

Learning Go

GO



<http://golang.org>

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Thanks to:
Go Authors
Google
Go Nuts mailing list



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Learning as we Go.

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1

Introduction

I am interested in this and hope to do something.

On adding complex numbers to Go
KEN THOMPSON

This is an introduction to the Go language from Google. Its aim is to provide a guide to this new and innovative language. What is Go? From the website [9]:

The Go programming language is an open source project to make programmers more productive. Go is expressive, concise, clean, and efficient. Its concurrency mechanisms make it easy to write programs that get the most out of multicore and networked machines, while its novel type system enables flexible and modular program construction. Go compiles quickly to machine code yet has the convenience of garbage collection and the power of run-time reflection. It's a fast, statically typed, compiled language that feels like a dynamically typed, interpreted language.

The intended audience of this book is people who are familiar with programming and know multiple programming languages, be it C[19], C++[29], Perl [21], Java [20], Erlang[18], Scala[1] or Haskell[10]. This is *not* a book which teaches you how to program, this is a book that just teaches you how to use Go.

As with learning new things, probably the best way to do this is to discover it for yourself by creating your own programs. Therefore includes each chapter a number of exercises (and answers) to acquaint you with the language. An exercise is numbered as **Q n** , where n is a number. After the exercise number another number in parentheses displays the difficulty of this particular assignment. This difficulty ranges from 0 to 9, where 0 is easy and 9 is difficult. Then a short name is given, for easier reference. For example

Q1. (1) A map function ...

introduces a question numbered **Q1** of a level 1 difficulty, concerning a `map()`-function. The answers are included after the exercises on a new page, except for those pesky unresolved research problems. The numbering and setup of the answers is identical to the exercises, except that an answer starts with **A n** , where the number n corresponds with the number of the exercise.

Go is a young language, where features are still being added or even *removed*. It may be possible that some text is outdated when you read it. Some exercise answers may become incorrect as Go continues to evolve. We will do our best to keep this document up to date with respect to the latest Go release. An effort has been made to create "future proof" code examples.

The following convention is used throughout this book:

- Code is displayed in *DejaVu Mono*;
- Keywords are displayed in **DejaVu Mono Bold**;
- Comments are displayed in *DejaVu Mono Italic*;

- Extra remarks in the code ← *Are displayed like this;*
- Longer remarks get a number – ❶ – with the explanation following;
- Line numbers are printed on the right side;
- Shell examples use a % as prompt;
- An emphasized paragraph is indented and has a vertical bar on the left.

Official documentation

There already is a substantial amount of documentation written about Go. The Go Tutorial [8], and the Effective Go document [3]. The website <http://golang.org/doc/> is a very good starting point for reading up on Go^a. Reading these documents is certainly not required, but is recommended.

Go comes with its own documentation in the form of a Go program called `godoc`. You can use it yourself to look in the online documentation. For instance, suppose we want to know more about the package `hash`. We would then give the command `godoc hash`. How to create your own package documentation is explained in chapter 4.

Getting Go

There are currently (2011) no packages for Go in any Linux distribution. The route to install Go is thus slightly longer than it could be. When Go stabilizes this situation will change. For now, you need to retrieve the code from the mercurial archive and compile Go yourself. For other Unix like systems the procedure is the same.

- First install Mercurial (to get the `hg` command). In Ubuntu/Debian/Fedora you must install the `mercurial` package;
- For building Go you need the packages: `bison`, `gcc`, `libc6-dev`, `ed`, `gawk` and `make`;
- Set the environment variable `GOROOT` to the root of your Go install:
% `export GOROOT=~/go`
- Then retrieve the Go source code:
% `hg clone -r release https://go.googlecode.com/hg/ $GOROOT`
- Set your `PATH` to so that the Shell can find the Go binaries:
% `export PATH=$GOROOT/bin:$PATH`
- Compile Go
% `cd $GOROOT/src`
% `./all.bash`

If all goes well, you should see the following:

^a<http://golang.org/doc/> itself is served by a Go program called `godoc`.

```
Installed Go for linux/amd64 in /home/gobook/go.  
Installed commands in /home/gobook/go/bin.  
The compiler is 6g.
```

You now have Go installed on your system and you can start playing.

Keeping up to date

New releases are announced on the Go Nuts mailing list [17]. To update an existing tree to the latest release, you can run:

```
% cd $GOROOT  
% hg pull  
% hg update release  
% cd src  
% ./all.bash
```

To see what you are running right now:

```
% cd $GOROOT  
% hg identify  
79997f0e5823 release/release.2010-10-20
```

That would be release.2010-10-20.

Origins

Go has its origins in Inferno [11] (which in turn was based upon Plan 9 [14]). Inferno included a language called Limbo [13]. Quoting from the Limbo paper:

Limbo is a programming language intended for applications running distributed systems on small computers. It supports modular programming, strong type checking at compile- and run-time, inter process communication over typed channels, automatic garbage collection, and simple abstract data types. It is designed for safe execution even on small machines without hardware memory protection.

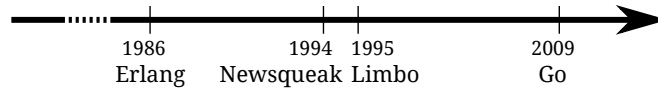
TODO
Ale language on Plan 9. Also presented in 1995.

One feature of Limbo that is included in Go is the support for cross compiling. Another feature Go inherited from Limbo is channels (see chapter 7). Again from the Limbo documentation.

[A channel] is a communication mechanism capable of sending and receiving objects of the specified type to another agent in the system. Channels may be used to communicate between local processes; using library procedures, they may be connected to named destinations. In either case send and receive operations may be directed to them.

The channels in Go are easier to use than those in Limbo. If we dig even deeper in the history of Go we also find references to "Newsqueak" [28], which pioneered the use of channel communication in a C-like language. Channel communication isn't unique to these languages, a big non-C-like language which also uses them is Erlang [18].

Figure 1.1. Chronology of Go



The whole of idea of using channels to communicate with other processes is called Communicating Sequential Processes (CSP) and was conceived by C. A. R. Hoare [23], who incidentally is the same man that invented QuickSort [24].

Go is the first C-like language that is widely available, runs on many different platforms and makes concurrency easy (or easier).

Exercises

Q1. (1) Documentation

1. Go's documentation can be read with the `godoc` program, which is included the Go distribution.
`godoc hash` gives information about the *hash* package. Reading the documentation on *container* gives the following result:

SUBDIRECTORIES

heap
list
ring
vector

With which `godoc` command can you read the documentation of *vector* contained in *container*?

Answers

A1. (1) Documentation

1. The package *vector* is in a *subdirectory* of *container*, so you will only need `godoc container/vector`.

Specific functions inside the "Go manual" can also be accessed. For instance the function `Printf` is described in *fmt*, but to only view the documentation concerning this function use: `godoc fmt Printf`.

You can even display the source code with: `godoc -src fmt Printf`.

2

Basics

In Go, the code does exactly what it says on the page.

Go Nuts mailing list
ANDREW GERRAND

There are a few things that make Go different from other languages.

Clean and Simple

Go strives to keep things small and beautiful, you should be able to do a lot in only a few lines of code;

Concurrent

Go makes it easy to "fire off" functions to be run as *very* lightweight threads. These threads are called goroutines ^a in Go;

Channels

Communication with these goroutines is done via channels [33][23];

Fast

Compilation is fast and execution is fast. The aim is to be as fast as C. Compilation time is measured in seconds;

Safe

Go has garbage collection, no more `free()` in Go, the language takes care of this;

Standard format

A Go program can be formatted in (almost) any way the programmers want, *but* an official format exist. The rule is very simple: The output of the filter `gofmt` is *the official* endorsed format.

Postfix types

Types are given *after* the variable name, thus `var a int`, instead of `int a`; as one would in C;

UTF-8

UTF-8 is everywhere, in strings *and* in the program code. Finally you can use $\Phi + 1$ in your source code;

Open Source

The Go license is completely open source, see the file `LICENSE` in the Go source code distribution;

Fun

Programming with Go should be fun!

Erlang [18] also shares some of the features of Go. Notable differences between Erlang and Go is that Erlang borders on being a functional language, where Go is an imperative one. And Erlang runs in a virtual machine, while Go is compiled. Go also has a much more Unix-like feeling to it.

^aYes, that sounds a lot like coroutines, but goroutines are slightly different as we will see in chapter 7.

Hello World

In the Go tutorial, Go is presented to the world in the typical manner: letting it print "Hello World" (Ken Thompson and Dennis Ritchie started this when they presented the C language in the nineteen seventies). We don't think we can do better, so here it is, "Hello World" in Go.

Listing 2.1. Hello world

```

package main ❶                                     1

import "fmt" // Implements formatted I/O. ❶          3

/* Print something */ ❷                             5
func main() { ❸                                     6
  ❹ fmt.Printf("Hello, world; or καλημέρα κόσμε; or こんにちは世界\n") 8
}                                                    9

```

Lets look at the program line by line.

- ❶ This first line is just required. All Go files start with **package** <something>, **package** *main* is required for a standalone executable;
- ❶ This says we need *"fmt"* in addition to *main*. A package other than *main* is commonly called a library, a familiar concept of many programming languages (see chapter 4). The line ends with a comment which is started with *//*;
- ❷ This is also a comment, but this one is enclosed in */** and **/*;
- ❸ Just as **package** *main* was required to be first, **import** may come next. In Go, **package** is always first, then **import**, then everything else. When your Go program is executed, the first function called will be *main.main()*, which mimics the behavior from C. Here we declare that function;
- ❹ On line 8 we call a function from the package *fmt* to print a string to the screen. The string is enclosed with *"* and may contain non-ASCII characters. Here we use Greek and Japanese.

Compiling and running code

The Go compiler is named <number>g, where the number is 6 for 64 bit Intel and 8 for 32 bit Intel. The linker has a similar naming scheme: <number>l. In this book we will use 6g and 6l for all the compiles. To compile the code from above, we use:

```
% 6g helloworld.go
```

And then we link it with 6l:

```
% 6l helloworld.6
```

And then we run it:

```
% ./6.out          ← The default name for a (64 bit) Go executable (8.out on 32 bits)
Hello, world; or καλημέρα κόσμε; or こんにちは世界
```

Using a Makefile

Another, less laborious (once setup), way to build a Go program, is to use a Makefile. The following one can be used to build helloworld:

Listing 2.2. Makefile for a program

```
include $(GOROOT)/src/Make.inc          1

TARG=helloworld                        3
GOFILES=\                              4
    helloworld.go\                      5

include $(GOROOT)/src/Make.cmd          7
```

At line 3 you specify the name for your compiled program and on line 5 you enumerate the source files. Now an invocation of `make` is enough to get your program compiled. Note that Go ships with the variant of `make`, called `gomake`, which is (currently) a small wrapper around GNU `make`. As the build system for Go programs may change in the future and `make` make go away, we use `gomake`. Note that this Makefile creates an executable named `helloworld`, not `6.out` or `8.out`.

Variables, types and keywords

In the next sections we will look at variables, basic types, keywords and control structures of our new language. Go has a C-like feel when it comes to its syntax. If you want to put two (or more) statements on one line, they must be separated with a semicolon (;). Normally you don't need the semicolon.

Go is different from other languages in that the type of a variable is specified *after* the variable name. So not: `int a`, but `a int`. When declaring a variable it is assigned the "natural" null value for the type. This means that after `var a int`, `a` has a value of 0. With `var s string`, `s` is assigned the zero string, which is "".

Declaring and assigning in Go is a two step process, but they may be combined. Compare the following pieces of code which have the same effect.

Listing 2.3. Declaration with =

```
var a int
var b bool
a = 15
b = false
```

Listing 2.4. Declaration with :=

```
a := 15
b := false
```

On the left we use the `var` keyword to declare a variable and *then* assign a value to it. The code on the right uses `:=` to do this in one step (this form may only be used *inside* functions). In that case the variable type is *deduced* from the value. A value of 15 indicates an `int`, a value of `false` tells Go that the type should be `bool`. Multiple `var` declarations may also be grouped, `const` and `import` also allow this. Note the use of parentheses:


```
var (
    x int
    b bool
)
```

Multiple variables of the same type can also be declared on a single line: `var x, y int`, makes `x` and `y` both `int` variables. You can also make use of parallel assignment:

```
a, b := 20, 16
```

Which makes `a` and `b` both integer variables and assigns 20 to `a` and 16 to `b`.

A special name for a variable is `_` (underscore). Any value assigned to it, is discarded. In this example we only assign the integer value of 35 to `b` and discard the value 34.

```
_, b := 34, 35
```

Declared, but otherwise unused variables are a compiler error in Go, the following code generates this error: `i declared and not used`

```
package main
func main() {
    var i int
}
```

Boolean types

A boolean type represents the set of boolean truth values denoted by the predeclared constants `true` and `false`. The boolean type is `bool`.

Numerical types

Go has the well known types such as `int`, this type has the appropriate length for your machine. Meaning that on a 32 bits machine they are 32 bits, and on a 64 bits machine they are 64 bits. Note: an `int` is either 32 or 64 bits, no other values are defined. Same goes for `uint`.

If you want to be explicit about the length you can have that too with `int32`, or `uint32`. The full list for (signed and unsigned) integers is `int8`, `int16`, `int32`, `int64` and `byte`, `uint8`, `uint16`, `uint32`, `uint64`. With `byte` being an alias for `uint8`. For floating point values there is `float32` and `float64` (there is no `float` type). A 64 bit integer or floating point value is *always* 64 bit, also on 32 bit architectures.

Note however that these types are all distinct and assigning variables which mix these types is a compiler error, like in the following code:

Listing 2.5. Familiar types are still distinct

```
package main                                     1

func main() {                                     3
    var a int                                     4
    var b int32                                   5
    a = 15                                        6
    b = a + a                                     7
    b = b + 5                                     8
}
```

`← Generic integer type`
`← 32 bits integer type`
`← Illegal mixing of these types`
`← 5 is a (typeless) constant, so this is OK`

Gives the error on the assignment on line 7:

```
types.go:7: cannot use a + a (type int) as type int32 in assignment
```

The assigned values may be denoted using octal, hexadecimal or the scientific notation: 077, 0xFF, 1e3 or 6.022e23 are all valid.

Constants

Constants in Go are just that — constant. They are created at compile time, and can only be numbers, strings or booleans; `const x = 42` makes `x` a constant. You can use `iota`^b to enumerate values.

```
const (
    a = iota
    b = iota
)
```

The first use of `iota` will yield 0, so `a` is equal to 0, whenever `iota` is used again on a new line its value is incremented with 1, so `b` has a value of 1.

You can even do the following, let Go repeat the use of `= iota`:

```
const (
    a = iota
    b           ← Implicitly b = iota
)
```

You may also explicitly type a constant, if you need that:

```
const (
    a = 0           ← Is an int now
    b string = "0"
)
```

Strings

An important other built in type is `string`. Assigning a string is as simple as:

```
s := "Hello World!"
```

Strings in Go are a sequence of UTF-8 characters enclosed in double quotes ("). If you use the single quote (') you mean one character (encoded in UTF-8) — which is *not* a `string` in Go.

Once assigned to a variable the string can not be changed anymore: strings in Go are immutable. For people coming from C, the following is not legal in Go:

```
var s string = "hello"
s[0] = 'c'      ← Change first char. to 'c', this is an error
```

To do this in Go you will need the following:

```
s := "hello"
c := []byte(s)  ❶
```

^bThe word [iota] is used in a common English phrase, 'not one iota', meaning 'not the slightest difference', in reference to a phrase in the New Testament: "until heaven and earth pass away, not an iota, not a dot, will pass from the Law." [35]

```

c[0] = 'c'           ❶
s2 := string(c)      ❷
fmt.Printf("%s\n", s2) ❸

```

- ❶ Convert `s` to an array of bytes, see chapter 5 section "Conversions" on page 59;
- ❷ Change the first element of this array;
- ❸ Create a *new* string `s2` with the alteration;
- ❹ print the string with `fmt.Printf`.

Multi-line strings

Due to the insertion of semicolons (see [3] section "Semicolons"), you need to be careful with using multiline strings. If you write:

```

s := "Starting part"
   + "Ending part"

```

This is transformed into:

```

s := "Starting part";
   + "Ending part";

```

Which is not valid syntax, you need to write:

```

s := "Starting part" +
     "Ending part"

```

Then Go will not insert the semicolons in the wrong places. Another way would be to use *raw* string literals by using back quotes: ```:

```

s := `Starting part
     Ending part`

```

Be aware that in this last example `s` now also contains the newline. Unlike *interpreted* string literals a raw string literal's value is composed of the *uninterpreted* characters between the quotes.

Complex numbers

Go has native support for complex numbers. If you use them you need a variable of the type `complex128` (64 bit imaginary part). If you want something smaller there is `complex64` – for a 32 bits imaginary part. Complex numbers are written as `re + imi`, where `re` is the real part, `im` is the imaginary part and `i` is the literal `'i'` ($\sqrt{-1}$). An example of using complex numbers:

```

var c complex64 = 5+5i; fmt.Printf("Value is: %v", c)

```

will print: (5+5i)

Operators and built-in functions

Go supports the normal set of numerical operations, table 2.1 lists the current ones and their relative precedence. They all associate from left to right.

Table 2.1. Operator precedence

Precedence	Operator(s)
Highest	* / % << >> & &^ + - ^ == != < <= > >= <- &&
Lowest	

+ - * / and % all do what you would expect, & | ^ and &^ are bit operators for **bitwise and**, **bitwise or**, **bitwise xor** and **bit clear** respectively. The && and || operators are logical **and** and logical **or**. Not listed in the table is the logical **not**: ! .

Although Go does not support operator overloading (or method overloading for that matter), some of the built-in operators *are* overloaded. For instance + can be used for integers, floats, complex numbers and strings (adding strings is concatenating them).

