

Design for Testability

Introduction to testing
integrated circuit designs

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2 lectures about testing circuits

■ PART 1: What to test

- Why is testing so important
- Which faults can happen
- Generation of test patterns

■ PART 2: How to test

- How to test a design
- How good is my test
- How to test hardware



Motivation

Design Phase	Error (mechanism), failure, bug, improper definition, ...	Approx. cost
Specification	Functionality, Performance, Testability, conflicting definitions, ..	1 €
Modeling, Design	Interoperability, Functionality, Performance,	1000 €
Prototype	Verification errors, Performance, ...	100 k€

Impl. Phase	Error (mechanism), failure, bug, improper implementation, ...	Approx. cost
Wafer	Yield, Silicon properties, Dust, Speed, Noise, ...	1 €
Chip	Packing, Bonding, ...	10 €
Module	Soldering, ESD impact, ...	100 €
Subsystem	Cable connections, Connectors, ...	1000 €
After deployment	Stress, Vibration, Reliability, ...	10 k€

- Example: Intel FDIV bug (1994)
 - Logic error not caught until > 1M units shipped
 - Recall cost \$450M



Motivation

- **Failure rate** for a full system

- Probability for a system failure

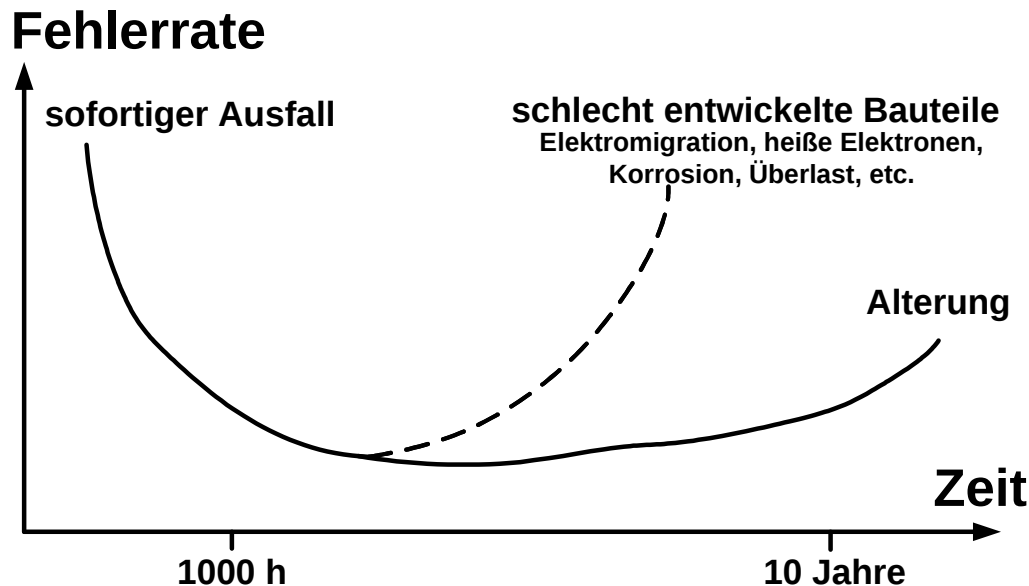
$$p(\text{System} | \text{faulty}) = 1 - p(\text{System} | \text{ok}) = 1 - \prod_{i=1}^N p_i(\text{component} | \text{ok})$$

- 70 components 99% OK → Full system is OK for 50%
- 70 components 90% OK → Full system is OK for 0%



Motivation

- **Reliability** of integrated circuits
 - “Bathtub” characteristic



- Number of components which fails within 1000 operation hours: 1...5 %
- Burn-In: Operating the ICs at 125° shortens this 1000 hours interval to about 24 hours.

Motivation

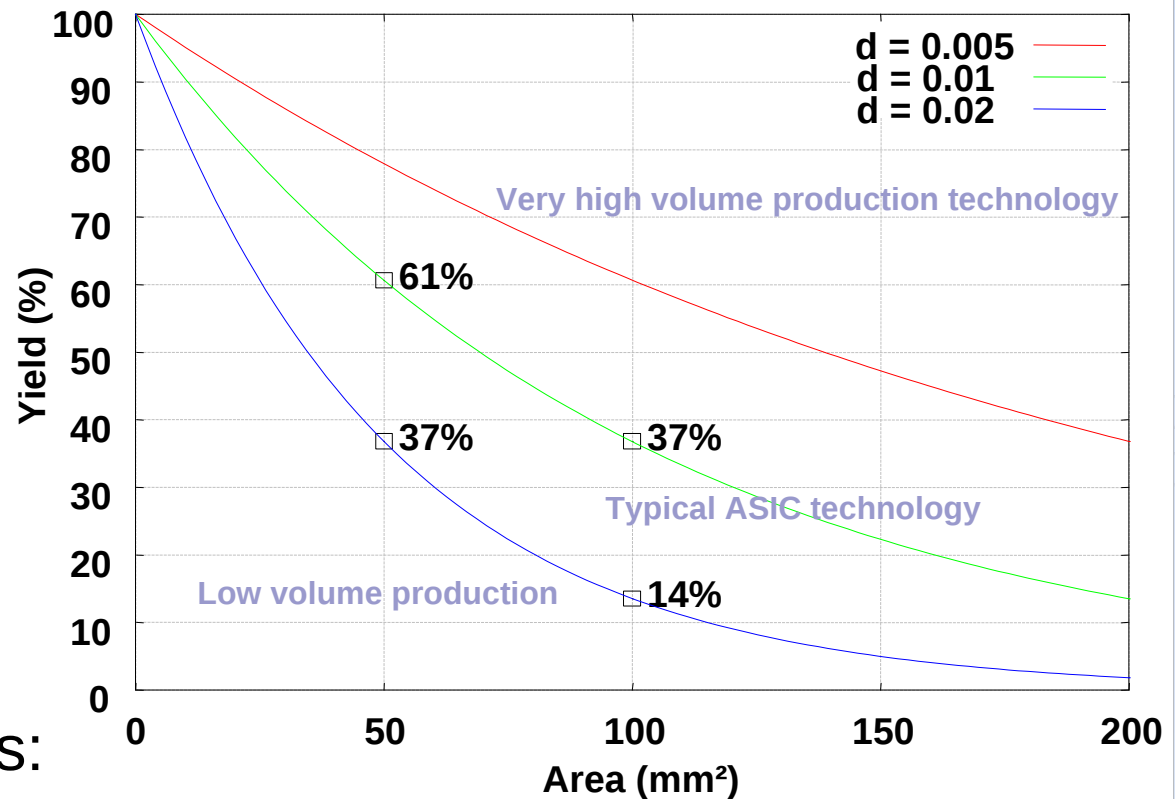
■ Yield:

$$100 \cdot e^{-d \cdot A}$$

d =defects/mm²

typically $d =$

0.005 ... 0.02



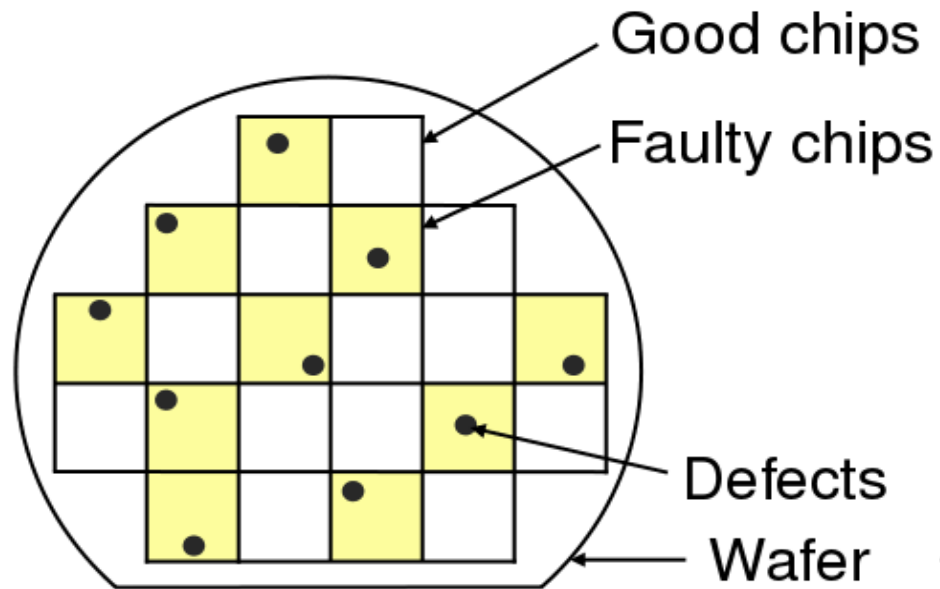
■ Yield improvements:

- Technology improvements
- Redundancy (e.g. memory)

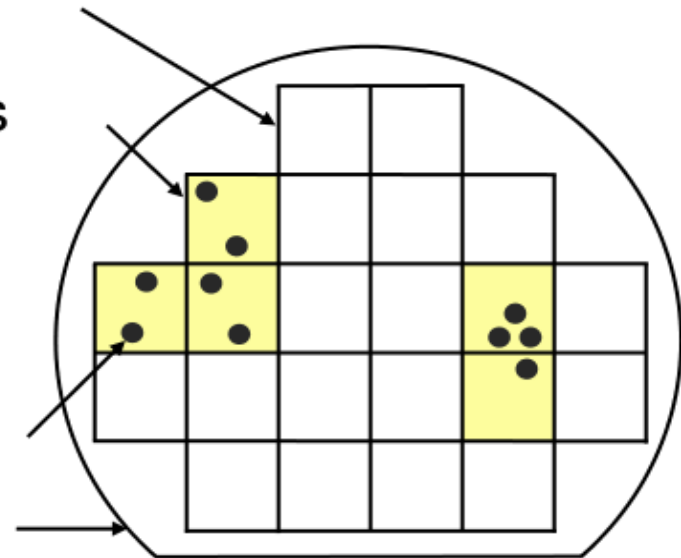
■ Cost: 100mm² → 50mm² chip with yields from figure

- $100\text{mm}^2/50\text{mm}^2 \times 0.61/0.37 = 3.4$ ($d=0.01$)
- $100\text{mm}^2/50\text{mm}^2 \times 0.37/0.14 = 5.4$ ($d=0.02$)

