

CS3640

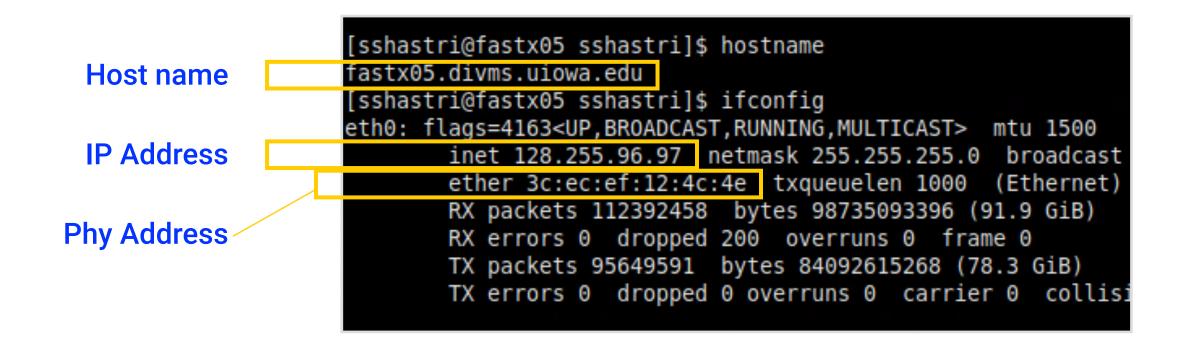
Link Layer (2): Addressing and Ethernet

Continuing from previous lecture

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The Three Network Identities



Address Resolution Protocol (ARP)

Network Working Group Request For Comments: 826 David C. Plummer (DCP@MIT-MC) November 1982

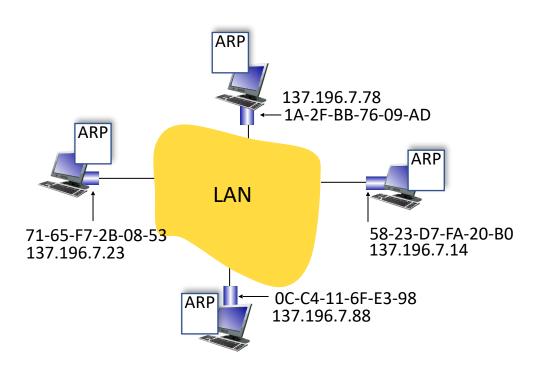
An Ethernet Address Resolution Protocol
-- or -Converting Network Protocol Addresses
to 48.bit Ethernet Address
for Transmission on
Ethernet Hardware

Abstract

The implementation of protocol P on a sending host S decides, through protocol P's routing mechanism, that it wants to transmit to a target host T located some place on a connected piece of 10Mbit Ethernet cable. To actually transmit the Ethernet packet a 48.bit Ethernet address must be generated. The addresses of

Address Resolution Protocol (ARP)

Given the IP address of an interface, how to determine its MAC address?

