

# **Chapter 4**

## **The Processor**

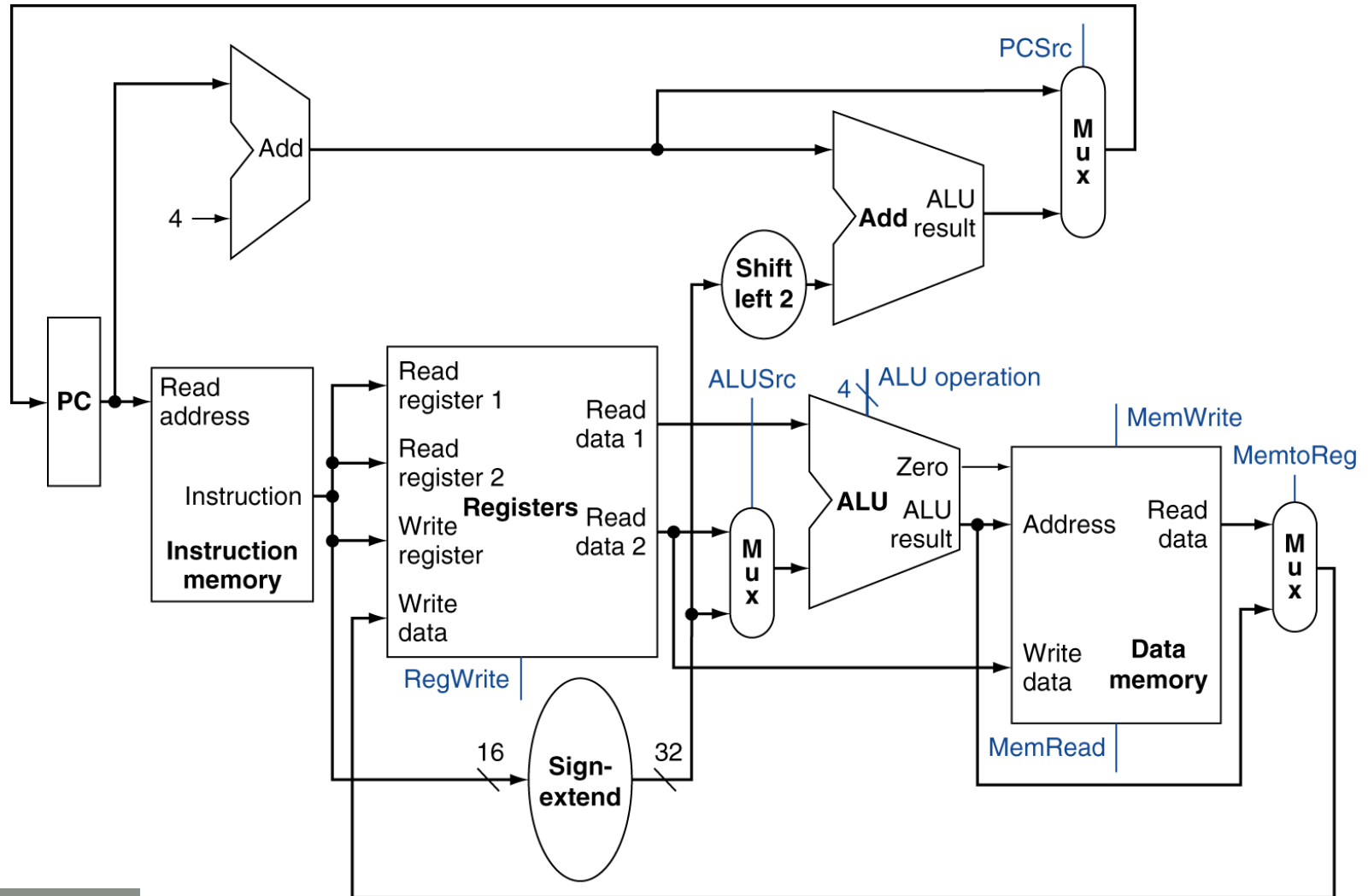
**Fall 2018**

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# Announcement

- Midterm exam
  - Place: this classroom
  - When: 16:00PM ~ 18:00PM on Oct. 16 (Tue.)
  - Coverage: all lecture notes and materials covered in class

