

CS3640

Overview (3): Performance & Protocols

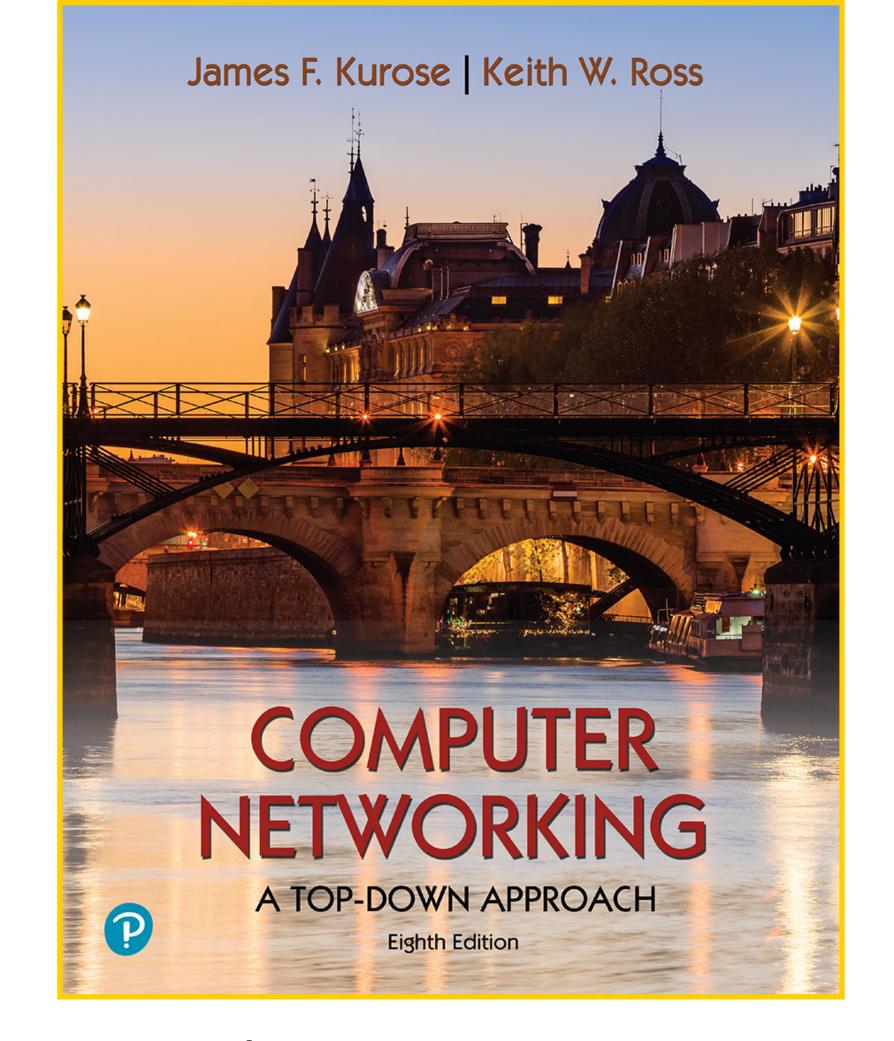
Prof. Supreeth Shastri

Computer Science
The University of Iowa

Lecture goals

Continuing our in-depth exploration into the structure and functioning of the Internet

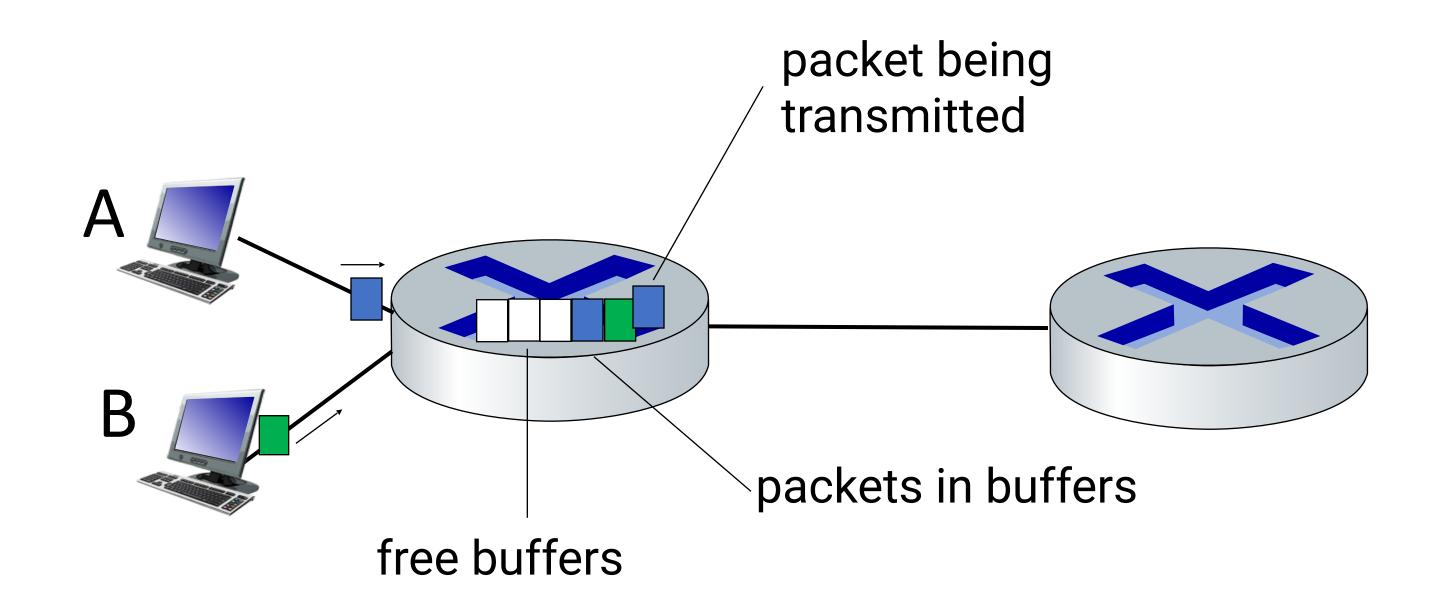
- Performance: delay and throughput
- Protocol architecture



