

A decorative graphic on the left side of the slide consisting of two overlapping parallelograms. The front one is blue and the back one is a light greenish-blue. They are positioned diagonally, with the blue one in front of the green one.

Introduction to Cybersecurity

Lecture #4: Reverse engineering I

Anton Semenko



CTF sections

- Web
- Crypto
- Stegano
- Reverse
- Pwn
- Ppc
- Recon
- Misc
- Forensics



CTF sections

- Web
- Crypto
- Stegano
- **Reverse**
- Pwn
- Ppc
- Recon
- Misc
- Forensics

What is it?

- Forward engineering

- Come up with idea
- Choose tools
- Create an architecture
- Develop application



- Reverse engineering

- You have a black box (desktop / mobile app)
- Study how this box works





Why?



Why?

Make better

Find hidden functions

Get +400

Adapt for new
hardware

Flex

Make your own

Find mistakes

Bugbounty

Harm (illegal)

Hardware reverse



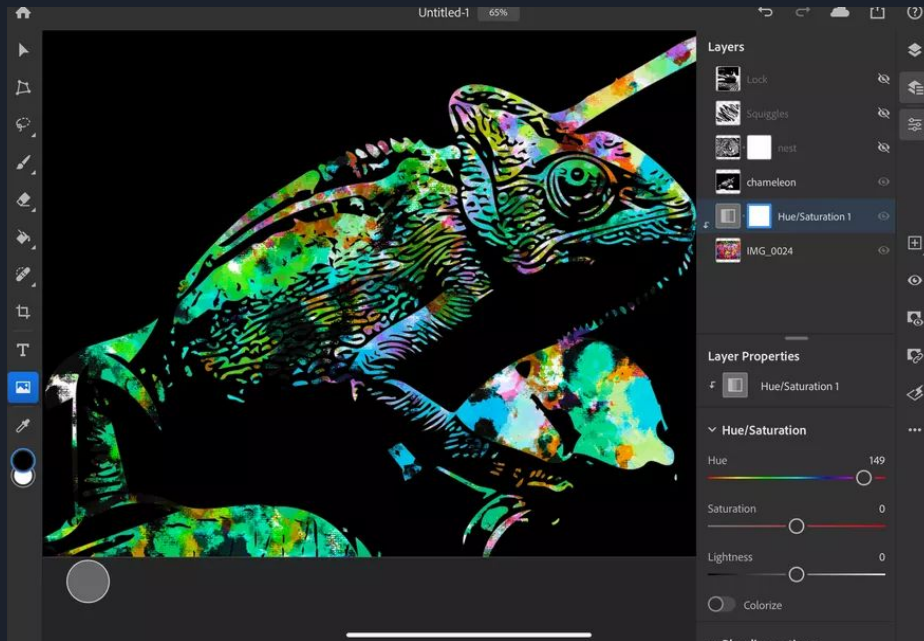
Hardware reverse



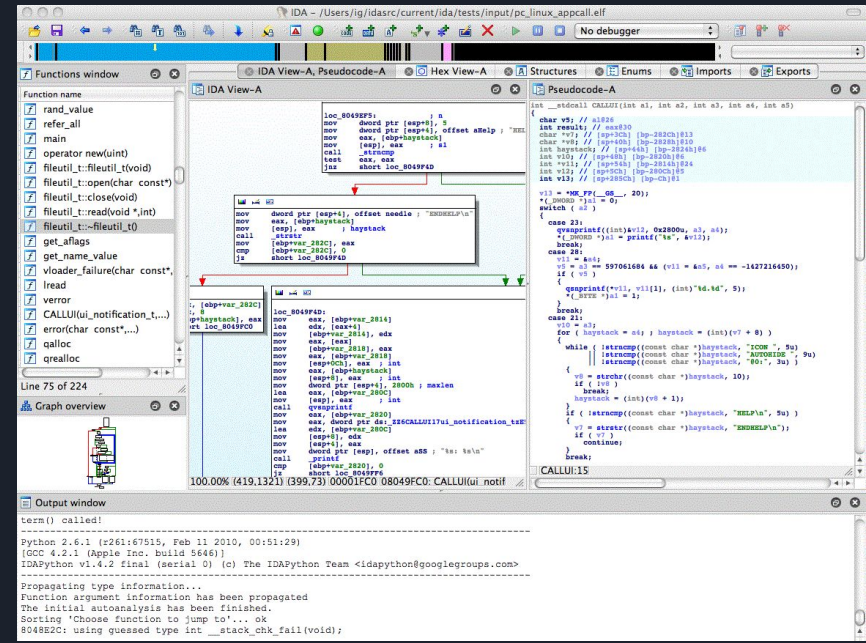
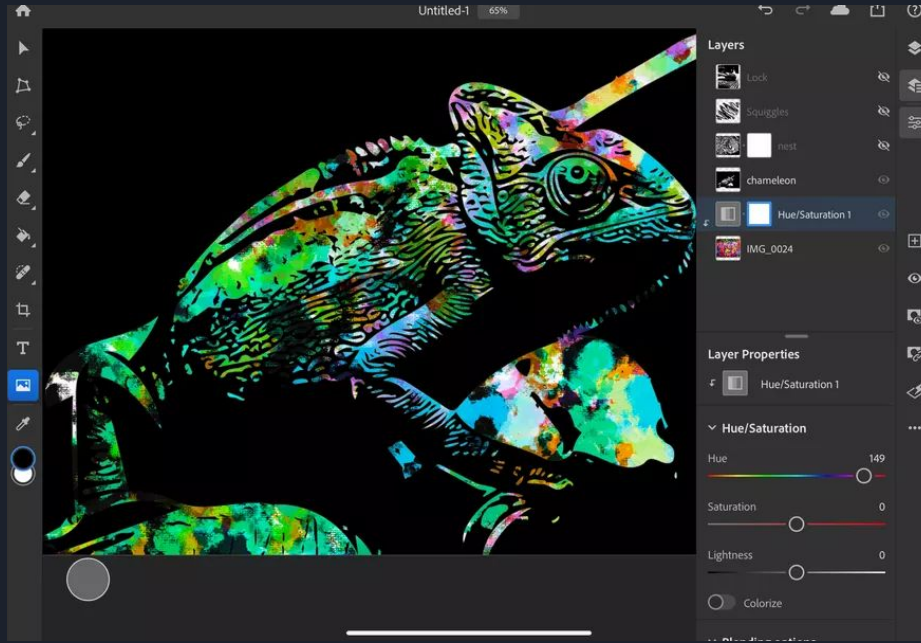
Hardware reverse



