DEEP COMPUTATIONS AND NIP

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ABSTRACT. This paper revisits and extends this bridge between functional analysis and model theory, emphasizing its relevance to the theoretical foundations of machine learning. We show that the compactness behavior of families of Baire class 1 functions mirrors the learnability conditions in the sense of *Probably Approximately Correct* (PAC) learning, and that the failure of compactness corresponds to the presence of infinite *Vapnik-Chervonenkis* (VC) dimension. From this perspective, Rosenthal compacta emerge as the natural topological counterpart of PAC-learnable concept classes, while NIP structures capture the precise boundary between analytical regularity and combinatorial intractability. These parallels suggest a unified framework linking compactness, definability, and learnability, proving how the topology of function spaces encodes the algorithmic and epistemic limits of prediction.

1. GENERAL TOPOLOGICAL PRELIMINARIES

The space of Baire class 1 functions occupies a central position at the intersection of topology, analysis, and logic. The Bourgain-Fremlin-Talagrand theorem characterizes relatively compact subsets of $\mathsf{B}_1(\mathsf{X})$ and, through the perspective of *Rosenthal compacta*, reveals a precise correspondence between topological compactness and model-theoretic tameness. Building on insights of Pierre Simon, this connection identifies the *Non-Independence Property* (NIP) as the exact combinatorial manifestation of the analytic compactness captured by the Bourgain-Fremlin-Talagrand theorem.

In this section we give preliminaries from General Topology and Function Space Theory. We include some of the proofs for completeness but a reader familiar with these topics may skip them.

A *Polish space* is a separable and completely metrizable topological space. The most important examples are the reals $\mathbb{R}$, the Cantor space $2^{\mathbb{N}}$ (the set of all infinite binary sequences, endowed with the product topology), and the Baire space $\mathbb{N}^{\mathbb{N}}$ (the set of all infinite sequences of naturals, also with the product topology). Countable products of Polish spaces are Polish; this includes spaces like $\mathbb{R}^{\mathbb{N}}$, the space of sequences of real numbers. A subspace of a Polish space is itself Polish if and only if it is a G_δ-set, that is, it can be written as the intersection of a countable family of open subsets; in particular, closed subsets and open subsets of Polish spaces are also Polish spaces.

In this work we talk a lot about subspaces, and so there is a pertinent subtlety of the definitions worth mentioning: *completely metrizable space* is not the same as *complete metric space*; for an illustrative example, notice that $(0, 1)$ is homeomorphic to the real line, and thus a Polish space (being Polish is a topological property), but with the metric inherited from the reals, as a subspace, $(0, 1)$ is **not** a complete metric space. In summary, a Polish space has its topology generated by *some* complete metric, but other metrics generating the same topology might not be. In practice, such as when studying descriptive set theory, one finds that we can often keep the metric implicit.

Given two topological spaces X and Y we denote by $B_1(X, Y)$ the set of all functions $f : X \to Y$ such that for all open $U \subseteq Y$, $f^{-1}[U]$ is an F_σ subset of X (that is, a countable union of closed sets); we call these types of functions *Baire class 1 functions*. When $Y = \mathbb{R}$ we simply denote this collection by $B_1(X)$. We endow $B_1(X, Y)$ with the topology of pointwise convergence (the topology inherited from the product topology of Y^X). By $C_p(X, Y)$ we denote the set of all continuous functions $f : X \to Y$ with the topology of pointwise convergence. Similarly, $C_p(X) := C_p(X, \mathbb{R})$. A natural question is, how do topological properties of X translate to $C_p(X)$ and vice versa? These questions, and in general the study of these spaces, are the concern of C_p-theory, a very active field of research in mathematics which was first studied by A. V. Arhangel'skiĭ around 1970's and 1980's. This field has found many exciting applications in model theory (see [HT21]) and functional analysis (see [IC20]). Good recent surveys on the topics include [HT23] and [Tka11]. We begin with the following:

Fact 1.1. *If X is metrizable, then $C_p(X, Y) \subseteq B_1(X, Y)$.*

The proof of the following fact (due to Baire) can be found in Section 10 of [Tod97].

Fact 1.2 (Baire)**.** *If X is a complete metric space, then the following are equivalent:*

(i) *f is a Baire class 1 function, that is, $f \in B_1(X)$.*
(ii) *f is a pointwise limit of continuous functions.*
(iii) *For every closed $F \subseteq X$, the restriction $f|_F$ has a point of continuity.*

Moreover, if X is Polish and $f \notin B_1(X)$, then there exists countable $D_0, D_1 \subseteq X$ and $\alpha < \beta$ such that $\overline{D_0} = \overline{D_1}$, $D_0 \subseteq f^{-1}(-\infty, \alpha]$ and $D_1 \subseteq f^{-1}[\beta, \infty)$.

Recall that a subset $L \subseteq X$ is *relatively compact* in X if the closure of L in X is compact. Relatively compact subsets of $B_1(X)$ (for X Polish space) have been objects of interest to many people working in Analysis and Topological Dynamics. We begin with the following well-known result. Recall that a set $A \subseteq \mathbb{R}^X$ of real-valued functions is *pointwise bounded* if for every $x \in X$ there is $M_x > 0$ such that $|f(x)| < M_x$ for all $f \in A$. We include the proof for the reader's convenience:

Lemma 1.3. *Let X be a Polish space and $A \subseteq B_1(X)$ be pointwise bounded. The following are equivalent:*

(i) *A is relatively compact in $B_1(X)$.*
(ii) *A is relatively countably compact in $B_1(X)$, i.e., every countable subset of A has an accumulation point in $B_1(X)$.*
(iii) *$\overline{A} \subseteq B_1(X)$, where $\overline{A}$ denotes the closure in $\mathbb{R}^X$.*

Proof. By definition, being pointwise bounded means that there is, for each $x \in X$, $M_x > 0$ such that, for every $f \in A$, $|f(x)| \leq M_x$.

(i)⇒(ii) holds in general.

