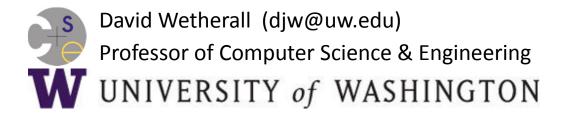
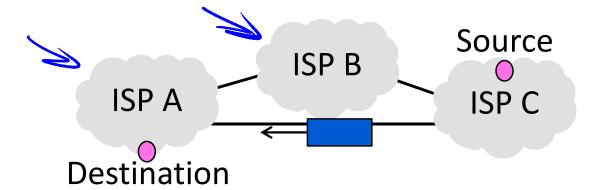
Introduction to Computer Networks

Routing with Policy (BGP) (§5.6.7)



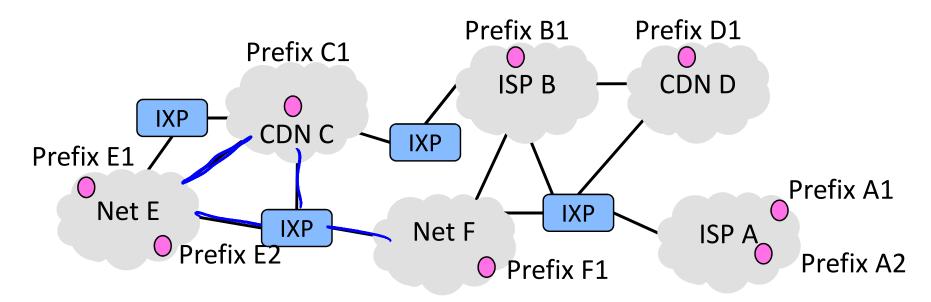
Topic

- How to route with multiple parties,
 each with their own routing policies
 - This is Internet-wide BGP routing



Structure of the Internet

- Networks (ISPs, CDNs, etc.) group hosts as IP prefixes
- Networks are richly interconnected, often using IXPs



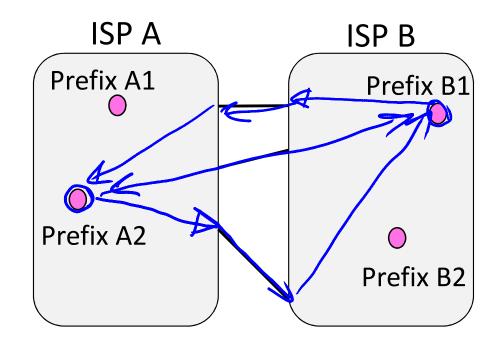
Internet-wide Routing Issues

- Two problems beyond routing within an individual network
- 1. Scaling to very large networks
 - Techniques of IP prefixes, hierarchy, prefix aggregation
- 2. Incorporating policy decisions
 - Letting different parties choose their routes to suit their own needs

Yikes!

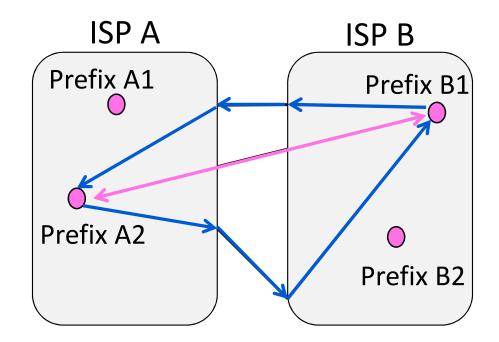
Effects of Independent Parties

- Each party selects routes to suit its own interests
 - e.g, shortest path in ISP
- What path will be chosen for A2→B1 and B1→A2?
 - What is the best path?



Effects of Independent Parties (2)

- Selected paths are longer than overall shortest path
 - And symmetric too!
- This is a consequence of independent goals and decisions, not hierarchy

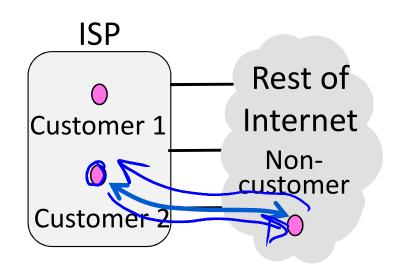


Routing Policies

- Capture the goals of different parties – could be anything
 - E.g., Internet2 only carries non-commercial traffic
- Common policies we'll look at:
 - ISPs give TRANSIT service to customers
 - ISPs give PEER service to each other

