

CS3640

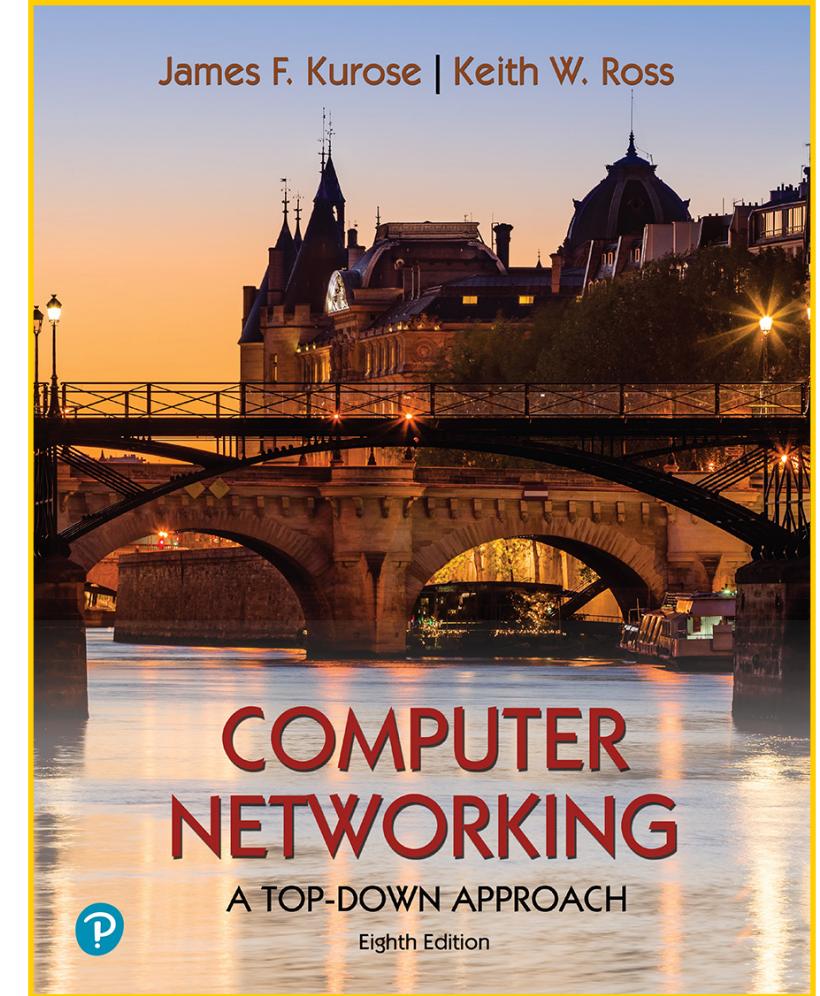
Overview (2): Network Edge & Core

Prof. Supreeth Shastri
Computer Science
The University of Iowa

Lecture goals

In-depth look into the structure and functioning of the Internet

- *Network Edge*
- *Network Core*
- *Packet Switching*

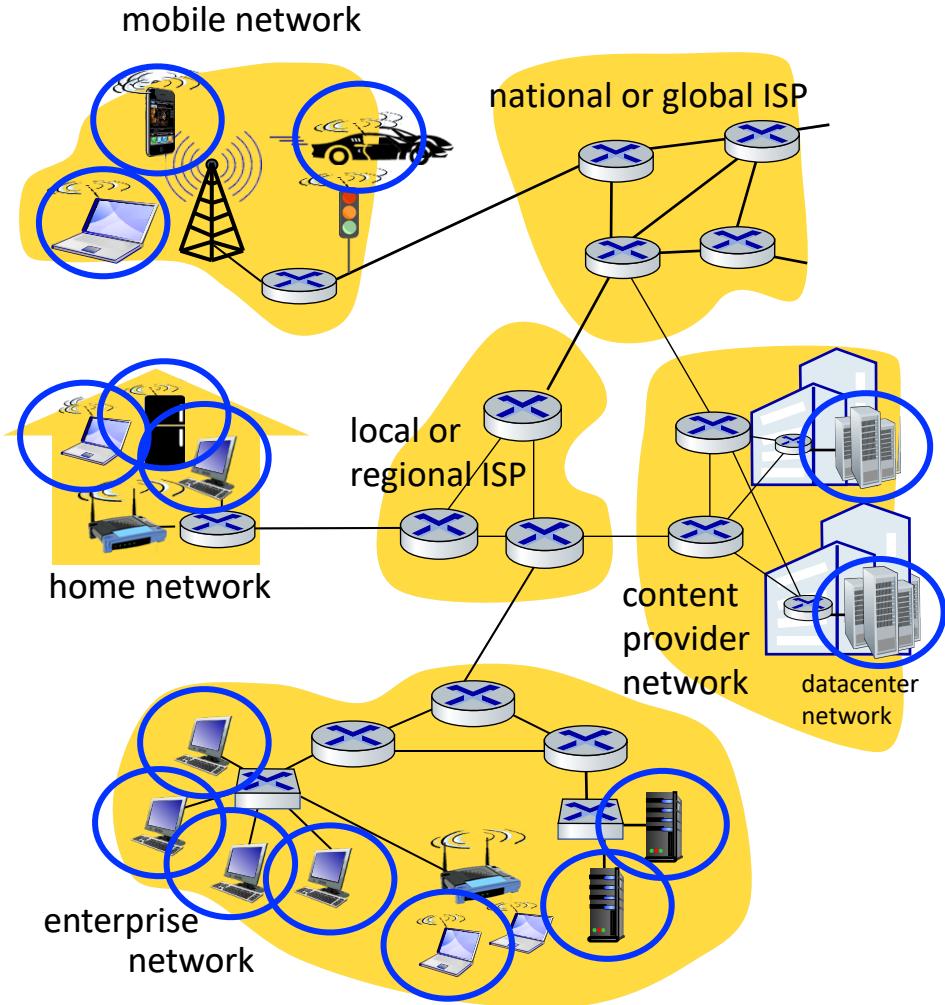


Chapter 1.2 - 1.3

A closer look at Internet structure

Network edge

- hosts: clients and servers
- servers often in data centers



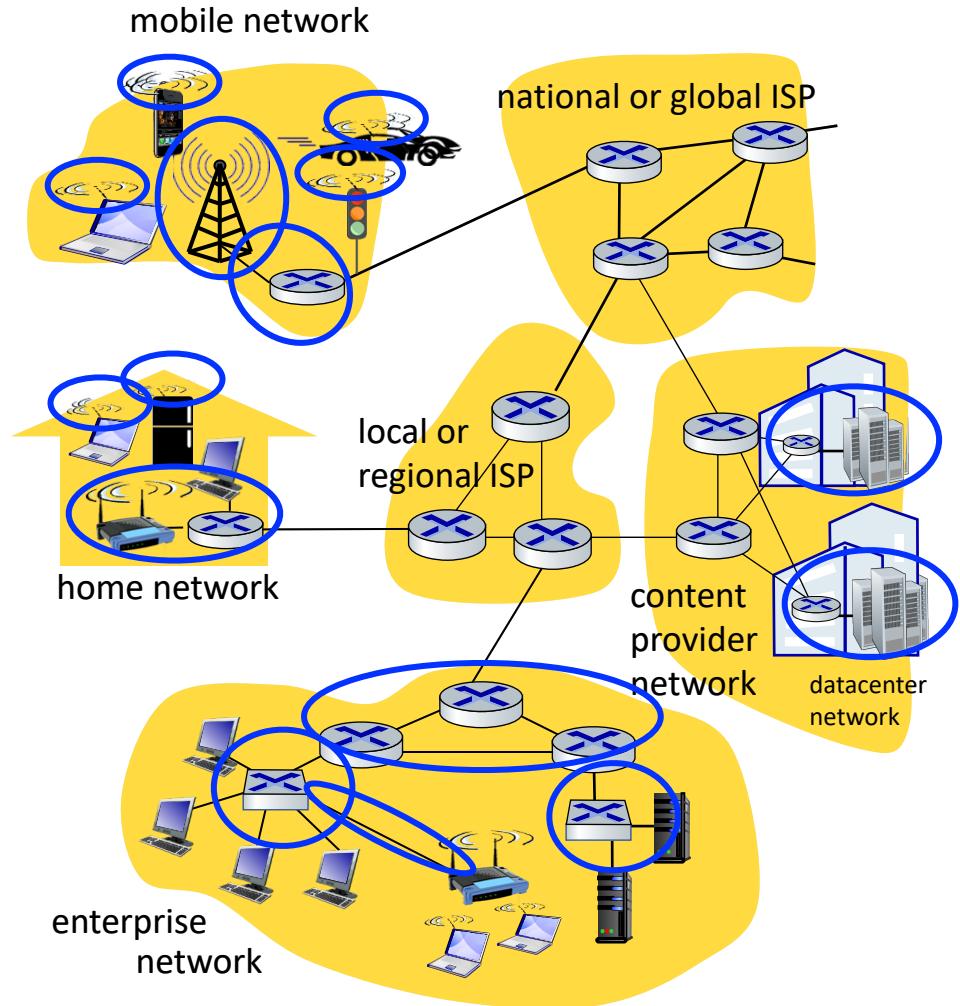
A closer look at Internet structure

Network edge

- hosts: clients and servers
- servers located in data centers

Access networks, physical media

- wired, wireless communication links



A closer look at Internet structure

Network edge

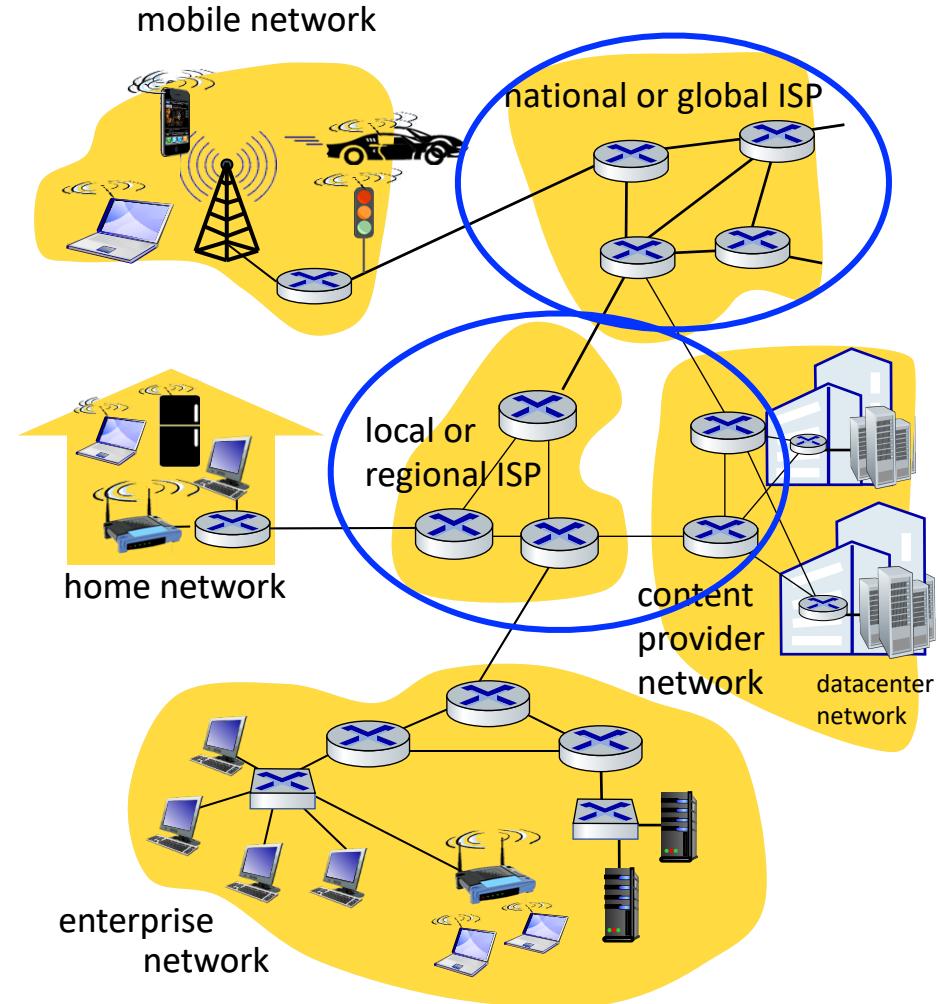
- hosts: clients and servers
- servers often in data centers