Software reverse



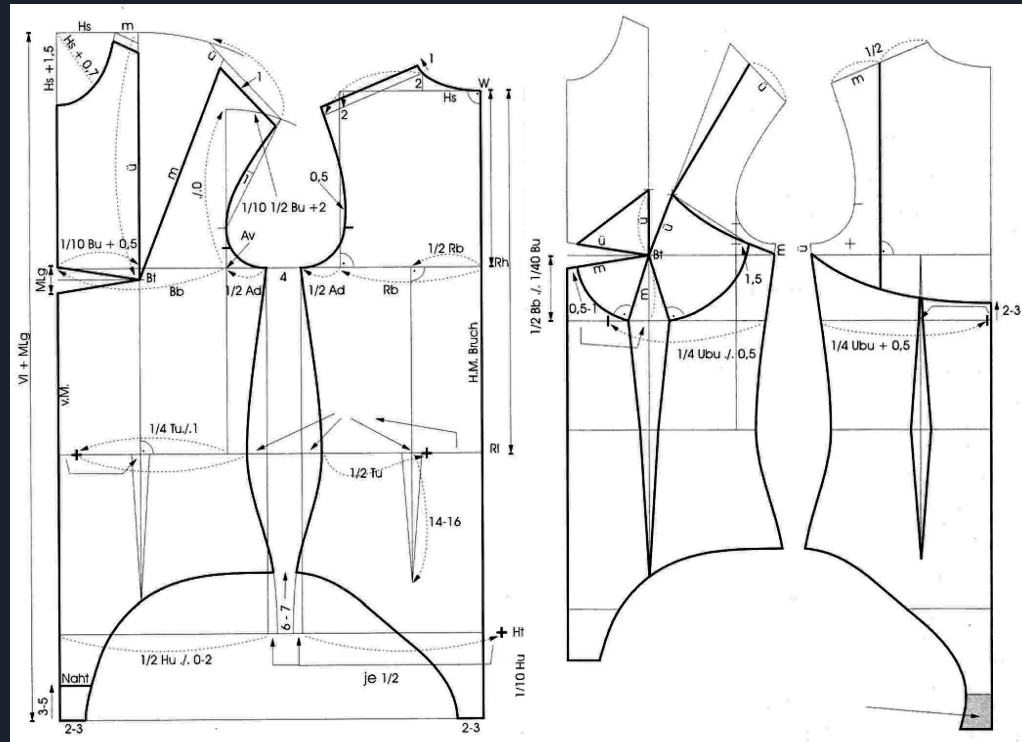
Software reverse



Clothes reverse



Clothes reverse



Clothes reverse





Time to play

<https://overthewire.org/wargames/bandit/>



Reverse is hard



What is the program?

- Code point of view
 - Operations
 - Custom functions
 - Library functions
 - System calls
- OS point of view
 - Executable and Linkable Format
 - Scripts (`#!/usr/local/bin/my_brainfuck_interp`, `#!/bin/zsh`)
 - Old formats (a.out, Common Object File Format)

ELF

- Other programs
- Dynamic libraries (*.so)
- Object files (*.o)

Executable file has only one entry point (*_start* or *main*), but the library has many (all exportable functions)





