rapido q a treap
// O construtor a partir do vector
// eh linear, todas as outras operacoes
// custam O(log(n)) amortizado
081 template < typename T > struct splay {
3 c 9
        struct node {
183
            node *ch[2], *p;
e4d
             int sz;
875
            T val, sub, lazy;
            bool rev:
aa6
da0
            node(T v) {
696
                 ch[0] = ch[1] = p = NULL;
a26
                 sz = 1;
1 e 4
                 sub = val = v;
c60
                 lazv = 0;
b67
                 rev = false;
```

```
48f
a9c
            void prop() {
0ec
                 if (lazy) {
                     val += lazy, sub += lazy*sz;
924
091
                     if (ch[0]) ch[0]->lazy += lazy;
                     if (ch[1]) ch[1]->lazy += lazy;
1 a 8
                }
a98
                 if (rev) {
1 bb
80a
                     swap(ch[0], ch[1]);
                     if (ch[0]) ch[0]->rev ^= 1;
628
adc
                     if (ch[1]) ch[1]->rev ^= 1;
30a
a32
                 lazv = 0, rev = 0;
6bf
            }
01e
            void update() {
0 c 3
                 sz = 1, sub = val;
с7с
                 for (int i = 0; i < 2; i++) if (ch[i]) {
05f
                     ch[i]->prop();
d5f
                     sz += ch[i] -> sz;
4 a 1
                     sub += ch[i]->sub;
6 c 1
                }
e98
b4a
        };
bb7
        node * root;
5d9
        splay() { root = NULL; }
9 b 1
        splay(node* x) {
4ea
            root = x:
32e
            if (root) root->p = NULL;
371
1 b 7
        splay(vector < T > v) { // O(n)}
950
            root = NULL:
806
            for (T i : v) {
                 node* x = new node(i);
2a0
bd1
                 x -> ch[0] = root;
37a
                 if (root) root->p = x;
4ea
                 root = x;
a0a
                 root ->update();
17с
            }
c6b
a9e
        splay(const splay& t) {
e62
             throw logic_error("Nao copiar a splay!");
        }
d4d
5ab
        \simsplav() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
```

```
e5d
                  node * x = q.back(); q.pop_back();
ee9
                  if (!x) continue;
73f
                  q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                  delete x;
             }
d3e
        }
a1c
73 с
         int size(node* x) { return x ? x->sz : 0; }
94f
         void rotate(node* x) { // x vai ficar em cima
d9b
             node *p = x->p, *pp = p->p;
ecf
             if (pp) pp->ch[pp->ch[1] == p] = x;
286
             bool d = p \rightarrow ch[0] == x;
d63
             p \rightarrow ch[!d] = x \rightarrow ch[d], x \rightarrow ch[d] = p;
bad
             if (p->ch[!d]) p->ch[!d]->p = p;
fc2
             x \rightarrow p = pp, p \rightarrow p = x;
             p->update(), x->update();
1ea
007
         node* splaya(node* x) {
6a0
a39
             if (!x) return x;
             root = x, x->update();
be6
3cf
             while (x->p) {
                  node *p = x->p, *pp = p->p;
d9b
359
                  if (!pp) return rotate(x), x; // zig
еЗс
                  if ((pp -> ch[0] == p)^(p -> ch[0] == x))
a2b
                      rotate(x), rotate(x); // zigzag
4b2
                  else rotate(p), rotate(x); // zigzig
028
ea5
             return x:
21a
        }
a7f
         node* find(int v) {
             if (!root) return NULL;
a2e
52f
             node *x = root;
6cd
             int key = 0;
31 e
             while (1) {
857
                  x -> prop();
                  bool d = key + size(x->ch[0]) < v;
ba1
877
                  if (\text{kev} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
15 e
                      if (d) key += size(x -> ch[0])+1;
30 e
                      x = x -> ch[d]:
a30
                  } else break;
3 c 3
152
             return splaya(x);
f 19
         int size() { return root ? root->sz : 0; }
сОс
c26
         void join(splay<T>& 1) { // assume que l < *this</pre>
690
             if (!size()) swap(root, l.root);
579
             if (!size() or !l.size()) return;
```

