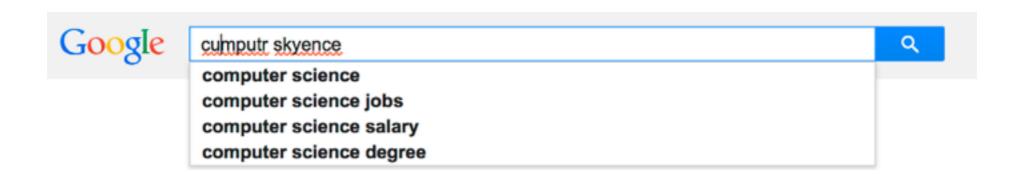
Spelling: Edit Distance & Memoization

Spelling



- Microsoft Word, Google Search, etc., can perform *spelling correction* have you ever wondered how they do that?!
- In order to determine if a word is misspelled, we need to be able to...
 - Compare words (misspelled with properly spelled words)
 - Pick relevant words to compare misspelled words to
- We will start by looking at how to compare words, then talk about how this can be used in a simple spelling corrector.

Edit Distance

- Problem: We want to determine similarity between strings.
- Question: What does it mean for two strings to be "close" to one another?
- Similarity Metric: "Edit Distance".
 - For two strings (s1 & s2), the edit distance is defined as the minimum # of editing operations that transform s1 into s2.
 - Possible operations
 - **Insert:** add a single character. If string s = uv, then inserting the char. x produces uxv. This can also be denoted $\varepsilon \rightarrow x$, using ε to denote the empty string.
 - **Delete:** delete a single character. e.g., uxv to uv $(x \rightarrow \varepsilon)$.
 - Replace: replace a single character x for a symbol $y \neq x$. e.g., uxv to uyv $(x \rightarrow y)$.
 - Swap: swaps two characters. e.g., uxyv to uyxv.
 - Note: this is useful for typos like "thier" vs. "their". (one op. to correct).
 - Assume "cost" of each operation is equal.

Edit Distance

Ex. editDistance("Virginia", "Vermont")

```
Virginia
Virginia (replace)
Verginia (replace)
Verminia (replace)
Vermonia (replace)
Vermonta (delete)
Vermont
```

- Thus, editDistance("Virginia", "Vermont") = 5.
- How did we know which characters to swap, delete, etc.?

EditDistance.java: Overview

- Contains two methods: naiveEditDistance() and editDistance() that both solve the problem recursively.
- Each function looks at the first character of s1 and s2.
- If they match, find the edit distance between the remaining strings.
- **If they don't match**, try *ALL* the possible operations to make the first characters match, then solve the remaining subproblems. The one that solves the problem in the fewest edits is the one we choose.
- Note: "match" and "replace" have the same subproblem(s) to solve, but their contributions to the edit distance differ by 1 (i.e., +0 for match, +1 for replace).

EditDistance.java (main)

• [Demo].

