

FYS4150 - Project 1: Differential Equations and Linear Algebra

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1 Introduction

Poisson's equation, given in general as

$$\nabla^2 u(\mathbf{r}) = f(\mathbf{r}) \quad (1)$$

where $f(\mathbf{r})$ is some known function, is an often-used equation in physics as it describes how a field (in this case u) is affected by some source (in this case f). If f is a complicated function this equation may not be analytically solvable, so one must use numerical methods instead, to find an approximate solution. In one dimension this can be done by using an approximate expression for the second derivative of u at discrete set of points; then equation 1 becomes a set of linear equations for the values of u at these points, which can be expressed as a matrix involving a tridiagonal matrix. Thus the problem of solving Poisson's equation numerically reduces to inverting such an equation numerically.

In this project I have examined three different methods for solving this matrix equation: One "brute-force" method which inverts the entire matrix numerically, and two which exploit the fact that the matrix is tridiagonal by storing data in vectors instead of a matrix, which reduces memory usage and speeds up the algorithm. One of these uses a general algorithm for row-reducing tridiagonal matrices, while the other takes advantage of the specific form of the matrix in question. In this text I will first outline these methods, and then compare their results, in particular the difference in CPU time used. I will also discuss the effects of round-off errors as the step size between the different data points is changed.

2 Methods

2.1 Turning the Poisson equation into a matrix equation

The target of this project was solving Poisson's equation,

$$u''(x) = -f(x), \quad (2)$$

for $f(x) = 100e^{-10x}$ for $x \in [0, 1]$ with the boundary conditions $u(0) = u(1) = 0$. This equation has an exact solution given by [1]

$$u(x) = (e^{-10} - 1)x + 1 - e^{-10x} \quad (3)$$

The first step is dividing this interval into a discrete set of points x_i for $i = 0, 1, 2, \dots, n+1$, where $x_0 = 0$ and $x_{n+1} = 1$. Thus the step size is given by $h = x_{i+1} - x_i = 1/(n+1)$. The second derivative can be approximated by [2]

$$u_i'' \approx \frac{u_{i-1} - 2u_i + u_{i+1}}{h^2} \quad (4)$$

(obtained by manipulating the Taylor expansion of $u(x)$) where the error is of order h^2 . Here $u_i \equiv u(x_i)$, etc. Inserting this into equation 2 gives

$$-u_{i-1} + 2u_i - u_{i+1} = g_i$$

where I have defined $g_i \equiv h^2 f(x_i)$. This is a set of linear equations for u_i ; it can thus be rewritten on matrix form as (using the fact that $u_0 = u_{n+1} = 0$, by the boundary conditions):

$$\mathbf{A}\mathbf{u} = \mathbf{g} \quad (5)$$

Here A is an $n \times n$ matrix whose non-zero elements on the i 'th row are $A_{i,i-1} = -1$, $A_{i,i} = 2$, $A_{i,i+1} = -1$ (except the first row which only has $A_{0,0} = 2$ and $A_{0,1} = -1$, and the n 'th row which only has $A_{n,n-1} = -1$ and $A_{n,n} = 2$). Thus it is a tridiagonal matrix with 2 on the diagonal and -1 on the elements directly adjacent to the diagonal. Furthermore $\mathbf{u}^T \equiv [u_1, u_2, \dots, u_n]$, and $\mathbf{g}^T \equiv [g_1, g_2, \dots, g_n]$.