(ii)⇒(iii) Assume that A is relatively countably compact in $B_1(X)$ and that $f \in \overline{A} \setminus B_1(X)$. By Fact 1.2, there are countable $D_0, D_1 \subseteq X$ with $\overline{D_0} = \overline{D_1}$, and $\alpha < \beta$ such that $D_0 \subseteq f^{-1}(-\infty, \alpha]$ and $D_1 \subseteq f^{-1}[\beta, \infty)$. We claim that there is a sequence $\{f_n\}_n \subseteq A$ such that for all $x \in D_0 \cup D_1$, $\lim_{n\to\infty} f_n(x) = f(x)$. Indeed, use the countability to enumerate $D_0 \cup D_1$ as $\{x_n\}_n$. Then find, for each positive n, $f_n \in A$ with $|f_n(x_i) - f(x_i)| < \frac{1}{n}$ for all $i \leq n$. The claim follows.

By relative countable compactness of A, there is an accumulation point $g \in B_1(X)$ of $\{f_n\}_n$. It is straightforward to show that since f and g agree on $D_0 \cup D_1$, g does not have a point of continuity on the closed set $\overline{D_0} = \overline{D_1}$, which contradicts Fact 1.2.

(iii)⇒(i) Suppose that $\overline{A} \subseteq B_1(x)$. Then $\overline{A} \cap B_1(X) = \overline{A}$ is a closed subset of $\prod_{x\in X}[-M_x, M_x]$; Tychonoff's theorem states that the product of compact spaces is always compact, and since closed subsets of compact spaces are compact, $\overline{A}$ must be compact, as desired. □

1.1. **From Rosenthal's dichotomy to NIP.** The fundamental idea that connects the rich theory here presented to real-valued computations is the concept of an *approximation*. In the reals, points of closure from some subset can always be approximated by points inside the set, via a convergent sequence. For more complicated spaces, such as $C_p(X)$, this fails in a remarkably intriguing way. Let us show an example that is actually the protagonist of a celebrated result. Consider the Cantor space $X = 2^{\mathbb{N}}$ and let $p_n(x) = x(n)$ define a continuous mapping $X \to \{0, 1\}$. Then one can show (see Chapter 1.1 of [Tod97] for details) that, perhaps surprisingly, the only continuous functions in the closure of $\{p_n\}_n$ are the functions p_n themselves; moreover, none of the subsequences of $\{p_n\}_n$ converge. In some sense, this example is the worst possible scenario for convergence. The topological space obtained from this closure is well-known. Topologists refer to it as the Čech-Stone compactification of the reals, or $\beta\mathbb{N}$ for short, and it is an important object of study in General Topology.

Theorem 1.4 (Rosenthal's Dichotomy). *If X is Polish and $\{f_n\} \subseteq C_p(X)$ is pointwise bounded, then either $\{f_n\}_n$ contains a convergence subsequence or a subsequence whose closure (in $\mathbb{R}^X$) is homeomorphic to $\beta\mathbb{N}$.*

In words, a pointwise bounded set of continuous functions will either contain a subsequence that converges or a subsequence those closure is essentially the same as the example mentioned in the previous paragraphs (the worst possible scenario). Note that in the preceding example, the functions are trivially pointwise bounded in $\mathbb{R}^X$ as the functions can only take values 0 and 1.

If we intend to generalize our results from $C_p(X)$ to the bigger space $B_1(X)$, we find a similar dichotomy. Either every point of closure of the set of functions will be a Baire class 1 function, or there is a sequence inside the set that behaves in the worst possible way (which in this context, is the IP!). The theorem is usually not phrased as a dichotomy but rather as an equivalence (with the NIP instead):

Theorem 1.5 (Bourgain-Fremlin-Talagrand, Theorem 4G in [BFT78]). *Let X be a Polish space and $A \subseteq C_p(X)$ be pointwise bounded. The following are equivalent:*

(i) A is relatively compact in $B_1(X)$, i.e., $\overline{A} \subseteq B_1(X)$.
(ii) For every $\{f_n\}_n \subseteq A$ and for every $\alpha < \beta$ there is $I \subseteq \mathbb{N}$ such that

$$\bigcap_{n\in I} f_n^{-1}(-\infty, \alpha] \cap \bigcap_{n\notin I} f_n^{-1}[\beta, \infty) = \emptyset.$$

Our goal now is to characterize relatively compact subsets of $B_1(X, Y)$ when $Y = \mathbb{R}^{\mathcal{P}}$ with $\mathcal{P}$ countable. Given $P \in \mathcal{P}$ we denote the *projection map* onto the

P-coordinate by $\pi_P : \mathbb{R}^{\mathcal{P}} \to \mathbb{R}$. From a high-level topological interpretation, the subsequent lemma states that, in this context, the spaces $\mathbb{R}$ and $\mathbb{R}^{\mathcal{P}}$ are really not that different, and that if we understand the Baire class 1 functions of one space, then we also understand the functions of both.

Lemma 1.6. *Let X be a Polish space and $\mathcal{P}$ be a countable set. Then, $f \in B_1(X, \mathbb{R}^{\mathcal{P}})$ if and only if $\pi_P \circ f \in B_1(X)$ for all $P \in \mathcal{P}$.*

Proof. Only one implication needs a proof. Suppose that $\pi_P \circ f \in B_1(X)$ for all $P \in \mathcal{P}$. Let V be a basic open subset of $\mathbb{R}^{\mathcal{P}}$. That is, there exists a finite $\mathcal{P}' \subseteq \mathcal{P}$ such that $V = \bigcap_{P \in \mathcal{P}'} \pi_P^{-1}[U_P]$ where U_P is open in $\mathbb{R}$. Finally,

$$f^{-1}[V] = \bigcap_{P \in \mathcal{P}'} (\pi_P \circ f)^{-1}[U_P]$$

is an F_σ set. Since $\mathcal{P}$ is countable, then $\mathbb{R}^{\mathcal{P}}$ is second countable so every open set U in $\mathbb{R}^{\mathcal{P}}$ is a countable union of basic open sets. Hence, $f^{-1}[U]$ is F_σ. □

Below we consider $\mathcal{P}$ with the discrete topology. For each $f : X \to \mathbb{R}^{\mathcal{P}}$ denote $\hat{f}(P, x) := \pi_P \circ f(x)$ for all $(P, x) \in \mathcal{P} \times X$. Similarly, for each $g : \mathcal{P} \times X \to \mathbb{R}$ denote $\check{g}(x)(P) := g(P, x)$. Given $A \subseteq (\mathbb{R}^{\mathcal{P}})^X$, we denote $\hat{A}$ as the set of all $\hat{f}$ such that $f \in A$.