Demo

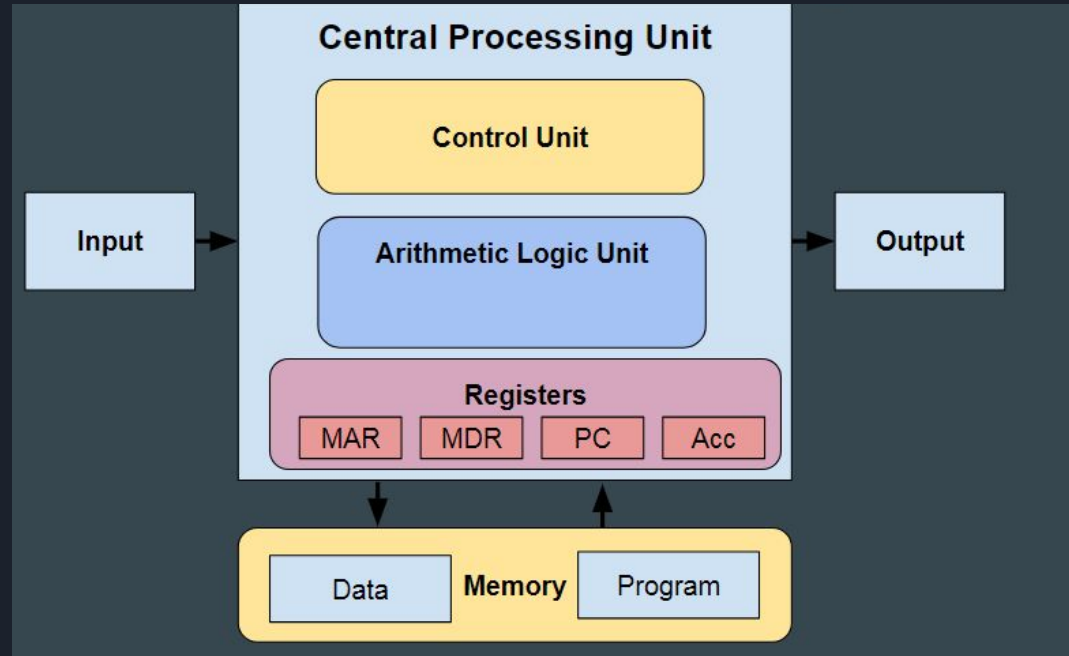


How computer works

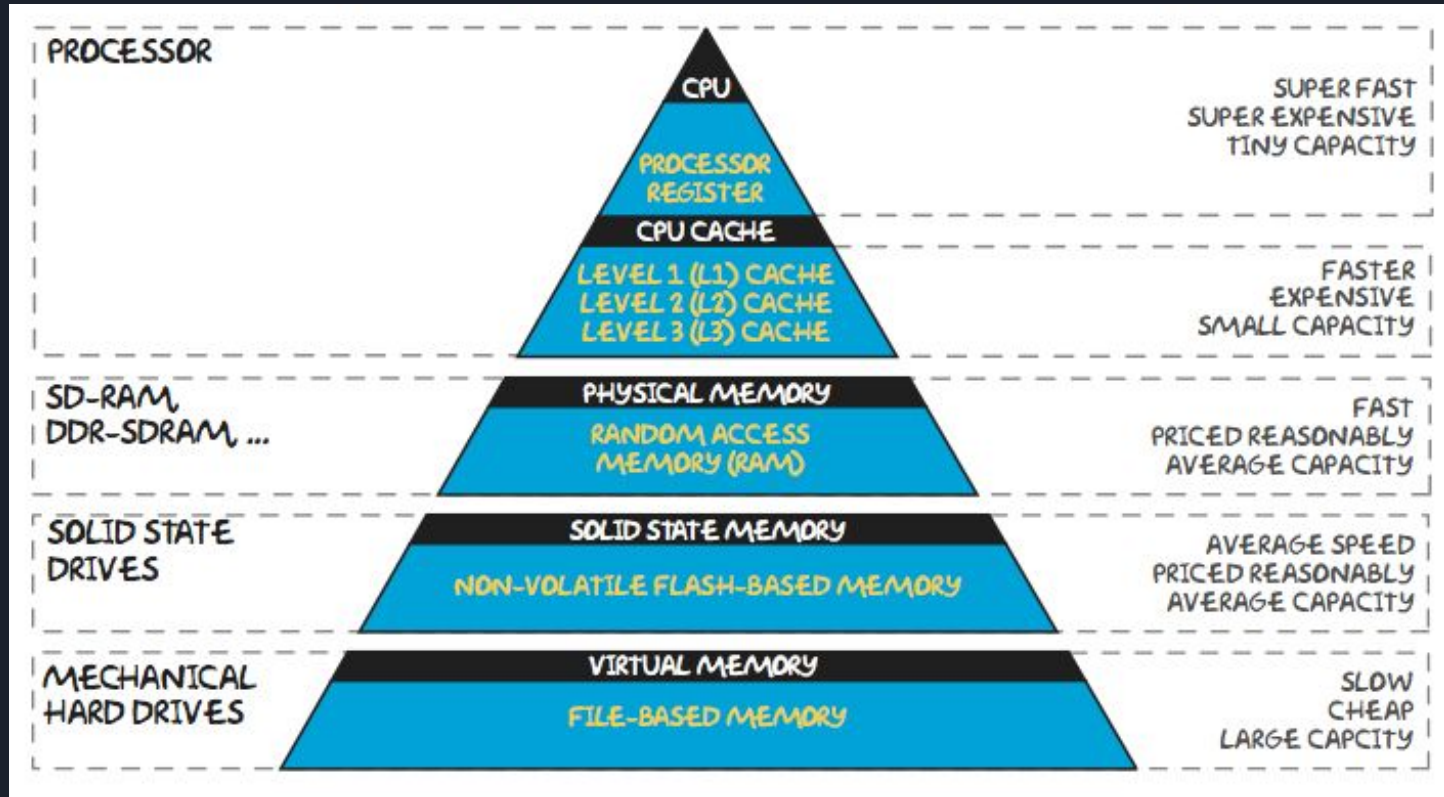
- Von Neumann architecture
- Memory hierarchy
- Process memory

