

## DEEP COMPUTATIONS AND NIP

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**ABSTRACT.** This paper revisits and extends a bridge between functional analysis and model theory, emphasizing its relevance to the theoretical foundations of machine learning. We show that the compactness behavior of families of Baire class 1 functions mirrors the learnability conditions in the sense of *Probably Approximately Correct* (PAC) learning, and that the failure of compactness corresponds to the presence of infinite *Vapnik-Chervonenkis* (VC) dimension. From this perspective, Rosenthal compacta emerge as the natural topological counterpart of PAC-learnable concept classes, while NIP vs. IP structures capture the precise boundary between analytical regularity and combinatorial intractability. These parallels suggest a unified framework linking compactness, definability, and learnability, exemplifying how the topology of function spaces encodes the algorithmic and epistemic limits of prediction.

## 1. INTRODUCTION

Suppose that  $A$  is a subset of the real line  $\mathbb{R}$  and that  $\overline{A}$  is its *closure*. It is a well-known fact that any point of closure of  $A$ , say  $x \in \overline{A}$ , can be *approximated* by points inside of  $A$ , in the sense that a sequence  $\{x_n\}_{n \in \mathbb{N}} \subseteq A$  must exist with the property that  $\lim_{n \rightarrow \infty} x_n = x$ . For most applications we wish to approximate objects more complicated than points, such as functions.

Suppose we wish to build a neural network that decides, given an 8 by 8 black-and-white image of a hand-written scribble, what single decimal digit the scribble represents. Maybe there exists  $f$ , a function representing an optimal solution to this classifier. Thus if  $X$  is the set of all (possible) images, then for  $I \in X$ ,  $f(I) \in \{0, 1, 2, \dots, 9\}$  is the “best” (or “good enough” for whatever deployment is needed) possible guess. Training the neural network involves approximating  $f$  until its guesses are within an acceptable error range. In general,  $f$  might be a function defined on a more complicated topological space  $X$ .

Often computers' viable operations are restricted (addition, subtraction, multiplication, division, etc.) and so we want to approximate a complicated function using simple functions (like polynomials). The problem is that, in contrast with mere points, functions in the closure of a set of functions need not be approximable (meaning the pointwise limit of a sequence of functions) by functions in the set.

Functions that are the pointwise limit of continuous functions are *Baire class 1 functions*, and the set of all of these is denoted by  $B_1(X)$ . Notice that these are not necessarily continuous themselves! A set of Baire class 1 functions,  $A$ , will be relatively compact if its closure consists of just Baire class 1 functions (we delay the formal definition of *relatively compact* until Section 2, but the fact mentioned here is sufficient). The Bourgain-Fremlin-Talagrand (BFT) theorem reveals a precise correspondence between relative compactness in  $B_1(X)$  and the model-theoretic

<sup>32</sup> notion of *Non-Independence Property* (NIP). This was realized by Pierre Simon in  
<sup>33</sup> [Sim15b].

<sup>34</sup> Simon's insight was to view definable families of functions as sets of real-valued  
<sup>35</sup> functions on type spaces and to interpret relative compactness in  $B_1(X)$  as a form  
<sup>36</sup> of "tame behavior" under ultrafilter limits. From this perspective, NIP theories are  
<sup>37</sup> those whose definable families behave like relatively compact sets of Baire class 1  
<sup>38</sup> functions, avoiding the wild,  $\beta\mathbb{N}$ -like configurations that witness instability. This  
<sup>39</sup> observation opened a new bridge between analysis and logic: topological compact-  
<sup>40</sup> ness corresponds to the absence of combinatorial independence. Simon's later de-  
<sup>41</sup> velopments connected these ideas to *Keisler measures* and *empirical averages*, al-  
<sup>42</sup> lowing tools from functional analysis to be used to study learnability and definable  
<sup>43</sup> types. This reinterpretation of model-theoretic tameness through the lens of the  
<sup>44</sup> BFT theorem has made NIP a central notion not only in stability theory but also  
<sup>45</sup> in contemporary connections with learning theory and ergodic analysis.

Historically, the notion of NIP arises from Shelah's foundational work on the classification theory of models. In his seminal book *Unstable Theories* [She78], Shelah introduced the independence property as a key dividing line within unstable structures, identifying the class of *stable* theories inside those in which this property fails. Fix a first-order formula  $\varphi(x, y)$  in a language  $L$  and a model  $M$  of an  $L$ -theory  $T$ . We say that  $\varphi(x, y)$  has the *independence property (IP)* in  $M$  if there is a sequence  $(c_i)_{i \in \mathbb{N}} \subseteq M^{|x|}$  such that for every  $S \subseteq \mathbb{N}$  there is  $a_S \in M^{|y|}$  with

$$\forall i \in \mathbb{N}, \quad M \models \varphi(c_i, a_S) \iff i \in S.$$

<sup>46</sup> The formula  $\varphi(x, y)$  has the IP if it does so in some model  $M$ , and the formula  
<sup>47</sup> has the *non-independence property (NIP)* if it does not have the IP. The latter  
<sup>48</sup> notion of NIP generalizes stability by forbidding the full combinatorial indepen-  
<sup>49</sup> dence pattern while allowing certain controlled forms of instability. Thus, Simon's  
<sup>50</sup> interpretation of the BFT theorem can be viewed as placing Shelah's dividing line  
<sup>51</sup> into a topological-analytic framework, connecting the earliest notions of stability  
<sup>52</sup> to compactness phenomena in spaces of Baire class 1 functions.

<sup>53</sup> One of the most important innovations in Machine Learning is the mathemati-  
<sup>54</sup> cal notion, introduced by Turing Awardee Leslie Valiant in the 1980s, of 'probably  
<sup>55</sup> approximately correct learning', or PAC-learning for short [BD19]. We give a stan-  
<sup>56</sup> dard but short overview of these concepts in the context that is relevant to this  
<sup>57</sup> work.

<sup>58</sup> Consider the following important idea in data classification. Suppose that  $A$  is  
<sup>59</sup> a set and that  $\mathcal{C}$  is a collection of sets. We say that  $\mathcal{C}$  *shatters*  $A$  if every subset  
<sup>60</sup> of  $A$  is of the form  $C \cap A$  for some  $C \in \mathcal{C}$ . For a classical geometric example, if  
<sup>61</sup>  $A$  is the set of four points on the Euclidean plane of the form  $(\pm 1, \pm 1)$ , then the  
<sup>62</sup> collection of all half-planes does not shatter  $A$ , the collection of all open balls does  
<sup>63</sup> not shatter  $A$ , but the collection of all convex sets shatters  $A$ . While  $A$  need not be  
<sup>64</sup> finite, it will usually be assumed to be so in Machine Learning applications. A finer  
<sup>65</sup> way to distinguish collections of sets that shatter a given set from those that do  
<sup>66</sup> not is by the *Vapnik-Chervonenkis dimension (VC-dimension)*, which is equal to  
<sup>67</sup> the cardinality of the largest finite set shattered by the collection, in case it exists,  
<sup>68</sup> or to infinity otherwise.

<sup>69</sup> A concrete illustration of these ideas appears when considering threshold clas-  
<sup>70</sup> sifiers on the real line. Let  $\mathcal{H}$  be the collection of all indicator functions  $h_t$  given

71 by  $h_t(x) = 1$  if  $x \leq t$  and  $h_t(x) = 0$  otherwise. Each  $h_t$  is a Baire class 1 function,  
 72 and the family  $\mathcal{H}$  is relatively compact in  $B_1(\mathbb{R})$ . In model-theoretic terms,  
 73  $\mathcal{H}$  is NIP, since no configuration of points and thresholds can realize the full inde-  
 74 pendence pattern of a binary matrix. By contrast, the family of parity functions  
 75  $\{x \mapsto (-1)^{\langle w, x \rangle} : w \in \{0, 1\}^n\}$  on  $\{0, 1\}^n$  (here  $\langle w, x \rangle$  is the usual vector dot product)  
 76 has the independence property and fails relative compactness in  $B_1(X)$ , capturing  
 77 the analytical meaning of instability. This dichotomy mirrors the behavior of con-  
 78 cept classes with finite versus infinite VC dimension in statistical learning theory.

Going back to the model theoretic framework, let

$$\mathcal{F}_\varphi(M) := \{\varphi(M, a) : a \in M^{|y|}\}$$

79 be the family of subsets of  $M^{|x|}$  defined by instances of the formula  $\varphi$ , where  
 80  $\varphi(M, a)$  is the set of  $|x|$ -tuples  $c$  in  $M$  for which  $M \models \varphi(c, a)$ . The fundamental  
 81 theorem of statistical learning states that a binary hypothesis class is PAC-learnable  
 82 if and only if it has finite VC-dimension, and the subsequent theorem connects the  
 83 rest of the concepts presented in this section.

84 **Theorem 1.1** (Laskowski). *The formula  $\varphi(x, y)$  has the NIP if and only if  $\mathcal{F}_\varphi(M)$   
 85 has finite VC-dimension.*

86 For two simple examples of formulas satisfying the NIP, consider first the lan-  
 87 guage  $L = \{<\}$  and the model  $M = (\mathbb{R}, <)$  of the reals with their usual linear order.  
 88 Take the formula  $\varphi(x, y)$  to mean  $x < y$ , then  $\varphi(M, a) = (-\infty, a)$ , and so  $\mathcal{F}_\varphi(M)$   
 89 is just the set of left open rays. The VC-dimension of this collection is 1, since it  
 90 can shatter a single point, but no two point set can be shattered since the rays are  
 91 downwards closed. Now in contrast, the collection of open intervals, given by the  
 92 formula  $\varphi(x; y_1, y_2) := (y_1 < x) \wedge (x < y_2)$ , has VC-dimension 2.

93 In this work, we study the corresponding notions of NIP (and hence PAC-  
 94 learnability) in the context of Compositional Computation Structures introduced  
 95 in [ADIW24].