# Full Datapath



# ALU Control

- ALU used for
  - Load/Store: F = add
  - Branch: F = subtract
  - R-type: F depends on funct field

ALU control	Function
0000	AND
0001	OR
0010	add
0110	subtract
0111	set-on-less-than
1100	NOR

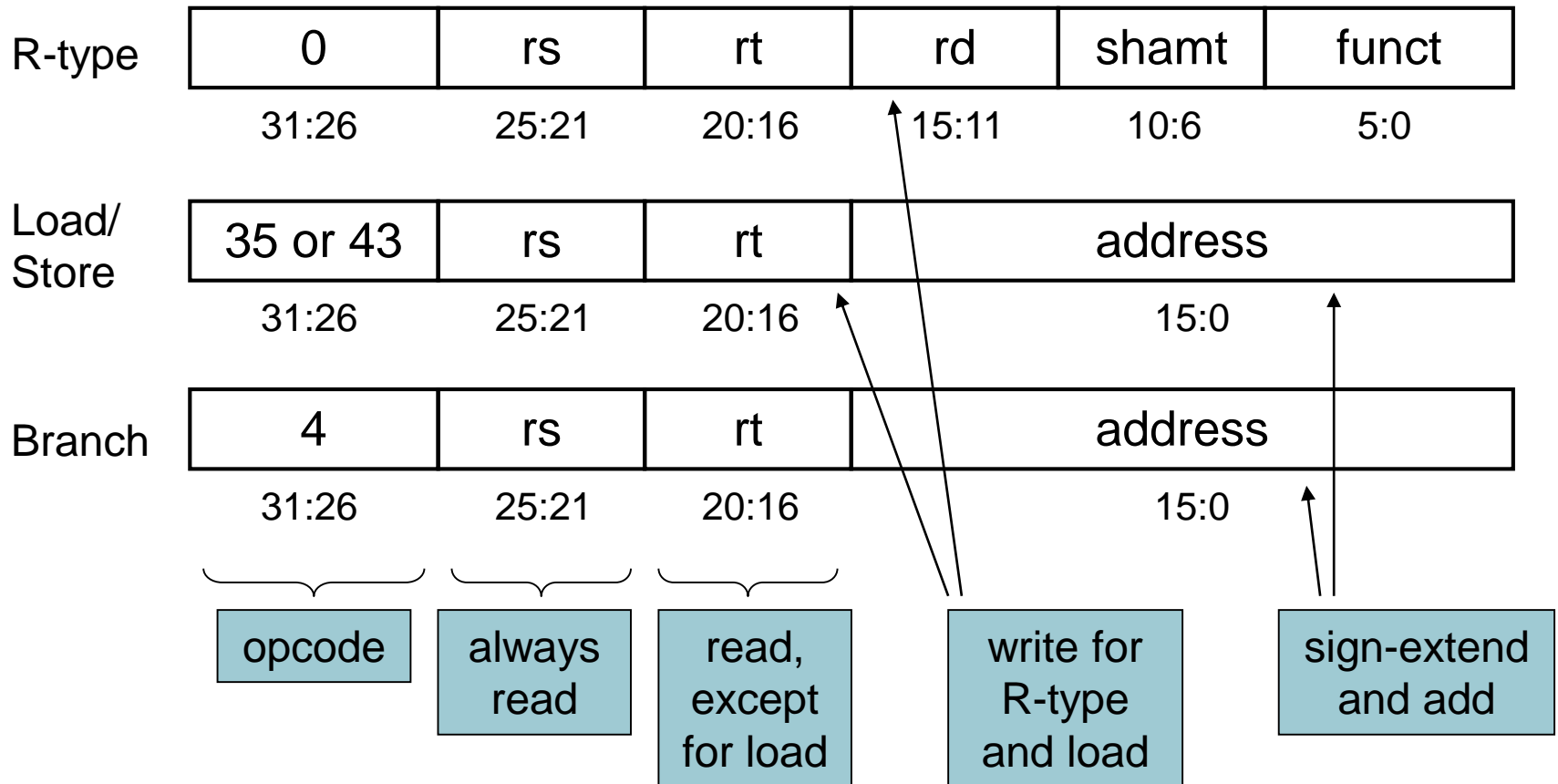
# ALU Control

- Assume 2-bit ALUOp derived from opcode
  - Combinational logic derives ALU control