Chip Yield



$$\text{Wafer yield} = 12/22 = 0.55$$

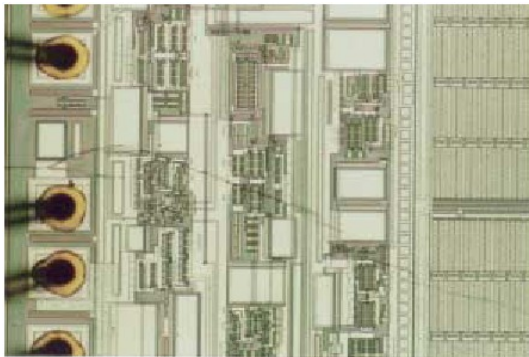


$$\text{Wafer yield} = 17/22 = 0.77$$

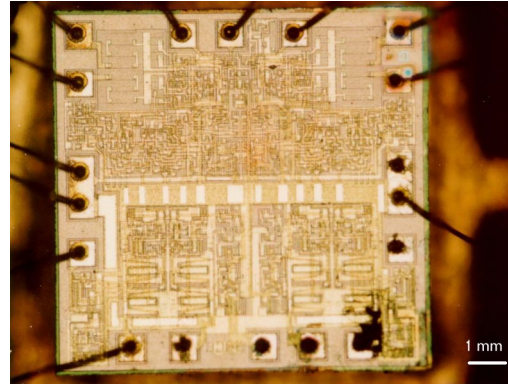
Defects



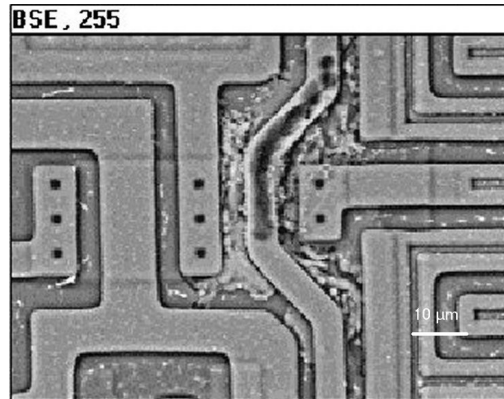
1) Kurzschluß zwischen Leiterbahnen durch Staub



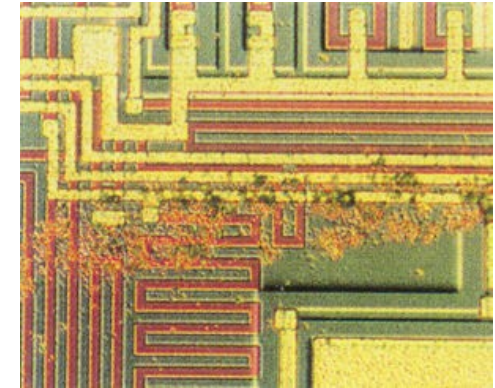
4) gebrochenes Die (mechanische Verspannungen)



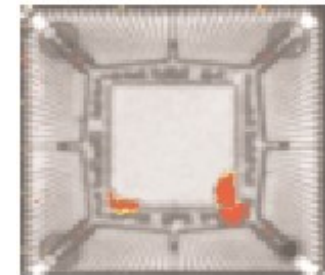
2) Elektrischer Durchschlag (rechts unten)



5) überlastete Leiterbahn



3) Verunreinigungen des Wafers (Diffundieren ins Polysilizium)



6) ins Gehäuse eingedrungener Schmutz (sonographisches Bild)

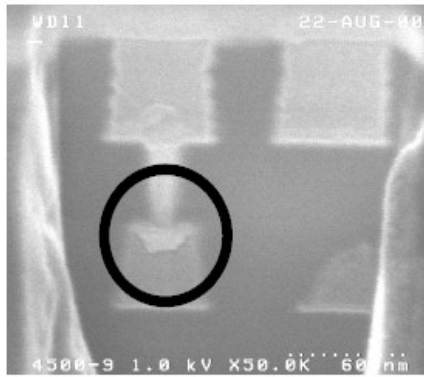
www.sony.co.jp/en/Products/SC-HP/PDF

www.era.co.uk/product

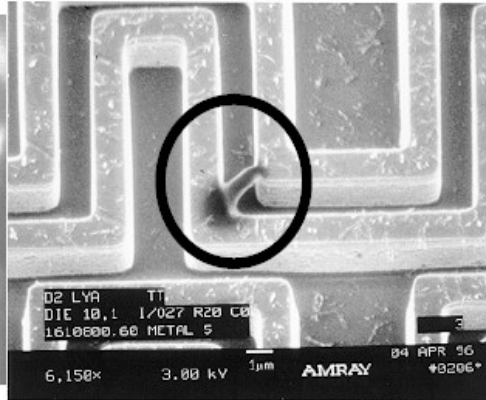
www.fabtech.org/features/tap/articles

www.empf.org/html

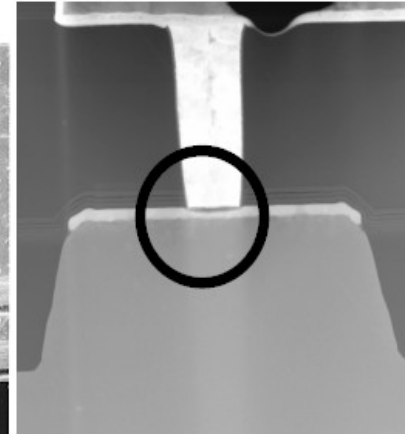
Defects



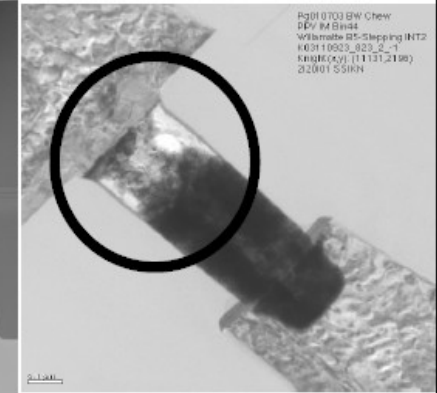
Metal 1 Shelving



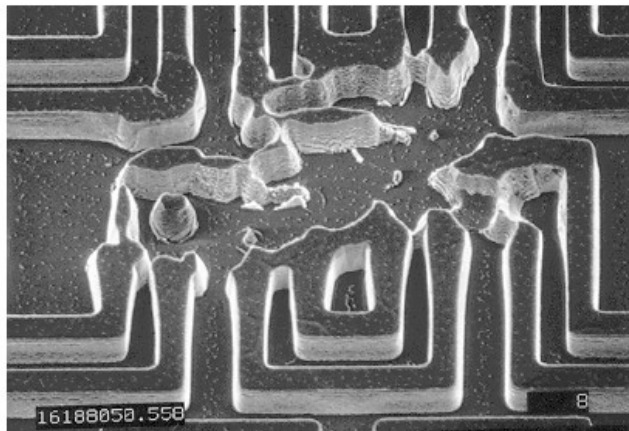
**Metal 5 film particle
(bridging defect)**



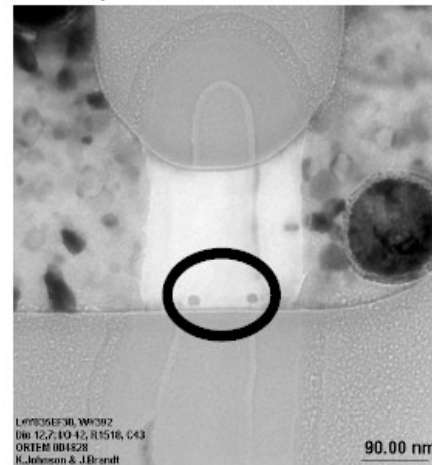
Open defect



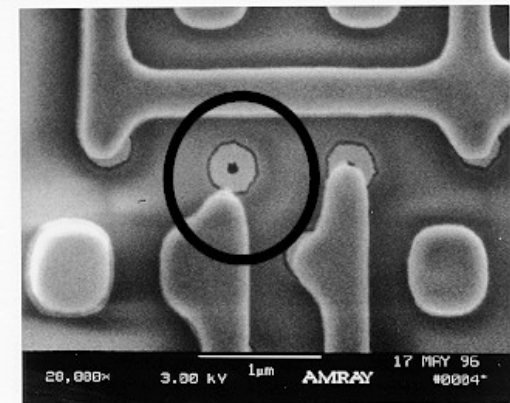
**Spongy Via2
(Infant mortality)**



**Metal 5 blocked etch
(patterning defect)**



**Spot defects
"Co" Defect under Gate**



**Metal 1 missing pattern
(open at contact)**

SEM images courtesy Intel Corporation

Terminology

■ Defect:

- Is an unintended difference between the implemented hardware and the design.
- Defects may be caused by disturbances during manufacturing or usage of the device.

■ Fault:

- The impact of a hardware defect to an abstracted functional level
- Fault indicates an event

■ Error:

- Wrong (unspecified) behavior of the defective system
- An error is caused by a fault
- Error indicates a state.



Terminology

■ **Controlability**

Effort for reaching a defined state in the system. This might be the initial state for a testsequence. Ease of forcing the output of a node to 0 or 1 by driving the inputs.

■ **Observability**

The ability to observe the state or logic values if internal nodes.
Effort for observing effects caused by a test sequence

Combinatorial logic is usually easy to observe and control. FSM can be very difficult, requiring many cycles to reach a desired initial test state.

■ **Testability**

- Ability to generate test sequences for a dedicated focus and/or to reach a required test coverage.
- Localization of faults. A test must be designed that the cause (component, signal, etc.) of a fault can be identified.
- Testability = Controlability + Observability



Terminology

- Fault Coverage

$$FC = \frac{N(\text{observable System errors})}{N(\text{possible System errors})} \cdot 100$$

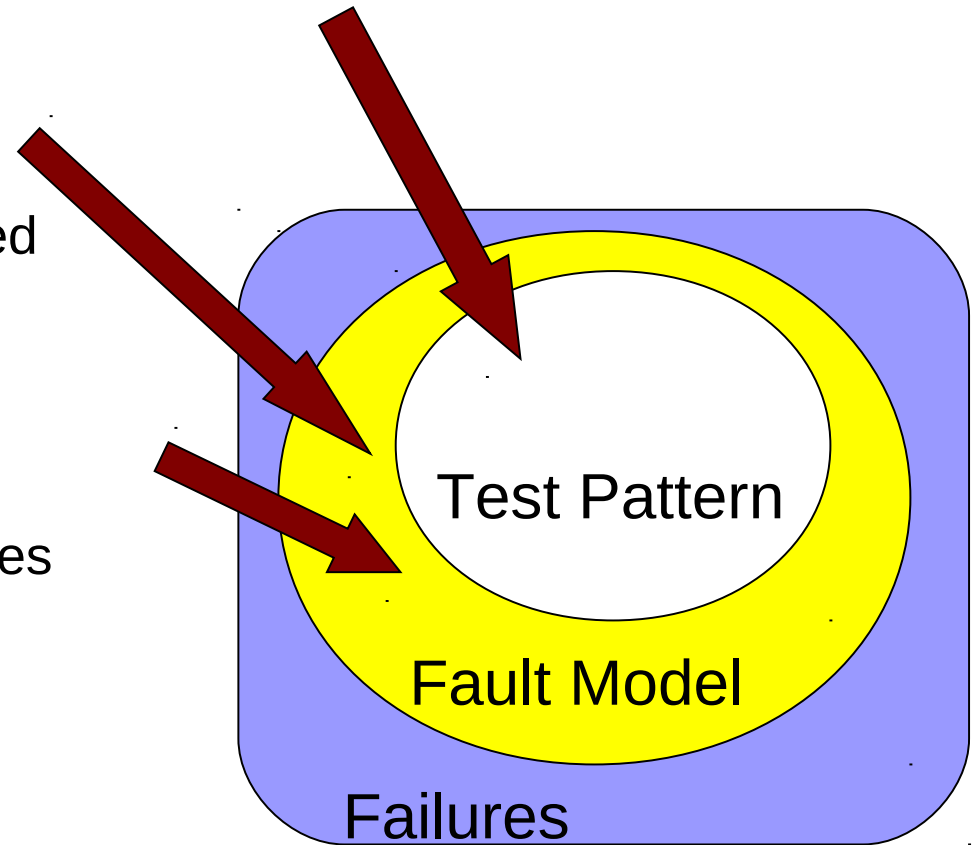
- Fault Model:

Hypothesis on which defects can occur at a used process technology

- Fault Model +

System structure:

Results in all possible cases a system can reach malformed behavior.

