High Performance Computing Term 4 2018/2019

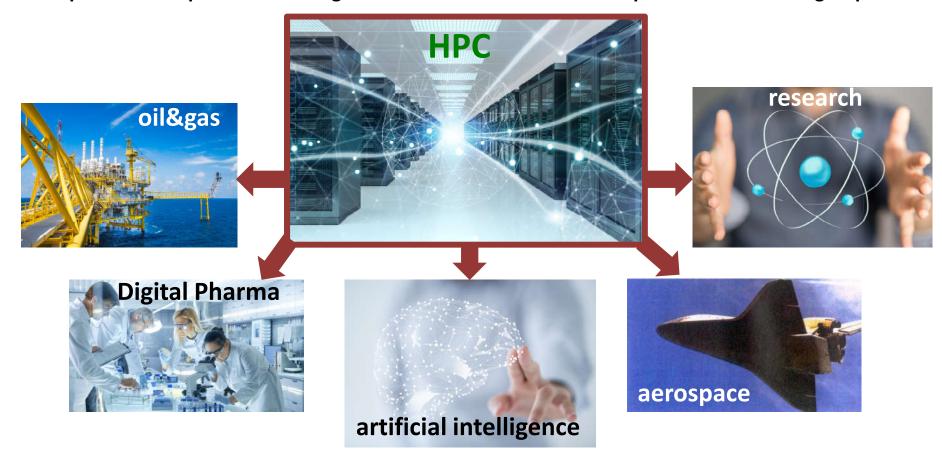
Lecture 2

Outline

- Applications of HPC and supercomputers
- Memory hierarchy
 - Matrix multiplication
- Building intuition: parallel processes
- Amdahl's law
- Flynn's taxonomy
- Shared memory and distributed memory systems
- Processes and threads
- Start with OpenMP

High Performance Computing is a strategic necessity

HPC provides competitive advantage for industrial and business companies and research groups.



High Performance Computing is a strategic necessity



¹ https://www.scientificamerican.com/article/cost-to-develop-new-pharmaceutical-drug-now-exceeds-2-

Laser plasma technology can revolutionize accelerators





- Large Hadron Collider: 27 km, candidate for Higgs boson. Limit of accelerating gradients.
- Relativistic Laser Plasma: 1 GeV in just 10 cm using powerful lasers interacting with plasma

More energy = increase in size. Cost? Size?

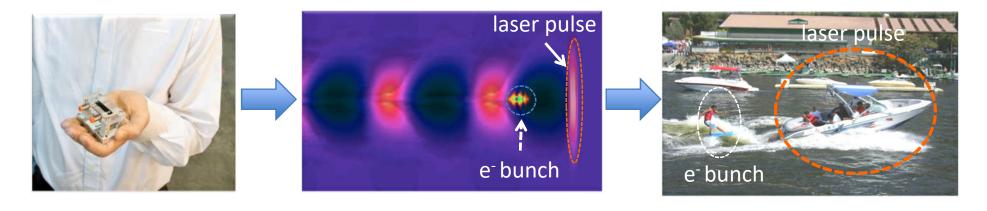




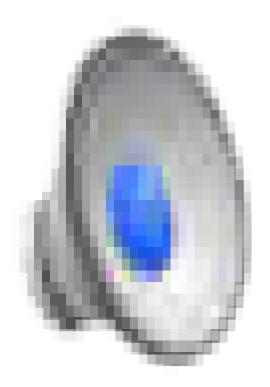


Laser plasma technology can revolutionize accelerators

Laser Plasma Acceleration = Surfing on a plasma wave



simulations by S. Kuschel



Most powerful industrial supercomputer is used for oil&gas exploration(Eni S.p.A.)

Supercomputer HPC4 installed at Eni is #15 in top500 (12 ПФлопс). For comparison, most powerful russian supercomputer: Lomonosov-2, 2.5 PFlops.