The map $\left(\mathbb{R}^{\mathcal{P}}\right)^X \to \mathbb{R}^{\mathcal{P} \times X}$ given by $f \mapsto \hat{f}$ is a homeomorphism and its inverse is given by $g \mapsto \check{g}$.

Lemma 1.7. *Let X be a Polish space and $\mathcal{P}$ be countable. Then, $f \in B_1(X, \mathbb{R}^{\mathcal{P}})$ if and only if $\hat{f} \in B_1(\mathcal{P} \times X)$.*

Proof. ($\Rightarrow$) Given an open set of reals U, we have that for every $P \in \mathcal{P}$, $f^{-1}[\pi_P^{-1}[U]]$ is F_σ by lemma 1.6. Given that $\mathcal{P}$ is a discrete countable space, we observe that

$$\hat{f}^{-1}[U] = \bigcup_{P \in \mathcal{P}} \left(\{P\} \times f^{-1}[\pi_P^{-1}[U]]\right)$$

is also an F_σ set. ($\Leftarrow$) By lemma 1.6 it suffices to show that $\pi_P \circ f \in B_1(X)$ for all $P \in \mathcal{P}$. Fix an open $U \subseteq \mathbb{R}$. Write $\hat{f}^{-1}[U] = \bigcup_{n \in \mathbb{N}} F_n$ where F_n is closed in $\mathcal{P} \times X$. Then,

$$(\pi_P \circ f)^{-1}[U] = \bigcup_{n \in \mathbb{N}} \{x \in X : (P, x) \in F_n\}$$

which is F_σ. □

1.2. **The Non-Independence Property.** One of the most important innovations in Machine Learning is the mathematical notion, introduced by Turing Awardee Leslie Valiant in the 1980s, of 'probably approximately correct learning', or PAC-learning for short [BD19]. We give a standard but short overview of these concepts in the context that is relevant to this work.

Consider the following important idea in data classification. Suppose that A is a set and that $\mathcal{C}$ is a collection of sets. We say that $\mathcal{C}$ *shatters* A if every subset of A is of the form $C \cap A$ for some $C \in \mathcal{C}$. For a classical geometric example, if A is the set of four points on the Euclidean plane of the form $(\pm 1, \pm 1)$, then the collection of all half-planes does not shatter A, the collection of all open balls does not shatter A, but the collection of all convex sets shatters A. While A need not be finite, it will usually be assumed to be so in Machine Learning applications. A finer way to distinguish collections of sets that shatter a given set from those that do not is by the *Vapnik-Chervonenkis dimension (VC-dimension)*, which is equal to the cardinality of the largest finite set shattered by the collection, in case it exists, or to infinity otherwise.

A concrete illustration of these ideas appears when considering threshold classifiers on the real line. Let $\mathcal{H} = \{h_t(x) = \mathbf{1}_{x \le t} : t \in \mathbb{R}\}$ be the class of indicator functions representing all possible threshold decisions. Each h_t is a Baire class 1 function, and the family $\mathcal{H}$ is relatively compact in $B_1(\mathbb{R})$. In model-theoretic terms, $\mathcal{H}$ is NIP, since no configuration of points and thresholds can realize the full independence pattern of a binary matrix. By contrast, the family of parity functions $\{x \mapsto (-1)^{\langle w, x \rangle} : w \in \{0,1\}^n\}$ on $\{0,1\}^n$ has the independence property and fails relative compactness in $B_1(X)$, capturing the analytical meaning of instability. This dichotomy mirrors the behavior of concept classes with finite versus infinite VC dimension in statistical learning theory.

Fix a first-order formula $\varphi(x, y)$ in a language L and a model M of an L-theory T, and let

$$\mathcal{F}_\varphi(M) := \{\varphi(M, a) : a \in M^{|y|}\}$$

be the family of subsets of $M^{|x|}$ defined by instances of the formula φ, where $M^{|x|}$ denoted the usual Cartesian power. We say that $\varphi(x, y)$ has the *independence property (IP)* in M if there is a sequence $(c_i)_{i \in \mathbb{N}} \subseteq M^{|x|}$ such that for every $S \subseteq \mathbb{N}$ there is $a_S \in M^{|y|}$ with

$$M \models \varphi(c_i, a_S) \quad \iff \quad \forall i \in \mathbb{N},\ i \in S.$$

The formula has the IP if it does so in some model, and the formula has the *non-independence property (NIP)* if it does not have the IP.

For two simple examples of formulas satisfying the NIP, consider first the language $L = \{<\}$ and the model $M = (\mathbb{R}, <)$ of the reals with their usual linear order. Take the formula $\varphi(x, y)$ to mean $x < y$, then $\varphi(M, a) = (-\infty, a)$, and so $\mathcal{F}_\varphi(M)$ is just the set of left open rays. The VC-dimension of this collection is 1, since it can shatter a single point, but no two point set can be shattered since the rays are downwards closed. Now in contrast, the collection of open intervals, given by the formula $\varphi(x; y_1, y_2) := (y_1 < x) \wedge (x < y_2)$, has VC-dimension 2.

The fundamental theorem of statistical learning states that a binary hypothesis class is PAC-learnable if and only if it has finite VC-dimension, and the subsequent theorem connects the rest of the concepts presented in this section.

Theorem 1.8 (Laskowski). *The formula $\varphi(x; y)$ has the NIP if and only if $\mathcal{F}_\varphi(M)$ has finite VC-dimension.*

We now direct our attention to a more general notion of the NIP. It can be interpreted as a sort of continuous version of the previously discrete one presented in the preceding paragraphs.