## 96 2. GENERAL TOPOLOGICAL PRELIMINARIES

97 In this section we give preliminaries from general topology and function space  
 98 theory. We include some of the proofs for completeness but a reader familiar with  
 99 these topics may skip them.

100 A *Polish space* is a separable and completely metrizable topological space. The  
 101 most important examples are the reals  $\mathbb{R}$ , the Cantor space  $2^\mathbb{N}$  (the set of all infinite  
 102 binary sequences, endowed with the product topology), and the Baire space  $\mathbb{N}^\mathbb{N}$  (the  
 103 set of all infinite sequences of naturals, also with the product topology). Countable  
 104 products of Polish spaces are Polish; this includes spaces like  $\mathbb{R}^\mathbb{N}$ , the space of  
 105 sequences of real numbers. A subspace of a Polish space is itself Polish if and only  
 106 if it is a  $G_\delta$ -set, that is, it can be written as the intersection of a countable family  
 107 of open subsets; in particular, closed subsets and open subsets of Polish spaces are  
 108 also Polish spaces.

109 In this work we talk a lot about subspaces, and so there is a pertinent subtlety  
 110 of the definitions worth mentioning: *completely metrizable space* is not the same  
 111 as *complete metric space*; for an illustrative example, notice that  $(0, 1)$  is home-  
 112 omorphic to the real line, and thus a Polish space (being Polish is a topological  
 113 property), but with the metric inherited from the reals, as a subspace,  $(0, 1)$  is **not**  
 114 a complete metric space. In summary, a Polish space has its topology generated by

115 *some* complete metric, but other metrics generating the same topology might not  
 116 be. In practice, such as when studying descriptive set theory, one finds that we can  
 117 often keep the metric implicit.

118 Given two topological spaces  $X$  and  $Y$  we denote by  $B_1(X, Y)$  the set of all func-  
 119 tions  $f : X \rightarrow Y$  such that for all open  $U \subseteq Y$ ,  $f^{-1}[U]$  is an  $F_\sigma$  subset of  $X$  (that  
 120 is, a countable union of closed sets); we call these types of functions *Baire class*  
 121 *1 functions*. When  $Y = \mathbb{R}$  we simply denote this collection by  $B_1(X)$ . We endow  
 122  $B_1(X, Y)$  with the topology of pointwise convergence (the topology inherited  
 123 from the product topology of  $Y^X$ ). By  $C_p(X, Y)$  we denote the set of all contin-  
 124 uous functions  $f : X \rightarrow Y$  with the topology of pointwise convergence. Similarly,  
 125  $C_p(X) := C_p(X, \mathbb{R})$ . A natural question is, how do topological properties of  $X$   
 126 translate to  $C_p(X)$  and vice versa? These questions, and in general the study of  
 127 these spaces, are the concern of  $C_p$ -theory, an active field of research in general  
 128 topology which was pioneered by A. V. Arhangel'skiĭ and his students in the 1970's  
 129 and 1980's. This field has found many exciting applications in model theory and  
 130 functional analysis. Good recent surveys on the topics include [HT23] and [Tka11].  
 131 We begin with the following:

132 **Fact 2.1.** *If all open subsets of  $X$  are  $F_\sigma$  (in particular if  $X$  is metrizable), then*  
 133  $C_p(X, Y) \subseteq B_1(X, Y)$ .

134 The proof of the following fact (due to Baire) can be found in Section 10 of  
 135 [Tod97].

136 **Fact 2.2** (Baire). *If  $X$  is a complete metric space, then the following are equivalent:*

- 137 (i)  *$f$  is a Baire class 1 function, that is,  $f \in B_1(X)$ .*
- 138 (ii)  *$f$  is a pointwise limit of continuous functions.*
- 139 (iii) *For every closed  $F \subseteq X$ , the restriction  $f|_F$  has a point of continuity.*

140 Moreover, if  $X$  is Polish and  $f \notin B_1(X)$ , then there exists countable  $D_0, D_1 \subseteq X$  and  
 141 reals  $a < b$  such that  $\overline{D_0} = \overline{D_1}$ ,  $D_0 \subseteq f^{-1}(-\infty, a]$  and  $D_1 \subseteq f^{-1}[b, \infty)$ .

142 A subset  $L \subseteq X$  is *relatively compact* in  $X$  if the closure of  $L$  in  $X$  is compact.  
 143 Relatively compact subsets of  $B_1(X)$  (for  $X$  Polish space) have been objects of  
 144 interest to many people working in Analysis and Topological Dynamics. We begin  
 145 with the following well-known result. Recall that a set  $A \subseteq \mathbb{R}^X$  of real-valued  
 146 functions is *pointwise bounded* if for every  $x \in X$  there is  $M_x > 0$  such that  $|f(x)| <$   
 147  $M_x$  for all  $f \in A$ . We include the proof for the reader's convenience:

148 **Lemma 2.3.** *Let  $X$  be a Polish space and  $A \subseteq B_1(X)$  be pointwise bounded. The*  
 149 *following are equivalent:*

- 150 (i)  *$A$  is relatively compact in  $B_1(X)$ .*
- 151 (ii)  *$A$  is relatively countably compact in  $B_1(X)$ , i.e., every countable subset of*  
 152  *$A$  has an accumulation point in  $B_1(X)$ .*
- 153 (iii)  *$\overline{A} \subseteq B_1(X)$ , where  $\overline{A}$  denotes the closure in  $\mathbb{R}^X$ .*

154 *Proof.* By definition, being pointwise bounded means that there is, for each  $x \in X$ ,  
 155  $M_x > 0$  such that, for every  $f \in A$ ,  $|f(x)| \leq M_x$ .

156 (i) $\Rightarrow$ (ii) holds in general.

157 (ii) $\Rightarrow$ (iii) Assume that  $A$  is relatively countably compact in  $B_1(X)$  and that  
 158  $f \in \overline{A} \setminus B_1(X)$ . By Fact 2.2, there are countable  $D_0, D_1 \subseteq X$  with  $\overline{D_0} = \overline{D_1}$ , and  
 159  $a < b$  such that  $D_0 \subseteq f^{-1}(-\infty, a]$  and  $D_1 \subseteq f^{-1}[b, \infty)$ . We claim that there is a

160 sequence  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  such that for all  $x \in D_0 \cup D_1$ ,  $\lim_{n \rightarrow \infty} f_n(x) = f(x)$ . Indeed,  
 161 use the countability to enumerate  $D_0 \cup D_1$  as  $\{x_n\}_{n \in \mathbb{N}}$ . Then find, for each positive  
 162  $n$ ,  $f_n \in A$  with  $|f_n(x_i) - f(x_i)| < \frac{1}{n}$  for all  $i \leq n$ . The claim follows.

163 By relative countable compactness of  $A$ , there is an accumulation point  $g \in$   
 164  $B_1(X)$  of  $\{f_n\}_{n \in \mathbb{N}}$ . It is straightforward to show that since  $f$  and  $g$  agree on  $D_0 \cup D_1$ ,  
 165  $g$  does not have a point of continuity on the closed set  $\overline{D_0} = \overline{D_1}$ , which contradicts  
 166 Fact 2.2.

167 (iii) $\Rightarrow$ (i) Suppose that  $\overline{A} \subseteq B_1(X)$ . Then  $\overline{A} \cap B_1(X) = \overline{A}$  is a closed subset of  
 168  $\prod_{x \in X} [-M_x, M_x]$ ; Tychonoff's theorem states that the product of compact spaces  
 169 is always compact, and since closed subsets of compact spaces are compact,  $\overline{A}$  must  
 170 be compact, as desired.  $\square$

171 **2.1. From Rosenthal's dichotomy to NIP.** The fundamental idea that con-  
 172 nects the rich theory here presented to real-valued computations is the concept  
 173 of an *approximation*. In the reals, points of closure from some subset can always  
 174 be approximated by points inside the set, via a convergent sequence. For more  
 175 complicated spaces, such as  $C_p(X)$ , this fails in a remarkably intriguing way. Let  
 176 us show an example that is actually the protagonist of a celebrated result. Con-  
 177 sider the Cantor space  $X = 2^{\mathbb{N}}$  and let  $p_n(x) = x(n)$  define a continuous mapping  
 178  $X \rightarrow \{0, 1\}$ . Then one can show (see Chapter 1.1 of [Tod97] for details) that, per-  
 179 haps surprisingly, the only continuous functions in the closure of  $\{p_n\}_{n \in \mathbb{N}}$  are the  
 180 functions  $p_n$  themselves; moreover, none of the subsequences of  $\{p_n\}_{n \in \mathbb{N}}$  converge.  
 181 In some sense, this example is the worst possible scenario for convergence. The  
 182 topological space obtained from this closure is well-known. Topologists refer to it  
 183 as the Stone-Čech compactification of the discrete space of natural numbers, or  $\beta\mathbb{N}$   
 184 for short, and it is an important object of study in general topology.

185 **Theorem 2.4** (Rosenthal's Dichotomy). *If  $X$  is Polish and  $\{f_n\} \subseteq C_p(X)$  is point-  
 186 wise bounded, then either  $\{f_n\}_{n \in \mathbb{N}}$  contains a convergent subsequence or a subse-  
 187 quence whose closure (in  $\mathbb{R}^X$ ) is homeomorphic to  $\beta\mathbb{N}$ .*

188 In other words, a pointwise bounded set of continuous functions will either con-  
 189 tain a subsequence that converges or a subsequence whose closure is essentially  
 190 the same as the example mentioned in the previous paragraphs (the worst possible  
 191 scenario). Note that in the preceding example, the functions are trivially pointwise  
 192 bounded in  $\mathbb{R}^X$  as the functions can only take values 0 and 1.