15	Eni S.p.A.	HPC4 - Proliant DL380 Gen10,	253,600	12,210.0	18,621.1	1,320
	Italy	Xeon Platinum 8160 24C				
		2.1GHz, Mellanox InfiniBand				
		EDR, NVIDIA Tesla P100				
		HPE				

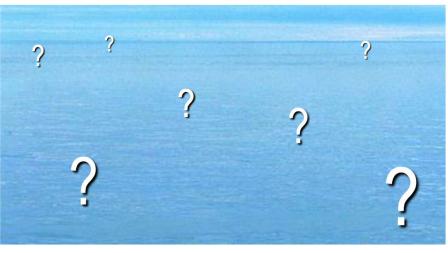


Much shorter oil&gas exploration by using supercomputers

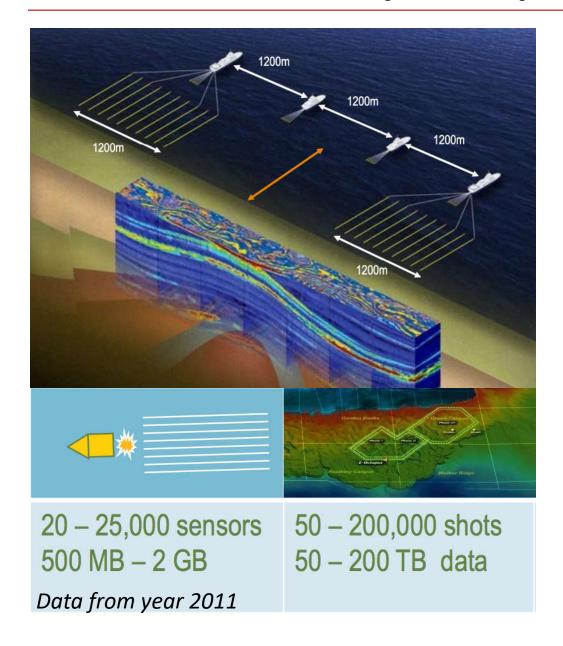
Seismic –Acoustic measurement Electromagnetic Gravity







Much shorter oil&gas exploration by using supercomputers



Creation of 3D seismic model:

~30000000 CPU core hours (~1000 years on a normal laptop)

Without the supercomputer:

- it is impossible to create a precise seismic model
- it is not clear where to drill for optimal oil output and the mistake can cost millions of dollars

New oil&gas deposits are more and more hard to reach.

Climate modelling

$$\frac{du}{dt} - \left(f + \frac{u}{a} \operatorname{tg} \varphi\right) v + \frac{1}{a \cos \varphi} \left(\frac{\partial \Phi}{\partial \lambda} + \frac{RT}{p_s} \frac{\partial p_s}{\partial \lambda}\right) = F_u,$$

$$\frac{dv}{dt} + \left(f + \frac{u}{a} \operatorname{tg} \varphi\right) u + \frac{1}{a} \left(\frac{\partial \Phi}{\partial \varphi} + \frac{RT}{p_s} \frac{\partial p_s}{\partial \varphi}\right) = F_v,$$

$$\frac{\partial \Phi}{\partial \sigma} = -\frac{RT}{\sigma},$$

$$\frac{\partial p_s}{\partial t} + \frac{1}{a \cos \varphi} \left(\frac{\partial p_s u}{\partial \lambda} + \frac{p_s v \cos \varphi}{\partial \varphi}\right) + \frac{\partial p_s \dot{\sigma}}{\partial \sigma} = 0,$$

$$\frac{dT}{dt} - \frac{RT}{c_p \sigma p_s} \left[p_s \dot{\sigma} + \sigma \left(\frac{\partial p_s}{\partial t} + \frac{u}{a \cos \varphi} \frac{\partial p_s}{\partial \lambda} + \frac{v}{a} \frac{\partial p_s}{\partial \varphi}\right)\right] = F_T + \epsilon,$$

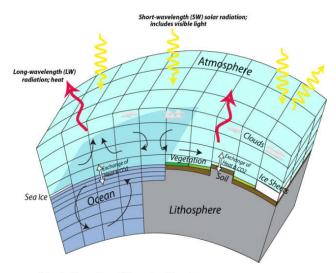
$$\frac{dq}{dt} = F_q - (C - E),$$

$$\frac{d}{dt} = \frac{\partial}{\partial t} + \frac{u}{a \cos \varphi} \frac{\partial}{\partial \lambda} + \frac{v}{a} \frac{\partial}{\partial \varphi} + \dot{\sigma} \frac{\partial}{\partial \sigma},$$