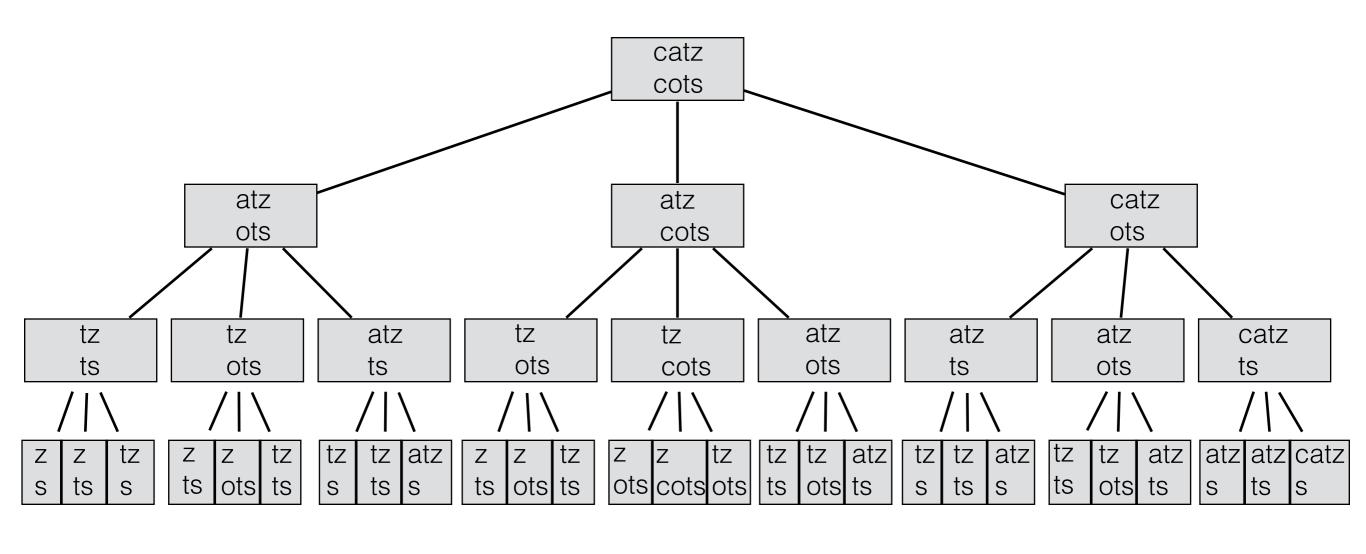
```
• s1 = "cats", s2 = "cotz"
```

•s1 = "Kate Blanchet", s2 = "Cate Blanchett"

EditDistance.java (naive)

```
public static int naiveEditDistance(String s1, String s2) {
                        // Edit distance if first char. match or do a replace
      int matchDist;
      int insertDist; // Edit distance if insert first char of s1 in front of s2.
      int deleteDist; // Edit distance if delete first char of s2.
      int swapDist;
                        // Edit distance for <u>twiddle</u> (first 2 char. must swap).
      if (s1.length() == 0)
          return s2.length(); // Insert the remainder of s2
      else if (s2.length() == 0)
          return s1.length(); // Delete the remainder of s1
      else {
                                                                               c<u>ats</u>
          matchDist = naiveEditDistance(s1.substring(1), s2.substring(1));
cats
          if (s1.charAt(0) != s2.charAt(0))
              matchDist++; // If first 2 char. don't match must replace
cots
          insertDist = naiveEditDistance(s1.substring(1), s2) + 1;
          deleteDist = naiveEditDistance(s1, s2.substring(1)) + 1;
cats
          if (s1.length() > 1 \&\& s2.length() > 1 \&\& s1.charAt(0) == s2.charAt(1) \&\& s1.charAt(1) == s2.charAt(0))
              swapDist = naiveEditDistance(s1.substring(2), s2.substring(2)) + 1;
cots
           else
              swapDist = Integer.MAX_VALUE; // Can't swap if first 2 char. don't match
                                                                                                 "swap"
          return Math.min(matchDist, Math.min(insertDist, Math.min(deleteDist, swapDist)));
  }
                                                                                                  (not relevant in
                                                                                                 this example)
```

EditDistance.java (naive)



Question: Can we do better?! Lots of repeated computations...

Note: this tree only shows match/replace, insert, and delete. The swap op. has been left out since it is a more specific, special case operation. It would, however, add a possible 4th branch to each recursive step.

Memoization

- Better approach: Memoization.
- "Dynamic Programming"; benefits/motivation:
 - A special type of algorithmic design pattern (simpler code)
 - Produces polynomial-time algorithms to solve problems that would seem to take exponential-time.

• Basic Idea:

Keep a matrix of subproblem answers and fill it in using an order that guarantees the subproblem's solution is known before you use it to fill in another place in the table.

- Key: Identify/reuse overlapping subproblems.
- Memory/Runtime tradeoff: use more memory, but save on computation (i.e., store results of already solved subproblems).

• Runtime:

• $O(m^*n)$ were m and n are the length of s1 and s2, respectively.

EditDistance.java (main)

EditDistance.java (memoized)

```
private Map<StringPair, Integer> solvedProblems;
public int memoizedEditDist(String s1, String s2) {
    solvedProblems = new HashMap<StringPair, Integer>();
    return editDist(s1, s2);
private int editDist(String s1, String s2) {
   else {
        StringPair pair = new StringPair(s1, s2);
        Integer result = solvedProblems.get(pair);
        if (result != null) // Did we find the subproblem in the map?
            return result; // If so, return the answer
        else {
            int dist = Math.min(matchDist, Math.min(insertDist, Math.min(deleteDist, swapDist)));
            solvedProblems.put(pair, dist); // Save the result for future
            return dist;
        }
```

StringPair.java

```
public class StringPair {
    private String s1, s2; // The pair of strings
    /**
     * Construct a new pair
     * @param str1 the first string
     * @param str2 the second string
     */
    public StringPair(String str1, String str2) {
        s1 = str1;
        s2 = str2;
    }
    public boolean equals(Object other) {
        StringPair otherPair = (StringPair) other;
        return s1.equals(otherPair.s1) && s2.equals(otherPair.s2);
    }
    public int hashCode() {
        return s1.hashCode() + 31 * s2.hashCode();
}
```

EditDistance.java (memoized)

