Constructions With Sets

The Clowder Project Authors

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- 000J This chapter develops some material relating to constructions with sets with an eye towards its categorical and higher-categorical counterparts to be introduced later in this work. In particular, it contains:
 - Explicit descriptions of the major types of co/limits in Sets, including in particular explicit descriptions of pushouts and coequalisers (see Definitions 2.4.1 and 2.5.1 and Remarks 2.4.3 and 2.5.3).
 - 2. A discussion of powersets as decategorifications of categories of presheaves (Remarks 4.1.2 and 4.3.2), including a (-1)-categorical analogue of un/straightening, described in Items 1 and 2 of Proposition 4.3.9 and Remark 4.3.11.
 - 3. A lengthy discussion of the adjoint triple

$$f_* \dashv f^{-1} \dashv f_! \colon \mathcal{P}(A) \xrightarrow{\rightleftarrows} \mathcal{P}(B)$$

of functors (morphisms of posets) between $\mathcal{P}(A)$ and $\mathcal{P}(B)$ induced by a map of sets $f: A \to B$, along with a discussion of the properties of f_* , f^{-1} , and $f_!$.

In line with the categorical viewpoint developed here, this adjoint triple may be described in terms of Kan extensions, and, as it turns out, it also shows up in some definitions and results in point-set topology, such as in e.g. notions of continuity for functions (??, ??).

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000K 1 Limits of Sets

000L 1.1 The Terminal Set

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DEFINITION 1.1.1 ► THE TERMINAL SET

The **terminal set** is the pair $(pt, \{!_A\}_{A \in Obj(Sets)})$ consisting of:

- · The Limit. The punctual set pt $\stackrel{\text{def}}{=} \{ \star \}$.
- · The Cone. The collection of maps

$$\{!_A : A \to \mathsf{pt}\}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

defined by

$$!_A(a) \stackrel{\text{def}}{=} \star$$

for each $a \in A$ and each $A \in Obj(Sets)$.

PROOF 1.1.2 ► PROOF OF DEFINITION 1.1.1

We claim that pt is the terminal object of Sets. Indeed, suppose we have a diagram of the form

$$A$$
 pt

in Sets. Then there exists a unique map $\phi \colon A \to \operatorname{pt}$ making the diagram

$$A \xrightarrow{\phi} \mathsf{pt}$$

commute, namely $!_A$.



000N 1.2 Products of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

000P

DEFINITION 1.2.1 ► THE PRODUCT OF A FAMILY OF SETS

The **product**¹ of $\{A_i\}_{i\in I}$ is the pair $(\prod_{i\in I} A_i, \{pr_i\}_{i\in I})$ consisting of:

· The Limit. The set $\prod_{i \in I} A_i$ defined by²

$$\prod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ f \in \operatorname{Sets} \left(I, \bigcup_{i \in I} A_i \right) \middle| \begin{array}{l} \text{for each } i \in I, \text{ we} \\ \text{have } f(i) \in A_i \end{array} \right\}.$$

· The Cone. The collection

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{pr}_{i}(f) \stackrel{\text{def}}{=} f(i)$$

for each $f \in \prod_{i \in I} A_i$ and each $i \in I$.

 2 Less formally, $\prod_{i\in I}A_i$ is the set whose elements are I-indexed collections $(a_i)_{i\in I}$ with $a_i\in A_i$ for each $i\in I$. The projection maps

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

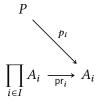
are then given by

$$\operatorname{pr}_i\left(\left(a_j\right)_{j\in I}\right)\stackrel{\text{def}}{=}a_i$$

for each $(a_j)_{j \in I} \in \prod_{i \in I} A_i$ and each $i \in I$.

PROOF 1.2.2 ► PROOF OF DEFINITION 1.2.1

We claim that $\prod_{i \in I} A_i$ is the categorical product of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi\colon P o \prod_{i \in I} A_i$ making the diagram

$$P$$

$$\phi \mid \exists ! \qquad p_i$$

$$\downarrow$$

$$\prod_{i \in I} A_i \xrightarrow{\operatorname{pr}_i} A_i$$

¹Further Terminology: Also called the **Cartesian product of** $\{A_i\}_{i\in I}$.

commute, being uniquely determined by the condition $\operatorname{pr}_i \circ \phi = p_i$ for each $i \in I$ via

$$\phi(x) = (p_i(x))_{i \in I}$$

for each $x \in P$.

000Q

PROPOSITION 1.2.3 ► PROPERTIES OF PRODUCTS OF FAMILIES OF SETS

Let $\{A_i\}_{i\in I}$ be a family of sets.

000R

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \prod_{i\in I}A_i$ defines a functor

$$\prod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

· Action on Objects. For each $(A_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}},\mathsf{Sets}))$, we have

$$\left[\prod_{i\in I}\right]((A_i)_{i\in I})\stackrel{\text{def}}{=}\prod_{i\in I}A_i$$

· Action on Morphisms. For each $(A_i)_{i \in I}, (B_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, the action on Hom-sets

$$\left(\prod_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}\colon \mathsf{Nat}((A_i)_{i\in I},(B_i)_{i\in I})\to \mathsf{Sets}\!\left(\prod_{i\in I}A_i,\prod_{i\in I}B_i\right)$$

of $\prod_{i\in I}$ at $((A_i)_{i\in I},(B_i)_{i\in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in $\operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I})$ to the map of sets

$$\prod_{i \in I} f_i \colon \prod_{i \in I} A_i \to \prod_{i \in I} B_i$$

defined by

$$\left[\prod_{i\in I} f_i\right] ((a_i)_{i\in I}) \stackrel{\text{def}}{=} (f_i(a_i))_{i\in I}$$

for each $(a_i)_{i \in I} \in \prod_{i \in I} A_i$.

PROOF 1.2.4 ► PROOF OF PROPOSITION 1.2.3

Item 1: Functoriality

This follows from ??, ?? of ??.

000S 1.3 Binary Products of Sets

Let A and B be sets.

000T DEFINITION 1.3.1 ► PRODUCTS OF SETS

The **product**¹ of A and B is the pair $(A \times B, \{pr_1, pr_2\})$ consisting of:

· The Limit. The set $A \times B$ defined by²

$$\begin{split} A \times B &\stackrel{\text{def}}{=} \prod_{z \in \{A,B\}} z \\ &\stackrel{\text{def}}{=} \{f \in \mathsf{Sets}(\{0,1\}, A \cup B) \mid \mathsf{we have} \ f(0) \in A \ \mathsf{and} \ f(1) \in B\} \\ &\cong \{\{\{a\}, \{a,b\}\} \in \mathcal{P}(\mathcal{P}(A \cup B)) \mid \mathsf{we have} \ a \in A \ \mathsf{and} \ b \in B\}. \end{split}$$

· The Cone. The maps

$$\operatorname{pr}_1 : A \times B \to A,$$

 $\operatorname{pr}_2 : A \times B \to B$

defined by

$$\operatorname{pr}_{1}(a,b) \stackrel{\text{def}}{=} a,$$
 $\operatorname{pr}_{2}(a,b) \stackrel{\text{def}}{=} b$

for each $(a, b) \in A \times B$.

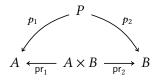
¹ Further Terminology: Also called the **Cartesian product of** A **and** B or the **binary Cartesian product of** A **and** B, for emphasis.

This can also be thought of as the $(\mathbb{E}_{-1}, \mathbb{E}_{-1})$ -tensor product of A and B.

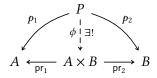
 $^{^2}$ In other words, $A \times B$ is the set whose elements are ordered pairs (a,b) with $a \in A$ and $b \in B$ as in Definition 3.4.1

PROOF 1.3.2 ▶ PROOF OF DEFINITION 1.3.1

We claim that $A \times B$ is the categorical product of A and B in Sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$.

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PROPOSITION 1.3.3 ► PROPERTIES OF PRODUCTS OF SETS

Let A, B, C, and X be sets.

000V

1. Functoriality. The assignments $A, B, (A, B) \mapsto A \times B$ define functors

$$A \times -:$$
 Sets \rightarrow Sets,
 $- \times B:$ Sets \rightarrow Sets,
 $-_1 \times -_2:$ Sets \times Sets \rightarrow Sets,

where -1×-2 is the functor where

· Action on Objects. For each $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$, we have

$$[-1 \times -2](A, B) \stackrel{\text{def}}{=} A \times B.$$

· Action on Morphisms. For each (A, B), $(X, Y) \in \mathsf{Obj}(\mathsf{Sets})$, the action on Hom-sets

$$\times_{(A,B),(X,Y)}$$
: $\mathsf{Sets}(A,X) \times \mathsf{Sets}(B,Y) \to \mathsf{Sets}(A \times B, X \times Y)$

of \times at ((A, B), (X, Y)) is defined by sending (f, g) to the function

$$f \times q : A \times B \to X \times Y$$

defined by

$$[f \times g](a,b) \stackrel{\text{def}}{=} (f(a),g(b))$$

for each $(a, b) \in A \times B$.

and where $A \times -$ and $- \times B$ are the partial functors of $-_1 \times -_2$ at $A, B \in Obj(Sets)$.

2. Adjointness. We have adjunctions

$$(A \times - + \mathsf{Hom}_{\mathsf{Sets}}(A, -)) \colon \mathsf{Sets} \underbrace{\bot}_{\mathsf{Hom}_{\mathsf{Sets}}(A, -)} \mathsf{Sets},$$

$$(- \times B + \mathsf{Hom}_{\mathsf{Sets}}(B, -)) \colon \mathsf{Sets} \underbrace{\bot}_{\mathsf{L}} \mathsf{Sets},$$

$$\mathsf{Hom}_{\mathsf{Sets}}(B, -)$$

witnessed by bijections

$$\operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A \times B, C) \cong \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A, \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(B, C)),$$

 $\operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A \times B, C) \cong \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(B, \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A, C)),$

natural in $A, B, C \in Obj(Sets)$.

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000X

3. Associativity. We have an isomorphism of sets

$$(A \times B) \times C \cong A \times (B \times C),$$

natural in $A, B, C \in Obj(Sets)$.

000Y

4. Unitality. We have isomorphisms of sets

$$pt \times A \cong A$$
,

$$A \times pt \cong A$$
,

natural in $A \in \mathsf{Obj}(\mathsf{Sets})$.

000Z

5. Commutativity. We have an isomorphism of sets

$$A \times B \cong B \times A$$
,

natural in $A, B \in Obj(Sets)$.

0010

6. Annihilation With the Empty Set. We have isomorphisms of sets

$$A \times \emptyset \cong \emptyset$$
,

$$\emptyset \times A \cong \emptyset$$
,

natural in $A \in Obj(Sets)$.

0011

7. Distributivity Over Unions. We have isomorphisms of sets

$$A \times (B \cup C) = (A \times B) \cup (A \times C),$$

$$(A \cup B) \times C = (A \times C) \cup (B \times C).$$

0012

8. Distributivity Over Intersections. We have isomorphisms of sets

$$A \times (B \cap C) = (A \times B) \cap (A \times C),$$

$$(A \cap B) \times C = (A \times C) \cap (B \times C).$$

0013

9. Middle-Four Exchange with Respect to Intersections. We have an isomorphism of sets

$$(A \times B) \cap (C \times D) \cong (A \cap B) \times (C \cap D).$$

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10. Distributivity Over Differences. We have isomorphisms of sets

$$A \times (B \setminus C) = (A \times B) \setminus (A \times C),$$

$$(A \setminus B) \times C = (A \times C) \setminus (B \times C),$$

natural in $A, B, C \in Obj(Sets)$.

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11. Distributivity Over Symmetric Differences. We have isomorphisms of sets

$$A \times (B \triangle C) = (A \times B) \triangle (A \times C),$$

$$(A \triangle B) \times C = (A \times C) \triangle (B \times C),$$

natural in $A, B, C \in \mathsf{Obj}(\mathsf{Sets})$.

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12. Symmetric Monoidality. The triple (Sets, \times , pt) is a symmetric monoidal category.

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13. Symmetric Bimonoidality. The quintuple (Sets, \coprod , \emptyset , \times , pt) is a symmetric bimonoidal category.

PROOF 1.3.4 ► PROOF OF PROPOSITION 1.3.3

Item 1: Functoriality

This follows from ??, ?? of ??.

Item 2: Adjointness

We prove only that there's an adjunction $- \times B \dashv \mathsf{Hom}_{\mathsf{Sets}}(B, -)$, witnessed by a bijection

$$Hom_{Sets}(A \times B, C) \cong Hom_{Sets}(A, Hom_{Sets}(B, C)),$$

natural in $B, C \in \mathsf{Obj}(\mathsf{Sets})$, as the proof of the existence of the adjunction $A \times - \dashv \mathsf{Hom}_{\mathsf{Sets}}(A, -)$ follows almost exactly in the same way.

· Map I. We define a map

$$\Phi_{B,C} \colon \mathsf{Hom}_{\mathsf{Sets}}(A \times B, C) \to \mathsf{Hom}_{\mathsf{Sets}}(A, \mathsf{Hom}_{\mathsf{Sets}}(B, C)),$$

by sending a function

$$\xi \colon A \times B \to C$$

to the function

$$\begin{split} \xi^{\dagger} \colon A \, &\to \, \mathsf{Hom}_{\mathsf{Sets}}(B,C), \\ a \, &\longmapsto \, \Big(\xi_a^{\dagger} \colon B \to C \Big), \end{split}$$

where we define

$$\xi_a^{\dagger}(b) \stackrel{\text{def}}{=} \xi(a,b)$$

for each $b \in B$. In terms of the $[\![a \mapsto f(a)]\!]$ notation of Sets, Notation 1.1.2, we have

$$\xi^{\dagger} \stackrel{\text{def}}{=} [\![a \mapsto [\![b \mapsto \xi(a, b)]\!]\!].$$

· Map II. We define a map

$$\Psi_{B,C}$$
: $\mathsf{Hom}_{\mathsf{Sets}}(A,\mathsf{Hom}_{\mathsf{Sets}}(B,C)), \to \mathsf{Hom}_{\mathsf{Sets}}(A \times B,C)$

given by sending a function

$$\xi \colon A \to \mathsf{Hom}_{\mathsf{Sets}}(B,C),$$

 $a \longmapsto (\xi_a \colon B \to C),$

to the function

$$\xi^{\dagger}: A \times B \to C$$

defined by

$$\xi^{\dagger}(a,b) \stackrel{\text{def}}{=} \operatorname{ev}_b(\operatorname{ev}_a(\xi))$$

$$\stackrel{\text{def}}{=} \operatorname{ev}_b(\xi_a)$$

$$\stackrel{\text{def}}{=} \xi_a(b)$$

for each $(a, b) \in A \times B$.

· Invertibility I. We claim that

$$\Psi_{A,B} \circ \Phi_{A,B} = \mathrm{id}_{\mathsf{Hom}_{\mathsf{Sets}}(A \times B,C)}.$$

Indeed, given a function $\xi \colon A \times B \to C$, we have

$$\begin{split} \left[\Psi_{A,B} \circ \Phi_{A,B} \right] (\xi) &= \Psi_{A,B} \big(\Phi_{A,B} (\xi) \big) \\ &= \Psi_{A,B} \big(\Phi_{A,B} \big(\left[(a,b) \mapsto \xi(a,b) \right] \big) \big) \\ &= \Psi_{A,B} \big(\left[a \mapsto \left[b \mapsto \xi(a,b) \right] \right] \big) \\ &= \Psi_{A,B} \big(\left[a' \mapsto \left[b' \mapsto \xi(a',b') \right] \right] \big) \\ &= \left[(a,b) \mapsto \operatorname{ev}_b \big(\left[a' \mapsto \left[b' \mapsto \xi(a',b') \right] \right] \big) \big) \right] \\ &= \left[(a,b) \mapsto \operatorname{ev}_b \big(\left[b' \mapsto \xi(a,b') \right] \big) \right] \\ &= \left[(a,b) \mapsto \xi(a,b) \right] \\ &= \xi. \end{split}$$

· Invertibility II. We claim that

$$\Phi_{A,B} \circ \Psi_{A,B} = \mathrm{id}_{\mathsf{HomSets}(A,\mathsf{HomSets}(B,C))}$$
.

Indeed, given a function

$$\xi \colon A \to \mathsf{Hom}_{\mathsf{Sets}}(B,C),$$

 $a \longmapsto (\xi_a \colon B \to C),$

we have

$$\begin{split} \left[\Phi_{A,B} \circ \Psi_{A,B} \right] (\xi) & \stackrel{\text{def}}{=} \Phi_{A,B} \big(\Psi_{A,B} (\xi) \big) \\ & \stackrel{\text{def}}{=} \Phi_{A,B} \big(\left[(a,b) \mapsto \xi_a(b) \right] \big) \\ & \stackrel{\text{def}}{=} \Phi_{A,B} \big(\left[(a',b') \mapsto \xi_{a'}(b') \right] \big) \\ & \stackrel{\text{def}}{=} \left[a \mapsto \left[b \mapsto \text{ev}_{(a,b)} \big(\left[(a',b') \mapsto \xi_{a'}(b') \right] \right) \right] \right] \\ & \stackrel{\text{def}}{=} \left[a \mapsto \left[b \mapsto \xi_a(b) \right] \right] \\ & \stackrel{\text{def}}{=} \left[a \mapsto \xi_a \right] \\ & \stackrel{\text{def}}{=} \xi. \end{split}$$

· Naturality for Φ , Part I. We need to show that, given a function $g \colon B \to B'$,

the diagram

commutes. Indeed, given a function

$$\xi: A \times B' \to C$$
,

we have

$$\begin{split} \big[\Phi_{B,C} \circ (\mathrm{id}_A \times g^*) \big] (\xi) &= \Phi_{B,C} ([\mathrm{id}_A \times g^*](\xi)) \\ &= \Phi_{B,C} (\xi(-_1, g(-_2))) \\ &= \big[\xi(-_1, g(-_2)) \big]^\dagger \\ &= \xi_{-_1}^\dagger (g(-_2)) \\ &= (g^*)_* \Big(\xi^\dagger \Big) \\ &= (g^*)_* \big(\Phi_{B',C} (\xi) \big) \\ &= \big[(g^*)_* \circ \Phi_{B',C} \big] (\xi). \end{split}$$

Alternatively, using the $[\![a\mapsto f(a)]\!]$ notation of Sets, Notation 1.1.2, we have

$$\begin{split} \left[\Phi_{B,C} \circ (\mathsf{id}_A \times g^*) \right] (\xi) &= \Phi_{B,C} (\left[\mathsf{id}_A \times g^* \right] (\xi)) \\ &= \Phi_{B,C} \left(\left[\mathsf{id}_A \times g^* \right] \left(\left[(a,b') \mapsto \xi(a,b') \right] \right) \right) \\ &= \Phi_{B,C} \left(\left[(a,b) \mapsto \xi(a,g(b)) \right] \right) \\ &= \left[a \mapsto \left[b \mapsto \xi(a,g(b)) \right] \right] \\ &= \left[a \mapsto g^* \left(\left[b' \mapsto \xi(a,b') \right] \right) \right] \\ &= (g^*)_* \left(\left[a \mapsto \left[b' \mapsto \xi(a,b') \right] \right] \right) \\ &= (g^*)_* \left(\Phi_{B',C} \left(\left[(a,b') \mapsto \xi(a,b') \right] \right) \right) \\ &= (g^*)_* \left(\Phi_{B',C} (\xi) \right) \\ &= \left[(g^*)_* \circ \Phi_{B',C} \right] (\xi). \end{split}$$

· Naturality for Φ , Part II. We need to show that, given a function $h\colon C\to C'$, the diagram

$$\begin{array}{cccc} \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A \times B, C) & \xrightarrow{\Phi_{B,C}} & \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A, \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(B, C)), \\ & & \downarrow & & \downarrow \\ h_* & & \downarrow & & \downarrow \\ \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A \times B, C') & \xrightarrow{\Phi_{B,C'}} & \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(A, \operatorname{\mathsf{Hom}}_{\mathsf{Sets}}(B, C')) \end{array}$$

commutes. Indeed, given a function

$$\xi: A \times B \to C$$

we have

$$\begin{split} \left[\Phi_{B,C} \circ h_* \right] (\xi) &= \Phi_{B,C} (h_*(\xi)) \\ &= \Phi_{B,C} \left(h_* ([[(a,b) \mapsto \xi(a,b)]] \right)) \\ &= \Phi_{B,C} ([[(a,b) \mapsto h(\xi(a,b))]]) \\ &= [[a \mapsto [[b \mapsto h(\xi(a,b))]]]) \\ &= [[a \mapsto h_* ([[b \mapsto \xi(a,b)]]])) \\ &= (h_*)_* ([[a \mapsto [[b \mapsto \xi(a,b)]]])) \\ &= (h_*)_* (\Phi_{B,C} ([[(a,b) \mapsto \xi(a,b)]])) \\ &= (h_*)_* (\Phi_{B,C} (\xi)) \\ &= [(h_*)_* \circ \Phi_{B,C}] (\xi). \end{split}$$

• Naturality for Ψ . Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Proposition 8.6.2 that Ψ is also natural in each argument.