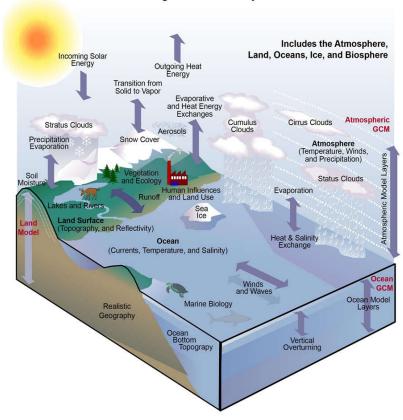
Weather modelling only:

- 100 years, 10 minutes interval ~ 5.3 * 10⁶ checkpoints
- 1° latitude and longitude grid of 40 layers ~ 2.6 * 10⁶ cells
- Each cell ~ 10 components to model
- Computational cost for one measurement ~ 1000 operations

Total: 10¹⁷ arithmetic operations, i. e. 3-30 hours on 1 TFlops

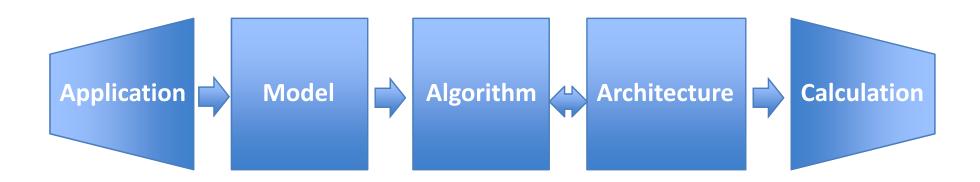


Modeling the Climate System



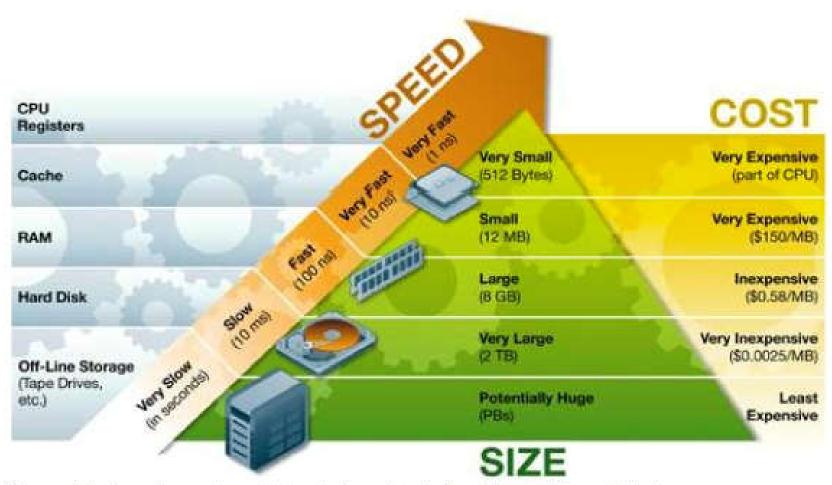
High Performance Computing is not only "computing": sometimes thinking outside the box and knowledge of other areas helps

Multi-disciplinary approach to HPC is necessary:

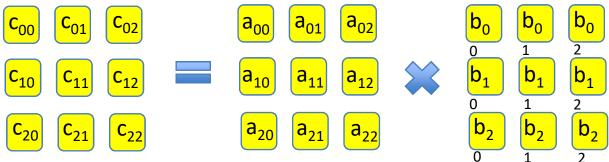


- Analytical solution to the problem can save computing time
- Clever model (reduced model) can save computing time
- Clever numerical methods fit to the certain architecture can save computing time

Extended Memory Hierarchy



Source: http://www.ts.avnet.com/uk/products and solutions/storage/hierarchy.html

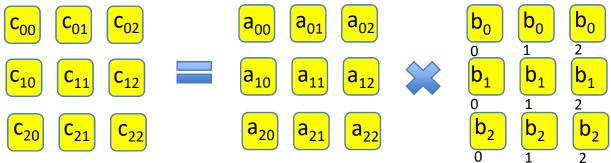


NxN matrices: A, B, C N >> 1

$$C = A * B$$

$$c_{ij} = \sum_{n=0}^{n=N-1} a_{in} b_{nj}$$

- 1. Write matmul in C
- 2. Exchange the order of for loops (i,j,n)->(n,i,j)
- 3. Measure execution time and if it changes try to explain why.
- 4. Try different compiler optimization options
- 5. Calculate the flops of a single CPU core of your machine
- 6. Imagine you have K workers how would you make this code parallel?
- 7. What if you use double * instead of double **?
- 8. Implement Hilbert curve for 2D matrix allocation.



