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# COMPLEXITY OF DEEP COMPUTATIONS VIA TOPOLOGY OF FUNCTION SPACES

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ABSTRACT. This paper revisits and extends a bridge between functional analysis and model theory, emphasizing its relevance to the theoretical foundations of machine learning. We show that the compactness behavior of families of Baire class 1 functions mirrors the learnability conditions in the sense of *Probably Approximately Correct* (PAC) learning, and that the failure of compactness corresponds to the presence of infinite *Vapnik-Chervonenkis* (VC) dimension. From this perspective, Rosenthal compacta emerge as the natural topological counterpart of PAC-learnable concept classes, while NIP vs. IP structures capture the precise boundary between analytical regularity and combinatorial intractability. These parallels suggest a unified framework linking compactness, definability, and learnability, exemplifying how the topology of function spaces encodes the algorithmic and epistemic limits of prediction.

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## 1. INTRODUCTION

8      In this paper we study limit behavior of real-valued computations as the value  
9 of certain parameters of the computation model tend towards infinity, or towards  
10 zero, or towards some other fixed value, e.g., the depth of a neural network tending  
11 to infinity, or the time interval between layers of the network tending toward zero.  
12 Recently, particular cases of this situation have attracted considerable attention  
13 in deep learning research (e.g., Neural Ordinary Differential Equations [CRBD],  
14 Physics-Informed Neural Networks [RPK19], deep equilibrium models [BKK], etc).  
15 In this paper, we combine ideas of topology and model theory to study these limit  
16 phenomena from a unified viewpoint.  
17      Informed by model theory, to each computation in a given computation model,  
18 we associate a continuous real-valued function, called the *type* of the computation,  
19 that describes the logical properties of this computation with respect to the rest of  
20 the model. This allows us to view computations in any given computational model  
21 as elements of a space of real-valued functions, which is called the *space of types*  
22 of the model. The idea of embedding models of theories into their type spaces is  
23 central in model theory. The embedding of computations into spaces of types allows  
24 us to utilize the vast theory of topology of function spaces, known as  $C_p$ -theory,  
25 to obtain results about complexity of topological limits of computations. As we  
26 shall indicate next, recent classification results for spaces of functions provide an  
27 elegant and powerful machinery to classify computations according to their levels  
28 of “tameness” or “wildness”, with the former corresponding roughly to polynomial  
29 approximability and the latter to exponential approximability. The viewpoint  
30 of spaces of types, which we have borrowed from model theory, thus becomes a  
31 “Rosetta stone” that allows us to interconnect various classification programs: In

topology, the classification of Rosenthal compacta pioneered by Todorćević [Tod99]; in logic, the classification of theories developed by Shelah [She90]; and in statistical learning, the notion PAC learning and VC dimension pioneered by Vapkins and Chervonenkis [VC74, VC71].

In a previous paper [ADIW24], we introduced the concept of limits of computations, which we called *ultracomputations* (given they arise as ultrafilter limits of standard computations) and *deep computations* (following usage in machine learning [BKK]). There is a technical difference between both designations, but in this paper, to simplify the nomenclature, we will ignore the difference and use only the term “deep computation”.

In [ADIW24], we proved a new “tame vs wild” (i.e., polynomial vs exponential) dichotomy for complexity of deep computations by invoking a classical result of Grothendieck from late 50s [Gro52]. Under our model-theoretic Rosetta stone, polynomial approximability in the sense of computation becomes identified with the notion of continuous extendability in the sense of topology, and with the notions of *stability* and *type definability* in the sense of model theory.

In this paper, we follow a more general approach, i.e., we view deep computations as pointwise limits of continuous functions. In topology, real-valued functions that arise as the pointwise limit of a sequence of continuous are called *functions of the first Baire class*, or *Baire class 1* functions, or *Baire-1* for short; Baire class 1 form a step above simple continuity in the hierarchy of functions studied in real analysis (Baire class 0 functions being continuous functions). Intuitively, Baire-1 functions represent functions with “controlled” discontinuities, so they are crucial in topology and set theory.

We prove a new “tame vs wild” Ramsey-theoretic dichotomy for complexity of general deep computations by invoking a famous paper by Bourgain, Fremlin and Talagrand from the late 70s [BFT78], and a new trichotomy for the class of “tame” deep computations by invoking an equally celebrated result of Todorćević, from the late 90s, for functions of the first Baire class [Tod99].

Todorćević’s trichotomy regards *Rosenthal compacta*; these are special classes of topological spaces, defined as compact spaces that can be embedded (homeomorphically identified as a subset) within the space of Baire class 1 functions on some Polish (separable, complete metric) space, under the pointwise convergence topology. Rosenthal compacta exhibit “topological tameness,” meaning they behave in relatively controlled ways, and since the late 70’s, they have played a crucial role for understanding complexity of structures of functional analysis, especially, Banach spaces. Todorćević’s trichotomy has been utilized to settle longstanding problems in topological dynamics and topological entriopy [GM22].

Through our Rosetta stone, Rosenthal compacta in topology correspond to the important concept of “No Independence Property” (known as “NIP”) in model theory, identified by Shelah [She71, She90], and to the concept of Probably Approximately Correct learning (known as “PAC learnability”) in statistical learning theory identified by Valiant [Val84].

Going beyond Todorćević’s trichotomy, we invoke a more recent heptachotomy for Rosenthal compacta obtained by Argyros, Dodos and Kanellopoulos [ADK08]. Argyros, Dodos and Kanellopoulos identified the fundamental “prototypes” of separable Rosenthal compacta, and proved that any non-metrizable separable Rosenthal compactum must contain a “canonical” embedding of one of these prototypes. They

showed that if a separable Rosenthal compactum is not hereditarily separable, it must contain an uncountable discrete subspace of the size of the continuum.

We believe that the results presented in this paper show practitioners of computation, or topology, or descriptive set theory, or model theory, how classification invariants used in their field translate into classification invariants of other fields. However, in the interest of accessibility, we do not assume previous familiarity with high-level topology or model theory, or computing. The only technical prerequisite of the paper is undergraduate-level topology. The necessary topological background beyond undergraduate topology is covered in section 2.

Throughout the paper, we focus on classical computation; however, by refining the model-theoretic tools, the results presented here can be extended to quantum computation and open quantum systems. This extension will be addressed in a forthcoming paper.

## 2. GENERAL TOPOLOGICAL PRELIMINARIES

In this section we give preliminaries from general topology and function space theory. We include some of the proofs for completeness, but the reader familiar with these topics may skip them.

Recall that a subset of a topological space is  $F_\sigma$  if it is a countable union of closed sets, and  $G_\delta$  if it is a countable intersection of closed sets. Note that in a metrizable space, every open set is  $F_\sigma$ ; equivalently, every closed set is  $G_\delta$ .

A *Polish space* is a separable and completely metrizable topological space. The most important examples are the reals  $\mathbb{R}$ , the Cantor space  $2^\mathbb{N}$  (the set of all infinite binary sequences, endowed with the product topology), and the Baire space  $\mathbb{N}^\mathbb{N}$  (the set of all infinite sequences of naturals, also with the product topology). Countable products of Polish spaces are Polish; this includes spaces like  $\mathbb{R}^\mathbb{N}$ , the space of sequences of real numbers.

In this paper, we shall discuss subspaces, and so there is a pertinent subtlety of the definitions worth mentioning: *completely metrizable space* is not the same as *complete metric space*; for an illustrative example, the interval  $(0, 1)$  with the metric inherited from the reals is not complete, but it is Polish since that is homeomorphic to the real line. Being Polish is a topological property.

The following result is a cornerstone of descriptive set theory, closely tied to the work of Waław Sierpiński and Kazimierz Kuratowski, with proofs often built upon their foundations and formalized later, notably, involving Stefan Mazurkiewicz's work on complete metric spaces.

**Fact 2.1.** *A subset  $A$  of a Polish space  $X$  is itself Polish in the subspace topology if and only if it is a  $G_\delta$  set. In particular, closed subsets and open subsets of Polish spaces are also Polish spaces.*

Given two topological spaces  $X$  and  $Y$  we denote by  $C_p(X, Y)$  the set of all continuous functions  $f : X \rightarrow Y$  endowed with the topology of pointwise convergence. When  $Y = \mathbb{R}$ , we denote this collection simply as  $C_p(X)$ . A natural question is, how do topological properties of  $X$  translate to  $C_p(X)$  and vice versa? These questions, and in general the study of these spaces, are the concern of  $C_p$ -theory, an active field of research in general topology which was pioneered by A. V. Arhangel'skiĭ and his students in the 1970's and 1980's. This field has found many applications in model theory and functional analysis. Recent surveys on the topics include [HT23] and [Tka11].

A *Baire class 1* function between topological spaces is a function that can be expressed as the pointwise limit of a sequence of continuous functions. If  $X$  and  $Y$  are topological spaces, the Baire class 1 functions  $f : X \rightarrow Y$  endowed with the topology of pointwise convergence is denoted  $B_1(X, Y)$ . As above, in the special case  $Y = \mathbb{R}$  we denote  $B_1(X, Y)$  as  $B_1(X)$ . Clearly,  $C_p(X, Y) \subseteq B_1(X, Y)$ . The Baire hierarchy of functions was introduced by French mathematician René-Louis Baire in his 1899 doctoral thesis, *Sur les fonctions de variables réelles*. His work moved away from the 19th-century preoccupation with "pathological" functions toward a constructive classification based on pointwise limits.

A topological space  $X$  is *perfectly normal* if it is normal and every closed subset of  $X$  is a  $G_\delta$  (equivalently, every open subset of  $X$  is a  $G_\delta$ ). Note that every metrizable space is perfectly normal.

The following fact was established by Baire in thesis. A proof can be found in Section 10 of [Tod97].