Go keywords

Table 2.2. Keywords in Go

break	default	func	interface	select
case	defer	go	map	struct
chan	else	goto	package	switch
const	fallthrough	if	range	type
continue	for	import	return	var

Table 2.2 lists all the keywords in Go. In the following paragraphs and chapters we will cover them. Some of these we have seen already.

- For **var** and **const** see section “Variables, types and keywords” on page 3;
- **package** and **import** are briefly touch upon in section “Hello World”. In chapter 4 they are documented in more detail.

Others deserve more text and have their own chapter/section:

- **func** is used to declare functions and methods;
- **return** is used to return from functions, for both **func** and **return** see chapter 3 for the details;
- **go** is used for concurrency (chapter 7);
- **select** used to choose from different types of communication, see chapter 7;
- **interface** see chapter 6;
- **struct** is used for abstract data types, see chapter 5;
- **type** also see chapter 5.

Control structures

There are only a few control structures in Go ^c. For instance there is no `do` or `while` loop, only a `for`. There is a (flexible) `switch` statement and `if` and `switch` accept an optional initialization statement like that of `for`. There also is something called a type switch and a multiway communications multiplexer, `select` (see chapter 7). The syntax is different (from that in C): parentheses are not required and the body must *always* be brace-delimited.

If

In Go an `if` looks like this:

```
if x > 0 {           ← { is mandatory
    return y
} else {
    return x
}
```

Mandatory braces encourage writing simple `if` statements on multiple lines. It is good style to do so anyway, especially when the body contains a control statement such as a `return` or `break`.

Since `if` and `switch` accept an initialization statement, it's common to see one used to set up a (local) variable.

```
if err := file.Chmod(0664); err != nil {    ← nil is like C's NULL
    log.Stderr(err)    ← Scope of err is limited to if's body
    return err
}
```

You can use the logical operators (see table 2.1) as you would normally do:

```
if true && true {
    println("true")
}
if ! false {
    println("true")
}
```

In the Go libraries, you will find that when an `if` statement doesn't flow into the next statement—that is, the body ends in `break`, `continue`, `goto`, or `return`, the unnecessary `else` is omitted.

```
f, err := os.Open(name, os.O_RDONLY, 0)
if err != nil {
    return err
}
doSomething(f)
```

This is an example of a common situation where code must analyze a sequence of error possibilities. The code reads well if the successful flow of control runs down the page, eliminating error cases as they arise. Since error cases tend to end in `return` statements, the resulting code needs no `else` statements.

^cThis section is copied from [3].

```
f, err := os.Open(name, os.O_RDONLY, 0)
if err != nil {
    return err
}
d, err := f.Stat()
if err != nil {
    return err
}
doSomething(f, d)
```

Syntax wise the following is illegal in Go:

```
if err != nil
{
    return err
}
```

← Must be on the same line as the if

See [3] section "Semicolons" for the deeper reasons behind this.

Ending with **if-then-else**

Note that if you end a function like this:

```
if err != nil {
    return err
} else {
    return nil
}
```

It will not compile. This is a bug in the Go compiler. See [16] for a extended problem description and hopefully a fix.

Goto

Go has a **goto** statement — use it wisely. With **goto** you jump to a label which must be defined within the current function. For instance a loop in disguise:

```
func myfunc() {
    i := 0
Here:    ← First word on a line ending with a colon is a label
    println(i)
    i++
    goto Here    ← Jump
}
```

The name of the label is case sensitive.

For

The Go **for** loop has three forms, only one of which has semicolons.

```
for init; condition; post { }    ← Like a C for

for condition { }                ← Like a while
```

```
for { } ← Like a C for(;;) (endless loop)
```

Short declarations make it easy to declare the index variable right in the loop.

```
sum := 0
for i := 0; i < 10; i++ {
    sum += i ← Short for sum = sum + i
} ← i ceases to exist after the loop
```

Finally, since Go has no comma operator and ++ and -- are statements not expressions, if you want to run multiple variables in a **for** you should use parallel assignment.

```
// Reverse a
for i, j := 0, len(a)-1; i < j; i, j = i+1, j-1 { ← Parallel assignment
    a[i], a[j] = a[j], a[i] ← Here too
}
```

Break and continue

With **break** you can quit loops early, **break** breaks the current loop.

```
for i := 0; i < 10; i++ {
    if i > 5 {
        break ← Stop this loop, making it only print 0 to 5
    }
    println(i)
}
```

With loops within loops you can specify a label after **break**. Making the label identify *which* loop to stop:

```
J: for j := 0; j < 5; j++ {
    for i := 0; i < 10; i++ {
        if i > 5 {
            break J ← Now it breaks the j-loop, not the i one
        }
        println(i)
    }
}
```

With **continue** you begin the next iteration of the loop, skipping any remaining code. In the same way as **break**, **continue** also accepts a label. The following prints 0 to 5.

```
for i := 0; i < 10; i++ {
    if i > 5 {
        continue ← Skip any remaining code
    }
    println(i)
}
```

Range

The keyword **range** can be used for loops. It can loop over slices, arrays, strings, maps and channels (see chapter 7). **range** is an iterator, that when called, returns a key-value pair from the thing it loops over. Depending on what that is, **range** returns different things.

When looping over a slice or array **range** returns the index in the slice as the key and value belonging to that index. Consider this code:

```
list := []string{"a", "b", "c", "d", "e", "f"} ❶
for k, v := range list { ❷
    // do what you want with k and v ❸
}
```

- ❶ Create a slice (see "Arrays, slices and maps" on page 14) of strings.
- ❷ Use **range** to loop over them. With each iteration **range** will return the index as **int** and the key as a **string**, starting with 0 and "a".
- ❸ k will have the value 0...5, and v will loop through "a"... "f".

You can also use **range** on strings directly. Then it will break out the individual Unicode characters^d and their start position, by parsing the UTF-8. The loop:

```
for pos, char := range "aΦx" {
    fmt.Printf("character '%c' starts at byte position %d\n", char, pos)
}
```

prints

```
character 'a' starts at byte position 0
character 'Φ' starts at byte position 1
character 'x' starts at byte position 3    ← Φ took 2 bytes
```

Switch

Go's **switch** is very flexible. The expressions need not be constants or even integers, the cases are evaluated top to bottom until a match is found, and if the **switch** has no expression it switches on **true**. It's therefore possible – and idiomatic – to write an **if-else-if-else** chain as a **switch**.

```
func unhex(c byte) byte {
    switch {
    case '0' <= c && c <= '9':
        return c - '0'
    case 'a' <= c && c <= 'f':
        return c - 'a' + 10
    case 'A' <= c && c <= 'F':
        return c - 'A' + 10
    }
    return 0
}
```

^dIn the UTF-8 world characters are sometimes called runes. Mostly, when people talk about characters, they mean 8 bit characters. As UTF-8 characters may be up to 32 bits the word rune is used.

There is no automatic fall through, you can however use **fallthrough** to make do just that.

Without **fallthrough**:

```
switch i {
    case 0: // empty case body
    case 1:
        f() // f is not called when i == 0!
}
```

And with:

```
switch i {
    case 0: fallthrough
    case 1:
        f() // f is called when i == 0!
}
```

With **default** you can specify an action when none of the other cases match.

```
switch i {
    case 0:
    case 1:
        f()
    default:
        g() // called when i is not 0 or 1
}
```

Cases can be presented in comma-separated lists.

```
func shouldEscape(c byte) bool {
    switch c {
    case ' ', '?', '&', '=', '#', '+': ← , as "or"
        return true
    }
    return false
}
```

Here's a comparison routine for byte arrays that uses two **switch** statements:

```
// Compare returns an integer comparing the two byte arrays
// lexicographically.
// The result will be 0 if a == b, -1 if a < b, and +1 if a > b
func Compare(a, b []byte) int {
    for i := 0; i < len(a) && i < len(b); i++ {
        switch {
        case a[i] > b[i]:
            return 1
        case a[i] < b[i]:
            return -1
        }
    }
    // String are equal except for possible tail
    switch {
    case len(a) < len(b):
        return -1
    }
```

```

    case len(a) > len(b):
        return 1
    }
    return 0    // Strings are equal
}

```

Built-in functions

A small number of functions are predefined, meaning you *don't* have to include any package to get access to them. Table 2.3 lists them all.

Table 2.3. Pre-defined functions in Go

close	new	panic	complex
closed	make	recover	real
len	append	print	imag
cap	copy	println	

`close` and `closed` are used in channel communication and the closing of those channels, see chapter 7 for more on this.

`len` and `cap` are used on a number of different types, `len` is used for returning the length of strings and the length of slices and arrays. See section "Arrays, slices and maps" for the details of slices and arrays and the function `cap`.

`new` is used for allocating memory for user defined data types. See section "Allocation with new" on page 55.

`make` is used for allocating memory for built-in types (maps, slices and channels). See section "Allocation with make" on page 55.

`copy` is used for copying slices. `append` is for concatenating slices. See section "Slices" in this chapter.

`panic` and `recover` are used for an *exception* mechanism. See the section "Panic and recovering" on page 31 for more.

`print` and `println` are low level printing functions that can be used without reverting to the `fmt` package. These are mainly used for debugging.

`complex`, `real` and `imag` all deal with complex numbers. Other than the simple example we gave, we will not further explain complex numbers.

Arrays, slices and maps

Storing multiple values in a list can be done by utilizing arrays, or their more flexible cousin: slices. A dictionary or hash type is also available, it is called a **map** in Go.

Arrays

An array is defined by: `[n]<type>`, where *n* is the length of the array and `<type>` is the stuff you want to store. Assigning, or indexing an element in the array is done with square brackets:

```
var arr [10]int
arr[0] = 42
arr[1] = 13
fmt.Printf("The first element is %d\n", arr[0])
```

Array types like `var arr = [10]int` have a fixed size. The size is *part* of the type. They can't grow, because then they would have a different type. Also arrays are values: Assigning one array to another *copies* all the elements. In particular, if you pass an array to a function, it will receive a copy of the array, not a pointer to it.

To declare an array you can use the following: `var a [3]int`, to initialize it to something else than zero, use a composite literal: `a := [3]int{1, 2, 3}` and this can be shortened to `a := [...]int{1, 2, 3}`, where Go counts the elements automatically. Note that all fields must be specified. So if you are using multidimensional arrays you have to do quite some typing:

```
a := [2][2]int{ [2]int{1,2}, [2]int{3,4} }
```

Which is the same as:

```
a := [2][2]int{ [...]int{1,2}, [...]int{3,4} }
```

When declaring arrays you *always* have to type something in between the square brackets, either a number or three dots (...) when using a composite literal. Since `release.2010-10-27` this syntax was further simplified. From the release notes:

The syntax for arrays, slices, and maps of composite literals has been simplified. Within a composite literal of array, slice, or map type, elements that are themselves composite literals may elide the type if it is identical to the outer literal's element type.

This means our example can become:

```
a := [2][2]int{ {1,2}, {3,4} }
```

Slices

A slice is similar to an array, but it can grow when new elements are added. A slice always refers to an underlying array. What makes slices different from arrays is that a slice is a pointer *to* an array; slices are reference types, which means that if you assign one slice to another, both refer to the same underlying array. For instance, if a function takes a slice argument, changes it makes to the elements of the slice will be visible to the caller, analogous to passing a pointer to the underlying array. With:

A composite literal allows you to assign a value directly to an array, slice or map.

Go release.2010-10-27 [7].

*Reference types are created with **make**.*

```
sl := make([]int, 10)
```

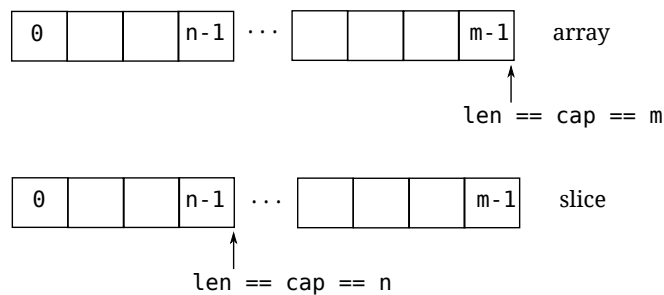
you create a slice which can hold ten elements. Note that the underlying array isn't specified. A slice is always coupled to an array that has a fixed size. For slices we define a capacity and a length. Figure 2.1 depicts the following Go code. First we create an array of m elements of the type `int`: `var array[m]int`

Next, we create a slice from this array: `slice := array[0:n]`

And now we have:

- `len(slice) == n == cap(slice) == n`;
- `len(array) == cap(array) == m`.

Figure 2.1. Array versus slice



Given an array, or another slice, a new slice is created via `a[I:J]`. This creates a new slice which refers to `a`, starts at index `I`, and ends before index `J`. It has length `J - I`.

```
// array[n:m], create a slice from array with elements n to m-1
a := [...]int{1, 2, 3, 4, 5} ❶
s1 := a[2:4] ❷
s2 := a[1:5] ❸
s3 := a[:] ❹
s4 := a[:4] ❺
s5 := s2[:]
```

- ❶ Define an array with 5 elements, from index 0 to 4;
- ❷ Create a slice with the elements from index 2 to 3, this contains: 3, 4;
- ❸ Create a slice with the elements from index 1 to 4, contains: 2, 3, 4, 5;
- ❹ Create a slice with all the elements of the array in it. This is a shorthand for: `a[0:len(a)]`;
- ❺ Create a slice with the elements from index 0 to 3, this is thus short for: `a[0:4]`, and yields: 1, 2, 3, 4;
- ❻ Create a slice from the slice `s2`, note that `s5` still refers to the array `a`.

In the code listed in 2.6 we dare to do the impossible on line 8 and try to allocate something beyond the capacity (maximum length of the underlying array) and we are greeted with a *runtime* error.

Listing 2.6. Arrays and slices

```

package main                                     1

func main() {                                     3
    var array [100]int    // Create array, index from 0 to 99    4
    slice := array[0:99]  // Create slice, index from 0 to 98    5

    slice[98] = 'a'      // OK                                     7
    slice[99] = 'a'      // Error: "throw: index out of range"    8
}                                                                9

```

If you want to extend a slice, there are a couple of built-in functions that make life easier: **append** and **copy**. From [5]:

*The function **append** appends zero or more values x to a slice s and returns the resulting slice, with the same type as s . If the capacity of s is not large enough to fit the additional values, **append** allocates a new, sufficiently large slice that fits both the existing slice elements and the additional values. Thus, the returned slice may refer to a different underlying array.*

```

s0 := []int{0, 0}
s1 := append(s0, 2)    ❶
s2 := append(s1, 3, 5, 7) ❷
s3 := append(s2, s0...) ❸

```

- ❶ append a single element, `s1 == []int{0, 0, 2};`
- ❷ append multiple elements, `s2 == []int{0, 0, 2, 3, 5, 7};`
- ❸ append a slice, `s3 == []int{0, 0, 2, 3, 5, 7, 0, 0}`. Note the three dots!

And

*The function **copy** copies slice elements from a source src to a destination dst and returns the number of elements copied. Source and destination may overlap. The number of arguments copied is the minimum of `len(src)` and `len(dst)`.*

```

var a = [...]int{0, 1, 2, 3, 4, 5, 6, 7}
var s = make([]int, 6)
n1 := copy(s, a[0:])    ← n1 == 6, s == []int{0, 1, 2, 3, 4, 5}
n2 := copy(s, s[2:])     ← n2 == 4, s == []int{2, 3, 4, 5, 4, 5}

```

Maps

Many other languages have a similar type built-in, Perl has hashes, Python has its dictionaries and C++ also has maps (in *lib*) for instance. In Go we have the **map** type. A **map** can be thought of as an array indexed by strings (in its most simple form). In the following listing we define a **map** which converts from a **string** (month abbreviation) to an **int** – the number of days in that month. The generic way to define a map is with: `map[<from type>]<to type>`

```
monthdays := map[string]int{
    "Jan": 31, "Feb": 28, "Mar": 31,
    "Apr": 30, "May": 31, "Jun": 30,
    "Jul": 31, "Aug": 31, "Sep": 30,
    "Oct": 31, "Nov": 30, "Dec": 31,    ← The comma here is required
}
```

Note to use **make** when only declaring a **map**: `monthdays := make(map[string]int)`

For indexing (searching) in the map, we use square brackets, for example suppose we want to print the number of days in December: `fmt.Printf("%d\n", monthdays["Dec"])` If you are looping over an array, slice, string, or map a **range** clause help you again, which returns the key and corresponding value with each invocation.

```
year := 0
for _, days := range monthdays {    ← Key is not used, hence _, days
    year += days
}
fmt.Printf("Numbers of days in a year: %d\n", year)
```

Adding elements to the **map** would be done as:

```
monthdays["Undecim"] = 30    ← Add a month
monthdays["Feb"]      = 29    ← Overwrite entry - for leap years
```

To test for existence, you would use the following[26]:

```
var value int
var present bool

value, present = monthdays["Jan"]    ← If exist, present has the value true
                                     ← Or better and more Go like
v, ok := monthdays["Jan"]           ← Hence, the "comma ok" form
```

And finally you can remove elements from the **map**:

```
monthdays["Mar"] = 0, false    ← Deletes "Mar", always rains anyway
```

Which looks a bit like the reverse of the "comma ok" form.

Exercises

Q2. (1) For-loop

1. Create a simple loop with the **for** construct. Make it loop 10 times and print out the loop counter with the *fmt* package.
2. Rewrite the loop from 1. to use **goto**. The keyword **for** may not be used.
3. Rewrite the loop again so that it fills an array and then prints that array to the screen.

Q3. (1) FizzBuzz

1. Solve this problem, called the Fizz-Buzz [31] problem:

Write a program that prints the numbers from 1 to 100. But for multiples of three print "Fizz" instead of the number and for the multiples of five print "Buzz". For numbers which are multiples of both three and five print "FizzBuzz".

