

CSCI 460 Operating Systems

Computer System & Operating System Overview

Professor Travis Peters Fall 2019



Goals for Today

Learning Objectives

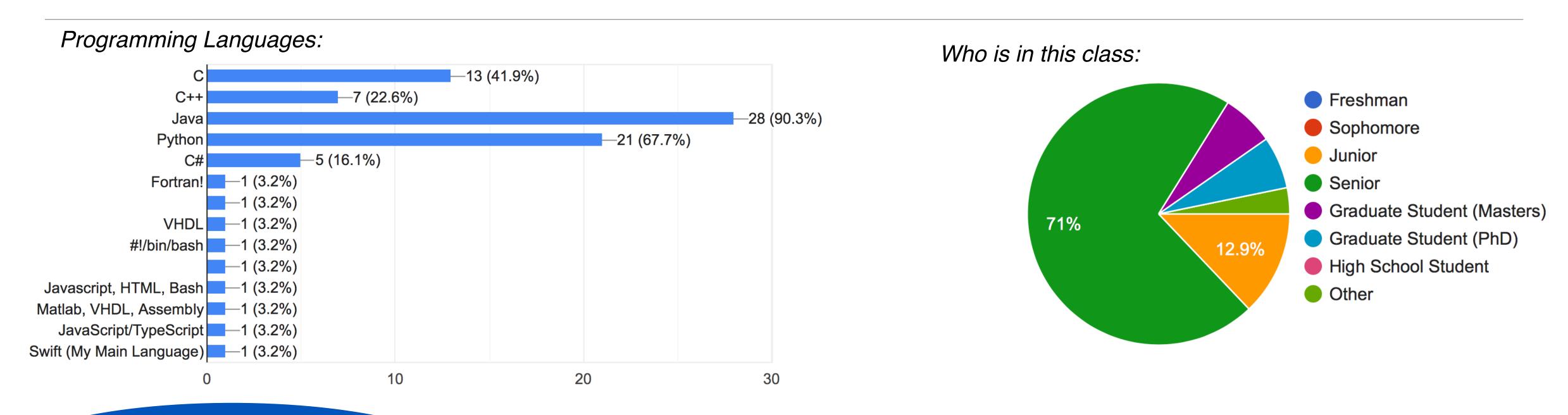
- Wrap up high-level computer system & OS concepts (Chapters 1 & 2)
- Connect some dots and pave the way for the rest of the course
 - not going deep into history, which was covered (enough) last week...
 - gory details in the text, of course...

Announcements

- Note taker?!! ;-)
- Grades for HW1 coming soon (restructuring things in D2L over the next week or so)
- More HW soon, I promise ;-)

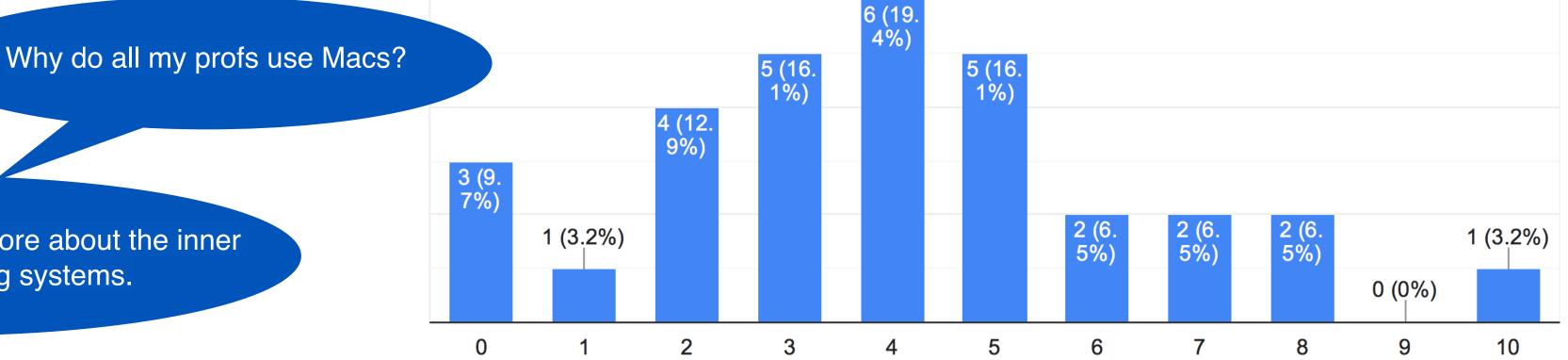


Questionnaire - The results are in (mostly...)