193 If we intend to generalize our results from  $C_p(X)$  to the bigger space  $B_1(X)$ , we  
 194 find a similar dichotomy. Either every point of closure of the set of functions will  
 195 be a Baire class 1 function, or there is a sequence inside the set that behaves in the  
 196 worst possible way (which in this context, is the IP!). The theorem is usually not  
 197 phrased as a dichotomy but rather as an equivalence (with the NIP instead):

198 **Theorem 2.5** (Bourgain-Fremlin-Talagrand, Theorem 4G in [BFT78]). *Let  $X$  be  
 199 a Polish space and  $A \subseteq C_p(X)$  be pointwise bounded. The following are equivalent:*

- 200 (i)  $A$  is relatively compact in  $B_1(X)$ , i.e.,  $\overline{A} \subseteq B_1(X)$ .  
 201 (ii) For every  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and for every  $a < b$  there is  $I \subseteq \mathbb{N}$  such that

$$\bigcap_{n \in I} f_n^{-1}(-\infty, a] \cap \bigcap_{n \notin I} f_n^{-1}[b, \infty) = \emptyset.$$

201 Our goal now is to characterize relatively compact subsets of  $B_1(X, Y)$  when  
 202  $Y = \mathbb{R}^P$  with  $P$  countable. Given  $P \in \mathcal{P}$  we denote the *projection map* onto the

203 P-coordinate by  $\pi_P : \mathbb{R}^P \rightarrow \mathbb{R}$ . From a high-level topological interpretation, the  
 204 subsequent lemma states that, in this context, the spaces  $\mathbb{R}$  and  $\mathbb{R}^P$  are really not  
 205 that different, and that if we understand the Baire class 1 functions of one space,  
 206 then we also understand the functions of both. In fact,  $\mathbb{R}$  and any other Polish  
 207 space is embeddable as a closed subspace of  $\mathbb{R}^P$ .

208 **Lemma 2.6.** *Let  $X$  be a Polish space and  $\mathcal{P}$  be a countable set. Then,  $f \in B_1(X, \mathbb{R}^P)$   
 209 if and only if  $\pi_P \circ f \in B_1(X)$  for all  $P \in \mathcal{P}$ .*

*Proof.* Only one implication needs a proof. Suppose that  $\pi_P \circ f \in B_1(X)$  for all  
 $P \in \mathcal{P}$ . Let  $V$  be a basic open subset of  $\mathbb{R}^P$ . That is, there exists a finite  $\mathcal{P}' \subseteq \mathcal{P}$   
 such that  $V = \bigcap_{P \in \mathcal{P}'} \pi_P^{-1}[U_P]$  where  $U_P$  is open in  $\mathbb{R}$ . Finally,

$$f^{-1}[V] = \bigcap_{P \in \mathcal{P}'} (\pi_P \circ f)^{-1}[U_P]$$

210 is an  $F_\sigma$  set. Since  $\mathcal{P}$  is countable,  $\mathbb{R}^P$  is second countable so every open set  $U$  in  
 211  $\mathbb{R}^P$  is a countable union of basic open sets. Hence,  $f^{-1}[U]$  is  $F_\sigma$ .  $\square$

212 Below we consider  $\mathcal{P}$  with the discrete topology. For each  $f : X \rightarrow \mathbb{R}^P$  denote  
 213  $\hat{f}(P, x) := \pi_P \circ f(x)$  for all  $(P, x) \in \mathcal{P} \times X$ . Similarly, for each  $g : \mathcal{P} \times X \rightarrow \mathbb{R}$  denote  
 214  $\check{g}(x)(P) := g(P, x)$ . Given  $A \subseteq (\mathbb{R}^P)^X$ , we denote  $\hat{A}$  as the set of all  $\hat{f}$  such that  
 215  $f \in A$ .

216 The map  $(\mathbb{R}^P)^X \rightarrow \mathbb{R}^{P \times X}$  given by  $f \mapsto \hat{f}$  is a homeomorphism and its inverse is  
 217 given by  $g \mapsto \check{g}$ .

218 **Lemma 2.7.** *Let  $X$  be a Polish space and  $\mathcal{P}$  be countable. Then,  $f \in B_1(X, \mathbb{R}^P)$  if  
 219 and only if  $\hat{f} \in B_1(\mathcal{P} \times X)$ .*

*Proof.* ( $\Rightarrow$ ) Given an open set of reals  $U$ , we have that for every  $P \in \mathcal{P}$ ,  $f^{-1}[\pi_P^{-1}[U]]$   
 is  $F_\sigma$  by Lemma 2.6. Given that  $\mathcal{P}$  is a discrete countable space, we observe that

$$\hat{f}^{-1}[U] = \bigcup_{P \in \mathcal{P}} (\{P\} \times f^{-1}[\pi_P^{-1}[U]])$$

is also an  $F_\sigma$  set. ( $\Leftarrow$ ) By lemma 2.6 it suffices to show that  $\pi_P \circ f \in B_1(X)$  for all  
 $P \in \mathcal{P}$ . Fix an open  $U \subseteq \mathbb{R}$ . Write  $\hat{f}^{-1}[U] = \bigcup_{n \in \mathbb{N}} F_n$  where  $F_n$  is closed in  $\mathcal{P} \times X$ .  
 Then,

$$(\pi_P \circ f)^{-1}[U] = \bigcup_{n \in \mathbb{N}} \{x \in X : (P, x) \in F_n\}$$

220 which is  $F_\sigma$ .  $\square$

221 We now direct our attention to a notion of the NIP that is more general than  
 222 the one from the introduction. It can be interpreted as a sort of continuous version  
 223 of the one presented in the preceding section.

**Definition 2.8.** We say that  $A \subseteq \mathbb{R}^X$  has the *Non-Independence Property* (NIP)  
 if and only if for every  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and for every  $a < b$  there are finite disjoint  
 sets  $E, F \subseteq \mathbb{N}$  such that

$$\bigcap_{n \in E} f_n^{-1}(-\infty, a] \cap \bigcap_{n \in F} f_n^{-1}[b, \infty) = \emptyset.$$

Note that if  $X$  is compact and  $A \subseteq C_p(X)$ , then  $A$  has the NIP if and only if for every  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and for every  $a < b$  there is  $I \subseteq \mathbb{N}$  such that

$$\bigcap_{n \in I} f_n^{-1}(-\infty, a] \cap \bigcap_{n \notin I} f_n^{-1}[b, \infty) = \emptyset.$$

Given  $A \subseteq Y^X$  and  $K \subseteq X$  we write  $A|_K := \{f|_K : f \in A\}$ , i.e., the set of all restrictions of functions in  $A$  to  $K$ . The following Theorem is a slightly more general version of Theorem 2.5.

**Theorem 2.9.** *Assume that  $\mathcal{P}$  is countable,  $X$  is a Polish space, and  $A \subseteq C_p(X, \mathbb{R}^\mathcal{P})$  is such that  $\pi_P \circ A$  is pointwise bounded for all  $P \in \mathcal{P}$ . The following are equivalent for every compact  $K \subseteq X$ :*

- (1)  $\overline{A|_K} \subseteq B_1(K, \mathbb{R}^\mathcal{P})$ .
- (2)  $\pi_P \circ A|_K$  has the NIP for every  $P \in \mathcal{P}$ .

*Proof.* (1)  $\Rightarrow$  (2). Let  $P \in \mathcal{P}$ . Fix  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and  $a < b$ . By (1) we have that  $A|_K \subseteq B_1(K, \mathbb{R}^\mathcal{P})$ . Applying the homeomorphism  $f \mapsto \hat{f}$  and using lemma 2.7 we get  $\overline{\hat{A}|_{\mathcal{P} \times K}} \subseteq B_1(\mathcal{P} \times K)$ . By Theorem 2.5, there is  $I \subseteq \mathbb{N}$  such that

$$(\mathcal{P} \times K) \cap \bigcap_{n \in I} \hat{f}_n^{-1}(-\infty, a] \cap \bigcap_{n \notin I} \hat{f}_n^{-1}[b, \infty) = \emptyset$$

Hence,

$$K \cap \bigcap_{n \in I} (\pi_P \circ f_n)^{-1}(-\infty, a] \cap \bigcap_{n \notin I} (\pi_P \circ f_n)^{-1}[b, \infty) = \emptyset$$

By compactness, there are finite  $E \subseteq I$  and  $F \subseteq \mathbb{N} \setminus I$  such that

$$K \cap \bigcap_{n \in E} (\pi_P \circ f_n)^{-1}(-\infty, a] \cap \bigcap_{n \in F} (\pi_P \circ f_n)^{-1}[b, \infty) = \emptyset$$

Thus,  $\pi_P \circ A|_L$  has the NIP.

(2)  $\Rightarrow$  (1) Fix  $f \in \overline{A|_K}$ . By lemma 2.6 it suffices to show that  $\pi_P \circ f \in B_1(K)$  for all  $P \in \mathcal{P}$ . By (2),  $\pi_P \circ A|_K$  has the NIP. Hence, by Theorem 2.5 we have  $\pi_P \circ A|_K \subseteq B_1(K)$ . But then  $\pi_P \circ f \in \pi_P \circ A|_K \subseteq B_1(K)$ .  $\square$

Lastly, a simple but significant result that helps understand the operation of restricting a set of functions to a specific subspace of the domain space  $X$ , of course in the context of the NIP, is that we may always assume that said subspace is closed. Concretely, whether we take its closure or not has no effect on the NIP:

**Lemma 2.10.** *Assume that  $X$  is Hausdorff and that  $A \subseteq C_p(X)$ . The following are equivalent for every  $L \subseteq X$ :*

- (i)  $A_L$  has the NIP.
- (ii)  $A|_{\overline{L}}$  has the NIP.