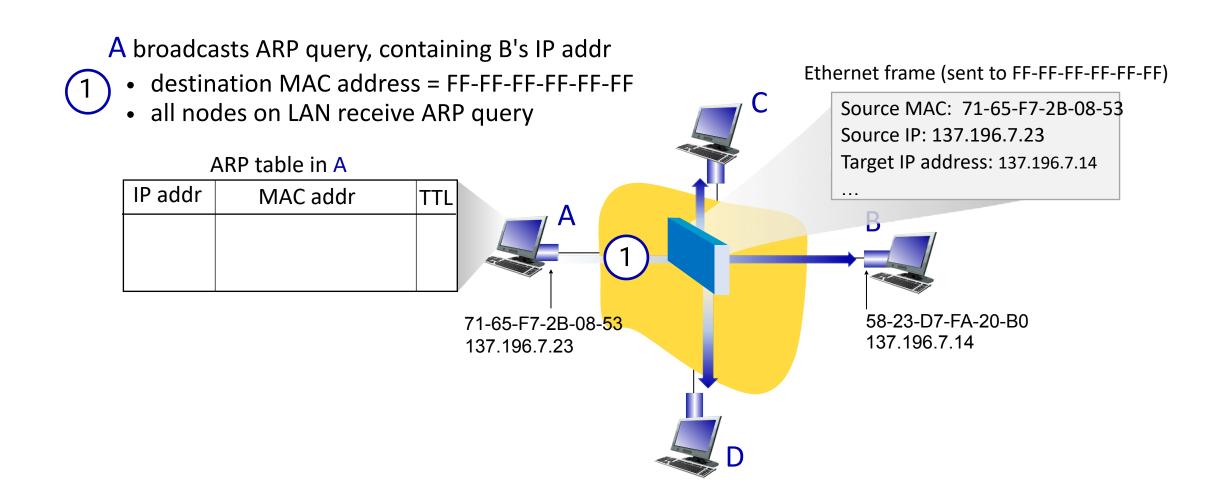


Every node (host, router) has an ARP table

- ARP table contains IP to MAC address mappings for nodes on the same LAN
- Three tuple: <IP address; MAC address; TTL>
- TTL (Time To Live): time after which address mapping will be forgotten (typically 20 min)

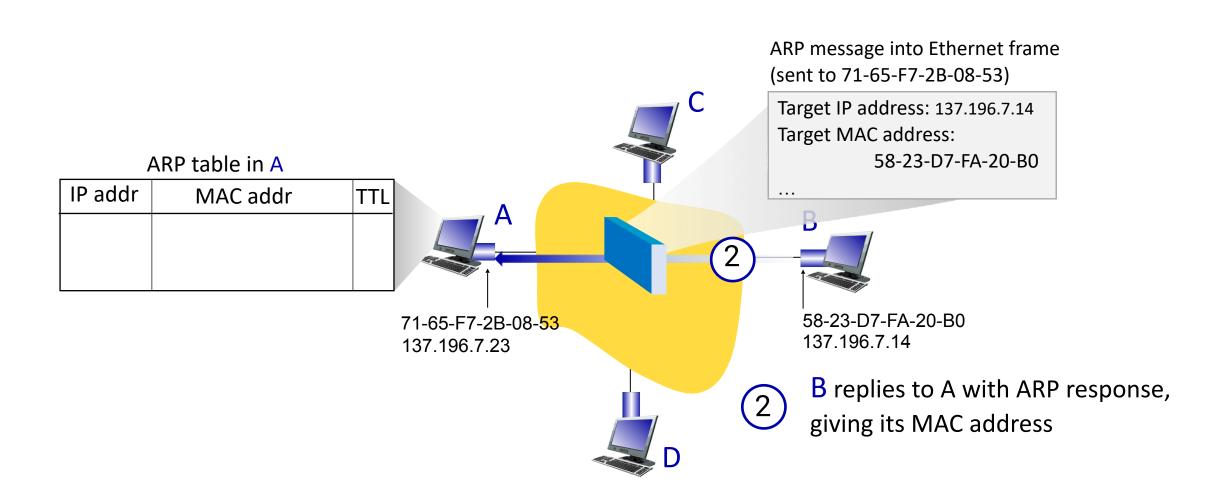
ARP in Action

E.g., A wants to send datagram to B B's MAC address not in A's ARP table, so A uses ARP to find B's MAC address