```
bee
            node* x = 1.root;
            while (1) {
31e
857
                 x->prop();
34d
                 if (!x->ch[1]) break;
bd8
                 x = x -> ch[1];
fa3
147
            1.splaya(x), root->prop(), root->update();
            x -> ch[1] = root, x -> ch[1] -> p = x;
42b
0aa
            root = 1.root, 1.root = NULL;
            root ->update();
a0a
7 e 6
5 e d
        node* split(int v) { // retorna os elementos < v</pre>
398
            if (v <= 0) return NULL;</pre>
            if (v >= size()) {
060
f87
                 node* ret = root;
950
                 root = NULL;
8 c 9
                 ret->update();
edf
                 return ret;
            }
d0f
adc
            find(v):
            node*1 = root -> ch[0];
a59
4df
            root -> ch[0] = NULL;
            if (1) 1->p = NULL;
5 a 3
a0a
            root ->update();
792
            return 1;
826
511
        T& operator [](int i) {
9 d 4
            find(i):
ae0
            return root -> val;
829
231
        void push_back(T v) { // 0(1)
            node* r = new node(v);
a01
0 de
            r \rightarrow ch[0] = root:
            if (root) root->p = r;
b11
            root = r, root->update();
b13
        }
315
b7a
        T query(int 1, int r) {
95f
            splay<T> M(split(r+1));
5ff
            splay < T > L (M.split(1));
d1c
            T ans = M.root->sub;
49с
            M. join(L), join(M);
ba7
            return ans;
ca3
        void update(int 1, int r, T s) {
41f
95f
            splay<T> M(split(r+1));
5ff
            splay < T > L (M.split(1));
            M.root->lazy += s;
996
```

```
49 с
             M.join(L), join(M);
9e9
8c1
        void reverse(int 1, int r) {
95 f
             splay <T> M(split(r+1));
5ff
             splay <T > L(M.split(1));
945
             M.root->rev ^= 1;
49 с
             M. join(L), join(M);
        }
c1a
2fb
        void erase(int 1, int r) {
95 f
             splay <T > M(split(r+1));
5ff
             splay <T> L(M.split(1));
dcc
             ioin(L);
68 e
a35 };
11.34 Split-Merge Set
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge, que custa O(log(N)) amortizado
// Usa O(\min(N, n \log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
2dc template < typename T, bool MULTI = false, typename SIZE_T = int > struct
    sms {
3 c 9
        struct node {
b19
             node *1, *r;
15 f
             SIZE T cnt:
658
             node() : l(NULL), r(NULL), cnt(0) {}
01e
             void update() {
a 0 1
                 cnt = 0;
d8a
                 if (1) cnt += 1->cnt;
e 49
                 if (r) cnt += r \rightarrow cnt;
74 d
            }
84f
        };
bb7
        node* root;
fd0
        T N:
f34
        sms() : root(NULL), N(0) {}
83b
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
5 e 1
        sms(const sms& t) : root(NULL), N(t.N) {
3af
             for (SIZE T i = 0; i < t.size(); i++) {</pre>
a0f
                 T at = t[i];
e6d
                 SIZE_T qt = t.count(at);
```