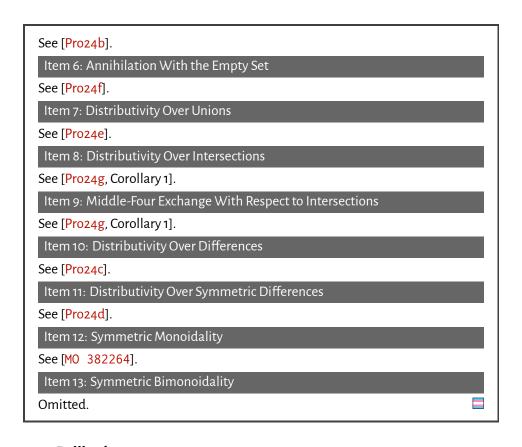
Item 3: Associativity

See [Pro24a].

Item 4: Unitality

Clear.

Item 5: Commutativity



0018 1.4 Pullbacks

Let A, B, and C be sets and let $f: A \to C$ and $g: B \to C$ be functions.

0019 DEFINITION 1.4.1 ► PULLBACKS OF SETS

The **pullback of** A **and** B **over** C **along** f **and** g^1 is the pair $(A \times_C B, \{\operatorname{pr}_1, \operatorname{pr}_2\})$ consisting of:

· The Limit. The set $A \times_C B$ defined by

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}.$$

· The Cone. The maps

$$\operatorname{pr}_1 : A \times_C B \to A,$$

 $\operatorname{pr}_2 : A \times_C B \to B$

defined by

$$\operatorname{pr}_{1}(a, b) \stackrel{\text{def}}{=} a,$$

 $\operatorname{pr}_{2}(a, b) \stackrel{\text{def}}{=} b$

for each $(a, b) \in A \times_C B$.

PROOF 1.4.2 ► PROOF OF DEFINITION 1.4.1

We claim that $A \times_C B$ is the categorical pullback of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pullback diagram commutes, i.e. that we have

$$f \circ \operatorname{pr}_{1} = g \circ \operatorname{pr}_{2}, \qquad A \times_{C} B \xrightarrow{\operatorname{pr}_{2}} B$$

$$\downarrow g$$

$$A \xrightarrow{f} C.$$

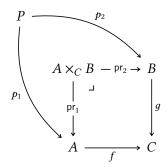
Indeed, given $(a, b) \in A \times_C B$, we have

$$\begin{split} [f \circ \mathsf{pr}_1](a,b) &= f(\mathsf{pr}_1(a,b)) \\ &= f(a) \\ &= g(b) \\ &= g(\mathsf{pr}_2(a,b)) \\ &= [g \circ \mathsf{pr}_2](a,b), \end{split}$$

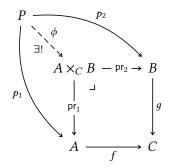
¹ Further Terminology: Also called the **fibre product of** A **and** B **over** C **along** f **and** g.

² Further Notation: Also written $A \times_{f,C,q} B$.

where f(a) = g(b) since $(a, b) \in A \times_C B$. Next, we prove that $A \times_C B$ satisfies the universal property of the pullback. Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi \colon P \to A \times_C B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$, where we note that $(p_1(x), p_2(x)) \in A \times B$ indeed lies in $A \times_C B$ by the condition

$$f\circ p_1=g\circ p_2,$$

which gives

$$f(p_1(x)) = g(p_2(x))$$

for each $x \in P$, so that $(p_1(x), p_2(x)) \in A \times_C B$.

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001A

EXAMPLE 1.4.3 ► **EXAMPLES OF PULLBACKS OF SETS**

Here are some examples of pullbacks of sets.

001B

1. Unions via Intersections. Let $A, B \subset X$. We have a bijection of sets

PROOF 1.4.4 ► PROOF OF EXAMPLE 1.4.3

Item 1: Unions via Intersections

Indeed, we have

$$A \times_{A \cup B} B \cong \{(x, y) \in A \times B \mid x = y\}$$

 $\cong A \cap B.$

This finishes the proof.

001C

PROPOSITION 1.4.5 ► PROPERTIES OF PULLBACKS OF SETS

Let A, B, C, and X be sets.

001D

1. Functoriality. The assignment $(A,B,C,f,g)\mapsto A\times_{f,C,g}B$ defines a functor

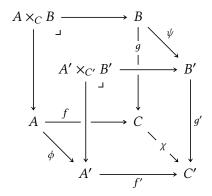
$$-_1 \times_{-_3} -_1 : \mathsf{Fun}(\mathcal{P}, \mathsf{Sets}) \to \mathsf{Sets},$$

where \mathcal{P} is the category that looks like this:



In particular, the action on morphisms of $-1 \times_{-3} -1$ is given by sending a

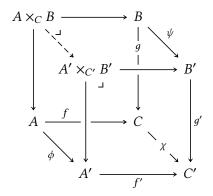
morphism



in Fun(\mathcal{P} , Sets) to the map $\xi \colon A \times_C B \xrightarrow{\exists !} A' \times_{C'} B'$ given by

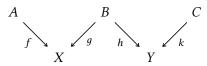
$$\xi(a,b) \stackrel{\text{def}}{=} (\phi(a), \psi(b))$$

for each $(a, b) \in A \times_C B$, which is the unique map making the diagram



commute.

2. Associativity. Given a diagram

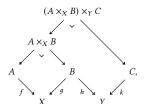


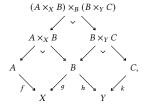
001E

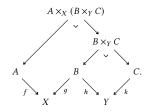
in Sets, we have isomorphisms of sets

$$(A \times_X B) \times_Y C \cong (A \times_X B) \times_B (B \times_Y C) \cong A \times_X (B \times_Y C),$$

where these pullbacks are built as in the diagrams







001F

3. Unitality. We have isomorphisms of sets



$$X \times_X A \cong A,$$

 $A \times_X X \cong A,$

$$\begin{array}{ccc} A & \xrightarrow{f} & X \\ \parallel & & \parallel \\ X & \xrightarrow{f} & X. \end{array}$$

001G

4. Commutativity. We have an isomorphism of sets

$$A \times_X B \cong B \times_X A$$

$$\begin{array}{ccc}
B \times_X A & \longrightarrow & A \\
\downarrow & & \downarrow f \\
B & \xrightarrow{g} & X.
\end{array}$$

001H

5. Annihilation With the Empty Set. We have isomorphisms of sets

$$\begin{array}{ccc}
\emptyset & \longrightarrow & \emptyset \\
\downarrow & \downarrow & \downarrow \\
A & \longrightarrow & X,
\end{array}$$

$$A \times_X \emptyset \cong \emptyset,$$

 $\emptyset \times_X A \cong \emptyset,$

$$\emptyset \longrightarrow \emptyset \qquad \qquad \emptyset \longrightarrow A$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad A \times_X \emptyset \cong \emptyset, \qquad \qquad \downarrow \qquad \downarrow f$$

$$A \longrightarrow_f X, \qquad \qquad \emptyset \longrightarrow X.$$

001J

6. Interaction With Products. We have an isomorphism of sets

$$A \times_{\mathsf{pt}} B \cong A \times B, \qquad A \times_{\mathsf{pt}} B \cong A \times B, \qquad A \xrightarrow{!_{A}} \mathsf{pt}.$$

001K

7. Symmetric Monoidality. The triple (Sets, \times_X , X) is a symmetric monoidal category.

PROOF 1.4.6 ► PROOF OF PROPOSITION 1.4.5

Item 1: Functoriality

This is a special case of functoriality of co/limits, ??, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pullback diagram.

Item 2: Associativity

Indeed, we have

```
(A \times_X B) \times_Y C \cong \{((a,b),c) \in (A \times_X B) \times C \mid h(b) = k(c)\}
\cong \{((a,b),c) \in (A \times B) \times C \mid f(a) = g(b) \text{ and } h(b) = k(c)\}
\cong \{(a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c)\}
\cong \{(a,(b,c)) \in A \times (B \times_Y C) \mid f(a) = g(b)\}
\cong A \times_X (B \times_Y C)
```

1.5 Equalisers 22

and

$$(A \times_X B) \times_B (B \times_Y C) \cong \left\{ ((a,b), (b',c)) \in (A \times_X B) \times (B \times_Y C) \mid b = b' \right\}$$

$$\cong \left\{ ((a,b), (b',c)) \in (A \times B) \times (B \times C) \mid f(a) = g(b), b = b', \\ \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a, (b, (b',c))) \in A \times (B \times (B \times C)) \mid f(a) = g(b), b = b', \\ \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a, ((b,b'),c)) \in A \times ((B \times B) \times C) \mid f(a) = g(b), b = b', \\ \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a, ((b,b'),c)) \in A \times ((B \times_B B) \times C) \mid f(a) = g(b) \text{ and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a, (b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c) \right\}$$

$$\cong A \times_X (B \times_Y C),$$

where we have used Item 3 for the isomorphism $B \times_B B \cong B$.

Item 3: Unitality

Indeed, we have

$$X \times_X A \cong \{(x, a) \in X \times A \mid f(a) = x\},\$$

$$A \times_X X \cong \{(a, x) \in X \times A \mid f(a) = x\},\$$

which are isomorphic to A via the maps $(x, a) \mapsto a$ and $(a, x) \mapsto a$.

Item 4: Commutativity

Clear.

Item 5: Annihilation With the Empty Set

Clear.

Item 6: Interaction With Products

Clear.

Item 7: Symmetric Monoidality

Omitted.



001L 1.5 Equalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

1.5 Equalisers

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001M

DEFINITION 1.5.1 ► EQUALISERS OF SETS

The **equaliser of** f **and** g is the pair (Eq(f,g),eq(f,g)) consisting of:

· The Limit. The set Eq(f, g) defined by

$$Eq(f,g) \stackrel{\text{def}}{=} \{ a \in A \, | \, f(a) = g(a) \}.$$

· The Cone. The inclusion map

$$eq(f,g): Eq(f,g) \hookrightarrow A.$$

PROOF 1.5.2 ► PROOF OF DEFINITION 1.5.1

We claim that ${\sf Eq}(f,g)$ is the categorical equaliser of f and g in Sets. First we need to check that the relevant equaliser diagram commutes, i.e. that we have

$$f \circ eq(f,g) = g \circ eq(f,g),$$

which indeed holds by the definition of the set ${\rm Eq}(f,g)$. Next, we prove that ${\rm Eq}(f,g)$ satisfies the universal property of the equaliser. Suppose we have a diagram of the form

$$\mathsf{Eq}(f,g) \xrightarrow{\mathsf{eq}(f,g)} A \xrightarrow{f} B$$

in Sets. Then there exists a unique map $\phi \colon E \to \operatorname{Eq}(f,g)$ making the diagram

$$\mathsf{Eq}(f,g) \xrightarrow{\mathsf{eq}(f,g)} A \xrightarrow{f} B$$

$$\downarrow \downarrow \downarrow \downarrow e$$

$$E$$

commute, being uniquely determined by the condition

$$\operatorname{eq}(f,g)\circ\phi=e$$

via

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in $\mathrm{Eq}(f,g)$ by the condition

$$f \circ e = g \circ e$$
,

which gives

$$f(e(x)) = g(e(x))$$

for each $x \in E$, so that $e(x) \in Eq(f, g)$.

001N

PROPOSITION 1.5.3 ► PROPERTIES OF EQUALISERS OF SETS

Let A, B, and C be sets.

001P

1. Associativity. We have isomorphisms of sets¹

$$\underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(g,h), g \circ \mathrm{eq}(g,h))}_{=\mathrm{Eq}(f \circ \mathrm{eq}(g,h), h \circ \mathrm{eq}(g,h))} \cong \underbrace{\mathrm{Eq}(f,g,h)}_{=\mathrm{Eq}(g \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))} \cong \underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}_{=\mathrm{Eq}(g \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}$$

where Eq(f, g, h) is the limit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets, being explicitly given by

$$Eq(f, q, h) \cong \{a \in A \mid f(a) = q(a) = h(a)\}.$$

001Q

2. Unitality. We have an isomorphism of sets

$$\operatorname{Eq}(f, f) \cong A$$
.

001R

3. Commutativity. We have an isomorphism of sets

$$\operatorname{Eq}(f,g) \cong \operatorname{Eq}(g,f)$$
.

001S

4. Interaction With Composition. Let

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have an inclusion of sets

$$\mathsf{Eq}(h \circ f \circ \mathsf{eq}(f,g), k \circ g \circ \mathsf{eq}(f,g)) \subset \mathsf{Eq}(h \circ f, k \circ g),$$

where Eq $(h \circ f \circ \text{eq}(f,g), k \circ g \circ \text{eq}(f,g))$ is the equaliser of the composition

$$\operatorname{Eq}(f,g) \overset{\operatorname{eq}(f,g)}{\hookrightarrow} A \overset{f}{\underset{q}{\Longrightarrow}} B \overset{h}{\underset{k}{\Longrightarrow}} C.$$

¹That is, the following three ways of forming "the" equaliser of (f, g, h) agree:

(a) Take the equaliser of (f, g, h), i.e. the limit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

(b) First take the equaliser of f and g, forming a diagram

$$\operatorname{Eq}(f,g) \overset{\operatorname{eq}(f,g)}{\hookrightarrow} A \overset{f}{\underset{g}{\rightrightarrows}} B$$

and then take the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{h}{\Longrightarrow}} B,$$

obtaining a subset

$$\mathsf{Eq}(f \circ \mathsf{eq}(f,g), h \circ \mathsf{eq}(f,g)) = \mathsf{Eq}(g \circ \mathsf{eq}(f,g), h \circ \mathsf{eq}(f,g))$$

of Eq(f, g).

(c) First take the equaliser of g and h, forming a diagram

$$\mathsf{Eq}(g,h) \overset{\mathsf{eq}(g,h)}{\hookrightarrow} A \overset{g}{\underset{h}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\mathsf{Eq}(g,h) \overset{\mathsf{eq}(g,h)}{\hookrightarrow} A \overset{f}{\underset{q}{\Rightarrow}} B,$$

obtaining a subset

$$\operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h)) = \operatorname{Eq}(f \circ \operatorname{eq}(g,h), h \circ \operatorname{eq}(g,h))$$

of Eq(g, h).

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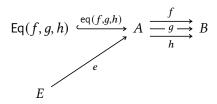
PROOF 1.5.4 ▶ PROOF OF PROPOSITION 1.5.3

Item 1: Associativity

We first prove that Eq(f, q, h) is indeed given by

$$Eq(f, q, h) \cong \{a \in A \mid f(a) = q(a) = h(a)\}.$$

Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi\colon E\to \operatorname{Eq}(f,g,h)$, uniquely determined by the condition

$$\operatorname{eq}(f,g)\circ\phi=e$$

being necessarily given by

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in Eq(f, g, h) by the condition

$$f \circ e = g \circ e = h \circ e$$
,

which gives

$$f(e(x)) = g(e(x)) = h(e(x))$$

for each $x \in E$, so that $e(x) \in Eq(f, g, h)$.

We now check the equalities

$$\mathsf{Eq}(f \circ \mathsf{eq}(g,h), g \circ \mathsf{eq}(g,h)) \cong \mathsf{Eq}(f,g,h) \cong \mathsf{Eq}(f \circ \mathsf{eq}(f,g), h \circ \mathsf{eq}(f,g)).$$

Indeed, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h)) &\cong \{x \in \operatorname{Eq}(g,h) \,|\, [f \circ \operatorname{eq}(g,h)](a) = [g \circ \operatorname{eq}(g,h)](a) \} \\ &\cong \{x \in \operatorname{Eq}(g,h) \,|\, f(a) = g(a) \} \\ &\cong \{x \in A \,|\, f(a) = g(a) \text{ and } g(a) = h(a) \} \\ &\cong \{x \in A \,|\, f(a) = g(a) = h(a) \} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Similarly, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g)) &\cong \{x \in \operatorname{Eq}(f,g) \,|\, [f \circ \operatorname{eq}(f,g)](a) = [h \circ \operatorname{eq}(f,g)](a)\} \\ &\cong \{x \in \operatorname{Eq}(f,g) \,|\, f(a) = h(a)\} \\ &\cong \{x \in A \,|\, f(a) = h(a) \text{ and } f(a) = g(a)\} \\ &\cong \{x \in A \,|\, f(a) = g(a) = h(a)\} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Item 2: Unitality

Clear.

Item 3: Commutativity

Clear.

Item 4: Interaction With Composition

Indeed, we have

$$\begin{split} \operatorname{Eq}(h \circ f \circ \operatorname{eq}(f,g), k \circ g \circ \operatorname{eq}(f,g)) & \cong \{a \in \operatorname{Eq}(f,g) \,|\, h(f(a)) = k(g(a))\} \\ & \cong \{a \in A \,|\, f(a) = g(a) \text{ and } h(f(a)) = k(g(a))\}. \end{split}$$

and

$$\mathsf{Eq}(h \circ f, k \circ g) \cong \{a \in A \,|\, h(f(a)) = k(g(a))\},\$$

and thus there's an inclusion from ${\rm Eq}(h\circ f\circ {\rm eq}(f,g),k\circ g\circ {\rm eq}(f,g))$ to ${\rm Eq}(h\circ f,k\circ g).$

001T 2 Colimits of Sets

001U 2.1 The Initial Set

001V

DEFINITION 2.1.1 ► THE INITIAL SET

The **initial set** is the pair $(\emptyset, \{\iota_A\}_{A \in \mathsf{Obj}(\mathsf{Sets})})$ consisting of:

- The Limit. The empty set Ø of Definition 3.1.1.
- · The Cone. The collection of maps

$$\{\iota_A \colon \emptyset \to A\}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

given by the inclusion maps from \emptyset to A.

PROOF 2.1.2 ▶ PROOF OF DEFINITION 2.1.1

We claim that \emptyset is the initial object of Sets. Indeed, suppose we have a diagram of the form

in Sets. Then there exists a unique map $\phi \colon \emptyset \to A$ making the diagram

$$\emptyset - \frac{\phi}{\exists !} \rightarrow A$$

commute, namely the inclusion map ι_A .

001W 2.2 Coproducts of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

001X DEFINITION 2.2.1 ➤ DISJOINT UNIONS OF FAMILIES

The **disjoint union of the family** $\{A_i\}_{i\in I}$ is the pair $(\coprod_{i\in I} A_i, \{\operatorname{inj}_i\}_{i\in I})$ consisting of:

· The Colimit. The set $\coprod_{i\in I} A_i$ defined by

$$\left[\prod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ (i, x) \in I \times \left(\bigcup_{i \in I} A_i \right) \middle| x \in A_i \right\}.$$

· The Cocone. The collection

$$\left\{\operatorname{inj}_i\colon A_i\to \coprod_{i\in I}A_i\right\}_{i\in I}$$

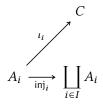
of maps given by

$$\operatorname{inj}_i(x) \stackrel{\text{def}}{=} (i, x)$$

for each $x \in A_i$ and each $i \in I$.

PROOF 2.2.2 ▶ PROOF OF DEFINITION 2.2.1

We claim that $\coprod_{i\in I} A_i$ is the categorical coproduct of $\{A_i\}_{i\in I}$ in Sets. Indeed, suppose we have, for each $i\in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi\colon\coprod_{i\in I}A_i o C$ making the diagram

$$A_{i} \xrightarrow{\inf_{i}} \coprod_{i \in I} A$$

commute, being uniquely determined by the condition $\phi \circ \operatorname{inj}_i = \iota_i$ for each $i \in I$ via

$$\phi((i,x)) = \iota_i(x)$$

for each $(i, x) \in \coprod_{i \in I} A_i$.

001Y

PROPOSITION 2.2.3 ► PROPERTIES OF COPRODUCTS OF FAMILIES OF SETS

Let $\{A_i\}_{i\in I}$ be a family of sets.

