

Fibred Sets

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This chapter contains a discussion of the un/straightening equivalence in the context of sets, as well as a general discussion of indexed and fibred sets. In particular, it contains:

1. A discussion of indexed sets (i.e. functors $K_{\text{disc}} \rightarrow \text{Sets}$ with K a set), constructions with them like dependent sums and dependent products, and their properties (???)
2. A discussion of fibred sets (i.e. maps of sets $X \rightarrow K$), constructions with them like dependent sums and dependent products, and their properties (???)
3. A discussion of the un/straightening equivalence for indexed and fibred sets (??).

Contents

1 Fibred Sets 00S1

1.1 Foundations 00S2

Let K be a set.

Definition 1.1.1.1. A K -**fibred set** is a pair (X, ϕ) consisting of¹

- *The Underlying Set.* A set X , called the **underlying set of** (X, ϕ) ;
- *The Fibration.* A map of sets $\phi: X \rightarrow K$.

¹*Further Terminology:* The **fibre of** (X, ϕ) **over** $x \in K$ is the set $\phi^{-1}(x)$ (also written ϕ_x) defined by

$$\phi^{-1}(x) \stackrel{\text{def}}{=} \text{pt} \times_{[x], K, \phi} X,$$

$$\begin{array}{ccc} \phi^{-1}(x) & \longrightarrow & X \\ \downarrow & \lrcorner & \downarrow \phi \\ \text{pt} & \xrightarrow{[x]} & K. \end{array}$$

1.2 Morphisms of Fibred Sets

Definition 1.2.1.1. A **morphism of K -fibred sets** from (X, ϕ) to (Y, ψ) is a function $f: X \rightarrow Y$ such that the diagram²

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \phi \searrow & & \swarrow \psi \\ & K & \end{array}$$

commutes.

1.3 The Category of Fibred Sets Over a Fixed Base

Definition 1.3.1.1. The **category of K -fibred sets** is the category $\text{FibSets}(K)$ defined as the slice category $\text{Sets}_{/K}$ of Sets over K :

$$\text{FibSets}(K) \stackrel{\text{def}}{=} \text{Sets}_{/K}.$$

Remark 1.3.1.2. In detail $\text{FibSets}(K)$ is a category where

- *Objects.* The objects of $\text{FibSets}(K)$ are pairs (X, ϕ) consisting of
 - *The Fibred Set.* A set X ;
 - *The Fibration.* A function $\phi: X \rightarrow K$;
- *Morphisms.* A morphism of $\text{FibSets}(K)$ from (X, ϕ) to (Y, ψ) is a function $f: X \rightarrow$

²*Further Terminology:* The **transport map associated to f at $x \in K$** is the function

$$f_x^*: \phi^{-1}(x) \rightarrow \psi^{-1}(x)$$

given by the dashed map in the diagram

$$\begin{array}{ccccc} \phi^{-1}(x) & \xrightarrow{\quad} & X & & \\ \downarrow & \dashrightarrow & \downarrow \phi & \searrow f & \\ \psi^{-1}(x) & \xrightarrow{\quad} & Y & & \\ \downarrow & \lrcorner & \downarrow \psi & & \\ \text{pt} & \xrightarrow{[x]} & K & & \\ \parallel & & \parallel & & \\ \text{pt} & \xrightarrow{[x]} & K. & & \end{array}$$

Y making the diagram

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ & \searrow \phi & \swarrow \psi \\ & X & \end{array}$$

commute;

- *Identities.* For each $(X, \phi) \in \text{Obj}(\text{FibSets}(K))$, the unit map

$$\mathbb{1}_{(X, \phi)}^{\text{FibSets}(K)} : \text{pt} \rightarrow \text{Hom}_{\text{FibSets}(K)}((X, \phi), (X, \phi))$$

of $\text{FibSets}(K)$ at (X, ϕ) is given by

$$\text{id}_{(X, \phi)}^{\text{FibSets}(K)} \stackrel{\text{def}}{=} \text{id}_X,$$

as witnessed by the commutativity of the diagram

$$\begin{array}{ccc} X & \xrightarrow{\text{id}_X} & X \\ & \searrow \phi & \swarrow \phi \\ & K & \end{array}$$

in Sets ;

- *Composition.* For each $\mathbf{X} = (X, \phi), \mathbf{Y} = (Y, \psi), \mathbf{Z} = (Z, \chi) \in \text{Obj}(\text{FibSets}(K))$, the composition map

$$\circ_{\mathbf{X}, \mathbf{Y}, \mathbf{Z}}^{\text{FibSets}(K)} : \text{Hom}_{\text{FibSets}(K)}(\mathbf{Y}, \mathbf{Z}) \times \text{Hom}_{\text{FibSets}(K)}(\mathbf{X}, \mathbf{Y}) \rightarrow \text{Hom}_{\text{FibSets}(K)}(\mathbf{X}, \mathbf{Z})$$

of $\text{FibSets}(K)$ at $(\mathbf{X}, \mathbf{Y}, \mathbf{Z})$ is defined by

$$\circ_{\mathbf{X}, \mathbf{Y}, \mathbf{Z}}^{\text{FibSets}(K)} \stackrel{\text{def}}{=} \circ_{X, Y, Z}^{\text{Sets}},$$

as witnessed by the commutativity of the diagram

$$\begin{array}{ccccc} X & \xrightarrow{f} & Y & \xrightarrow{g} & Z \\ & \searrow \phi & \downarrow \psi & \swarrow \chi & \\ & & K & & \end{array}$$

in Sets .

1.4 The Category of Fibred Sets

Definition 1.4.1.1. The **category of fibred sets** is the category $\mathbf{FibSets}$ defined as the Grothendieck construction of the functor $\mathbf{FibSets}: \mathbf{Sets}^{\text{op}} \rightarrow \mathbf{Cats}$ of ??:

$$\mathbf{FibSets} \stackrel{\text{def}}{=} \int^{\mathbf{Sets}} \mathbf{FibSets}.$$

Remark 1.4.1.2. In detail, the **category of fibred sets** is the category $\mathbf{FibSets}$ where

- *Objects.* The objects of $\mathbf{FibSets}$ are pairs $(K, (X, \phi_X))$ consisting of
 - *The Base Set.* A set K ;
 - *The Fibred Set.* A K -fibred set $\phi_X: X \rightarrow K$;
- *Morphisms.* A morphism of $\mathbf{FibSets}$ from $(K, (X, \phi_X))$ to $(K', (Y, \phi_Y))$ is a pair (ϕ, f) consisting of
 - *The Base Map.* A map of sets $\phi: K \rightarrow K'$;
 - *The Morphism of Fibred Sets.* A morphism of K -fibred sets

$$f: (X, \phi_X) \rightarrow \phi_Y^*(Y), \quad \begin{array}{ccc} X & \xrightarrow{f} & Y \times_{K'} K \\ & \searrow \phi_X & \swarrow \text{pr}_2 \\ & & K \end{array}$$

- *Identities.* For each $(K, X) \in \text{Obj}(\mathbf{FibSets})$, the unit map

$$\mu_{(K,X)}^{\mathbf{FibSets}}: \text{pt} \rightarrow \mathbf{FibSets}((K, X), (K, X))$$

of $\mathbf{FibSets}$ at (K, X) is defined by

$$\text{id}_{(K,X)}^{\mathbf{FibSets}} \stackrel{\text{def}}{=} (\text{id}_K, \sim),$$

where \sim is the isomorphism $X \rightarrow X \times_K K$ as in the diagram

$$\begin{array}{ccc} X & \xrightarrow{\sim} & X \times_K K \\ & \searrow \phi_X & \swarrow \text{pr}_2 \\ & & K \end{array}$$