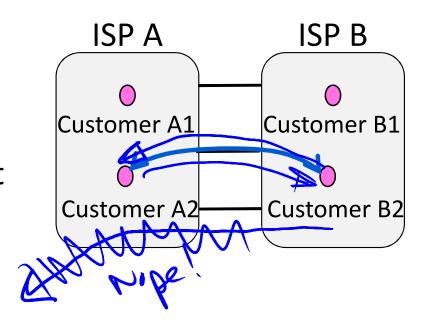
Routing Policies – Transit

- One party (customer) gets TRANSIT service from another party (ISP)
 - ISP accepts traffic for customer from the rest of Internet
 - ISP sends traffic from customer to the rest of Internet
 - Customer pays ISP for the privilege



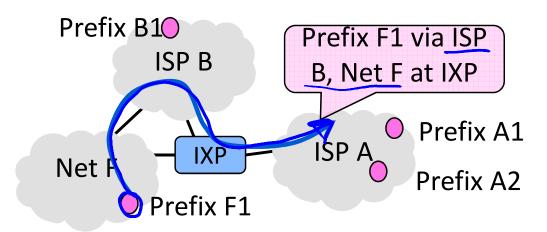
Routing Policies – Peer

- Both party (ISPs in example) get
 PEER service from each other
 - Each ISP accepts traffic from the other ISP only for their customers
 - ISPs do not carry traffic to the rest of the Internet for each other
 - ISPs don't pay each other



Routing with BGP (Border Gateway Protocol)

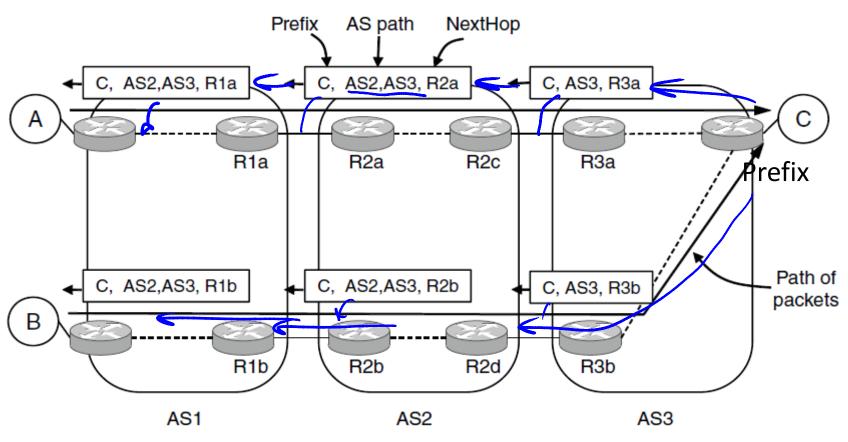
- BGP is the <u>interdomain</u> routing protocol used in the Internet
 - Path vector, a kind of distance vector



Routing with BGP (2)

- Different parties like ISPs are called AS (Autonomous <u>Systems</u>)
- Border routers of ASes announce
 BGP routes to each other
- Route announcements contain an IP prefix, path vector, next hop
 - Path vector is list of ASes on the way to the prefix; list is to find loops
- Route announcements move in the opposite direction to traffic

Routing with BGP (3)



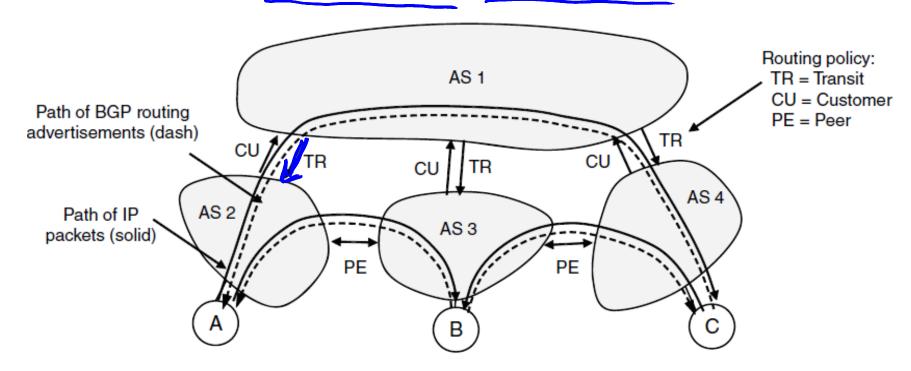
Routing with BGP (4)

Policy is implemented in two ways:

- 1. Border routers of ISP announce paths only to other parties who may use those paths
 - Filter out paths others can't use
- 2. Border routers of ISP select the best path of the ones they hear in any, non-shortest way