2.2 Row-reducing a tridiagonal matrix

The problem is now finding the reduces row echelon form of the matrix

$$[A; \mathbf{g}] = \begin{bmatrix} b_1 & c_1 & 0 & \dots & \dots & \dots & g_1 \\ a_1 & b_2 & c_2 & \dots & \dots & \dots & g_2 \\ \dots & a_2 & b_3 & c_3 & \dots & \dots & g_3 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & a_{n-2} & b_{n-1} & c_{n-1} & g_{n-1} \\ \dots & \dots & \dots & \dots & a_{n-1} & b_n & g_n \end{bmatrix} \quad (6)$$

For the moment, I am not making any assumptions about the values of a_i , b_i or c_i . Renaming b_1 and g_1 to \tilde{b}_1 and \tilde{g}_1 respectively, and then subtracting $a_1/\tilde{b}_1 \times (\text{row } 1)$ from the second row leaves

$$\begin{bmatrix} \tilde{b}_1 & c_1 & 0 & \dots & \dots & \dots & \tilde{g}_1 \\ 0 & \tilde{b}_2 & c_2 & \dots & \dots & \dots & \tilde{g}_2 \\ \dots & a_2 & b_3 & c_3 & \dots & \dots & g_3 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & a_{n-2} & b_{n-1} & c_{n-1} & g_{n-1} \\ \dots & \dots & \dots & \dots & a_{n-1} & b_n & g_n \end{bmatrix} \quad (7)$$

where $\tilde{b}_2 = b_2 - (a_1/\tilde{b}_1)c_1$ and $\tilde{g}_2 = g_2 - (a_1/b_1)\tilde{g}_1$. The matrix can be brought to an upper triangular form by repeating this process, i.e

$$\begin{bmatrix} \tilde{b}_1 & c_1 & 0 & \dots & \dots & \dots & \tilde{g}_1 \\ 0 & \tilde{b}_2 & c_2 & \dots & \dots & \dots & \tilde{g}_2 \\ \dots & 0 & \tilde{b}_3 & c_3 & \dots & \dots & \tilde{g}_3 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & 0 & \tilde{b}_{n-1} & c_{n-1} & \tilde{g}_{n-1} \\ \dots & \dots & \dots & \dots & 0 & \tilde{b}_n & \tilde{g}_n \end{bmatrix} \quad (8)$$

where $\tilde{b}_i = b_i - (a_{i-1}/\tilde{b}_{i-1})c_{i-1}$ and $\tilde{g}_i = g_i - (a_{i-1}/\tilde{b}_{i-1})\tilde{g}_{i-1}$ for $i = 2, 3, \dots, n$. This matrix can be brought to reduced row echelon form by dividing the lower row by \tilde{b}_n (the n 'th element of the solution is then given by $u_n = \tilde{g}_n/\tilde{b}_n$), and subtracting c_{n-1} times this line from the line above it; dividing this line by \tilde{b}_{n-1} then gives the $(n-1)$ 'th element of the solution; $u_{n-1} = (\tilde{g}_{n-1} - c_{n-1}u_n)/\tilde{b}_{n-1}$. Repeating this process brings the matrix to the form

$$\begin{bmatrix} 1 & 0 & 0 & \dots & \dots & \dots & u_1 \\ 0 & 1 & 0 & \dots & \dots & \dots & u_2 \\ \dots & 0 & 1 & 0 & \dots & \dots & u_3 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & 0 & 1 & 0 & u_{n-1} \\ \dots & \dots & \dots & \dots & 0 & 1 & u_n \end{bmatrix} \quad (9)$$

where $u_i = (\tilde{g}_i - c_i u_{i+1})/\tilde{b}_i$ for $i = 1, 2, \dots, n-1$ (note that this is not the full solution as the interval is divided into $n+2$ points; however the end points u_0 and u_{n+1} are both known from the boundary conditions to be zero). Given vectors \mathbf{a} , \mathbf{b} , \mathbf{c} and \mathbf{g} , the solution \mathbf{u} can thus be found by the following two-step algorithm:

Forward substitution (general):

$$\tilde{b}_1 = b_1; \quad \tilde{g}_1 = g_1 \quad (10)$$

For $i = 2, 3, \dots, n$:

$$\tilde{b}_i = b_i - \frac{a_{i-1}}{\tilde{b}_{i-1}}c_{i-1} \quad (11)$$

$$\tilde{g}_i = g_i - \frac{a_{i-1}}{\tilde{b}_{i-1}}\tilde{g}_{i-1} \quad (12)$$