- *Composition.* For each $\mathbf{X} = (K, X), \mathbf{Y} = (K', Y), \mathbf{Z} = (K'', Z) \in \text{Obj}(\mathbf{FibSets})$, the composition map

$$\circ_{\mathbf{X}, \mathbf{Y}, \mathbf{Z}}^{\mathbf{FibSets}}: \mathbf{FibSets}(\mathbf{Y}, \mathbf{Z}) \times \mathbf{FibSets}(\mathbf{X}, \mathbf{Y}) \rightarrow \mathbf{FibSets}(\mathbf{X}, \mathbf{Z})$$

of FibSets at $(\mathbf{X}, \mathbf{Y}, \mathbf{Z})$ is defined by

$$g \circ_{\mathbf{X}, \mathbf{Y}, \mathbf{Z}}^{\text{FibSets}} f \stackrel{\text{def}}{=} (g \times_{K'} \text{id}_K) \circ f$$

as in the diagram

$$\begin{array}{ccccc}
 X & \xrightarrow{f} & Y \times_{K'} K & \xrightarrow{g \times_{K'} \text{id}_K} & \overbrace{(Z \times_{K''} K') \times_{K'} K}^{\cong Z \times_{K''} K} \\
 & \searrow \phi_X & \downarrow \text{pr}_2 & \nearrow \text{pr}_2 & \\
 & & K & &
 \end{array}$$

for each $f \in \text{Obj}(\text{FibSets}(\mathbf{X}, \mathbf{Y}))$ and each $g \in \text{Obj}(\text{FibSets}(\mathbf{Y}, \mathbf{Z}))$.

2 Constructions With Fibred Sets

2.1 Change of Base

Let $f: K \rightarrow K'$ be a function and let (X, ϕ_X) be a K' -fibred set.

Definition 2.1.1.1. The **change of base of (X, ϕ_X) to K** is the K -fibred set $f^*(X)$ defined by

$$f^*(X) \stackrel{\text{def}}{=} (K \times_{K'} X, \text{pr}_1), \quad
 \begin{array}{ccc}
 f^*(X) & \xrightarrow{\text{pr}_2} & X \\
 \text{pr}_1 \downarrow & \lrcorner & \downarrow \phi_X \\
 K & \xrightarrow{f} & K'.
 \end{array}$$

Proposition 2.1.1.2. The assignment $X \mapsto f^*(X)$ defines a functor

$$f^*: \text{FibSets}(K') \rightarrow \text{FibSets}(K),$$

where

- *Action on Objects.* For each $(X, \phi_X) \in \text{Obj}(\text{FibSets}(K'))$, we have

$$f^*(X, \phi_X) \stackrel{\text{def}}{=} f^*(X);$$

- *Action on Morphisms.* For each $(X, \phi_X), (Y, \phi_Y) \in \text{Obj}(\text{FibSets}(K'))$, the action

on Hom-sets

$$f_{X,Y}^* : \text{Hom}_{\text{FibSets}(K')}(X, Y) \rightarrow \text{Hom}_{\text{FibSets}(K)}(f^*(X), f^*(Y))$$

of f^* at $((X, \phi_X), (Y, \phi_Y))$ is the map sending a morphism of K' -fibred sets

$$g : (X, \phi_X) \rightarrow (Y, \phi_Y)$$

to the morphism of K -fibred sets given by the dashed morphism in the diagram

$$\begin{array}{ccccc}
 f^*(X) & \xrightarrow{\quad} & X & & \\
 \downarrow & \searrow \lrcorner & \downarrow \phi_X & \searrow g & \\
 & & f^*(Y) & \xrightarrow{\quad} & Y \\
 & & \downarrow & & \downarrow \phi_Y \\
 K & \xrightarrow{f} & K' & & \\
 \parallel & & \parallel & & \\
 K & \xrightarrow{f} & K' & &
 \end{array}$$

Proof. Omitted. □

Proposition 2.1.1.3. The assignment $K \mapsto \text{FibSets}(K)$ defines a functor

$$\text{FibSets} : \text{Sets}^{\text{op}} \rightarrow \text{Cats},$$

where

- *Action on Objects.* For each $K \in \text{Obj}(\text{Sets})$, we have

$$[\text{FibSets}](K) \stackrel{\text{def}}{=} \text{FibSets}(K);$$

- *Action on Morphisms.* For each $K, K' \in \text{Obj}(\text{Sets})$, the action on Hom-sets

$$\text{Sets}_{/(-)|K,K'} : \text{Sets}^{\text{op}}(K, K') \rightarrow \text{Fun}(\text{FibSets}(K), \text{FibSets}(K'))$$

of $\text{Sets}_{/(-)}$ at (K, K') is the map sending a map of sets $f : K \rightarrow K'$ to the functor

$$\text{Sets}_{/f} : \text{FibSets}(K') \rightarrow \text{FibSets}(K)$$

defined by

$$\text{Sets}_{/f} \stackrel{\text{def}}{=} f^*.$$

Proof. Omitted. □

2.2 Dependent Sums

Let $f: K \rightarrow K'$ be a function and let (X, ϕ_X) be a K -fibred set.

Definition 2.2.1.1. The **dependent sum**³ of (X, ϕ_X) is the K' -fibred set $\Sigma_f(X)$ ⁴ defined by

$$\begin{aligned}\Sigma_f(X) &\stackrel{\text{def}}{=} (\Sigma_f(X), \Sigma_f(\phi_X)) \\ &\stackrel{\text{def}}{=} (X, f \circ \phi_X).\end{aligned}$$

Proposition 2.2.1.2. Let $f: K \rightarrow K'$ be a function.

1. *Functoriality.* The assignment $X \mapsto \Sigma_f(X)$ defines a functor

$$\Sigma_f: \text{FibSets}(K) \rightarrow \text{FibSets}(K'),$$

where

- *Action on Objects.* For each $(X, \phi_X) \in \text{Obj}(\text{FibSets}(K))$, we have

$$\Sigma_f(X, \phi_X) \stackrel{\text{def}}{=} (\Sigma_f(X), \Sigma_f(\phi_X));$$

- *Action on Morphisms.* For each $(X, \phi_X), (Y, \phi_Y) \in \text{Obj}(\text{FibSets}(K))$, the action on Hom-sets

$$\Sigma_f|_{X,Y}: \text{Hom}_{\text{FibSets}(K)}(X, Y) \rightarrow \text{Hom}_{\text{FibSets}(K')}(\Sigma_f(X), \Sigma_f(Y))$$

of Σ_f at $((X, \phi_X), (Y, \phi_Y))$ is the map sending a morphism of K -fibred sets

$$g: (X, \phi_X) \rightarrow (Y, \phi_Y)$$

to the morphism of K' -fibred sets defined by

$$\Sigma_f(g) \stackrel{\text{def}}{=} g.$$

2. *Interaction With Fibres.* We have a bijection of sets

$$\Sigma_f(\phi_X)^{-1}(k') \cong \coprod_{k \in f^{-1}(k')} \phi_X^{-1}(k)$$

for each $k' \in K'$.

³The name “dependent sum” comes from the fact that the fibre $\Sigma_f(\phi_X)^{-1}(x)$ of $\Sigma_f(X)$ at $x \in K'$ is given by

$$\Sigma_f(\phi_X)^{-1}(x) \cong \coprod_{y \in f^{-1}(x)} \phi_X^{-1}(y);$$

see ?? of ??.

⁴*Further Notation:* Also written $f_*(X)$.

Proof. ??, Functoriality: Omitted.

??, Interaction With Fibres: Indeed, we have

$$\begin{aligned}\Sigma_f(\phi_X)^{-1}(k') &\stackrel{\text{def}}{=} \mathbf{pt} \times_{[k'], K', f \circ \phi_X} X \\ &\cong \{x \in X \mid f(\phi_X(x)) = k'\} \\ &\cong \coprod_{k \in f^{-1}(k')} \{x \in X \mid \phi_X(x) = k\} \\ &\cong \coprod_{k \in f^{-1}(k')} \phi_X^{-1}(k)\end{aligned}$$

for each $k' \in K'$. □

2.3 Dependent Products

Let $f: K \rightarrow K'$ be a function and let (X, ϕ_X) be a K -fibred set.