*Proof.* It suffices to show that (i)  $\Rightarrow$  (ii). Suppose that (ii) does not hold, i.e., that there are  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and  $a < b$  such that for all finite disjoint  $E, F \subseteq \mathbb{N}$ :

$$\overline{L} \cap \bigcap_{n \in E} f_n^{-1}(-\infty, a] \cap \bigcap_{n \in F} f_n^{-1}[b, \infty) \neq \emptyset.$$

Pick  $a' < b'$  such that  $a < a' < b' < b$ . Then, for any finite disjoint  $E, F \subseteq \mathbb{N}$  we can choose

$$x \in \overline{L} \cap \bigcap_{n \in E} f_n^{-1}(-\infty, a') \cap \bigcap_{n \in F} f_n^{-1}(b', \infty)$$

By definition of closure:

$$L \cap \bigcap_{n \in E} f_n^{-1}(-\infty, a'] \cap \bigcap_{n \in F} f_n^{-1}[b', \infty) \neq \emptyset.$$

244 This contradicts (i). □

### 245 3. NIP IN THE CONTEXT OF COMPOSITIONAL COMPUTATION STRUCTURES

246 In this section, we study what the NIP tell us in the context of deep computations as defined in [ADIW24]. We say a structure  $(L, \mathcal{P}, \Gamma)$  is a *Compositional 247 Computation Structure* (CCS) if  $L \subseteq \mathbb{R}^{\mathcal{P}}$  is a subspace of  $\mathbb{R}^{\mathcal{P}}$ , with the pointwise 248 convergence topology, and  $\Gamma \subseteq L^L$  is a semigroup under composition. The motivation 249 for CCS comes from (continuous) model theory, where  $\mathcal{P}$  is a fixed collection 250 of predicates and  $L$  is a (real-valued) structure. Every point in  $L$  is identified with 251 its “type”, which is the tuple of all values the point takes on the predicates from 252  $\mathcal{P}$ , i.e., an element of  $\mathbb{R}^{\mathcal{P}}$ . In this context, elements of  $\mathcal{P}$  are called *features*. In the 253 discrete model theory framework, one views the space of complete-types as a sort of 254 compactification of the structure  $L$ . In this context, we don’t want to consider only 255 points in  $L$  (realized types) but in its closure  $\bar{L}$  (possibly unrealized types). The 256 problem is that the closure  $\bar{L}$  is not necessarily compact, an assumption that turns 257 out to be very useful in the context of continuous model theory. To bypass this 258 problem in a framework for deep computations, Alva, Dueñez, Iovino and Walton 259 introduced in [ADIW24] the concept of *shards*, which essentially consists in covering 260 (a large fragment) of the space  $\bar{L}$  by compact, and hence pointwise-bounded, 261 subspaces (shards). We shall give the formal definition next.

263 A *sizer* is a tuple  $r_{\bullet} = (r_p)_{p \in \mathcal{P}}$  of positive real numbers indexed by  $\mathcal{P}$ . Given a 264 sizer  $r_{\bullet}$ , we define the  $r_{\bullet}$ -shard as:

$$L[r_{\bullet}] = L \cap \prod_{p \in \mathcal{P}} [-r_p, r_p]$$

265 For an illustrative example, we can frame Newton’s polynomial root approximation 266 method in the context of a CCS (see Example 5.6 of [ADIW24] for details) as 267 follows. Begin by considering the extended complex numbers  $\hat{\mathbb{C}} := \mathbb{C} \cup \{\infty\}$  with 268 the usual Riemann sphere topology that makes it into a compact space (where 269 unbounded sequences converge to  $\infty$ ). In fact, not only is this space compact 270 but it is covered by the shard given by the sizer  $(1, 1, 1)$  (the unit sphere is con- 271 tained in the cube  $[-1, 1]^3$ ). The space  $\hat{\mathbb{C}}$  is homeomorphic to the usual unit 272 sphere  $S^2 := \{(x, y, z) : x^2 + y^2 + z^2 = 1\}$  of  $\mathbb{R}^3$ , by means of the stereographic 273 projection and its inverse  $\hat{\mathbb{C}} \rightarrow S^2$ . This function is regarded as a triple of pred- 274 icates  $x, y, z : \hat{\mathbb{C}} \rightarrow [-1, 1]$  where each will map an extended complex number to 275 its corresponding real coordinate on the cube  $[-1, 1]^3$ . Now fix the cubic com- 276 plex polynomial  $p(s) := s^3 - 1$ , and consider the map which performs one step 277 in Newton’s method at a particular (extended) complex number  $s$ , for finding 278 a root of  $p$ ,  $\gamma_p : \hat{\mathbb{C}} \rightarrow \hat{\mathbb{C}}$ . The explicit inner workings of  $\gamma_p$  are irrelevant for 279 this example, except for the fact that it is a continuous mapping. It follows that 280  $(S^2, \{x, y, z\}, \{\gamma_p^k : k \in \mathbb{N}\})$  is a CCS. The idea is that repeated applications of 281  $\gamma_p(s), \gamma_p \circ \gamma_p(s), \gamma_p \circ \gamma_p \circ \gamma_p(s), \dots$  would approximate a root of  $p$  provided  $s$  was 282 a good enough initial guess.

283     The  $r_\bullet$ -type-shard is defined as  $\mathcal{L}[r_\bullet] = \overline{L[r_\bullet]}$  and  $\mathcal{L}_{sh}$  is the union of all type-  
284     shards. Notice that  $\mathcal{L}_{sh}$  is not necessarily equal to  $\mathcal{L} = \overline{L}$ , unless  $\mathcal{P}$  is countable  
285     (see [ADIW24]). A *transition* is a map  $f : L \rightarrow L$ , in particular, every element  
286     in the semigroup  $\Gamma$  is a transition (these are called *realized computations*). In  
287     practice, one would like to work with “definable” computations, i.e., ones that can  
288     be described by a computer. In this topological framework, being continuous is an  
289     expected requirement. However, as in the case of complete-types in model theory,  
290     we will work with “unrealized computations”, i.e., maps  $f : \mathcal{L}_{sh} \rightarrow \mathcal{L}_{sh}$ . Note that  
291     continuity of a computation does not imply that it can be continuously extended  
292     to  $\mathcal{L}_{sh}$ .

293     Suppose that a transition map  $f : L \rightarrow \mathcal{L}$  can be extended continuously to a  
294     map  $\mathcal{L} \rightarrow \mathcal{L}$ . Then, the Stone-Weierstrass theorem implies that the feature  $P \circ f$   
295     (here  $P$  is a fixed predicate, and the feature is hence continuous) can be uniformly  
296     approximated by polynomials on the compact set  $\mathcal{L}[r_\bullet]$ . Theorem 2.2 [ADIW24]  
297     of formalizes the converse of this fact, in the sense that transitions maps that  
298     are not continuously extendable in this fashion cannot be obtained from simple  
299     constructions involving predicates. Under this framework, the features  $P \circ f$  of such  
300     transitions  $f$  are not approximable by polynomials, and so they are understood as  
301     “non-computable” since, again, we expect the operations computers carry out to be  
302     determined by elementary algebra corresponding to polynomials (namely addition  
303     and multiplication). Therefore it is crucial we assume some extendibility conditions.

304     We say that the CCS  $(L, \mathcal{P}, \Gamma)$  satisfies the *Extendibility Axiom* if for all  $\gamma \in \Gamma$ ,  
305     there is  $\tilde{\gamma} : \mathcal{L}_{sh} \rightarrow \mathcal{L}_{sh}$  such that for every sizer  $r_\bullet$  there is an  $s_\bullet$  such that  
306      $\tilde{\gamma}|_{\mathcal{L}[r_\bullet]} : \mathcal{L}[r_\bullet] \rightarrow \mathcal{L}[s_\bullet]$  is continuous. For a deeper discussion about this axiom, we  
307     refer the reader to [ADIW24].

308     A collection  $R$  of sizers is called *exhaustive* if  $\mathcal{L}_{sh} = \bigcup_{r_\bullet \in R} \mathcal{L}[r_\bullet]$ . We say that  
309      $\Delta \subseteq \Gamma$  is  $R$ -*confined* if  $\gamma|_{L[r_\bullet]} : L[r_\bullet] \rightarrow L[r_\bullet]$  for every  $r_\bullet \in R$  and  $\gamma \in \Delta$ . Elements in  
310      $\Delta$  are called *real-valued computations* (in this article we will refer to them simply as  
311     *computations*) and elements in  $\Delta \subseteq \mathcal{L}_{sh}^L$  are called (real-valued) *deep computations*  
312     or *ultracomputations*. By  $\tilde{\Delta}$  we denote the set of all extensions  $\tilde{\gamma}$  for  $\gamma \in \Delta$ . For a  
313     more complete description of this framework, we refer the reader to [ADIW24].

314     **3.1. NIP and Baire-1 definability of deep computations.** Under what con-  
315     ditions are deep computations Baire class 1, and thus well-behaved according to  
316     our framework, on type-shards? The next Theorem says that, again under the  
317     assumption that  $\mathcal{P}$  is countable, the space of deep computations is a Rosenthal  
318     compactum (when restricted to shards) if and only if the set of computations has  
319     the NIP on features. Hence, we can import the theory of Rosenthal compacta into  
320     this framework of deep computations.

321     **Theorem 3.1.** *Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom with  $\mathcal{P}$   
322     countable. Let  $R$  be an exhaustive collection of sizers. Let  $\Delta \subseteq \Gamma$  be  $R$ -confined. The  
323     following are equivalent.*

- 324        (1)  $\tilde{\Delta}|_{\mathcal{L}[r_\bullet]} \subseteq B_1(\mathcal{L}[r_\bullet], \mathcal{L}[r_\bullet])$  for all  $r_\bullet \in R$ .  
                (2)  $\pi_P \circ \Delta|_{L[r_\bullet]}$  has the NIP for all  $P \in \mathcal{P}$  and  $r_\bullet \in R$ , that is, for all  $P \in \mathcal{P}$ ,  
                 $r_\bullet \in R$ ,  $a < b$ ,  $\{\gamma_n\}_{n \in \mathbb{N}} \subseteq \Delta$  there are finite disjoint  $E, F \subseteq \mathbb{N}$  such that

$$L[r_\bullet] \cap \bigcap_{n \in E} (\pi_P \circ \gamma_n)^{-1}(-\infty, a] \cap \bigcap_{n \in F} (\pi_P \circ \gamma_n)^{-1}[b, \infty) = \emptyset.$$

325     Moreover, if any (hence all) of the preceding conditions hold, then every deep  
 326     computation  $f \in \bar{\Delta}$  can be extended to a Baire-1 function on shards, i.e., there is  
 327      $\tilde{f} : \mathcal{L}_{sh} \rightarrow \mathcal{L}_{sh}$  such that  $\tilde{f}|_{\mathcal{L}[r_\bullet]} \in B_1(\mathcal{L}[r_\bullet], \mathcal{L}[r_\bullet])$  for all  $r_\bullet \in R$ . In particular, on  
 328     each shard every deep computation is the pointwise limit of a countable sequence of  
 329     computations.