I've heard from software engineers that taking OS is a very helpful class for once you get out into the field

Sys Programming Skillzzz:

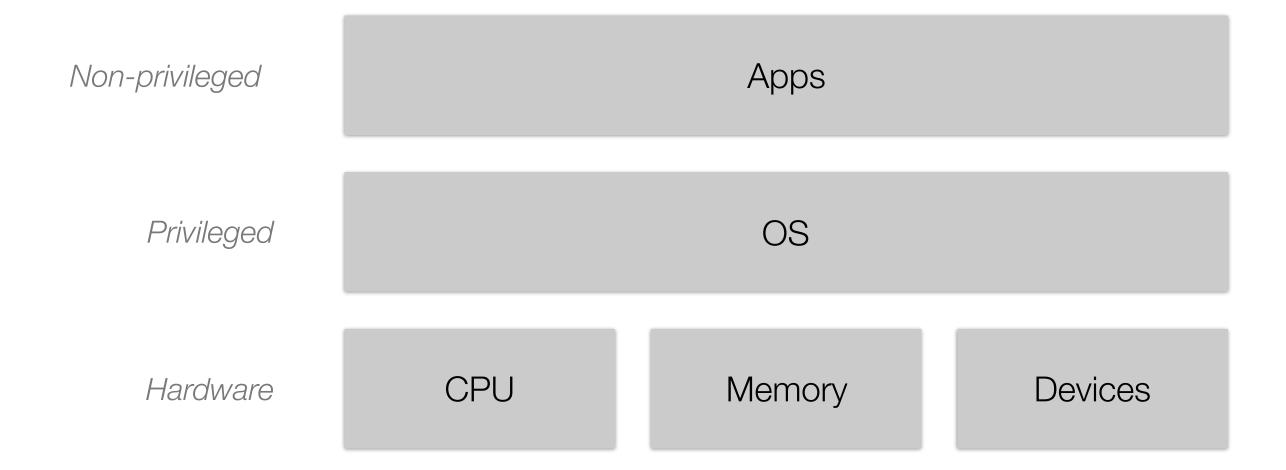


I did not like the other course offerings

I'm just interested to learn more about the inner workings of operating systems.