Access networks, physical media

- wired, wireless communication links

Network core

- interconnected routers
- network of networks



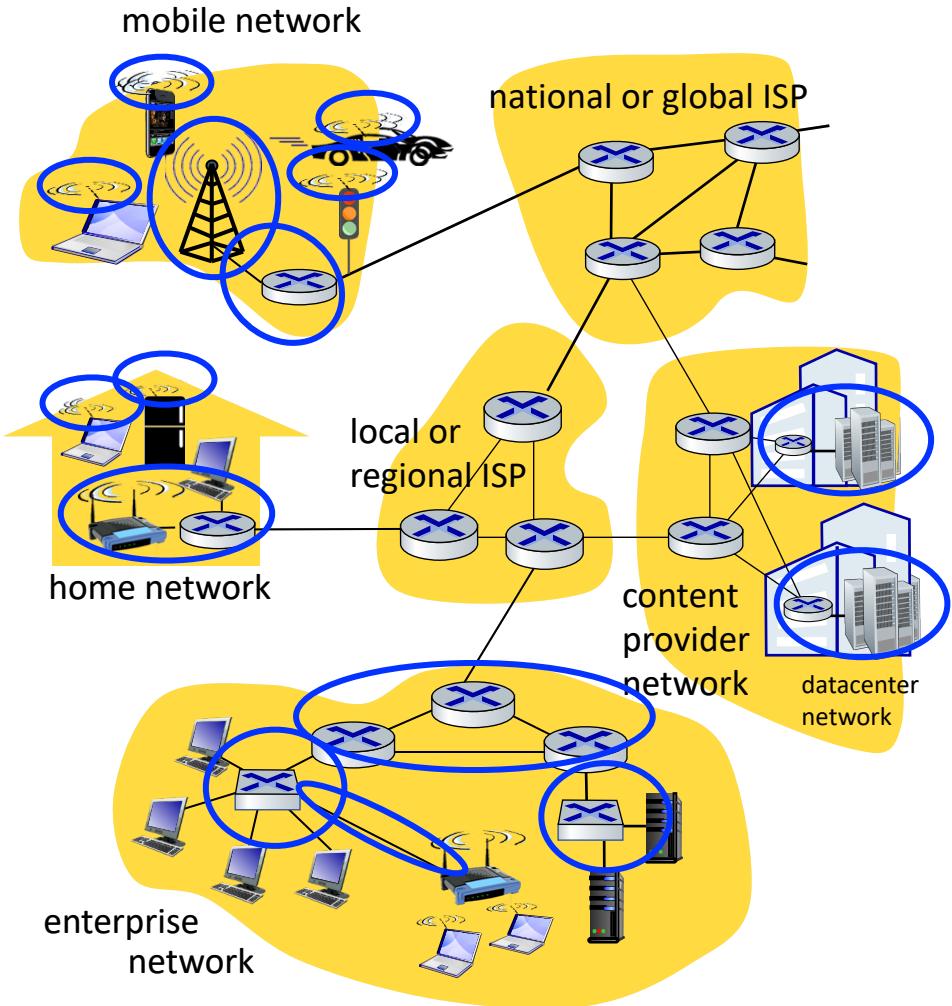
Access networks

Q: How are edge devices connected to the network core?

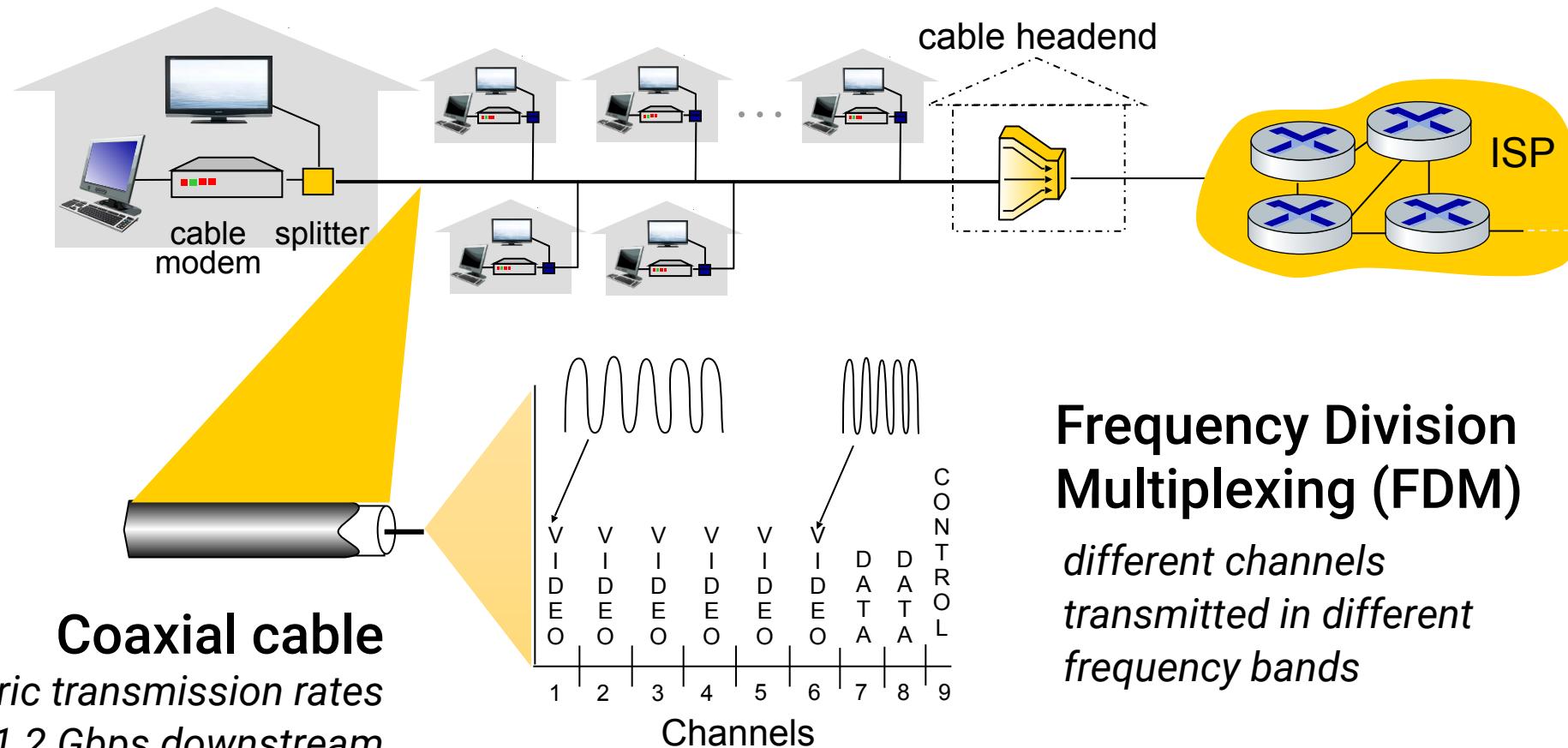
- **cable** based access
- **telephone line** based access
- **wireless or cellular** access

Access network characteristics

- *Transmission rate of access networks*
- *Dedicated or shared access among hosts*

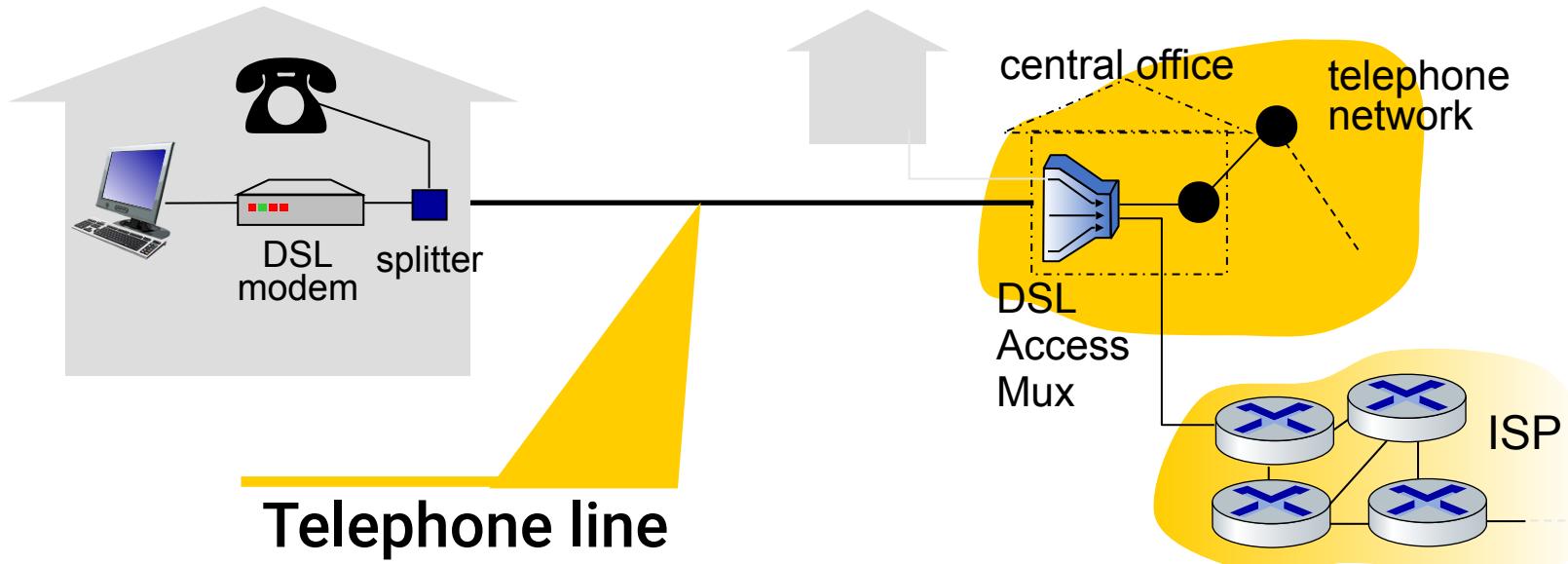


Access technology: cable-based



shared access

Access technology: digital subscriber line (DSL)



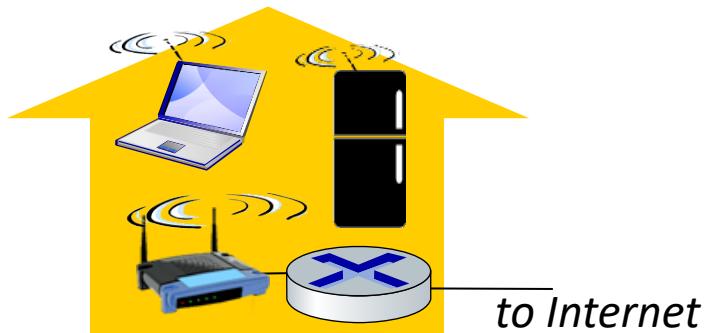
*Transmission rates could be
symmetric or asymmetric
Typical range: 1 - 52 Mbps*

dedicated access

Access technology: wireless

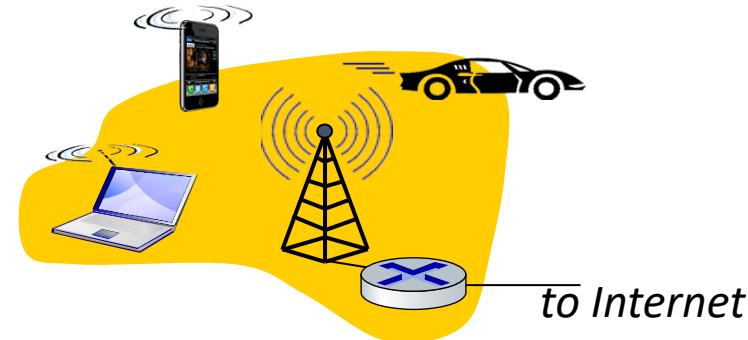
Wireless local area networks (WLANs)

- 802.11b/g/n (WiFi)
- Within or around building (~100 ft)
- 11 - 450 Mbps transmission rate



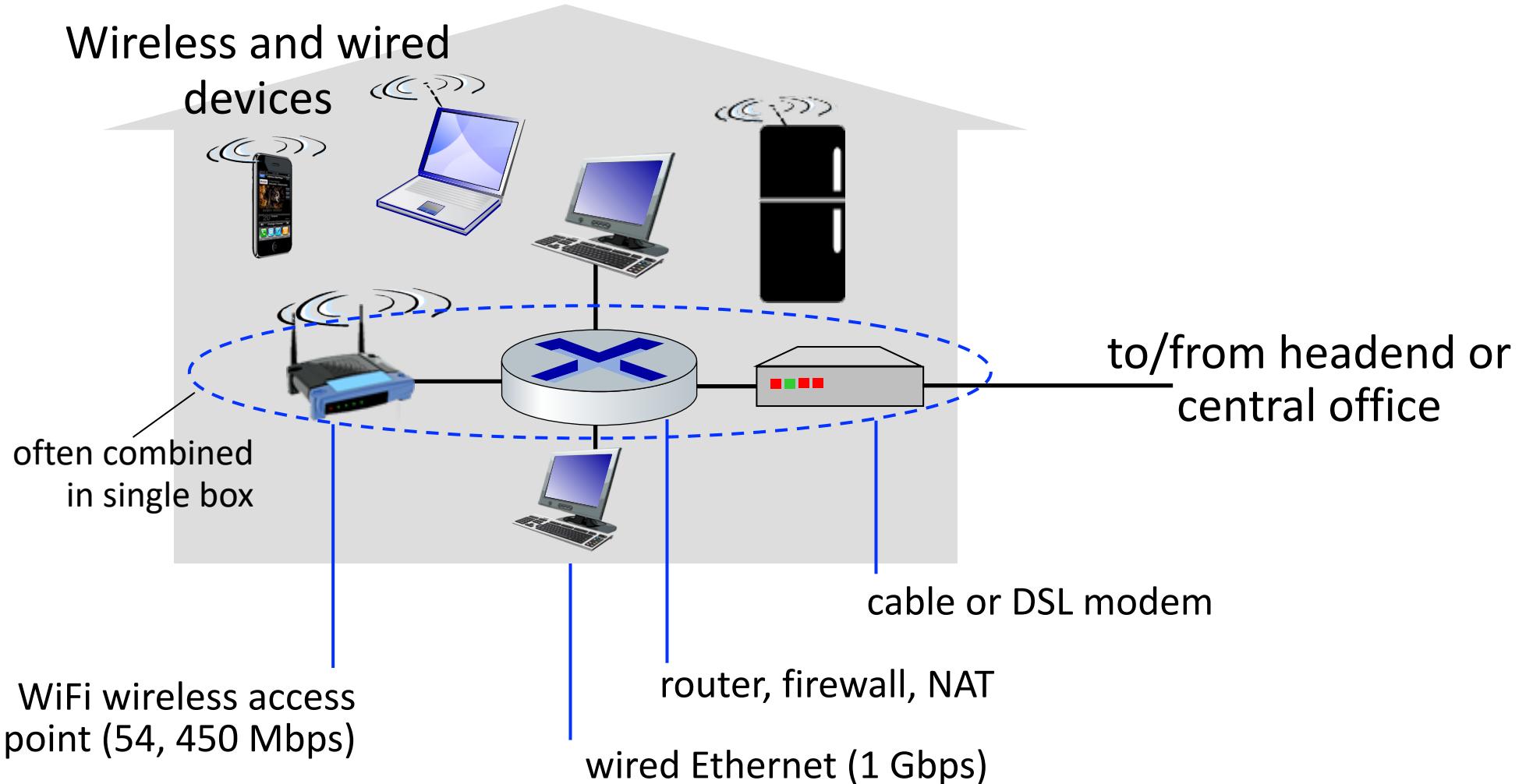
Wide-area cellular networks

- Provided by cellular operators
- 3G and 4G cellular networks
- Range of 10's miles
- 10's Mbps transmission rate