NxN matrices: A, B, C



 $c_{ij} = \sum_{n=0}^{n=N-1} a_{in} b_{nj}$

for i from 0 to N-1:

for j from 0 to N-1:

for n from 0 to N-1:

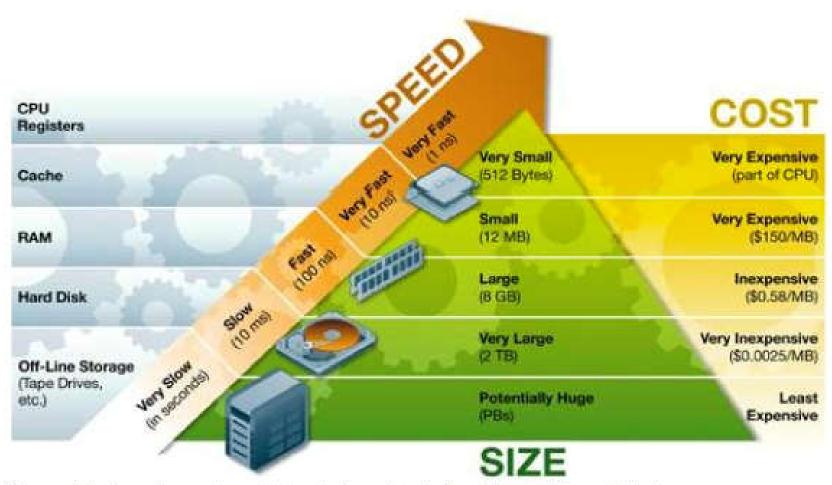
C[i][j]+=A[i][n]*B[n][j]

goo.gl/vQeUD8

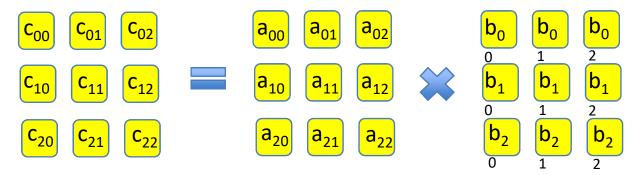
- 1. Write matmul in C
- 2. Exchange the order of for loops (i,j,n)->(n,i,j)
- 3. Measure execution time and if it changes try to explain why.
- 4. Try different compiler optimization options
- 5. Calculate the flops of a single CPU core of your machine
- 6. Imagine you have K workers how would you make this code parallel?
- 7. What if you use double * instead of double **?
- 8. Implement Hilbert curve for 2D matrix allocation.

double ** ptr double ** ptr N arrays of M double each ptr[0][0] ptr[0][M-1] ptr[0][1] ptr[0] ptr[1][0] ptr[1][1] ptr[1][M-1] ptr[1] ptr[2] ptr[2][0] ptr[2][1] ptr[2][M-1] ptr[N-1] ptr[N-1][0] ptr[N-1][1] ptr[N-1][M-1] Array of N double *

Extended Memory Hierarchy



Source: http://www.ts.avnet.com/uk/products and solutions/storage/hierarchy.html

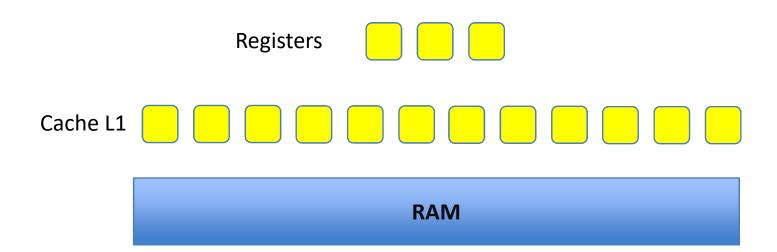


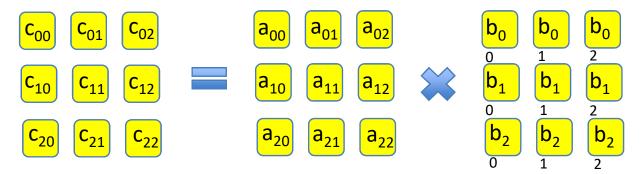
for i from 0 to N-1:

for j from 0 to N-1:

for n from 0 to N-1:

C[i][j]+=A[i][n]*B[n][j]





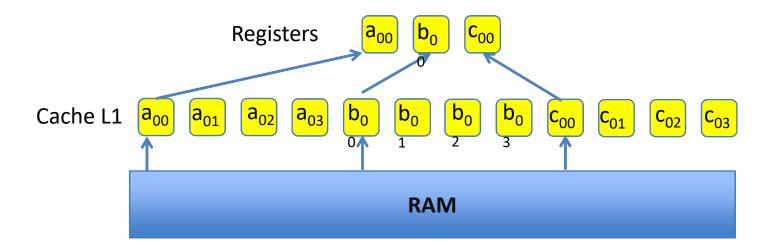
for i from 0 to N-1:

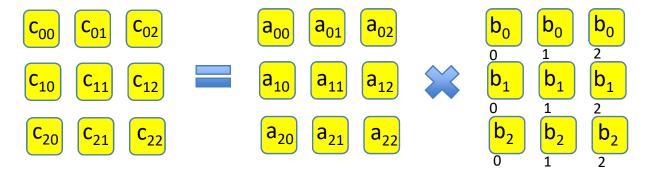
for j from 0 to N-1:

for n from 0 to N-1:

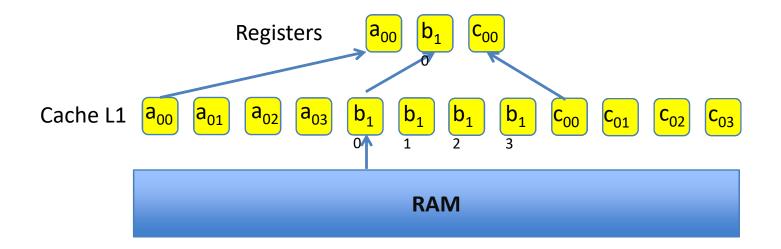
C[i][j]+=A[i][n]*B[n][j]

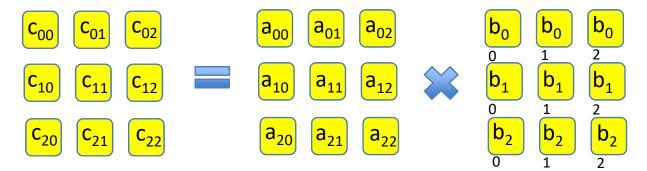
i=0 j=0 n=0 C[0][0]+=A[0][0]*B[0][0]





```
\begin{array}{lll} \text{for i from 0 to N-1:} & & \text{i=0} \\ & \text{for j from 0 to N-1:} & \text{j=0} \\ & \text{for n from 0 to N-1:} & \text{n=1} \\ & \text{C[i][j]+=A[i][n]*B[n][j]} & \text{C[0][0]+=A[0][1]*B[1][0]} \end{array}
```





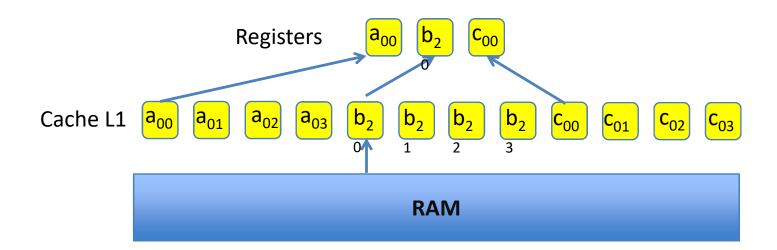
for i from 0 to N-1:

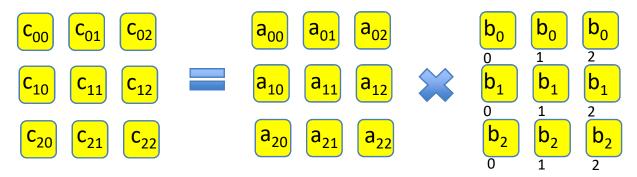
for j from 0 to N-1:

for n from 0 to N-1:

C[i][j]+=A[i][n]*B[n][j]

i=0 j=0 n=2 C[0][0]+=A[0][2]*B[2][0]



