

# Rental Bike Share Problem

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## 1 Introduction

RentalBike is a bike sharing company specialized in short-term bicycle rentals along various areas of a city, offering different categories of bicycles. Clients may rent bicycles for any duration and may deliver the rented bicycles in a place different from where the rental started. However, the delivery place must be in one of the areas currently operated by RentalBike. Since clients are able to change the bicycles' original position, RentalBike must often redistribute the surplus of bicycles according to the expected demand of each area. The decision of how many bicycles will be redistributed from one area to another is based on expected profits for each possible volume of bicycles to be moved. Therefore, before starting the relocation of bicycles, it is necessary to optimize the volume of bicycles to be moved from one area to the others. The following sections describe the problem to be solved. Read the input descriptions and assumptions before doing the tasks.

## 2 Formulation Problem

This section presents the proposed formulation for the RentalBike optimization problem.

### 2.1 Definition

- $I \rightarrow$  Set of bicycle categories.
- $J \rightarrow$  Set of destination areas.
- $s_i \rightarrow$  Number of surplus bicycles of category  $i$  in the origin area.
- $c_i \rightarrow$  Space occupied by a bicycle of category  $i$  in the truck.
- $T \rightarrow$  Total truck capacity.
- $p_{ij,k} \rightarrow$  Expected profit from relocating the  $k$ -th bicycle of category  $i$  to area  $j$ .

- $x_{i,j} \rightarrow$  Decision variable representing the number of bicycles of category  $i$  sent to area  $j$ , where  $x_{i,j}$  is a non-negative integer.

## 2.2 Objective Function

The objective function is given by the maximization of the total profit from the redistribution of bicycles:

$$\max \sum_{i \in I} \sum_{j \in J} \sum_{k=1}^{x_{i,j}} p_{ij,k} \quad (1)$$

## 2.3 Decision Vars

This section describes all decision variable. As quantidades de bicicletas enviadas para cada área devem ser inteiras e não negativas:

$$x_{i,j} \in \mathbb{Z}^+, \quad \forall i \in I, \forall j \in J \quad (2)$$

## 2.4 Parameters

- $s_i \rightarrow$  Quantity of bicycles of category  $c$  available (surplus) in the origin area.
- $c_i \rightarrow$  Space occupied by a bicycle of category  $i$  in the truck.
- $p_{ij,k} \rightarrow$  Expected profit from relocating the  $k$ -th bicycle of category  $i$  to area  $j$ .
- $x_{i,j} \rightarrow$  Decision variable representing the number of bicycles of category  $i$  sent to area  $j$ , where  $x_{i,j}$  is a non-negative integer.

# 3 Constraints

This section describes all the constraints of the model.

## 3.1 Truck Capacity

The total space occupied by the transported bicycles cannot exceed the truck's capacity:

$$\sum_{i \in I} \sum_{j \in J} c_i x_{i,j} \leq T \quad (3)$$

### 3.2 Constraint of available bicycles

The number of bicycles shipped from each category cannot exceed the available stock.

$$\sum_{i \in I} x_{ij} \leq s_i, \quad \forall i \in I \quad (4)$$

### 3.3 Mathematical Model

The problem involves the redistribution of bicycles of different categories from a specific area to others, aiming to maximize the expected profit from the re-allocation while subject to the truck's capacity constraint. Thus, the complete model can be given by:

$$\max \sum_{i \in I} \sum_{j \in J} \sum_{k=1}^{x_{i,j}} p_{ij,k} \quad (5)$$

$$\sum_{i \in I} \sum_{j \in J} c_i x_{i,j} \leq T \quad (6)$$

$$\sum_{i \in I} x_{ij} \leq s_i, \quad \forall i \in I \quad (7)$$

$$x_{i,j} \in \mathbb{Z}^+, \quad \forall i \in I, \forall j \in J \quad (8)$$

### 3.4 Computational Aspects

In this study, the computational implementation was carried out using the Python programming language 3.9 and Gurobi 12.0.1 as solver. The simulations were performed on an AMD Ryzen 7 series processor, 3.2 GHz, with 32 GB of RAM.

## 4 Methodology

This section describes the methods and techniques used to address the research problem and achieve the study's objectives.

### 4.1 Problem Description

In this study, we investigate the bicycles relocation problem. The objective is to determine how many bicycles of each category should be moved to each area so that the sum of expected profit is maximized.

## 4.2 Data Collection and Processing

The data used in this study was kindly provided by SOT (Shipping Optimization Team) of Mercado Livre. The data preprocessing included:

- **Handling of Missing Values:** Imputation was performed by replacing missing values with zero.