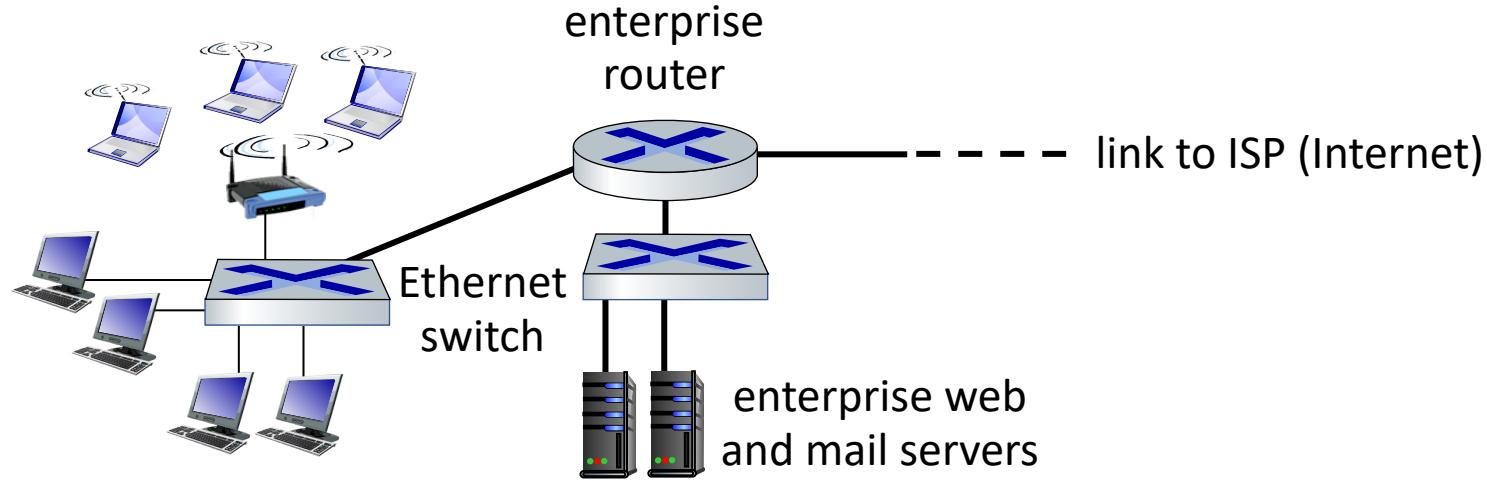


shared access

Access networks: residential



Access networks: enterprise



- *Companies, universities, government agencies etc.,*
- *Mix of wired (e.g., Ethernet) and wireless (e.g., WiFi) technologies*
- *Connected with a hierarchy of switches and routers*

Access networks: data center networks

- Hundreds to thousands of servers connected to each other and to Internet
- High-bandwidth wired links (1 to 100s Gbps)
- Switches and routers organized in intricate **topologies** to maximize bandwidth and minimize latency
 - **Spine - leaf topology**
 - **Three layer topology**

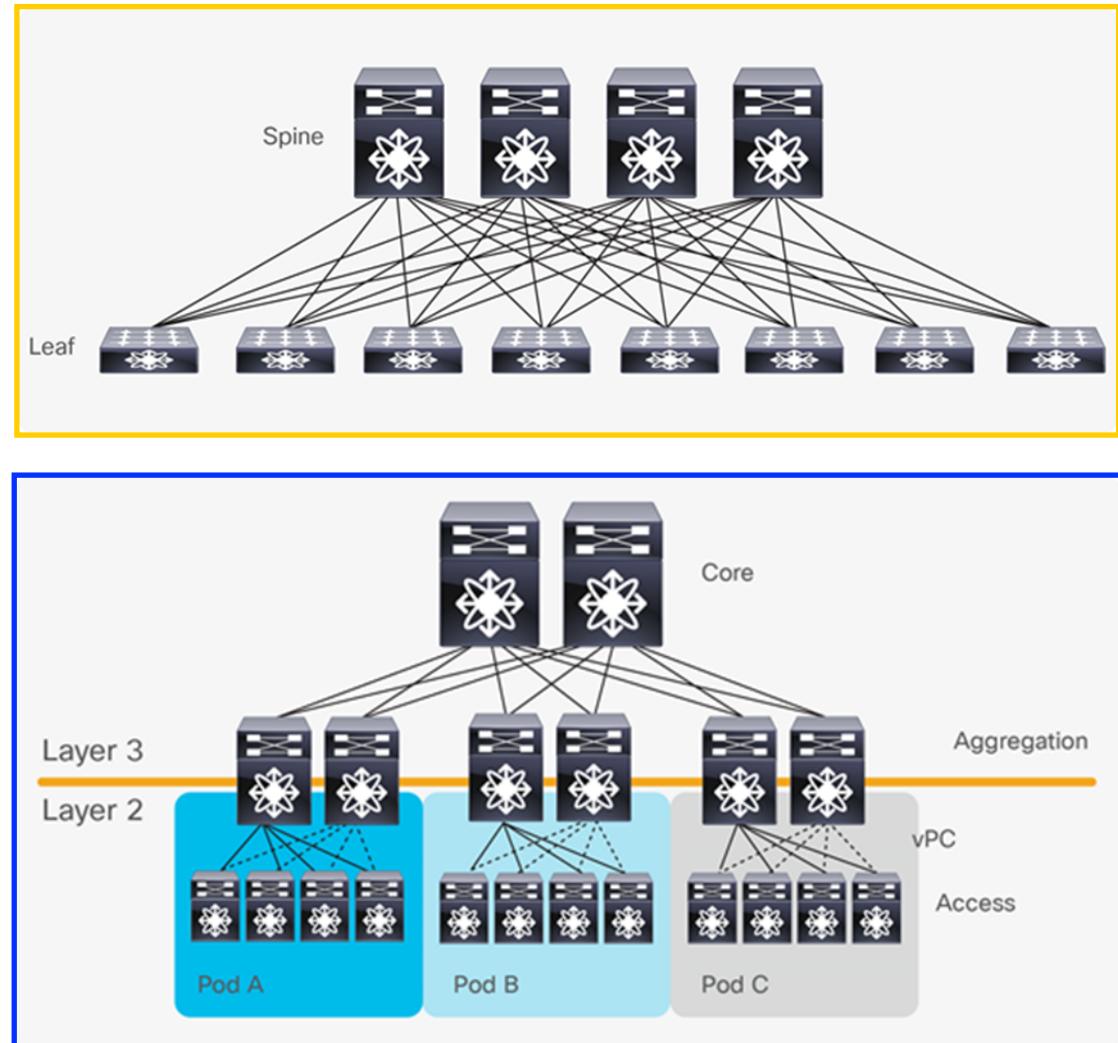
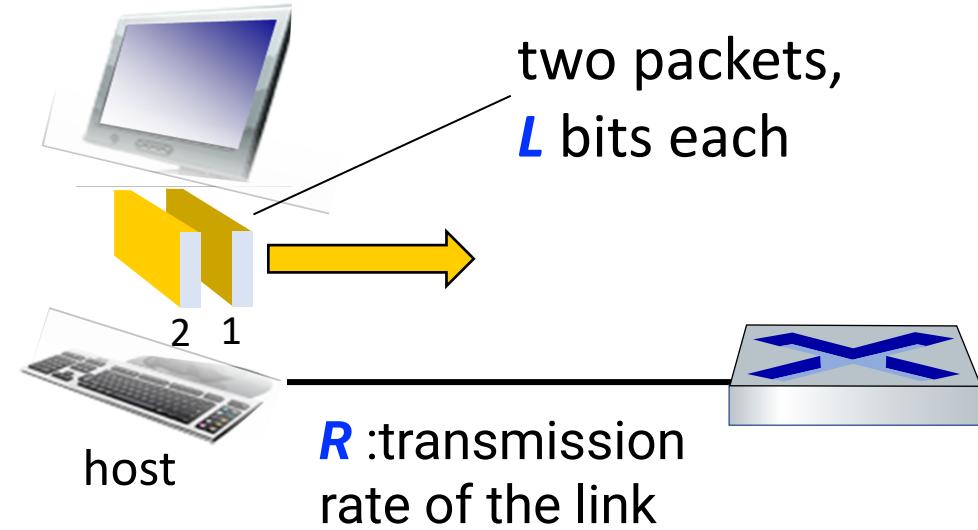


Photo courtesy: Cisco Systems

Host perspective: sending packets of data

Packet sending function

- takes application message
- breaks into smaller chunks, known as **packets**, of **length L** bits
- transmits packet into access network at **transmission rate R** (aka bandwidth)

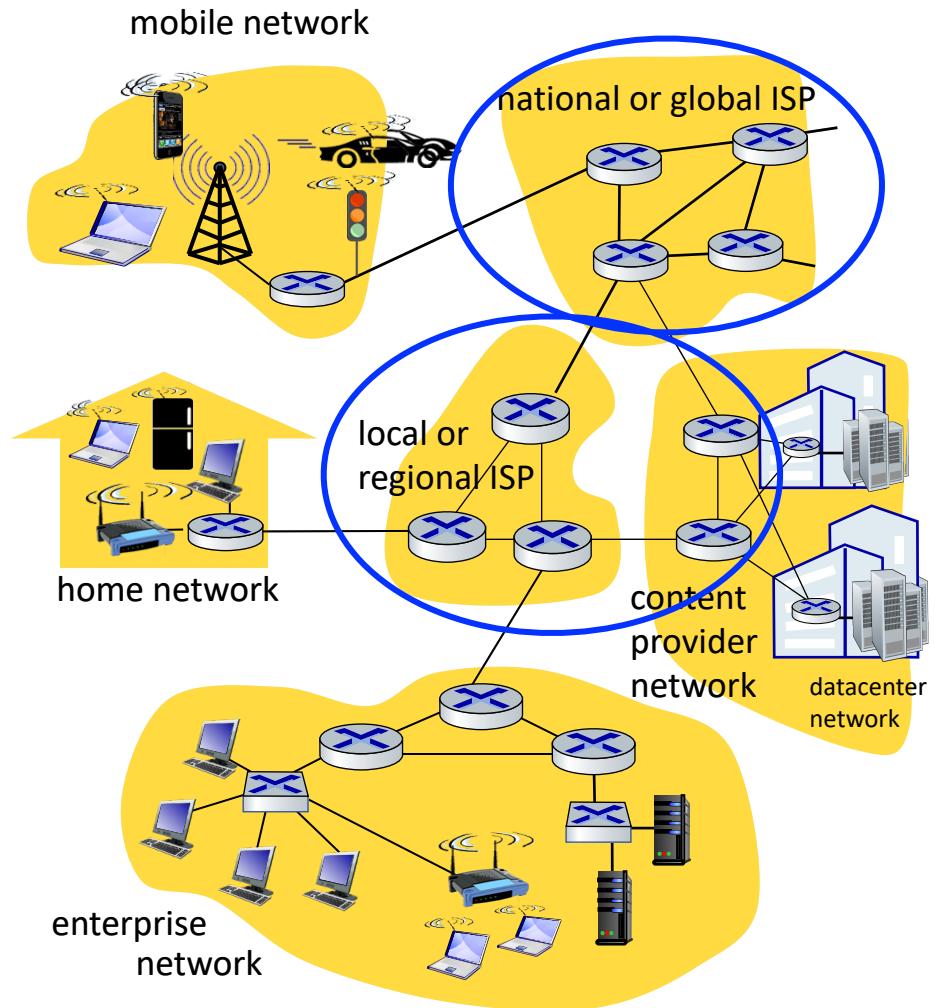


$$\text{packet transmission delay} = \frac{\text{time needed to transmit } L\text{-bit packet into link}}{R \text{ (bits/sec)}}$$