Definition 2.3.1.1. The **dependent product**⁶ of (X, ϕ_X) is the K' -fibred set $\Pi_f(X)$ ⁶ consisting of⁷

- *The Underlying Set.* The set $\Pi_f(X)$ defined by

$$\Pi_f(X) \stackrel{\text{def}}{=} \coprod_{k' \in K'} \prod_{k \in f^{-1}(k')} \phi_X^{-1}(k);$$

- *The Fibration.* The map of sets

$$\Pi_f(\phi_X): \Pi_f(X) \rightarrow K'$$

defined by sending an element of

$$\Pi_f(X) \stackrel{\text{def}}{=} \coprod_{k' \in K'} \prod_{k \in f^{-1}(k')} \phi_X^{-1}(k)$$

to its index k' in K' .

Example 2.3.1.2. Here are some examples of dependent products of sets.

⁵The name “dependent product” comes from the fact that the fibre $\Pi_f(\phi_X)^{-1}(k')$ of $\Pi_f(X)$ at $k' \in K'$ is given by

$$\Pi_f(\phi_X)^{-1}(k') \cong \prod_{k \in f^{-1}(k')} \phi_X^{-1}(k);$$

see ?? of ??.

⁶*Further Notation:* Also written $f_!(X)$.

⁷We can also define dependent products via the internal **Hom** in $\mathbf{FibSets}(K')$; see ?? of ??.

1. *Spaces of Sections.* Let $K = X, K' = \text{pt}$, let $\phi: E \rightarrow X$ be a map of sets, and write $!_X: X \rightarrow \text{pt}$ for the terminal map from X to pt . We have a bijection of sets

$$\begin{aligned} \Pi_{!_X}((E, \phi)) &\cong \Gamma_X(\phi) \\ &\stackrel{\text{def}}{=} \{h \in \text{Sets}(X, E) \mid \phi \circ h = \text{id}_X\}. \end{aligned}$$

2. *Function Spaces.* Let $K = K' = \text{pt}$ and write $!_X: X \rightarrow \text{pt}$ and $!_Y: Y \rightarrow \text{pt}$ for the terminal maps from X and Y to pt . We have a bijection of sets

$$\text{Sets}(X, Y) \cong \Pi_{!_X}(!_X^*(Y, !_Y)).$$

Proof. ??, Spaces of Sections: Indeed, we have

$$\begin{aligned} \Pi_{!_X}((E, \phi)) &\stackrel{\text{def}}{=} \coprod_{\star \in \text{pt}} \prod_{k \in !_X^{-1}(\star)} \phi_X^{-1}(k) \\ &= \prod_{x \in X} \phi_X^{-1}(x) \\ &\cong \{h \in \text{Sets}(X, E) \mid \phi_X \circ h = \text{id}_X\} \\ &\stackrel{\text{def}}{=} \Gamma_X(\phi). \end{aligned}$$

??, Function Spaces: Indeed, we have

$$\begin{aligned} \Pi_{!_X}(!_X^*(Y, !_Y)) &\stackrel{\text{def}}{=} \Pi_{!_X}(X \times_{!_X, \text{pt}, !_Y} Y) \\ &\stackrel{\text{def}}{=} \coprod_{\star \in \text{pt}} \prod_{x \in !_X^{-1}(\star)} \text{pr}_1^{-1}(x) \\ &= \prod_{x \in X} Y \\ &\cong \text{Sets}(X, Y). \end{aligned}$$

This finishes the proof. □

Proposition 2.3.1.3. Let $f: K \rightarrow K'$ be a function.

1. *Functoriality.* The assignment $X \mapsto \Pi_f(X)$ defines a functor

$$\Pi_f: \text{FibSets}(K) \rightarrow \text{FibSets}(K'),$$

where

- *Action on Objects.* For each $(X, \phi_X) \in \text{Obj}(\text{FibSets}(K))$, we have

$$\Pi_f(X, \phi_X) \stackrel{\text{def}}{=} \Pi_f(X);$$

- *Action on Morphisms.* For each $(X, \phi_X), (Y, \phi_Y) \in \text{Obj}(\text{FibSets}(K))$, the action on Hom-sets

$$\Pi_f|_{X,Y}: \text{Hom}_{\text{FibSets}(K)}(X, Y) \rightarrow \text{Hom}_{\text{FibSets}(K')}(\Pi_f(X), \Pi_f(Y))$$

of Π_f at $((X, \phi_X), (Y, \phi_Y))$ is the map sending a morphism of K -fibred sets

$$\xi: (X, \phi_X) \rightarrow (Y, \phi_Y), \quad \begin{array}{ccc} X & \xrightarrow{\xi} & Y \\ \phi_X \searrow & & \swarrow \phi_Y \\ & K & \end{array}$$

to the morphism

$$\Pi_f(\xi): (\Pi_f(X), \Pi_f(\phi_X)) \rightarrow (\Pi_f(Y), \Pi_f(\phi_Y)) \quad \begin{array}{ccc} \Pi_f(X) & \xrightarrow{\Pi_f(\xi)} & \Pi_f(Y) \\ \Pi_f(\phi_X) \searrow & & \swarrow \Pi_f(\phi_Y) \\ & K & \end{array}$$

of K' -fibred sets given by⁸

$$[\Pi_f(\xi)]((x_k)_{k \in f^{-1}(k')}) \stackrel{\text{def}}{=} (\xi(x_k))_{k \in f^{-1}(k')}$$

for each $(x_k)_{k \in f^{-1}(k')} \in \Pi_f(X)$.

2. *Interaction With Fibres.* We have a bijection of sets

$$\Pi_f(\phi_X)^{-1}(k') \cong \prod_{k \in f^{-1}(k')} \phi_X^{-1}(k);$$

for each $k' \in K'$.

3. *Construction Using the Internal Hom.* We have

$$\Pi_f(X, \phi_X) = (K' \times_{\text{Hom}_{\text{FibSets}(K')}((K, f), (K, f))} \text{Hom}_{\text{FibSets}(K')}((K, f), (X, f \circ \phi_X)), \text{pr}_1),$$

⁸Note that we indeed have $\xi(x_k) \in \phi_Y^{-1}(k)$, since

$$\begin{aligned} \phi_Y(\xi(x_k)) &= [\phi_Y \circ \xi](x_k) \\ &= \phi_X(x_k) \\ &= k, \end{aligned}$$

where we have used that ξ is a morphism of K -fibred sets for the second equality.

forming a pullback diagram

$$\begin{array}{ccc}
 \Pi_f(X, \phi_X) & \xrightarrow{\text{pr}_2} & \mathbf{Hom}_{\mathbf{FibSets}(K')}((K, f), (X, f \circ \phi_X)) \\
 \text{pr}_1 \downarrow & \lrcorner & \downarrow (\phi_X)_* \\
 K' & \xrightarrow{I} & \mathbf{Hom}_{\mathbf{FibSets}(K')}((K, f), (K, f)),
 \end{array}$$

where the bottom map is given by

$$I(k') \stackrel{\text{def}}{=} \text{id}_{f^{-1}(k')}$$

for each $k' \in K'$ and where $\mathbf{Hom}_{\mathbf{FibSets}(K')}$ denotes the internal Hom of $\mathbf{FibSets}(K')$ of ??.

4. *Internal Homs via Dependent Products.* We have

00SX

$$\mathbf{Hom}_{\mathbf{FibSets}(K)}(X, Y) \cong \Pi_{\phi_X}(\phi_X^*(Y)).$$

Proof. ??, Functoriality: Omitted.

??, *Interaction With Fibres:* Clear.

