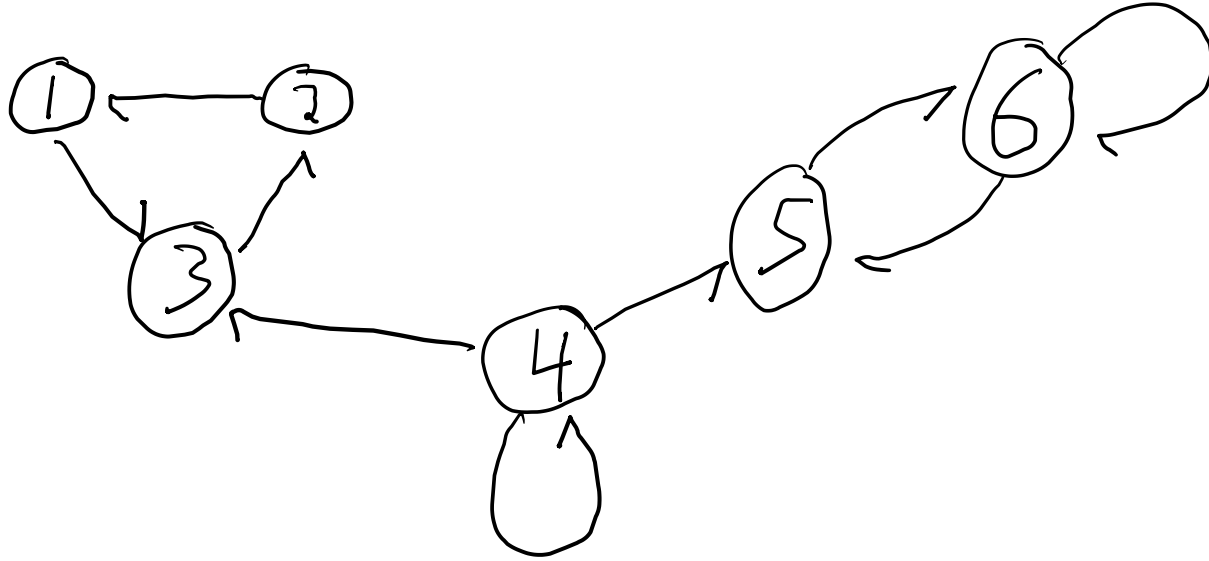


# Lecture 5

## Graph properties of MC

- Communicating classes (and accessibility)
- Periodicity
- Transience/recurrence

Ex)



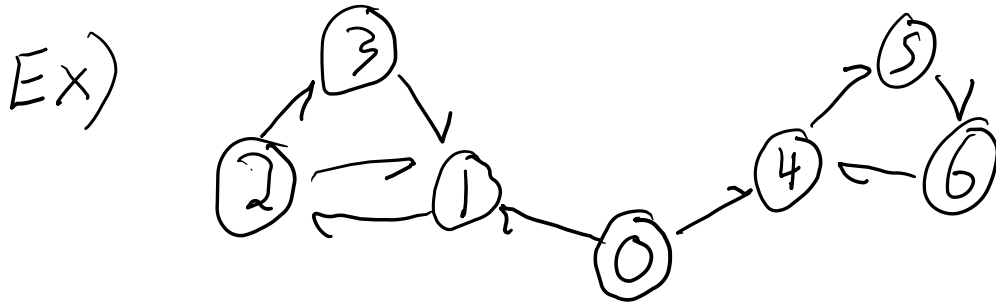
We will learn to make the following kind of classifications:

Communicating classes:

- $\{1, 2, 3\} \leftarrow$  recurrent, period 3
- $\{4\} \leftarrow$  transient, aperiodic
- $\{5, 6\} \leftarrow$  recurrent, aperiodic

## Defs:

- State  $s_i$  is accessible from  $s_j$  if  $p_{i,j}^n > 0$  for some  $n \geq 0$ .
- States  $s_i$  and  $s_j$  communicate if  $s_i$  is accessible from  $s_j$  and  $s_j$  is accessible from  $s_i$ . This is denoted by  $i \leftrightarrow j$ .