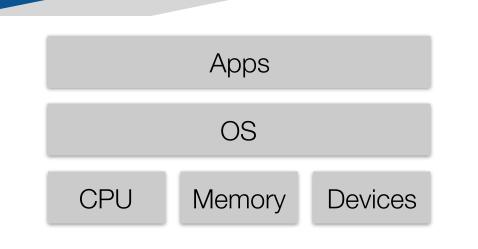


Basic Model

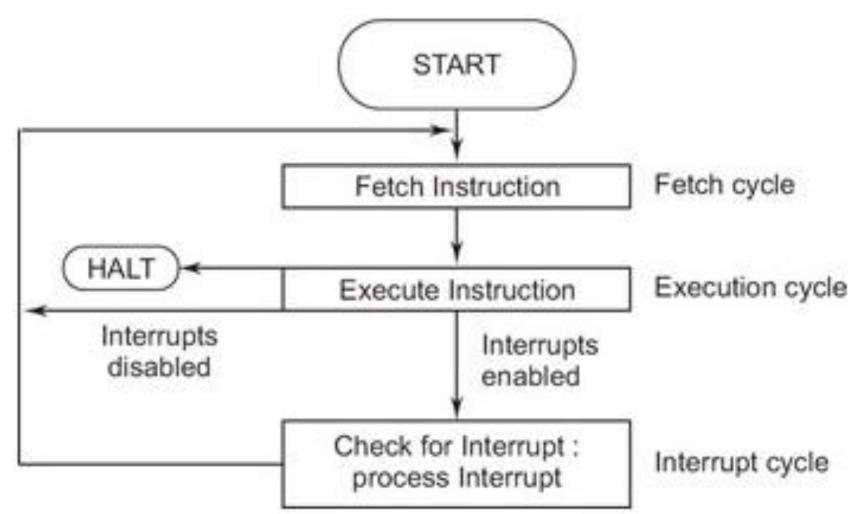


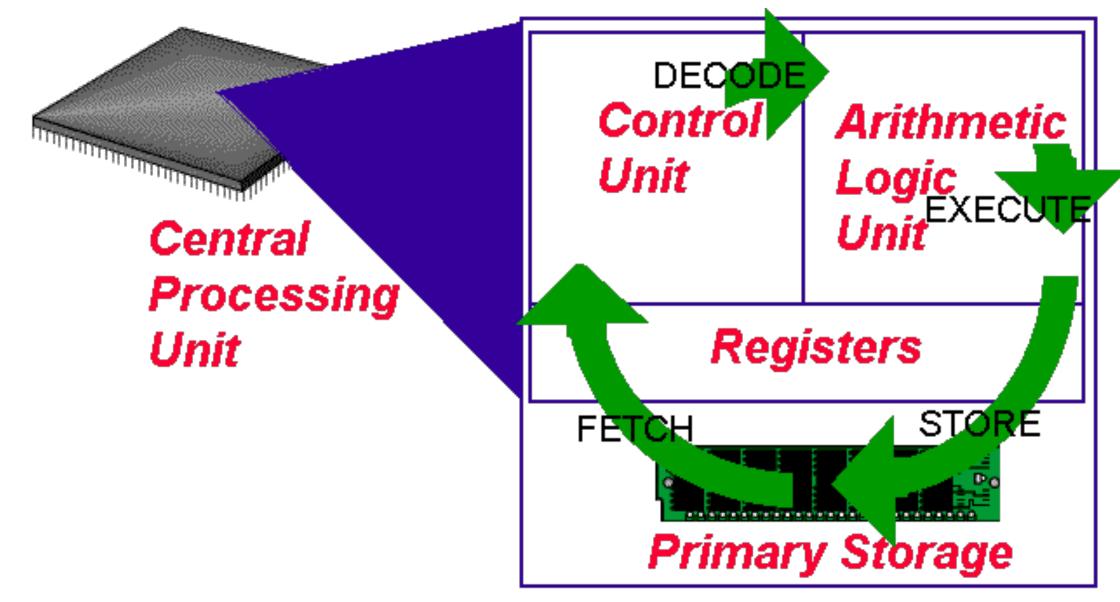
CSCI 460: Operating Systems Montana State University





- Instruction Exectuion / Instruction Cycle
 - FETCH the next instruction, DECODE it, EXECUTE it, and STORE the result.
- Interrupts (interrupt the processor!)
 - E.g., program int. (e.g., illegal instruction, out-of-bounds access), timer, I/O, hardware failure
 - Run corresponding "handler"





CSCI 460: Operating Systems

Montana State University



Apps
OS
CPU Memory Devices

Memory

The Memory Hierarchy

want fast access to a large amount of memory

locality of reference: memory access tends to cluster (e.g., loops, subroutines)

Cache Memory

keep as much as possible, as close as possible...
...but still support large amounts of memory.
Multiple levels of cache... need cache coherency.

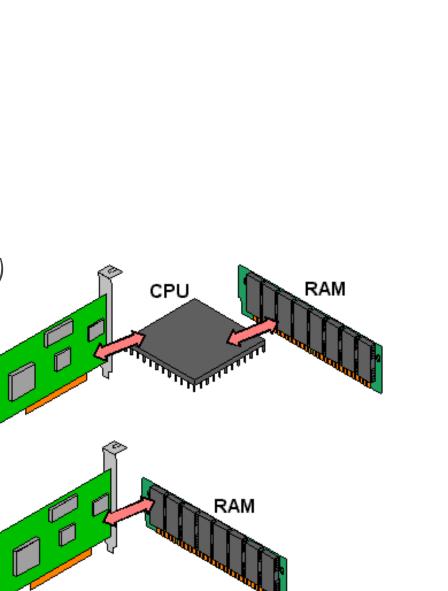
Direct Memory Access (DMA)

