Chapter 4

The Processor

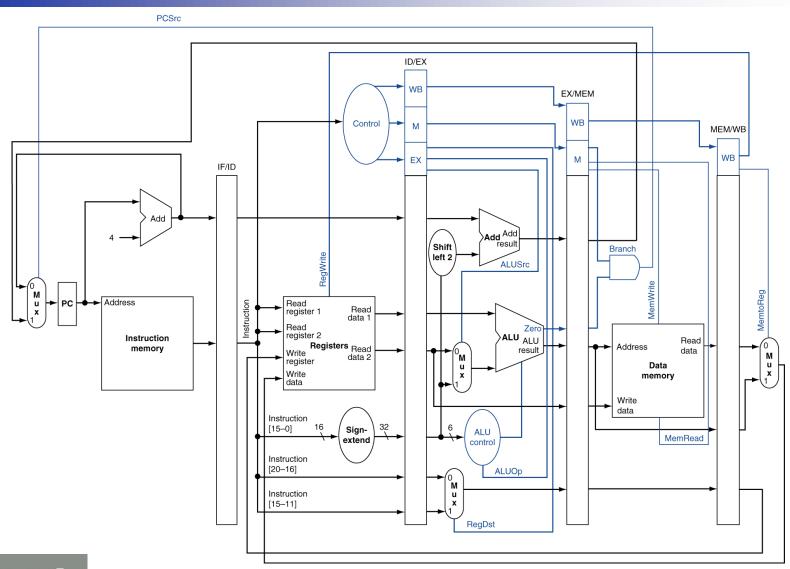
Fall 2018

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Announcement

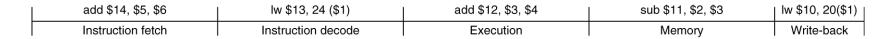
- Project #2
 - Due on Nov. 9
 - Studying pipelining with simulator
- Cheating on HW#1
 - Copying from Internet
 - Copying friends' HW
 - Modifying friends' HW
 - Changing variable names
 - Using slightly different algorithms
 - Start HW by yourselves and discuss later

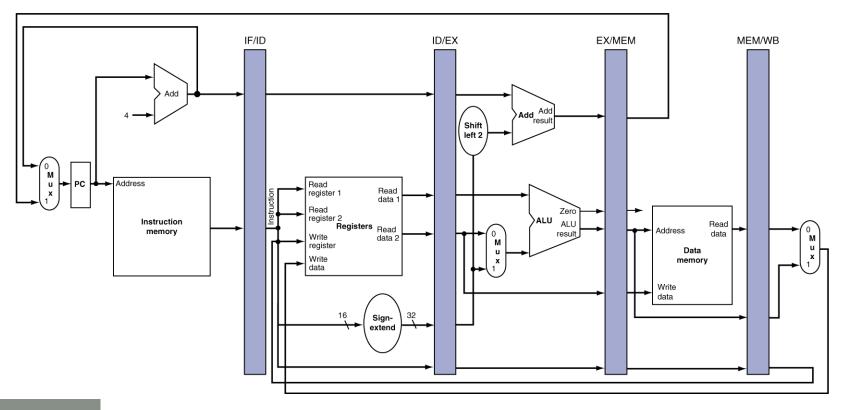
Pipelined Control



Single-Cycle Pipeline Diagram

Different instructions in stages





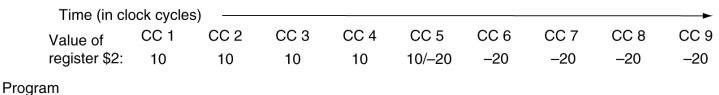
Data Hazards in ALU Instructions

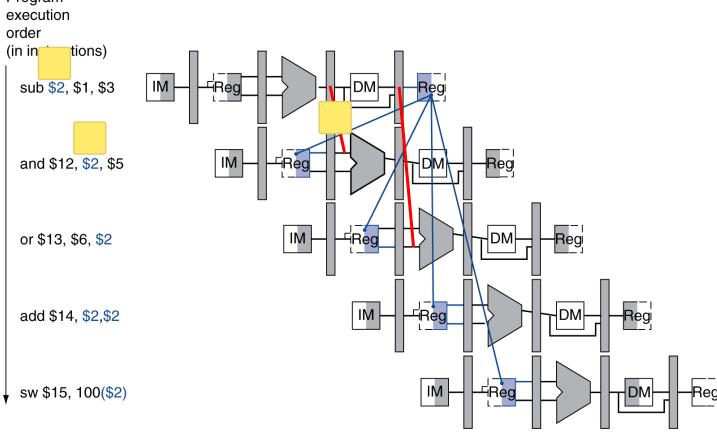
Consider this sequence:

```
sub $2, $1,$3
and $12,$2,$5
or $13,$6,$2
add $14,$2,$2
sw $15,100($2)
```

- We can resolve hazards with forwarding
 - How do we detect when to forward?