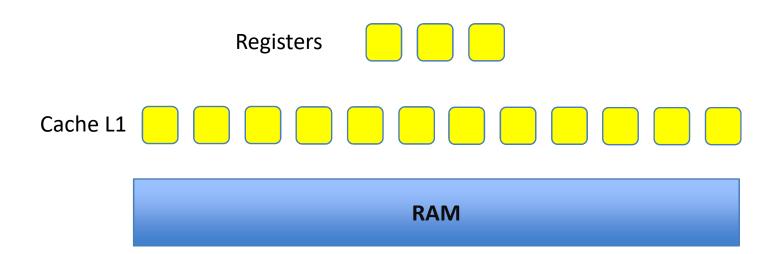


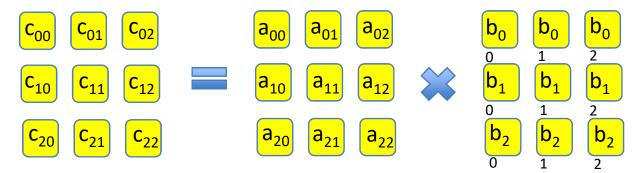
for n from 0 to N-1:

for i from 0 to N-1:

for j from 0 to N-1:

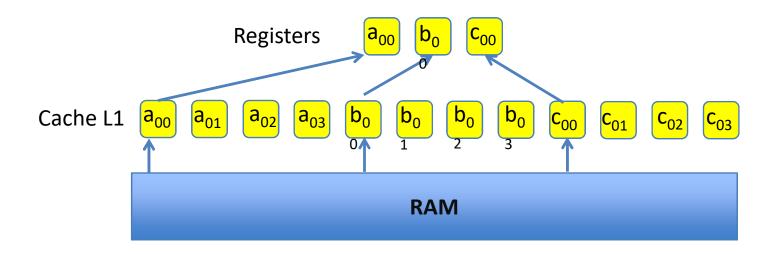
C[i][j]+=A[i][n]*B[n][j]

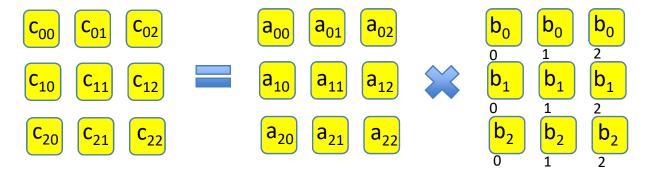




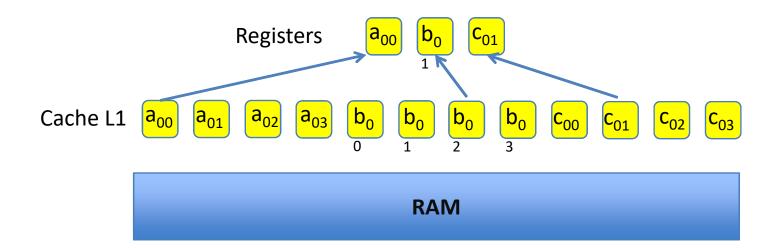
```
for n from 0 to N-1:
    for i from 0 to N-1:
        for j from 0 to N-1:
        C[i][j]+=A[i][n]*B[n][j]
```

n=0 i=0 j=0 C[0][0]+=A[0][0]*B[0][0]





```
\begin{array}{lll} \text{for n from 0 to N-1:} & & n=0 \\ & \text{for i from 0 to N-1:} & & i=0 \\ & \text{for j from 0 to N-1:} & & j=1 \\ & & C[i][j]+=A[i][n]*B[n][j] & & C[0][1]+=A[0][0]*B[0][1] \end{array}
```



Parallel intuition

• Examples of parallelism in real world?

Parallel intuition How to really understand concurrency.



Building a house



Say, we have 4 workers building a house

Construction will be faster than if it would be done by 1 worker

After each floor, workers have to synchronize

Connect their parts - communication

Can't build 7th floor before 3rd

Amdahl's law

Speed-up of a parallel program/work:

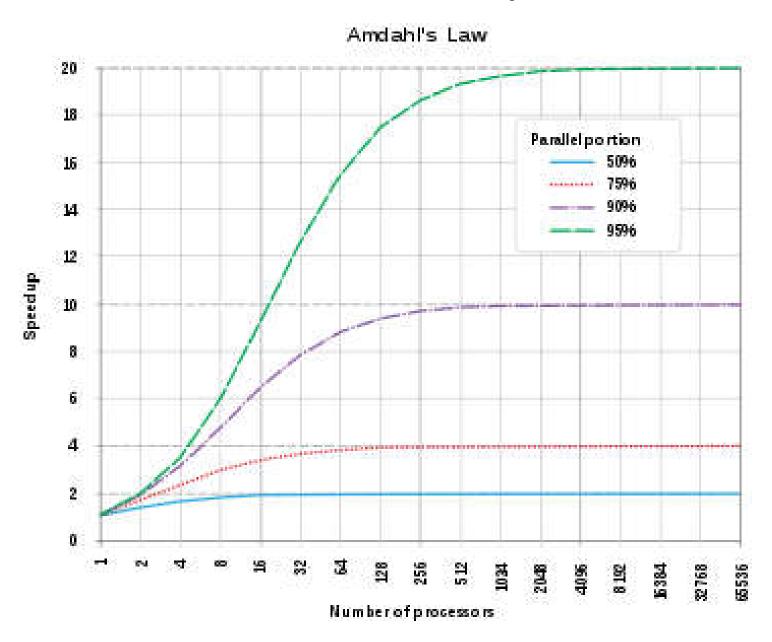


$$S = \frac{1}{(1-f) + \frac{f}{n}}$$

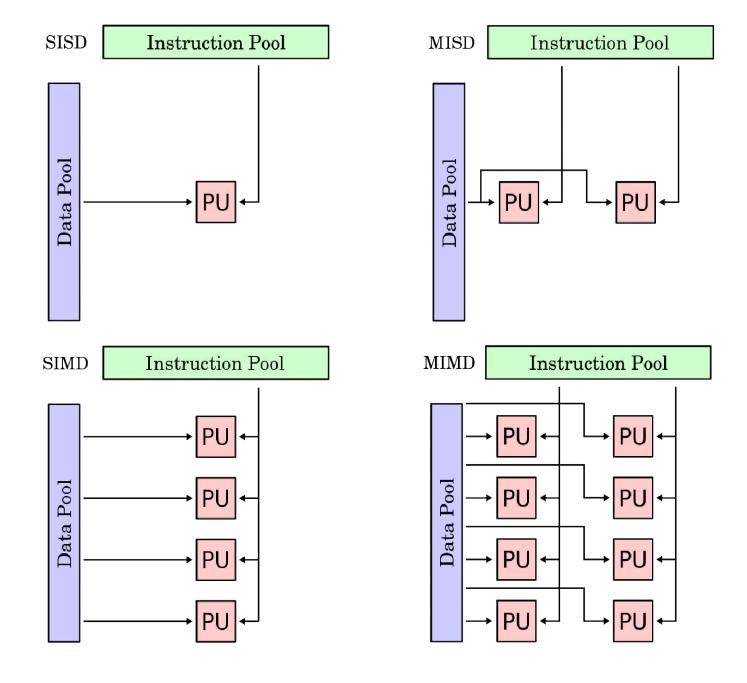
f – is a fraction of the work that can be done in parallel

n – number of parallel workers

Amdahl's law implication



Flynn's taxonomy



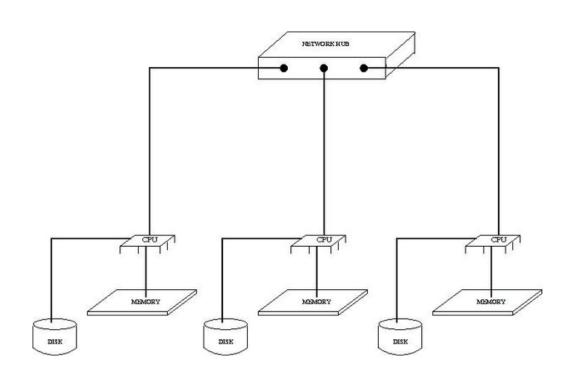
Shared memory and distributed memory

shared memory

System Bus or Crossbar Switch

Memory

distributed memory



Ways to work parallel

- Scientific software (Matlab, Mathematica, ANSYS, OpenFoam, Tensorflo w, PyTorch, etc)
- 2. Libraries (Math Kernel Library, CuDNN, CuBlas, Thrust, Boost)
- 3. Parallel instruments (OpenMP, MPI, Cuda, OpenCL)

Shared memory systems: processes and threads



Shared memory systems: processes and threads

