

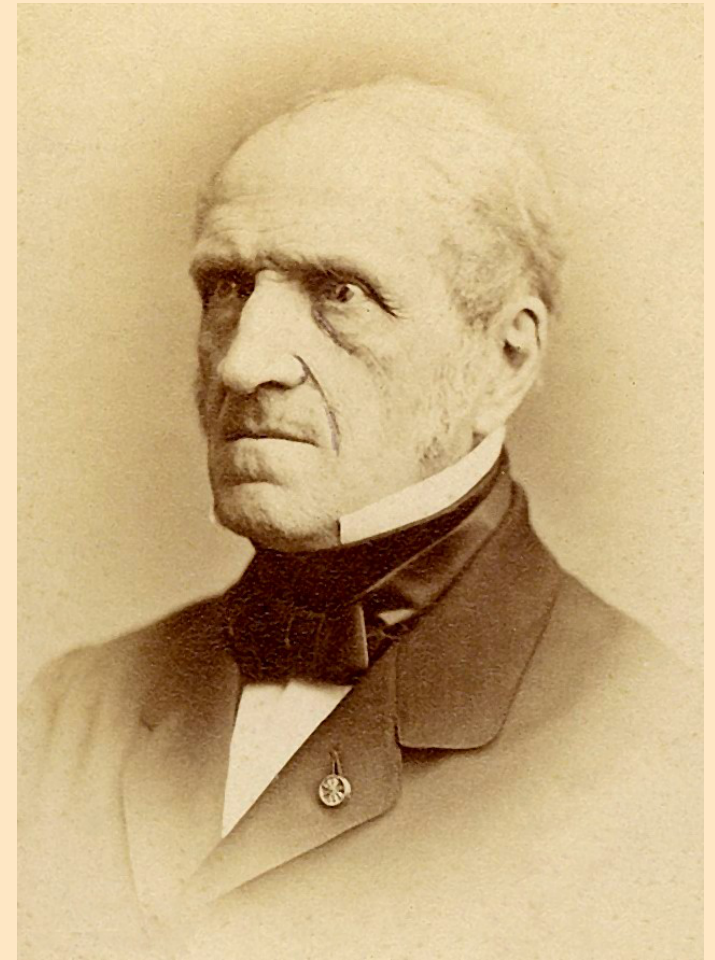
On this day...

the mathematician

Michel Chasles

would have

his 229th birthday



Famous for Chasles' identity:

$$\overrightarrow{ab} + \overrightarrow{bc} = \overrightarrow{ac}$$

A whole career
reduced to this!



Computing Numerically the Homology of a Semialgebraic Set

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Why do we care?

Central problem
in computational
semi alg. geometry
(with applications)

THE THEOREM:

There is an algorithm Hom that given a q -tuple $p \in \prod_{i=1}^q \mathbb{R}[X_1, \dots, X_n]_{\leq d_i}$ of polynomials and a semialgebraic formula Φ over p of size s , it computes the homology groups of

$$S_{\text{aff}}(p, \Phi) \subseteq \mathbb{R}^n$$

semialg. set
defined by (p, Φ)

in time

$$\mathcal{O}(qs(128nD\bar{K}_{\text{aff}}(p))^{10n(n+2)})$$

THE THEOREM: (Homogeneous Version)

There is an algorithm Hom that given a q -tuple $\mathcal{g} \in \prod_{i=1}^q \mathbb{R}[X_0, \dots, X_n]_{d_i}$ of homogeneous polynomials and a semialgebraic formula Φ over \mathcal{g} of size s , it computes the homology groups of

$$S(\mathcal{g}, \Phi) \subseteq S^n$$

spherical
semialg. set
def. by (\mathcal{g}, Φ)

in time

$$\mathcal{O}(qs(128nD\bar{H}(\mathcal{g}))^{10n(n+2)})$$

THE THEOREM: (Homogeneous Version +)

There is an algorithm Hom_ℓ that given a q -tuple $\mathcal{g} \in \prod_{i=1}^q \mathbb{R}[X_0, \dots, X_n]_{d_i}$ of homogeneous polynomials and a semialgebraic formula

Φ over \mathcal{g} of size s , it computes the

first ℓ homology groups of

$$S(\mathcal{g}, \Phi) \subseteq S^n$$

spherical
semialg. set
def. by (\mathcal{g}, Φ)

in time

$$\mathcal{O}(qs(128nD\bar{H}(\mathcal{g}))^{10n(\ell+2)})$$

THE PROB. THEOREM: (Homogeneous Ver)

Given a random KSS q -tuple F ,

i.e. $F_i = \sum (d_i^\alpha)^{1/2} c_{i,\alpha} X^\alpha$ with

the $c_{i,\alpha} \sim N(0,1)$ i.i.d., then Hom_e
runs with

$$\mathcal{O}\left(s(qD)^{\mathcal{O}(n^2 e)} q^{10n(e+1)}\right)$$

with prob $\geq 1 - 1/q$

Also smoothed analysis

and more robust probabilistic models

Single exponential in n
with high probability!

