



THE UNIVERSITY OF ARIZONA
UASouth

CYBV 471 Assembly Programming for
Security Professionals
Week 9

Control Structures: Branching and Looping

Agenda



➤ **Control Structures**

- Compare instruction
- Unconditional jump
- Conditional jump
- Test operation with conditional jump
- if-then-else
- For loop instruction
- While loop instruction
- Do while loop
- Do while loop vs while loop
- Repeat string instructions (string encoding)

Compare Instruction: CMP



- CMP instruction subtracts one operand from another but doesn't store the result anywhere

```
cmp    a, b    ; compare a with b,  
               ; which computes a-b
```

- CMP instruction will modify the FLAGS register based on values of a and b
- Which FLAGS?

CMP With **Unsigned** Integers (**+ values**)



```
cmp    a, b    ; compare a with b,  
               ; which computes a-b
```

- CMP instruction will modify the values of ZF (Zero Flag) and CF (Carry Flag) based on value of operands a and b
 - ZF: set to 1 if result is 0
 - CF: set if we have overflow, have a carry, or borrow conditions

If $a = b$: ZF is set (=1), CF is not set (=0)

If $a > b$: ZF is not set (=0), CF is not set (=0)

If $a < b$: ZF is not set (=0), **CF** is set (borrow) (=1)



CMP With Signed Integers (-/+ values)

```
cmp    a, b    ; compare a with b,  
               ; which computes a-b
```

- CMP instruction will modify the values of **ZF** (Zero Flag), **OF** (Overflow Flag), and **SF** (Sign Flag) based on value of operands a and b
 - ZF: set to 1 if result is 0
 - OF: set to 1 if the result overflows or underflows
 - SF: set to 1 if the result is negative

If $a = b$: **ZF** is set (=1), **OF** is not set (=0), **SF** is not set (=0)

If $a > b$: **ZF** is not set (=0), and **OF = SF**

If $a < b$: **ZF** is not set (=0), and **OF \neq SF**

Understand the relationship between OF and SF Flags for Signed Integers



`cmp a, b ; compute (a - b)`

- If $a = b$: ZF is set ($=1$), OF is not set ($=0$), SF is not set ($=0$)
- If $a > b$: ZF is not set ($=0$), and $OF = SF$
 - If there is no overflow ($OF=0$), the (correct) result is positive ($SF=0$)
 - If there is an overflow ($OF=1$), the (incorrect) result is negative ($SF=1$)
 - Therefore, in both cases $SF = OF$
- If $a < b$: ZF is not set ($=0$), and $OF \neq SF$
 - If there is no overflow ($OF=0$), the (correct) result is negative ($SF=1$)
 - If there is overflow ($OF=1$), the (incorrect) result is positive ($SF=0$)
 - The correct answer should be negative.
 - Therefore, in both cases $SF \neq OF$

Examples: Values of OF and SF Flags for Signed Integers



Example: Assume $a = 80h$ (-128d)?

$80h = 1000\ 0000 \rightarrow$ 2nd complement \rightarrow value = (-128d)

$b = 23h > 0010\ 0011 \rightarrow$ value = 35d

(a < b, one byte)

$a - b = a + (-b) = 80h + (2's\ complement\ of\ b) = 80h + DDh (=15Dh)$

$15Dh = 0001\ 0101\ 1101.$

$15Dh >$ for two bytes, we have $0101\ 1101$ (Drop 1 because we have one byte) = 5Dh, (SF=0)

-128d - 35d should be negative.

Result is $0101\ 1101$ which is positive. Therefore, **OF=1 (incorrect)** (SF \neq OF)

Example: Assume $a = F3h$ (-13d), $b = 23h$ (+35d)

(a < b, one byte)

$a - b = a + (-b) = F3h + (2's\ complement\ of\ b) = F3h + DDh = D0h$

$D0h = 1101\ 0000 \rightarrow$ value?? (= - 2's complement of D0h) = - 48d (correct)

D0h is negative and we have no overflow (in range). So, SF=1 and **OF=0** (SF \neq OF)

Examples: Values of “OF” and “SF” Flags for Signed Integers



Example: Assume : $a = \text{F3h}$ (-13d), $b = \text{82h}$ (-126d)

($a > b$, one byte)

$$a - b = a + (-b) = \text{F3h} + \text{7Eh} = \text{171h}$$

We have overflow condition:

$\text{171h} = 0001\ 0111\ 0001$. We have only one byte.