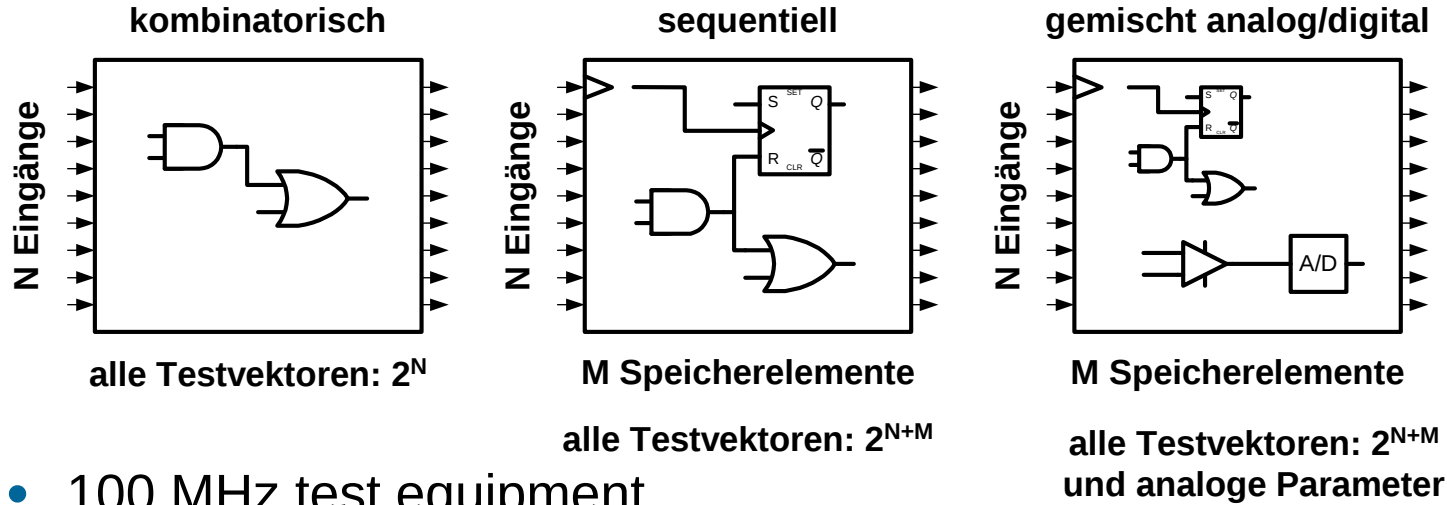


Overview

- Testing is a very important phase of the development and design of a new IC
- Testing is cost intensive and complex
 - Cost for testing is about **2/3 of total ASIC development cost**
- Testing might be a bottleneck for IC fabrication
- Test time is cost intensive and must be reduced
 - Trade-off between coverage (number of testcases) and test time.
- Test-effort
 - Sometimes complex and expensive test equipment is required (analog testing, RF testing, ...)
 - External test equipment for digital functionalities
100 k€ to 10 M€; 32k – 100M testvectors
 - On-chip test equipment → chip area

Overview

- What is being tested?



- 100 MHz test equipment
 - $N=32 \rightarrow$ testtime = 40 seconds
 - $N=63 \rightarrow$ testtime = 6000 years
- Topology of the DUT (device under test) must be used to reduce the number of testcases.
- Synchronization of analog and digital Signals

Testability Measures

- TMEAS [Stephenson & Grason, 1976]
- SCOAP [Goldstein, 1979]
- TESTSCREEN [Kovijanic 1979]
- ...

- Importance of testability measures
 - Guide designers to improve testability in their designs
 - Test generation algorithms using heuristics apply testability measures to speed up the test generation process (e.g. in search decisions)

SCOAP

- Sandia Controllability Observability Analysis Program
- Using integers to reflect the difficulty of controlling and observing the internal nodes.
- Higher numbers indicate more difficult to control or observe.
- Applicable to both combinational & sequential circuits.
- SCOAP measurements on each line
 - Combinational 0-controlability, CC0(I)
 - Combinational 1-controlability, CC1(I)
 - Combinational observability, CO (I)
 - Sequential 0-controlability, SC0(I)
 - Sequential 1-controlability, SC1(I)
 - Sequential observability, SO(I)



SCOAP

- Combinational measures related to number of signals to be manipulated to control or observe 1
- Sequential measures indicating number of clock cycles needed to control or observe signals on line 1
- Controlability ranging from 1 to ∞
- Observability ranging from 0 to ∞
- High measures indicating difficulties with controlling or observing the given line



SCOAP – combinatorial Controlability

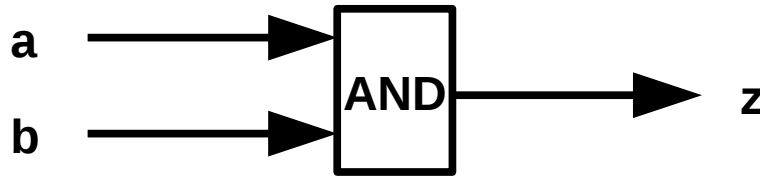
- Computed at propagating from inputs to outputs
- Step1: For all primary inputs set all $CC0 = 1$ and all $CC1 = 1$
- Step 2: Traverse in level order (Level of logic gate: max distance of its logic inputs) through circuit towards primary outputs updating controllability measures
- Step 3: For each traversed logic gate:
 - If logic output produced by setting only one input to controlling value then
output controllability = min(input controllability)+1
 - If logic output only obtained by setting all inputs to non controlling values then
output controllability = sum(input controllabilities)+1
 - If possible to control output by multiple input sets (XOR: “01” or “10” cause output 1) then
output controllability = min(controllabilites of input sets)+1

SCOAP – combinatorial Observability

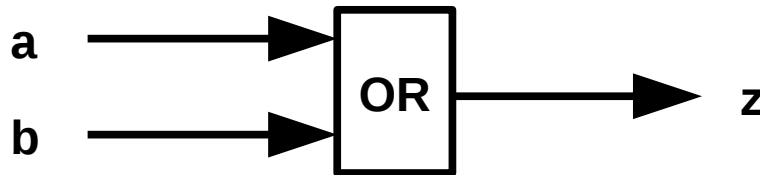
- Computed at propagating from outputs to inputs
- Step1: For all primary outputs set all CO= 0
- Step 2: Traverse through circuit towards primary inputs updating observability measures
- Step 3: For each traversed logic gate:
 - **CO = observability of output +
difficulty in setting all inputs to non-controlling values +
1 to accommodate for logic depth**
 - No distinction between 0 and 1 observability



SCOAP - Examples

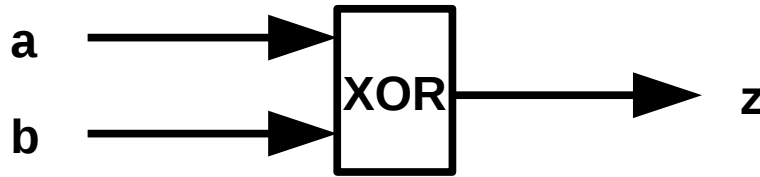


- $CC0(z) = \min(CC0(a), CC0(b)) + 1$
- $CC1(z) = CC1(a) + CC1(b) + 1$
- $CO(a) = CO(z) + CC1(b) + 1$
- $CO(b) = CO(z) + CC1(a) + 1$

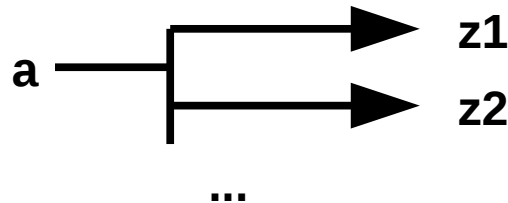


- $CC0(z) = CC0(a) + CC0(b) + 1$
- $CC1(z) = \min(CC1(a), CC1(b)) + 1$
- $CO(a) = CO(z) + CC0(b) + 1$
- $CO(b) = CO(z) + CC0(a) + 1$