001Z

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \coprod_{i\in I}A_i$ defines a functor

$$\coprod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

 $\cdot \ \mathit{Action on Objects}. \ \mathsf{For each} \ (A_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets})), \mathsf{we have}$

$$\left[\bigsqcup_{i \in I} \right] ((A_i)_{i \in I}) \stackrel{\text{def}}{=} \bigsqcup_{i \in I} A_i$$

· Action on Morphisms. For each $(A_i)_{i \in I}, (B_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, the action on Hom-sets

$$\left(\coprod_{i \in I} \right)_{(A_i)_{i \in I}, (B_i)_{i \in I}} \colon \mathsf{Nat}((A_i)_{i \in I}, (B_i)_{i \in I}) \to \mathsf{Sets}\!\left(\coprod_{i \in I} A_i, \coprod_{i \in I} B_i \right)$$

of $\coprod_{i \in I}$ at $((A_i)_{i \in I}, (B_i)_{i \in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i \in I}, (B_i)_{i \in I})$ to the map of sets

$$\coprod_{i \in I} f_i \colon \coprod_{i \in I} A_i \to \coprod_{i \in I} B_i$$

defined by

$$\left[\bigsqcup_{i \in I} f_i \right] (i, a) \stackrel{\text{def}}{=} f_i(a)$$

for each $(i, a) \in \coprod_{i \in I} A_i$.

PROOF 2.2.4 ► PROOF OF PROPOSITION 2.2.3

Item 1: Functoriality

This follows from ??, ?? of ??.

0020 2.3 Binary Coproducts

Let A and B be sets.

0021 DEFINITION 2.3.1 ➤ COPRODUCTS OF SETS

The **coproduct**¹ **of** A **and** B is the pair $(A \coprod B, \{inj_1, inj_2\})$ consisting of:

· The Colimit. The set $A \coprod B$ defined by

$$A \coprod B \stackrel{\text{def}}{=} \coprod_{z \in \{A, B\}} z$$
$$\cong \{(0, a) \mid a \in A\} \cup \{(1, b) \mid b \in B\}.$$

· The Cocone. The maps

$$\begin{aligned} &\inf_1\colon A\to A\coprod B,\\ &\inf_2\colon B\to A\coprod B,\end{aligned}$$

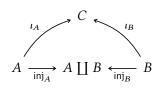
given by

$$\operatorname{inj}_{1}(a) \stackrel{\text{def}}{=} (0, a),$$
 $\operatorname{inj}_{2}(b) \stackrel{\text{def}}{=} (1, b),$

for each $a \in A$ and each $b \in B$.

PROOF 2.3.2 ► PROOF OF DEFINITION 2.3.1

We claim that $A \coprod B$ is the categorical coproduct of A and B in Sets. Indeed, suppose we have a diagram of the form



¹ Further Terminology: Also called the **disjoint union of** A **and** B, or the **binary disjoint union of** A **and** B, for emphasis.

in Sets. Then there exists a unique map $\phi\colon A\coprod B\to C$ making the diagram

$$A \xrightarrow[\text{inj}_{A}]{C} \downarrow_{l_{B}} \downarrow_{l_{B}}$$

$$A \xrightarrow[\text{inj}_{A}]{A} A \coprod B \xrightarrow[\text{inj}_{B}]{B}$$

commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_A = \iota_A,$$

 $\phi \circ \operatorname{inj}_B = \iota_B$

via

0023

$$\phi(x) = \begin{cases} \iota_A(a) & \text{if } x = (0, a), \\ \iota_B(b) & \text{if } x = (1, b) \end{cases}$$

for each $x \in A \coprod B$.

0022 PROPOSITION 2.3.3 ► PROPERTIES OF COPRODUCTS OF SETS

Let A, B, C, and X be sets.

1. Functoriality. The assignment $A, B, (A, B) \mapsto A \coprod B$ defines functors

$$A \coprod -: \mathsf{Sets} \to \mathsf{Sets},$$

 $- \coprod B: \mathsf{Sets} \to \mathsf{Sets},$
 $-_1 \coprod -_2: \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets},$

where $-_1 \coprod -_2$ is the functor where

· Action on Objects. For each $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$, we have

$$[-_1 \coprod -_2](A, B) \stackrel{\text{def}}{=} A \coprod B.$$

· Action on Morphisms. For each $(A, B), (X, Y) \in \mathsf{Obj}(\mathsf{Sets})$, the action on Hom-sets

$$\coprod_{(A,B),(X,Y)} : \mathsf{Sets}(A,X) \times \mathsf{Sets}(B,Y) \to \mathsf{Sets}(A \coprod B,X \coprod Y)$$

of \coprod at ((A, B), (X, Y)) is defined by sending (f, g) to the function

$$f \coprod g: A \coprod B \to X \coprod Y$$

defined by

$$[f \coprod g](x) \stackrel{\text{def}}{=} \begin{cases} (0, f(a)) & \text{if } x = (0, a), \\ (1, g(b)) & \text{if } x = (1, b), \end{cases}$$

for each $x \in A \coprod B$.

and where $A \coprod -$ and $- \coprod B$ are the partial functors of $-_1 \coprod -_2$ at $A, B \in$ Obj(Sets).

2. Associativity. We have an isomorphism of sets

$$(A \coprod B) \coprod C \cong A \coprod (B \coprod C),$$

natural in $A, B, C \in Obj(Sets)$.

3. Unitality. We have isomorphisms of sets

$$A \coprod \emptyset \cong A,$$

 $\emptyset \coprod A \cong A,$

natural in $A \in Obj(Sets)$.

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4. Commutativity. We have an isomorphism of sets

$$A \mid A \mid A \cong B \mid A$$

natural in $A, B \in Obj(Sets)$.

5. Symmetric Monoidality. The triple (Sets, \coprod , \emptyset) is a symmetric monoidal category.

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PROOF 2.3.4 ► PROOF OF PROPOSITION 2.3.3
Item 1: Functoriality
This follows from ??, ?? of ??.
Item 2: Associativity
Clear.
Item 3: Unitality
Clear.
Item 4: Commutativity
Clear.
Item 5: Symmetric Monoidality
Omitted.

0028 2.4 Pushouts

Let A, B, and C be sets and let $f: C \to A$ and $g: C \to B$ be functions.

0029 DEFINITION 2.4.1 ➤ PUSHOUTS OF SETS

The **pushout of** A **and** B **over** C **along** f **and** g^1 is the pair A^2 $A = A^2$ $A = A^2$ ($A = A^2$ $A = A^2$

· The Colimit. The set $A \coprod_C B$ defined by

$$A \coprod_{C} B \stackrel{\text{def}}{=} A \coprod_{C} B/\sim_{C},$$

where \sim_C is the equivalence relation on $A \coprod B$ generated by $(0, f(c)) \sim_C (1, g(c))$.

· The Cocone. The maps

$$\operatorname{inj}_1 : A \to A \coprod_C B,$$

 $\operatorname{inj}_2 : B \to A \coprod_C B$

given by

$$inj_1(a) \stackrel{\text{def}}{=} [(0, a)]$$

$$\mathsf{inj}_2(b) \stackrel{\mathsf{def}}{=} [(1,b)]$$

2.4 Pushouts 35

for each $a \in A$ and each $b \in B$.

¹ Further Terminology: Also called the **fibre coproduct of** A **and** B **over** C **along** f **and** g.

PROOF 2.4.2 ► PROOF OF DEFINITION 2.4.1

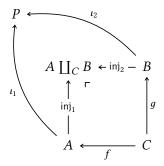
We claim that $A \coprod_C B$ is the categorical pushout of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pushout diagram commutes, i.e. that we have

$$\begin{aligned} & A \coprod_C B \xleftarrow{\operatorname{inj}_2} & B \\ & \operatorname{inj}_1 \circ f = \operatorname{inj}_2 \circ g, & & \operatorname{inj}_1 \\ & A \xleftarrow{f} & C. \end{aligned}$$

Indeed, given $c \in C$, we have

$$\begin{split} [\inf_1 \circ f](c) &= \inf_1(f(c)) \\ &= [(0, f(c))] \\ &= [(1, g(c))] \\ &= \inf_2(g(c)) \\ &= [\inf_2 \circ g](c), \end{split}$$

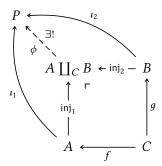
where [(0,f(c))] = [(1,g(c))] by the definition of the relation \sim on $A \coprod B$. Next, we prove that $A \coprod CB$ satisfies the universal property of the pushout. Suppose we have a diagram of the form



² Further Notation: Also written $A \coprod_{f,C,q} B$.

2.4 Pushouts 36

in Sets. Then there exists a unique map $\phi: A \coprod_C B \to P$ making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_1 = \iota_1,$$

 $\phi \circ \operatorname{inj}_2 = \iota_2$

via

$$\phi(x) = \begin{cases} \iota_1(a) & \text{if } x = [(0, a)], \\ \iota_2(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, where the well-definedness of ϕ is guaranteed by the equality $\iota_1 \circ f = \iota_2 \circ g$ and the definition of the relation \sim on $A \coprod B$ as follows:

1. Case 1: Suppose we have x = [(0, a)] = [(0, a')] for some $a, a' \in A$. Then, by Remark 2.4.3, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a').$$

2. Case 2: Suppose we have x = [(1, b)] = [(1, b')] for some $b, b' \in B$. Then, by Remark 2.4.3, we have a sequence

$$(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b').$$

3. Case 3: Suppose we have x = [(0, a)] = [(1, b)] for some $a \in A$ and $b \in B$. Then, by Remark 2.4.3, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b).$$

2.4 Pushouts 37

In all these cases, we declare $x \sim' y$ iff there exists some $c \in C$ such that x = (0, f(c)) and y = (1, g(c)) or x = (1, g(c)) and y = (0, f(c)). Then, the equality $\iota_1 \circ f = \iota_2 \circ g$ gives

$$\begin{split} \phi([x]) &= \phi([(0, f(c))]) \\ &\stackrel{\text{def}}{=} \iota_1(f(c)) \\ &= \iota_2(g(c)) \\ &\stackrel{\text{def}}{=} \phi([(1, g(c))]) \\ &= \phi([y]), \end{split}$$

with the case where x=(1,g(c)) and y=(0,f(c)) similarly giving $\phi([x])=\phi([y])$. Thus, if $x\sim' y$, then $\phi([x])=\phi([y])$. Applying this equality pairwise to the sequences

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a'),$$

 $(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b'),$
 $(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b)$

gives

$$\phi([(0, a)]) = \phi([(0, a')]),$$

$$\phi([(1, b)]) = \phi([(1, b')]),$$

$$\phi([(0, a)]) = \phi([(1, b)]),$$

showing ϕ to be well-defined.

002A REMARK 2.4.3 ► UNWINDING DEFINITION 2.4.1

In detail, by Equivalence Relations and Apartness Relations, Construction 4.2.2, the relation \sim of Definition 2.4.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

- · We have $a, b \in A$ and a = b;
- · We have $a, b \in B$ and a = b;
- · There exist $x_1, \ldots, x_n \in A \coprod B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where

2.4 Pushouts

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we declare $x \sim' y$ if one of the following conditions is satisfied:

1. There exists $c \in C$ such that x = (0, f(c)) and y = (1, q(c)).

2. There exists $c \in C$ such that x = (1, q(c)) and y = (0, f(c)).

That is: we require the following condition to be satisfied:

(★) There exist $x_1, ..., x_n \in A \coprod B$ satisfying the following conditions:

- 1. There exists $c_0 \in C$ satisfying one of the following conditions:
 - (a) We have $a = f(c_0)$ and $x_1 = g(c_0)$.
 - (b) We have $a = q(c_0)$ and $x_1 = f(c_0)$.
- 2. For each $1 \le i \le n-1$, there exists $c_i \in C$ satisfying one of the following conditions:
 - (a) We have $x_i = f(c_i)$ and $x_{i+1} = g(c_i)$.
 - (b) We have $x_i = g(c_i)$ and $x_{i+1} = f(c_i)$.
- 3. There exists $c_n \in C$ satisfying one of the following conditions:
 - (a) We have $x_n = f(c_n)$ and $b = g(c_n)$.
 - (b) We have $x_n = g(c_n)$ and $b = f(c_n)$.

002B EXAMPLE 2.4.4 ► EXAMPLES OF PUSHOUTS OF SETS

002C

002D

Here are some examples of pushouts of sets.

1. Wedge Sums of Pointed Sets. The wedge sum of two pointed sets of Pointed Sets, Definition 3.3.1 is an example of a pushout of sets.

2. Intersections via Unions. Let $A, B \subset X$. We have a bijection of sets

$$A \cup B \cong A \coprod_{A \cap B} B, \qquad A \longleftarrow B$$

$$A \longleftarrow A \cap B$$

Pushouts 2.4 39

PROOF 2.4.5 ▶ PROOF OF EXAMPLE 2.4.4

Item 1: Wedge Sums of Pointed Sets

Follows by definition.

Item 2: Intersections via Unions

Indeed, $A \coprod_{A \cap B} B$ is the quotient of $A \coprod B$ by the equivalence relation obtained by declaring $(0, a) \sim (1, b)$ iff $a = b \in A \cap B$, which is in bijection with $A \cup B$ via the map with $[(0, a)] \mapsto a$ and $[(1, b)] \mapsto b$.

002E PROPOSITION 2.4.6 ► PROPERTIES OF PUSHOUTS OF SETS

Let A, B, C, and X be sets.

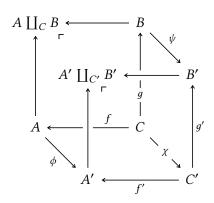
1. Functoriality. The assignment $(A,B,C,f,g)\mapsto A\coprod_{f,C,g} B$ defines a functor

$$-_1 \coprod_{-_3} -_1 : \operatorname{\mathsf{Fun}}(\mathcal{P},\operatorname{\mathsf{Sets}}) \to \operatorname{\mathsf{Sets}},$$

where \mathcal{P} is the category that looks like this:



In particular, the action on morphisms of $-_1 \coprod_{-_3} -_1$ is given by sending a morphism



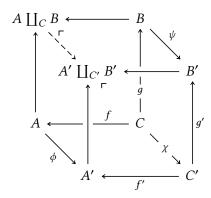
002F

2.4 Pushouts 40

in Fun(\mathcal{P} , Sets) to the map $\xi \colon A \coprod_C B \xrightarrow{\exists !} A' \coprod_{C'} B'$ given by

$$\xi(x) \stackrel{\text{def}}{=} \begin{cases} \phi(a) & \text{if } x = [(0, a)], \\ \psi(b) & \text{if } x = [(1, b)] \end{cases}$$

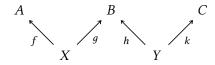
for each $x \in A \coprod_C B$, which is the unique map making the diagram



commute.

002G

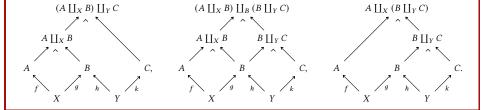
2. Associativity. Given a diagram



in Sets, we have isomorphisms of sets

$$(A \coprod_X B) \coprod_Y C \cong (A \coprod_X B) \coprod_B (B \coprod_Y C) \cong A \coprod_X (B \coprod_Y C),$$

where these pullbacks are built as in the diagrams



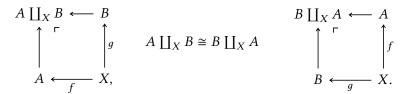
002H

3. Unitality. We have isomorphisms of sets



002J

4. Commutativity. We have an isomorphism of sets



002K

5. Interaction With Coproducts. We have

$$A \coprod_{\emptyset} B \cong A \coprod_{B}, \qquad A \coprod_{\iota_{A}} B \longleftarrow B$$

$$A \coprod_{\iota_{A}} B \longleftarrow_{\iota_{A}} B$$

$$A \longleftarrow_{\iota_{A}} \emptyset.$$

002L

6. Symmetric Monoidality. The triple (Sets, \coprod_X , X) is a symmetric monoidal category.

PROOF 2.4.7 ▶ PROOF OF PROPOSITION 2.4.6

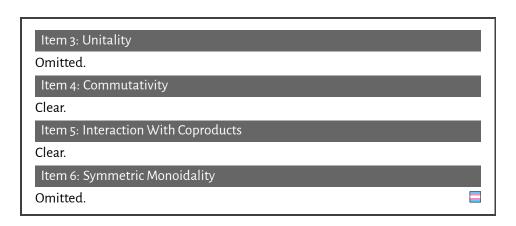
Item 1: Functoriality

This is a special case of functoriality of co/limits, ??, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pushout diagram.

Item 2: Associativity

Omitted.

2.5 Coequalisers 42



002M 2.5 Coequalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

002N DEFINITION 2.5.1 ➤ COEQUALISERS OF SETS

The **coequaliser of** f **and** g is the pair (CoEq(f,g), coeq(f,g)) consisting of:

· The Colimit. The set CoEq(f, g) defined by

$$CoEq(f,g) \stackrel{\text{def}}{=} B/\sim$$
,

where \sim is the equivalence relation on B generated by $f(a) \sim g(a)$.

· The Cocone. The map

$$coeq(f,g): B \to CoEq(f,g)$$

given by the quotient map $\pi \colon B \twoheadrightarrow B/\sim$ with respect to the equivalence relation generated by $f(a) \sim g(a)$.

PROOF 2.5.2 ▶ PROOF OF DEFINITION 2.5.1

We claim that $\operatorname{CoEq}(f,g)$ is the categorical coequaliser of f and g in Sets. First we need to check that the relevant coequaliser diagram commutes, i.e. that we have

$$\operatorname{coeq}(f,g)\circ f=\operatorname{coeq}(f,g)\circ g.$$

2.5 Coequalisers

Indeed, we have

$$[\operatorname{coeq}(f,g) \circ f](a) \stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](f(a))$$

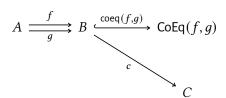
$$\stackrel{\text{def}}{=} [f(a)]$$

$$= [g(a)]$$

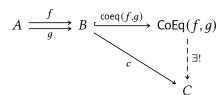
$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](g(a))$$

$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g) \circ g](a)$$

for each $a \in A$. Next, we prove that CoEq(f,g) satisfies the universal property of the coequaliser. Suppose we have a diagram of the form



in Sets. Then, since c(f(a)) = c(g(a)) for each $a \in A$, it follows from Equivalence Relations and Apartness Relations, Items 4 and 5 of Proposition 5.2.3 that there exists a unique map $CoEq(f,g) \xrightarrow{\exists !} C$ making the diagram



commute.

002P REMARK 2.5.3 ► UNWINDING DEFINITION 2.5.1

In detail, by Equivalence Relations and Apartness Relations, Construction 4.2.2, the relation \sim of Definition 2.5.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

· We have a = b;

- There exist $x_1, \ldots, x_n \in B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - 1. There exists $z \in A$ such that x = f(z) and y = g(z).
 - 2. There exists $z \in A$ such that x = q(z) and y = f(z).

That is: we require the following condition to be satisfied:

- (★) There exist $x_1, ..., x_n ∈ B$ satisfying the following conditions:
 - 1. There exists $z_0 \in A$ satisfying one of the following conditions:
 - (a) We have $a = f(z_0)$ and $x_1 = g(z_0)$.
 - (b) We have $a = g(z_0)$ and $x_1 = f(z_0)$.
 - 2. For each $1 \le i \le n-1$, there exists $z_i \in A$ satisfying one of the following conditions:
 - (a) We have $x_i = f(z_i)$ and $x_{i+1} = g(z_i)$.
 - (b) We have $x_i = g(z_i)$ and $x_{i+1} = f(z_i)$.
 - 3. There exists $z_n \in A$ satisfying one of the following conditions:
 - (a) We have $x_n = f(z_n)$ and $b = g(z_n)$.
 - (b) We have $x_n = g(z_n)$ and $b = f(z_n)$.

0020 EXAMPLE 2.5.4 ► EXAMPLES OF COEQUALISERS OF SETS

Here are some examples of coequalisers of sets.

1. Quotients by Equivalence Relations. Let R be an equivalence relation on a set X. We have a bijection of sets

$$X/\sim_R \cong \mathsf{CoEq}\bigg(R \hookrightarrow X \times X \overset{\mathsf{pr}_1}{\underset{\mathsf{pr}_2}{\Longrightarrow}} X\bigg).$$

PROOF 2.5.5 ▶ PROOF OF EXAMPLE 2.5.4

Item 1: Quotients by Equivalence Relations

See [Pro24z].

002R

002S PROPOSITION 2.5.6 ➤ PROPERTIES OF COEQUALISERS OF SETS

Let A, B, and C be sets.

1. Associativity. We have isomorphisms of sets¹

$$\underbrace{\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ f, \mathsf{coeq}(f,g) \circ h)}_{=\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ g, \mathsf{coeq}(f,g) \circ h)} \cong \mathsf{CoEq}(f,g,h) \cong \underbrace{\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g, \mathsf{coeq}(g,h) \circ h)}_{=\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h)}$$

where CoEq(f, g, h) is the colimit of the diagram

$$A \xrightarrow{f \atop g \xrightarrow{h}} B$$

in Sets.