**Fact 2.2** (Baire). *If  $X$  is perfectly normal, then the following conditions are equivalent for a function  $f : X \rightarrow \mathbb{R}$ :*

- $f$  is a Baire class 1 function, that is,  $f \in B_1(X)$ .
- $f^{-1}[U]$  is an  $F_\sigma$  subset of  $X$  whenever  $U \subseteq \mathbb{R}$  is open.
- $f$  is a pointwise limit of continuous functions.
- For every closed  $F \subseteq X$ , the restriction  $f|_F$  has a point of continuity.

Moreover, if  $X$  is Polish and  $f \notin B_1(X)$ , then there exists countable  $D_0, D_1 \subseteq X$  and reals  $a < b$  such that

$$D_0 \subseteq f^{-1}(-\infty, a], \quad D_1 \subseteq f^{-1}[b, \infty), \quad \overline{D_0} = \overline{D_1}.$$

A subset  $L$  of a topological space  $X$  is *relatively compact* in  $X$  if the closure of  $L$  in  $X$  is compact. Relatively compact subsets of  $B_1(X)$  (for  $X$  Polish space) have been objects of interest for researchers in Analysis and Topological Dynamics. We begin with the following well-known result. Recall that a set  $A \subseteq \mathbb{R}^X$  of real-valued functions is *pointwise bounded* if for every  $x \in X$  there is  $M_x > 0$  such that  $|f(x)| < M_x$  for all  $f \in A$ . We include a proof for the reader's convenience:

**Lemma 2.3.** *Let  $X$  be a Polish space and  $A \subseteq B_1(X)$  be pointwise bounded. The following are equivalent:*

- (i)  $A$  is relatively compact in  $B_1(X)$ .
- (ii)  $A$  is relatively countably compact in  $B_1(X)$ , i.e., every countable subset of  $A$  has an accumulation point in  $B_1(X)$ .
- (iii)  $\overline{A} \subseteq B_1(X)$ , where  $\overline{A}$  denotes the closure in  $\mathbb{R}^X$ .

*Proof.* Since  $A$  is pointwise bounded, for each  $x \in X$ , fix  $M_x > 0$  such that  $|f(x)| \leq M_x$  for every  $f \in A$ .

(i)  $\Rightarrow$  (ii) holds in general.

(ii)  $\Rightarrow$  (iii) Assume that  $A$  is relatively countably compact in  $B_1(X)$  and that  $f \in \overline{A} \setminus B_1(X)$ . By Fact 2.2, there are countable  $D_0, D_1 \subseteq X$  with  $\overline{D_0} = \overline{D_1}$ , and  $a < b$  such that  $D_0 \subseteq f^{-1}(-\infty, a]$  and  $D_1 \subseteq f^{-1}[b, \infty)$ . We claim that there is a sequence  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  such that  $\lim_{n \rightarrow \infty} f_n(x) = f(x)$  for all  $x \in D_0 \cup D_1$ . Indeed, use the countability to enumerate  $D_0 \cup D_1$  as  $\{x_n\}_{n \in \mathbb{N}}$ . Then for each positive  $n$  find  $f_n \in A$  with  $|f_n(x_i) - f(x_i)| < \frac{1}{n}$  for all  $i \leq n$ . The claim follows.

By relative countable compactness of  $A$ , there is an accumulation point  $g \in B_1(X)$  of  $\{f_n\}_{n \in \mathbb{N}}$ . It is straightforward to show that since  $f$  and  $g$  agree on  $D_0 \cup D_1$ ,

172  $g$  does not have a point of continuity on the closed set  $\overline{D_0} = \overline{D_1}$ , which contradicts  
 173 Fact 2.2.

174 (iii) $\Rightarrow$ (i) Suppose that  $\overline{A} \subseteq B_1(X)$ . Then  $\overline{A} \cap B_1(X) = \overline{A}$  is a closed subset of  
 175  $\prod_{x \in X} [-M_x, M_x]$ ; Tychonoff's theorem states that the product of compact spaces  
 176 is always compact, and since closed subsets of compact spaces are compact,  $\overline{A}$  must  
 177 be compact, as desired.  $\square$

178 **2.1. From Rosenthal's dichotomy to Shelah's NIP.** The fundamental idea  
 179 that connects the rich theory here presented to real-valued computations is the  
 180 concept of an *approximation*. In the reals, points of closure from some subset  
 181 can always be approximated by points inside the set, via a convergent sequence.  
 182 For more complicated spaces, such as  $C_p(X)$ , this fails in remarkable ways. To  
 183 see an example, consider the Cantor space  $X = 2^{\mathbb{N}}$ , and for each  $n \in \mathbb{N}$  define  
 184  $p_n : X \rightarrow \{0, 1\}$  by  $p_n(x) = x(n)$  for each  $x \in X$ . Then  $p_n$  is continuous for each  $n$ ,  
 185 but one can show (see Chapter 1.1 of [Tod97] for details) that the only continuous  
 186 functions in the closure of  $\{p_n\}_{n \in \mathbb{N}}$  are the functions  $p_n$  themselves; moreover, none  
 187 of the subsequences of  $\{p_n\}_{n \in \mathbb{N}}$  converge. In some sense, this example is the worst  
 188 possible scenario for convergence. The topological space obtained from this closure  
 189 is well-known: it is the *Stone-Ćech compactification* of the discrete space of natural  
 190 numbers, or  $\beta\mathbb{N}$  for short, and it is an important object of study in general topology.

191 The following theorem, established by Haskell Rosenthal in 1974, is fundamental  
 192 in functional analysis, and describes a sharp division in the behavior of sequences  
 193 within a Banach space:

194 **Theorem 2.4** (Rosenthal's Dichotomy, 1974). *If  $X$  is Polish and  $\{f_n\} \subseteq C_p(X)$   
 195 is pointwise bounded, then either  $\{f_n\}_{n \in \mathbb{N}}$  contains a convergent subsequence or a  
 196 subsequence whose closure (in  $\mathbb{R}^X$ ) is homeomorphic to  $\beta\mathbb{N}$ .*

197 In other words, a pointwise bounded set of continuous functions either contains  
 198 a convergent subsequence, or a subsequence whose closure is essentially the same as  
 199 the example mentioned in the previous paragraphs (the “wildest” possible scenario).  
 200 Note that in the preceding example, the functions are trivially pointwise bounded  
 201 in  $\mathbb{R}^X$  as the functions can only take values 0 and 1.

202 As we go from  $C_p(X)$  to the larger space  $B_1(X)$ , we find a similar dichotomy.  
 203 Either every point of closure of the set of functions will be a Baire class 1 function,  
 204 or there is a sequence inside the set that behaves in the wildest possible way. The  
 205 theorem is usually not phrased as a dichotomy but rather as an equivalence:

206 **Theorem 2.5** (Bourgain-Fremlin-Talagrand, Theorem 4G in [BFT78]). *Let  $X$  be  
 207 a Polish space and  $A \subseteq C_p(X)$  be pointwise bounded. The following are equivalent:*

- 208 (i)  $A$  is relatively compact in  $B_1(X)$ , i.e.,  $\overline{A} \subseteq B_1(X)$ .  
 (ii) For every  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and every  $a < b$  there is  $I \subseteq \mathbb{N}$  such that

$$\bigcap_{n \in I} f_n^{-1}(-\infty, a] \cap \bigcap_{n \notin I} f_n^{-1}[b, \infty) = \emptyset.$$

209 **Definition 2.6.** We shall say that a set  $A \subseteq \mathbb{R}^X$  has the *Independence Property*, or  
 210 IP for short, if it satisfies the following condition: There exists every  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$   
 211 and  $a < b$  such that for every pair of disjoint sets  $E, F \subseteq \mathbb{N}$ , we have

$$\bigcap_{n \in E} f_n^{-1}(-\infty, a] \cap \bigcap_{n \in F} f_n^{-1}[b, \infty) \neq \emptyset.$$

212 If  $A$  satisfies the negation of this condition, we will say that  $A$  *satisfies NIP*, or  
 213 that has the NIP.

*Remark 2.7.* Note that if  $X$  is compact and  $A \subseteq C_p(X)$ , then  $A$  has the NIP if and only if for every  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and for every  $a < b$  there is  $I \subseteq \mathbb{N}$  such that

$$\bigcap_{n \in I} f_n^{-1}(-\infty, a] \cap \bigcap_{n \notin I} f_n^{-1}[b, \infty) = \emptyset.$$

214 The Independence Property was first isolated by Saharon Shelah in model theory  
 215 as a dividing line between theories whose models are analyzable, or “tame” (corre-  
 216 sponding to NIP) theories of models are “unclassifiable” or “wild” ((corresponding  
 217 to IP). See [She71, She90].

218 **2.2. The special case  $Y = \mathbb{R}^{\mathcal{P}}$  with  $\mathcal{P}$  countable.** Our goal now is to characterize  
 219 relatively compact subsets of  $B_1(X, Y)$  for the particular case when  $Y = \mathbb{R}^{\mathcal{P}}$  with  $\mathcal{P}$   
 220 countable. Given  $P \in \mathcal{P}$  we denote the projection map onto the  $P$ -coordinate by  
 221  $\pi_P : \mathbb{R}^{\mathcal{P}} \rightarrow \mathbb{R}$ . From a high-level topological interpretation, the subsequent lemma  
 222 states that, in this context, the spaces  $\mathbb{R}$  and  $\mathbb{R}^{\mathcal{P}}$  are really not that different,  
 223 and that if we understand the Baire class 1 functions of one space, then we also  
 224 understand the functions of both.

