Lecture 07

CSE 1201: Data Structure

Insertion Sort

Md. Shahid Uz Zaman Dept. of CSE, RUET

The Sorting Problem

Input: A sequence of *n* numbers $[a_1, a_2, ..., a_n]$.

Output: A permutation or reordering $[a'_1, a'_2, \dots, a'_n]$ of the input

sequence such that $a'_1 \le a'_2 \le ... \le a'_n$.

An instance of the Sorting Problem:

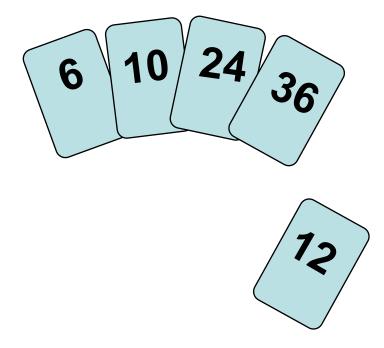
Input: A sequence of 6 number [31, 41, 59, 26, 41, 58].

Expected output for given instance:

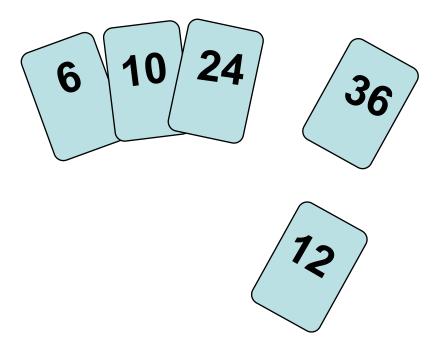
Expected

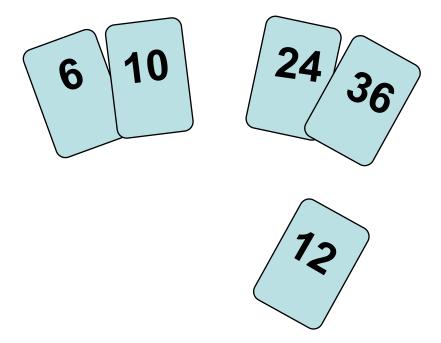
Output: The permutation of the input [26, 31, 41, 41, 58, 59].

- Idea: like sorting a hand of playing cards
 - Start with an empty left hand and the cards facing down on the table.
 - Remove one card at a time from the table, and insert it into the correct position in the left hand
 - compare it with each of the cards already in the hand, from right to left
 - The cards held in the left hand are sorted
 - these cards were originally the top cards of the pile on the table



To insert 12, we need to make room for it by moving first 36 and then 24.





