

# **OS Security**

Professor Travis Peters
CSCI 460 Operating Systems
Fall 2019



# Will Peteroy

Prep Resources (Read Before Class!)

### Reading / Prep (Read Before Class)

- https://msrc-blog.microsoft.com/2010/12/08/on-the-effectiveness-of-dep-and-aslr/
- https://msrc-blog.microsoft.com/2013/08/12/mitigating-the-ldrhotpatchroutine-depaslr-bypass-with-ms13-063/
- <a href="https://arstechnica.com/information-technology/2019/08/armed-with-ios-0days-hackers-indiscriminately-infected-iphones-for-two-years/">https://arstechnica.com/information-technology/2019/08/armed-with-ios-0days-hackers-indiscriminately-infected-iphones-for-two-years/</a>

### Case Study 1 - Effectiveness of DEP and ASLR / Bypasses

- https://msrc-blog.microsoft.com/2010/12/08/on-the-effectiveness-of-dep-and-aslr/
- https://msrc-blog.microsoft.com/2013/08/12/mitigating-the-ldrhotpatchroutine-depaslr-bypass-with-ms13-063/

### Case Study 2 - Apple - Extreme Odays in 2019

- <a href="https://arstechnica.com/information-technology/2019/08/armed-with-ios-0days-hackers-indiscriminately-infected-iphones-for-two-years/">https://arstechnica.com/information-technology/2019/08/armed-with-ios-0days-hackers-indiscriminately-infected-iphones-for-two-years/</a>
- <a href="https://googleprojectzero.blogspot.com/2019/08/a-very-deep-dive-into-ios-exploit.html">https://googleprojectzero.blogspot.com/2019/08/a-very-deep-dive-into-ios-exploit.html</a>



### A Guest Lecture w/

Will Peteroy — CEO @ icebrg.io

11/04/2019

https://www.traviswpeters.com/cs460/notes/cs460-22-os-security-will-peteroy.pdf



### A Guest Lecture w/

# Daniel Pagan — OS Enthusiast & Pen Tester @ IBM X-Force Red

11/13/2019

https://www.traviswpeters.com/cs460/notes/cs460-23-os-security-daniel-pagan.pdf



# Tying Up Some Loose Ends in OS

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## Threats

- Intruders
  - Masquerader Non-authorized actor that tries to impersonate a legitimate user / exploit a legit. user's access privileges (outsider)
  - Misfeasor A legit. user that abuses their inside access to carry out nefarious deeds (insider)
  - Clandestine User Any entity that seizes privileged control of a system and evades detection
- Malicious Software ("Malware")
  - Trojan horse code that misuses its environment
    - · classic examples: NetBus, Back Orifice (BO), Back Orifice 2000 (B02K)
  - Backdoor specific user identifier or password that bypasses normal security procedures
    - · classic examples: LiteBot, Remote Connection (RedNeck)
  - Stack and Buffer Overflow exploits overflow bug to obtain ability to execute arbitrary code
    - · classic examples: Internet worm (fingerd), Code Red (IIS), SQLSlammer (MS SQL Server)
  - Spyware SW that installs itself on a computer, reports personal info or activities
    - classic examples: adware, stealware