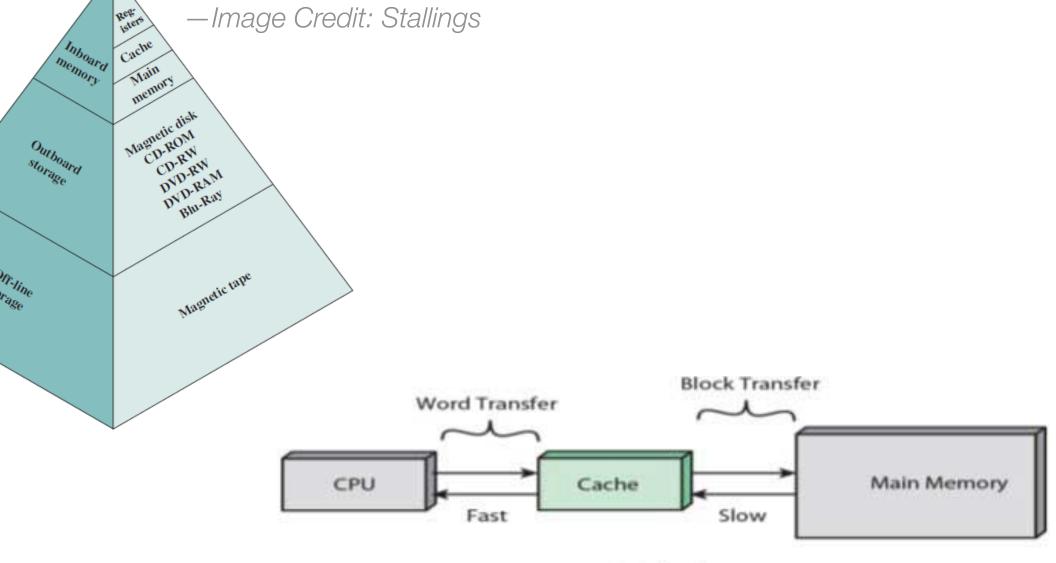
programmed I/O (active; processor polls I/O device) **vs.**

interrupt-driven I/O (assist; I/O device does work; interrupt processor to help)

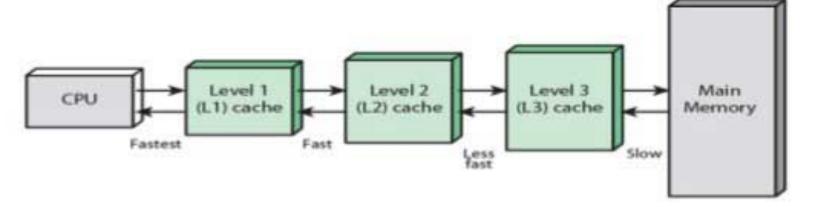
VS

DMA I/O (<u>delegate</u>; DMA module is given OP, DEV, ADDR, #WORDS; sent INT when complete)





