

Comparison of Radix Sort and Tim Sort

An Abstract

Tonmoy

CSE, RUET

January 23, 2024

Presentation Outline

- 1 Abstract
- 2 Introduction to Sorting Algorithms
 - Radix Sort Overview
 - Tim Sort Overview
- 3 Background Study
 - Radix Sort with Counting Sort
 - Counting Sort
 - Tim Sort with Insertion Sort and Merge Sort
 - Insertion Sort
 - Merge Sort
- 4 Implementation of Radix Sort
- 5 Implementation of Tim Sort
- 6 Results and Analysis
- 7 Results and Analysis
- 8 Conclusion
- 9 References

Abstract

[8] This paper presents a comparative analysis of Radix Sort and Tim Sort, two popular sorting algorithms. Radix Sort, a linear time, non-comparative algorithm, is explored with an enhanced version incorporating Counting Sort. Tim Sort, a hybrid algorithm combining Merge Sort and Insertion Sort, is detailed with a focus on its adaptability. The study aims to provide insights into their strengths, weaknesses, and optimal use cases.

Introduction to Sorting Algorithms

Radix Sort Overview

Radix Sort is a linear time, non-comparative sorting algorithm that operates on integers by processing individual digits. It exhibits excellent performance for fixed-size integer keys and is known for its simplicity and efficiency in certain scenarios.[7]

The time complexity of Radix Sort is linear, $O(kn)$, where k is the number of digits in the input integers.

Introduction to Sorting Algorithms

Tim Sort Overview

Tim Sort, a hybrid sorting algorithm derived from Merge Sort and Insertion Sort, is designed for real-world data. Its ability to adapt to various scenarios, along with its efficiency in handling partially ordered data, makes it a popular choice in practice.

The average-case time complexity of Tim Sort is $O(n \log n)$, where n is the size of the input array.[3]

Background Study

Radix Sort Pseudocode with Counting Sort[3]

Algorithm 1: Radix Sort with Counting Sort

Data: Array A , Number of digits d

```
1 for  $i \leftarrow 1$  to  $d$  do  
2   | Use Counting Sort to sort array  $A$  on digit  $i$ ;  
3 end
```

Background Study

Counting Sort Pseudocode

Algorithm 2: Counting Sort

Data: Array A , Range of elements k

Result: Sorted array A

```
1  $C \leftarrow$  array of size  $k$  initialized with zeros;  
2 for  $i \leftarrow 1$  to  $n$  do  
3    $C[A[i]] \leftarrow C[A[i]] + 1$ ;  
4 end  
5 for  $i \leftarrow 1$  to  $k$  do  
6    $C[i] \leftarrow C[i] + C[i - 1]$ ;  
7 end  
8 for  $i \leftarrow n$  to 1 do  
9    $B[C[A[i]]] \leftarrow A[i]$ ;  
10   $C[A[i]] \leftarrow C[A[i]] - 1$ ;  
11 end
```

Background Study

Tim Sort Pseudocode with Insertion Sort and Merge Sort[5]

Algorithm 3: Tim Sort with Insertion Sort and Merge Sort

Data: Array A

Result: Sorted array A

- 1 Divide the array into small blocks using Insertion Sort;
 - 2 **while** *blocks to merge* **do**
 - 3 Merge adjacent blocks using Merge Sort;
 - 4 **end**
-

Background Study

Insertion Sort Pseudocode

Algorithm 4: Insertion Sort

Data: Array A **Result:** Sorted array A

```
1 for  $i \leftarrow 2$  to  $n$  do
2    $key \leftarrow A[i];$ 
3    $j \leftarrow i - 1;$ 
4   while  $j > 0$  and  $A[j] > key$  do
5     Swap  $A[j + 1]$  and  $A[j];$ 
6      $j \leftarrow j - 1;$ 
7   end
8    $A[j + 1] \leftarrow key;$ 
9 end
```

Background Study

Merge Sort Pseudocode (Part 1)

Algorithm 5: Merge Sort

Data: Array A , Indices p , q , r **Result:** Sorted array A

```
1  $n_1 \leftarrow q - p + 1, n_2 \leftarrow r - q;$ 
2 Create arrays  $L[1..n_1 + 1]$  and  $R[1..n_2 + 1];$ 
3 for  $i \leftarrow 1$  to  $n_1$  do
4    $L[i] \leftarrow A[p + i - 1];$ 
5 end
6 for  $j \leftarrow 1$  to  $n_2$  do
7    $R[j] \leftarrow A[q + j];$ 
8 end
9  $L[n_1 + 1] \leftarrow \infty, j \leftarrow 1;$ 
10  $R[n_2 + 1] \leftarrow \infty, i \leftarrow 1;$ 
```

Background Study

Merge Sort Pseudocode (Part 2)

Algorithm 5: Merge Sort (continued)

