

# Data Flow Criteria

# Data Flow Testing -- Additional Criteria

Data flow testing

```
graph TD; A[Data flow testing] --> B[Static data flow]; A --> C[Dynamic data]; B --> D[Involves actual program execution (similar to CFG) but paths are identified based on data flow testing criteria.]; C --> E["• Identify potential defects (anomalies)  
• Analyze source code  
• Do not execute code (dead code)"];
```

Static data flow

Involves actual program execution (similar to CFG) but paths are identified based on **data flow testing criteria**.

Dynamic data

- Identify potential defects (**anomalies**)
- Analyze source code
- Do not execute code (dead code)

# Data Flow Anomaly

There are different possible two letter combinations for 'd', 'k' and 'u'. Some are bugs, some are suspicious and some are okay

- **dd**: Probably harmless, but suspicious. Why define the object twice without an intervening usage?
- **dk**: Probably a bug,. Why define the object without using it?
- **du**: The normal case. The object is defined and then used.
- **kd**: Normal situation. An object is killed and then redefined.
- **kk**: Harmless but probably buggy. Did you want to be sure it was really killed?
- **ku**: A bug. The object does not exist.
- **ud**: Usually not a bug because the language permits reassignment at almost any time.
- **uu**: Normal situation

# Limitation of Static Anomaly Detection

- **Arrays:** reference to array is done at the level of the element in the array while defining or killing the array is done at the object level
- **Dead variable:** There no general solution to prove that the variable is dead. Only at a given point in time that is possible
- **Dynamic subroutine or function names**
- **Concurrency**
- **Interrupts and incidents**

# Data Flow Anomaly Examples

$y = f1(x);$

$y = f2(z);$

Type: **dd**: Defined and then defined again  
Why this happened?

## Four interpretations of Example 1

The first statement is redundant.

The first statement has a fault -- the intended one might be:  $w = f1(x)$ .

The second statement has a fault – the intended one might be:  $v = f2(z)$ .

There is a missing statement in between the two:  $v = f3(x)$ .

# Def and Use Definition

- Occurrences of variables
- **A *definition (def)*** is a location where a value for a variable is stored into memory (assignment, input, etc.).
  - $X = 0;$
- **Use:** This occurs when the value is fetched from the memory location of the variable. There are **two forms** of uses of a variable.
  - Computation use (c-use)
    - Example:  $x = 2 * y;$
  - Predicate use (p-use)
    - Example: `if (y > 100) { ... }`

# Data flow testing criteria

- Data flow testing criteria use the fact that values are carried from defs to uses. We call these *du-pairs* (they are also known as *definition-use*, *def-use*, and *du* associations in the testing literature).
- The idea of data flow criteria is to exercise du-pairs in various ways

# Example :Def and Use

Given the following code statements :

1.  $X = y + 3;$
2.  $X = 10;$
3. ....
4. ....
5.  $\text{If } (x < y) \{$

Write all definitions and uses for the three statements?

$\text{Def}(1) = \{x\}$                        $\text{use}(1) = \{y\}$

$\text{Def}(2) = \{x\}$

$\text{use}(5) = \{x, y\}$



# Example :Def and Use

1. $x = y + 3$	$\text{def}(1) = \{x\}$	$\text{use}(1) = \{y\}$
2. $y = 10;$	$\text{def}(2) = \{y\}$	
3. ...		
4. ...		
5. $X = x2 + y$	$\text{def}(5) = \{x\}$	$\text{use}(5) = \{x, y\}$
6. If $(x > y)$ {		$\text{use}(6) = \{x, y\}$

- For the path [1 2 3 4 ] the definition of x has not change  $\rightarrow$  we call it definition clear path (dc)
- For the path [1 2 3 4 5 6] the definition of x changed (redefined)  $\rightarrow$  we do not have a clear path

# Example :Def and Use

1.

2. If ( $x < 1$ ) {

3.      $y = z + 1$

4. }

def(3)={y}

use(2) = {x}

use(2) = {z}  $\rightarrow$  c-used

# Def Occurrences

**def** may occur for variable  $x$  in the following situations:

1.  $x$  appears on the left side of an assignment statement
2.  $x$  is an actual parameter in a call site and its value is changed within the method
3.  $x$  is a formal parameter of a method (an implicit def when the method begins execution)
4.  $x$  is an input to the program

# Use Occurrences

A use may occur for variable  $x$  in the following situations:

1.  $x$  appears on the right side of an assignment statement
2.  $x$  appears in a conditional test (note that such a test is always associated with at least two edges)
3.  $x$  is an actual parameter to a method
4.  $x$  is an output of the program
5.  $x$  is an output of a method in a return statement or returned as a parameter

## *def-clear*

- An important concept when discussing data flow criteria is that a def of a variable may or may not reach a particular use. The most obvious reason that a def of a variable  $v$  at location  $li$  (a location could be a node or an edge) will not reach a use at location  $lj$  is because no path goes from  $li$  to  $lj$
  - A more subtle reason is that the variable's value may be changed by another def before it reaches the use
- No location between  $li$  and  $lj$  changes the value.

# Data Flow Graph

- A data flow graph is a directed graph constructed as follows.
  - A sequence of **definitions** and **c-uses** is associated with each **node** of the graph.
  - A set of **p-uses** is associated with each **edge** of the graph.
  - The entry node has a definition of each edge parameter and each nonlocal variable used in the program.
  - The exit node has an undefinition of each local variable.

# Motivation for Data Flow Graph

- A program unit accepts inputs, performs computations, assigns new values to variables, and returns results.
- One can visualize of “flow” of data values from one statement to another.
- A data value produced in one statement is expected to be used later.
- The memory location for a variable is accessed in a “desirable” way.
- Verify the correctness of data values “defined” (i.e. generated)
- Observe that all the “uses” of the value produce the desired results.

```

public int pat (char[] subject, char[] pattern)
{
    // Post: if pattern is not a substring of subject, return -1
    //       else return (zero-based) index where the pattern (first)
    //       starts in subject

