Kushnirenko's Fewnomials, the Number of Real Zeros & Condition Numbers



The University of Texas at San Antonio[™]

Josué Tonelli-Cuto

UT San Antonio

josue.tonelli.cueto@bizkaia.eu

Elias Tsigaridas

Inria Paris & IMJ-PRG elias.tsigaridas@inria.fr



What is a fewnomial?

A fewnomial is a polynomial with few monomials.

The polynomial

$$1 - T + 7T^5 + T^{1978}$$

has degree 1978, but, despite this, it has very few monomials, so it is an example of a fewnomial.

Why don't we say oligonomial?

polynomial — poly-nomial — много-члены много члены — мало члены (many monomials) (few monomials) \rightarrow мало-члены \rightarrow few-nomial \rightarrow fewnomial

Kushnirenko Hypotheses

Kushnirenko Hypothesis I

Topological complexity (e.g., Betti numbers) of the zero set of a real polynomial system can be controlled by the complexity of such a system (e.g., number of monomials) rather than by the degree (or the Newton polygon)

Kushnirenko Hypothesis II

The number of real zeros of a general real polynomial system

$$f_1(X_1,\ldots,X_n)=\cdots=f_n(X_1,\ldots,X_n)=0$$

can be bounded from above by the total number of nonzero terms of f_1, \ldots, f_n .

Kushnirenko Hypothesis III (corrected)

The number of nondegenerate real zeros of a real polynomial system

$$f_1(X_1, ..., X_n) = \cdots = f_n(X_1, ..., X_n) = 0$$

can be bounded from above by

$$O\left(\prod_{i=1}^n \# \operatorname{supp}(f_i)\right).$$

Kushnirenko Hypothesis IV

The number of nondegenerate real zeros of a real polynomial system

$$f_1(X_1, X_2) = f_2(X_1, X_2) = 0$$

can be bounded by $Z(\deg f_1, \# \operatorname{supp} f_2)$, where Z is some universal function.

Why hypotheses? In some Slavic languages, "hypothesis" is also used as "conjecture"

Examples

Descartes' rule of signs: A real univariate polynomial *f* has at most 1 + 2# supp \boldsymbol{f} real roots.

For a generic choice of coefficients, the system

$$\begin{cases} \mathbf{\alpha}_1 + \mathbf{\beta}_1 \mathbf{X} + \mathbf{\gamma}_1 \mathbf{Y} + \mathbf{\delta}_1 \mathbf{X} \mathbf{Y} \mathbf{Z}^d = 0 \\ \mathbf{\alpha}_2 + \mathbf{\beta}_2 \mathbf{X} + \mathbf{\gamma}_2 \mathbf{Y} + \mathbf{\delta}_2 \mathbf{X} \mathbf{Y} \mathbf{Z}^d = 0 \\ \mathbf{\alpha}_3 + \mathbf{\beta}_3 \mathbf{X} + \mathbf{\gamma}_3 \mathbf{Y} + \mathbf{\delta}_3 \mathbf{X} \mathbf{Y} \mathbf{Z}^d = 0 \end{cases}$$

has always d complex solutions, but at most 2 real solutions.

More on fewnomial history...

Kushnirenko's Letter to Prof. Sottile, and Appendix F of J. Tonelli-Cueto (2019), Condition and Homology in Semialgebraic Geometry, PhD thesis, Technische Universität Berlin.

Some Milestones

Sevastyanov (1978): Kushnirenko hypothesis IV is true, but original version of hypothesis III is false.

Khovanskii (1991): Kushnirenko hypotheses I and II are

Khovanskii's Theorem (1991)

The number of nondegenerate positive zeros of a fewnomial system $f_1 = \cdots = f_n = 0$ in n variables is at most

$$2^{\binom{t-1}{2}}(n+1)^{t-1}$$

where $t = \# \left(\bigcup_{i=1}^n \operatorname{supp} f_i \right)$.

Bihan, Sottile (2007): Improvement of Khovanskii's bound. Bihan, Sottile, Rojas (2008): Improvement of Khovanskii's bound for number of connected components

Bihan, Sottile (2009): Improvement of Khovanskii's bound for sum of Betti numbers

Bihan, Sottile (2007)

The number of nondegenerate positive zeros of a fewnomial system $f_1 = \cdots = f_n = 0$ in n variables is at most

$$3 \cdot 2^{\binom{t-n-1}{2}} \boldsymbol{n}^{t-n-1}$$

where $t = \# \left(\bigcup_{i=1}^n \operatorname{supp} f_i \right)$.