Data: Array A , Indices p , q , r **Result:** Sorted array A (continued)

```
1 for  $k \leftarrow p$  to  $r$  do
2   if  $L[i] \leq R[j]$  then
3      $A[k] \leftarrow L[i];$ 
4      $i \leftarrow i + 1;$ 
5   else
6      $A[k] \leftarrow R[j];$ 
7      $j \leftarrow j + 1;$ 
```

Implementation of Radix Sort[2]

- **Input:** Array A
- **Output:** Sorted array A using Radix Sort
- Find the maximum number of digits, d , in elements of A
- **For** i from 1 to d :
 - Use Counting Sort to sort array A based on the i -th digit

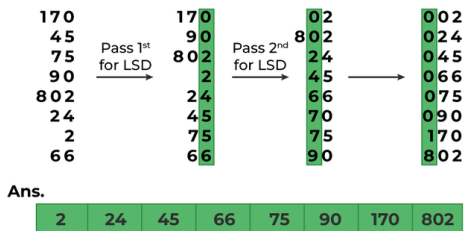


Figure: Radix Sort Visualization

Implementation of Tim Sort

- **Input:** Array A
- **Output:** Sorted array A using Tim Sort
- Divide the array into small blocks using Insertion Sort
- **While** there are blocks to merge:
 - Merge adjacent blocks using Merge Sort

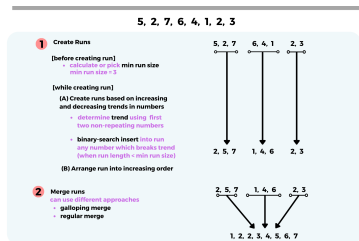


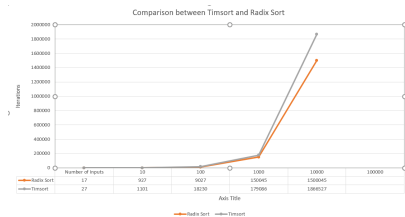
Figure: Tim Sort Visualization

Comparison Table

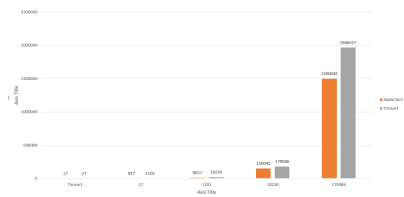
Table: Comparison between Timsort and Radix Sort

Topics	Radix Sort	Timsort
Time Complexity	$O(nk)$	$O(n \log n)$
Space Complexity	$O(n + k)$	$O(n)$
Application	Sorting items of fixed-length. Sorting strings where product of the length of the largest item is and number of item is not too large.	[6]Timsort is a hybrid sorting algorithm derived from merge sort and insertion sort. Where is a balance between time and space complexity is desired.

Time complexity



(a) Chart Comparison



(b) Bar Comparison

Figure: Comparison between Timsort and Radix Sort.

Time complexity

Table: Comparison between Radix Sort's and Timsort's iterations

Number of Inputs	Radix Sort	Timsort
10	17	27
100	927	1101
1000	9027	18230
10000	150045	179086
100000	1500045	1666527

- 1 Radix Sort's time complexity is almost similar to $O(n * k)$, where n stands for the number of items in the data set, and k is the length of the longest number in the data set.
- 2 Timsort's time complexity is almost similar to $O(n \log n)$, where n stands for the number of items in the data set[4].

Conclusion

- In conclusion, Radix Sort and Timsort are both effective sorting algorithms, each with its own strengths and use cases.

Radix Sort:

- Radix Sort demonstrates efficiency in scenarios where the range of input values is not significantly larger than the number of elements.
- With a time complexity of $O(n \cdot k)$, where n is the number of elements and k is the length of the longest digit.[1]

Timsort:

- Timsort, a hybrid sorting algorithm derived from merge sort and insertion sort, exhibits a time complexity of $O(n \log n)$.

References I

- [1] Khalid Suleiman et al. Al-Kharabsheh. Review on sorting algorithms: A comparative study. *International Journal of Computer Science and Security (IJCSS)*, 2013.
- [2] Arne Anderson and S. N. Implementing radix sort. *Experimental Algorithmic (JEA)*, 1998.
- [3] J. Hammad. A comparative study between various sorting algorithms. 15(March):6, 2015. March.
- [4] Md. Khairullah. Enhancing worst sorting algorithms. *International Journal of Advanced Science and Technology*, 2013.
- [5] S. K. P. Megha Jain. Bitonic sorting algorithm: A review. Volume 113 - No. 13:2, 2015. March.
- [6] S. B. Sahni. Gpu radix sort for. *IEEE*, 2011.

References II

- [7] Avinash Shukla and A. K. K. S. Review of radix sort & proposed modified radix sort for heterogeneous data set in distributed computing environment. *Journal Name*, 2(5):6, 2012.
- [8] A. K. Verma. A new approach for sorting list to reduce execution time. page 7, 2013. 13 October.