```
a43
                insert(at, qt);
                                                                             8d9
                                                                                         else at->r = insert(at->r, idx, qt, m+1, r);
                i += qt-1;
                                                                             cff
f42
                                                                                         return at -> update(), at;
1 e 9
            }
                                                                             83b
                                                                                     void insert(T v, SIZE_T qt=1) { // insere 'qt' ocorrencias de
ea8
        }
                                                                             cf7
a96
        sms(initializer list<T> v) : sms() { for (T i : v) insert(i): }
2 dd
        \simsms() {
                                                                                          if (qt <= 0) return erase(v, -qt);</pre>
                                                                             882
                                                                                         assert(v >= 0):
609
            vector<node*> a = {root};
                                                                             72b
            while (q.size()) {
402
                                                                             f 52
                                                                                         expand(v);
e5d
                 node* x = q.back(); q.pop_back();
                                                                             5e9
                                                                                         root = insert(root, v, qt, 0, N);
                if (!x) continue;
                                                                                     }
ee9
                                                                             f62
1 c7
                q.push_back(x->1), q.push_back(x->r);
bf0
                delete x:
                                                                             f06
                                                                                     node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
653
            }
                                                                             28 c
                                                                                         if (!at) return at:
                                                                             54b
                                                                                         if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt - qt;
f0d
        }
                                                                             4e6
                                                                                         else {
        friend void swap(sms& a, sms& b) {
                                                                             841
                                                                                              T m = 1 + (r-1)/2;
fdc
49e
            swap(a.root, b.root), swap(a.N, b.N);
                                                                             281
                                                                                              if (idx \le m) at \rightarrow 1 = erase(at \rightarrow 1, idx, qt, 1, m);
984
                                                                                              else at->r = erase(at->r, idx, qt, m+1, r);
                                                                             ba1
83e
        sms& operator =(const sms& v) {
                                                                             7b4
                                                                                              at -> update();
768
            sms tmp = v:
                                                                                         }
                                                                             d3d
420
                                                                                         if (!at->cnt) delete at, at = NULL;
            swap(tmp, *this);
                                                                             135
357
            return *this;
                                                                                         return at:
                                                                             ce6
e9b
                                                                             e1f
d06
        SIZE_T size() const { return root ? root->cnt : 0; }
                                                                             43 d
                                                                                     void erase(T v, SIZE_T qt=1) { // remove 'qt' ocorrencias de
17f
        SIZE T count(node * x) const { return x ? x->cnt : 0: }
                                                                              , ,, ,
        void clear() {
75a
                                                                             9 c 3
                                                                                         if (v < 0 \text{ or } v > N \text{ or } !at) return:
                                                                                         if (qt < 0) insert(v, -qt);</pre>
0a0
            sms tmp;
                                                                             9dc
            swap(*this, tmp);
4ac
                                                                             b1d
                                                                                         root = erase(root, v. qt. 0, N);
fcb
        }
                                                                             b32
a06
        void expand(T v) {
                                                                             8d6
                                                                                     void erase_all(T v) { // remove todos os 'v'
bc3
            for (: N < v: N = 2*N+1) if (root) {
                                                                             347
                                                                                         if (v < 0 \text{ or } v > N) return:
63c
                node* nroot = new node();
                                                                             9f2
                                                                                         root = erase(root, v, numeric_limits < SIZE_T >:: max(), 0, N);
956
                                                                                     }
                nroot ->1 = root;
                                                                             569
897
                root = nroot;
                                                                                     SIZE_T count(node * at, T a, T b, T 1, T r) const {
                root ->update();
a0a
                                                                             0fe
            }
dd9
                                                                             61 b
                                                                                         if (!at or b < l or r < a) return 0;
9f0
        }
                                                                             0fe
                                                                                         if (a <= l and r <= b) return at->cnt;
                                                                             841
                                                                                         T m = 1 + (r-1)/2;
b14
        node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
                                                                             84a
                                                                                         return count(at->1, a, b, 1, m) + count(at->r, a, b, m+1,
1 a 4
            if (!at) at = new node();
                                                                                r);
            if (1 == r) {
893
                                                                             4e6
435
                at->cnt += qt;
                                                                                     SIZE_T count(T v) const { return count(root, v, v, 0, N); }
                                                                             0a9
                if (!MULTI) at->cnt = 1;
                                                                                     SIZE_T order_of_key(T v) { return count(root, 0, v-1, 0, N); }
beb
                                                                             ffc
                                                                                     SIZE_T lower_bound(T v) { return order_of_key(v); }
ce6
                return at:
                                                                             df2
a53
            T m = 1 + (r-1)/2:
                                                                                     const T operator [](SIZE_T i) const { // i-esimo menor elemento
841
                                                                             e 68
                                                                                         assert(i >= 0 and i < size());
            if (idx \le m) at->1 = insert(at->1, idx, qt, 1, m);
                                                                             809
a02
```