The value is $0111\ 0001 = 71\text{h} = (+113\text{d})$ (correct answer)

The result is correct answer. Therefore $\text{OF} = 0$

71 is positive ($\text{SF} = 0$). Therefore, in both cases $\text{SF} = \text{OF}$

Example: $a = 70\text{h}$ (112d), $b = \text{D8h}$ (-40d)

($a > b$, one byte)

$$a - b = a + (-b) = 70\text{h} + 28\text{h} = 98\text{h}, \quad (152 \text{ is } > 127, \text{ overflow})$$

$98\text{h} = 1001\ 1000 = \text{negative}$. (2^{nd} complement value) = -104d (incorrect)

Therefore $\text{OF} = 1$

$98\text{h} = 1001\ 1000 \rightarrow \text{SF} = 1$

So, $\text{SF} = 1$ and $\text{OF} = 1$

Summary: CMP for Unsigned and Signed Integers



	cmp a,b	ZF	CF	OF	SF
unsigned	a = b	1	0		
	a < b	0	1		
	a > b	0	0		
signed	a = b	1		0	0
	a < b	0		v v (0 or 1)	!v
	a > b	0		v	v

Jumps and Branches



- In the normal program flow, we execute instructions one by one (next instruction fashion)
- Assembly language has instructions that allow you to modify the order in which instructions are executed.
 - i.e., we not always execute the next instruction
 - Unconditional jumps
 - Conditional jumps

Unconditional Jump

The “JMP” Instruction



- JMP instruction allows you to “jump” to a **code label**

- Example:

instruction-1
instruction-2

jmp label-1 ; then jump to Label-1
instruction ; will never be executed
instruction ; will never be executed

...

Label-1:

instruction-1
instruction-2

...

- JMP instruction will change the content of the EIP register (program counter) to execute sub1 label
- The JMP instruction is an **unconditional branch (jump)**
- The content of the FLAGS register doesn't affect “JMP” instruction

Conditional Jump Instructions



- These instructions jump to an address in the code segment (i.e., a label) based on the content of the **FLAGS** register
- **Conditional jumps use the flags** to determine whether to jump or to proceed to the next instruction.
- **FLAGS** register is updated by
 - All instructions that perform arithmetic operations
 - The **CMP** (compare) instruction
 - **CMP** that sets some flags in the **FLAEG** register
- Example:

```
instruction-1
instruction-2
Arithmetic instruction or CMP instruction ; may modify FLAGS register
Conditional jump label-1 ; jump if condition is satisfied
instruction ; (no jump) continue from here if condition is not satisfied
instruction
...
```

label1:

```
instruction-1
instruction-2
```

Conditional Jump Instructions



- There is a large set of conditional jump instructions that act based on bits in the FLAGS register
- The simple jump instructions jumps (or not) depending on the value of **one** of the following flags:
 - ZF, OF, SF, CF, PF
 - PF: Parity Flag
 - Set to 0 if there is odd number of ones in the lower 8-bit of the “result”
 - Set to 1 if there is even number of ones in the lower 8-bit of the “result”
 - Examples: Result = 11011111 10001000 (we have 2 ones, PF=1)
 - Result = 11011111 10001100 (we have 3 ones, PF=0)

Conditional Jump Instructions



- The simple jump instructions jumps (or not) depending on the value of **one** of the following flags: ZF, OF, SF, CF, PF

JZ jumps (branches) if ZF is set (1)

JNZ branches if ZF is unset

JO branches if OF is set

JNO branches if OF is unset

JS branches if SF is set

JNS branches if SF is unset

JC branches if CF is set

JNC branches if CF is unset

JP branches if PF is set

JNP branches if PF is unset

Conditional Jump Based on Unsigned Comparison



Instruction	branches if
JE x, y	$x = y$
JNE	$x \neq y$
JB, JNAE	$x < y$
JBE, JNA	$x \leq y$
JA, JNBE	$x > y$
JAE, JNB	$x \geq y$

Mnemonic	Description
JA	Jump if above (if $leftOp > rightOp$)
JNBE	Jump if not below or equal (same as JA)
JAE	Jump if above or equal (if $leftOp \geq rightOp$)
JNB	Jump if not below (same as JAE)
JB	Jump if below (if $leftOp < rightOp$)
JNAE	Jump if not above or equal (same as JB)
JBE	Jump if below or equal (if $leftOp \leq rightOp$)
JNA	Jump if not above (same as JBE)