Chapter 1.4 - 1.5

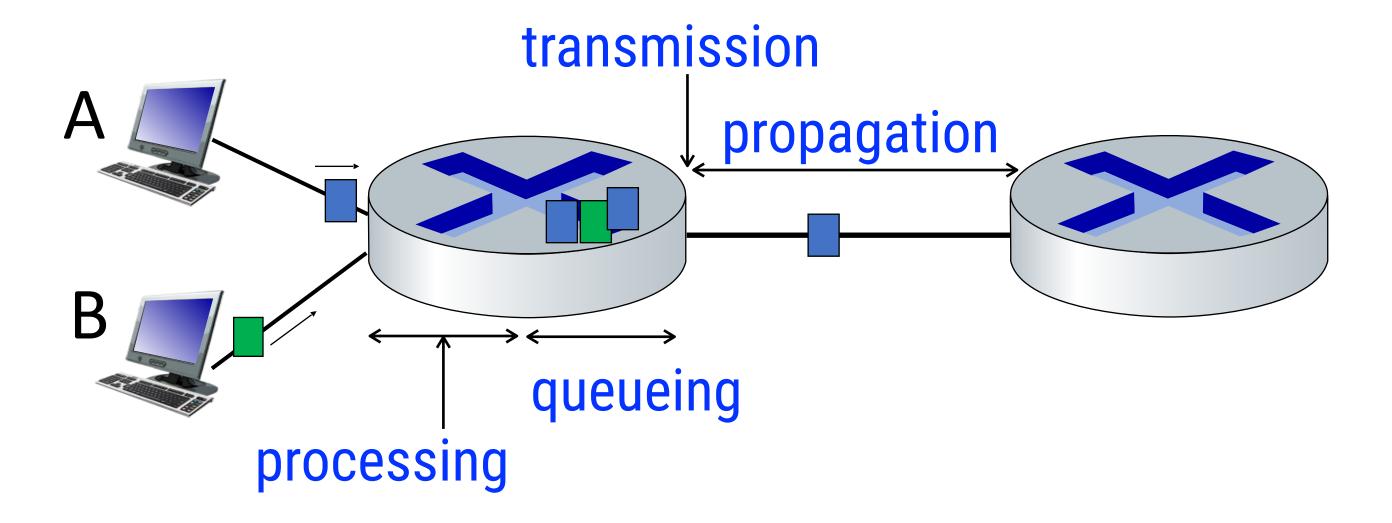


Queueing and loss in packet switching



- packets are queued in when arrival rate for a link exceeds its output capacity
- packet are dropped when memory to hold queued packets fills up
 - lost packet may be retransmitted by previous node, by source, or not at all

Packet delay: four sources



$$d_{\text{total}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{proc}: processing

- check bit errors; determine output link
- typically < microseconds

d_{queue}: queueing

- time waiting at output link for transmission
- depends on router's congestion level

d_{trans}: transmission

- L: packet length (bits)
- R: link transmission rate (bps)
- $d_{trans} = L/R$

d_{prop}: propagation

- d: length of physical link
- s: propagation speed (~2x10⁸ m/sec)
- $d_{prop} = d/s$

Understanding traffic intensity

- a: avg. arrival rate (packets/sec)
- L: avg. packet length (bits/packet)
- R: link transmission rate (bits/sec)

traffic intensity

 $\frac{L \cdot a}{R}$ arrival rate of bits



La/R ~ 0 Avg delay $\approx none$

La/R \rightarrow 1 Avg delay depends on arrival distribution



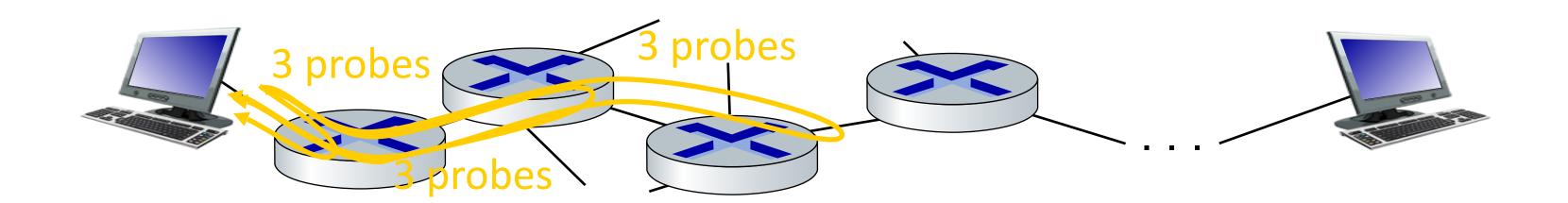
La/R > 1
Avg delay \rightarrow *infinity*

Quantifying delays in the "real" Internet

traceroute: a tool that provides delay measurement from source to router along end-end Internet path towards destination.