```
c43
             node * at = root:
             T 1 = 0, r = N;
4 a 5
40c
             while (1 < r) {
841
                 T m = 1 + (r-1)/2;
5 c 2
                 if (count(at->1) > i) at = at->1, r = m;
4 e 6
b4a
                     i -= count(at->1);
                      at = at -> r: 1 = m+1:
ded
fa6
                 }
             }
41a
792
             return 1;
67f
        }
78c
        node* merge(node* 1, node* r) {
347
             if (!1 or !r) return 1 ? 1 : r;
504
             if (!1\rightarrow 1 \text{ and } !1\rightarrow r) \{ // \text{ folha} \}
599
                 if (MULTI) 1 \rightarrow cnt + r \rightarrow cnt;
55d
                 delete r:
792
                 return 1;
92c
             }
             1 - > 1 = merge(1 - > 1, r - > 1), 1 - > r = merge(1 - > r, r - > r);
f58
f4f
             1->update(), delete r;
             return 1:
792
06a
        }
f59
        void merge(sms& s) { // mergeia dois sets
068
             if (N > s.N) swap(*this, s);
             expand(s.N);
785
             root = merge(root, s.root);
938
             s.root = NULL;
ee2
2f6
        }
dc6
         node* split(node*& x, SIZE_T k) {
7ca
             if (k <= 0 or !x) return NULL:
             node* ret = new node();
6 d 0
             if (!x->1 \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
386
4 e 6
             else {
85e
                 if (k \le count(x->1)) ret->1 = split(x->1, k);
4 e 6
06f
                      ret->r = split(x->r, k - count(x->1));
cfd
                      swap(x->1, ret->1);
63b
                 ret->update(), x->update();
674
379
             if (!x->cnt) delete x, x = NULL;
d5b
edf
             return ret;
f18
02b
         void split(SIZE_T k, sms& s) { // pega os 'k' menores
```

```
e 63
            s.clear():
6e5
            s.root = split(root, min(k, size()));
еЗс
            s.N = N;
9a6
        }
        // pega os menores que 'k'
        void split_val(T k, sms& s) { split(order_of_key(k), s); }
131
2d2 };
11.35 SQRT Tree
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entag a query eh praticamente O(1))
//
// build - O(n log log n)
// query - O(log log n)
97a namespace sqrtTree {
052
        int n. *v:
        int pref[4][MAX], sulf[4][MAX], getl[4][MAX], entre[4][MAX],
e c.7
   sz[4];
5f7
        int op(int a, int b) { return min(a, b); }
        inline int getblk(int p, int i) { return (i-getl[p][i])/sz[p];
c72
  }
2 c 6
        void build(int p, int l, int r) {
            if (1+1 >= r) return;
bc8
368
            for (int i = 1; i <= r; i++) getl[p][i] = 1;</pre>
f 16
            for (int L = 1; L <= r; L += sz[p]) {
                int R = min(L+sz[p]-1, r);
191
89 с
                pref[p][L] = v[L], sulf[p][R] = v[R];
                for (int i = L+1; i <= R; i++) pref[p][i] =
59f
   op(pref[p][i-1], v[i]);
                for (int i = R-1; i >= L; i--) sulf[p][i] = op(v[i],
d9a
   sulf[p][i+1]);
221
                build(p+1, L, R);
c7b
            for (int i = 0; i <= sz[p]; i++) {</pre>
695
                int at = entre[p][l+i*sz[p]+i] = sulf[p][l+i*sz[p]];
ca5
759
                for (int j = i+1; j \le sz[p]; j++)
   entre[p][l+i*sz[p]+j] = at =
23a
                        op(at, sulf[p][1+j*sz[p]]);
c 5 1
            }
861
8b0
        void build(int n2, int* v2) {
```