225 **Lemma 2.8.** *Let  $X$  be a Polish space and  $\mathcal{P}$  be a countable set. Then,  $f \in B_1(X, \mathbb{R}^{\mathcal{P}})$   
 226 if and only if  $\pi_P \circ f \in B_1(X)$  for all  $P \in \mathcal{P}$ .*

*Proof.* Only one implication needs a proof. Suppose that  $\pi_P \circ f \in B_1(X)$  for all  $P \in \mathcal{P}$ . Let  $V$  be a basic open subset of  $\mathbb{R}^{\mathcal{P}}$ . That is, there exists a finite  $\mathcal{P}' \subseteq \mathcal{P}$  such that  $V = \bigcap_{P \in \mathcal{P}'} \pi_P^{-1}[U_P]$  where  $U_P$  is open in  $\mathbb{R}$ . Then,

$$f^{-1}[V] = \bigcap_{P \in \mathcal{P}'} (\pi_P \circ f)^{-1}[U_P]$$

227 is an  $F_\sigma$  set. Since  $\mathcal{P}$  is countable,  $\mathbb{R}^{\mathcal{P}}$  is second countable so every open set  $U$  in  
 228  $\mathbb{R}^{\mathcal{P}}$  is a countable union of basic open sets. Hence,  $f^{-1}[U]$  is  $F_\sigma$ .  $\square$

229 Below we consider  $\mathcal{P}$  with the discrete topology. For each  $f : X \rightarrow \mathbb{R}^{\mathcal{P}}$  denote  
 230  $\hat{f}(P, x) := \pi_P \circ f(x)$  for all  $(P, x) \in \mathcal{P} \times X$ . Similarly, for each  $g : \mathcal{P} \times X \rightarrow \mathbb{R}$  denote  
 231  $\check{g}(x)(P) := g(P, x)$ . Given  $A \subseteq (\mathbb{R}^{\mathcal{P}})^X$ , we denote  $\hat{A}$  as the set of all  $\hat{f}$  such that  
 232  $f \in A$ . Note that the map  $(\mathbb{R}^{\mathcal{P}})^X \rightarrow \mathbb{R}^{\mathcal{P} \times X}$  given by  $f \mapsto \hat{f}$  is a homeomorphism  
 233 and its inverse is given by  $g \mapsto \check{g}$ .

234 **Lemma 2.9.** *Let  $X$  be a Polish space and  $\mathcal{P}$  be countable. Then,  $f \in B_1(X, \mathbb{R}^{\mathcal{P}})$  if  
 235 and only if  $\hat{f} \in B_1(\mathcal{P} \times X)$ .*

*Proof.*  $(\Rightarrow)$  By Lemma 2.8, given an open set of reals  $U$ , we have  $f^{-1}[\pi_P^{-1}[U]]$  is  $F_\sigma$  for every  $P \in \mathcal{P}$ . Given that  $\mathcal{P}$  is a discrete countable space, we observe that

$$\hat{f}^{-1}[U] = \bigcup_{P \in \mathcal{P}} (\{P\} \times f^{-1}[\pi_P^{-1}[U]])$$

236 is an  $F_\sigma$  as well.

$(\Leftarrow)$  By lemma 2.8 it suffices to show that  $\pi_P \circ f \in B_1(X)$  for all  $P \in \mathcal{P}$ . Fix an open  $U \subseteq \mathbb{R}$ . Write  $\hat{f}^{-1}[U] = \bigcup_{n \in \mathbb{N}} F_n$  where  $F_n$  is closed in  $\mathcal{P} \times X$ . Then,

$$(\pi_P \circ f)^{-1}[U] = \bigcup_{n \in \mathbb{N}} \{x \in X : (P, x) \in F_n\}$$

237 which is  $F_\sigma$ . □

238 Given  $A \subseteq Y^X$  and  $K \subseteq X$  we write  $A|_K := \{f|_K : f \in A\}$ , i.e., the set of  
 239 all restrictions of functions in  $A$  to  $K$ . The following Theorem is a slightly more  
 240 general version of Theorem 2.5.

241 **Theorem 2.10.** *Assume that  $\mathcal{P}$  is countable,  $X$  is a Polish space, and  $A \subseteq C_p(X, \mathbb{R}^{\mathcal{P}})$   
 242 is such that  $\pi_P \circ A$  is pointwise bounded for all  $P \in \mathcal{P}$ . The following are equivalent  
 243 for every compact  $K \subseteq X$ :*

- 244 (1)  $\overline{A|_K} \subseteq B_1(K, \mathbb{R}^{\mathcal{P}})$ .  
 245 (2)  $\pi_P \circ A|_K$  has the NIP for every  $P \in \mathcal{P}$ .

*Proof.* (1) $\Rightarrow$ (2). Let  $P \in \mathcal{P}$ . Fix  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and  $a < b$ . By (1), we have  
 $\overline{A|_K} \subseteq B_1(K, \mathbb{R}^{\mathcal{P}})$ . Applying the homeomorphism  $f \mapsto \hat{f}$  and using lemma 2.9 we  
 get  $\hat{A}|_{\mathcal{P} \times K} \subseteq B_1(\mathcal{P} \times K)$ . By Theorem 2.5, there is  $I \subseteq \mathbb{N}$  such that

$$(\mathcal{P} \times K) \cap \bigcap_{n \in I} \hat{f}_n^{-1}(-\infty, a] \cap \bigcap_{n \notin I} \hat{f}_n^{-1}[b, \infty) = \emptyset$$

Hence,

$$K \cap \bigcap_{n \in I} (\pi_P \circ f_n)^{-1}(-\infty, a] \cap \bigcap_{n \notin I} (\pi_P \circ f_n)^{-1}[b, \infty) = \emptyset$$

By compactness, there are finite  $E \subseteq I$  and  $F \subseteq \mathbb{N} \setminus I$  such that

$$K \cap \bigcap_{n \in E} (\pi_P \circ f_n)^{-1}(-\infty, a] \cap \bigcap_{n \in F} (\pi_P \circ f_n)^{-1}[b, \infty) = \emptyset$$

246 Thus,  $\pi_P \circ A|_L$  has the NIP.

247 (2) $\Rightarrow$ (1) Fix  $f \in \overline{A|_K}$ . By lemma 2.8 it suffices to show that  $\pi_P \circ f \in B_1(K)$   
 248 for all  $P \in \mathcal{P}$ . By (2),  $\pi_P \circ A|_K$  has the NIP. Hence, by Theorem 2.5 we have  
 249  $\overline{\pi_P \circ A|_K} \subseteq B_1(K)$ . But then  $\pi_P \circ f \in \overline{\pi_P \circ A|_K} \subseteq B_1(K)$ . □

250 Lastly, a simple but significant result that helps understand the operation of  
 251 restricting a set of functions to a specific subspace of the domain space  $X$ , of course  
 252 in the context of the NIP, is that we may always assume that said subspace is  
 253 closed. Concretely, whether we take its closure or not has no effect on the NIP:

254 **Lemma 2.11.** *Assume that  $X$  is Hausdorff and that  $A \subseteq C_p(X)$ . The following  
 255 are equivalent for every  $L \subseteq X$ :*

- 256 (i)  $A_L$  has the NIP.  
 257 (ii)  $A|_{\overline{L}}$  has the NIP.

*Proof.* It suffices to show that (i) $\Rightarrow$ (ii). Suppose that (ii) does not hold, i.e., that  
 there are  $\{f_n\}_{n \in \mathbb{N}} \subseteq A$  and  $a < b$  such that for all finite disjoint  $E, F \subseteq \mathbb{N}$ :

$$\overline{L} \cap \bigcap_{n \in E} f_n^{-1}(-\infty, a] \cap \bigcap_{n \in F} f_n^{-1}[b, \infty) \neq \emptyset.$$

Pick  $a' < b'$  such that  $a < a' < b' < b$ . Then, for any finite disjoint  $E, F \subseteq \mathbb{N}$  we  
 can choose

$$x \in \overline{L} \cap \bigcap_{n \in E} f_n^{-1}(-\infty, a') \cap \bigcap_{n \in F} f_n^{-1}(b', \infty)$$

By definition of closure:

$$L \cap \bigcap_{n \in E} f_n^{-1}(-\infty, a'] \cap \bigcap_{n \in F} f_n^{-1}[b', \infty) \neq \emptyset.$$

258 This contradicts (i). □

### 259 3. NIP IN THE CONTEXT OF COMPOSITIONAL COMPUTATION STRUCTURES

260 In this section, we study what the NIP tell us in the context of deep compu-  
 261 tations as defined in [ADIW24]. We say a structure  $(L, \mathcal{P}, \Gamma)$  is a *Compositional*  
 262 *Computation Structure* (CCS) if  $L \subseteq \mathbb{R}^{\mathcal{P}}$  is a subspace of  $\mathbb{R}^{\mathcal{P}}$ , with the pointwise  
 263 convergence topology, and  $\Gamma \subseteq L^L$  is a semigroup under composition. The motiva-  
 264 tion for CCS comes from (continuous) model theory, where  $\mathcal{P}$  is a fixed collection  
 265 of predicates and  $L$  is a (real-valued) structure. Every point in  $L$  is identified with  
 266 its “type”, which is the tuple of all values the point takes on the predicates from  
 267  $\mathcal{P}$ , i.e., an element of  $\mathbb{R}^{\mathcal{P}}$ . In this context, elements of  $\mathcal{P}$  are called *features*. In the  
 268 discrete model theory framework, one views the space of complete-types as a sort of  
 269 compactification of the structure  $L$ . In this context, we don’t want to consider only  
 270 points in  $L$  (realized types) but in its closure  $\bar{L}$  (possibly unrealized types). The  
 271 problem is that the closure  $\bar{L}$  is not necessarily compact, an assumption that turns  
 272 out to be very useful in the context of continuous model theory. To bypass this  
 273 problem in a framework for deep computations, Alva, Dueñez, Iovino and Walton  
 274 introduced in [ADIW24] the concept of *shards*, which essentially consists in cover-  
 275 ing (a large fragment) of the space  $\bar{L}$  by compact, and hence pointwise-bounded,  
 276 subspaces (shards). We shall give the formal definition next.