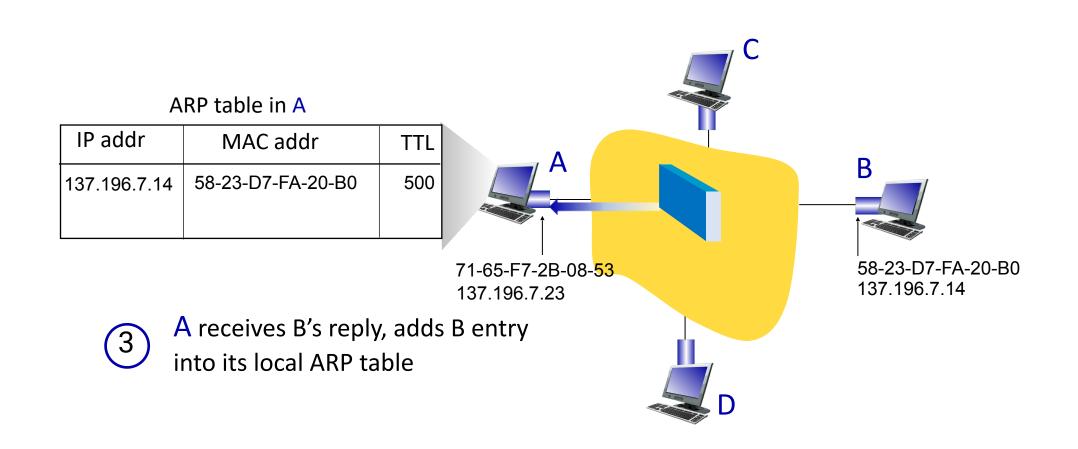
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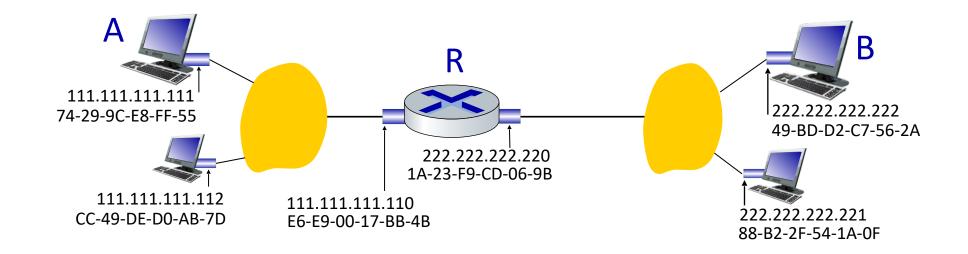
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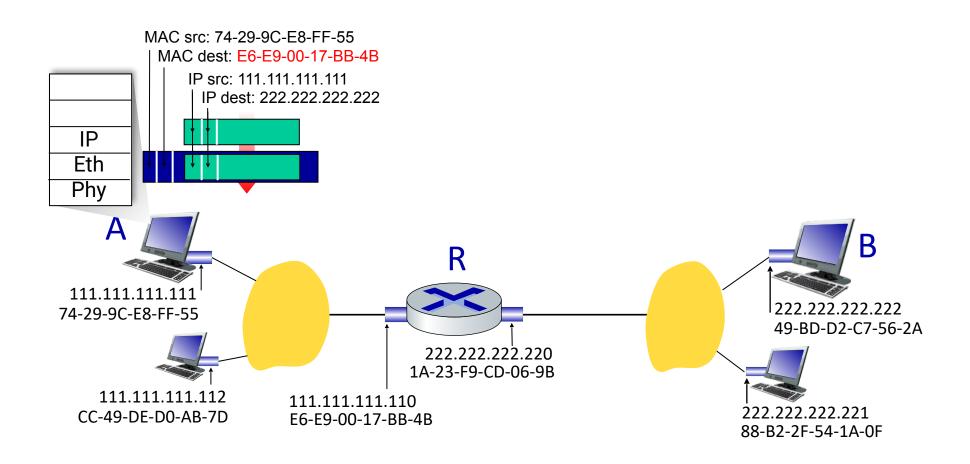


E.g., A wants to send datagram to B via R

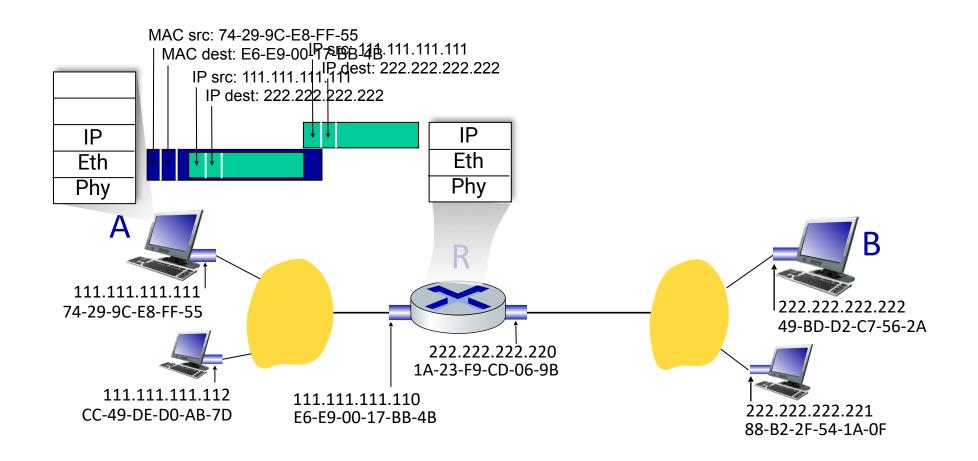
A knows B's IP address; A also knows R's IP address and MAC address (how?)