Backward substitution (general):

$$u_n = \frac{\tilde{g}_n}{\tilde{b}_n} \quad (13)$$

For $i = n-1, n-2, \dots, 1$:

$$u_i = \frac{\tilde{g}_i - c_i u_{i+1}}{\tilde{b}_i} \quad (14)$$

When programming, this requires $8n$ floating point operations (if the factor a_{i-1}/\tilde{b}_{i-1} is calculated only once per i) (not exactly, but for large n the error is negligible). I implemented this algorithm in C++ (with $a_i = c_i = -1$, $b_i = 2$ and $g_i = h^2 f(x_i)$ with f as specified above) for $n = 10, 100, 1000, \dots, 10^7$ and wrote the solution (for $n = 10$, $n = 100$ and $n = 1000$), maximum relative error and time elapsed running the algorithm to files.

2.3 Specialization of the algorithm

In this case all $a_i = c_i = -1$, and $b_i = 2$. Thus the algorithm as described above can be specialized; first of all, $\tilde{b}_i = (i+1)/i$ (this can be proved by induction, from $\tilde{b}_i = 2 - 1/\tilde{b}_{i-1}$). If these are pre-calculated they will not contribute to the number of floating point operations of the algorithm. This can also improve precision, as all elements are calculated directly as opposed to from repeated subtractions in the general method; for large n these repeated subtractions can lead to round-off errors building up. A specialized version of the previously described algorithm is then:

Forward substitution (specialized):

$$\tilde{b}_i = \frac{i+1}{i} \text{ for } i = 1, 2, \dots, n; \quad \tilde{g}_1 = g_1 \quad (15)$$

For $i = 2, 3, \dots, n$:

$$\tilde{g}_i = g_i + \frac{\tilde{g}_{i-1}}{\tilde{b}_{i-1}} \quad (16)$$

Backward substitution (specialized):

$$u_n = \frac{\tilde{g}_n}{\tilde{b}_n} \quad (17)$$

For $i = n-1, n-2, \dots, 1$:

$$u_i = \frac{\tilde{g}_i + u_{i+1}}{\tilde{b}_i} \quad (18)$$

Not counting the pre-calculated variables, this requires $4n$ floating point operations. Naively one would then expect this algorithm to be roughly twice as fast as the general one; however this is not the case, as the number of memory reads/writes is the same for both cases. Thus this algorithm should be faster, but not quite twice as fast. I implemented this algorithm in C++ for the same parameters as the general one, and compared the results (in particular the run-time) to those from the general algorithm.

2.4 Solution by LU decomposition

For the present problem the algorithms as described above are the most efficient; by exploiting the tridiagonal nature of the matrix they only require a number of floating point operations which is proportional to n . They also use relatively little memory; as only the non-zero elements of the matrix are stored the number of stored values is also proportional to n . However, for a more general matrix equation $\mathbf{A}\mathbf{u} = \mathbf{g}$ more generally applicable methods must be used, such as LU

decomposition. In such methods the entire $n \times n$ matrix is used, so that the number of stored values is proportional to n^2 .

In LU decomposition, the matrix A is written as a product of a lower and an upper triangular matrix, $A = LU$, where [2]

$$L_{ij} = \begin{cases} l_{ij} & \text{for } i > j \\ 1 & \text{for } i = j \\ 0 & \text{for } i < j \end{cases} \quad (19)$$

$$U_{ij} = \begin{cases} u_{ij} & \text{for } i \leq j \\ 0 & \text{for } i > j \end{cases} \quad (20)$$

This decomposition is possible if A is non-singular, i.e. $\det(A) \neq 0$. When this decomposition is obtained the matrix equation $A\mathbf{u} = \mathbf{g}$ can be solved in two steps (which are easier to solve than the original one due to the shape of the L and U matrices): First by solving $L\mathbf{y} = \mathbf{g}$ for \mathbf{y} , then solving $U\mathbf{u} = \mathbf{y}$ [2]. I used this method in C++ (using the library Armadillo[3][4] for LU decomposition and solving the equations) for $n = 10, \dots, 10^4$ (higher powers of 10 are not possible; $n = 10^5$ would require 10^{10} numbers to be stored, i.e. roughly 80GB of memory), and compared the CPU time with the previous algorithms.