For all i

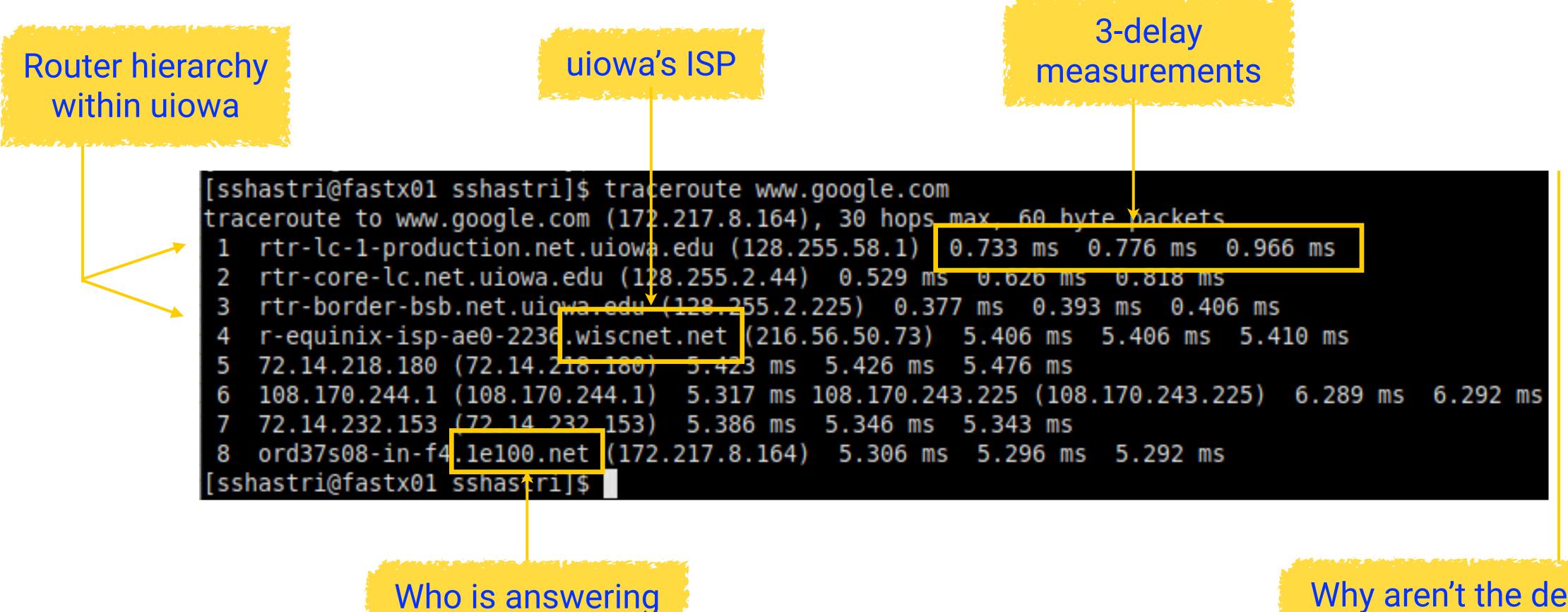
- (i) send three packets that will reach router i on path towards destination (with time-to-live field set to i)
- (ii) router i will return packets to sender
- (iii) sender measures time interval between transmission and reply



traceroute demo

Quantifying delays in the "real" Internet

traceroute: from fastx01.divms.uiowa.edu to www.google.com



for google?

Why aren't the delays strictly increasing?

