

# Chapter 10

## Other Public-Key Cryptosystems



정보보안

Chapter 10  
Other Public-Key Cryptosystems

$$n = pq$$
$$y = g^x \bmod p$$

# Diffie-Hellman Key Exchange 공개키 암호

- History Public key encryption

- The scheme was first published by Whitfield Diffie and Martin Hellman in 1976, although it had been separately invented a few years earlier within GCHQ (The Government Communications Headquarters) but was kept classified.

- Key exchange [만약식이 없는데 키를 교환] 두사람만이 알수있는 암호 생성

- The Diffie–Hellman key exchange method allows two parties that have no prior knowledge of each other to jointly establish a shared secret key over an insecure communications channel.

- The (Computational) Diffie-Hellman problem

- Let  $g$  be a generator of some group  $G$ . For randomly chosen integers  $x$  and  $y$ , the DHP is stated informally as follows.
  - Given an element  $g$  and the values  $g^x$  and  $g^y$ , what is the value of  $g^{xy}$ ?
- Over many groups, the DHP is almost as hard as the DLP.

$g^{x1}, g^{x2} \rightarrow g^{x1x2}$  맞출수 있나?  $g^{283} \text{ mod } P, g^{182}$  283 안알려주고 결과도 안알려주고 맞추게 하기

# Diffie-Hellman Key Exchange

- Scheme

- Alice and Bob agree on a cyclic group  $G$  of (prime) order  $q$  and a generator  $g \in G$ .  
(Note: A cyclic group of prime order is recommended.)  
 $0 \sim q-1$
- Alice chooses a random integer  $1 \leq a < q$  and sends  $pk_A = g^a$  to Bob.
- Bob chooses a random integer  $1 \leq b < q$  and sends  $pk_B = g^b$  to Alice.
- Alice computes  $(pk_B)^a = (g^b)^a = g^{ab}$ .  
숫자가 크기 때문에 0이 나올 확률 거의 없음.  
등호는 맞으나 대나.
- Bob computes  $(pk_A)^b = (g^a)^b = g^{ab}$ .

Alice

Public key: PK

Bob

$g, P$ 는 공개된 값

$$pk_A = g^a \bmod P$$

$$\xrightarrow{pk_A}$$

$$pk_B = g^b \bmod P$$

$$\xleftarrow{pk_B}$$

$$\begin{aligned} \text{BOB은} \\ (pk_A)^b \\ = (g^a)^b \\ = g^{ab} \end{aligned}$$

$a$ 는 BOB도 모름  
 $g^b$  해서  $B$ 를 보냄  
결과만 알. 서로  $a, b$ 는  
모름

# Diffie-Hellman Key Exchange

- Example

1. Alice and Bob agree on  $G = Z_{23}^*$  ( $p = 23, q = 22$ ) and  $g = 5$ .

2. Alice chooses  $a = 6$  and sends  $pk_A = 5^6 \bmod 23 = 8$ . - Alice 가 보낸

3. Bob chooses  $b = 15$  and sends  $pk_B = 5^{15} \bmod 23 = 19$ . - Bob 이 보낸

4. Alice computes  $(pk_B)^a = 19^6 \bmod 23 = 2$ .

5. Bob computes  $(pk_A)^b = 8^{15} \bmod 23 = 2$ .

두사람만 2라는 것을 알.

$5^{22} \equiv 1$  generator

$5^6$  인지  $5^{15}$  인지 모름  
실제로는 10진수로 천자리

# ElGamal Encryption Diff, hellman 의해 생성. 실용적

- Key generation

- Alice chooses a cyclic group  $G$  of (prime) order  $q$  and a generator  $g \in G$ .
- Alice chooses a private key  $x \in \{1, \dots, q-1\}$  randomly and computes her public key  $y = g^x$ .  $x$ 는 비밀인 알고리즘.  $b$ 를 한번 쓰고 다시 쓰는 것  $y^r = g^{xr}$

- Encryption

- Bob chooses a random integer  $r \in \{1, \dots, q-1\}$  and computes  $c_1 = g^r$ .
- To encrypt a message  $m \in G$ , Bob calculates  $c_2 = my^r (= mg^{xr})$ .
- Bob sends the ciphertext  $(c_1, c_2)$  to Alice.

- Decryption pk를 다시 보내는게 아닌 한번 쓰고 다시 쓰는 것을 뺀

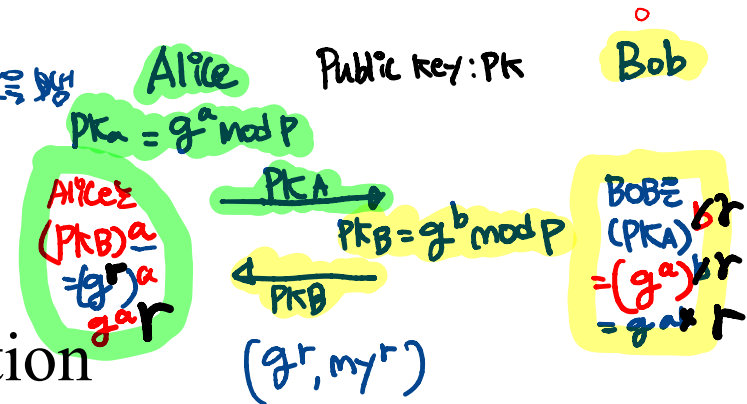
- Alice computes  $c_2 (c_1^x)^{-1} = mg^{xr} (g^{rx})^{-1} = m$ .