SCOAP - Examples

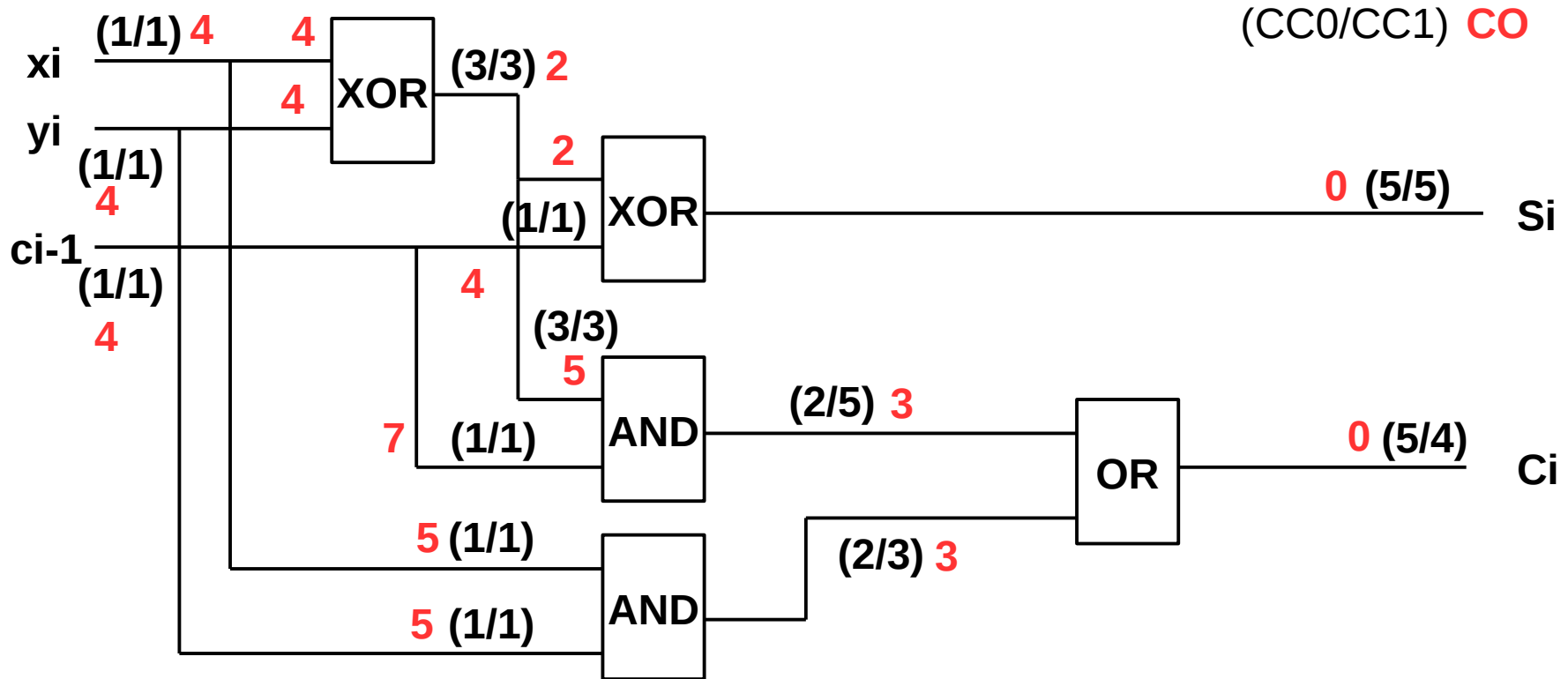


- $CC0(z) = \min(CC0(a)+CC0(b), CC1(a)+CC1(b)) + 1$
- $CC1(z) = \min(CC1(a)+CC0(b), CC0(a)+CC1(b)) + 1$
- $CO(a) = CO(z) + \min(CC0(b), CC1(b)) + 1$
- $CO(b) = CO(z) + \min(CC0(a), CC1(a)) + 1$



- $CO(a) = \min(CO(z1), CO(z2), \dots)$

SCOAP – Example Adder



- $CC0(si) = 5$ $CC1(si) = 5$ $CO(si) = 0$
- $CC0(ci) = 5$ $CC1(ci) = 4$ $CO(ci) = 0$
- $CC0(xi) = 1$ $CC1(xi) = 1$ $CO(xi) = 4$
- $CC0(yi) = 1$ $CC1(yi) = 1$ $CO(yi) = 4$
- $CC0(ci-1) = 1$ $CC1(ci-1) = 1$ $CO(ci-1) = 4$

Basic Models for defects

■ **Hard Faults** (Catastrophic faults)

- Large deviations of the designed system behavior. May be electrical characteristics, structural deviations, mechanical faults, etc.
- Stuck at short: Unspecified electrical connections inside a component, at component interconnects, etc.
- Stuck at open: Unspecified disconnection.
- Punch through: Unspecified electrical connection caused by a high voltage impact.

■ **Soft faults** (Deviations inside a specified tolerance)

- Variability in electrical, mechanical, structural characteristics. Process fluctuations
- Silicon (material) characteristics: Carrier mobility, doping concentrations, etc.



Fault Models

- Layout level fault model

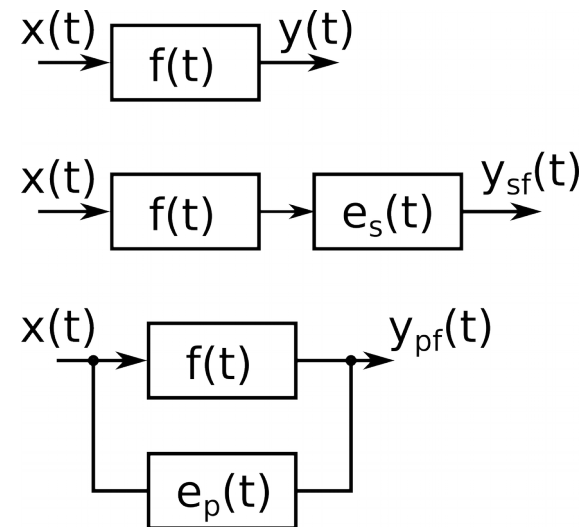
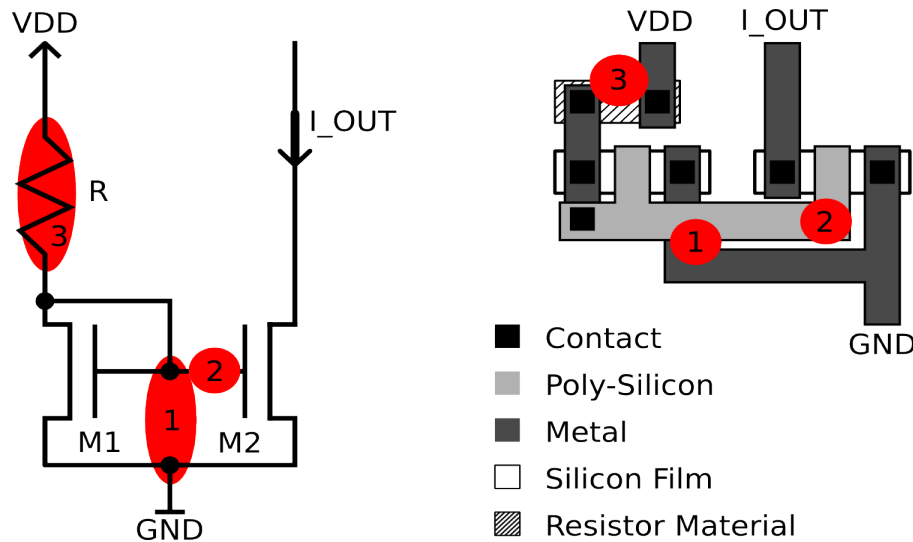
Schematic as a unique model – layout as a realization

- Behavioral-Functional fault model

Serial and parallel fault modeling

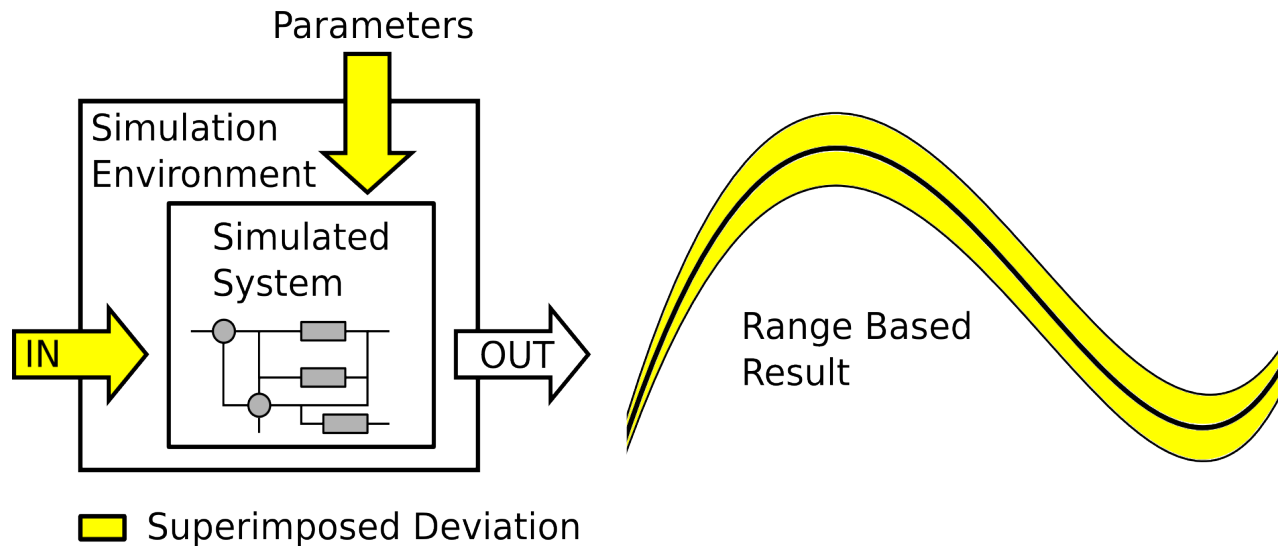
- Parametric fault model

Deviation modeling – range based analysis of the system



Fault Models

- Layout level fault model
Schematic as a unique model – layout as a realization
- Behavioral-Functional fault model
Serial and parallel fault modeling
- Parametric fault model
Deviation modeling – range based analysis of the system



Fault Types

- Permanent faults

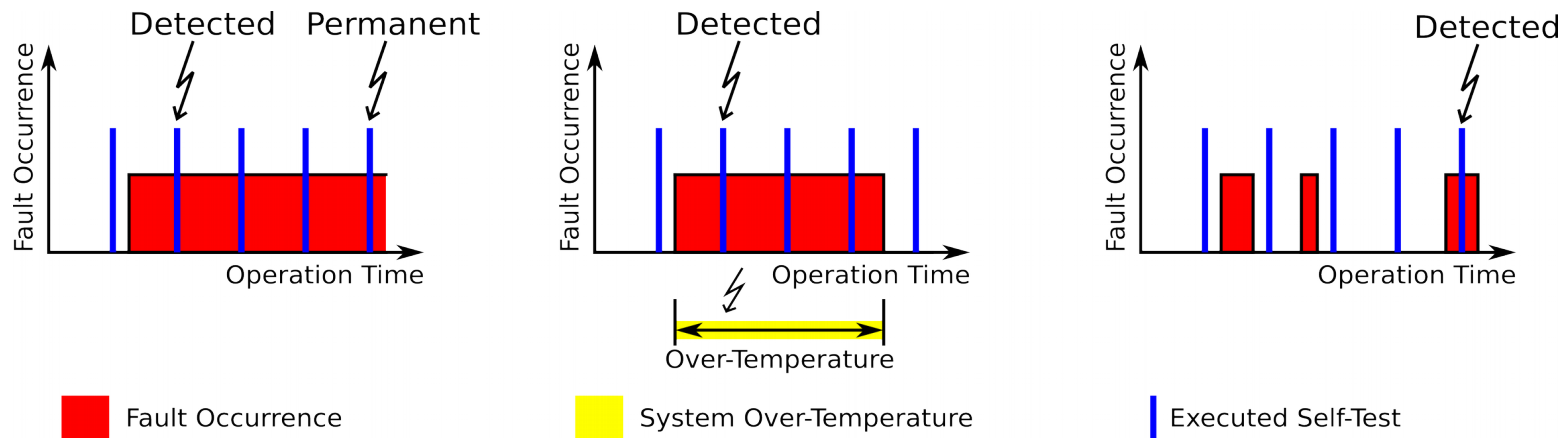
Errors at the production or during operation. Can be detected by cyclic self tests