330     Proof. Since  $\mathcal{P}$  is countable, then  $\mathcal{L}[r_\bullet] \subseteq \mathbb{R}^{\mathcal{P}}$  is Polish. Also, the Extendibility  
 331     Axiom implies that  $\pi_P \circ \bar{\Delta}|_{\mathcal{L}[r_\bullet]}$  is a pointwise bounded set of continuous functions  
 332     for all  $P \in \mathcal{P}$ . Hence, Theorem 2.9 and Lemma 2.10 prove the equivalence of (1)  
 333     and (2). If (1) holds and  $f \in \bar{\Delta}$ , then write  $f = \underline{Ulim}_i \tilde{y}_i$  as an ultralimit. Define  
 334      $\tilde{f} := \underline{Ulim}_i \tilde{y}_i$ . Hence, for all  $r_\bullet \in R$  we have  $\tilde{f}|_{\mathcal{L}[r_\bullet]} \in \bar{\Delta}|_{\mathcal{L}[r_\bullet]} \subseteq B_1(\mathcal{L}[r_\bullet], \mathcal{L}[r_\bullet])$ . That  
 335     every deep computation is a pointwise limit of a countable sequence of computations  
 336     follows from the fact that for a Polish space  $X$  every compact subset of  $B_1(X)$  is  
 337     Fréchet-Urysohn (that is, a space where topological closures coincide with sequential  
 338     closures, see Theorem 3F in [BFT78] or Theorem 4.1 in [Deb13]).  $\square$

339     Given a countable set  $\Delta$  of computations satisfying the NIP on features and  
 340     shards (condition (2) of Theorem 3.1) we have that  $\bar{\Delta}|_{\mathcal{L}[r_\bullet]}$  (for a fixed sizer  $r_\bullet$ ) is  
 341     a separable *Rosenthal compactum* (compact subset of  $B_1(P \times \mathcal{L}[r_\bullet])$ ). The work of  
 342     Todorčević ([Tod99]) and Argyros, Dodos, Kanellopoulos ([ADK08]) culminates in  
 343     a trichotomy theorem for separable Rosenthal Compacta. Inspired by the work of  
 344     Glasner and Megrelishvili ([GM22]), we are interested to see how this allows us to  
 345     classify and obtain different levels of PAC-learnability (NIP).

346     Recall that a topological space  $X$  is *hereditarily separable* (HS) if every subspace  
 347     is separable and that  $X$  is *first countable* if every point in  $X$  has a countable local  
 348     basis. Every separable metrizable space is hereditarily separable and it is a result  
 349     of R. Pol that every hereditarily separable Rosenthal compactum is first countable  
 350     (see section 10 in [Deb13]). This suggests the following definition:

351     **Definition 3.2.** Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom and  $R$   
 352     be an exhaustive collection of sizers. Let  $\Delta \subseteq \Gamma$  be an  $R$ -confined countable set of  
 353     computations satisfying the NIP on shards and features (condition (2) in Theorem  
 354     3.1). We say that  $\Delta$  is:

- 355       (i)  $NIP_1$  if  $\bar{\Delta}|_{\mathcal{L}[r_\bullet]}$  is first countable for every  $r_\bullet \in R$ .
- 356       (ii)  $NIP_2$  if  $\bar{\Delta}|_{\mathcal{L}[r_\bullet]}$  is hereditarily separable for every  $r_\bullet \in R$ .
- 357       (iii)  $NIP_3$  if  $\bar{\Delta}|_{\mathcal{L}[r_\bullet]}$  is metrizable for every  $r_\bullet \in R$ .

358     Observe that  $NIP_3 \Rightarrow NIP_2 \Rightarrow NIP_1 \Rightarrow NIP$ . A natural question that would  
 359     continue this work is to find examples of CCS that separate these levels of NIP. In  
 360     [Tod99], Todorčević isolates 3 canonical examples of Rosenthal compacta that wit-  
 361     ness the failure of the converse implications above. We now present some separable  
 362     and non-separable examples of Rosenthal compacta:

- 363       (1) *Alexandroff compactification of a discrete space of size continuum.* For each  
 364          $a \in 2^\mathbb{N}$  consider the map  $\delta_a : 2^\mathbb{N} \rightarrow \mathbb{R}$  given by  $\delta_a(x) = 1$  if  $x = a$  and  
 365          $\delta_a(x) = 0$  otherwise. Let  $A(2^\mathbb{N}) = \{\delta_a : a \in 2^\mathbb{N}\} \cup \{0\}$ , where  $0$  is the zero  
 366         map. Notice that  $A(2^\mathbb{N})$  is a compact subset of  $B_1(2^\mathbb{N})$ , in fact  $\{\delta_a : a \in 2^\mathbb{N}\}$   
 367         is a discrete subspace of  $B_1(2^\mathbb{N})$  and its pointwise closure is precisely  $A(2^\mathbb{N})$ .  
 368         Hence, this is a Rosenthal compactum which is not first countable. Notice  
 369         that this space is also not separable.

- 370    (2) *Extended Alexandroff compactification.* For each finite binary sequence  $s \in$   
 371     $2^{<\mathbb{N}}$ , let  $v_s : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  be given by  $v_s(x) = 1$  if  $x$  extends  $s$  and  $v_s(x) =$   
 372    0 otherwise. Let  $\hat{A}(2^{\mathbb{N}})$  be the pointwise closure of  $\{v_s : s \in 2^{<\mathbb{N}}\}$ , i.e.,  
 373     $\hat{A}(2^{\mathbb{N}}) = A(2^{\mathbb{N}}) \cup \{v_s : s \in 2^{<\mathbb{N}}\}$ . Note that this space is a separable  
 374    Rosenthal compactum which is not first countable.  
 375    (3) *Split Cantor.* Let  $<$  be the lexicographic order in the space of infinite  
 376    binary sequences, i.e.,  $2^{\mathbb{N}}$ . For each  $a \in 2^{\mathbb{N}}$  let  $f_a^- : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  be given by  
 377     $f_a^-(x) = 1$  if  $x < a$  and  $f_a^-(x) = 0$  otherwise. Let  $f_a^+ : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  be given  
 378    by  $f_a^+(x) = 1$  if  $x \leq a$  and  $f_a^+(x) = 0$  otherwise. The split Cantor is the  
 379    space  $S(2^{\mathbb{N}}) = \{f_a^- : a \in 2^{\mathbb{N}}\} \cup \{f_a^+ : a \in 2^{\mathbb{N}}\}$ . This is a separable Rosenthal  
 380    compactum. One example of a countable dense subset is the set of all  $f_a^+$   
 381    and  $f_a^-$  where  $a$  is an infinite binary sequence that is eventually constant.  
 382    Moreover, it is hereditarily separable but it is not metrizable.  
 383    (4) *Alexandroff Duplicate.* Let  $K$  be any compact metric space and consider  
 384    the Polish space  $X = C(K) \sqcup K$ , i.e., the disjoint union of  $C(K)$  (with its  
 385    supremum norm topology) and  $K$ . For each  $a \in K$  define  $g_a^0, g_a^1 : X \rightarrow \mathbb{R}$  as  
 follows:

$$g_a^0(x) = \begin{cases} x(a), & x \in C(K) \\ 0, & x \in K \end{cases}$$

$$g_a^1(x) = \begin{cases} x(a), & x \in C(K) \\ \delta_a(x), & x \in K \end{cases}$$

383    Let  $D(K) = \{g_a^0 : a \in K\} \cup \{g_a^1 : a \in K\}$ . Notice that  $D(K)$  is a first  
 384    countable Rosenthal compactum. It is not separable if  $K$  is uncountable.  
 385    The interesting case will be when  $K = 2^{\mathbb{N}}$ .

- 386    (5) *Extended Alexandroff Duplicate of the split Cantor.* For each finite binary  
 387    sequence  $t \in 2^{<\mathbb{N}}$  let  $a_t \in 2^{\mathbb{N}}$  be the sequence starting with  $t$  and ending  
 388    with 0's and let  $b_t \in 2^{\mathbb{N}}$  be the sequence starting with  $t$  and ending with  
 389    1's. Define  $h_t : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  by

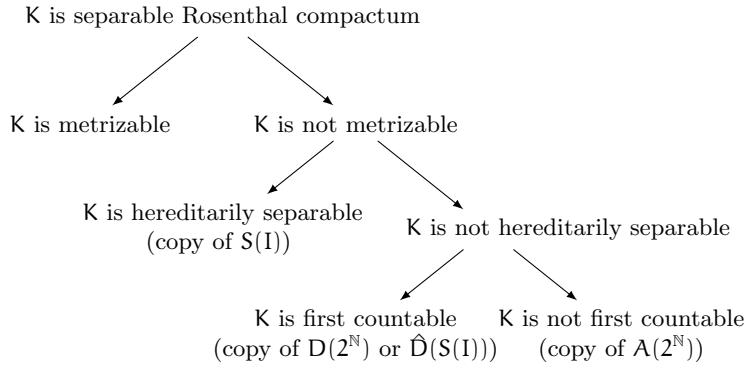
$$h_t(x) = \begin{cases} 0, & x < a_t \\ 1/2, & a_t \leq x \leq b_t \\ 1, & b_t < x \end{cases}$$

386    Let  $\hat{D}(S(2^{\mathbb{N}}))$  be the pointwise closure of the set  $\{h_t : t \in 2^{<\mathbb{N}}\}$ . Hence,  
 387     $\hat{D}(S(2^{\mathbb{N}}))$  is a separable first countable Rosenthal compactum which is not  
 388    hereditarily separable. In fact, it contains an uncountable discrete subspace  
 389    (see Theorem 5 in [Tod99]).