```

```

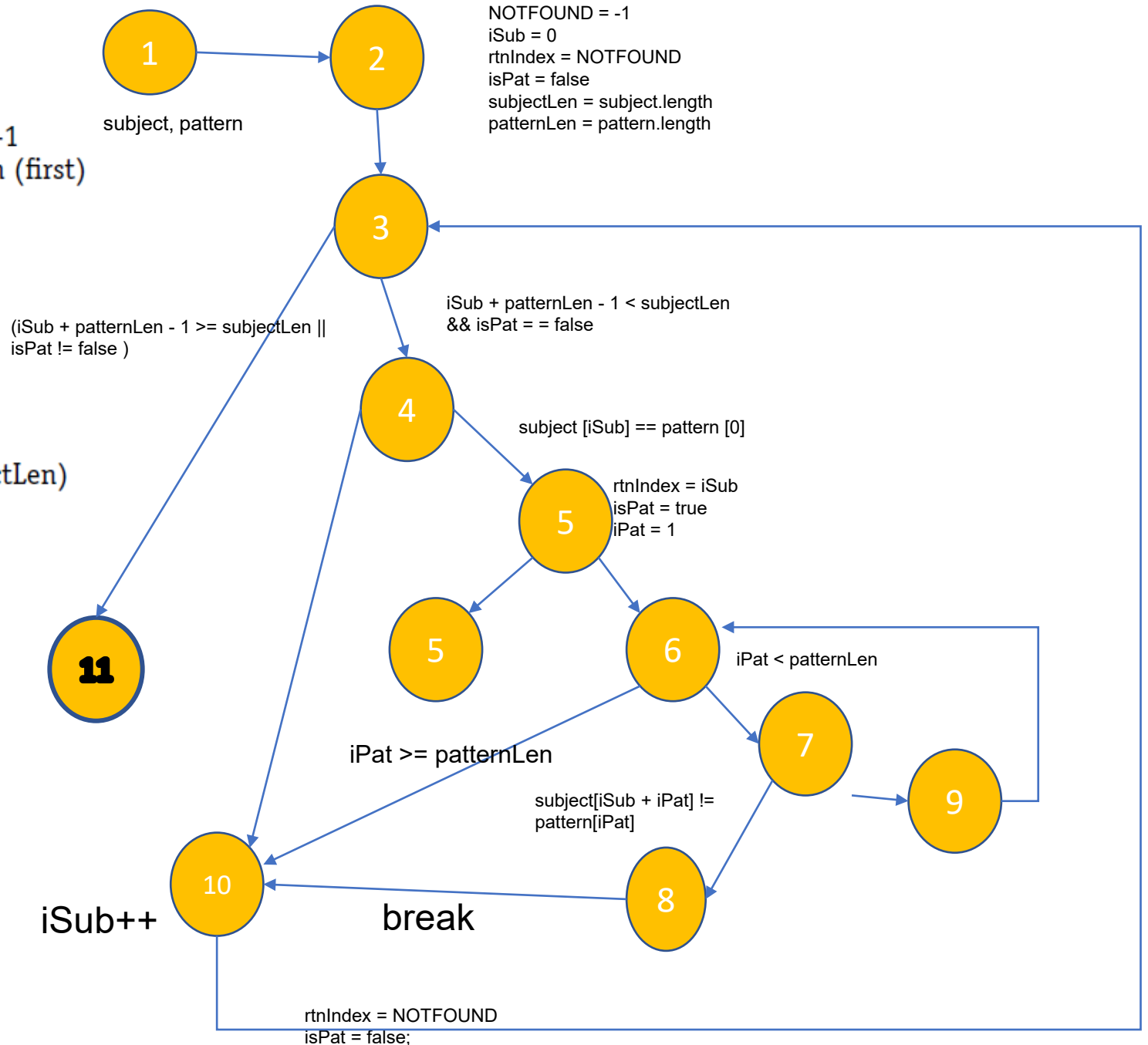
    final int NOTFOUND = -1;
    int iSub = 0, rtnIndex = NOTFOUND;
    boolean isPat = false;
    int subjectLen = subject.length;
    int patternLen = pattern.length;

```

```

    while (isPat == false && iSub + patternLen - 1 < subjectLen)
    {
        if (subject[iSub] == pattern[0])
        {
            rtnIndex = iSub; // Starting at zero
            isPat = true;
            for (int iPat = 1; iPat < patternLen; iPat++)
            {
                if (subject[iSub + iPat] != pattern[iPat])
                {
                    rtnIndex = NOTFOUND;
                    isPat = false;
                    break; // out of for loop
                }
            }
            iSub++;
        }
    }
    return (rtnIndex);

```





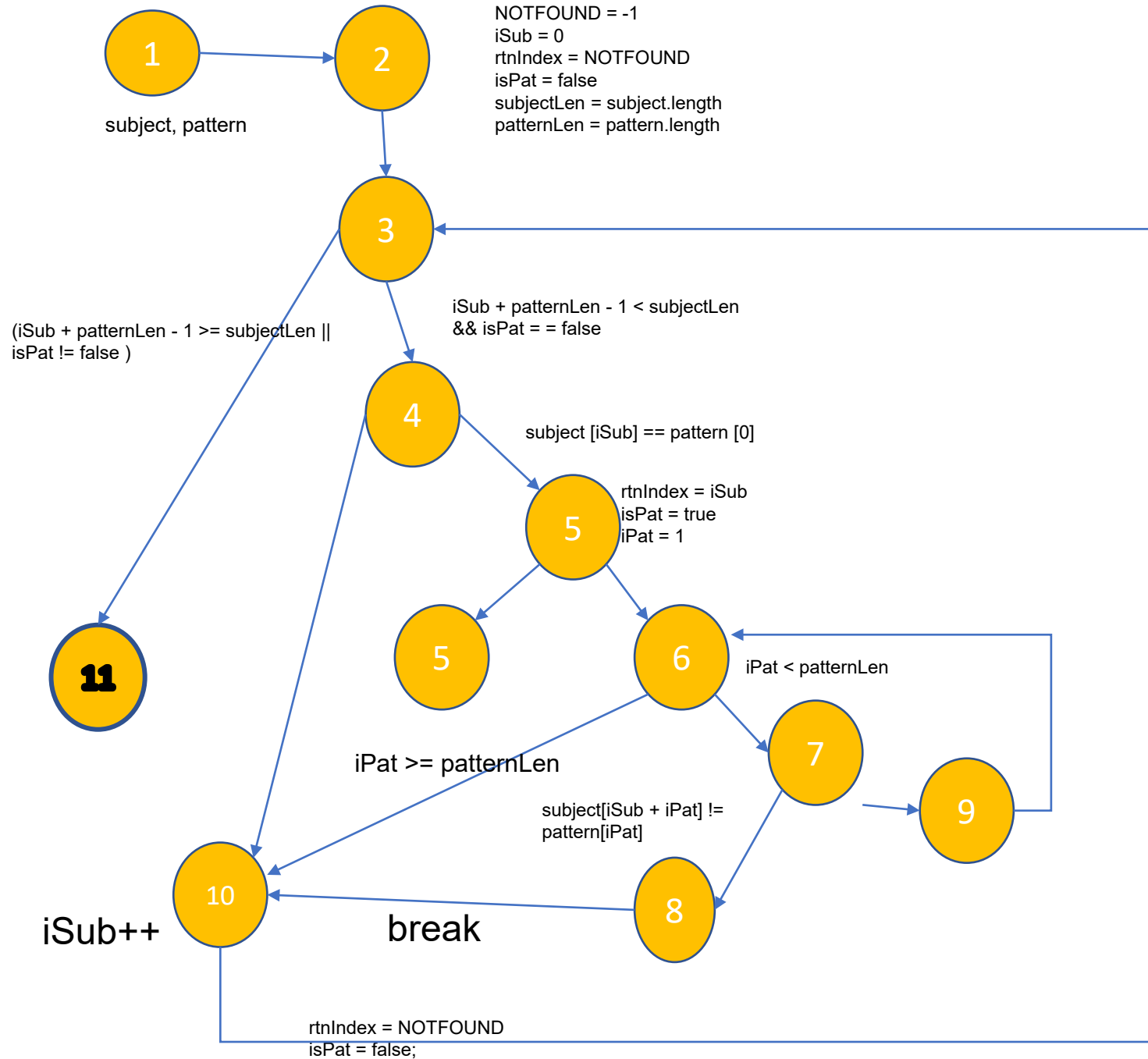
def(1) = { subject, pattern }

def(2) = {NOTFOUND, iSub, rtnIndex,isPat,  
subjectLen, atternLen }

use(2) = {subject , pattern }

use (3,11) = use (3,4) = { iSub, patternLen, subjectLen, isPat  
}

use(4,10) = use(4,5) = { subject, iSub, pattern }



```

public int pat (char[] subject, char[] pattern)
{
    // Post: if pattern is not a substring of subject, return -1
    //       else return (zero-based) index where the pattern (first)
    //       starts in subject

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    final int NOTFOUND = -1;
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    int patternLen = pattern.length;

