

Relational Database's Beta ranges over names

$\tau : \beta := T_{\{1\}} \text{ as } N_{\{1\}}, \dots, T_{\{k\}} \text{ as } N_{\{k\}};$  for  $k > 0, \tau = (\tau_{\{1\}} \text{ to } \tau_{\{k\}}), \beta = (N_{\{1\}} \text{ to } N_{\{k\}})$

Relational Database's Alpha ranges over terms

Each Entry has k terms:

$\alpha : \{\beta\}' := t_{\{1\}} \text{ as } \{N_{\{1\}}\}', \dots, t_{\{m\}} \text{ as } \{N_{\{k\}}\}';$  for:  $m > 0, \alpha = (\tau_{\{1\}} \text{ to } \tau_{\{m\}}), \{\beta\}' = (\{N_{\{1\}}\}' \text{ to } \{N_{\{k\}}\}')$

$n * m * k$  values associated this way

Of course none of these values can be zero for a proper database to form.

One of the problems with SQL Standard is that the semantics of queries is non com-positional, semantically, a query can behave differently depending on the context in which it occurs.

Conditions for our SQL Query language are important to outline:

Our conditions all boil down to boolean states, a conditional is either True, False, or a predicate that evaluates to a boolean True, False.

$\theta := \text{True} | \text{False} | P(t_{\{1\}} \text{ to } t_{\{k\}}),$  Where  $P$  is a collection of predicates

The predicate states are typical in any boolean system. Membership vs. negated membership,  $t$  is [NOT] NULL

A query evaluation where (assuming  $t$  was a member)  $t$ -bar represents the conglomeration of values from  $t$ . So our value either, is/is not in  $Q$ , this assumes  $Q$  exists,  $\bar{t}$  is [NOT] in  $Q$  | Exists  $Q$

Theta has union, intersection, and negation as part of it's conditional functionality,

$\theta$  [AND]  $\theta$

[NOT]  $\theta$

For SQL Queries, it is different from what we have seen before:

$Q := \text{SELECT}[\text{DISTINCT}] \alpha : \{\beta\}' \text{ from } \tau : \beta \text{ where } \theta$

Given our conditional theta, we search the entries terms (Alpha:Beta') from the full database entries (Tau:Beta). This can either be distinct or non-distinct selection.

Then,

$|\text{SELECT}[\text{DISTINCT}] * \text{from}(\tau : \beta) \text{ where } \theta|$

In this usage, '\*' simply implies the return of all the entries terms from the full database, so if you grabbed a dog → You'd get Color, Name, Size or whatever your terms happened to be.

As mentioned before with the conditionals, SQL queries have specific clauses for different combinational structures:

$Q(\text{UNION} | \text{INTERCEPT} | \text{EXCEPT})[\text{ALL}]Q$

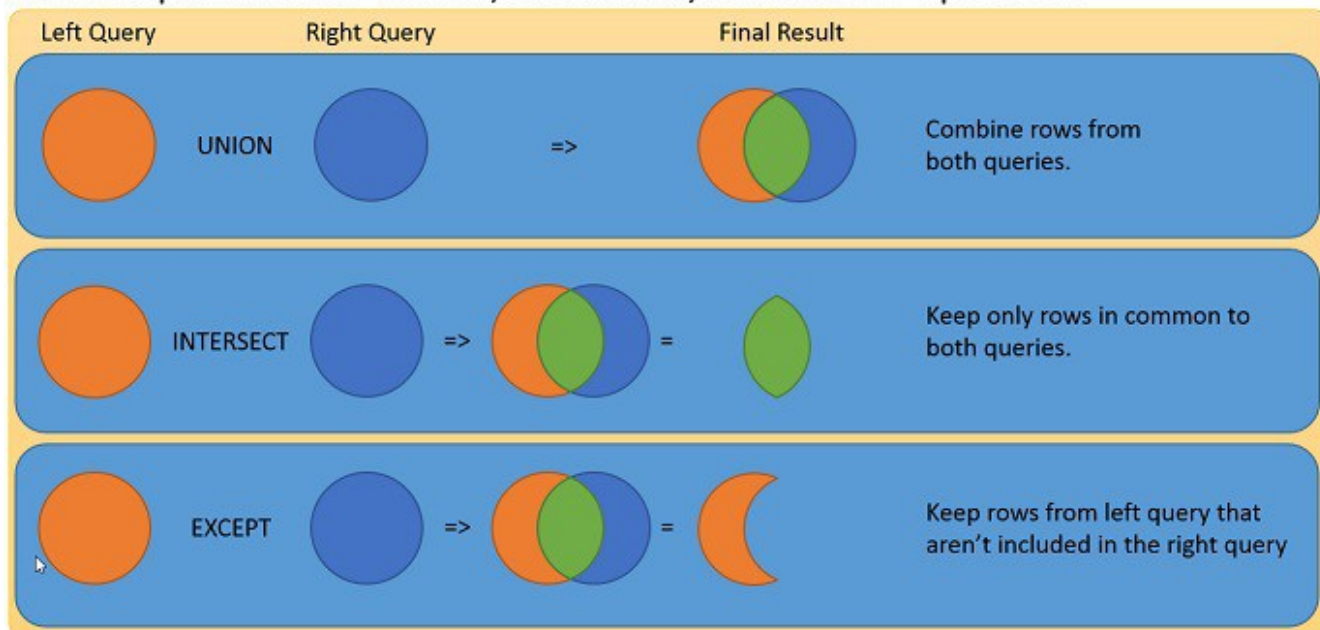
The SQL Union operator returns rows from both databases. If used by itself, UNION returns a distinct list of rows. Using UNION ALL, returns all rows from both databases