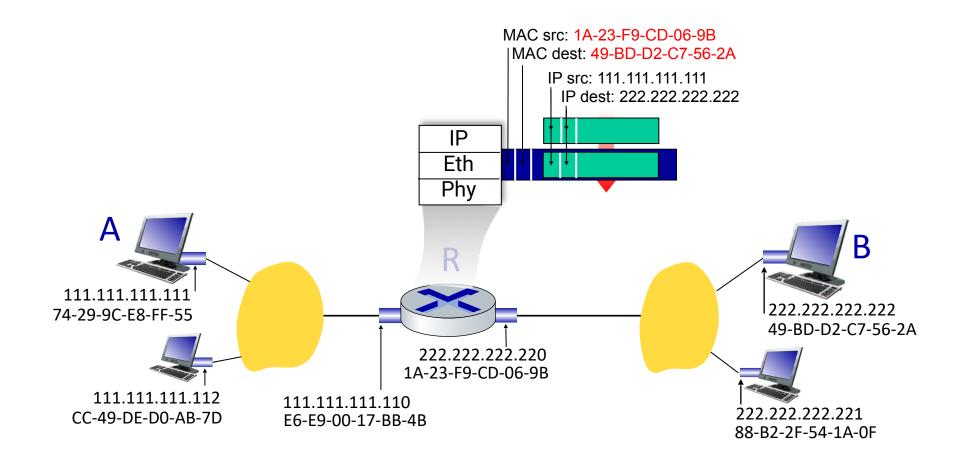
- A creates IP datagram with IP source A, destination B
- A creates link-layer frame containing A-to-B IP datagram,
 then puts R's MAC address is the frame's destination



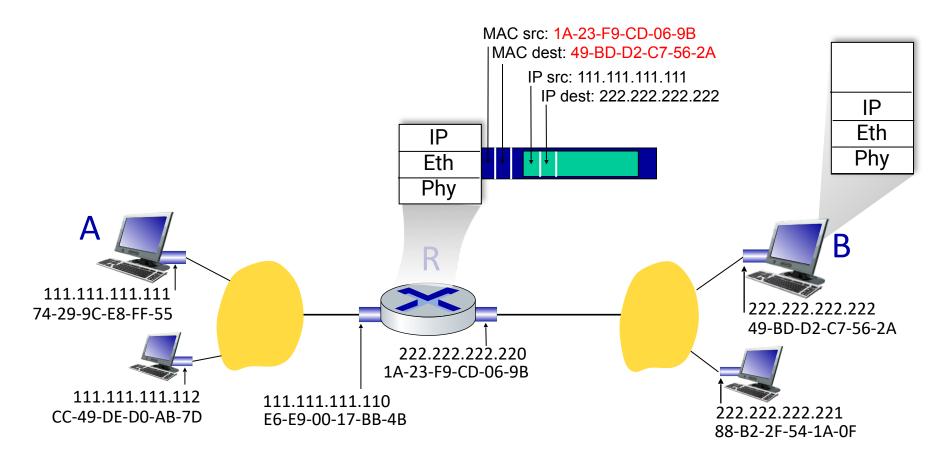
- frame sent from A to R
- frame received at R, datagram removed, passed up to IP



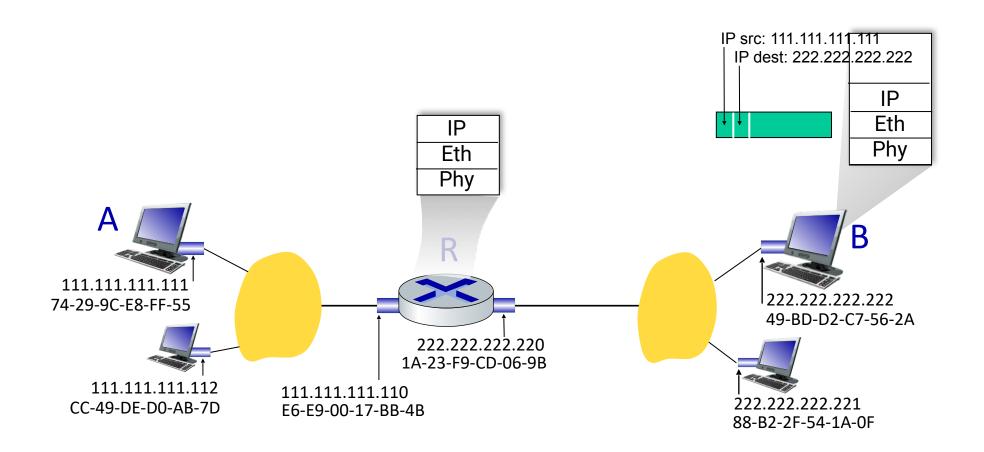
- R determines outgoing interface, passes datagram with IP source A, destination B to link layer
- R creates link-layer frame containing A-to-B IP datagram. Frame destination address: B's MAC address