- Intermittent faults

Correlated with environmental conditions. Additional measurement of system parameters to differ from a permanent fault

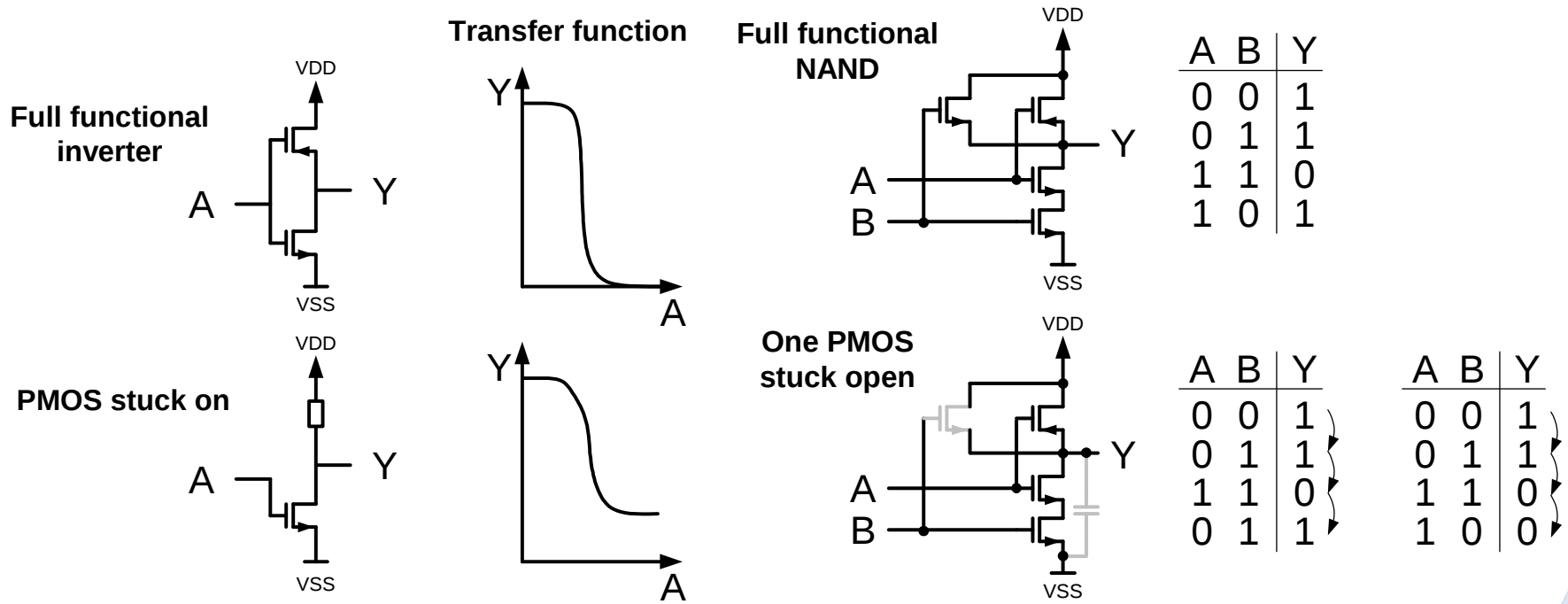
- Transient faults

completely random manner. Hard to detect if fault time is shorter than self test period.



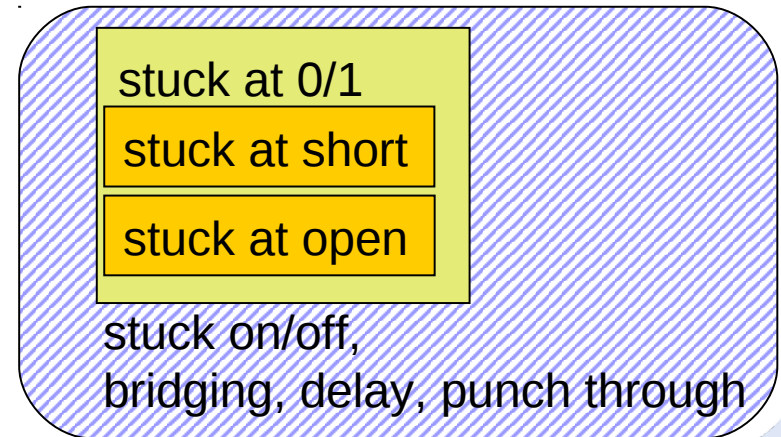
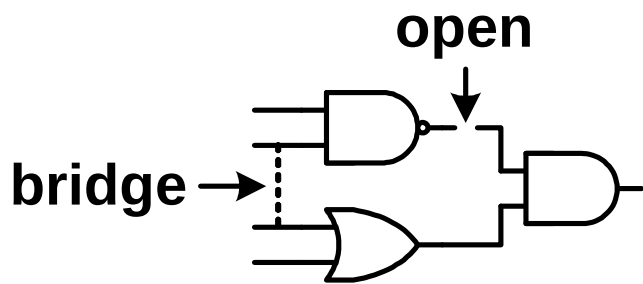
CMOS stuck at faults

- It is assumed that a fault in a logic gate results in one of its inputs or the output is fixed to either logic 0 or 1
- A single line in the structure is connected to 0 or 1



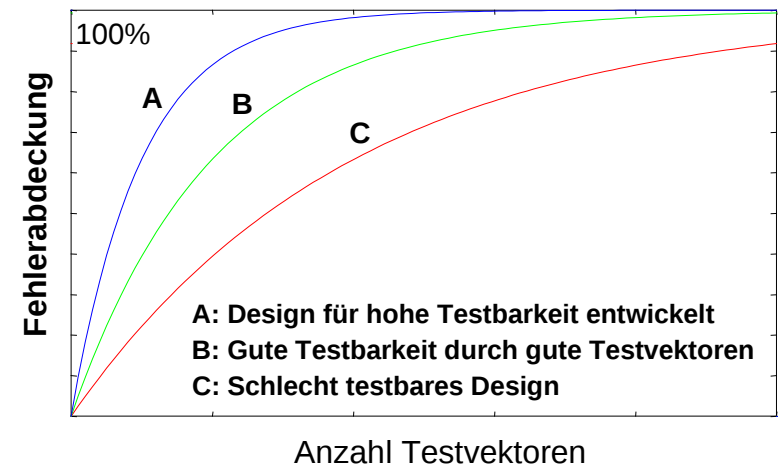
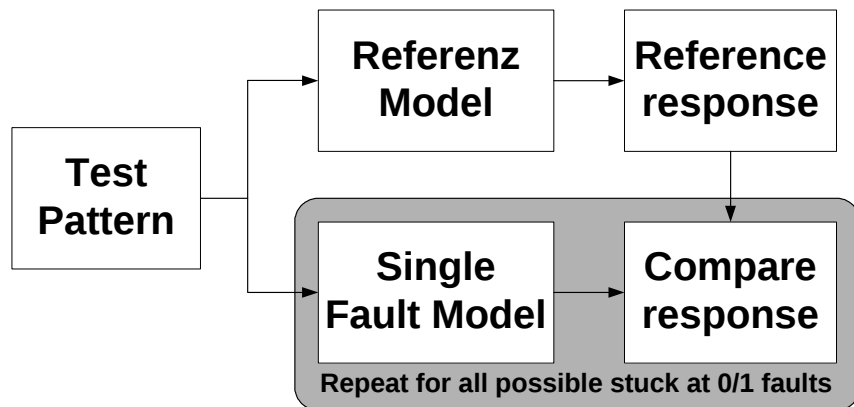
CMOS stuck at faults

- Complexity compared to a parametric fault model is reduced.
- For each component (gate) in the netlist connections to 0 and 1 are considered.
- Reduced effort for fault simulations and achieving a requested coverage.
- For VLSI designs sometimes other parametric fault models are too computation intensive.
- Covers 80-90 % of classical defects.



Fault simulation

- For the calculation of the required test coverage
- Returns just detectable faults
- Conclusions regarding the quality of testvectors
- For complex designs long simulation runs



Generation of testvectors

- Testcases are designed by a Test team not by the designers.
 - Require a functional description of the DUT
 - Use script languages for modeling test scenarios.
- Design of testvectors is guided by CAE systems and fault simulation software.
- Parts of the testvectors can be derived from functional simulation and system verification processes.



Test and Test Set

- A **test** t for a fault α in circuit C (realizing function f) is an input combination X for which the output of C is not equal when α is present than when it is not.

$$f(X) \text{ XOR } f_{\alpha}(X) = 1$$

- A **test set** T is a set of t for a class of faults A ,

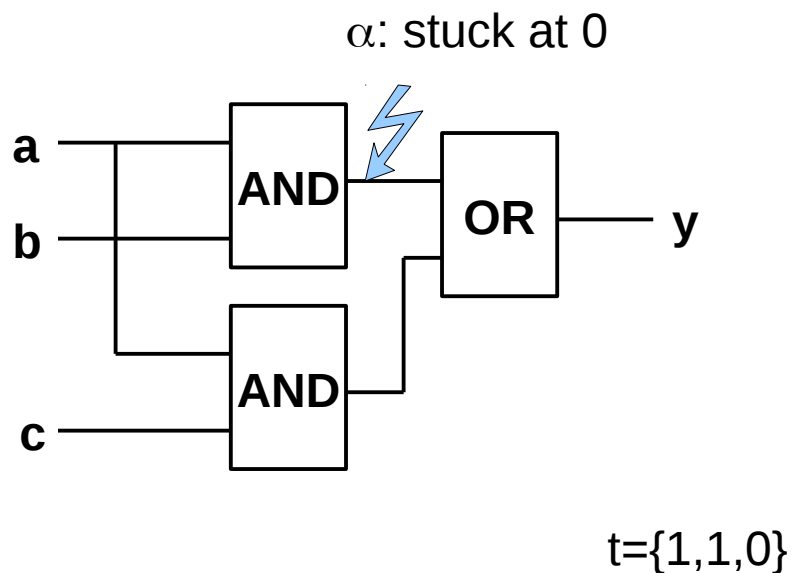
$$\forall \alpha \in A, \exists t \in T$$