```

```

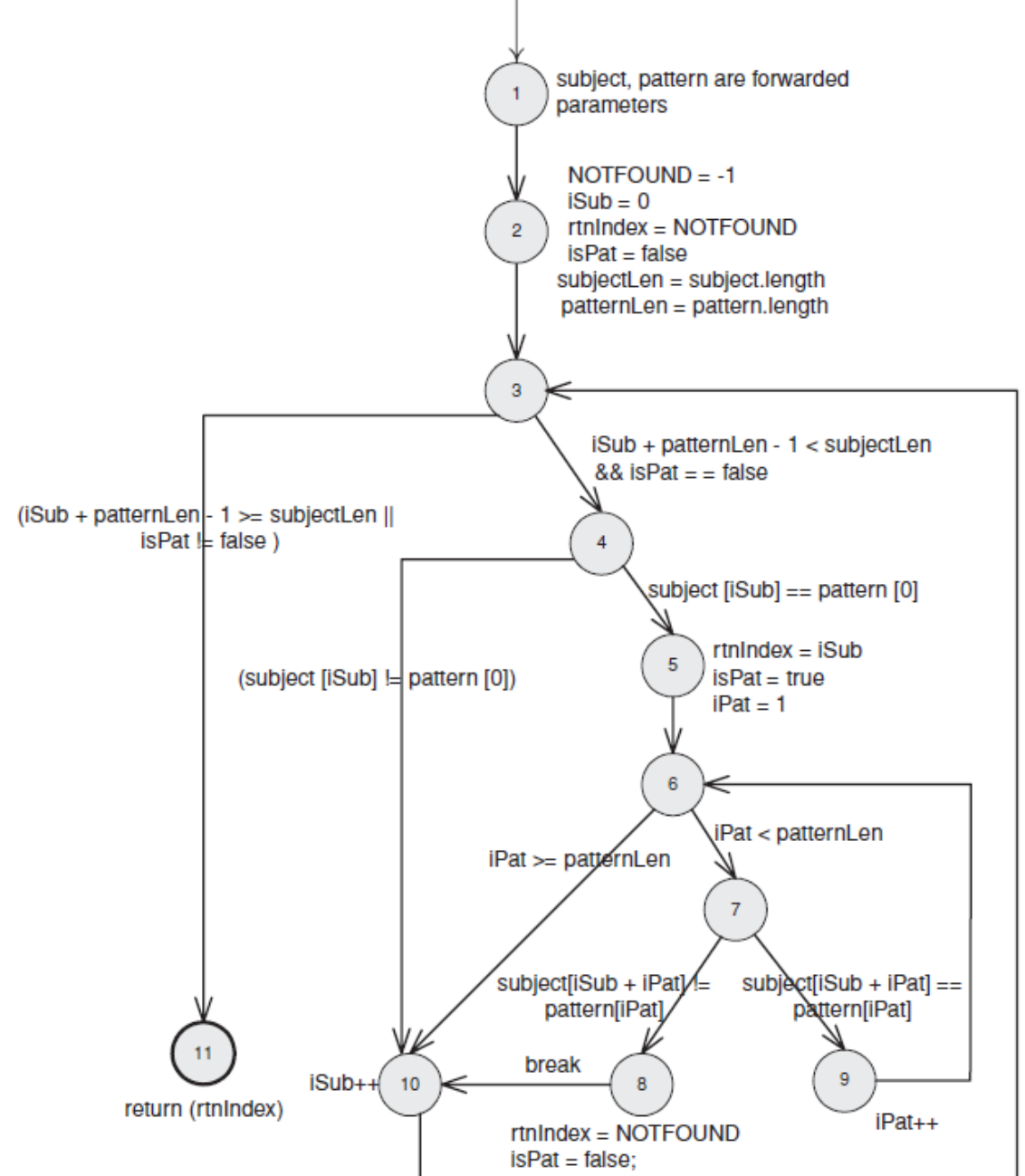
    while (isPat == false && iSub + patternLen - 1 < subjectLen)
    {

```

```

        if (subject[iSub] == pattern[0])
        {
            rtnIndex = iSub; // Starting at zero
            isPat = true;
            for (int iPat = 1; iPat < patternLen; iPat++)
            {
                if (subject[iSub + iPat] != pattern[iPat])
                {
                    rtnIndex = NOTFOUND;
                    isPat = false;
                    break; // out of for loop
                }
            }
            iSub++;
        }
        return (rtnIndex);
    }
}

```



# Data Flow Example

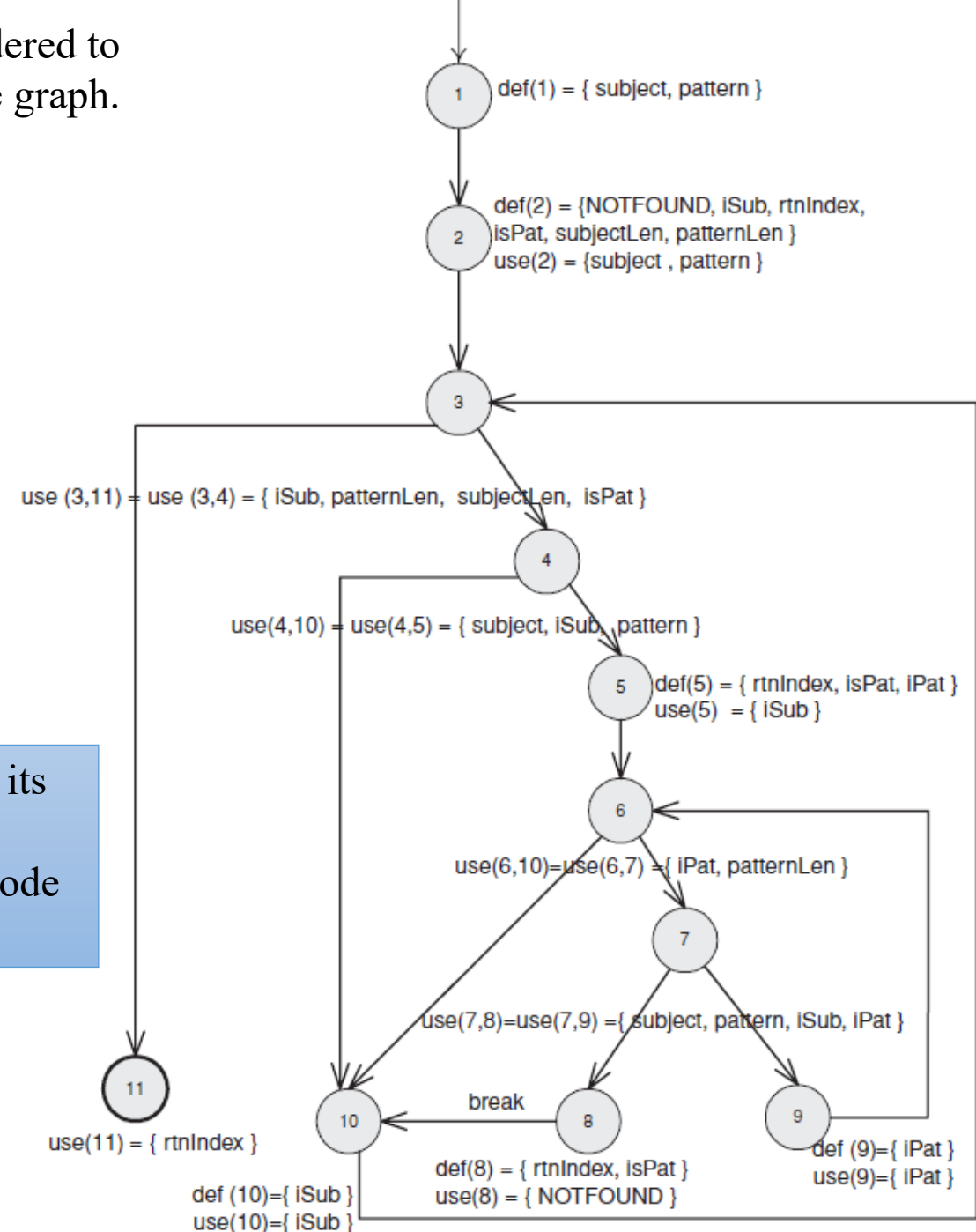
Note that the parameters (*subject* and *pattern*) are considered to be *explicitly defined* by the first node in the graph.

def set of node 1 is  $def(1) = \{subject, pattern\}$ .

if  $subject[i\ Sub] == pattern[0]$  result in uses of each of the associated variables for both edges in the decision. That is,  $use(4, 10) \equiv use(4,5) \equiv \{subject, i\ Sub, pattern\}$ .