True or False:

- All states are accessible from 0 (T or F)
- All states communicate with 0 (T or F)
- $1 \leftrightarrow 2$  (T or F)
- $1 \leftrightarrow 4$  (T or F)

**Proposition:** Communication is an equivalence relation, i.e., it is

- (i) • **Reflexive:**  $i \leftrightarrow i$  for all states  $i$
- (ii) • **Symmetric:**  $i \leftrightarrow j \Rightarrow j \leftrightarrow i$
- (iii) • **Transitive:**  $i \leftrightarrow j$  and  $j \leftrightarrow k \Rightarrow i \leftrightarrow k$

Proof: (i)  $P_{ii}^0 = 1$  so  $i \leftrightarrow i$  ( $n=0$ )

(ii)  $i \leftrightarrow j$  then  $\exists m, n$   $P_{ij}^n > 0$  and  $P_{ji}^m > 0 \Rightarrow j \leftrightarrow i$

(iii) Suppose  $\exists m \geq 0$   $P_{ij}^m > 0$  and  $\exists n \geq 0$   $P_{jk}^n > 0$

and let's show  $P_{ik}^{m+n} > 0$  (so  $i \rightarrow k$ )

Using the Chapman Kolmogorov equation:

$$P_{ik}^{m+n} = \sum_{l \in S} \underbrace{P_{il}^m}_{\geq 0} \cdot \underbrace{P_{lk}^n}_{\geq 0} \geq \underbrace{P_{ij}^m}_{\geq 0} \cdot \underbrace{P_{jk}^n}_{\geq 0} > 0$$

$$\Rightarrow i \rightarrow k$$

and similarly one can show  $k \rightarrow i \Rightarrow \boxed{i \leftrightarrow k}$   ~~$\square$~~

Consequence (from equivalence relation in general)

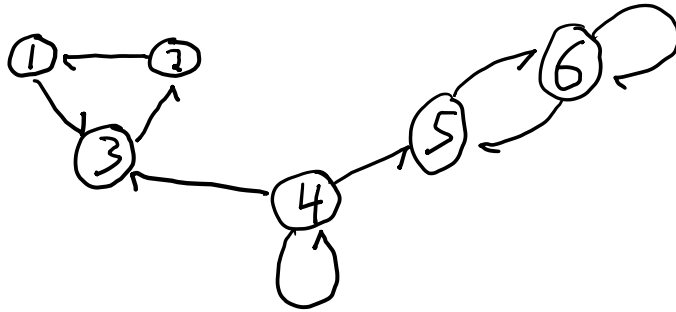
We can partition (split) the state space into communicating classes

# Partitioning

Any equivalence relation on a set partitions the set into *equivalence classes*.

Thus, the state space always partitions into communicating classes.

Ex:



$$S = \{ \underbrace{1, 2, 3}_{\text{red}}, \underbrace{4}_{\text{green}}, \underbrace{5, 6}_{\text{blue}} \}$$

$$\underline{C_1} = \{1, 2, 3\} \quad \underline{C_2} = \{4\}$$

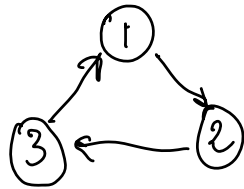
$$\underline{C_3} = \{5, 6\}$$

**Def:** A MC is irreducible if it has only one communicating class.

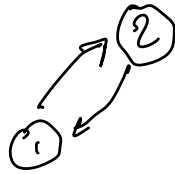
**Def:** The period of state  $s_i$ , denoted  $d(i)$ , is the greatest common divisor of  $\{n \in \{1, 2, \dots\} : p_{i,i}^n > 0\}$ .

I.e., starting in  $s_i$  the MC can only return to  $s_i$  at numbers of steps that are multiples of  $d$ , and  $d$  is the largest such integer.

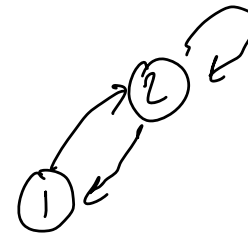
Examples: What is  $d(1)$ ?