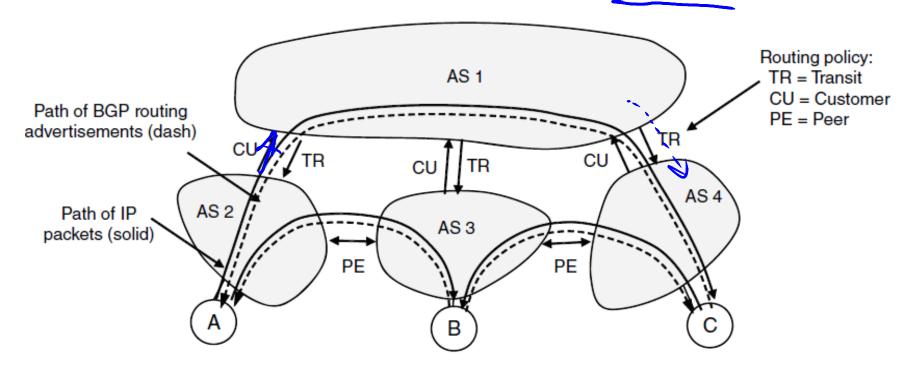
Routing with BGP (5)

• TRANSIT: AS1 says [B, (AS1, AS3)], [C, (AS1, AS4)] to AS2



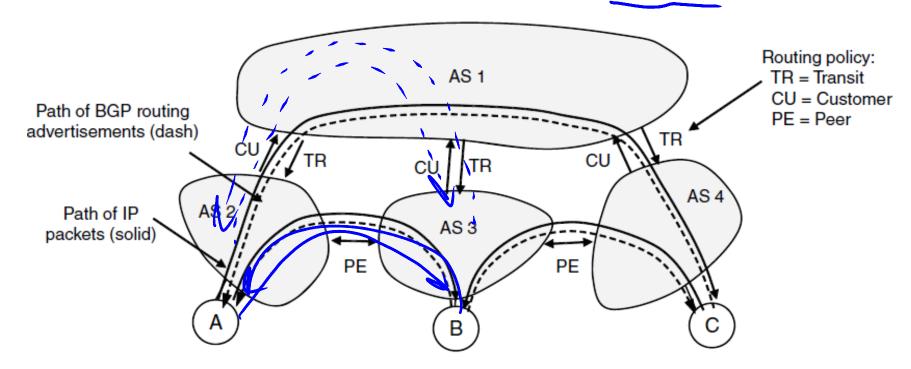
Routing with BGP (6)

CUSTOMER (other side of TRANSIT): AS2 says [A, (AS2)] to AS1



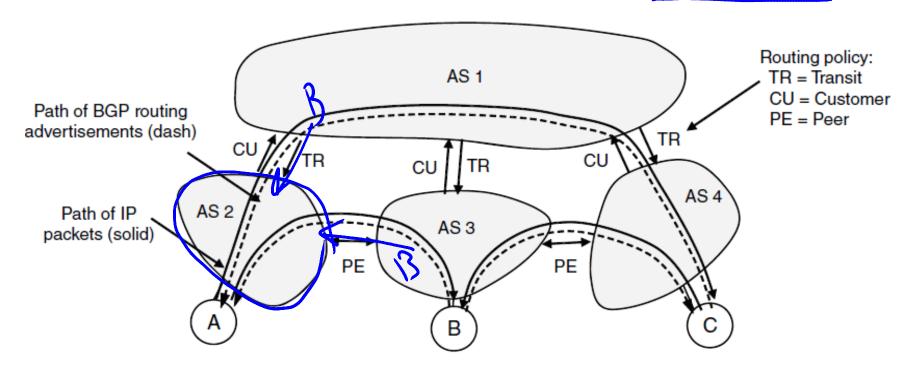
Routing with BGP (7)

PEER: AS2 says [A, (AS2)] to AS3, AS3 says [B, (AS3)] to AS2



Routing with BGP (8)

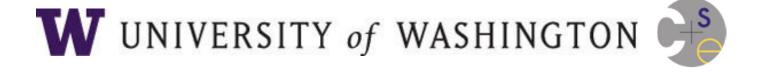
AS2 hears two routes to B (via AS1, AS3) and chooses AS3 (Free!)



BGP Thoughts

- Much more beyond basics to explore!
- Policy is a substantial factor
 - Can we even be independent decisions will be sensible overall?
- Other important factors:
 - Convergence effects
 - How well it scales
 - Integration with intradomain routing
 - And more ...

END



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