The parameter *subject* is used at node 2 (with a reference to its *length* attribute) and at edges (4, 5), (4, 10), (7, 8), and (7, 9), thus du-paths exist from node 1 to node 2 and from node 1 to each of those four edges.

$iPat++$ , which is equivalent to  $iPat = iPat + 1$ .



# Du-path

- A *du-path* with respect to a variable  $v$  is a simple path that is def-clear with respect to  $v$  from a node  $ni$  for which  $v$  is in  $def(ni)$  to a node  $nj$  for which  $v$  is in  $use(nj)$ .
- Note that a du-path is always associated with a specific **variable  $v$** , a du-path **always has to be simple**, and there may be intervening uses on the path
- **The test criteria for data flow will be defined as sets of du-paths.** This makes the criteria quite simple, but first we need to categorize the du-paths into **several groups**

# Data Flow Coverage Criteria

- Data flow coverage criteria will be defined as sets of du-paths
  - Du-paths will check for definitions of variable reaching their uses
1. Grouping du-paths as per definitions
  2. Grouping du-paths as per definitions and uses

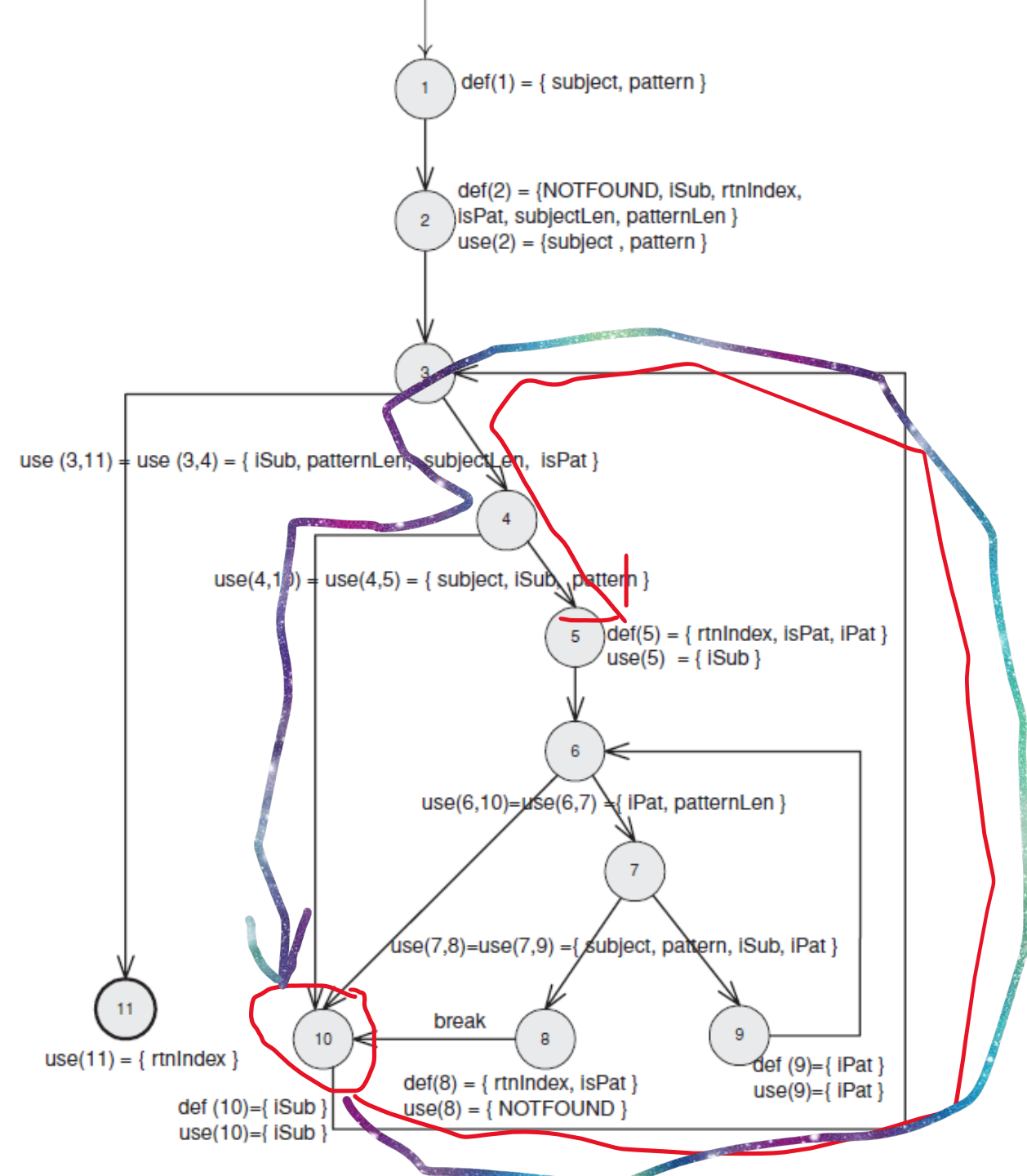
# Grouping du-paths as per definitions

- Consider all du-paths with respect to a given variable in a given node
- The def-path set  $du(n_i, v)$  is the set of du-paths with respect to variable  $v$  that start at node  $n_i$
- We do not group du-paths by uses

The first grouping of du-paths is according to definitions. Specifically, consider all of the du-paths with respect to a given variable defined in a given node

Let the *def-path* set  $du(n_i, v)$  be the set of du-paths with respect to variable  $v$  that start at node  $n_i$

- The def-path set for the use of *isub* at node 10 is:
- $du(10, iSub) = \{[10, 3, 4], [10, 3, 4, 5], [10, 3, 4, 5, 6, 7, 8], [10, 3, 4, 5, 6, 7, 9], [10, 3, 4, 5, 6, 10], [10, 3, 4, 5, 6, 7, 8, 10], [10, 3, 4, 10], [10, 3, 11]\}$
- This def-path set can be broken up into the following def-pair sets:
- $du(10, 4, iSub) = is\{[10, 3, 4]\}$
- $du(10, 5, iSub) = \{[10, 3, 4, 5]\}$
- $du(10, 8, iSub) = \{[10, 3, 4, 5, 6, 7, 8]\}$
- $du(10, 9, iSub) = \{[10, 3, 4, 5, 6, 7, 9]\}$
- $du(10, 10, iSub) = \{[10, 3, 4, 5, 6, 10], [10, 3, 4, 5, 6, 7, 8, 10], [10, 3, 4, 10]\}$
- $du(10, 11, iSub) = \{[10, 3, 11]\}$



