# CSE 150 Foundations of Computer Science: Honors, Fall 2016

Assignment #3
Part A: Oct. 24, 2016
Part B: Oct. 31, 2016

Please format your homework in LATEX. That includes any tables or equations.

The goal of this homework is to improve your formal proof techniques and give you more insights on the issues of countability. The first exercise is solved for you as an example. Please be sure to read, understand and use the definitions below; you need to be precise when you prove any statement, and the only way to be precise is to use the definitions to show your results.

**Note:** Part B is not due until after the exam. However we strongly encourage you to do and understand the problems on induction ahead of time.

### Definition Recap:

Function  $f: X \to Y$  is said to be **one-to-one**, or *injective*, if for all a, b in X: f(a) = f(b) if and only if a = b.

Function  $f: X \to Y$  is said to be **onto**, or *surjective*, if for all b in Y there exists an a in X such that f(a) = b.

Function F is a **bijection** (or is *bijective*) if it is both one-to-one and onto (injective and surjective).

#### Example 1

Show that  $\mathbb{N} - \{1\}$  is countable.

#### Proof:

Let  $X = \mathbb{N} - \{1\}$ . We show that X is countable by showing a bijection between X and N.

Let  $f: X \to \mathbb{N}$ . Set f(n) = n - 1. We show that f is a bijection.

First observe that f is one-to-one. This is true because if f(a) = f(b) for some  $a, b \in X$ , then a - 1 = b - 1, implying that a = b.

Now we show that f is onto. We need to show that for all  $a \in \mathbb{N}$  there exists b such that f(b) = a and b is in X. Let b = a + 1. Then f(b) = f(a + 1) = a + 1 - 1 = a. Since  $a \in \mathbb{N}$ ,  $a \ge 1$  and  $a + 1 \ge 2$ ; therefore, b = a + 1 =is in X.

Thus, f is a bijection.  $\square$ 

### Part A

### Problem 2

Show that every subset of  $\mathbb{N}$  is countable.

*Hint:* Use the fact that any nonempty set of positive integers has a least element. If X is a subset of N, show that  $f(n) = \min\{X - \bigcup_{i=1}^{n-1} \{f(i)\}\}$  gives the desired bijection for infinite X.

#### Problem 3

Show that the set of all integers  $(\mathbb{Z})$  is countable.

# Problem 4

Show that the set of all *finite* subsets of a countable set is countable.

## Problem 5

Show that the following are statements are equivalent and true (Statements P and Q are equivalent if  $P \Leftrightarrow Q$ ):

- $\mathbb{N} \times \mathbb{N}$  is countable
- Union of countably many countable sets is countable.
- $\mathbb{Q}$  is countable.

Hint: look at  $\mathbb{Z}^2 = \mathbb{Z} \times \mathbb{Z}$  (also known as the integer lattice). Start at the origin and try walking the grid in such a way that any point in it will be reached after a finite number of steps. One example of such walk would walking  $\mathbb{Z}^2$  in a spiral starting from the origin (see Figure 1); try to come up with other kinds of walks on  $\mathbb{Z}^2$  or  $\mathbb{N}^2$ . Then the walk will be an enumeration of the points in  $\mathbb{Z} \times \mathbb{Z}$  or  $\mathbb{N} \times \mathbb{N}$ , since it yields a bijection  $\phi(n) = (\text{the point you reach after } n \text{ steps})$ 

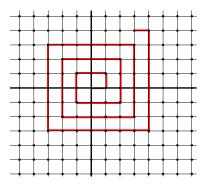


Figure 1: Lattice Walk on  $\mathbb{Z} \times \mathbb{Z}$ 

### Problem 6

A submarine is moving along the integer number line at a constant speed s so that at each hour it is on an integer number. It started moving at time 0 at some position s. If t is the (whole) number of hours elapsed since the submarine started moving, then its position is given by the equation s0 are integers.

You are working at Rocket Pizza delivery and you are to deliver pizza to the submarine. At each hour you can drop pizza on any number on the integer line. If the submarine is there at that time, then you have delivered the pizza and your job is done (you will be notified as soon as it happens).

The problem is that you don't know where the submarine is, you cannot see it, you don't know where it started and how fast it is moving (i.e., you don't know values of s and b - classified top secret data). The upside is that you have infinite number of pizzas.

Show that you can deliver pizza in a finite amount of time.

# Problem 7

The goal of this quiz/question is to make sure you will have no problems with countability in the future.

