

# Sets

The Clowder Project Authors

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This chapter (will eventually) contain material on axiomatic set theory, as well as a couple other things.

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## 1 Sets and Functions

### 1.1 Functions

**Definition 1.1.1.1.** A **function** is a functional and total relation.

**Notation 1.1.1.2.** Throughout this work, we will sometimes denote a function  $f: X \rightarrow Y$  by

$$f \stackrel{\text{def}}{=} \llbracket x \mapsto f(x) \rrbracket.$$

1. For example, given a function

$$\Phi: \text{Hom}_{\text{Sets}}(X, Y) \rightarrow K$$

taking values on a set of functions such as  $\text{Hom}_{\text{Sets}}(X, Y)$ , we will sometimes also write

$$\Phi(f) \stackrel{\text{def}}{=} \Phi(\llbracket x \mapsto f(x) \rrbracket).$$

2. This notational choice is based on the lambda notation

$$f \stackrel{\text{def}}{=} (\lambda x. f(x)),$$

but uses a “ $\mapsto$ ” symbol for better spacing and double brackets instead of either:

- (a) Square brackets  $[x \mapsto f(x)]$ ;
- (b) Parentheses  $(x \mapsto f(x))$ ;

hoping to improve readability when dealing with e.g.:

- (a) Equivalence classes, cf.:
  - i.  $\llbracket [x] \mapsto f([x]) \rrbracket$
  - ii.  $\llbracket [x] \mapsto f([x]) \rrbracket$
  - iii.  $(\lambda [x]. f([x]))$
- (b) Function evaluations, cf.:
  - i.  $\Phi(\llbracket x \mapsto f(x) \rrbracket)$
  - ii.  $\Phi((x \mapsto f(x)))$
  - iii.  $\Phi((\lambda x. f(x)))$

3. We will also sometimes write  $-_1, -_2$ , etc. for the arguments of a function. Some examples include:

- (a) Writing  $f(-_1)$  for a function  $f: A \rightarrow B$ .
- (b) Writing  $f(-_1, -_2)$  for a function  $f: A \times B \rightarrow C$ .
- (c) Given a function  $f: A \times B \rightarrow C$ , writing

$$f(a, -): B \rightarrow C$$

for the function  $\llbracket b \mapsto f(a, b) \rrbracket$ .

(d) Denoting a composition of the form

$$A \times B \xrightarrow{\phi \times \text{id}_B} A' \times B \xrightarrow{f} C$$

by  $f(\phi(-_1), -_2)$ .

4. Finally, given a function  $f: A \rightarrow B$ , we write

$$\text{ev}_a(f) \stackrel{\text{def}}{=} f(a)$$

for the value of  $f$  at some  $a \in A$ .

For an example of the above notations being used in practice, see the proof of the adjunction

$$(A \times - \dashv \text{Hom}_{\mathbf{Sets}}(A, -)): \mathbf{Sets} \begin{array}{c} \xrightarrow{A \times -} \\ \perp \\ \xleftarrow{\text{Hom}_{\mathbf{Sets}}(A, -)} \end{array} \mathbf{Sets},$$

stated in [Constructions With Sets, Item 2](#) of [Proposition 1.3.1.2](#).

## 2 The Enrichment of Sets in Classical Truth Values

### 2.1 $(-2)$ -Categories

**Definition 2.1.1.1.** A  $(-2)$ -category is the “necessarily true” truth value.<sup>1,2,3</sup>

### 2.2 $(-1)$ -Categories

**Definition 2.2.1.1.** A  $(-1)$ -category is a classical truth value.

**Remark 2.2.1.2.** <sup>4</sup> $(-1)$ -categories should be thought of as being “categories enriched in  $(-2)$ -categories”, having a collection of objects and, for each pair of objects, a Hom-object  $\text{Hom}(x, y)$  that is a  $(-2)$ -category (i.e. trivial).

Therefore, a  $(-1)$ -category  $C$  is either ([\[BS10, pp. 33–34\]](#)):

1. *Empty*, having no objects;

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<sup>1</sup>Thus, there is only one  $(-2)$ -category.