# Security Goal Prevent attempts to gain unauthorized privileges (or at least detect it...)



# General Countermeasures / Defenses

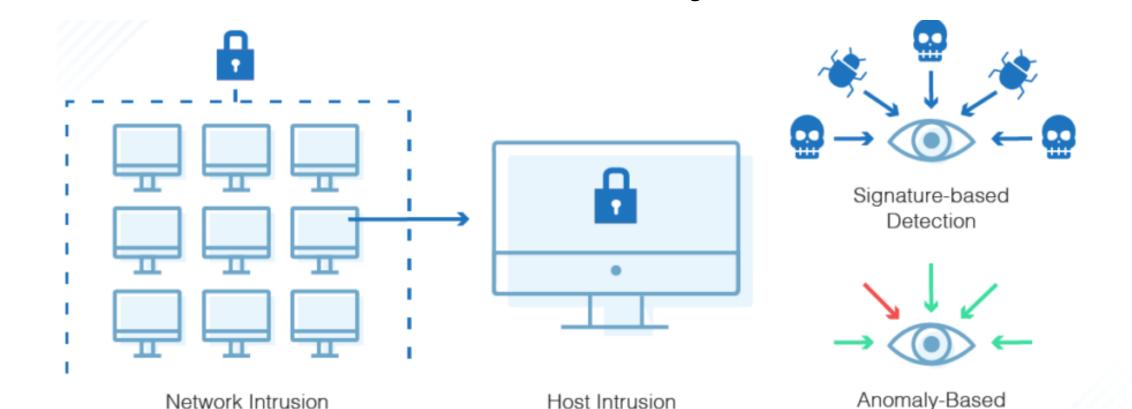


(Multi-Factor) Authentication

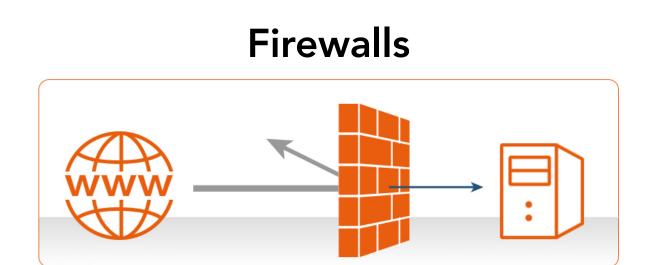
VS.

Verification

& Access Control (later)



Intrusion Detection Systems (IDS)





### "Hardening" the OS/system

Update and install patches

Detection

- · Remove unnecessary services, applications, and protocols
- · Configure appropriate users, groups, and permissions + resource controls

Detection

Detection

- Install additional security tools/controls
- Test
- (Repeat)



# Exploitation Example: The Buffer Overflow

į	<pre>int main(int argc, char *argv[]) {   int valid = FALSE;   char str1[8];   char str2[8];</pre>
3	<pre>next_tag(str1); gets(str2); if (strncmp(str1, str2, 8) == 0)     valid = TRUE; printf("buffer1: str1(%s), str2(%s), valid(%d)\n", str1, str2, valid); }</pre>

#### (a) Basic buffer overflow C code

```
$ cc -g -o buffer1 buffer1.c
$ ./buffer1
START
buffer1: str1(START), str2(START), valid(1)
$ ./buffer1
EVILINPUTVALUE
buffer1: str1(TVALUE), str2(EVILINPUTVALUE), valid(0)
$ ./buffer1
BADINPUTBADINPUT
buffer1: str1(BADINPUT), str2(BADINPUTBADINPUT), valid(1)
```

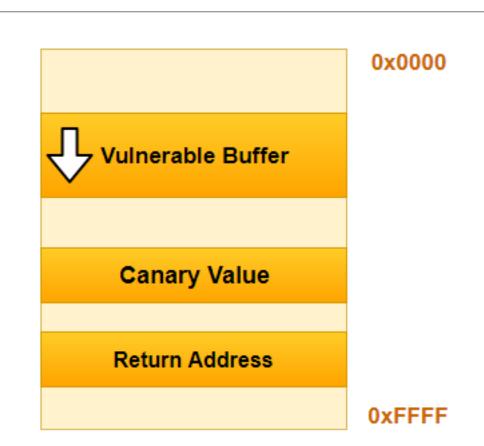
#### (b) Basic buffer overflow example runs

Memory Address	Before gets(str2)	After gets(str2)	Contains Value of
• • •	• • •	• • •	
bffffbf4	34fcffbf 4	34fcffbf 3	argv
bffffbf0	01000000	01000000	argc
bffffbec	c6bd0340	c6bd0340	return addr
bffffbe8	08fcffbf	08fcffbf	old base ptr
bffffbe4	0000000	01000000	valid
bffffbe0	80640140	00640140	
bffffbdc	. d . @ 54001540	. d . @ 4e505554	str1[4-7]
bffffbd8	т@ 53544152	N P U T 42414449	str1[0-3]
bffffbd4	S T A R 00850408	B A D I 4e505554	str2[4-7]
		NPUT	-
bffffbd0	30561540 0 V . @	42414449 B A D I	str2[0-3]
• • •		• • •	