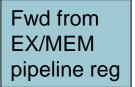
Dependencies & Forwarding





Detecting the Need to Forward

- Pass register numbers along pipeline
 - e.g., ID/EX.RegisterRs = register number for Rs sitting in ID/EX pipeline register
- ALU operand register numbers in EX stage are given by
 - ID/EX.RegisterRs, ID/EX.RegisterRt
- Data hazards when
 - 1a. EX/MEM.RegisterRd = ID/EX.RegisterRs
 - 1b. EX/MEM.RegisterRd = ID/EX.RegisterRt
 - 2a. MEM/WB.RegisterRd = ID/EX.RegisterRs
 - 2b. MEM/WB.RegisterRd = ID/EX.RegisterRt

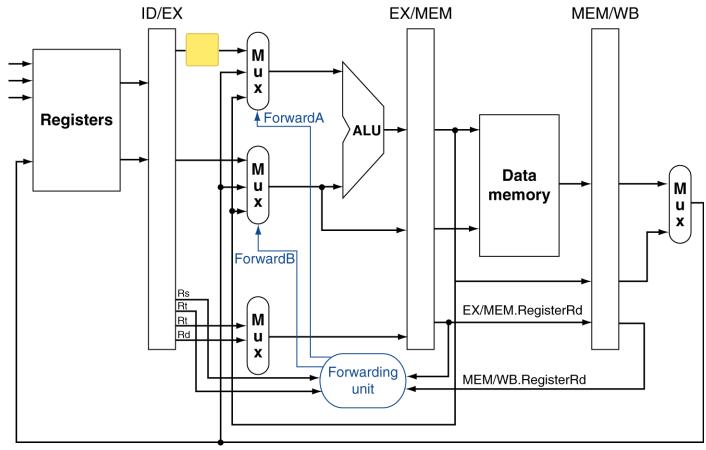


Fwd from MEM/WB pipeline reg

Detecting the Need to Forward

- But only if forwarding instruction will write to a register!
 - EX/MEM.RegWrite, MEM/WB.RegWrite
- And only if Rd for that instruction is not \$zero
 - EX/MEM.RegisterRd ≠ 0,
 MEM/WB.RegisterRd ≠ 0

Forwarding Paths



b. With forwarding



Forwarding Conditions

- EX hazard
 - if (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0) and (EX/MEM.RegisterRd = ID/EX.RegisterRs)) ForwardA = 10
 - if (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0) and (EX/MEM.RegisterRd = ID/EX.RegisterRt)) ForwardB = 10
- MEM hazard
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0) and (MEM/WB.RegisterRd = ID/EX.RegisterRs)) ForwardA = 01
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0) and (MEM/WB.RegisterRd = ID/EX.RegisterRt)) ForwardB = 01

Double Data Hazard

Consider the sequence:

```
add $1,$1,$2
add $1,$1,$3
add $1,$1,$4
```

- Both hazards occur
 - Want to use the most recent
- Revise MEM hazard condition
 - Only fwd if EX hazard condition isn't true

Revised Forwarding Condition

- MEM hazard
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0) and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0) and (EX/MEM.RegisterRd = ID/EX.RegisterRs))

and (MEM/WB.RegisterRd = ID/EX.RegisterRs))