390    **Theorem 3.3** (Todorčević's Trichotomy, [Tod99], Theorem 3 in [ADK08]). *Let  $K$*   
 391    *be a separable Rosenthal Compactum.*

- 392    (i) *If  $K$  is hereditarily separable but non-metrizable, then  $S(2^{\mathbb{N}})$  embeds into  $K$ .*  
 393    (ii) *If  $K$  is first countable but not hereditarily separable, then either  $D(2^{\mathbb{N}})$  or*  
 394     $\hat{D}(S(2^{\mathbb{N}}))$  *embeds into  $K$ .*  
 395    (iii) *If  $K$  is not first countable, then  $A(2^{\mathbb{N}})$  embeds into  $K$ .*

396    In other words, we have the following classification:



397

398 Lastly, the definitions provided here for  $NIP_i$  ( $i = 1, 2, 3$ ) are topological.

399 **Question 3.4.** Is there a non-topological characterization for  $NIP_i$ ,  $i = 1, 2, 3$ ?

400 More can be said about the nature of the embeddings in Todorčević's Trichotomy.  
 401 Given a separable Rosenthal compactum  $K$ , there is typically more than one countable dense subset of  $K$ . We can view a separable Rosenthal compactum as the accumulation points of a countable family of pointwise bounded real-valued functions.  
 402 The choice of the countable families is not important when a bijection between them can be lifted to a homeomorphism of their closures. To be more precise:

406 **Definition 3.5.** Given a Polish space  $X$ , a countable set  $I$  and two pointwise bounded families  $\{f_i : i \in I\} \subseteq \mathbb{R}^X$ ,  $\{g_i : i \in I\} \subseteq \mathbb{R}^X$  indexed by  $I$ . We say that  $\{f_i : i \in I\}$  and  $\{g_i : i \in I\}$  are *equivalent* if and only if the map  $f_i \mapsto g_i$  is extended to a homeomorphism from  $\overline{\{f_i : i \in I\}}$  to  $\overline{\{g_i : i \in I\}}$ .

410 Notice that in the separable examples discussed before  $(\hat{A}(2^N), S(2^N)$  and  $\hat{D}(S(2^N))$ ),  
 411 the countable dense subsets are indexed by the binary tree  $2^{<\mathbb{N}}$ . This choice of index is useful because the Ramsey theory of perfect subsets of the Cantor space  $2^{\mathbb{N}}$   
 413 can be imported to analyze the behavior of the accumulation points. Since  $2^{<\mathbb{N}}$  is  
 414 countable, we can always choose this index for the countable dense subsets. This  
 415 is done in [ADK08].

416 **Definition 3.6.** Given a Polish space  $X$  and a pointwise bounded family  $\{f_t : t \in 2^{<\mathbb{N}}\}$ . We say that  $\{f_t : t \in 2^{<\mathbb{N}}\}$  is *minimal* if and only if for every dyadic subtree  $\{s_t : t \in 2^{<\mathbb{N}}\}$  of  $2^{<\mathbb{N}}$ ,  $\{f_{s_t} : t \in 2^{<\mathbb{N}}\}$  is equivalent to  $\{f_t : t \in 2^{<\mathbb{N}}\}$ .

419 One of the main results in [ADK08] is that there are (up to equivalence) seven  
 420 minimal families of Rosenthal compacta and that for every relatively compact  $\{f_t : t \in 2^{<\mathbb{N}}\} \subseteq B_1(X)$  there is a dyadic subtree  $\{s_t : t \in 2^{<\mathbb{N}}\}$  such that  $\{f_{s_t} : t \in 2^{<\mathbb{N}}\}$   
 422 is equivalent to one of the minimal families. We shall describe the minimal families  
 423 next. We will follow the same notation as in [ADK08]. For any node  $t \in 2^{<\mathbb{N}}$ , we  
 424 denote by  $t^\frown 0^\infty$  ( $t^\frown 1^\infty$ ) the infinite binary sequence starting with  $t$  and ending  
 425 in 0's (1's). Fix a regular dyadic subtree  $R = \{s_t : t \in 2^{<\mathbb{N}}\}$  of  $2^{<\mathbb{N}}$  (i.e., a dyadic  
 426 subtree such that every level of  $R$  is contained in a level of  $2^{<\mathbb{N}}$ ) with the property  
 427 that for all  $s, s' \in R$ ,  $s^\frown 0^\infty \neq s'^\frown 0^\infty$  and  $s^\frown 1^\infty \neq s'^\frown 1^\infty$ . Given  $t \in 2^{<\mathbb{N}}$ , let  
 428  $v_t$  be the characteristic function of the set  $\{x \in 2^{\mathbb{N}} : x \text{ extends } t\}$ . Let  $<$  be the  
 429 lexicographic order in  $2^{\mathbb{N}}$ . Given  $a \in 2^{\mathbb{N}}$ , let  $f_a^+ : 2^{\mathbb{N}} \rightarrow \{0, 1\}$  be the characteristic  
 430 function of  $\{x \in 2^{\mathbb{N}} : a \leq x\}$  and let  $f_a^- : 2^{\mathbb{N}} \rightarrow \{0, 1\}$  be the characteristic function of

431     $\{x \in 2^{\mathbb{N}} : a < x\}$ . Given two maps  $f, g : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  we denote by  $(f, g) : 2^{\mathbb{N}} \sqcup 2^{\mathbb{N}} \rightarrow \mathbb{R}$   
 432    the function which is  $f$  on the first copy of  $2^{\mathbb{N}}$  and  $g$  on the second copy of  $2^{\mathbb{N}}$ .

- 433    (1)  $D_1 = \{\frac{1}{|t|+1}v_t : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_1} = A(2^{\mathbb{N}})$ .
- 434    (2)  $D_2 = \{s_t^\frown 0^\infty : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_2} = 2^{\leq\mathbb{N}}$ .
- 435    (3)  $D_3 = \{f_{s_t^\frown 0^\infty}^+ : t \in 2^{<\mathbb{N}}\}$ . This is a discrete in  $\overline{D_3} = S(2^{\mathbb{N}})$ .
- 436    (4)  $D_4 = \{f_{s_t^\frown 1^\infty}^- : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_4} = S(2^{\mathbb{N}})$ .
- 437    (5)  $D_5 = \{v_t : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_5} = \hat{A}(2^{\mathbb{N}})$ .
- 438    (6)  $D_6 = \{(v_{s_t}, s_t^\frown 0^\infty) : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_6} = \hat{D}(2^{\mathbb{N}})$ .
- 439    (7)  $D_7 = \{(v_{s_t}, x_{s_t^\frown 0^\infty}^+) : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_7} = \hat{D}(S(2^{\mathbb{N}}))$

440    **Theorem 3.7** (Heptacotomy of minimal families, Theorem 2 in [ADK08]). *Let  
 441     $X$  be Polish. For every relatively compact  $\{f_t : t \in 2^{<\mathbb{N}}\} \subseteq B_1(X)$ , there exists  $i =  
 442    1, 2, \dots, 7$  and a regular dyadic subtree  $\{s_t : t \in 2^{<\mathbb{N}}\}$  of  $2^{<\mathbb{N}}$  such that  $\{f_{s_t} : t \in 2^{<\mathbb{N}}\}$   
 443    is equivalent to  $D_i$ . Moreover, all  $D_i$  are minimal and mutually non-equivalent.*

444    **3.2. NIP and definability by universally measurable functions.** We now  
 445    turn to the question: what happens when  $\mathcal{P}$  is uncountable? Notice that the  
 446    countability assumption is crucial in the proof of Theorem 2.9 essentially because it  
 447    makes  $\mathbb{R}^{\mathcal{P}}$  a Polish space. For the uncountable case, we may lose Baire-1 definability  
 448    so we shall replace  $B_1(X)$  by a bigger class. Recall that the purpose of studying the  
 449    class of Baire-1 functions is that a pointwise limit of continuous functions is not  
 450    necessarily continuous. In [BFT78], J. Bourgain, D.H. Fremlin and M. Talagrand  
 451    characterized the Non-Independence Property of a set of continuous functions with  
 452    various notions of compactness in function spaces containing  $C(X)$ , such as  $B_1(X)$ .  
 453    In this section we will replace  $B_1(X)$  with the larger space  $M_r(X)$  of universally  
 454    measurable functions. The development of this section is based on Theorem 2F in  
 455    [BFT78]. We now give the relevant definitions. Readers with little familiarity with  
 456    measure theory can review the appendix for standard definitions appearing in this  
 457    subsection.

458    Given a Hausdorff space  $X$  and a measurable space  $(Y, \Sigma)$ , we say that  $f : X \rightarrow Y$   
 459    is *universally measurable* (with respect to  $\Sigma$ ) if  $f^{-1}(E)$  is universally measurable  
 460    for every  $E \in \Sigma$ , i.e.,  $f^{-1}(E)$  is  $\mu$ -measurable for every Radon probability measure  
 461     $\mu$  on  $X$ . When  $Y = \mathbb{R}$  we will always take  $\Sigma = \mathcal{B}(\mathbb{R})$ , the Borel  $\sigma$ -algebra of  $\mathbb{R}$ .  
 462    In that case, a function  $f : X \rightarrow \mathbb{R}$  is universally measurable if and only if  $f^{-1}(U)$   
 463    is  $\mu$ -measurable for every Radon probability measure  $\mu$  on  $X$  and every open set  
 464     $U \subseteq \mathbb{R}$ . Following [BFT78], the collection of all universally measurable real-valued  
 465    functions will be denoted by  $M_r(X)$ . In the context of deep computations, we will  
 466    be interested in transition maps from a state space  $L \subseteq \mathbb{R}^{\mathcal{P}}$  to itself. There are two  
 467    natural  $\sigma$ -algebras one can consider in the product space  $\mathbb{R}^{\mathcal{P}}$ : the Borel  $\sigma$ -algebra,  
 468    i.e., the  $\sigma$ -algebra generated by open sets in  $\mathbb{R}^{\mathcal{P}}$ ; and the cylinder  $\sigma$ -algebra, i.e.,  
 469    the  $\sigma$ -algebra generated by Borel cylinder sets or equivalently basic open sets in  
 470     $\mathbb{R}^{\mathcal{P}}$ . Note that when  $\mathcal{P}$  is countable, both  $\sigma$ -algebras coincide but in general the  
 471    cylinder  $\sigma$ -algebra is strictly smaller. We will use the cylinder  $\sigma$ -algebra to define  
 472    universally measurable maps  $f : \mathbb{R}^{\mathcal{P}} \rightarrow \mathbb{R}^{\mathcal{P}}$ . The reason for this choice is because of  
 473    the following characterization:

474    **Lemma 3.8.** *Let  $X$  be a Hausdorff space and  $Y = \prod_{i \in I} Y_i$  be any product of  
 475    measurable spaces  $(Y_i, \Sigma_i)$  for  $i \in I$ . Let  $\Sigma_Y$  be the cylinder  $\sigma$ -algebra generated by  
 476    the measurable spaces  $(Y_i, \Sigma_i)$ . Let  $f : X \rightarrow Y$ . The following are equivalent:*

- 477        (i)  $f : X \rightarrow Y$  is universally measurable (with respect to  $\Sigma_Y$ ).  
 478        (ii)  $\pi_i \circ f : X \rightarrow Y_i$  is universally measurable (with respect to  $\Sigma_i$ ) for all  $i \in I$ .