The test set for a fault α is $T_{\alpha} = f \text{ XOR } f_{\alpha}$



ATPG – from truth table

- Consider a single fault and evaluate the truth table for good and faulty circuit

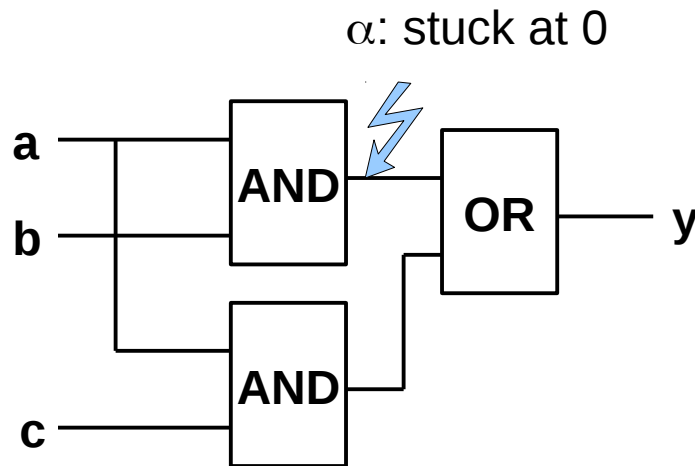


a	b	c	f	$f\alpha$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	0	0
1	0	0	0	0
1	0	1	1	1
1	1	0	1	0
1	1	1	1	1

- T for all stuck at faults at each line results analogously where $T = \min\{t_1, t_2, \dots\}$

ATPG – from boolean equations

- Consider a single fault and evaluate boolean equations for good and faulty circuit



$$f = (a \wedge b) \vee (a \wedge c), f_{\alpha} = a \wedge c$$

$$T_{\alpha} = f \text{ XOR } f_{\alpha}$$

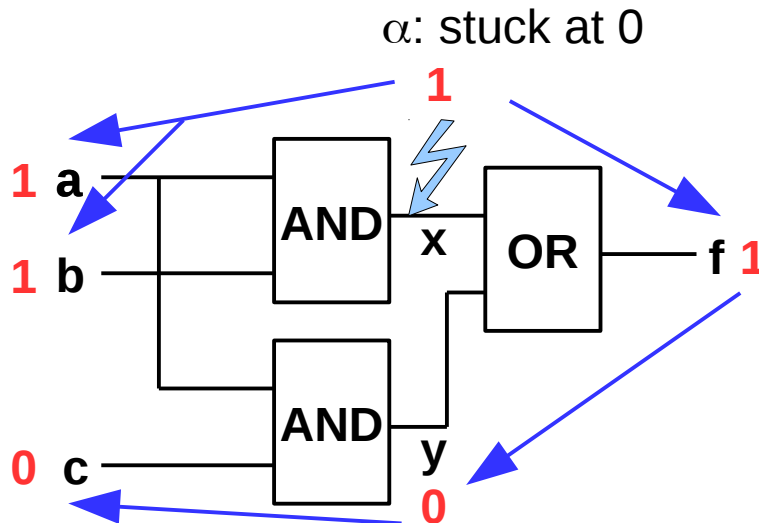
$$T_{\alpha} = (((a \wedge b) \vee (a \wedge c)) \wedge \neg(a \wedge c)) \vee (\neg((a \wedge b) \vee (a \wedge c)) \wedge (a \wedge c)) = 1$$

$$T_{\alpha} = a \wedge b \wedge \neg c = 1$$

$$T_{\alpha} = \{1, 1, 0\}$$

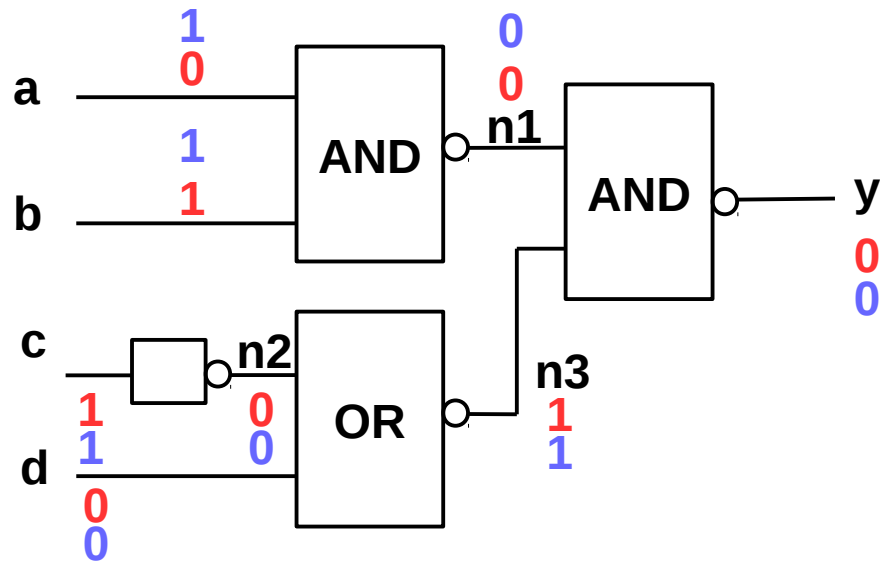
- Complexity is high since it needs to compute the faulty function for each fault

ATPG – from circuit structure



- **Fault Activation:** For observing the fault α we need to set the line x to 1.
- **Line Justification:** We set the input of previous gate(s) in a way to satisfy fault activation.
- **Fault propagation:** Propagate the fault on line x to the output f of the circuit.

ATPG – Example



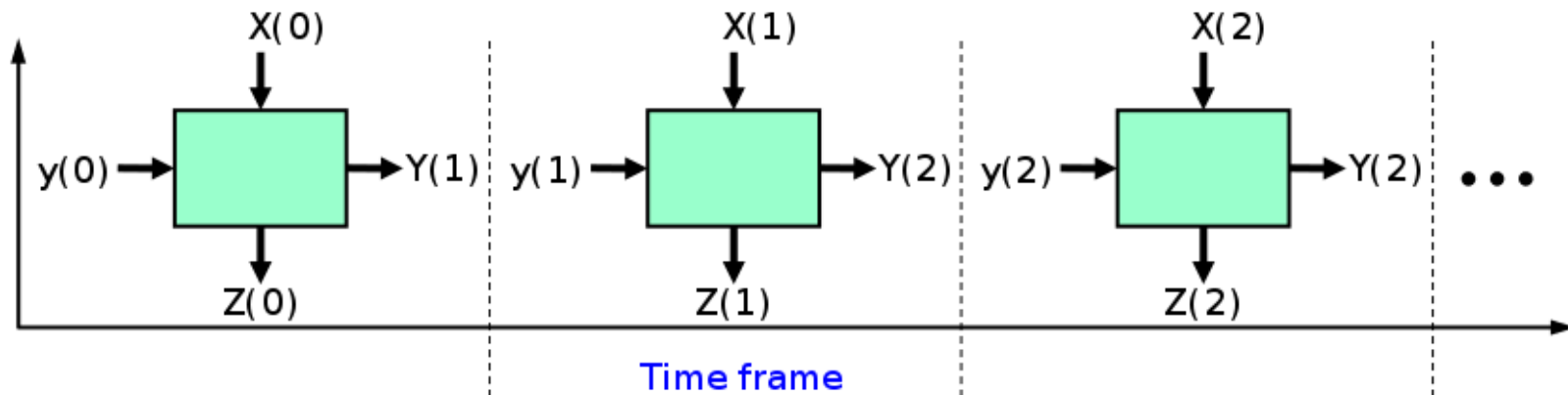
$t=\{a,b,c,d\}$

	stuck at 1	stuck at 0
a	$\{0,1,1,0\}$	$\{1,1,1,0\}$
b	$\{1,0,1,0\}$	$\{1,1,1,0\}$
c	$\{0,1,0,0\}$	$\{0,1,1,0\}$
d	$\{0,1,1,0\}$	$\{0,1,1,1\}$
n1	$\{1,1,1,0\}$	$\{0,1,1,0\}$
n2	$\{0,1,1,0\}$	$\{0,1,0,0\}$
n3	$\{0,1,0,1\}$	$\{0,1,1,0\}$
y	$\{0,1,1,0\}$	$\{1,1,1,0\}$

$$T=\min(t_1,t_2,t_3 \dots) = \{0,1,1,0\}, \{1,1,1,0\}, \{1,0,1,0\}, \{0,1,0,0\}, \{0,1,1,1\}, \{0,1,0,1\}$$

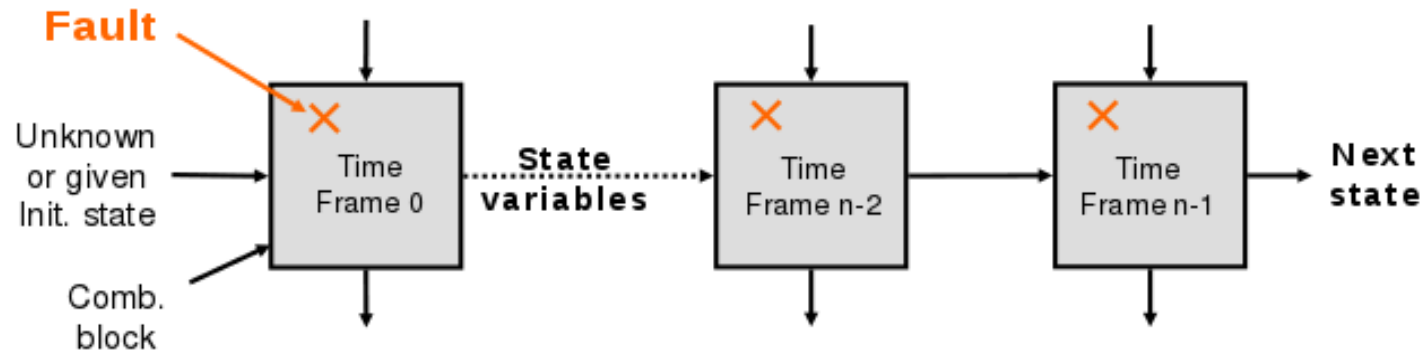
ATPG for sequential circuits

- Combinatorial test generations are all reduced to a SAT problem (see next lecture)
- Sequential circuits are computed by sequential operation of the combinatorial logic of the corresponding FSMs.
- Concept of time frame expansion



ATPG time frame expansion

- For FSM test pattern generation
- When a single stuck at fault is present a real sequential circuit, it will appear as a multiple fault in each unfolded timeframe.
- It may happen that conflicting logical states appear.
- D-Algorithm for each timeframe.