	С	a	t	S
C	0	1	2	3
0	1	1	2	3
t	2	2	1	3
Z	3	3	3	2

- [Do Demo again w/ memoized code].
 - *s1* = "catz", *s2* = "cots"
 - •s1 = "Kate Blanchet", s2 = "Cate Blanchett"

Spelling Corrector

- Some Notes:
 - A link to the spell corrector algorithm we will look at linked from the notes.
 - Initial implementation in Python ported to Java (Spelling.java).
- Basic Idea for how it works:
 - 1. Use some *measure* to compute distance between two words.
 - Edit Distance is a reasonable measure; other measures work too.
 - 2. Only look at words with minimum edit distance (i.e., E.D. = 1).
 - 3. Pick most common word in English
 - Determined by "big.txt" (compiled from books and lists of most frequent words in Wiktionary, etc.).
 - "bigger.txt" also includes the Unix spelling dictionary (to incorporate more rarely used words).

Algorithm:

- Find all words of minimum edit distance.
- Return the one that occurs most frequently in big.txt.

• [Demo].

```
/**
 * Original version read a single word to correct from the command line. It
 * is commented out below
 *
 * @throws IOException
 */
public static void main(String args[]) throws IOException {
    //Spelling corrector = new Spelling("inputs/big.txt");
    Spelling corrector = new Spelling("inputs/bigger.txt");
    Scanner input = new Scanner(System.in);
    System.out.println("Enter words to correct");
    String word = input.next();
    while (true) {
        System.out.println(word + " is corrected to " + corrector.correct(word));
        word = input.next();
    }
}
```

```
private HashMap<String, Integer> nWords;
/**
* Constructs a new spell corrector. Builds up a map of correct words with
* their frequencies, based on the words in the given file.
* @param file the text to process
* @throws IOException
*/
public Spelling(String file) throws IOException {
   nWords = new HashMap<String, Integer>();
   BufferedReader in = new BufferedReader(new FileReader(file));
   // This pattern matches any word character (letters or digits)
   Pattern p = Pattern.compile("\\w+");
   for (String temp = ""; temp != null; temp = in.readLine()) {
       Matcher m = p.matcher(temp.toLowerCase());
        /*
        * find looks for next match for pattern p (in this case a word). True if found.
         * group then returns the last thing matched. The ? is a conditional expression.
       while (m.find())
           nWords.put((temp = m.group()), nWords.containsKey(temp) ? nWords.get(temp) + 1 : 1);
   in.close();
}
```

```
public String correct(String word) {
   // If in the dictionary, return it as correctly spelled
   if (nWords.containsKey(word))
       return word;
   HashMap<Integer, String> andidates new HashMap<Integer, String>();
       Find all things edit distance 1 that are in the dictionary. Also remember
       their frequency count from nWords.
       (Note if equal frequencies the last one will be the one remembered.)
   for (String s : list)
       if (nWords.containsKey(s))
           candidates.put(nWords.get(s), s);
   // If found something edit distance 1 return the most frequent word
   if (candidates.size() > 0)
       return candidates.get(Collections.max(candidates.keySet()));
    * Find all things edit distance 1 from everything of edit distance 1.
    * These will be all things of edit distance 2 (plus original word). Remember frequencies
   for (String s : list)
       for (String w : edits(s))
           if (nWords.containsKey(w))
               candidates.put(nWords.get(w), w);
    * If found something edit distance 2 return the most frequent word.
    * If not return the word with a "?" prepended. (Original just returned
    * the word.)
    */
   return candidates.size() > 0 ? candidates.get(Collections.max(candidates.keySet())) : "?" + word;
```

```
/**
 * Constructs a list of all words within edit distance 1 of the given word.
 * @param word the word to construct the list from
 * @return a list of words with in edit distance 1 of word
 */
private ArrayList<String> edits(String word) {
    ArrayList<String> result = new ArrayList<String>();
   // All deletes of a single letter
   for (int i = 0; i < word.length(); ++i)
        result.add(word.substring(0, i) + word.substring(i + 1));
   // All swaps of adjacent letters
   for (int i = 0; i < word.length() - 1; ++i)
        result.add(word.substring(0, i) + word.substring(i + 1, i + 2) + word.substring(i, i + 1) +
                                                                                         word.substring(i + 2));
   // All replacements of a letter
    for (int i = 0; i < word.length(); ++i)
        for (char c = 'a'; c <= 'z'; ++c)
            result.add(word.substring(0, i) + String.valueOf(c) + word.substring(i + 1));
    // All insertions of a letter
   for (int i = 0; i <= word.length(); ++i)
        for (char c = 'a'; c <= 'z'; ++c)
            result.add(word.substring(0, i) + String.valueOf(c) + word.substring(i));
    return result;
}
```