```
680
            n = n2, v = v2;
            for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
44c
c50
            build(0, 0, n-1);
940
        }
        int query(int 1, int r) {
9 e 3
792
            if (1+1 >= r) return 1 == r ? v[1] : op(v[1], v[r]);
1ba
            int p = 0;
            while (getblk(p, 1) == getblk(p, r)) p++;
4ba
9 e 4
            int ans = sulf[p][1], a = getblk(p, 1)+1, b = getblk(p, 1)+1
   r)-1;
8bf
            if (a \le b) ans = op(ans, entre[p][get1[p][1]+a*sz[p]+b]);
            return op(ans, pref[p][r]);
dea
589
8ff }
11.36 Treap
// Todas as operacoes custam
// O(log(n)) com alta probabilidade, exceto meld
```

#### // meld custa O(log^2 n) amortizado com alta prob., // e permite unir duas treaps sem restricao adicional // Na pratica, esse meld tem constante muito boa e // o pior caso eh meio estranho de acontecer 878 mt19937 rng((int) chrono::steady\_clock::now().time\_since\_epoch().count()); aa1 template < typename T > struct treap { struct node { 3 c.9 b19 node \*1, \*r; 284 int p, sz; 36d T val, mi; 4 c7 node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v), mi(v) {} void update() { 01e a26 sz = 1: mi = val; d6e bd7 if (1) sz += 1->sz, mi = min(mi, 1->mi); if (r) sz += r->sz, mi = min(mi, r->mi); a54 } 660 c1b }; bb7 node \* root;

84b

2 d8

treap() { root = NULL; }

treap(const treap& t) {

```
465
             throw logic_error("Nao copiar a treap!");
        }
1e9
        \simtreap() {
cec
609
             vector < node *> q = {root};
402
             while (q.size()) {
                  node* x = q.back(); q.pop_back();
e5d
ee9
                  if (!x) continue:
1 c 7
                  q.push_back(x->1), q.push_back(x->r);
bf0
                  delete x;
653
             }
50e
        }
73 c
         int size(node* x) { return x ? x->sz : 0; }
         int size() { return size(root): }
b2b
bcf
         void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
             if (!1 or !r) return void(i = 1 ? 1 : r);
80e
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
             else join(1, r\rightarrow 1, r\rightarrow 1), i = r;
fa0
bda
             i->update();
        }
671
есе
         void split(node* i, node*& 1, node*& r, T v) {
26a
             if (!i) return void(r = l = NULL);
f 05
             if (i-\forall val < v) split(i-\forall r, i-\forall r, r, v), l = i;
807
             else split(i \rightarrow 1, l, i \rightarrow 1, v), r = i;
             i->update();
bda
2cd
        }
3fc
         void split_leq(node* i, node*& 1, node*& r, T v) {
26a
             if (!i) return void(r = l = NULL);
181
             if (i\rightarrow val \leftarrow v) split_leq(i\rightarrow r, i\rightarrow r, r, v), l=i;
58f
             else split_leg(i \rightarrow l, l, i \rightarrow l, v), r = i;
             i->update();
bda
70 f
        }
         int count(node* i, T v) {
e 13
6b4
             if (!i) return 0:
352
             if (i->val == v) return 1;
8d0
             if (v < i->val) return count(i->1, v);
4 d 0
             return count(i->r, v);
5 e 6
26 d
         void index_split(node* i, node*& 1, node*& r, int v, int key =
   0) {
             if (!i) return void(r = l = NULL);
26a
             if (key + size(i->1) < v) index_split(i->r, i->r, r, v,
c10
   key+size(i->1)+1), l = i;
             else index_split(i->1, 1, i->1, v, key), r = i;
e5a
             i->update();
bda
        }
ccf
         int count(T v) {
a1f
```