# Defs and uses at each node in the CFG for TestPat

node	def	use
1	{subject, pattern}	
2	{NOTFOUND, isPat, iSub, rtnIndex, subjectLen, patternLen}	{subject, pattern}
3		
4		
5	{rtnIndex, isPat, iPat}	{iSub}
6		
7		
8	{rtnIndex, isPat}	{NOTFOUND}
9	{iPat}	{iPat}
10	{iSub}	{iSub}
11		{rtnIndex}



# du-paths for each variable in TestPat

edge	use
(1, 2)	
(2, 3)	
(3, 4)	{iSub, patternLen, subjectLen, isPat}
(3, 11)	{iSub, patternLen, subjectLen, isPat}
(4, 5)	{subject, iSub, pattern}
(4, 10)	{subject, iSub, pattern}
(5, 6)	
(6, 7)	{iPat, patternLen}
(6, 10)	{iPat, patternLen}
(7, 8)	{subject, iSub, iPat, pattern}
(7, 9)	{subject, iSub, iPat, pattern}
(8, 10)	
(9, 6)	
(10, 3)	

# Grouping du-paths as per definitions and uses

- A def-pair set  $du(n_i, n_j, v)$  is the set of du-paths with respect to variable  $v$  that start at node  $n_i$  and end at node  $n_j$
- A def-per set collect together all the simple ways to get from a given definition to a given use
- A def-pair for a def at node  $n_i$  is the union of all the def-path set for the def.  $du(n_i, v) = \bigcup_j du(n_i, n_j, v)$

# Definitions

**CRITERION 2.9 All-Defs Coverage (ADC):** For each def-pair set  $S = du(n, v)$ ,  $TR$  contains at least one path  $d$  in  $S$ .

**CRITERION 2.10 All-Uses Coverage (AUC):** For each def-pair set  $S = du(n_i, n_j, v)$ ,  $TR$  contains at least one path  $d$  in  $S$ .

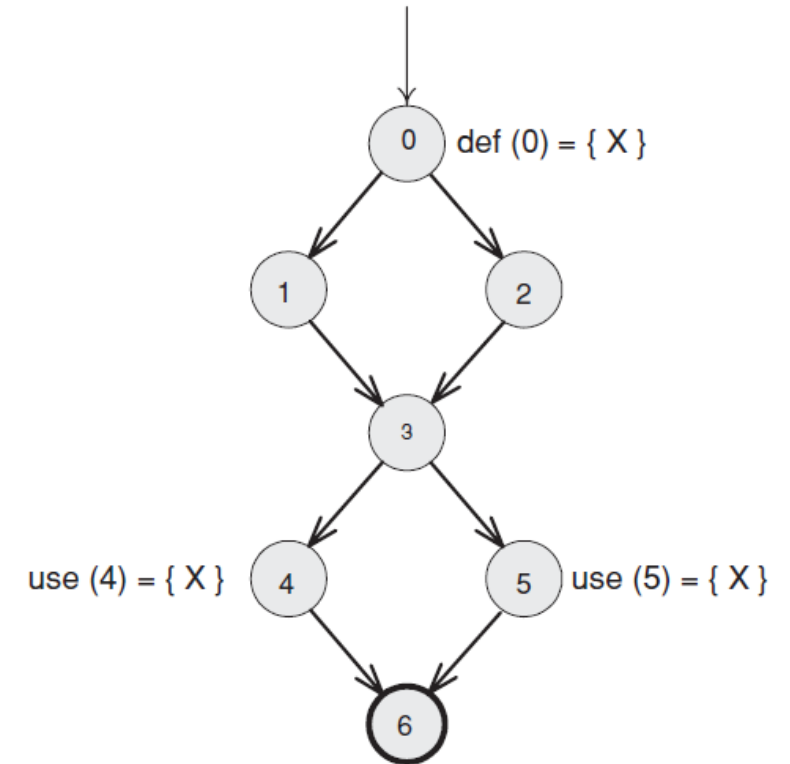
**CRITERION 2.11 All-du-Paths Coverage (ADUPC):** For each def-pair set  $S = du(n_i, n_j, v)$ ,  $TR$  contains every path  $d$  in  $S$ .

All-defs
0-1-3-4

All-uses
0-1-3-4
0-1-3-5

All-du-paths
0-1-3-4
0-1-3-5
0-2-3-4
0-2-3-5

ADC: [0 1 3 4] or [0 2 3 5]



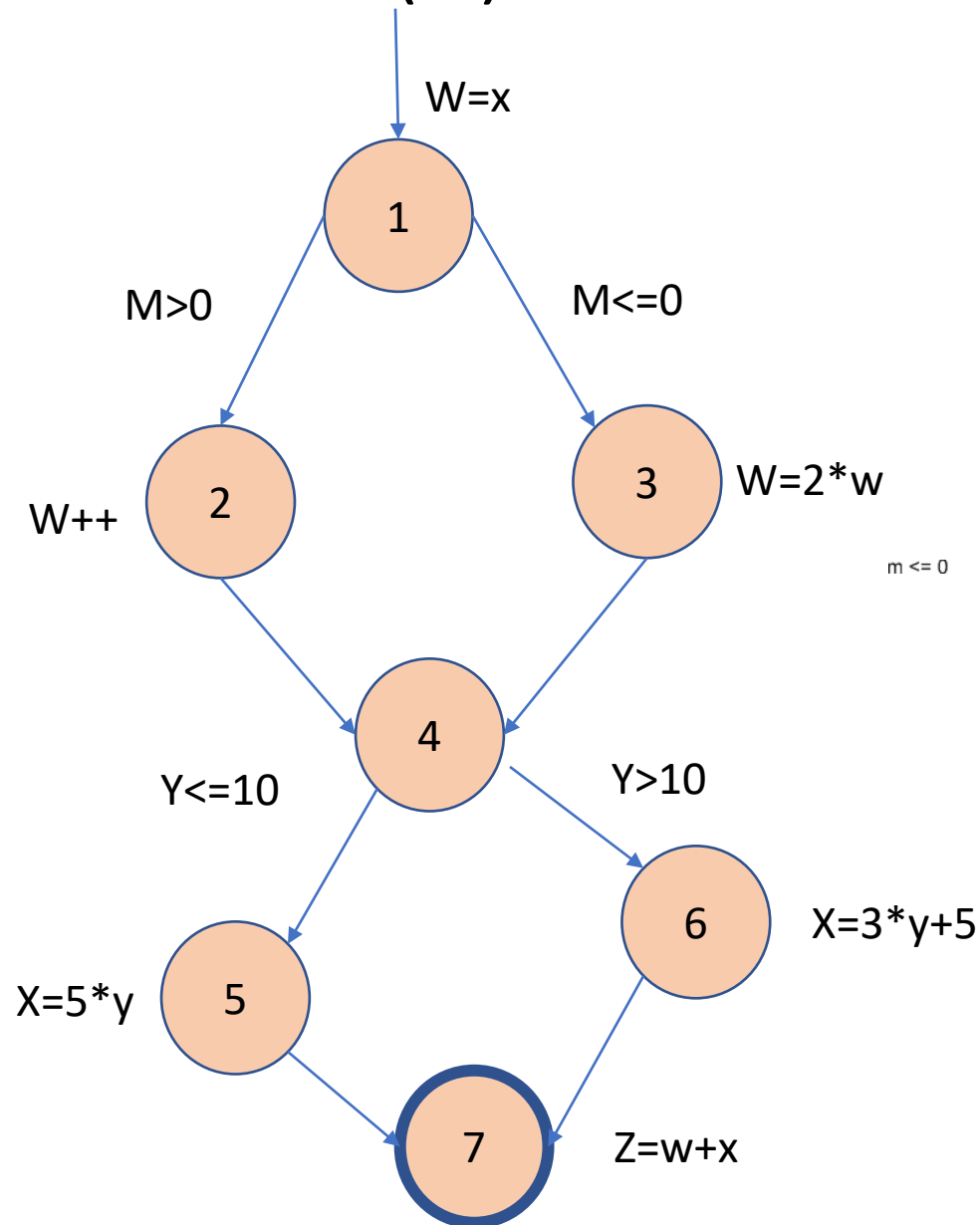
# Exercise

Use the following program fragment for questions a-e below.