Extend Euclid's

a 대신  $x$

- Decisional Diffie-Hellman assumption

- The DDH assumption states that the probability distributions  $(g^a, g^b, g^{ab})$  and  $(g^a, g^b, g^c)$  are computationally indistinguishable, where  $a, b$ , and  $c$  are randomly chosen from  $Z_q$ .



# ElGamal Encryption

- Example

(Key generation)

- Alice chooses  $G = Z_{19}^*$  ( $p = 19$ ,  $q = 18$ ) and  $g = 10$ .
- Alice chooses the private key  $x = 5$  and computes the public key  $y = g^x = 10^5 \bmod 19 = 3$ .

(Encryption) *rand number*

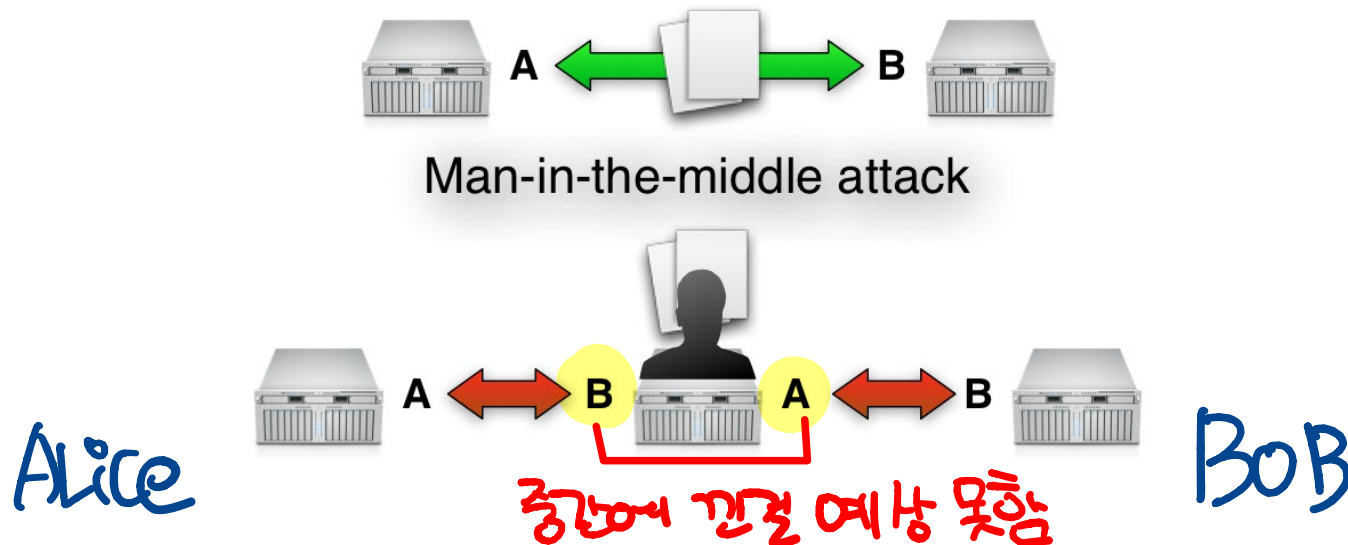
- Bob chooses  $\hat{x} = 6$  and computes  $c_1 = g^r = 10^6 \bmod 19 = 11$ .
- To encrypt  $m = 17$ , Bob calculates  $c_2 = my^r = 17 \cdot 3^6 \bmod 19 = 5$ .
- Bob sends the ciphertext  $(c_1, c_2) = (11, 5)$  to Alice.

(Decryption)

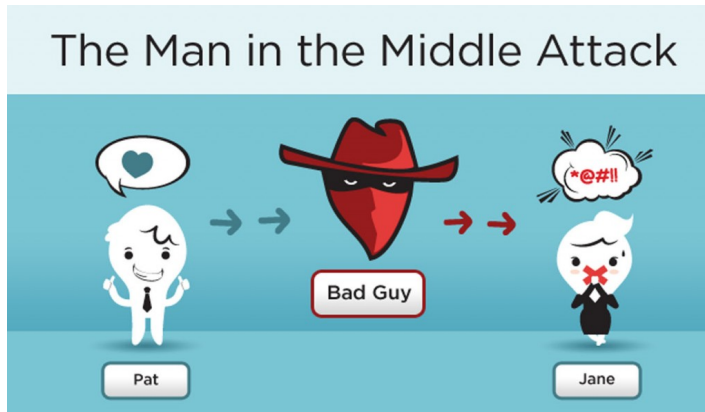
- Alice computes  $m = c_2 (c_1^x)^{-1} = 5 \cdot (11^5)^{-1} \bmod 19 = 17$   
where  $5 \cdot (11^5)^{-1} \equiv 5 \cdot (7)^{-1} \equiv 5 \cdot 11 \equiv 17 \pmod{19}$ .

# Man-in-the-Middle Attack

- In cryptography and computer security, a **man-in-the-middle attack** (often abbreviated to MITM, MitM, MIM, MiM or MITMA) is an attack where the attacker secretly relays and possibly alters the communication between two parties who believe they are directly communicating with each other.



# Man-in-the-Middle Attack



DH

Alice Carol Bob  
 $g^a$   $C$   $g^b$   
 $\rightarrow g^a$   $\leftarrow b$   
 $g^c$   $g^c$   
 $\leftarrow g^a$   $g^b \rightarrow$

Alice와 Bob은 서로 갖

는 것을 알. Carol은 모름 (1-2-3)