2. Unitality. We have an isomorphism of sets

$$CoEq(f, f) \cong B$$
.

3. Commutativity. We have an isomorphism of sets

$$CoEq(f, q) \cong CoEq(q, f)$$
.

4. Interaction With Composition. Let

$$A \underset{g}{\overset{f}{\Longrightarrow}} B \underset{k}{\overset{h}{\Longrightarrow}} C$$

be functions. We have a surjection

$$\mathsf{CoEq}(h \circ f, k \circ g) \twoheadrightarrow \mathsf{CoEq}(\mathsf{coeq}(h, k) \circ h \circ f, \mathsf{coeq}(h, k) \circ k \circ g)$$

exhibiting CoEq(coeq(h,k) \circ h \circ f, coeq(h,k) \circ k \circ g) as a quotient of CoEq(h \circ f, k \circ g) by the relation generated by declaring $h(y) \sim k(y)$ for each $y \in B$.

002U

002T

002V

002W

¹That is, the following three ways of forming "the" coequaliser of (f, g, h) agree:

(a) Take the coequaliser of (f,g,h), i.e. the colimit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

(b) First take the coequaliser of f and g, forming a diagram

$$A \underset{q}{\overset{f}{\Rightarrow}} B \overset{\mathsf{coeq}(f,g)}{\twoheadrightarrow} \mathsf{CoEq}(f,g)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\Longrightarrow} B \stackrel{\mathsf{coeq}(f,g)}{\twoheadrightarrow} \mathsf{CoEq}(f,g),$$

obtaining a quotient

 $\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ f, \mathsf{coeq}(f,g) \circ h) = \mathsf{CoEq}(\mathsf{coeq}(f,g) \circ g, \mathsf{coeq}(f,g) \circ h)$

of CoEq(f, g)

(c) First take the coequaliser of g and h, forming a diagram

$$A \stackrel{g}{\underset{h}{\Longrightarrow}} B \stackrel{\mathsf{coeq}(g,h)}{\twoheadrightarrow} \mathsf{CoEq}(g,h)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{\mathsf{coeq}(g,h)}{\twoheadrightarrow} \mathsf{CoEq}(g,h),$$

obtaining a quotient

 ${\sf CoEq}({\sf coeq}(g,h)\circ f, {\sf coeq}(g,h)\circ g) = {\sf CoEq}({\sf coeq}(g,h)\circ f, {\sf coeq}(g,h)\circ h)$ of ${\sf CoEq}(g,h).$

PROOF 2.5.7 ► PROOF OF PROPOSITION 2.5.6

Item 1: Associativity

Omitted.

Item 2: Unitality

Clear.

Item 3: Commutativity

Clear.

Item 4: Interaction With Composition

Omitted.

002X 3 Operations With Sets

002Y 3.1 The Empty Set

002Z DEFINITION 3.1.1 ► THE EMPTY SET

The **empty set** is the set \emptyset defined by

$$\emptyset \stackrel{\text{def}}{=} \{ x \in X \mid x \neq x \},\$$

where A is the set in the set existence axiom, ?? of ??.

0030 3.2 Singleton Sets

Let X be a set.

0031 DEFINITION 3.2.1 ➤ SINGLETON SETS

The **singleton set containing** X is the set $\{X\}$ defined by

$$\{X\} \stackrel{\mathrm{def}}{=} \{X,X\},$$

where $\{X, X\}$ is the pairing of X with itself (Definition 3.3.1).

0032 3.3 Pairings of Sets

Let X and Y be sets.

0033 DEFINITION 3.3.1 ► PAIRINGS OF SETS

The **pairing of** X **and** Y is the set $\{X, Y\}$ defined by

$${X, Y} \stackrel{\text{def}}{=} {x \in A \mid x = X \text{ or } x = Y},$$

3.4 Ordered Pairs 48

where A is the set in the axiom of pairing, ?? of ??.

0034 3.4 Ordered Pairs

Let *A* and *B* be sets.

0035 DEFINITION 3.4.1 ► ORDERED PAIRS

The **ordered pair associated to** A **and** B is the set (A, B) defined by

$$(A, B) \stackrel{\text{def}}{=} \{ \{A\}, \{A, B\} \}.$$

0036 PROPOSITION 3.4.2 ► PROPERTIES OF ORDERED PAIRS

Let A and B be sets.

0037

0039

- 1. Uniqueness. Let A, B, C, and D be sets. The following conditions are equivalent:
- 0038 (a) We have (A, B) = (C, D).
 - (b) We have A = C and B = D.

PROOF 3.4.3 ► PROOF OF PROPOSITION 3.4.2

Item 1: Uniqueness

See [Cie97, Theorem 1.2.3].

003A 3.5 Sets of Maps

Let *A* and *B* be sets.

003B DEFINITION 3.5.1 ► SETS OF MAPS

The **set of maps from** A **to** B^1 is the set $Hom_{Sets}(A,B)^2$ whose elements are the functions from A to B.

¹ Further Terminology: Also called the **Hom set from** A **to** B.

² Further Notation: Also written Sets(A, B).

003C

PROPOSITION 3.5.2 ► PROPERTIES OF SETS OF MAPS

Let *A* and *B* be sets.

003D

1. Functoriality. The assignments $X,Y,(X,Y)\mapsto \operatorname{Hom}_{\mathsf{Sets}}(X,Y)$ define functors

$$\mathsf{Hom}_{\mathsf{Sets}}(X,-)\colon \mathsf{Sets} \to \mathsf{Sets},$$
 $\mathsf{Hom}_{\mathsf{Sets}}(-,Y)\colon \mathsf{Sets}^\mathsf{op} \to \mathsf{Sets},$ $\mathsf{Hom}_{\mathsf{Sets}}(-_1,-_2)\colon \mathsf{Sets}^\mathsf{op} \times \mathsf{Sets} \to \mathsf{Sets}.$

PROOF 3.5.3 ► PROOF OF PROPOSITION 3.5.2

Item 1: Functoriality

This follows from Categories, Items 2 and 5 of Proposition 1.6.2.

003E 3.6 Unions of Families

Let $\{A_i\}_{i\in I}$ be a family of sets.

003F

DEFINITION 3.6.1 ► Unions of Families

The **union of the family** $\{A_i\}_{i\in I}$ is the set $\bigcup_{i\in I}A_i$ defined by

$$\bigcup_{i \in I} A_i \stackrel{\text{def}}{=} \{x \in F \mid \text{there exists some } i \in I \text{ such that } x \in A_i\},$$

where F is the set in the axiom of union, ?? of ??.

003G 3.7 Binary Unions

Let A and B be sets.

003H

DEFINITION 3.7.1 ► BINARY UNIONS

The **union**¹ of A and B is the set $A \cup B$ defined by

$$A \cup B \stackrel{\text{def}}{=} \bigcup_{z \in \{A,B\}} z.$$

¹ Further Terminology: Also called the **binary union of** A **and** B, for emphasis.

003J PROPOSITION 3.7.2 ► PROPERTIES OF BINARY UNIONS

Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cup V$ define functors

$$U \cup -: (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$- \cup V : (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$-_1 \cup -_2 : (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset),$$

where $-_1 \cup -_2$ is the functor where

· Action on Objects. For each $(U,V) \in \mathcal{P}(X) \times \mathcal{P}(X)$, we have

$$[-_1 \cup -_2](U, V) \stackrel{\text{def}}{=} U \cup V.$$

· Action on Morphisms. For each pair of morphisms

$$\iota_U \colon U \hookrightarrow U',$$

 $\iota_V \colon V \hookrightarrow V'$

of $\mathcal{P}(X) \times \mathcal{P}(X)$, the image

$$\iota_U \cup \iota_V \colon U \cup V \hookrightarrow U' \cup V'$$

of (ι_U, ι_V) by \cup is the inclusion

$$U \cup V \subset U' \cup V'$$

i.e. where we have

$$(\star)$$
 If $U \subset U'$ and $V \subset V'$, then $U \cup V \subset U' \cup V'$.

and where $U \cup -$ and $- \cup V$ are the partial functors of $-_1 \cup -_2$ at $U, V \in \mathcal{P}(X)$.

003K

003L

2. Via Intersections and Symmetric Differences. We have an equality of sets

$$U \cup V = (U \triangle V) \triangle (U \cap V)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

003M

3. Associativity. We have an equality of sets

$$(U \cup V) \cup W = U \cup (V \cup W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

003N

4. Unitality. We have equalities of sets

$$U \cup \emptyset = U$$
,

$$\emptyset \cup U = U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

003P

5. Commutativity. We have an equality of sets

$$U \cup V = V \cup U$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

003Q

6. Idempotency. We have an equality of sets

$$U \cup U = U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

003R

7. Distributivity Over Intersections. We have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

003S

8. Interaction With Characteristic Functions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

003T

9. Interaction With Characteristic Functions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

003U

10. Interaction With Powersets and Semirings. The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

PROOF 3.7.3 ► PROOF OF PROPOSITION 3.7.2

Item 1: Functoriality

See [Pro24an].

Item 2: Via Intersections and Symmetric Differences

See [Pro24ay].

Item 3: Associativity

See [Pro24ba].

Item 4: Unitality

This follows from [Pro24bd] and Item 5.

Item 5: Commutativity

See [Pro24bb].

Item 6: Idempotency

See [Pro24am].

Item 7: Distributivity Over Intersections

See [Pro24az].

Item 8: Interaction With Characteristic Functions I

See [Pro24k].

Item 9: Interaction With Characteristic Functions II

See [Pro24k].

Item 10: Interaction With Powersets and Semirings

This follows from Items 3 to 6 and Items 3 to 5, 7 and 8 of Proposition 3.9.2.

003V 3.8 Intersections of Families

Let ${\mathcal F}$ be a family of sets.

003W

DEFINITION 3.8.1 ► INTERSECTIONS OF FAMILIES

The intersection of a family $\mathcal F$ of sets is the set $\bigcap_{X\in\mathcal F} X$ defined by

$$\bigcap_{X\in\mathcal{F}}X\stackrel{\mathrm{def}}{=} \bigg\{z\in\bigcup_{X\in\mathcal{F}}X\,\bigg|\, \text{for each}\, X\in\mathcal{F}, \text{ we have }z\in X\bigg\}.$$

Binary Intersections 003X **3.9**

Let *X* and *Y* be sets.

003Y **DEFINITION 3.9.1** ► BINARY INTERSECTIONS

The **intersection**¹ **of** X **and** Y is the set $X \cap Y$ defined by

$$X \cap Y \stackrel{\mathsf{def}}{=} \bigcap_{z \in \{X,Y\}} z.$$

¹ Further Terminology: Also called the **binary intersection of** X **and** Y, for emphasis.

003Z

PROPOSITION 3.9.2 ► PROPERTIES OF BINARY INTERSECTIONS

Let *X* be a set.

0040

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \cap -: (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$- \cap V \colon (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$-_1 \cap -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset),$$

where $-_1 \cap -_2$ is the functor where

· Action on Objects. For each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(X)$, we have

$$[-_1 \cap -_2](U, V) \stackrel{\mathsf{def}}{=} U \cap V.$$

· Action on Morphisms. For each pair of morphisms

$$\iota_U \colon U \hookrightarrow U',$$

 $\iota_V \colon V \hookrightarrow V'$

of
$$\mathcal{P}(X) \times \mathcal{P}(X)$$
, the image

$$\iota_{U} \cap \iota_{V} : U \cap V \hookrightarrow U' \cap V'$$

of (ι_U, ι_V) by \cap is the inclusion

$$U \cap V \subset U' \cap V'$$

i.e. where we have

$$(\star)$$
 If $U \subset U'$ and $V \subset V'$, then $U \cap V \subset U' \cap V'$.

and where $U \cap -$ and $- \cap V$ are the partial functors of $-_1 \cap -_2$ at $U, V \in$ $\mathcal{P}(X)$.

2. Adjointness. We have adjunctions

$$(-\cap V \dashv \mathsf{Hom}_{\mathcal{P}(X)}(V,-)): \mathcal{P}(X) \xrightarrow{-\cap V} \mathcal{P}(X)$$

where

$$\text{Hom}_{\mathcal{P}(X)}(-_1, -_2) \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

is the bifunctor defined by1

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,V)\stackrel{\mathsf{def}}{=} (X\setminus U)\cup V$$

witnessed by bijections

$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, \operatorname{Hom}_{\mathcal{P}(X)}(V, W)),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, \operatorname{Hom}_{\mathcal{P}(X)}(U, W)),$

natural in $U, V, W \in \mathcal{P}(X)$, i.e. where:

(a) The following conditions are equivalent:

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0044

0045

- i. We have $U \cap V \subset W$.
- ii. We have $U \subset \operatorname{Hom}_{\mathcal{P}(X)}(V, W)$.
- iii. We have $U \subset (X \setminus V) \cup W$.
- (b) The following conditions are equivalent:
 - i. We have $V \cap U \subset W$.
 - ii. We have $V \subset \mathbf{Hom}_{\mathcal{P}(X)}(U, W)$.
 - iii. We have $V \subset (X \setminus U) \cup W$.
- 3. Associativity. We have an equality of sets

$$(U \cap V) \cap W = U \cap (V \cap W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

4. Unitality. Let X be a set and let $U \in \mathcal{P}(X)$. We have equalities of sets

$$X \cap U = U$$
,

$$U \cap X = U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

5. Commutativity. We have an equality of sets

$$U \cap V = V \cap U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

6. Idempotency. We have an equality of sets

$$U \cap U = U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

7. Distributivity Over Unions. We have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$

$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

0047

8. Annihilation With the Empty Set. We have an equality of sets

$$\emptyset\cap X=\emptyset,$$

$$X \cap \emptyset = \emptyset$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

0048

9. Interaction With Characteristic Functions I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

0049

10. Interaction With Characteristic Functions II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

004A

11. Interaction With Powersets and Monoids With Zero. The quadruple $((\mathcal{P}(X),\emptyset),\cap,X)$ is a commutative monoid with zero.

004B

12. Interaction With Powersets and Semirings. The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

¹For intuition regarding the expression defining $\mathbf{Hom}_{\mathcal{P}(X)}(U,V)$, see Remark 3.9.4.

PROOF 3.9.3 ► PROOF OF PROPOSITION 3.9.2

Item 1: Functoriality

See [Pro24al].

Item 2: Adjointness

See [MSE 267469].

Item 3: Associativity

See [Pro24s].

Item 4: Unitality

This follows from [Pro24w] and Item 5.

3.10 Differences 58

Item 5: Commutativity

See [Pro24t].

Item 6: Idempotency

See [Pro24ak].

Item 7: Distributivity Over Unions

See [Pro24aj].

Item 8: Annihilation With the Empty Set

This follows from [Pro24u] and Item 5.

Item 9: Interaction With Characteristic Functions I

See [Pro24h].

Item 10: Interaction With Characteristic Functions II

See [Pro24h].

Item 11: Interaction With Powersets and Monoids With Zero

This follows from Items 3 to 5 and 8.

Item 12: Interaction With Powersets and Semirings

This follows from Items 3 to 6 and Items 3 to 5, 7 and 8 of Proposition 3.9.2.

004C REMARK 3.9.4 \blacktriangleright Intuition for the Internal Hom of $\mathcal{P}(X)$

Since intersections are the products in $\mathcal{P}(X)$ (Item 1 of Proposition 4.3.3), the left adjoint $\operatorname{Hom}_{\mathcal{P}(X)}(U,V)$ may be thought of as a function type [U,V]. Then, under the Curry–Howard correspondence, the function type [U,V] corresponds to implication $U \Longrightarrow V$, which is logically equivalent to the statement $\neg U \lor V$. This in turn corresponds to the set $U^{\mathsf{c}} \lor V = (X \setminus U) \cup V$.

004D 3.10 Differences

Let *X* and *Y* be sets.

3.10 Differences 59

004E DEFINITION 3.10.1 ➤ DIFFERENCES

The **difference of** X **and** Y is the set $X \setminus Y$ defined by

$$X \setminus Y \stackrel{\text{def}}{=} \{ a \in X \mid a \notin Y \}.$$

004F Proposition 3.10.2 ▶ Properties of Differences

Let X be a set.

004G

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \setminus -: (\mathcal{P}(X), \supset) \to (\mathcal{P}(X), \subset),$$
$$- \setminus V: (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$-_1 \setminus -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \supset) \to (\mathcal{P}(X), \subset),$$

where $-_1 \setminus -_2$ is the functor where

· Action on Objects. For each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(X)$, we have

$$[-_1 \setminus -_2](U,V) \stackrel{\mathsf{def}}{=} U \setminus V.$$

· Action on Morphisms. For each pair of morphisms

$$\iota_A \colon A \hookrightarrow B,$$

 $\iota_U \colon U \hookrightarrow V$

of $\mathcal{P}(X) \times \mathcal{P}(X)$, the image

$$\iota_U \setminus \iota_V : A \setminus V \hookrightarrow B \setminus U$$

of (ι_U, ι_V) by \ is the inclusion

$$A \setminus V \subset B \setminus U$$

i.e. where we have

 (\star) If $A \subset B$ and $U \subset V$, then $A \setminus V \subset B \setminus U$.

and where $U \setminus -$ and $- \setminus V$ are the partial functors of $-_1 \setminus -_2$ at $U, V \in \mathcal{P}(X)$.

004H

2. De Morgan's Laws. We have equalities of sets

$$X \setminus (U \cup V) = (X \setminus U) \cap (X \setminus V),$$

$$X \setminus (U \cap V) = (X \setminus U) \cup (X \setminus V)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

004J

3. Interaction With Unions I. We have equalities of sets

$$U \setminus (V \cup W) = (U \setminus V) \cap (U \setminus W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

004K

4. Interaction With Unions II. We have equalities of sets

$$(U\setminus V)\cup W=(U\cup W)\setminus (V\setminus W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

004L

5. Interaction With Unions III. We have equalities of sets

$$U \setminus (V \cup W) = (U \cup W) \setminus (V \cup W)$$
$$= (U \setminus V) \setminus W$$
$$= (U \setminus W) \setminus V$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

004M

6. Interaction With Unions IV. We have equalities of sets

$$(U \cup V) \setminus W = (U \setminus W) \cup (V \setminus W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

004N

7. Interaction With Intersections. We have equalities of sets

$$(U \setminus V) \cap W = (U \cap W) \setminus V$$
$$= U \cap (W \setminus V)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

004P

0040

8. Interaction With Complements. We have an equality of sets

$$U \setminus V = U \cap V^{c}$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. We have an equality of sets

$$U \setminus V = U \triangle (U \cap V)$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

004R 10. Triple Differences. We have

$$U \setminus (V \setminus W) = (U \cap W) \cup (U \setminus V)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

11. Left Annihilation. We have

$$\emptyset \setminus U = \emptyset$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

12. Right Unitality. We have

$$U \setminus \emptyset = U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

004U 13. Invertibility. We have

$$U \setminus U = \emptyset$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

004V 14. Interaction With Containment. The following conditions are equivalent:

- (a) We have $V \setminus U \subset W$.
- (b) We have $V \setminus W \subset U$.

15. Interaction With Characteristic Functions. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

004S

004T

004W 004X

004Y

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PROOF 3.10.3 ► PROOF OF PROPOSITION 3.10.2 Item 1: Functoriality See [Pro24ad] and [Pro24ah]. Item 2: De Morgan's Laws See [Pro24m]. Item 3: Interaction With Unions I See [Pro24n]. Item 4: Interaction With Unions II Omitted. Item 5: Interaction With Unions III See [Pro24ai]. Item 6: Interaction With Unions IV See [Pro24ac]. Item 7: Interaction With Intersections See [Pro24v]. Item 8: Interaction With Complements See [Pro24aa]. Item 9: Interaction With Symmetric Differences See [Pro24ab]. Item 10: Triple Differences See [Pro24ag]. Item 11: Left Annihilation Clear. Item 12: Right Unitality See [Pro24ae]. Item 13: Invertibility See [Pro24af]. Item 14: Interaction With Containment Omitted.

Item 15: Interaction With Characteristic Functions

See [Pro24i].

004Z 3.11 Complements

Let X be a set and let $U \in \mathcal{P}(X)$.

0050 DEFINITION 3.11.1 ► COMPLEMENTS

The **complement of** U is the set U^{c} defined by

$$U^{\mathsf{c}} \stackrel{\text{def}}{=} X \setminus U$$

$$\stackrel{\text{def}}{=} \{ a \in X \mid a \notin U \}.$$

0051 PROPOSITION 3.11.2 ▶ PROPERTIES OF COMPLEMENTS

Let *X* be a set.

