Concurrent Programming

Exercise Booklet 9: Promela and Spin¹

Solutions to selected exercises (\Diamond) are provided at the end of this document. Important: You should first try solving them before looking at the solutions. You will otherwise learn **nothing**.

1 Basic Promela

Exercise 1. (\Diamond) Implement the following entry/exit protocol (Attempt I) seen in class, in Promela. Use the Promela code in the slides as an aid in understanding Promela syntax.

```
global int turn = 1;

thread P: {
    while (true) {
        await (turn==1);
        turn = 2;
        turn = 1;
    }
}
```

Exercise 2.

Draw the transition system of the following two programs separately and then compare them (note: you can use SpinSpider as a guide):

```
% Program 1
                                % Program 2
     byte state = 1;
                                  byte state = 1;
     active proctype A(){
                                  active proctype A(){
                                       (state == 1) ->
       atomic {
          (state == 1) ->
                                         state = state+1
            state = state+1 6
                                  }
6
                                  active proctype B() {
     }
                                      (state == 1) ->
     active proctype B() {
                                         state = state-1
       atomic {
         (state == 1) ->
            state = state-1
12
```

Exercise 3. Check whether the following algorithm guarantees mutual exclusion by adding an auxiliary variable critical and appropriate statements. Do you recognize this algorithm?

```
bool flag[2]
bool turn

4 active [2] proctype user()
{
    flag[_pid] = true
    turn = _pid
```

¹Sources include: http://www.cs.toronto.edu/~chechik/courses01/csc2108/lectures/spin.2up.pdf

```
s (flag[1-_pid] == false || turn == 1-_pid)

10 crit: skip // critical section

12 flag[_pid] = false
}
```

Exercise 4. What happens in the previous algorithm if you exchange the line flag[_pid] = true with the line turn = _pid?

Exercise 5. Consider the following simplified presentation of the Bakery Algorithm for two processes:

```
global int np,nq =0;
thread P: {
                              2 thread Q: {
  while (true) {
                                   while (true) {
   // non-critical section
                                    // non-critical section
   np = nq + 1;
                                   nq = np + 1;
   await nq==0 or np \le nq;
                                    await np == 0 or nq < np;
   // CRITICAL SECTION
                                    // CRITICAL SECTION
   np = 0;
                                   nq = 0;
   // non-critical section
                                   // non-critical section
                                  }
}
```

- 1. Encode it in Promela.
- 2. Show that it does not guarantee mutual exclusion.
- 3. Add the command np = 1 in thread P just before line 4 and add nq=1 in thread Q just before line 4. Show, using Spin, that the resulting program does enjoy mutual exclusion.

Exercise 6. The following extension to 3 processes of Dekker's algorithm is known to deadlock.

```
global int turn=0;
  global bool[] flags = {false, false, false};
  thread (id): {
     int left = (id+2)\%3;
     int right = (id+1)%3;
6
     flags[id] = true;
     while (flags[left] || flags[right])
       if (turn == left) {
         flags[id] = false;
10
         await (turn==id);
         flags[id] = true;
12
       }
     }
14
     // Critical Section
     turn = right;
     flags[id]=false;
18
```

Encode it in Promela and verify where it deadlocks (there are multiple traces that lead to deadlock, exhibiting one suffices). Here is some code that you can start from.

```
byte turn;
2 bool flags[3];
```

```
proctype P() {
   byte myId = _pid-1;
   /* complete here */
}

init {
   turn=0;
   byte i;
   for(i:0..2) { flags[i] = false; }

atomic {
    for (i:0..2) { run P(); }
}
}
```

Exercise 7. (\lozenge)

- 1. Define a general semaphore with operations acquire(sem) and release(sem), where sem is a numeric shared global variable. Since there are no functions in Promela, you can declare macros using inline. Hint: use atomic and blocking expressions.
- 2. Check that your solution is correct by providing an example of its use.

Exercise 8. Implement the Bar exercise from eb5, in Promela, using the semaphore encodings from Exercise. 7. Use Spin to show that the required invariant is upheld, i.e. that only one man goes in for every two women, by inserting an appropriate assertion. In your example use 20 women and 20 men threads.

2 Channels

Exercise 9. (\lozenge) What does the following program print?

```
proctype A(chan q1) {
     chan q2;
     q1?q2;
     q2!123
   proctype B(chan qforb) {
     int x;
     qforb?x;
     printf("x=%d\n",x)
10
   init {
     chan qname = [1] of { chan };
12
     chan qforb = [1] of { int };
     run A(qname);
     run B(qforb);
     qname!qforb
16
   }
```

Exercise 10. (\Diamond) Implement a binary semaphore in Promela using message passing and synchronous communication. You should have two proctypes, semaphore and user. Here is the code for the user:

```
#define acquire 0
#define release 1
   chan sema = [0] of { bit }; /* synchronous channel */
   proctype semaphore() {
     /* complete this */
6
   proctype user() {
     dο
     :: sema?acquire;
10
        /* crit. sect */
        sema!release;
        /* non-crit. sect. */
12
     o d
14 }
   init {
    run semaphore();
    run user();
    run user();
   }
```

Exercise 11. Implement the turnstile example from class using channels. Here is some code to get you started:

```
mtype { bump , read };
   chan counter = [0] of { mtype, chan }
  chan reply[2] = [0] of { byte }
  proctype Counter() {
      /* complete */
  proctype Turnstile() {
     /* complete */
11
  init {
     byte x;
     atomic {
15
      run Counter();
      run Turnstile();
17
       run Turnstile();
19
     _nr_pr==2;
     counter!read(reply[0]);
21
     reply[0]?x;
     printf("%d\n",x) /* should print 20 */
23
   }
```

3 Solutions to Selected Exercises

Answer to exercise 1

```
byte turn=1;

active proctype P() {
   do
     :: turn==1;
```

```
printf("P went in \n");
    turn = 2

od

active proctype Q() {

do
    :: turn==2;
    printf("Q went in \n");
    turn = 1
    od
}
```

Answer to exercise 7

```
inline acquire(sem) {
   atomic {
    sem > 0;
    sem --
}
}

inline release(sem) {
   sem++
}
```

Answer to exercise 9

- Init sends qforb to A on qname
- A sends 123 to qforb
- B receives 123 on qforb and prints it

Answer to exercise 10

```
#define acquire 0
   #define release 1
   chan sema = [0] of { bit };
   proctype semaphore() {
     byte count = 1;
5
     :: (count == 1) \rightarrow sema!acquire; count = 0
     :: (count == 0) -> sema?release; count = 1
     o d
   }
11
   proctype user() {
13
     :: sema?acquire;
        /* crit. sect */
15
       sema!release;
        /* non-crit. sect. */
17
     o d
19
   init {
    run semaphore();
21
     run user();
     run user();
```