??, *Construction Using the Internal Hom:* Using the explicit formula for pullbacks of sets given in **Constructions With Sets, Definition 1.3.1.1**, we see that the pullback

$$K' \times_{\mathbf{Hom}_{\mathbf{FibSets}(K')}((K, f), (K, f))} \mathbf{Hom}_{\mathbf{FibSets}(K')}((K, f), (X, f \circ \phi_X))$$

is given by

$$\left\{ (k', h) \in \coprod_{k' \in K'} \text{Sets}(f^{-1}(k'), \phi_X^{-1}(f^{-1}(k'))) \mid \phi_X \circ h = \text{id}_{f^{-1}(k')} \right\},$$

which is isomorphic to

$$\coprod_{k' \in K'} \{h \in \text{Sets}(f^{-1}(k'), \phi_X^{-1}(f^{-1}(k'))) \mid \phi_X \circ h = \text{id}_{f^{-1}(k')}\}.$$

We claim that

$$\{h \in \text{Sets}(f^{-1}(k'), \phi_X^{-1}(f^{-1}(k'))) \mid \phi_X \circ h = \text{id}_{f^{-1}(k')}\} \cong \prod_{k \in f^{-1}(k')} \phi_X^{-1}(k),$$

so that the pullback is indeed given by $\Pi_f(X)$. There are two cases:

1. If $f^{-1}(k') = \emptyset$, then there is only one map from $f^{-1}(k')$ to $\phi_X^{-1}(f^{-1}(k'))$ (the inclusion), so $\text{Sets}(f^{-1}(k'), \phi_X^{-1}(f^{-1}(k'))) \cong \text{pt}$. Since products indexed by the empty set are isomorphic to pt , the isomorphism follows.
2. Otherwise, by the condition $\phi_X \circ h = \text{id}_{f^{-1}(k')}$, it follows that, for each $k \in f^{-1}(k')$, we must have

$$\phi_X(h(k)) = k,$$

and thus $h(k) \in \phi_X^{-1}(k)$. Therefore, a map from $f^{-1}(k')$ to $\phi_X^{-1}(f^{-1}(k'))$ consists of a choice of an element from $\phi_X^{-1}(k)$ for each $k \in f^{-1}(k')$, which is precisely given by an element of the product $\prod_{k \in f^{-1}(k')} \phi_X^{-1}(k)$, showing the bijection to be true.

??, *Internal Homs via Dependent Products*: Indeed we have

$$\begin{aligned}
 \Pi_{\phi_X}(\phi_X^*(Y)) &\stackrel{\text{def}}{=} \Pi_{\phi_X}(X \times_K Y) \\
 &\stackrel{\text{def}}{=} \coprod_{k \in K} \prod_{x \in \phi_X^{-1}(k)} \text{pr}_1^{-1}(x) \\
 &\cong \coprod_{k \in K} \prod_{x \in \phi_X^{-1}(k)} \{y \in Y \mid \phi_X(x) = \phi_Y(y)\} \\
 &\cong \coprod_{k \in K} \prod_{x \in \phi_X^{-1}(k)} \{y \in Y \mid k = \phi_Y(y)\} \\
 &\cong \coprod_{k \in K} \prod_{x \in \phi_X^{-1}(k)} \phi_Y^{-1}(k) \\
 &\cong \coprod_{k \in K} \text{Sets}(\phi_X^{-1}(k), \phi_Y^{-1}(k)) \\
 &\stackrel{\text{def}}{=} \mathbf{Hom}_{\text{FibSets}(K)}(X, Y).
 \end{aligned}$$

This finishes the proof. □

2.4 Internal Homs 00SY

Let K be a set and let (X, ϕ_X) and (Y, ϕ_Y) be K -fibred sets.

Definition 2.4.1.1. The **internal Hom of fibred sets from (X, ϕ_X) to (Y, ϕ_Y)** is the fibred set $\mathbf{Hom}_{\text{FibSets}(K)}(X, Y)$ consisting of

- *The Underlying Set.* The set $\mathbf{Hom}_{\text{FibSets}(K)}(X, Y)$ defined by

$$\mathbf{Hom}_{\text{FibSets}(K)}(X, Y) \stackrel{\text{def}}{=} \coprod_{k \in K} \text{Sets}(\phi_X^{-1}(k), \phi_Y^{-1}(k));$$

- *The Fibration.* The map of sets⁹

$$\phi_{\mathbf{Hom}_{\mathbf{FibSets}(K)}(X,Y)} : \underbrace{\mathbf{Hom}_{\mathbf{FibSets}(K)}(X,Y)}_{\coprod_{k \in K} \mathbf{Sets}(\phi_X^{-1}(k), \phi_Y^{-1}(k))} \rightarrow K$$

defined by sending a map $f: \phi_X^{-1}(k) \rightarrow \phi_Y^{-1}(k)$ to its index $k \in K$.

Proof. Omitted. □

Proposition 2.4.1.2. Let K be a set and let (X, ϕ_X) and (Y, ϕ_Y) be K -fibred sets.

1. *Functoriality.* Let (X, ϕ_X) and (Y, ϕ_Y) be K -fibred sets.

- (a) The assignment $X \mapsto \mathbf{Hom}_{\mathbf{FibSets}(K)}(X, Y)$ defines a functor

$$\mathbf{Hom}_{\mathbf{FibSets}(K)}(X, -) : \mathbf{FibSets}(K) \rightarrow \mathbf{FibSets}(K).$$

- (b) The assignment $Y \mapsto \mathbf{Hom}_{\mathbf{FibSets}(K)}(X, Y)$ defines a functor

$$\mathbf{Hom}_{\mathbf{FibSets}(K)}(-, Y) : \mathbf{FibSets}(K)^{\text{op}} \rightarrow \mathbf{FibSets}(K).$$

- (c) The assignment $(X, Y) \mapsto \mathbf{Hom}_{\mathbf{FibSets}(K)}(X, Y)$ defines a functor

$$\mathbf{Hom}_{\mathbf{FibSets}(K)}(-, -) : \mathbf{FibSets}(K)^{\text{op}} \times \mathbf{FibSets}(K) \rightarrow \mathbf{FibSets}(K).$$

2. *Internal Homs via Dependent Products.* We have 00T2

$$\mathbf{Hom}_{\mathbf{FibSets}(K)}(X, Y) \cong \Pi_{\phi_X}(\phi_X^*(Y)).$$

Proof. ??, *Functoriality:* Omitted.

??, *Internal Homs via Dependent Products:* This was proved in ?? of ??. □

2.5 Adjointness for Fibred Sets

Let $f: K \rightarrow K'$ be a map of sets.

⁹The fibres of the internal \mathbf{Hom} of $\mathbf{FibSets}(K)$ are precisely the sets $\mathbf{Sets}(\phi_X^{-1}(k), \phi_Y^{-1}(k))$, i.e. we have

$$\phi_{\mathbf{Hom}_{\mathbf{FibSets}(K)}(X,Y)}|_k \cong \mathbf{Sets}(\phi_X^{-1}(k), \phi_Y^{-1}(k))$$

for each $k \in K$.

Proposition 2.5.1.1. We have a triple adjunction [2.5.1](#)

$$(\Sigma_f \dashv f^* \dashv \Pi_f): \text{FibSets}(K) \begin{array}{c} \xrightarrow{\Sigma_f} \\ \perp \\ \xleftarrow{\Pi_f} \end{array} f^* \text{FibSets}(K').$$

We offer two proofs. The first uses the corresponding adjunction for indexed sets ([Indexed Sets, Proposition 4.5.1.1](#)) and the un/straightening equivalence together with its compatibility with dependent sums and products to “transfer” the adjunction to fibred sets, while the second is a direct proof.

