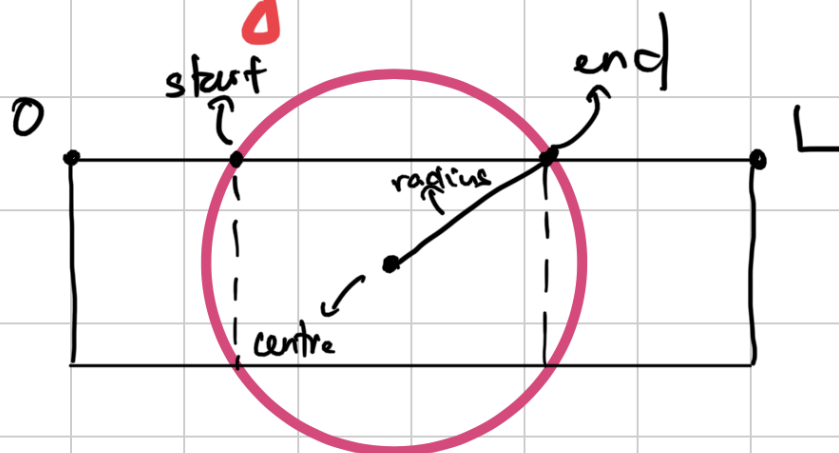


1.3 Algorithm



so that we always "stay ahead"

Recursion Idea:

- Sort Sprinkler by start point (natural order), then by radius (reverse order).

- Given n sprinklers:

$\text{MaxArea}(c, k)$ = the max area (starting from 0.0)

that can be covered by k sprinklers in the first c sprinklers ($k \leq c$)

e.g: $\text{MaxArea}(2, 4) = 3.0$

↳ The max area that can be covered by any 2 sprinklers in the first 4 sprinklers is 3.0

- Recursion:

k sprinklers in the first $c-1$ sprinkler
 ? \Rightarrow not include c^{th} sprinkler

$$\text{MaxArea}(c, k) = \text{Max} \left(\text{MaxArea}(c-1, k), \text{MaxArea}(c, k-1) + \text{Sprinkler } c^{\text{th}} \right)$$

- Recursion:

k sprinklers in the first $c-1$ sprinklers
 \Rightarrow not include c^{th} sprinkler

$$\text{MaxArea}(c, k) = \text{Max} \left(\text{MaxArea}(c-1, k), \text{MaxArea}(c, k-1) + \text{Sprinkler } c^{\text{th}} \right)$$

"exchange"
 only include c^{th} sprinkler
 if it gives better result

$k-1$ sprinklers in the first c sprinklers
 (may or may not include c^{th} sprinkler)

- Use memoization technique

$L = 5, 0$

c	Sprinklers	k	1	2	3
1	$[0, 3]$		3	3	3
2	$[0, 2]$		3	3	3
3	$[2, 4]$		3	4	4
4	$[4, 5]$		3	4	5

sorted
 start
 end

$$\text{table}[3, 2] = \text{Max} \left(\text{table}[2, 2], \text{table}[3, 1] + 3^{\text{rd}} \text{ sprinkler} \right)$$

$$= \text{Max} \left(\text{area from } 0 \rightarrow 3, \text{area from } 0 \rightarrow 3 + \text{area from } 2 \rightarrow 4 \right)$$

contiguous \rightarrow area from $0 \rightarrow 4$

	LU, 2J	0	3	5
2	[0, 2]	3	3	3
3	[2, 4]	3	4	4
4	[4, 5]	3	4	5

sorted (vertical arrow pointing down)
start (arrow pointing to [0, 2])
end (arrow pointing to [4, 5])

$$\text{table}[3, 2] = \text{Max} \left(\text{table}[2, 2], \text{table}[3, 1] + 3^{\text{rd}} \text{ sprinkler} \right)$$

$$= \text{Max} \left(\text{area from } 0 \rightarrow 3, \text{area from } 0 \rightarrow 3 + \text{area from } 2 \rightarrow 4 \right)$$

contiguous \rightarrow area from $0 \rightarrow 4$

$$= \text{area from } 0 \rightarrow 4$$

$$= 4$$

$$\text{table}[4, 2] = \text{Max} \left(\text{table}[3, 2], \text{table}[4, 1] + 4^{\text{th}} \text{ sprinkler} \right)$$

$$= \text{Max} \left(\text{area from } 0 \rightarrow 4, \text{area from } 0 \rightarrow 3 + \text{area from } 4 \rightarrow 5 \right)$$

not contiguous \Rightarrow area from $0 \rightarrow 3$

$$= \text{area from } 0 \rightarrow 4$$

ALGORITHM

* find Minimum (Sprinklers, L, W):

Sort Sprinklers()

Verify All Sprinklers Can Cover the Garden()

table = array [no. of sprinklers] [no. of sprinklers]

Populate the 1st column & 1st row with the first sprinkler's end.

for i in range(1, no. of sprinkler):

table[0, i] = 1st sprinkler's end.

table[i, 0] = 1st sprinkler's end

if sprinkler's end \geq L:

return 1

for k in range(2, no. of sprinkler):

for c in range(2, no. of sprinkler):

a = table[c-1, k]

b = table[c, k-1] + sprinklers[c]

table[c, k] = max(a, b)

if table[c, k] \geq L:

return k

return 0;