- (a) Draw a control flow graph for this program fragment. Use the node numbers given above.
- (b) Which nodes have defs for variable *w*?
- (c) Which nodes have uses for variable *w*?
- (d) Are there any du-paths with respect to variable *w* from node 1 to node 7? If not, explain why not. If any exist, show one.
- (e) Enumerate all of the du-paths for variables *w* and *x*.

```
w = x;           // node 1
if (m > 0)
{
    w++;          // node 2
}
else
{
    w=2*w;        // node 3
}
// node 4 (no executable statement)
if (y <= 10)
{
    x = 5*y;      // node 5
}
else
{
    x = 3*y+5;    // node 6
}
z = w + x;       // node 7
```

# Solution for (a): Draw a control ow graph

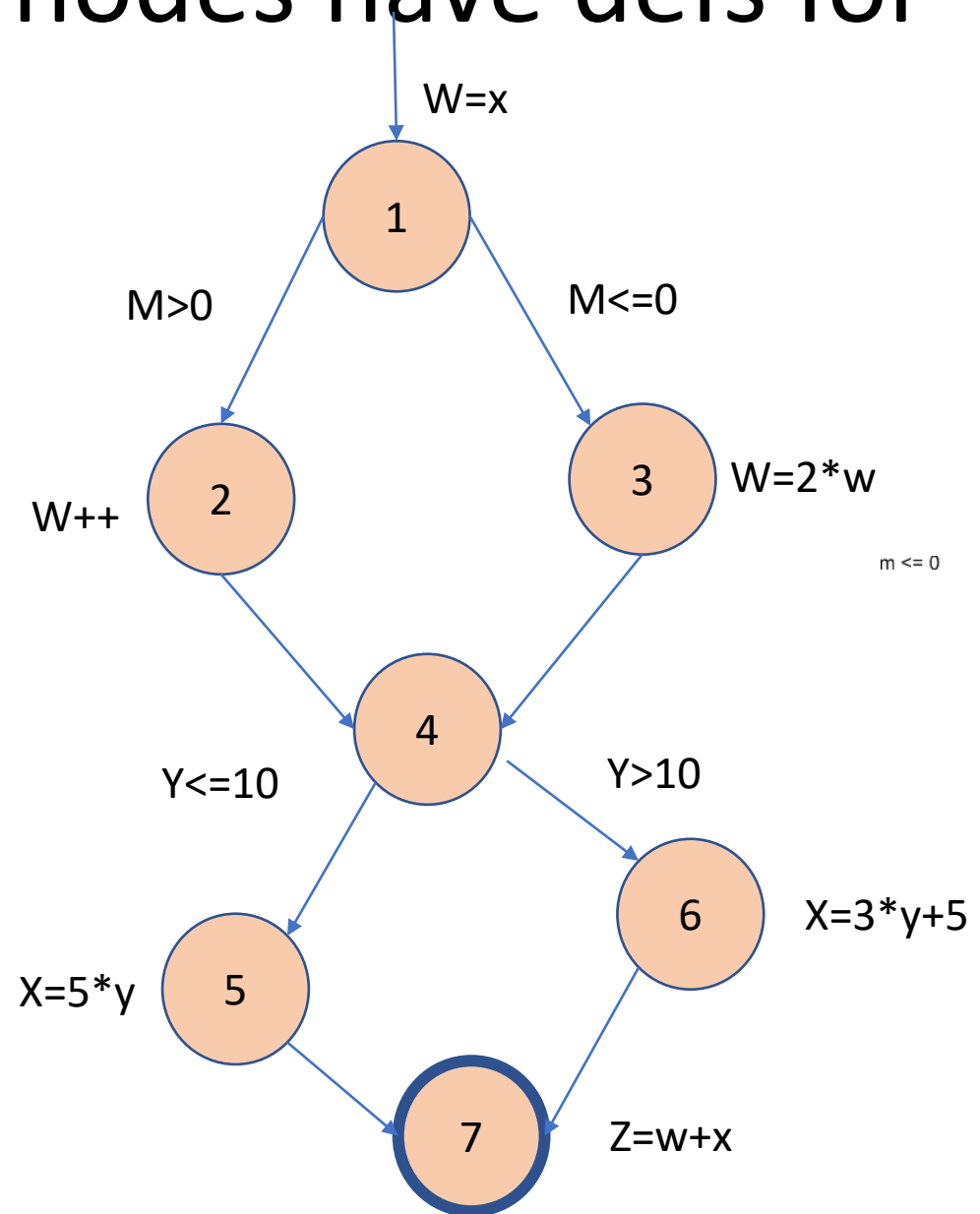


```
w = x;           // node 1
if (m > 0)
{
    w++;          // node 2
}
else
{
    w = 2 * w;    // node 3
}
// node 4 (no executable statement)
if (y <= 10)
{
    x = 5 * y;    // node 5
}
else
{
    x = 3 * y + 5; // node 6
}
z = w + x;       // node 7
```