The algorithm for obtaining the LU decomposition of a matrix requires $\sim n^3$ floating point operations, and solving a matrix equation with an LU decomposed matrix requires $\sim n^2$ such operations[2]. Thus the CPU time used for this algorithm should roughly be proportional to n^3 .

3 Results

The numerical solution (obtained from the general tridiagonal matrix algorithm), compared to the exact solution from equation 3 is shown for $n = 1$, $n = 10$ and $n = 1000$ in figure 1. The CPU time used and the maximum relative error (i.e. the maximum value of $|(u_i^* - u_i)/u_i|$ where \mathbf{u}^* and \mathbf{u} are the numerical and the exact solutions, respectively) in the different solution methods for different values of n is shown in table 1. The results here are largely as expected; the CPU time of both tridiagonal solvers are roughly proportional to n , while for the LU solver it is proportional to somewhere between n^2 and n^3 for larger n ; this is not the case for $n = 10$, probably due to some "minimum time" to run the different functions; due to the relatively small size of this matrix it is likely that the run-time cannot be reduced much more. This is likely also the reason why the CPU time still does not increase as quickly as one would expect for $\sim n^3$ floating point operations even for large n . One can also see a general trend indicating that the specialized version of the tridiagonal algorithm is somewhat faster than the general one, but not quite twice as fast, as I discussed previously.

3.1 Round-off errors

The errors in table 1 also agree with the previous discussion; for small n the errors are identical, i.e. all methods give the same results, however as n grows large the specialized tridiagonal algorithm gives significantly more precise results. This is, as I discussed in the previous section, due to the fact that in the

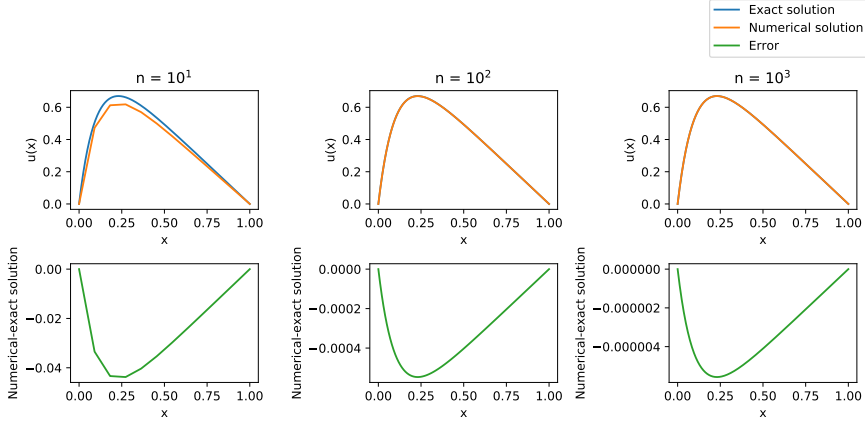


Figure 1: Numerical solution (from the general tridiagonal matrix algorithm) for $n = 1$, $n = 10$ and $n = 1000$, compared to the exact solution.