Proof. The Adjunction $\Sigma_f \dashv f^*$: The adjunction

$$(\Sigma_f \dashv f^*): \text{ISets}(K) \begin{array}{c} \xrightarrow{\Sigma_f} \\ \perp \\ \xleftarrow{f^*} \end{array} \text{ISets}(K')$$

of [Indexed Sets, Proposition 4.5.1.1](#) gives a unit and counit of the form

$$\begin{aligned} \eta: \text{id}_{\text{ISets}(K)} &\Longrightarrow \Sigma_f \circ f^*, \\ \epsilon: f^* \circ \Sigma_f &\Longrightarrow \text{id}_{\text{ISets}(K')}. \end{aligned}$$

With these in hand, we construct natural transformations

$$\begin{aligned} \eta': \text{id}_{\text{FibSets}(K)} &\Longrightarrow \Sigma_f \circ f^*, \\ \epsilon': f^* \circ \Sigma_f &\Longrightarrow \text{id}_{\text{FibSets}(K')}. \end{aligned}$$

as follows:

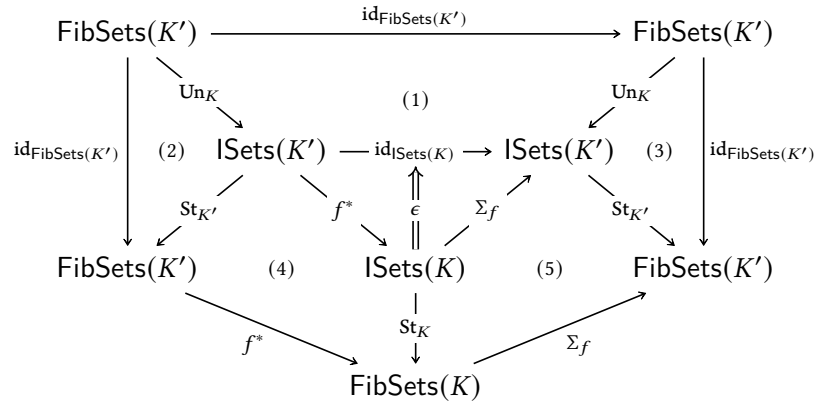
1. *The Unit.* We define $\eta': \text{id}_{\text{FibSets}(K)} \Longrightarrow \Sigma_f \circ f^*$ as the pasting of the diagram

$$\begin{array}{ccccc}
 & & \text{FibSets}(K') & & \\
 & \nearrow \Sigma_f & \uparrow \text{St}_{K'} & \nwarrow f^* & \\
 \text{FibSets}(K) & (1) & \text{ISets}(K') & (2) & \text{FibSets}(K) \\
 \uparrow \text{id}_{\text{FibSets}(K)} & \nwarrow \text{St}_K & \nearrow \Sigma_f & \nwarrow f^* & \nearrow \text{St}_K \\
 & (3) & \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) & (4) \\
 & \nearrow \text{Un}_K & & & \nwarrow \text{Un}_K & \\
 \text{FibSets}(K) & & & & & \text{FibSets}(K) \\
 & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & & & &
 \end{array}$$

where:

- (a) Subdiagram (1) commutes by Un/Straightening for Indexed and Fibred Sets, ?? of ??.
- (b) Subdiagram (2) commutes by Un/Straightening for Indexed and Fibred Sets, ?? of ??.
- (c) Subdiagram (3) commutes by Un/Straightening for Indexed and Fibred Sets, ??.
- (d) Subdiagram (4) commutes by Un/Straightening for Indexed and Fibred Sets, ??.
- (e) Subdiagram (5) commutes by unitality of composition.

2. *The Counit.* We define $\epsilon' : f^* \circ \Sigma_f \Rightarrow \text{id}_{\text{FibSets}(K')}$ as the pasting of the diagram



where:

- (a) Subdiagram (1) commutes by unitality of composition.
- (b) Subdiagram (2) commutes by Un/Straightening for Indexed and Fibred Sets, ??.
- (c) Subdiagram (3) commutes by Un/Straightening for Indexed and Fibred Sets, ??.
- (d) Subdiagram (4) commutes by Un/Straightening for Indexed and Fibred Sets, ?? of ??.
- (e) Subdiagram (5) commutes by Un/Straightening for Indexed and Fibred Sets, ?? of ??.

Next, we prove the left triangle identity,

$$\begin{array}{ccc}
 \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') \\
 \Sigma_f \nearrow & \uparrow \eta & \downarrow f^* \\
 \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K) \\
 & \uparrow \epsilon & \downarrow \Sigma_f
 \end{array}
 =
 \begin{array}{ccc}
 \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') \\
 \Sigma_f \nearrow & \uparrow \text{id}_{\Sigma_f} & \downarrow \Sigma_f \\
 \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K)
 \end{array}$$

whose left side in our case looks like this:

$$\begin{array}{ccccc}
 & & \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') \\
 & \Sigma_f \nearrow & \uparrow \text{St}_{K'} & \downarrow f^* & \downarrow \text{Un}_{K'} & \downarrow \text{Un}_{K'} & \downarrow \text{id}_{\text{FibSets}(K')} \\
 \text{FibSets}(K) & & \text{ISets}(K') & & \text{ISets}(K') & \xrightarrow{\text{id}_{\text{ISets}(K')}} & \text{ISets}(K') \\
 \uparrow \text{id}_{\text{FibSets}(K)} & \uparrow \text{St}_K & \uparrow \Sigma_f & \uparrow \eta & \uparrow f^* & \uparrow \Sigma_f & \uparrow \text{St}_{K'} \\
 \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) \\
 \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K \\
 \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K)
 \end{array}$$

It can be rearranged into

$$\begin{array}{ccccc}
 & & \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') \\
 & \Sigma_f \nearrow & \uparrow \text{St}_{K'} & \downarrow f^* & \downarrow \text{Un}_{K'} & \downarrow \text{Un}_{K'} & \downarrow \text{id}_{\text{FibSets}(K')} \\
 \text{FibSets}(K) & & \text{ISets}(K') & \xrightarrow{\text{id}_{\text{ISets}(K')}} & \text{ISets}(K') & \xrightarrow{\text{id}_{\text{ISets}(K')}} & \text{ISets}(K') \\
 \uparrow \text{id}_{\text{FibSets}(K)} & \uparrow \text{St}_K & \uparrow \Sigma_f & \uparrow \eta & \uparrow f^* & \uparrow \Sigma_f & \uparrow \text{St}_{K'} \\
 \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) \\
 \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K \\
 \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K)
 \end{array}$$

(1) (2) (3)

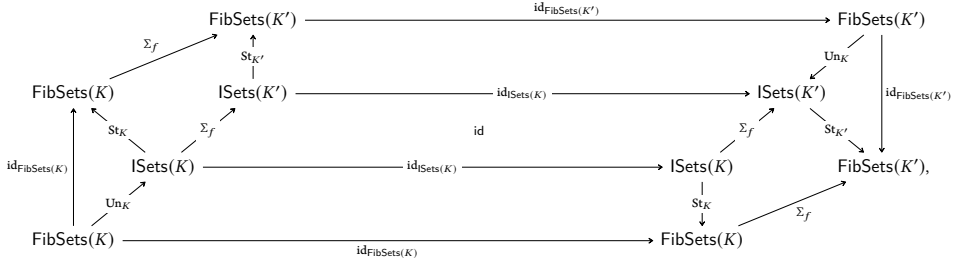
where:

1. Subdiagram (1) commutes by Un/Straightening for Indexed and Fibred Sets, ??.
2. Subdiagram (2) commutes by unitality of composition.
3. Subdiagram (3) commutes by Un/Straightening for Indexed and Fibred Sets, ??.