## 4.3 Mathematical Formulation

The problem was initially modeled as an optimization problem. Subsequently, heuristic approaches were used to solve the same problem. The formulation can be expressed as follows:

$$\max \sum_{i \in I} \sum_{j \in J} \sum_{k=1}^{x_{i,j}} p_{ij,k} \quad (9)$$

s.t

$$\sum_{i \in I} \sum_{j \in J} c_i x_{i,j} \leq T \quad (10)$$

$$\sum_{i \in I} x_{ij} \leq s_i, \quad \forall i \in I \quad (11)$$

$$x_{i,j} \in \mathbb{Z}^+, \quad \forall i \in I, \forall j \in J \quad (12)$$

Onde:

- $I \rightarrow$  Set of bicycle categories.
- $J \rightarrow$  Set of destination areas.
- $s_i \rightarrow$  Number of surplus bicycles of category  $i$  in the origin area.
- $c_i \rightarrow$  Space occupied by a bicycle of category  $i$  in the truck.
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- $x_{i,j} \rightarrow$  Decision variable representing the number of bicycles of category  $i$  sent to area  $j$ , where  $x_{i,j}$  is a non-negative integer.

## 4.4 Algorithms and Methods Used

To solve the problem, a mathematical model was initially formulated in its exact form. Then, algorithmic approaches were implemented to validate and compare the obtained results. The main steps of the method are:

1. **Implementation of the Exact Algorithm** – First, the exact algorithm is implemented considering the capacity constraint suggested in the challenge.
2. **Implementation of the Exact Algorithm without Capacity Constraint** – Aiming for a sensitivity analysis, the exact algorithm is re-implemented considering possible scenarios for the capacity constraint.
3. **Implementation of Approximate Algorithm 1** – The first heuristic algorithm is implemented. The objective is to compare it with the exact implementation and use it as a baseline for other heuristic algorithms.
4. **Implementation of the Following Approximate Algorithms** – Subsequently, two more approximate algorithms are implemented (totaling three algorithms) for comparison with the baseline and the exact algorithms.

## 4.5 Tools and Development Environment

The implementation and experimentation were conducted using:

- **Programming Language:** Python 3.9.
- **Libraries Used:** NumPy, Pandas, Gurobi 12.0.1.
- **Computational Environment:** The algorithms were initially implemented in a Jupyter notebook, and later, scripts were created for each implemented algorithm.

## 4.6 Performance Evaluation

To evaluate the results, we used the following metrics:

- **Objective Function Value:** The value of the objective function directly reflects the amount of profit obtained in the optimization process.
- **Allocation by Category:** Although a more subjective metric, the allocation by category reflects the coherence of resource allocation within the optimization process.

The quantitative results will be compared and presented in the following sections.

## 5 Results

The following section presents the results obtained from the proposed algorithms, providing a comprehensive analysis of its performance across different scenarios. The discussion of these findings aims to contextualize their significance within the broader literature, emphasizing both the strengths and potential limitations of the proposed methods.

Table 1 presents the results obtained by the exact algorithm. Considering the problem in its exact form and with the capacity constraint "relaxed," the algorithm achieved a value of 78,866.16 in its objective function. This value corresponds to a total optimized profit of USD 78.866.16.

Category	Area 1	Area 2	Area 3	Area 4	Area 5	Area 6	Area 7
Child	0.0	0.0	91.0	103.0	0.0	78.0	0.0
Adult	6.0	0.0	92.0	97.0	0.0	75.0	0.0
Electric	0.0	0.0	110.0	72.0	0.0	97.0	0.0
Racing	0.0	0.0	93.0	103.0	0.0	71.0	0.0
Mountain	0.0	0.0	104.0	108.0	0.0	70.0	0.0
Tricycle	0.0	0.0	106.0	109.0	0.0	64.0	0.0

Table 1: Bicycle Reallocation by Category and Area - Exact Algorithm

Next, we present the allocation values achieved by the VNS heuristic. The result obtained by the algorithm reached an objective function value of approximately 74,132.8581. This value corresponds to a total optimized profit of USD 74.132.85. Table 2 presents the allocation results by category and area.

Category	Area 1	Area 2	Area 3	Area 4	Area 5	Area 6	Area 7
Child	10	0	79	103	0	80	0
Adult	43	0	55	97	0	75	0
Electric	47	0	109	77	0	46	0
Racing	85	0	45	105	0	32	0
Mountain	44	0	94	100	0	44	0
Tricycle	32	0	106	109	0	32	0

Table 2: Bicycle Reallocation by Category and Area - VNS

The results obtained by the tabu search heuristic are presented in Table 3. The result obtained by the algorithm reached an objective function value of approximately 78,063.9850. This value corresponds to a total optimized profit of USD 78.063.98.

<b>Category</b>	<b>Area 1</b>	<b>Area 2</b>	<b>Area 3</b>	<b>Area 4</b>	<b>Area 5</b>	<b>Area 6</b>	<b>Area 7</b>
Child	0	0	91	103	0	78	0
Adult	0	0	91	97	0	82	0
Electric	14	0	91	77	0	97	0
Racing	0	0	93	103	0	71	0
Mountain	0	0	100	103	0	79	0
Tricycle	10	0	106	95	0	68	0

Table 3: Bicycle Reallocation by Category and Area - Tabu Search

<b>Algorithm</b>	<b>Exact Algorithm</b>	<b>Tabu Search</b>	<b>VNS</b>
<b>Obj. Value</b>	78.866.16	78.063.9850	74.132.85

Table 4: Comparison of Objective Values Obtained