

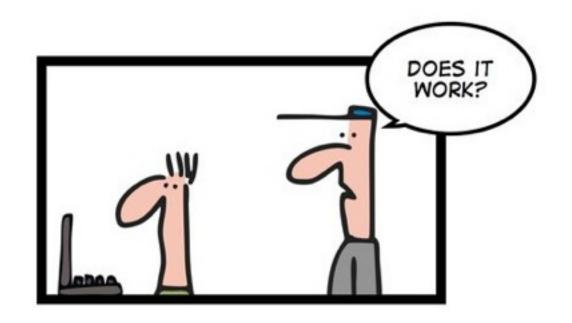
Concurrency (Part II): Mutual Exclusion, Synchronization, Deadlock, and Starvation

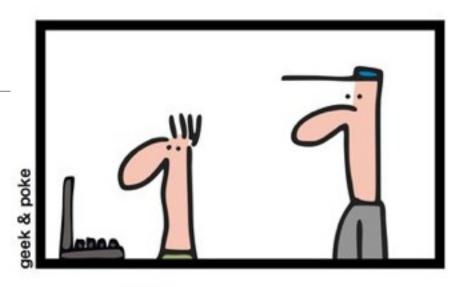
Professor Travis Peters
CSCI 460 Operating Systems
Fall 2019

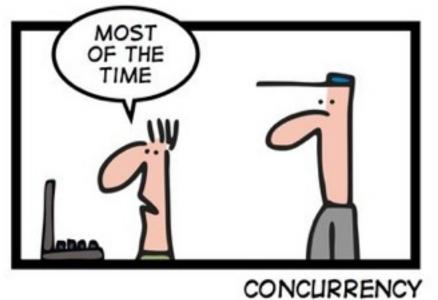
Some slides & figures adapted from Stallings instructor resources.

Some slides adapted from Adam Bates's F'18 CS423 course @ UIUC https://courses.engr.illinois.edu/cs423/sp2018/schedule.html

SIMPLY EXPLAINED







-http://www.datamation.com/news/tech-comics-quantum-physics-2.html



Goals for Today

Learning Objectives

- Dive a bit deeper into core topics in concurrency
- Discuss common mechanisms for achieving mutual exclusion & synchronization

Announcements

- Schedule Updates
 - Exam #1 delayed…
 - · ...career fair...
 - topics to include everything up through concurrency (at least... possibly scheduling as well...)
 - Exam #2 topics will include:
 - (Scheduling)
 - Memory Management & Virtual Memory
 - File Systems & I/O
 - Selected OS Security Topics
 - No Exam #3!
- Homework 2 (Chapters 3-4) DUE 10/11 trust me, you don't need that long!
- · Homework 3 (Chapters 5-6) DUE 10/18
- Programming Assignment 1 (Concurrency + C programming/pthreads) should be posted by Friday



Recap: Concurrency & Solutions for Mutual Exclusion

To acheive correct & meaningful solutions to concurrency problems, mutual exclusion is a must!

Software Support

- · Assume elementary mutual exclusion at the memory access level; serialized by "memory arbiter"
- · Decker's Algorithm, Peterson's Algorithm