Network core and switching

The network core

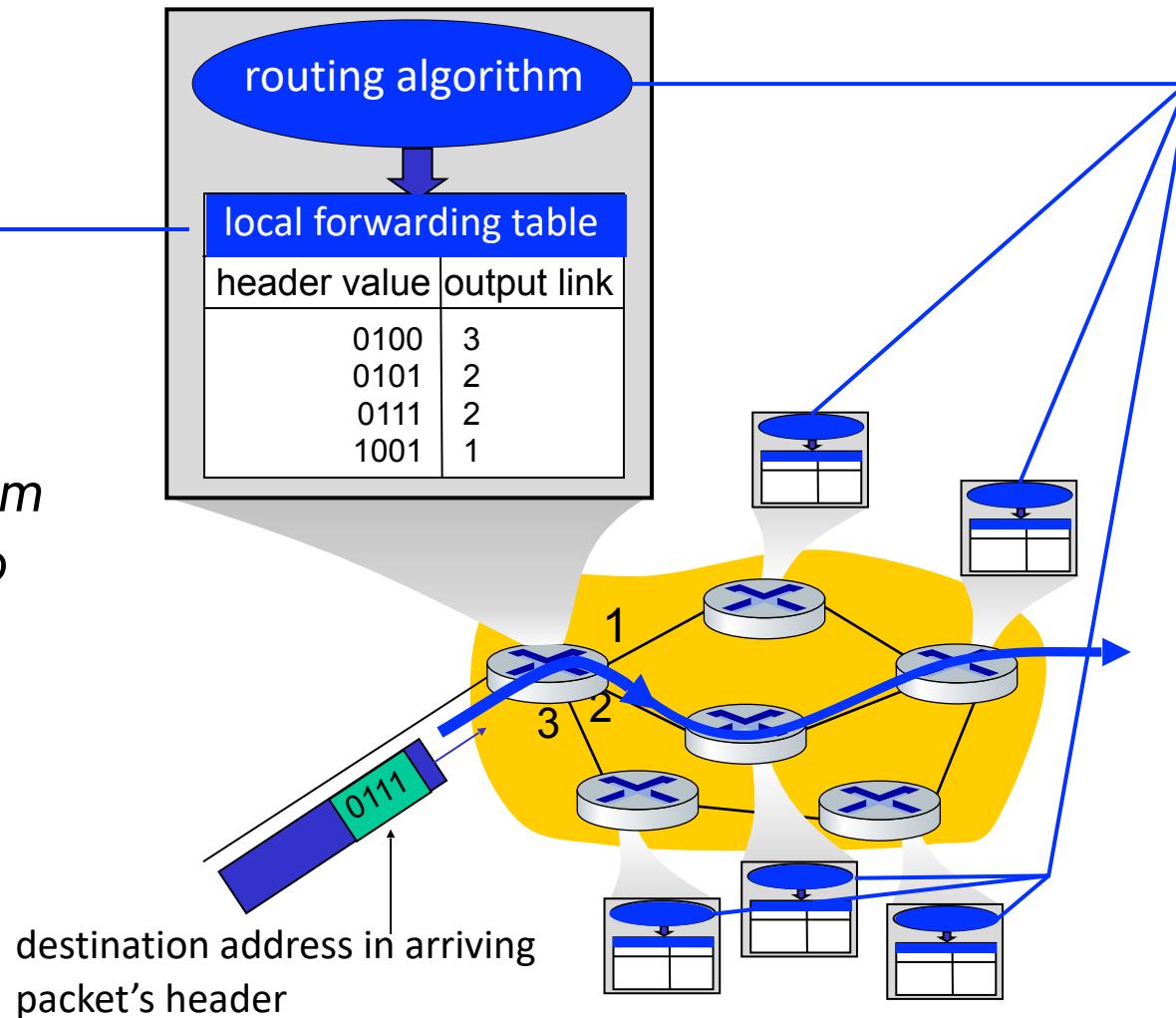
- it is a mesh of interconnected routers
- it performs packet-switching
 - ▶ hosts *break application-layer messages into packets*
 - ▶ *network forwards packets from one router to the next, across links on path from source to destination*



Two key network-core functions

Forwarding

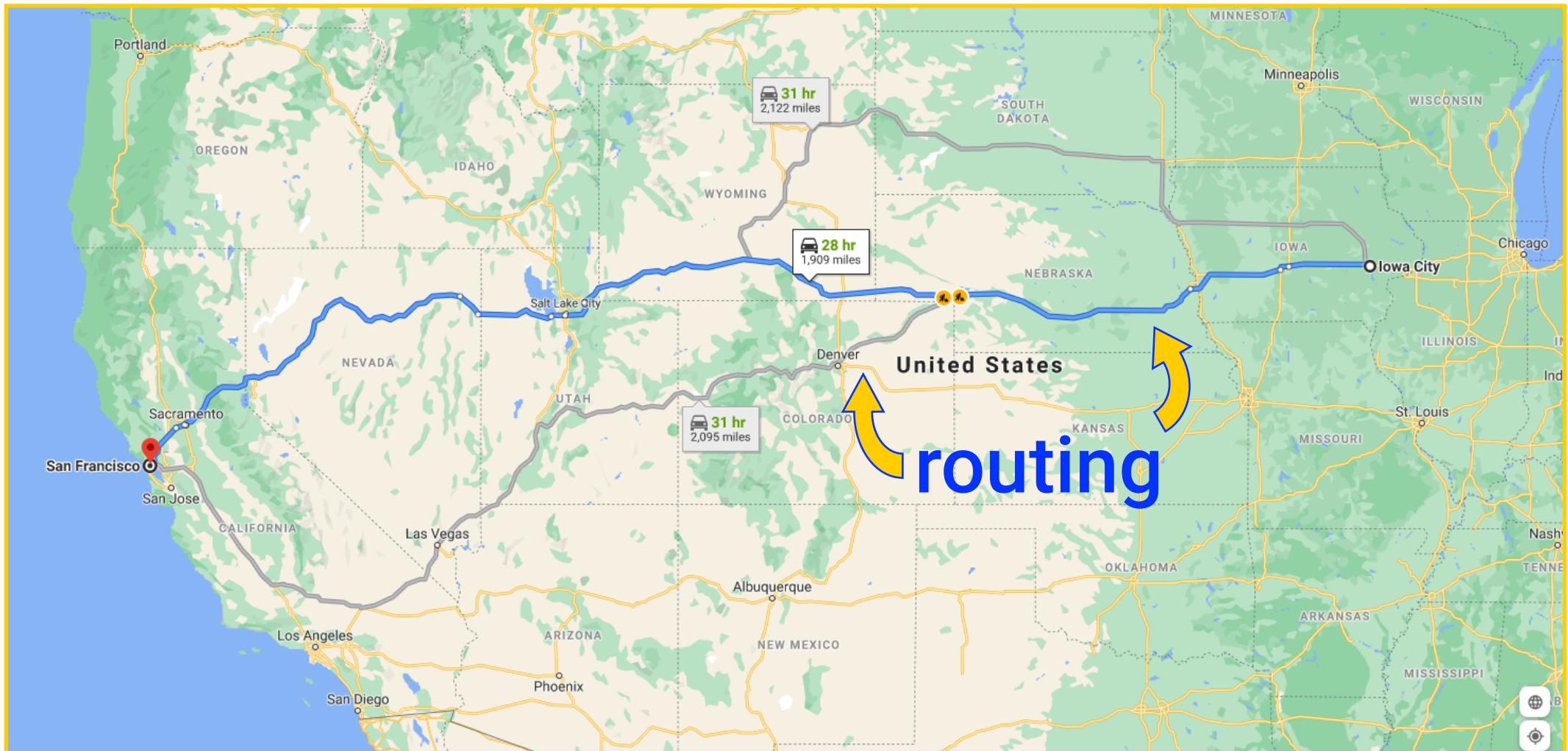
- aka “switching”
- **local action:** move arriving packets from router’s *input link* to appropriate router *output link*



Routing

- **global action:** determine source-destination paths taken by packets
- **routing algorithms**

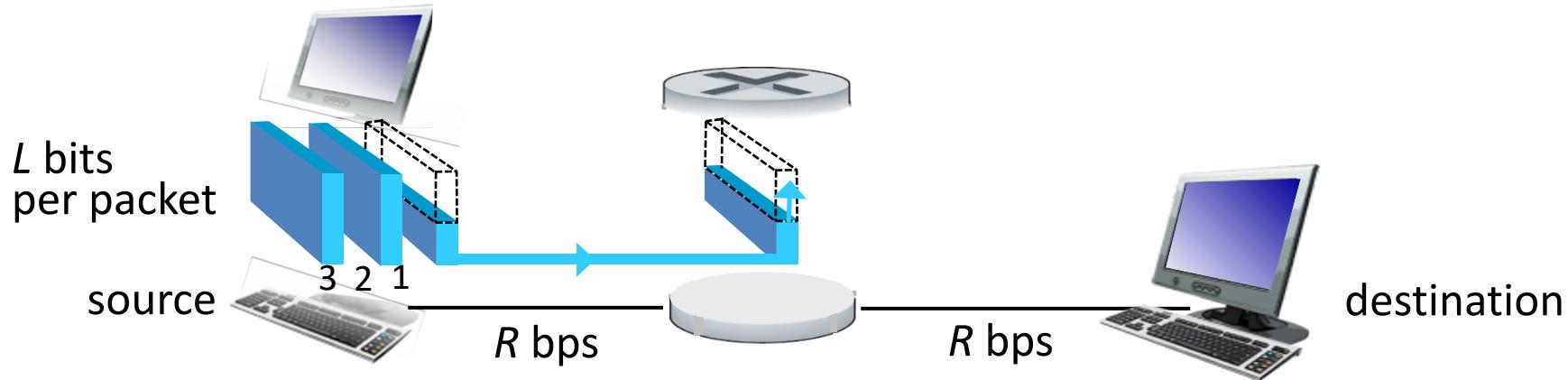
Routing vs. forwarding



Routing vs. forwarding



Packet-switching: store-and-forward



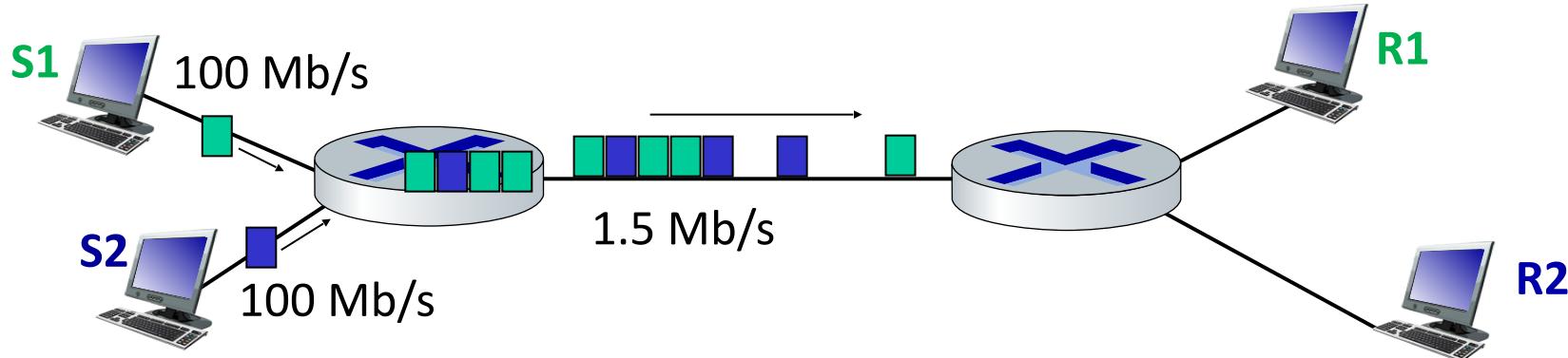
packet transmission delay: takes L/R seconds to transmit (push out) L -bit packet into link at R bps

store and forward: entire packet must arrive at router before it can be transmitted on next link

One-hop example

- $L = 10 \text{ Kbits}$
- $R = 100 \text{ Mbps}$
- one-hop transmission delay = 0.1 msec