Quantifying delays in the "real" Internet

traceroute: from fastx01.divms.uiowa.edu to www.wimbledon.org

uiowa is using more then one ISP

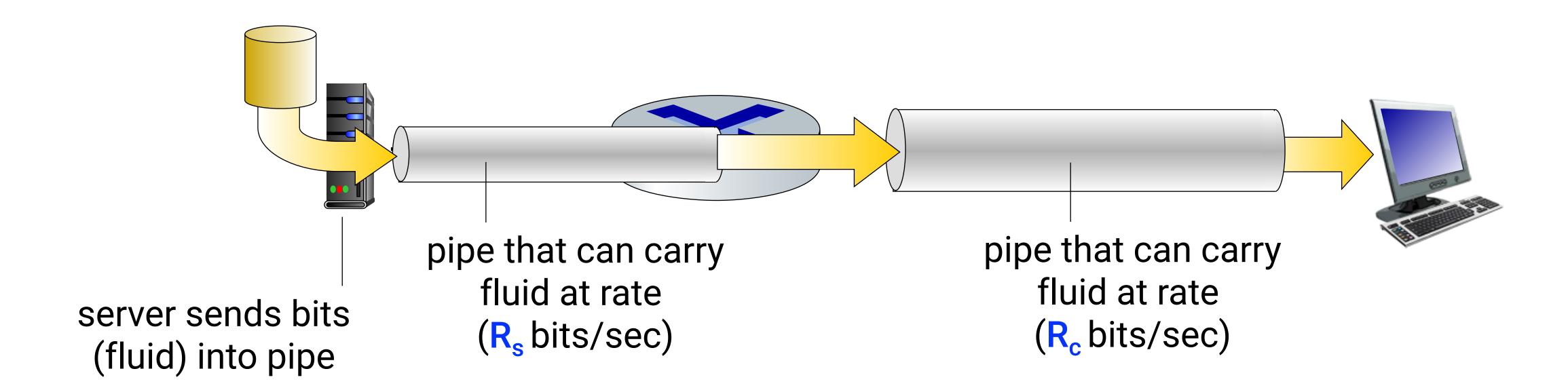
```
[sshastri@fastx01 sshastri]$ traceroute www.wimbledon.org
traceroute to www.wimbledon.org (104.114.79.50), 30 hops max, 60 byte packets
1 rtr-lc-1-production.net.uiowa.edu (128.255.58.1) 0.921 ms 0.921 ms 0.906 ms
2 rtr-core-lc.net.uiowa.edu (128.255.2.44) 0.610 ms 0.872 ms 1.027 ms
3 rtr-border-bsb.net.uiowa.edu (128.255.2.225) 0.415 ms 0.344 ms 0.404 ms
4 et-5-1-5-102.cr1-min1.ip4.gtt.net (208.116.156.121) 8.262 ms 8.274 ms 8.261 ms
5 ae19.cr9-chi1.ip4.gtt.net (141.136.108.189) 16.754 ms 27.958 ms 16.725 ms
6 ip4.gtt.net (98.124.183.18) 25.264 ms 24.924 ms 24.907 ms
7 ae3.ctl-ord3.netarch.akamai.com (23.203.151.229) 17.680 ms 17.637 ms 15.612 ms
8 a104-114-79-50.deploy.static.akamaitechnologies.com (104.114.79.50) 11.331 ms 11.332 ms 11.329 ms
```

CDN! No cross Atlantic traffic Delays are longer than google.com

Throughput

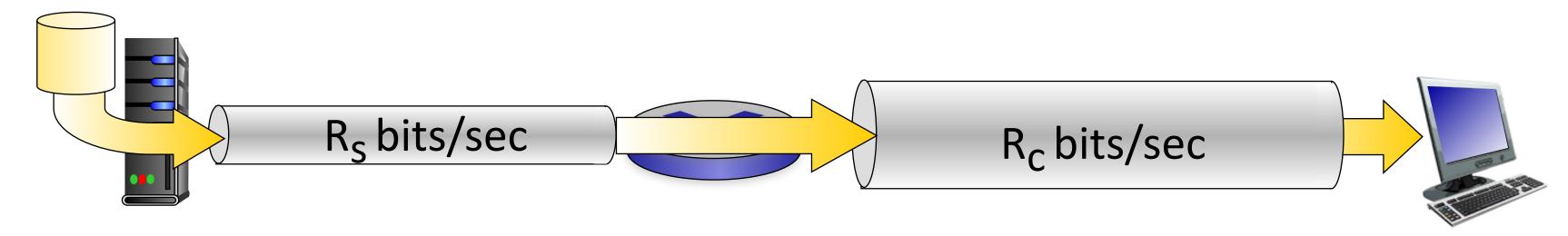
Rate (measured in bits/sec) at which bits are being sent from sender to receiver

- instantaneous: rate at a given point in time
- average: rate over a longer period of time

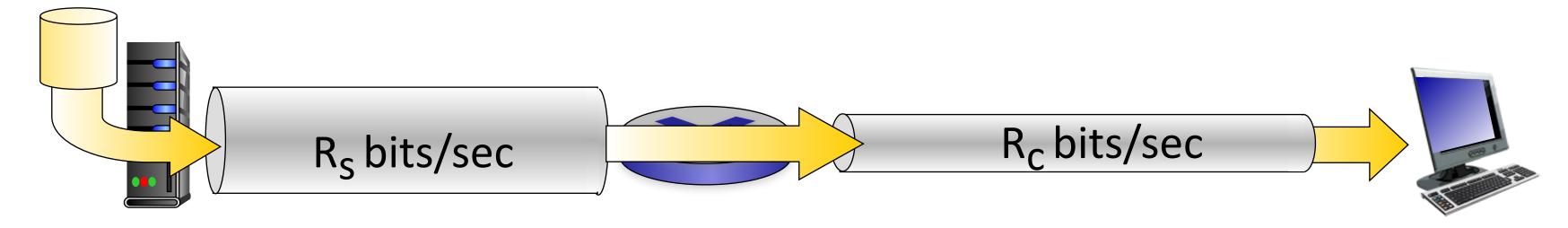


Throughput: bottleneck link

 $R_s < R_c$ What is average end-end throughput?