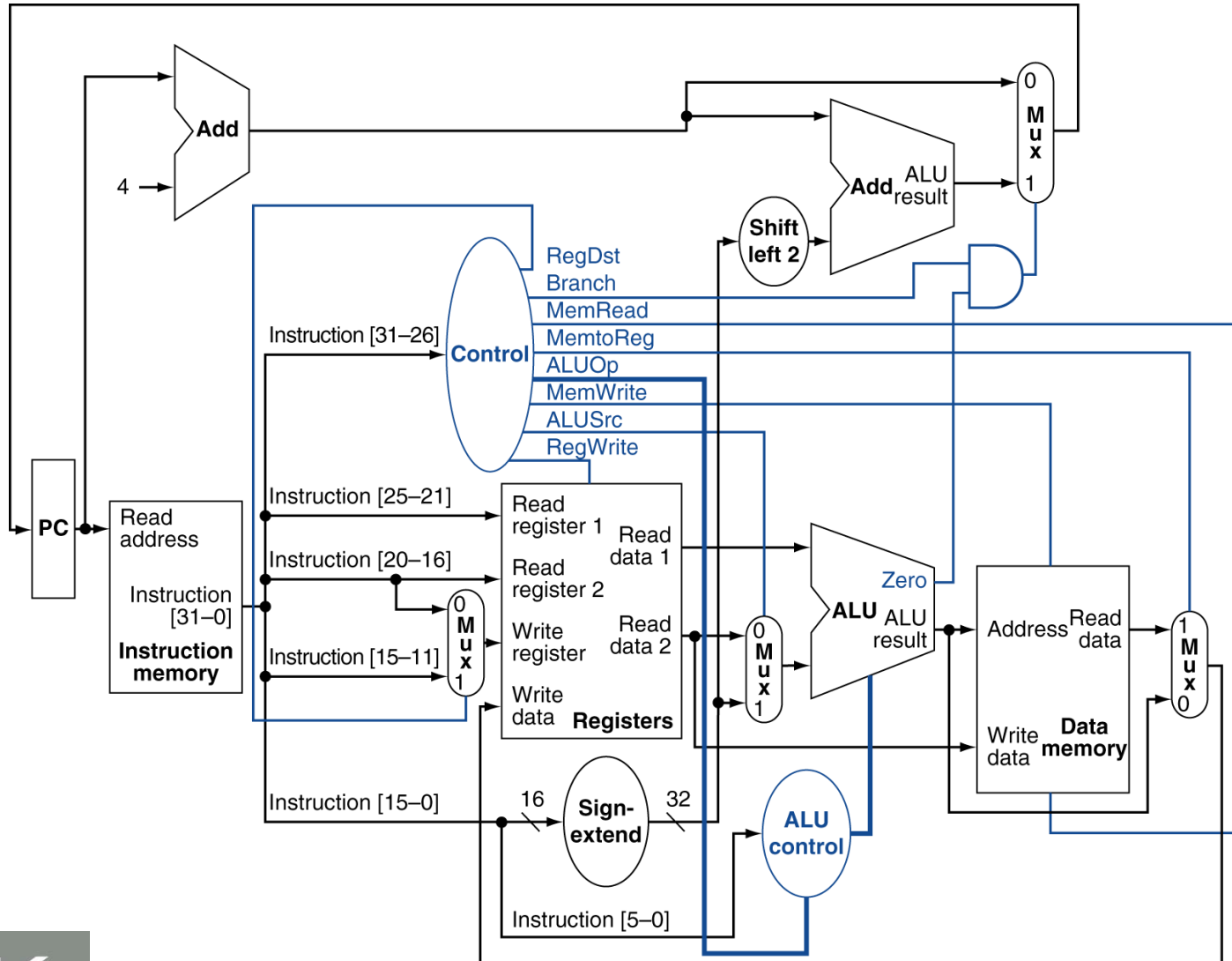
opcode	ALUOp	Operation	funct	ALU function	ALU control
lw	00	load word	XXXXXX	add	0010
sw	00	store word	XXXXXX	add	0010
beq	01	branch equal	XXXXXX	subtract	0110
R-type	10	add	100000	add	0010
		subtract	100010	subtract	0110
		AND	100100	AND	0000
		OR	100101	OR	0001
		set-on-less-than	101010	set-on-less-than	0111