277 A *sizer* is a tuple  $\mathbf{r}_{\bullet} = (r_p)_{p \in \mathcal{P}}$  of positive real numbers indexed by  $\mathcal{P}$ . Given a  
 278 sizer  $\mathbf{r}_{\bullet}$ , we define the  $\mathbf{r}_{\bullet}$ -*shard* as:

$$L[\mathbf{r}_{\bullet}] = L \cap \prod_{p \in \mathcal{P}} [-r_p, r_p]$$

279 For an illustrative example, we can frame Newton’s polynomial root approxima-  
 280 tion method in the context of a CCS (see Example 5.6 of [ADIW24] for details) as  
 281 follows. Begin by considering the extended complex numbers  $\hat{\mathbb{C}} := \mathbb{C} \cup \{\infty\}$  with  
 282 the usual Riemann sphere topology that makes it into a compact space (where  
 283 unbounded sequences converge to  $\infty$ ). In fact, not only is this space compact  
 284 but it is covered by the shard given by the sizer  $(1, 1, 1)$  (the unit sphere is con-  
 285 tained in the cube  $[-1, 1]^3$ ). The space  $\hat{\mathbb{C}}$  is homeomorphic to the usual unit  
 286 sphere  $S^2 := \{(x, y, z) : x^2 + y^2 + z^2 = 1\}$  of  $\mathbb{R}^3$ , by means of the stereographic  
 287 projection and its inverse  $\hat{\mathbb{C}} \rightarrow S^2$ . This function is regarded as a triple of pred-  
 288 icates  $x, y, z : \hat{\mathbb{C}} \rightarrow [-1, 1]$  where each will map an extended complex number to  
 289 its corresponding real coordinate on the cube  $[-1, 1]^3$ . Now fix the cubic com-  
 290 plex polynomial  $p(s) := s^3 - 1$ , and consider the map which performs one step  
 291 in Newton’s method at a particular (extended) complex number  $s$ , for finding  
 292 a root of  $p$ ,  $\gamma_p : \hat{\mathbb{C}} \rightarrow \hat{\mathbb{C}}$ . The explicit inner workings of  $\gamma_p$  are irrelevant for  
 293 this example, except for the fact that it is a continuous mapping. It follows that  
 294  $(S^2, \{x, y, z\}, \{\gamma_p^k : k \in \mathbb{N}\})$  is a CCS. The idea is that repeated applications of  
 295  $\gamma_p(s), \gamma_p \circ \gamma_p(s), \gamma_p \circ \gamma_p \circ \gamma_p(s), \dots$  would approximate a root of  $p$  provided  $s$  was  
 296 a good enough initial guess.



297 The  $\mathbf{r}_\bullet$ -type-shard is defined as  $\mathcal{L}[\mathbf{r}_\bullet] = \overline{\mathcal{L}[\mathbf{r}_\bullet]}$  and  $\mathcal{L}_{\text{sh}}$  is the union of all type-  
 298 shards. Notice that  $\mathcal{L}_{\text{sh}}$  is not necessarily equal to  $\mathcal{L} = \overline{\mathcal{L}}$ , unless  $\mathcal{P}$  is countable  
 299 (see [ADIW24]). A *transition* is a map  $f : \mathbf{L} \rightarrow \mathbf{L}$ , in particular, every element  
 300 in the semigroup  $\Gamma$  is a transition (these are called *realized computations*). In  
 301 practice, one would like to work with “definable” computations, i.e., ones that can  
 302 be described by a computer. In this topological framework, being continuous is an  
 303 expected requirement. However, as in the case of complete-types in model theory,  
 304 we will work with “unrealized computations”, i.e., maps  $f : \mathcal{L}_{\text{sh}} \rightarrow \mathcal{L}_{\text{sh}}$ . Note that  
 305 continuity of a computation does not imply that it can be continuously extended  
 306 to  $\mathcal{L}_{\text{sh}}$ .

307 Suppose that a transition map  $f : \mathbf{L} \rightarrow \mathcal{L}$  can be extended continuously to  
 308 a map  $\mathcal{L} \rightarrow \mathcal{L}$ . Then, the Stone-Weierstrass theorem implies that the feature  
 309  $\pi_{\mathbf{P}} \circ f$  (here  $\mathbf{P}$  is a fixed predicate, and the feature is hence continuous) can be  
 310 uniformly approximated by polynomials on the compact set  $\mathcal{L}[\mathbf{r}_\bullet]$ . Theorem 2.2 in  
 311 [ADIW24] formalizes the converse of this fact, in the sense that transitions maps  
 312 that are not continuously extendable in this fashion cannot be obtained from simple  
 313 constructions involving predicates. Under this framework, the features  $\pi_{\mathbf{P}} \circ f$  of such  
 314 transitions  $f$  are not approximable by polynomials, and so they are understood as  
 315 “non-computable” since, again, we expect the operations computers carry out to be  
 316 determined by elementary algebra corresponding to polynomials (namely addition  
 317 and multiplication). Therefore it is crucial we assume some extendibility conditions.

318 We say that the CCS  $(\mathbf{L}, \mathcal{P}, \Gamma)$  satisfies the *Extendibility Axiom* if for all  $\gamma \in \Gamma$ ,  
 319 there is  $\tilde{\gamma} : \mathcal{L}_{\text{sh}} \rightarrow \mathcal{L}_{\text{sh}}$  such that for every sizer  $\mathbf{r}_\bullet$  there is an  $\mathbf{s}_\bullet$  such that  
 320  $\tilde{\gamma}|_{\mathcal{L}[\mathbf{r}_\bullet]} : \mathcal{L}[\mathbf{r}_\bullet] \rightarrow \mathcal{L}[\mathbf{s}_\bullet]$  is continuous. For a deeper discussion about this axiom, we  
 321 refer the reader to [ADIW24].

322 A collection  $\mathbf{R}$  of sizers is called *exhaustive* if  $\mathcal{L}_{\text{sh}} = \bigcup_{\mathbf{r}_\bullet \in \mathbf{R}} \mathcal{L}[\mathbf{r}_\bullet]$ . We say that  
 323  $\Delta \subseteq \Gamma$  is  *$\mathbf{R}$ -confined* if  $\gamma|_{\mathcal{L}[\mathbf{r}_\bullet]} : \mathcal{L}[\mathbf{r}_\bullet] \rightarrow \mathcal{L}[\mathbf{r}_\bullet]$  for every  $\mathbf{r}_\bullet \in \mathbf{R}$  and  $\gamma \in \Delta$ . Elements in  
 324  $\Delta$  are called *real-valued computations* (in this article we will refer to them simply as  
 325 *computations*) and elements in  $\overline{\Delta} \subseteq \mathcal{L}_{\text{sh}}^{\mathbf{L}}$  are called (real-valued) *deep computations*  
 326 or *ultracomputations*. By  $\tilde{\Delta}$  we denote the set of all extensions  $\tilde{\gamma}$  for  $\gamma \in \Delta$ . For a  
 327 more complete description of this framework, we refer the reader to [ADIW24].

328 **3.1. NIP and Baire-1 definability of deep computations.** Under what con-  
 329 ditions are deep computations Baire class 1, and thus well-behaved according to  
 330 our framework, on type-shards? The next Theorem says that, again under the  
 331 assumption that  $\mathcal{P}$  is countable, the space of deep computations is a Rosenthal  
 332 compactum (when restricted to shards) if and only if the set of computations has  
 333 the NIP on features. Hence, we can import the theory of Rosenthal compacta into  
 334 this framework of deep computations.

335 **Theorem 3.1.** *Let  $(\mathbf{L}, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom with  $\mathcal{P}$*   
 336 *countable. Let  $\mathbf{R}$  be an exhaustive collection of sizers. Let  $\Delta \subseteq \Gamma$  be  $\mathbf{R}$ -confined. The*  
 337 *following are equivalent.*

- 338 (1)  $\tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]} \subseteq \mathbf{B}_1(\mathcal{L}[\mathbf{r}_\bullet], \mathcal{L}[\mathbf{r}_\bullet])$  for all  $\mathbf{r}_\bullet \in \mathbf{R}$ .  
 (2)  $\pi_{\mathbf{P}} \circ \Delta|_{\mathcal{L}[\mathbf{r}_\bullet]}$  has the NIP for all  $\mathbf{P} \in \mathcal{P}$  and  $\mathbf{r}_\bullet \in \mathbf{R}$ , that is, for all  $\mathbf{P} \in \mathcal{P}$ ,  
 $\mathbf{r}_\bullet \in \mathbf{R}$ ,  $\mathbf{a} < \mathbf{b}$ ,  $\{\gamma_n\}_{n \in \mathbb{N}} \subseteq \Delta$  there are finite disjoint  $\mathbf{E}, \mathbf{F} \subseteq \mathbb{N}$  such that

$$\mathcal{L}[\mathbf{r}_\bullet] \cap \bigcap_{n \in \mathbf{E}} (\pi_{\mathbf{P}} \circ \gamma_n)^{-1}(-\infty, \mathbf{a}] \cap \bigcap_{n \in \mathbf{F}} (\pi_{\mathbf{P}} \circ \gamma_n)^{-1}[\mathbf{b}, \infty) = \emptyset.$$

Moreover, if any (hence all) of the preceding conditions hold, then every deep computation  $f \in \bar{\Delta}$  can be extended to a Baire-1 function on shards, i.e., there is  $\tilde{f} : \mathcal{L}_{\text{sh}} \rightarrow \mathcal{L}_{\text{sh}}$  such that  $\tilde{f}|_{\mathcal{L}[\mathbf{r}_\bullet]} \in B_1(\mathcal{L}[\mathbf{r}_\bullet], \mathcal{L}[\mathbf{r}_\bullet])$  for all  $\mathbf{r}_\bullet \in \mathbf{R}$ . In particular, on each shard every deep computation is the pointwise limit of a countable sequence of computations.