0052

1. Functoriality. The assignment $U \mapsto U^{c}$ defines a functor

$$(-)^{c} : \mathcal{P}(X)^{op} \to \mathcal{P}(X),$$

where

· Action on Objects. For each $U \in \mathcal{P}(X)$, we have

$$[(-)^{\mathsf{c}}](U) \stackrel{\text{def}}{=} U^{\mathsf{c}}.$$

· Action on Morphisms. For each morphism $\iota_U\colon U\hookrightarrow V$ of $\mathcal{P}(X)$, the image

$$\iota_{IJ}^{\mathsf{c}} \colon V^{\mathsf{c}} \hookrightarrow U^{\mathsf{c}}$$

of ι_U by $(-)^{\mathsf{c}}$ is the inclusion

$$V^{\mathsf{c}} \subset U^{\mathsf{c}}$$

i.e. where we have

 (\star) If $U \subset V$, then $V^{c} \subset U^{c}$.

0053

2. De Morgan's Laws. We have equalities of sets

$$(U \cup V)^{c} = U^{c} \cap V^{c},$$

$$(U \cap V)^{c} = U^{c} \cup V^{c}$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

0054

3. Involutority. We have

$$(U^{\mathsf{c}})^{\mathsf{c}} = U$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

0055

4. Interaction With Characteristic Functions. We have

$$\chi_{U^{c}} = 1 - \chi_{U}$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

PROOF 3.11.3 ► PROOF OF PROPOSITION 3.11.2

Item 1: Functoriality

This follows from Item 1 of Proposition 3.10.2.

Item 2: De Morgan's Laws

See [Pro24m].

Item 3: Involutority

See [Pro24l].

Item 4: Interaction With Characteristic Functions

Clear.



0056 3.12 Symmetric Differences

Let A and B be sets.

0057 DEFINITION 3.12.1 ► SYMMETRIC DIFFERENCES

The **symmetric difference of** A **and** B is the set $A \triangle B$ defined by

$$A \triangle B \stackrel{\text{def}}{=} (A \setminus B) \cup (B \setminus A).$$

0058 PROPOSITION 3.12.2 ➤ PROPERTIES OF SYMMETRIC DIFFERENCES

Let X be a set.

0059

005B

005C

005D

1. Lack of Functoriality. The assignment $(U,V)\mapsto U\bigtriangleup V$ need not define functors

$$U \triangle -: (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$- \triangle V : (\mathcal{P}(X), \subset) \to (\mathcal{P}(X), \subset),$$
$$-_{1} \triangle -_{2} : (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

2. Via Unions and Intersections. We have

$$U \vartriangle V = (U \cup V) \setminus (U \cap V)$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

3. Associativity. We have²

$$(U \triangle V) \triangle W = U \triangle (V \triangle W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

4. Commutativity. We have

$$U \triangle V = V \triangle U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

5. Unitality. We have

$$U \triangle \emptyset = U,$$
$$\emptyset \triangle U = U$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

005E

6. Invertibility. We have

$$U \vartriangle U = \emptyset$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

005F

7. Interaction With Unions. We have

$$(U \triangle V) \cup (V \triangle T) = (U \cup V \cup W) \setminus (U \cap V \cap W)$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

005G

8. Interaction With Complements I. We have

$$U \triangle U^{c} = X$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

005H

9. Interaction With Complements II. We have

$$U \mathbin{\vartriangle} X = U^{\mathsf{c}},$$

$$X \triangle U = U^{\mathsf{c}}$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

005J

10. Interaction With Complements III. We have

$$U^{c} \triangle V^{c} = U \triangle V$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

005K

11. "Transitivity". We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle W$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

005L

12. The Triangle Inequality for Symmetric Differences. We have

$$U \triangle W \subset U \triangle V \cup V \triangle W$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

005M

13. Distributivity Over Intersections. We have

$$U \cap (V \triangle W) = (U \cap V) \triangle (U \cap W),$$

$$(U \triangle V) \cap W = (U \cap W) \triangle (V \cap W)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V, W \in \mathcal{P}(X)$.

005N

14. Interaction With Characteristic Functions. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \mod 2$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

005P

15. Bijectivity. Given $A, B \subset \mathcal{P}(X)$, the maps

$$A \triangle -: \mathcal{P}(X) \to \mathcal{P}(X),$$

 $- \triangle B: \mathcal{P}(X) \to \mathcal{P}(X)$

are bijections with inverses given by

$$(A \triangle -)^{-1} = - \cup (A \cap -),$$

 $(- \triangle B)^{-1} = - \cup (B \cap -).$

Moreover, the map

$$C \mapsto C \triangle (A \triangle B)$$

is a bijection of $\mathcal{P}(X)$ onto itself sending A to B and B to A.

005Q

16. Interaction With Powersets and Groups. Let X be a set.

005R

(a) The quadruple $(\mathcal{P}(X), \triangle, \emptyset, \mathrm{id}_{\mathcal{P}(X)})$ is an abelian group. ³

005S

(b) Every element of $\mathcal{P}(X)$ has order 2 with respect to \triangle , and thus $\mathcal{P}(X)$ is a Boolean group (i.e. an abelian 2-group).

005T

17. Interaction With Powersets and Vector Spaces I. The pair $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ consisting of

- · The group $\mathcal{P}(X)$ of ??;
- · The map $\alpha_{\mathcal{P}(X)} : \mathbb{F}_2 \times \mathcal{P}(X) \to \mathcal{P}(X)$ defined by

$$0\cdot U\stackrel{\mathrm{def}}{=}\emptyset,$$

$$1 \cdot U \stackrel{\text{def}}{=} U$$
:

is an \mathbb{F}_2 -vector space.

- 18. Interaction With Powersets and Vector Spaces II. If X is finite, then:
 - (a) The set of singletons sets on the elements of X forms a basis for the \mathbb{F}_2 -vector space $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ of Item 17.
 - (b) We have

$$\dim(\mathcal{P}(X)) = \#\mathcal{P}(X).$$

19. Interaction With Powersets and Rings. The quintuple $(\mathcal{P}(X), \triangle, \cap, \emptyset, X)$ is a commutative ring.⁴

005U

005V

$$\boxed{\bigcup_{U \land V}} = \boxed{\bigcup_{U \lor V}} \setminus \boxed{\bigcup_{U \circlearrowleft V}}$$

²Illustration:



³Here are some examples:

i. When $X = \emptyset$, we have an isomorphism of groups between $\mathcal{P}(\emptyset)$ and the trivial group:

$$\left(\mathcal{P}(\emptyset), \triangle, \emptyset, \mathsf{id}_{\mathcal{P}(\emptyset)}\right) \cong \mathsf{pt}.$$

ii. When $X=\operatorname{pt}$, we have an isomorphism of groups between $\mathcal{P}(\operatorname{pt})$ and $\mathbb{Z}_{/2}$:

$$\left(\mathcal{P}(\mathsf{pt}), \triangle, \emptyset, \mathsf{id}_{\mathcal{P}(\mathsf{pt})}\right) \cong \mathbb{Z}_{/2}.$$

iii. When $X=\{0,1\}$, we have an isomorphism of groups between $\mathcal{P}(\{0,1\})$ and $\mathbb{Z}_{/2}\times\mathbb{Z}_{/2}$:

$$\left(\mathcal{P}(\{0,1\}), \triangle, \emptyset, \mathsf{id}_{\mathcal{P}(\{0,1\})}\right) \cong \mathbb{Z}_{/2} \times \mathbb{Z}_{/2}.$$

4 Warning: The analogous statement replacing intersections by unions (i.e. that the quintuple $(\mathcal{P}(X), \Delta, \cup, \emptyset, X)$ is a ring) is false, however. See [Pro24aw] for a proof.

¹Illustration:

PROOF 3.12.3 ► PROOF OF PROPOSITION 3.12.2

Item 1: Lack of Functoriality

Omitted.

Item 2: Via Unions and Intersections

See [Pro240].

Item 3: Associativity

See [Pro24ao].

Item 4: Commutativity

See [Pro24ap].

Item 5: Unitality

This follows from Item 4 and [Pro24at].

Item 6: Invertibility

See [Pro24av].

Item 7: Interaction With Unions

See [Pro24bc].

Item 8: Interaction With Complements I

See [Pro24as].

Item 9: Interaction With Complements II

This follows from Item 4 and [Pro24ax].

Item 10: Interaction With Complements III

See [Pro24aq].

Item 11: "Transitivity"

We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle (V \triangle (V \triangle W))$$
 (by Item 3)

$$= U \triangle ((V \triangle V) \triangle W)$$
 (by Item 6)

$$= U \triangle (\emptyset \triangle W)$$
 (by Item 5)

Item 12: The Triangle Inequality for Symmetric Differences

This follows from Items 2 and 11. Item 13: Distributivity Over Intersections See [Pro24r]. Item 14: Interaction With Characteristic Functions See [Pro24j]. Item 15: Bijectivity Clear. Item 16: Interaction With Powersets and Groups Item 16a follows from Items 3 to 6, while Item 16b follows from Item 6. Item 17: Interaction With Powersets and Vector Spaces I Item 18: Interaction With Powersets and Vector Spaces II Omitted. Item 19: Interaction With Powersets and Rings This follows from Items 8 and 11 of Proposition 3.9.2 and Items 13 and 16.2 ¹Reference: [Pro24ar]. ²Reference: [Pro24au].

005W 4 Powersets

005X 4.1 Characteristic Functions

Let X be a set.

005Z

005Y DEFINITION 4.1.1 ► CHARACTERISTIC FUNCTIONS

Let $U \subset X$ and let $x \in X$.

1. The **characteristic function of** U^1 is the function²

$$\chi_U \colon X \to \{\mathsf{t},\mathsf{f}\}$$

defined by

$$\chi_U(x) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in U, \\ \text{false} & \text{if } x \notin U \end{cases}$$

for each $x \in X$.

2. The **characteristic function of** x is the function ³

$$\chi_x \colon X \to \{\mathsf{t},\mathsf{f}\}$$

defined by

$$\chi_x \stackrel{\text{def}}{=} \chi_{\{x\}},$$

i.e. by

$$\chi_x(y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $y \in X$.

3. The **characteristic relation on** X^4 is the relation⁵

$$\chi_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

on X defined by 6

$$\chi_X(x,y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $x, y \in X$.

4. The **characteristic embedding** of X into $\mathcal{P}(X)$ is the function

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

defined by

$$\chi_{(-)}(x) \stackrel{\text{def}}{=} \chi_x$$

for each $x \in X$.

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¹ Further Terminology: Also called the **indicator function of** U.

⁷The name "characteristic *embedding*" comes from the fact that there is an analogue of fully faithfulness for $\chi_{(-)}$: given a set X, we have

$$\operatorname{Hom}_{\mathcal{P}(X)}(\chi_{x},\chi_{y})=\chi_{X}(x,y),$$

for each $x, y \in X$.

0063 REMARK 4.1.2 ➤ CHARACTERISTIC FUNCTIONS AS DECATEGORIFICATIONS OF PRESHEAVES

The definitions in Definition 4.1.1 are decategorifications of co/presheaves, representable co/presheaves, Hom profunctors, and the Yoneda embedding:¹

0064 1. A function

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$$f: X \to \{\mathsf{t},\mathsf{f}\}$$

is a decategorification of a presheaf

$$\mathcal{F} \colon \mathcal{C}^{\mathsf{op}} \to \mathsf{Sets},$$

with the characteristic functions χ_U of the subsets of X being the primordial examples (and, in fact, all examples) of these.

2. The characteristic function

$$\chi_X \colon X \to \{\mathsf{t},\mathsf{f}\}$$

of an element x of X is a decategorification of the representable presheaf

$$h_X \colon C^{\mathsf{op}} \to \mathsf{Sets}$$

of an object x of a category C.

3. The characteristic relation

$$\chi_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

of X is a decategorification of the Hom profunctor

$$\operatorname{Hom}_C(-1,-2): C^{\operatorname{op}} \times C \to \operatorname{Sets}$$

of a category C.

² Further Notation: Also written $\chi_X(U,-)$ or $\chi_X(-,U)$.

³ Further Notation: Also written χ^x , $\chi_X(x,-)$, or $\chi_X(-,x)$.

⁴ Further Terminology: Also called the **identity relation on** X.

⁵ Further Notation: Also written χ_{-2}^{-1} , or \sim_{id} in the context of relations.

⁶As a subset of $X \times X$, the relation χ_X corresponds to the diagonal $\Delta_X \subset X \times X$ of X.

4. The characteristic embedding

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

of X into $\mathcal{P}(X)$ is a decategorification of the Yoneda embedding

of a category C into PSh(C).

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- 5. There is also a direct parallel between unions and colimits:
 - · An element of $\mathcal{P}(X)$ is a union of elements of X, viewed as one-point subsets $\{x\} \in \mathcal{P}(A)$.
 - · An object of PSh(C) is a colimit of objects of C, viewed as representable presheaves $h_X \in Obj(PSh(C))$.

$$(-)_{\mbox{disc}} \colon \mbox{Sets} \hookrightarrow \mbox{Cats},$$
 $(-)_{\mbox{disc}} \colon \{t,f\}_{\mbox{disc}} \hookrightarrow \mbox{Sets}$

of sets into categories and of classical truth values into sets. For instance, in this approach the characteristic function

$$\chi_X \colon X \to \{\mathsf{t},\mathsf{f}\}$$

of an element x of X, defined by

$$\chi_X(y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $y \in X$, is recovered as the representable presheaf

$$\operatorname{Hom}_{X_{\operatorname{disc}}}(-,x)\colon X_{\operatorname{disc}} \to \operatorname{Sets}$$

of the corresponding object x of X_{disc} , defined on objects by

$$\operatorname{Hom}_{X_{\operatorname{disc}}}(y,x) \stackrel{\text{def}}{=} \begin{cases} \operatorname{pt} & \text{if } x = y, \\ \emptyset & \text{if } x \neq y \end{cases}$$

for each $y \in \text{Obj}(X_{\text{disc}})$.

¹These statements can be made precise by using the embeddings

0069 Proposition 4.1.3 ➤ Properties of Characteristic Functions

Let *X* be a set.

006A

006C

006D

006E

006F

1. The Inclusion of Characteristic Relations Associated to a Function. Let $f: A \to B$ be a function. We have an inclusion 1

$$\chi_B \circ (f \times f) \subset \chi_A, \qquad A \times A \xrightarrow{f \times f} B \times B$$

$$\chi_A \searrow \chi_A \qquad \chi_A \downarrow \chi_B$$

$$\{t, f\}.$$

006B 2. Interaction With Unions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

3. Interaction With Unions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $X \in \mathsf{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

4. Interaction With Intersections I. We have

$$\chi_{U\cap V}=\chi_U\chi_V$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

5. Interaction With Intersections II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

6. Interaction With Differences. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

006G

7. Interaction With Complements. We have

$$\chi_{U^c} = 1 - \chi_U$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U \in \mathcal{P}(X)$.

006H

8. Interaction With Symmetric Differences. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \mod 2$$

for each $X \in \text{Obj}(\mathsf{Sets})$ and each $U, V \in \mathcal{P}(X)$.

006J

9. Interaction Between the Characteristic Embedding and Morphisms. Let $f: X \to Y$ be a map of sets. The diagram

$$f_* \circ \chi_X = \chi_{X'} \circ f, \qquad \chi_X \downarrow \qquad \qquad \downarrow \chi_{X'} \downarrow$$

commutes.

PROOF 4.1.4 ► PROOF OF PROPOSITION 4.1.3

Item 1: The Inclusion of Characteristic Relations Associated to a Function

The inclusion $\chi_B(f(a), f(b)) \subset \chi_A(a, b)$ is equivalent to the statement "if a = b, then f(a) = f(b)", which is true.

Item 2: Interaction With Unions I

This is a repetition of Item 8 of Proposition 3.7.2 and is proved there.

Item 3: Interaction With Unions II

This is a repetition of Item 9 of Proposition 3.7.2 and is proved there.

¹This is the 0-categorical version of Categories, Definition 4.4.1.

Item 4: Interaction With Intersections I

This is a repetition of Item 9 of Proposition 3.9.2 and is proved there.

Item 5: Interaction With Intersections II

This is a repetition of Item 10 of Proposition 3.9.2 and is proved there.

Item 6: Interaction With Differences

This is a repetition of Item 15 of Proposition 3.10.2 and is proved there.

Item 7: Interaction With Complements

This is a repetition of Item 4 of Proposition 3.11.2 and is proved there.

Item 8: Interaction With Symmetric Differences

This is a repetition of Item 14 of Proposition 3.12.2 and is proved there.

Item 9: Interaction Between the Characteristic Embedding and Morphisms

Indeed, we have

$$[f_* \circ \chi_X](x) \stackrel{\text{def}}{=} f_*(\chi_X(x))$$

$$\stackrel{\text{def}}{=} f_*(\{x\})$$

$$= \{f(x)\}$$

$$\stackrel{\text{def}}{=} \chi_{X'}(f(x))$$

$$\stackrel{\text{def}}{=} [\chi_{X'} \circ f](x),$$

for each $x \in X$, showing the desired equality.

006K 4.2 The Yoneda Lemma for Sets

Let X be a set and let $U \subset X$ be a subset of X.

006L Proposition 4.2.1 ➤ The Yoneda Lemma for Sets

We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_U) = \chi_U(x)$$

for each $x \in X$, giving an equality of functions

$$\chi_{\mathcal{P}(X)}(\chi_{(-)},\chi_U)=\chi_U.$$

PROOF 4.2.2 ▶	PROOF OF	PROPOSITION	4.2.1
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Clear.

006M

COROLLARY 4.2.3 ► THE CHARACTERISTIC EMBEDDING IS FULLY FAITHFUL

The characteristic embedding is fully faithful, i.e., we have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_y)=\chi_X(x,y)$$

for each $x, y \in X$.

PROOF 4.2.4 ► PROOF OF COROLLARY 4.2.3

This follows from Proposition 4.2.1.

006N 4.3 Powersets

Let X be a set.

006P

DEFINITION 4.3.1 ► POWERSETS

The **powerset of** X is the set $\mathcal{P}(X)$ defined by

$$\mathcal{P}(X) \stackrel{\text{def}}{=} \{ U \in P \mid U \subset X \},\$$

where P is the set in the axiom of powerset, ?? of ??.

006Q

REMARK 4.3.2 ► POWERSETS AS DECATEGORIFICATIONS OF CO/PRESHEAF CATEGORIES

The powerset of a set is a decategorification of the category of presheaves of a category: while¹

• The powerset of a set X is equivalently (Items 1 and 2 of Proposition 4.3.9) the set

$$Sets(X, \{t, f\})$$

of functions from X to the set $\{t,f\}$ of classical truth values.

 \cdot The category of presheaves on a category C is the category

$$\operatorname{Fun}(\mathcal{C}^{\operatorname{op}},\operatorname{\mathsf{Sets}})$$

of functors from C^{op} to the category Sets of sets.

¹This parallel is based on the following comparison:

· A category is enriched over the category

$$Sets \stackrel{\text{def}}{=} Cats_0$$

of sets (i.e. "0-categories"), with presheaves taking values on it.

· A set is enriched over the set

$$\{t, f\} \stackrel{\text{def}}{=} \mathsf{Cats}_{-1}$$

of classical truth values (i.e. "(-1)-categories"), with characteristic functions taking values on it.

006R

PROPOSITION 4.3.3 ► PROPERTIES OF POWERSETS: AS CATEGORIES

Let X be a set.

006S

- 1. *Co/Completeness*. The (posetal) category (associated to) $(\mathcal{P}(X), \subset)$ is complete and cocomplete:
 - (a) *Products*. The products in $\mathcal{P}(X)$ are given by intersection of subsets.
 - (b) *Coproducts*. The coproducts in $\mathcal{P}(X)$ are given by union of subsets.
 - (c) Co/Equalisers. Being a posetal category, $\mathcal{P}(X)$ only has at most one morphisms between any two objects, so co/equalisers are trivial.

006T

2. Cartesian Closedness. The category $\mathcal{P}(X)$ is Cartesian closed with internal Hom

$$\operatorname{Hom}_{\mathcal{P}(X)}(-_1, -_2) \colon \mathcal{P}(X)^{\operatorname{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

given by1

$$\operatorname{Hom}_{\mathcal{P}(X)}(U,V) \stackrel{\text{def}}{=} (X \setminus U) \cup V$$

for each $U, V \in \text{Obj}(\mathcal{P}(X))$.

 $^{^1}$ For intuition regarding the expression defining $\mathbf{Hom}_{\mathcal{P}(X)}(U,V)$, see Remark 3.9.4.