<sup>2</sup>A  $(-n)$ -category for  $n = 3, 4, \dots$  is also the “necessarily true” truth value, coinciding with a  $(-2)$ -category.

<sup>3</sup>For motivation, see [\[BS10, p. 13\]](#).

<sup>4</sup>For more motivation, see [\[BS10, p. 13\]](#).

2. *Contractible*, having a collection of objects  $\{a, b, c, \dots\}$ , but with  $\text{Hom}_C(a, b)$  being a  $(-2)$ -category (i.e. trivial) for all  $a, b \in \text{Obj}(C)$ , forcing all objects of  $C$  to be uniquely isomorphic to each other.

As such, there are only two  $(-1)$ -categories, up to equivalence:

- The  $(-1)$ -category false (the empty one);
- The  $(-1)$ -category true (the contractible one).

**Definition 2.2.1.3.** The **poset of truth values**<sup>5</sup> is the poset  $(\{\text{true}, \text{false}\}, \preceq)$  consisting of

- *The Underlying Set.* The set  $\{\text{true}, \text{false}\}$  whose elements are the truth values true and false.
- *The Partial Order.* The partial order

$$\preceq: \{\text{true}, \text{false}\} \times \{\text{true}, \text{false}\} \rightarrow \{\text{true}, \text{false}\}$$

on  $\{\text{true}, \text{false}\}$  defined by<sup>6</sup>

$$\begin{aligned} \text{false} &\preceq \text{false} \stackrel{\text{def}}{=} \text{true}, \\ \text{true} &\preceq \text{false} \stackrel{\text{def}}{=} \text{false}, \\ \text{false} &\preceq \text{true} \stackrel{\text{def}}{=} \text{true}, \\ \text{true} &\preceq \text{true} \stackrel{\text{def}}{=} \text{true}. \end{aligned}$$

**Notation 2.2.1.4.** We also write  $\{t, f\}$  for the poset  $\{\text{true}, \text{false}\}$ .

**Proposition 2.2.1.5.** The poset of truth values  $\{t, f\}$  is Cartesian closed with product given by<sup>7</sup>

$$\begin{aligned} t \times t &= t, \\ t \times f &= f, \\ f \times t &= f, \\ f \times f &= f, \end{aligned}$$

<sup>5</sup>*Further Terminology:* Also called the **poset of  $(-1)$ -categories**.

<sup>6</sup>This partial order coincides with logical implication.

<sup>7</sup>Note that  $\times$  coincides with the “and” operator, while  $\mathbf{Hom}_{\{t, f\}}$  coincides with the logical

and internal Hom  $\mathbf{Hom}_{\{t,f\}}$  given by the partial order of  $\{t, f\}$ , i.e. by

$$\mathbf{Hom}_{\{t,f\}}(t, t) = t,$$

$$\mathbf{Hom}_{\{t,f\}}(t, f) = f,$$

$$\mathbf{Hom}_{\{t,f\}}(f, t) = t,$$

$$\mathbf{Hom}_{\{t,f\}}(f, f) = t.$$

*Proof. Existence of Products:* We claim that the products  $t \times t$ ,  $t \times f$ ,  $f \times t$ , and  $f \times f$  satisfy the universal property of the product in  $\{t, f\}$ . Indeed, consider the diagrams

$$\begin{array}{ccccccc} & & P_1 & & & & \\ & p_1^1 \swarrow & \downarrow \exists! & \searrow p_2^1 & & & \\ t & \xleftarrow{\text{pr}_1} & t \times t & \xrightarrow{\text{pr}_2} & t & & \\ & \underbrace{\hspace{2cm}}_{=t} & & & & & \end{array} \quad \begin{array}{ccccccc} & & P_2 & & & & \\ & p_1^2 \swarrow & \downarrow \exists! & \searrow p_2^2 & & & \\ t & \xleftarrow{\text{pr}_1} & t \times f & \xrightarrow{\text{pr}_2} & f & & \\ & \underbrace{\hspace{2cm}}_{=f} & & & & & \end{array} \quad \begin{array}{ccccccc} & & P_3 & & & & \\ & p_1^3 \swarrow & \downarrow \exists! & \searrow p_2^3 & & & \\ f & \xleftarrow{\text{pr}_1} & f \times t & \xrightarrow{\text{pr}_2} & t & & \\ & \underbrace{\hspace{2cm}}_{=f} & & & & & \end{array} \quad \begin{array}{ccccccc} & & P_4 & & & & \\ & p_1^4 \swarrow & \downarrow \exists! & \searrow p_2^4 & & & \\ f & \xleftarrow{\text{pr}_1} & f \times f & \xrightarrow{\text{pr}_2} & f & & \\ & \underbrace{\hspace{2cm}}_{=f} & & & & & \end{array}$$