Von Neumann architecture

- Processing unit
- Control unit
- Memory unit
- Input / output devices



Memory hierarchy





x64 VS x32?

x32 processor registers

General-purpose
registers

EAX
ECX
EDX
EBX

Control and Status
'Flags' register

EFLAGS

Pointer and Index
registers

ESP
EBP
ESI
EDI

Instruction-Pointer
register

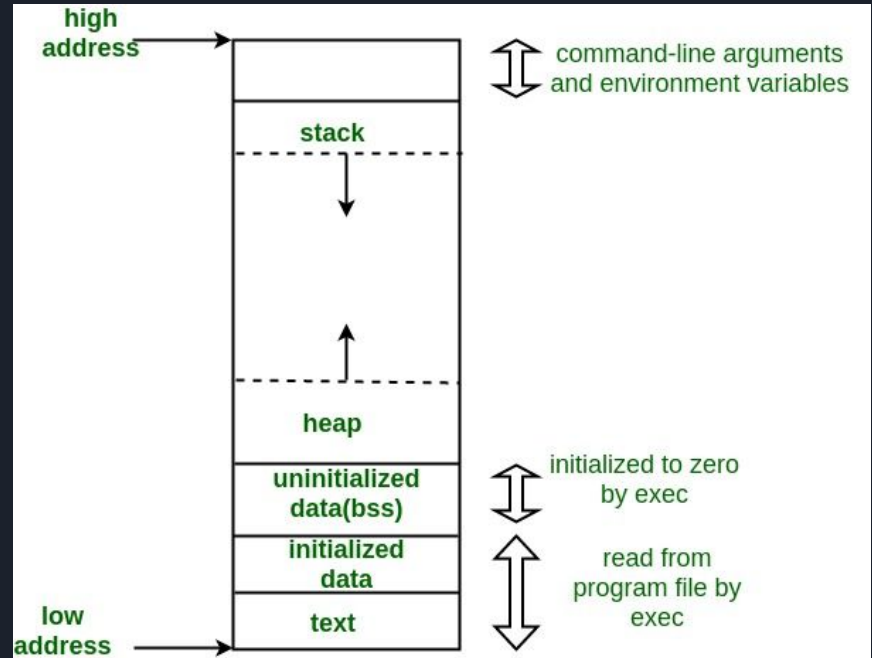
EIP

Segment/selector
registers

ES
CS
SS
DS
FS
GS

Memory layout

- Static
 - .bss, .data
- Dynamic
 - heap
- Automatic
 - stack





Demo



stack VS heap

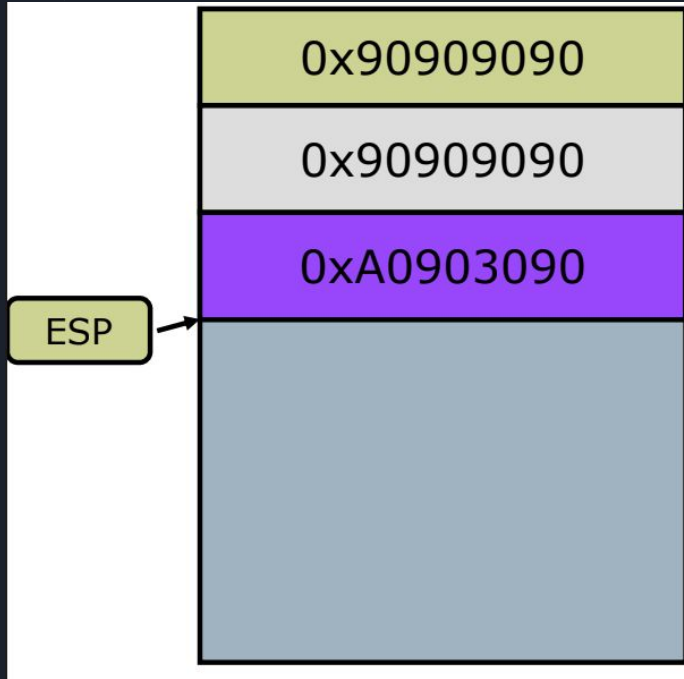


Stack

- Last-in-first-out (LIFO) structure
- Contains temporary (local) function variables