PROOF 4.3.4 ► PROOF OF PROPOSITION 4.3.3

Item 1: Co/Completeness

Clear.

Item 2: Cartesian Closedness

This follows from Item 2 of Proposition 3.9.2.

006U

PROPOSITION 4.3.5 ► PROPERTIES OF POWERSETS: FUNCTORIALITY AND ADJOINTNESS

Let *X* be a set.

006V

1. Functoriality I. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_* : \mathsf{Sets} \to \mathsf{Sets}$$
,

where

· Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_*(A) \stackrel{\text{def}}{=} \mathcal{P}(A)$$
.

· Action on Morphisms. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{*|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of \mathcal{P}_* at (A,B) is the map defined by by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}_*(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_*(f) \stackrel{\text{def}}{=} f_*,$$

as in Definition 4.4.1.

006W

2. Functoriality II. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}^{-1}$$
: Sets^{op} \rightarrow Sets.

where

· Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}^{-1}(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

· Action on Morphisms. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{A.B}^{-1} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(B),\mathcal{P}(A))$$

of \mathcal{P}^{-1} at (A,B) is the map defined by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}^{-1}(f) \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by

$$\mathcal{P}^{-1}(f) \stackrel{\text{def}}{=} f^{-1}$$
,

as in Definition 4.5.1.

3. Functoriality III. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_1 \colon \mathsf{Sets} \to \mathsf{Sets}$$

where

· Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_!(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

· Action on Morphisms. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{!|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of $\mathcal{P}_!$ at (A,B) is the map defined by by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}_!(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_!(f) \stackrel{\text{def}}{=} f_!,$$

as in Definition 4.6.1.

006X

006Y

4. Adjointness I. We have an adjunction

$$\big(\mathcal{P}^{-1}\dashv\mathcal{P}^{-1,\mathsf{op}}\big)\colon\quad\mathsf{Sets}^{\mathsf{op}}\underbrace{\overset{\mathcal{P}^{-1}}{\bot}}_{\mathcal{P}^{-1,\mathsf{op}}}\mathsf{Sets},$$

witnessed by a bijection

$$\underbrace{\mathsf{Sets}^{\mathsf{op}}(\mathcal{P}(A), B)}_{\overset{\mathsf{def}}{=}\mathsf{Sets}(B, \mathcal{P}(A))} \cong \mathsf{Sets}(A, \mathcal{P}(B)),$$

natural in $A \in Obj(Sets)$ and $B \in Obj(Sets^{op})$.

006Z

5. Adjointness II. We have an adjunction

$$(\operatorname{\mathsf{Gr}} \dashv \mathcal{P}_*) \colon \operatorname{\mathsf{Sets}} \underbrace{\overset{\operatorname{\mathsf{Gr}}}{\vdash}}_{\mathcal{P}_*} \operatorname{\mathsf{Rel}},$$

witnessed by a bijection of sets

$$Rel(Gr(A), B) \cong Sets(A, \mathcal{P}(B))$$

natural in $A \in \text{Obj}(\mathsf{Sets})$ and $B \in \text{Obj}(\mathsf{Rel})$, where Gr is the graph functor of Constructions With Relations, Item 1 of Proposition 3.1.2 and \mathcal{P}_* is the functor of Constructions With Relations, Proposition 4.5.1.

PROOF 4.3.6 ► PROOF OF PROPOSITION 4.3.5

Item 1: Functoriality I

This follows from Items 3 and 4 of Proposition 4.4.6.

Item 2: Functoriality II

This follows Items 3 and 4 of Proposition 4.5.5.

Item 3: Functoriality III

This follows Items 3 and 4 of Proposition 4.6.8.

Item 4: Adjointness I

We have

```
Sets^{op}(\mathcal{P}(A), B) \stackrel{\text{def}}{=} Sets(B, \mathcal{P}(A))
\cong Sets(B, Sets(A, \{t, f\})) \quad \text{(by Item 1 of Proposition 4.3.9)}
\cong Sets(A \times B, \{t, f\}) \quad \text{(by Item 2 of Proposition 1.3.3)}
\cong Sets(A, Sets(B, \{t, f\})) \quad \text{(by Item 2 of Proposition 1.3.3)}
\cong Sets(A, \mathcal{P}(B)) \quad \text{(by Item 1 of Proposition 4.3.9)}
```

with all bijections natural in A and B (where we use Item 2 of Proposition 4.3.9 here).

Item 5: Adjointness II

We have

with all bijections natural in A (where we use Item 2 of Proposition 4.3.9 here). Explicitly, this isomorphism is given by sending a relation $R: Gr(A) \to B$ to the map $R^{\dagger}: A \to \mathcal{P}(B)$ sending a to the subset R(a) of B, as in Relations, Remark 1.1.4. Naturality in B is then the statement that given a relation $R: B \to B'$, the diagram

commutes, which follows from Constructions With Relations, Remark 4.1.2.

0070 PROPOSITION 4.3.7 ► PROPERTIES OF POWERSETS: MONOIDALITY

Let *X* be a set.

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1. Symmetric Strong Monoidality With Respect to Coproducts I. The powerset functor \mathcal{P}_* of Item 1 of Proposition 4.3.5 has a symmetric strong monoidal structure

$$\left(\mathcal{P}_*,\mathcal{P}_*^{\coprod},\mathcal{P}_{*|\mathbb{1}}^{\coprod}\right)\colon (\mathsf{Sets},\mathsf{X},\mathsf{pt})\to (\mathsf{Sets},\underline{\coprod},\emptyset)$$

being equipped with isomorphisms

$$\begin{split} \mathcal{P}^{\coprod}_{*|X,Y} \colon \mathcal{P}(X) \times \mathcal{P}(Y) &\xrightarrow{\cong} \mathcal{P}(X \coprod Y), \\ \mathcal{P}^{\coprod}_{*|\mathfrak{I}} \colon \mathsf{pt} &\xrightarrow{\cong} \mathcal{P}(\emptyset), \end{split}$$

natural in $X, Y \in Obj(Sets)$.

2. Symmetric Strong Monoidality With Respect to Coproducts II. The powerset functor \mathcal{P}^{-1} of Item 2 of Proposition 4.3.5 has a symmetric strong monoidal structure

$$\left(\mathcal{P}^{-1},\mathcal{P}^{-1|\coprod},\mathcal{P}_{\mathbb{1}}^{-1|\coprod}\right)\colon(\mathsf{Sets}^{\mathsf{op}},\mathsf{x}^{\mathsf{op}},\mathsf{pt})\to(\mathsf{Sets},\coprod,\emptyset)$$

being equipped with isomorphisms

$$\mathcal{P}_{X,Y}^{-1|\coprod} : \mathcal{P}(X) \times \mathcal{P}(Y) \xrightarrow{\cong} \mathcal{P}(X \coprod Y),$$
$$\mathcal{P}_{1}^{-1|\coprod} : \mathsf{pt} \xrightarrow{\cong} \mathcal{P}(\emptyset),$$

natural in $X, Y \in Obj(Sets)$.

3. Symmetric Strong Monoidality With Respect to Coproducts III. The powerset functor $\mathcal{P}_!$ of Item 3 of Proposition 4.3.5 has a symmetric strong monoidal structure

$$\left(\mathcal{P}_!,\mathcal{P}_!^{\coprod},\mathcal{P}_{!\mid \mathbb{1}}^{\coprod}\right) \colon (\mathsf{Sets},\mathsf{x},\mathsf{pt}) \to (\mathsf{Sets}, \underline{\coprod}, \emptyset)$$

being equipped with isomorphisms

$$\begin{split} \mathcal{P}^{\coprod}_{!|X,Y} \colon \mathcal{P}(X) \times \mathcal{P}(Y) &\xrightarrow{\cong} \mathcal{P}(X \coprod Y), \\ \mathcal{P}^{\coprod}_{!|\mathfrak{1}} \colon \mathsf{pt} &\xrightarrow{\cong} \mathcal{P}(\emptyset), \end{split}$$

natural in $X, Y \in Obj(Sets)$.

4. Symmetric Lax Monoidality With Respect to Products I. The powerset functor \mathcal{P}_* of Item 1 of Proposition 4.3.5 has a symmetric lax monoidal structure

$$\left(\mathcal{P}_*,\mathcal{P}_*^\otimes,\mathcal{P}_{*|\mathbb{1}}^\otimes\right)\colon(\mathsf{Sets},\mathsf{x},\mathsf{pt})\to(\mathsf{Sets},\mathsf{x},\mathsf{pt})$$

being equipped with morphisms

$$\mathcal{P}_{*|X,Y}^{\times} \colon \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y),$$
$$\mathcal{P}_{*|\mathfrak{I}}^{\times} \colon \mathsf{pt} \to \mathcal{P}(\mathsf{pt}),$$

natural in $X, Y \in \mathsf{Obj}(\mathsf{Sets})$, where

· The map $\mathcal{P}_{*|X,Y}^{\times}$ is given by

$$\mathcal{P}_{*|X,Y}^{\times}(U,V) \stackrel{\text{def}}{=} U \times V$$

for each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$,

· The map $\mathcal{P}_{*|_{\mathbb{I}}}^{\times}$ is given by

$$\mathcal{P}_{*|1}^{\times}(\star) = \mathsf{pt}.$$

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5. Symmetric Lax Monoidality With Respect to Products II. The powerset functor \mathcal{P}^{-1} of Item 2 of Proposition 4.3.5 has a symmetric lax monoidal structure

$$\left(\mathcal{P}^{-1},\mathcal{P}^{-1|\otimes},\mathcal{P}_{\mathbb{1}}^{-1|\otimes}\right)\colon(\mathsf{Sets}^{\mathsf{op}},\mathsf{x}^{\mathsf{op}},\mathsf{pt})\to(\mathsf{Sets},\mathsf{x},\mathsf{pt})$$

being equipped with morphisms

$$\mathcal{P}_{X,Y}^{-1|\times} \colon \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y),$$
$$\mathcal{P}_{1}^{\times} \colon \mathsf{pt} \to \mathcal{P}(\emptyset),$$

natural in $X, Y \in Obj(Sets)$, defined as in Item 4.

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6. Symmetric Lax Monoidality With Respect to Products III. The powerset functor \mathcal{P}_1 of Item 3 of Proposition 4.3.5 has a symmetric lax monoidal structure

$$\left(\mathcal{P}_!, \mathcal{P}_!^{\otimes}, \mathcal{P}_{!|1}^{\otimes}\right) \colon (\mathsf{Sets}, \mathsf{x}, \mathsf{pt}) \to (\mathsf{Sets}, \mathsf{x}, \mathsf{pt})$$

being equipped with morphisms

$$\mathcal{P}_{!|X,Y}^{\times} \colon \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y),$$
$$\mathcal{P}_{!|\mathfrak{I}}^{\times} \colon \mathsf{pt} \to \mathcal{P}(\emptyset),$$

natural in $X, Y \in Obj(Sets)$, defined as in Item 4.

PROOF 4.3.8 ► PROOF OF PROPOSITION 4.3.7

Item 1: Symmetric Strong Monoidality With Respect to Coproducts I

The isomorphism

$$\mathcal{P}^{\coprod}_{*|X,Y} \colon \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \coprod Y)$$

is given by sending $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$ to $U \coprod V$, with inverse given by sending a subset S of $X \coprod Y$ to the pair $(S_X, S_Y) \in \mathcal{P}(X) \times \mathcal{P}(Y)$ with

$$S_X \stackrel{\text{def}}{=} \{ x \in X \mid (0, x) \in S \}$$

 $S_Y \stackrel{\text{def}}{=} \{ y \in Y \mid (1, y) \in S \}.$

The isomorphism pt $\cong \mathcal{P}(\emptyset)$ is given by $\star \mapsto \emptyset \in \mathcal{P}(\emptyset)$.

Naturality for the isomorphism $\mathcal{P}^{\coprod}_{*|X,Y}$ is the statement that, given maps of sets $f\colon X\to X'$ and $g\colon Y\to Y'$, the diagram

$$\mathcal{P}(X) \times \mathcal{P}(Y) \xrightarrow{f_* \times g_*} \mathcal{P}(X') \times \mathcal{P}(Y')$$

$$\downarrow \wr \qquad \qquad \downarrow \downarrow \qquad \qquad \downarrow \qquad \qquad$$

commutes, which is clear, as it acts on elements as

$$(U,V) \longmapsto (f_*(U),g_*(V))$$

$$\downarrow \qquad \qquad \downarrow$$

$$U \coprod V \longmapsto (f \coprod g)_*(U \coprod V) = f_*(U) \coprod g_*(V).$$

where we are using Item 7 of Proposition 4.4.4.

Finally, monoidality, unity, and symmetry of \mathcal{P}_* as a monoidal functor follow by checking the commutativity of the relevant diagrams on elements.

Item 2: Symmetric Strong Monoidality With Respect to Coproducts II

The proof is similar to Item 1, and is hence omitted.

Item 3: Symmetric Strong Monoidality With Respect to Coproducts III

The proof is similar to Item 1, and is hence omitted.

Item 4: Symmetric Lax Monoidality With Respect to Products I

Naturality for the morphism $\mathcal{P}_{*|X,Y}^{\times}$ is the statement that, given maps of sets $f\colon X\to X'$ and $g\colon Y\to Y'$, the diagram

$$\mathcal{P}(X) \times \mathcal{P}(Y) \xrightarrow{f_* \times g_*} \mathcal{P}(X') \times \mathcal{P}(Y')$$

$$\downarrow \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad$$

commutes, which is clear, as it acts on elements as

$$(U,V) \longmapsto (f_*(U),g_*(V))$$

$$\downarrow \qquad \qquad \downarrow$$

$$U \times V \longmapsto (f \times g)_*(U \times V) = f_*(U) \times g_*(V),$$

where we are using Item 8 of Proposition 4.4.4.

Finally, monoidality, unity, and symmetry of \mathcal{P}_* as a monoidal functor follow by checking the commutativity of the relevant diagrams on elements.

Item 5: Symmetric Lax Monoidality With Respect to Products II

The proof is similar to Item 4, and is hence omitted.

Item 6: Symmetric Lax Monoidality With Respect to Products III

The proof is similar to Item 4, and is hence omitted.

0077 PROPOSITION 4.3.9 ➤ PROPERTIES OF POWERSETS: AS SETS OF FUNCTIONS/RELATIONS

Let X be a set.

0078

1. Powersets as Sets of Functions I. The assignment $U\mapsto \chi_U$ defines a bijection

$$\chi_{(-)} \colon \mathcal{P}(X) \xrightarrow{\cong} \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}),$$

for each $X \in Obj(Sets)$.

2. Powersets as Sets of Functions II. The bijection

$$\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

of Item 1 is natural in $X \in \mathsf{Obj}(\mathsf{Sets})$, refining to a natural isomorphism of functors

$$\mathcal{P}^{-1} \cong \mathsf{Sets}(-, \{t, f\}).$$

3. Powersets as Sets of Relations. We have bijections

$$\mathcal{P}(X) \cong \mathsf{Rel}(\mathsf{pt}, X),$$

$$\mathcal{P}(X) \cong \text{Rel}(X, \text{pt}),$$

natural in $X \in Obj(Sets)$.

PROOF 4.3.10 ► PROOF OF PROPOSITION 4.3.9

Item 1: Powersets as Sets of Functions I

Indeed, the inverse of $\chi_{(-)}$ is given by sending a function $f: X \to \{\mathsf{t}, \mathsf{f}\}$ to the subset U_f of $\mathcal{P}(X)$ defined by

$$U_f \stackrel{\text{def}}{=} \{ x \in X \mid f(x) = \text{true} \},$$

i.e. by $U_f=f^{-1}({\rm true})$. That $\chi_{(-)}$ and $f\mapsto U_f$ are inverses is then straightforward to check.

Item 2: Powersets as Sets of Functions II

We need to check that, given a function $f: X \to Y$, the diagram

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

$$\chi_{(-)} \downarrow \chi \qquad \qquad \downarrow \chi_{(-)}$$

$$\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{f^*} \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

commutes, i.e. that for each $V \in \mathcal{P}(Y)$, we have

$$\chi_V \circ f = \chi_{f^{-1}(V)}$$
.

007A

0079

And indeed, we have

$$\begin{split} [\chi_V \circ f](v) &\stackrel{\text{def}}{=} \chi_V(f(v)) \\ &= \begin{cases} \text{true} & \text{if } f(v) \in V, \\ \text{false} & \text{otherwise} \end{cases} \\ &= \begin{cases} \text{true} & \text{if } v \in f^{-1}(V), \\ \text{false} & \text{otherwise} \end{cases} \\ &\stackrel{\text{def}}{=} \chi_{f^{-1}(V)}(v) \end{split}$$

for each $v \in V$.

Item 3: Powersets as Sets of Relations

Indeed, we have

$$\mathsf{Rel}(\mathsf{pt}, X) \stackrel{\mathsf{def}}{=} \mathcal{P}(\mathsf{pt} \times X)$$
$$\cong \mathcal{P}(X)$$

and

$$\mathsf{Rel}(X,\mathsf{pt}) \stackrel{\mathsf{def}}{=} \mathcal{P}(X \times \mathsf{pt})$$

 $\cong \mathcal{P}(X),$

where we have used Item 4 of Proposition 1.3.3.

007B

REMARK 4.3.11 ▶ Powersets as Sets of Functions and Un/Straightening

The bijection

$$\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

of Item 1 of Proposition 4.3.9, which

- · Takes a subset $U \hookrightarrow X$ of X and straightens it to a function $\chi_U \colon X \to \{\text{true}, \text{false}\};$
- · Takes a function $f: X \to \{\text{true}, \text{false}\}$ and unstraightens it to a subset $f^{-1}(\text{true}) \hookrightarrow X$ of X;

may be viewed as the (-1)-categorical version of the un/straightening isomorphism for indexed and fibred sets

$$\underbrace{\mathsf{FibSets}(X)}_{\overset{\text{def}}{=}\mathsf{Sets}_{/X}} \cong \underbrace{\mathsf{ISets}(X)}_{\overset{\text{def}}{=}\mathsf{Fun}(X_{\mathsf{disc}},\mathsf{Sets})}$$

of ??, ??, where we view:

- · Subsets $U \hookrightarrow X$ as analogous to X-fibred sets $\phi_X \colon A \to X$.
- · Functions $f: X \to \{t, f\}$ as analogous to X-indexed sets $A: X_{disc} \to \mathsf{Sets}$.

007C PROPOSITION 4.3.12 ➤ PROPERTIES OF POWERSETS: AS FREE COCOMPLETIONS

Let *X* be a set.

007D

- 1. Universal Property. The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of
 - · The powerset $\mathcal{P}(X)$ of X;
 - · The characteristic embedding $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$;

satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - A cocomplete poset (Y, \preceq) ;
 - A function $f: X \to Y$;

there exists a unique cocontinuous morphism of posets

$$(\mathcal{P}(X),\subset) \xrightarrow{\exists !} (Y,\preceq)$$

making the diagram



commute.

007E

2. Adjointness. We have an adjunction¹

$$(\mathcal{P} \dashv \overline{\Xi})$$
: Sets $\stackrel{\mathcal{P}}{\underset{\Xi}{\longleftarrow}}$ Pos^{cocomp.}

witnessed by a bijection

$$\mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq)) \cong \mathsf{Sets}(X,Y),$$

natural in $X \in \mathsf{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \mathsf{Obj}(\mathsf{Pos}^{\mathsf{cocomp.}})$, where the maps witnessing this bijection are given by

· The map

$$\chi_X^* \colon \mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq)) \to \mathsf{Sets}(X,Y)$$

defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a cocontinuous morphism of posets $f\colon \mathcal{P}(X) \to Y$ to the composition

$$X \stackrel{\chi_X}{\hookrightarrow} \mathcal{P}(X) \stackrel{f}{\longrightarrow} Y.$$

· The map

$$\mathsf{Lan}_{\chi_X} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq))$$

is given by sending a function $f\colon X\to Y$ to its left Kan extension along χ_X ,

$$\operatorname{Lan}_{\chi_X}(f) \colon \mathcal{P}(X) \to Y, \qquad \begin{array}{c} \mathcal{P}(X) \\ \chi_X & \downarrow \\ X \xrightarrow{f} & Y. \end{array}$$

Moreover, $\operatorname{Lan}_{\chi_{\!X}}(f)$ can be explicitly computed by

$$\begin{split} \big[\mathsf{Lan}_{\chi_X}(f) \big](U) &\cong \int_{-x \in X}^{x \in X} \chi_{\mathcal{P}(X)}(\chi_x, U) \odot f(x) \\ &\cong \int_{-x \in X}^{x \in X} \chi_U(x) \odot f(x) \qquad \text{(by Proposition 4.2.1)} \\ &\cong \bigvee_{x \in X} (\chi_U(x) \odot f(x)) \end{split}$$

for each $U \in \mathcal{P}(X)$, where:

- \lor is the join in (Y, \preceq) .
- We have

true
$$\odot f(x) \stackrel{\text{def}}{=} f(x)$$
,
false $\odot f(x) \stackrel{\text{def}}{=} \varnothing_Y$,

where \emptyset_Y is the minimal element of (Y, \preceq) .