Here:

1. If  $P_1 = t$ , then  $p_1^1 = p_2^1 = \text{id}_t$ , and there's indeed a unique morphism from  $P_1$  to  $t$  making the diagram commute, namely  $\text{id}_t$ ;
2. If  $P_1 = f$ , then  $p_1^1 = p_2^1$  are given by the unique morphism from  $f$  to  $t$ , and there's indeed a unique morphism from  $P_1$  to  $t$  making the diagram commute, namely the unique morphism from  $f$  to  $t$ ;
3. If  $P_2 = t$ , then there is no morphism  $p_2^2$ .
4. If  $P_2 = f$ , then  $p_1^2$  is the unique morphism from  $f$  to  $t$  while  $p_2^2 = \text{id}_f$ , and there's indeed a unique morphism from  $P_2$  to  $f$  making the diagram commute, namely  $\text{id}_f$ ;
5. The proof for  $P_3$  is similar to the one for  $P_2$ ;
6. If  $P_4 = t$ , then there is no morphism  $p_1^4$  or  $p_2^4$ .
7. If  $P_4 = f$ , then  $p_1^4 = p_2^4 = \text{id}_f$ , and there's indeed a unique morphism from  $P_4$  to  $f$  making the diagram commute, namely  $\text{id}_f$ .

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implication operator.

*Cartesian Closedness:* We claim there's a bijection

$$\mathrm{Hom}_{\{t,f\}}(A \times B, C) \cong \mathrm{Hom}_{\{t,f\}}(A, \mathbf{Hom}_{\{t,f\}}(B, C))$$

natural in  $A, B, C \in \{t, f\}$ . Indeed:

- For  $(A, B, C) = (t, t, t)$ , we have

$$\begin{aligned} \mathrm{Hom}_{\{t,f\}}(t \times t, t) &\cong \mathrm{Hom}_{\{t,f\}}(t, t) \\ &= \{\mathrm{id}_{\mathrm{true}}\} \\ &\cong \mathrm{Hom}_{\{t,f\}}(t, t) \\ &\cong \mathrm{Hom}_{\{t,f\}}(t, \mathbf{Hom}_{\{t,f\}}(t, t)). \end{aligned}$$

- For  $(A, B, C) = (t, t, f)$ , we have

$$\begin{aligned} \mathrm{Hom}_{\{t,f\}}(t \times t, f) &\cong \mathrm{Hom}_{\{t,f\}}(t, f) \\ &= \emptyset \\ &\cong \mathrm{Hom}_{\{t,f\}}(t, f) \\ &\cong \mathrm{Hom}_{\{t,f\}}(t, \mathbf{Hom}_{\{t,f\}}(t, f)). \end{aligned}$$

- For  $(A, B, C) = (t, f, t)$ , we have

$$\begin{aligned} \mathrm{Hom}_{\{t,f\}}(t \times f, t) &\cong \mathrm{Hom}_{\{t,f\}}(f, t) \\ &\cong \mathrm{pt} \\ &\cong \mathrm{Hom}_{\{t,f\}}(f, t) \\ &\cong \mathrm{Hom}_{\{t,f\}}(f, \mathbf{Hom}_{\{t,f\}}(f, t)). \end{aligned}$$

- For  $(A, B, C) = (t, f, f)$ , we have

$$\begin{aligned} \mathrm{Hom}_{\{t,f\}}(t \times f, f) &\cong \mathrm{Hom}_{\{t,f\}}(f, f) \\ &\cong \{\mathrm{id}_{\mathrm{false}}\} \\ &\cong \mathrm{Hom}_{\{t,f\}}(f, f) \\ &\cong \mathrm{Hom}_{\{t,f\}}(t, \mathbf{Hom}_{\{t,f\}}(f, f)). \end{aligned}$$