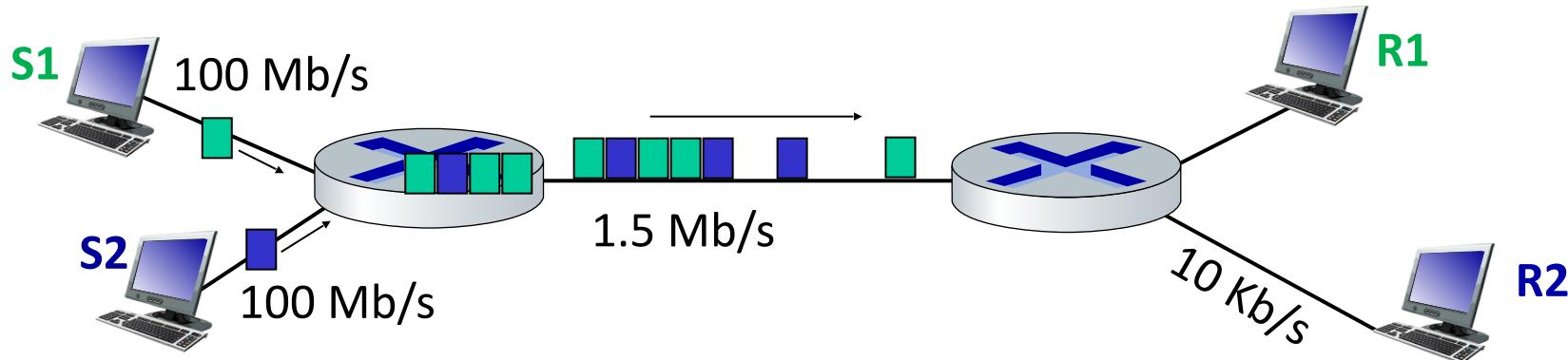
Packet-switching: queueing



Queueing occurs when
work arrives faster than it
can be serviced



Packet-switching: queueing



Packet queuing and loss

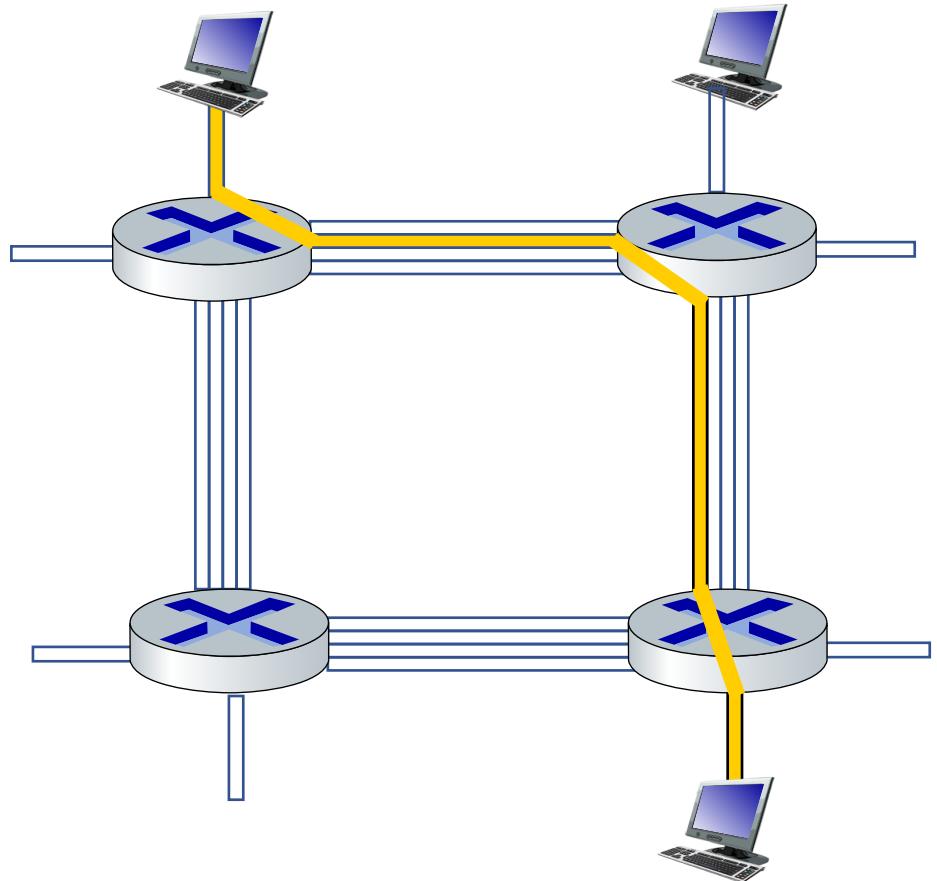
if arrival rate to link exceeds its transmission for some period of time,

- packets will queue, waiting to be transmitted on output link
- packets can be dropped (lost) if memory (buffer) in router fills up

An alternative: circuit switching

End-end resources are allocated and reserved for the “call” between source and destination

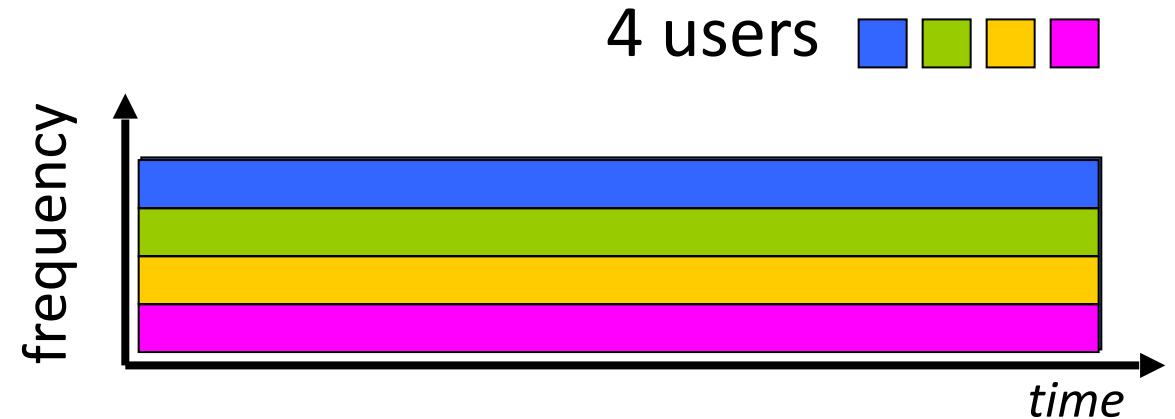
- each link has four circuits
- our call is allocated 2nd circuit in top link and 1st circuit in right link
- dedicated resources: no sharing
- guaranteed performance
- circuit segments remain idle if not used by call
- commonly used in traditional telephone networks



Circuit switching: FDM and TDM

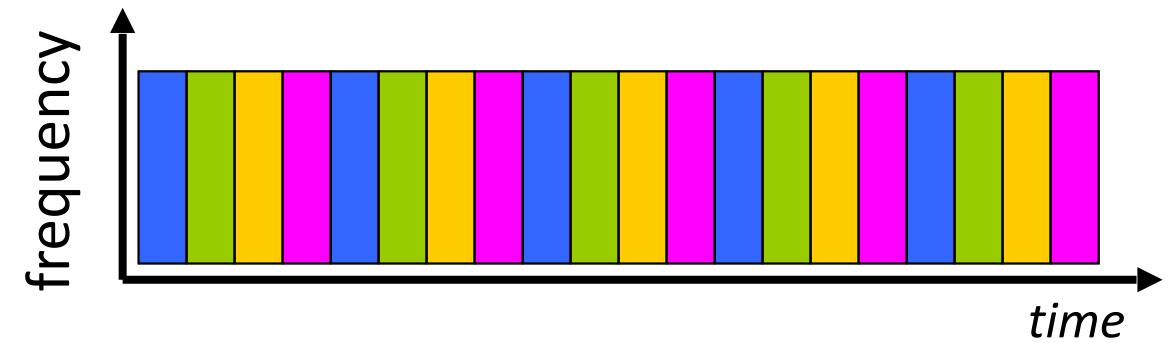
Frequency Division Multiplexing (FDM)

- optical, electromagnetic frequencies divided into frequency bands
- each call allocated its own band, can transmit at max rate of that narrow band



Time Division Multiplexing (TDM)

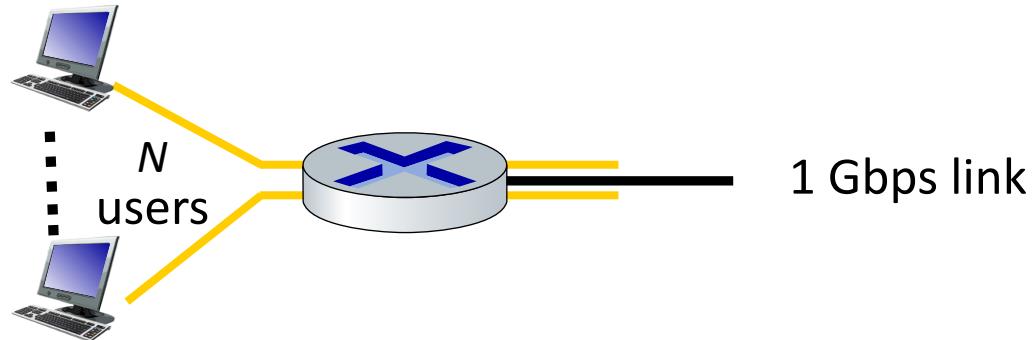
- time divided into slots
- each call allocated periodic slots, can transmit at maximum rate of (wider) frequency band during its time slots



Packet switching vs. circuit switching

N users, each of them is

- active 10% of time
- use 100 Mb/s when “active”



How many users can use this network under circuit-switching and packet switching?

10

Circuit switching

>10

Packet switching

With **35** users
probability that **>10**
are active at the same time is

< 0.0004

Packet switching vs. circuit switching

Is packet switching a “slam dunk winner”?

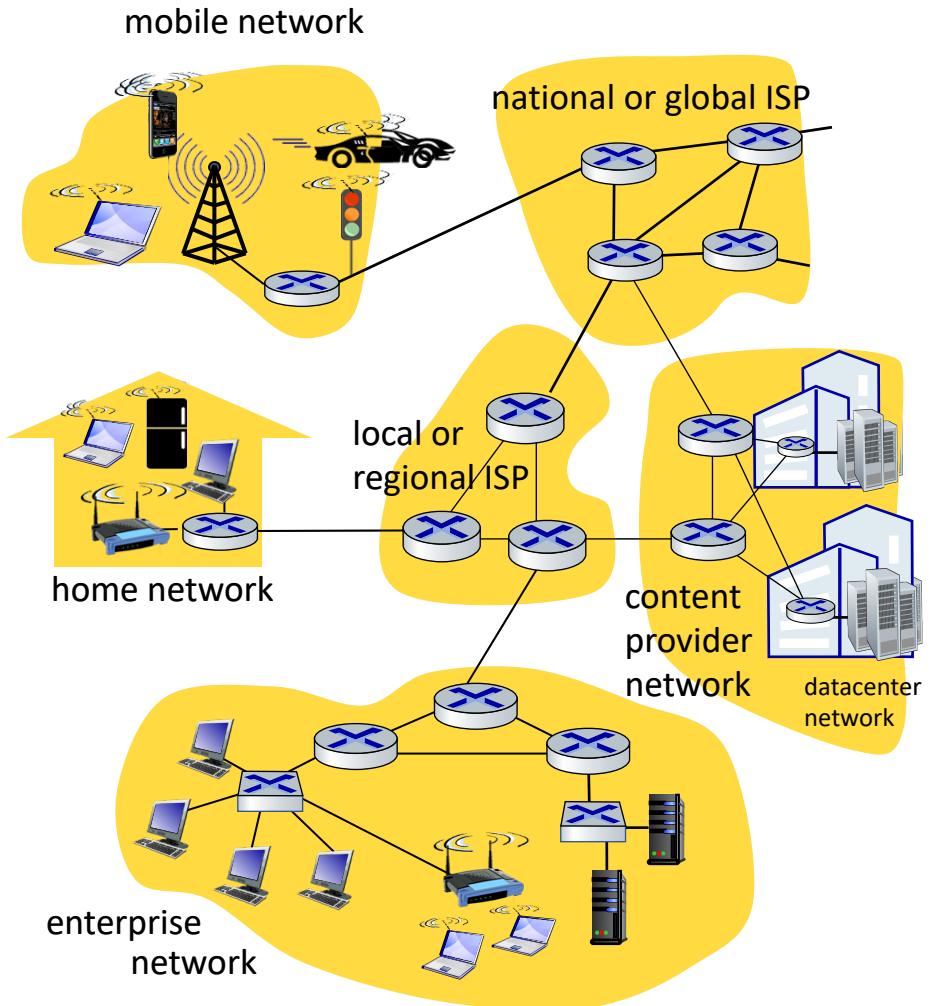
- Great for “bursty” data
 - *efficient resource sharing via statistical multiplexing*
 - *simpler call setup with no a priori reservations*
- Congestion could occur
 - *packets could be delayed or lost due to buffer overflow*
 - *protocols are needed for reliable data transfer, congestion control*

Is there a way to get circuit-like behavior with packet switching?