In the following table, F stands for Finite, I for Infinitely Countable, C for Countable, U for uncountable. Put a checkmark next to the strongest statement you can make about the resulting set (i.e. correct answer for  $F \cup F$  is finite, even though it is countable as well). Check ? whenever there's not enough information to decide, i.e. there can be different cases with no 'strongest' answer.

Set	F	Ι	С	U	?	*	Set	F	Ι	С	U	?
$F \cup F$						*	$C \cup F$					
$F \cup I$						*	$C \cup I$					
$F \cup C$						*	$C \cup C$					
$F \cup U$						*	$C \cup U$					
$F \cap F$						*	$C \cap F$					
$F \cap I$						*	$C \cap I$					
$F \cap C$						*	$C \cap C$					
$F \cap U$						*	$C \cap U$					
F - F						*	C-F					
F-I						*	C-I					
F-C						*	C-C					
F-U						*	C-U					
$F \times F$						*	$C \times F$					
$\overline{F \times I}$						*	$C \times I$					
$F \times C$						*	$C \times C$					
$F \times U$						*	$C \times U$					
$I \cup F$						*	$U \cup F$					
$I \cup I$						*	$U \cup I$					
$I \cup C$						*	$U \cup C$					
$I \cup U$						*	$U \cup U$					
$I \cap F$						*	$U \cap F$					
$I \cap I$						*	$U \cap I$					
$I \cap C$						*	$U \cap C$					
$I \cap U$						*	$U \cap U$					
I-F						*	U-F					
I-I						*	U-I					
I-C						*	U-C					
I-U						*	U - U					
$I \times F$						*	$U \times F$					
$I \times I$						*	$U \times I$					
$I \times C$						*	$U \times C$					
$I \times U$						*	$U \times U$					

#### Part B

#### Problem 8

You have 10000 kilograms of pickles. Pickles are 99 percent water by volume. Water comprises 100 percent of the mass of the pickle. Time goes by, and you observe that some water has evaporated. Now water comprises only 98 percent of the volume. What is the weight of the pickles now?

Some of the following problems require the use of mathematical induction. The following problem is solved for you as an example:

# Example 9

Show that 
$$1 + 2 + 3 + ... + n = \frac{n(n+1)}{2}$$

**Proof:** proof by induction.

Let P(n) be the predicate that  $1+2+3+\ldots+n=\frac{n(n+1)}{2}$ 

Base case:  $1 = \frac{1 \cdot 2}{2}$ ; hence P(0) holds.

Inductive assumption: Let P(n) hold for n = k, i.e.  $1 + 2 + 3 + ... + k = \frac{k(k+1)}{2}$  for some number k.

Inductive step: We show that given the inductive assumption, P(k+1) holds. Observe that

$$1 + 2 + 3 + \dots + k + (k+1) = \frac{k(k+1)}{2} + (k+1) = \frac{k(k+1) + 2(k+1)}{2} = \frac{(k+1)(k+2)}{2}$$

(where the second step comes from inductive assumption, and the rest follows by simplification).

Therefore, P(n) holds for all n by the principle of mathematical induction, thus proving the theorem.  $\square$ 

# Problem 10

Prove the following using mathematical induction:

- 1.  $2n \le 2^n$
- 2.  $1+3+5+...+(2n-1)=n^2$
- 3.  $1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{(n)(n+1)(2n+1)}{6}$

### Problem 11

You have an  $n \times m$  bar of chocolate. Your goal is to separate all of the squares of chocolate. The way that you can break the chocolate is to take a single piece of chocolate (connected component of squares) and break it along one horizontal or vertical line. What is the minimum number of breaks necessary? Please prove your answer.

# Problem 12

You have an  $n \times n$  checkerboard with an initial set of checkers placed on it. You are allowed to add additional checkers under the following conditions: You can place a checker on a square if two or more neighboring squares also have checkers on them. Neighboring cells are those above, below, to the left and to the right, as shown in Figure 2. As we showed in class, there are initial configurations of n checkers that enable the entire board to be covered. Prove that no configuration of n-1 checkers can let you cover the board.

Hint: this is not induction on the number of pieces.

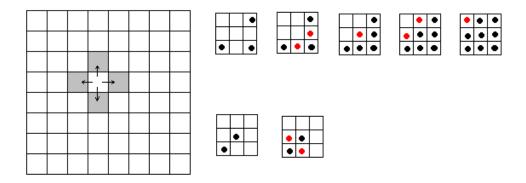


Figure 2: Neighboring cells and checkerboard evolution.