Q4. (1) Strings

1. Create a Go program that prints the following (up to 100 characters):

```
A
AA
AAA
AAAA
AAAAA
AAAAAA
AAAAAAA
...
```

2. Create a program that counts the numbers of characters in this string:
`asSASA ddd dsjkdsjs dk`
Make it also output the number of bytes in that string. *Hint.* Check out the `utf8` package.
3. Extend the program from the previous question to replace the three runes at position 4 with `'abc'`.
4. Write a Go program that reverses a string, so `"foobar"` is printed as `"raboof"`. *Hint.* Unfortunately you need to know about conversion, skip ahead to section "Conversions" on page 59.

Q5. (4) Average

1. Give the code that calculates the average of a `float64` slice. In a later exercise (Q6) you will make it into a function.

Answers

A2. (1) For-loop

1. There are a multitude of possibilities, one of the solutions could be:

Listing 2.7. Simple for loop

```
package main

import "fmt"

func main() {
    for i := 0; i < 10; i++ {    ← See section For on page 9
        fmt.Printf("%d\n", i)
    }
}
```

Lets compile this and look at the output.

```
% 6g for.go && 6l -o for for.6
% ./for
0
1
.
.
.
9
```

2. Rewriting the loop results in code that should look something like this (only showing the main-function):

```
func main() {
    i := 0                ← Define our loop variable
I:                      ← Define a label
    fmt.Printf("%d\n", i)
    i++
    if i < 10 {
        goto I          ← Jump back to the label
    }
}
```

3. The following might one possible solution:

Listing 2.8. For loop with an array

```
func main() {
    var arr [10]int      ← Create an array with 10 elements
    for i := 0; i < 10; i++ {
        arr[i] = i      ← Fill it one by one
    }
    fmt.Printf("%v", arr)    ← With %v Go prints the type
}
```


You could even do this in one fell swoop by using a composite literal:

```
a := [...]int{0,1,2,3,4,5,6,7,8,9}    ← With [...] you let Go count
fmt.Printf("%v\n", a)
```

A3. (1) FizzBuzz

1. A possible solution to this simple problem is the following program.

Listing 2.9. Fizz-Buzz

```
package main

import "fmt"

func main() {
    const (
        FIZZ = 3 ❶
        BUZZ = 5
    )
    var p bool ❶
    for i := 1; i < 100; i++ { ❷;
        p = false
        if i%FIZZ == 0 { ❸
            fmt.Printf("Fizz")
            p = true
        }
        if i%BUZZ == 0 { ❹
            fmt.Printf("Buzz")
            p = true
        }
        if !p { ❺
            fmt.Printf("%v", i)
        }
        fmt.Println() ❻
    }
}
```

- ❶ Define two constants to make the code more readable. See section "Constants";
- ❶ Holds if we already printed something;
- ❷ for-loop, see section "For"
- ❸ If divisible by FIZZ, print "Fizz";
- ❹ And if divisible by BUZZ, print "Buzz". Note that we have also taken care of the FizzBuzz case;
- ❺ If neither FIZZ nor BUZZ printed, print the value;
- ❻ Format each output on a new line.

A4. (1) Strings

1. This program is a solution:

Listing 2.10. Strings

```
package main

import "fmt"

func main() {
    str := "A"
    for i := 0; i < 100; i++ {
        fmt.Printf("%s\n", str)
        str = str + "A"    ← String concatenation
    }
}
```

2. To answer this question we need some help of the *utf8* package. First we check the documentation with `godoc utf8 | less`. When we read the documentation we notice `func RuneCount(p []byte)int`. Secondly we can convert *string* to a **byte** slice with

```
str := "hello"
b    := []byte(str)    ← Conversion, see page 59
```

Putting this together leads to the following program.

Listing 2.11. Runes in strings

```
package main

import (
    "fmt"
    "utf8"
)

func main() {
    str := "dsjdkshdjsdh...js"
    fmt.Printf("String %s\nLenght: %d, Runes: %d\n", str,
        len([]byte(str)), utf8.RuneCount([]byte(str)))
}
```

3. Reversing a string can be done as follows. We start from the left (i) and the right (j) and swap the characters as we see them:

Listing 2.12. Reverse a string

```
import "fmt"

func main() {
    s := "foobar"
    a := []byte(s)    ← Again a conversion
    // Reverse a
    for i, j := 0, len(a)-1; i < j; i, j = i+1, j-1 {
        a[i], a[j] = a[j], a[i]    ← Parallel assignment
    }
    fmt.Printf("%s\n", string(a))    ← Convert it back
}
```

A5. (4) Average

1. The following code calculates the average.

```
sum := 0.0
switch len(xs) {
case 0:    ❶
    ave = 0
default:    ❷
    for _, v := range xs {
        sum += v
    }
    ave = sum / float64(len(xs))    ❸
}
```

- ❶ If the length is zero, we return 0;
- ❷ Otherwise we calculate the average;
- ❸ We have to convert the value to a **float64** to make the division work.

3

Functions

I'm always delighted by the light touch and stillness of early programming languages. Not much text; a lot gets done. Old programs read like quiet conversations between a well-spoken research worker and a well-studied mechanical colleague, not as a debate with a compiler. Who'd have guessed sophistication bought such noise?


RICHARD P. GABRIEL

Functions are the basic building blocks in Go programs; all interesting stuff happens in them. A function is declared as follows:

Listing 3.1. A function declaration

```
type mytype int    ← New type, see chapter 5

func (p mytype) funcname(q int) (r,s int) { return 0,0 }
```



- ❶ The keyword **func** is used to declare a function;
- ❷ A function can be defined to work on a specific type, a more common name for such a function is method. This part is called a *receiver* and it is optional. This will be used in chapter 6;
- ❸ *funcname* is the name of your function;
- ❹ The variable *q* of type **int** is the input parameter. The parameters are passed *pass-by-value* meaning they are copied. But be aware that reference types (slices, channels, maps and interfaces) are *pass-by-reference* even though you do not see the pointers directly in the code;
- ❺ The variables *r* and *s* are the named return parameters for this function. Note that functions in Go can have multiple return values. See section "Multiple return values" on page 26. If you want the return parameters not to be named you only give the types: **(int,int)**. If you have only one value to return you may omit the parentheses. If your function is a subroutine and does not have anything to return you may omit this entirely;
- ❻ This is the function's body, note that **return** is a statement so the braces around the parameter(s) are optional.

Here are a two examples, the left is a function without a return value, the one on the right is a simple function that returns its input.

```
func subroutine(in int) {
    return
}

func identity(in int) int {
    return in
}
```

Functions can be declared in any order you wish, the compiler scans the entire file before execution. So function prototyping is a thing of the past in Go. Go disallows nested functions. You can however work around this by using anonymous functions, see section “Functions as values” on page 30 in this chapter.

Recursive functions just work as in other languages:

Listing 3.2. Recursive function

```
func rec(i int) {
    if i == 10 {
        return
    }
    rec(i+1)
    fmt.Printf("%d ", i)
}
```

This prints 9 8 7 6 5 4 3 2 1 0.

Scope

Variables declared outside any functions are global in Go, those defined in functions are local to those functions. If names overlap — a local variable is declared with the same name as a global one — the local variable hides the global one when the current function is executed.

Listing 3.3. Local scope

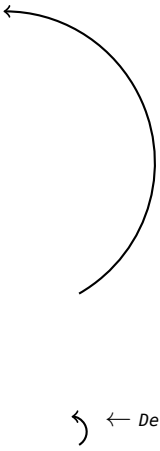
```
package main

var a = 6

func main() {
    p()
    q()
    p()
}

func p() {
    println(a)
}

func q() {
    a := 5
    println(a)
}
```



← Definition

Listing 3.4. Global scope


```
package main

var a = 6

func main() {
    p()
    q()
    p()
}

func p() {
    println(a)
}

func q() {
    a = 5
    println(a)
}
```



← Assignment

In listing 3.3 we introduce a local variable `a` in the function `q()`. This local `a` is only

visible in `q()`. That is why the code will print: 656. In listing 3.4 no new variables are introduced, there is only a global `a`. Assigning a new value to `a` will be globally visible. This code will print: 655

In the following example we call `g()` from `f()`:

Listing 3.5. Scope when calling functions from functions

```
package main

var a int

func main() {
    a = 5
    println(a)
    f()
}

func f() {
    a := 6
    println(a)
    g()
}

func g() {
    println(a)
}
```

The printout will be: 565. A *local* variable is *only* valid when we are executing the function in which it is defined.

Multiple return values

One of Go's unusual features is that functions and methods can return multiple values (Python can do this too). This can be used to improve on a couple of clumsy idioms in C programs: in-band error returns (such as -1 for EOF) and modifying an argument. In Go, `Write` returns a count and an error: "Yes, you wrote some bytes but not all of them because you filled the device". The signature of `*File.Write` in package `os` is:

```
func (file *File) Write(b []byte) (n int, err Error)
```

and as the documentation says, it returns the number of bytes written and a non-nil `Error` when `n != len(b)`. This is a common style in Go.

A similar approach obviates the need to pass a pointer to a return value to simulate a reference parameter. Here's a simple-minded function to grab a number from a position in a byte array, returning the number and the next position.

```
func nextInt(b []byte, i int) (int, int) {
    x := 0
    // Naively assume everything is a number
    for ; i < len(b); i++ {
        x = x*10 + int(b[i])-'0'
    }
    return x, i
}
```

```
}

```

You could use it to scan the numbers in an input array a like this:

```
a := []byte{'1', '2', '3', '4'}
var x int
for i := 0; i < len(a); {           ← No i++
    x, i = nextInt(a, i)
    println(x)
}
```

Without having tuples as a native type, multiple return values is the next best thing to have. You can return precisely what you want without overloading the domain space with special values to signal errors.

Named result parameters

The return or result parameters of a Go function can be given names and used as regular variables, just like the incoming parameters. When named, they are initialized to the zero values for their types when the function begins; if the function executes a **return** statement with no arguments, the current values of the result parameters are used as the returned values. Using this features enables you (again) to do more with less code ^a.

The names are not mandatory but they can make code shorter and clearer: *they are documentation*. If we name the results of `nextInt` it becomes obvious which returned `int` is which.

```
func nextInt(b []byte, pos int) (value, nextPos int) { /* ... */ }
```

Because named results are initialized and tied to an unadorned **return**, they can simplify as well as clarify. Here's a version of `io.ReadFull` that uses them well:

```
func ReadFull(r Reader, buf []byte) (n int, err os.Error) {
    for len(buf) > 0 && err == nil {
        var nr int
        nr, err = r.Read(buf)
        n += nr
        buf = buf[nr:len(buf)]
    }
    return
}
```

In the following example we declare a simple function which calculates the factorial value of a value `x`.

Some text in this section comes from [15].

```
func Factorial(x int) int {           ← func Factorial(x int) (int) is also OK
    if x == 0 {
        return 1
    } else {
        return x * Factorial(x - 1)
    }
}
```

So you could also write factorial as:

^aThis is a motto of Go; "Do more with less code".


```
func Factorial(x int) (result int) {  
    if x == 0 {  
        result = 1  
    } else {  
        result = x * Factorial(x - 1)  
    }  
    return  
}
```

When we use named result values, the code is shorter and easier to read. You can also write a function with multiple return values:

```
func fib(n) (val, pos int) {    ← Both ints  
    if n == 0 {  
        val = 1  
        pos = 0  
    } else if n == 1 {  
        val = 1  
        pos = 1  
    } else {  
        v1, _ := fib(n-1)  
        v2, _ := fib(n-2)  
        val = v1 + v2  
        pos = n  
    }  
    return  
}
```

Deferred code

Suppose you have a function in which you open a file and perform various writes and reads on it. In such a function there are often spots where you want to return early. If you do that, you will need to close the file descriptor you are working on. This often leads to the following code:

Listing 3.6. Without defer

```
func ReadWrite() bool {  
    file.Open("file")  
    // Do your thing  
    if failureX {  
        file.Close()  
        return false  
    }  
  
    if failureY {  
        file.Close()  
        return false  
    }  
    file.Close()  
    return true  
}
```

```
}
```

Here a lot of code is repeated. To overcome this Go has the **defer** statement. After **defer** you specify a function which is called just *before* a return from the function is executed.

The code above could be rewritten as follows. This makes the function more readable, shorter and puts the Close right next to the Open.

Listing 3.7. With defer

```
func ReadWrite() bool {
    file.Open("file")
    defer file.Close()    ← file.Close() is the function
    // Do your thing
    if failureX {
        return false      ← Close() is now done automatically
    }
    if failureY {
        return false      ← And here too
    }
    return true
}
```

You can put multiple functions on the "deferred list", like this example from [3]:

```
for i := 0; i < 5; i++ {
    defer fmt.Printf("%d ", i)
}
```

Deferred functions are executed in LIFO order, so the above code prints: 4 3 2 1 0.

With defer you can even change return values, provided that you are using named result parameters and a function literal^b, i.e:

Listing 3.8. Function literal

```
defer func() {
    /* ... */
}()    ← () is needed here
```

Or this example which makes it easier to understand why and where you need the braces:

Listing 3.9. Function literal with parameters

```
defer func(x int) {
    /* ... */
}(5)    ← Give the input variable x the value 5
```

In that (unnamed) function you can access any named return parameter:

Listing 3.10. Access return values within defer

```
func f() (ret int) {    ← ret is initialized with zero
    defer func() {
        ret++          ← Increment ret with 1
    }()
    return 0            ← 1 not 0 will be returned!
}
```

^bA function literal is sometimes called a closure.

Variadic parameters

Functions that take variadic parameters are functions that have a variable number of parameters. To do this, you first need to declare your function to take variadic arguments:

```
func myfunc(arg ...int) {}
```

The `arg ... int` instructs Go to see this as a function that takes a variable number of arguments. Note that these arguments all have the type `int`. Inside your function's body the variable `arg` is a slice of ints:

```
for _, n := range arg {
    fmt.Printf("And the number is: %d\n", n)
}
```

If you don't specify the type of the variadic argument it defaults to the empty interface `interface{}` (see chapter 6"). Suppose we have another variadic function called `myfunc2`, the following example shows how to pass the variadic arguments to it:

```
func myfunc(arg ...int) {
    myfunc2(arg...)      ← Pass it as-is
    myfunc2(arg[:2]...)  ← Slice it
}
```

Functions as values

As with almost everything in Go, functions are also *just* values. They can be assigned to variables as follows:

Listing 3.11. Anonymous function

```
func main() {
    a := func() {          ← Define a nameless function and assign to a
        println("Hello")
    }                      ← No () here
    a()                    ← Call the function
}
```

If we use `fmt.Printf("%T\n", a)` to print the type of `a`, it prints `func()`.

Functions-as-values may also be used in other places, like in maps. Here we convert from integers to functions:

Listing 3.12. Functions as values in maps

```
var xs = map[int]func() int{
    1: func() int { return 10 },
    2: func() int { return 20 },
    3: func() int { return 30 },    ← Mandatory ,
    /* ... */
}
```

Or you can write a function that takes a function as its parameter, for example a `Map` function that works on `int` slices. This is left as an exercise for the reader, see exercise Q12 on page 33.

Callbacks and closures

With functions as values they are easy to pass to functions, from where they can be used as callbacks. First define a function that does "something" with an integer value:

```
func printit(x int) {           ← Function returns nothing
    fmt.Print("%v\n", x)        ← Just print it
}
```

The signature of this function is: **func** printit(**int**), or without the function name: **func**(**int**). To create a new function that uses this one as a callback we need to use this signature:

```
func callback(y int, f func(int)) {    ← f will hold the function
    f(y)                                ← Call the callback f with y
}
```

We've already seen some use of closures in section "Deferred code", but there is more to tell. When you define a closure, i.e. when you start using a function literal you still have access to the (local) variables defined in the current function.

```
// Define some local vars
// This code is WAY to complex, but illustrates it nicely
frameSquare := func(x, y int) {
    // closure effortlessly passes local variables to callback
    if thickFrame {
        // draw 3 x 3 pixel block for thicker rectangle
        for x0 := x - 1; x0 <= x+1; x0++ {
            for y0 := y - 1; y0 <= y+1; y0++ {
                rgba.Set(x0, y0, redColor)
            }
        }
    } else {
        rgba.Set(x, y, blueColor)
    }
}
```

If you want to do this by *not* using a closure and defining a completely new function you need to pass all vars to the function.

TODO
Gist is there, but
needs much better
wording and code

Panic and recovering

Go does not have an exception mechanism, like that in Java for instance: you can not throw exceptions. Instead it uses a panic-and-recover mechanism. It is worth remembering that you should use this as a last resort, your code will not look, or be, better if it is littered with panics. It's a powerful tool: use it wisely. So, how do you use it.

The following description was taken from [2]:

Panic

is a built-in function that stops the ordinary flow of control and begins panicking. When the function **F** calls **panic**, execution of **F** stops, any deferred functions in **F** are executed normally, and then **F** returns to its caller. To the caller, **F** then behaves like

a call to **panic**. The process continues up the stack until all functions in the current goroutine have returned, at which point the program crashes.

Panics can be initiated by invoking `panic` directly. They can also be caused by *run-time errors*, such as out-of-bounds array accesses.

Recover

is a built-in function that regains control of a panicking goroutine. *Recover* is *only* useful inside *deferred* functions.

During normal execution, a call to `recover` will return `nil` and have no other effect. If the current goroutine is panicking, a call to `recover` will capture the value given to `panic` and resume normal execution.