479        *Proof.* (i) $\Rightarrow$ (ii) is clear since the projection maps  $\pi_i$  are measurable and the com-  
 480        position of measurable functions is measurable. To prove (ii) $\Rightarrow$ (i), suppose that  
 481         $C = \prod_{i \in I} C_i$  is a measurable cylinder and let  $J$  be the finite set of  $i \in I$  such that  
 482         $C_i \neq Y_i$ . Then,  $C = \bigcap_{i \in J} \pi_i^{-1}(C_i)$  so  $f^{-1}(C) = \bigcap_{i \in J} (\pi_i \circ f)^{-1}(C_i)$  is a universally  
 483        measurable set by assumption.  $\square$

484        The previous lemma says that a transition map is universally measurable if and  
 485        only if it is universally measurable on all its features. In other words, we can check  
 486        measurability of a transition just by checking measurability in all its features. We  
 487        will denote by  $M_r(X, \mathbb{R}^P)$  the collection of all universally measurable functions  
 488         $f : X \rightarrow \mathbb{R}^P$  (with respect to the cylinder  $\sigma$ -algebra), endowed with the topology of  
 489        pointwise convergence.

490        **Definition 3.9.** Let  $(L, \mathcal{P}, \Gamma)$  be a CCS. We say that a transition  $f : L \rightarrow L$  is  
 491        *universally measurable shard-definable* if and only if there exists  $\tilde{f} : \mathcal{L}_{sh} \rightarrow \mathcal{L}_{sh}$   
 492        extending  $f$  such that for every sizer  $r_\bullet$  there is a sizer  $s_\bullet$  such that the restriction  
 493         $\tilde{f}|_{\mathcal{L}[r_\bullet]} : \mathcal{L}[r_\bullet] \rightarrow \mathcal{L}[s_\bullet]$  is universally measurable, i.e.  $\pi_P \circ \tilde{f}|_{\mathcal{L}[r_\bullet]} : \mathcal{L}[r_\bullet] \rightarrow [-s_P, s_P]$   
 494        is  $\mu$ -measurable for every Radon probability measure  $\mu$  on  $\mathcal{L}[r_\bullet]$ .

495        We will need the following result about NIP and universally measurable func-  
 496        tions:

497        **Theorem 3.10** (Bourgain-Fremlin-Talagrand, Theorem 2F in [BFT78]). *Let  $X$  be a  
 498        Hausdorff space and  $A \subseteq C(X)$  be pointwise bounded. The following are equivalent:*

- 499        (i)  $\overline{A} \subseteq M_r(X)$ .  
 500        (ii) For every compact  $K \subseteq X$ ,  $A|_K$  has the NIP.  
 501        (iii) For every Radon measure  $\mu$  on  $X$ ,  $A$  is relatively countably compact in  
 502         $\mathcal{L}^0(X, \mu)$ , i.e., every countable subset of  $A$  has an accumulation point in  
 503         $\mathcal{L}^0(X, \mu)$ .

504        Theorem 2.5 immediately yields the following.

505        **Theorem 3.11.** *Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom. Let  $R$   
 506        be an exhaustive collection of sizers. Let  $\Delta \subseteq \Gamma$  be  $R$ -confined. If  $\pi_P \circ \Delta|_{L[r_\bullet]}$  has  
 507        the NIP for all  $P \in \mathcal{P}$  and all  $r_\bullet \in R$ , then every deep computation is universally  
 508        measurable shard-definable.*

509        *Proof.* By the Extendibility Axiom, Theorem 2.5 and lemma 2.10 we have that  
 510         $\pi_P \circ \tilde{\Delta}|_{\mathcal{L}[r_\bullet]} \subseteq M_r(\mathcal{L}[r_\bullet])$  for all  $r_\bullet \in R$  and  $P \in \mathcal{P}$ . Let  $f \in \overline{\Delta}$  be a deep computation.  
 511        Write  $f = \mathcal{U} \lim_i \gamma_i$  as an ultralimit of computations in  $\Delta$ . Define  $\tilde{f} := \mathcal{U} \lim_i \tilde{\gamma}_i$ .  
 512        Then, for all  $r_\bullet \in R$  and  $P \in \mathcal{P}$   $\pi_P \circ \tilde{\gamma}_i|_{\mathcal{L}[r_\bullet]} \in M_r(\mathcal{L}[r_\bullet])$  for all  $i$  so  $\pi_P \circ f|_{\mathcal{L}[r_\bullet]} \in$   
 513         $\pi_P \circ \tilde{\Delta}|_{\mathcal{L}[r_\bullet]} \subseteq M_r(\mathcal{L}[r_\bullet])$ .  $\square$

514        **Question 3.12.** Under the same assumptions of the previous Theorem, suppose  
 515        that every deep computation of  $\Delta$  is universally measurable shard-definable. Must  
 516         $\pi_P \circ \Delta|_{L[r_\bullet]}$  have the NIP for all  $P \in \mathcal{P}$  and all  $r_\bullet \in R$ ?

517 **3.3. Talagrand stability and definability by universally measurable func-**  
 518 **tions.** There is another notion closely related to NIP, introduced by Talagrand  
 519 in [Tal84] while studying Pettis integration. Suppose that  $X$  is a compact Haus-  
 520 dorff space and  $A \subseteq \mathbb{R}^X$ . Let  $\mu$  be a Radon probability measure on  $X$ . Given a  
 521  $\mu$ -measurable set  $E \subseteq X$ , a positive integer  $k$  and real numbers  $a < b$ . we write:

$$D_k(A, E, a, b) = \bigcup_{f \in A} \{x \in E^{2k} : f(x_{2i}) \leq a, f(x_{2i+1}) \geq b \text{ for all } i < k\}$$

522 We say that  $A$  is *Talagrand  $\mu$ -stable* if and only if for every  $\mu$ -measurable  
 523 set  $E \subseteq X$  of positive measure and for every  $a < b$  there is  $k \geq 1$  such that  
 524  $(\mu^{2k})^*(D_k(A, E, a, b)) < (\mu(E))^{2k}$ . Notice that we work with the outer measure  
 525 because it is not necessarily true that the sets  $D_k(A, E, a, b)$  are  $\mu$ -measurable.  
 526 This is certainly the case when  $A$  is a countable set of continuous (or  $\mu$ -measurable)  
 527 functions.

528 The following lemma establishes that Talagrand stability is a way to ensure that  
 529 deep computations are definable by measurable functions. We include the proof for  
 530 the reader's convenience.

531 **Lemma 3.13.** *If  $A$  is Talagrand  $\mu$ -stable, then  $\bar{A}$  is also Talagrand  $\mu$ -stable and*  
 532  $\bar{A} \subseteq \mathcal{L}^0(X, \mu)$ .

533 *Proof.* First, observe that a subset of a  $\mu$ -stable set is  $\mu$ -stable. To show that  $\bar{A}$   
 534 is  $\mu$ -stable, observe that  $D_k(\bar{A}, E, a, b) \subseteq D_k(A, E, a', b')$  where  $a < a' < b' <$   
 535  $b$  and  $E$  is a  $\mu$ -measurable set with positive measure. It suffices to show that  
 536  $\bar{A} \subseteq \mathcal{L}^0(X, \mu)$ . Suppose that there exists  $f \in \bar{A}$  such that  $f \notin \mathcal{L}^0(X, \mu)$ . By a  
 537 characterization of measurable functions (see 413G in [Fre03]), there exists a  $\mu$ -  
 538 measurable set  $E$  of positive measure and  $a < b$  such that  $\mu^*(P) = \mu^*(Q) = \mu(E)$   
 539 where  $P = \{x \in E : f(x) \leq a\}$  and  $Q = \{x \in E : f(x) \geq b\}$ . Then, for any  $k \geq 1$ :  
 540  $(P \times Q)^k \subseteq D_k(\{f\}, E, a, b)$  so  $(\mu^{2k})^*(D_k(\{f\}, E, a, b)) = (\mu^*(P)\mu^*(Q))^k = (\mu(E))^{2k}$ .  
 541 Thus,  $\{f\}$  is not  $\mu$ -stable, but we argued before that a subset of a  $\mu$ -stable set must  
 542 be  $\mu$ -stable.  $\square$

543 We say that  $A$  is *universally Talagrand stable* if  $A$  is Talagrand  $\mu$ -stable for  
 544 every Radon probability measure  $\mu$  on  $X$ . A similar argument as before, yields the  
 545 following:

546 **Theorem 3.14.** *Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom. If*  
 547  $\pi_P \circ \Delta|_{L[r_\bullet]}$  *is universally Talagrand stable for all  $P \in \mathcal{P}$  and all sizers  $r_\bullet$ , then*  
 548 *every deep computation is universally measurable sh-definable.*

549 It is then natural to ask: what is the relationship between Talagrand stability  
 550 and the NIP? The following dichotomy will be useful.