Conditional Jump Based on Signed Comparison



Instruction	branches if
JE	$x = y$
JNE	$x \neq y$
JL, JNGE	$x < y$
JLE, JNG	$x \leq y$
JG, JNLE	$x > y$
JGE, JNL	$x \geq y$

Mnemonic	Description
JG	Jump if greater (if $leftOp > rightOp$)
JNLE	Jump if not less than or equal (same as JG)
JGE	Jump if greater than or equal (if $leftOp \geq rightOp$)
JNL	Jump if not less (same as JGE)
JL	Jump if less (if $leftOp < rightOp$)
JNGE	Jump if not greater than or equal (same as JL)
JLE	Jump if less than or equal (if $leftOp \leq rightOp$)
JNG	Jump if not greater (same as JLE)

Conditional Jump for Unsigned/Signed Comparison



cmp x, y			
signed		Unsigned	
Instruction	branches if	Instruction	branches if
JE	$x = y$	JE	$x = y$
JNE	$x \neq y$	JNE	$x \neq y$
JL, JNGE	$x < y$	JB, JNAE	$x < y$
JLE, JNG	$x \leq y$	JBE, JNA	$x \leq y$
JG, JNLE	$x > y$	JA, JNBE	$x > y$
JGE, JNL	$x \geq y$	JAE, JNB	$x \geq y$

Summary Conditional Jump Instructions



- Here are the most common conditional jump instructions

Mnemonic	Description	Flags
JZ	Jump if zero	ZF = 1
JNZ	Jump if not zero	ZF = 0
JC	Jump if carry	CF = 1
JNC	Jump if not carry	CF = 0
JO	Jump if overflow	OF = 1
JNO	Jump if not overflow	OF = 0
JS	Jump if signed	SF = 1
JNS	Jump if not signed	SF = 0
JP	Jump if parity (even)	PF = 1
JNP	Jump if not parity (odd)	PF = 0

Test operation with conditional jump

- Test instruction performs a nondestructive **AND operation** between each pair of matching **bits** in two operands
- No operands are modified, but the Zero flag (**ZF**) is affected.
- Example-1: jump to a label-1 if **either** bit 0 or bit 1 in AL is set (1)

```
test al,00000011b    ; zf = 1  if al has xxxxxx00
                        ; zf = 0 if al has xxxxxx01 (or 10 or 11)
jnz label-1           ; jump if zf is not set (zf=0)
```

Example-2: jump to a label-1 if neither bit 0 **nor** bit 1 in AL is set (1)

```
test al,00000011b    ; zf = 1 if al has xxxxxx00
                        ; zf = 0 if al has xxxxxx01 (or 10 or 11)
jz  label-1           ; jump if zf = 1
                        ; no jump if zf is 0
```

Control Structures: if-then-else



A generic if-then-else construct:

```
if (condition) then
{
    execute block-code-1 ; if condition is true
}
else
{
    execute else-block-code-2 ; if condition is false
}
next instruction
```

Translation into x86 assembly:

```
; instructions to set flags (e.g., cmp ...)
jyy else-block ; jyy branches to else-block if-condition is false
```

if-block:

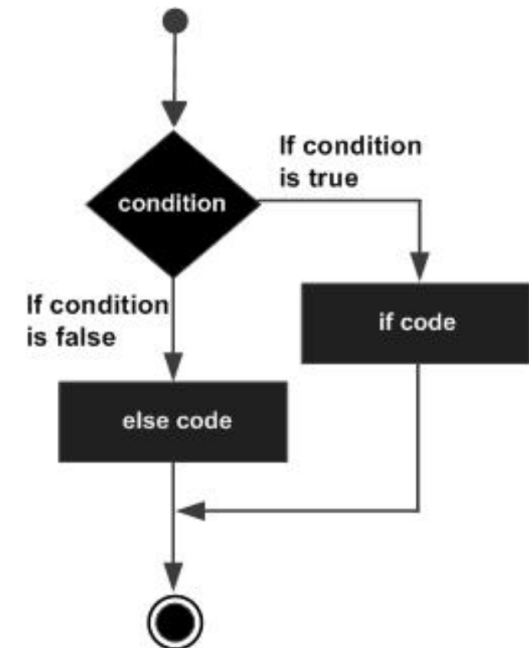
```
; code of block-code-1
jmp endif
```

else-block:

```
; code of block-code2
```

endif:

```
next instruction
```



Control Structures: if-then-else



Example:

```
if (val == val2)
```

```
    x = 1;
```

```
else
```

```
    x = 2
```

```
Next instruction
```