- R determines outgoing interface, passes datagram with IP source A, destination B to link layer
- R creates link-layer frame containing A-to-B IP datagram. Frame destination address: B's MAC address
- transmits link-layer frame



- B receives frame, extracts IP datagram destination B
- B passes datagram up protocol stack to IP





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Link Layer (3): A Day in The Life of a Packet

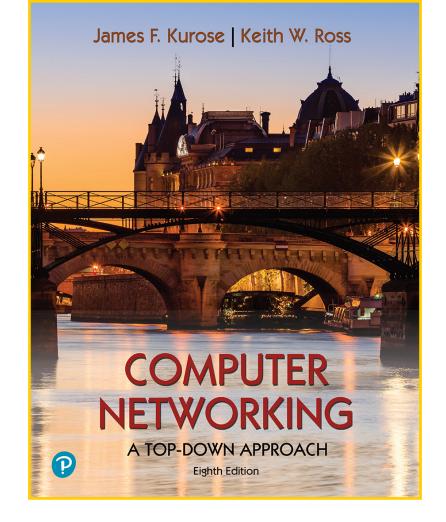
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Lecture Goals

Retrospective: a day in the life of a web page request

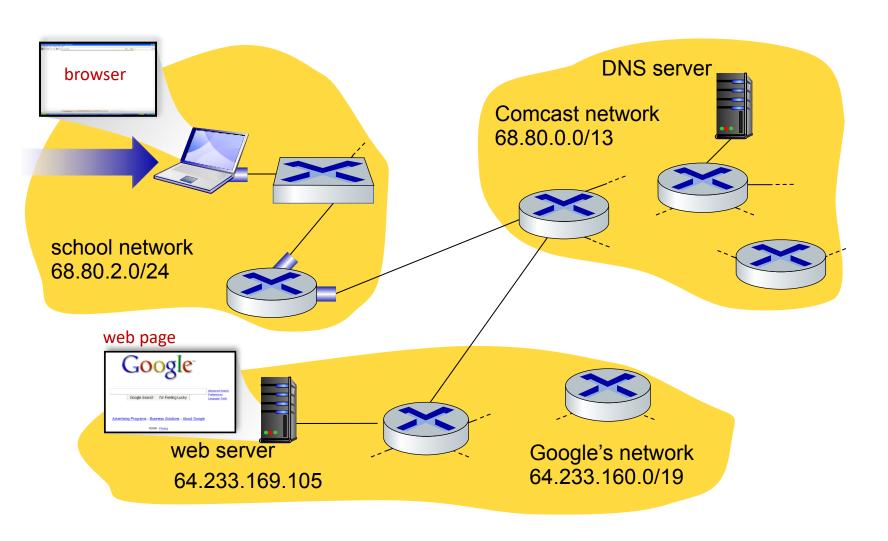
- touches all five layers of the stack
- utilizes more than a dozen protocols
- end-to-end flow of control and data