Definition 1.9. We say that $A \subseteq \mathbb{R}^X$ has the *Non-Independence Property* (NIP) if and only if for every $\{f_n\}_n \subseteq A$ and for every $\alpha < \beta$ there are finite disjoint sets $E, F \subseteq \mathbb{N}$ such that

$$\bigcap_{n \in E} f_n^{-1}(-\infty, \alpha] \cap \bigcap_{n \in F} f_n^{-1}[\beta, \infty) = \emptyset.$$

Note that if X is compact and $A \subseteq C_p(X)$, then A has the NIP if and only if for every $\{f_n\}_n \subseteq A$ and for every $\alpha < \beta$ there is $I \subseteq \mathbb{N}$ such that

$$\bigcap_{n\in I} f_n^{-1}(-\infty,\alpha] \cap \bigcap_{n\notin I} f_n^{-1}[\beta,\infty) = \emptyset.$$

Given $A \subseteq Y^X$ and $K \subseteq X$ we write $A|_K := \{f|_K : f \in A\}$, i.e., the set of all restrictions of functions in A to K. The following Theorem is a slightly more general version of Theorem 2.7.

Theorem 1.10. *Assume that $\mathcal{P}$ is countable, X is a Polish space, and $A \subseteq C_p(X,\mathbb{R}^{\mathcal{P}})$ be such that $\pi_P \circ A$ is pointwise bounded for all $P \in \mathcal{P}$. The following are equivalent for every compact $K \subseteq X$:*

(1) $\overline{A|_K} \subseteq B_1(K,\mathbb{R}^{\mathcal{P}})$.

(2) $\pi_P \circ A|_K$ *has the NIP for every* $P \in \mathcal{P}$.

Proof. (1)⇒(2). Let $P \in \mathcal{P}$. Fix $\{f_n\}_n \subseteq A$ and $\alpha < \beta$. By (1) we have that $\overline{A|_K} \subseteq B_1(K,\mathbb{R}^{\mathcal{P}})$. Applying the homeomorphism $f \mapsto \hat{f}$ and using lemma 1.7 we get $\overline{\hat{A}|_{\mathcal{P}\times K}} \subseteq B_1(\mathcal{P}\times K)$. By Theorem 2.7, there is $I \subseteq \mathbb{N}$ such that

$$(\mathcal{P}\times K) \cap \bigcap_{n\in I} \hat{f_n}^{-1}(-\infty,\alpha] \cap \bigcap_{n\notin I} \hat{f_n}^{-1}[\beta,\infty) = \emptyset$$

Hence,

$$K \cap \bigcap_{n\in I} (\pi_P \circ f_n)^{-1}(-\infty,\alpha] \cap \bigcap_{n\notin I} (\pi_P \circ f_n)^{-1}[\beta,\infty) = \emptyset$$

By compactness, there are finite $E \subseteq I$ and $F \subseteq \mathbb{N}\setminus I$ such that

$$K \cap \bigcap_{n\in E} (\pi_P \circ f_n)^{-1}(-\infty,\alpha] \cap \bigcap_{n\in F} (\pi_P \circ f_n)^{-1}[\beta,\infty) = \emptyset$$

Thus, $\pi_P \circ A|_L$ has the NIP.

(2)⇒(1) Fix $f \in \overline{A|_K}$. By lemma 1.6 it suffices to show that $\pi_P \circ f \in B_1(K)$ for all $P \in \mathcal{P}$. By (2), $\pi_P \circ A|_K$ has the NIP. Hence, by Theorem 2.7 we have $\overline{\pi_P \circ A|_K} \subseteq B_1(K)$. But then $\pi_P \circ f \in \overline{\pi_P \circ A|_K} \subseteq B_1(K)$. □

Lastly, a simple but significant result that helps understand the operation of restricting a set of functions to a specific subspace of the domain space X, of course in the context of the NIP, is that we may always assume that said subspace is closed. Concretely, whether we take its closure or not has no effect on the NIP:

Lemma 1.11. *Assume that X is Hausdorff and that $A \subseteq C_p(X)$. The following are equivalent for every $L \subseteq X$:*

(i) A_L *has the NIP.*

(ii) $A|_{\overline{L}}$ *has the NIP.*

Proof. It suffices to show that (i)⇒(ii). Suppose that (ii) does not hold, i.e., that there are $\{f_n\}_n \subseteq A$ and $\alpha < \beta$ such that for all finite disjoint $E, F \subseteq \mathbb{N}$:

$$\overline{L} \cap \bigcap_{n\in E} f_n^{-1}(-\infty,\alpha] \cap \bigcap_{n\in F} f_n^{-1}[\beta,\infty) \neq \emptyset.$$

Pick $\alpha' < \beta'$ such that $\alpha < \alpha' < \beta' < \beta$. Then, for any finite disjoint $E, F \subseteq \mathbb{N}$ we can choose

$$x \in \overline{L} \cap \bigcap_{n\in E} f_n^{-1}(-\infty,\alpha') \cap \bigcap_{n\in F} f_n^{-1}(\beta',\infty)$$

By definition of closure:

$$L \cap \bigcap_{n\in E} f_n^{-1}(-\infty, \alpha'] \cap \bigcap_{n\in F} f_n^{-1}[\beta', \infty) \neq \emptyset.$$

This contradicts (i). □

2. NIP in the context of Compositional Computation Structures

In this section, we study what the NIP tell us in the context of deep computations as defined in [ADIW24]. We say a structure $(L, \mathcal{P}, \Gamma)$ is a *Compositional Computation Structure* (CCS) if $L \subseteq \mathbb{R}^{\mathcal{P}}$ is a subspace of $\mathbb{R}^{\mathcal{P}}$, with the pointwise convergence topology, and $\Gamma \subseteq L^L$ is a semigroup under composition. The motivation of CCS comes from (continuous) Model Theory, where $\mathcal{P}$ is a fixed collection of predicates and L is a (real-valued) structure. Every point in L is identified with its "type", which is the tuple of all values the point takes on the predicates from $\mathcal{P}$, i.e., an element of $\mathbb{R}^{\mathcal{P}}$. In this context, elements of $\mathcal{P}$ are called *features.* In the discrete Model Theory framework, one view the space of complete-types as a compactification of the structure L. In this context, we don't want to consider only points in L (realized types) but in its closure $\overline{L}$ (possibly unrealized types). The problem is that the closure $\overline{L}$ is not necessarily compact, an assumption that turns out very useful in the context of continuous Model Theory. To bypass this problem in the framework for deep computations, Alva, Dueñez, Iovino and Walton introduced in [ADIW24] the concept of shards, which essentially consists in covering (a large fragment) of the space $\overline{L}$ by compact pointwise-bounded subspaces (shards). We shall give the formal definition next.