# The Main Control Unit

- Control signals derived from instruction

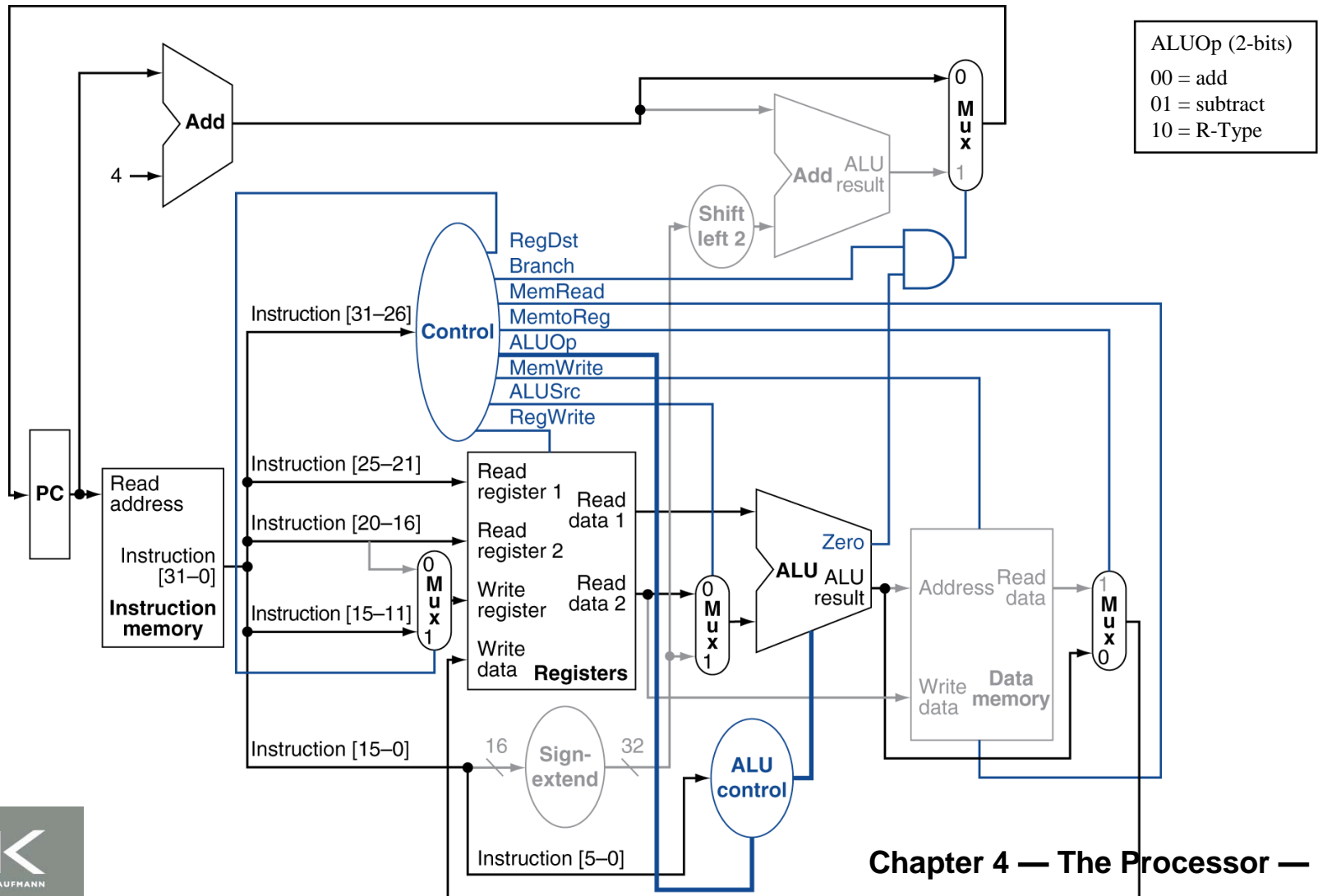


# Datapath With Control



# R-Type Instruction

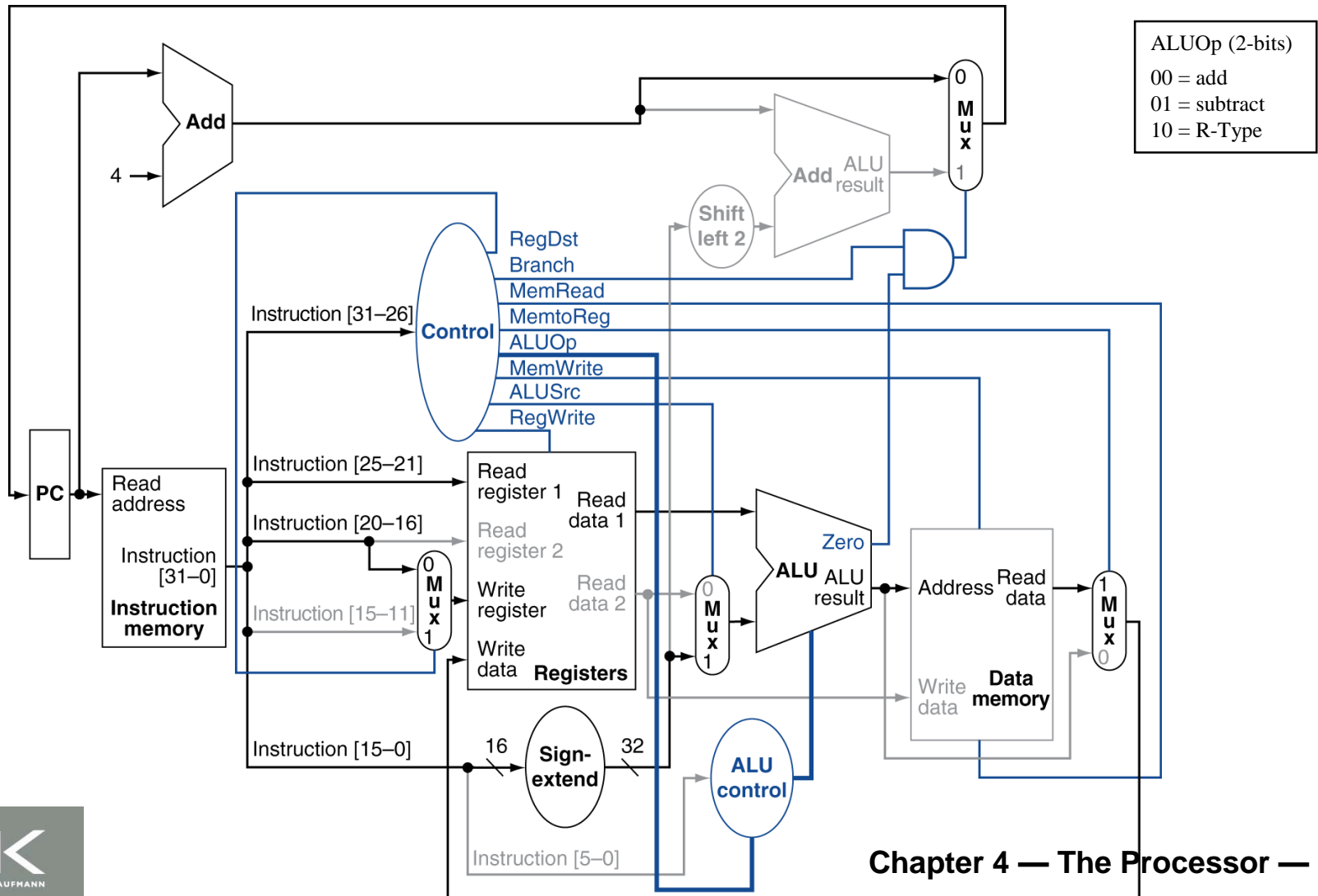
Instruction	RegDst	ALUSrc	MemtoReg	RegWrite	MemRead	MemWrite	Branch	ALUOp1	ALUOp0
<b>R-Format</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>
lw	0	1	1	1	1	0	0	0	0
sw	X	1	X	0	0	1	0	0	0
beq	X	0	X	0	0	0	1	0	1