*Proof.* Since  $\mathcal{P}$  is countable, then  $\mathcal{L}[\mathbf{r}_\bullet] \subseteq \mathbb{R}^{\mathcal{P}}$  is Polish. Also, the Extendibility Axiom implies that  $\pi_{\mathcal{P}} \circ \tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]}$  is a pointwise bounded set of continuous functions for all  $\mathbf{P} \in \mathcal{P}$ . Hence, Theorem 2.10 and Lemma 2.11 prove the equivalence of (1) and (2). If (1) holds and  $f \in \bar{\Delta}$ , then write  $f = \mathcal{U}\lim_i \gamma_i$  as an ultralimit. Define  $\tilde{f} := \mathcal{U}\lim_i \tilde{\gamma}_i$ . Hence, for all  $\mathbf{r}_\bullet \in \mathbf{R}$  we have  $\tilde{f}|_{\mathcal{L}[\mathbf{r}_\bullet]} \in \overline{\tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]}} \subseteq B_1(\mathcal{L}[\mathbf{r}_\bullet], \mathcal{L}[\mathbf{r}_\bullet])$ . That every deep computation is a pointwise limit of a countable sequence of computations follows from the fact that for a Polish space  $X$  every compact subset of  $B_1(X)$  is Fréchet-Urysohn (that is, a space where topological closures coincide with sequential closures, see Theorem 3F in [BFT78] or Theorem 4.1 in [Deb13]).  $\square$

Given a countable set  $\Delta$  of computations satisfying the NIP on features and shards (condition (2) of Theorem 3.1) we have that  $\overline{\tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]}}$  (for a fixed sizer  $\mathbf{r}_\bullet$ ) is a separable *Rosenthal compactum* (compact subset of  $B_1(\mathcal{P} \times \mathcal{L}[\mathbf{r}_\bullet])$ ). The work of Todorčević ([Tod99]) and Argyros, Dodos, Kanellopoulos ([ADK08]) culminates in a trichotomy theorem for separable Rosenthal Compacta. Inspired by the work of Glasner and Megrelishvili ([GM22]), we are interested to see how this allows us to classify and obtain different levels of PAC-learnability (NIP).

Recall that a topological space  $X$  is *hereditarily separable* (HS) if every subspace is separable and that  $X$  is *first countable* if every point in  $X$  has a countable local basis. Every separable metrizable space is hereditarily separable and it is a result of R. Pol that every hereditarily separable Rosenthal compactum is first countable (see section 10 in [Deb13]). This suggests the following definition:

**Definition 3.2.** Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom and  $\mathbf{R}$  be an exhaustive collection of sizers. Let  $\Delta \subseteq \Gamma$  be an  $\mathbf{R}$ -confined countable set of computations satisfying the NIP on shards and features (condition (2) in Theorem 3.1). We say that  $\Delta$  is:

- (i) NIP<sub>1</sub> if  $\overline{\tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]}}$  is first countable for every  $\mathbf{r}_\bullet \in \mathbf{R}$ .
- (ii) NIP<sub>2</sub> if  $\overline{\tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]}}$  is hereditarily separable for every  $\mathbf{r}_\bullet \in \mathbf{R}$ .
- (iii) NIP<sub>3</sub> if  $\overline{\tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_\bullet]}}$  is metrizable for every  $\mathbf{r}_\bullet \in \mathbf{R}$ .

Observe that  $\text{NIP}_3 \Rightarrow \text{NIP}_2 \Rightarrow \text{NIP}_1 \Rightarrow \text{NIP}$ . A natural question that would continue this work is to find examples of CCS that separate these levels of NIP. In [Tod99], Todorčević isolates 3 canonical examples of Rosenthal compacta that witness the failure of the converse implications above. We now present some separable and non-separable examples of Rosenthal compacta:

- (1) *Alexandroff compactification of a discrete space of size continuum.* For each  $\mathbf{a} \in 2^{\mathbb{N}}$  consider the map  $\delta_{\mathbf{a}} : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  given by  $\delta_{\mathbf{a}}(\mathbf{x}) = 1$  if  $\mathbf{x} = \mathbf{a}$  and  $\delta_{\mathbf{a}}(\mathbf{x}) = 0$  otherwise. Let  $A(2^{\mathbb{N}}) = \{\delta_{\mathbf{a}} : \mathbf{a} \in 2^{\mathbb{N}}\} \cup \{0\}$ , where  $0$  is the zero map. Notice that  $A(2^{\mathbb{N}})$  is a compact subset of  $B_1(2^{\mathbb{N}})$ , in fact  $\{\delta_{\mathbf{a}} : \mathbf{a} \in 2^{\mathbb{N}}\}$  is a discrete subspace of  $B_1(2^{\mathbb{N}})$  and its pointwise closure is precisely  $A(2^{\mathbb{N}})$ . Hence, this is a Rosenthal compactum which is not first countable. Notice that this space is also not separable.

- 384 (2) *Extended Alexandroff compactification.* For each finite binary sequence  $s \in$   
 385  $2^{<\mathbb{N}}$ , let  $v_s : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  be given by  $v_s(x) = 1$  if  $x$  extends  $s$  and  $v_s(x) =$   
 386  $0$  otherwise. Let  $\hat{A}(2^{\mathbb{N}})$  be the pointwise closure of  $\{v_s : s \in 2^{<\mathbb{N}}\}$ , i.e.,  
 387  $\hat{A}(2^{\mathbb{N}}) = A(2^{\mathbb{N}}) \cup \{v_s : s \in 2^{<\mathbb{N}}\}$ . Note that this space is a separable  
 388 Rosenthal compactum which is not first countable.
- 389 (3) *Split Cantor.* Let  $<$  be the lexicographic order in the space of infinite  
 390 binary sequences, i.e.,  $2^{\mathbb{N}}$ . For each  $a \in 2^{\mathbb{N}}$  let  $f_a^- : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  be given by  
 391  $f_a^-(x) = 1$  if  $x < a$  and  $f_a^-(x) = 0$  otherwise. Let  $f_a^+ : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  be given  
 392 by  $f_a^+(x) = 1$  if  $x \leq a$  and  $f_a^+(x) = 0$  otherwise. The split Cantor is the  
 393 space  $S(2^{\mathbb{N}}) = \{f_a^- : a \in 2^{\mathbb{N}}\} \cup \{f_a^+ : a \in 2^{\mathbb{N}}\}$ . This is a separable Rosenthal  
 394 compactum. One example of a countable dense subset is the set of all  $f_a^+$   
 395 and  $f_a^-$  where  $a$  is an infinite binary sequence that is eventually constant.  
 396 Moreover, it is hereditarily separable but it is not metrizable.
- (4) *Alexandroff Duplicate.* Let  $K$  be any compact metric space and consider  
 the Polish space  $X = C(K) \sqcup K$ , i.e., the disjoint union of  $C(K)$  (with its  
 supremum norm topology) and  $K$ . For each  $a \in K$  define  $g_a^0, g_a^1 : X \rightarrow \mathbb{R}$  as  
 follows:

$$g_a^0(x) = \begin{cases} x(a), & x \in C(K) \\ 0, & x \in K \end{cases}$$

$$g_a^1(x) = \begin{cases} x(a), & x \in C(K) \\ \delta_a(x), & x \in K \end{cases}$$

- 397 Let  $D(K) = \{g_a^0 : a \in K\} \cup \{g_a^1 : a \in K\}$ . Notice that  $D(K)$  is a first  
 398 countable Rosenthal compactum. It is not separable if  $K$  is uncountable.  
 399 The interesting case will be when  $K = 2^{\mathbb{N}}$ .
- (5) *Extended Alexandroff Duplicate of the split Cantor.* For each finite binary  
 sequence  $t \in 2^{<\mathbb{N}}$  let  $a_t \in 2^{\mathbb{N}}$  be the sequence starting with  $t$  and ending  
 with 0's and let  $b_t \in 2^{\mathbb{N}}$  be the sequence starting with  $t$  and ending with  
 1's. Define  $h_t : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  by

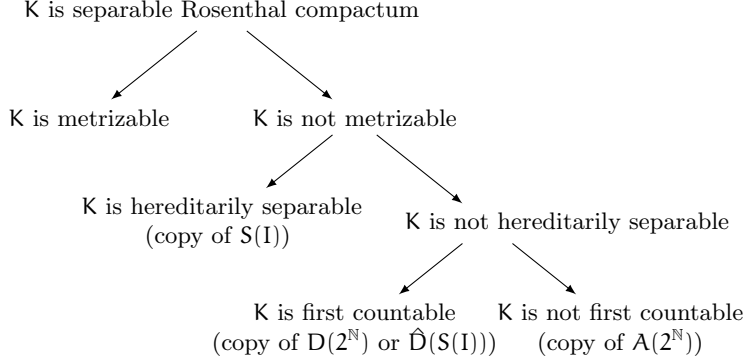
$$h_t(x) = \begin{cases} 0, & x < a_t \\ 1/2, & a_t \leq x \leq b_t \\ 1, & b_t < x \end{cases}$$

- 400 Let  $\hat{D}(S(2^{\mathbb{N}}))$  be the pointwise closure of the set  $\{h_t : t \in 2^{<\mathbb{N}}\}$ . Hence,  
 401  $\hat{D}(S(2^{\mathbb{N}}))$  is a separable first countable Rosenthal compactum which is not  
 402 hereditarily separable. In fact, it contains an uncountable discrete subspace  
 403 (see Theorem 5 in [Tod99]).

404 **Theorem 3.3** (Todorćević's Trichotomy, [Tod99], Theorem 3 in [ADK08]). *Let  $K$*   
 405 *be a separable Rosenthal Compactum.*

- 406 (i) *If  $K$  is hereditarily separable but non-metrizable, then  $S(2^{\mathbb{N}})$  embeds into  $K$ .*  
 407 (ii) *If  $K$  is first countable but not hereditarily separable, then either  $D(2^{\mathbb{N}})$  or*  
 408  *$\hat{D}(S(2^{\mathbb{N}}))$  embeds into  $K$ .*  
 409 (iii) *If  $K$  is not first countable, then  $A(2^{\mathbb{N}})$  embeds into  $K$ .*

410 In other words, we have the following classification:



411

412 Lastly, the definitions provided here for  $NIP_i$  ( $i = 1, 2, 3$ ) are topological.

413 **Question 3.4.** Is there a non-topological characterization for  $NIP_i$ ,  $i = 1, 2, 3$ ?