A *sizer* is a tuple $r_\bullet = (r_P)_{P\in\mathcal{P}}$ of positive real numbers indexed by $\mathcal{P}$. Given a sizer $r_\bullet$, we define the $r_\bullet$-*shard* as:

$$L[r_\bullet] = L \cap \prod_{P\in\mathcal{P}} [-r_P, r_P]$$

For an illustrative example, we can frame Newton's polynomial root approximation method in the context of a CCS (see Example 5.6 of [ADIW24] for details) as follows. Begin by considering the extended complex numbers $\hat{\mathbb{C}} := \mathbb{C} \cup \{\infty\}$ with the usual Riemann sphere topology that makes it into a compact space (where unbounded sequences converge to ∞). In fact, not only is this space compact but is covered by the shard given by the sizer $(1, 1, 1)$ (the unit sphere is contained in the cube $[-1, 1]^3$). The space $\hat{\mathbb{C}}$ is homeomorphic to the usual unit sphere $S^2 := \{(x, y, z) : x^2 + y^2 + z^2 = 1\}$ of $\mathbb{R}^3$, by means of the stereographic projection and its inverse $\hat{\mathbb{C}} \to S^2$. This function is regarded as a triple of predicates $x, y, z : \hat{\mathbb{C}} \to [-1, 1]$ where each will map an extended complex number to its corresponding real coordinate on the unit sphere. Now fix the cubic complex polynomial $p(s) := s^3 - 1$, and consider the map which performs one step in Newton's method at a particular (extended) complex number s, for finding a root of p, $\gamma_p : \hat{\mathbb{C}} \to \hat{\mathbb{C}}$. The explicit inner workings of γ_p are irrelevant for this example, except for the fact that it is a continuous mapping. It follows that $(S^3, \{x, y, z\}, \{\gamma_p^k : k \in \mathbb{N}\})$ is a CCS. The idea is that repeated applications of $\gamma_p(s), \gamma_p \circ \gamma_p(s), \gamma_p \circ \gamma_p \circ \gamma_p(s), \ldots$ would approximate a root of p provided s was a good enough initial guess.

The $r_\bullet$-type-shard is defined as $\mathcal{L}[r_\bullet] = \overline{L[r_\bullet]}$ and $\mathcal{L}_{sh}$ is the union of all type-shards. Notice that $\mathcal{L}_{sh}$ is not necessarily equal to $\mathcal{L} = \overline{L}$, unless $\mathcal{P}$ is countable (see [ADIW24]). A *transition* is a map $f : L \to L$, in particular, every element in the semigroup Γ is a transition (these are called realized computations). In practice, one would like to work with "definable" computations, i.e., ones that can be described by a computer. In this topological framework, being continuous is an expected requirement. However, as in the case of complete-types in Model Theory, we will work with "unrealized computations", i.e., maps $f : \mathcal{L}_{sh} \to \mathcal{L}_{sh}$. Note that continuity of a computation does not imply that it can be continuously extended to $\mathcal{L}_{sh}$. The Extendibility Axiom (introduced in [ADIW24]) is a reasonable assumption made to work with a nice space of computations:

We say that the CCS $(L, \mathcal{P}, \Gamma)$ satisfies the *Extendibility Axiom* if for all $\gamma \in \Gamma$, there is $\tilde{\gamma} : \mathcal{L}_{sh} \to \mathcal{L}_{sh}$ such that for every sizer $r_\bullet$ there is an $s_\bullet$ such that $\tilde{\gamma}|_{\mathcal{L}[r_\bullet]} : \mathcal{L}[r_\bullet] \to \mathcal{L}[s_\bullet]$ is continuous.

A collection R of sizers is called *exhaustive* if $\mathcal{L}_{sh} = \bigcup_{r_\bullet \in R} \mathcal{L}[r_\bullet]$. We say that $\Delta \subseteq \Gamma$ is R-confined if $\gamma|_{L[r_\bullet]} : L[r_\bullet] \to L[r_\bullet]$ for every $r_\bullet \in R$ and $\gamma \in \Delta$. Elements in Δ are called *real-valued computations* (in this article we will refer to them simply as *computations*) and elements in $\overline{\Delta} \subseteq \mathcal{L}_{sh}^L$ are called (real-valued) *deep computations* or *ultracomputations*. By $\tilde{\Delta}$ we denote the set of all extensions $\tilde{\gamma}$ for $\gamma \in \Delta$. For a more complete description of this framework, we refer the reader to [ADIW24].

2.1. **NIP and Baire-1 definability of deep computations.** Under what conditions are deep computations Baire class 1, and thus well-behaved according to our framework, on type-shards? The next Theorem says that, again under the assumption that $\mathcal{P}$ is countable, the space of deep computations is a Rosenthal compactum (on shards) if and only if the set of computations has the NIP on features. Hence, we can import the theory of Rosenthal compacta into this framework of deep computations.

Theorem 2.1. *Let $(L, \mathcal{P}, \Gamma)$ be a CCS satisfying the Extendibility Axiom with $\mathcal{P}$ countable. Let R be an exhaustive collection of sizers. Let $\Delta \subseteq \Gamma$ be R-confined. The following are equivalent.*

(1) $\overline{\tilde{\Delta}|_{\mathcal{L}[r_\bullet]}} \subseteq B_1(\mathcal{L}[r_\bullet], \mathcal{L}[r_\bullet])$ for all $r_\bullet \in R$.
(2) $\pi_P \circ \Delta|_{L[r_\bullet]}$ has the NIP for all $P \in \mathcal{P}$ and $r_\bullet \in R$, that is, for all $P \in \mathcal{P}$, $r_\bullet \in R$, $\alpha < \beta$, $\{\gamma_n\}_n \subseteq \Delta$ there are finite disjoint $E, F \subseteq \mathbb{N}$ such that

$$L[r_\bullet] \cap \bigcap_{n \in E} (\pi_P \circ \gamma_n)^{-1}(-\infty, \alpha] \cap \bigcap_{n \in F} (\pi_P \circ \gamma_n)^{-1}[\beta, \infty) = \emptyset.$$

Moreover, if any (hence all) of the preceding conditions hold, then every deep computation $f \in \overline{\Delta}$ can be extended to a Baire-1 function on shards, i.e., there is $\tilde{f} : \mathcal{L}_{sh} \to \mathcal{L}_{sh}$ such that $\tilde{f}|_{\mathcal{L}[r_\bullet]} \in B_1(\mathcal{L}[r_\bullet], \mathcal{L}[r_\bullet])$ for all $r_\bullet \in R$. In particular, every deep computation is the pointwise limit of a countable sequence of computations on every shard.