PROOF 4.3.13 ► PROOF OF PROPOSITION 4.3.7

Item 1: Universal Property

This is a rephrasing of Item 2.

Item 2: Adjointness

We claim we have adjunction \mathcal{P} \dashv 忘, witnessed by a bijection

$$\mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq)) \cong \mathsf{Sets}(X,Y),$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \text{Obj}(\mathsf{Pos}^{\mathsf{cocomp.}})$.

· Map I. We define a map

$$\Phi_{X,Y} \colon \mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq)) \to \mathsf{Sets}(X,Y)$$

as in the statement, by

$$\Phi_{X,Y}(f) \stackrel{\mathsf{def}}{=} f \circ \chi_X$$

¹In this sense, $\mathcal{P}(A)$ is the free cocompletion of A. (Note that, despite its name, however, this is not an idempotent operation, as we have $\mathcal{P}(\mathcal{P}(A)) \neq \mathcal{P}(A)$.)

for each
$$f \in \mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq)).$$

· Map II. We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq))$$

as in the statement, by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \mathsf{Lan}_{\chi_X}(f), \qquad X \xrightarrow{\chi_X} \downarrow \mathsf{Lan}_{\chi_X}(f)$$

$$X \xrightarrow{f} Y,$$

for each $f \in \mathsf{Sets}(X, Y)$.

· Invertibility I. We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathsf{id}_{\mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq))}.$$

Indeed, given a cocontinuous morphism of posets

$$\xi \colon (\mathcal{P}(X), \subset) \to (Y, \preceq),$$

we have

$$\begin{split} \left[\Psi_{X,Y} \circ \Phi_{X,Y}\right] (\xi) &\stackrel{\text{def}}{=} \Psi_{X,Y} \big(\Phi_{X,Y}(\xi)\big) \\ &\stackrel{\text{def}}{=} \Psi_{X,Y} (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} \operatorname{Lan}_{\chi_X} (\xi \circ \chi_X) \\ &\cong \bigvee_{x \in X} \chi_{(-)}(x) \odot \xi(\chi_X(x)) \\ &\stackrel{\text{clm}}{=} \xi, \end{split}$$

where indeed

$$\begin{split} \left[\bigvee_{x \in X} \chi_{(-)}(x) \odot \xi(\chi_X(x))\right] (U) &\stackrel{\text{def}}{=} \bigvee_{x \in X} \chi_U(x) \odot \xi(\chi_X(x)) \\ &= \left(\bigvee_{x \in U} \chi_U(x) \odot \xi(\chi_X(x))\right) \vee \left(\bigvee_{x \in X \setminus U} \chi_U(x) \odot \xi(\chi_X(x))\right) \\ &= \left(\bigvee_{x \in U} \xi(\chi_X(x))\right) \vee \left(\bigvee_{x \in X \setminus U} \varnothing_Y\right) \\ &= \bigvee_{x \in U} \xi(\chi_X(x)) \\ &\stackrel{(\dagger)}{=} \xi \left(\bigvee_{x \in U} \chi_X(x)\right) \\ &= \xi(U) \end{split}$$

for each $U\in\mathcal{P}(X)$, where we have used that ξ is cocontinuous for the equality $\stackrel{(\dagger)}{=}$.

· Invertibility II. We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathrm{Sets}(X,Y)}$$
.

Indeed, given a function $f: X \to Y$, we have

$$\begin{split} \left[\Phi_{X,Y} \circ \Psi_{X,Y}\right](f) &\stackrel{\text{def}}{=} \Phi_{X,Y}\big(\Psi_{X,Y}(f)\big) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}\big(\text{Lan}_{\chi_X}(f)\big) \\ &\stackrel{\text{def}}{=} \text{Lan}_{\chi_X}(f) \circ \chi_X \\ &\stackrel{\text{clm}}{=} f, \end{split}$$

where indeed

$$\begin{split} \left[\mathsf{Lan}_{\chi_X}(f) \circ \chi_X \right] (x) & \stackrel{\mathsf{def}}{=} \bigvee_{y \in X} \chi_{\{x\}}(y) \odot f(y) \\ & = \left(\chi_{\{x\}}(x) \odot f(x) \right) \vee \left(\bigvee_{y \in X \setminus \{x\}} \chi_{\{x\}}(y) \odot f(y) \right) \\ & = f(x) \vee \left(\bigvee_{y \in X \setminus \{x\}} \varnothing_Y \right) \\ & = f(x) \vee \varnothing_Y \\ & = f(x) \end{split}$$

for each $x \in X$.

· Naturality for Φ , Part I. We need to show that, given a function $f\colon X\to X'$, the diagram

$$\begin{array}{ccc} \mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X'),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} \mathsf{Sets}(X',Y) \\ & & & \downarrow f^* \\ & & & \downarrow f^* \end{array}$$

$$\mathsf{Pos}^{\mathsf{cocomp.}}((\mathcal{P}(X),\subset),(Y,\preceq)) \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y)$$

commutes. Indeed, given a cocontinuous morphism of posets

$$\xi \colon (\mathcal{P}(X'), \subset) \to (Y, \preceq),$$

we have

$$\begin{split} \left[\Phi_{X,Y}\circ\mathcal{P}_*(f)^*\right](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}\big(\mathcal{P}_*(f)^*(\xi)\big) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\xi\circ f_*) \\ &\stackrel{\text{def}}{=} (\xi\circ f_*)\circ\chi_X \\ &= \xi\circ (f_*\circ\chi_X) \\ &\stackrel{(\dagger)}{=} \xi\circ (\chi_{X'}\circ f) \\ &= (\xi\circ\chi_{X'})\circ f \\ &\stackrel{\text{def}}{=} \Phi_{X',Y}(\xi)\circ f \\ &\stackrel{\text{def}}{=} f^*\big(\Phi_{X',Y}(\xi)\big) \\ &\stackrel{\text{def}}{=} \left[f^*\circ\Phi_{X',Y}\right](\xi), \end{split}$$

where we have used Item 9 of Proposition 4.1.3 for the equality $\stackrel{(\dagger)}{=}$ above.

· Naturality for Φ , Part II. We need to show that, given a cocontinuous morphism of posets

$$q: (Y, \preceq_Y) \to (Y', \preceq_{Y'}),$$

the diagram

$$\begin{array}{ccc} \mathsf{Pos^{\mathsf{cocomp.}}}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} & \mathsf{Sets}(X,Y) \\ & & & \downarrow g_* \\ & & & \downarrow g_* \end{array}$$

$$\mathsf{Pos^{\mathsf{cocomp.}}}((\mathcal{P}(X),\subset),(Y',\preceq)) \xrightarrow{\Phi_{X,Y'}} & \mathsf{Sets}(X,Y')$$

commutes. Indeed, given a cocontinuous morphism of posets

$$\xi \colon (\mathcal{P}(X), \subset) \to (Y, \preceq),$$

we have

$$\begin{split} \left[\Phi_{X,Y'} \circ g_*\right](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_*(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} g_*(\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} \left[g_* \circ \Phi_{X,Y}\right](\xi). \end{split}$$

• Naturality for Ψ . Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Proposition 8.6.2 that Ψ is also natural in each argument.

This finishes the proof.

007F 4.4 Direct Images

Let A and B be sets and let $f: A \rightarrow B$ be a function.

007G DEFINITION 4.4.1 ► DIRECT IMAGES

The **direct image function associated to** f is the function

$$f_* \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by^{1,2}

$$f_*(U) \stackrel{\text{def}}{=} f(U)$$

$$\stackrel{\text{def}}{=} \left\{ b \in B \middle| \begin{array}{l} \text{there exists some } a \in U \\ \text{such that } b = f(a) \end{array} \right\}$$

$$= \left\{ f(a) \in B \mid a \in U \right\}$$

for each $U \in \mathcal{P}(A)$.

¹ Further Terminology: The set f(U) is called the **direct image of** U **by** f.

²We also have

$$f_*(U) = B \setminus f_!(A \setminus U);$$

see Item 9 of Proposition 4.4.4.

007H NOTATION 4.4.2 ► FURTHER NOTATION FOR DIRECT IMAGES

Sometimes one finds the notation

$$\exists_f \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

for f_* . This notation comes from the fact that the following statements are equivalent, where $b \in B$ and $U \in \mathcal{P}(A)$:

- · We have $b \in \exists_f(U)$.
- · There exists some $a \in U$ such that f(a) = b.

007J REMARK 4.4.3 ► UNWINDING DEFINITION 4.4.1

Identifying subsets of A with functions from A to $\{ \text{true}, \text{false} \}$ via $\frac{1}{2}$ and $\frac{2}{2}$ of Proposition 4.3.9, we see that the direct image function associated to f is equivalently the function

$$f_* \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$f_*(\chi_U) \stackrel{\text{def}}{=} \operatorname{Lan}_f(\chi_U)$$

$$= \operatorname{colim}\left(\left(f \stackrel{\rightarrow}{\times} (\underline{-_1})\right) \stackrel{\text{pr}}{\twoheadrightarrow} A \stackrel{\chi_U}{\longrightarrow} \{\mathsf{t},\mathsf{f}\}\right)$$

$$= \operatorname{colim}_{\substack{a \in A \\ f(a) = -_1}} (\chi_U(a))$$

$$= \bigvee_{\substack{a \in A \\ f(a) = -_1}} (\chi_U(a)),$$

where we have used ??, ?? for the second equality. In other words, we have

$$[f_*(\chi_U)](b) = \bigvee_{\substack{a \in A \\ f(a) = b}} (\chi_U(a))$$

$$= \begin{cases} \text{true} & \text{if there exists some } a \in A \text{ such} \\ & \text{that } f(a) = b \text{ and } a \in U, \end{cases}$$

$$= \begin{cases} \text{true} & \text{if there exists some } a \in U \\ & \text{such that } f(a) = b, \end{cases}$$

$$\text{false} & \text{otherwise}$$

for each $b \in B$.

007K PROPOSITION 4.4.4 ➤ PROPERTIES OF DIRECT IMAGES I

Let $f: A \to B$ be a function.

1. Functoriality. The assignment $U\mapsto f_*(U)$ defines a functor

$$f_* \colon (\mathcal{P}(A), \subset) \to (\mathcal{P}(B), \subset)$$

where

· Action on Objects. For each $U \in \mathcal{P}(A)$, we have

$$[f_*](U) \stackrel{\text{def}}{=} f_*(U).$$

· Action on Morphisms. For each $U, V \in \mathcal{P}(A)$:

$$(\star)$$
 If $U \subset V$, then $f_*(U) \subset f_*(V)$.

2. Triple Adjointness. We have a triple adjunction

$$(f_* \dashv f^{-1} \dashv f_!): \mathcal{P}(A) \leftarrow f^{-1} - \mathcal{P}(B),$$

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007M

witnessed by bijections of sets

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(B)}(f_*(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(A)}(U,f^{-1}(V)),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(A)}(f^{-1}(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(A)}(U,f_!(V)),$

natural in $U \in \mathcal{P}(A)$ and $V \in \mathcal{P}(B)$ and (respectively) $V \in \mathcal{P}(A)$ and $U \in \mathcal{P}(B)$, i.e. where:

- (a) The following conditions are equivalent:
 - i. We have $f_*(U) \subset V$.
 - ii. We have $U \subset f^{-1}(V)$.
- (b) The following conditions are equivalent:
 - i. We have $f^{-1}(U) \subset V$.
 - ii. We have $U \subset f_!(V)$.
- 3. Preservation of Colimits. We have an equality of sets

$$f_*\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f_*(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(A)^{ imes I}$. In particular, we have equalities

$$f_*(U) \cup f_*(V) = f_*(U \cup V),$$

$$f_*(\emptyset) = \emptyset,$$

natural in $U, V \in \mathcal{P}(A)$.

4. Oplax Preservation of Limits. We have an inclusion of sets

$$f_*\left(\bigcap_{i\in I}U_i\right)\subset\bigcap_{i\in I}f_*(U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(A)^{\times I}$. In particular, we have inclusions

$$f_*(U \cap V) \subset f_*(U) \cap f_*(V),$$

 $f_*(A) \subset B,$

natural in $U, V \in \mathcal{P}(A)$.

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007Q

5. Symmetric Strict Monoidality With Respect to Unions. The direct image function of Item1 has a symmetric strict monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}) \colon (\mathcal{P}(A), \cup, \emptyset) \to (\mathcal{P}(B), \cup, \emptyset),$$

being equipped with equalities

$$f_{*|U,V}^{\otimes} \colon f_{*}(U) \cup f_{*}(V) \xrightarrow{=} f_{*}(U \cup V),$$
$$f_{*|1}^{\otimes} \colon \emptyset \xrightarrow{=} \emptyset,$$

natural in $U, V \in \mathcal{P}(A)$.

007R

6. Symmetric Oplax Monoidality With Respect to Intersections. The direct image function of Item 1 has a symmetric oplax monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|\mathbb{1}}^{\otimes}) \colon (\mathcal{P}(A), \cap, A) \to (\mathcal{P}(B), \cap, B),$$

being equipped with inclusions

$$f_{*|U,V}^{\otimes} \colon f_{*}(U \cap V) \hookrightarrow f_{*}(U) \cap f_{*}(V),$$
$$f_{*|1}^{\otimes} \colon f_{*}(A) \hookrightarrow B,$$

natural in $U, V \in \mathcal{P}(A)$.

007S

7. Interaction With Coproducts. Let $f:A\to A'$ and $g:B\to B'$ be maps of sets. We have

$$(f \coprod g)_*(U \coprod V) = f_*(U) \coprod g_*(V)$$

for each $U \in \mathcal{P}(A)$ and each $V \in \mathcal{P}(B)$.

007T

8. Interaction With Products. Let $f:A\to A'$ and $g:B\to B'$ be maps of sets. We have

$$(f \times g)_*(U \times V) = f_*(U) \times g_*(V)$$

for each $U \in \mathcal{P}(A)$ and each $V \in \mathcal{P}(B)$.

007U

9. Relation to Direct Images With Compact Support. We have

$$f_*(U) = B \setminus f_!(A \setminus U)$$

for each $U \in \mathcal{P}(A)$.

PROOF 4.4.5 ► PROOF OF PROPOSITION 4.4.4

Item 1: Functoriality

Clear.

Item 2: Triple Adjointness

This follows from Remark 4.4.3, Remark 4.5.2, Remark 4.6.3, and ??, ?? of ??.

Item 3: Preservation of Colimits

This follows from Item 2 and ??, ?? of ??.1

Item 4: Oplax Preservation of Limits

The inclusion $f_*(A) \subset B$ is clear. See [Pro24p] for the other inclusions.

Item 5: Symmetric Strict Monoidality With Respect to Unions

This follows from Item 3.

Item 6: Symmetric Oplax Monoidality With Respect to Intersections

This follows from Item 4.

Item 7: Interaction With Coproducts

Clear.

Item 8: Interaction With Products

Clear.

Item 9: Relation to Direct Images With Compact Support

Applying Item 9 of Proposition 4.6.6 to $A \setminus U$, we have

$$f_!(A \setminus U) = B \setminus f_*(A \setminus (A \setminus U))$$
$$= B \setminus f_*(U).$$

Taking complements, we then obtain

$$f_*(U) = B \setminus (B \setminus f_*(U)),$$

= $B \setminus f_!(A \setminus U),$

which finishes the proof.

¹See also [Pro24q].

007V

PROPOSITION 4.4.6 ► PROPERTIES OF DIRECT IMAGES II

Let $f: A \to B$ be a function.

007W

1. Functionality I. The assignment $f\mapsto f_*$ defines a function

$$(-)_{*|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B)).$$

007X

2. Functionality II. The assignment $f \mapsto f_*$ defines a function

$$(-)_{*|A|B}$$
: Sets $(A, B) \to \mathsf{Pos}((\mathcal{P}(A), \subset), (\mathcal{P}(B), \subset)).$

007Y

3. Interaction With Identities. For each $A \in Obj(Sets)$, we have

$$(\mathrm{id}_A)_* = \mathrm{id}_{\mathcal{P}(A)}.$$

007Z

4. Interaction With Composition. For each pair of composable functions $f:A\to B$ and $g:B\to C$, we have

$$(g \circ f)_* = g_* \circ f_*,$$

$$\mathcal{P}(A) \xrightarrow{f_*} \mathcal{P}(B)$$

$$\downarrow^{g_*}$$

$$\mathcal{P}(C).$$

PROOF 4.4.7 ► PROOF OF PROPOSITION 4.4.6

Item 1: Functionality I

Clear.

Item 2: Functionality II

Clear.

Item 3: Interaction With Identities

This follows from Remark 4.4.3 and ??, ?? of ??.

Item 4: Interaction With Composition

This follows from Remark 4.4.3 and ??, ?? of ??.

0080 4.5 Inverse Images

Let A and B be sets and let $f: A \rightarrow B$ be a function.

0081 DEFINITION 4.5.1 ► INVERSE IMAGES

The **inverse image function associated to** f is the function¹

$$f^{-1} \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by²

$$f^{-1}(V) \stackrel{\text{def}}{=} \{a \in A \mid \text{we have } f(a) \in V\}$$

for each $V \in \mathcal{P}(B)$.

0082 REMARK 4.5.2 ► UNWINDING DEFINITION 4.5.1

Identifying subsets of B with functions from B to {true, false} via Items 1 and 2 of Proposition 4.3.9, we see that the inverse image function associated to f is equivalently the function

$$f^* \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by

$$f^*(\chi_V) \stackrel{\text{def}}{=} \chi_V \circ f$$

for each $\chi_V \in \mathcal{P}(B)$, where $\chi_V \circ f$ is the composition

$$A \xrightarrow{f} B \xrightarrow{\chi_V} \{\text{true}, \text{false}\}$$

in Sets.

0083 Proposition 4.5.3 ► Properties of Inverse Images I

Let $f: A \rightarrow B$ be a function.

¹Further Notation: Also written $f^* : \mathcal{P}(B) \to \mathcal{P}(A)$.

² Further Terminology: The set $f^{-1}(V)$ is called the **inverse image of** V by f.

1. Functoriality. The assignment $V \mapsto f^{-1}(V)$ defines a functor

$$f^{-1} \colon (\mathcal{P}(B), \subset) \to (\mathcal{P}(A), \subset)$$

where

· Action on Objects. For each $V \in \mathcal{P}(B)$, we have

$$[f^{-1}](V) \stackrel{\text{def}}{=} f^{-1}(V).$$

· Action on Morphisms. For each $U, V \in \mathcal{P}(B)$:

$$(\star)$$
 If $U \subset V$, then $f^{-1}(U) \subset f^{-1}(V)$.

2. Triple Adjointness. We have a triple adjunction

$$(f_* \dashv f^{-1} \dashv f_!): \mathcal{P}(A) \leftarrow f^{-1} - \mathcal{P}(B),$$

witnessed by bijections of sets

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(B)}(f_*(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(A)}(U,f^{-1}(V)),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(A)}(f^{-1}(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(A)}(U,f_!(V)),$

natural in $U \in \mathcal{P}(A)$ and $V \in \mathcal{P}(B)$ and (respectively) $V \in \mathcal{P}(A)$ and $U \in \mathcal{P}(B)$, i.e. where:

- (a) The following conditions are equivalent:
 - i. We have $f_*(U) \subset V$;
 - ii. We have $U \subset f^{-1}(V)$;
- (b) The following conditions are equivalent:
 - i. We have $f^{-1}(U) \subset V$.
 - ii. We have $U \subset f_!(V)$.

0085

3. Preservation of Colimits. We have an equality of sets

$$f^{-1}\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(B)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cup f^{-1}(V) = f^{-1}(U \cup V),$$

 $f^{-1}(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(B)$.

0087

4. Preservation of Limits. We have an equality of sets

$$f^{-1}\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(B)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cap f^{-1}(V) = f^{-1}(U \cap V),$$

 $f^{-1}(B) = A,$

natural in $U, V \in \mathcal{P}(B)$.