```
e06
             return count(root, v);
                                                                              3 c 9
                                                                                       struct node {
980
                                                                              b19
                                                                                           node *1, *r;
c27
        void insert(T v) {
                                                                              284
                                                                                           int p, sz;
980
            if (count(v)) return:
                                                                              875
                                                                                           T val, sub, lazy;
031
            node *L. *R:
                                                                              aa6
                                                                                           bool rev:
d42
             split(root, L, R, v);
                                                                              8dc
                                                                                           node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
585
             node * at = new node(v);
                                                                                  sub(v), lazy(0), rev(0) {}
59f
                                                                                           void prop() {
             join(L, at, L);
                                                                              а9с
a28
             join(L, R, root);
                                                                              0ec
                                                                                               if (lazy) {
37 c
        }
                                                                              924
                                                                                                   val += lazy, sub += lazy*sz;
26b
        void erase(T v) {
                                                                              b87
                                                                                                   if (1) 1->lazy += lazy;
df9
            node *L, *M, *R;
                                                                              d3b
                                                                                                   if (r) r->lazy += lazy;
b6b
             split_leq(root, M, R, v), split(M, L, M, v);
                                                                              cea
                                                                                               }
f 1 7
            if (M) delete M;
                                                                              1bb
                                                                                               if (rev) {
f38
            M = NULL;
                                                                              e4f
                                                                                                   swap(1, r);
a28
                                                                              dc8
                                                                                                   if (1) 1->rev ^= 1;
            join(L, R, root);
b92
        }
                                                                              f2f
                                                                                                   if (r) r->rev ^= 1;
e77
        void meld(treap& t) { // segmented merge
                                                                              3e5
             node *L = root, *R = t.root;
4a6
                                                                              a 32
                                                                                               lazv = 0, rev = 0;
950
             root = NULL:
                                                                                           }
                                                                              ca6
             while (L or R) {
6 b 1
                                                                              01e
                                                                                           void update() {
fe2
                 if (!L or (L and R and L->mi > R->mi)) std::swap(L, R);
                                                                              0 c 3
                                                                                               sz = 1, sub = val;
                 if (!R) join(root, L, root), L = NULL;
5 e 1
                                                                              a09
                                                                                               if (1) 1 - > prop(), sz += 1 - > sz, sub += 1 - > sub;
3 c 9
                 else if (L->mi == R->mi) {
                                                                              095
                                                                                               if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub;
a76
                     node * LL:
                                                                              360
                                                                                          }
439
                     split(L, LL, L, R->mi+1);
                                                                              d37
                                                                                      };
359
                     delete LL;
2 a 3
                } else {
                                                                              bb7
                                                                                       node* root;
a76
                     node * LL;
537
                     split(L, LL, L, R->mi);
                                                                              84b
                                                                                       treap() { root = NULL; }
dbb
                     join(root, LL, root);
                                                                              2 d 8
                                                                                       treap(const treap& t) {
f4f
                 }
                                                                              465
                                                                                           throw logic_error("Nao copiar a treap!");
            }
                                                                                      }
576
                                                                              1e9
689
            t.root = NULL:
                                                                                      \simtreap() {
                                                                              cec
8 e 7
                                                                              609
                                                                                           vector < node *> q = {root};
651 };
                                                                              402
                                                                                           while (q.size()) {
                                                                              e5d
                                                                                               node* x = q.back(); q.pop_back();
                                                                              ee9
                                                                                               if (!x) continue;
                                                                              1 c 7
                                                                                               q.push_back(x->1), q.push_back(x->r);
11.37 Treap Implicita
                                                                              bf0
                                                                                               delete x;
                                                                              653
                                                                                          }
// Todas as operacoes custam
                                                                                      }
                                                                              50e
// O(log(n)) com alta probabilidade
                                                                              73 с
                                                                                       int size(node* x) { return x ? x->sz : 0; }
878 mt19937 rng((int)
                                                                              b2b
                                                                                       int size() { return size(root); }
    chrono::steady_clock::now().time_since_epoch().count());
                                                                              bcf
                                                                                       void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
                                                                              986
                                                                                           if (!l or !r) return void(i = 1 ? 1 : r);
aa1 template < typename T > struct treap {
```

