tast (Semple Efficient) RL New setting : acting online every step get to make a decision (take an action) and observe a next state & reward Often important to learn "quickly" multiple definitions of interest quickly can mean minimizing # of samples (actions) to learn an optimal policy quickly may mean that empirical performance (reward obtained) goes up rapidly as a function of the # of data points we will consider/discuss several objectives but key idea is that care about Impact of amount of data on performance sample efficiency not computational efficiency A lot of work or computationally efficient algorithms When is fast learning / sample efficiency important & when is computational efficiency bay? (of course, nice to have both) - computational efficiency robotics, consider self-driving vehicle driving at 60 mph. by the time I second has passed already moved over 26 meters! quickly computing next action meded (or, operating hierarchically, e.g. may have meta-actions or mini-policies and Icarn & plan over these) simulated domains : video games - sample efficiency when obtaining data is hard / costly, avoiding or limiting bad decisions is important most applications that involve RL agents interacting w/people education; what problem to give a student part healthcare; what treatment to give a petient consumer marketing: what ad (or set of eds) robotics in extreme/remote environments eig. Mars scientific exploration by robots

ection sprague

Types of algorithms

value function based

Evaluation criteria
how do we know if an algorithm ienables fast learning?
Many possible objectives (& these can yield different
algorithms)

1) Empirical performence

- completive reward after fixed # of time steps (decisions)
- how quickly (# time steps) can learn optimal policy (note: may or may not be executing that policy even if tearned it, consider E-greedy exploration approaches)
- slope of increasing reward: rapidly increase average reward obtaining as a function of the # of timesteps (decisions)
- 2) PAC: on all but N steps/actions, take an action whose state-action value is near optimal N is a polynomial function of the MDP parameters

Knows What It Knows (KWIK) is a related criteria

3) Regret. Compare total accomulated reward to expected reward from making optimal decisions. Can compare in terms of reward could have obtained if made optimal decisions from the start vs decisions made or for each state reached making actual decisions made, what's the difference in reward from the action taken vs the action on should have taken.

-> Why are these different? What is PAC guaranteeing?

regret bounds tend to be in terms of rates
e.g. is regret growing limarly, sublimarly as a func of

U T

15 Kps

4) Bayesian optimality given mittal uncurtainty over MDP parameters maximize expected sum of comulative rewards If acting forever almost definitely will still went to learn optimal policy for states encounter for finite horizon It where His short-ish, less clear valve of information what is the value of information if only making I more decision?

Today: Probably Approximately Correct (PAC) RL 'Early ideas pronund by Keans & Singh (1998, 2002). Brafman & Tennenholtz (2002), Sham Kakade (2003) A lot of later work, esp by Michael Littman & his group (Alux Strehl, Lihong Li, Tom Welch, ...) my group elso active in this space defria RL algorithm A is PAC-MDP if on all but N steps, the algorithm's non-stationary policy at time to At (note: this policy is computely defined by the algorithm Adhistory up to t) is at wast 6-optimal from the current state VA+ (s+) ≥ V*(s+)-E with probability at hast 1-8, where

N= polynomial func of (151,1A1,1/6,1/8,1/1-7).

Why do we consider only a high probability bound? quen finite data l'experience in a stochastic environment, could always yield experience that makes it impossible to identify the optimal TI W/perfect confidence

number of mistakes (times take potentially poor decisions) is bounded

* note; doesn't say when (on what time step) this mistakes will occur

Why might not all mistakes occur at the start?

