## CMPT 360 Assignment Two due: in Canvas at 11:59PM on November 13, 2023.

This assignment offers opportunities to construct, verify, find counterexamples, and implement dynamic-programming.

Submit your solutions to the following problems. Upload them as two files:

- 1. a PDF of your written answers
- 2. a tar archive (named abc123-a3.tar) of your directory (named abc123-a3) containing your solution to first coding problem. Please consult assignment 0 for detailed instructions on the content, format, naming, and development techniques. Recall that the absence of a git commit log showing incremental development yields a zero grade even if the code works correctly.

## 1 Problems

1. (35 pts) Let G=(V,E) be an undirected graph with |V|=n. A set  $I\subseteq V$  of the vertices is *independent* if no pair of vertices  $(u,v)\in I$  has an edge in G:

$$\forall u, v \in I \subseteq V \qquad (u, v) \notin E$$

Finding independent sets is difficult in general, but it can be done efficiently if the graph is simple enough.

The graph G=(V,E) is a path if it's nodes can be arranged in a sequence

$$V = \{v_1, v_2, v_3, \dots, v_n\}$$

where each vertex appears in the sequence exactly once, and the edges E of G are only those between consecutive vertices in the sequence:

$$\forall (u, v) \in E$$
  $\exists i \in [1, ..., n]$  with  $u = v_i and v = v_{i+1}$ 

Each vertex will have an associated weight  $w \in \mathbb{N} = \{0, 1, \ldots\}$ . An example of a path is:

```
\begin{array}{c} \mathrm{path}\; G = (V,E)\; \mathrm{with} \\ \\ V = \left\{ \begin{array}{c} v_1,\, v_2,\, v_3,\, v_4,\, v_5 \end{array} \right\} \\ E = \left\{ \begin{array}{c} (v_1,\, v_2),\, (v_2,\, v_3),\, (v_3,\, v_4),\, (v_4,\, v_5) \end{array} \right\} \\ \\ \mathrm{and\; weights}\; \left\{ \begin{array}{c} v_1 \mapsto 1, \;\; v_2 \mapsto 8,\, v_3 \mapsto 6,\, v_4 \mapsto 3,\, v_5 \mapsto 1 \end{array} \right\} \\ \hline \\ 1 \\ \hline \end{array}
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The goal is to find an independent set in a path with maximum total weight.

• (5 pts) Give an example path where the following greedy algorithm does not find the maximum weight independent set:

```
\label{eq:continuity} $ \begin{tabular}{ll} $//$ given $ G = (V,E) $ \\ S &= \varnothing $ \\ \hline \end{tabular} $ \begin{tabular}{ll} $with maximum $ weight $ \\ S &= S \cup \{v_i\} $ \\ V &= V \setminus \{v_{i-1}, v_i, v_{i+1}\} $ \\ \hline \end{tabular} $ \begin{tabular}{ll} $cone $ \\ \end{tabular} $ \begin{tabular}{ll} $v \in V \\ \end{tabular} $ \begin{tabular}{ll} $v_{i+1}$ \\ \hline \end{tabular} $ \begin{tabular}{ll} $cone $ \\ \end{tabular} $ \begin{tabular}{ll} $v \in V \\ \end{tabular} $ \begin{tabular}{
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You need to give the independent set identified by the algorithm and show that it is not maximal.

• (5 pts) Give an example path where the following algorithm also does not fint the maximum weigh independent set:

```
// largest independent sets 
// given G=(V,E) 
S_{even} = \{ v_i : i \in [1, ..., n] \text{ where } i \% 2 = 0 \} 
S_{odd} = \{ v_i : i \in [1, ..., n] \text{ where } i \% 2 = 1 \} 
if total_weight(S_{even}) > total_weight(S_{odd}) 
then return S_{even} 
then return S_{odd}
```

You need to give the independent set identified by the algorithm and show that it is not maximal.