(a) Single cache



- Image Credit: Stallings

CSCI 460: Operating Systems

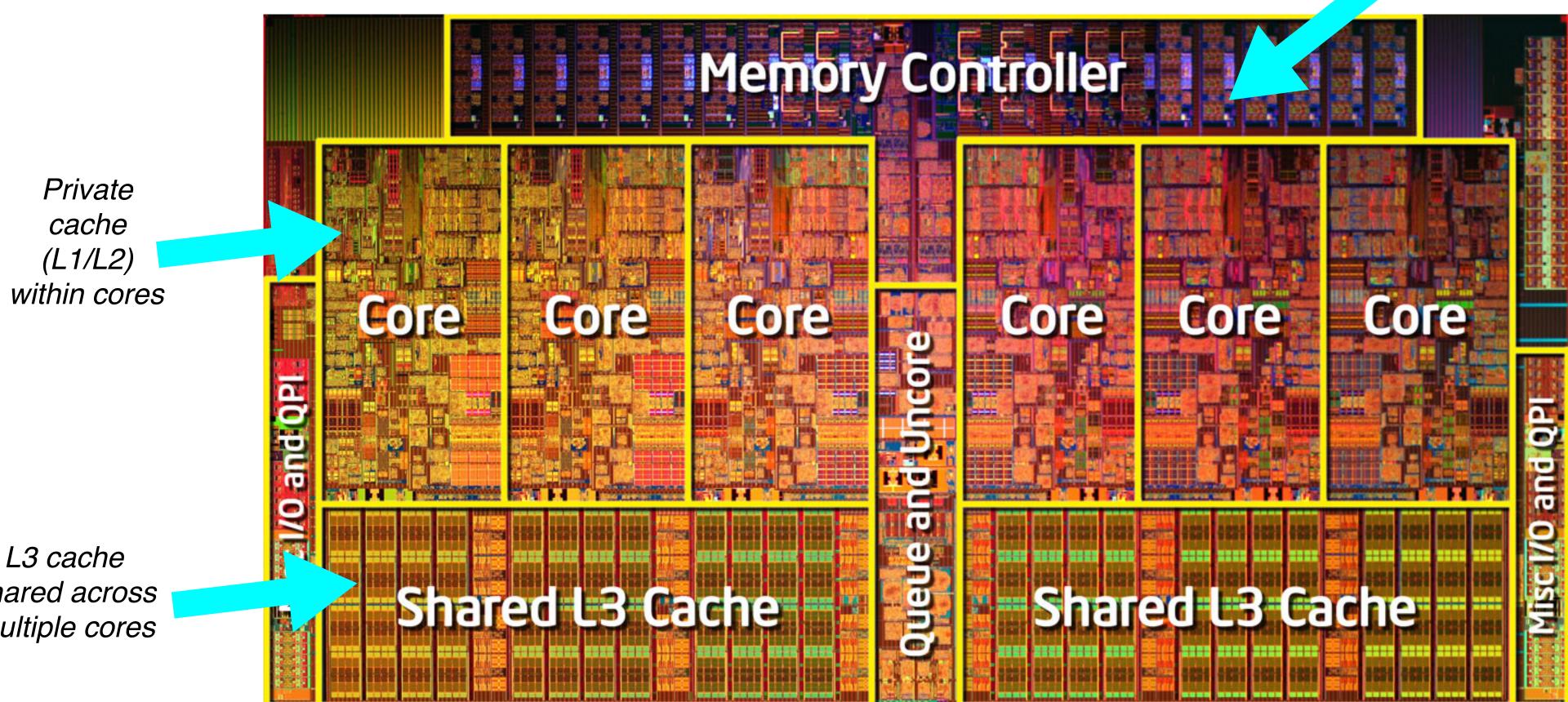
Montana State University



Apps OS Devices Memory

Organization IRL

Access to RAM and Disk via Memory Controller



shared across multiple cores

- Image Credit: Stallings

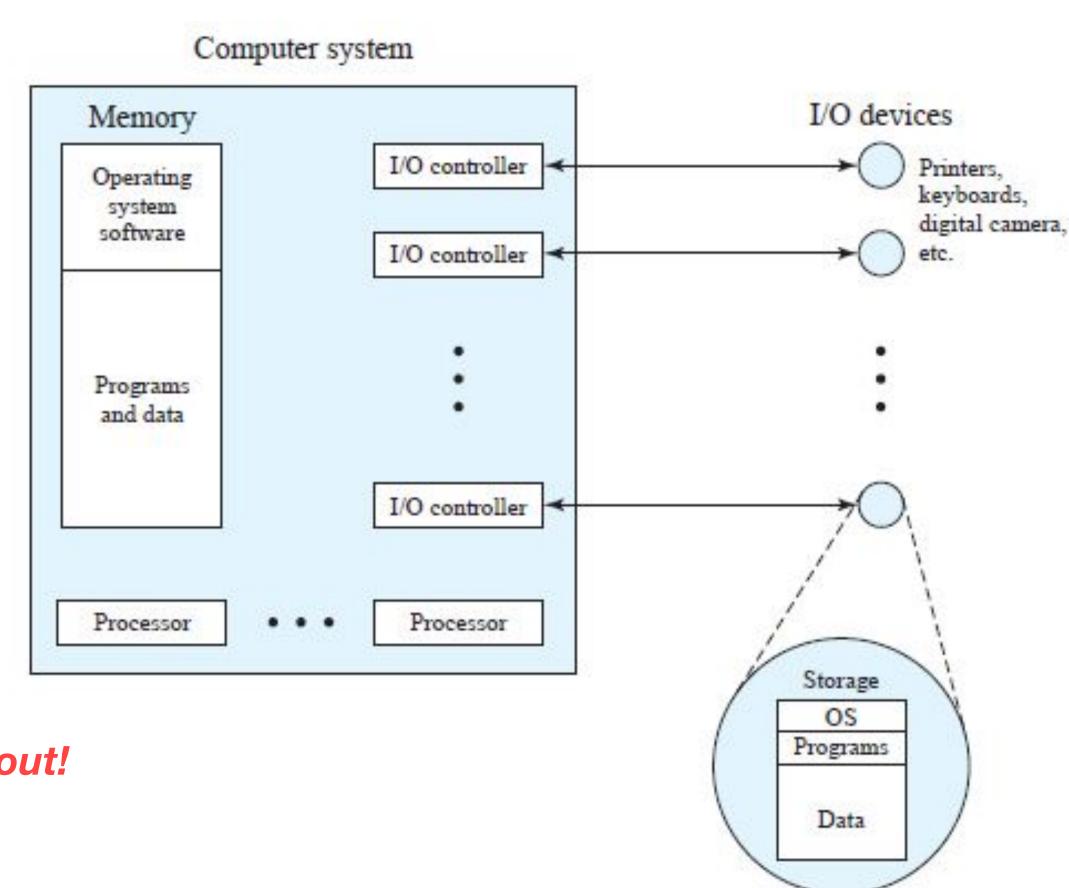
CSCI 460: Operating Systems Montana State University