stert so 1= p=10-10 7 53

can bern to act near optimally here, likely long before visit

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example PAC algorithm: 12-max
      mentain set of "known" state-action pairs
"unknown" " "
      initially all soa pairs are unknown
      intuitively siz pair becomes known when have a good estimate
         of the model parameters for that so pairs
          (can also formulate related intuition for modul free algorithms)
       define a MDP M w/some state & action spaces of 7
         for all s, a pairs in K+ (known set)
             set R(sia) = S; r; (sia) /# times sur sia (eig. emplicial avq
                  T(1/s,a) = empirical estim of trans model
         for all (s,a) & K+ (unknown)
                 R(s,a) = Rmax
                  Î(sls, 2) = 1 (self loop)
      compute fix for A
      act using A*
      Kup counts of # visits to each (sia) pair
      if counts(s,a) ≥ m, add (s,a) to known set K+
intuitively what does this algorithm do?
optimism under uncurtainty
with correct setting of m, R-max is a PAC elgorithm
proof approach (Strehl et al UAI 2006)
    generic proof that can be used for related algorithms
     assume RL alg Amaintains Q(sia)
     let Q+(s, a) be is estimate immediately before taking the action
     A is a "grady" alg if at = aigmaxa Qt (stia)
    - define known s-a MDP MK = < SU 8505, A, TK, RK, T>
      related to a MDP M = < S, A, T, R, p > Let K be the set of
      known so pairs and Q be a set of state-action values
      (not necessarily related to valves of Mor Mx).
       VSIZEK, RK(SIZ)=R(SIZ) TK(1/SIZ)=T(1/SIZ) } true model
       R(so,)=0 T(solso,)=1 (sink state)
       Ysia &K RK(sia) = Q(sia) TK(solsia)=1
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Proposition 1. Let A(E, 8) be any grady learning alg sit. for every time step t there exist a set of Kt state-action pairs. Kt=Ktri unless during timestapt an unknown (sia) pair is visited or a Q valve is updated. MK+ is the known s-a MDP and TI+(s)= argmaxa Q+(s,a). Suppose that for any inputs e & 8, w/prob = 1-8 the following holds # sia, t 1) Optimism: Q+(s,a) = Q*(s,a)- E 2) Accuracy: V+(s)-VMK+(s) = E 3) Learning complexity bounded: total # of updates of Q estimates and the # of times visit an unknown (sia) pair is bounded by \$(E,8). Then when A(E, 8) is executed on any MOP M, it will follow a 46-optimal T from its current state on all but $O\left(\frac{f(\epsilon, \delta)}{\epsilon(1-\gamma)^2} \ln \frac{1}{\delta} \ln \frac{1}{\epsilon(1-\gamma)}\right) + \lim_{\delta \to \infty} \frac{1}{\delta} \ln \frac$ Proof sketch, Assume Recoil Define D= Top In Elley) VT (s,D)-VT (s) = E finite housen value e-close to on housen value if housen long enough (keeins & Singh 2002) let At be the current Tof the agent (nonstationary) let whee the event that by executing At starting in state of the 2 occurs 1) the algorithm successfully updates Q(sia) for some (sia) 2) visit an unknown (sia) pair Lover D skps can bound value of At in true MDP by prob of W VM (St, D) > VMK+ (St, D) (1- Pr(W)) + Pr(W). Value when Working note At= TT+ unless update Q+(s1=) M and Mkt are identical on all (sia) & Kt so get exact same behavior & rewards unless event wo happens (lest kin pos, V= 17) > VMKI (St.D) - Pr (W) 1/1-7 ≥ VMK+ (S+)- E - Pr(W) 1/1-7 defr of D = V(st)-2E - Pr(w)/1-p by accuracy condition = 1+(s+)-3E - Pr(w)11-7 by optimism 8-1 / Lold w/prob >1-8

IF Pr(W) < E(1-7) VAT (ST) > VAT (ST, D) = V+(ST)-4E

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proof of Prop. I continued
   if Pr(W) > E(1-7)
        want to know # of D-intervals until # of w events
           = { (e.s) (since know from cond 3 this bounds
            the # of Wevents)
        to get upper bound, treat each D-interval as iid
            opportunity for w to occur w/prob = E(1-7)
 lumme 56 (Lihong Li, PhD thusis, 2009) ilet X1... Xm be &
    seg of m indep Bernoulli triels, each w/a success
     prob of at least M: E[Xi] = M > 0, Then for any
     K & IN and SE (0,1), with prob 21-8.
            X, # X2 + 1. Xm ≥ K
           m = = (k+ In 1/8).
 applying lemma 56 we get that if the # of intervals is
         = = ( } ( } ( ( ( ( ( ) + | n | 1 ) )
 then whigh prob all & (E, E) W events will have occured.
  since each interval is O-steps this wields
         = 20 3(E, 8) In (1/6) = since In (1/6) >1
 so # time steps on which Pr(w) > E(1-7) is upper bounded by
        20 5 (E.8) In (1/8) In (E/1-71)
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how do we achieve the required preconditions?

simulation lemma & bounds on how far empirical model peram

estimates can be from true parameters (see earlier lecture)

yield accuracy requirement

yield accuracy requirement

also yield criteria to make (sia) pair known

pigeonhole principle bounds # times can visit unknown (sia) pair optimism comes from simul lemma plus using optimistic optimism comes from simul lemma plus using optimistic estimate for all unknown (sia) pairs

what is final sample complexity (# of steps not necessarily near optimal?) for R-mex*

Are these practical? No Can we do better? Yes.

- Lettimore & Hutter (ALT 2012), discounted MDP

$$\begin{array}{ll}
& \text{NSSQ} \\
& \text{NSSQ} \\
& \text{VSSQ}
\end{array}$$

$$\begin{array}{ll}
& \text{NSSQ} = \# \text{Non-zero Significans} \\
& \text{VSSQ} = \# \text{Non-zero Significans} \\
& \text{VSSQ} = \# \text{Non-zero Significans}
\end{array}$$

much better (by = (1-x)s) than prior bounds in terms of ety

can be worse in state dup

- Denn & Brunskill (NIPS 2015). Finite horizon MDP episodic

lower bound

1 of key insights of the above approaches; not all (sia) pairs
are equally likely to be visited nor influence value force egrally

Bounds still impractical

optimism under uncertainty (goes back, sans theory, at but ideas are useful least to early 1990s Kaelbling)
quantifying "enough" info

many extensions, altrinatives

don't have to be binary in "knowness" use info about state-action experience but still build in rep of uncertainty

erg. confidina bounds