Exercises

Q6. (4) Average

1. Write a function that calculates the average of a `float64` slice.

Q7. (3) Integer ordering

1. Write a function that returns its (two) parameters in the right, numerical (ascending) order:

$f(7, 2) \rightarrow 2, 7$

$f(2, 7) \rightarrow 2, 7$

Q8. (4) Scope

1. What is wrong with the following program?

```
package main                                1

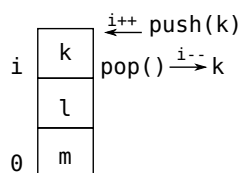
import "fmt"                                3

func main() {                                5
    for i := 0; i < 10; i++ {                6
        fmt.Printf("%v\n", i)                7
    }                                         8
    fmt.Printf("%v\n", i)                    9
}
```

Q9. (5) Stack

1. Create a simple stack which can hold a fixed amount of ints. It does not have to grow beyond this limit. Define both a `push` – put something on the stack – and a `pop` – retrieve something from the stack – function. The stack should be a LIFO (last in, first out) stack.

Figure 3.1. A simple LIFO stack



2. Bonus. Write a `String` method which converts the stack to a string representation. This way you can print the stack using: `fmt.Printf("My stack %v\n", stack)`
The stack in the figure could be represented as: `[0:m] [1:l] [2:k]`

Q10. (5) Var args

1. Write a function that takes a variable numbers of ints and prints each integer on a separate line

Q11. (5) Fibonacci

1. The Fibonacci sequence starts as follows: 1, 1, 2, 3, 5, 8, 13, ... Or in mathematical terms: $x_1 = 1; x_2 = 1; x_n = x_{n-1} + x_{n-2} \quad \forall n > 2$.
Write a function that takes an `int` value and gives that many terms of the Fibonacci sequence.

Q12. (4) Map function A `map()`-function is a function that takes a function and a list. The function is applied to each member in the list and a new list containing these calculated values is returned. Thus:

$$\text{map}(f(), (a_1, a_2, \dots, a_{n-1}, a_n)) = (f(a_1), f(a_2), \dots, f(a_{n-1}), f(a_n))$$

1. Write a simple `map()`-function in Go. It is sufficient for this function only to work for ints.
2. Expand your code to also work on a list of strings.

Q13. (3) Minimum and maximum

1. Write a function that calculates the maximum value in an `int` slice (`[]int`).
2. Write a function that calculates the minimum value in a `int` slice (`[]int`).

Q14. (5) Bubble sort

1. Write a function that performs Bubble sort on slice of ints. From [32]:

It works by repeatedly stepping through the list to be sorted, comparing each pair of adjacent items and swapping them if they are in the wrong order. The pass through the list is repeated until no swaps are needed, which indicates that the list is sorted. The algorithm gets its name from the way smaller elements "bubble" to the top of the list.

[32] also gives an example in pseudo code:

```
procedure bubbleSort( A : list of sortable items )
do
  swapped = false
  for each i in 1 to length(A) - 1 inclusive do:
    if A[i-1] > A[i] then
      swap( A[i-1], A[i] )
      swapped = true
    end if
  end for
while swapped
end procedure
```

Q15. (6) Functions that return functions

1. Write a function that returns a function that performs a `+2` on integers. Name the function `plusTwo`. You should then be able to do the following:

```
p := plusTwo()  
fmt.Printf("%v\n", p(2))
```

Which should print 4. See section Callbacks and closures on page 31 for information about this topic.

2. Generalize the function from 1, and create a `plusX(x)` which returns a function that add `x` to an integer.

Answers

A6. (4) Average

1. The following function calculates the average.

Listing 3.13. Average function in Go

```
func average(xs []float64) (ave float64) { ❶
    sum := 0.0
    switch len(xs) {
    case 0: ❷
        ave = 0
    default: ❸
        for _, v := range xs {
            sum += v
        }
        ave = sum / float64(len(xs)) ❹
    }
    return ❺
}
```

- ❶ We use a named return parameter;
- ❷ If the length is zero, we return 0;
- ❸ Otherwise we calculate the average;
- ❹ We have to convert the value to a **float64** to make the division work;
- ❺ We have an average, return it.

A7. (3) Integer ordering

1. Here we can use the multiple return values (section "Multiple return values") from Go:

```
func order(a, b int) (int, int) {
    if a > b {
        return b, a
    }
    return a, b
}
```

A8. (4) Scope

1. The program does not even compile, because **i** on line 9 is not defined: **i** is only defined within the **for**-loop. To fix this the function **main()** should become:

```
func main() {
    var i int
    for i = 0; i < 10; i++ {
        fmt.Printf("%v\n", i)
    }
}
```



```

        fmt.Printf("%v\n", i)
    }

```

Now `i` is defined outside the `for`-loop and still visible afterwards. This code will print the numbers 0 through 10.

A9. (5) Stack

1. First we define a new type that represents a stack; we need an array (to hold the keys) and an index, which points to the last element. Our small stack can only hold 10 elements.

```

type stack struct {    ← stack is not exported
    i    int
    data [10]int
}

```

Next we need the push and pop functions to actually use the thing. *First we show the wrong solution!* In Go data passed to functions is *passed-by-value* meaning a copy is created and given to the function. The first stab for the function push could be:

```

func (s stack) push(k int) {    ← Works on copy of argument
    if s.i+1 > 9 {
        return
    }
    s.data[s.i] = k
    s.i++
}

```

The function works on the `s` which is of the type `stack`. To use this we just call `s.push(50)`, to push the integer 50 on the stack. But the push function gets a copy of `s`, so it is *not* working the *real* thing. Nothing gets pushed to our stack this way, for example the following code:

```

var s stack    ← make s a simple stack variable
s.push(25)
fmt.Printf("stack %v\n", s);
s.push(14)
fmt.Printf("stack %v\n", s);

```

prints:

```

stack [0:0]
stack [0:0]

```

To solve this we need to give the function push a pointer to the stack. This means we need to change push from

```
func (s stack)push(k int) → func (s *stack)push(k int)
```

We should now use `new()` (see "Allocation with new" in chapter 5) to create a *pointer* to a newly allocated `stack`, so line 1 from the example above needs to be `s := new(stack)`

And our two functions become:

```

func (s *stack) push(k int) {
    s.data[s.i] = k
}

```

```

        s.i++
    }

    func (s *stack) pop() int {
        s.i--
        return s.data[s.i]
    }

```

Which we then use as follows

```

func main() {
    var s stack
    s.push(25)
    s.push(14)
    fmt.Printf("stack %v\n", s)
}

```

2. While this was a bonus question, having the ability to print the stack was very valuable when writing the code for this exercise. According to the Go documentation `fmt.Printf("%v")` can print any value (`%v`) that satisfies the `Stringer` interface. For this to work we only need to define a `String()` function for our type:

Listing 3.14. stack.String()

```

func (s stack) String() string {
    var str string
    for i := 0; i <= s.i; i++ {
        str = str + "[" +
            strconv.Itoa(i) + ":" + strconv.Itoa(s.data[i]) + "]"
    }
    return str
}

```

A10. (5) Var args

1. For this we need the `...`-syntax to signal we define a function that takes an arbitrary number of arguments.

Listing 3.15. A function with variable number of arguments

```

package main

import "fmt"

func main() {
    printthem(1, 4, 5, 7, 4)
    printthem(1, 2, 4)
}

func printthem(numbers ... int) {    ← numbers is now a slice of ints

```

```

        for _, d := range numbers {
            fmt.Printf("%d\n", d)
        }
    }
}

```

A11. (5) Fibonacci

1. The following program calculates the Fibonacci numbers.

Listing 3.16. Fibonacci function in Go

```

package main

import "fmt"

func fibonacci(value int) []int {
    x := make([]int, value) ❶
    x[0], x[1] = 1, 1 ❷
    for n := 2; n < value; n++ {
        x[n] = x[n-1] + x[n-2] ❸
    }
    return x ❹
}

func main() {
    for _, term := range fibonacci(10) { ❺
        fmt.Printf("%v ", term)
    }
}

```

- ❶ We create an **array** to hold the integers up to the value given in the function call;
- ❷ Starting point of the Fibonacci calculation;
- ❸ $x_n = x_{n-1} + x_{n-2}$;
- ❹ Return the *entire* array;
- ❺ Using the **range** keyword we "walk" the numbers returned by the fibonacci function. Here up to 10. And we print them.

A12. (4) Map function

Listing 3.17. A Map function

```

1. func Map(f func(int) int, l []int) []int {
    j := make([]int, len(l))
    for k, v := range l {
        j[k] = f(v)
    }
}

```

```

        return j
    }

    func main() {
        m := []int{1, 3, 4}
        f := func(i int) int {
            return i * i
        }
        fmt.Printf("%v", (Map(f, m)))
    }

```

2. Answer to question but now with strings

A13. (3) Minimum and maximum

1. This is a function for calculating a maximum:

```

func max(l []int) (max int) { ❶
    max = l[0]
    for _, v := range l { ❷
        if v > max { ❸
            max = v
        }
    }
    return ❹
}

```

- ❶ We use a named return parameter;
- ❷ Loop over `l`. The index of the element is not important;
- ❸ If we find a new maximum, remember it;
- ❹ A "lone" return, the current value of `max` is now returned.

2. This is a function for calculating a minimum, that is almost identical to `max`:

```

func min(l []int) (min int) {
    min = l[0]
    for _, v := range l {
        if v < min {
            min = v
        }
    }
    return
}

```

The interested reader may combine `max` and `min` into one function with a selector that lets you choose between the minimum or the maximum, or one that returns both values.

A14. (5) Bubble sort

1. The Bubble sort isn't terribly efficient, for n elements it scales $O(n^2)$. See Quick sort (ref XXX) for a better sorting algorithm.
But Bubble sort is easy to implement, the following is an example.

Listing 3.18. Bubble sort

```

func main() {
    n := []int{5, -1, 0, 12, 3, 5}
    fmt.Printf("unsorted %v\n", n)
    // even though it's call by value, a slice is a
    // reference type, so the underlying array is changed!
    bubblesort(n)
    fmt.Printf("sorted %v\n", n)
}

func bubblesort(n []int) {
    for i := 0; i < len(n) - 1; i++ {
        for j := i + 1; j < len(n); j++ {
            if n[j] < n[i] {
                n[i], n[j] = n[j], n[i]
            }
        }
    }
}

```

A15. (6) Functions that return functions

```

1. func main() {
    p2 := plusTwo()
    fmt.Printf("%v\n", p2(2))
}

func plusTwo() func(int) int { ❶
    return func(x int) int { return x + 2 } ❷
}

```

❶ Define a new function that returns a function. See how you can just write down what you mean;

❷ Function literals at work, we define the +2-function right there in the return statement.

2. Here we use a closure:

```

func plusX(x int) func(int) int { ❶
    return func(y int) int { return x + y } ❷
}

```

❶ Again define a function that returns a function;

❷ Use the *local* variable *x* in the function literal.

4

Packages

Answer to whether there is a bitwise negation operator.

KEN THOMPSON

Packages are a collection of functions and data, they are like the Perl packages[12]. You declare a package with the **package** keyword. The filename does not have to match the package name. The convention for package names is to use lowercase characters. Go packages may consist of multiple files, but they share the **package** <name> line. Lets define a package *even* in the file *even.go*.

Listing 4.1. A small package

```
package even           ← Start our own namespace

func Even(i int) bool { ← Exported function
    return i % 2 == 0
}

func odd(i int) bool {  ← Private function
    return i % 2 == 1
}
```

Names that start with a capital letter are *exported* and may be used outside your package, more on that later. We can now use the package as follows in our own program *myeven.go*:

Listing 4.2. Use of the even package

```
package main

import (                ❶
    "./even"            ❷
    "fmt"                ❸
)

func main() {
    i := 5
    fmt.Printf("Is %d even? %v\n", i, even.Even(i)) ❹
}
```

- ❶ Import the following packages;
- ❷ The *local* package *even* is imported here;
- ❸ The official *fmt* package gets imported;
- ❹ Use the function from the *even* package. The syntax for accessing a function from a package is <package>.Function().

Now we just need to compile and link, first the package, then `myeven.go` and then link it:

```
% 6g even.go      ← The package
% 6g myeven.go    ← Our program
% 6l -o myeven myeven.6      ← Linking stage
```

And test it:

```
% ./myeven
Is 5 even? false
```

In Go, a function from a package is exported (visible outside the package, i.e. public) when the first letter of the function name is a capital, hence the function name `Even`. If we change our `myeven.go` on line 7 to using to unexported function `even.odd`:

```
fmt.Printf("Is %d even? %v\n", i, even.odd(i))
```

We get an error when compiling, because we are trying to use a *private* function:

```
myeven.go:7: cannot refer to unexported name even.odd
myeven.go:7: undefined: even.odd
```

To summarize:

- Public functions have a name starting with a *capital* letter;
- Private function have a name starting with a *lowercase* letter.

This convention also holds true for other names (new types, global variables) you define in a package. Note that the term "capital" is not limited to US ASCII, it extends into the entire Unicode range. So captial Greek, Coptic, etc. is OK too.

Building a package

To create a package that other people can use (by just using `import "even"`) we first need to create a directory to put the package files in.

```
% mkdir even
% cp even.go even/
```

Next we can use the following `Makefile`, which is adapted for our *even* package.

Listing 4.3. Makefile for a package

```
include $(GOROOT)/src/Make.inc      1

TARG=even                          3
GOFILES=\                           4
    even.go\                         5

include $(GOROOT)/src/Make.pkg      7
```

Note that on line 3 we define our *even* package and on the lines 4 and 5 we enter the files which make up the package. Also note that this is *not* the *same* `Makefile` setup as we

used in section "Using a Makefile" in chapter 2. The last line with the `include` statement is different.

If we now issue `gomake`, a file named `"_go_.6"`, a directory named `"_obj/"` and a file inside `"_obj/"` called `"even.a"` is created. The file `even.a` is a static library which holds the compiled Go code. With `gomake install` the package (well, only the `even.a`) is installed in the *official* package directory:

```
% make install
cp _obj/even.a $GOROOT/pkg/linux_amd64/even.a
```

After installing we can change our `import` from `"./even"` to `"even"`.

But what if you do not want to install packages in the official Go directory, or maybe you do not have the permission to do so? When using 6/8g you can use the `-I`-flag to specify an alternate package directory. With this flag you can leave your import statement as-is (`import "even"`) and continue development in your own directory. So the following commands should build (and link) our `myeven` program with our package.

```
% 6g -I even/_obj myeven.go    ← Building
% 6l -L even/_obj myeven.6    ← Linking
```

Imported, but otherwise unused packages are an error.

Identifiers

Names are as important in Go as in any other language. In some cases they even have semantic effect: for instance, the visibility of a name outside a package is determined by whether its first character is upper case. It's therefore worth spending a little time talking about naming conventions in Go programs.

The convention that is used was to leave well-known legacy not-quite-words alone rather than try to figure out where the capital letters go. `Atoi`, `Getwd`, `Chmod`. Camelcasing works best when you have whole words to work with: `ReadFile`, `NewWriter`, `MakeSlice`.

Package names

When a package is imported (with `import`), the package name becomes an accessor for the contents. After

```
import "bytes"
```

the importing package can talk about `bytes.Buffer`. It's helpful if everyone using the package can use the same name to refer to its contents, which implies that the package name should be good: short, concise, evocative. By convention, packages are given lower case, single-word names; there should be no need for underscores or mixedCaps. Err on the side of brevity, since everyone using your package will be typing that name. And don't worry about collisions a priori. The package name is only the default name for imports. With the above import you can use `bytes.Buffer`. With

```
import bar "bytes"
```

it becomes `bar.Buffer`. So it does need not be unique across all source code, and in the rare case of a collision the importing package can choose a different name to use locally. In any case, confusion is rare because the file name in the import determines just which package is being used.

Another convention is that the package name is the base name of its source directory; the package in `src/pkg/container/vector` is imported as `container/vector` but has name `vector`, not `container_vector` and not `containerVector`.

The importer of a package will use the name to refer to its contents, so exported names in the package can use that fact to avoid stutter. For instance, the buffered reader type in the `bufio` package is called `Reader`, not `BufReader`, because users see it as `bufio.Reader`, which is a clear, concise name. Moreover, because imported entities are always addressed with their package name, `bufio.Reader` does not conflict with `io.Reader`. Similarly, the function to make new instances of `ring.Ring`—which is the definition of a constructor in Go—would normally be called `NewRing`, but since **`Ring`** is the only type exported by the package, and since the package is called `ring`, it's called just `New`. Clients of the package see that as `ring.New`. Use the package structure to help you choose good names.

Another short example is `once.Do`; `once.Do(setup)` reads well and would not be improved by writing `once.DoOrWaitUntilDone(setup)`. Long names don't automatically make things more readable. If the name represents something intricate or subtle, it's usually better to write a helpful doc comment than to attempt to put all the information into the name.

Finally, the convention in Go is to use `MixedCaps` or `mixedCaps` rather than underscores to write multiword names.

Documenting packages

Every package should have a *package comment*, a block comment preceding the **`package`** clause. For multi-file packages, the package comment only needs to be present in one file, and any one will do. The package comment should introduce the package and provide information relevant to the package as a whole. It will appear first on the godoc page and should set up the detailed documentation that follows. An example from the official *regexp* package:

This text is copied from [3].

```
/*
    The regexp package implements a simple library for
    regular expressions.

    The syntax of the regular expressions accepted is:

    regexp:
        concatenation '|' concatenation
*/
package regexp
```

Each defined (and exported) function should have a small line of text documenting the behavior of the function, from the *fmt* package:

```
// Printf formats according to a format specifier and writes to standard output.
func Printf(format string, a ...interface{}) (n int, errno os.Error) {
```

Testing packages

In Go it is customary to write (unit) tests for your package. Writing tests involves the *testing* package and the program `gotest`. Both have excellent documentation. When you include

tests with your package keep in mind that has to be build using the standard Makefile (see section "Building a package").