ForwardA = 01

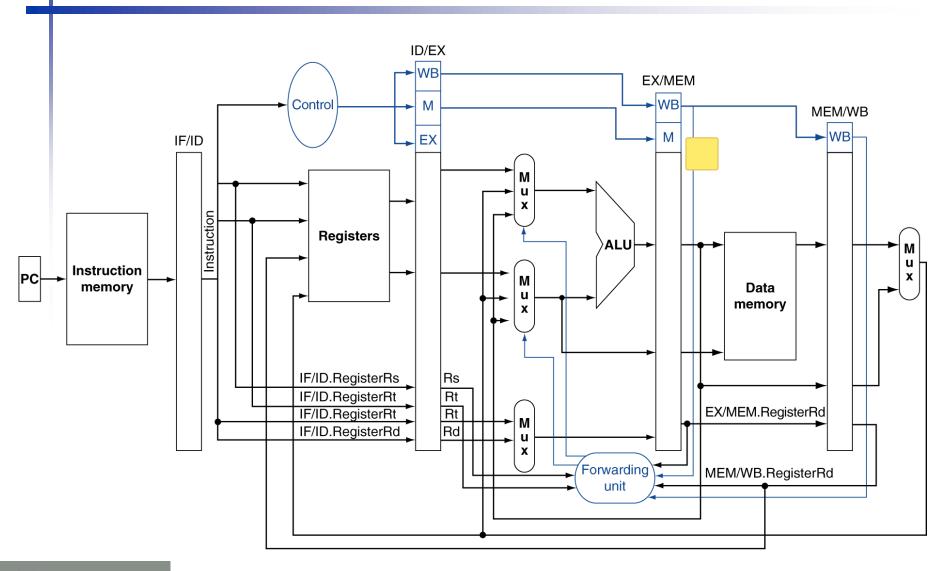
if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)

```
and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd = ID/EX.RegisterRt))
```

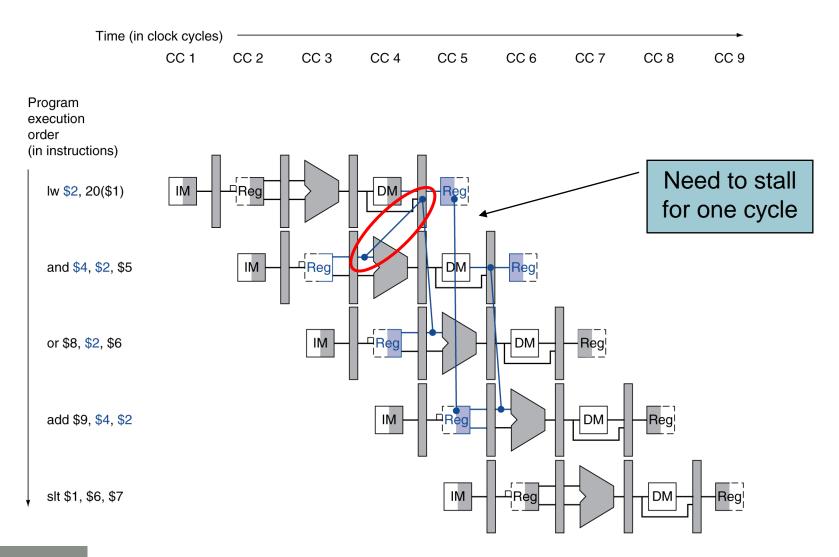
and (MEM/WB.RegisterRd = ID/EX.RegisterRt))