ATPG – D-Algorithm

- Structural method which performs a boolean search using the circuit topology [Roth'66]
- Use a 5-valued logic representing states in the good and faulty circuit. Implies that only one circuit is necessary.
 - $0 \rightarrow 0$ in true circuit, 0 in faulty circuit
 - $1 \rightarrow 1$ in true circuit, 1 in faulty circuit
 - $D \rightarrow 1$ in true circuit, 0 in faulty circuit
 - $D' \rightarrow 0$ in true circuit, 1 in faulty circuit
 - $X \rightarrow$ unknown value in either true or faulty circuit
- Goal: Find an assignment, an input sequence that causes a D or D' at one of the output(s)



ATPG – D-Algorithm

■ Termination rules of the Algorithm:

- If an output is D or D', the stuck at error is propagated to an output
→ an input sequence is found for the fault
- If $k > 4^n$, where n is the number of FF of the circuit, then the circuit is redundant.

A and B

A\B	0	1	X	D	D'
0	0	0	0	0	0
1	0	1	X	D	D'
X	0	X	X	X	X
D	0	D	X	D	0
D'	0	D'	X	0	D'

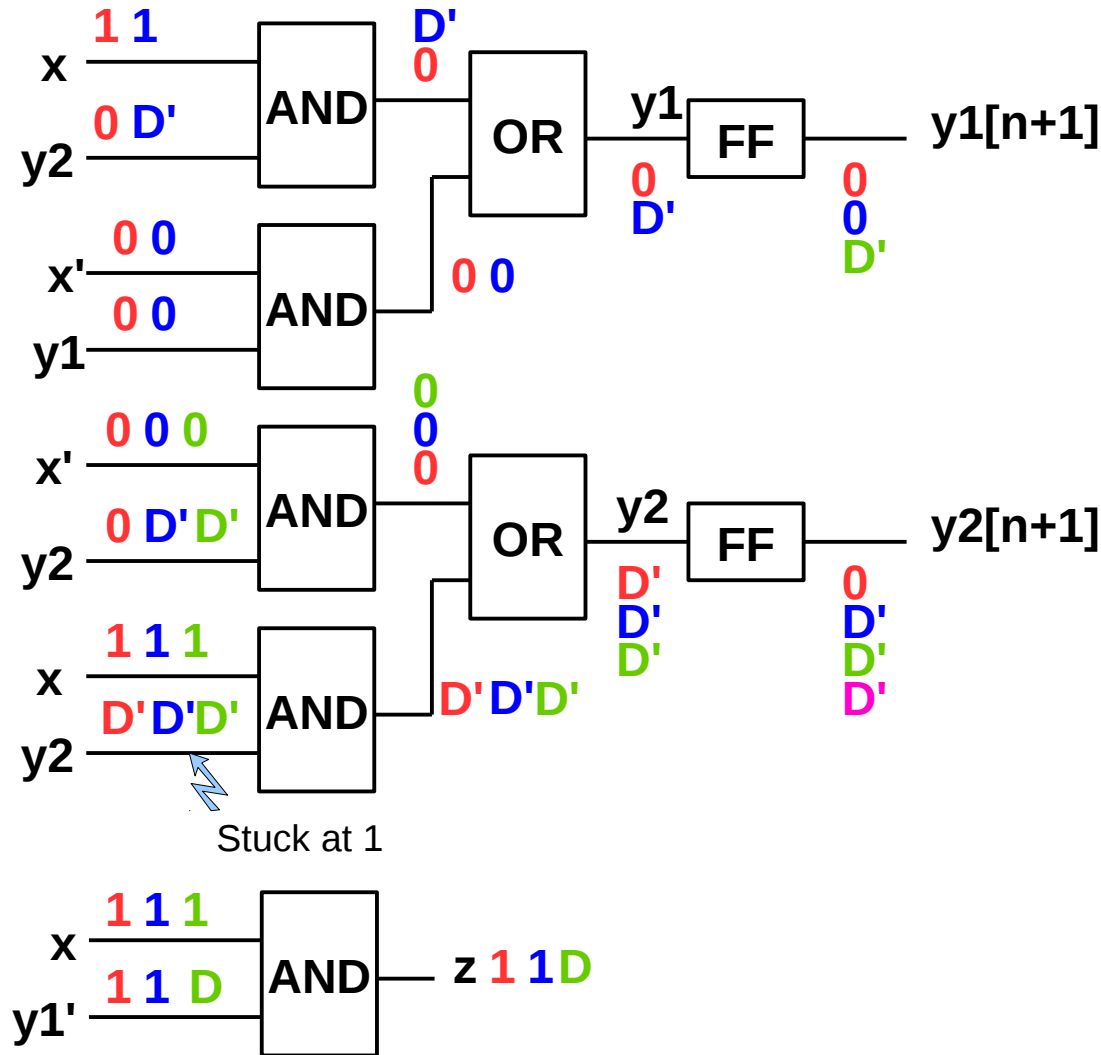
A or B

A\B	0	1	X	D	D'
0	0	1	X	D	D'
1	1	1	1	1	1
X	X	1	X	X	X
D	D	1	X	D	1
D'	D'	1	X	1	D'

not A

A	A'
0	1
1	0
X	X
D	D'
D'	D

ATPG – Example



Initial state of the FF
 $y_1 = 0, y_2 = 0$

t0
t1
t2
t3

Resulting test sequence
 $T = [1, 1, 1]$

ATPG – Example

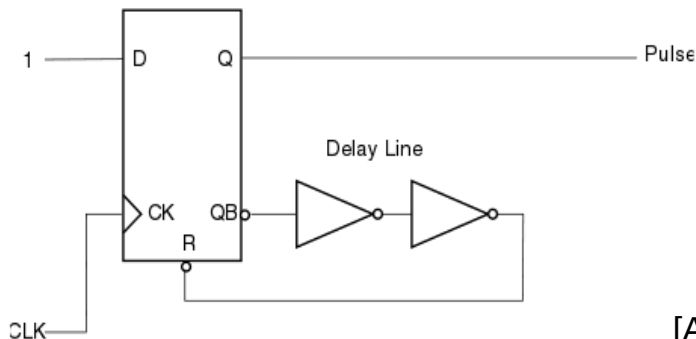
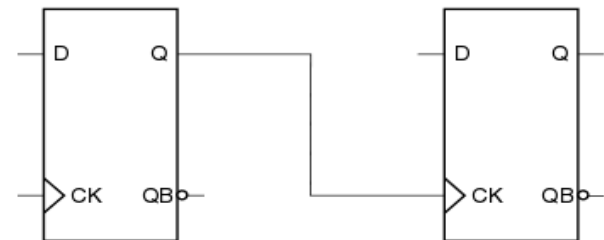
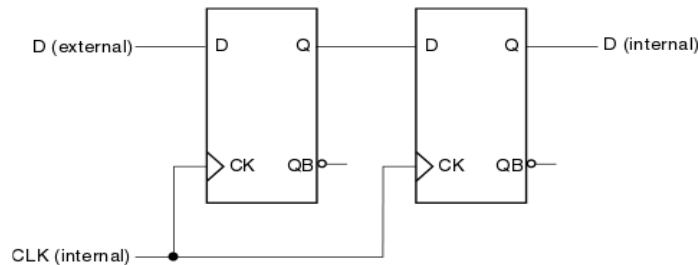
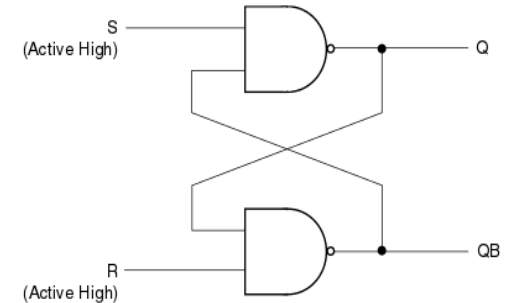
- T0:
 - Initial state $y_1=y_2=0$; $y_2=0$ and $y_2=1 \rightarrow D'$
 - Set $x(0)=1 \rightarrow y_2(1)=D'$ and $z(0)=1$
 - $z \neq D$ or $D' \rightarrow$ continue
- T1:
 - Set $x(1)=1 \rightarrow y_1(2)=y_2(2)=D'$ and $z(1)=1$
 - $z \neq D$ or $D' \rightarrow$ continue
- T2:
 - $X(2)=1$; $y_1(2)=D' \rightarrow z(2)=D$
- Resulting test sequence is $\{1, 1, 1\}$. Under this test sequence the stuck at fault can be observed at the circuit output after 3 clock cycles.



ASIC design for Testability

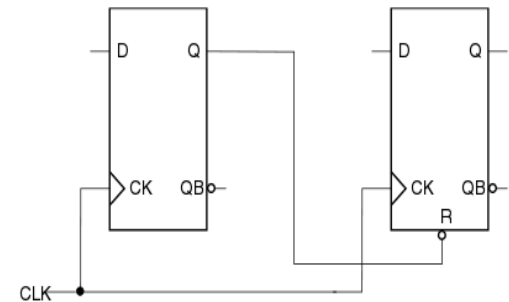
■ Synchronous design

- Use D FF, avoid latches
- Capture (synchronize) asynchronous inputs
- A FF must not drive reset
- Do not generate pulses based on delay line behavior



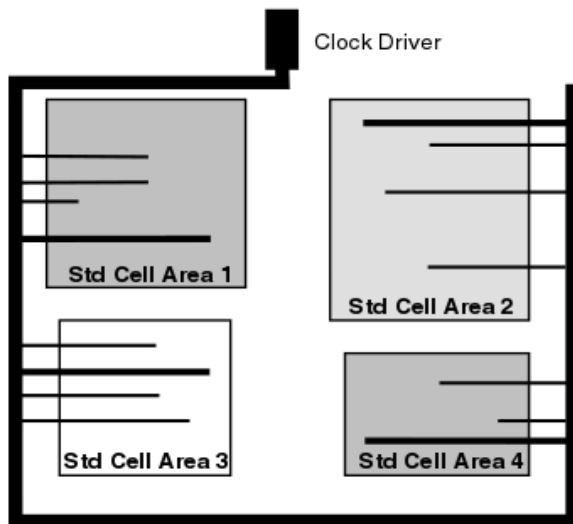
NO !