The testing itself is carried out with `gotest`. The `gotest` program run all the test functions. Without any defined tests for our *even* package a `gomake test`, yields:

```
% gomake test
no test files found (*_test.go)
make: *** [test] Error 2
```

Let us fix this by defining a test in a test file. Test files reside in the package directory and are named `*_test.go`. Those test files are just like other Go program, but `gotest` will only execute the test functions. Each test function has the same signature and its name should start with `Test`:

```
func TestXxx(t *testing.T)    ← Test<Capital>restOftheName
```

When writing test you will need to tell `gotest` that a test has failed or was successful. A successful test function just returns. When the test fails you can signal this with the following functions [6]. These are the most important ones (see `godoc testing` for more):

```
func (t *T) Fail()
```

`Fail` marks the test function as having failed but continues execution.

```
func (t *T) FailNow()
```

`FailNow` marks the test function as having failed and stops its execution. Execution will continue at the next test. So any other test in *this* file are skipped too.

```
func (t *T) Log(args ...interface{})
```

`Log` formats its arguments using default formatting, analogous to `Print()`, and records the text in the error log.

```
func (t *T) Fatal(args ...interface{})
```

`Fatal` is equivalent to `Log()` followed by `FailNow()`.

Putting all this together we can write our test. First we pick a name: `even_test.go`. Then we add the following contents:

Listing 4.4. Test file for even package

```
package even                                                    1

import "testing"                                                3

func TestEven(t *testing.T) {                                  5
    if true != Even(2) {                                        6
        t.Log("2 should be even!")                            7
        t.Fail()                                              8
    }                                                         9
}                                                            10
```

Note that we use `package even` on line 1, the tests fall in the same namespace as the package we are testing. This not only convenient, but also allows tests of unexported function and structures. We then import the *testing* package and on line 5 we define the only test function in this file. The displayed Go code should not hold any surprises: we check if the `Even` function works OK. Now, the moment we've been waiting for, executing the test:

```
% gomake test
6g -o _gotest_.6 even.go even_test.go
rm -f _test/even.a
gopack grc _test/even.a _gotest_.6
PASS
```

Our test ran and reported PASS. Success! To show how a failed test look we redefine our test function:

```
// Entering the twilight zone
func TestEven(t *testing.T) {
    if ! Even(2) {
        t.Log("2 should be odd!")
        t.Fail()
    }
}
```

We now get:

```
--- FAIL: even.TestEven
    2 should be odd!
FAIL
make: *** [test] Error 1
```

And you can act accordingly (by fixing the test for instance).

Writing new packages should go hand in hand with writing (some) documentation and test functions. It will make your code better and it shows that you really put in the effort.

Useful packages

The standard Go repository includes a huge number of packages and it is even possible to install more along side your current Go installation. It is very enlightening to browse the `$GOROOT/src/pkg` directory and look at the packages. We cannot comment on each package, but the following are worth a mention: ^a

fmt

Package *fmt* implements formatted I/O with functions analogous to C's `printf` and `scanf`. The format verbs are derived from C's but are simpler. Some verbs (%-sequences) that can be used:

`%v`

The value in a default format. when printing structs, the plus flag (`%+v`) adds field names;

`%#v`

a Go-syntax representation of the value.

`%T`

a Go-syntax representation of the type of the value;

^aThe descriptions are copied from the packages' godoc. Extra remarks are type set in italic.

io

This package provides basic interfaces to I/O primitives. Its primary job is to wrap existing implementations of such primitives, such as those in package `os`, into shared public interfaces that abstract the functionality, plus some other related primitives.

bufio

This package implements buffered I/O. It wraps an `io.Reader` or `io.Writer` object, creating another object (Reader or Writer) that also implements the interface but provides buffering and some help for textual I/O.

sort

The *sort* package provides primitives for sorting arrays and user-defined collections.

strconv

The *strconv* package implements conversions to and from string representations of basic data types.

os

The *os* package provides a platform-independent interface to operating system functionality. The design is Unix-like.

flag

The *flag* package implements command-line flag parsing. See “Command line arguments” on page 87.

json

The *json* package implements encoding and decoding of JSON objects as defined in RFC 4627.

template

Data-driven templates for generating textual output such as HTML.

Templates are executed by applying them to a data structure. Annotations in the template refer to elements of the data structure (typically a field of a struct or a key in a map) to control execution and derive values to be displayed. The template walks the structure as it executes and the “cursor” `@` represents the value at the current location in the structure.

http

The *http* package implements parsing of HTTP requests, replies, and URLs and provides an extensible HTTP server and a basic HTTP client.

unsafe

The *unsafe* package contains operations that step around the type safety of Go programs. *Normally you don't need this package.*

reflect

The *reflect* package implements run-time reflection, allowing a program to manipulate objects with arbitrary types. The typical use is to take a value with static type `interface{}` and extract its dynamic type information by calling `Typeof`, which returns an object with interface type `Type`. That contains a pointer to a struct of type `*StructType`, `*IntType`, etc. representing the details of the underlying type. A type switch or type assertion can reveal which. See chapter 6, section “Introspection”.

exec

The *exec* package runs external commands.

Exercises

Q16. (2) Stack as package

1. See the Q9 exercise. In this exercise we want to create a separate package for that code.
Create a proper package for your stack implementation, Push, Pop and the **Stack** type need to be exported.
2. Which official Go package could also be used for a (FIFO) stack?

Q17. (7) Calculator

1. Create a reverse polish calculator. Use your stack package.
2. Bonus. Rewrite your calculator to use the the package you found for question 2.

TODO

Write exercise for testing the stack package. (MG)

Answers

A16. (2) Stack as package

1. There are a few details that should be changed to make a proper package for our stack. First, the exported function should begin with a capital letter and so should **Stack**. So the full package (including the `String`) function becomes

Listing 4.5. Stack in a package

```
package stack

import (
    "strconv"
)

type Stack struct {
    i    int
    data [10]int
}

func (s *Stack) Push(k int) {
    s.data[s.i] = k
    s.i++
}

func (s *Stack) Pop() (ret int) {
    s.i--
    ret = s.data[s.i]
}

func (s *Stack) String() string {
    var str string
    for i := 0; i < s.i; i++ {
        str = str + "[" + strconv.Itoa(i) + ":"
            + strconv.Itoa(s.data[i]) + "]"
    }
    return str
}
```

A17. (7) Calculator

1. This is one answer

Listing 4.6. A (rpn) calculator

```
package main

import (
    "bufio"
    "os"
    "strconv"
    "fmt"
)
```



```

    )

    var reader *bufio.Reader = bufio.NewReader(os.Stdin)
    var st = new(Stack)

    type Stack struct {
        i    int
        data [10]int
    }

    func (s *Stack) push(k int) {
        if s.i+1 > 9 {
            return
        }
        s.data[s.i] = k
        s.i++
    }

    func (s *Stack) pop() (ret int) {
        s.i--
        if s.i < 0 {
            s.i = 0
            return 0
        }
        ret = s.data[s.i]
        return ret
    }

    func (s *Stack) String() string {
        var str string
        for i := 0; i < s.i; i++ {
            str = str + "[" +
                strconv.Itoa(i) + ":" + strconv.Itoa(s.data[i]) + "]"
        }
        return str
    }

    func main() {
        for {
            s, err := reader.ReadString('\n')
            var token string
            if err != nil {
                return
            }
            for _, c := range s {
                switch {
                case c >= '0' && c <= '9':
                    token = token + string(c)
                case c == ' ':
                    r, _ := strconv.Atoi(token)
                    st.push(r)
                    token = ""
                case c == '+':
                    fmt.Printf("%d\n", st.pop()+st.pop())
                case c == '*':
                    fmt.Printf("%d\n", st.pop()*st.pop())
                case c == '-':
                    p := st.pop()
                    q := st.pop()
                    fmt.Printf("%d\n", q-p)
                case c == 'q':
                    return
                default:
                    //error
                }
            }
        }
    }
}

```

2. The *container/vector* package would be a good candidate. It even comes with Push and Pop functions *predefined*. The changes needed to our program are *minimal* to say the least, the following unified diff shows the differences:

```
--- calc.go      2010-05-16 10:19:13.886855818 +0200
+++ calcvec.go   2010-05-16 10:13:35.000000000 +0200
@@ -5,11 +5,11 @@
     "os"
     "strconv"
     "fmt"
-    "./stack"
+    "container/vector"
 )

 var reader *bufio.Reader = bufio.NewReader(os.Stdin)
-var st = new(Stack)
+var st = new(vector.IntVector)

 func main() {
     for {
```

Only two lines need to be changed. *Nice.*

5

Beyond the basics

Go has pointers but not pointer arithmetic. You cannot use a pointer variable to walk through the bytes of a string.

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Go has pointers. There is however no pointer arithmetic, so they act more like references than pointers that you may know from C. Pointers are useful. Remember that when you call a function in Go, the variables are *pass-by-value*. So, for efficiency and the possibility to modify a passed value *in* functions we have pointers.

Just like in C you declare a pointer by prefixing the type with an `'*'`: `var p *int`. Now `p` is a pointer to an integer value. All newly declared variables are assigned their zero value and pointers are no difference. A newly declared, or just a pointer that points to nothing has a `nil`-value. In other languages this is often called a `NULL` pointer in Go it is just `nil`. To make a pointer point to something you can use the address-of operator (`&`), which we do on line 5:

Listing 5.1. Use of a pointer

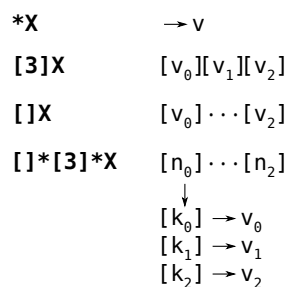
```
var p *int
fmt.Printf("%v", p)    ← Prints nil

var i int             ← Declare integer variable i
p = &i                ← Make p point to i

fmt.Printf("%v", p)    ← Prints something like 0x7ff96b81c000a
```

More general: `*X` is a pointer to an `X`; `[3]X` is an array of three `X`s. The types are therefore really easy to read just read out the names of the type modifiers: `[]` declares an array slice; `'*'` declares a pointer; `[size]` declares an array. So `[]* [3]*X` is an array slice of pointers to arrays of three pointers to `X`s (also see figure 5.1).

Figure 5.1. Pointers and types, the values `v` all have type `X`



Dereferencing a pointer is done by prefixing the pointer variable with `'*'`:

Listing 5.2. Dereferencing a pointer

```

p = &i           ← Take the address of i
*p = 8           ← Change the value of i
fmt.Printf("%v\n", *p) ← Prints 8
fmt.Printf("%v\n", i)  ← Idem

```

As said, there is no pointer arithmetic, so if you write: `*p++`, it is interpreted as `(*p)++`: first dereference and then increment the value.

Allocation

Go has garbage collection, meaning that you don't have to worry about memory allocation and deallocation. Of course almost every language since 1980 has this, but it is nice to see garbage collection in a C-like language.

Go has two allocation primitives, **new** and **make**. They do different things and apply to different types, which can be confusing, but the rules are simple. The following sections show how to handle allocation in Go and hopefully clarifies the somewhat artificial distinction between **new** and **make**.

Allocation with new

The built-in function **new** is essentially the same as its namesakes in other languages: `new(T)` allocates zeroed storage for a new item of type **T** and returns its address, a value of type `*T`. In Go terminology, it returns a pointer to a newly allocated zero value of type **T**. This is important to remember:

| **new** returns *pointers*.

This means a user of the data structure can create one with **new** and get right to work. For example, the documentation for `bytes.Buffer` states that "the zero value for `Buffer` is an empty buffer ready to use." Similarly, `sync.Mutex` does not have an explicit constructor or `Init` method. Instead, the zero value for a `sync.Mutex` is defined to be an unlocked mutex.

The zero-value-is-useful property works transitively. Consider this type declaration. See section "Defining your own types" on page 57.

```

type SyncedBuffer struct {
    lock    sync.Mutex
    buffer  bytes.Buffer
}

```

Values of type **SyncedBuffer** are also ready to use immediately upon allocation or just declaration. In this snippet, both `p` and `v` will work correctly without further arrangement.

```

p := new(SyncedBuffer) ← Type *SyncedBuffer
var v SyncedBuffer      ← Type SyncedBuffer

```

Allocation with make

Back to allocation. The built-in function `make(T, args)` serves a purpose different from `new(T)`. It creates slices, maps, and channels only, and it returns an initialized (not zero) value of type **T**, not `*T`. The reason for the distinction is that these three types are, under

the covers, references to data structures that must be initialized before use. A slice, for example, is a three-item descriptor containing a pointer to the data (inside an array), the length, and the capacity; until those items are initialized, the slice is `nil`. For slices, maps, and channels, `make` initializes the internal data structure and prepares the value for use.

| `make` returns initialized (non zero) *values*.

For instance, `make([]int, 10, 100)` allocates an array of 100 ints and then creates a slice structure with length 10 and a capacity of 100 pointing at the first 10 elements of the array. In contrast, `new([]int)` returns a pointer to a newly allocated, zeroed slice structure, that is, a pointer to a `nil` slice value.

These examples illustrate the difference between `new` and `make`.

```
var p *[]int = new([]int)      // allocates slice structure; *p == nil
                               // rarely useful
var v []int = make([]int, 100) // v refers to a new array of 100 ints

// Unnecessarily complex:
var p *[]int = new([]int)
*p = make([]int, 100, 100)

// Idiomatic:
v := make([]int, 100)
```

Remember that `make` applies only to maps, slices and channels and does not return a pointer. To obtain an explicit pointer allocate with `new`.

Constructors and composite literals

Sometimes the zero value isn't good enough and an initializing constructor is necessary, as in this example taken from the package `os`.

```
func NewFile(fd int, name string) *File {
    if fd < 0 {
        return nil
    }
    f := new(File)
    f.fd = fd
    f.name = name
    f.dirinfo = nil
    f.nepipe = 0
    return f
}
```

There's a lot of boiler plate in there. We can simplify it using a composite literal, which is an expression that creates a new instance each time it is evaluated.

```
func NewFile(fd int, name string) *File {
    if fd < 0 {
        return nil
    }
    f := File{fd, name, nil, 0}    ← Create a new File
    return &f                     ← Return the address of f
}
```

It is OK to return the address of a local variable; the storage associated with the variable survives after the function returns.

In fact, taking the address of a composite literal allocates a fresh instance each time it is evaluated, so we can combine these last two lines.^a

```
return &File{fd, name, nil, 0}
```

The fields of a composite literal are laid out in order and must all be present. However, by labeling the elements explicitly as field:value pairs, the initializers can appear in any order, with the missing ones left as their respective zero values. Thus we could say

```
return &File{fd: fd, name: name}
```

As a limiting case, if a composite literal contains no fields at all, it creates a zero value for the type. The expressions `new(File)` and `&File{}` are equivalent.

Composite literals can also be created for arrays, slices, and maps, with the field labels being indices or map keys as appropriate. In these examples, the initializations work regardless of the values of `Enone`, `Eio`, and `Einval`, as long as they are distinct.

```
ar := [...]string {Enone: "no error", Eio: "Eio", Einval: "invalid
argument"}
sl := []string {Enone: "no error", Eio: "Eio", Einval: "invalid
argument"}
ma := map[int]string{Enone: "no error", Eio: "Eio", Einval: "invalid
argument"}
```

Defining your own types

Of course Go allows you to define new types, it does this with the `type` keyword:

```
type foo int
```

Creates a new type `foo` which acts like an `int`. Creating more sophisticated types is done with the `struct` keyword. An example would be when we want record somebody's name (`string`) and age (`int`) in a single structure and make it a new type:

Listing 5.3. Structures

```
package main

import "fmt"

type NameAge struct {
    name string    ← Not exported
    age  int       ← Not exported, Age would be exported
}

func main() {
    a := new(NameAge)
    a.name = "Pete"
    a.age = 42
}
```

^aTaking the address of a composite literal tells the compiler to allocate it on the heap, not the stack.

```
        fmt.Printf("%v\n", a)
    }
```

Apropos, the output of `fmt.Printf("%v\n", a)` is

```
&{Pete 42}
```

That is nice! Go knows how to print your structure. If you only want to print one, or a few, fields of the structure you'll need to use `.<field name>`. For example to only print the name:

```
fmt.Printf("%s", a.name)    ← %s formats a string
```

More on structure fields

Each item in a structure is called a field. A struct with no fields:

```
struct {}
```

One with five fields:

```
struct {
    x, y int
    _ float64    ← padding
    A *[]int
    F func()
}
```

If you omit the name for a field, you create an anonymous field, for instance:

```
struct {
    T1          // field name is T1
    *T2         // field name is T2
    P.T3        // field name is T3
    x, y int    // field names are x and y
}
```

Note the field names that start with a capital letter are exported, i.e. can be set or read from other packages. Field names that start with a lowercase are private to the current package. The same goes for functions defined in packages, see chapter 4.

Methods

If you create functions that works on your newly defined type, you can take two routes:

1. Create a function that takes the type as an argument.

```
func doSomething(in1 *NameAge, in2 int) { /* ... */ }
```

This is (you might have guessed) a *function call*.

2. Create a function that works on the type (see *receiver* in listing 3.1):

```
func (in1 *NameAge) doSomething(in2 int) { /* ... */ }
```

This is a *method call*, which can be used as:

```
var n *NameAge
n.doSomething(2)
```

But keep the following in mind, this is quoted from [5]:

If x is addressable and $\&x$'s method set contains m , $x.m()$ is shorthand for $(\&x).m()$.

In the above case this means that the following is *not* an error:

```
var n NameAge           ← Not a pointer
n.doSomething(2)
```

Here Go will search the method list for `n` of type `NameAge`, come up empty and will then *also* search the method list for the type `*NameAge` and will translate this call to `(&n).doSomething(2)`.

Conversions

Sometimes you want to convert a type to another type. This is possible in Go, but there are some rules. For starters, converting from one value to another is done by functions and not all conversions are allowed.