Operating System Overview

- Objectives of an OS
 - Covenience, Efficient, Modular
 vs.
 Referee, Illusionist, Glue
- Major Achievements
 - The Process
 - Memory Management
 - Security
 - Scheduling & Resource Management

These topics are what most of the rest of the class is about!



- Image Credit: Stallings



OS Objectives/Roles

The textbook offers one way of looking at the role of the OS. Here is another:

Referee

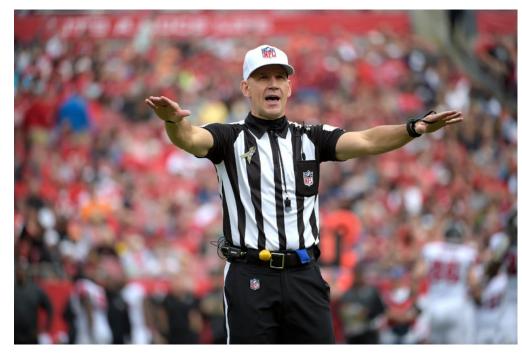
- manage resources between apps
- isolate apps and users from one another
- facilitate communications between apps/users

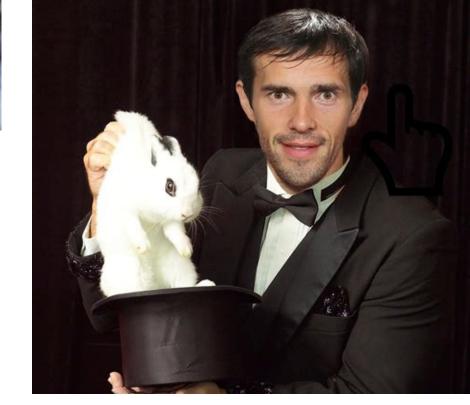
Illusionist

- make apps believe they have the whole machine to themselves
- create appearance of infinite processors and memory
- abstract away complexity of storage, network communications, etc.

Glue

manage hardware so apps can be machine-agnostic; provide a set of common services