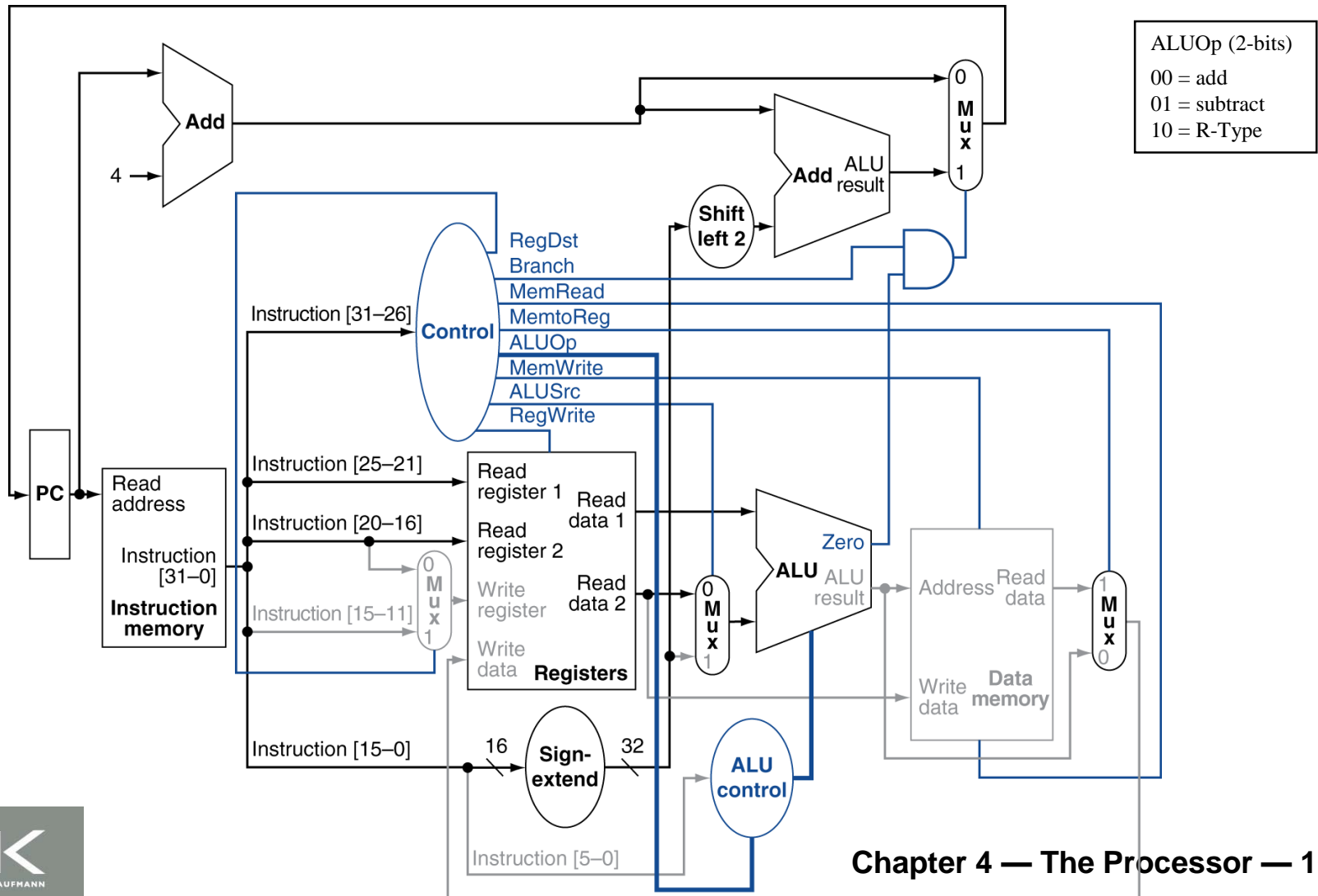
# Load Instruction

Instruction	RegDst	ALUSrc	MemtoReg	RegWrite	MemRead	MemWrite	Branch	ALUOp1	ALUOp0
R-Format	1	0	0	1	0	0	0	1	0
lw	0	1	1	1	1	0	0	0	0
sw	X	1	X	0	0	1	0	0	0
beq	X	0	X	0	0	0	1	0	1

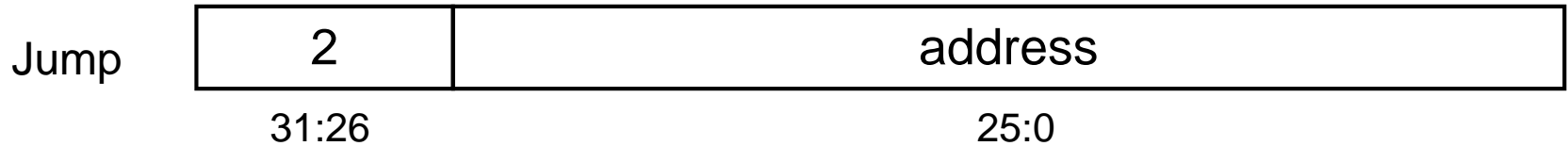


# Branch-on-Equal Instruction

Instruction	RegDst	ALUSrc	MemtoReg	RegWrite	MemRead	MemWrite	Branch	ALUOp1	ALUOp0
R-Format	1	0	0	1	0	0	0	1	0
lw	0	1	1	1	1	0	0	0	0
sw	X	1	X	0	0	1	0	0	0
beq	X	0	X	0	0	0	1	0	1