Table 5.1. Valid conversions, `float64` works the same as `float32`

From	xb []byte	xi []int	s string	f float32	i int
To					
[]byte	×		[]byte(s)		
[]int		×	[]int(s)		
string	string(xb)	string(xi)	×		
float32				×	float32(i)
int				int(f)	×

- From a **string** to a slice of bytes or ints.

```
mystring := "hello this is string"
```

```
byteslice := []byte(mystring)
```

Converts to a **byte** slice, each **byte** contains the integer value of the corresponding byte in the string. Note that as strings in Go are encoded in UTF-8 some characters in the string may end up in 1, 2, 3 or 4 bytes.

```
intslice := []int(mystring)
```

Converts to an **int** slice, each **int** contains a Unicode code point. Every character from the string is corresponds to one integer.

- From a slice of bytes or ints to a **string**.


```

b := []byte{'h','e','l','l','o'}    ← Composite literal
s := string(b)
i := []int{257,1024,65}
r := string(i)

```

For numeric values the following conversion are defined:

- Convert to a integer with a specific (bit) length: `uint8(int)`;
- From floating point to an integer value: `int(float32)`. This discards the fraction part from the floating point value;
- The other way around: `float32(int)`;

User defined types and conversions

How can you convert between the types you have defined yourself? We create two types here `Foo` and `Bar`, where `Bar` is an alias for `Foo`:

```

type foo struct { int }    ← anonymous struct field
type bar foo               ← bar is an alias for foo

```

Then we:

```

var b bar = bar{1}        ← Declare b to be a bar
var f foo = b              ← Assign b to f

```

Which fails on the last line with:

cannot use b (type bar) as type foo in assignment

This can be fixed with a conversion:

```

var f foo = foo(b)

```

Note the converting structures that are not identical in their fields is more difficult. Also note that converting `b` to a plain `int` also fails; an integer is not the same as a structure containing an integer.

Exercises

Q18. (6) Map function with interfaces

1. Use the answer from exercise Q12, but now make it generic using interfaces.

Q19. (6) Pointers

1. Suppose we have defined the following structure:

```

type Person struct {
    name string
    age  int
}

```

What is the difference between the following two lines?

```

var p1 Person
p2 := new(Person)

```

2. What is the difference between the following two allocations?

```
func Set(t *T) {
    x = t
}
```

and

```
func Set(t T) {
    x = &t
}
```

Q20. (6) Linked List

1. Make use of the package *container/list* to create a (double) linked list. Push the values 1, 2 and 4 to the list and then print it.
2. Create your own linked list implementation. And perform the same actions as in question 1

Q21. (6) Cat

1. Write a program which mimics Unix program cat. For those who don't know this program, the following invocation displays the contents of the file *blah*:

```
% cat blah
```
2. Make it support the *n* flag, where each line is numbered.

Q22. (8) Method calls

1. Suppose we have the following program:

```
package main

import "container/vector"

func main() {
    k1 := vector.IntVector{}
    k2 := &vector.IntVector{}
    k3 := new(vector.IntVector)
    k1.Push(2)
    k2.Push(3)
    k3.Push(4)
}
```

What are the types of *k1*, *k2* and *k3*?

2. Now, this program compiles and runs OK. All the *Push* operations work even though the variables are of a different type. The documentation for *Push* says:

*func (p *IntVector) Push(x int) Push appends x to the end of the vector.*

So the receiver has to be of type ***IntVector**, why does the code above work then?

Answers

A18. (6) Map function with interfaces

Listing 5.4. A generic map function in Go

```
1. package main
   import "fmt"

   /* define the empty interface as a type */
   type e interface{}

   func mult2(f e) e {
       switch f.(type) {
       case int:
           return f.(int) * 2
       case string:
           return f.(string) + f.(string) + f.(string) + f.(string)
       }
       return f
   }

   func Map(n []e, f func(e) e) []e {
       m := make([]e, len(n))
       for k, v := range n {
           m[k] = f(v)
       }
       return m
   }

   func main() {
       m := []e{1, 2, 3, 4}
       s := []e{"a", "b", "c", "d"}
       mf := Map(m, mult2)
       sf := Map(s, mult2)
       fmt.Printf("%v\n", mf)
       fmt.Printf("%v\n", sf)
   }
```

A19. (6) Pointers

1. In first line: `var p1 Person` allocates a *Person-value* to `p1`. The type of `p1` is **Person**. The second line: `p2 := new(Person)` allocates memory and assigns a *pointer* to `p2`. The type of `p2` is ***Person**.
2. In the second function, `x` points to a new (heap-allocated) variable `t` which contains a copy of whatever the actual argument value is.
In the first function, `x` points to the same thing that `t` does, which is the same thing that the actual argument points to.

So in the second function, we have an "extra" variable containing a copy of the interesting value.

A20. (6) Linked List

- 1.
- 2.

A21. (6) Cat

1. The following is implementation of cat which also supports a n flag to number each line.

Listing 5.5. A cat program

```
package main

❶
import (
    "os"
    "fmt"
    "bufio"
    "flag"
)

var numberFlag = flag.Bool("n", false, "number each line") ❶

❷
func cat(r *bufio.Reader) {
    i := 1
    for {
        buf, e := r.ReadBytes('\n') ❸
        if e == os.EOF { ❹
            break
        }
        if *numberFlag { ❺
            fmt.Fprintf(os.Stdout, "%5d %s", i, buf)
            i++
        } else { ❻
            fmt.Fprintf(os.Stdout, "%s", buf)
        }
    }
    return
}

func main() {
    flag.Parse()
    if flag.NArg() == 0 {
        cat(bufio.NewReader(os.Stdin))
    }
}
```

```

    for i := 0; i < flag.NArg(); i++ {
        f, e := os.Open(flag.Arg(i), os.O_RDONLY, 0)
        if e != nil {
            fmt.Fprintf(os.Stderr, "%s: error reading from\n",
                os.Args[0], flag.Arg(i), e.String())
            continue
        }
        cat(bufio.NewReader(f))
    }
}

```

- ❶ Include all the packages we need;
- ❷ Define a new flag "n", which defaults to off. Note that we get the help for free;
- ❸ Start the function that actually reads the file's contents and displays it;
- ❹ Read one line at the time;
- ❺ Or stop if we hit the end;
- ❻ If we should number each line, print the line number and then the line itself;
- ❼ Otherwise we could just print the line.

A22. (8) Method calls

1. The type of `k1` is **vector.IntVector**. Why? We use a composite literal (the `{}`), so we get a value of that type back. The variable `k2` is of ***vector.IntVector**, because we take the address (`&`) of the composite literal. And finally `k3` has also the type ***vector.IntVector**, because `new` returns a pointer to the type.
2. The answer is given in [5] in the section "Calls", where among other things it says:

A method call `x.m()` is valid if the method set of (the type of) `x` contains `m` and the argument list can be assigned to the parameter list of `m`. If `x` is addressable and `&x`'s method set contains `m`, `x.m()` is shorthand for `(&x).m()`.

In other words because `k1` is addressable and ***vector.IntVector** *does* have the `Push` method, the call `k1.Push(2)` is translated by Go into `(&k1).Push(2)` which makes the type system happy again (and you too — now you know this).^b

^bAlso see section "Methods" in this chapter.

6

Interfaces

I have this phobia about having my body
penetrated surgically. You know what I mean?

eXistenZ
TED PIKUL

The following text is from [30]. Written by Ian Lance Taylor — one of the authors of Go.

In Go, the word *interface* is overloaded to mean several different things. Every type has an interface, which is the *set of methods defined* for that type. This bit of code defines a struct type **S** with one field, and defines two methods for **S**.

Listing 6.1. Defining a struct and methods on it

```
type S struct { i int }
func (p *S) Get() int { return p.i }
func (p *S) Put(v int) { p.i = v }
```

You can also define an interface type, which is simply a set of methods. This defines an interface **I** with two methods:

```
type I interface {
    Get() int
    Put(int)
}
```

| An interface type is a set of methods.

S is a valid *implementation* for interface **I**, because it defines the two methods which **I** requires. Note that this is true even though there is no explicit declaration that **S** implements **I**.

A Go program can use this fact via yet another meaning of interface, which is an interface value:

```
func f(p I) { ❶
    fmt.Println(p.Get()) ❷
    p.Put(1) ❸
}
```

❶ Declare a function that takes an interface type as the argument;

❷ As **p** implements interface **I** is *must* have the `Get()` method;

❸ Same holds for the `Put()` method.

Here the variable **p** holds a value of interface type. Because **S** implements **I**, we can call **f** passing in a pointer to a value of type **S**:

```
var s S; f(&s)
```

The reason we need to take the address of `s`, rather than a value of type `S`, is because we defined the methods on `s` to operate on pointers, see the code above in listing 6.1. This is not a requirement — we could have defined the methods to take values — but then the `Put` method would not work as expected.

The fact that you do not need to declare whether or not a type implements an interface means that Go implements a form of duck typing[34]. This is not pure duck typing, because when possible the Go compiler will statically check whether the type implements the interface. However, Go does have a purely dynamic aspect, in that you can convert from one interface type to another. In the general case, that conversion is checked at runtime. If the conversion is invalid — if the type of the value stored in the existing interface value does not satisfy the interface to which it is being converted — the program will fail with a runtime error.

Interfaces in Go are similar to ideas in several other programming languages: pure abstract virtual base classes in C++, typeclasses in Haskell or duck typing in Python. However there is no other language which combines interface values, static type checking, dynamic runtime conversion, and no requirement for explicitly declaring that a type satisfies an interface. The result in Go is powerful, flexible, efficient, and easy to write.

Which is what?

Lets define another type that also implements the interface `I`:

Listing 6.2. Another type that implements I

```
type R struct { i int }
func (p *R) Get() int { return p.i }
func (p *R) Put(v int) { p.i = v }
```

The function `f` can now accept variables of type `R` and `S`. Suppose you need to know the actual type in the function `f`. In Go you can figure that out by using a type switch.

```
func f(p I) {
    switch t := p.(type) { ❶
        case *S: ❷
        case *R: ❸
        case S: ❹
        case R: ❺
        default: ❻
    }
}
```

❶ The type switch. Use `(type)` in a `switch` statement. We store the type in the variable `t`;

❷ The actual type of `p` is a pointer to `S`;

❸ The actual type of `p` is a pointer to `R`;

❹ The actual type of `p` is a `S`;

❺ The actual type of `p` is a `R`;

⑤ It's another type that implements **I**.

Note that a type switch is the only way to figure out what type of an interface variable is. Using (**type**) outside a **switch** is illegal.

Empty interface

Since every type satisfies the empty interface: **interface{}**. We can create a generic function which has an empty interface as its argument:

Listing 6.3. A function with a empty interface argument

```
func g(any interface{}) int {
    return any.(I).Get()
}
```

The **return any.(I).Get()** is the tricky bit in this function. The value **any** has type **interface{}**, meaning no guarantee of any methods at all: it could contain any type. The **.(I)** is a type assertion which converts **any** to an interface of type **I**. If we have that type we can invoke the **Get()** function. So if we create a new variable of the type ***S**, we can just call **g()**, because ***S** also implements the empty interface.

```
s = new(S)
fmt.Println(g(s));
```

The call to **g** will work fine and will print 0. If we however invoke **g()** with a value that does not implement **I** we have a problem:

Listing 6.4. Failing to implement an interface

```
i := 5           ← Make i a "lousy" int
fmt.Println(g(i))
```

This compiles OK, but when we run this we get slammed with:

```
panic: interface conversion: int is not main.I: missing method Get
Which is completely true, the built-in type int does not have a Get() method.
```

Checking for interfaces

In your code you want to prevent these kind of errors, therefore Go provides you with a way to check if a variable implements a certain interface, again this uses a type assertion, but now in an **if** statement.

```
if ok := any.(I); ok {
    /* any implements the interface I */
}
```

Methods

Methods are functions that have a receiver (see chapter 3). You can define methods on any type (except on non-local types, this includes built-in types: the type **int** can not have methods). You can however make a new integer type with its own methods. For example:

```
type Foo int
```

```
func (self Foo) Emit() {
    fmt.Printf("%v", self)
}

type Emitter interface {
    Emit()
}
```

Doing this on non-local (types defined in other packages) types yields:

Listing 6.5. Failure extending built-in types

```
func (i int) Emit() {
    fmt.Printf("%d", i)
}
```

cannot define new methods
on non-local type int

Listing 6.6. Failure extending non-local types

```
func (a *net.AddrError) Emit() {
    fmt.Printf("%v", a)
}
```

cannot define new methods
on non-local type net.AddrError

Methods on interface types

An interface defines a set of methods. A method contains the actual code. In other words, an interface is the definition and the methods are the implementation. So a receiver can not be defined for interface types, doing so results in a `invalid receiver type ...` compiler error. The authoritative word from the language spec [5]:

*The receiver type must be of the form T or $*T$ where T is a type name. T is called the receiver base type or just base type. The base type must not be a pointer or interface type and must be declared in the same package as the method.*

Pointers to interfaces

Creating a pointer to an interface value is a useless action in Go. It is in fact illegal to create a pointer to an interface value. The release notes for that release that made them illegal leave no room for doubt:

The language change is that uses of pointers to interface values no longer automatically dereference the pointer. A pointer to an interface value is more often a beginner's bug than correct code.

From the [4]. If not for this restriction, this code:

```
var buf bytes.Buffer
io.Copy(buf, os.Stdin)
```

Would copy standard input into a copy of `buf`, not into `buf` itself. This is almost never the desired behavior.

Interface names

By convention, one-method interfaces are named by the method name plus the `-er` suffix: `Reader`, `Writer`, `Formatter` etc.

TODO
*Interfaces is not a
pointer, not reference
type*

Go release.2010-10-13.

There are a number of such names and it's productive to honor them and the function names they capture. Read, Write, Close, Flush, String and so on have canonical signatures and meanings. To avoid confusion, don't give your method one of those names unless it has the same signature and meaning. Conversely, if your type implements a method with the same meaning as a method on a well-known type, give it the same name and signature; call your string-converter method String not ToString.

Text copied from [3].

A sorting example

Recall the Bubblesort exercise (Q14), where we sorted an array of integers:

```
func bubblesort(n []int) {
    for i := 0; i < len(n)-1; i++ {
        for j := i + 1; j < len(n); j++ {
            if n[j] < n[i] {
                n[i], n[j] = n[j], n[i]
            }
        }
    }
}
```

A version that sorts strings is identical except for the signature of the function:

```
func bubblesortString(n []string) { /* ... */ }
```

Using this approach would lead to two functions, one for each type. By using interfaces we can make this more generic.

Lets create a new function that will sort both strings and integers, something along the lines of the non-working example:

```
func sort(i []interface{}) { ❶
    switch i.(type) {        ❷
        case string:
            // ...
        case int:
            // ...
    }
    ❸
}
```

- ❶ Our function will receive a slice of empty interfaces;
- ❷ Using a type switch we find out what the actual type is of the input;
- ❸ And then sort accordingly;
- ❹ Return the sorted slice.

But when we call this function with `sort([]int{1, 4, 5})`, it fails with: cannot use i (type []int) as type []interface in function argument

This is because Go can not easily convert to a *slice* of interfaces. Just converting to an interface is easy, but to a slice is much more costly. To keep a long story short: Go does not

The full mailing list discussion on this subject can be found at [25].

(implicitly) convert slices for you.

So what is the Go way of creating such a "generic" function? Instead of doing the type inference our selves with a type switch, we let Go do it implicitly: The following steps are required:

1. Define an interface type (called **Sorter** here) with a number of methods needed for sorting. We will at least need a function to get the length of the slice, a function to compare two values and a swap function;

```
type Sorter interface {
    Len() int           ← len() as a method
    Less(i, j int) bool ← p[j] < p[i] as a method
    Swap(i, j int)      ← p[i], p[j] = p[j], p[i] as a method
}
```

2. Define new types for the slices we want to sort. Note that we declare slice types;

```
type Xi []int
type Xs []string
```

3. Implementation of the methods of the **Sorter** interface. For integers:

```
func (p Xi) Len() int           { return len(p) }
func (p Xi) Less(i int, j int) bool { return p[j] < p[i] }
func (p Xi) Swap(i int, j int)   { p[i], p[j] = p[j], p[i] }
```

And for strings:

```
func (p Xs) Len() int           { return len(p) }
func (p Xs) Less(i int, j int) bool { return p[j] < p[i] }
func (p Xs) Swap(i int, j int)   { p[i], p[j] = p[j], p[i] }
```

4. Write a *generic* Sort function that works on the **Sorter** interface.

```
func Sort(x Sorter) { ❶
    for i := 0; i < x.Len() - 1; i++ { ❷
        for j := i + 1; j < x.Len(); j++ {
            if x.Less(i, j) {
                x.Swap(i, j)
            }
        }
    }
}
```

❶ x is now of the Sorter type;

❷ Using the defined functions, we implement Bubblesort.

We can now use you generic Sort function as follows:

```
ints := Xi{44, 67, 3, 17, 89, 10, 73, 9, 14, 8}
strings := Xs{"nut", "ape", "elephant", "zoo", "go"}
```

```
Sort(ints)
fmt.Printf("%v\n", ints)
Sort(strings)
fmt.Printf("%v\n", strings)
```

Introspection

Type assertion.

In a program, you can discover the dynamic type of an interface variable by using a **switch**. Such a type assertion uses the syntax of a type assertion with the keyword **type** inside the parentheses. If the switch declares a variable in the expression, the variable will have the corresponding type in each clause.

Listing 6.7. Dynamically find out the type

```
package main

type PersonAge struct { ❶
    name string
    age  int
}

type PersonShoe struct { ❶
    name      string
    shoesize  int
}

func main() {
    p1 := new(PersonAge)
    p2 := new(PersonShoe)
    WhichOne(p1)
    WhichOne(p2)
}

func WhichOne(x interface{}) { ❷
    switch t := x.(type) { ❸
    case *PersonAge: ❹
        println("Age person")
    case *PersonShoe:
        println("Shoe person")
    }
}
```

- ❶ First we define two structures as a new type, PersonAge;
- ❶ And PersonShoe;
- ❷ This function must accept *both* types as valid input, so we use the empty Interface, which every type implements;
- ❸ The type switch: (type);

- ④ When allocated with `new` it's a pointer. So we check for `*PersonAge`. If `WhichOne()` was called with a non pointer value, we should check for `PersonAge`.