Translation into x86 assembly:

```
mov eax, var1
```

```
cmp eax, var2
```

```
jne L1      ; condition false
```

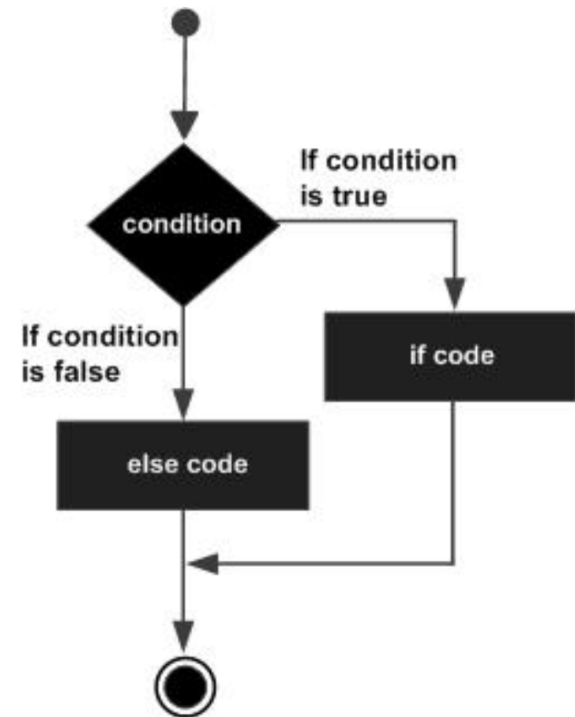
```
mov x, 1    ; condition true
```

```
jmp L2      ; end of if-else
```

```
L1: mov x, 2
```

```
L2:
```

```
Next instruction
```



Control Structures: if-then (no else)



A generic if-then-else construct:

```
if (condition) then
{
  execute block-code-1 ; if condition is true
}
next instruction
```

Translation into x86 assembly:

```
; instructions to set flags (e.g., cmp ...)
```

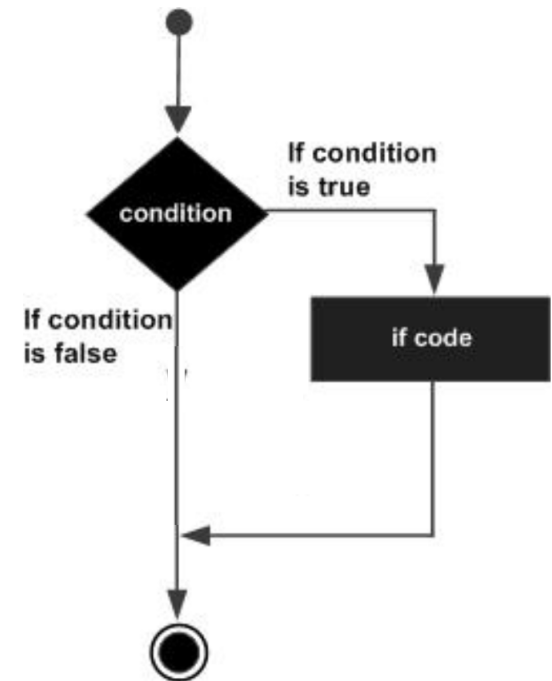
```
jxx endif ; jxx branches to endif if condition is false
```

if-block:

```
; code of block-code-1
```

endif:

```
next instruction
```



Control Structures: if-then (no else)