And then, it can be rearranged into

$$\begin{array}{ccccc}
 & & \text{FibSets}(K') & \xrightarrow{\text{id}_{\text{FibSets}(K')}} & \text{FibSets}(K') \\
 & \Sigma_f \nearrow & \uparrow \text{St}_{K'} & \downarrow f^* & \downarrow \text{Un}_{K'} \\
 \text{FibSets}(K) & & \text{ISets}(K') & \xrightarrow{\text{id}_{\text{ISets}(K')}} & \text{ISets}(K') \\
 \uparrow \text{id}_{\text{FibSets}(K)} & \uparrow \text{St}_K & \uparrow \Sigma_f & \uparrow \eta & \uparrow f^* \\
 \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) & \xrightarrow{\text{id}_{\text{ISets}(K)}} & \text{ISets}(K) \\
 \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K & \downarrow \text{Un}_K \\
 \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K) & \xrightarrow{\text{id}_{\text{FibSets}(K)}} & \text{FibSets}(K)
 \end{array}$$

which by the left triangle identity for (η, ϵ) , becomes



finishing the proof of the left triangle identity. The proof of the right triangle identity is similar, and is thus omitted.

The Adjunction $f^ \dashv \Pi_f$:* This proof is similar to the proof of the adjunction $\Sigma_f \dashv f^*$, and is thus omitted. \square

We proceed to the direct proof of ??.

Proof. The Adjunction $\Sigma_f \dashv f^$:* We claim there's a bijection

$$\text{Hom}_{\text{FibSets}(K')}(\Sigma_f(X), Y) \cong \text{Hom}_{\text{FibSets}(K)}(X, f^*(Y)),$$

natural in $(X, \phi_X) \in \text{Obj}(\text{FibSets}(K))$ and $(Y, \phi_Y) \in \text{Obj}(\text{FibSets}(K'))$:

- *Map I.* We define a map

$$\Phi_{X,Y} : \text{Hom}_{\text{FibSets}(K')}(\Sigma_f(X), Y) \rightarrow \text{Hom}_{\text{FibSets}(K)}(X, f^*(Y)),$$

by sending a morphism

$$\xi : \Sigma_f(X) \rightarrow Y, \quad \begin{array}{ccc} X & \xrightarrow{\xi} & Y \\ \phi_X \searrow & & \nearrow \phi_Y \\ & K & \\ & f \searrow & \\ & & K' \end{array}$$

of K' -fibred sets to the morphism

$$\xi^\dagger : X \rightarrow f^*(Y), \quad \begin{array}{ccc} X & \xrightarrow{\xi^\dagger} & K \times_{K'} Y \\ \phi_X \searrow & & \nearrow \text{pr}_1 \\ & & K' \end{array}$$

of K -fibred sets given by the dashed morphism in the diagram

$$\begin{array}{ccccc}
 X & & & & \\
 \downarrow \phi_X & \searrow \xi & & & \\
 & K \times_{K'} Y & \xrightarrow{\text{pr}_2} & Y & \\
 & \downarrow \text{pr}_1 & \lrcorner & \downarrow \phi_Y & \\
 & K & \xrightarrow{f} & K' &
 \end{array}$$

- *Map II.* We define a map

$$\Psi_{X,Y} : \text{Hom}_{\text{FibSets}(K)}(X, f^*(Y)) \rightarrow \text{Hom}_{\text{FibSets}(K')}(\Sigma_f(X), Y),$$

given by sending a map

$$\xi : X \rightarrow f^*(Y), \quad \begin{array}{ccc} X & \xrightarrow{\xi} & K \times_{K'} Y \\ \phi_X \searrow & & \swarrow \text{pr}_1 \\ & K' & \end{array}$$

of K' -fibred sets to the map

$$\xi^\dagger : \Sigma_f(X) \rightarrow Y, \quad \begin{array}{ccc} X & \xrightarrow{\xi^\dagger} & Y \\ \phi_X \searrow & & \swarrow \phi_Y \\ & K & \\ & \searrow f & \\ & K' & \end{array}$$

of K -fibred sets given by

$$\xi^\dagger \stackrel{\text{def}}{=} \text{pr}_2 \circ \xi,$$

where we indeed have

$$\begin{aligned}
 \phi_Y \circ (\text{pr}_2 \circ \xi) &= (\phi_Y \circ \text{pr}_2) \circ \xi \\
 &= (f \circ \text{pr}_1) \circ \xi && \text{(by the pullback square of } K \times_{K'} Y) \\
 &= f \circ (\text{pr}_1 \circ \xi) \\
 &= f \circ \phi_X. && \text{(since } \xi \text{ is a morphism of } K' \text{-fibred sets)}
 \end{aligned}$$

- *Naturality I.* We need to show that, given a morphism

$$\alpha: (X, \phi_X) \rightarrow (X', \phi_{X'})$$

of K -fibred sets, the diagram

$$\begin{array}{ccc} \mathrm{Hom}_{\mathrm{FibSets}(K')}(\Sigma_f(X'), Y) & \xrightarrow{\Phi_{X',Y}} & \mathrm{Hom}_{\mathrm{FibSets}(K)}(X', f^*(Y)), \\ \downarrow \Sigma_f(\alpha)^* & & \downarrow \alpha^* \\ \mathrm{Hom}_{\mathrm{FibSets}(K')}(\Sigma_f(X), Y) & \xrightarrow{\Phi_{X,Y}} & \mathrm{Hom}_{\mathrm{FibSets}(K)}(X, f^*(Y)) \end{array}$$

commutes. Indeed, given a morphism

$$\xi: \Sigma_f(X') \rightarrow Y, \quad \begin{array}{ccc} X' & \xrightarrow{\xi} & Y \\ \phi_{X'} \downarrow & & \downarrow \phi_Y \\ K & \xrightarrow{f} & K' \end{array}$$

of K' -fibred-sets, the map $\Phi_{X',Y}(\xi) \circ \alpha$ is the composition, coloured in **vermillion**, of the dashed arrow with α in the diagram

while $\Phi_{X,Y}(\xi \circ \Sigma_f(\alpha))$ is given by the dashed arrow, coloured in **blue**. Since both the **blue arrow** and the **vermillion arrow** make the outer pullback diagram commute, it follows from the universal property of the pullback that they must be equal, i.e. that

$$\Phi_{X',Y}(\xi) \circ \alpha = \Phi_{X,Y}(\xi \circ \Sigma_f(\alpha)),$$

showing that the naturality diagram above indeed commutes.

- $$\beta: (Y, \phi_Y) \rightarrow (Y', \phi_{Y'})$$

of K -fibred sets, the diagram

$$\begin{array}{ccc} \mathrm{Hom}_{\mathrm{FibSets}(K')}(\Sigma_f(X), Y) & \xrightarrow{\Phi_{X,Y}} & \mathrm{Hom}_{\mathrm{FibSets}(K)}(X, f^*(Y)), \\ \beta_* \downarrow & & \downarrow f^*(\beta)_* \\ \mathrm{Hom}_{\mathrm{FibSets}(K')}(\Sigma_f(X), Y') & \xrightarrow{\Phi_{X,Y'}} & \mathrm{Hom}_{\mathrm{FibSets}(K)}(X, f^*(Y')) \end{array}$$

commutes. Indeed, given a morphism

$$\xi: \Sigma_f(X') \rightarrow Y,$$

$$\begin{array}{ccc} X' & \xrightarrow{\xi} & Y \\ \phi_{X'} \searrow & & \nearrow \phi_Y \\ & K & \\ & \searrow f & \\ & & K' \end{array}$$

of K' -fibred-sets, the map $f^*(\beta) \circ \Phi_{X,Y}(\xi)$ is the composition, coloured in **vermillion**, of the dashed arrow from X to $K \times_{K'} Y$ with the dashed arrow from $K \times_{K'} Y$ to $K \times_{K'} Y'$ in the diagram

The diagram shows a commutative diagram with the following components:

- Objects:** X , $K \times_{K'} Y$, Y , K , $K \times_{K'} Y'$, Y' , K , K' .
- Maps:**
 - $\phi_X: X \rightarrow K$ (solid black arrow)
 - $\xi: X \rightarrow Y$ (solid black arrow)
 - $\beta: Y \rightarrow Y'$ (solid black arrow)
 - $\phi_Y: Y \rightarrow K$ (solid black arrow)
 - $\phi_{Y'}: Y' \rightarrow K'$ (solid black arrow)
 - $f: K \rightarrow K$ (solid black arrow)
 - $f: K \rightarrow K'$ (solid black arrow)
 - $\eta: K \times_{K'} Y \rightarrow K \times_{K'} Y'$ (solid black arrow)
 - $\gamma: Y \rightarrow K \times_{K'} Y'$ (solid black arrow)
 - $\gamma': Y' \rightarrow K \times_{K'} Y'$ (solid black arrow)
 - $\eta': K \times_{K'} Y' \rightarrow K$ (solid black arrow)
- Isomorphisms:**
 - $\eta: K \times_{K'} Y \rightarrow K \times_{K'} Y'$ (double line)
 - $\eta': K \times_{K'} Y' \rightarrow K$ (double line)
- Existence of Maps (dashed arrows):**
 - Orange dashed arrow from X to $K \times_{K'} Y$ labeled $\exists!$.
 - Blue dashed arrow from X to $K \times_{K'} Y'$ labeled $\exists!$.

while $\Phi_{X,Y'}(\beta \circ \xi)$ is given by the dashed arrow from X to $K \times_{K'} Y'$, coloured in **blue**. Since both the **blue arrow** and the **vermillion arrow** make the outer pullback

diagram for $K \times_{K'} Y'$ commute, it follows from the universal property of the pullback that they must be equal, i.e. that

$$f^*(\beta) \circ \Phi_{X,Y}(\xi) = \Phi_{X,Y'}(\beta \circ \xi),$$

showing that the naturality diagram above indeed commutes.

- *Invertibility I.* We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \text{id}_{\text{HomFibSets}(K')(\Sigma_f(X), Y)}.$$

Indeed, $\Phi_{X,Y}$ sends a map

$$\xi: \Sigma_f(X) \rightarrow Y,$$

of K' -fibred sets to the dashed morphism in the diagram

and $\Psi_{X,Y}$ then postcomposes that map with pr_2 , which, by the commutativity of the diagram above, is ξ again, showing the claimed equality to be true.

- *Invertibility II.* We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \text{id}_{\text{HomFibSets}(K)(X, f^*(Y))}.$$

Indeed, $\Psi_{X,Y}$ sends a map

$$\xi: X \rightarrow f^*(Y),$$

of K' -fibred sets to $\text{pr}_2 \circ \xi$, which is then sent by $\Phi_{X,Y}$ to the dashed morphism in the diagram

$$\begin{array}{ccccc}
 X & & & & \\
 \downarrow \phi_X & \searrow \exists! & & \xrightarrow{\text{pr}_2 \circ \xi} & Y \\
 & & K \times_{K'} Y & \xrightarrow{\text{pr}_2} & Y \\
 & & \downarrow \text{pr}_1 & \lrcorner & \downarrow \phi_Y \\
 & & K & \xrightarrow{f} & K'
 \end{array}$$

(†)

which, by the commutativity of the subdiagram marked with (†), is given by ξ again, showing the claimed equality to be true.

The Adjunction $f^ \dashv \Pi_f$:* We claim there's a bijection

$$\text{Hom}_{\text{FibSets}(K)}(f^*(X), Y) \cong \text{Hom}_{\text{FibSets}(K')}(X, \Pi_f(Y))$$

natural in $(X, \phi_X) \in \text{Obj}(\text{FibSets}(K'))$ and $(Y, \phi_Y) \in \text{Obj}(\text{FibSets}(K))$:

1. *Map I.* We define a map

$$\Phi_{X,Y}: \text{Hom}_{\text{FibSets}(K)}(f^*(X), Y) \rightarrow \text{Hom}_{\text{FibSets}(K')}(X, \Pi_f(Y))$$

defined as follows. Given a morphism

$$\xi: f^*(X) \rightarrow Y, \quad \begin{array}{ccc} K \times_{K'} X & \xrightarrow{\xi} & Y \\ \text{pr}_1 \searrow & & \nearrow \phi_Y \\ & K & \end{array}$$

of K -fibred sets, where

$$\begin{aligned}
 f^*(X) &\stackrel{\text{def}}{=} K \times_{K'} X \\
 &\stackrel{\text{def}}{=} \{(k, x) \in K \times X \mid f(k) = \phi_X(x)\},
 \end{aligned}$$

we construct a morphism

$$\xi^\dagger: X \rightarrow \Pi_f(Y), \quad \begin{array}{ccc} X & \xrightarrow{\xi^\dagger} & \Pi_f(Y) \\ \phi_X \searrow & & \nearrow \Pi_f(\phi_Y) \\ & K' & \end{array}$$

of K' -fibred sets, where

$$\Pi_f(Y) \stackrel{\text{def}}{=} \coprod_{k' \in K'} \prod_{k \in f^{-1}(k')} \phi_Y^{-1}(k),$$

by defining

$$\xi^\dagger(x) \stackrel{\text{def}}{=} (\xi(k, x))_{k \in f^{-1}(\phi_X(x))}$$

for each $x \in X$. There are two things to be checked here:

- We have $\xi(k, x) \in \phi_Y^{-1}(\phi_X(x))$ since $\phi_Y(\xi(k, x)) = \phi_X(x)$ as ξ is a morphism of K -fibred sets.
- The map ξ^\dagger is indeed a morphism of K' -fibred sets, i.e. we have

$$\Pi_f(\phi_Y) \circ \xi^\dagger = \phi_X,$$

since

$$[\Pi_f(\phi_Y)]\left((\xi(k, x))_{k \in f^{-1}(\phi_X(x))}\right) = \phi_X(x)$$

for each $x \in X$.

2. *Map II.* We define a map

$$\Psi_{X,Y} : \text{Hom}_{\text{FibSets}(K')}(X, \Pi_f(Y)) \rightarrow \text{Hom}_{\text{FibSets}(K)}(f^*(X), Y)$$

as follows. Given a morphism

$$\xi : X \rightarrow \Pi_f(Y), \quad \begin{array}{ccc} X & \xrightarrow{\xi} & \Pi_f(Y) \\ \phi_X \searrow & & \swarrow \Pi_f(\phi_Y) \\ & K' & \end{array}$$

of K' -fibred sets, where

$$\Pi_f(Y) \stackrel{\text{def}}{=} \coprod_{k' \in K'} \prod_{k \in f^{-1}(k')} \phi_Y^{-1}(k),$$

we construct a morphism

$$\xi^\dagger : f^*(X) \rightarrow Y, \quad \begin{array}{ccc} K \times_{K'} X & \xrightarrow{\xi^\dagger} & Y \\ \text{pr}_1 \searrow & & \swarrow \phi_Y \\ & K & \end{array}$$

of K -fibred sets, where

$$\begin{aligned} f^*(X) &\stackrel{\text{def}}{=} K \times_{K'} X \\ &\stackrel{\text{def}}{=} \{(k, x) \in K \times X \mid f(k) = \phi_X(x)\}, \end{aligned}$$

by defining

$$\xi^\dagger(k, x) \stackrel{\text{def}}{=} \xi(x)_k$$

for each $(k, x) \in f^*(X)$, where $\xi(x)_k$ is the k th component of $\xi(x) = (y_k)_{k \in f^{-1}(k')}$. We also need to check that ξ^\dagger is a morphism of K -fibred sets, i.e. that

$$\phi_Y \circ \xi^\dagger = \text{pr}_1,$$

or

$$\phi_Y(\xi^\dagger(k, x)) = k,$$

for each $(k, x) \in f^*(X)$, which is clear, since $\xi^\dagger(k, x) \in \phi_Y^{-1}(k)$ by definition.