Proof. Since $\mathcal{P}$ is countable, then $\mathcal{L}[r_\bullet] \subseteq \mathbb{R}^{\mathcal{P}}$ is Polish. Also, the Extendibility Axiom implies that $\pi_P \circ \tilde{\Delta}|_{\mathcal{L}[r_\bullet]}$ is a pointwise bounded set of continuous functions for all $P \in \mathcal{P}$. Hence, Theorem 1.10 and Lemma 1.11 prove the equivalence of (1) and (2). If (1) holds and $f \in \overline{\Delta}$, then write $f = \mathcal{U}\lim_i \gamma_i$ as an ultra-limit. Define $\tilde{f} := \mathcal{U}\lim_i \tilde{\gamma}_i$. Hence, for all $r_\bullet \in R$ we have $\tilde{f}|_{\mathcal{L}[r_\bullet]} \in \overline{\tilde{\Delta}|_{\mathcal{L}[r_\bullet]}} \subseteq B_1(\mathcal{L}[r_\bullet], \mathcal{L}[r_\bullet])$. That every deep computation is a pointwise limit of a countable sequence of computations follows from the fact that every compact subset of $B_1(X)$ is Frechet-Urysohn (see Theorem 3F in [BFT78] or Theorem 4.1 in [Deb13]). $\square$

Given a countable set Δ of computations satisfying the NIP on features and shards (condition (2) of Theorem 2.1) we have that $\overline{\tilde{\Delta}_{\mathcal{L}[r_\bullet]}}$ (for a fixed sizer $r_\bullet$) is a separable Rosenthal Compactum (compact subset of $B_1(P \times \mathcal{L}[r_\bullet])$. The work of Todorčević ([Tod99]) and Argyros, Dodos, Kanellopoulos ([ADK08]) culminates in a trichotomy Theorem for separable Rosenthal Compacta. Inspired by the work of Glasner and Megrelishvili ([GM22]), we are interested to see how this allows us to classify and obtain different levels of PAC-learnability (NIP).

Recall that a topological space X is *hereditarily separable* (HS) if every subspace is separable and that X is *first countable* if every point in X has a countable local basis. It is a result of R. Pol that every metrizable separable Rosenthal compactum is hereditarily separable and that every hereditarily separable Rosenthal compactum is first countable (see section 10 in [Deb13]). This suggests the following definition:

Definition 2.2. Let $(L, \mathcal{P}, \Gamma)$ be a CCS satisfying the Extendibility Axiom and R be an exhaustive collection of sizers. Let $\Delta \subseteq \Gamma$ be an R-confined countable set of computations satisfying the NIP on shards and features (condition (2) in Theorem 2.1). We say that Δ is:

(i) NIP_1 if $\overline{\tilde{\Delta}|_{\mathcal{L}[r_\bullet]}}$ is first countable for every $r_\bullet \in R$.
(ii) NIP_2 if $\overline{\tilde{\Delta}|_{\mathcal{L}[r_\bullet]}}$ is hereditarily separable for every $r_\bullet \in R$.
(iii) NIP_3 if $\overline{\tilde{\Delta}|_{\mathcal{L}[r_\bullet]}}$ is metrizable for every $r_\bullet \in R$.

Observe that $\mathsf{NIP}_3 \Rightarrow \mathsf{NIP}_2 \Rightarrow \mathsf{NIP}_1 \Rightarrow \mathsf{NIP}$. A natural question that would continue this work is to find examples of CCS that separate these levels of NIP. In [Tod99], Todorčević isolates 3 canonical examples of Rosenthal compacta that witness the failure of the converse implications above.

Theorem 2.3 (Todorčević's Trichotomy, [Tod99], Theorem 3 in [ADK08])**.** *Let K be a separable Rosenthal Compactum.*

(i) If K is hereditarily separable but non-metrizable, then $S(2^{\mathbb{N}})$ embeds into K.
(ii) If K is first countable but not hereditarily separable, then either $D(2^{\mathbb{N}})$ or $D(S(2^{\mathbb{N}}))$ embeds into K.
(iii) If K is not first countable, then $A(2^{\mathbb{N}})$ embeds into K.

In other words, we have the following classification:

$$\begin{tikzcd}[column sep=tiny]
& \text{K is separable Rosenthal compactum} \arrow[dl] \arrow[dr] & & & \\
\text{K is metrizable} & & \text{K is not metrizable} \arrow[dl] \arrow[dr] & & \\
& \begin{array}{c}\text{K is HS}\\ \text{(copy of $S(I)$)}\end{array} & & \text{K is not HS} \arrow[dl] \arrow[dr] & \\
& & \begin{array}{c}\text{K is first countable}\\ \text{(copy of $D(2^{\mathbb{N}})$ or $D(S(I))$)}\end{array} & & \begin{array}{c}\text{K is not first countable}\\ \text{(copy of $A(2^{\mathbb{N}})$)}\end{array}
\end{tikzcd}$$

Lastly, the definitions provided here for NIP_i $(i = 1, 2, 3)$ are topological.

Question 2.4. Is there a non-topological characterization for NIP_i, $i = 1, 2, 3$?