ForwardB = 01

Datapath with Forwarding



Load-Use Data Hazard



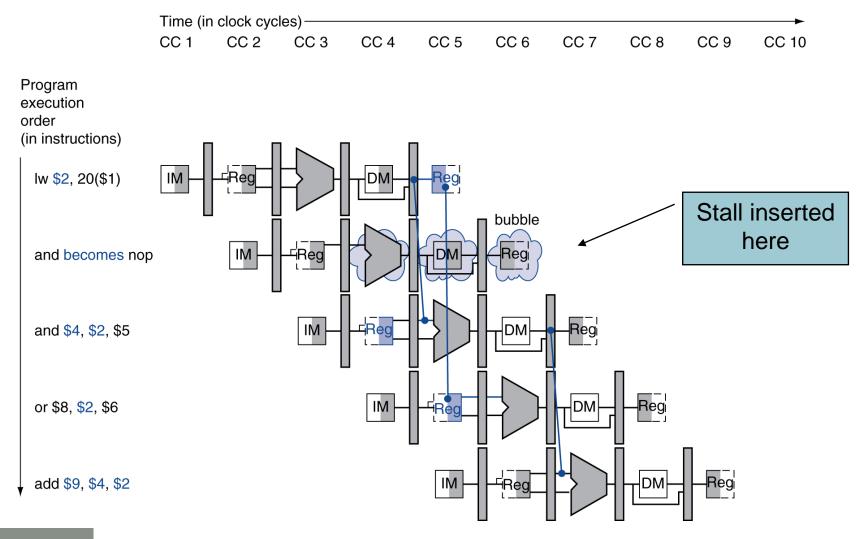
Load-Use Hazard Detection

- Check when using instruction is decoded in ID stage
- ALU operand register numbers in ID stage are given by
 - IF/ID.RegisterRs, IF/ID.RegisterRt
- Load-use hazard when
 - ID/EX.MemRead and ((ID/EX.RegisterRt = IF/ID.RegisterRs) or (ID/EX.RegisterRt = IF/ID.RegisterRt))
- If detected, stall and insert bubble

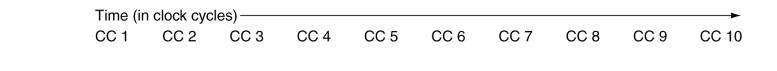
How to Stall the Pipeline

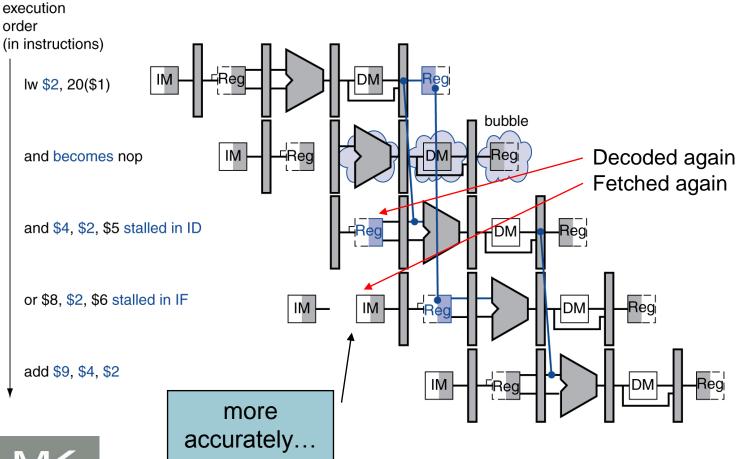
- Force control values in ID/EX register to 0
 - EX, MEM and WB do nop (no-operation)
- Prevent update of PC and IF/ID registers
 - Using instruction is decoded again
 - Following instruction is fetched again
 - 1-cycle stall allows MEM to read data for \(\)\rm \
 - Can subsequently forward to EX stage

Stall/Bubble in the Pipeline



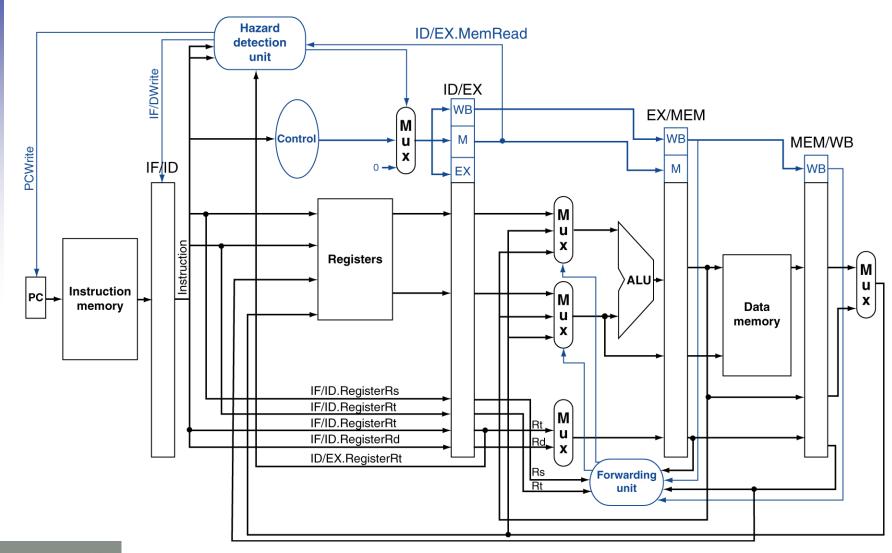
Stall/Bubble in the Pipeline





Program

Datapath with Hazard Detection



Stalls and Performance

The BIG Picture

- Stalls reduce performance
 - But are required to get correct results
- Compiler can arrange code to avoid hazards and stalls
 - Requires knowledge of the pipeline structure