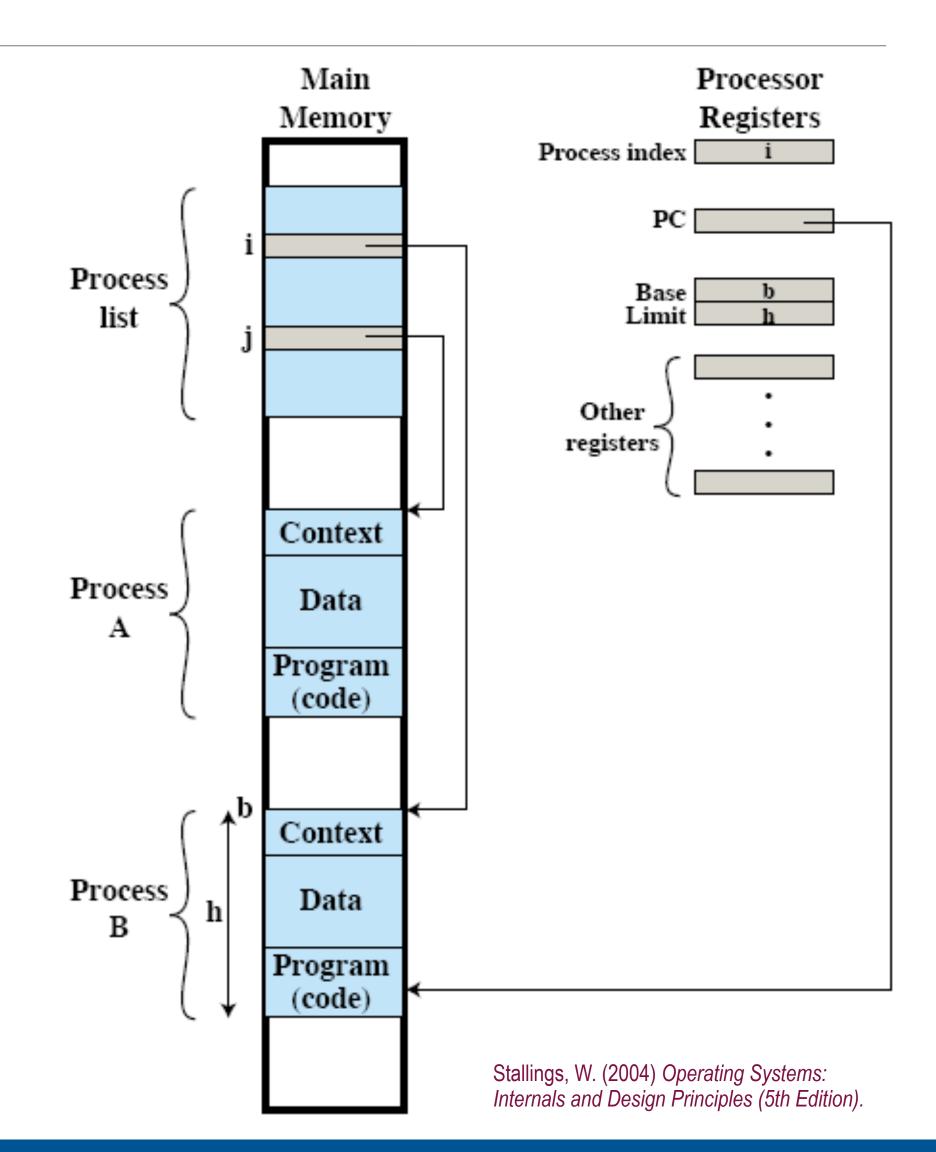


- Thanks to Adam Bates for the nice analogies: https://courses.engr.illinois.edu/cs423/sp2018/



The Process

- The notion of a *Process*....
 - · arose from multiprogramming, time sharing, real-time transactions;
 - is more general than a "job"; many definitions; e.g., an instance of a program running on a computer;
 - consists of
 - (1) an executable program (code),
 - · (2) assocated data,
 - (3) **execution context** (info the OS needs to manage the process)
 - is realized as nothing more than a data structure!
 - A thread = cooperative execution within a process; use shared context.
- Process Switching
 - Interrupt = save context (e.g., PC and other registers)
 - execute interrupt handler
 - resume processing (same or different process)
 - states (simply put) = executing or awaiting execution





Memory Management

OS responsibities w.r.t. MM

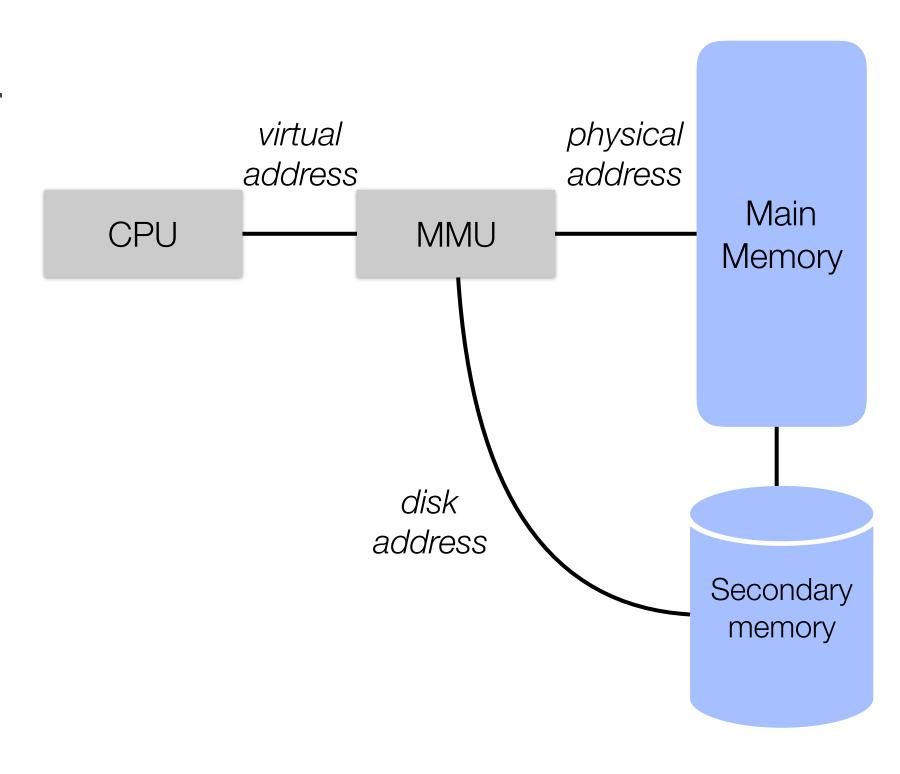
• Process Isolation + Automatic Allocation/Mgmt. + Modular + Protection & Access Control + Long-Term Storage