# Solution for: (b) Which nodes have defs for variable w?

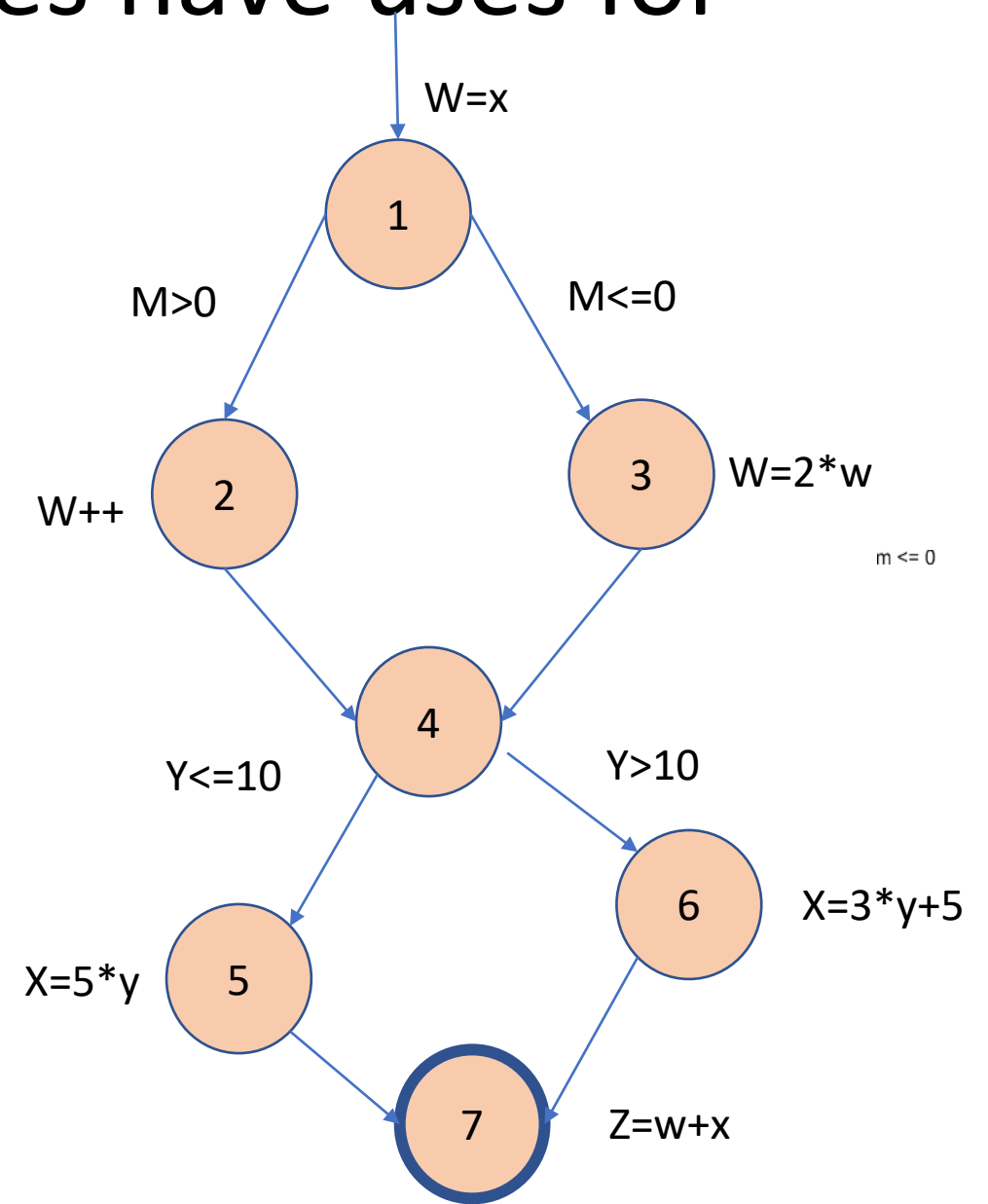
- $\text{Def}(1): \{w\}$
- $\text{Def}(3): \{w\}$
- $\text{Def}(2): \{w\}$

So nodes 1,2,3 have def for variable w



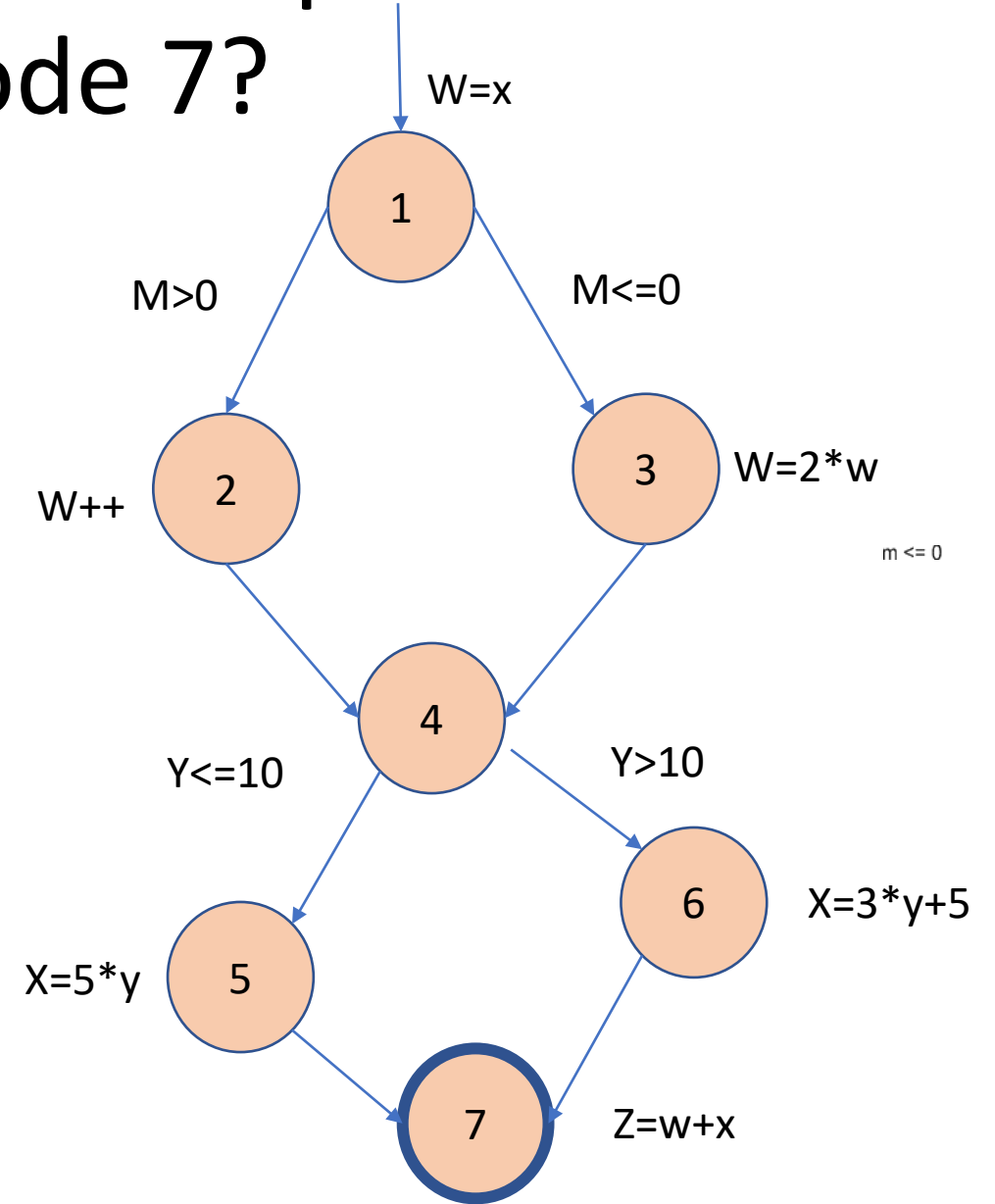
# Solution for:(c) Which nodes have uses for variable w?

- Nodes 2,3,7 have uses for variable w



(d) Are there any du-paths with respect to variable  $w$  from node 1 to node 7?