- For  $(A, B, C) = (f, t, t)$ , we have

$$\begin{aligned} \mathrm{Hom}_{\{t,f\}}(f \times t, t) &\cong \mathrm{Hom}_{\{t,f\}}(f, t) \\ &\cong \mathrm{pt} \\ &\cong \mathrm{Hom}_{\{t,f\}}(f, t) \\ &\cong \mathrm{Hom}_{\{t,f\}}(f, \mathbf{Hom}_{\{t,f\}}(t, t)). \end{aligned}$$

- For  $(A, B, C) = (f, t, f)$ , we have

$$\begin{aligned}
 \mathrm{Hom}_{\{t,f\}}(f \times t, f) &\cong \mathrm{Hom}_{\{t,f\}}(f, f) \\
 &\cong \{\mathrm{id}_{\mathrm{false}}\} \\
 &\cong \mathrm{Hom}_{\{t,f\}}(f, f) \\
 &\cong \mathrm{Hom}_{\{t,f\}}(f, \mathbf{Hom}_{\{t,f\}}(t, f)).
 \end{aligned}$$

- For  $(A, B, C) = (f, f, t)$ , we have

$$\begin{aligned}
 \mathrm{Hom}_{\{t,f\}}(f \times f, t) &\cong \mathrm{Hom}_{\{t,f\}}(f, t) \\
 &\cong \mathrm{pt} \\
 &\cong \mathrm{Hom}_{\{t,f\}}(f, t) \\
 &\cong \mathrm{Hom}_{\{t,f\}}(f, \mathbf{Hom}_{\{t,f\}}(f, t)).
 \end{aligned}$$

- For  $(A, B, C) = (f, f, f)$ , we have

$$\begin{aligned}
 \mathrm{Hom}_{\{t,f\}}(f \times f, f) &\cong \mathrm{Hom}_{\{t,f\}}(f, f) \\
 &= \{\mathrm{id}_{\mathrm{false}}\} \\
 &\cong \mathrm{Hom}_{\{t,f\}}(f, f) \\
 &\cong \mathrm{Hom}_{\{t,f\}}(f, \mathbf{Hom}_{\{t,f\}}(f, f)).
 \end{aligned}$$

The proof of naturality is omitted. □

## 2.3 0-Categories

**Definition 2.3.1.1.** A 0-category is a poset.<sup>8</sup>

**Definition 2.3.1.2.** A 0-groupoid is a 0-category in which every morphism is invertible.<sup>9</sup>

## 2.4 Tables of Analogies Between Set Theory and Category Theory

Here we record some analogies between notions in set theory and category theory. Note that the analogies relating to presheaves relate equally well to copresheaves, as the opposite  $X^{\mathrm{op}}$  of a set  $X$  is just  $X$  again.

Basics:

<sup>8</sup>Motivation: A 0-category is precisely a category enriched in the poset of  $(-1)$ -categories.

<sup>9</sup>That is, a set.

SET THEORY	CATEGORY THEORY
Enrichment in $\{\text{true}, \text{false}\}$	Enrichment in Sets
Set $X$	Category $C$
Element $x \in X$	Object $X \in \text{Obj}(C)$
Function	Functor
Function $X \rightarrow \{\text{true}, \text{false}\}$	Functor $C \rightarrow \text{Sets}$
Function $X \rightarrow \{\text{true}, \text{false}\}$	Presheaf $C^{\text{op}} \rightarrow \text{Sets}$

Powersets and categories of presheaves:

SET THEORY	CATEGORY THEORY
Powerset $\mathcal{P}(X)$	Presheaf category $\text{PSh}(C)$
Characteristic function $\chi_{\{x\}}$	Representable presheaf $h_X$
Characteristic embedding $\chi_{(-)} : X \hookrightarrow \mathcal{P}(X)$	Yoneda embedding $\mathcal{Y} : C^{\text{op}} \hookrightarrow \text{PSh}(C)$
Characteristic relation $\chi_X(-_1, -_2)$	Hom profunctor $\text{Hom}_C(-_1, -_2)$
The Yoneda lemma for sets $\text{Hom}_{\mathcal{P}(X)}(\chi_x, \chi_U) = \chi_U(x)$	The Yoneda lemma for categories $\text{Nat}(h_X, \mathcal{F}) \cong \mathcal{F}(X)$
The characteristic embedding is fully faithful, $\text{Hom}_{\mathcal{P}(X)}(\chi_x, \chi_y) = \chi_X(x, y)$	The Yoneda embedding is fully faithful, $\text{Nat}(h_X, h_Y) \cong \text{Hom}_C(X, Y)$
Subsets are unions of their elements $U = \bigcup_{x \in U} \{x\}$ or $\chi_U = \text{colim}_{\chi_x \in \text{Sets}(U, \{\text{t}, \text{f}\})} (\chi_x)$	Presheaves are colimits of representables, $\mathcal{F} \cong \text{colim}_{h_X \in \int_C \mathcal{F}} (h_X)$

Categories of elements:



SET THEORY	CATEGORY THEORY
Assignment $U \mapsto \chi_U$	Assignment $\mathcal{F} \mapsto \int_C \mathcal{F}$ (the category of elements)
Assignment $U \mapsto \chi_U$ giving an isomorphism $\mathcal{P}(X) \cong \text{Sets}(X, \{t, f\})$	Assignment $\mathcal{F} \mapsto \int_C \mathcal{F}$ giving an equivalence $\text{PSh}(C) \stackrel{\text{eq.}}{\cong} \text{DFib}(C)$

Functions between powersets and functors between presheaf categories:

SET THEORY	CATEGORY THEORY
Direct image function $f_*: \mathcal{P}(X) \rightarrow \mathcal{P}(Y)$	Inverse image functor $f^{-1}: \text{PSh}(C) \rightarrow \text{PSh}(\mathcal{D})$
Inverse image function $f^{-1}: \mathcal{P}(Y) \rightarrow \mathcal{P}(X)$	Direct image functor $f_*: \text{PSh}(\mathcal{D}) \rightarrow \text{PSh}(C)$
Direct image with compact support function $f_!: \mathcal{P}(X) \rightarrow \mathcal{P}(Y)$	Direct image with compact support functor $f_!: \text{PSh}(C) \rightarrow \text{PSh}(\mathcal{D})$

Relations and profunctors:

SET THEORY	CATEGORY THEORY
Relation $R: X \times Y \rightarrow \{t, f\}$	Profunctor $\mathfrak{p}: \mathcal{D}^{\text{op}} \times C \rightarrow \text{Sets}$
Relation $R: X \rightarrow \mathcal{P}(Y)$	Profunctor $\mathfrak{p}: C \rightarrow \text{PSh}(\mathcal{D})$
Relation as a cocontinuous morphism of posets $R: (\mathcal{P}(X), \subset) \rightarrow (\mathcal{P}(Y), \subset)$	Profunctor as a colimit-preserving functor $\mathfrak{p}: \text{PSh}(C) \rightarrow \text{PSh}(\mathcal{D})$

## Appendices

## A Other Chapters

### Sets

1. [Sets](#)
2. [Constructions With Sets](#)
3. [Pointed Sets](#)
4. [Tensor Products of Pointed Sets](#)

### Relations

5. [Relations](#)

6. [Constructions With Relations](#)

7. [Equivalence Relations and Apartness Relations](#)

### Category Theory

8. [Categories](#)

### Bicategories

9. [Types of Morphisms in Bicategories](#)

## References

- [BS10] John C. Baez and Michael Shulman. “Lectures on  $n$ -Categories and Cohomology”. In: *Towards higher categories*. Vol. 152. IMA Vol. Math. Appl. Springer, New York, 2010, pp. 1–68. doi: [10.1007/978-1-4419-1524-5\\_1](https://doi.org/10.1007/978-1-4419-1524-5_1). URL: [https://doi.org/10.1007/978-1-4419-1524-5\\_1](https://doi.org/10.1007/978-1-4419-1524-5_1) (cit. on p. 3).