- Possible but complicated (*we have three dedicated lectures studying TCP*)

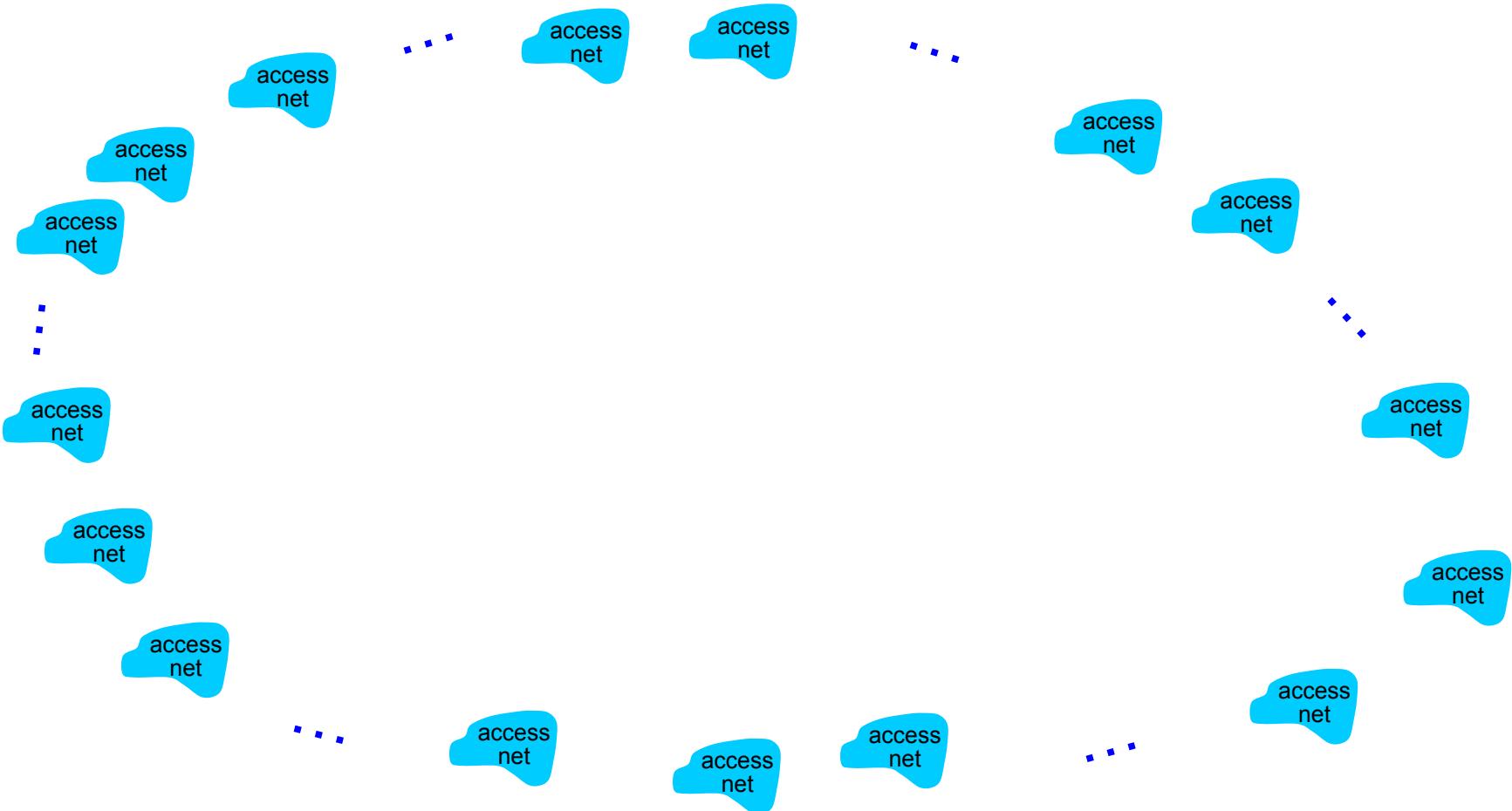
Internet structure: a network of networks

- hosts connect to Internet via access **Internet Service Providers (ISPs)**
- these access ISPs in turn must be interconnected
 - ▶ *to allow that any pair of hosts (anywhere!) to send packets to each other*
- resulting network of networks is complex
 - ▶ *evolution driven by economics, national policies*



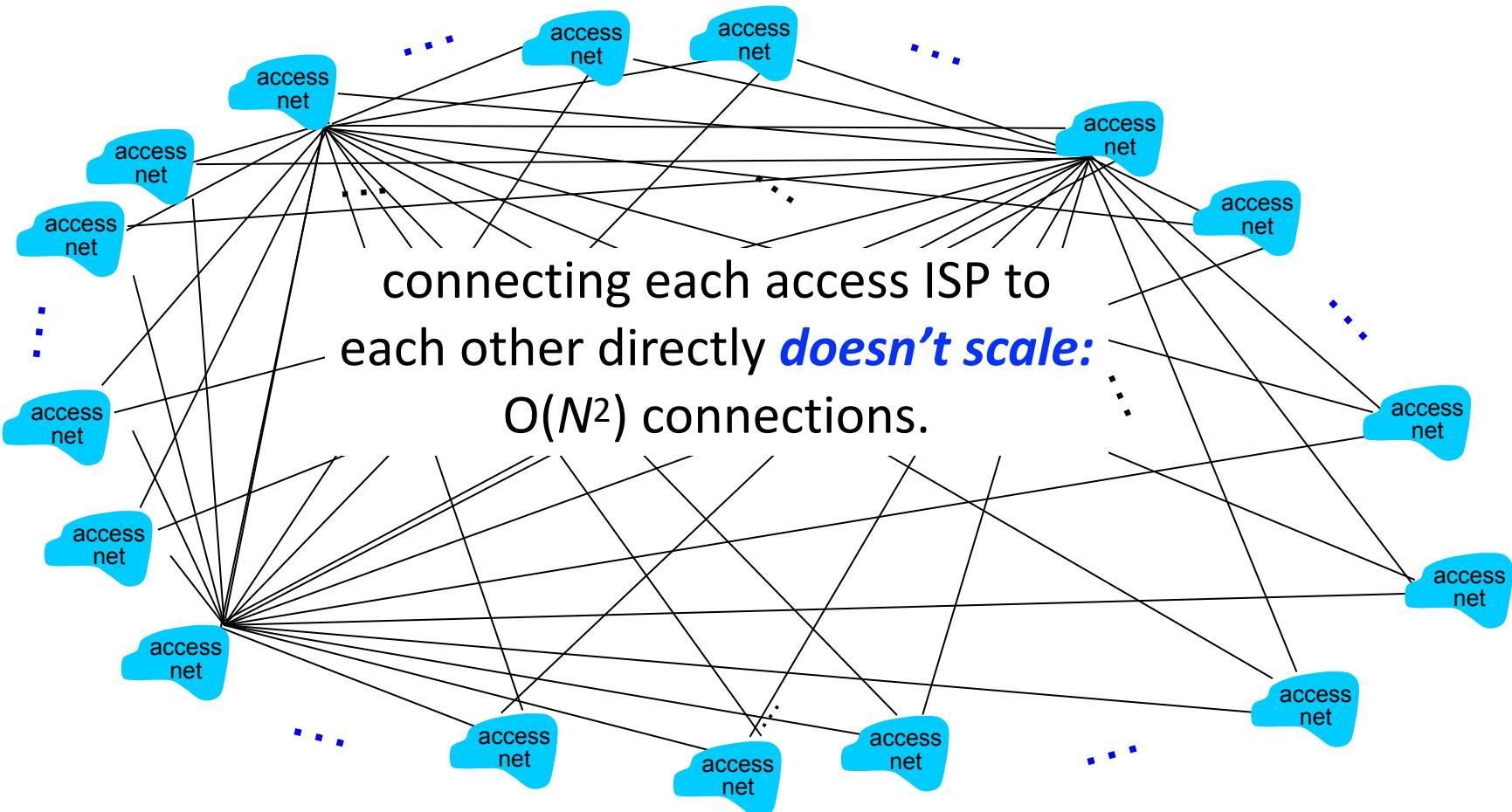
Internet structure: a network of networks

How to connect *millions* of access ISPs together?



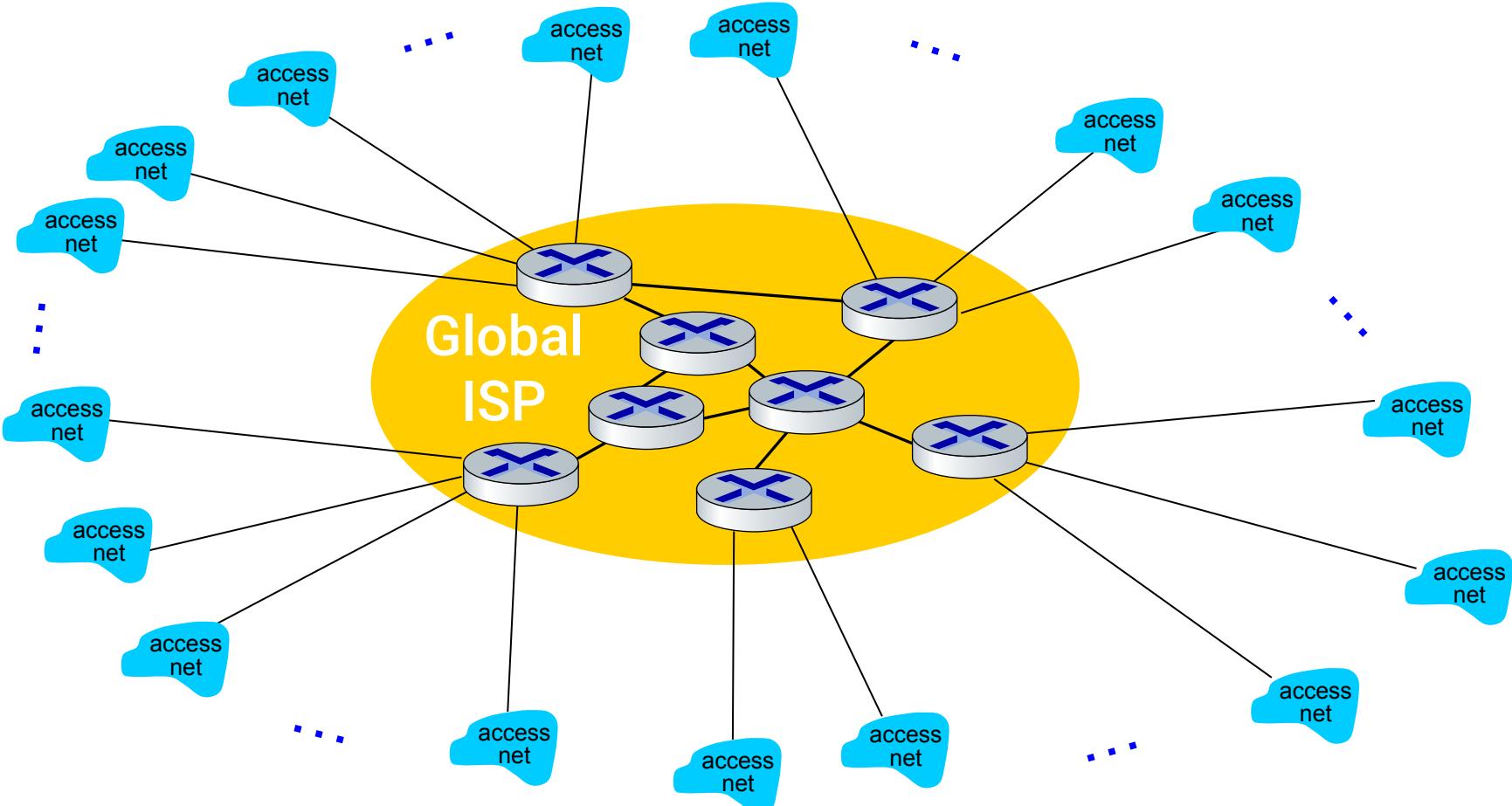
Internet structure: a network of networks

How to connect *millions* of access ISPs together?



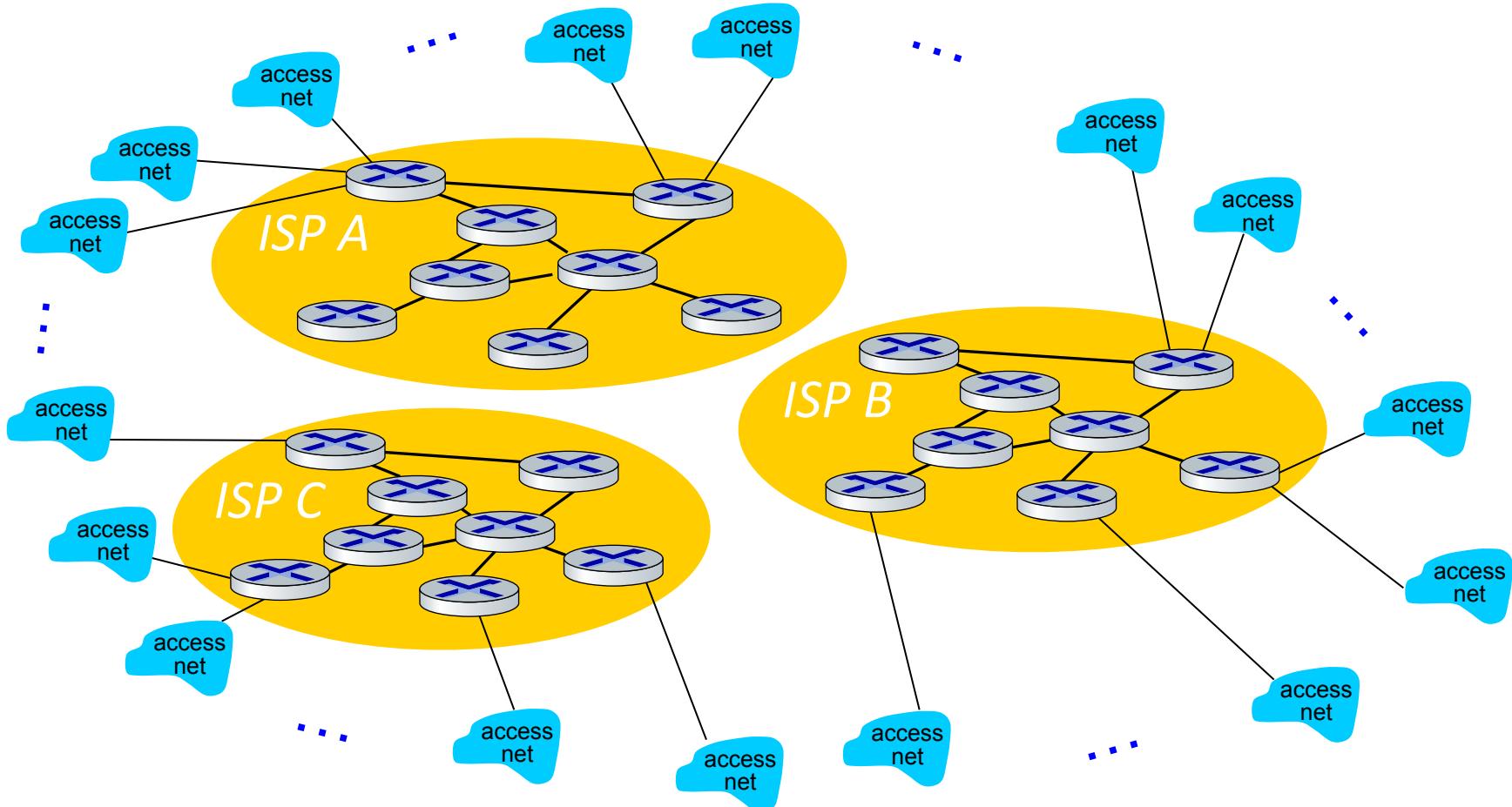
Internet structure: a network of networks

