The background is a light gray color. It features several stylized, light gray leaf shapes scattered across the right side. Some leaves are simple outlines, while others are filled with a light gray gradient. There are also small, solid light gray circles scattered around the leaves. The text is centered on the left side of the image.

# Assignment 1

## DATA STRUCTURES & ALGORITHMS

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Create a design specification for data structures, explaining the valid operations that can be carried out on the structures.

# Identify the Data Structures



## 1. Array

Definition: An array is a collection of elements of the same type, stored consecutively in memory. Each element can be accessed through an index.

Characteristics:

Fixed size (in many programming languages).

Access elements by index in  $O(1)$  time.

## 2. Stack

Definition: Stack is a data structure that follows the LIFO (Last In, First Out) principle, meaning the last element added will be taken out first.

Characteristics:

Operate only on the top element of the stack.

Push (add) and Pop (take out) operations take place in  $O(1)$  time.

## 3. Tree

Definition: A tree is a hierarchical data structure, each node has a value and can have many children. A Binary Tree is a common type of tree, in which each node has a maximum of two children.

Characteristics:

There are many ways to traverse a tree: Preorder, Inorder, Postorder.

The insertion and search operations have a complexity ranging from  $O(\log n)$  to  $O(n)$  depending on whether the tree is balanced or not.

## 4. Hash Table

Definition: A hash table uses a hash function to map a key to a specific location in the table. This allows for fast searches, insertions, and deletions.

Characteristics:

The average search and insertion time is  $O(1)$  but can be as bad as  $O(n)$  in the case of key conflicts.



# Define the Operations

## 2.1 Stack Operations

Push: Add an element to the top of the stack.

Pop: Remove and return the top element from the stack.

Peek/Top: Return the top element without removing it.

isEmpty: Check if the stack is empty.

size: Return the number of elements in the stack.

## 2.2 Array Operations

Access: Retrieve an element at a specified index.

Insert: Add an element at a specified index, shifting subsequent elements.

Delete: Remove an element at a specified index, shifting subsequent elements left.

Search: Find the index of a specified element.

Update: Change the value of an element at a specified index.

## 2.3 Tree Operations

Insert: Add a new node to the tree based on its value.

Delete: Remove a node from the tree.

Search: Find a node with a specified value in the tree.

Traversal: Visit all nodes in a specific order (pre-order, in-order, post-order).

Update: Change the value of a specified node.

## 2.4 Hash Table Operations

Insert: Add a key-value pair to the hash table.

Delete: Remove a key-value pair based on the key.

Search: Retrieve the value associated with a specified key.

Update: Change the value of an existing key.

isEmpty: Check if the hash table is empty.

# Define Pre- and Post-conditions

## 2.1 Stack Pre- and Post-conditions

Push(value)

- Pre-condition: The stack is not full (for bounded stacks).
- Post-condition: The stack size increases by 1, and the top element becomes the pushed value.

Pop()

- Pre-condition: The stack is not empty.
- Post-condition: The stack size decreases by 1, and the returned value is the previous top element.

Peek()

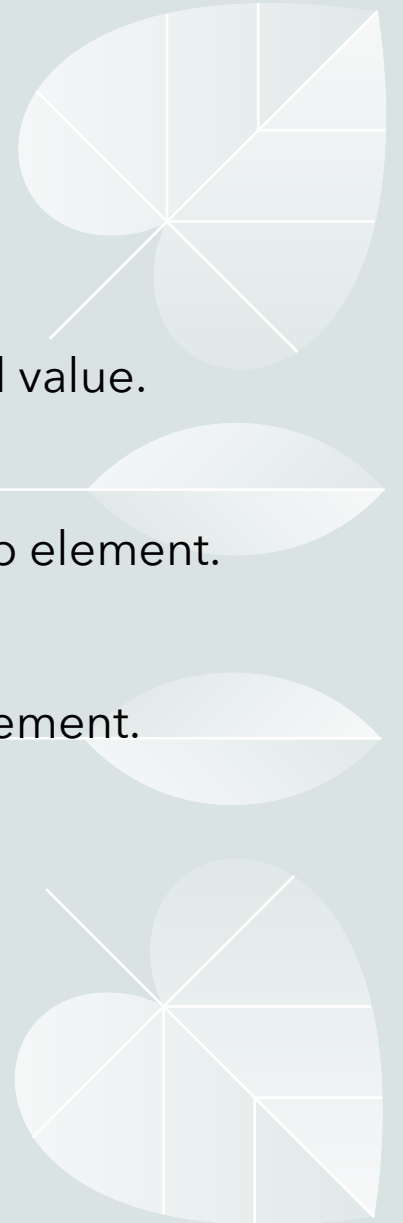
- Pre-condition: The stack is not empty.
- Post-condition: The top element remains the same, and the returned value is the top element.

isEmpty()

- Pre-condition: None.
- Post-condition: Returns true if the stack size is 0, otherwise false.

size()

- Pre-condition: None.
- Post-condition: Returns the current number of elements in the stack.



## 2.2 Array Pre- and Post-conditions

### Access(index)

- Pre-condition: The index is within bounds ( $0 \leq \text{index} < \text{array length}$ ).
- Post-condition: Returns the element at the specified index.

### Insert(index, value)

- Pre-condition: The index is within bounds for insertion ( $0 \leq \text{index} \leq \text{array length}$ ).
- Post-condition: The value is added at the specified index, and subsequent elements are shifted to the right.