414 More can be said about the nature of the embeddings in Todorćević's Trichotomy.  
 415 Given a separable Rosenthal compactum  $K$ , there is typically more than one countable  
 416 dense subset of  $K$ . We can view a separable Rosenthal compactum as the accumulation  
 417 points of a countable family of pointwise bounded real-valued functions.  
 418 The choice of the countable families is not important when a bijection between  
 419 them can be lifted to a homeomorphism of their closures. To be more precise:

420 **Definition 3.5.** Given a Polish space  $X$ , a countable set  $I$  and two pointwise  
 421 bounded families  $\{f_i : i \in I\} \subseteq \mathbb{R}^X$ ,  $\{g_i : i \in I\} \subseteq \mathbb{R}^X$  indexed by  $I$ . We say that  
 422  $\{f_i : i \in I\}$  and  $\{g_i : i \in I\}$  are *equivalent* if and only if the map  $f_i \mapsto g_i$  is extended  
 423 to a homeomorphism from  $\overline{\{f_i : i \in I\}}$  to  $\overline{\{g_i : i \in I\}}$ .

424 Notice that in the separable examples discussed before ( $\hat{A}(2^{\mathbb{N}})$ ,  $S(2^{\mathbb{N}})$  and  $\hat{D}(S(2^{\mathbb{N}}))$ )  
 425 the countable dense subsets are indexed by the binary tree  $2^{<\mathbb{N}}$ . This choice of index  
 426 is useful because the Ramsey theory of perfect subsets of the Cantor space  $2^{\mathbb{N}}$   
 427 can be imported to analyze the behavior of the accumulation points. Since  $2^{<\mathbb{N}}$  is  
 428 countable, we can always choose this index for the countable dense subsets. This  
 429 is done in [ADK08].

430 **Definition 3.6.** Given a Polish space  $X$  and a pointwise bounded family  $\{f_t : t \in$   
 431  $2^{<\mathbb{N}}\}$ . We say that  $\{f_t : t \in 2^{<\mathbb{N}}\}$  is *minimal* if and only if for every dyadic subtree  
 432  $\{s_t : t \in 2^{<\mathbb{N}}\}$  of  $2^{<\mathbb{N}}$ ,  $\{f_{s_t} : t \in 2^{<\mathbb{N}}\}$  is equivalent to  $\{f_t : t \in 2^{<\mathbb{N}}\}$ .

433 One of the main results in [ADK08] is that there are (up to equivalence) seven  
 434 minimal families of Rosenthal compacta and that for every relatively compact  $\{f_t :$   
 435  $t \in 2^{<\mathbb{N}}\} \subseteq B_1(X)$  there is a dyadic subtree  $\{s_t : t \in 2^{<\mathbb{N}}\}$  such that  $\{f_{s_t} : t \in 2^{<\mathbb{N}}\}$   
 436 is equivalent to one of the minimal families. We shall describe the minimal families  
 437 next. We will follow the same notation as in [ADK08]. For any node  $t \in 2^{<\mathbb{N}}$ , we  
 438 denote by  $t \smallfrown 0^\infty$  ( $t \smallfrown 1^\infty$ ) the infinite binary sequence starting with  $t$  and ending  
 439 in 0's (1's). Fix a regular dyadic subtree  $R = \{s_t : t \in 2^{<\mathbb{N}}\}$  of  $2^{<\mathbb{N}}$  (i.e., a dyadic  
 440 subtree such that every level of  $R$  is contained in a level of  $2^{<\mathbb{N}}$ ) with the property  
 441 that for all  $s, s' \in R$ ,  $s \smallfrown 0^\infty \neq s' \smallfrown 0^\infty$  and  $s \smallfrown 1^\infty \neq s' \smallfrown 1^\infty$ . Given  $t \in 2^{<\mathbb{N}}$ , let  
 442  $v_t$  be the characteristic function of the set  $\{x \in 2^{\mathbb{N}} : x \text{ extends } t\}$ . Let  $<$  be the  
 443 lexicographic order in  $2^{\mathbb{N}}$ . Given  $a \in 2^{\mathbb{N}}$ , let  $f_a^+ : 2^{\mathbb{N}} \rightarrow \{0, 1\}$  be the characteristic  
 444 function of  $\{x \in 2^{\mathbb{N}} : a \leq x\}$  and let  $f_a^- : 2^{\mathbb{N}} \rightarrow \{0, 1\}$  be the characteristic function of

445  $\{x \in 2^{\mathbb{N}} : a < x\}$ . Given two maps  $f, g : 2^{\mathbb{N}} \rightarrow \mathbb{R}$  we denote by  $(f, g) : 2^{\mathbb{N}} \sqcup 2^{\mathbb{N}} \rightarrow \mathbb{R}$   
 446 the function which is  $f$  on the first copy of  $2^{\mathbb{N}}$  and  $g$  on the second copy of  $2^{\mathbb{N}}$ .

- 447 (1)  $D_1 = \{\frac{1}{|t|+1}v_t : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_1} = A(2^{\mathbb{N}})$ .  
 448 (2)  $D_2 = \{s_t \widehat{0}^\infty : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_2} = 2^{\leq \mathbb{N}}$ .  
 449 (3)  $D_3 = \{f_{s_t \widehat{0}^\infty}^+ : t \in 2^{<\mathbb{N}}\}$ . This is a discrete in  $\overline{D_3} = S(2^{\mathbb{N}})$ .  
 450 (4)  $D_4 = \{f_{s_t \widehat{1}^\infty}^- : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_4} = S(2^{\mathbb{N}})$ .  
 451 (5)  $D_5 = \{v_t : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_5} = \widehat{A}(2^{\mathbb{N}})$ .  
 452 (6)  $D_6 = \{(v_{s_t}, s_t \widehat{0}^\infty) : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_6} = \widehat{D}(2^{\mathbb{N}})$ .  
 453 (7)  $D_7 = \{(v_{s_t}, x_{s_t \widehat{0}^\infty}^+) : t \in 2^{<\mathbb{N}}\}$ . This is discrete in  $\overline{D_7} = \widehat{D}(S(2^{\mathbb{N}}))$

454 **Theorem 3.7** (Heptacotomy of minimal families, Theorem 2 in [ADK08]). *Let*  
 455  *$X$  be Polish. For every relatively compact  $\{f_t : t \in 2^{<\mathbb{N}}\} \subseteq B_1(X)$ , there exists  $i =$*   
 456  *$1, 2, \dots, 7$  and a regular dyadic subtree  $\{s_t : t \in 2^{<\mathbb{N}}\}$  of  $2^{<\mathbb{N}}$  such that  $\{f_{s_t} : t \in 2^{<\mathbb{N}}\}$*   
 457 *is equivalent to  $D_i$ . Moreover, all  $D_i$  are minimal and mutually non-equivalent.*

458 **3.2. NIP and definability by universally measurable functions.** We now  
 459 turn to the question: what happens when  $\mathcal{P}$  is uncountable? Notice that the count-  
 460 ability assumption is crucial in the proof of Theorem 2.10 essentially because it  
 461 makes  $\mathbb{R}^{\mathcal{P}}$  a Polish space. For the uncountable case, we may lose Baire-1 definabil-  
 462 ity so we shall replace  $B_1(X)$  by a bigger class. Recall that the purpose of studying  
 463 the class of Baire-1 functions is that a pointwise limit of continuous functions is not  
 464 necessarily continuous. In [BFT78], J. Bourgain, D.H. Fremlin and M. Talagrand  
 465 characterized the Non-Independence Property of a set of continuous functions with  
 466 various notions of compactness in function spaces containing  $C(X)$ , such as  $B_1(X)$ .  
 467 In this section we will replace  $B_1(X)$  with the larger space  $M_r(X)$  of universally  
 468 measurable functions. The development of this section is based on Theorem 2F in  
 469 [BFT78]. We now give the relevant definitions. Readers with little familiarity with  
 470 measure theory can review the appendix for standard definitions appearing in this  
 471 subsection.

472 Given a Hausdorff space  $X$  and a measurable space  $(Y, \Sigma)$ , we say that  $f : X \rightarrow Y$   
 473 is *universally measurable* (with respect to  $\Sigma$ ) if  $f^{-1}(E)$  is universally measurable  
 474 for every  $E \in \Sigma$ , i.e.,  $f^{-1}(E)$  is  $\mu$ -measurable for every Radon probability measure  
 475  $\mu$  on  $X$ . When  $Y = \mathbb{R}$  we will always take  $\Sigma = \mathcal{B}(\mathbb{R})$ , the Borel  $\sigma$ -algebra of  $\mathbb{R}$ .  
 476 In that case, a function  $f : X \rightarrow \mathbb{R}$  is universally measurable if and only if  $f^{-1}(U)$   
 477 is  $\mu$ -measurable for every Radon probability measure  $\mu$  on  $X$  and every open set  
 478  $U \subseteq \mathbb{R}$ . Following [BFT78], the collection of all universally measurable real-valued  
 479 functions will be denoted by  $M_r(X)$ . In the context of deep computations, we will  
 480 be interested in transition maps from a state space  $L \subseteq \mathbb{R}^{\mathcal{P}}$  to itself. There are two  
 481 natural  $\sigma$ -algebras one can consider in the product space  $\mathbb{R}^{\mathcal{P}}$ : the Borel  $\sigma$ -algebra,  
 482 i.e., the  $\sigma$ -algebra generated by open sets in  $\mathbb{R}^{\mathcal{P}}$ ; and the cylinder  $\sigma$ -algebra, i.e.,  
 483 the  $\sigma$ -algebra generated by Borel cylinder sets or equivalently basic open sets in  
 484  $\mathbb{R}^{\mathcal{P}}$ . Note that when  $\mathcal{P}$  is countable, both  $\sigma$ -algebras coincide but in general the  
 485 cylinder  $\sigma$ -algebra is strictly smaller. We will use the cylinder  $\sigma$ -algebra to define  
 486 universally measurable maps  $f : \mathbb{R}^{\mathcal{P}} \rightarrow \mathbb{R}^{\mathcal{P}}$ . The reason for this choice is because of  
 487 the following characterization:

488 **Lemma 3.8.** *Let  $X$  be a Hausdorff space and  $Y = \prod_{i \in I} Y_i$  be any product of*  
 489 *measurable spaces  $(Y_i, \Sigma_i)$  for  $i \in I$ . Let  $\Sigma_Y$  be the cylinder  $\sigma$ -algebra generated by*  
 490 *the measurable spaces  $(Y_i, \Sigma_i)$ . Let  $f : X \rightarrow Y$ . The following are equivalent:*

- 491 (i)  $f : X \rightarrow Y$  is universally measurable (with respect to  $\Sigma_Y$ ).  
 492 (ii)  $\pi_i \circ f : X \rightarrow Y_i$  is universally measurable (with respect to  $\Sigma_i$ ) for all  $i \in I$ .