[ATMEL ASIC design guidelines]



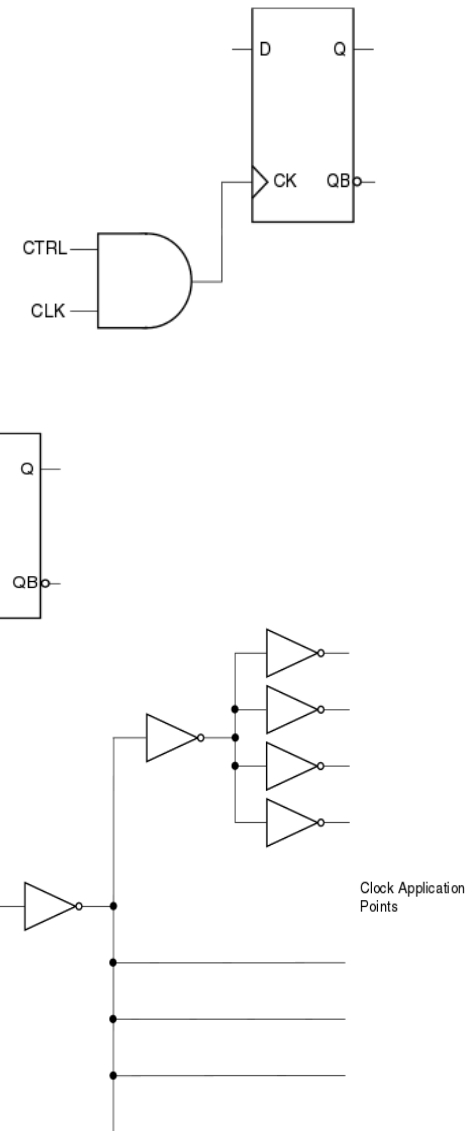
ASIC design for Testability

- Clock signals
 - Avoid unequal depth of clock buffering
 - Take care on a good clock tree signal distribution
 - No gated clocking
 - No double edge clocking



NO !

[ATMEL ASIC design guidelines]



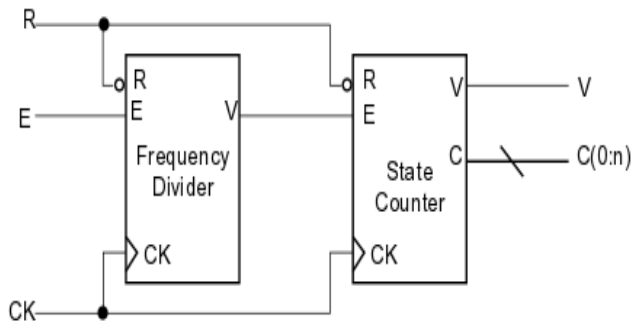
ASIC design for Testability

- Chain of counters:

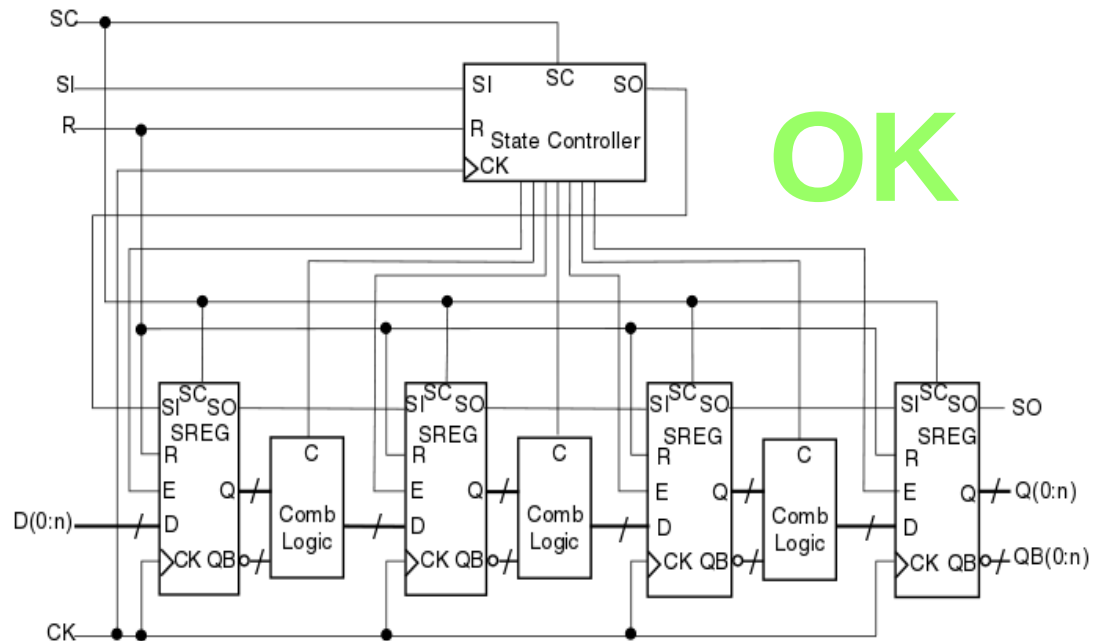
The first counter is not directly observable, and the second is not directly controllable

- Use BIST with compiled Megacells
- Form scan path, from scan input si to scan output so

NO !



[ATMEL ASIC design guidelines]

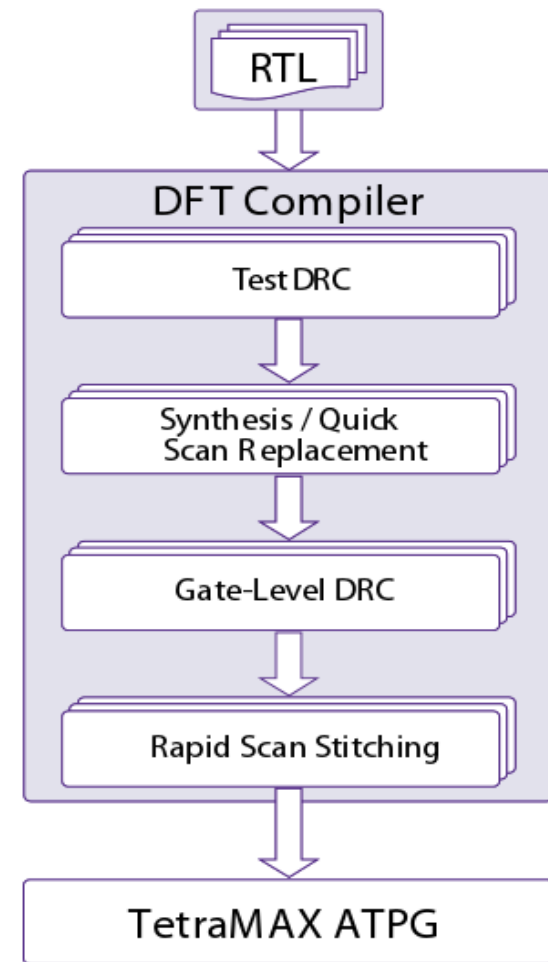


OK

Software Support

■ SYNOPSYS DFT Compiler

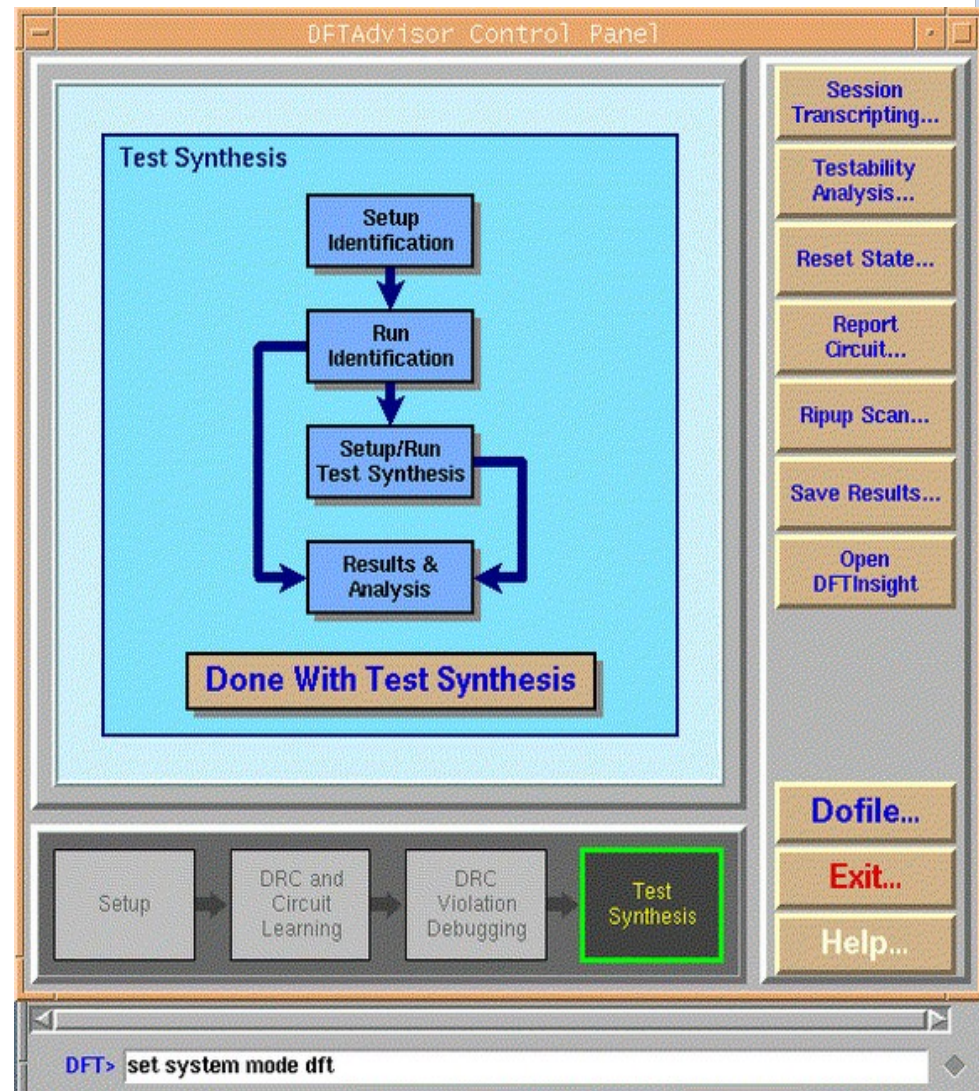
- Delivers scan DFT transparently within Synopsys' synthesis flows
- Conduct in-depth testability analysis at RTL level
- Enables rapid implementation of the most effective test structures at the hierarchical block level.
- TetraMAX ATPG:
 - PrimeTime interface selects critical timing path
 - Tester-ready patterns with complete timing
 - Easy-to-use flow with graphical support for analysis and debug



Software Support

■ MENTOR DFT ADVISOR

- Supports both full and partial-scan identification and insertion
- Reads most standard gate-level netlist
- Contains a powerful design rules checker
- After inserting the scan circuitry, you can then generate test vectors for your design by using Fastscan



Thank you for your attention

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