Koiran, Portier, Tavenas (2015): Sevastyanov's lost theorem is reproven

Bürgisser, Ergür, Tonelli-Cueto (2018): A probabilistic Kushnirenko hypothesis III is proven

Bürgisser, Ergür, Tonelli-Cueto (2018)

Let $f_i = \sum_{\alpha \in A} f_{i,\alpha} X^{\alpha}$ be random polynomials in n variables such that the $f_{i,\alpha}$ are independent centered Gaussian whose variance only depends on α . Then the expected number of nondegenerate positive zeros of $\mathfrak{f}_1 = \cdots = \mathfrak{f}_n =$ 0 is at most

Probabilistic Approach

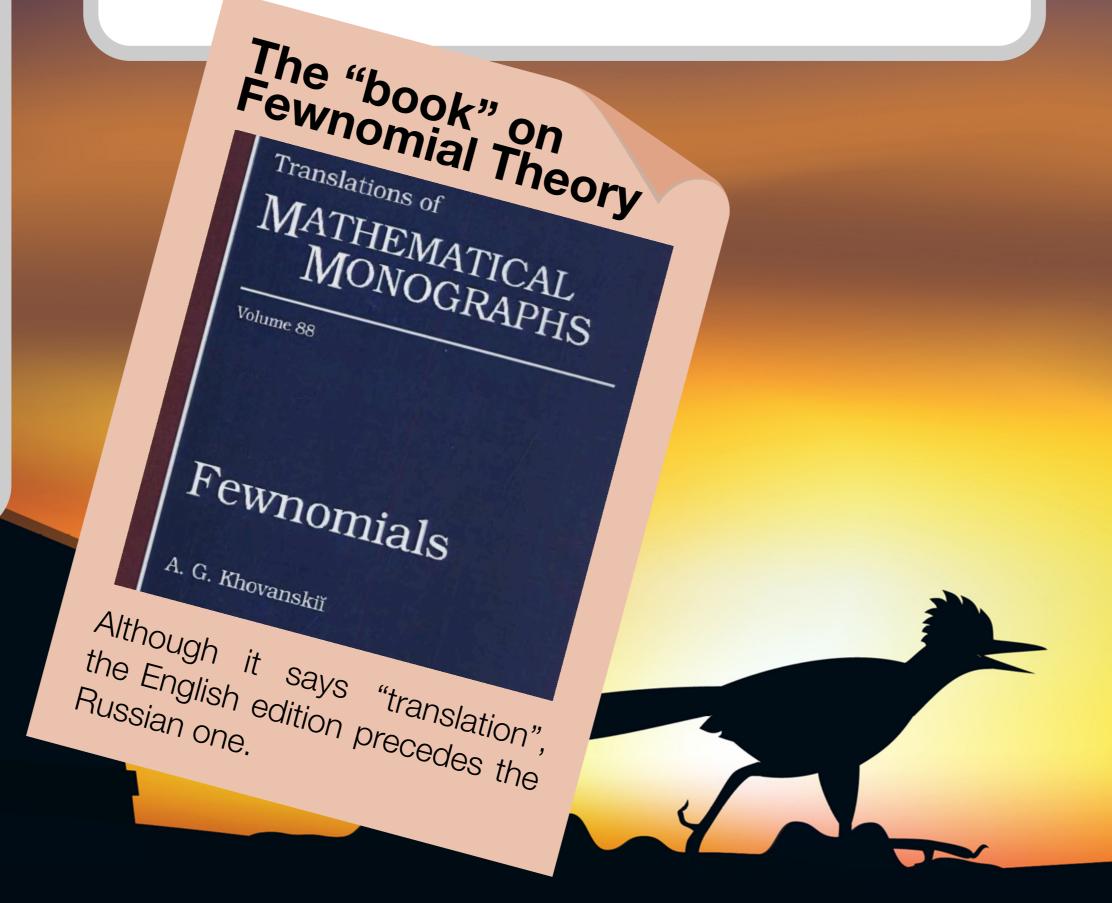
Probabilistic Reformulation

Let $f_i = \sum_{\alpha \in A} f_{i,\alpha} X^{\alpha}$ be random polynomials in n variables with some absolutely continuous distribution whose density does not vanish. Then