493 *Proof.* (i) $\Rightarrow$ (ii) is clear since the projection maps  $\pi_i$  are measurable and the com-  
 494 position of measurable functions is measurable. To prove (ii) $\Rightarrow$ (i), suppose that  
 495  $C = \prod_{i \in I} C_i$  is a measurable cylinder and let  $J$  be the finite set of  $i \in I$  such that  
 496  $C_i \neq Y_i$ . Then,  $C = \bigcap_{i \in J} \pi_i^{-1}(C_i)$  so  $f^{-1}(C) = \bigcap_{i \in J} (\pi_i \circ f)^{-1}(C_i)$  is a universally  
 497 measurable set by assumption.  $\square$

498 The previous lemma says that a transition map is universally measurable if and  
 499 only if it is universally measurable on all its features. In other words, we can check  
 500 measurability of a transition just by checking measurability in all its features. We  
 501 will denote by  $M_r(X, \mathbb{R}^{\mathcal{P}})$  the collection of all universally measurable functions  
 502  $f : X \rightarrow \mathbb{R}^{\mathcal{P}}$  (with respect to the cylinder  $\sigma$ -algebra), endowed with the topology of  
 503 pointwise convergence.

504 **Definition 3.9.** Let  $(L, \mathcal{P}, \Gamma)$  be a CCS. We say that a transition  $f : L \rightarrow L$  is  
 505 *universally measurable shard-definable* if and only if there exists  $\tilde{f} : \mathcal{L}_{sh} \rightarrow \mathcal{L}_{sh}$   
 506 extending  $f$  such that for every sizer  $\mathbf{r}_{\bullet}$  there is a sizer  $\mathbf{s}_{\bullet}$  such that the restriction  
 507  $\tilde{f}|_{\mathcal{L}[\mathbf{r}_{\bullet}]} : \mathcal{L}[\mathbf{r}_{\bullet}] \rightarrow \mathcal{L}[\mathbf{s}_{\bullet}]$  is universally measurable, i.e.  $\pi_P \circ \tilde{f}|_{\mathcal{L}[\mathbf{r}_{\bullet}]} : \mathcal{L}[\mathbf{r}_{\bullet}] \rightarrow [-s_P, s_P]$   
 508 is  $\mu$ -measurable for every Radon probability measure  $\mu$  on  $\mathcal{L}[\mathbf{r}_{\bullet}]$ .

509 We will need the following result about NIP and universally measurable func-  
 510 tions:

511 **Theorem 3.10** (Bourgain-Fremlin-Talagrand, Theorem 2F in [BFT78]). *Let  $X$  be a*  
 512 *Hausdorff space and  $A \subseteq C(X)$  be pointwise bounded. The following are equivalent:*

- 513 (i)  $\overline{A} \subseteq M_r(X)$ .  
 514 (ii) For every compact  $K \subseteq X$ ,  $A|_K$  has the NIP.  
 515 (iii) For every Radon measure  $\mu$  on  $X$ ,  $A$  is relatively countably compact in  
 516  $\mathcal{L}^0(X, \mu)$ , i.e., every countable subset of  $A$  has an accumulation point in  
 517  $\mathcal{L}^0(X, \mu)$ .

518 Theorem 2.5 immediately yields the following.

519 **Theorem 3.11.** *Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom. Let  $R$*   
 520 *be an exhaustive collection of sizers. Let  $\Delta \subseteq \Gamma$  be  $R$ -confined. If  $\pi_P \circ \Delta|_{\mathcal{L}[\mathbf{r}_{\bullet}]}$  has*  
 521 *the NIP for all  $P \in \mathcal{P}$  and all  $\mathbf{r}_{\bullet} \in R$ , then every deep computation is universally*  
 522 *measurable shard-definable.*

523 *Proof.* By the Extendibility Axiom, Theorem 2.5 and lemma 2.11 we have that  
 524  $\pi_P \circ \tilde{\Delta}|_{\mathcal{L}[\mathbf{r}_{\bullet}]} \subseteq M_r(\mathcal{L}[\mathbf{r}_{\bullet}])$  for all  $\mathbf{r}_{\bullet} \in R$  and  $P \in \mathcal{P}$ . Let  $f \in \overline{\Delta}$  be a deep computation.  
 525 Write  $f = \mathcal{U} \lim_i \gamma_i$  as an ultralimit of computations in  $\Delta$ . Define  $\tilde{f} := \mathcal{U} \lim_i \tilde{\gamma}_i$ .  
 526 Then, for all  $\mathbf{r}_{\bullet} \in R$  and  $P \in \mathcal{P}$   $\pi_P \circ \tilde{\gamma}_i|_{\mathcal{L}[\mathbf{r}_{\bullet}]} \in M_r(\mathcal{L}[\mathbf{r}_{\bullet}])$  for all  $i$  so  $\pi_P \circ \tilde{f}|_{\mathcal{L}[\mathbf{r}_{\bullet}]} \in$   
 527  $\overline{M_r(\mathcal{L}[\mathbf{r}_{\bullet}])} \subseteq M_r(\mathcal{L}[\mathbf{r}_{\bullet}])$ .  $\square$

528 **Question 3.12.** Under the same assumptions of the previous Theorem, suppose  
 529 that every deep computation of  $\Delta$  is universally measurable shard-definable. Must  
 530  $\pi_P \circ \Delta|_{\mathcal{L}[\mathbf{r}_{\bullet}]}$  have the NIP for all  $P \in \mathcal{P}$  and all  $\mathbf{r}_{\bullet} \in R$ ?

531 **3.3. Talagrand stability and definability by universally measurable func-**  
 532 **tions.** There is another notion closely related to NIP, introduced by Talagrand  
 533 in [Tal84] while studying Pettis integration. Suppose that  $X$  is a compact Haus-  
 534 dorff space and  $A \subseteq \mathbb{R}^X$ . Let  $\mu$  be a Radon probability measure on  $X$ . Given a  
 535  $\mu$ -measurable set  $E \subseteq X$ , a positive integer  $k$  and real numbers  $a < b$ , we write:

$$D_k(A, E, a, b) = \bigcup_{f \in A} \{x \in E^{2k} : f(x_{2i}) \leq a, f(x_{2i+1}) \geq b \text{ for all } i < k\}$$

536 We say that  $A$  is *Talagrand  $\mu$ -stable* if and only if for every  $\mu$ -measurable  
 537 set  $E \subseteq X$  of positive measure and for every  $a < b$  there is  $k \geq 1$  such that  
 538  $(\mu^{2k})^*(D_k(A, E, a, b)) < (\mu(E))^{2k}$ . Notice that we work with the outer measure  
 539 because it is not necessarily true that the sets  $D_k(A, E, a, b)$  are  $\mu$ -measurable.  
 540 This is certainly the case when  $A$  is a countable set of continuous (or  $\mu$ -measurable)  
 541 functions.

542 The following lemma establishes that Talagrand stability is a way to ensure that  
 543 deep computations are definable by measurable functions. We include the proof for  
 544 the reader's convenience.

545 **Lemma 3.13.** *If  $A$  is Talagrand  $\mu$ -stable, then  $\overline{A}$  is also Talagrand  $\mu$ -stable and*  
 546  $\overline{A} \subseteq \mathcal{L}^0(X, \mu)$ .

547 *Proof.* First, observe that a subset of a  $\mu$ -stable set is  $\mu$ -stable. To show that  $\overline{A}$   
 548 is  $\mu$ -stable, observe that  $D_k(\overline{A}, E, a, b) \subseteq D_k(A, E, a', b')$  where  $a < a' < b' <$   
 549  $b$  and  $E$  is a  $\mu$ -measurable set with positive measure. It suffices to show that  
 550  $\overline{A} \subseteq \mathcal{L}^0(X, \mu)$ . Suppose that there exists  $f \in \overline{A}$  such that  $f \notin \mathcal{L}^0(X, \mu)$ . By a  
 551 characterization of measurable functions (see 413G in [Fre03]), there exists a  $\mu$ -  
 552 measurable set  $E$  of positive measure and  $a < b$  such that  $\mu^*(P) = \mu^*(Q) = \mu(E)$   
 553 where  $P = \{x \in E : f(x) \leq a\}$  and  $Q = \{x \in E : f(x) \geq b\}$ . Then, for any  $k \geq 1$ :  
 554  $(P \times Q)^k \subseteq D_k(\{f\}, E, a, b)$  so  $(\mu^{2k})^*(D_k(\{f\}, E, a, b)) = (\mu^*(P)\mu^*(Q))^k = (\mu(E))^{2k}$ .  
 555 Thus,  $\{f\}$  is not  $\mu$ -stable, but we argued before that a subset of a  $\mu$ -stable set must  
 556 be  $\mu$ -stable.  $\square$

557 We say that  $A$  is *universally Talagrand stable* if  $A$  is Talagrand  $\mu$ -stable for  
 558 every Radon probability measure  $\mu$  on  $X$ . A similar argument as before, yields the  
 559 following:

560 **Theorem 3.14.** *Let  $(L, \mathcal{P}, \Gamma)$  be a CCS satisfying the Extendibility Axiom. If  $\pi_P \circ$   
 561  $\Delta|_{L[r_\bullet]}$  is universally Talagrand stable for all  $P \in \mathcal{P}$  and all sizers  $r_\bullet$ , then every  
 562 deep computation is universally measurable sh-definable.*

563 It is then natural to ask: what is the relationship between Talagrand stability  
 564 and the NIP? The following dichotomy will be useful.