### Delete(index)

- Pre-condition: The index is within bounds ( $0 \leq \text{index} < \text{array length}$ ).
- Post-condition: The element at the specified index is removed, and subsequent elements are shifted to the left.

### Search(value)

- Pre-condition: None.
- Post-condition: Returns the index of the value if found; otherwise, returns -1.

### Update(index, value)

- Pre-condition: The index is within bounds ( $0 \leq \text{index} < \text{array length}$ ).
- Post-condition: The value at the specified index is updated to the new value.

## 2.3 Tree Pre- and Post-conditions

Insert(value)

- Pre-condition: None.
- Post-condition: A new node with the specified value is added to the tree.

Delete(value)

- Pre-condition: The value exists in the tree.
- Post-condition: The node with the specified value is removed from the tree.

Search(value)

- Pre-condition: None.
- Post-condition: Returns the node with the specified value if found; otherwise, returns null.

Traversal(order)

- Pre-condition: None.
- Post-condition: All nodes are visited in the specified order.

Update(oldValue, newValue)

- Pre-condition: The oldValue exists in the tree.
- Post-condition: The value of the specified node is updated to newValue.

## 2.4 Hash Table Pre- and Post-conditions

Insert(key, value)

- Pre-condition: The key does not already exist in the hash table (for unique keys).
- Post-condition: A new key-value pair is added to the hash table.

Delete(key)

- Pre-condition: The key exists in the hash table.
- Post-condition: The key-value pair associated with the specified key is removed.

Search(key)

- Pre-condition: None.
- Post-condition: Returns the value associated with the specified key if found; otherwise, returns null.

Update(key, value)

- Pre-condition: The key exists in the hash table.
- Post-condition: The value associated with the specified key is updated to the new value.

isEmpty()

- Pre-condition: None.
- Post-condition: Returns true if the hash table contains no key-value pairs; otherwise, returns false.



# Discuss Time and Space Complexity

## 2.1 Stack Complexity

### •Push(value)

- **Time Complexity:**  $O(1)$  - Adding an element to the top is done in constant time.
- **Space Complexity:**  $O(n)$  - The space complexity is dependent on the number of elements stored in the stack.

### •Pop()

- **Time Complexity:**  $O(1)$  - Removing the top element is done in constant time.
- **Space Complexity:**  $O(n)$  - The space complexity remains dependent on the number of elements.

### •Peek()

- **Time Complexity:**  $O(1)$  - Accessing the top element is done in constant time.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •isEmpty()

- **Time Complexity:**  $O(1)$  - Checking if the stack is empty is done in constant time.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •size()

- **Time Complexity:**  $O(1)$  - Returning the size is done in constant time.
- **Space Complexity:**  $O(1)$  - No additional space is used.

## 2.2 Array Complexity

### •Access(index)

- **Time Complexity:**  $O(1)$  - Accessing an element by index is done in constant time.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •Insert(index, value)

- **Time Complexity:**  $O(n)$  - Inserting requires shifting elements, leading to linear time complexity.
- **Space Complexity:**  $O(n)$  - The space complexity is dependent on the number of elements in the array.

### •Delete(index)

- **Time Complexity:**  $O(n)$  - Deleting requires shifting elements, leading to linear time complexity.
- **Space Complexity:**  $O(n)$  - The space complexity is dependent on the number of elements in the array.

### •Search(value)

- **Time Complexity:**  $O(n)$  - In the worst case, all elements need to be checked.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •Update(index, value)

- **Time Complexity:**  $O(1)$  - Updating an element by index is done in constant time.
- **Space Complexity:**  $O(1)$  - No additional space is used.



## 2.3 Tree Complexity

### •Insert(value)

- **Time Complexity:**  $O(h)$  - Where  $h$  is the height of the tree; in a balanced tree, this is  $O(\log n)$ , while in an unbalanced tree, it can be  $O(n)$ .
- **Space Complexity:**  $O(n)$  - The space complexity depends on the number of nodes.

### •Delete(value)

- **Time Complexity:**  $O(h)$  - Similar to insert; it depends on the height of the tree.
- **Space Complexity:**  $O(n)$  - The space complexity is dependent on the number of nodes.

### •Search(value)

- **Time Complexity:**  $O(h)$  - Where  $h$  is the height of the tree.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •Traversal(order)

- **Time Complexity:**  $O(n)$  - All nodes are visited.
- **Space Complexity:**  $O(h)$  - Space used in recursion stack ( $h$  is the height).

### •Update(oldValue, newValue)

- **Time Complexity:**  $O(h)$  - Depends on the height of the tree.
- **Space Complexity:**  $O(1)$  - No additional space is used.

## 2.4 Hash Table Complexity

### •Insert(key, value)

- **Time Complexity:**  $O(1)$  on average (amortized),  $O(n)$  in the worst case (due to collisions).
- **Space Complexity:**  $O(n)$  - The space complexity is dependent on the number of key-value pairs.

### •Delete(key)

- **Time Complexity:**  $O(1)$  on average,  $O(n)$  in the worst case.
- **Space Complexity:**  $O(n)$  - The space complexity is dependent on the number of key-value pairs.

### •Search(key)

- **Time Complexity:**  $O(1)$  on average,  $O(n)$  in the worst case.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •Update(key, value)

- **Time Complexity:**  $O(1)$  on average,  $O(n)$  in the worst case.
- **Space Complexity:**  $O(1)$  - No additional space is used.

### •isEmpty()

- **Time Complexity:**  $O(1)$  - Checking if the hash table is empty is done in constant time.
- **Space Complexity:**  $O(1)$  - No additional space is used.