Improves state of the art!

Algorithm at a glance

0. Condition number estimation
1. Reduction to Lax Case
2. Reduction to Basic Cases
3. Basic case

CONDITION NUMBER

$$\bar{\kappa}(\delta, x) := \max_{\substack{I \subseteq [n] \\ \#I \leq n+1}} \frac{\|\delta_I\|}{\sqrt{\|\delta_I(x)\|^2 + \sigma_{|I|}(\Delta_I^{-1/2} D_x \delta_I)}}$$

where $\delta_I := (\delta_i)_{i \in I}$ $D_x \delta_I: T_x S^n \rightarrow \mathbb{R}^I$
 $\Delta = \text{diag}(d_i)_{i \in I}$ tangent map
 $\sigma_{|I|}$ $|I|$ -sing. val.

$$\bar{\kappa}(\delta) := \max_{x \in S^n} \bar{\kappa}(\delta, x) \in [1, \infty)$$

Interpretation of $\overline{K}(g)$

$$\overline{K}(g) < \infty$$

$\Leftrightarrow: \forall i, \mathcal{Z}_g(g_i)$ smooth hypersur.

& $\forall I \subseteq [g],$

$\bigcap_{i \in I} \mathcal{Z}_g(g_i)$ transversal

Note: If $|I| \geq n+1$, then $\bigcap_{i \in I} \mathcal{Z}_g(g_i)$ Transversal is empty

ESTIMATION OF $\bar{\mathcal{K}}(g)$

Prop. $S^n \ni x \mapsto 1/\bar{\mathcal{K}}(g, x)$ is D -Lipschitz

Cor. If $G \subseteq S^n$ satisfies

$$d_H(G, S^n) < \varepsilon,$$

& $\max \{ \bar{\mathcal{K}}(g, x) \mid x \in G \} \leq \varepsilon < 1/2$, then

$$\bar{\mathcal{K}}(g) \leq 2 \max \{ \bar{\mathcal{K}}(g, x) \mid x \in G \}$$

REDUCTION TO LAX CASE I

Gabrielov-Vorobjov Construction:

$$\Gamma B_{\delta, \varepsilon}(\mathcal{g}, \underline{\Phi}) = S(\mathcal{g}, \tilde{\underline{\Phi}})$$

$$\text{where in } \tilde{\underline{\Phi}} \quad \begin{cases} \mathcal{g}_i = 0 \rightsquigarrow |\mathcal{g}_i| \leq \varepsilon \|\mathcal{g}_i\|_w \\ \mathcal{g}_i > 0 \rightsquigarrow \mathcal{g}_i \geq \delta \|\mathcal{g}_i\|_w \\ \mathcal{g}_i < 0 \rightsquigarrow \mathcal{g}_i \leq -\delta \|\mathcal{g}_i\|_w \end{cases}$$

$$\Gamma B_{\underline{\delta}, \underline{\varepsilon}}(\mathcal{g}, \underline{\Phi}) = \bigcup_i \Gamma B_{\delta_i, \varepsilon_i}(\mathcal{g}, \underline{\Phi})$$

Gabrielov-Vorobjov Theorem:

IF $0 \ll \varepsilon_1 \ll \delta_1 \ll \varepsilon_2 \ll \delta_2 \ll \dots \ll \delta_m \ll 1$,

then $\exists H_k(\Gamma B_{\underline{\delta}, \underline{\varepsilon}}(\mathcal{g}, \underline{\Phi})) \rightarrow H_k(S(\mathcal{g}, \tilde{\underline{\Phi}}))$: iso for $k \leq m-1$
& surj. for $k = m$

REDUCTION TO LAX CASE II

Gabrielov-Vorobjov Construction:

$$\Gamma B_{\delta, \varepsilon}(\mathcal{F}, \underline{\Phi}) = S(\mathcal{F}, \tilde{\underline{\Phi}})$$

$$\text{where in } \tilde{\underline{\Phi}} \quad \begin{cases} \mathcal{F}_i = 0 \rightsquigarrow |\mathcal{F}_i| \leq \varepsilon \|\mathcal{F}_i\|_w \\ \mathcal{F}_i > 0 \rightsquigarrow \mathcal{F}_i \geq \delta \|\mathcal{F}_i\|_w \\ \mathcal{F}_i < 0 \rightsquigarrow \mathcal{F}_i \leq -\delta \|\mathcal{F}_i\|_w \end{cases}$$

$$\Gamma B_{\underline{\delta}, \underline{\varepsilon}}(\mathcal{F}, \underline{\Phi}) = \bigcup_i \Gamma B_{\delta_i, \varepsilon_i}(\mathcal{F}, \underline{\Phi})$$