565 **Lemma 3.15** (Fremlin's Dichotomy, 463K in [Fre03]). *If  $(X, \Sigma, \mu)$  is a perfect  $\sigma$ -*  
 566 *finite measure space (in particular, for  $X$  compact and  $\mu$  a Radon probability measure*  
 567 *on  $X$ ) and  $\{f_n : n \in \mathbb{N}\}$  be a sequence of real-valued measurable functions on  $X$ , then*  
 568 *either:*

- 569 (i)  $\{f_n : n \in \mathbb{N}\}$  has a subsequence that converges  $\mu$ -almost everywhere, or
- 570 (ii)  $\{f_n : n \in \mathbb{N}\}$  has a subsequence with no  $\mu$ -measurable accumulation point in  
 571  $\mathbb{R}^X$ .

572 The preceding lemma can be considered as the measure theoretic version of  
 573 Rosenthal's Dichotomy. Combining this dichotomy with the Theorem 3.10 we get  
 574 the following result:

575 **Theorem 3.16.** *Let  $X$  be a Hausdorff space and  $A \subseteq C(X)$  be pointwise bounded.*  
 576 *The following are equivalent:*

- 577 (i)  $\overline{A} \subseteq M_r(X)$ .
- 578 (ii) *For every compact  $K \subseteq X$ ,  $A|_K$  has the NIP.*
- 579 (iii) *For every Radon measure  $\mu$  on  $X$ ,  $A$  is relatively countably compact in*  
 580  *$\mathcal{L}^0(X, \mu)$ , i.e., every countable subset of  $A$  has an accumulation point in*  
 581  *$\mathcal{L}^0(X, \mu)$ .*
- 582 (iv) *For every Radon measure  $\mu$  on  $X$  and every sequence  $\{f_n : n \in \mathbb{N}\} \subseteq A$ ,*  
 583 *there is a subsequence that converges  $\mu$ -almost everywhere.*

584 *Proof.* Notice that the equivalence (i)-(iii) is Theorem 3.10. Notice that the equiv-  
 585 alence of (iii) and (iv) is Fremlin's Dichotomy Theorem.  $\square$

586 **Lemma 3.17.** *Let  $X$  be a compact Hausdorff space and  $A \subseteq C(X)$  be pointwise*  
 587 *bounded. If  $A$  is universally Talagrand stable, then  $A$  has the NIP.*

588 *Proof.* By Theorem 3.10, it suffices to show that  $A$  is relatively countably compact  
 589 in  $\mathcal{L}^0(X, \mu)$  for all Radon probability measure  $\mu$  on  $X$ . Since  $A$  is Talagrand  $\mu$ -stable  
 590 for any such  $\mu$ , then  $\overline{A} \subseteq \mathcal{L}^0(X, \mu)$ . In particular,  $A$  is relatively countably compact  
 591 in  $\mathcal{L}^0(X, \mu)$ .  $\square$

592 **Question 3.18.** Is the converse true?

593 There is a delicate point in this question, as it may be sensitive to set-theoretic  
 594 axioms (even assuming countability of  $A$ ).

595 **Theorem 3.19** (Talagrand, Theorem 9-3-1(a) in [Tal84]). *Let  $X$  be a compact*  
 596 *Hausdorff space and  $A \subseteq M_r(X)$  be countable and pointwise bounded. Assume that*  
 597  *$[0, 1]$  is not the union of  $< \mathfrak{c}$  closed measure zero sets. If  $A$  has the NIP, then  $A$  is*  
 598 *universally Talagrand stable.*

599 **Theorem 3.20** (Fremlin, Shelah, [FS93]). *It is consistent that there exists a count-*  
 600 *able pointwise bounded set of Lebesgue measurable functions with the NIP which is*  
 601 *not Talagrand stable with respect to Lebesgue measure.*

## 602 APPENDIX: MEASURE THEORY

603 Given a set  $X$ , a collection  $\Sigma$  of subsets of  $X$  is called a  $\sigma$ -algebra if  $\Sigma$  contains  
 604  $X$  and is closed under complements and countable unions. Hence, for example, a  
 605  $\sigma$ -algebra is also closed under countable intersections. Intuitively, a  $\sigma$ -algebra is  
 606 a collection of sets in which we can define a  $\sigma$ -additive measure. We call sets in  
 607 a  $\sigma$ -algebra  $\Sigma$  *measurable sets* and the pair  $(X, \Sigma)$  a measurable space. If  $X$  is a  
 608 topological space, there is a natural  $\sigma$ -algebra of subsets of  $X$ , namely the *Borel*  
 609  *$\sigma$ -algebra*  $\mathcal{B}(X)$ , i.e., the smallest  $\sigma$ -algebra containing all open subsets of  $X$ . Given  
 610 two measurable spaces  $(X, \Sigma_X)$  and  $(Y, \Sigma_Y)$ , we say that a function  $f : X \rightarrow Y$  is  
 611 *measurable* if and only if  $f^{-1}(E) \in \Sigma_X$  for every  $E \in \Sigma_Y$ . In particular, we say that  
 612  $f : X \rightarrow \mathbb{R}$  is measurable if  $f^{-1}(E) \in \Sigma_X$  for all  $E \in \mathcal{B}(\mathbb{R})$  (equivalently,  $E$  open in  
 613  $\mathbb{R}$ ).



Given a measurable space  $(X, \Sigma)$ , a  $\sigma$ -additive measure is a non-negative function  $\mu : \Sigma \rightarrow \mathbb{R}$  with the property that  $\mu(\emptyset) = 0$  and  $\mu(\bigcup_{n=0}^{\infty} A_n) = \sum_{n=0}^{\infty} \mu(A_n)$  whenever  $\{A_n : n \in \mathbb{N}\} \subseteq \Sigma$  is pairwise disjoint. We call  $(X, \Sigma, \mu)$  a *measure space*. A  $\sigma$ -additive measure is called a *probability measure* if  $\mu(X) = 1$ . A measure  $\mu$  is *complete* if for every  $A \subseteq B \in \Sigma$ ,  $\mu(B) = 0$  implies  $A \in \Sigma$ . In words, subsets of measure-zero sets are always measurable (and hence, by the monotonicity of  $\mu$ , have measure zero as well). A measure  $\mu$  is  $\sigma$ -finite if  $X = \bigcup_{n=1}^{\infty} X_n$  where  $\mu(X_n) < \infty$  for all  $n \in \mathbb{N}$  (i.e.,  $X$  can be decomposed into countably many finite measure sets). A measure  $\mu$  is *perfect* if for every measurable  $f : X \rightarrow \mathbb{R}$  and every measurable set  $E$  with  $\mu(E) > 0$ , there exists a compact  $K \subseteq f(E)$  such that  $\mu(f^{-1}(K)) > 0$ . We say that a property  $\phi(x)$  about  $x \in X$  holds  $\mu$ -almost everywhere if  $\mu(\{x \in X : \phi(x) \text{ does not hold}\}) = 0$ .

A special example of the preceding concepts is that of a *Radon measure*. If  $X$  is a Hausdorff topological space, then a measure  $\mu$  on the Borel sets of  $X$  is called a *Radon measure* if

- for every open set  $U$ ,  $\mu(U)$  is the supremum of  $\mu(K)$  over all compact  $K \subseteq U$ , that is, the measure of open sets may be approximated via compact sets; and
- every point of  $X$  has a neighborhood  $U \ni x$  for which  $\mu(U)$  is finite.

Perhaps the most famous example of a Radon measure on  $\mathbb{R}$  is the Lebesgue measure of Borel sets. If  $X$  is finite,  $\mu(A) := |A|$  (the cardinality of  $A$ ) defines a Radon measure on  $X$ . Every Radon measure is perfect (see 451A, 451B and 451C in [Fre03]).

While not immediately obvious, sets can be measurable according to one measure, but non-measurable according to another. Given a measure space  $(X, \Sigma, \mu)$  we say that a set  $E \subseteq X$  is  $\mu$ -measurable if there are  $A, B \in \Sigma$  such that  $A \subseteq E \subseteq B$  and  $\mu(B \setminus A) = 0$ . The set of all  $\mu$ -measurable sets is a  $\sigma$ -algebra containing  $\Sigma$  and it is denoted by  $\Sigma_\mu$ . A set  $E \subseteq X$  is *universally measurable* if it is  $\mu$ -measurable for every Radon probability measure on  $X$ . It follows that Borel sets are universally measurable. We say that  $f : X \rightarrow \mathbb{R}$  is  $\mu$ -measurable if  $f^{-1}(E) \in \Sigma_\mu$  for all  $E \in \mathcal{B}(\mathbb{R})$  (equivalently,  $E$  open in  $\mathbb{R}$ ). The set of all  $\mu$ -measurable functions is denoted by  $\mathcal{L}^0(X, \mu)$ .

Recall that if  $\{X_i : i \in I\}$  is a collection of topological spaces indexed by some set  $I$ , then the product space  $X := \prod_{i \in I} X_i$  is endowed with the topology generated by *cylinders*, that is, sets of the form  $\prod_{i \in I} U_i$  where each  $U_i$  is open in  $X_i$ , and  $U_i = X_i$  except for finitely many indices  $i \in I$ . If each space is measurable, say we pair  $X_i$  with a  $\sigma$ -algebra  $\Sigma_i$ , then there are multiple ways to interpret the product space  $X$  as a measurable space, but the interpretation we care about in this paper is the so called *cylinder  $\sigma$ -algebra*, as used in Lemma 3.8. Namely, let  $\Sigma$  be the  $\sigma$ -algebra generated by sets of the form

$$\prod_{i \in I} C_i, \quad C_i \in \Sigma_i, \quad C_i = X_i \text{ for all but finitely many } i \in I.$$

We remark that when  $I$  is uncountable and  $\Sigma_i = \mathcal{B}(X_i)$  for all  $i \in I$ , then  $\Sigma$  is, in general, strictly **smaller** than  $\mathcal{B}(X)$ .

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