The following is another example of performing a type switch, but this time checking for more (built-in) types:

Listing 6.8. A more generic type switch

```
switch t := interfaceValue.(type) {    ← The type switch
case bool:
    fmt.Printf("boolean %t\n", t)
case int:
    fmt.Printf("integer %d\n", t)
case *bool:
    fmt.Printf("pointer to boolean %t\n", *t)
case *int:
    fmt.Printf("pointer to integer %d\n", *t)
default:
    fmt.Printf("unexpected type %T", t) // %T prints type
}
```

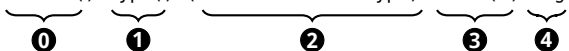
Introspection and reflection

In the following example we want to look at the "tag" (here named "namestr") defined in the type definition of `Person`. To do this we need the `reflect` package (there is no other way in Go). Keep in mind that looking at a tag means going back the *type* definition. So we use the `reflect` package to figure out the type of the variable and *then* access the tag.

Listing 6.9. Introspection using reflection

```
type Person struct {
    name string "namestr"    ← "namestr" is the tag
    age  int
}

p1 := new(Person)          ← new returns a pointer to Person
ShowTag(p1)                ← ShowTag() is now called with this pointer

func ShowTag(i interface{}) {
    switch t := reflect.NewValue(i).(type) {    ← Type assert on reflect value
    case *reflect.PtrValue:                    ← Hence the case for *reflect.PtrValue
        tag := t.Elem().Type().(*reflect.StructType).Field(0).Tag
        

```

- ① We are dealing with a `PtrValue` and according to the documentation^a:

```
func (v *PtrValue) Elem() Value
Elem returns the value that v points to. If v is a nil pointer, Elem returns
a nil Value.
```

^agodoc reflect

So on `t` we use `Elem()` to get the value the pointer points to.

- ❶ On `Value` we use the function `Type()` which returns `reflect.Type`. We need to get the type because there is where the tag is defined;
- ❷ So now we have a `reflect.Type`:

*...which returns an object with interface type `Type`. That contains a pointer to a struct of type `*StructType`, `*IntType`, etc. representing the details of the underlying type. A type switch or type assertion can reveal which.*

So we can access your specific type as a member of this struct. Which we do with `(*reflect.StructType)`;

- ❸ A `StructType` has a number of methods, one of which is `Field(n)` which returns the n^{th} field of a structure. The type of this return is a `StructField`;
- ❹ The struct `StructField` has a `Tag` member which returns the tag-name as a string. So on the 0^{th} field we can unleash `.Tag` to access this name: `Field(0).Tag`. This *finally* gives us `namestr`.

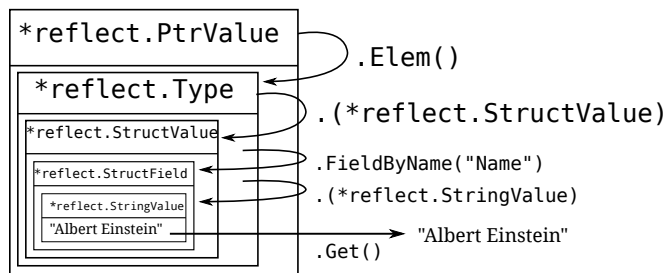
To make the difference between looking at types and values more clear, take a look at the following code:

Listing 6.10. Reflection and the type and value

```
func show(i interface{}) {
    switch t := i.(type) {
    case *Person:
        r := reflect.NewValue(i)      ← Enter the world of reflection
        tag := ❶
        r.(*reflect.PtrValue).Elem().Type().(*reflect.StructType).Field(0).tag
        nam := ❷
        r.(*reflect.PtrValue).Elem().(*reflect.StructValue).Field(0).(*reflect.StringValue).Get()
    }
}
```

- ❶ Here we want to get the "tag", which means going for the type. Thus we need `Elem().Type().(*reflect.StructType)` to get to it;
- ❷ Now we want to get access to the *value* of one of the members and we employ `Elem().(*reflect.StructValue)` to get to it. Now we have arrived at the structure. Then we go to the first field `Field(0)`, tell `reflect` is a `*reflect.StringValue` and invoke the `Get()` method on it.

Figure 6.1. Peeling away the layers using reflection. Going from a `*Person` via `*reflect.PtrValue` using the methods described in `godoc reflect` to get the actual `string` contained deep within.



Reflection works by peeling off layers once you have got your hands on a `Value` in the reflection world.

Setting a value works similarly as getting a value, but only works on *exported* members. Again some code:

Listing 6.11. Reflect with private member

```
type Person struct {
    name string "namestr"
    age  int
}

func Set(i interface{}) {
    switch t := i.(type) {
    case *Person:
        r := reflect.NewValue(i)
        r.(*reflect.PtrValue).Elem().
            (*reflect.StructValue).
            FieldByName("name").
            (*reflect.StringValue).
            Set("Albert Einstein")
    }
}
```

Listing 6.12. Reflect with public member

```
type Person struct {
    Name string "namestr" ←
    age  int
}

func Set(i interface{}) {
    switch t := i.(type) {
    case *Person:
        r := reflect.NewValue(i)
        r.(*reflect.PtrValue).Elem().
            (*reflect.StructValue).
            FieldByName("Name"). ←
            (*reflect.StringValue).
            Set("Albert Einstein")
    }
}
```

The code on the left compiles and runs, but when you run it, you are greeted with a stack trace and a *runtime* error:

```
panic: cannot set value obtained via unexported struct field
```

The code on the right works OK and sets the member `Name` to "Albert Einstein". Of course this only works when you call `Set()` with a pointer argument.

Exercises

Q23. (6) Interfaces and compilation

1. The code in listing 6.4 on page 68 compiles OK — as stated in the text. But when you run it you'll get a runtime error, so something is wrong. Why does the code

compile cleanly then?

Q24. (5) Pointers and reflection

1. One of the last paragraphs in section "Introspection and reflection" on page 73, has the following words:

The code on the right works OK and sets the member Name to "Albert Einstein". Of course this only works when you call Set() with a pointer argument.

Why is this the case?

Q25. (1) Interfaces and min-max

1. Make min-max generic so that it works for both integers and strings as in the example in section "A sorting example".

Answers

A23. (6) Interfaces and compilation

1. The code compiles because an integer type implements the empty interface and that is the check that happens at compile time.

A proper way to fix this, is to test if such an empty interface can be converted and if so, call the appropriate method. The Go code that defines the function `g` in listing 6.3 – repeated here:

```
func g(any interface{}) int { return any.(I).Get() }
```

Should be changed to become:

```
func g(any interface{}) int {  
    if v, ok := any.(I); ok { // Check if any can be converted  
        return v.Get()          // If so invoke Get()  
    }  
    return -1                  // Just so we return anything  
}
```

If `g()` is called now there are no run-time errors anymore. The idiom used is called "comma ok" in Go.

A24. (5) Pointers and reflection

1. When called with a non-pointer argument the variable is a copy (call-by-value). So you are doing the reflection voodoo on a copy. And thus you are *not* changing the original value, but only this copy.

A25. (1) Interfaces and min-max

- 1.

7

Concurrency

- Parallelism is about performance;
- Concurrency is about program design.

Google IO 2010

ROBE PIKE

In this chapter we will show off Go's ability for concurrent programming using channels and goroutines. Goroutines are the central entity in Go's ability for concurrency. But what is a goroutine? From [3]:

They're called goroutines because the existing terms — threads, coroutines, processes, and so on — convey inaccurate connotations. A goroutine has a simple model: it is a function executing in parallel with other goroutines in the same address space. It is lightweight, costing little more than the allocation of stack space. And the stacks start small, so they are cheap, and grow by allocating (and freeing) heap storage as required.

A goroutine is a normal function, except that you start it with the keyword **go**.

```
ready("Tee", 2)      ← Normal function call
go ready("Tee", 2)    ← ready() started as goroutine
```

The following idea for a program was taken from [27]. We run a function as two goroutines, the goroutines wait for an amount of time and then print something to the screen. On the lines 14 and 15 we start the goroutines. The `main` function waits long enough, so that both goroutines will have printed their text. Right now we wait for 5 seconds (`time.Sleep()` counts in ns) on line 17, but in fact we have no idea how long we should wait until all goroutines have exited.

Listing 7.1. Go routines in action

```
func ready(w string, sec int64) {           8
    time.Sleep(sec * 1e9)                   9
    fmt.Println(w, "is ready!")             10
}                                           11

func main() {                               13
    go ready("Tee", 2)                      14
    go ready("Coffee", 1)                   15
    fmt.Println("I'm waiting")             16
    time.Sleep(5 * 1e9)                    17
}                                           18
```

Listing 7.1 outputs:

```
I'm waiting      ← right away
Coffee is ready!  ← after 1 second
Tee is ready!     ← after 2 seconds
```

If we did not wait for the goroutines (i.e. remove line 17) the program would be terminated immediately and any running goroutines would *die with it*. To fix this we need some kind of mechanism which allows us to communicate with the goroutines. This mechanism is available to us in the form of channels. A channel can be compared to a two-way pipe in Unix shells: you can send to and receive values from it. Those values can only be of a specific type: the type of the channel. If we define a channel, we must also define the type of the values we can send on the channel. Note that we must use **make** to create a channel:

```
ci := make(chan int)
cs := make(chan string)
cf := make(chan interface{})
```

Makes *ci* a channel on which we can send and receive integers, makes *cs* a channel for strings and *cf* a channel for types that satisfy the empty interface. Sending on a channel and receiving from it, is done with the same operator: **<-**. Depending on the operands it figures out what to do:

```
ci <- 1      ← Send the integer 1 to the channel ci
<-ci        ← Receive an integer from the channel ci
i := <-ci    ← Receive from the channel ci and storing it in i
```

Lets put this to use.

Listing 7.2. Go routines and a channel

```
var c chan int ❶

func ready(w string, sec int) {
    time.Sleep(int64(sec) * 1e9)
    fmt.Println(w, "is ready!")
    c <- 1 ❷
}

func main() {
    c = make(chan int) ❸
    go ready("Tee", 2) ❹
    go ready("Coffee", 1)
    fmt.Println("I'm waiting, but not too long")
    <-c ❺
    <-c ❻
}
```

- ❶ Declare *c* to be a variable that is a channel of ints. That is: this channel can move integers. Note that this variable is global so that the goroutines have access to it;
- ❷ Send the integer 1 on the channel *c*;
- ❸ Initialize *c*;
- ❹ Start the goroutines with the keyword **go**;
- ❺ Wait until we receive a value from the channel. Note that the value we receive is discarded;

5 Two goroutines, two values to receive.

There is still some remaining ugliness; we have to read twice from the channel (lines 14 and 15). This is OK in this case, but what if we don't know how many goroutines we started? This is where another Go built-in comes in: **select**. With **select** you can (among other things) listen for incoming data on a channel.

Using **select** in our program does not really make it shorter, because we run too few goroutines. We remove the lines 14 and 15 and replace them with the following:

Listing 7.3. Using **select**

```

L: for {                                     14
    select {                                 15
        case <-c:                           16
            i++                             17
            if i > 1 {                       18
                break L                     19
            }                               20
    }                                       21
}                                         22

```

Make it run in parallel

While our goroutines were running concurrent, they were not running in parallel. When you do not tell Go anything there can only be one goroutine running at a time. With `runtime.GOMAXPROCS(n)` you can set the number of goroutines that can run in parallel. From the documentation:

GOMAXPROCS sets the maximum number of CPUs that can be executing simultaneously and returns the previous setting. If $n < 1$, it does not change the current setting. This call will go away when the scheduler improves.

If you do not want to change any source code you can also set an environment variable `GOMAXPROCS` to the desired value.

More on channels

When you create a channel in Go with `ch := make(chan bool)`, an unbuffered channel for bools is created. What does this mean for your program? For one, if you read (`value := <-ch`) it will block until there is data to receive. Secondly anything sending (`ch<-5`) will block until there is somebody to read it. Unbuffered channels make a perfect tool for synchronising multiple goroutines.

But Go allows you to specify the buffer size of a channel, which is quite simple how many elements a channel can hold. `ch := make(chan bool, 4)`, creates a buffered channels of bools that can hold 4 elements. The first 4 elements in this channels are written without any blocking. When you write the 5th element, your code *will* block, until another goroutine reads some elements from the channel to make room.

Although reads from channels block, you can perform a non-blocking read with the following syntax:

```
x, ok = <-ch
```

TODO
Still need to test this.

Where `ok` is set to `true` when there was something to read (otherwise it is `false`). And if that was the case `x` get the value read from the channel. In conclusion, the following is true in Go:

$$\text{ch} := \text{make}(\text{chan type}, \text{value}) \begin{cases} \text{value} == 0 & \rightarrow \text{unbuffered (blocking)} \\ \text{value} > 0 & \rightarrow \text{buffered (non-blocking) up to value elements} \end{cases}$$

Closing channels

Wrapping functions in goroutines

See `godns/resolver.go` and how to wrap the `Query()` function in a goroutine with error handling and such.

Exercises

Q26. (4) Channels

1. Modify the program you created in exercise Q2 to use channels, in other words, the function called in the body should now be a goroutine and communication should happen via channels. You should not worry yourself on how the goroutine terminates.
2. There are a few annoying issues left if you resolve question 1. One of the problems is that the goroutine isn't neatly cleaned up when `main.main()` exits. And worse, due to a race condition between the exit of `main.main()` and `main.shower()` not all numbers are printed. It should print up until 9, but sometimes it prints only to 8. Adding a second quit-channel you can remedy both issues. Do this.^a

Q27. (7) Fibonacci II

1. This is the same exercise as the one given page 33 in exercise 11. For completeness the complete question:

The Fibonacci sequence starts as follows: 1, 1, 2, 3, 5, 8, 13, ... Or in mathematical terms: $x_1 = 1; x_2 = 1; x_n = x_{n-1} + x_{n-2} \quad \forall n > 2$.

*Write a function that takes an **int** value and gives that many terms of the Fibonacci sequence.*

But now the twist: You must use channels.

^aYou will need the `select` statement.

Answers

A26. (4) Channels

1. A possible program is:

Listing 7.4. Channels in Go

```

package main                                     1

import "fmt"                                       3

func main() {                                     5
    ch := make(chan int)                          6
    go shower(ch)                                  7
    for i := 0; i < 10; i++ {                      8
        ch <- i                                    9
    }                                              10
}                                                  11

func shower(c chan int) {                         13
    for {                                         14
        j := <-c                                  15
        fmt.Printf("%d\n", j)                    16
    }                                             17
}                                                  18

```

We start off in the usual way, then at line 6 we create a new channel of ints. In the next line we fire off the function `shower` with the `ch` variable as its argument, so that we may communicate with it. Next we start our for-loop (lines 8-10) and in the loop we send (with `<-`) our number to the function (now a goroutine) `shower`. In the function `shower` we wait (as this blocks) until we receive a number (line 15). Any received number is printed (line 16) and then continue the endless loop started on line 14.

2. An answer is

Listing 7.5. Adding an extra quit channel

```

package main                                     1

import "fmt"                                       3

func main() {                                     5
    ch := make(chan int)                          6
    quit := make(chan bool)                       7
    go shower(ch, quit)                           8
    for i := 0; i < 10; i++ {                      9
        ch <- i                                    10
    }                                              11
    quit <- false // or true, does not matter    12
}

```



```

    }
    13

    func shower(c chan int, quit chan bool) {
    15
        for {
    16
            select {
    17
                case j := <-c:
    18
                    fmt.Printf("%d\n", j)
    19
                case <-quit:
    20
                    break
    21
            }
    22
        }
    23
    }
    24

```

On line 20 we read from the quit channel and we discard the value we read. We could have used `q := <-quit`, but then we would have used the variable only once — which is illegal in Go. Another trick you might have pulled out of your hat may be: `_ = <-quit`. This is valid in Go, but the Go idiom favors the one given on line 20.

A27. (7) Fibonacci II

1. The following program calculates the Fibonacci numbers using channels.

Listing 7.6. A Fibonacci function in Go

```

package main
import "fmt"

func dup3(in <-chan int) (<-chan int, <-chan int, <-chan int) {
    a, b, c := make(chan int, 2), make(chan int, 2), make(chan int, 2)
    go func() {
        for {
            x := <-in
            a <- x
            b <- x
            c <- x
        }
    }()
    return a, b, c
}

func fib() <-chan int {
    x := make(chan int, 2)
    a, b, out := dup3(x)
    go func() {
        x <- 0
        x <- 1
        <-a
        for {
            x <- <-a+<-b

```

```
        }
    }()
    return out
}

func main() {
    x := fib()
    for i := 0; i < 10; i++ {
        fmt.Println(<-x)
    }
}

// See sdh33b.blogspot.com/2009/12/fibonacci-in-go.html
```

8

Communication

Good communication is as stimulating as black coffee, and just as hard to sleep after.

ANNE MORROW LINDBERGH

In this chapter we are going to look at the building blocks in Go for communicating with the outside world.

Files

Reading from (and writing to) files is easy in Go. This program only uses the *os* package to read data from the file */etc/passwd*.

