

# Fundamentals of Optimization Homework 2

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(4.1)

$$\begin{aligned}
 & \min_{s.t.} -2x_1 + x_5 - 3x_3^- - 2x_4^+ + 2x_4^- + 0x_5 + 0x_6 \\
 & 2x_1 + x_5 - x_3^- - x_4^+ + x_4^- = 1 \\
 & x_1 + 2x_5 - 3x_3^- - x_4^+ + x_4^- + x_5 = -2 \\
 & x_1 - 2x_5 + x_3^- + 2x_4^+ - 2x_4^- + x_6 = 4 \\
 & x_1 \geq 0, x_5 \geq 0, x_3^- \geq 0, x_4^+ \geq 0, x_4^- \geq 0, x_5 \geq 0, x_6 \geq 0
 \end{aligned}$$

(5.1)

$P \in P_1$ , and from definition 6.4 we know  $P_1$  is bounded, so  $P$  is bounded, too. In this case, the bounded polyhedron is polytope

(5.2)

counterexample:

$$P_1 = x_1 \geq -1, x_1 \leq 1$$

$$P_2 = x_2 \geq -1, x_2 \leq 1$$

then  $P \in P_1 \cup P_2$  is a square (bounded) so it is a polytope but neither of  $P_1, P_2$  is polytope actually

In other word, several 'unbounded' polyhedron can bound into a polytope with constraints increasing (adding extra polyhedron with extra constraints, so that the original unbounded constraints may become bounded)

(6.1)

1. prove "P does not contain a vertex  $\rightarrow \text{rank}(A) < n$ "

Since P does not contain a vertex, it contains a 'line'. Suppose  $\hat{x} \in P$ ,  $\exists d \neq 0$ ,  $\lambda \in R$ ,  $\hat{x} + \lambda d \in P$  into  $Ax \leq b$ , we have  $Ad = 0$  &  $d$  is non-zero, so that A could not be full-rank then

2. prove " $\text{rank}(A) < n \rightarrow$  P does not contain a vertex"

If  $\text{rank}(A) < n$ , then A is not a full column rank matrix, and from proposition 13.1, there would be no basic solution, so no vertex any more