Virtual Memory

- · address memory *logically*, without regard to the actual, *physical* memory.
- (recall "paging") virtual address = page # + offset within the page.
- MM maps between virtual address and real (physical) addresses.
- pages can be in memory (or not!) how is that possible?!

File Systems

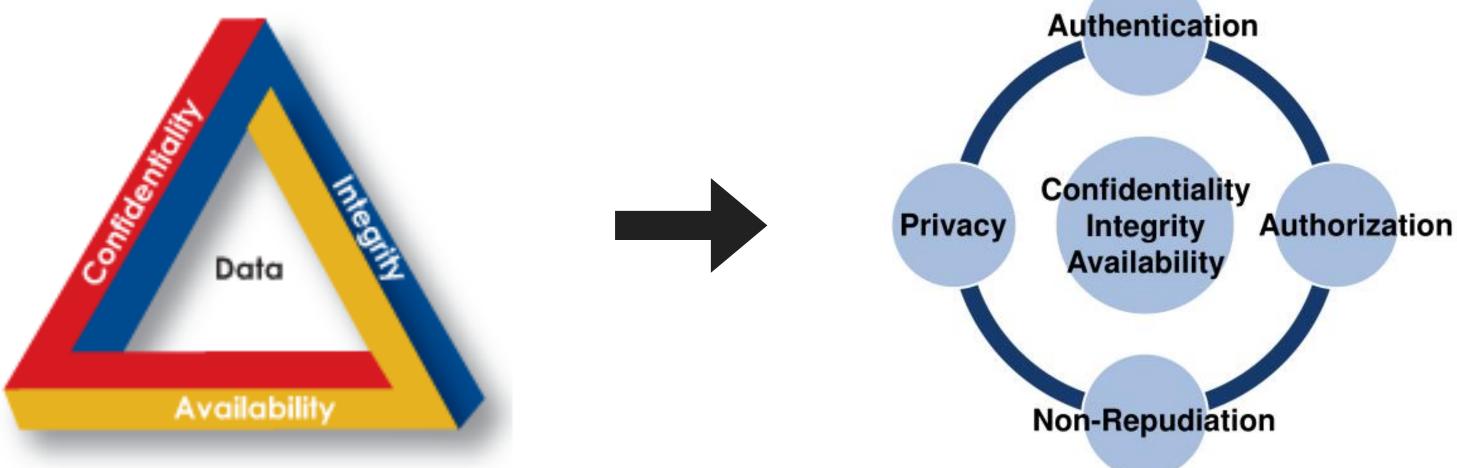
- good for persisting information for extended periods of time
- **files** = named objects





Information Protection & Security

- · CIA
 - Confidentiality protection from unauthorized access (e.g., snooping)
 - Integrity protection from unauthorized modification (e.g., tampering)
 - Availability protection from interruption (e.g., denial of service)
- Also...
 - Authenticity verification of the origin/identity
 - +...



traditional infosec model...

...increasingly common to include other goals

CSCI 460: Operating Systems Montana State University



Scheduling & Resource Management

- Manage Resources!
 - Manage system resources (Main Memory, I/O devices, Processors)
 - Keep these resources *utilized* (i.e., schedule processes to utilize them)
 - · ...all while being *fair*, *responsive*, and *efficient*
- Operations-Research Problems & Lots of Maths
 - -How do we do things good?
 - Data Structures: A bunch of queues, and lists, and (mostly) other simple data structures
 - · Algorithms: A scheduler or dispatcher to pick which process runs next
 - round-robin = everyone gets some time in turn. Other approaches?