Example:

```
if (val == val2)
```

```
    x = 1;
```

```
Next instruction
```

Translation into x86 assembly:

```
    mov eax, var1
```

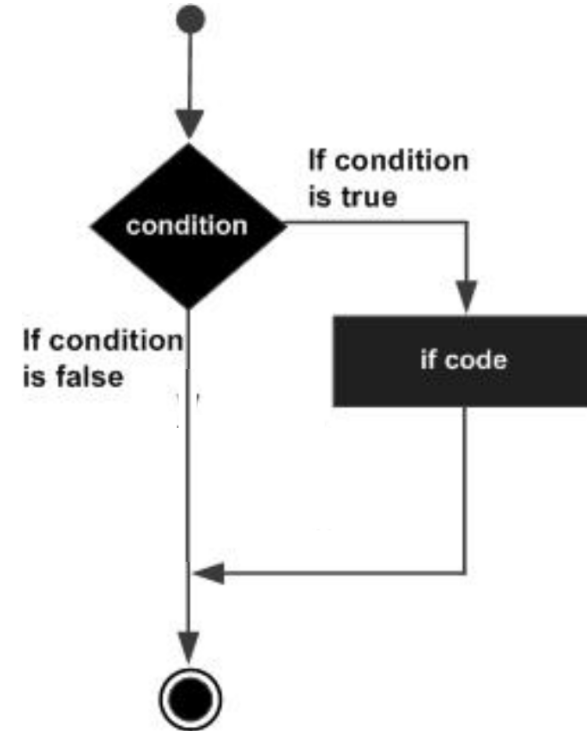
```
    cmp eax, var2
```

```
    jne endif      ; condition false
```

```
L1: mov x, 1
```

```
endif:
```

```
Next instruction
```





For Loop

Assume we have would like to implement the following c-code

```
sum = 0
for (i = 0; i <=10 ; i++)
    sum = sum + i
```

Translation into x86 assembly:

```
mov eax, 0      ; store sum value in eax
mov ecx, 0      ; store i value in ecx
```

loop_start:

```
cmp ecx, 10     ; compare i and 10
jg loop_end     ; if (i>10) go loop_end
add eax, ecx     ; if i < 10, sum = sum + i
inc ecx         ; i = i + 1
jmp loop_start  ; start over, go to loop_start
```

loop_end:

For Loop



Assume we have would like to implement the following c-code

```
sum = 0
for (i = 0; i <=10 ; i++)
{
    sum = sum + i
    several code lines
}
```

Translation into x86 assembly:

```
mov eax, 0 ; store sum value in eax
mov ecx, 0 ; store i ( loop index) value in ecx
; use other registers for other variables as needed
```

loop1_start:

```
cmp ebx, 10      ; compare i and 10;
jg loop1_end     ; jump greater, if (I >10) go loop_end
add eax, ebx     ; if i <= 10, sum = sum + i
inc ecx          ; i = i + 1
several code lines ; use other registers
jmp loop1_start  ; start over, goto loop_start
```

loop1_end:



For Loop

The previous loop can be written **in a reverse**

```
sum = 0
for (i = 10; i >= 0 ; i--)
    sum = sum + i
```

Translation into x86 assembly:

```
mov eax, 0 ; store sum value in eax
mov ecx, 10 ; store i (loop index) value in ecx
```

loop1_start:

```
cmp ecx, 0          ; compare i and 0  (LOOP CONDITION)
jnl loop1_end       ; if (i < 0) go loop_end (CHECK CONDITION)
add eax, ecx        ; if i > 0, sum = sum + i
dec ecx             ; i = i - 1  (LOOP CONTRO)
jmp loop1_start      ; start over, goto loop_start
```

loop1_end:

```
; END
```

Using the `loop` Instruction



The previous loop can be written in a reverse

```
sum = 0
for (i = 10; i > 10 ; i--)
    sum = sum + i
    other codes
```

Translation into x86 assembly:

```
mov eax, 0 ; store sum value in eax
mov ecx, 10 ; store i (loop index) value in ecx
```

loop1_start:

```
add eax, ecx
    other codes
loop loop1_end: ; check the required conditions and decrement the index
```

NOTE-1: Store loop index in ecx

Make sure you have the reversed loop (decrement i)



Do While loop

A generic do while loop construct:

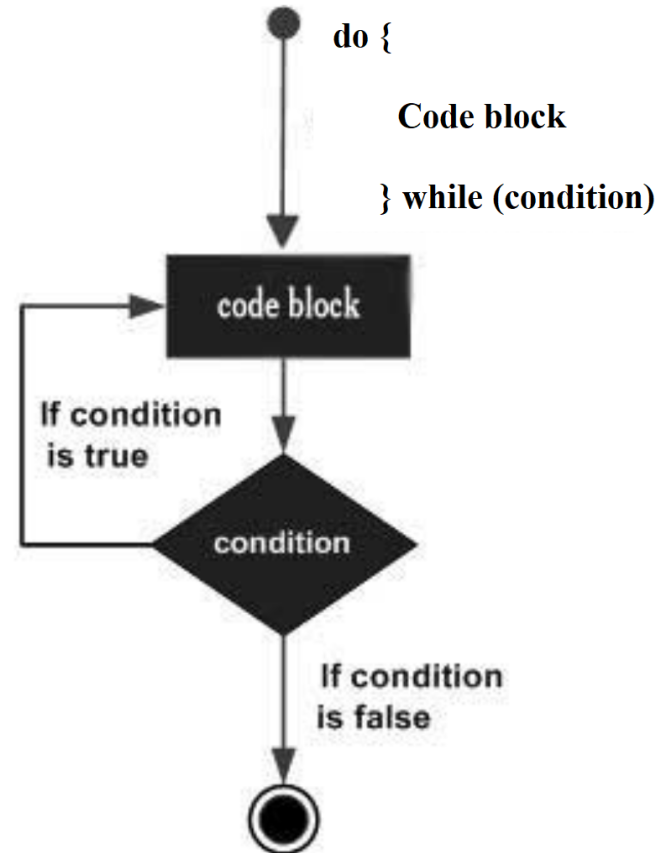
```
do
{
    execute block-code ; if condition is true
} while (condition)

    next instruction
```

Translation into x86 assembly:

```
do:
    block-code
    ; instructions to set flags (e.g., cmp ...)
    jxx do ; branches if condition is true

    next instruction
```





Do While loop

A generic do while loop construct:

```
do
{
    execute block-code      ; if condition is true
} while (condition)

next instruction
```

The block-code will be executed at least one time

The code will be repeated if the condition is true



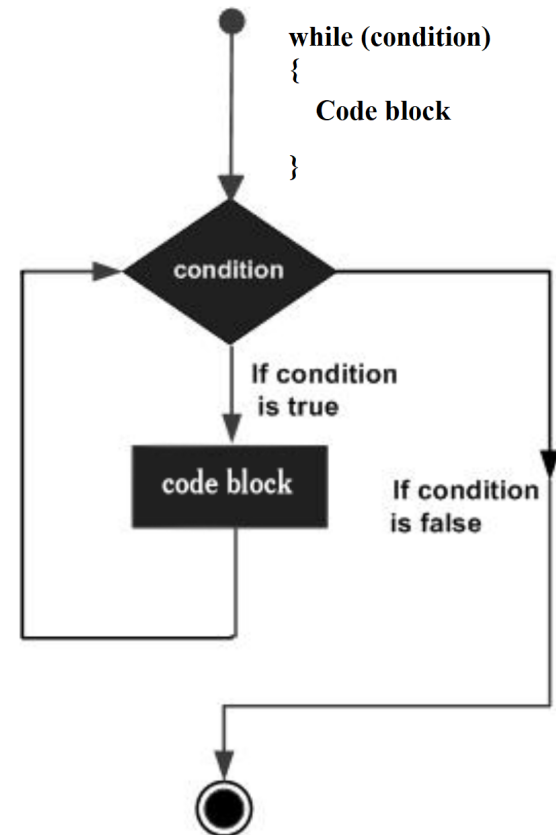
While loop

A generic while loop construct:

```
while (condition)
{
    execute block-code ; if condition is true
}
next instruction
```

Translation into x86 assembly:

```
while:
    ; instructions to set flags (e.g., cmp ...)
    jxx end_while ; branches if condition is false
    block-code
    jmp while
end_while:
    next instruction
```



Do While loop vs While loop

A generic do while loop construct:

```
do
{
    execute block-code      ; if condition is true
} while (condition)

next instruction
```

The block-code will be executed at least one time

The code will be repeated if the condition is true

```
while (condition)
{
    execute block-code      ; if condition is true
}
```

The block-code will be executed **only** if condition is true



Putting It All Together

You should know:

➤ **Control Structures**

- Compare instruction
- Unconditional jump
- Conditional jump
- Test operation with conditional jump
- if-then-else
- Loop instructions
- Repeat (while)string instructions



Questions?

Coming Next Week
C-Stream I/O

Week 9 Assignments



- **Learning Materials**

- 1- Week 9 Presentation

- 2- Read pages 298-303 (ch. 10): Duntermann, Jeff. Assembly Language Step by Step, Programming with Linux

- **Assignment**

- 1- Complete “Lab 9” by coming Sunday 11:59 PM.