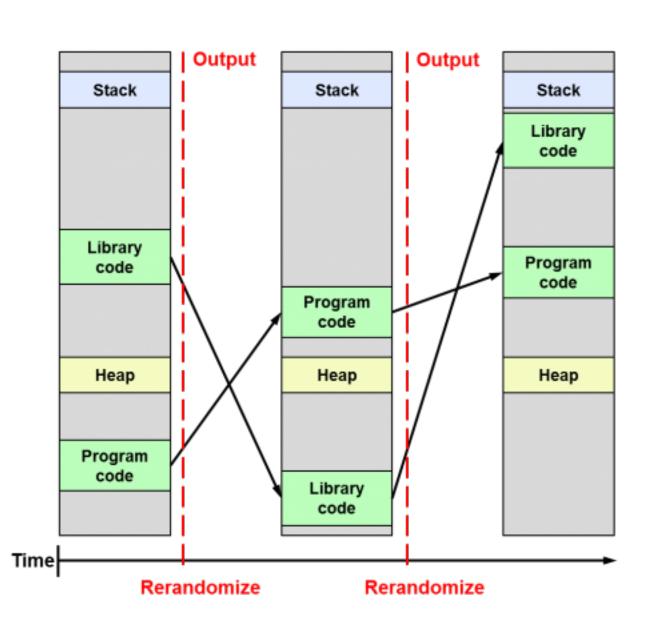


# Overflow Countermeasures / Defenses

- Compile-Time Defenses
  - High-Level Programming Languages
  - Scrutinize code open source vs. closed source...
  - Language Extensions use safe, validated implementations of libraries libsafe
  - Stack Protection stackguard (gcc); "canaries"  $\rightarrow$  favor aborting/crashing programs over being pwned



- Run-Time Defenses
  - Executable Address Space Protection nx-bit on x86
  - ASLR
  - Guard Pages



# Access Control For each protected resource, store a list of **who** is permitted to do **what**

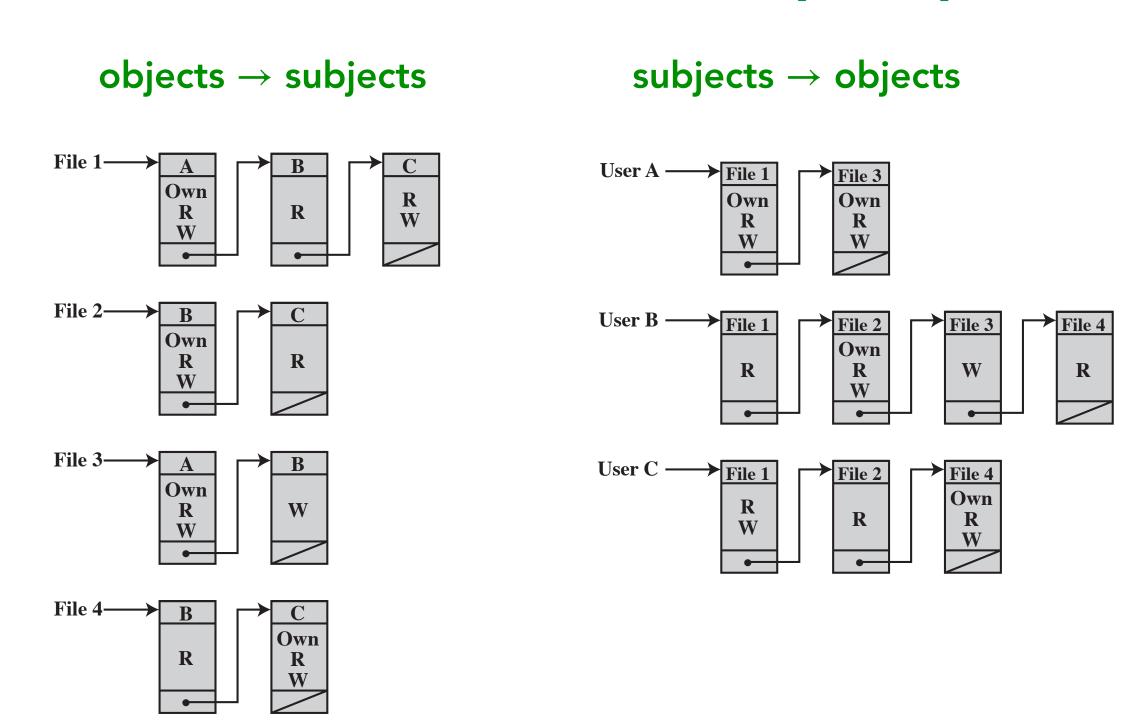
### **Access Matrix**

#### **OBJECTS**

		File 1	File 2	File 3	File 4
	User A	Own Read Write		Own Read Write	
SUBJECTS	User B	Read	Own Read Write	Write	Read
	User C	Read Write	Read		Own Read Write

