Comparison of Radix Sort and Tim Sort

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Presentation Outline

Abstract

[8]This paper presents a comparative analysis of Radix Sort and Tim Sort, two popular sorting algorithms. Radix Sort, a linear time, non-comparative algorithm, is explored with an enhanced version incorporating Counting Sort. Tim Sort, a hybrid algorithm combining Merge Sort and Insertion Sort, is detailed with a focus on its adaptability. The study aims to provide insights into their strengths, weaknesses, and optimal use cases.

Introduction to Sorting Algorithms

Radix Sort Overview

Radix Sort is a linear time, non-comparative sorting algorithm that operates on integers by processing individual digits. It exhibits excellent performance for fixed-size integer keys and is known for its simplicity and efficiency in certain scenarios.[7]

The time complexity of Radix Sort is linear, $O(kn)$, where k is the number of digits in the input integers.

Introduction to Sorting Algorithms

Tim Sort Overview

Tim Sort, a hybrid sorting algorithm derived from Merge Sort and Insertion Sort, is designed for real-world data. Its ability to adapt to various scenarios, along with its efficiency in handling partially ordered data, makes it a popular choice in practice.

The average-case time complexity of Tim Sort is $O(n \log n)$, where n is the size of the input array.[3]

Background Study

Radix Sort Pseudocode with Counting Sort[3]

Algorithm 1: Radix Sort with Counting Sort

Data: Array A, Number of digits d

1 for $i \leftarrow 1$ *to* d **do**

2 $\quad$ | $\quad$ Use Counting Sort to sort array A on digit i;

3 end

Background Study
Counting Sort Pseudocode

Algorithm 2: Counting Sort

Data: Array A, Range of elements k

Result: Sorted array A

1   $C \leftarrow$ array of size k initialized with zeros;

2 **for** $i \leftarrow 1$ *to* n **do**

3 | $C[A[i]] \leftarrow C[A[i]] + 1$;

4 **end**

5 **for** $i \leftarrow 1$ *to* k **do**

6 | $C[i] \leftarrow C[i] + C[i-1]$;

7 **end**

8 **for** $i \leftarrow n$ *to* 1 **do**

9 | $B[C[A[i]]] \leftarrow A[i]$;

10 | $C[A[i]] \leftarrow C[A[i]] - 1$;

11 **end**

Comparison of Radix Sort and Tim Sort
Background Study
  Tim Sort with Insertion Sort and Merge Sort

Background Study
Tim Sort Pseudocode with Insertion Sort and Merge Sort[5]

Algorithm 3: Tim Sort with Insertion Sort and Merge Sort

Data: Array A

Result: Sorted array A

1 Divide the array into small blocks using Insertion Sort;

2 while *blocks to merge* **do**

3 | Merge adjacent blocks using Merge Sort;

4 end

Background Study
Insertion Sort Pseudocode

Algorithm 4: Insertion Sort

Data: Array A

Result: Sorted array A

1 for $i \leftarrow 2$ *to* n **do**

2 $key \leftarrow A[i]$;

3 $j \leftarrow i - 1$;

4   **while** $j > 0$ *and* $A[j] > key$ **do**

5   Swap $A[j + 1]$ and $A[j]$;

6 $j \leftarrow j - 1$;

7   **end**

8 $A[j + 1] \leftarrow key$;

9 end

Background Study

Merge Sort Pseudocode (Part 1)

Algorithm 5: Merge Sort

Data: Array A, Indices p, q, r

Result: Sorted array A

1   $n_1 \leftarrow q - p + 1, n_2 \leftarrow r - q$;

2   Create arrays $L[1..n_1 + 1]$ and $R[1..n_2 + 1]$;

3   **for** $i \leftarrow 1$ *to* n_1 **do**

4 | $L[i] \leftarrow A[p + i - 1]$;

5   **end**

6   **for** $j \leftarrow 1$ *to* n_2 **do**

7 | $R[j] \leftarrow A[q + j]$;

8   **end**

9   $L[n_1 + 1] \leftarrow \infty, j \leftarrow 1$;