Option: connect each access ISP to one global transit ISP



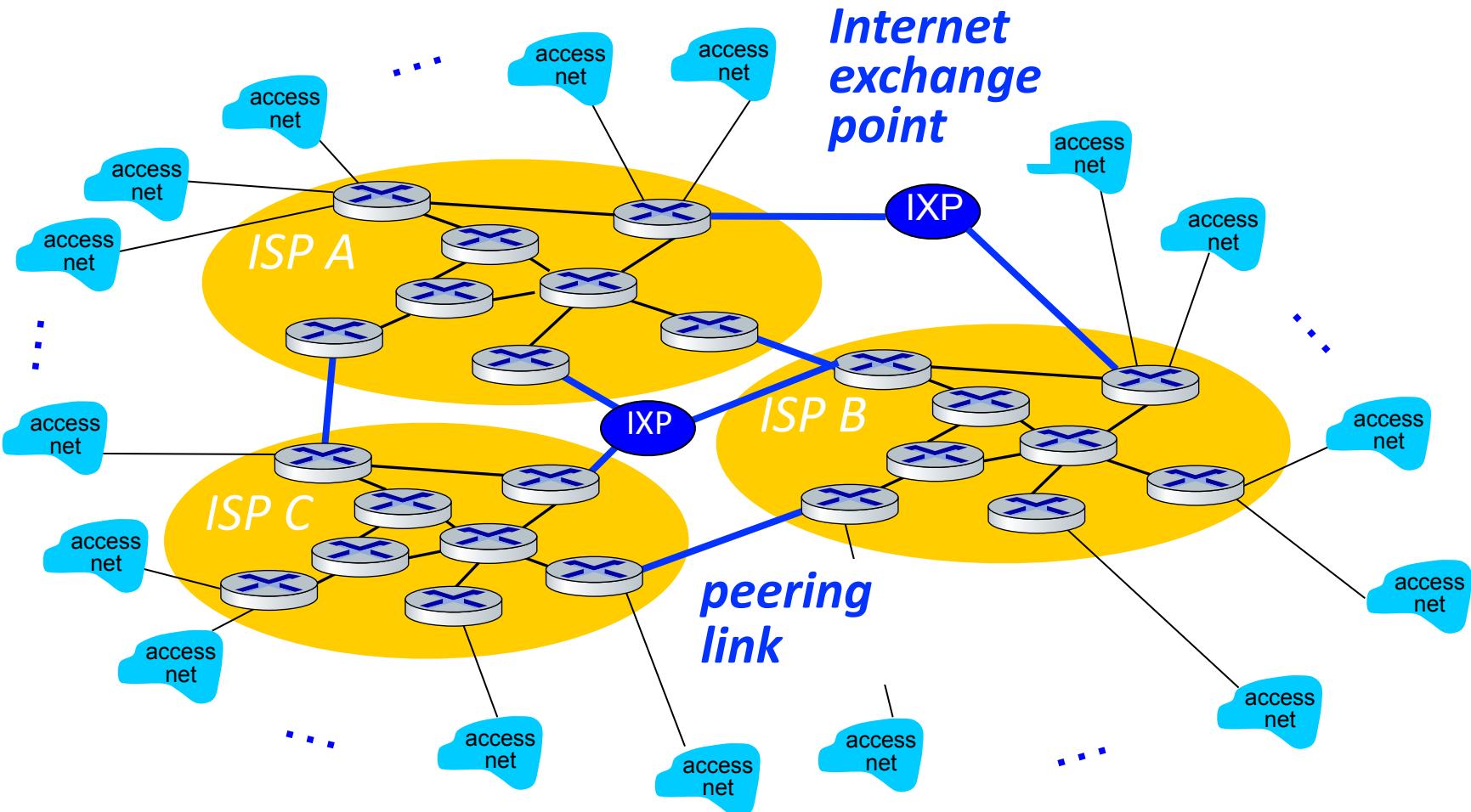
Internet structure: a network of networks

But if one global ISP is viable business, there will be competitors....



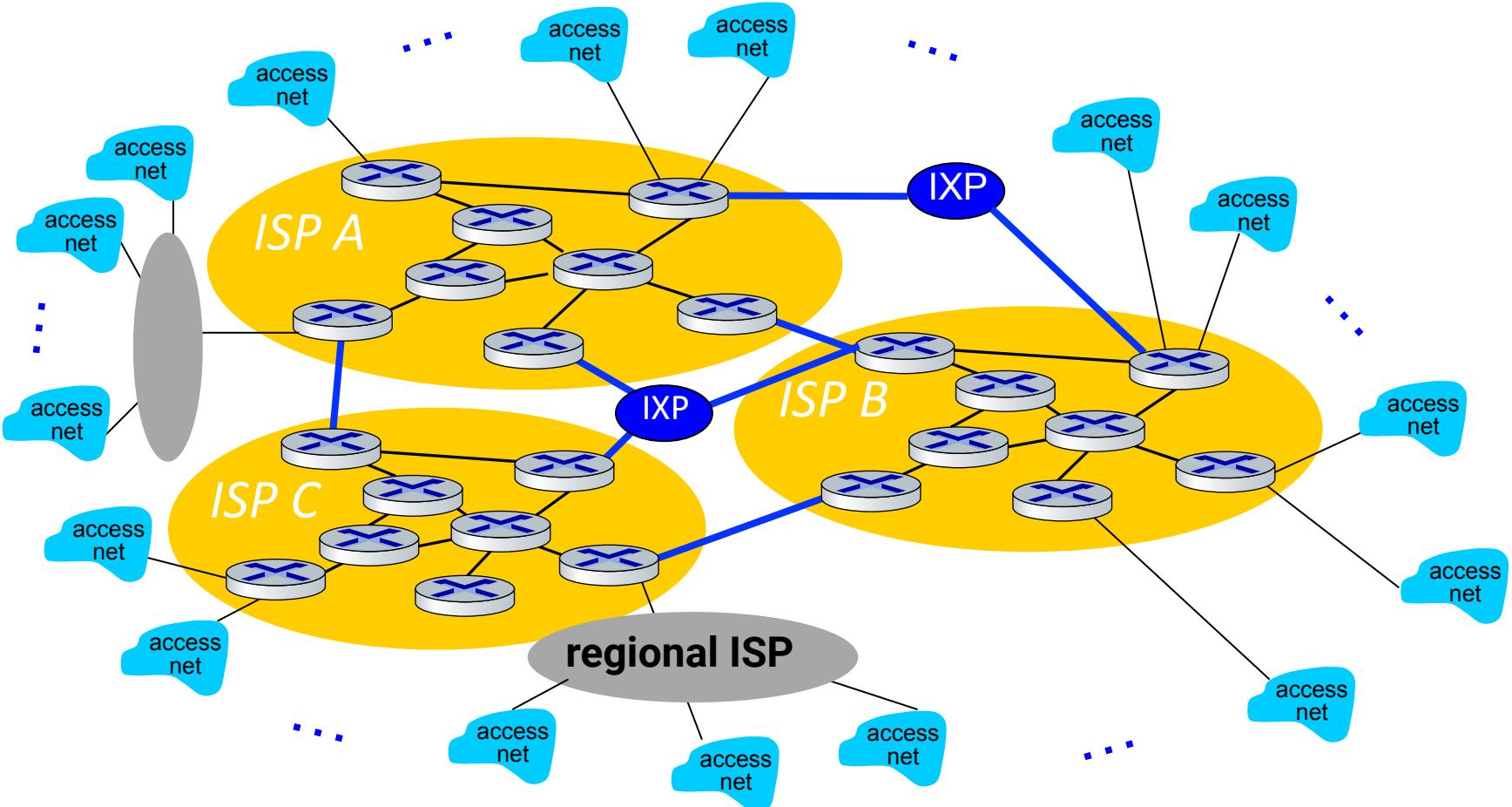
Internet structure: a network of networks

But if one global ISP is viable business, there will be competitors...
who will want to be connected



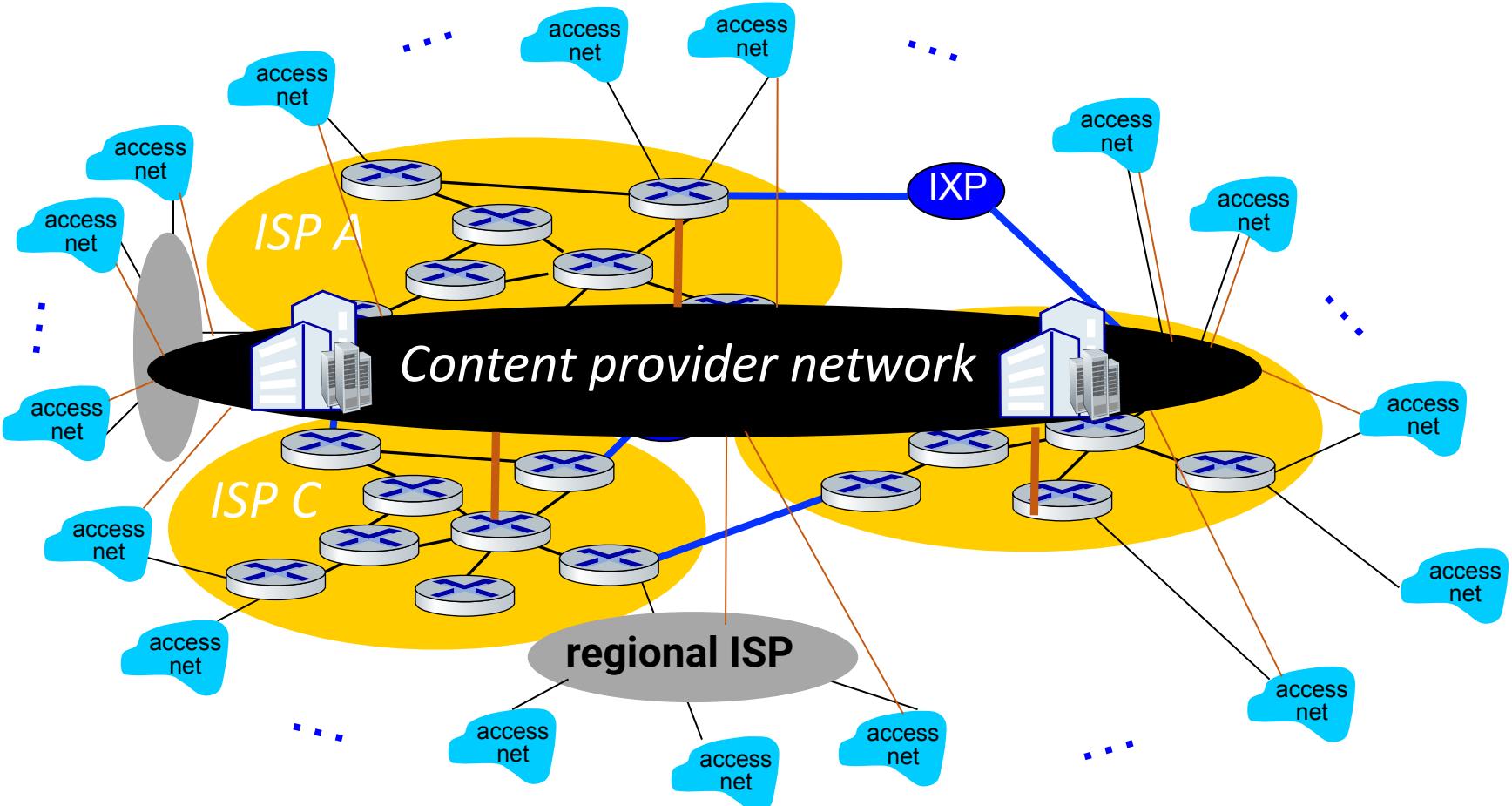
Internet structure: a network of networks