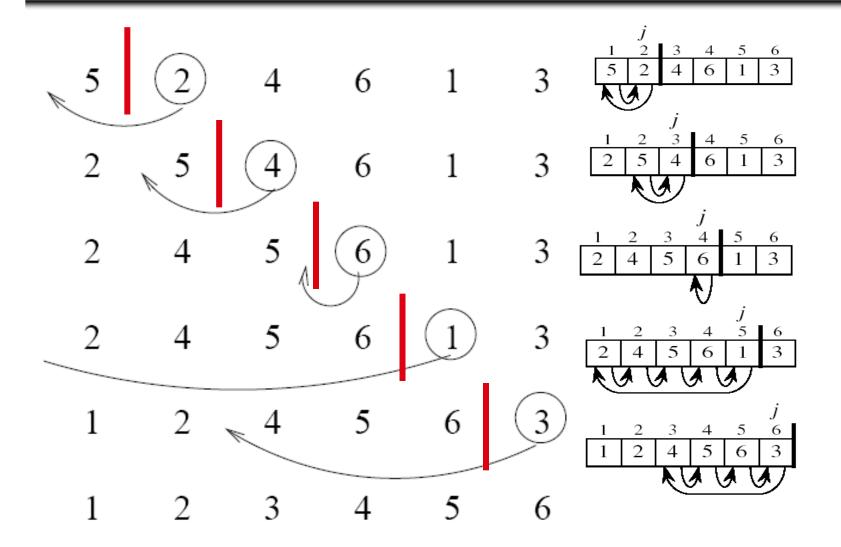
input array

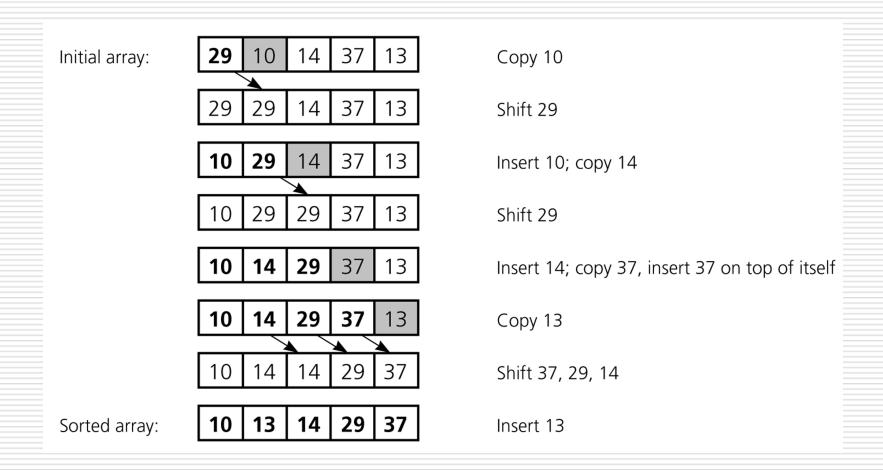
5 2 4 6 1

sorted

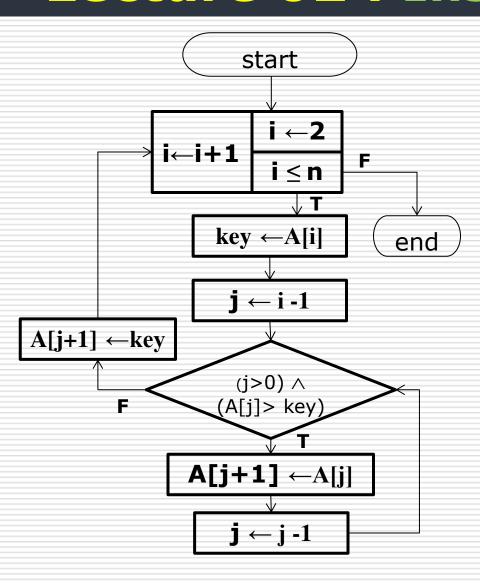
at each iteration, the array is divided in two sub-arrays:

left sub-array right sub-array unsorted





An insertion sort of an array of five integers

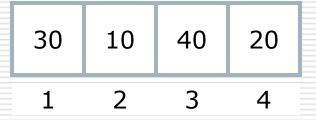


```
InsertionSort(A, n) {
  for i = 2 to n \{
     key = A[i]
     j = i - 1;
     while (j > 0) and
  (A[j] > key) {
       A[j+1] = A[j]
       j = j - 1
     A[j+1] = key
```

Flowchart

Coding

```
InsertionSort(A, n) {
  for i = 2 to n {
    key = A[i]
     j = i - 1;
     while (j > 0) and (A[j] > key) {
          A[j+1] = A[j]
          j = j - 1
     A[j+1] = key
```



```
i = \emptyset j = \emptyset key = \emptyset

A[j] = \emptyset A[j+1] = \emptyset
```



```
InsertionSort(A, n) {
    for i = 2 to n {
        key = A[i]
        j = i - 1;
        while (j > 0) and (A[j] > key) {
            A[j+1] = A[j]
            j = j - 1
        }
        A[j+1] = key
    }
}
```

```
30 10 40 20

1 2 3 4

↑ ↑

j i
```

```
i = 2 j = 1 key = 10
A[j] = 30 A[j+1] = 10
```

```
InsertionSort(A, n) {
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}
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            j = j - 1
        }
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    }
}
```

An Example: Insertion Sort

```
30 30 40 20
0 1 2 3 4
1 i
```

David Luebke

$$i = 2$$
 $j = 0$ key = 10
A[j] = \emptyset A[j+1] = 30

```
InsertionSort(A, n) {
    for i = 2 to n {
        key = A[i]
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            j = j - 1
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        A[j+1] = key
    }
}
```

```
10 30 40 20
0 1 2 3 4
↑ ↑
i
```

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i = 2 j = 0 key = 10

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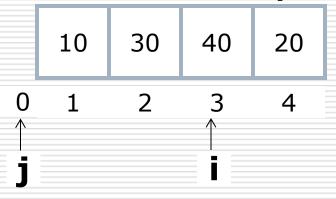
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10 30 40 20
0 1 2 3 4
1 i
```