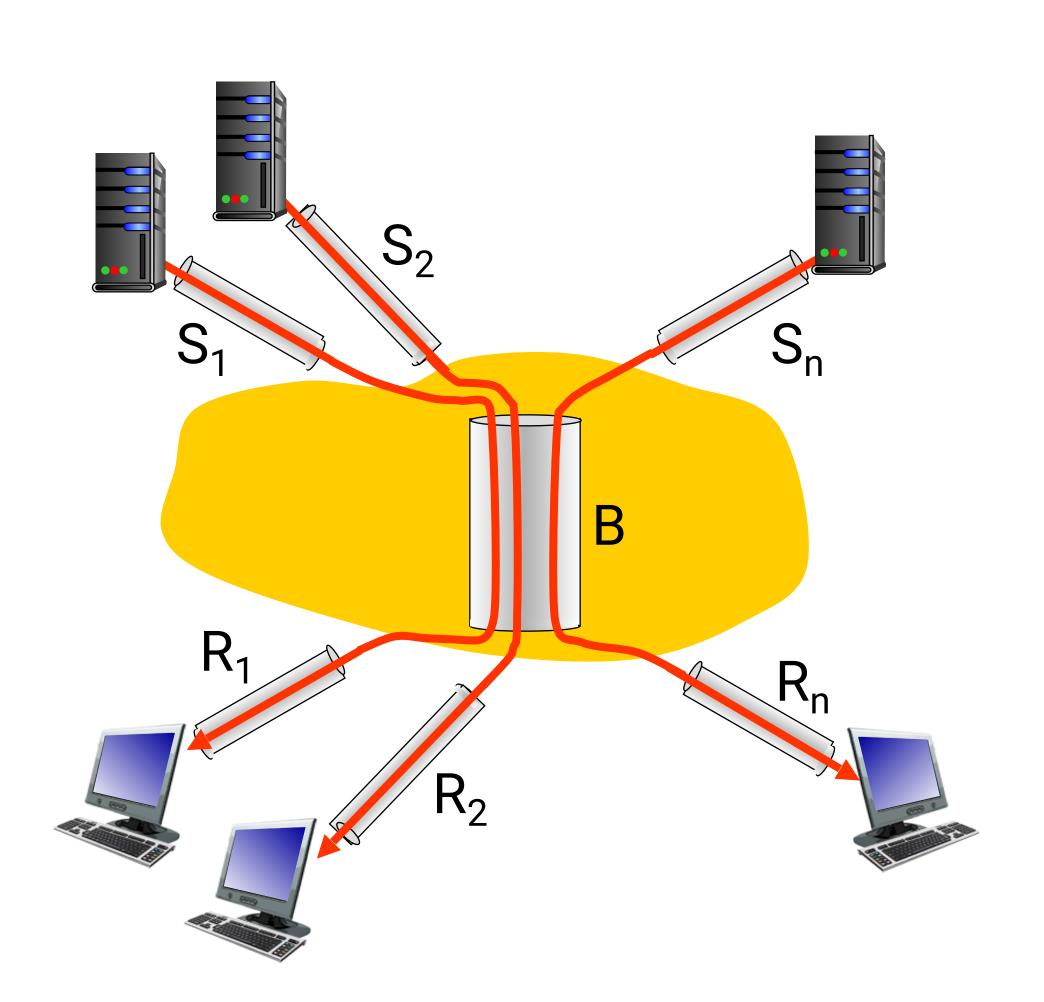
 $R_s > R_c$ What is average end-end throughput?



bottleneck link

link on end-end path that constrains the end-end throughput

Throughput: network scenario



- n connections fairly share the backbone link (B bits/sec)
- End-end throughput for connection $k = min(S_k, R_k, B/n)$
- in practice: S_k or R_k is often the bottleneck

Protocol layering

What is layering?

An approach to designing complex systems

allows identifying system's components and explicitly defining their relationship

Modularization eases maintenance and updating of system

change in layer's service implementation: transparent to rest of system

Why is layering useful for computer networks?

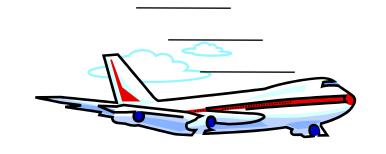
Computer networks have multitude of components interacting with each other

host devices, routers, links, protocols, applications, policies, and so on

Internet is arguably the largest engineered system ever created by humans!

A layered system: air travel

a complex system involving people, goods, airplanes, airports, and services



end-to-end transfer of person plus baggage

ticketing (purchase)
baggage (check)
gates (load)
airplane takeoff

ticketing (complain)
baggage (claim)
gates (unload)
airplane landing

airplane routing

A layered system: air travel

each layer implements a service and relies on services provided by layer below

ticket (purchase)	ticketing service	ticket (complain)	
baggage (check)	baggage service	baggage (claim)	
gates (load)	gate service	gates (unload)	
airplane takeoff	runway service	airplane landing	
routing	routing service	routing	