Listing 8.1. Reading from a file (unbufferd)

```
package main                                1

import "os"                                  3

func main() {                                5
    buf := make([]byte, 1024)                6
    f, _ := os.Open("/etc/passwd", os.O_RDONLY, 0666) 7
    defer f.Close()                          8
    for {                                    9
        n, _ := f.Read(buf)                  10
        if n == 0 { break }                  11
        os.Stdout.Write(buf[0:n])            12
    }                                         13
}                                             14
```

If you want to use buffered IO there is the *bufio* package:

Listing 8.2. Reading from a file (bufferd)

```
package main                                1

import ( "os"; "bufio" )                    3

func main() {                                5
    buf := make([]byte, 1024)                6
    f, _ := os.Open("/etc/passwd", os.O_RDONLY, 0666) 7
    defer f.Close()                          8
    r := bufio.NewReader(f)                  9
    w := bufio.NewWriter(os.Stdout)          10
    defer w.Flush()                          11
    for {                                    12
        n, _ := r.Read(buf)                  13
        if n == 0 { break }                  14
    }
```

```

        w.Write(buf[0:n])           15
    }                               16
}                                  17

```

On line 9 we create a **bufio.Reader** from **f** which is of type ***File**. **NewReader** expects an **io.Reader**, so you might think this will fail. But it doesn't. An **io.Reader** is defined as:

```

type Reader interface {
    Read(p []byte) (n int, err os.Error)
}

```

So *anything* that has such a **Read()** function implements this interface. And from listing 8.1 (line 10) we can see that ***File** indeed does so.

Command line arguments

Arguments from the command line are available inside your program via the string slice **os.Args**, provided you have imported the package **os**. The **flag** package has a more sophisticated interface, and also provides a way to parse flags. Take this example from a little DNS query tool:

```

var dnssec *bool = flag.Bool("dnssec", false, "Request DNSSEC records") ❶
var port *string = flag.String("port", "53", "Set the query port")       ❶
flag.Usage = func() { ❷
    fmt.Fprintf(os.Stderr, "Usage: %s [@server] [qtype] [qclass] [name\n    ...]\n", os.Args[0])
    flag.PrintDefaults() ❸
}
flag.Parse() ❹

```

- ❶ Define a bool flag, **-dnssec**. The variable must be a pointer otherwise the package can not set its value;
- ❶ Idem, but for a port option;
- ❷ Slightly redefine the **Usage** function, to be a little more verbose;
- ❸ For every flag given, **PrintDefaults** will output the help string;
- ❹ Parse the flags and fill the variables.

Executing commands

The **exec** package has function to run external commands, and it the premier way to execute commands from within a Go program. We start commands with the **Run** function:

```

func Run(argv0 string, argv, envv []string, dir string, stdin, stdout,
    stderr int) (p *Cmd, err os.Error)

```

*Run starts the binary program running with arguments **argv** and environment **envv**. It returns a pointer to a new **Cmd** representing the command or an error.*

Lets execute `ls -l`:

```
import "exec"

cmd, err := exec.Run("/bin/ls", []string{"ls", "-l"}, nil, "", exec.
    DevNull, exec.DevNull, exec.DevNull)
```

In the `os` package we find the `StartProcess` function. This is another way (but more low level) to start executables.^a The prototype for `StartProcess` is:

```
func StartProcess(name string, argv []string, envv []string, dir string,
    fd []*File) (pid int, err Error)
```

With the following documentation:

StartProcess starts a new process with the program, arguments, and environment specified by name, argv, and envv. The fd array specifies the file descriptors to be set up in the new process: fd[0] will be Unix file descriptor 0 (standard input), fd[1] descriptor 1, and so on. A nil entry will cause the child to have no open file descriptor with that index. If dir is not empty, the child chdirs into the directory before execing the program.

Suppose we want to execute `ls -l` again:

```
import "os"

pid, err := os.StartProcess("/bin/ls", []string{"ls", "-l"}, nil, "", []*
    os.File{ os.Stdin, os.Stdout, os.Stderr})
defer os.Wait(pid, os.WNOHANG)    ← Otherwise you create a zombie
```

Note that `os.Wait` (among other things) returns the exit code, with:

```
w := os.Wait(pid, os.WNOHANG)
e := w.WaitStatus.ExitStatus()    ← ExitStatus() returns an integer
```

Networking

All network related types and functions can be found in the package `net`. One of the most important functions in there is `Dial`. When you `Dial` into a remote system the function returns a `Conn` interface type, which can be used to send and receive information. The function `Dial` neatly abstracts away the network family and transport. So IPv4 or IPv6, TCP or UDP can all share a common interface.

Dialing a remote system (port 80) over TCP, then UDP and lastly TCP over IPv6 looks like this:^b

```
conn, err := Dial("tcp", "", "192.0.32.10:80")
conn, err := Dial("udp", "", "192.0.32.10:80")
conn, err := Dial("tcp", "", "[2620:0:2d0:200::10]:80")
    ← Brackets are mandatory
```

And with `conn` you can do read/write .

^aThere is talk on the go-nuts mailing list about separating `Fork` and `Exec`.

^bIn case you are wondering, 192.0.32.10 and 2620:0:2d0:200::10 are `www.example.org`.

TODO
dkls

TODO
Write echo server

Netchan: networking and channels

Exercises

Q28. (8) Processes

1. Write a program that takes a list of all running processes and prints how many child processes each parent has spawned. The output should look like:

```
Pid 0 has 2 children: [1 2]
Pid 490 has 2 children: [1199 26524]
Pid 1824 has 1 child: [7293]
```

- For acquiring the process list, you'll need to capture the output of `ps -e -opid,ppid,comm`. This output looks like:

```
PID  PPID  COMMAND
9024  9023  zsh
19560  9024  ps
```

- If a parent has one child you must print `child`, if there are more than one print `children`;
- The process list must be numerically sorted, so you start with pid 0 and work your way up.

Here is a Perl version to help you on your way (or to create complete and utter confusion).

Listing 8.3. Processes in Perl

```
#!/usr/bin/perl -l
my (%child, $pid, $parent);
my @ps=`ps -e -opid,ppid,comm`; # Capture the output from `ps`
foreach (@ps[1..$#ps]) {        # Discard the header line
    ($pid, $parent, undef) = split; # Split the line, discard 'comm'
    push @{$child{$parent}}, $pid; # Save the child PIDs on a list
}
# Walk through the sorted PPIDs
foreach (sort { $a <=> $b } keys %child) {
    print "Pid ", $_, " has ", @{$child{$_}}+0, " child", # Print
        them
        @{$child{$_}} == 1 ? " : " : "ren: ", "[@{$child{$_}}]";
}
}
```

Q29. (5) Word and letter count

1. Write a small program that reads text from standard input and performs the following actions:
 1. Count the number of characters (including spaces);
 2. Count the number of words;
 3. Count the numbers of lines.

In other words implement `wc(1)` (check your local manual page), however you only have to read from standard input.

Q30. (4) Uniq

1. Write a Go program that mimics the function of the Unix `uniq` command. This program should work as follows, given a list with the following items:

'a' 'b' 'a' 'a' 'a' 'c' 'd' 'e' 'f' 'g'

it should print only those item which don't have the same successor:

'a' 'b' 'a' 'c' 'd' 'e' 'f'

Listing 8.6 is a Perl implementation of the algorithm.

Listing 8.6. uniq(1) in Perl

```
#!/usr/bin/perl
my @a = qw/a b a a a c d e f g/;
print my $first = shift @a;
foreach (@a) {
    if ($first ne $_) { print; $first = $_; }
}
```

Q31. (9) Quine A *Quine* is a program that prints itself.

1. Write a Quine in Go.

Q32. (9) Number cruncher

- Pick six (6) random numbers from this list:

1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 25, 50, 75, 100

Numbers may be picked multiple times;

- Pick one (1) random number (i) in the range: 1 . . . 1000;
- Tell how, by combining the first 6 numbers (or a subset thereof) with the operators +, −, * and /, you can make i ;

An example. We have picked the numbers: 1, 6, 7, 8, 8 and 75. And i is 977. This can be done in many different ways, one way is:

$$((((1 * 6) * 8) + 75) * 8) - 7 = 977$$

or

$$(8 * (75 + (8 * 6))) - (7/1) = 977$$

1. Implement a number cruncher that works like that. Make it print the solution in a similar format (i.e. output should be infix with parenthesis) as used above.
2. Calculate *all* possible solutions and show them (or only show how many there are). In the example above there are 544 ways to do it.

Answers

A28. (8) Processes

1. There is lots of stuff to do here. We can divide our program up in the following sections:

1. Starting `ps` and capturing the output;
2. Parsing the output and saving the child PIDs for each PPID;
3. Sorting the PPID list;
4. Printing the sorted list to the screen

In the solution presented below, we've opted to use *container/vector* to hold the PIDs. This "list" grows automatically.

The function `atoi` (lines 19 through 22) is defined to get rid of the multiple return values of the original `strconv.Atoi`, so that it can be used inside function calls that only accept one argument, as we do on lines 45, 47 and 50.

A possible program is:

Listing 8.4. Processes in Go

```

package main                                     1

import (                                           3
    "os"                                           4
    "fmt"                                          5
    "sort"                                         6
    "bufio"                                       7
    "strings"                                     8
    "strconv"                                     9
    "container/vector"                            10
)                                                  11

const (                                           13
    PID = iota                                    14
    PPID                                           15
)                                                  16

func atoi(s string) (x int) {                    18
    x, _ = strconv.Atoi(s)                        19
    return                                         20
}                                                  21

func main() {                                     23
    pr, pw, _ := os.Pipe()                        24
    defer pr.Close()                              25
    r := bufio.NewReader(pr)                      26
    w := bufio.NewWriter(os.Stdout)              27
    defer w.Flush()                               28
    pid, _ := os.StartProcess("/bin/ps", []string{"ps", "-e", "- 29
        opid,ppid,comm"}, nil, "", []*os.File{nil, pw, nil})

```



```

defer os.Wait(pid, os.WNOHANG)      30
pw.Close()                          31

child := make(map[int]*vector.IntVector)  33
s, ok := r.ReadString('\n') // Discard the header line  34
s, ok = r.ReadString('\n')           35
for ok == nil {                      36
    f := strings.Fields(s)           37
    if _, present := child[atoi(f[PPID])]; !present {  38
        v := new(vector.IntVector)  39
        child[atoi(f[PPID])] = v  40
    }                                41
    // Save the child PIDs on a vector  42
    child[atoi(f[PPID])].Push(atoi(f[PID]))  43
    s, ok = r.ReadString('\n')       44
}                                    45

// Sort the PIDs                      47
schild := make([]int, len(child))      48
i := 0                                49
for k, _ := range child {             50
    schild[i] = k                     51
    i++                              52
}                                    53
sort.SortInts(schild)                54
// Walk through the sorted list       55
for _, ppid := range schild {         56
    fmt.Printf("Pid %d has %d child", ppid, child[ppid].  57
        Len())
    if child[ppid].Len() == 1 {       58
        fmt.Printf(": %v\n", []int(*child[ppid]))  59
    } else {                          60
        fmt.Printf("ren: %v\n", []int(*child[ppid]))  61
    }                                62
}                                    63
}                                    64

```

A29. (5) Word and letter count

1. The following program is an implementation of `wc(1)`.

Listing 8.5. wc(1) in Go

```

package main

import (
    "os"
    "fmt"
    "bufio"
    "strings"

```

```

)

func main() {
    var chars, words, lines int
    r := bufio.NewReader(os.Stdin) ❶
    for {
        switch s, ok := r.ReadString('\n'); true { ❶
        case ok != nil: ❷
            fmt.Printf("%d %d %d\n", chars, words, lines);
            return
        default: ❸
            chars += len(s)
            words += len(strings.Fields(s))
            lines++
        }
    }
}

```

- ❶ Start a new reader that reads from standard input;
- ❶ Read a line from the input;
- ❷ If we received an error, we assume it was because of a EOF. So we print the current values;
- ❸ Otherwise we count the charaters, words and increment the lines.

A30. (4) Uniq

1. The following is a uniq implementation in Go.

Listing 8.7. uniq(1) in Go

```

package main

import "fmt"

func main() {
    list := []string{"a", "b", "a", "a", "c", "d", "e", "f"}
    first := list[0]

    fmt.Printf("%s ", first)
    for _, v := range list[1:] {
        if first != v {
            fmt.Printf("%s ", v)
            first = v
        }
    }
}

```

A31. (9) Quine

1. The following Quine is from Russ Cox:

```
/* Go quine */
package main
import "fmt"
func main() {
    fmt.Printf("%s%C%s%C\n", q, 0x60, q, 0x60)
}
var q = `/* Go quine */
package main
import "fmt"
func main() {
    fmt.Printf("%s%C%s%C\n", q, 0x60, q, 0x60)
}
var q = `
```

A32. (9) Number cruncher

1. The following is one possibility. It uses recursion and backtracking to get an answer.

Listing 8.8. Number cruncher

```
package main

import (
    "fmt"
    "strconv"
    "container/vector"
    "flag"
)

const (
    _ = 1000 * iota
    ADD
    SUB
    MUL
    DIV
    MAXPOS = 11
)

var mop = map[int]string{
    ADD: "+",
    SUB: "-",
    MUL: "*",
    DIV: "/",
}

var (
    ok bool
    value int
)

type Stack struct {
    i int
    data [MAXPOS]int
}

func (s *Stack) Reset() {
    s.i = 0
}

func (s *Stack) Len() int {
    return s.i
}

func (s *Stack) Push(k int) {
    s.data[s.i] = k
    s.i++
}
```

```

func (s *Stack) Pop() int {
    s.i--
    return s.data[s.i]
}

var found int
var stack = new(Stack)

func main() {
    flag.Parse()
    list := []int{1, 6, 7, 8, 8, 75, ADD, SUB, MUL, DIV}
    // list := []int{1, 6, 7, ADD, SUB, MUL, DIV
    magic, ok := strconv.Atoi(flag.Arg(0))
    if ok != nil {
        return
    }
    f := make([]int, MAXPOS)
    solve(f, list, 0, magic)
}

func solve(form, numberop []int, index, magic int) {
    var tmp int
    for i, v := range numberop {
        if v == 0 {
            goto NEXT
        }

        if v < ADD {
            // it is a number, save it
            tmp = numberop[i]
            numberop[i] = 0
        }
        form[index] = v
        value, ok = rpncalc(form[0 : index+1])

        if ok && value == magic {
            if v < ADD {
                numberop[i] = tmp // reset and go on
            }
            found++
            fmt.Printf("%s = %d  %#d\n", rpnmstr(form[0:index+1]), value, found)
            //goto NEXT
        }

        if index == MAXPOS-1 {
            if v < ADD {
                numberop[i] = tmp // reset and go on
            }
            goto NEXT
        }
        solve(form, numberop, index+1, magic)
        if v < ADD {
            numberop[i] = tmp // reset and go on
        }
    }
    NEXT:
}

// convert rpn to nice infix notation and string
// the r must be valid rpn form
func rpnmstr(r []int) (ret string) {
    s := new(vector.StringVector)
    for k, t := range r {
        switch t {
        case ADD, SUB, MUL, DIV:
            a := s.Pop()
            b := s.Pop()
            if k == len(r)-1 {
                s.Push(b + mop[t] + a)
            } else {
                s.Push("(" + b + mop[t] + a + ")")
            }
        default:
            s.Push(strconv.Itoa(t))
        }
    }
    for _, v := range *s {
        ret += v
    }
    return
}

// return result from the rpn form.
// if the expression is not valid, ok is false

```

```

func rpnCalc(r []int) (int, bool) {
    stack.Reset()
    for _, t := range r {
        switch t {
            case ADD, SUB, MUL, DIV:
                if stack.Len() < 2 {
                    return 0, false
                }
                a := stack.Pop()
                b := stack.Pop()
                if t == ADD {
                    stack.Push(b + a)
                }
                if t == SUB {
                    // disallow negative subresults
                    if b-a < 0 {
                        return 0, false
                    }
                    stack.Push(b - a)
                }
                if t == MUL {
                    stack.Push(b * a)
                }
                if t == DIV {
                    if a == 0 {
                        return 0, false
                    }
                    // disallow fractions
                    if b%a != 0 {
                        return 0, false
                    }
                    stack.Push(b / a)
                }
            default:
                stack.Push(t)
        }
    }
    if stack.Len() == 1 { // there is only one!
        return stack.Pop(), true
    }
    return 0, false
}

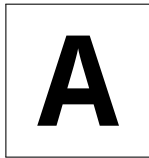
```

2. When starting permrec we give 977 as the first argument:

```

% ./permrec 977
1+((6+7)*75)+(8/8) = 977 #1
...
((75+(8*6))*8)-7 = 977 #542
(((75+(8*6))*8)-7)*1 = 977 #543
(((75+(8*6))*8)-7)/1 = 977 #544

```

Colophon

This work was created with \LaTeX . The main text is set in the Google Droid fonts. All type-writer text is typeset in DejaVu Mono.

Contributors

The following people have helped to make this book what it is today.

- Miek Gieben `<miek@miek.nl>;`
- JC van Winkel.

Help with proof reading, checking exercises and text improvements (no particular order): *Filip Zaludek, Jonathan Kans, Jaap Akkerhuis, Mayuresh Kathe, Makoto Inoue, Ben Bullock, Bob Cunningham, Dan Kortschak, Sonia Keys, Babu Sreekanth, Haiping Fan, Cecil New, Andrey Mirtchovski, and Russel Winder.*

Miek Gieben

Miek Gieben has a master's degree in Computer Science from the Radboud University Nijmegen (Netherlands).

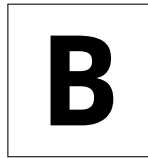
He is involved in the development and now the deployment of the DNSSEC protocol – the successor of the DNS and as such co-authored [22].

After playing with the language Erlang, Go was the first concurrent language that actually stuck with him.

He fills his spare time with coding in, and writing of Go. He is the maintainer of the Go DNS library: <https://github.com/miekg/godns>.

He maintains a personal blog on <http://www.miek.nl> and tweets under the name @miekg. The postings and tweets may sometimes actually have to do something with Go.

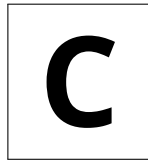




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