2.2. **NIP and definability by universally measurable functions.** We now turn to the question: what happens when $\mathcal{P}$ is uncountable? Notice that the countability assumption is crucial in the proof of Theorem 1.10 essentially because it makes $\mathbb{R}^{\mathcal{P}}$ a Polish space. For the uncountable case, we may lose Baire-1 definability so we shall replace $B_1(X)$ by a bigger class. Recall that the purpose of studying the class of Baire-1 functions is that a pointwise limit of continuous functions is not necessarily continuous. In [BFT78], J. Bourgain, D.H. Fremlin and M. Talagrand characterized the Non-Independence Property of a set of continuous functions with various notions of compactness in function spaces containing $C(X)$, such as $B_1(X)$. In this section we will replace $B_1(X)$ with the larger space $M_r(X)$ of universally measurable functions. The development of this section is based on Theorem 2F in [BFT78]. We now give the relevant definitions. Readers with little familiarity with measure theory can review the appendix for standard definitions appearing in this subsection.

Given a Hausdorff space X and a measurable space (Y, Σ), we say that $f : X \to Y$ is *universally measurable* (with respect to Σ) if $f^{-1}(E)$ is universally measurable for every $E \in \Sigma$, i.e., $f^{-1}(E)$ is μ-measurable for every Radon Probability measure μ on X. When $Y = \mathbb{R}$ we will always take $\Sigma = \mathcal{B}(\mathbb{R})$, the Borel σ-algebra of $\mathbb{R}$. In that case, a function $f : X \to \mathbb{R}$ is universally measurable if and only if $f^{-1}(U)$ is μ-measurable for every Radon Probability measure μ on X and every open set $U \subseteq \mathbb{R}$. Following [BFT78], the collection of all universally measurable real-valued functions will be denoted by $M_r(X)$. In the context of deep computations, we will be interested in transition maps from a state space $L \subseteq \mathbb{R}^{\mathcal{P}}$ to itself. There are two natural σ-algebras one can consider in the product space $\mathbb{R}^{\mathcal{P}}$: the Borel σ-algebra, i.e., the σ-algebra generated by open sets in $\mathbb{R}^{\mathcal{P}}$; and the cylinder σ-algebra, i.e., the σ-algebra generated by Borel cylinder sets or equivalently basic open sets in $\mathbb{R}^{\mathcal{P}}$. Note that when $\mathcal{P}$ is countable, both σ-algebras coincide but in general the cylinder σ-algebra is strictly smaller. We will use the cylinder σ-algebra to define universally measurable maps $f : \mathbb{R}^{\mathcal{P}} \to \mathbb{R}^{\mathcal{P}}$. The reason for this choice is because of the following characterization:

Lemma 2.5. *Let X be a Hausdorff space and $Y = \prod_{i\in I} Y_i$ be any product of measurable spaces (Y_i, Σ_i) for $i \in I$. Let Σ_Y be the cylinder σ-algebra generated by the measurable spaces (Y_i, Σ_i). Let $f : X \to Y$. The following are equivalent:*

(i) $f : X \to Y$ *is universally measurable (with respect to Σ_Y).*
(ii) $\pi_i \circ f : X \to Y_i$ *is universally measurable (with respect to Σ_i) for all $i \in I$.*

Proof. (i)⇒(ii) is clear since the projection maps π_i are measurable and the composition of measurable functions is measurable. To prove (ii)⇒(i), suppose that $C = \prod_{i\in I} C_i$ is a measurable cylinder and let J be the finite set of $i \in I$ such that $C_i \neq Y_i$. Then, $C = \bigcap_{i\in J} \pi_i^{-1}(C_i)$ so $f^{-1}(C) = \bigcap_{i\in J}(\pi_i \circ f)^{-1}(C_i)$ is a universally measurable set by assumption. □

The previous lemma says that a transition map is universally measurable if and only if it is universally measurable on all its features. In other words, we can check measurability of a transition just by checking measurability in all its features. We will denote by $M_r(X, \mathbb{R}^{\mathcal{P}})$ the collection of all universally measurable functions $f : X \to \mathbb{R}^{\mathcal{P}}$ (with respect to the cylinder σ-algebra), endowed with the topology of pointwise convergence.

Definition 2.6. Let $(\mathsf{L}, \mathcal{P}, \Gamma)$ be a CCS. We say that a transition $f : \mathsf{L} \to \mathsf{L}$ is **universally measurable shard-definable** if and only if there exists $\tilde{f} : \mathcal{L}_{sh} \to \mathcal{L}_{sh}$ extending f such that for every sizer $r_\bullet$ there is a sizer $s_\bullet$ such that the restriction $\tilde{f}|_{\mathcal{L}[r_\bullet]} : \mathcal{L}[r_\bullet] \to \mathcal{L}[s_\bullet]$ is universally measurable, i.e. $\pi_P \circ \tilde{f}|_{\mathcal{L}[r_\bullet]} : \mathcal{L}[r_\bullet] \to [-s_P, s_P]$ is μ-measurable for every Radon probability measure μ on $\mathcal{L}[r_\bullet]$.

We will need the following result about NIP and universally measurable functions:

Theorem 2.7 (Bourgain-Fremlin-Talagrand, Theorem 2F in [BFT78])**.** *Let X be a Hausdorff space and $A \subseteq C(X)$ be pointwise bounded. The following are equivalent:*

(i) $\overline{A} \subseteq M_r(X)$.
(ii) For every compact $K \subseteq X$, $A|_K$ has the NIP.

Pierre Simon made a major contribution to the understanding of the Bourgain-Fremlin-Talagrand theorem within model theory by showing that its analytic content coincides with the model-theoretic notion of the *Non-Independence Property* (NIP). In his work on NIP theories (see [**?**]), Simon demonstrated that the same combinatorial configurations forbidden by clause (ii) of Theorem 2.7 are exactly those that correspond to the independence property in logic. The classical BFT theorem thus provides a topological-functional characterization of model-theoretic tameness.

Simon's insight was to view definable families of functions as sets of real-valued functions on type spaces and to interpret relative compactness in $B_1(X)$ as a form of "tame behavior" under ultrafilter limits. From this perspective, NIP theories are those whose definable families behave like relatively compact sets of Baire class 1 functions, avoiding the wild, $\beta\mathbb{N}$-like configurations that characterize instability. This observation opened a new bridge between analysis and logic: topological compactness corresponds to the absence of combinatorial independence.

Simon's later developments connected these ideas to Keisler measures and empirical averages, allowing tools from functional analysis to study learnability and definable types. This reinterpretation of model-theoretic tameness through the lens of the BFT theorem has made NIP a central notion not only in stability theory but also in contemporary connections with learning theory and ergodic analysis.