$$\sup_{\boldsymbol{\ell}} \left(\mathbb{E}_{\boldsymbol{\mathfrak{f}}} \# \mathcal{Z}_{\boldsymbol{r}}(\boldsymbol{\mathfrak{f}})^{\boldsymbol{\ell}} \right)^{\frac{1}{\boldsymbol{\ell}}} = \max_{\boldsymbol{f}} \# \mathcal{Z}_{\boldsymbol{r}}(\boldsymbol{f})$$

where $\mathcal{Z}_r(f)$ is the set of nondegenerate real zeros of f_1 = $\cdots = f_n = 0.$

Proof idea: A fewnomial system with maximum number of nondegenerate zeros is stable.



Condition Number

Let $f = (f_1, \dots, f_n)$ be a real polynomial system in n variables with f_i of degree at most d_i , its condition number is

$$\mathtt{C}(f) := \sup_{\mathbf{x} \in [-1,1]} \frac{\|f\|}{\max\{\|f(\mathbf{x})\|, \|\mathsf{D}_{\mathbf{x}}f^{-1}\Delta\|\}}$$

where $\Delta := \operatorname{diag}(\boldsymbol{d}_1, \ldots, \boldsymbol{d}_n)$

How to think about the condition number?

Roughly, $\mathbf{C}(\mathbf{f})$ is the inverse of the distance to the discriminant variety.

A new bound!

MAIN THEOREM (T.-C., Ts.; '22 +)

Let $f = (f_1, \ldots, f_n)$ be a real polynomial system in nvariables. Then

 $\#\mathcal{Z}(f, [-1, 1]^n) \le O(\log \mathbf{D} \max\{n \log \mathbf{D}, \log \mathbf{C}(f)\})^n$

where **D** is the maximum degree.

Corollary:

Well-posed Real Polynomial Systems

HAVE FEW REAL ZEROS

Probabilistic Consequences

PROB. THEOREM (T.-C., Ts.; '22 +)

Let $\mathfrak{f} = (\mathfrak{f}_1, \dots, \mathfrak{f}_n)$ be a random real fewnomial system in nvariables whose coefficients are independent and uniformly distributed in [-1, 1]. Then

$$\mathbb{E}_{\mathfrak{f}} \# \mathcal{Z}(\mathfrak{f}, \mathbb{R}^n_+)^{\ell} \leq O\left(n\ell \log^2 \mathbf{D}\right)^{n\ell}$$

where **D** is the maximum degree.

Corollary:

FEWNOMIAL SYSTEMS WITH MANY ZEROS

ARE VERY IMPROBABLE

We can cover a wide range of probabilistic assumptions

Algorithmic Consequences

PROOF IS FULLY CONSTRUCTIBLE!

Issue: Computing C(f) is expensive

ALG. THEOREM (T.-C., Ts.; '22 +)

There is a explicit partition \mathcal{B} of $[-1,1]^n$ into $O(\log \mathbf{D})^n$ boxes such that for all real polynomial system f = (f_1,\ldots,f_n) in n variables of degree at most \mathbf{D} and all $\mathbf{B}\in\mathcal{B}$, there is a polynomial

 $\Phi_{f,B}$

of degree $O(\max\{n \log \mathbf{D}, \log \mathbf{C}(f)\})$ such that

$$\#\mathcal{Z}(f,\mathsf{B}) \leq \#\mathcal{Z}(\varphi_{f,\mathsf{B}},\mathbb{R}^n).$$

Moreover, every real zero of f in B has a zero of $\phi_{f,B}$ that converges quadratically to it under Newton's method.

Proof idea: Well-conditioned polynomials are fast converging Taylor series

On the sphere...

We can cover also random polynomial system with the Weyl scaling. However, in this case, we are unable to cover fewnomials, and we get probabilistic bounds of the form $O\left(\sqrt{\mathbf{D}}\log\mathbf{D}\right)^{"}$ appears in the bound.