Quantitative Gabrielov-Vorobjov Theorem:

$$\text{If } 0 < \varepsilon_1 < \delta_1 < \varepsilon_2 < \dots < \delta_m < 1/\sqrt{2}K(\mathcal{F})$$

then $\exists H_K(\Gamma B_{\underline{\delta}, \underline{\varepsilon}}(\mathcal{F}, \underline{\Phi})) \rightarrow H_K(S(\mathcal{F}, \underline{\Phi}))$: iso for $K \leq m-1$
& surj. for $K = m$

Reduction to basic case I

Main Idea:

- \mathcal{X}_i^∞ sample of $S(\mathcal{L}; \infty, 0)$
- If for all basis $\phi = \bigwedge_{i \in I} (\mathcal{L}; \infty; 0)$

$\bigcap \mathcal{X}_i^\infty$ good sample of $S(\mathcal{L}, \phi)$,
then for all semialg. formula Φ ,

$$\Phi(\text{Simp}(\mathcal{X}_i^\infty)) \simeq S(\mathcal{L}, \Phi)$$

Reduction to basic case II

Ingredients:

Explicit Functorial Nerve Theorem

$$\begin{array}{l} \pi: \check{C}_\varepsilon(x) \rightarrow B_\varepsilon(x) \text{ homology-equiv.} \\ \nearrow \check{C}_{\text{ech}} \quad \Sigma t: [x_i] \mapsto \Sigma t: x_i \end{array}$$

Homological Inclusion-Exclusion Transfer

$$\left. \begin{array}{l} \delta: X \rightarrow Y \text{ cont.} \\ X = \bigcup X_i, Y = \bigcup Y_i \\ \delta(X_i) \subseteq Y_i + \text{Tech.} \end{array} \right\} \begin{array}{l} \forall I, \delta: \bigwedge_{i \in I} X_i \rightarrow \bigwedge_{i \in I} Y_i \\ \text{homology equiv.} \\ \Rightarrow \delta: X \rightarrow Y \text{ hom. equiv.} \end{array}$$

Also hom. iso up to k and sur; For $(k+1)$ -homology

Basic Case TDA ingredients

Niyogi-Smale-Weinberger Thm

$$\exists d_H(X, X) < \varepsilon < \frac{1}{2} \gamma(X) \Rightarrow B_\varepsilon(x) \overset{\sim}{\hookrightarrow} X$$

\uparrow
Hausdorff
distance

\uparrow
hom. equiv

reach

$$\gamma(X) := \inf \{ r > 0 \mid \exists p \in \mathbb{R}^n, x, \tilde{x} \in X \text{ } d(p, X) = d(p, x) = d(p, \tilde{x}) = r \text{ } \& \text{ } x \neq \tilde{x} \}$$

Attali-Lientier-Salinas Thm

$$\exists d_H(X, X) < \varepsilon < \frac{1}{5} \gamma(X) \Rightarrow VR_\varepsilon(x) \overset{\sim}{\hookrightarrow} \check{C}_\varepsilon(x)$$

\uparrow
Vietoris-Rips

\uparrow
reach

Basic Case Approx Results

Reach Bound

$$\gamma(Y \wedge \bigwedge X_i) \geq \min_I \gamma(Y \wedge \bigwedge_{i \in I} \partial X_i)$$

Reduction to reach of boundaries

$$\phi \text{ basic semialg. form} \Rightarrow \exists D^{3/2} \bar{\kappa}(\delta) \gamma(S(\delta, \phi)) > 1$$

Condition-based bound of reach

Sampling Thm $g \in S^n$, $d_H(g, S^n) < r$

If $\sqrt{2} D^{1/2} \bar{\kappa}(\delta) r < 1$, then for all lax formulas

$$\Phi, \quad d_H(g \cap S_{D^{1/2}r}(\delta, \Phi), S(\delta, \Phi)) < \sqrt{2} D^{1/2} \bar{\kappa}(\delta)$$

Relaxation: $\delta \geq 0 \rightsquigarrow \delta \geq -D^{1/2}r, \dots$

Improvements

- Using $\|g\|_\infty := \max_i \max_{x \in S^n} |g_i(x)|$

instead of $\|g\|_w$ reduces in one n
the exponent of prob. run-time

- Can we implement it?

Eskervik asko

bere arretagatik!