The five layer architecture of the Internet

application transport network link physical

- Application: supporting network applications.
 E.g., HTTP, IMAP, SMTP, DNS
- Transport: process to process data transfer.
 E.g.,TCP, UDP
- Network: routing of datagrams from source machine to destination. E.g., IP, IPv6
- Link: deliver data between neighboring network elements. E.g., Ethernet, 802.11 (WiFi)
- Physical: bits "on the wire". E.g., 10BASE-T

Protocol layering and services

application

transport

network

link

physical

Application exchanges messages to implement some application service using *services* of transport layer

 $H_t M - M$

Transport-layer protocol transfers M from one *process* to another, using services of network layer

- transport-layer protocol encapsulates applicationlayer message, M, with transport layer-layer header H_t to create a transport-layer segment
- H_t used by transport layer protocol to implement its service

application

transport

network

link

physical





Protocol layering and services

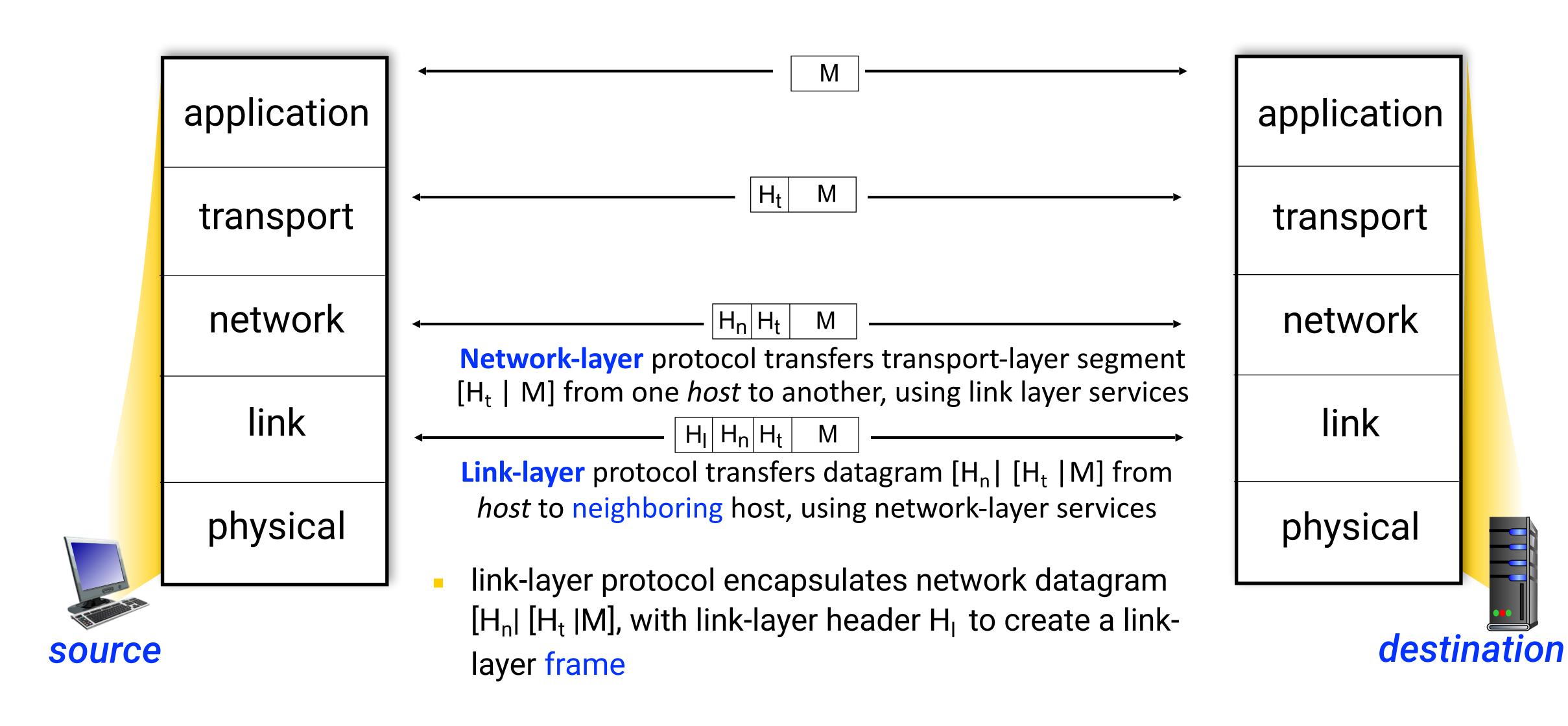
M application transport Transport-layer protocol transfers M from one *process* to another, using services of network layer network $H_n|H_t$ Network-layer protocol transfers transport-layer segment [H_t | M] from one *host* to another, using link layer services link network-layer protocol encapsulates transportlayer segment [H_t | M] with network layer-layer physical header H_n to create a network-layer datagram. H_n used by network layer protocol to implement its service