```
i = 3 j = 0 key = 10

A[j] = \emptyset A[j+1] = 10
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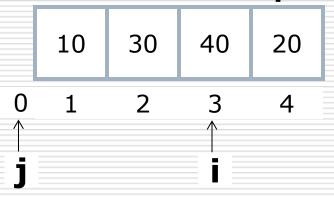
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```
\Rightarrow
```

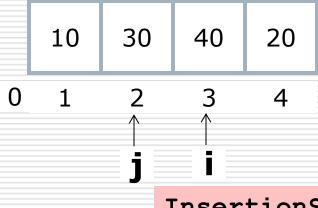


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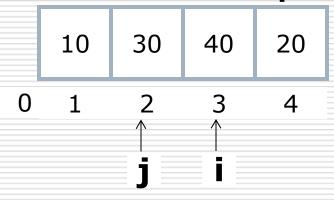


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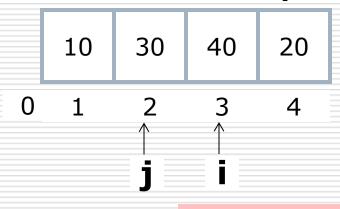
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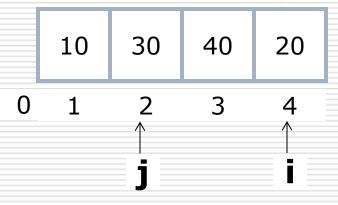
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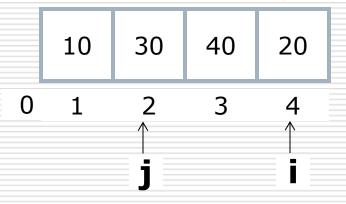


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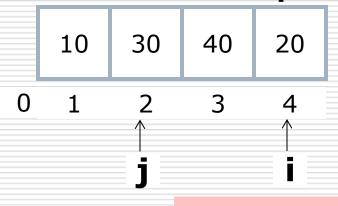


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```
\Longrightarrow
```

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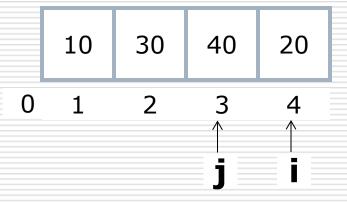


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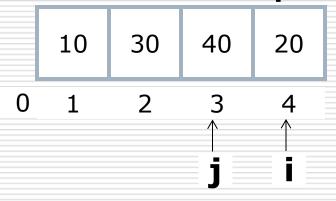
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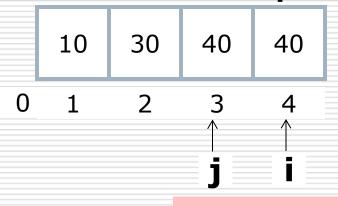
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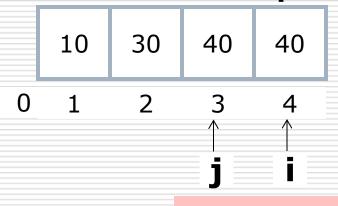
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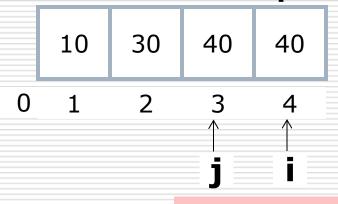


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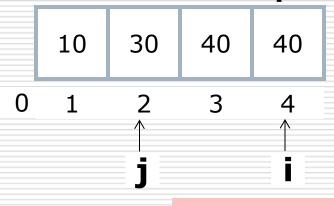


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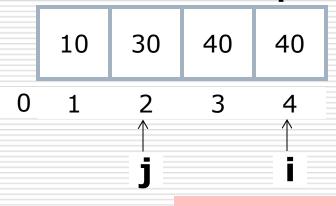




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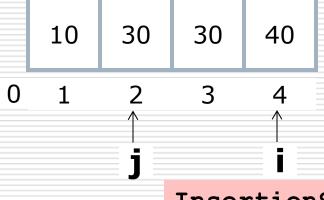
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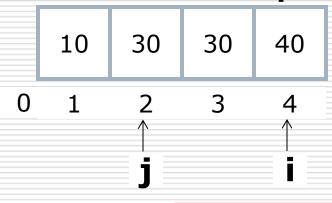
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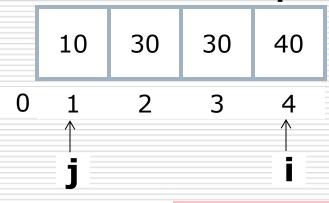