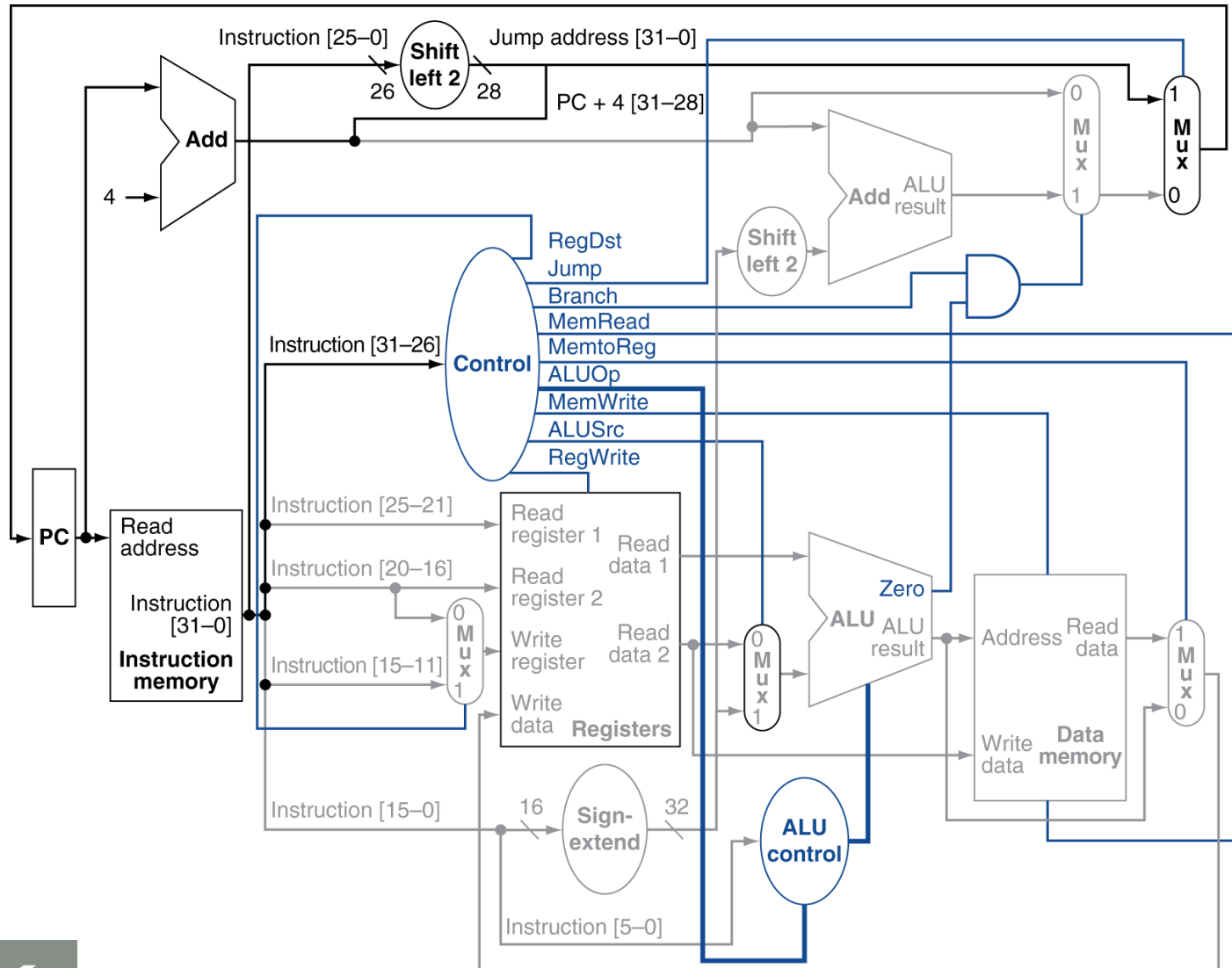


# Implementing Jumps



- Jump uses word address
- Update PC with concatenation of
  - Top 4 bits of old PC
  - 26-bit jump address
  - 00
- Need an extra control signal decoded from opcode

# Datapath With Jumps Added



# Performance Issues

- Longest delay determines clock period
  - Critical path: load instruction
  - Instruction memory → register file → ALU → data memory → register file
- Not feasible to vary clock period for different instructions
- Violates design principle
  - Making the common case fast
- We will improve performance by pipelining