

1. (a) 8 项 2 比较

$$C_8^2 = 28 \quad ??$$

如何简化



$$\sqrt{n} < n^{1.5} \quad \text{3rd, 8th}$$

$n^{1/2}$ 因为 $\frac{1}{2} < 1.5$

good observation

8 项简化 n^α 形式

$$n, n^2, n^3, n^4 \quad \text{不需要 } C_4^2$$

尝试去转化

$$\textcircled{1} 4^{2n} = 16^n$$

exponential

$$\alpha^n$$

$$O(\alpha^n) \geq$$

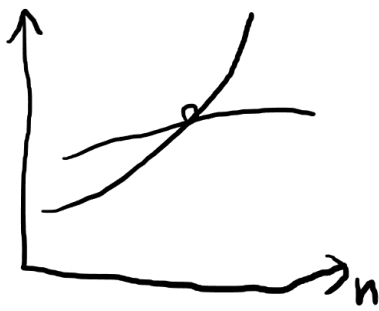
polynomial

$$n^\alpha$$

$$O(n^\alpha)$$

α is constant

$$\alpha^n \notin O(n^\alpha) \quad \& \quad n^\alpha \in O(\alpha^n)$$



写证明不用

$$\textcircled{2} \quad n^n \geq \alpha^n \geq 16^n$$

$$\textcircled{3} \quad n^{1.5}$$

$$\textcircled{4} \quad n^3 - n^2 \log n \in O(n^3)$$

$$\textcircled{5} \quad 9^{\log_3 n} = n^2$$

$$\textcircled{6} \quad \log_{10} 2^{(n^3)} = n^3 \log_{10} 2$$

$$\textcircled{7} \quad n! \geq \alpha^n$$

$$\textcircled{8} \quad n^{\frac{1}{2}}$$

$\alpha \cdot n \cdot \alpha \rightarrow \infty$

$$\textcircled{4} \text{ and } \textcircled{6} \quad \textcircled{4} \simeq \textcircled{6}$$

Hence

$$\sqrt{n} \leq n^{1.5} \leq 9^{\log_3 n} \leq \textcircled{4} \sim \textcircled{6} \leq 4^{2^n} \leq n! \leq n^n$$

(b) $g(n) = 2n$	closed-form	open
$g(n) = g(\frac{n}{2})$	recursive	$x^2 + y^2 = 1$

$$g(3) = 3, \quad g(n) = 3g(\frac{n}{3}) + n$$

$$g(5) \quad \times \quad g(6) \quad \times$$

$$n \in \{3^k : k \in \mathbb{Z}\}$$

$$g(3), g(9) = 9 + 9 = 18$$

$$g(27) = 3 \times 18 + 27 = 81$$

思考. $g(n) = 3g(\frac{n}{3}) + n$

$$= 3(3g(\frac{n}{9}) + n) + n$$

$$= 3^2 g(\frac{n}{3^2}) + 3n + n$$

$$\text{令 } n = 3^x$$

$$g(n) = g(3^x) = G(x) = 3g(3^{x-1})$$

$$= 3G(x-1) + 3^x$$

$$x \in \mathbb{Z}$$

$$\begin{aligned}
 a_n &= 3a_{n-1} + 3^n, \quad n=1, \quad G(1)=g(3)=3 \\
 &= 3(3a_{n-2} + 3^{n-1}) + 3^n \\
 &= 3^2 a_{n-2} + 2 \cdot 3^n \\
 &\vdots \\
 &= 3^{n-1} a_1 + (n-1) \cdot 3^n
 \end{aligned}$$

$$G(x) = 3^{x-1} \cdot 3 + (x-1) \cdot 3^x = x 3^x$$

$$g(n) = n \log_3 n$$

□

2. 0 1 1 2 3 5

$5 = 2 + 3$? 不能连续 $5 = 5$
 m

$4 = 2 + 2$? 不能重复 $4 = 1 + 3$
 m

任意 $m \geq 0$ 0 1 2 3 4 5
数学归纳法 A_0 \checkmark induction

if $A_{n-1} \checkmark$, to prove A_n .

Hence, A_n holds for every n .

Assume $m = F_{i_1} + \dots + F_{i_k}$ Induction
to prove $m+1$

也可以 Assume $\forall x < m$ holds,

to prove m

Strong Induction

$\triangle \triangle \triangle \triangle \triangle$ m
 holds

$\forall k: m = F_1 \dots F_k \quad F \leq m$

Fibonacci $F_0 \ F_1 \ \dots \ F_k \ F_{k+1}$
 $\leq m <$

$\exists k, \text{ s.t. } F_k \leq m < F_{k+1}$

$F_{k+1} = F_k + F_{k-1} > m \Leftrightarrow m - F_k < F_{k-1}$
 $\Delta \parallel m^*$

$m = F_k + ?$

$m = F_k + m^* \quad m^* < F_{k-1}$

hypothesis $m^* = \sum_i F_{i_j} \quad \text{holds}$

$(F_{i_j}) + F_k$? \nwarrow not consecutive
 no duplicate

3. $f(n) = n$, $g(n) = n^2$

hint :

Extension: $f(n) \in O(g(n))$

is it true that $\log f(n) \in O(\log g(n))$
?