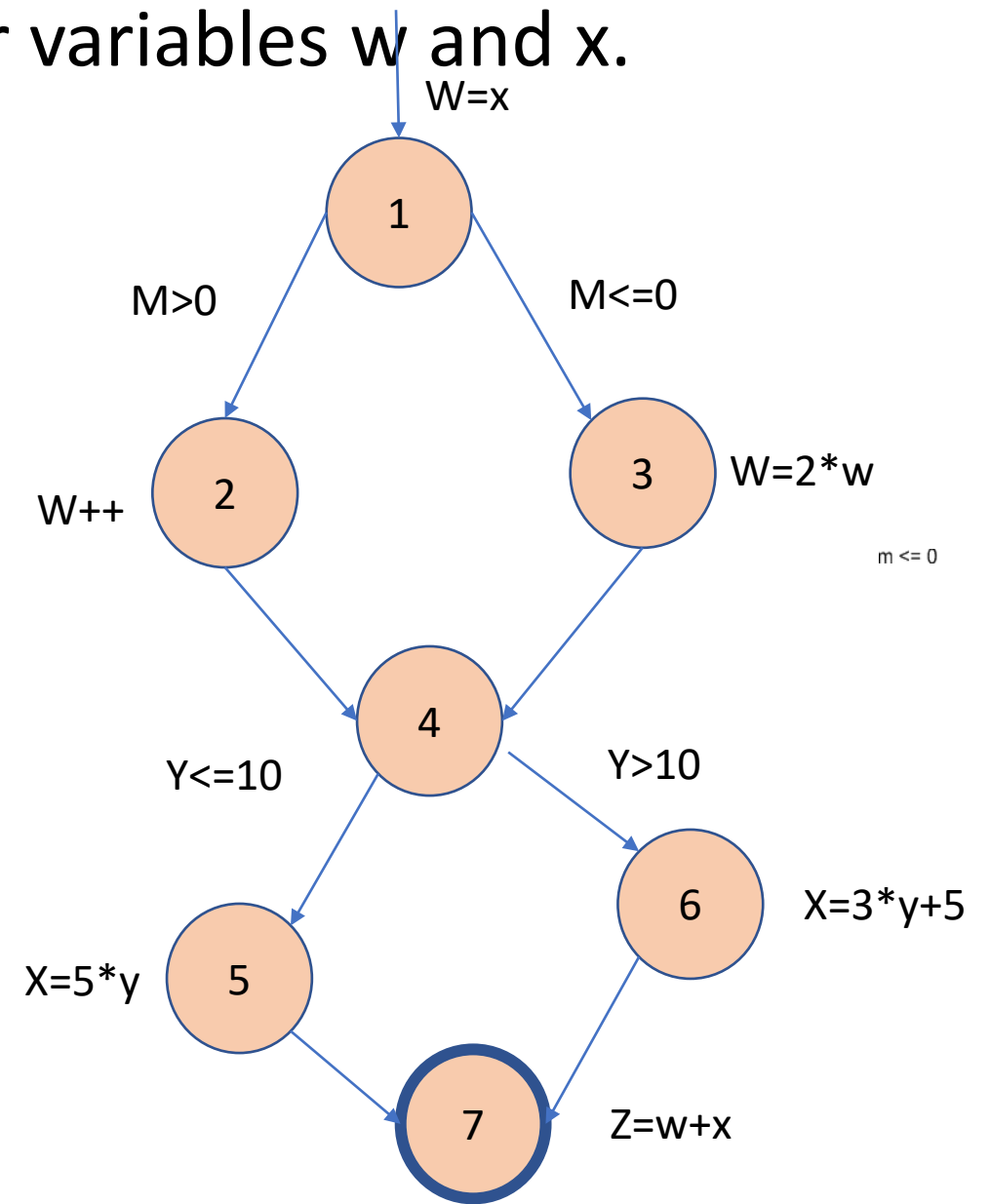
- Node 1 have def
- Node 7 has use but in node 2 we have redefine so there is no clear path from 1 to 7





(e) Enumerate all of the du-paths for variables  $w$  and  $x$ .

- The du-path should be clear
- For variable  $w$ 
  - [1 2] [1 3]
  - [2 4 5 7] [2 4 6 7]
  - [3 4 5 7] [3 4 6 7]
- For variable  $x$ 
  - [5 7] [6 7]



# Test Coverage Metrics -Recap

There are many different ways to measure test coverage, but some common metrics include:

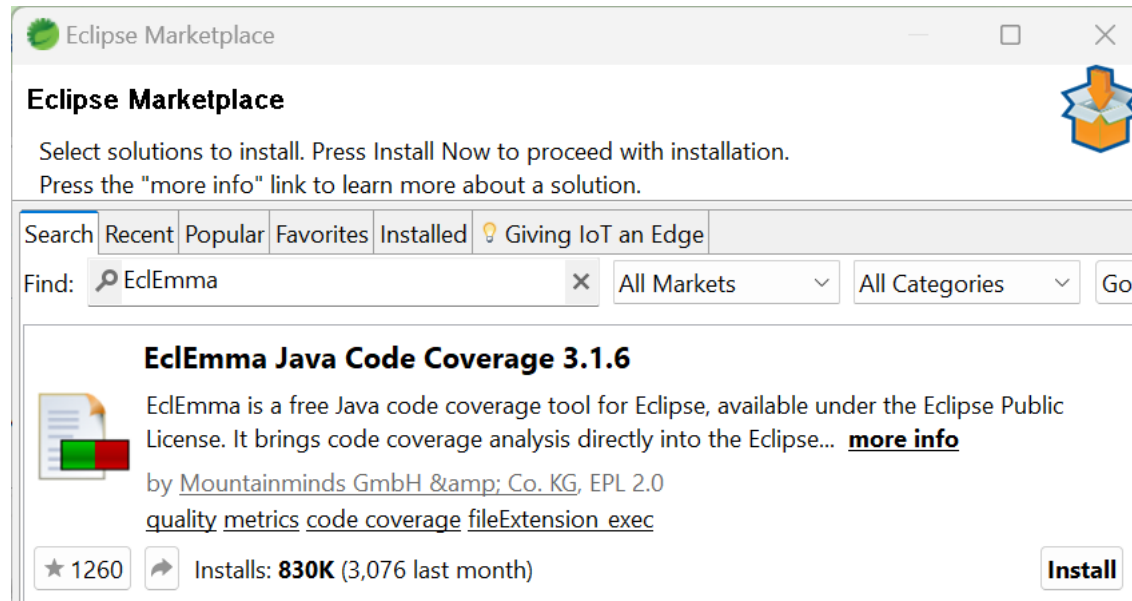
- **Lines of code covered:** This metric simply measures the number of lines of code that are covered by tests. This is a good starting point, but it doesn't give the whole picture, since some lines of code are more important than others.
- **Blocks covered:** This metric measures the number of blocks of code (e.g., if-statements, for-loops, etc.) that are covered by tests. This is a more granular metric than lines of code, but it still doesn't give the whole picture.
- **Functions/methods covered:** This metric measures the number of functions or methods that are covered by tests. This is a more granular metric than blocks, but it still doesn't give the whole picture.
- **Statements covered:** This metric measures the number of statements that are covered by tests. This is the most granular metric, but it still doesn't give the whole picture.
- **Conditions covered:** This metric measures the number of conditions (e.g., if-statements, boolean expressions, etc.) that are covered by tests.

# Code Coverage Tools

- JaCoCo

<https://www.eclemma.org/jacoco/>

Download the plugin into your STS or Eclipse . Go to help → marketplace then search for *EclEmma*










# Example

```
class CurrencyNameTest {  
    @Test  
    void test() {  
        BorwaserTest coverage = new BorwaserTest();  
  
        String currency1= coverage.getCurrencyName("USD");  
        assertEquals("American Dollar",currency1);  
    }  
}
```

```
public String getCurrencyName(String code) {  
    switch (code) {  
        case "USD":  
            return "American Dollar";  
        case "JOD" :  
            return "Jordanian Dinar";  
        default:  
            return "Wrong Currency";  
    }  
}
```

# Testing Results

Coverage ×					
Element		Coverage	Covered Instructions	Missed Instr...	Total Instructions
▼ selenium		 86.2 %	25	4	29
▼ src		 86.2 %	25	4	29
▼ chrome		 86.2 %	25	4	29
▼ CaseCoverageTest.java		 71.4 %	10	4	14
▼ CaseCoverageTest		 71.4 %	10	4	14
• getCurrencyName(String)		 63.6 %	7	4	11
> CurrencyNameTest.java		 100.0 %	15	0	15

Refine the test cases to achieve 100% coverage

# Reading Material

- **Control Flow Analysis for Java Methods**
  - <https://www.jacoco.org/jacoco/trunk/doc/flow.html>
- This article is part of the course material