Step: $\vdash \varphi \to \bigvee_{\varphi_l \Rightarrow^{\exists} \varphi_r \in \mathcal{S}} \exists Free Vars(\varphi_l).\varphi_l$ $\vdash ((\varphi \land \varphi_l) \neq \bot_{Cfg}) \land \varphi_r \to \varphi' \quad \text{for each } \varphi_l \Rightarrow^{\exists} \varphi_r \in \mathcal{S}$ $S. \mathcal{A} \vdash_C \varphi \Rightarrow^{\forall} \varphi'$

 $\frac{\varphi \Rightarrow^{Q} \varphi' \in S \cup \mathcal{A} \qquad \psi \text{ is FOL formula (logical frame)}}{S \mathcal{A} + V + V}$

$$\mathcal{S}, \mathcal{A} \vdash_{\mathcal{C}} \varphi \land \psi \Rightarrow^{\mathcal{Q}} \varphi' \land \psi$$

Reflexivity: $\frac{\cdot}{S, \mathcal{A} \vdash \varphi \Rightarrow^{Q} \varphi}$

TRANSITIVITY:

$$\frac{S, \mathcal{A} \vdash_{C} \varphi_{1} \Rightarrow^{Q} \varphi_{2}}{S, \mathcal{A} \vdash_{C} \varphi_{1} \Rightarrow^{Q} \varphi_{3}} \xrightarrow{S, \mathcal{A} \vdash_{C} \varphi_{1} \Rightarrow^{Q} \varphi_{3}}$$

Consequence:

$$\frac{\models \varphi_1 \to \varphi_1' \qquad \mathcal{S}, \mathcal{A} \vdash_{\mathcal{C}} \varphi_1' \Rightarrow^{\mathcal{Q}} \varphi_2' \qquad \models \varphi_2' \to \varphi_2}{\mathcal{S}, \mathcal{A} \vdash_{\mathcal{C}} \varphi_1 \Rightarrow^{\mathcal{Q}} \varphi_2}$$

CASE ANALYSIS:

$$\frac{S, \mathcal{A} \vdash_{C} \varphi_{1} \Rightarrow^{Q} \varphi \qquad S, \mathcal{A} \vdash_{C} \varphi_{2} \Rightarrow^{Q} \varphi}{S, \mathcal{A} \vdash_{C} \varphi_{1} \lor \varphi_{2} \Rightarrow^{Q} \varphi}$$

Abstraction:

$$\frac{S, \mathcal{A} \vdash_{C} \varphi \Rightarrow^{Q} \varphi' \qquad X \cap FreeVars(\varphi') = \emptyset}{S, \mathcal{A} \vdash_{C} \exists \mathsf{X} \varphi \Rightarrow^{Q} \varphi'}$$

CIRCULARITY:

$$\frac{\mathcal{S},\mathcal{A} \vdash_{C \cup \{\varphi \Rightarrow \mathcal{Q}_{\varphi'}\}} \varphi \Rightarrow^{\mathcal{Q}} \varphi'}{\mathcal{S},\mathcal{A} \vdash_{C} \varphi \Rightarrow^{\mathcal{Q}} \varphi'}$$