0088

5. Symmetric Strict Monoidality With Respect to Unions. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$\left(f^{-1}, f^{-1, \otimes}, f_{\mathbb{1}}^{-1, \otimes}\right) \colon (\mathcal{P}(B), \cup, \emptyset) \to (\mathcal{P}(A), \cup, \emptyset),$$

being equipped with equalities

$$\begin{split} f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cup f^{-1}(V) &\stackrel{=}{\to} f^{-1}(U \cup V), \\ f_{\parallel}^{-1,\otimes} \colon \emptyset &\stackrel{=}{\to} f^{-1}(\emptyset), \end{split}$$

natural in $U, V \in \mathcal{P}(B)$.

6. Symmetric Strict Monoidality With Respect to Intersections. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$\left(f^{-1}, f^{-1,\otimes}, f_{\mathbb{1}}^{-1,\otimes}\right) \colon (\mathcal{P}(B), \cap, B) \to (\mathcal{P}(A), \cap, A),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cap f^{-1}(V) \xrightarrow{=} f^{-1}(U \cap V),$$

$$f_{\parallel}^{-1,\otimes} \colon A \xrightarrow{=} f^{-1}(B),$$

natural in $U, V \in \mathcal{P}(B)$.

008A

7. Interaction With Coproducts. Let $f:A\to A'$ and $g:B\to B'$ be maps of sets. We have

$$(f \coprod g)^{-1}(U' \coprod V') = f^{-1}(U') \coprod g^{-1}(V')$$

for each $U' \in \mathcal{P}(A')$ and each $V' \in \mathcal{P}(B')$.

008B

8. Interaction With Products. Let $f:A\to A'$ and $g:B\to B'$ be maps of sets. We have

$$(f \times g)^{-1}(U' \times V') = f^{-1}(U') \times g^{-1}(V')$$

for each $U' \in \mathcal{P}(A')$ and each $V' \in \mathcal{P}(B')$.

PROOF 4.5.4 ► PROOF OF PROPOSITION 4.5.3

Item 1: Functoriality

Clear.

Item 2: Triple Adjointness

This follows from Remark 4.4.3, Remark 4.5.2, Remark 4.6.3, and ??, ?? of ??.

Item 3: Preservation of Colimits

This follows from Item 2 and ??, ?? of ??.1

Item 4: Preservation of Limits

This follows from Item 2 and ??, ?? of ??.²

Item 5: Symmetric Strict Monoidality With Respect to Unions

This follows from Item 3.

Item 6: Symmetric Strict Monoidality With Respect to Intersections

This follows from Item 4.

Item 7: Interaction With Coproducts

Clear.

Item 8: Interaction With Products

Clear.

008D

008E

008F

008G

¹See also [Pro24y].

²See also [Pro24x].

008C PROPOSITION 4.5.5 ► PROPERTIES OF INVERSE IMAGES II

Let $f: A \to B$ be a function.

1. Functionality I. The assignment $f \mapsto f^{-1}$ defines a function

$$(-)_{A,B}^{-1}$$
: Sets $(A,B) \to \text{Sets}(\mathcal{P}(B),\mathcal{P}(A))$.

2. Functionality II. The assignment $f\mapsto f^{-1}$ defines a function

$$(-)^{-1}_{AB}$$
: Sets $(A, B) \to \mathsf{Pos}((\mathcal{P}(B), \subset), (\mathcal{P}(A), \subset)).$

3. Interaction With Identities. For each $A \in Obj(Sets)$, we have

$$\operatorname{id}_A^{-1} = \operatorname{id}_{\mathcal{P}(A)}.$$

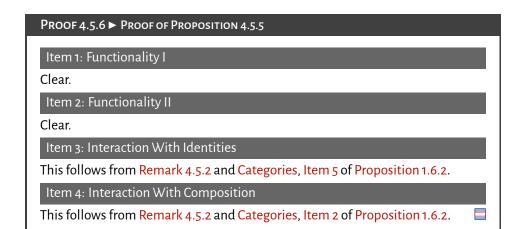
4. Interaction With Composition. For each pair of composable functions $f\colon A\to B$ and $g\colon B\to C$, we have

$$(g \circ f)^{-1} = f^{-1} \circ g^{-1},$$

$$\mathcal{P}(C) \xrightarrow{g^{-1}} \mathcal{P}(B)$$

$$(g \circ f)^{-1} \qquad \qquad \downarrow_{f^{-1}}$$

$$\mathcal{P}(A)$$



008H 4.6 Direct Images With Compact Support

Let *A* and *B* be sets and let $f: A \rightarrow B$ be a function.

008J DEFINITION 4.6.1 ➤ DIRECT IMAGES WITH COMPACT SUPPORT

The direct image with compact support function associated to f is the function

$$f_! \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by^{1,2}

$$f_!(U) \stackrel{\text{def}}{=} \left\{ b \in B \middle| \begin{array}{l} \text{for each } a \in A, \text{ if we have} \\ f(a) = b, \text{ then } a \in U \end{array} \right\}$$

$$= \left\{ b \in B \middle| \text{ we have } f^{-1}(b) \subset U \right\}$$

for each $U \in \mathcal{P}(A)$.

$$f_!(U) = B \setminus f_*(A \setminus U);$$

see Item 9 of Proposition 4.6.6.

 $^{^{1}}$ Further Terminology: The set $f_{!}(U)$ is called the **direct image with compact support of** U **by** f.

²We also have

008K

NOTATION 4.6.2 ► FURTHER NOTATION FOR DIRECT IMAGES WITH COMPACT SUPPORT

Sometimes one finds the notation

$$\forall_f \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

for f_* . This notation comes from the fact that the following statements are equivalent, where $b \in B$ and $U \in \mathcal{P}(A)$:

- · We have $b \in \forall_f(U)$.
- · For each $a \in A$, if b = f(a), then $a \in U$.

008L

REMARK 4.6.3 ► UNWINDING DEFINITION 4.6.1

Identifying subsets of A with functions from A to {true, false} via Items 1 and 2 of Proposition 4.3.9, we see that the direct image with compact support function associated to f is equivalently the function

$$f_! \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\begin{split} f_!(\chi_U) &\stackrel{\text{def}}{=} \mathsf{Ran}_f(\chi_U) \\ &= \mathsf{lim}\Big(\Big(\underbrace{(-_1)}_{} \stackrel{\rightarrow}{\times} f\Big) \stackrel{\mathsf{pr}}{\twoheadrightarrow} A \xrightarrow{\chi_U} \{\mathsf{true}, \mathsf{false}\}\Big) \\ &= \lim_{\substack{a \in A \\ f(a) = -_1}} (\chi_U(a)) \\ &= \bigwedge_{\substack{a \in A \\ f(a) = -_1}} (\chi_U(a)). \end{split}$$

where we have used ??, ?? for the second equality. In other words, we have

$$[f_!(\chi_U)](b) = \bigwedge_{\substack{a \in A \\ f(a) = b}} (\chi_U(a))$$

$$= \begin{cases} \text{true} & \text{if, for each } a \in A \text{ such that} \\ f(a) = b, \text{ we have } a \in U, \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } f^{-1}(b) \subset U \\ \text{false} & \text{otherwise} \end{cases}$$

for each $b \in B$.

008M

DEFINITION 4.6.4 ► THE IMAGE AND COMPLEMENT PARTS OF $f_!$

Let U be a subset of A.^{1,2}

008N

1. The image part of the direct image with compact support $f_!(U)$ of U is the set $f_!, \text{im}(U)$ defined by

$$\begin{split} f_{\mathrm{l,im}}(U) &\stackrel{\mathrm{def}}{=} f_{\mathrm{!}}(U) \cap \mathrm{Im}(f) \\ &= \left\{ b \in B \,\middle|\, \begin{aligned} &\text{we have } f^{-1}(b) &\subset \\ &U \text{ and } f^{-1}(b) \neq \emptyset \end{aligned} \right\}. \end{split}$$

008P

2. The complement part of the direct image with compact support $f_!(U)$ of U is the set $f_!(U)$ defined by

$$\begin{split} f_{!,\mathsf{cp}}(U) &\stackrel{\mathsf{def}}{=} f_!(U) \cap (B \setminus \mathsf{Im}(f)) \\ &= B \setminus \mathsf{Im}(f) \\ &= \left\{ b \in B \,\middle|\, \begin{aligned} \mathsf{we have} \, f^{-1}(b) &\subset \\ U \, \mathsf{and} \, f^{-1}(b) &= \emptyset \end{aligned} \right\} \\ &= \left\{ b \in B \,\middle|\, f^{-1}(b) &= \emptyset \right\}. \end{split}$$

$$f_!(U) = f_{!,\mathsf{im}}(U) \cup f_{!,\mathsf{cp}}(U),$$

¹Note that we have

as

$$\begin{split} f_!(U) &= f_!(U) \cap B \\ &= f_!(U) \cap (\operatorname{Im}(f) \cup (B \setminus \operatorname{Im}(f))) \\ &= (f_!(U) \cap \operatorname{Im}(f)) \cup (f_!(U) \cap (B \setminus \operatorname{Im}(f))) \\ &\stackrel{\text{def}}{=} f_{!,\operatorname{im}}(U) \cup f_{!,\operatorname{cp}}(U). \end{split}$$

 2 In terms of the meet computation of $f_!(U)$ of Remark 4.6.3, namely

$$f_!(\chi_U) = \bigwedge_{\substack{a \in A \\ f(a) = -1}} (\chi_U(a)),$$

we see that $f_{i,im}$ corresponds to meets indexed over nonempty sets, while $f_{i,cp}$ corresponds to meets indexed over the empty set.

008Q

EXAMPLE 4.6.5 ► **EXAMPLES OF DIRECT IMAGES WITH COMPACT SUPPORT**

Here are some examples of direct images with compact support.

1. The Multiplication by Two Map on the Natural Numbers. Consider the function $f: \mathbb{N} \to \mathbb{N}$ given by

$$f(n) \stackrel{\text{def}}{=} 2n$$

for each $n \in \mathbb{N}$. Since f is injective, we have

$$f_{!,\text{im}}(U) = f_*(U)$$

 $f_{!,\text{cp}}(U) = \{\text{odd natural numbers}\}$

for any $U \subset \mathbb{N}$.

2. Parabolas. Consider the function $f: \mathbb{R} \to \mathbb{R}$ given by

$$f(x) \stackrel{\text{def}}{=} x^2$$

for each $x \in \mathbb{R}$. We have

$$f_{!,\mathsf{cp}}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}$. Moreover, since $f^{-1}(x) = \{-\sqrt{x}, \sqrt{x}\}$, we have e.g.:

$$f_{!,\text{im}}([0,1]) = \{0\},\$$

$$f_{!,im}([-1,1]) = [0,1],$$

$$f_{!,\mathsf{im}}([1,2]) = \emptyset,$$

$$f_{!,\text{im}}([-2,-1] \cup [1,2]) = [1,4].$$

3. Circles. Consider the function $f: \mathbb{R}^2 \to \mathbb{R}$ given by

$$f(x,y) \stackrel{\text{def}}{=} x^2 + y^2$$

for each $(x, y) \in \mathbb{R}^2$. We have

$$f_{!,\mathsf{cp}}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}^2$, and since

$$f^{-1}(r) = \begin{cases} \text{a circle of radius } r \text{ about the origin} & \text{if } r > 0, \\ \{(0,0)\} & \text{if } r = 0, \\ \emptyset & \text{if } r < 0, \end{cases}$$

we have e.g.:

$$f_{!,\text{im}}([-1,1] \times [-1,1]) = [0,1],$$

$$f_{!,\text{im}}(([-1,1] \times [-1,1]) \setminus [-1,1] \times \{0\}) = \emptyset.$$

008R

PROPOSITION 4.6.6 ► PROPERTIES OF DIRECT IMAGES WITH COMPACT SUPPORT I

Let $f: A \to B$ be a function.

008S

1. Functoriality. The assignment $U \mapsto f_!(U)$ defines a functor

$$f_! \colon (\mathcal{P}(A), \subset) \to (\mathcal{P}(B), \subset)$$

where

· Action on Objects. For each $U \in \mathcal{P}(A)$, we have

$$[f_!](U) \stackrel{\text{def}}{=} f_!(U).$$

· Action on Morphisms. For each $U, V \in \mathcal{P}(A)$:

$$(\star)$$
 If $U \subset V$, then $f_!(U) \subset f_!(V)$.

008T

2. Triple Adjointness. We have a triple adjunction

$$(f_* \dashv f^{-1} \dashv f_!): \mathcal{P}(A) \leftarrow f^{-1} - \mathcal{P}(B),$$

witnessed by bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(B)}(f_*(U), V) \cong \operatorname{Hom}_{\mathcal{P}(A)}(U, f^{-1}(V)),$$

 $\operatorname{Hom}_{\mathcal{P}(A)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(A)}(U, f_!(V)),$

natural in $U \in \mathcal{P}(A)$ and $V \in \mathcal{P}(B)$ and (respectively) $V \in \mathcal{P}(A)$ and $U \in \mathcal{P}(B)$, i.e. where:

- (a) The following conditions are equivalent:
 - i. We have $f_*(U) \subset V$.
 - ii. We have $U \subset f^{-1}(V)$.
- (b) The following conditions are equivalent:
 - i. We have $f^{-1}(U) \subset V$.
 - ii. We have $U \subset f_!(V)$.

008U

3. Lax Preservation of Colimits. We have an inclusion of sets

$$\bigcup_{i\in I} f_!(U_i) \subset f_!\left(\bigcup_{i\in I} U_i\right),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(A)^{\times I}$. In particular, we have inclusions

$$f_!(U) \cup f_!(V) \hookrightarrow f_!(U \cup V),$$

 $\emptyset \hookrightarrow f_!(\emptyset),$

natural in $U, V \in \mathcal{P}(A)$.

008V

4. Preservation of Limits. We have an equality of sets

$$f_! \left(\bigcap_{i \in I} U_i \right) = \bigcap_{i \in I} f_! (U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(A)^{\times I}$. In particular, we have equalities

$$f^{-1}(U \cap V) = f_!(U) \cap f^{-1}(V),$$

$$f_!(A) = B,$$

natural in $U, V \in \mathcal{P}(A)$.

5. Symmetric Lax Monoidality With Respect to Unions. The direct image with compact support function of Item1 has a symmetric lax monoidal structure

$$\left(f_!, f_!^{\otimes}, f_{!|1}^{\otimes}\right) \colon (\mathcal{P}(A), \cup, \emptyset) \to (\mathcal{P}(B), \cup, \emptyset),$$

being equipped with inclusions

$$f_{!|U,V}^{\otimes} \colon f_{!}(U) \cup f_{!}(V) \hookrightarrow f_{!}(U \cup V),$$
$$f_{!|\mathfrak{1}}^{\otimes} \colon \emptyset \hookrightarrow f_{!}(\emptyset),$$

natural in $U, V \in \mathcal{P}(A)$.

6. Symmetric Strict Monoidality With Respect to Intersections. The direct image function of Item 1 has a symmetric strict monoidal structure

$$\left(f_!, f_!^{\otimes}, f_{!|1}^{\otimes}\right) \colon (\mathcal{P}(A), \cap, A) \to (\mathcal{P}(B), \cap, B),$$

being equipped with equalities

$$f_{!|U,V}^{\otimes} : f_{!}(U \cap V) \xrightarrow{=} f_{!}(U) \cap f_{!}(V),$$
$$f_{!|\mathbb{1}}^{\otimes} : f_{!}(A) \xrightarrow{=} B,$$

natural in $U, V \in \mathcal{P}(A)$.

7. Interaction With Coproducts. Let $f:A\to A'$ and $g:B\to B'$ be maps of sets. We have

$$(f \coprod g)_!(U \coprod V) = f_!(U) \coprod g_!(V)$$

for each $U \in \mathcal{P}(A)$ and each $V \in \mathcal{P}(B)$.

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8. Interaction With Products. Let $f:A\to A'$ and $g:B\to B'$ be maps of sets. We have

$$(f \times g)_!(U \times V) = f_!(U) \times g_!(V)$$

for each $U \in \mathcal{P}(A)$ and each $V \in \mathcal{P}(B)$.

0090

9. Relation to Direct Images. We have

$$f_!(U) = B \setminus f_*(A \setminus U)$$

for each $U \in \mathcal{P}(A)$.

0091

10. Interaction With Injections. If f is injective, then we have

$$\begin{split} f_{!,\text{im}}(U) &= f_*(U), \\ f_{!,\text{cp}}(U) &= B \setminus \text{Im}(f), \\ f_!(U) &= f_{!,\text{im}}(U) \cup f_{!,\text{cp}}(U) \\ &= f_*(U) \cup (B \setminus \text{Im}(f)) \end{split}$$

for each $U \in \mathcal{P}(A)$.

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11. Interaction With Surjections. If f is surjective, then we have

$$\begin{split} f_{!,\text{im}}(U) &\subset f_*(U), \\ f_{!,\text{cp}}(U) &= \emptyset, \\ f_!(U) &\subset f_*(U) \end{split}$$

for each $U \in \mathcal{P}(A)$.

PROOF 4.6.7 ► PROOF OF PROPOSITION 4.6.6

Item 1: Functoriality

Clear.

Item 2: Triple Adjointness

This follows from Remark 4.4.3, Remark 4.5.2, Remark 4.6.3, and ??, ?? of ??.

Item 3: Lax Preservation of Colimits

Omitted.

Item 4: Preservation of Limits

This follows from Item 2 and ??, ?? of ??.

Item 5: Symmetric Lax Monoidality With Respect to Unions

This follows from Item 3.

Item 6: Symmetric Strict Monoidality With Respect to Intersections

This follows from Item 4.

Item 7: Interaction With Coproducts

Clear.

Item 8: Interaction With Products

Clear.

Item 9: Relation to Direct Images

We claim that $f_!(U) = B \setminus f_*(A \setminus U)$.

· The First Implication. We claim that

$$f_!(U) \subset B \setminus f_*(A \setminus U).$$

Let $b \in f_!(U)$. We need to show that $b \notin f_*(A \setminus U)$, i.e. that there is no $a \in A \setminus U$ such that f(a) = b.

This is indeed the case, as otherwise we would have $a \in f^{-1}(b)$ and $a \notin U$, contradicting $f^{-1}(b) \subset U$ (which holds since $b \in f_!(U)$).

Thus $b \in B \setminus f_*(A \setminus U)$.

· The Second Implication. We claim that

$$B \setminus f_*(A \setminus U) \subset f_!(U)$$
.

Let $b \in B \setminus f_*(A \setminus U)$. We need to show that $b \in f_!(U)$, i.e. that $f^{-1}(b) \subset U$.

Since $b \notin f_*(A \setminus U)$, there exists no $a \in A \setminus U$ such that b = f(a), and hence $f^{-1}(b) \subset U$.

Thus $b \in f_!(U)$.

This finishes the proof of Item 9.

Item 10: Interaction With Injections

Clear.

Item 11: Interaction With Surjections

Clear.

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0093 PROPOSITION 4.6.8 ► PROPERTIES OF DIRECT IMAGES WITH COMPACT SUPPORT II

Let $f: A \to B$ be a function.

1. Functionality I. The assignment $f \mapsto f_!$ defines a function

$$(-)_{!|A,B}$$
: Sets $(A,B) \to \text{Sets}(\mathcal{P}(A),\mathcal{P}(B))$.

2. Functionality II. The assignment $f \mapsto f_!$ defines a function

$$(-)_{1|A|B}$$
: Sets $(A, B) \to \mathsf{Pos}((\mathcal{P}(A), \subset), (\mathcal{P}(B), \subset))$.

3. Interaction With Identities. For each $A \in Obj(Sets)$, we have

$$(id_A)_! = id_{\mathcal{P}(A)}.$$

4. Interaction With Composition. For each pair of composable functions $f:A\to B$ and $g:B\to C$, we have

$$(g \circ f)_{!} = g_{!} \circ f_{!}, \qquad \mathcal{P}(A) \xrightarrow{f_{!}} \mathcal{P}(B)$$

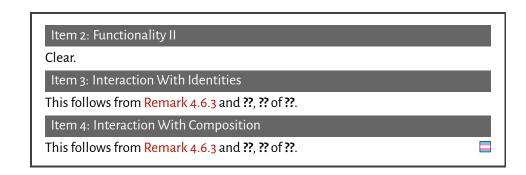
$$\downarrow g_{!}$$

$$\mathcal{P}(C)$$

PROOF 4.6.9 ► PROOF OF PROPOSITION 4.6.8

Item 1: Functionality I

Clear.



Appendices

A Other Chapters

Sets

- 1. Sets
- 2. Constructions With Sets
- 3. Pointed Sets
- 4. Tensor Products of Pointed Sets
- Relations
 - 5. Relations

- 6. Constructions With Relations
- 7. Equivalence Relations and Apartness Relations

Category Theory

8. Categories

Bicategories

9. Types of Morphisms in Bicategories

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