David Luebke

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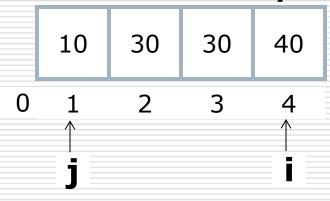
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```



```
i = 4 j = 1 key = 20

A[j] = 10 A[j+1] = 30
```

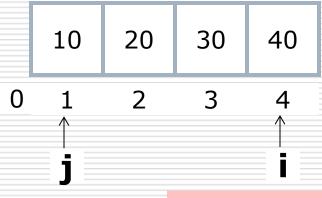
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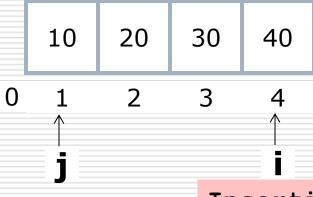
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i = 4 j = 1 key = 20

A[j] = 10 A[j+1] = 20
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```
Done!
```

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            j = j - 1
        }
        A[j+1] = key
    }
}
```

```
Statement
```

Cost Times

```
InsertionSort(A, n) {
   for i = 2 to n \{
                                               C₁ n
                                               c_2 (n-1)
      key = A[i]
                                               c_3 (n-1)
      i = i - 1;
      while (j > 0) and (A[j] > key) { C_4 Tx
                                               C_5 Ty
             A[j+1] = A[j]
                                               c<sub>6</sub> Ty
             j = j - 1
                                               c_7 (n-1)
      A[j+1] = key
```

```
i = 2, 3, 4,....., n

Tx = 1+1, 2+1, 3+1,....(n-1)+1 = 2+3+4+...+n = [n(n+1)/2] - 1

Ty = 1, 2, 3,.... (n-1) = 1+2+3+...+(n-1)= n(n-1)/2

Total Time = c_1*n+(c_2+c_3)*(n-1)+c_4*Tx+(c_5+c_6)*Ty+c_7*(n-1)
```

Best-case Analysis

Total Time=
$$c_1*n+(c_2+c_3)*(n-1)+c_4*Tx+(c_5+c_6)*Ty+c_7*(n-1)$$

$$T(n) = c_1 n + c_2 (n - 1) + c_3 (n - 1) + c_4 (n - 1) + c_5 * 0 + c_6 * 0 + c_7 (n - 1)$$

$$= (c_1 + c_2 + c_3 + c_4 + c_7) n + (c_1 + c_2 + c_3 + c_4 + c_7)$$

$$= an + b$$

We can express this running time as an + b for *constants* a and b that depend on the statement costs c_i ; it is thus a *linear function* of n.

- Least amount of (time) resource ever needed by algorithm
- Achieved when incoming list is already sorted in increasing order
- Inner loop is never iterated

Worse-case Analysis

i = 2, 3, 4,....., n
Tx = 1+1, 2+1, 3+1,....(n-1)+1 = 2+3+4+..+n =
$$[n(n+1)/2] - 1$$

Ty = 1, 2, 3,.... (n-1) = 1+2+3+..+(n-1)= $n(n-1)/2$
Total Time= $c_1*n+(c_2+c_3)*(n-1)+c_4*Tx+(c_5+c_6)*Ty+c_7*(n-1)$

$$T(n) = c_1 n + c_2 (n - 1) + c_3 (n - 1) + c_4 \left[\frac{n(n + 1)}{2} - 1 \right] + c_5 \left[\frac{n(n - 1)}{2} \right]$$

$$+ c_6 \left[\frac{n(n - 1)}{2} \right] + c_7 (n - 1)$$

$$= \left(\frac{c_4}{2} + \frac{c_5}{2} + \frac{c_6}{2} \right) n^2 + (c_1 + c_2 + c_3 + \frac{c_4}{2} - \frac{c_5}{2} - \frac{c_6}{2} + c_7) n - (c_2 + c_3 + c_4 + c_7)$$

$$= an^2 + bn + c$$

We can express this worst-case running time as $an^2 + bn + c$ for constants a, b, and c that again depend on the statement costs c_i ; it is thus a *quadratic function* of n.

- Greatest amount of (time) resource ever needed by algorithm
- Achieved when incoming list is in reverse order
- Inner loop is iterated the maximum number of times