10   $R[n_2 + 1] \leftarrow \infty, i \leftarrow 1$;

Background Study
Merge Sort Pseudocode (Part 2)

Algorithm 5: Merge Sort (continued)

Data: Array A, Indices p, q, r

Result: Sorted array A (continued)

1 **for** $k \leftarrow p$ *to* r **do**

2   **if** $L[i] \leq R[j]$ **then**

3 $A[k] \leftarrow L[i]$;

4 $i \leftarrow i + 1$;

5   **else**

6 $A[k] \leftarrow R[j]$;

7 $j \leftarrow j + 1$;

Implementation of Radix Sort[2]

- **Input:** Array A
- **Output:** Sorted array A using Radix Sort
- Find the maximum number of digits, d, in elements of A
- **For** i from 1 to d:
    - Use Counting Sort to sort array A based on the i-th digit
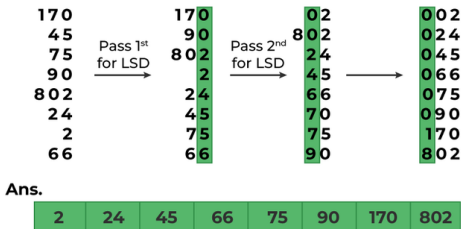

Figure: Radix Sort Visualization

Implementation of Tim Sort

- **Input:** Array A
- **Output:** Sorted array A using Tim Sort
- Divide the array into small blocks using Insertion Sort
- **While** there are blocks to merge:
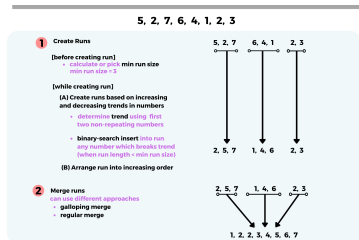 - Merge adjacent blocks using Merge Sort


Figure: Tim Sort Visualization

Comparison Table

Table: Comparison between Timsort and Radix Sort

Topics	Radix Sort	Timsort
Time Complexity	$O(nk)$	$O(n \log n)$
Space Complexity	$O(n + k)$	$O(n)$
Application	Sorting items of fixed-length. Sorting strings where product of the length of the largest item is and number of item is not too large.	[6]Timsort is a hybrid sorting algorithm derived from merge sort and insertion sort. Where is a balance between time and space complexity is desired.

Time complexity
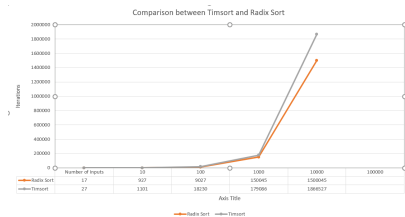
(a) Chart Comparison
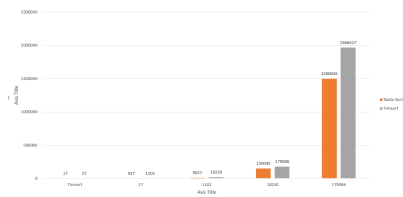
(b) Bar Comparison

Figure: Comparison between Timsort and Radix Sort.

Time complexity

Table: Comparison between Radix Sort's and Timsort's iterations

Number of Inputs	Radix Sort	Timsort
10	17	27
100	927	1101
1000	9027	18230
10000	150045	179086
100000	1500045	1666527

1. Radix Sort's time complexity is almost similar to $O(n * k)$, where n stands for the number of items in the data set, and k is the length of the longest number in the data set.

2. Timsort's time complexity is almost similar to $O(n\log n)$, where n stands for the number of items in the data set[4].

Conclusion

- In conclusion, Radix Sort and Timsort are both effective sorting algorithms, each with its own strengths and use cases.
  Radix Sort:
  - Radix Sort demonstrates efficiency in scenarios where the range of input values is not significantly larger than the number of elements.
  - With a time complexity of $O(n \cdot k)$, where n is the number of elements and k is the length of the longest digit.[1]

  Timsort:
  - Timsort, a hybrid sorting algorithm derived from merge sort and insertion sort, exhibits a time complexity of $O(n \log n)$.

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