- Managed and optimized by CPU (not by you)
- Functions clear stack after exit^{*}

^{*}For more information read about [calling conventions](#)

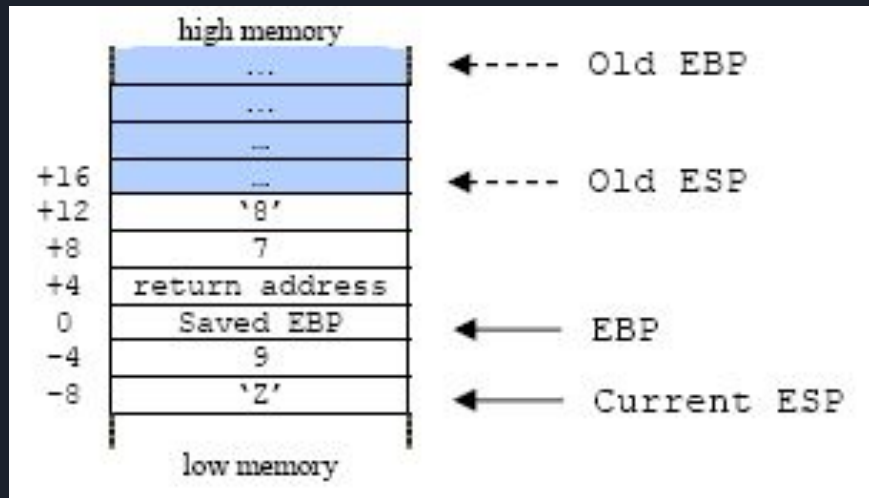
Stack: in-depth view



- Push - put element on the top of the stack
- Pop - take top value from stack
- ESP - pointer to the top of the stack

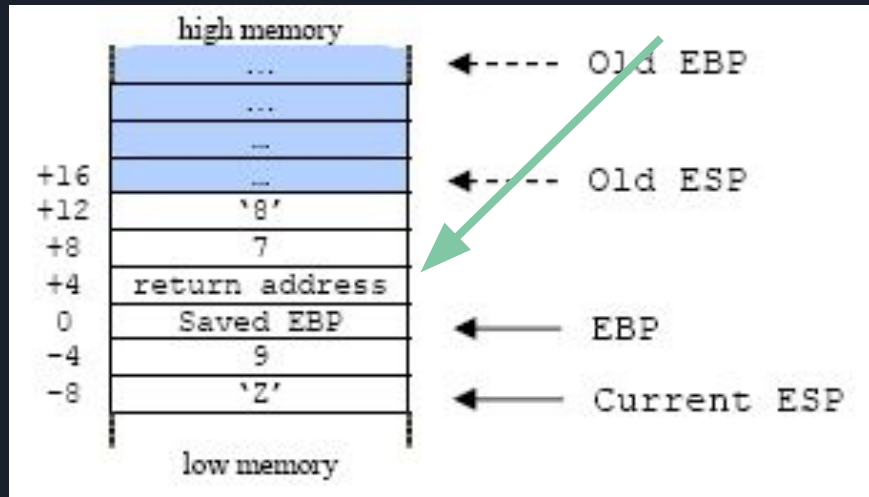
Stack: in-depth view

```
1 #include <stdio.h>
2
3 int MyFunc(int parameter1, char parameter2) {
4     int local1 = 9;
5     char local2 = 'Z';
6     return 0;
7 }
8
9 int main(int argc, char **argv) {
10     MyFunc(7, '8');
11     return 0;
12 }
```