... and regional networks will arise to connect access nets to ISPs

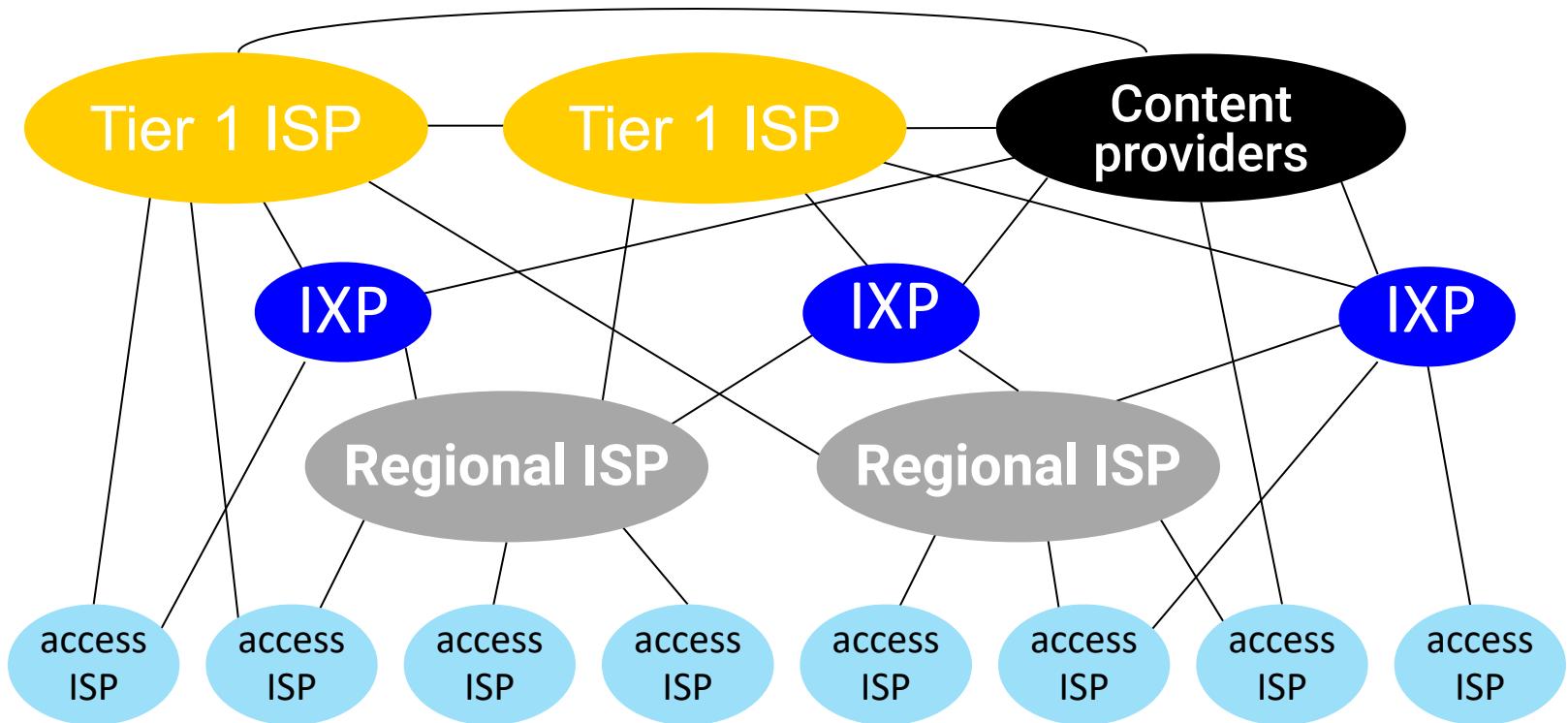


Internet structure: a network of networks

... and content providers (e.g., Google, Akamai, Amazon) may run their own network



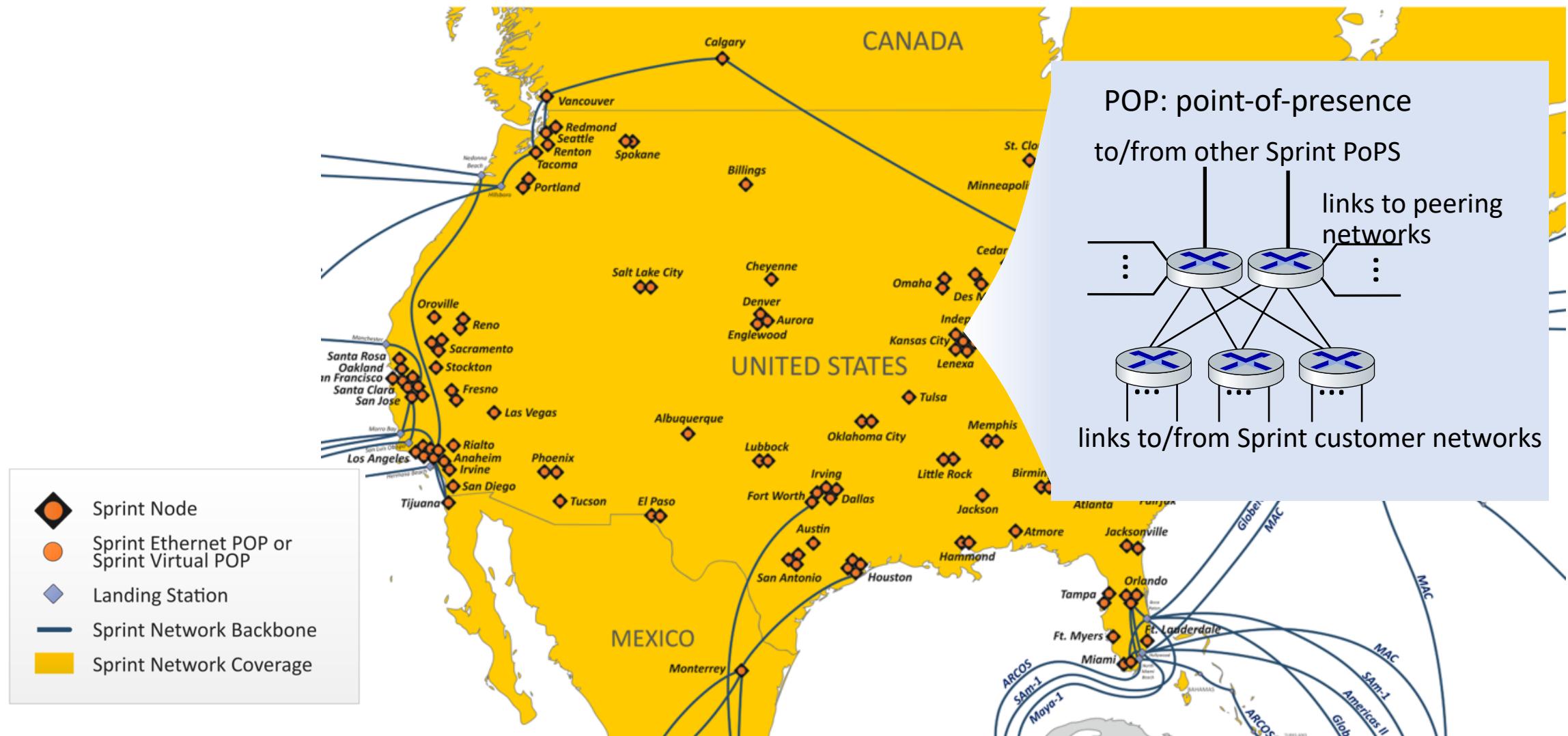
Internet structure: a network of networks



At the center: a small number of well-connected large networks

- “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T) for national & international coverage
- content provider networks (e.g., Google, Akamai, Amazon): private networks that connect data centers to Internet, often bypassing tier-1 and regional ISPs

Tier-1 ISP network: Sprint (circa 2019)



Content provider network: Google (circa 2017)

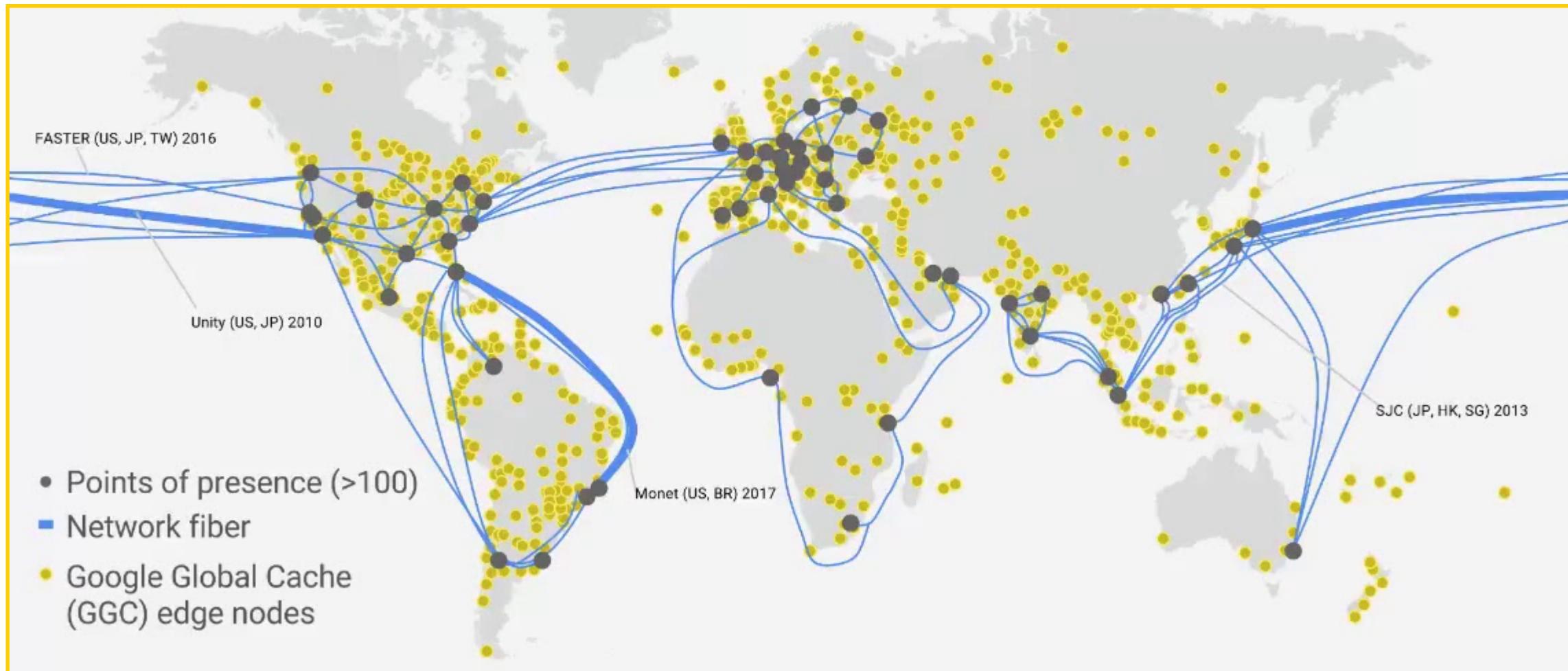


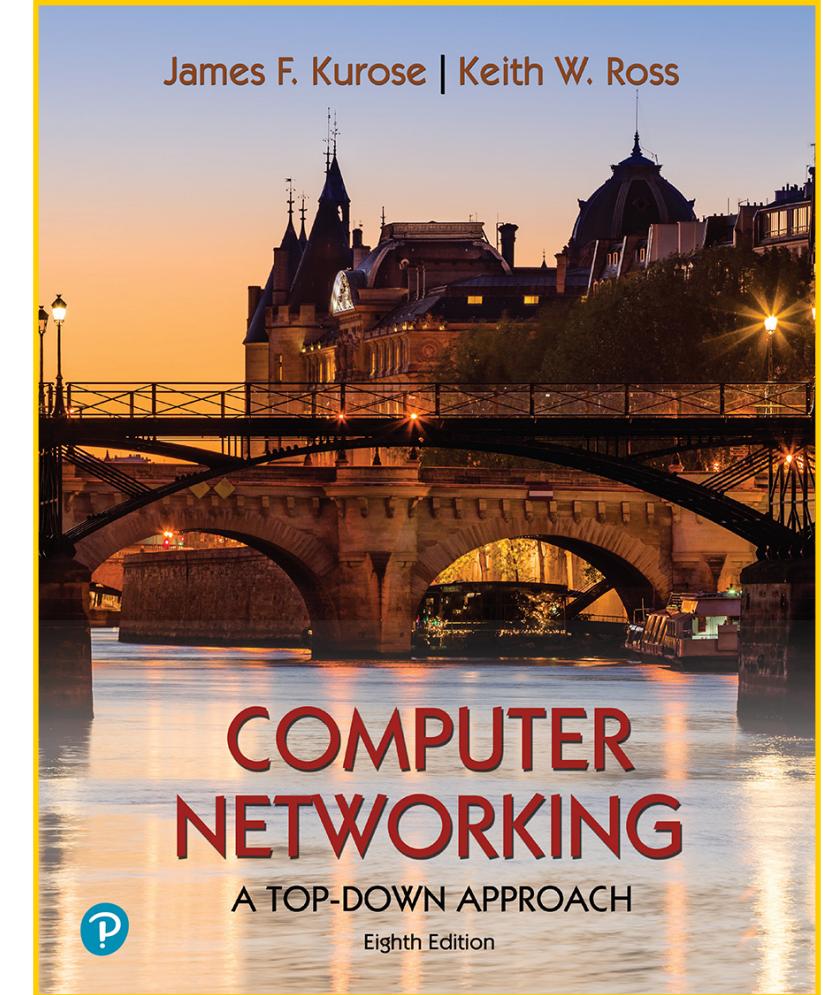
Image courtesy: Google cloud

James F. Kurose | Keith W. Ross

Next Lecture

*Continuing our in-depth exploration
into the structure and functioning
of the Internet*

- *Network performance*
- *Protocol architecture*



Chapter 1.4 - 1.5

Spot Quiz (ICON)