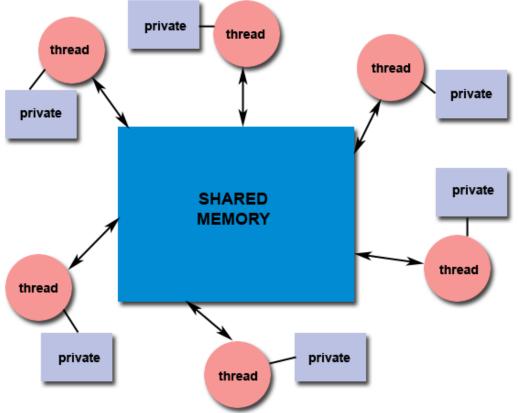
Hardware Support

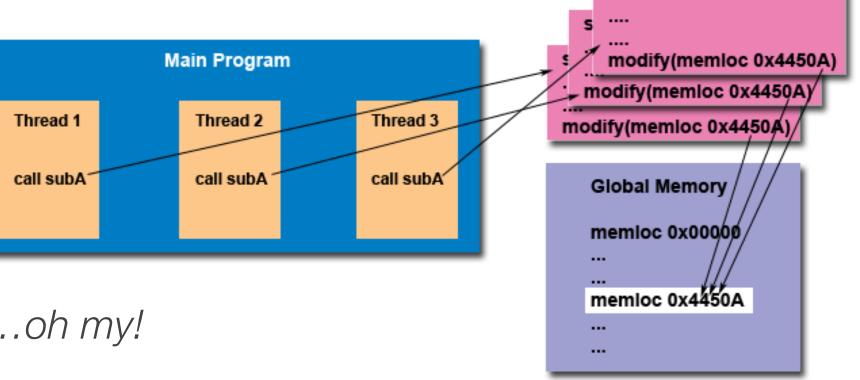
- Interrupt Disabling
 - · Disadvantages: inhibits processor's ability to interleave processes; doesn't work across processors.
- Special Machine Instructions
 - Compare&Swap: compare values => if values are the same, swap!
 - Exchange (XCHG): exchanges the contents of a register w/ that of a memory location
 - · Advantages: simple & easy to implement; can be used on multi-processor machines
 - · Disadvantages: possibly expensive busy-waiting; starvation & deadlock are still possibe

Programming Language Mechanisms

- Examples using pthreads
- · Semaphores, Monitors, Condition Variables, Message Passing, Mutexes (Locks), ...oh my!







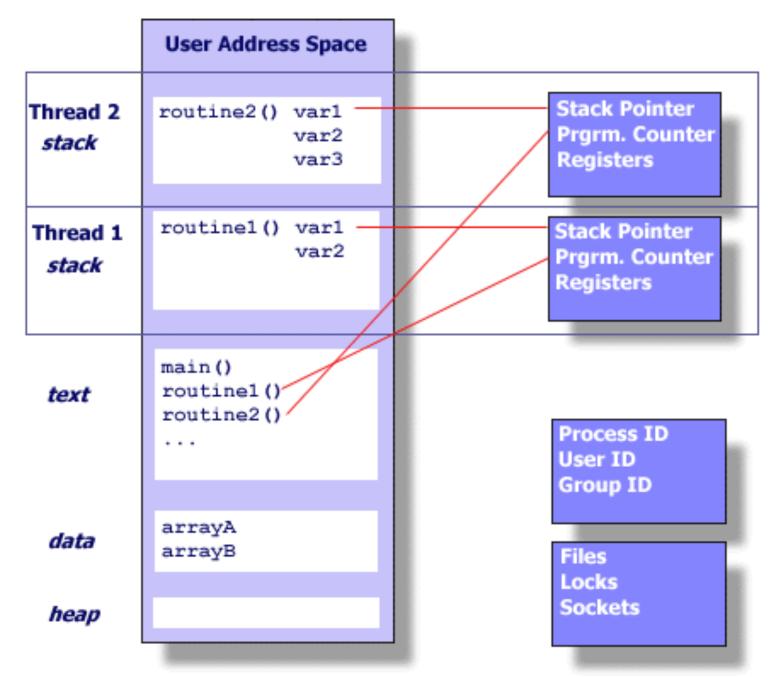
-https://computing.llnl.gov/tutorials/pthreads/



Programming w/ pthreads

A good write-up on pthreads: https://computing.llnl.gov/tutorials/pthreads/

- pthreads (POSIX Threads)
 - · defines a set of C programming language types, functions and constants.
 - is implemented with a pthread.h header and a thread library.
 - · includes mutexes, condition variables, etc.
- E.g., A typical sequence in the use of a **mutex** is as follows:
 - → Create and initialize a mutex variable
 - → Create and start several threads
 - → Several threads attempt to lock the mutex (only 1 succeeds and "owns" the lock)
 - →The owner thread performs some set of actions
 - ◆The owner unlocks the mutex
 - →Another thread acquires the mutex and repeats the process
 - → Finally, the mutex is destroyed



- https://computing.llnl.gov/tutorials/pthreads/

NOTE: When several threads compete for a mutex, the losers block at that call (there does exist a non-blocking call: "trylock")

NOTE: It is the programmer's responsibility to make sure **every thread** that needs to use a mutex does so. For example, if 10 threads are updating the same data, but only one uses a mutex, the data can still be corrupted!



Programming w/ pthreads

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Most of class featured in-class demo:

- Code Walk-Through threads01.c & threads02.c>
- · What happens if I vary the number of threads?
 - 1 vs. 2 vs. 5 vs. 10 vs.
- What happens if the target number is small?
 - 100 vs. 100000000000
- What happens if we don't wait for the threads to complete?
- · What happens if I compile & run on *different machines* (e.g., naitive machine/OS vs. my local Linux VM)?
- Why are these things happening?!
 - NOTE: make sure you (re-)build the executable whereever you are going to run it....
 - objdump -d t1
- · With proper (simple) synchronization, why does the program appear to execute *slower*?