source

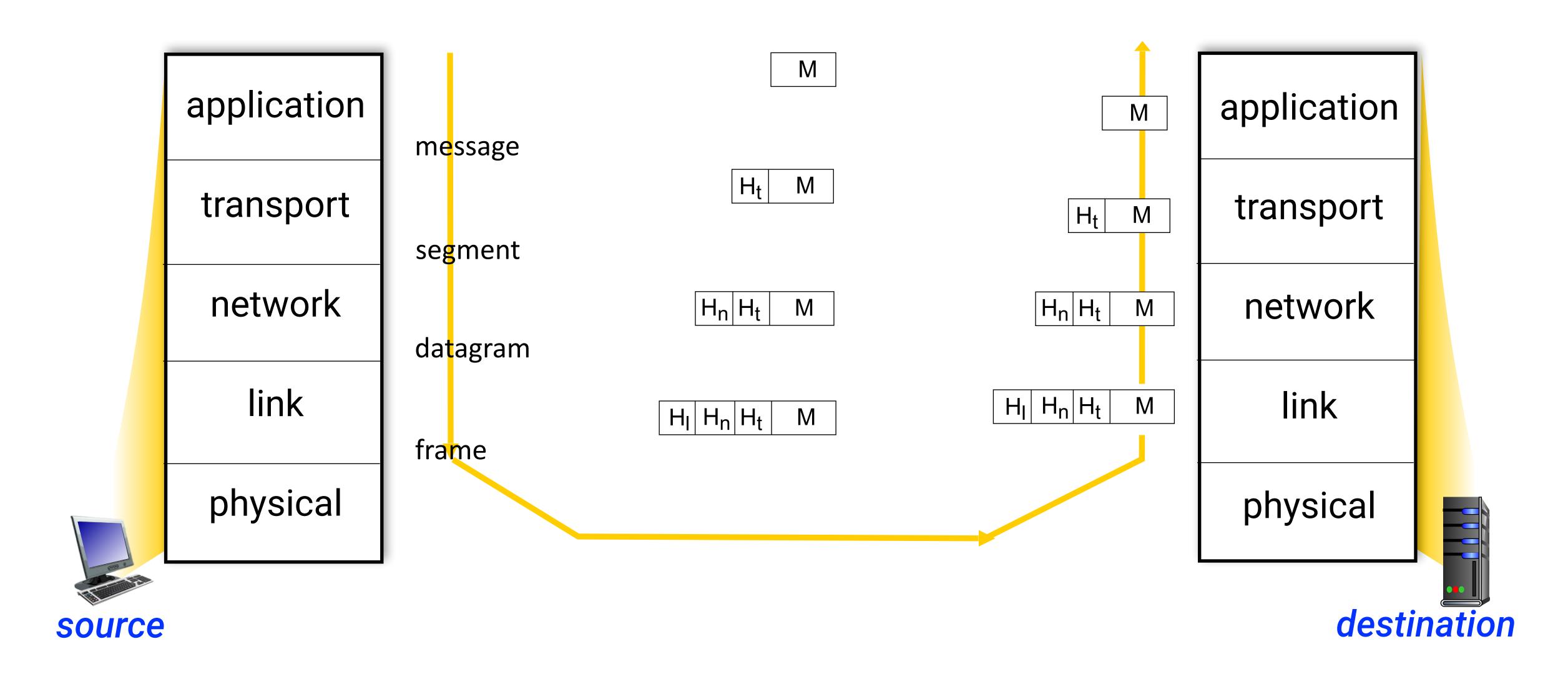
application transport network link physical