551 **Lemma 3.15** (Fremlin's Dichotomy, 463K in [Fre03]). *If  $(X, \Sigma, \mu)$  is a perfect  $\sigma$ -*  
 552 *finite measure space (in particular, for  $X$  compact and  $\mu$  a Radon probability measure*  
 553 *on  $X$ ) and  $\{f_n : n \in \mathbb{N}\}$  be a sequence of real-valued measurable functions on  $X$ , then*  
 554 *either:*

- 555 (i)  $\{f_n : n \in \mathbb{N}\}$  has a subsequence that converges  $\mu$ -almost everywhere, or
- 556 (ii)  $\{f_n : n \in \mathbb{N}\}$  has a subsequence with no  $\mu$ -measurable accumulation point in  
 557  $\mathbb{R}^X$ .

558     The preceding lemma can be considered as the measure theoretic version of  
 559     Rosenthal's Dichotomy. Combining this dichotomy with the Theorem 3.10 we get  
 560     the following result:

561     **Theorem 3.16.** *Let  $X$  be a Hausdorff space and  $A \subseteq C(X)$  be pointwise bounded.  
 562     The following are equivalent:*

- 563       (i)  $\overline{A} \subseteq M_r(X)$ .
- 564       (ii) *For every compact  $K \subseteq X$ ,  $A|_K$  has the NIP.*
- 565       (iii) *For every Radon measure  $\mu$  on  $X$ ,  $A$  is relatively countably compact in  
                    $L^0(X, \mu)$ , i.e., every countable subset of  $A$  has an accumulation point in  
                    $L^0(X, \mu)$ .*
- 566       (iv) *For every Radon measure  $\mu$  on  $X$  and every sequence  $\{f_n : n \in \mathbb{N}\} \subseteq A$ ,  
                   there is a subsequence that converges  $\mu$ -almost everywhere.*

570     *Proof.* Notice that the equivalence (i)-(iii) is Theorem 3.10. Notice that the equivalence of (iii) and (iv) is Fremlin's Dichotomy Theorem.  $\square$

572     **Lemma 3.17.** *Let  $X$  be a compact Hausdorff space and  $A \subseteq C(X)$  be pointwise bounded. If  $A$  is universally Talagrand stable, then  $A$  has the NIP.*

574     *Proof.* By Theorem 3.10, it suffices to show that  $A$  is relatively countably compact  
 575     in  $L^0(X, \mu)$  for all Radon probability measure  $\mu$  on  $X$ . Since  $A$  is Talagrand  $\mu$ -stable  
 576     for any such  $\mu$ , then  $\overline{A} \subseteq L^0(X, \mu)$ . In particular,  $A$  is relatively countably compact  
 577     in  $L^0(X, \mu)$ .  $\square$

578     **Question 3.18.** Is the converse true?

579     There is a delicate point in this question, as it may be sensitive to set-theoretic  
 580     axioms (even assuming countability of  $A$ ).

581     **Theorem 3.19** (Talagrand, Theorem 9-3-1(a) in [Tal84]). *Let  $X$  be a compact  
 582     Hausdorff space and  $A \subseteq M_r(X)$  be countable and pointwise bounded. Assume that  
 583      $[0, 1]$  is not the union of  $< c$  closed measure zero sets. If  $A$  has the NIP, then  $A$  is  
 584     universally Talagrand stable.*

585     **Theorem 3.20** (Fremlin, Shelah, [FS93]). *It is consistent that there exists a countable  
 586     pointwise bounded set of Lebesgue measurable functions with the NIP which is  
 587     not Talagrand stable with respect to Lebesgue measure.*

## 588 APPENDIX: MEASURE THEORY

589     Given a set  $X$ , a collection  $\Sigma$  of subsets of  $X$  is called a  $\sigma$ -algebra if  $\Sigma$  contains  
 590      $X$  and is closed under complements and countable unions. Hence, for example, a  
 591      $\sigma$ -algebra is also closed under countable intersections. Intuitively, a  $\sigma$ -algebra is  
 592     a collection of sets in which we can define a  $\sigma$ -additive measure. We call sets in  
 593     a  $\sigma$ -algebra  $\Sigma$  measurable sets and the pair  $(X, \Sigma)$  a measurable space. If  $X$  is a  
 594     topological space, there is a natural  $\sigma$ -algebra of subsets of  $X$ , namely the Borel  
 595      $\sigma$ -algebra  $\mathcal{B}(X)$ , i.e., the smallest  $\sigma$ -algebra containing all open subsets of  $X$ . Given  
 596     two measurable spaces  $(X, \Sigma_X)$  and  $(Y, \Sigma_Y)$ , we say that a function  $f : X \rightarrow Y$  is  
 597     measurable if and only if  $f^{-1}(E) \in \Sigma_X$  for every  $E \in \Sigma_Y$ . In particular, we say that  
 598      $f : X \rightarrow \mathbb{R}$  is measurable if  $f^{-1}(E) \in \Sigma_X$  for all  $E \in \mathcal{B}(\mathbb{R})$  (equivalently,  $E$  open in  
 599      $\mathbb{R}$ ).

Given a measurable space  $(X, \Sigma)$ , a  $\sigma$ -additive measure is a non-negative function  $\mu : \Sigma \rightarrow \mathbb{R}$  with the property that  $\mu(\emptyset) = 0$  and  $\mu(\bigcup_{n=0}^{\infty} A_n) = \sum_{n=0}^{\infty} \mu(A_n)$  whenever  $\{A_n : n \in \mathbb{N}\} \subseteq \Sigma$  is pairwise disjoint. We call  $(X, \Sigma, \mu)$  a *measure space*. A  $\sigma$ -additive measure is called a *probability measure* if  $\mu(X) = 1$ . A measure  $\mu$  is *complete* if for every  $A \subseteq B \in \Sigma$ ,  $\mu(B) = 0$  implies  $A \in \Sigma$ . In words, subsets of measure-zero sets are always measurable (and hence, by the monotonicity of  $\mu$ , have measure zero as well). A measure  $\mu$  is  $\sigma$ -finite if  $X = \bigcup_{n=1}^{\infty} X_n$  where  $\mu(X_n) < \infty$  for all  $n \in \mathbb{N}$  (i.e.,  $X$  can be decomposed into countably many finite measure sets). A measure  $\mu$  is *perfect* if for every measurable  $f : X \rightarrow \mathbb{R}$  and every measurable set  $E$  with  $\mu(E) > 0$ , there exists a compact  $K \subseteq f(E)$  such that  $\mu(f^{-1}(K)) > 0$ . We say that a property  $\phi(x)$  about  $x \in X$  holds  $\mu$ -almost everywhere if  $\mu(\{x \in X : \phi(x) \text{ does not hold}\}) = 0$ .

A special example of the preceding concepts is that of a *Radon measure*. If  $X$  is a Hausdorff topological space, then a measure  $\mu$  on the Borel sets of  $X$  is called a *Radon measure* if

- for every open set  $U$ ,  $\mu(U)$  is the supremum of  $\mu(K)$  over all compact  $K \subseteq U$ , that is, the measure of open sets may be approximated via compact sets; and
- every point of  $X$  has a neighborhood  $U \ni x$  for which  $\mu(U)$  is finite.

Perhaps the most famous example of a Radon measure on  $\mathbb{R}$  is the Lebesgue measure of Borel sets. If  $X$  is finite,  $\mu(A) := |A|$  (the cardinality of  $A$ ) defines a Radon measure on  $X$ . Every Radon measure is perfect (see 451A, 451B and 451C in [Fre03]).

While not immediately obvious, sets can be measurable according to one measure, but non-measurable according to another. Given a measure space  $(X, \Sigma, \mu)$  we say that a set  $E \subseteq X$  is  $\mu$ -measurable if there are  $A, B \in \Sigma$  such that  $A \subseteq E \subseteq B$  and  $\mu(B \setminus A) = 0$ . The set of all  $\mu$ -measurable sets is a  $\sigma$ -algebra containing  $\Sigma$  and it is denoted by  $\Sigma_{\mu}$ . A set  $E \subseteq X$  is *universally measurable* if it is  $\mu$ -measurable for every Radon probability measure on  $X$ . It follows that Borel sets are universally measurable. We say that  $f : X \rightarrow \mathbb{R}$  is  $\mu$ -measurable if  $f^{-1}(E) \in \Sigma_{\mu}$  for all  $E \in \mathcal{B}(\mathbb{R})$  (equivalently,  $E$  open in  $\mathbb{R}$ ). The set of all  $\mu$ -measurable functions is denoted by  $\mathcal{L}^0(X, \mu)$ .

Recall that if  $\{X_i : i \in I\}$  is a collection of topological spaces indexed by some set  $I$ , then the product space  $X := \prod_{i \in I} X_i$  is endowed with the topology generated by *cylinders*, that is, sets of the form  $\prod_{i \in I} U_i$  where each  $U_i$  is open in  $X_i$ , and  $U_i = X_i$  except for finitely many indices  $i \in I$ . If each space is measurable, say we pair  $X_i$  with a  $\sigma$ -algebra  $\Sigma_i$ , then there are multiple ways to interpret the product space  $X$  as a measurable space, but the interpretation we care about in this paper is the so called *cylinder  $\sigma$ -algebra*, as used in Lemma 3.8. Namely, let  $\Sigma$  be the  $\sigma$ -algebra generated by sets of the form

$$\prod_{i \in I} C_i, \quad C_i \in \Sigma_i, \quad C_i = X_i \text{ for all but finitely many } i \in I.$$

We remark that when  $I$  is uncountable and  $\Sigma_i = \mathcal{B}(X_i)$  for all  $i \in I$ , then  $\Sigma$  is, in general, strictly **smaller** than  $\mathcal{B}(X)$ .

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