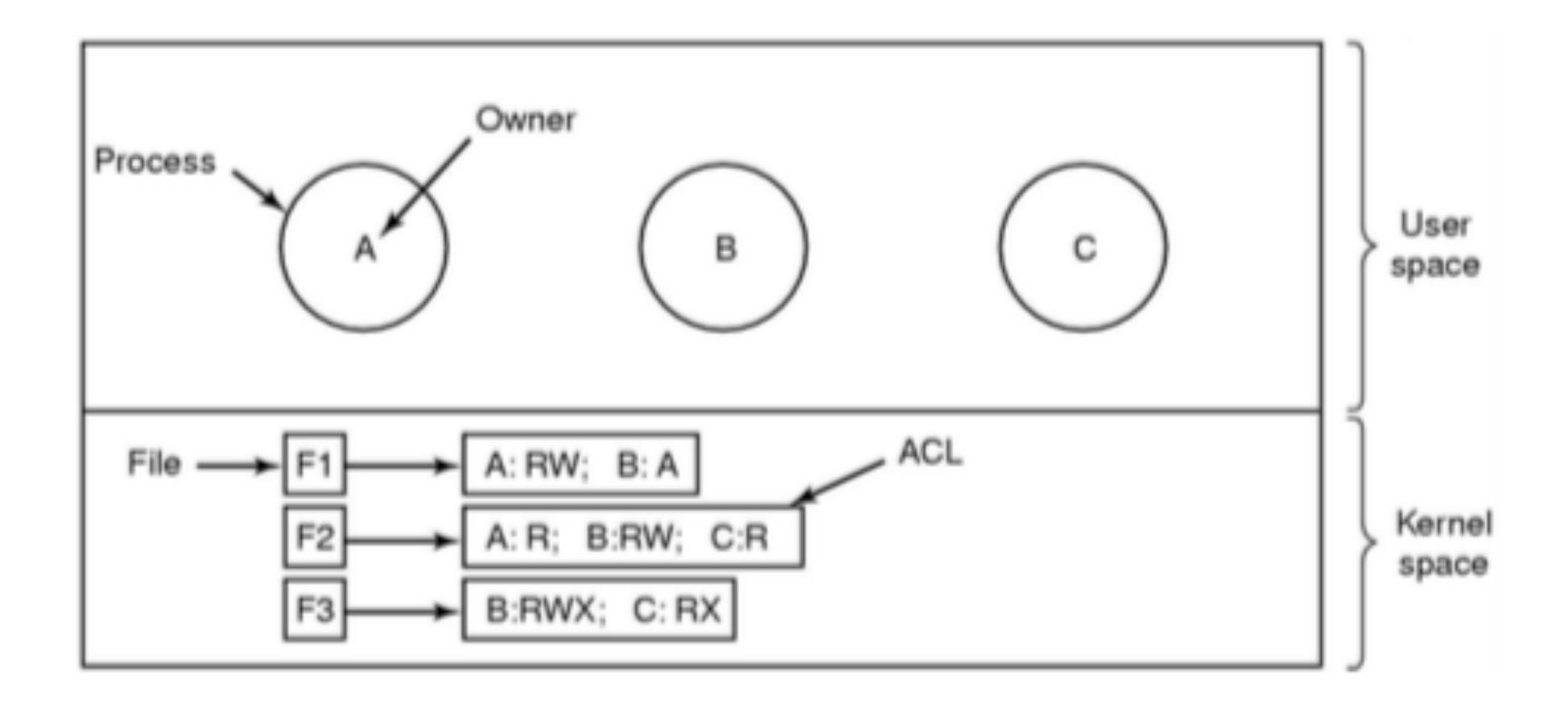
For each subject → easy to look up their capabilities For each object → easy to look up authorized users