3. *Naturality I.* We need to show that, given a morphism

$$\alpha: (X, \phi_X) \rightarrow (X', \phi_{X'})$$

of K' -fibred sets, the diagram

$$\begin{array}{ccc} \text{Hom}_{\text{FibSets}(K')} (f^*(X'), Y) & \xrightarrow{\Phi_{X', Y}} & \text{Hom}_{\text{FibSets}(K)} (X', \Pi_f(Y)) \\ f^*(\alpha)^* \downarrow & & \downarrow \alpha^* \\ \text{Hom}_{\text{FibSets}(K')} (f^*(X), Y) & \xrightarrow{\Phi_{X, Y}} & \text{Hom}_{\text{FibSets}(K)} (X, \Pi_f(Y)) \end{array}$$

commutes. Indeed, given a morphism $\xi: f^*(X') \rightarrow Y$ of K' -fibred sets, we have

$$\begin{aligned} [[\Phi_{X, Y} \circ f^*(\alpha)](\xi)](x) &\stackrel{\text{def}}{=} [\Phi_{X, Y}(\xi \circ f^*(\alpha))](x) \\ &\stackrel{\text{def}}{=} ([\xi \circ f^*(\alpha)](k, x))_{k \in f^{-1}(\phi_X(x))} \\ &\stackrel{\text{def}}{=} (\xi(k, \alpha(x)))_{k \in f^{-1}(\phi_X(x))} \\ &\stackrel{\text{def}}{=} \alpha^*((\xi(k, x))_{k \in f^{-1}(\phi_X(x))}) \\ &\stackrel{\text{def}}{=} \alpha^*(\xi^\dagger(x)) \\ &\stackrel{\text{def}}{=} [[\alpha^* \circ \Phi_{X, Y}](\xi)](x) \end{aligned}$$

for each $x \in X$.

4. *Naturality II.* We need to show that, given a morphism

$$\beta: (Y, \phi_Y) \rightarrow (Y', \phi_{Y'})$$

of K -fibred sets, the diagram

$$\begin{array}{ccc} \mathrm{Hom}_{\mathrm{FibSets}(K')}(f^*(X), Y) & \xrightarrow{\Phi_{X,Y}} & \mathrm{Hom}_{\mathrm{FibSets}(K)}(X, \Pi_f(Y)) \\ \beta_* \downarrow & & \downarrow \Pi_f(\beta)_* \\ \mathrm{Hom}_{\mathrm{FibSets}(K')}(f^*(X), Y') & \xrightarrow{\Phi_{X,Y'}} & \mathrm{Hom}_{\mathrm{FibSets}(K)}(X, \Pi_f(Y')) \end{array}$$

commutes. Indeed, given a morphism $\xi: X \rightarrow \Pi_f(Y)$ of K -fibred sets, we have

$$\begin{aligned} [[\Phi_{X,Y'} \circ \beta_*](\xi)](x) &\stackrel{\mathrm{def}}{=} [\Phi_{X,Y'}(\beta \circ \xi)](x) \\ &\stackrel{\mathrm{def}}{=} [\Phi_{X,Y'}(\beta \circ \xi)](x) \\ &\stackrel{\mathrm{def}}{=} ([\beta \circ \xi](k, x))_{k \in f^{-1}(\phi_X(x))} \\ &\stackrel{\mathrm{def}}{=} (\beta(\xi(k, x)))_{k \in f^{-1}(\phi_X(x))} \\ &\stackrel{\mathrm{def}}{=} \Pi_f(\beta)_*((\xi(k, x))_{k \in f^{-1}(\phi_X(x))}) \\ &\stackrel{\mathrm{def}}{=} [\Pi_f(\beta)_* \circ \xi^\dagger](x) \\ &\stackrel{\mathrm{def}}{=} [\Pi_f(\beta)_* \circ \Phi_{X,Y'}(\xi)](x) \end{aligned}$$

for each $x \in X$.

5. *Invertibility I.* We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathrm{Hom}_{\mathrm{FibSets}(K')}(f^*(X), Y)}.$$

Indeed, given a morphism $\xi: f^*(X') \rightarrow Y$ of K' -fibred sets, we have

$$\begin{aligned} [[\Psi_{X,Y} \circ \Phi_{X,Y}](\xi)](k, x) &\stackrel{\mathrm{def}}{=} [\Psi_{X,Y}(\Phi_{X,Y}(\xi))](k, x) \\ &\stackrel{\mathrm{def}}{=} ([\Phi_{X,Y}(\xi)](x))_k \\ &\stackrel{\mathrm{def}}{=} ((\xi(k_1, x))_{k_1 \in f^{-1}(\phi_X(x))})_k \\ &\stackrel{\mathrm{def}}{=} \xi(k, x) \end{aligned}$$

for each $(k, x) \in f^*(X)$, and thus the stated equality follows.

6. *Invertibility II.* We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathrm{Hom}_{\mathrm{FibSets}(K)}(X, \Pi_f(Y))}.$$

Indeed, given a morphism $\xi: X \rightarrow \Pi_f(Y)$ of K -fibred sets, write

$$\xi(x) = (y_k)_{k \in f^{-1}(k'_x)}.$$

We then have

$$\begin{aligned} [[\Phi_{X,Y} \circ \Psi_{X,Y}](\xi)](x) &\stackrel{\text{def}}{=} [\Phi_{X,Y}(\Psi_{X,Y}(\xi))](x) \\ &\stackrel{\text{def}}{=} ([\Psi_{X,Y}(\xi)](k_1, x))_{k_1 \in f^{-1}(\phi_X(x))} \\ &\stackrel{\text{def}}{=} ((\xi(x))_{k_1})_{k_1 \in f^{-1}(\phi_X(x))} \\ &\stackrel{\text{def}}{=} ((y_k)_{k \in f^{-1}(k'_x)})_{k_1 \in f^{-1}(\phi_X(x))} \\ &\stackrel{\text{def}}{=} (y_{k_1})_{k_1 \in f^{-1}(\phi_X(x))} \\ &= (y_{k_1})_{k_1 \in f^{-1}(k'_x)} \\ &= (y_k)_{k \in f^{-1}(k'_x)} \\ &\stackrel{\text{def}}{=} \xi(x) \end{aligned}$$

for each $x \in X$, where the equality $\phi_X(x) = k'_x$ follows from the fact that ξ is a morphism of K' -fibred sets. Thus the stated equality follows.

This finishes the proof. \square

Appendices

A Other Chapters

Sets

1. **Sets**
2. **Constructions With Sets**
3. **Pointed Sets**
4. **Tensor Products of Pointed Sets**
5. **Relations**
6. **Spans**
7. **Posets**

Indexed and Fibred Sets

7. **Indexed Sets**

8. **Fibred Sets**

9. **Un/Straightening for Indexed and Fibred Sets**

Category Theory

11. **Categories**
12. **Types of Morphisms in Categories**
13. **Adjunctions and the Yoneda Lemma**
14. **Constructions With Categories**
15. **Kan Extensions**

Bicategories

- 17. Bicategories
- 18. Internal Adjunctions

Internal Category Theory

- 19. Internal Categories

Cyclic Stuff

- 20. The Cycle Category

Cubical Stuff

- 21. The Cube Category

Globular Stuff

- 22. The Globe Category

Cellular Stuff

- 23. The Cell Category

Monoids

- 24. Monoids
- 25. Constructions With Monoids

Monoids With Zero

- 26. Monoids With Zero
- 27. Constructions With Monoids With Zero

Groups

- 28. Groups
- 29. Constructions With Groups

Hyper Algebra

- 30. Hypermonoids

- 31. Hypergroups

- 32. Hypersemirings and Hyperrings

- 33. Quantaes

Near-Rings

- 34. Near-Semirings

- 35. Near-Rings

Real Analysis

- 36. Real Analysis in One Variable

- 37. Real Analysis in Several Variables

Measure Theory

- 38. Measurable Spaces

- 39. Measures and Integration

Probability Theory

- 39. Probability Theory

Stochastic Analysis

- 40. Stochastic Processes, Martingales, and Brownian Motion

- 41. Itô Calculus

- 42. Stochastic Differential Equations

Differential Geometry

- 43. Topological and Smooth Manifolds

Schemes

- 44. Schemes