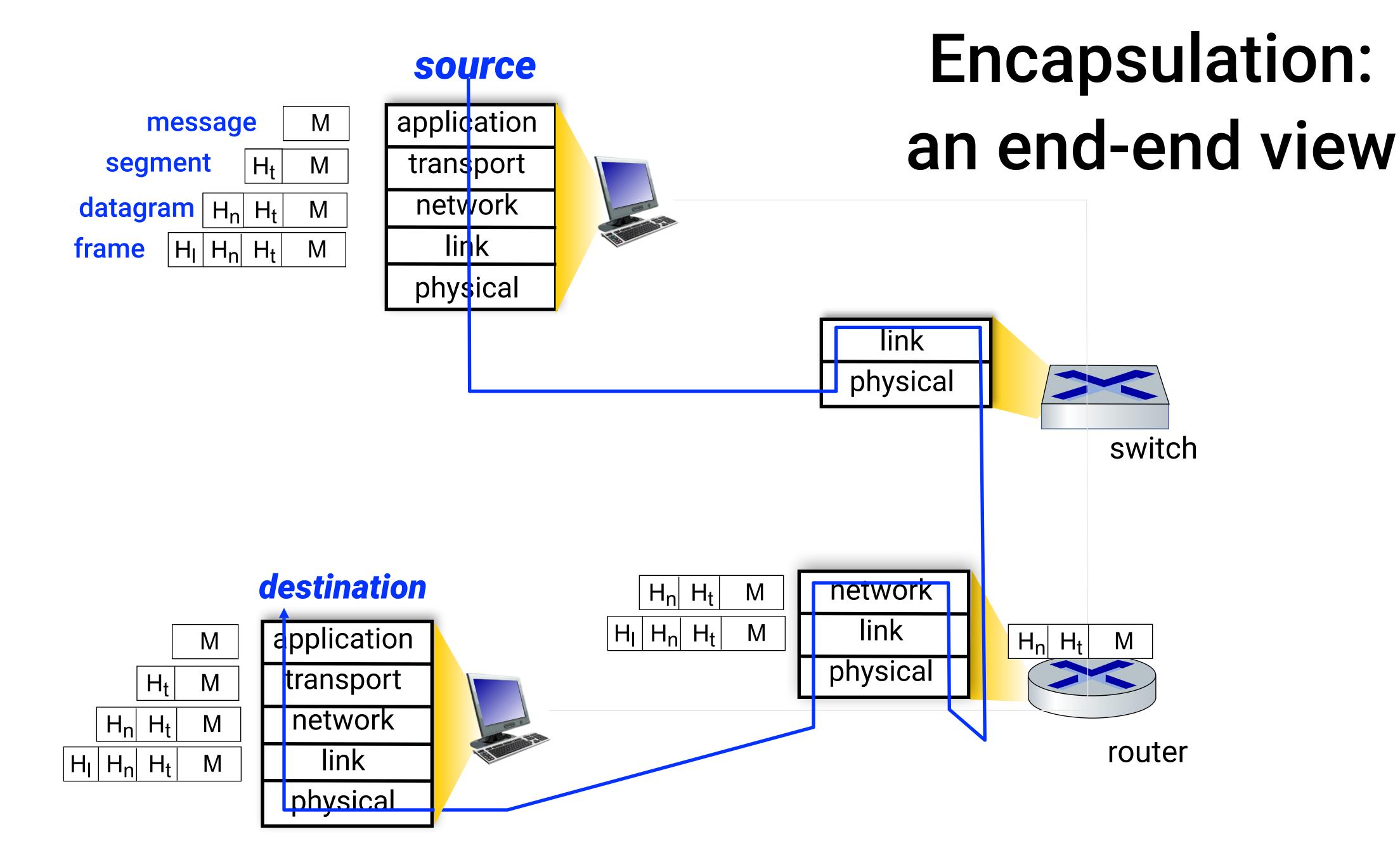
destination

Protocol layering and services



Protocol encapsulation

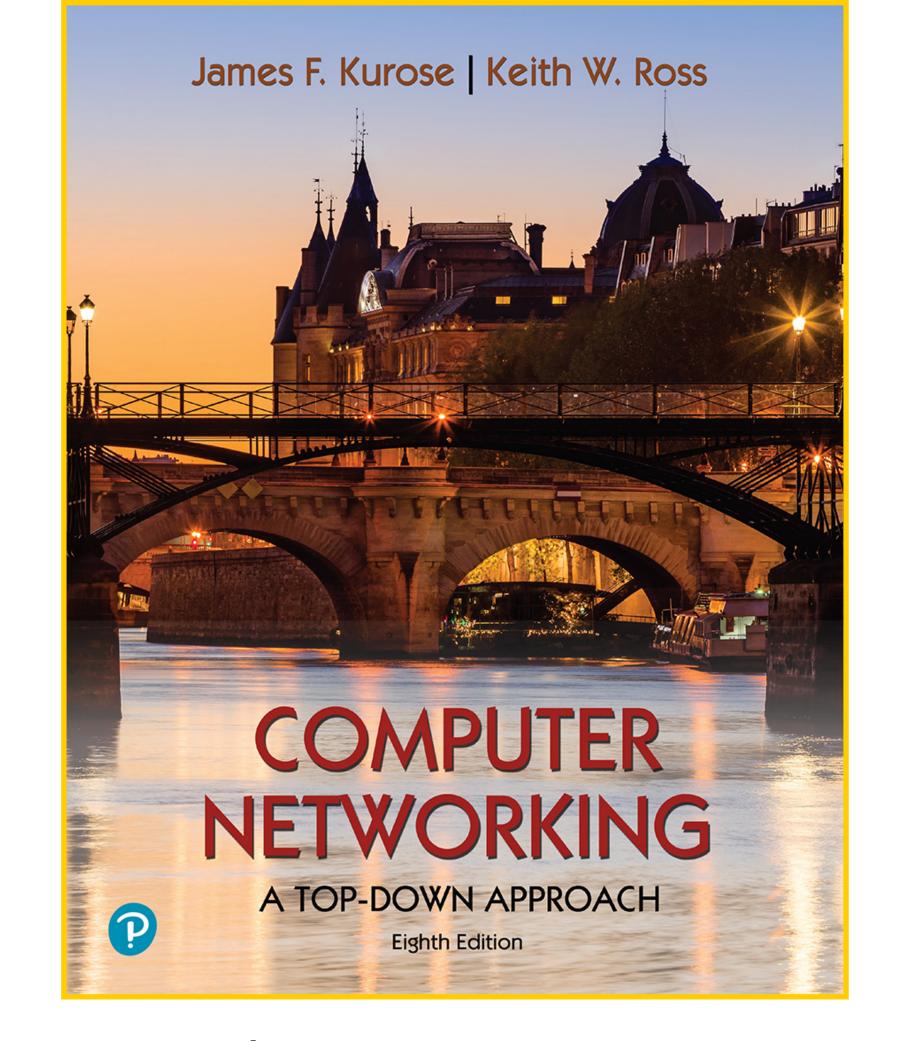




Next Lecture

Continuing our in-depth exploration into the structure and functionality of the Internet

- Network security
- Internet history and evolution



Chapter 1.6 - 1.7



Spot Quiz (ICON)