Stack: in-depth view

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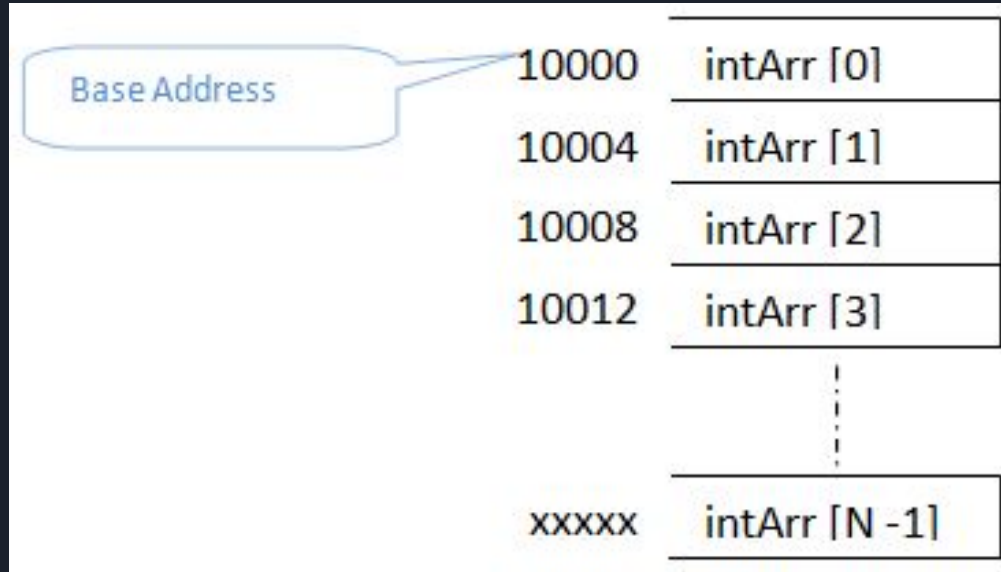
Demo



What's wrong?

```
int func(char *user) {  
    fprintf(stdout, user);  
}
```

Arrays in stack





Heap

- Managed by you
- Global variables
- malloc(), calloc()
- YOU are responsible to free() the memory you allocated
- If you fail to free the space, you've got **memory leak**

For hardcore reversers:

<https://www.contextis.com/en/resources/white-papers/glibc-adventures-the-forgotten-chunks>