Table 1: Comparison of the CPU time used and the maximum relative error for the general and specialized tridiagonal matrix algorithms, and the general LU decomposition method, for various numbers of integration points $n + 2$.

n	Tridiagonal (general)		Tridiagonal (special)		LU decomposition	
	t [s]	Rel. error	t [s]	Rel. err.	t [s]	Rel. err.
10	2.32e-7	0.0661	3.23e-7	0.0661	1.10e-4	0.0661
10 ²	1.95e-6	8.17e-4	1.84e-6	8.17e-4	7.99e-4	8.17e-4
10 ³	2.39e-5	8.32e-6	2.12e-5	8.32e-6	0.0594	8.32e-6
10 ⁴	3.07e-4	8.33e-8	2.66e-4	8.33e-8	13.3	8.33e-8
10 ⁵	3.13e-3	1.44e-9	2.61e-3	8.33e-10	-	-
10 ⁶	0.0384	8.40e-7	0.0278	6.87e-11	-	-
10 ⁷	0.369	2.98e-6	0.266	8.13e-10	-	-

specialized version all the \tilde{b}_i 's are calculated exactly, before the main part of the algorithm, whereas in the general version they are calculated in the main part of the algorithm, by repeated subtractions. When this process is repeated a very large number of times small round-off errors in the values of \tilde{b}_i can begin to add up, which by extension leads to errors in the solution.

The sizes of the errors for small n also behave as expected from the approximation for the second derivative I used; the leading term of the error in equation 4 is of order $h^2[2]$ ($h = 1/(n + 1)$), and one can indeed see that for small n the maximum relative error increases roughly by a factor 100 every time n increases by a factor of 10. However as n grows larger this trend stops, and the error starts to rise again. To examine this, I compared the logarithm of the maximum relative error (i.e. the maximum value of $\epsilon_i = \log_{10} |(u_i^* - u_i)/u_i|$) to the logarithm of h , for the specialized tridiagonal algorithm. The results are shown in table 2 and figure 2.

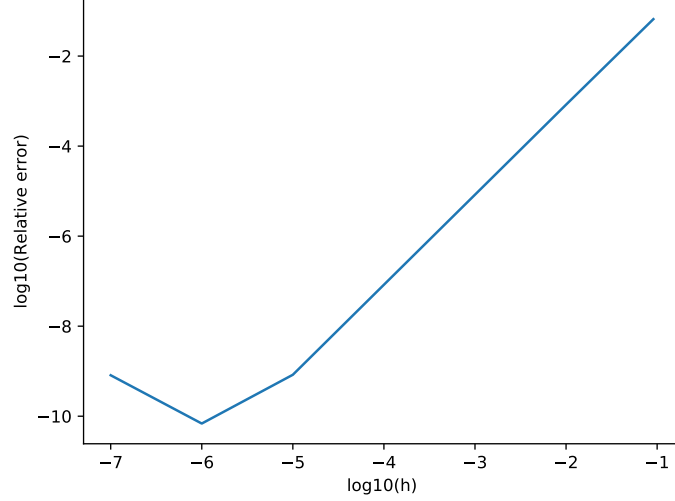


Figure 2: $\epsilon \equiv \max(\epsilon_i)$ (maximum value of the logarithm of the relative error from the specialized tridiagonal algorithm) plotted against the logarithm of the step size h .

Table 2: $\epsilon \equiv \max(\epsilon_i)$ (maximum value of the logarithm of the relative error from the specialized tridiagonal algorithm) vs the logarithm of the step size h .

$\log_{10}(h)$	ϵ
-1.0414	-1.1797
-2.0043	-3.0880
-3.0004	-5.0801
-4.0000	-7.0793
-5.0000	-9.0791
-6.0000	-10.1628
-7.0000	-9.0900

These results show that the logarithm of the relative error is indeed proportional to the logarithm of the step size, with a slope of 2, for relatively large h . As h gets smaller, however, the error starts to increase. This happens because of round-off error; as the number of integration points increases the difference between the elements of the vector \mathbf{g} , and between those of $\tilde{\mathbf{g}}$, become very small. When floating point operations are performed on these (especially subtraction - the difference between two values is given by a small number of significant digits - but this does not occur in this case) small round-off errors can easily arise, and when a very large number of such operations are performed these errors can add up and become significant. As was seen in table 1, this effect is more apparent for the other solution methods, which as I discussed earlier is because the effect of round-off errors adding up also affects the values of $\tilde{\mathbf{b}}$ in those methods.

References

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