The SQL Intercept operator returns rows that the two databases share in common, as well as the ones that are unique in between them.

The SQL Except operator returns rows from the first databases that aren't also in the second database.

Quick Visualization I found:

## Visual Explanation of UNION, INTERSECT, and EXCEPT operators



Source: <https://277dfx2bm2883ohl6u2g3l59-wpengine.netdna-ssl.com/wp-content/uploads/2014/10/UnionIntersectExcept.jpg>

Small Steps:

If  $Q$  steps to some  $E$ , then  $Q$  exists, and the  $E$  it steps to is non-empty

$t \vdash Q \overset{\{n\}}{\rightarrow} E \vdash \mid \text{ EXISTS } \vdash Q \overset{\{n\}}{\rightarrow} \vdash \mid \text{ NOT } \vdash \text{ EMPTY}(E)$

It is possible if our base value, say  $b$  is already boolean,  $T$  or  $F$ , then  $b$  steps to  $b$  in any number of steps  
 $b \overset{\{n\}}{\rightarrow} b \mid b = \text{ TRUE } \vdash \text{ or } \vdash \text{ FALSE}$

If  $t$  IS NULL, it steps to null  $\hat{t}$  if  $t$  stepped to  $\hat{t}$  before:

$t \vdash \text{ IS } \vdash \text{ NULL } \overset{\{n\}}{\rightarrow} \text{ null}(\hat{t}) \mid \text{ IF } \vdash t \overset{\{n\}}{\rightarrow} \hat{t}$

If  $t$  IS NOT NULL, it steps to null  $\hat{t}$  if  $t$  stepped to  $\hat{t}$  before:

$t \vdash \text{ IS } \vdash \text{ NOT } \vdash \text{ NULL } \overset{\{n\}}{\rightarrow} \text{ NOT } \vdash \text{ null}(\hat{t}) \mid \text{ IF } \vdash t \overset{\{n\}}{\rightarrow} \hat{t}$

We can step the whole predicate base to  $\hat{t}$  if entry  $I$  steps to  $\hat{I}$  and entry  $I'$  steps to  $\hat{I}'$ :

$P(t_1, \dots, t_i) \overset{\{n\}}{\rightarrow} P(\hat{t}_1, \dots, \hat{t}_i) \mid \text{ IF } \vdash t_i \overset{\{n\}}{\rightarrow} \hat{t}_i, t_i' \overset{\{n\}}{\rightarrow} \hat{t}_i'$

While in the state with conditionals ( $\theta$ ) evaluating to  $\theta'$  we have:

$\mid \text{ For } \vdash \Theta_{\{1\}} \overset{\{n\}}{\rightarrow} \Theta'_{\{1\}}, \Theta_{\{2\}} \overset{\{n\}}{\rightarrow} \Theta_{\{2\}} \mid$

$\mid \Theta_{\{1\}} \vdash \text{ AND } \vdash \Theta_{\{2\}} \overset{\{n\}}{\rightarrow} \Theta_{\{1\}} \wedge \Theta_{\{2\}} \mid$

$\mid \Theta_{\{1\}} \vdash \text{ OR } \vdash \Theta_{\{2\}} \overset{\{n\}}{\rightarrow} \Theta_{\{1\}} \vee \Theta_{\{2\}} \mid$

$\mid \text{ NOT } \vdash \Theta_{\{1\}} \overset{\{n\}}{\rightarrow} \text{ NOT } \vdash \Theta_{\{1\}} \mid$

Such that the union of  $\theta_1$  and  $\theta_2$  can evaluate to the union of  $\theta_1'$  and  $\theta_2'$  over  $n$  steps. The same applies for intersection and NOT, but not EXCEPT