```
161
            1->prop(), r->prop();
                                                                              3c9 struct node {
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
                                                                             b19
                                                                                      node *1. *r:
fa0
            else join(1, r\rightarrow 1, r\rightarrow 1), i = r;
                                                                             f 14
                                                                                      ll sz, val, sub;
                                                                                      node(ll v) : l(NULL), r(NULL), sz(1), val(v), sub(v) {}
            i->update();
bda
                                                                              304
b57
                                                                                      node(node* x) : 1(x->1), r(x->r), sz(x->sz), val(x->val),
                                                                              c12
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
a20
                                                                                 sub(x->sub) {}
26a
            if (!i) return void(r = l = NULL);
                                                                              01e
                                                                                      void update() {
c89
            i->prop();
                                                                              0 c 3
                                                                                          sz = 1, sub = val;
5bd
            if (\text{key} + \text{size}(i->1) < v) split(i->r, i->r, r, v,
                                                                              77 e
                                                                                          if (1) sz += 1->sz, sub += 1->sub;
   key+size(i->1)+1), l = i;
                                                                                          if (r) sz += r->sz, sub += r->sub;
                                                                              d6e
219
             else split(i\rightarrow 1, l, i\rightarrow 1, v, key), r = i;
                                                                              124
                                                                                          sub %= MOD;
bda
            i->update();
                                                                              472
                                                                                      }
d37
                                                                              95f }:
231
        void push_back(T v) {
                                                                              bc9 ll size(node* x) { return x ? x->sz : 0; }
2 e 0
             node* i = new node(v);
7ab
             join(root, i, root);
                                                                              761 void update(node* x) { if (x) x->update(); }
46b
                                                                              828 node* copy(node* x) { return x ? new node(x) : NULL; }
        T query(int 1, int r) {
b7a
df9
            node *L, *M, *R;
                                                                              b02 node* join(node* 1, node* r) {
             split(root, M, R, r+1), split(M, L, M, 1);
                                                                                      if (!1 or !r) return 1 ? copy(1) : copy(r);
dca
                                                                              e1f
d43
            T ans = M->sub;
                                                                              48b
                                                                                      node* ret:
                                                                                      if (rng() % (size(1) + size(r)) < size(1)) {</pre>
69d
            join(L, M, M), join(M, R, root);
                                                                              49 f
ba7
            return ans;
                                                                              7eb
                                                                                          ret = copy(1);
1 f 7
                                                                              cc1
                                                                                          ret -> r = join(ret -> r, r);
41f
        void update(int 1, int r, T s) {
                                                                              784
                                                                                      } else {
df9
            node *L. *M. *R:
                                                                              4 c 5
                                                                                          ret = copy(r);
             split(root, M, R, r+1), split(M, L, M, 1);
dca
                                                                              551
                                                                                          ret ->1 = join(1, ret ->1);
8f6
            M \rightarrow lazy += s;
                                                                              7a0
69d
                                                                              74 f
                                                                                      return update(ret), ret;
            join(L, M, M), join(M, R, root);
29f
                                                                              2cc }
        void reverse(int 1. int r) {
8 c 1
df9
            node *L, *M, *R;
                                                                              723 void split(node* x, node*& 1, node*& r, 11 v, 11 key = 0) {
dca
             split(root, M, R, r+1), split(M, L, M, 1);
                                                                              421
                                                                                      if (!x) return void(l = r = NULL);
66a
            M -> rev ^= 1:
                                                                              b4b
                                                                                      if (kev + size(x->1) < v) {
             join(L, M, M), join(M, R, root);
                                                                              72 f
                                                                                          1 = copy(x);
69d
                                                                                          split(1->r, 1->r, r, v, key+size(1->1)+1);
ea8
        }
                                                                              d70
139 };
                                                                              710
                                                                                      } else {
                                                                              303
                                                                                          r = copv(x);
                                                                              417
                                                                                          split(r\rightarrow 1, 1, r\rightarrow 1, v, key);
                                                                              3d1
11.38 Treap Persistent Implicita
                                                                                      update(1), update(r);
                                                                              da2
                                                                              666 }
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
                                                                             f9e vector<node*> treap;
6cf mt19937_64 rng((int)
                                                                              139 void init(const vector<ll>& v) {
   chrono::steady_clock::now().time_since_epoch().count());
                                                                              bbd
                                                                                      treap = {NULL};
```

#### 11.39 Wavelet Tree