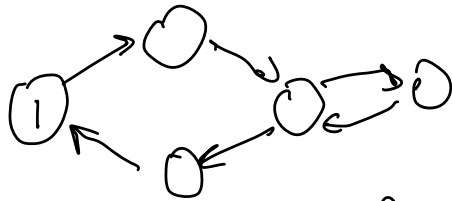
$$d(1) = 3 (=d(2)=d(3))$$



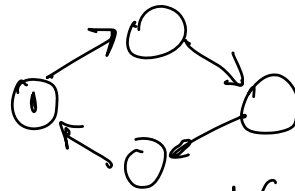
$$d(1) = 2$$



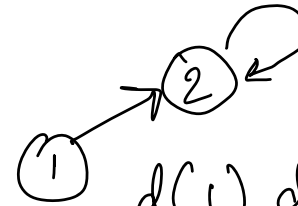
$$\begin{aligned} 121 \quad n=2 \\ 1221 \quad n=3 \\ \Rightarrow \begin{cases} d(1) | 2 \\ d(1) | 3 \end{cases} \Rightarrow d(1)=1 \end{aligned}$$



$$d(1) = 2 \quad \{n \mid p_{1,1}^n > 0\} = \{4, 6, 8, \dots\}$$



$$d(1) = 4$$

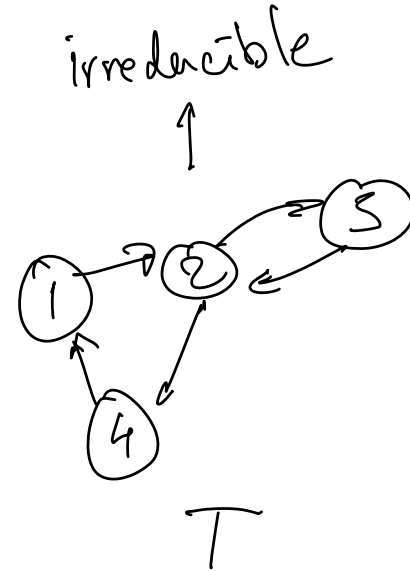
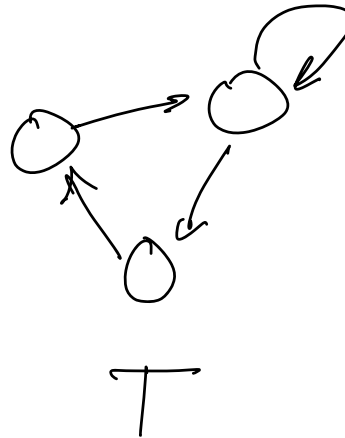
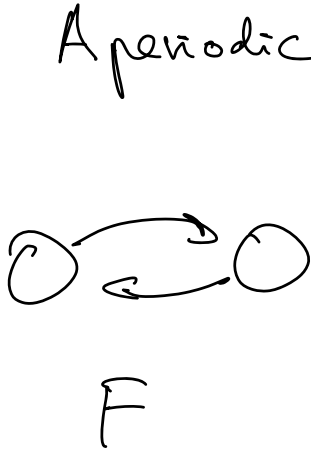


$d(1)$  does not exist

Remark: A state can have no period

**Def:** We say that a state is aperiodic if it has period 1, and we say that an MC is aperiodic if all states have period 1.

Examples:



$$\left( \begin{array}{l} d(2) | 2 \\ d(2) | 3 \end{array} \right) \Rightarrow d(2) = 1$$



**Proposition:** Period is a class property, i.e., all states in the same communicating class have the same period.

Proof:

↳ Remark: In practice, to study the periodicity of a M-C, we just need to find the periodicity associated with 1 state in each communicating class.

↳ Suppose  $i \leftrightarrow j$  ( $\exists m \mid P_{ij}^m > 0$   
 $\exists n \mid P_{ji}^n > 0$ ) and let's show that  $d(i) = d(j)$

• First, we show  $d(i) \leq d(j)$

$$P_{ii}^{m+n} = \sum_k P_{ik}^m P_{ki}^n \geq \underbrace{P_{ij}^m}_{>0} \underbrace{P_{ji}^n}_{>0} > 0 \Rightarrow d(i) \mid m+n$$