Chapter 6.7



A day in the life: scenario

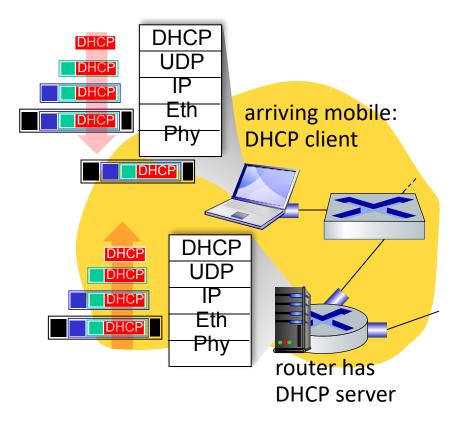


 arriving mobile client attaches to network

requests web page: www.google.com

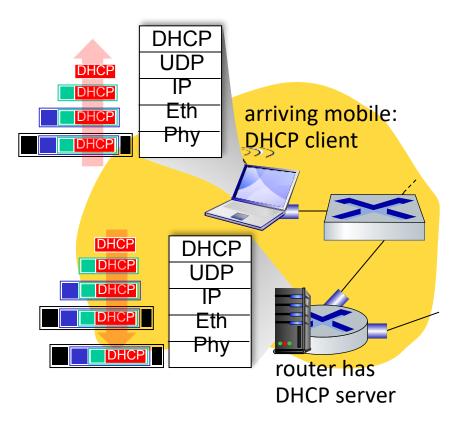


A day in the life: connecting to the Internet



- connecting laptop needs to get its own IP address, address of first-hop router, address of DNS server: use DHCP
- DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in 802.3 Ethernet
- Ethernet frame broadcast (dest: FF:FF:FF:FF:FF)
 on LAN, received at the router running DHCP server
- Ethernet demuxed to IP demuxed to UDP demuxed to DHCP

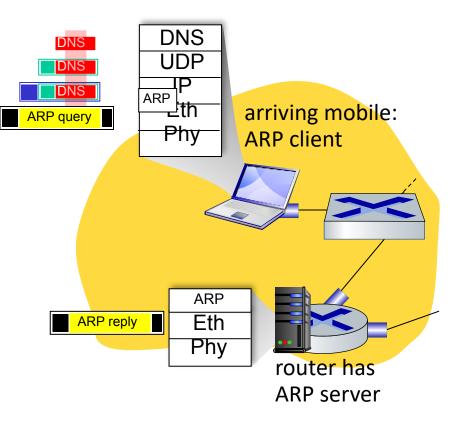
A day in the life: connecting to the Internet



- DHCP server formulates DHCP ACK containing client's IP address, IP address of first-hop router, name and IP address of DNS server
- encapsulation at DHCP server, frame forwarded through LAN, demultiplexing at client
- DHCP client receives DHCP ACK reply

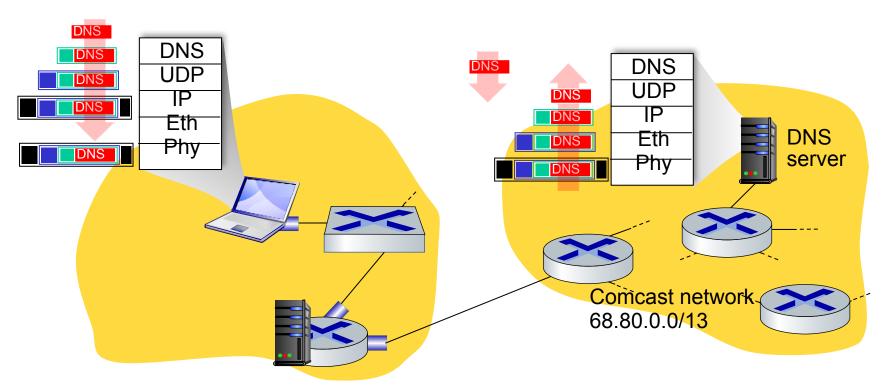
Client now has IP address, knows name and address of **DNS server** and IP address of its **first-hop router**

A day in the life: ARP (before DNS, before HTTP)



- before sending HTTP request, need IP address of www.google.com: use DNS
- DNS query created, encapsulated in UDP, encapsulated in IP, encapsulated in Ethernet. However, to send frame to router, need MAC address of router interface: use ARP
- ARP query broadcast, received by router, which replies with ARP reply giving MAC address of router interface
- client now knows MAC address of first hop router, so can now send frame containing DNS query