Historically, the notion of NIP arises from Shelah's foundational work on the classification theory of models. In his seminal book *Unstable Theories* [She78], Shelah introduced the independence property as a key dividing line within unstable structures, identifying the class of stable theories as those in which this property fails. The later notion of NIP generalizes stability by forbidding the full combinatorial independence pattern while allowing certain controlled forms of instability. Thus, Simon's interpretation of the BFT theorem can be viewed as placing Shelah's dividing line into a topological-analytic framework, connecting the earliest notions of stability to compactness phenomena in spaces of Baire class 1 functions.

Theorem 2.7 immediately yields the following.

Theorem 2.8. *Let $(\mathsf{L}, \mathcal{P}, \Gamma)$ be a CCS satisfying the Extendibility Axiom. Let R be an exhaustive collection of sizers. Let $\Delta \subseteq \Gamma$ be R-confined. If $\pi_P \circ \Delta|_{\mathsf{L}[r_\bullet]}$ has the NIP for all $P \in \mathcal{P}$ and all $r_\bullet \in R$, then every deep computation is universally measurable shard-definable.*

Proof. By the Extendibility Axiom, Theorem 2.7 and lemma 1.11 we have that $\pi_P \circ \overline{\tilde{\Delta}|_{\mathcal{L}[r_\bullet]}} \subseteq M_r(\mathcal{L}[r_\bullet])$ for all $r_\bullet \in R$ and $P \in \mathcal{P}$. Let $f \in \overline{\Delta}$ be a deep computation. Write $f = \mathcal{U}\lim_i \gamma_i$ as an ultralimit of computations in Δ. Define $\tilde{f} := \mathcal{U}\lim_i \tilde{\gamma}_i$. Then, for all $r_\bullet \in R$ and $P \in \mathcal{P}$ $\pi_P \circ \tilde{\gamma}_i|_{\mathcal{L}[r_\bullet]} \in M_r(\mathcal{L}[r_\bullet])$ for all i so $\pi_P \circ f|_{\mathcal{L}[r_\bullet]} \in \overline{\pi_P \circ \tilde{\Delta}|_{\mathcal{L}[r_\bullet]}} \subseteq M_r(\mathcal{L}[r_\bullet])$. □

Question 2.9. Under the same assumptions of the previous Theorem, suppose that every deep computation of Δ is universally measurable shard-definable. Must $\pi_P \circ \Delta|_{L[r_\bullet]}$ have the NIP for all $P \in \mathcal{P}$ and all $r_\bullet \in R$?

2.3. **Talagrand stability and definability by universally measurable functions.** There is another notion closely related to NIP, introduced by Talagrand while studying Pettis integration. Suppose that X is a compact Hausdorff space and $A \subseteq \mathbb{R}^X$. Let μ be a Radon measure on X. Given a μ-measurable set $E \subseteq X$, a positive integer k and real numbers $\alpha < \beta$. we write:

$$D_k(A, E, \alpha, \beta) = \bigcup_{f \in A} \{x \in E^{2k} : f(x_{2i}) \leq \alpha,\ f(x_{2i+1}) \geq \beta \ \text{ for all } i < k\}$$

We say that A is Talagrand μ-stable if and only if for every μ-measurable set $E \subseteq X$ of positive measure and for every $\alpha < \beta$ there is $k \geq 1$ such that $(\mu^{2k})^*(D_k(A, E, \alpha, \beta)) < (\mu(E))^{2k}$.

Lemma 2.10. *If A is Talagrand μ-stable, then $\overline{A}$ is also Talagrand μ-stable and $\overline{A} \subseteq \mathcal{L}^0(X, \mu)$.*

We say that A is universally Talagrand stable if A is Talagrand μ-stable for every Radon measure μ on X. A similar argument as before, yields the following:

Theorem 2.11. *Let $(L, \mathcal{P}, \Gamma)$ be a CCS satisfying the Extendibility Axiom. If $\pi_P \circ \Delta|_{L[r_\bullet]}$ is universally Talagrand stable for all $P \in \mathcal{P}$ and all sizers $r_\bullet$, then every deep computation is universally measurable sh-definable.*

It is then natural to ask, what is the relationship between Talagrand stability and the NIP? We know that Fremlin' Dichotomy imply:

Lemma 2.12. *Let X be a compact Hausdorff space and $A \subseteq C(X)$ be pointwise bounded. If A is universally Talagrand stable, then A has the NIP.*

Question 2.13. Is the converse true?

There is a delicate point in this question, as it may be sensitive to set-theoretic axioms (even assuming countability of A).

Theorem 2.14 (Talagrand)**.** *Let X be a compact Hausdorff space and $A \subseteq M_r(X)$ be countable and pointwise bounded. Assume that $[0,1]$ is not the union of $< \mathfrak{c}$ closed measure zero sets. If A has the NIP, then A is universally Talagrand stable.*

Theorem 2.15 (Fremlin, Shelah)**.** *It is consistent that there exists a countable pointwise bounded set of Lebesgue measurable functions with the NIP which is not Talagrand stable with respect to Lebesgue measure.*

Appendix: Measure Theory

Given a set X, a collection Σ of subsets of X is called a **σ-algebra** if Σ contains X and is closed under complements and countable unions. Hence, for example, a σ-algebra is also closed under countable intersections. Intuitively, a σ-algebra is a collection of sets in which we can define a σ-additive measure. We call sets in a σ-algebra Σ **measurable sets** and the pair (X, Σ) a measurable space. If X is a topological space, there is a natural σ-algebra of subsets of X, namely the **Borel σ-algebra** $\mathcal{B}(X)$, i.e., the smallest σ-algebra containing all open subsets of X. Given a measurable space (X, Σ), a σ-additive measure is a non-negative function $\mu : \Sigma \to \mathbb{R}$

(Define Topological measurable spaces, example of topological spaces with Borel sigma-algebra, Radon probability measures on Hausdorff spaces, universally measurable sets, products of topological measurable spaces and the cylinder sigma-algebra, universally measurable functions when the co-domain is a topological measurable space)

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