$C-K$

Now consider  $l \mid p_{ij}^l > 0$  (we know that such  $l$  exists since  $d(i)$  exists)

$$p_{ii}^{m+n+l} \geq \underbrace{p_{ij}^m}_{>0} \underbrace{p_{\bar{i}\bar{j}}^l}_{>0} \underbrace{p_{j\bar{i}}^n}_{>0} > 0 \Rightarrow d(i) \mid m+n+l$$

$$\text{So } d(i) \mid m+n+l - (m+n) \Rightarrow d(i) \mid l \Rightarrow \boxed{d(i) \leq d(j)}$$

↓  
 $d(i)$  is a common divisor

~~$X$~~  =  $\{n \mid p_{\bar{i}\bar{j}}^n > 0\}$        $d(j)$  is the gcd of  $X$

By symmetry we also have  $d(j) \geq d(i)$

$$\text{So } \boxed{d(i) = d(j)}$$

~~$X$~~

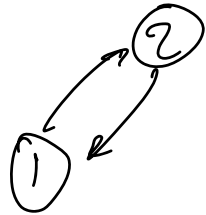
**Def:** Let  $f_i := P(X_n = i \text{ for some } n \geq 1 | X_0 = i)$ . Then

- if  $f_i = 1$ , State  $s_i$  is recurrent,
- if  $f_i < 1$ , State  $s_i$  is transient.

In words:  $f_i = 1$  if you are always sure to visit  $i$  again

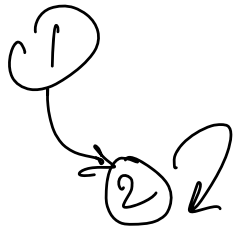
**Proposition:** Recurrence and transience are class properties.

Examples:



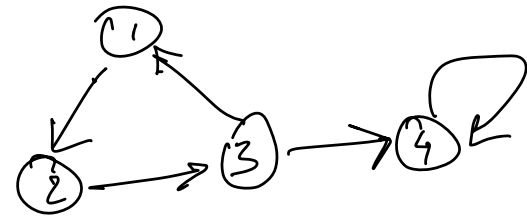
① recurrent  
or

$\{1, 2\}$  transient  
( $f_1 \neq P_{11}^2 = 1$ )



② recurrent

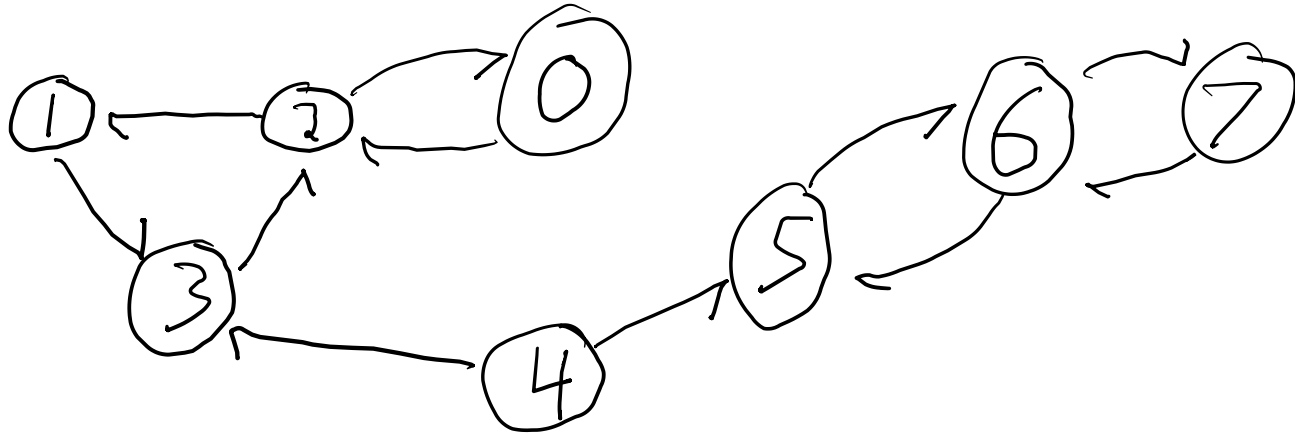
① transient ( $f_1 = 0$ )



$\{1, 2, 3\}$  transient

④ recurrent

Ex)



Communication classes :  $\{0, 1, 2, 3\}$ ,  $\{4\}$ ,  $\{5, 6, 7\}$

$\downarrow$                        $\downarrow$                        $\downarrow$

recurrent                      transient                      recurrent