# **Access Control List (ACL)**



# Access Control (cont.)

# Example: Use of access control lists to manage file access in UNIX





# Access Control (cont.)

#### DAC MAC RBAC Roles have useful meaning and Each subject can perform Each subject can perform action(s) on objects action(s) on objects context for access permissions subjects & objects each have subjects & objects each have Example: student, lab associated security attributes associated security attributes manager, professor, dean, president Permissions based on Permissions based on Each **subject** is mapped to a user/group user/group role(s); each role has permission to perform actions **Discretionary** → users can grant **Mandatory** $\rightarrow$ sec. policy is others access similar to its own centrally controlled; users do not on **objects**. permissions have the ability to override the policy (e.g., cannot grant others access to files that would otherwise be restricted) Subject can Action to Object Subject can grant other Subject Subject can Action to Object Object can be Action by Subject Subject is a Role which has Permission of Action to Object

### ABAC

Policies expressed as a complex Boolean rule set that can evaluate many different attributes

#### Examples:

Subject Attributes (age, clearance, dept. role, job title)

Action Attributes (read, write, view, approve)

Resource Attributes (object type, sensitivity, location)

Contextual Attributes (time, location)

Subject who is xxx can Action to Object which is xxx in Environment

Some examples drawn from: https://dinolai.com/notes/others/authorization-models-acl-dac-mac-rbac-abac.html

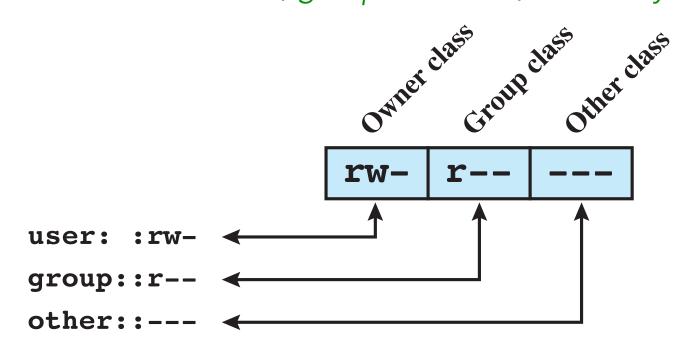


# Some Examples of Access Control IRL

### **Basic UNIX File Access Control**

Every file has an owner and group owner.

Define r/w/x permissions for owner, group members, and everyone else (other)



setuid — program runs with permission of principle who installed it / owns it.

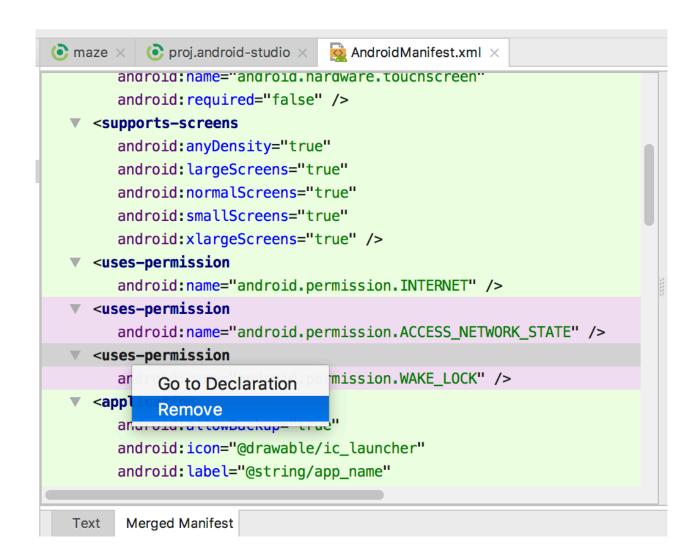
 $\rightarrow$  if owner == superuser, and setuid bit is set, then program executes with superuser privileges!

(Try it!) Use stat to look at file details

# **Smartphones & App Permissions**

There exists a mechanism for apps to declare needed permissions and request permissions.

E.g., AndroidManifest.xml





# Access Control looking forward...

### Basic Access Control

UNIX, ACL, various capability systems

### Aggregated Access Matrix

TE, RBAC, groups and attributes, parameterized

### Plus Domain Transitions

• DTE, SELinux, Java

### Lattice Access Control Models

Bell-LaPadula, Biba, Denning

### Predicate Models

ASL, OASIS, domain-specific models, many others

### Safety Models

· Take-grant, Schematic Protection Model, Typed Access Matrix