Reverse levels

- Source code analysis



- Tracing



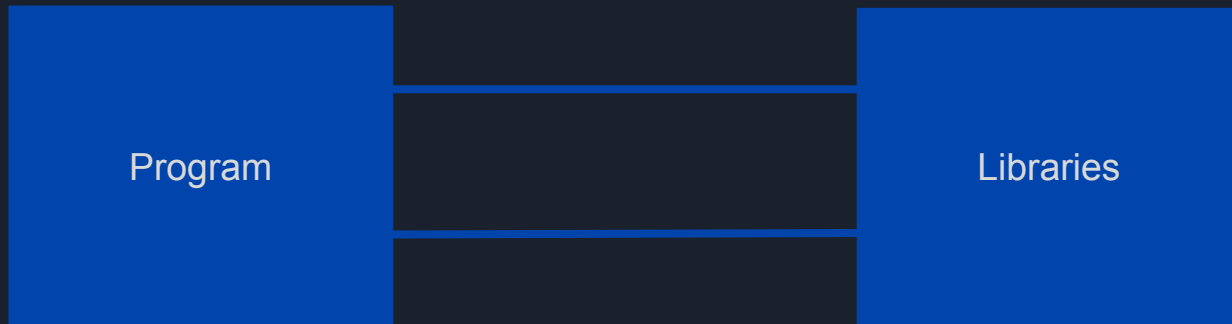
- Static analysis



- Binary patching

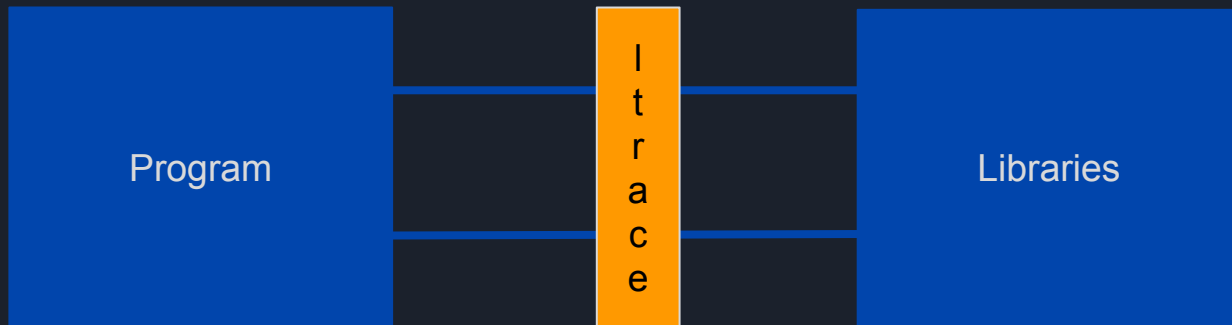


Tracing: functions interception





Tracing: functions interception

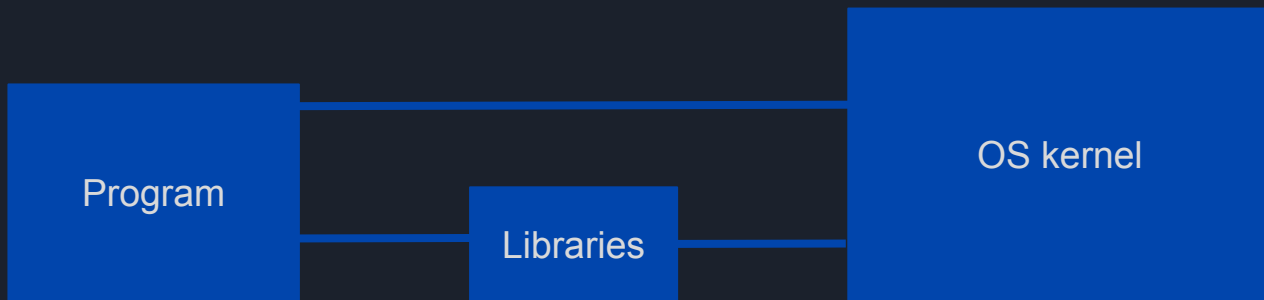




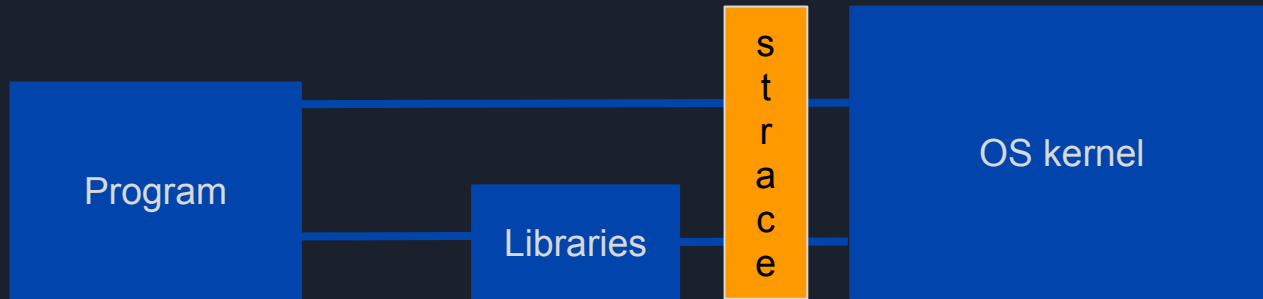
Demo



Tracing programs: syscalls interception



Tracing programs: syscalls interception





Demo