```
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, v]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - 0(log(sigma))
// sum - 0(log(sigma))
// sumk - O(log(sigma))
597 int n. v[MAX]:
578 vector < int > esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
f8d void build(int b = 0, int e = n, int p = 1, int l = MINN, int r =
    MAXN) {
58f
        int m = (1+r)/2; esq[p].push_back(0); pref[p].push_back(0);
f2f
        for (int i = b; i < e; i++) {
6 b 9
            esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
26f
            pref[p].push_back(pref[p].back()+v[i]);
206
8ce
        if (1 == r) return;
        int m2 = stable_partition(v+b, v+e, [=](int i){return i <=
3a7
    m; \}) - v;
347
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
Ofb }
540 int count(int i, int j, int x, int y, int p = 1, int l = MINN, int
    r = MAXN)  {
2ad
        if (y < 1 \text{ or } r < x) \text{ return } 0;
4db
       if (x \le 1 \text{ and } r \le y) \text{ return } j-i;
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
```

```
0a5
        return count(ei, ej, x, y, 2*p, l, m)+count(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
3cf }
f62 int kth(int i, int j, int k, int p=1, int l = MINN, int r = MAXN) {
        if (1 == r) return 1;
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
585
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
28b
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
8b6 }
f2c int sum(int i, int j, int x, int y, int p = 1, int l = MINN, int r
   = MAXN)
2ad
        if (y < 1 \text{ or } r < x) \text{ return } 0;
2a9
        if (x <= l and r <= y) return pref[p][j]-pref[p][i];</pre>
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x, y,
   2*p+1. m+1. r):
b6d }
b84 int sumk(int i, int j, int k, int p = 1, int l = MINN, int r =
   MAXN) {
8a1
        if (1 == r) return 1*k;
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
50 c
        if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);
4c9
        return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej, k-(ej-ei),
   2*p+1, m+1, r):
940 }
```

- 12 data tructures
- 13 misc

### 14 Extra

#### 14.1 stress.sh

```
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fi
    echo $i
done
```

#### 14.2 hash.sh

```
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp l1 l2
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P -fpreprocessed | tr
    -d '[:space:]' | md5sum | cut -c-6
```

## 14.3 rand.cpp

```
mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution<int> uid(1, r);
    return uid(rng);
}
```

# 14.4 timer.cpp

```
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
    const time_point start;
    timer(): start(now()) {}
    int operator()() {
        return duration_cast < milliseconds > (now() - start).count();
};
14.5 makefile
CXX = g++
CXXFLAGS = -fsanitize = address, undefined -fno-omit-frame-pointer -g
   -Wall -Wshadow -std=c++17 -Wno-unused-result -Wno-sign-compare
   -Wno-char-subscripts #-fuse-ld=gold
14.6 vimrc
189 "" {
d79 set ts=4 sw=4 mouse=a nu ai si undofile
7c9 function H(1)
        return system("sed '/^#/d' | cpp -dD -P -fpreprocessed | tr -d
   '[:space:]' | md5sum", a:1)
Obe endfunction
329 function P() range
dd8
        for i in range(a:firstline, a:lastline)
            let l = getline(i)
ссс
139
            call cursor(i, len(l))
            echo H(getline(search('{}'[1], 'bc', i) ? searchpair('{',
   '', '}', 'bn') : i, i))[0:2] 1
bf9
        endfor
Obe endfunction
90e vmap < C-H > : call P() < CR >
de2 "" }
14.7 pragma.cpp
// Otimizacoes agressivas, pode deixar mais rapido ou mais devagar
#pragma GCC optimize("Ofast")
// Auto explicativo
#pragma GCC optimize("unroll-loops")
```

```
// Vetorizacao
#pragma GCC target("avx2")
// Para operacoes com bits
#pragma GCC target("bmi,bmi2,popcnt,lzcnt")
14.8 debug.cpp
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
    if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
    } while (s.size() and s[0] != ',');
    cerr << " = " << h;
    debug_out(s, line, t...);
}
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__, __VA_ARGS__)
#else
#define debug(...) 42
#endif
14.9 template.cpp
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
typedef long long 11;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
int main() {
    exit(0);
```

# 14.10 fastIO.cpp

```
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
       if (ch == '-') break;
       if (ch >= '0' && ch <= '9') break;
        ch = getchar();
   }
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
   if (minus) return -result;
    else return result;
```