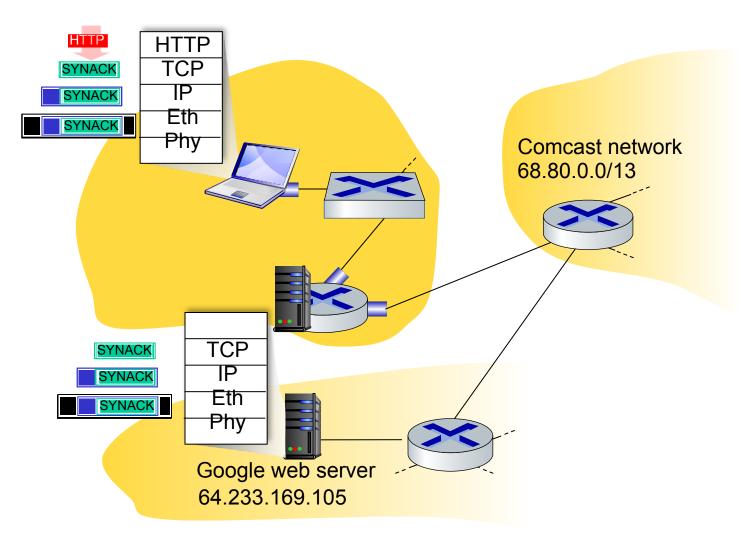
A day in the life: using DNS



- demuxed to DNS
- DNS replies to client with IP address of www.google.com

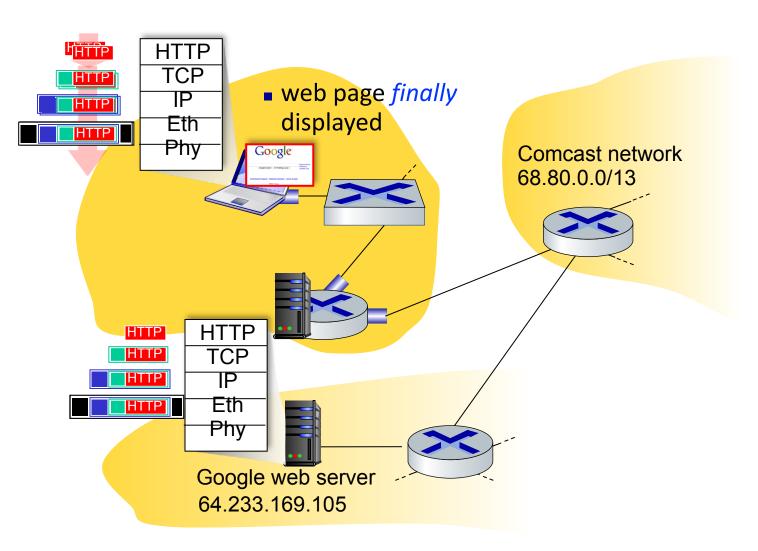
- IP datagram containing DNS query forwarded via LAN switch from client to first-hop router
- IP datagram forwarded from campus network into Comcast network, routed (tables created by RIP, OSPF, and/or BGP routing protocols) to DNS server

A day in the life: TCP connection carrying HTTP



- To send HTTP request, client first opens TCP socket to web server
- TCP SYN segment (step 1 in TCP 3way handshake) inter-domain routed to web server
- web server responds with TCP SYNACK (step 2 in TCP 3-way handshake)
- TCP connection established!

A day in the life: HTTP request/reply

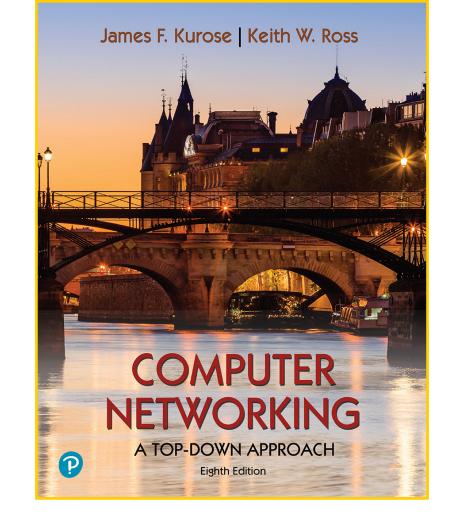


- HTTP request sent into TCP socket
- IP datagram containing HTTP request routed to www.google.com
- web server responds with HTTP reply (containing web page)
- IP datagram containing HTTP reply routed back to client

Next two lectures

Research topics to expand our horizon and get a taste of the state-of-the-art in networking

- Software-Defined Networking
- Cloud Computing



Chapters 4.4, 5.5



Spot Quiz (ICON)